Oracle® Solaris Tunable Parameters Reference Manual

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Contents

Preface

The *Oracle Solaris Tunable Parameters Reference Manual* provides reference information about Oracle Solaris OS kernel and network tunable parameters. This manual does not provide tunable parameter information about desktop systems or Java environments.

This manual contains information for both SPARC based and x86 based systems.

Note – This Oracle Solaris release supports systems that use the SPARC and x86 families of processor architectures. The supported systems appear in the *Oracle Solaris Hardware Compatibility List* at `http://www.oracle.com/webfolder/technetwork/hcl/index.html`. This document cites any implementation differences between the platform types.

In this document these x86 terms mean the following:

- "x86" refers to the larger family of 64-bit and 32-bit x86 compatible products.
- "x64" relates specifically to 64-bit x86 compatible CPUs.
- "32-bit x86" points out specific 32-bit information about x86 based systems.

Who Should Use This Book

This book is intended for experienced Oracle Solaris system administrators who might need to change kernel tunable parameters in certain situations. For guidelines on changing Oracle Solaris tunable parameters, refer to "Tuning an Oracle Solaris System" on page 23.

How This Book Is Organized

The following table describes the chapters and appendixes in this book.

Chapter	Description
Chapter 1, "Overview of Oracle Solaris System Tuning"	An overview of tuning an Oracle Solaris system. Also provides a description of the format used in the book to describe the kernel tunables.

Chapter	Description
Chapter 2, "Oracle Solaris Kernel Tunable Parameters"	A description of Oracle Solaris kernel tunables such as kernel memory, file system, process size, and paging parameters.
Chapter 3, "NFS Tunable Parameters"	A description of NFS tunables such as caching symbolic links, dynamic retransmission, and RPC security parameters.
Chapter 4, "Internet Protocol Suite Tunable Parameters"	A description of TCP/IP tunables such as IP forwarding, source routing, and buffer-sizing parameters.
Chapter 5, "Network Cache and Accelerator Tunable Parameters"	A description of tunable parameters for the Network Cache and Accelerator (NCA).
Chapter 6, "System Facility Parameters"	A description of parameters used to set default values of certain system facilities. Changes are made by modifying files in the `/etc/default` directory.
Appendix A, "Tunable Parameters Change History"	A history of parameters that have changed or are now obsolete.
Appendix B, "Revision History for This Manual"	A history of this manual's revisions including the current Oracle Solaris release.

Other Resources for Oracle Solaris Tuning Information

This table describes other resources for Oracle Solaris tuning information.

Tuning Resource	For More Information
Online performance tuning information	`http://www.solarisinternals.com/si/index.php`
In-depth technical white papers	`http://www.oracle.com/technetwork/server-storage/solaris/overview/index.html`

Access to Oracle Support

Oracle customers have access to electronic support through My Oracle Support. For information, visit `http://www.oracle.com/pls/topic/lookup?ctx=acc&id=info` or visit `http://www.oracle.com/pls/topic/lookup?ctx=acc&id=trs` if you are hearing impaired.

Typographic Conventions

The following table describes the typographic conventions that are used in this book.

TABLE P–1 Typographic Conventions

Typeface	Description	Example
`AaBbCc123`	The names of commands, files, and directories, and onscreen computer output	Edit your `.login` file. Use `ls -a` to list all files. `machine_name% you have mail.`
`AaBbCc123`	What you type, contrasted with onscreen computer output	`machine_name%` **`su`** `Password:`
aabbcc123	Placeholder: replace with a real name or value	The command to remove a file is `rm` *filename*.
AaBbCc123	Book titles, new terms, and terms to be emphasized	Read Chapter 6 in the *User's Guide*. A *cache* is a copy that is stored locally. Do *not* save the file. **Note:** Some emphasized items appear bold online.

Shell Prompts in Command Examples

The following table shows the default UNIX system prompt and superuser prompt for shells that are included in the Oracle Solaris OS. Note that the default system prompt that is displayed in command examples varies, depending on the Oracle Solaris release.

TABLE P–2 Shell Prompts

Shell	Prompt
Bash shell, Korn shell, and Bourne shell	`$`
Bash shell, Korn shell, and Bourne shell for superuser	`#`
C shell	`machine_name%`
C shell for superuser	`machine_name#`

◆ ◆ ◆ CHAPTER 1

Overview of Oracle Solaris System Tuning

This section provides overview information about the format of the tuning information in this manual. This section also describes the different ways to tune an Oracle Solaris system.

What's New in Oracle Solaris System Tuning?

This section describes new or changed parameters in the Oracle Solaris 10 release.

- **Oracle Solaris 10 8/11**: The `rstchown` parameter that was previously set in the `/etc/system` file is obsolete. If you set this parameter in the `/etc/system` file, it is ignored.

  This parameter has been replaced by the ZFS `rstchown` file system property and a general file system mount option. For more information, see *Oracle Solaris ZFS Administration Guide* and `mount(1M)`.

- **Oracle Solaris 10 8/11**: This release includes the `ngroups_max` parameter description. For more information, see "`ngroups_max`" on page 46.

- **Solaris 10 10/09**: This release includes the `zfs_arc_min` and `zfs_arc_max` parameter descriptions. For more information, see "`zfs_arc_min`" on page 33 and "`zfs_arc_max`" on page 33.

  For additional information about tuning ZFS file systems, see the following site:

  `http://www.solarisinternals.com/wiki/index.php/ZFS_Evil_Tuning_Guide`

- **Solaris 10 10/09**: Memory locality group parameters are provided in this release. For more information about these parameters, see "Locality Group Parameters" on page 88.
- **Solaris 10 5/08**: The translation storage buffers parameters in the "SPARC System Specific Parameters" on page 84 section have been revised to provide better information. In this release, the following parameters have changed:
  - "default_tsb_size" on page 86
  - "enable_tsb_rss_sizing" on page 87
  - "tsb_rss_factor" on page 87
- **Solaris 10 8/07**: Parameter information was updated to include sun4v systems. For more information, see the following references:
  - "maxphys" on page 63
  - "tmpfs:tmpfs_maxkmem" on page 75
  - "SPARC System Specific Parameters" on page 84
- **Solaris 10 8/07**: The IP instances project enables you to configure a zone as an exclusive-IP zone and assign exclusive access of some LANs or VLANs to that zone.

  The previous behavior of shared-IP zones remains the default behavior. The exclusive-IP zone means that all aspects of the TCP/IP state and policy are per exclusive-IP zone, including TCP/IP tunable parameters.

  The introduction of the IP instances feature means that the following TCP parameters can only be set in the global zone because they require the `PRIV_SYS_NET_CONFIG` privilege:

  - "ip_squeue_fanout" on page 133
  - "ip_squeue_worker_wait" on page 149

  The other TCP, IP, and SCTP parameters and route metrics only require the `PRIV_SYS_IP_CONFIG` privilege. Each exclusive-IP zone controls its own set of these parameters. For shared-IP zones, TCP, IP, SCTP, and route parameters are controlled by the global zone since the settings of these parameters are shared between the global zone and all shared IP zones.

  For more information about using IP instances in Solaris zones, see *System Administration Guide: Oracle Solaris Containers-Resource Management and Oracle Solaris Zones*.

Oracle Solaris System Tuning in the Solaris 10 Release

This section describes significant tuning enhancements in the Oracle Solaris 10 release.

- "SCTP Tunable Parameters" on page 23

Default Stack Size

A new parameter, `default_stksize`, specifies the default stack size of all threads, kernel or user. The `lwp_default_stksize` parameter is still available, but it does not affect all kernel stacks. If `default_stksize` is set, it overrides `lwp_default_stksize`. For more information, see "`default_stksize`" on page 34.

System V IPC Configuration

In the Oracle Solaris 10 release, all System V IPC facilities are either automatically configured or can be controlled by resource controls. Facilities that can be shared are memory, message queues, and semaphores.

Resource controls allow IPC settings to be made on a per-project or per-process basis on the local system or in a name service environment.

In previous Solaris releases, IPC facilities were controlled by kernel tunables. You had to modify the `/etc/system` file and reboot the system to change the default values for these facilities.

Because the IPC facilities are now controlled by resource controls, their configuration can be modified while the system is running.

Many applications that previously required system tuning to function might now run without tuning because of increased defaults and the automatic allocation of resources.

The following table identifies the now obsolete IPC tunables and the possible resource controls that could be used as replacements. An important distinction between the obsolete IPC tunables and resource controls is that the IPC tunables were set on a system-wide basis and the resource controls are set on a per-project or per-process basis.

Resource Control	Obsolete Tunable	Old Default Value	Maximum Value	New Default Value
`process.max-msg-qbytes`	`msgsys:msginfo_msgmnb`	4096	`ULONG_MAX`	65536
`process.max-msg-messages`	`msgsys:msginfo_msgtql`	40	`UINT_MAX`	8192
`process.max-sem-ops`	`semsys:seminfo_semopm`	10	`INT_MAX`	512
`process.max-sem-nsems`	`semsys:seminfo_semmsl`	25	`SHRT_MAX`	512
`project.max-shm-memory`	`shmsys:shminfo_shmmax*`	0x800000	`UINT64_MAX`	1/4 of physical memory

Resource Control	Obsolete Tunable	Old Default Value	Maximum Value	New Default Value
project.max-shm-ids	shmsys:shminfo_shmmni	100	2^{24}	128
project.max-msg-ids	msgsys:msginfo_msgmni	50	2^{24}	128
project.max-sem-ids	semsys:seminfo_semmni	10	2^{24}	128

* Note that the project.max-shm-memory resource control limits the total amount of shared memory of one project, whereas previously, the shmsys:shminfo_shmmax parameter limited the size of a single shared memory segment.

For more detailed descriptions of the resource controls, see "Available Resource Controls" in *System Administration Guide: Oracle Solaris Containers-Resource Management and Oracle Solaris Zones*.

Obsolete parameters can still be included in the /etc/system file on an Oracle Solaris system. If so, the parameters are used to initialize the default resource control values as in previous Oracle Solaris releases. For more information, see "Parameters That Are Obsolete or Have Been Removed" on page 184. However, using the obsolete parameters is not recommended.

The following related parameters have been removed. If these parameters are included in the /etc/system file on an Oracle Solaris system, the parameters are commented out.

```
semsys:seminfo_semmns
semsys:seminfo_semmnu
semsys:seminfo_semume
semsys:seminfo_semmap
shmsys:shminfo_shmmin
msgsys:msginfo_msgseg
msgsys:msginfo_msgmax
semsys:seminfo_semvmx
semsys:seminfo_semaem
semsys:seminfo_semusz
shmsys:shminfo_shmseg
msgsys:msginfo_msgmap
msgsys:msginfo_msgssz
```

For the current list of available resource controls, see rctladm(1M). For information about configuring resource controls, see project(4), and Chapter 6, "Resource Controls (Overview)," in *System Administration Guide: Oracle Solaris Containers-Resource Management and Oracle Solaris Zones*.

NFSv4 Parameters

The following parameters for the NFSv4 protocol are included in the Oracle Solaris 10 release:

- "nfs:nfs4_pathconf_disable_cache" on page 94
- "nfs:nfs4_cots_timeo" on page 97
- "nfs:nfs4_do_symlink_cache" on page 99
- "nfs:nfs4_lookup_neg_cache" on page 102
- "nfs:nfs4_max_threads" on page 105
- "nfs:nfs4_nra" on page 107
- "nfs:nfs4_bsize" on page 112
- "nfs:nfs4_async_clusters" on page 115
- "nfs:nfs4_max_transfer_size" on page 119

For information about NFSv4 parameters, see "NFS Module Parameters" on page 94.

New and Changed TCP/IP Parameters

The following IP parameters are available in the Oracle Solaris 10 release:

- "ip_squeue_worker_wait" on page 149
- "ip_squeue_fanout" on page 133
- "ipcl_conn_hash_size" on page 148

The following TCP parameters are available in the Oracle Solaris 10 release:

- "tcp_rst_sent_rate_enabled" on page 145
- "tcp_rst_sent_rate" on page 146
- "tcp_mdt_max_pbufs" on page 146

The following TCP/IP parameters are obsolete in this Oracle Solaris release.

- ipc_tcp_conn_hash_size
- tcp_compression_enabled
- tcp_conn_hash_size
- ip_forwarding
- ip6_forwarding
- xxx_forwarding

IP Forwarding Changes

In this Oracle Solaris release, IP forwarding is enabled or disabled by using the `routeadm` command or the `ifconfig` commands instead of setting the following tunable parameters with the `ndd` command:

- ip_forwarding
- ip6_forwarding

- `xxx_forwarding`

Using the `routeadm` command and the `ifconfig` command instead of the `ndd` command to set IP forwarding provides the following advantages:

- All settings are persistent across reboots
- The new `ifconfig router` and `-router` commands can be placed in the `/etc/hostname.`*interface* files, along with other `ifconfig` commands that are run when the interface is initially configured.

To enable IPv4 or IPv6 packet forwarding on all interfaces of a system, you would use the following commands:

```
# routeadm -e ipv4-forwarding

# routeadm -e ipv6-forwarding
```

To disable IPv4 or IPv6 packet forwarding on all interfaces of a system, you would use the following commands:

```
# routeadm -d ipv4-forwarding

# routeadm -d ipv6-forwarding
```

In previous Solaris releases, you would enable IPv4 or IPv6 packet forwarding on all interfaces of a system as follows:

```
# ndd -set /dev/ip ip_forwarding 1

# ndd -set /dev/ip ip6_forwarding 1
```

In previous Solaris releases, you would disable IPv4 or IPv6 packet forwarding on all interfaces of a system as follows:

```
# ndd -set /dev/ip ip_forwarding 0

# ndd -set /dev/ip ip6_forwarding 0
```

If you want to enable IP forwarding on a specific IPv4 interface or IPv6 interface, you would use syntax similar to the following for your interface. The `bge0` interface is used an as example.

```
# ifconfig bge0 router

# ifconfig bge0 inet6 router
```

If you want to disable IP forwarding on a specific IPv4 interface or IPv6 interface, you would use syntax similar to the following for your interface. The `bge0` interface is used an as example.

```
# ifconfig bge0 -router
```

```
# ifconfig bge0 inet6 -router
```

Previously, IP forwarding was enabled on a specific interface as follows:

```
# ndd -set /dev/ip bge0:ip_forwarding 1
```

```
# ndd -set /dev/ip bge0:ip_forwarding 1
```

Previously, IP forwarding on a specific interface was disabled as follows:

```
# ndd -set /dev/ip ip_forwarding 0
```

```
# ndd -set /dev/ip ip6_forwarding 0
```

If you want any of the preceding `routeadm` settings to take effect on the running system, use the following command:

```
# routeadm -u
```

For more information, see `routeadm(1M)` and `ifconfig(1M)`.

SPARC: Translation Storage Buffer (TSB) Parameters

New parameters for tuning Translation Storage Buffer (TSB) are included in the Oracle Solaris 10 release. For information about TSB parameters, see "SPARC System Specific Parameters" on page 84.

SCTP Tunable Parameters

Stream Control Transmission Protocol (SCTP), a reliable transport protocol that provides services similar to the services provided by TCP, is provided in this Oracle Solaris release. For more information about SCTP tunable parameters, see "SCTP Tunable Parameters" on page 156.

Tuning an Oracle Solaris System

The Oracle Solaris OS is a multi-threaded, scalable UNIX operating system that runs on SPARC and x86 processors. It is self-adjusting to system load and demands minimal tuning. In some cases, however, tuning is necessary. This book provides details about the officially supported kernel tuning options available for the Oracle Solaris OS.

The Solaris kernel is composed of a core portion, which is always loaded, and a number of loadable modules that are loaded as references are made to them. Many variables referred to in the kernel portion of this guide are in the core portion. However, a few variables are located in loadable modules.

A key consideration in system tuning is that setting system parameters (or system variables) is often the least effective action that can be done to improve performance. Changing the behavior of the application is generally the most effective tuning aid available. Adding more physical memory and balancing disk I/O patterns are also useful. In a few rare cases, changing one of the variables described in this guide will have a substantial effect on system performance.

Remember that one system's `/etc/system` settings might not be applicable, either wholly or in part, to another system's environment. Carefully consider the values in the file with respect to the environment in which they will be applied. Make sure that you understand the behavior of a system before attempting to apply changes to the system variables that are described here.

We recommend that you start with an empty `/etc/system` file when moving to a new Oracle Solaris release. As a first step, add only those tunables that are required by in-house or third-party applications. Any tunables that involve System V IPC (semaphores, shared memory, and message queues) have been modified in the Oracle Solaris 10 release and should be changed in your environment. For more information, see "System V IPC Configuration" on page 19. After baseline testing has been established, evaluate system performance to determine if additional tunable settings are required.

Caution – The tunable parameters described in this book can and do change from Oracle Solaris release to Oracle Solaris release. Publication of these tunable parameters does not preclude changes to the tunable parameters and their descriptions without notice.

Tuning Format of Tunable Parameters Descriptions

The format for the description of each tunable parameter is as follows:

- Parameter Name
- Description
- Data Type
- Default
- Range
- Units
- Dynamic?
- Validation
- Implicit
- When to Change
- Zone Configuration
- Commitment Level
- Change History

Parameter Name — Is the exact name that is typed in the `/etc/system` file, or found in the `/etc/default/`*facility* file.

	Most parameters names are of the form *parameter* where the parameter name does not contain a colon (:). These names refer to variables in the core portion of the kernel. If the name does contain a colon, the characters to the left of the colon reference the name of a loadable module. The name of the parameter within the module consists of the characters to the right of the colon. For example: *module_name*:*variable*
Description	Briefly describes what the parameter does or controls.
Data Type	Indicates the signed or unsigned short integer or long integer with the following distinctions: ■ On a system that runs a 32-bit kernel, a long integer is the same size as an integer. ■ On a system that runs a 64-bit kernel, a long integer is twice the width in bits as an integer. For example, an unsigned integer = 32 bits, an unsigned long integer = 64 bits.
Data Type	Indicates the signed or unsigned short integer or long integer. A long integer is twice the width in bits as an integer. For example, an unsigned integer = 32 bits, an unsigned long integer = 64 bits.
Units	(Optional) Describes the unit type.
Default	What the system uses as the default value.
Range	Specifies the possible range allowed by system validation or the bounds of the data type. ■ MAXINT – A shorthand description for the maximum value of a signed integer (2,147,483,647) ■ MAXUINT – A shorthand description for the maximum value of an unsigned integer (4,294,967,295)
Dynamic?	Yes, if the parameter can be changed on a running system with the `mdb` or `kmdb` debugger. No, if the parameter is a boot time initialization only.
Validation	Checks that the system applies to the value of the variable either as specified in the `/etc/system` file or the default value, as well as when the validation is applied.
Implicit	(Optional) Provides unstated constraints that might exist on the parameter, especially in relation to other parameters.
When to Change	Explains why someone might want to change this value. Includes error messages or return codes.

Zone Configuration	Identifies whether the parameter can be set in a exclusive-IP zone or must be set in the global zone. None of the parameters can be set in shared-IP zones.
Commitment Level	Identifies the stability of the interface. Many of the parameters in this manual are still evolving and are classified as unstable. For more information, see `attributes(5)`.
Change History	(Optional) Contains a link to the Change History appendix, if applicable.

Tuning the Oracle Solaris Kernel

The following table describes the different ways tunable parameters can be applied.

Apply Tunable Parameters in These Ways	For More Information
Modify the `/etc/system` file	"`/etc/system` File" on page 26
Use the kernel debugger (`kmdb`)	"`kmdb` Command" on page 27
Use the modular debugger (`mdb`)	"`mdb` Command" on page 27
Use the `ndd` command to set TCP/IP parameters	Chapter 4, "Internet Protocol Suite Tunable Parameters"
Modify the `/etc/default` files	"Tuning NCA Parameters" on page 167

/etc/system File

The `/etc/system` file provides a static mechanism for adjusting the values of kernel parameters. Values specified in this file are read at boot time and are applied. Any changes that are made to the file are not applied to the operating system until the system is rebooted.

One pass is made to set all the values before the configuration parameters are calculated.

Example—Setting a Parameter in /etc/system

The following `/etc/system` entry sets the ZFS ARC maximum (`zfs_arc_max`) to 30 GB.

```
set zfs:zfs_arc_max = 0x780000000
```

Recovering From an Incorrect Value

Make a copy of the `/etc/system` file before modifying it so that you can easily recover from incorrect value. For example:

```
# cp /etc/system /etc/system.good
```

If a value specified in the `/etc/system` file causes the system to become unbootable, you can recover with the following command:

```
ok boot -a
```

This command causes the system to ask for the name of various files used in the boot process. Press the Return key to accept the default values until the name of the `/etc/system` file is requested. When the `Name of system file [/etc/system]:` prompt is displayed, type the name of the good `/etc/system` file or `/dev/null`:

```
Name of system file [/etc/system]: /etc/system.good
```

If `/dev/null` is specified, this path causes the system to attempt to read from `/dev/null` for its configuration information. Because this file is empty, the system uses the default values. After the system is booted, the `/etc/system` file can be corrected.

For more information on system recovery, see *System Administration Guide: Basic Administration*.

kmdb Command

`kmdb` is a interactive kernel debugger with the same general syntax as `mdb`. An advantage of interactive kernel debugger is that you can set breakpoints. When a breakpoint is reached, you can examine data or step through the execution of kernel code.

`kmdb` can be loaded and unloaded on demand. You do not have to reboot the system to perform interactive kernel debugging, as was the case with `kadb`.

For more information, see `kmdb(1)`.

mdb Command

The modular debugger, `mdb`, is unique among Solaris debuggers because it is easily extensible. A programming API is available that allows compilation of modules to perform desired tasks within the context of the debugger.

`mdb` also includes a number of desirable usability features, including command-line editing, command history, built-in output pager, syntax checking, and command pipelining. `mdb` is the recommended post-mortem debugger for the kernel.

For more information, see mdb(1).

Example–Using mdb to Display Information

Display a high-level view of a system's memory usage. For example:

```
# mdb -k
Loading modules: [ unix genunix specfs dtrace zfs sd pcisch sockfs ip hook neti sctp arp
usba fcp fctl md lofs cpc random crypto fcip nca logindmux ptm ufs sppp nfs ]
> ::memstat
Page Summary                Pages                MB  %Tot
------------     ----------------  ----------------  ----
Kernel                      95193               743   37%
ZFS File Data               96308               752   38%
Anon                        28132               219   11%
Exec and libs                1870                14    1%
Page cache                   1465                11    1%
Free (cachelist)             4242                33    2%
Free (freelist)             28719               224   11%

Total                      255929              1999
Physical                   254495              1988
> $q
```

For more information on using the modular debugger, see the *Solaris Modular Debugger Guide*.

When using either `kmdb` or `mdb` debugger, the module name prefix is not required. After a module is loaded, its symbols form a common name space with the core kernel symbols and any other previously loaded module symbols.

For example, `ufs:ufs_WRITES` would be accessed as `ufs_WRITES` in each debugger (assuming the UFS module is loaded). The `ufs:` prefix is required when set in the `/etc/system` file.

Special Oracle Solaris tune and var Structures

Oracle Solaris tunable parameters come in a variety of forms. The `tune` structure defined in the`/usr/include/sys/tuneable.h` file is the runtime representation of `tune_t_fsflushr`, `tune_t_minarmem`, and `tune_t_flkrec`. After the kernel is initialized, all references to these variables are found in the appropriate field of the `tune` structure.

The proper way to set parameters for this structure at boot time is to initialize the special parameter that corresponds to the desired field name. The system initialization process then loads these values into the `tune` structure.

A second structure into which various tunable parameters are placed is the `var` structure named `v`. You can find the definition of a `var` structure in the `/usr/include/sys/var.h` file. The runtime representation of variables such as `autoup` and `bufhwm` is stored here.

Do not change either the `tune` or `v` structure on a running system. Changing any field in these structures on a running system might cause the system to panic.

Viewing Oracle Solaris System Configuration Information

Several tools are available to examine system configuration information. Some tools require superuser privilege. Other tools can be run by a non-privileged user. Every structure and data item can be examined with the kernel debugger by using `mdb` on a running system or by booting under `kmdb`.

For more information, see `mdb(1)` or `kadb(1M)`.

sysdef Command

The `sysdef` command provides the values of memory and process resource limits, and portions of the `tune` and `v` structures. For example, the `sysdef` "Tunable Parameters" section from a SPARC system with 16 GB of memory is as follows:

```
20840448        maximum memory allowed in buffer cache (bufhwm)
   15898        maximum number of processes (v.v_proc)
      99        maximum global priority in sys class (MAXCLSYSPRI)
   15893        maximum processes per user id (v.v_maxup)
      30        auto update time limit in seconds (NAUTOUP)
      25        page stealing low water mark (GPGSLO)
       1        fsflush run rate (FSFLUSHR)
      25        minimum resident memory for avoiding deadlock (MINARMEM)
      25        minimum swapable memory for avoiding deadlock (MINASMEM)
```

For more information, see `sysdef(1M)`.

kstat Utility

`kstats` are data structures maintained by various kernel subsystems and drivers. They provide a mechanism for exporting data from the kernel to user programs without requiring that the program read kernel memory or have superuser privilege. For more information, see `kstat(1M)` or `kstat(3KSTAT)`.

◆ ◆ ◆ CHAPTER 2

Oracle Solaris Kernel Tunable Parameters

This chapter describes most of the Oracle Solaris kernel tunable parameters.

Where to Find Tunable Parameter Information

Tunable Parameter	For Information
NFS tunable parameters	Chapter 3, "NFS Tunable Parameters"
Internet Protocol Suite tunable parameters	Chapter 4, "Internet Protocol Suite Tunable Parameters"
Network Cache and Accelerator (NCA) tunable parameters	Chapter 5, "Network Cache and Accelerator Tunable Parameters"

General Kernel and Memory Parameters

This section describes general kernel parameters that are related to physical memory and stack configuration.

physmem

Description	Modifies the system's configuration of the number of physical pages of memory after the Oracle Solaris OS and firmware are accounted for.
Data Type	Unsigned long
Default	Number of usable pages of physical memory available on the system, not counting the memory where the core kernel and data are stored
Range	1 to amount of physical memory on system
Units	Pages
Dynamic?	No
Validation	None
When to Change	Whenever you want to test the effect of running the system with less physical memory. Because this parameter does *not* take into account the memory used by the core kernel and data, as well as various other data structures allocated early in the startup process, the value of `physmem` should be less than the actual number of pages that represent the smaller amount of memory.
Commitment Level	Unstable

zfs_arc_min

Description	Determines the minimum size of the ZFS Adaptive Replacement Cache (ARC). See also "zfs_arc_max" on page 33.
Data Type	Unsigned Integer (64-bit)
Default	1/32nd of physical memory or 64 MB, whichever value is larger.
Range	64 MB to zfs_arc_max
Units	Bytes
Dynamic?	No
Validation	Yes, the range is validated.
When to Change	When a system's workload demand for memory fluctuates, the ZFS ARC caches data at a period of weak demand and then shrinks at a period of strong demand. However, ZFS does not shrink below the value of zfs_arc_min. The default value of zfs_arc_min is 12% of memory on large memory systems and so, can be a significant amount of memory. If a workload's highest memory usage requires more than 88% of system memory, consider tuning this parameter.
Commitment Level	Unstable
Change History	For information, see "zfs_arc_min (Solaris 10 Releases)" on page 180.

zfs_arc_max

Description	Determines the maximum size of the ZFS Adaptive Replacement Cache (ARC). See also "zfs_arc_min" on page 33.
Data Type	Unsigned Integer (64-bit)
Default	Three-fourths of memory on systems with less than 4 GB of memory physmem minus 1 GB on systems with greater than 4 GB of memory
Range	64 MB to physmem
Units	Bytes
Dynamic?	No
Validation	Yes, the range is validated.
When to Change	If a future memory requirement is significantly large and well defined, you might consider reducing the value of this parameter to cap the ARC so that it does not compete with the memory requirement. For

example, if you know that a future workload requires 20% of memory, it makes sense to cap the ARC such that it does not consume more than the remaining 80% of memory.

Commitment Level: Unstable

Change History: For information, see "zfs_arc_max (Solaris 10 Releases)" on page 180.

default_stksize

Description: Specifies the default stack size of all threads. No thread can be created with a stack size smaller than `default_stksize`. If `default_stksize` is set, it overrides `lwp_default_stksize`. See also "lwp_default_stksize" on page 35.

Data Type: Integer

Default:

- 3 x `PAGESIZE` on SPARC systems
- 2 x `PAGESIZE` on x86 systems
- 5 x `PAGESIZE` on AMD64 systems

Range: Minimum is the default values:

- 3 x `PAGESIZE` on SPARC systems
- 2 x `PAGESIZE` on x86 systems

Maximum is 32 times the default value.

Units: Bytes in multiples of the value returned by the `getpagesize` parameter. For more information, see `getpagesize(3C)`.

Dynamic?: Yes. Affects threads created after the variable is changed.

Validation: Must be greater than or equal to 8192 and less than or equal to 262,144 (256 x 1024). Also must be a multiple of the system page size. If these conditions are not met, the following message is displayed:

```
Illegal stack size, Using N
```

The value of *N* is the default value of `default_stksize`.

When to Change: When the system panics because it has run out of stack space. The best solution for this problem is to determine why the system is running out of space and then make a correction.

Increasing the default stack size means that almost every kernel thread will have a larger stack, resulting in increased kernel memory consumption for no good reason. Generally, that space will be unused.

The increased consumption means other resources that are competing for the same pool of memory will have the amount of space available to them reduced, possibly decreasing the system's ability to perform work. Among the side effects is a reduction in the number of threads that the kernel can create. This solution should be treated as no more than an interim workaround until the root cause is remedied.

Commitment Level: Unstable

lwp_default_stksize

Description: Specifies the default value of the stack size to be used when a kernel thread is created, and when the calling routine does not provide an explicit size to be used.

Data Type: Integer

Default:

- 8192 for x86 platforms
- 24,576 for SPARC platforms
- 20,480 for AMD64 platforms

Range: Minimum is the default values:

- 3 x `PAGESIZE` on SPARC systems
- 2 x `PAGESIZE` on x86 systems
- 5 x `PAGESIZE` on AMD64 systems

Maximum is 32 times the default value.

Units: Bytes in multiples of the value returned by the `getpagesize` parameter. For more information, see `getpagesize(3C)`.

Dynamic?: Yes. Affects threads created after the variable is changed.

Validation: Must be greater than or equal to 8192 and less than or equal to 262,144 (256 x 1024). Also must be a multiple of the system page size. If these conditions are not met, the following message is displayed:

`Illegal stack size, Using` *N*

The value of *N* is the default value of `lwp_default_stksize`.

When to Change: When the system panics because it has run out of stack space. The best solution for this problem is to determine why the system is running out of space and then make a correction.

	Increasing the default stack size means that almost every kernel thread will have a larger stack, resulting in increased kernel memory consumption for no good reason. Generally, that space will be unused. The increased consumption means other resources that are competing for the same pool of memory will have the amount of space available to them reduced, possibly decreasing the system's ability to perform work. Among the side effects is a reduction in the number of threads that the kernel can create. This solution should be treated as no more than an interim workaround until the root cause is remedied.
Commitment Level	Unstable
Change History	For information, see "`lwp_default_stksize` (Solaris 10 Releases)" on page 180.

logevent_max_q_sz

Description	Maximum number of system events allowed to be queued and waiting for delivery to the `syseventd` daemon. Once the size of the system event queue reaches this limit, no other system events are allowed on the queue.
Data Type	Integer
Default	5000
Range	0 to MAXINT
Units	System events
Dynamic?	Yes
Validation	The system event framework checks this value every time a system event is generated by `ddi_log_sysevent` and `sysevent_post_event`. For more information, see `ddi_log_sysevent(9F)` and `sysevent_post_event(3SYSEVENT)`.
When to Change	When error log messages indicate that a system event failed to be logged, generated, or posted.
Commitment Level	Unstable

segkpsize

Description	Specifies the amount of kernel pageable memory available. This memory is used primarily for kernel thread stacks. Increasing this

	number allows either larger stacks for the same number of threads or more threads. This parameter can only be set on a system running a 64-bit kernel. A system running a 64-bit kernel uses a default stack size of 24 KB.
Data Type	Unsigned long
Default	64-bit kernels, 2 GB 32-bit kernels, 512 MB
Range	64-bit kernels, 512 MB to 24 GB
Units	8-KB pages
Dynamic?	No
Validation	Value is compared to minimum and maximum sizes (512 MB and 24 GB for 64-bit systems). If smaller than the minimum or larger than the maximum, it is reset to 2 GB. A message to that effect is displayed. The actual size used in creation of the cache is the lesser of the value specified in `segkpsize` after the validation checking or 50 percent of physical memory.
When to Change	Required to support large numbers of processes on a system. The default size of 2 GB, assuming at least 1 GB of physical memory is present. This default size allows creation of 24-KB stacks for more than 87,000 kernel threads. The size of a stack in a 64-bit kernel is the same, whether the process is a 32-bit process or a 64-bit process. If more than this number is needed, `segkpsize` can be increased, assuming sufficient physical memory exists.
Commitment Level	Unstable

noexec_user_stack

Description	Enables the stack to be marked as nonexecutable, which helps make buffer-overflow attacks more difficult. an Oracle Solaris system running a 64-bit kernel makes the stacks of all 64-bit applications nonexecutable by default. Setting this parameter is necessary to make 32-bit applications nonexecutable on systems running 64-bit or 32-bit kernels.

Note – This parameter is only effective on 64–bit SPARC and AMD64 architectures.

Data Type	Signed integer
Default	0 (disabled)
Range	0 (disabled) or 1 (enabled)
Units	Toggle (on/off)
Dynamic?	Yes. Does not affect currently running processes, only processes created after the value is set.
Validation	None
When to Change	Should be enabled at all times unless applications are deliberately placing executable code on the stack without using `mprotect` to make the stack executable. For more information, see `mprotect(2)`.
Commitment Level	Unstable
Change History	For information, see "`noexec_user_stack` (Solaris 10 Releases)" on page 180.

fsflush and Related Parameters

This section describes `fsflush` and related tunables.

fsflush

The system daemon, `fsflush`, runs periodically to do three main tasks:

1. On every invocation, `fsflush` flushes dirty file system pages over a certain age to disk.
2. On every invocation, `fsflush` examines a portion of memory and causes modified pages to be written to their backing store. Pages are written if they are modified and if they do not meet one of the following conditions:
   - Pages are kernel page
   - Pages are free
   - Pages are locked
   - Pages are associated with a swap device
   - Pages are currently involved in an I/O operation

The net effect is to flush pages from files that are mapped with `mmap` with write permission and that have actually been changed.

Pages are flushed to backing store but left attached to the process using them. This will simplify page reclamation when the system runs low on memory by avoiding delay for writing the page to backing store before claiming it, if the page has not been modified since the flush.

3. `fsflush` writes file system metadata to disk. This write is done every *n*th invocation, where *n* is computed from various configuration variables. See "`tune_t_fsflushr`" on page 39 and "`autoup`" on page 40 for details.

The following features are configurable:

- Frequency of invocation (`tune_t_fsflushr`)
- Whether memory scanning is executed (`dopageflush`)
- Whether file system data flushing occurs (`doiflush`)
- The frequency with which file system data flushing occurs (`autoup`)

For most systems, memory scanning and file system metadata synchronizing are the dominant activities for `fsflush`. Depending on system usage, memory scanning can be of little use or consume too much CPU time.

tune_t_fsflushr

Description	Specifies the number of seconds between `fsflush` invocations
Data Type	Signed integer
Default	1
Range	1 to MAXINT
Units	Seconds
Dynamic?	No
Validation	If the value is less than or equal to zero, the value is reset to 1 and a warning message is displayed. This check is done only at boot time.
When to Change	See the `autoup` parameter.
Commitment Level	Unstable

autoup

Description	Along with `tune_t_flushr`, `autoup` controls the amount of memory examined for dirty pages in each invocation and frequency of file system synchronizing operations. The value of `autoup` is also used to control whether a buffer is written out from the free list. Buffers marked with the `B_DELWRI` flag (which identifies file content pages that have changed) are written out whenever the buffer has been on the list for longer than *autoup* seconds. Increasing the value of `autoup` keeps the buffers in memory for a longer time.
Data Type	Signed integer
Default	30
Range	1 to MAXINT
Units	Seconds
Dynamic?	No
Validation	If `autoup` is less than or equal to zero, it is reset to 30 and a warning message is displayed. This check is done only at boot time.
Implicit	`autoup` should be an integer multiple of `tune_t_fsflushr`. At a minimum, `autoup` should be at least 6 times the value of `tune_t_fsflushr`. If not, excessive amounts of memory are scanned each time `fsflush` is invoked. The total system pages multiplied by `tune_t_fsflushr` should be greater than or equal to `autoup` to cause memory to be checked if `dopageflush` is non-zero.
When to Change	Here are several potential situations for changing `autoup`, `tune_t_fsflushr`, or both: ■ Systems with large amounts of memory – In this case, increasing `autoup` reduces the amount of memory scanned in each invocation of `fsflush`. ■ Systems with minimal memory demand – Increasing both `autoup` and `tune_t_fsflushr` reduces the number of scans made. `autoup` should be increased also to maintain the current ratio of `autoup` / `tune_t_fsflushr`.

- Systems with large numbers of transient files (for example, mail servers or software build machines) – If large numbers of files are created and then deleted, `fsflush` might unnecessarily write data pages for those files to disk.

Commitment Level: Unstable

dopageflush

Description: Controls whether memory is examined for modified pages during `fsflush` invocations. In each invocation of `fsflush`, the number of physical memory pages in the system is determined. This number might have changed because of a dynamic reconfiguration operation. Each invocation scans by using this algorithm: total number of pages x `tune_t_fsflushr` / `autoup` pages

Data Type: Signed integer

Default: 1 (enabled)

Range: 0 (disabled) or 1 (enabled)

Units: Toggle (on/off)

Dynamic?: Yes

Validation: None

When to Change: If the system page scanner rarely runs, which is indicated by a value of 0 in the `sr` column of `vmstat` output.

Commitment Level: Unstable

Change History: For information, see "`dopageflush` (Solaris 10 Releases)" on page 180.

doiflush

Description: Controls whether file system metadata syncs will be executed during `fsflush` invocations. This synchronization is done every *N*th invocation of `fsflush` where *N*= `(autoup / tune_t_fsflushr)`. Because this algorithm is integer division, if `tune_t_fsflushr` is greater than `autoup`, a synchronization is done on every invocation of `fsflush` because the code checks to see if its iteration counter is greater than or equal to *N*. Note that *N* is computed once on invocation of `fsflush`. Later changes to `tune_t_fsflushr` or `autoup` have no effect on the frequency of synchronization operations.

Data Type	Signed integer
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Units	Toggle (on/off)
Dynamic?	Yes
Validation	None
When to Change	When files are frequently modified over a period of time and the load caused by the flushing perturbs system behavior. Files whose existence, and therefore consistency of state, does not matter if the system reboots are better kept in a TMPFS file system (for example, `/tmp`). Inode traffic can be reduced on systems by using the `mount -noatime` option. This option eliminates inode updates when the file is accessed. For a system engaged in realtime processing, you might want to disable this option and use explicit application file synchronizing to achieve consistency.
Commitment Level	Unstable

Process-Sizing Parameters

Several parameters (or variables) are used to control the number of processes that are available on the system and the number of processes that an individual user can create. The foundation parameter is `maxusers`. This parameter drives the values assigned to `max_nprocs` and `maxuprc`.

maxusers

Description	Originally, `maxusers` defined the number of logged in users the system could support. When a kernel was generated, various tables were sized based on this setting. Current Oracle Solaris releases do much of its sizing based on the amount of memory on the system. Thus, much of the past use of `maxusers` has changed. A number of subsystems that are still derived from `maxusers`: ■ The maximum number of processes on the system ■ The number of quota structures held in the system ■ The size of the directory name look-up cache (DNLC)

Data Type	Signed integer
Default	Lesser of the amount of memory in MB or 2048
Range	1 to 2048, based on physical memory if not set in the `/etc/system` file 1 to 4096, if set in the `/etc/system` file
Units	Users
Dynamic?	No. After computation of dependent parameters is done, `maxusers` is never referenced again.
Validation	None
When to Change	When the default number of user processes derived by the system is too low. This situation is evident when the following message displays on the system console: `out of processes` You might also change this parameter when the default number of processes is too high, as in these situations: ■ Database servers that have a lot of memory and relatively few running processes can save system memory when the default value of `maxusers` is reduced. ■ If file servers have a lot of memory and few running processes, you might reduce this value. However, you should explicitly set the size of the DNLC. See "`ncsize`" on page 65. ■ If compute servers have a lot of memory and few running processes, you might reduce this value.
Commitment Level	Unstable

reserved_procs

Description	Specifies the number of system process slots to be reserved in the process table for processes with a UID of root (0). For example, `fsflush` has a UID of root (0).
Data Type	Signed integer
Default	5
Range	5 to MAXINT
Units	Processes
Dynamic?	No. Not used after the initial parameter computation.

Validation	Any `/etc/system` setting is honored.
Commitment Level	Unstable
When to Change	Consider increasing to 10 + the normal number of UID 0 (root) processes on system. This setting provides some cushion should it be necessary to obtain a root shell when the system is otherwise unable to create user-level processes.

pidmax

Description	Specifies the value of the largest possible process ID. `pidmax` sets the value for the `maxpid` variable. Once `maxpid` is set, `pidmax` is ignored. `maxpid` is used elsewhere in the kernel to determine the maximum process ID and for validation checking. Any attempts to set `maxpid` by adding an entry to the `/etc/system` file have no effect.
Data Type	Signed integer
Default	30,000
Range	266 to 999,999
Units	Processes
Dynamic?	No. Used only at boot time to set the value of `pidmax`.
Validation	Yes. Value is compared to the value of `reserved_procs` and 999,999. If less than `reserved_procs` or greater than 999,999, the value is set to 999,999.
Implicit	`max_nprocs` range checking ensures that `max_nprocs` is always less than or equal to this value.
When to Change	Required to enable support for more than 30,000 processes on a system.
Commitment Level	Unstable

max_nprocs

Description	Specifies the maximum number of processes that can be created on a system. Includes system processes and user processes. Any value specified in `/etc/system` is used in the computation of `maxuprc`.

This value is also used in determining the size of several other system data structures. Other data structures where this parameter plays a role are as follows:

- Determining the size of the directory name lookup cache (if `ncsize` is not specified)
- Allocating disk quota structures for UFS (if `ndquot` is not specified)
- Verifying that the amount of memory used by configured system V semaphores does not exceed system limits
- Configuring Hardware Address Translation resources for x86 platforms.

Data Type	Signed integer
Default	10 + (16 x `maxusers`)
Range	266 to value of `maxpid`
Dynamic?	No
Validation	Yes. The value is compared to `maxpid` and set to `maxpid` if it is larger. On x86 platforms, an additional check is made against a platform-specific value. `max_nprocs` is set to the smallest value in the triplet (`max_nprocs`, `maxpid`, platform value). Both SPARC and x86 platforms use 65,534 as the platform value.
When to Change	Changing this parameter is one of the steps necessary to enable support for more than 30,000 processes on a system.
Commitment Level	Unstable
Change History	For information, see "`max_nprocs` (Solaris 10 Releases)" on page 179.

maxuprc

Description	Specifies the maximum number of processes that can be created on a system by any one user.
Data Type	Signed integer
Default	`max_nprocs` - `reserved_procs`
Range	1 to `max_nprocs` - `reserved_procs`
Units	Processes
Dynamic?	No

Validation	Yes. This value is compared to `max_nprocs` - `reserved_procs` and set to the smaller of the two values.
When to Change	When you want to specify a hard limit for the number of processes a user can create that is less than the default value of however many processes the system can create. Attempting to exceed this limit generates the following warning messages on the console or in the messages file: `out of per-user processes for uid` *N*
Commitment Level	Unstable

ngroups_max

Description	Specifies the maximum number of supplemental groups per process.
Data Type	Signed integer
Default	16
Range	0 to 1024
Units	Groups
Dynamic?	No
Validation	No
When to Change	When you want to increase the maximum number of groups. Keep in mind that if a particular user is assigned to more than 16 groups, the user might experience problems with `AUTH_SYS` credentials in an NFS environment.
Commitment Level	Unstable

Paging-Related Parameters

The Solaris OS uses a demand paged virtual memory system. As the system runs, pages are brought into memory as needed. When memory becomes occupied above a certain threshold and demand for memory continues, paging begins. Paging goes through several levels that are controlled by certain parameters.

The general paging algorithm is as follows:

- A memory deficit is noticed. The page scanner thread runs and begins to walk through memory. A two-step algorithm is employed:

1. A page is marked as unused.
2. If still unused after a time interval, the page is viewed as a subject for reclaim.

   If the page has been modified, a request is made to the pageout thread to schedule the page for I/O. Also, the page scanner continues looking at memory. Pageout causes the page to be written to the page's backing store and placed on the free list. When the page scanner scans memory, no distinction is made as to the origin of the page. The page might have come from a data file, or it might represent a page from an executable's text, data, or stack.

- As memory pressure on the system increases, the algorithm becomes more aggressive in the pages it will consider as candidates for reclamation and in how frequently the paging algorithm runs. (For more information, see "`fastscan`" on page 54 and "`slowscan`" on page 54.) As available memory falls between the range `lotsfree` and `minfree`, the system linearly increases the amount of memory scanned in each invocation of the pageout thread from the value specified by `slowscan` to the value specified by `fastscan`. The system uses the `desfree` parameter to control a number of decisions about resource usage and behavior.

The system initially constrains itself to use no more than 4 percent of one CPU for `pageout` operations. As memory pressure increases, the amount of CPU time consumed in support of `pageout` operations linearly increases until a maximum of 80 percent of one CPU is consumed. The algorithm looks through some amount of memory between `slowscan` and `fastscan`, then stops when one of the following occurs:

- Enough pages have been found to satisfy the memory shortfall.
- The planned number of pages have been looked at.
- Too much time has elapsed.

If a memory shortfall is still present when pageout finishes its scan, another scan is scheduled for 1/4 second in the future.

The configuration mechanism of the paging subsystem was changed. Instead of depending on a set of predefined values for `fastscan`, `slowscan`, and `handspreadpages`, the system determines the appropriate settings for these parameters at boot time. Setting any of these parameters in the `/etc/system` file can cause the system to use less than optimal values.

Caution – Remove all tuning of the VM system from the `/etc/system` file. Run with the default settings and determine if it is necessary to adjust any of these parameters. Do not set either `cachefree` or `priority_paging`.

Dynamic reconfiguration (DR) for CPU and memory is supported. A system in a DR operation that involves the addition or deletion of memory recalculates values for the relevant parameters, unless the parameter has been explicitly set in `/etc/system`. In that case, the value specified in `/etc/system` is used, unless a constraint on the value of the variable has been violated. In this case, the value is reset.

lotsfree

Description	Serves as the initial trigger for system paging to begin. When this threshold is crossed, the page scanner wakes up to begin looking for memory pages to reclaim.
Data Type	Unsigned long
Default	The greater of 1/64th of physical memory or 512 KB
Range	The minimum value is 512 KB or 1/64th of physical memory, whichever is greater, expressed as pages using the page size returned by `getpagesize`. For more information, see`getpagesize(3C)`. The maximum value is the number of physical memory pages. The maximum value should be no more than 30 percent of physical memory. The system does not enforce this range, other than that described in the Validation section.
Units	Pages
Dynamic?	Yes, but dynamic changes are lost if a memory-based DR operation occurs.
Validation	If `lotsfree` is greater than the amount of physical memory, the value is reset to the default.
Implicit	The relationship of `lotsfree` being greater than `desfree`, which is greater than `minfree`, should be maintained at all times.
When to Change	When demand for pages is subject to sudden sharp spikes, the memory algorithm might be unable to keep up with demand. One workaround is to start reclaiming memory at an earlier time. This solution gives the paging system some additional margin. A rule of thumb is to set this parameter to 2 times what the system needs to allocate in a few seconds. This parameter is workload dependent. A DBMS server can probably work fine with the default settings. However, you might need to adjust this parameter for a system doing heavy file system I/O. For systems with relatively static workloads and large amounts of memory, lower this value. The minimum acceptable value is 512 KB, expressed as pages using the page size returned by `getpagesize`.
Commitment Level	Unstable

desfree

Description	Specifies the preferred amount of memory to be free at all times on the system.
Data Type	Unsigned integer
Default	`lotsfree` / 2
Range	The minimum value is 256 KB or 1/128th of physical memory, whichever is greater, expressed as pages using the page size returned by `getpagesize`. The maximum value is the number of physical memory pages. The maximum value should be no more than 15 percent of physical memory. The system does not enforce this range other than that described in the Validation section.
Units	Pages
Dynamic?	Yes, unless dynamic reconfiguration operations that add or delete memory occur. At that point, the value is reset to the value provided in the `/etc/system` file or calculated from the new physical memory value.
Validation	If `desfree` is greater than `lotsfree`, `desfree` is set to `lotsfree` / 2. No message is displayed.
Implicit	The relationship of `lotsfree` being greater than `desfree`, which is greater than `minfree`, should be maintained at all times.
Side Effects	Several side effects can arise from increasing the value of this parameter. When the new value nears or exceeds the amount of available memory on the system, the following can occur: ■ Asynchronous I/O requests are not processed, unless available memory exceeds `desfree`. Increasing the value of `desfree` can result in rejection of requests that otherwise would succeed. ■ NFS asynchronous writes are executed as synchronous writes. ■ The swapper is awakened earlier, and the behavior of the swapper is biased towards more aggressive actions. ■ The system might not prefault as many executable pages into the system. This side effect results in applications potentially running slower than they otherwise would.
When to Change	For systems with relatively static workloads and large amounts of memory, lower this value. The minimum acceptable value is 256 KB, expressed as pages using the page size returned by `getpagesize`.

Commitment Level	Unstable

minfree

Description	Specifies the minimum acceptable memory level. When memory drops below this number, the system biases allocations toward allocations necessary to successfully complete pageout operations or to swap processes completely out of memory. Either allocation denies or blocks other allocation requests.
Data Type	Unsigned integer
Default	`desfree` / 2
Range	The minimum value is 128 KB or 1/256th of physical memory, whichever is greater, expressed as pages using the page size returned by `getpagesize`. The maximum value is the number of physical memory pages. The maximum value should be no more than 7.5 percent of physical memory. The system does not enforce this range other than that described in the Validation section.
Units	Pages
Dynamic?	Yes, unless dynamic reconfiguration operations that add or delete memory occur. At that point, the value is reset to the value provided in the `/etc/system` file or calculated from the new physical memory value.
Validation	If `minfree` is greater than `desfree`, `minfree` is set to `desfree` / 2. No message is displayed.
Implicit	The relationship of `lotsfree` being greater than `desfree`, which is greater than `minfree`, should be maintained at all times.
When to Change	The default value is generally adequate. For systems with relatively static workloads and large amounts of memory, lower this value. The minimum acceptable value is 128 KB, expressed as pages using the page size returned by `getpagesize`.
Commitment Level	Unstable

throttlefree

Description	Specifies the memory level at which blocking memory allocation requests are put to sleep, even if the memory is sufficient to satisfy the request.
Data Type	Unsigned integer
Default	`minfree`
Range	The minimum value is 128 KB or 1/256th of physical memory, whichever is greater, expressed as pages using the page size returned by `getpagesize`. The maximum value is the number of physical memory pages. The maximum value should be no more than 4 percent of physical memory. The system does not enforce this range other than that described in the Validation section.
Units	Pages
Dynamic?	Yes, unless dynamic reconfiguration operations that add or delete memory occur. At that point, the value is reset to the value provided in the `/etc/system` file or calculated from the new physical memory value.
Validation	If `throttlefree` is greater than `desfree`, `throttlefree` is set to `minfree`. No message is displayed.
Implicit	The relationship of `lotsfree` is greater than `desfree`, which is greater than `minfree`, should be maintained at all times.
When to Change	The default value is generally adequate. For systems with relatively static workloads and large amounts of memory, lower this value. The minimum acceptable value is 128 KB, expressed as pages using the page size returned by `getpagesize`. For more information, see `getpagesize(3C)`.
Commitment Level	Unstable

pageout_reserve

Description	Specifies the number of pages reserved for the exclusive use of the pageout or scheduler threads. When available memory is less than this value, nonblocking allocations are denied for any processes other than pageout or the scheduler. Pageout needs to have a small pool of

	memory for its use so it can allocate the data structures necessary to do the I/O for writing a page to its backing store.
Data Type	Unsigned integer
Default	`throttlefree` / 2
Range	The minimum value is 64 KB or 1/512th of physical memory, whichever is greater, expressed as pages using the page size returned by `getpagesize(3C)`. The maximum is the number of physical memory pages. The maximum value should be no more than 2 percent of physical memory. The system does not enforce this range, other than that described in the Validation section.
Units	Pages
Dynamic?	Yes, unless dynamic reconfiguration operations that add or delete memory occur. At that point, the value is reset to the value provided in the `/etc/system` file or calculated from the new physical memory value.
Validation	If `pageout_reserve` is greater than `throttlefree` / 2, `pageout_reserve` is set to `throttlefree` / 2. No message is displayed.
Implicit	The relationship of `lotsfree` being greater than `desfree`, which is greater than `minfree`, should be maintained at all times.
When to Change	The default value is generally adequate. For systems with relatively static workloads and large amounts of memory, lower this value. The minimum acceptable value is 64 KB, expressed as pages using the page size returned by `getpagesize`.
Commitment Level	Unstable

pages_pp_maximum

Description	Defines the number of pages that must be unlocked. If a request to lock pages would force available memory below this value, that request is refused.
Data Type	Unsigned long
Default	The greater of (`tune_t_minarmem` + 100 and [4% of memory available at boot time + 4 MB])
Range	Minimum value enforced by the system is `tune_t_minarmem` + 100. The system does not enforce a maximum value.

Units	Pages
Dynamic?	Yes, unless dynamic reconfiguration operations that add or delete memory occur. At that point, the value is reset to the value provided in the `/etc/system` file or was calculated from the new physical memory value.
Validation	If the value specified in the `/etc/system` file or the calculated default is less than `tune_t_minarmem` + 100, the value is reset to `tune_t_minarmem` + 100. No message is displayed if the value from the `/etc/system` file is increased. Validation is done only at boot time and during dynamic reconfiguration operations that involve adding or deleting memory.
When to Change	When memory-locking requests fail or when attaching to a shared memory segment with the `SHARE_MMU` flag fails, yet the amount of memory available seems to be sufficient. Excessively large values can cause memory locking requests (`mlock`, `mlockall`, and `memcntl`) to fail unnecessarily. For more information, see `mlock(3C)`, `mlockall(3C)`, and `memcntl(2)`.
Commitment Level	Unstable

tune_t_minarmem

Description	Defines the minimum available resident (not swappable) memory to maintain necessary to avoid deadlock. Used to reserve a portion of memory for use by the core of the OS. Pages restricted in this way are not seen when the OS determines the maximum amount of memory available.
Data Type	Signed integer
Default	25
Range	1 to physical memory
Units	Pages
Dynamic?	No
Validation	None. Large values result in wasted physical memory.
When to Change	The default value is generally adequate. Consider increasing the default value if the system locks up and debugging information indicates that no memory was available.

Commitment Level	Unstable

fastscan

Description	Defines the maximum number of pages per second that the system looks at when memory pressure is highest.
Data Type	Signed integer
Default	After the system is booted, fastscan is set to 64 MB. Then, this value is automatically reset to the number of pages that the scanner can scan in one second by using 10% of a CPU. If this derived value is more than half the system's physical memory, the default value is limited to half the system's physical memory.
Range	64 MB to half the system's physical memory
Units	Pages
Dynamic?	Yes, unless dynamic reconfiguration operations that add or delete memory occur. At that point, the value is reset to the value provided by `/etc/system` or calculated from the new physical memory value.
Validation	The maximum value is the lesser of 64 MB and 1/2 of physical memory.
When to Change	When more aggressive scanning of memory is preferred during periods of memory shortfall, especially when the system is subject to periods of intense memory demand or when performing heavy file I/O.
Commitment Level	Unstable

slowscan

Description	Defines the minimum number of pages per second that the system looks at when attempting to reclaim memory.
Data Type	Signed integer
Default	The smaller of 1/20th of physical memory in pages and 100.
Range	1 to `fastscan` / 2
Units	Pages
Dynamic?	Yes, unless dynamic reconfiguration operations that add or delete memory occur. At that point, the value is reset to the value provided in the `/etc/system` file or calculated from the new physical memory value.

Validation	If `slowscan` is larger than `fastscan` / 2, `slowscan` is reset to `fastscan` / 2. No message is displayed.
When to Change	When more aggressive scanning of memory is preferred during periods of memory shortfall, especially when the system is subject to periods of intense memory demand.
Commitment Level	Unstable

min_percent_cpu

Description	Defines the minimum percentage of CPU that `pageout` can consume. This parameter is used as the starting point for determining the maximum amount of time that can be consumed by the page scanner.
Data Type	Signed integer
Default	4
Range	1 to 80
Units	Percentage
Dynamic?	Yes
Validation	None
When to Change	Increasing this value on systems with multiple CPUs and lots of memory, which are subject to intense periods of memory demand, enables the pager to spend more time attempting to find memory.
Commitment Level	Unstable

handspreadpages

Description	The Oracle Solaris OS uses a two-handed clock algorithm to look for pages that are candidates for reclaiming when memory is low. The first hand of the clock walks through memory marking pages as unused. The second hand walks through memory some distance after the first hand, checking to see if the page is still marked as unused. If so, the page is subject to being reclaimed. The distance between the first hand and the second hand is `handspreadpages`.
Data Type	Unsigned long
Default	`fastscan`
Range	1 to maximum number of physical memory pages on the system

Units	Pages
Dynamic?	Yes. This parameter requires that the kernel `reset_hands` parameter also be set to a non-zero value. Once the new value of `handspreadpages` has been recognized, `reset_hands` is set to zero.
Validation	The value is set to the lesser of either the amount of physical memory and the `handspreadpages` *value*.
When to Change	When you want to increase the amount of time that pages are potentially resident before being reclaimed. Increasing this value increases the separation between the hands, and therefore, the amount of time before a page can be reclaimed.
Commitment Level	Unstable

pages_before_pager

Description	Defines part of a system threshold that immediately frees pages after an I/O completes instead of storing the pages for possible reuse. The threshold is `lotsfree` + `pages_before_pager`. The NFS environment also uses this threshold to curtail its asynchronous activities as memory pressure mounts.
Data Type	Signed integer
Default	200
Range	1 to amount of physical memory
Units	Pages
Dynamic?	No
Validation	None
When to Change	You might change this parameter when the majority of I/O is done for pages that are truly read or written once and never referenced again. Setting this variable to a larger amount of memory keeps adding pages to the free list. You might also change this parameter when the system is subject to bursts of severe memory pressure. A larger value here helps maintain a larger cushion against the pressure.
Commitment Level	Unstable

maxpgio

Description	Defines the maximum number of page I/O requests that can be queued by the paging system. This number is divided by 4 to get the actual maximum number used by the paging system. This parameter is used to throttle the number of requests as well as to control process swapping.
Data Type	Signed integer
Default	40
Range	1 to a variable maximum that depends on the system architecture, but mainly by the I/O subsystem, such as the number of controllers, disks, and disk swap size
Units	I/0s
Dynamic?	No
Validation	None
Implicit	The maximum number of I/O requests from the pager is limited by the size of a list of request buffers, which is currently sized at 256.
When to Change	Increase this parameter to page out memory faster. A larger value might help to recover faster from memory pressure if more than one swap device is configured or if the swap device is a striped device. Note that the existing I/O subsystem should be able to handle the additional I/O load. Also, increased swap I/O could degrade application I/O performance if the swap partition and application files are on the same disk.
Commitment Level	Unstable
Change History	For information, see "`maxpgio` (Solaris 10 Releases)" on page 180.

Swapping-Related Parameters

Swapping in the Oracle Solaris OS is accomplished by the swapfs pseudo file system. The combination of space on swap devices and physical memory is treated as the pool of space available to support the system for maintaining backing store for anonymous memory. The system attempts to allocate space from disk devices first, and then uses physical memory as backing store. When swapfs is forced to use system memory for backing store, limits are enforced to ensure that the system does not deadlock because of excessive consumption by swapfs.

swapfs_reserve

Description	Defines the amount of system memory that is reserved for use by system (UID = 0) processes.
Data Type	Unsigned long
Default	The smaller of 4 MB and 1/16th of physical memory
Range	The minimum value is 4 MB or 1/16th of physical memory, whichever is smaller, expressed as pages using the page size returned by `getpagesize`. The maximum value is the number of physical memory pages. The maximum value should be no more than 10 percent of physical memory. The system does not enforce this range, other than that described in the Validation section.
Units	Pages
Dynamic?	No
Validation	None
When to Change	Generally not necessary. Only change when recommended by a software provider, or when system processes are terminating because of an inability to obtain swap space. A much better solution is to add physical memory or additional swap devices to the system.
Commitment Level	Unstable

swapfs_minfree

Description	Defines the desired amount of physical memory to be kept free for the rest of the system. Attempts to reserve memory for use as swap space by any process that causes the system's perception of available memory to fall below this value are rejected. Pages reserved in this manner can only be used for locked-down allocations by the kernel or by user-level processes.
Data Type	Unsigned long
Default	The larger of 2 MB and 1/8th of physical memory
Range	1 to amount of physical memory
Units	Pages
Dynamic?	No

Validation	None
When to Change	When processes are failing because of an inability to obtain swap space, yet the system has memory available.
Commitment Level	Unstable

Kernel Memory Allocator

The Oracle Solaris kernel memory allocator distributes chunks of memory for use by clients inside the kernel. The allocator creates a number of caches of varying size for use by its clients. Clients can also request the allocator to create a cache for use by that client (for example, to allocate structures of a particular size). Statistics about each cache that the allocator manages can be seen by using the `kstat -c kmem_cache` command.

Occasionally, systems might panic because of memory corruption. The kernel memory allocator supports a debugging interface (a set of flags), that performs various integrity checks on the buffers. The kernel memory allocator also collects information on the allocators. The integrity checks provide the opportunity to detect errors closer to where they actually occurred. The collected information provides additional data for support people when they try to ascertain the reason for the panic.

Use of the flags incurs additional overhead and memory usage during system operations. The flags should only be used when a memory corruption problem is suspected.

kmem_flags

Description

The Oracle Solaris kernel memory allocator has various debugging and test options.

Five supported flag settings are described here.

Flag	Setting	Description
`AUDIT`	`0x1`	The allocator maintains a log that contains recent history of its activity. The number of items logged depends on whether `CONTENTS` is also set. The log is a fixed size. When space is exhausted, earlier records are reclaimed.

Flag	Setting	Description
TEST	0x2	The allocator writes a pattern into freed memory and checks that the pattern is unchanged when the buffer is next allocated. If some portion of the buffer is changed, then the memory was probably used by a client that had previously allocated and freed the buffer. If an overwrite is identified, the system panics.
REDZONE	0x4	The allocator provides extra memory at the end of the requested buffer and inserts a special pattern into that memory. When the buffer is freed, the pattern is checked to see if data was written past the end of the buffer. If an overwrite is identified, the kernel panics.
CONTENTS	0x8	The allocator logs up to 256 bytes of buffer contents when the buffer is freed. This flag requires that AUDIT also be set. The numeric value of these flags can be logically added together and set by the /etc/system file.
LITE	0x100	Does minimal integrity checking when a buffer is allocated and freed. When enabled, the allocator checks that the redzone has not been written into, that a freed buffer is not being freed again, and that the buffer being freed is the size that was allocated. Do not combine this flag with any other flags.

Data Type	Signed integer
Default	0 (disabled)
Range	0 (disabled) or 1 - 15 or 256 (0x100)
Dynamic?	Yes. Changes made during runtime only affect new kernel memory caches. After system initialization, the creation of new caches is rare.
Validation	None
When to Change	When memory corruption is suspected
Commitment Level	Unstable

General Driver Parameters

moddebug

Description — When this parameter is enabled, messages about various steps in the module loading process are displayed.

Data Type — Signed integer

Default — 0 (messages off)

Range — Here are the most useful values:

- 0x80000000 – Prints `[un] loading...` message. For every module loaded, messages such as the following appear on the console and in the `/var/adm/messages` file:

```
Apr 20 17:18:04 neo genunix: [ID 943528 kern.notice] load 'sched/TS_DPTBL' id 15
loaded @ 0x7be1b2f8/0x19c8380 size 176/2096
Apr 20 17:18:04 neo genunix: [ID 131579 kern.notice] installing TS_DPTBL,
module id 15.
```

- 0x40000000 – Prints detailed error messages. For every module loaded, messages such as the following appear on the console and in the `/var/adm/messages` file:

```
Apr 20 18:30:00 neo unix: Errno = 2
Apr 20 18:30:00 neo unix: kobj_open: vn_open of /platform/sun4v/kernel/exec/sparcv9/intpexec fails
Apr 20 18:30:00 neo unix: Errno = 2
Apr 20 18:30:00 neo unix: kobj_open: '/kernel/exec/sparcv9/intpexec'
Apr 20 18:30:00 neo unix:  vp = 60015777600
Apr 20 18:30:00 neo unix: kobj_close: 0x60015777600
Apr 20 18:30:00 neo unix: kobj_open: vn_open of /platform/SUNW,Sun-Fire-T200/kernel/exec/sparcv9
/intpexec fails,
Apr 20 18:30:00 neo unix: Errno = 2
Apr 20 18:30:00 neo unix: kobj_open: vn_open of /platform/sun4v/kernel/exec/sparcv9/intpexec fails
```

- 0x20000000 - Prints even more detailed messages. This value doesn't print any additional information beyond what the `0x40000000` flag does during system boot. However, this value does print additional information about releasing the module when the module is unloaded.

These values can be added together to set the final value.

Dynamic? — Yes

Validation — None

When to Change: When a module is either not loading as expected, or the system seems to hang while loading modules. Note that when `0x40000000` is set, system boot is slowed down considerably by the number of messages written to the console.

Commitment Level: Unstable

ddi_msix_alloc_limit

Description: x86 only: This parameter controls the number of Extended Message Signaled Interrupts (MSI-X) that a device instance can allocate. Due to an existing system limitation, the default value is 2. You can increase the number of MSI-X interrupts that a device instance can allocate by increasing the value of this parameter. This parameter can be set either by editing the `/etc/system` file or by setting it with `mdb` before the device driver attach occurs.

Data Type: Signed integer

Default: 2

Range: 1 to 16

Dynamic?: Yes

Validation: None

When to Change: To increase the number of MSI-X interrupts that a device instance can allocate. However, if you increase the number of MSI-X interrupts that a device instance can allocate, adequate interrupts might not be available to satisfy all allocation requests. If this happens, some devices might stop functioning or the system might fail to boot. Reduce the value or remove the parameter in this case.

Commitment Level: Unstable

Change History: For information, see "`ddi_msix_alloc_limit` (Solaris 10 Releases)" on page 179.

General I/O Parameters

maxphys

Description	Defines the maximum size of physical I/O requests. If a driver encounters a request larger than this size, the driver breaks the request into `maxphys` sized chunks. File systems can and do impose their own limit.
Data Type	Signed integer
Default	131,072 (sun4u or sun4v) or 57,344 (x86). The `sd` driver uses the value of 1,048,576 if the drive supports wide transfers. The `ssd` driver uses 1,048,576 by default.
Range	Machine-specific page size to `MAXINT`
Units	Bytes
Dynamic?	Yes, but many file systems load this value into a per-mount point data structure when the file system is mounted. A number of drivers load the value at the time a device is attached to a driver-specific data structure.
Validation	None
When to Change	When doing I/O to and from raw devices in large chunks. Note that a DBMS doing OLTP operations issues large numbers of small I/Os. Changing `maxphys` does not result in any performance improvement in that case. You might also consider changing this parameter when doing I/O to and from a UFS file system where large amounts of data (greater than 64 KB) are being read or written at any one time. The file system should be optimized to increase contiguity. For example, increase the size of the cylinder groups and decrease the number of inodes per cylinder group. UFS imposes an internal limit of 1 MB on the maximum I/O size it transfers.
Commitment Level	Unstable
Change History	For information, see "`maxphys` (Solaris 10 Releases)" on page 180.

rlim_fd_max

Description	Specifies the "hard" limit on file descriptors that a single process might have open. Overriding this limit requires superuser privilege.
Data Type	Signed integer
Default	65,536
Range	1 to `MAXINT`
Units	File descriptors
Dynamic?	No
Validation	None
When to Change	When the maximum number of open files for a process is not enough. Other limitations in system facilities can mean that a larger number of file descriptors is not as useful as it might be. For example: ■ A 32-bit program using standard I/O is limited to 256 file descriptors. A 64-bit program using standard I/O can use up to 2 billion descriptors. Specifically, standard I/O refers to the `stdio(3C)` functions in `libc(3LIB)`. ■ `select` is by default limited to 1024 descriptors per `fd_set`. For more information, see `select(3C)`. A 32-bit application code can be recompiled with a larger `fd_set` size (less than or equal to 65,536). A 64-bit application uses an `fd_set` size of 65,536, which cannot be changed. An alternative to changing this on a system wide basis is to use the `plimit(1)` command. If a parent process has its limits changed by `plimit`, all children inherit the increased limit. This alternative is useful for daemons such as `inetd`.
Commitment Level	Unstable

rlim_fd_cur

Description	Defines the "soft" limit on file descriptors that a single process can have open. A process might adjust its file descriptor limit to any value up to the "hard" limit defined by `rlim_fd_max` by using the `setrlimit()` call or by issuing the `limit` command in whatever shell it is running. You do not require superuser privilege to adjust the limit to any value less than or equal to the hard limit.

Data Type	Signed integer
Default	256
Range	1 to `MAXINT`
Units	File descriptors
Dynamic?	No
Validation	Compared to `rlim_fd_max`. If `rlim_fd_cur` is greater than `rlim_fd_max`, `rlim_fd_cur` is reset to `rlim_fd_max`.
When to Change	When the default number of open files for a process is not enough. Increasing this value means only that it might not be necessary for a program to use `setrlimit` to increase the maximum number of file descriptors available to it.
Commitment Level	Unstable

General File System Parameters

ncsize

Description	Defines the number of entries in the directory name look-up cache (DNLC). This parameter is used by UFS, NFS, and ZFS to cache elements of path names that have been resolved. The DNLC also caches negative look-up information, which means it caches a name not found in the cache.
Data Type	Signed integer
Default	(4 x (`v.v_proc` + `maxusers`) + 320) + (4 x (`v.v_proc` + `maxusers`) + 320 / 100
Range	0 to `MAXINT`
Units	DNLC entries
Dynamic?	No
Validation	None. Larger values cause the time it takes to unmount a file system to increase as the cache must be flushed of entries for that file system during the unmount process.
When to Change	You can use the `kstat -n dnlcstats` command to determine when entries have been removed from the DNLC because it was too small.

The sum of the `pick_heuristic` and the `pick_last` parameters represents otherwise valid entries that were reclaimed because the cache was too small.

Excessive values of `ncsize` have an immediate impact on the system because the system allocates a set of data structures for the DNLC based on the value of `ncsize`. A system running a 32-bit kernel allocates 36-byte structures for `ncsize`, while a system running a 64-bit kernel allocates 64-byte structures for `ncsize`. The value has a further effect on UFS and NFS, unless `ufs_ninode` and `nfs:nrnode` are explicitly set.

Commitment Level: Unstable

Change History: For information, see "`ncsize` (Solaris 10 Releases)" on page 181.

dnlc_dir_enable

Description: Enables large directory caching

Note – This parameter has no effect on NFS or ZFS file systems.

Data Type: Unsigned integer

Default: 1 (enabled)

Range: 0 (disabled) or 1 (enabled)

Dynamic?: Yes, but do not change this tunable dynamically. You can enable this parameter if it was originally disabled. Or, you can disable this parameter if it was originally enabled. However, enabling, disabling, and then enabling this parameter might lead to stale directory caches.

Validation: No

When to Change: Directory caching has no known problems. However, if problems occur, then set `dnlc_dir_enable` to 0 to disable caching.

Commitment Level: Unstable

dnlc_dir_min_size

Description: Specifies the minimum number of entries cached for one directory.

Note – This parameter has no effect on NFS or ZFS file systems.

Data Type	Unsigned integer
Default	40
Range	0 to MAXUINT (no maximum)
Units	Entries
Dynamic?	Yes, this parameter can be changed at any time.
Validation	None
When to Change	If performance problems occur with caching small directories, then increase `dnlc_dir_min_size`. Note that individual file systems might have their own range limits for caching directories. For instance, UFS limits directories to a minimum of `ufs_min_dir_cache` bytes (approximately 1024 entries), assuming 16 bytes per entry.
Commitment Level	Unstable

dnlc_dir_max_size

Description — Specifies the maximum number of entries cached for one directory.

Note – This parameter has no effect on NFS or ZFS file systems.

Data Type	Unsigned integer
Default	MAXUINT (no maximum)
Range	0 to MAXUINT
Dynamic?	Yes, this parameter can be changed at any time.
Validation	None
When to Change	If performance problems occur with large directories, then decrease `dnlc_dir_max_size`.
Commitment Level	Unstable

segmap_percent

Description	Defines the maximum amount of memory that is used for the fast-access file system cache. This pool of memory is subtracted from the free memory list.
Data Type	Unsigned integer
Default	12 percent of free memory at system startup time
Range	2 MB to 100 percent of `physmem`
Units	% of physical memory
Dynamic?	No
Validation	None
When to Change	If heavy file system activity is expected, and sufficient free memory is available, you should increase the value of this parameter.
Commitment Level	Unstable

UFS Parameters

bufhwm and bufhwm_pct

Description	Defines the maximum amount of memory for caching I/O buffers. The buffers are used for writing file system metadata (superblocks, inodes, indirect blocks, and directories). Buffers are allocated as needed until the amount of memory (in KB) to be allocated exceed `bufhwm`. At this point, metadata is purged from the buffer cache until enough buffers are reclaimed to satisfy the request. For historical reasons, `bufhwm` does not require the `ufs:` prefix.
Data Type	Signed integer
Default	2 percent of physical memory
Range	80 KB to 20 percent of physical memory, or 2 TB, whichever is less. Consequently, `bufhwm_pct` can be between 1 and 20.
Units	`bufhwm`: KB `bufhwm_pct`: percent of physical memory

Dynamic?

No. `bufhwm` and `bufhwm_pct` are only evaluated at system initialization to compute hash bucket sizes. The limit in bytes calculated from these parameters is then stored in a data structure that adjusts this value as buffers are allocated and deallocated.

Attempting to adjust this value without following the locking protocol on a running system can lead to incorrect operation.

Modifying `bufhwm` or `bufhwm_pct` at runtime has no effect.

Validation

If `bufhwm` is less than its lower limit of 80 KB or greater than its upper limit (the lesser of 20 percent of physical memory, 2 TB, or one quarter (1/4) of the maximum amount of kernel heap), it is reset to the upper limit. The following message appears on the system console and in the `/var/adm/messages` file if an invalid value is attempted:

"`binit:` *`bufhwm`* `(value attempted) out of range (`*`range start..range end`*`). Using` *`N`* `as default.`"

"Value attempted" refers to the value specified in the `/etc/system` file or by using a kernel debugger. *N* is the value computed by the system based on available system memory.

Likewise, if `bufhwm_pct` is set to a value that is outside the allowed range of 1 percent to 20 percent, it is reset to the default of 2 percent. And, the following message appears on the system console and in the `/var/adm/messages` file:

"`binit:` *`bufhwm_pct`*`(value attempted) out of range(0..20). Using 2 as default.`"

If both `bufhwm` or `bufhwm_pct` are set to non-zero values, `bufhwm` takes precedence.

When to Change

Because buffers are only allocated as they are needed, the overhead from the default setting is the required allocation of control structures for the buffer hash headers. These structures consume 52 bytes per potential buffer on a 32-bit kernel and 96 bytes per potential buffer on a 64-bit kernel.

On a 512-MB 64-bit kernel, the number of hash chains calculates to 10316 / 32 == 322, which scales up to next power of 2, 512. Therefore, the hash headers consume 512 x 96 bytes, or 48 KB. The hash header allocations assume that buffers are 32 KB.

The amount of memory, which has not been allocated in the buffer pool, can be found by looking at the `bfreelist` structure in the kernel with a

kernel debugger. The field of interest in the structure is `b_bufsize`, which is the possible remaining memory in bytes. Looking at it with the `buf` macro by using the `mdb` command:

```
# mdb -k
Loading modules: [ unix krtld genunix ip nfs ipc ]
> bfreelist::print "struct buf" b_bufsize
b_bufsize = 0x225800
```

The default value for `bufhwm` on this system, with 6 GB of memory, is 122277. You cannot determine the number of header structures used because the actual buffer size requested is usually larger than 1 KB. However, some space might be profitably reclaimed from control structure allocation for this system.

The same structure on a 512-MB system shows that only 4 KB of 10144 KB has not been allocated. When the `biostats kstat` is examined with `kstat -n biostats`, it is determined that the system had a reasonable ratio of `buffer_cache_hits` to `buffer_cache_lookups` as well. As such, the default setting is reasonable for that system.

Commitment Level	Unstable

ndquot

Description	Defines the number of quota structures for the UFS file system that should be allocated. Relevant only if quotas are enabled on one or more UFS file systems. Because of historical reasons, the `ufs:` prefix is not needed.
Data Type	Signed integer
Default	((`maxusers` x 40) / 4) + `max_nprocs`
Range	0 to MAXINT
Units	Quota structures
Dynamic?	No
Validation	None. Excessively large values hang the system.
When to Change	When the default number of quota structures is not enough. This situation is indicated by the following message displayed on the console or written in the message log: `dquot table full`
Commitment Level	Unstable

ufs_ninode

Description

Specifies the number of inodes to be held in memory. Inodes are cached globally for UFS, not on a per-file system basis.

A key parameter in this situation is `ufs_ninode`. This parameter is used to compute two key limits that affect the handling of inode caching. A high watermark of `ufs_ninode / 2` and a low watermark of `ufs_ninode / 4` are computed.

When the system is done with an inode, one of two things can happen:

- The file referred to by the inode is no longer on the system so the inode is deleted. After it is deleted, the space goes back into the inode cache for use by another inode (which is read from disk or created for a new file).
- The file still exists but is no longer referenced by a running process. The inode is then placed on the idle queue. Any referenced pages are still in memory.

When inodes are idled, the kernel defers the idling process to a later time. If a file system is a logging file system, the kernel also defers deletion of inodes. Two kernel threads handle this deferred processing. Each thread is responsible for one of the queues.

When the deferred processing is done, the system drops the inode onto either a delete queue or an idle queue, each of which has a thread that can run to process it. When the inode is placed on the queue, the queue occupancy is checked against the low watermark. If the queue occupancy exceeds the low watermark, the thread associated with the queue is awakened. After the queue is awakened, the thread runs through the queue and forces any pages associated with the inode out to disk and frees the inode. The thread stops when it has removed 50 percent of the inodes on the queue at the time it was awakened.

A second mechanism is in place if the idle thread is unable to keep up with the load. When the system needs to find a vnode, it goes through the `ufs_vget` routine. The *first* thing `vget` does is check the length of the idle queue. If the length is above the high watermark, then it takes two inodes off the idle queue and "idles" them (flushes pages and frees inodes). `vget` does this *before* it gets an inode for its own use.

The system does attempt to optimize by placing inodes with no in-core pages at the head of the idle list and inodes with pages at the end of the

	idle list. However, the system does no other ordering of the list. Inodes are always removed from the front of the idle queue. The only time that inodes are removed from the queues as a whole is when a synchronization, unmount, or remount occur. For historical reasons, this parameter does not require the `ufs:` prefix.
Data Type	Signed integer
Default	`ncsize`
Range	0 to MAXINT
Units	Inodes
Dynamic?	Yes
Validation	If `ufs_ninode` is less than or equal to zero, the value is set to `ncsize`.
When to Change	When the default number of inodes is not enough. If the `maxsize reached` field as reported by `kstat -n inode_cache` is larger than the `maxsize` field in the `kstat`, the value of `ufs_ninode` might be too small. Excessive inode idling can also be a problem. You can identify excessive inode idling by using `kstat -n inode_cache` to look at the `inode_cache kstat`. `Thread idles` are inodes idled by the background threads while `vget idles` are idles by the requesting process before using an inode.
Commitment Level	Unstable

ufs_WRITES

Description	If `ufs_WRITES` is non-zero, the number of bytes outstanding for writes on a file is checked. See `ufs_HW` to determine whether the write should be issued or deferred until only `ufs_LW` bytes are outstanding. The total number of bytes outstanding is tracked on a per-file basis so that if the limit is passed for one file, it won't affect writes to other files.
Data Type	Signed integer
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Units	Toggle (on/off)
Dynamic?	Yes

Validation	None
When to Change	When you want UFS write throttling turned off entirely. If sufficient I/O capacity does not exist, disabling this parameter can result in long service queues for disks.
Commitment Level	Unstable

ufs_LW and ufs_HW

Description	`ufs_HW` specifies the number of bytes outstanding on a single file barrier value. If the number of bytes outstanding is greater than this value and `ufs_WRITES` is set, then the write is deferred. The write is deferred by putting the thread issuing the write to sleep on a condition variable. `ufs_LW` is the barrier for the number of bytes outstanding on a single file below which the condition variable on which other sleeping processes are toggled. When a write completes and the number of bytes is less than `ufs_LW`, then the condition variable is toggled, which causes all threads waiting on the variable to awaken and try to issue their writes.
Data Type	Signed integer
Default	8 x 1024 x 1024 for `ufs_LW` and 16 x 1024 x 1024 for `ufs_HW`
Range	0 to MAXINT
Units	Bytes
Dynamic?	Yes
Validation	None
Implicit	`ufs_LW` and `ufs_HW` have meaning only if `ufs_WRITES` is not equal to zero. `ufs_HW` and `ufs_LW` should be changed together to avoid needless churning when processes awaken and find that either they cannot issue a write (when `ufs_LW` and `ufs_HW` are too close) or they might have waited longer than necessary (when `ufs_LW` and `ufs_HW` are too far apart).
When to Change	Consider changing these values when file systems consist of striped volumes. The aggregate bandwidth available can easily exceed the current value of `ufs_HW`. Unfortunately, this parameter is not a per-file system setting.

	You might also consider changing this parameter when `ufs_throttles` is a non-trivial number. Currently, `ufs_throttles` can only be accessed with a kernel debugger.
Commitment Level	Unstable

freebehind

Description	Enables the `freebehind` algorithm. When this algorithm is enabled, the system bypasses the file system cache on newly read blocks when sequential I/O is detected during times of heavy memory use.
Data Type	Boolean
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes
Validation	None
When to Change	The `freebehind` algorithm can occur too easily. If no significant sequential file system activity is expected, disabling `freebehind` makes sure that all files, no matter how large, will be candidates for retention in the file system page cache. For more fine-grained tuning, see `smallfile`.
Commitment Level	Unstable

smallfile

Description	Determines the size threshold of files larger than this value are candidates for no cache retention under the `freebehind` algorithm. Large memory systems contain enough memory to cache thousands of 10-MB files without making severe memory demands. However, this situation is highly application dependent. The goal of the `smallfile` and `freebehind` parameters is to reuse cached information, without causing memory shortfalls by caching too much.
Data Type	Signed integer
Default	0

Range	0 to 2,147,483,647
Dynamic?	Yes
Validation	None
When to Change	Increase `smallfile` if an application does sequential reads on medium-sized files and can most likely benefit from buffering, and the system is not otherwise under pressure for free memory. Medium-sized files are 32 KB to 2 GB in size.
Commitment Level	Unstable

TMPFS Parameters

tmpfs:tmpfs_maxkmem

Description	Defines the maximum amount of kernel memory that TMPFS can use for its data structures (tmpnodes and directory entries).
Data Type	Unsigned long
Default	One page or 4 percent of physical memory, whichever is greater.
Range	Number of bytes in one page (8192 for sun4u or sun4v systems, 4096 for all other systems) to 25 percent of the available kernel memory at the time TMPFS was first used.
Units	Bytes
Dynamic?	Yes
Validation	None
When to Change	Increase if the following message is displayed on the console or written in the messages file: `tmp_memalloc: tmpfs over memory limit` The current amount of memory used by TMPFS for its data structures is held in the `tmp_kmemspace` field. This field can be examined with a kernel debugger.
Commitment Level	Unstable
Change History	For information, see "`tmpfs:tmpfs_maxkmem` (Solaris 10 Releases)" on page 181.

tmpfs:tmpfs_minfree

Description	Defines the minimum amount of swap space that TMPFS leaves for the rest of the system.
Data Type	Signed long
Default	256
Range	0 to maximum swap space size
Units	Pages
Dynamic?	Yes
Validation	None
When to Change	To maintain a reasonable amount of swap space on systems with large amounts of TMPFS usage, you can increase this number. The limit has been reached when the console or messages file displays the following message: *fs-name*: `File system full, swap space limit exceeded`
Commitment Level	Unstable

Pseudo Terminals

Pseudo terminals, `ptys`, are used for two purposes in Oracle Solaris software:

- Supporting remote logins by using the `telnet`, `rlogin`, or `rsh` commands
- Providing the interface through which the X Window system creates command interpreter windows

The default number of pseudo-terminals is sufficient for a desktop workstation. So, tuning focuses on the number of `ptys` available for remote logins.

The default number of `ptys` is now based on the amount of memory on the system. This default should be changed only to restrict or increase the number of users who can log in to the system.

Three related variables are used in the configuration process:

- `pt_cnt` – Default maximum number of `ptys`.
- `pt_pctofmem` – Percentage of kernel memory that can be dedicated to `pty` support structures. A value of zero means that no remote users can log in to the system.
- `pt_max_pty` – Hard maximum for number of `ptys`.

pt_cnt has a default value of zero, which tells the system to limit logins based on the amount of memory specified in pct_pctofmem, unless pt_max_pty is set. If pt_cnt is non-zero, ptys are allocated until this limit is reached. When that threshold is crossed, the system looks at pt_max_pty. If pt_max_pty has a non-zero value, it is compared to pt_cnt. The pty allocation is allowed if pt_cnt is less than pt_max_pty. If pt_max_pty is zero, pt_cnt is compared to the number of ptys supported based on pt_pctofmem. If pt_cnt is less than this value, the pty allocation is allowed. Note that the limit based on pt_pctofmem only comes into play if both pt_cnt and ptms_ptymax have default values of zero.

To put a hard limit on ptys that is different than the maximum derived from pt_pctofmem, set pt_cnt and ptms_ptymax in /etc/system to the preferred number of ptys. The setting of ptms_pctofmem is not relevant in this case.

To dedicate a different percentage of system memory to pty support and let the operating system manage the explicit limits, do the following:

- Do not set pt_cnt or ptms_ptymax in /etc/system.
- Set pt_pctofmem in /etc/system to the preferred percentage. For example, set pt_pctofmem=10 for a 10 percent setting.

Note that the memory is not actually allocated until it is used in support of a pty. Once memory is allocated, it remains allocated.

pt_cnt

Description	The number of available /dev/pts entries is dynamic up to a limit determined by the amount of physical memory available on the system. pt_cnt is one of three variables that determines the minimum number of logins that the system can accommodate. The default maximum number of /dev/pts devices the system can support is determined at boot time by computing the number of pty structures that can fit in a percentage of system memory (see pt_pctofmem). If pt_cnt is zero, the system allocates up to that maximum. If pt_cnt is non-zero, the system allocates to the greater of pt_cnt and the default maximum.
Data Type	Unsigned integer
Default	0
Range	0 to maxpid
Units	Logins/windows
Dynamic?	No
Validation	None

When to Change	When you want to explicitly control the number of users who can remotely log in to the system.
Commitment Level	Unstable

pt_pctofmem

Description	Specifies the maximum percentage of physical memory that can be consumed by data structures to support `/dev/pts` entries. A system running a 64-bit kernel consumes 176 bytes per `/dev/pts` entry. A system running a 32-bit kernel consumes 112 bytes per `/dev/pts` entry.
Data Type	Unsigned integer
Default	5
Range	0 to 100
Units	Percentage
Dynamic?	No
Validation	None
When to Change	When you want to either restrict or increase the number of users who can log in to the system. A value of zero means that no remote users can log in to the system.
Commitment Level	Unstable

pt_max_pty

Description	Defines the maximum number of `ptys` the system offers
Data Type	Unsigned integer
Default	0 (Uses system-defined maximum)
Range	0 to MAXUINT
Units	Logins/windows
Dynamic?	Yes
Validation	None
Implicit	Should be greater than or equal to `pt_cnt`. Value is not checked until the number of `ptys` allocated exceeds the value of `pt_cnt`.

When to Change	When you want to place an absolute ceiling on the number of logins supported, even if the system could handle more based on its current configuration values.
Commitment Level	Unstable

STREAMS Parameters

nstrpush

Description	Specifies the number of modules that can be inserted into (pushed onto) a STREAM.
Data Type	Signed integer
Default	9
Range	9 to 16
Units	Modules
Dynamic?	Yes
Validation	None
When to Change	At the direction of your software vendor. No messages are displayed when a STREAM exceeds its permitted push count. A value of `EINVAL` is returned to the program that attempted the push.
Commitment Level	Unstable

strmsgsz

Description	Specifies the maximum number of bytes that a single system call can pass to a STREAM to be placed in the data part of a message. Any `write` exceeding this size is broken into multiple messages. For more information, see `write(2)`.
Data Type	Signed integer
Default	65,536
Range	0 to 262,144
Units	Bytes

Dynamic?	Yes
Validation	None
When to Change	When `putmsg` calls return `ERANGE`. For more information, see `putmsg(2)`.
Commitment Level	Unstable

strctlsz

Description	Specifies the maximum number of bytes that a single system call can pass to a STREAM to be placed in the control part of a message
Data Type	Signed integer
Default	1024
Range	0 to MAXINT
Units	Bytes
Dynamic?	Yes
Validation	None
When to Change	At the direction of your software vendor. `putmsg(2)` calls return `ERANGE` if they attempt to exceed this limit.
Commitment Level	Unstable

System V Message Queues

System V message queues provide a message-passing interface that enables the exchange of messages by queues created in the kernel. Interfaces are provided in the Oracle Solaris environment to enqueue and dequeue messages. Messages can have a type associated with them. Enqueueing places messages at the end of a queue. Dequeuing removes the first message of a specific type from the queue or the first message if no type is specified.

For information about System V message queues in the Oracle Solaris 10 release, see "System V IPC Configuration" on page 19.

For detailed information on tuning these system resources, see Chapter 6, "Resource Controls (Overview)," in *System Administration Guide: Oracle Solaris Containers-Resource Management and Oracle Solaris Zones*.

For legacy information about the obsolete System V message queues, see "Parameters That Are Obsolete or Have Been Removed" on page 184.

System V Semaphores

System V semaphores provide counting semaphores in the Oracle Solaris OS. A *semaphore* is a counter used to provide access to a shared data object for multiple processes. In addition to the standard set and release operations for semaphores, System V semaphores can have values that are incremented and decremented as needed (for example, to represent the number of resources available). System V semaphores also provide the ability to do operations on a group of semaphores simultaneously as well as to have the system undo the last operation by a process if the process dies.

For information about the changes to semaphore resources in the Oracle Solaris 10 release, see "System V IPC Configuration" on page 19.

For detailed information about using the new resource controls in the Oracle Solaris 10 release, see Chapter 6, "Resource Controls (Overview)," in *System Administration Guide: Oracle Solaris Containers-Resource Management and Oracle Solaris Zones*.

For legacy information about the obsolete System V semaphore parameters, see "Parameters That Are Obsolete or Have Been Removed" on page 184.

System V Shared Memory

System V shared memory allows the creation of a segment by a process. Cooperating processes can attach to the memory segment (subject to access permissions on the segment) and gain access to the data contained in the segment. This capability is implemented as a loadable module. Entries in the `/etc/system` file must contain the `shmsys:` prefix..

A special kind of shared memory known as *intimate shared memory* (ISM) is used by DBMS vendors to maximize performance. When a shared memory segment is made into an ISM segment, the memory for the segment is locked. This feature enables a faster I/O path to be followed and improves memory usage. A number of kernel resources describing the segment are then shared between all processes that attach to the segment in ISM mode.

For information about the changes to shared memory resources in the Oracle Solaris 10 release, see "System V IPC Configuration" on page 19.

For detailed information about using the new resource controls in the Oracle Solaris 10 release, see Chapter 6, "Resource Controls (Overview)," in *System Administration Guide: Oracle Solaris Containers-Resource Management and Oracle Solaris Zones*.

For legacy information about the obsolete System V shared memory parameters, see "Parameters That Are Obsolete or Have Been Removed" on page 184.

segspt_minfree

Description	Identifies pages of system memory that cannot be allocated for ISM shared memory.
Data Type	Unsigned long
Default	5 percent of available system memory when the first ISM segment is created
Range	0 to 50 percent of physical memory
Units	Pages
Dynamic?	Yes
Validation	None. Values that are too small can cause the system to hang or performance to severely degrade when memory is consumed with ISM segments.
When to Change	On database servers with large amounts of physical memory using ISM, the value of this parameter can be decreased. If ISM segments are not used, this parameter has no effect. A maximum value of 128 MB (0x4000) is almost certainly sufficient on large memory machines.
Commitment Level	Unstable

Scheduling

rechoose_interval

Description	Specifies the number of clock ticks before a process is deemed to have lost all affinity for the last CPU it ran on. After this interval expires, any CPU is considered a candidate for scheduling a thread. This parameter is relevant only for threads in the timesharing class. Real-time threads are scheduled on the first available CPU.
Data Type	Signed integer
Default	3
Range	0 to MAXINT

Dynamic?	Yes
Validation	None
When to Change	When caches are large, or when the system is running a critical process or a set of processes that seem to suffer from excessive cache misses not caused by data access patterns. Consider using the processor set capabilities or processor binding before changing this parameter. For more information, see `psrset(1M)` or `pbind(1M)`.
Commitment Level	Unstable

Timers

hires_tick

Description	When set, this parameter causes the Oracle Solaris OS to use a system clock rate of 1000 instead of the default value of 100.
Data Type	Signed integer
Default	0
Range	0 (disabled) or 1 (enabled)
Dynamic?	No. Causes new system timing variable to be set at boot time. Not referenced after boot.
Validation	None
When to Change	When you want timeouts with a resolution of less than 10 milliseconds, and greater than or equal to 1 millisecond.
Commitment Level	Unstable

timer_max

Description	Specifies the number of POSIX timers available.
Data Type	Signed integer
Default	32
Range	0 to MAXINT

Dynamic?	No. Increasing the value can cause a system crash.
Validation	None
When to Change	When the default number of timers offered by the system is inadequate. Applications receive an `EAGAIN` error when executing `timer_create` system calls.
Commitment Level	Unstable

SPARC System Specific Parameters

consistent_coloring

Description

The ability to use different page placement policies on the UltraSPARC (`sun4u`) platform is available. A page placement policy attempts to allocate physical page addresses to maximize the use of the L2 cache. Whatever algorithm is chosen as the default algorithm, that algorithm can potentially provide less optimal results than another algorithm for a particular application set. This parameter changes the placement algorithm selected for all processes on the system.

Based on the size of the L2 cache, memory is divided into bins. The page placement code allocates a page from a bin when a page fault first occurs on an unmapped page. The page chosen depends on which of the three possible algorithms are used:

- Page coloring – Various bits of the virtual address are used to determine the bin from which the page is selected. `consistent_coloring` is set to zero to use this algorithm. No per-process history exists for this algorithm.
- Virtual addr=physical address – Consecutive pages in the program selects pages from consecutive bins. `consistent_coloring` is set to 1 to use this algorithm. No per-process history exists for this algorithm.
- Bin-hopping – Consecutive pages in the program generally allocate pages from every other bin, but the algorithm occasionally skips more bins. `consistent_coloring` is set to 2 to use this algorithm. Each process starts at a randomly selected bin, and a per-process memory of the last bin allocated is kept.

Dynamic? Yes

Validation	None. Values larger than 2 cause a number of `WARNING: AS_2_BIN: bad consistent coloring value` messages to appear on the console. The system hangs immediately thereafter. A power-cycle is required to recover.
When to Change	When the primary workload of the system is a set of long-running high-performance computing (HPC) applications. Changing this value might provide better performance. File servers, database servers, and systems with a number of active processes (for example, compile or time sharing servers) do not benefit from changes.
Commitment Level	Unstable

tsb_alloc_hiwater_factor

Description	Initializes `tsb_alloc_hiwater` to impose an upper limit on the amount of physical memory that can be allocated for translation storage buffers (TSBs) as follows: `tsb_alloc_hiwater` = physical memory (bytes) / `tsb_alloc_hiwater_factor` When the memory that is allocated to TSBs is equal to the value of `tsb_alloc_hiwater`, the TSB memory allocation algorithm attempts to reclaim TSB memory as pages are unmapped. Exercise caution when using this factor to increase the value of `tsb_alloc_hiwater`. To prevent system hangs, the resulting high water value must be considerably lower than the value of `swapfs_minfree` and `segspt_minfree`.
Data Type	Integer
Default	32
Range	1 to MAXINIT Note that a factor of 1 makes all physical memory available for allocation to TSBs, which could cause the system to hang. A factor that is too high will not leave memory available for allocation to TSBs, decreasing system performance.
Dynamic?	Yes
Validation	None

When to Change: Change the value of this parameter if the system has many processes that attach to very large shared memory segments. Under most circumstances, tuning of this variable is not necessary.

Commitment Level: Unstable

default_tsb_size

Description: Selects size of the initial translation storage buffers (TSBs) allocated to all processes.

Data Type: Integer

Default: Default is 0 (8 KB), which corresponds to 512 entries

Range: Possible values are:

Value	Description
0	8 KB
1	16 KB
3	32 KB
4	128 KB
5	256 KB
6	512 KB
7	1 Mbyte

Dynamic?: Yes

Validation: None

When to Change: Generally, you do not need to change this value. However, doing so might provide some advantages if the majority of processes on the system have a larger than average working set, or if resident set size (RSS) sizing is disabled.

Commitment Level: Unstable

Change History: For information, see "`default_tsb_size` (Solaris 10 Releases)" on page 181.

enable_tsb_rss_sizing

Description	Enables a resident set size (RSS) based TSB sizing heuristic.
Data Type	Boolean
Default	1 (TSBs can be resized)
Range	0 (TSBs remain at `tsb_default_size`) or 1 (TSBs can be resized) If set to 0, then `tsb_rss_factor` is ignored.
Dynamic?	Yes
Validation	Yes
When to Change	Can be set to 0 to prevent growth of the TSBs. Under most circumstances, this parameter should be left at the default setting.
Commitment Level	Unstable
Change History	For information, see "`enable_tsb_rss_sizing` (Solaris 10 Releases)" on page 181.

tsb_rss_factor

Description	Controls the RSS to TSB span ratio of the RSS sizing heuristic. This factor divided by 512 yields the percentage of the TSB span which must be resident in memory before the TSB is considered as a candidate for resizing.
Data Type	Integer
Default	384, resulting in a value of 75%. Thus, when the TSB is 3/4 full, its size will be increased. Note that some virtual addresses typically map to the same slot in the TSB. Therefore, conflicts can occur before the TSB is at 100% full.
Range	0 to 512
Dynamic?	Yes
Validation	None
When to Change	If the system is experiencing an excessive number of traps due to TSB misses, for example, due to virtual address conflicts in the TSB, you might consider decreasing this value toward 0. For example, changing `tsb_rss_factor` to 256 (effectively, 50%) instead of 384 (effectively, 75%) might help eliminate virtual address

	conflicts in the TSB in some cases, but will use more kernel memory, particularly on a heavily loaded system. TSB activity can be monitored with the `trapstat -T` command.
Commitment Level	Unstable
Change History	For information, see "`tsb_rss_factor` (Solaris 10 Releases)" on page 181.

Locality Group Parameters

This section provides generic memory tunables, which apply to any SPARC or x86 system that uses a Non-Uniform Memory Architecture (NUMA).

lpg_alloc_prefer

Description	Controls a heuristic for allocation of large memory pages when the requested page size is not immediately available in the local memory group, but could be satisfied from a remote memory group. By default, the Oracle Solaris OS allocates a remote large page if local free memory is fragmented, but remote free memory is not. Setting this parameter to 1 indicates that additional effort should be spent attempting to allocate larger memory pages locally, potentially moving smaller pages around to coalesce larger pages in the local memory group.
Data Type	Boolean
Default	0 (Prefer remote allocation if local free memory is fragmented and remote free memory is not)
Range	0 (Prefer remote allocation if local free memory is fragmented and remote free memory is not) 1 (Prefer local allocation whenever possible, even if local free memory is fragmented and remote free memory is not)
Dynamic?	No
Validation	None
When to Change	This parameter might be set to 1 if long-running programs on the system tend to allocate memory that is accessed by a single program, or if memory that is accessed by a group of programs is known to be

running in the same locality group (lgroup). In these circumstances, the extra cost of page coalesce operations can be amortized over the long run of the programs.

This parameter might be left at the default value (0) if multiple programs tend to share memory across different locality groups, or if pages tend to be used for short periods of time. In these circumstances, quick allocation of the requested size tends to be more important than allocation in a particular location.

TLB miss activity might be observed by using the `trapstat -T` command.

Commitment Level: Uncommitted

lgrp_mem_default_policy

Description: This variable reflects the default memory allocation policy used by the Oracle Solaris OS. This variable is an integer, and its value should correspond to one of the policies listed in the `sys/lgrp.h` file.

Data Type: Integer

Default: 1, `LGRP_MEM_POLICY_NEXT` indicating that memory allocation defaults to the home lgroup of the thread performing the memory allocation.

Range: Possible values are:

Value	Description	Comment
0	`LGRP_MEM_POLICY_DEFAULT`	use system default policy
1	`LGRP_MEM_POLICY_NEXT`	next to allocating thread's home lgroup
2	`LGRP_MEM_POLICY_RANDOM_PROC`	randomly across process
3	`LGRP_MEM_POLICY_RANDOM_PSET`	randomly across processor set
4	`LGRP_MEM_POLICY_RANDOM`	randomly across all lgroups
5	`LGRP_MEM_POLICY_ROUNDROBIN`	round robin across all lgroups
6	`LGRP_MEM_POLICY_NEXT_CPU`	near next CPU to touch memory

Dynamic?: No

Validation: None

When to Change	For applications that are sensitive to memory latencies due to allocations that occur from remote versus local memory on systems that use NUMA.
Commitment Level	Uncommitted

lgrp_mem_pset_aware

Description	If a process is running within a user processor set, this variable determines whether *randomly* placed memory for the process is selected from among all the lgroups in the system or only from those lgroups that are spanned by the processors in the processor set. For more information about creating processor sets, see `psrset(1M)`.
Data Type	Boolean
Default	0, the Oracle Solaris OS selects memory from all the lgroups in the system
Range	■ 0, the Oracle Solaris OS selects memory from all the lgroups in the system (default) ■ 1, try selecting memory only from those lgroups that are spanned by the processors in the processor set. If the first attempt fails, memory can be allocated in any lgroup.
Dynamic?	No
Validation	None
When to Change	Setting this value to a value of one (1) might lead to more reproducible performance when processor sets are used to isolate applications from one another.
Commitment Level	Uncommitted

Solaris Volume Manager Parameters

md_mirror:md_resync_bufsz

Description	Sets the size of the buffer used for resynchronizing RAID 1 volumes (mirrors) as the number of 512-byte blocks in the buffer. Setting larger values can increase resynchronization speed.
Data Type	Integer
Default	The default value is 128. Larger systems could use higher values to increase mirror resynchronization speed.
Range	128 to 2048
Units	Blocks (512 bytes)
Dynamic?	No
Validation	None
When to Change	If you use Solaris Volume Manager RAID 1 volumes (mirrors), and you want to increase the speed of mirror resynchronizations. Assuming that you have adequate memory for overall system performance, you can increase this value without causing other performance problems. If you need to increase the speed of mirror resynchronizations, increase the value of this parameter incrementally (using 128-block increments) until performance is satisfactory. On fairly large or new systems, a value of 2048 seems to be optimal. High values on older systems might hang the system.
Commitment Level	Unstable

md:mirrored_root_flag

Description	Overrides Solaris Volume Manager requirements for replica quorum and forces Solaris Volume Manager to start if any valid state database replicas are available. The default value is disabled, which requires that a majority of all replicas are available and synchronized before Solaris Volume Manager will start.
Data Type	Boolean values

Default	0 (disabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	No
Validation	None
When to Change	Use of this parameter is not supported. Some people using Solaris Volume Manager accept the risk of enabling this parameter if all three of the following conditions apply: ■ When root (/) or other system-critical file systems are mirrored ■ Only two disks or controllers are available ■ An unattended reboot of the system is required If this parameter is enabled, the system might boot with a stale replica that inaccurately represents the system state (including which mirror sides are good or in Maintenance state). This situation could result in data corruption or system corruption. Change this parameter only if system availability is more important than data consistency and integrity. Closely monitor the system for any failures. You can mitigate the risk by keeping the number of failed, Maintenance, or hot-swapped volumes as low as possible. For more information about state database replicas, see Chapter 6, "State Database (Overview)," in *Solaris Volume Manager Administration Guide*.
Commitment Level	Unstable

◆ ◆ ◆ CHAPTER 3

NFS Tunable Parameters

This section describes the NFS tunable parameters.

Where to Find Tunable Parameter Information

Tunable Parameter	For Information
Oracle Solaris kernel tunables	Chapter 2, "Oracle Solaris Kernel Tunable Parameters"
Internet Protocol Suite tunable parameters	Chapter 4, "Internet Protocol Suite Tunable Parameters"
Network Cache and Accelerator (NCA) tunable parameters	Chapter 5, "Network Cache and Accelerator Tunable Parameters"

Tuning the NFS Environment

You can define NFS parameters in the `/etc/system` file, which is read during the boot process. Each parameter includes the name of its associated kernel module. For more information, see "Tuning an Oracle Solaris System" on page 23.

Caution – The names of the parameters, the modules that they reside in, and the default values can change between releases. Check the documentation for the version of the active SunOS release before making changes or applying values from previous releases.

NFS Module Parameters

This section describes parameters related to the NFS kernel module.

nfs:nfs3_pathconf_disable_cache

Description	Controls the caching of `pathconf` information for NFS Version 3 mounted file systems.
Data Type	Integer (32-bit)
Default	0 (caching enabled)
Range	0 (caching enabled) or 1 (caching disabled)
Units	Boolean values
Dynamic?	Yes
Validation	None
When to Change	The `pathconf` information is cached on a per file basis. However, if the server can change the information for a specific file dynamically, use this parameter to disable caching. There is no mechanism for the client to validate its cache entry.
Commitment Level	Unstable

nfs:nfs4_pathconf_disable_cache

Description	Controls the caching of `pathconf` information for NFS Version 4 mounted file systems.
Data Type	Integer (32-bit)
Default	0 (caching enabled)
Range	0 (caching enabled) or 1 (caching disabled)
Units	Boolean values
Dynamic?	Yes

Validation	None
When to Change	The `pathconf` information is cached on a per file basis. However, if the server can change the information for a specific file dynamically, use this parameter to disable caching. There is no mechanism for the client to validate its cache entry.
Commitment Level	Unstable

nfs:nfs_allow_preepoch_time

Description	Controls whether files with incorrect or *negative* time stamps should be made visible on the client. Historically, neither the NFS client nor the NFS server would do any range checking on the file times being returned. The over-the-wire timestamp values are unsigned and 32-bits long. So, all values have been legal. However, on a system running a 32-bit Solaris kernel, the timestamp values are signed and 32-bits long. Thus, it would be possible to have a timestamp representation that appeared to be prior to January 1, 1970, or *pre-epoch*. The problem on a system running a 64-bit Solaris kernel is slightly different. The timestamp values on the 64-bit Solaris kernel are signed and 64-bits long. It is impossible to determine whether a time field represents a full 32-bit time or a negative time, that is, a time prior to January 1, 1970. It is impossible to determine whether to sign extend a time value when converting from 32 bits to 64 bits. The time value should be sign extended if the time value is truly a negative number. However, the time value should not be sign extended if it does truly represent a full 32-bit time value. This problem is resolved by simply disallowing full 32-bit time values.
Data Type	Integer (32-bit)
Default	0 (32-bit time stamps disabled)
Range	0 (32-bit time stamps disabled) or 1 (32-bit time stamps enabled)
Units	Boolean values
Dynamic?	Yes

Validation	None
When to Change	Even during normal operation, it is possible for the timestamp values on some files to be set very far in the future or very far in the past. If access to these files is preferred using NFS mounted file systems, set this parameter to 1 to allow the timestamp values to be passed through unchecked.
Commitment Level	Unstable

nfs:nfs_cots_timeo

Description	Controls the default RPC timeout for NFS version 2 mounted file systems using connection-oriented transports such as TCP for the transport protocol.
Data Type	Signed integer (32-bit)
Default	600 (60 seconds)
Range	0 to 2^{31} - 1
Units	10th of seconds
Dynamic?	Yes, but the RPC timeout for a file system is set when the file system is mounted. To affect a particular file system, unmount and mount the file system after changing this parameter.
Validation	None
When to Change	TCP does a good job ensuring requests and responses are delivered appropriately. However, if the round-trip times are very large in a particularly slow network, the NFS version 2 client might time out prematurely. Increase this parameter to prevent the client from timing out incorrectly. The range of values is very large, so increasing this value too much might result in situations where a retransmission is not detected for long periods of time.
Commitment Level	Unstable

nfs:nfs3_cots_timeo

Description	Controls the default RPC timeout for NFS version 3 mounted file systems using connection-oriented transports such as TCP for the transport protocol.

Data Type	Signed integer (32-bit)
Default	600 (60 seconds)
Range	0 to 2^{31} - 1
Units	10th of seconds
Dynamic?	Yes, but the RPC timeout for a file system is set when the file system is mounted. To affect a particular file system, unmount and mount the file system after changing this parameter.
Validation	None
When to Change	TCP does a good job ensuring requests and responses are delivered appropriately. However, if the round-trip times are very large in a particularly slow network, the NFS version 3 client might time out prematurely. Increase this parameter to prevent the client from timing out incorrectly. The range of values is very large, so increasing this value too much might result in situations where a retransmission is not detected for long periods of time.
Commitment Level	Unstable

nfs:nfs4_cots_timeo

Description	Controls the default RPC timeout for NFS version 4 mounted file systems using connection-oriented transports such as TCP for the transport protocol. The NFS Version 4 protocol specification disallows retransmission over the same TCP connection. Thus, this parameter primarily controls how quickly the client responds to certain events, such as detecting a forced unmount operation or detecting how quickly the server fails over to a new server.
Data Type	Signed integer (32-bit)
Default	600 (60 seconds)
Range	0 to 2^{31} - 1
Units	10th of seconds
Dynamic?	Yes, but this parameter is set when the file system is mounted. To affect a particular file system, unmount and mount the file system after changing this parameter.

Validation	None
When to Change	TCP does a good job ensuring requests and responses are delivered appropriately. However, if the round-trip times are very large in a particularly slow network, the NFS version 4 client might time out prematurely. Increase this parameter to prevent the client from timing out incorrectly. The range of values is very large, so increasing this value too much might result in situations where a retransmission is not detected for long periods of time.
Commitment Level	Unstable

nfs:nfs_do_symlink_cache

Description	Controls whether the contents of symbolic link files are cached for NFS version 2 mounted file systems.
Data Type	Integer (32–bit)
Default	1 (caching enabled)
Range	0 (caching disabled) or 1 (caching enabled)
Units	Boolean values
Dynamic?	Yes
Validation	None
When to Change	If a server changes the contents of a symbolic link file without updating the modification timestamp on the file or if the granularity of the timestamp is too large, then changes to the contents of the symbolic link file might not be visible on the client for extended periods. In this case, use this parameter to disable the caching of symbolic link contents. Doing so makes the changes immediately visible to applications running on the client.
Commitment Level	Unstable

nfs:nfs3_do_symlink_cache

Description	Controls whether the contents of symbolic link files are cached for NFS version 3 mounted file systems.
Data Type	Integer (32-bit)

Default	1 (caching enabled)
Range	0 (caching disabled) or 1 (caching enabled)
Units	Boolean values
Dynamic?	Yes
Validation	None
When to Change	If a server changes the contents of a symbolic link file without updating the modification timestamp on the file or if the granularity of the timestamp is too large, then changes to the contents of the symbolic link file might not be visible on the client for extended periods. In this case, use this parameter to disable the caching of symbolic link contents. Doing so makes the changes immediately visible to applications running on the client.
Commitment Level	Unstable

nfs:nfs4_do_symlink_cache

Description	Controls whether the contents of symbolic link files are cached for NFS version 4 mounted file systems.
Data Type	Integer (32-bit)
Default	1 (caching enabled)
Range	0 (caching disabled) or 1 (caching enabled)
Units	Boolean values
Dynamic?	Yes
Validation	None
When to Change	If a server changes the contents of a symbolic link file without updating the modification timestamp on the file or if the granularity of the timestamp is too large, then changes to the contents of the symbolic link file might not be visible on the client for extended periods. In this case, use this parameter to disable the caching of symbolic link contents. Doing so makes the changes immediately visible to applications running on the client.
Commitment Level	Unstable

nfs:nfs_dynamic

Description	Controls whether a feature known as *dynamic retransmission* is enabled for NFS version 2 mounted file systems using connectionless transports such as UDP. This feature attempts to reduce retransmissions by monitoring server response times and then adjusting RPC timeouts and read- and write- transfer sizes.
Data Type	Integer (32-bit)
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes, but this parameter is set per file system at mount time. To affect a particular file system, unmount and mount the file system after changing this parameter.
Validation	None
When to Change	Do not change this parameter.
Commitment Level	Unstable

nfs:nfs3_dynamic

Description	Controls whether a feature known as *dynamic retransmission* is enabled for NFS version 3 mounted file systems using connectionless transports such as UDP. This feature attempts to reduce retransmissions by monitoring server response times and then adjusting RPC timeouts and read- and write- transfer sizes.
Data Type	Integer (32-bit)
Default	0 (disabled)
Range	0 (disabled) or 1 (enabled)
Units	Boolean values
Dynamic?	Yes, but this parameter is set per file system at mount time. To affect a particular file system, unmount and mount the file system after changing this parameter.
Validation	None
When to Change	Do not change this parameter.
Commitment Level	Unstable

nfs:nfs_lookup_neg_cache

Description	Controls whether a negative name cache is used for NFS version 2 mounted file systems. This negative name cache records file names that were looked up, but not found. The cache is used to avoid over-the-network look-up requests made for file names that are already known to not exist.
Data Type	Integer (32-bit)
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Units	Boolean values
Dynamic?	Yes
Validation	None
When to Change	For the cache to perform correctly, negative entries must be strictly verified before they are used. This consistency mechanism is relaxed slightly for read-only mounted file systems. It is assumed that the file system on the server is not changing or is changing very slowly, and that it is okay for such changes to propagate slowly to the client. The consistency mechanism becomes the normal attribute cache mechanism in this case. If file systems are mounted read-only on the client, but are expected to change on the server and these changes need to be seen immediately by the client, use this parameter to disable the negative cache. If you disable the `nfs:nfs_disable_rddir_cache` parameter, you should probably also disable this parameter. For more information, see "`nfs:nfs_disable_rddir_cache`" on page 111.
Commitment Level	Unstable

nfs:nfs3_lookup_neg_cache

Description	Controls whether a negative name cache is used for NFS version 3 mounted file systems. This negative name cache records file names that were looked up, but were not found. The cache is used to avoid over-the-network look-up requests made for file names that are already known to not exist.
Data Type	Integer (32-bit)

Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Units	Boolean values
Dynamic?	Yes
Validation	None
When to Change	For the cache to perform correctly, negative entries must be strictly verified before they are used. This consistency mechanism is relaxed slightly for read-only mounted file systems. It is assumed that the file system on the server is not changing or is changing very slowly, and that it is okay for such changes to propagate slowly to the client. The consistency mechanism becomes the normal attribute cache mechanism in this case. If file systems are mounted read-only on the client, but are expected to change on the server and these changes need to be seen immediately by the client, use this parameter to disable the negative cache. If you disable the `nfs:nfs_disable_rddir_cache` parameter, you should probably also disable this parameter. For more information, see "`nfs:nfs_disable_rddir_cache`" on page 111.
Commitment Level	Unstable

nfs:nfs4_lookup_neg_cache

Description	Controls whether a negative name cache is used for NFS version 4 mounted file systems. This negative name cache records file names that were looked up, but were not found. The cache is used to avoid over-the-network look-up requests made for file names that are already known to not exist.
Data Type	Integer (32-bit)
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Units	Boolean values
Dynamic?	Yes
Validation	None

When to Change: For the cache to perform correctly, negative entries must be strictly verified before they are used. This consistency mechanism is relaxed slightly for read-only mounted file systems. It is assumed that the file system on the server is not changing or is changing very slowly, and that it is okay for such changes to propagate slowly to the client. The consistency mechanism becomes the normal attribute cache mechanism in this case.

If file systems are mounted read-only on the client, but are expected to change on the server and these changes need to be seen immediately by the client, use this parameter to disable the negative cache.

If you disable the `nfs:nfs_disable_rddir_cache` parameter, you should probably also disable this parameter. For more information, see "`nfs:nfs_disable_rddir_cache`" on page 111.

Commitment Level: Unstable

nfs:nfs_max_threads

Description: Controls the number of kernel threads that perform asynchronous I/O for the NFS version 2 client. Because NFS is based on RPC and RPC is inherently synchronous, separate execution contexts are required to perform NFS operations that are asynchronous from the calling thread.

The operations that can be executed asynchronously are read for read-ahead, readdir for readdir read-ahead, write for putpage and pageio operations, commit, and inactive for cleanup operations that the client performs when it stops using a file.

Data Type: Integer (16-bit)

Default: 8

Range: 0 to 2^{15} - 1

Units: Threads

Dynamic?: Yes, but this parameter is set per file system at mount time. To affect a particular file system, unmount and mount the file system after changing this parameter.

Validation: None

When to Change: To increase or reduce the number of simultaneous I/O operations that are outstanding at any given time. For example, for a very low bandwidth network, you might want to decrease this value so that the

NFS client does not overload the network. Alternately, if the network is very high bandwidth, and the client and server have sufficient resources, you might want to increase this value. Doing so can more effectively utilize the available network bandwidth, and the client and server resources.

Commitment Level: Unstable

nfs:nfs3_max_threads

Description: Controls the number of kernel threads that perform asynchronous I/O for the NFS version 3 client. Because NFS is based on RPC and RPC is inherently synchronous, separate execution contexts are required to perform NFS operations that are asynchronous from the calling thread.

The operations that can be executed asynchronously are read for read-ahead, readdir for readdir read-ahead, write for putpage and pageio requests, and commit.

Data Type: Integer (16-bit)

Default: 8

Range: 0 to 2^{15} - 1

Units: Threads

Dynamic?: Yes, but this parameter is set per file system at mount time. To affect a particular file system, unmount and mount the file system after changing this parameter.

Validation: None

When to Change: To increase or reduce the number of simultaneous I/O operations that are outstanding at any given time. For example, for a very low bandwidth network, you might want to decrease this value so that the NFS client does not overload the network. Alternately, if the network is very high bandwidth, and the client and server have sufficient resources, you might want to increase this value. Doing so can more effectively utilize the available network bandwidth, and the client and server resources.

Commitment Level: Unstable

nfs:nfs4_max_threads

Description	Controls the number of kernel threads that perform asynchronous I/O for the NFS version 4 client. Because NFS is based on RPC and RPC is inherently synchronous, separate execution contexts are required to perform NFS operations that are asynchronous from the calling thread. The operations that can be executed asynchronously are read for read-ahead, write-behind, directory read-ahead, and cleanup operations that the client performs when it stops using a file.
Data Type	Integer (16-bit)
Default	8
Range	0 to 2^{15} - 1
Units	Threads
Dynamic?	Yes, but this parameter is set per file system at mount time. To affect a particular file system, unmount and mount the file system after changing this parameter.
Validation	None
When to Change	To increase or reduce the number of simultaneous I/O operations that are outstanding at any given time. For example, for a very low bandwidth network, you might want to decrease this value so that the NFS client does not overload the network. Alternately, if the network is very high bandwidth, and the client and server have sufficient resources, you might want to increase this value. Doing so can more effectively utilize the available network bandwidth, and the client and server resources.
Commitment Level	Unstable

nfs:nfs_nra

Description	Controls the number of read-ahead operations that are queued by the NFS version 2 client when sequential access to a file is discovered. These read-ahead operations increase concurrency and read throughput. Each read-ahead request is generally for one logical block of file data.
Data Type	Integer (32-bit)
Default	4

Range	0 to 2^{31} - 1
Units	Logical blocks.
Dynamic?	Yes
Validation	None
When to Change	To increase or reduce the number of read-ahead requests that are outstanding for a specific file at any given time. For example, for a very low bandwidth network or on a low memory client, you might want to decrease this value so that the NFS client does not overload the network or the system memory. Alternately, if the network is very high bandwidth, and the client and server have sufficient resources, you might want to increase this value. Doing so can more effectively utilize the available network bandwidth, and the client and server resources.
Commitment Level	Unstable

nfs:nfs3_nra

Description	Controls the number of read-ahead operations that are queued by the NFS version 3 client when sequential access to a file is discovered. These read-ahead operations increase concurrency and read throughput. Each read-ahead request is generally for one logical block of file data.
Data Type	Integer (32-bit)
Default	4
Range	0 to 2^{31} - 1
Units	Logical blocks. (See "nfs:nfs3_bsize" on page 112.)
Dynamic?	Yes
Validation	None
When to Change	To increase or reduce the number of read-ahead requests that are outstanding for a specific file at any given time. For example, for a very low bandwidth network or on a low memory client, you might want to decrease this value so that the NFS client does not overload the network or the system memory. Alternately, if the network is very high bandwidth and the client and server have sufficient resources, you might want to increase this value. Doing so can more effectively utilize the available network bandwidth, and the client and server resources.
Commitment Level	Unstable

Change History — For information, see "nfs:nfs3_nra (Solaris 10 Releases)" on page 182.

nfs:nfs4_nra

Description	Controls the number of read-ahead operations that are queued by the NFS version 4 client when sequential access to a file is discovered. These read-ahead operations increase concurrency and read throughput. Each read-ahead request is generally for one logical block of file data.
Data Type	Integer (32-bit)
Default	4
Range	0 to 2^{31} - 1
Units	Logical blocks. (See "nfs:nfs4_bsize" on page 112.)
Dynamic?	Yes
Validation	None
When to Change	To increase or reduce the number of read-ahead requests that are outstanding for a specific file at any given time. For example, for a very low bandwidth network or on a low memory client, you might want to decrease this value so that the NFS client does not overload the network or the system memory. Alternately, if the network is very high bandwidth, and the client and server have sufficient resources, you might want to increase this value. Doing so can more effectively utilize the available network bandwidth, and the client and server resources.
Commitment Level	Unstable

nfs:nrnode

Description	Controls the size of the `rnode` cache on the NFS client. The `rnode`, used by both NFS version 2, 3, and 4 clients, is the central data structure that describes a file on the NFS client. The `rnode` contains the file handle that identifies the file on the server. The `rnode` also contains pointers to various caches used by the NFS client to avoid network calls to the server. Each `rnode` has a one-to-one association with a `vnode`. The `vnode` caches file data.

	The NFS client attempts to maintain a minimum number of rnodes to attempt to avoid destroying cached data and metadata. When an rnode is reused or freed, the cached data and metadata must be destroyed.
Data Type	Integer (32-bit)
Default	The default setting of this parameter is 0, which means that the value of nrnode should be set to the value of the ncsize parameter. Actually, any non positive value of nrnode results in nrnode being set to the value of ncsize.
Range	1 to 2^{31} - 1
Units	rnodes
Dynamic?	No. This value can only be changed by adding or changing the parameter in the /etc/system file, and then rebooting the system.
Validation	The system enforces a maximum value such that the rnode cache can only consume 25 percent of available memory.
When to Change	Because rnodes are created and destroyed dynamically, the system tends to settle upon a *nrnode*-size cache, automatically adjusting the size of the cache as memory pressure on the system increases or as more files are simultaneously accessed. However, in certain situations, you could set the value of nrnode if the mix of files being accessed can be predicted in advance. For example, if the NFS client is accessing a few very large files, you could set the value of nrnode to a small number so that system memory can cache file data instead of rnodes. Alternately, if the client is accessing many small files, you could increase the value of nrnode to optimize for storing file metadata to reduce the number of network calls for metadata. Although it is not recommended, the rnode cache can be effectively disabled by setting the value of nrnode to 1. This value instructs the client to only cache 1 rnode, which means that it is reused frequently.
Commitment Level	Unstable

nfs:nfs_shrinkreaddir

Description	Some older NFS servers might incorrectly handle NFS version 2 READDIR requests for more than 1024 bytes of directory information. This problem is due to a bug in the server implementation. However, this parameter contains a workaround in the NFS version 2 client.

When this parameter is enabled, the client does not generate a `READDIR` request for larger than 1024 bytes of directory information. If this parameter is disabled, then the over-the-wire size is set to the lesser of either the size passed in by using the `getdents` system call or by using `NFS_MAXDATA`, which is 8192 bytes. For more information, see `getdents(2)`.

Data Type	Integer (32-bit)
Default	0 (disabled)
Range	0 (disabled) or 1 (enabled)
Units	Boolean values
Dynamic?	Yes
Validation	None
When to Change	Examine the value of this parameter if an older NFS version 2 only server is used and interoperability problems occur when the server tries to read directories. Enabling this parameter might cause a slight decrease in performance for applications that read directories.
Commitment Level	Unstable

nfs:nfs3_shrinkreaddir

Description

Some older NFS servers might incorrectly handle NFS version 3 `READDIR` requests for more than 1024 bytes of directory information. This problem is due to a bug in the server implementation. However, this parameter contains a workaround in the NFS version 3 client.

When this parameter is enabled, the client does not generate a `READDIR` request for larger than 1024 bytes of directory information. If this parameter is disabled, then the over-the-wire size is set to the minimum of either the size passed in by using the `getdents` system call or by using `MAXBSIZE`, which is 8192 bytes. For more information, see `getdents(2)`.

Data Type	Integer (32-bit)
Default	0 (disabled)
Range	0 (disabled) or 1 (enabled)
Units	Boolean values
Dynamic?	Yes

Validation	None
When to Change	Examine the value of this parameter if an older NFS version 3 only server is used and interoperability problems occur when the server tries to read directories. Enabling this parameter might cause a slight decrease in performance for applications that read directories.
Commitment Level	Unstable

nfs:nfs_write_error_interval

Description	Controls the time duration in between logging `ENOSPC` and `EDQUOT` write errors received by the NFS client. This parameter affects NFS version 2, 3, and 4 clients.
Data Type	Long integer (32 bits on 32-bit platforms and 64 bits on 64-bit platforms)
Default	5 seconds
Range	0 to 2^{31} - 1 on 32-bit platforms 0 to 2^{63} - 1 on 64-bit platforms
Units	Seconds
Dynamic?	Yes
Validation	None
When to Change	Increase or decrease the value of this parameter in response to the volume of messages being logged by the client. Typically, you might want to increase the value of this parameter to decrease the number of `out of space` messages being printed when a full file system on a server is being actively used.
Commitment Level	Unstable

nfs:nfs_write_error_to_cons_only

Description	Controls whether NFS write errors are logged to the system console and `syslog` or to the system console only. This parameter affects messages for NFS version 2, 3, and 4 clients.
Data Type	Integer (32-bit)
Default	0 (system console and `syslog`)

Range	0 (system console and `syslog`) or 1 (system console)
Units	Boolean values
Dynamic?	Yes
Validation	None
When to Change	Examine the value of this parameter to avoid filling up the file system containing the messages logged by the `syslogd` daemon. When this parameter is enabled, messages are printed on the system console only and are not copied to the `syslog` messages file.
Commitment Level	Unstable

nfs:nfs_disable_rddir_cache

Description	Controls the use of a cache to hold responses from `READDIR` and `READDIRPLUS` requests. This cache avoids over-the-wire calls to the server to retrieve directory information.
Data Type	Integer (32-bit)
Default	0 (caching enabled)
Range	0 (caching enabled) or 1 (caching disabled)
Units	Boolean values
Dynamic?	Yes
Validation	None
When to Change	Examine the value of this parameter if interoperability problems develop due to a server that does not update the modification time on a directory when a file or directory is created in it or removed from it. The symptoms are that new names do not appear in directory listings after they have been added to the directory or that old names do not disappear after they have been removed from the directory. This parameter controls the caching for NFS version 2, 3, and 4 mounted file systems. This parameter applies to all NFS mounted file systems, so caching cannot be disabled or enabled on a per file system basis. If you disable this parameter, you should also disable the following parameters to prevent bad entries in the DNLC negative cache: ■ "`nfs:nfs_lookup_neg_cache`" on page 101 ■ "`nfs:nfs3_lookup_neg_cache`" on page 101

- "nfs:nfs4_lookup_neg_cache" on page 102

Commitment Level: Unstable

nfs:nfs3_bsize

Description: Controls the logical block size used by the NFS version 3 client. This block size represents the amount of data that the client attempts to read from or write to the server when it needs to do an I/O.

Data Type: Unsigned integer (32-bit)

Default: 32,768 (32 KB)

Range: 0 to 2^{31} - 1

Units: Bytes

Dynamic?: Yes, but the block size for a file system is set when the file system is mounted. To affect a particular file system, unmount and mount the file system after changing this parameter.

Validation: None. Setting this parameter too low or too high might cause the system to malfunction. Do not set this parameter to anything less than `PAGESIZE` for the specific platform. Do not set this parameter too high because it might cause the system to hang while waiting for memory allocations to be granted.

When to Change: Examine the value of this parameter when attempting to change the maximum data transfer size. Change this parameter in conjunction with the `nfs:nfs3_max_transfer_size` parameter. If larger transfers are preferred, increase both parameters. If smaller transfers are preferred, then just reducing this parameter should suffice.

Commitment Level: Unstable

nfs:nfs4_bsize

Description: Controls the logical block size used by the NFS version 4 client. This block size represents the amount of data that the client attempts to read from or write to the server when it needs to do an I/O.

Data Type: Unsigned integer (32-bit)

Default: 32,768 (32 KB)

Range: 0 to 2^{31} - 1

Units	Bytes
Dynamic?	Yes, but the block size for a file system is set when the file system is mounted. To affect a particular file system, unmount and mount the file system after changing this parameter.
Validation	None. Setting this parameter too low or too high might cause the system to malfunction. Do not set this parameter to anything less than `PAGESIZE` for the specific platform. Do not set this parameter too high because it might cause the system to hang while waiting for memory allocations to be granted.
When to Change	Examine the value of this parameter when attempting to change the maximum data transfer size. Change this parameter in conjunction with the `nfs:nfs4_max_transfer_size` parameter. If larger transfers are preferred, increase both parameters. If smaller transfers are preferred, then just reducing this parameter should suffice.
Commitment Level	Unstable

nfs:nfs_async_clusters

Description	Controls the mix of asynchronous requests that are generated by the NFS version 2 client. The four types of asynchronous requests are read-ahead, putpage, pageio, and readdir-ahead. The client attempts to round-robin between these different request types to attempt to be fair and not starve one request type in favor of another. However, the functionality in some NFS version 2 servers such as write gathering depends upon certain behaviors of existing NFS Version 2 clients. In particular, this functionality depends upon the client sending out multiple WRITE requests at about the same time. If one request is taken out of the queue at a time, the client would be defeating this server functionality designed to enhance performance for the client. Thus, use this parameter to control the number of requests of each request type that are sent out before changing types.
Data Type	Unsigned integer (32-bit)
Default	1
Range	0 to 2^{31} - 1
Units	Asynchronous requests

Dynamic?	Yes, but the cluster setting for a file system is set when the file system is mounted. To affect a particular file system, unmount and mount the file system after changing this parameter.
Validation	None. However, setting the value of this parameter to 0 causes all of the queued requests of a particular request type to be processed before moving on to the next type. This effectively disables the fairness portion of the algorithm.
When to Change	To increase the number of each type of asynchronous request that is generated before switching to the next type. Doing so might help with server functionality that depends upon clusters of requests coming from the client.
Commitment Level	Unstable

nfs:nfs3_async_clusters

Description	Controls the mix of asynchronous requests that are generated by the NFS version 3 client. The five types of asynchronous requests are read-ahead, putpage, pageio, readdir-ahead, and commit. The client attempts to round-robin between these different request types to attempt to be fair and not starve one request type in favor of another. However, the functionality in some NFS version 3 servers such as write gathering depends upon certain behaviors of existing NFS version 3 clients. In particular, this functionality depends upon the client sending out multiple WRITE requests at about the same time. If one request is taken out of the queue at a time, the client would be defeating this server functionality designed to enhance performance for the client. Thus, use this parameter to control the number of requests of each request type that are sent out before changing types.
Data Type	Unsigned integer (32-bit)
Default	1
Range	0 to 2^{31} - 1
Units	Asynchronous requests
Dynamic?	Yes, but the cluster setting for a file system is set when the file system is mounted. To affect a particular file system, unmount and mount the file system after changing this parameter.

Validation	None. However, setting the value of this parameter to 0 causes all of the queued requests of a particular request type to be processed before moving on to the next type. This value effectively disables the fairness portion of the algorithm.
When to Change	To increase the number of each type of asynchronous operation that is generated before switching to the next type. Doing so might help with server functionality that depends upon clusters of operations coming from the client.
Commitment Level	Unstable

nfs:nfs4_async_clusters

Description	Controls the mix of asynchronous requests that are generated by the NFS version 4 client. The six types of asynchronous requests are read-ahead, putpage, pageio, readdir-ahead, commit, and inactive. The client attempts to round-robin between these different request types to attempt to be fair and not starve one request type in favor of another. However, the functionality in some NFS version 4 servers such as write gathering depends upon certain behaviors of existing NFS version 4 clients. In particular, this functionality depends upon the client sending out multiple WRITE requests at about the same time. If one request is taken out of the queue at a time, the client would be defeating this server functionality designed to enhance performance for the client. Thus, use this parameter to control the number of requests of each request type that are sent out before changing types.
Data Type	Unsigned integer (32-bit)
Default	1
Range	0 to 2^{31} - 1
Units	Asynchronous requests
Dynamic?	Yes, but the cluster setting for a file system is set when the file system is mounted. To affect a particular file system, unmount and mount the file system after changing this parameter.
Validation	None. However, setting the value of this parameter to 0 causes all of the queued requests of a particular request type to be processed before moving on to the next type. This effectively disables the fairness portion of the algorithm.

When to Change: To increase the number of each type of asynchronous request that is generated before switching to the next type. Doing so might help with server functionality that depends upon clusters of requests coming from the client.

Commitment Level: Unstable

nfs:nfs_async_timeout

Description: Controls the duration of time that threads, which execute asynchronous I/O requests, sleep with nothing to do before exiting. When there are no more requests to execute, each thread goes to sleep. If no new requests come in before this timer expires, the thread wakes up and exits. If a request does arrive, a thread is woken up to execute requests until there are none again. Then, the thread goes back to sleep waiting for another request to arrive, or for the timer to expire.

Data Type: Integer (32-bit)

Default: 6000 (1 minute expressed as 60 sec * 100Hz)

Range: 0 to 2^{31} - 1

Units: Hz. (Typically, the clock runs at 100Hz.)

Dynamic?: Yes

Validation: None. However, setting this parameter to a non positive value causes these threads exit as soon as there are no requests in the queue for them to process.

When to Change: If the behavior of applications in the system is known precisely and the rate of asynchronous I/O requests can be predicted, it might be possible to tune this parameter to optimize performance slightly in either of the following ways:

- By making the threads expire more quickly, thus freeing up kernel resources more quickly
- By making the threads expire more slowly, thus avoiding thread create and destroy overhead

Commitment Level: Unstable

nfs:nacache

Description	Tunes the number of hash queues that access the file access cache on the NFS client. The file access cache stores file access rights that users have with respect to files that they are trying to access. The cache itself is dynamically allocated. However, the hash queues used to index into the cache are statically allocated. The algorithm assumes that there is one access cache entry per active file and four of these access cache entries per hash bucket. Thus, by default, the value of this parameter is set to the value of the `nrnode` parameter.
Data Type	Integer (32-bit)
Default	The default setting of this parameter is 0. This value means that the value of `nacache` should be set to the value of the `nrnode` parameter.
Range	1 to 2^{31} - 1
Units	Access cache entries
Dynamic?	No. This value can only be changed by adding or changing the parameter in the `/etc/system` file, and then rebooting system.
Validation	None. However, setting this parameter to a negative value will probably cause the system to try to allocate a very large set of hash queues. While trying to do so, the system is likely to hang.
When to Change	Examine the value of this parameter if the basic assumption of one access cache entry per file would be violated. This violation could occur for systems in a timesharing mode where multiple users are accessing the same file at about the same time. In this case, it might be helpful to increase the expected size of the access cache so that the hashed access to the cache stays efficient.
Commitment Level	Unstable

nfs:nfs3_jukebox_delay

Description	Controls the duration of time that the NFS version 3 client waits to transmit a new request after receiving the `NFS3ERR_JUKEBOX` error from a previous request. The `NFS3ERR_JUKEBOX` error is generally returned from the server when the file is temporarily unavailable for some reason. This error is generally associated with hierarchical storage, and CD or tape jukeboxes.
Data Type	Long integer (32 bits on 32-bit platforms and 64 bits on 64-bit platforms)

Default	1000 (10 seconds expressed as 10 sec * 100Hz)
Range	0 to 2^{31} - 1 on 32-bit platforms 0 to 2^{63} - 1 on 64-bit platforms
Units	Hz. (Typically, the clock runs at 100Hz.)
Dynamic?	Yes
Validation	None
When to Change	Examine the value of this parameter and perhaps adjust it to match the behaviors exhibited by the server. Increase this value if the delays in making the file available are long in order to reduce network overhead due to repeated retransmissions. Decrease this value to reduce the delay in discovering that the file has become available.
Commitment Level	Unstable

nfs:nfs3_max_transfer_size

Description	Controls the maximum size of the data portion of an NFS version 3 `READ`, `WRITE`, `READDIR`, or `READDIRPLUS` request. This parameter controls both the maximum size of the request that the server returns as well as the maximum size of the request that the client generates.
Data Type	Integer (32-bit)
Default	1,048,576 (1 Mbyte)
Range	0 to 2^{31} - 1
Units	Bytes
Dynamic?	Yes, but this parameter is set per file system at mount time. To affect a particular file system, unmount and mount the file system after changing this parameter.
Validation	None. However, setting the maximum transfer size on the server to 0 is likely to cause clients to malfunction or just decide not to attempt to talk to the server. There is also a limit on the maximum transfer size when using NFS over the UDP transport. UDP has a hard limit of 64 KB per datagram. This 64 KB must include the RPC header as well as other NFS information, in addition to the data portion of the request. Setting the limit too high might result in errors from UDP and communication problems between the client and the server.

When to Change: To tune the size of data transmitted over the network. In general, the `nfs:nfs3_bsize` parameter should also be updated to reflect changes in this parameter.

For example, when you attempt to increase the transfer size beyond 32 KB, update `nfs:nfs3_bsize` to reflect the increased value. Otherwise, no change in the over-the-wire request size is observed. For more information, see "`nfs:nfs3_bsize`" on page 112.

If you want to use a smaller transfer size than the default transfer size, use the `mount` command's `-wsize` or `-rsize` option on a per-file system basis.

Commitment Level: Unstable

nfs:nfs4_max_transfer_size

Description: Controls the maximum size of the data portion of an NFS version 4 `READ`, `WRITE`, `READDIR`, or `READDIRPLUS` request. This parameter controls both the maximum size of the request that the server returns as well as the maximum size of the request that the client generates.

Data Type: Integer (32-bit)

Default: 32, 768 (32 KB)

Range: 0 to 2^{31} - 1

Units: Bytes

Dynamic?: Yes, but this parameter is set per file system at mount time. To affect a particular file system, unmount and mount the file system after changing this parameter.

Validation: None. However, setting the maximum transfer size on the server to 0 is likely to cause clients to malfunction or just decide not to attempt to talk to the server.

There is also a limit on the maximum transfer size when using NFS over the UDP transport. For more information on the maximum for UDP, see "`nfs:nfs3_max_transfer_size`" on page 118.

When to Change: To tune the size of data transmitted over the network. In general, the `nfs:nfs4_bsize` parameter should also be updated to reflect changes in this parameter.

For example, when you attempt to increase the transfer size beyond 32 KB, update `nfs:nfs4_bsize` to reflect the increased value. Otherwise, no change in the over-the-wire request size is observed. For more information, see "`nfs:nfs4_bsize`" on page 112.

If you want to use a smaller transfer size than the default transfer size, use the `mount` command's `-wsize` or `-rsize` option on a per-file system basis.

Commitment Level: Unstable

nfs:nfs3_max_transfer_size_clts

Description: Controls the maximum size of the data portion of an NFS version 3 `READ`, `WRITE`, `READDIR`, or `READDIRPLUS` request over UDP. This parameter controls both the maximum size of the request that the server returns as well as the maximum size of the request that the client generates.

Data Type: Integer (32-bit)

Default: 32, 768 (32 KB)

Range: 0 to 2^{31} - 1

Units: Bytes

Dynamic?: Yes, but this parameter is set per file system at mount time. To affect a particular file system, unmount and mount the file system after changing this parameter.

Validation: None. However, setting the maximum transfer size on the server to 0 is likely to cause clients to malfunction or just decide not to attempt to talk to the server.

When to Change: Do not change this parameter.

Commitment Level: Unstable

nfs:nfs3_max_transfer_size_cots

Description: Controls the maximum size of the data portion of an NFS version 3 `READ`, `WRITE`, `READDIR`, or `READDIRPLUS` request over TCP. This parameter controls both the maximum size of the request that the server returns as well as the maximum size of the request that the client generates.

Data Type	Integer (32-bit)
Default	1,048,576 bytes
Range	0 to 2^{31} - 1
Units	Bytes
Dynamic?	Yes, but this parameter is set per file system at mount time. To affect a particular file system, unmount and mount the file system after changing this parameter.
Validation	None. However, setting the maximum transfer size on the server to 0 is likely to cause clients to malfunction or just decide not to attempt to talk to the server.
When to Change	Do not change this parameter unless transfer sizes larger than 1 Mbyte are preferred.
Commitment Level	Unstable

nfssrv Module Parameters

This section describes NFS parameters for the `nfssrv` module.

nfssrv:nfs_portmon

Description	Controls some security checking that the NFS server attempts to do to enforce integrity on the part of its clients. The NFS server can check whether the source port from which a request was sent was a *reserved port*. A reserved port has a number less than 1024. For BSD-based systems, these ports are reserved for processes being run by root. This security checking can prevent users from writing their own RPC-based applications that defeat the access checking that the NFS client uses.
Data Type	Integer (32-bit)
Default	0 (security checking disabled)
Range	0 (security checking disabled) or 1 (security checking enabled)
Units	Boolean values
Dynamic?	Yes
Validation	None

When to Change: Use this parameter to prevent malicious users from gaining access to files by using the NFS server that they would not ordinarily have access to. However, the *reserved port* notion is not universally supported. Thus, the security aspects of the check are very weak. Also, not all NFS client implementations bind their transport endpoints to a port number in the reserved range. Thus, interoperability problems might result if the security checking is enabled.

Commitment Level: Unstable

nfssrv:rfs_write_async

Description: Controls the behavior of the NFS version 2 server when it processes `WRITE` requests. The NFS version 2 protocol mandates that all modified data and metadata associated with the `WRITE` request reside on stable storage before the server can respond to the client. NFS version 2 `WRITE` requests are limited to 8192 bytes of data. Thus, each `WRITE` request might cause multiple small writes to the storage subsystem. This can cause a performance problem.

One method to accelerate NFS version 2 `WRITE` requests is to take advantage of a client behavior. Clients tend to send `WRITE` requests in batches. The server can take advantage of this behavior by clustering together the different `WRITE` requests into a single request to the underlying file system. Thus, the data to be written to the storage subsystem can be written in fewer, larger requests. This method can significantly increase the throughput for `WRITE` requests.

Data Type: Integer (32-bit)

Default: 1 (clustering enabled)

Range: 0 (clustering disabled) or 1 (clustering enabled)

Units: Boolean values

Dynamic?: Yes

Validation: None

When to Change: Some very small NFS clients, particularly PC clients, might not batch `WRITE` requests. Thus, the behavior required from the clients might not exist. In addition, the clustering in the NFS version 2 server might just add overhead and slow down performance instead of increasing it.

Commitment Level: Unstable

rpcmod Module Parameters

This section describes NFS parameters for the `rpcmod` module.

rpcmod:clnt_max_conns

Description	Controls the number of TCP connections that the NFS client uses when communicating with each NFS server. The kernel RPC is constructed so that it can multiplex RPCs over a single connection. However, multiple connections can be used, if preferred.
Data Type	Integer (32-bit)
Default	1
Range	1 to 2^{31} - 1
Units	Connections
Dynamic?	Yes
Validation	None
When to Change	In general, one connection is sufficient to achieve full network bandwidth. However, if TCP cannot utilize the bandwidth offered by the network in a single stream, then multiple connections might increase the throughput between the client and the server. Increasing the number of connections doesn't come without consequences. Increasing the number of connections also increases kernel resource usage needed to keep track of each connection.
Commitment Level	Unstable

rpcmod:clnt_idle_timeout

Description	Controls the duration of time on the client that a connection between the client and server is allowed to remain idle before being closed.
Data Type	Long integer (32 bits on 32-bit platforms and 64 bits on 64-bit platforms)
Default	300,000 milliseconds (5 minutes)
Range	0 to 2^{31} - 1 on 32-bit platforms 0 to 2^{63} - 1 on 64-bit platforms

Units	Milliseconds
Dynamic?	Yes
Validation	None
When to Change	Use this parameter to change the time that idle connections are allowed to exist on the client before being closed. You might might want to close connections at a faster rate to avoid consuming system resources.
Commitment Level	Unstable

rpcmod:svc_idle_timeout

Description	Controls the duration of time on the server that a connection between the client and server is allowed to remain idle before being closed.
Data Type	Long integer (32 bits on 32-bit platforms and 64 bits on 64-bit platforms)
Default	360,000 milliseconds (6 minutes)
Range	0 to 2^{31} - 1 on 32-bit platforms 0 to 2^{63} - 1 on 64-bit platforms
Units	Milliseconds
Dynamic?	Yes
Validation	None
When to Change	Use this parameter to change the time that idle connections are allowed to exist on the server before being closed. You might want to close connections at a faster rate to avoid consuming system resources.
Commitment Level	Unstable

rpcmod:svc_default_stksize

Description	Sets the size of the kernel stack for kernel RPC service threads.
Data Type	Integer (32-bit)
Default	The default value is 0. This value means that the stack size is set to the system default.
Range	0 to 2^{31} - 1
Units	Bytes

Dynamic?	Yes, for all new threads that are allocated. The stack size is set when the thread is created. Therefore, changes to this parameter do not affect existing threads but are applied to all new threads that are allocated.
Validation	None
When to Change	Very deep call depths can cause the stack to overflow and cause red zone faults. The combination of a fairly deep call depth for the transport, coupled with a deep call depth for the local file system, can cause NFS service threads to overflow their stacks. Set this parameter to a multiple of the hardware `pagesize` on the platform.
Commitment Level	Unstable

rpcmod:svc_default_max_same_xprt

Description	Controls the maximum number of requests that are processed for each transport endpoint before switching transport endpoints. The kernel RPC works by having a pool of service threads and a pool of transport endpoints. Any one of the service threads can process requests from any one of the transport endpoints. For performance, multiple requests on each transport endpoint are consumed before switching to a different transport endpoint. This approach offers performance benefits while avoiding starvation.
Data Type	Integer (32-bit)
Default	8
Range	0 to 2^{31} - 1
Units	Requests
Dynamic?	Yes, but the maximum number of requests to process before switching transport endpoints is set when the transport endpoint is configured into the kernel RPC subsystem. Changes to this parameter only affect new transport endpoints, not existing transport endpoints.
Validation	None
When to Change	Tune this parameter so that services can take advantage of client behaviors such as the clustering that accelerate NFS version 2 `WRITE` requests. Increasing this parameter might result in the server being better able to take advantage of client behaviors.
Commitment Level	Unstable

rpcmod:maxdupreqs

Description	Controls the size of the duplicate request cache that detects RPC- level retransmissions on connectionless transports. This cache is indexed by the client network address and the RPC procedure number, program number, version number, and transaction ID. This cache avoids processing retransmitted requests that might not be idempotent.
Data Type	Integer (32-bit)
Default	1024
Range	1 to 2^{31} - 1
Units	Requests
Dynamic?	The cache is dynamically sized, but the hash queues that provide fast access to the cache are statically sized. Making the cache very large might result in long search times to find entries in the cache. Do not set the value of this parameter to 0. This value prevents the NFS server from handling non idempotent requests.
Validation	None
When to Change	Examine the value of this parameter if false failures are encountered by NFS clients. For example, if an attempt to create a directory fails, but the directory is actually created, perhaps that retransmitted `MKDIR` request was not detected by the server. The size of the cache should match the load on the server. The cache records non idempotent requests and so only needs to track a portion of the total requests. The cache does need to hold the information long enough to be able to detect a retransmission by the client. Typically, the client timeout for connectionless transports is relatively short, starting around 1 second and increasing to about 20 seconds.
Commitment Level	Unstable

rpcmod:cotsmaxdupreqs

Description	Controls the size of the duplicate request cache that detects RPC- level retransmissions on connection-oriented transports. This cache is indexed by the client network address and the RPC procedure number, program number, version number, and transaction ID. This cache avoids processing retransmitted requests that might not be idempotent.

Data Type	Integer (32–bit)
Default	1024
Range	1 to 2^{31} - 1
Units	Requests
Dynamic?	Yes
Validation	The cache is dynamically sized, but the hash queues that provide fast access to the cache are statically sized. Making the cache very large might result in long search times to find entries in the cache. Do not set the value of this parameter to 0. It prevents the NFS server from handling non-idempotent requests.
When to Change	Examine the value of this parameter if false failures are encountered by NFS clients. For example, if an attempt to create a directory fails, but the directory is actually created, it is possible that a retransmitted `MKDIR` request was not detected by the server. The size of the cache should match the load on the server. The cache records non-idempotent requests and so only needs to track a portion of the total requests. It does need to hold the information long enough to be able to detect a retransmission on the part of the client. Typically, the client timeout for connection oriented transports is very long, about 1 minute. Thus, entries need to stay in the cache for fairly long times.
Commitment Level	Unstable

◆ ◆ ◆ CHAPTER 4

Internet Protocol Suite Tunable Parameters

This chapter describes various Internet Protocol suite parameters, such as TCP, IP, UDP, and SCTP.

Where to Find Tunable Parameter Information

Tunable Parameter	For Information
Oracle Solaris kernel tunables	Chapter 2, "Oracle Solaris Kernel Tunable Parameters"
NFS tunable parameters	Chapter 3, "NFS Tunable Parameters"
Network Cache and Accelerator (NCA) tunable parameters	Chapter 5, "Network Cache and Accelerator Tunable Parameters"

Overview of Tuning IP Suite Parameters

For new information about IP forwarding, see "New and Changed TCP/IP Parameters" on page 21.

You can set all of the tuning parameters described in this chapter by using the `ndd` command except for the following parameters:

- "`ipcl_conn_hash_size`" on page 148

- "ip_squeue_worker_wait" on page 149

These parameters can only be set in the /etc/system file.

Use the following syntax to set TCP/IP parameters by using the ndd command:

```
# ndd -set driver parameter
```

For more information, see ndd(1M).

Although the SMF framework provides a method for managing system services, ndd commands are still included in system startup scripts. For more information on creating a startup script, see "Using Run Control Scripts" in *System Administration Guide: Basic Administration*.

IP Suite Parameter Validation

All parameters described in this section are checked to verify that they fall in the parameter range. The parameter's range is provided with the description for each parameter.

Internet Request for Comments (RFCs)

Internet protocol and standard specifications are described in RFC documents. You can get copies of RFCs from ftp://ftp.rfc-editor.org/in-notes. Browse RFC topics by viewing the rfc-index.txt file at this site.

IP Tunable Parameters

ip_icmp_err_interval and ip_icmp_err_burst

Description	Controls the rate of IP in generating IPv4 or IPv6 ICMP error messages. IP generates only up to ip_icmp_err_burst IPv4 or IPv6 ICMP error messages in any ip_icmp_err_interval. The ip_icmp_err_interval parameter protects IP from denial of service attacks. Setting this parameter to 0 disables rate limiting. It does not disable the generation of error messages.
Default	100 milliseconds for ip_icmp_err_interval 10 error messages for ip_icmp_err_burst
Range	0 – 99,999 milliseconds for ip_icmp_err_interval 1 – 99,999 error messages for ip_icmp_err_burst

Dynamic?	Yes
When to Change	If you need a higher error message generation rate for diagnostic purposes.
Commitment Level	Unstable

ip_respond_to_echo_broadcast and ip6_respond_to_echo_multicast

Description	Controls whether IPv4 or IPv6 responds to a broadcast ICMPv4 echo request or a multicast ICMPv6 echo request.
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes
When to Change	If you do not want this behavior for security reasons, disable it.
Commitment Level	Unstable

ip_send_redirects and ip6_send_redirects

Description	Controls whether IPv4 or IPv6 sends out ICMPv4 or ICMPv6 redirect messages.
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes
When to Change	If you do not want this behavior for security reasons, disable it.
Commitment Level	Unstable

ip_forward_src_routed and ip6_forward_src_routed

Description	Controls whether IPv4 or IPv6 forwards packets with source IPv4 routing options or IPv6 routing headers.
Default	0 (disabled)
Range	0 (disabled) or 1 (enabled)

Dynamic?	Yes
When to Change	Keep this parameter disabled to prevent denial of service attacks.
Commitment Level	Unstable
Change History	For information, see "`ip_forward_src_routed` and `ip6_forward_src_routed` (Solaris 10 Releases)" on page 182.

ip_addrs_per_if

Description	Defines the maximum number of logical interfaces associated with a real interface.
Default	256
Range	1 to 8192
Dynamic?	Yes
When to Change	Do not change the value. If more logical interfaces are required, you might consider increasing the value. However, recognize that this change might have a negative impact on IP's performance.
Commitment Level	Unstable

ip_strict_dst_multihoming and ip6_strict_dst_multihoming

Description	Determines whether a packet arriving on a non forwarding interface can be accepted for an IP address that is not explicitly configured on that interface. If `ip_forwarding` is enabled, or *xxx*`:ip_forwarding` for the appropriate interfaces is enabled, then this parameter is ignored, because the packet is actually forwarded. Refer to RFC 1122, 3.3.4.2.
Default	0 (loose multihoming)
Range	0 = Off (loose multihoming) 1 = On (strict multihoming)
Dynamic?	Yes
When to Change	If a machine has interfaces that cross strict networking domains (for example, a firewall or a VPN node), set this parameter to 1.

Commitment Level	Unstable

ip_multidata_outbound

Description	Enables the network stack to send more than one packet at one time to the network device driver during transmission. Enabling this parameter reduces the per-packet processing costs by improving host CPU utilization, network throughput, or both. This parameter now controls the use of multidata transmit (MDT) for transmitting IP fragments. For example, when sending out a UDP payload larger than the link MTU. When this tunable is enabled, IP fragments of a particular upper-level protocol, such as UDP, are delivered in batches to the network device driver. Disabling this feature results in both TCP and IP fragmentation logic in the network stack to revert back to sending one packet at a time to the driver. The MDT feature is only effective for device drivers that support this feature. See also "tcp_mdt_max_pbufs" on page 146.
Default	1 (Enabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes
When to Change	If you do not want this parameter enabled for debugging purposes or for any other reasons, disable it.
Commitment Level	Unstable
Change History	For information, see "ip_multidata_outbound (Solaris 10 Releases)" on page 182.

ip_squeue_fanout

Description	Determines the mode of associating TCP/IP connections with squeues A value of 0 associates a new TCP/IP connection with the CPU that creates the connection. A value of 1 associates the connection with multiple squeues that belong to different CPUs. The number of squeues that are used to fanout the connection is based upon "ip_soft_rings_cnt" on page 134

Default	0
Range	0 or 1
Dynamic?	Yes
When to Change	Consider setting this parameter to 1 to spread the load across all CPUs in certain situations. For example, when the number of CPUs exceed the number of NICs, and one CPU is not capable of handling the network load of a single NIC, change this parameter to 1.
Zone Configuration	This parameter can only be set in the global zone.
Commitment Level	Unstable
Change History	For information, see "`ip_squeue_fanout` (Solaris 10 11/06 Release)" on page 182.

ip_soft_rings_cnt

Description	Determines the number of squeues to be used to fanout the incoming TCP/IP connections.

Note – The incoming traffic is placed on one of the rings. If the ring is overloaded, packets are dropped. For every packet that gets dropped, the kstat dls counter, `dls_soft_ring_pkt_drop`, is incremented.

Default	2
Range	0 - nCPUs, where nCPUs is the maximum number of CPUs in the system
Dynamic?	No. The interface should be plumbed again when changing this parameter.
When to Change	Consider setting this parameter to a value greater than 2 on systems that have 10 Gbps NICs and many CPUs.
Zone Configuration	This parameter can only be set in the global zone.
Commitment Level	Obsolete
Change History	For information, see "`ip_soft_rings_cnt` (Solaris 10 11/06 Release)" on page 182.

IP Tunable Parameters With Additional Cautions

Changing the following parameters is not recommended.

ip_ire_pathmtu_interval

Description	Specifies the interval in milliseconds when IP flushes the path maximum transfer unit (PMTU) discovery information, and tries to rediscover PMTU. Refer to RFC 1191 on PMTU discovery.
Default	10 minutes
Range	5 seconds to 277 hours
Dynamic?	Yes
When to Change	Do not change this value.
Commitment Level	Unstable

ip_icmp_return_data_bytes and ip6_icmp_return_data_bytes

Description	When IPv4 or IPv6 sends an ICMPv4 or ICMPv6 error message, it includes the IP header of the packet that caused the error message. This parameter controls how many extra bytes of the packet beyond the IPv4 or IPv6 header are included in the ICMPv4 or ICMPv6 error message.
Default	64 bytes (`ip_icmp_return_data_bytes`) 1280 bytes (`ip6_icmp_return_data_bytes`)
Range	8 to 65,536 bytes
Dynamic?	Yes
When to Change	Do not change the value. Including more information in an ICMP error message might help in diagnosing network problems. If this feature is needed, increase the value.
Commitment Level	Unstable

TCP Tunable Parameters

tcp_deferred_ack_interval

Description	Specifies the time-out value for the TCP-delayed acknowledgment (ACK) timer for hosts that are not directly connected. Refer to RFC 1122, 4.2.3.2.
Default	100 milliseconds
Range	1 millisecond to 1 minute
Dynamic?	Yes
When to Change	Do not increase this value to more than 500 milliseconds. Increase the value under the following circumstances: ■ Slow network links (less than 57.6 Kbps) with greater than 512 bytes maximum segment size (MSS) ■ The interval for receiving more than one TCP segment is short
Commitment Level	Unstable

tcp_local_dack_interval

Description	Specifies the time-out value for TCP-delayed acknowledgment (ACK) timer for hosts that are directly connected. Refer to RFC 1122, 4.2.3.2.
Default	50 milliseconds
Range	10 milliseconds to 500 milliseconds
Dynamic?	Yes
When to Change	Do not increase this value to more than 500 milliseconds. Increase the value under the following circumstances: ■ Slow network links (less than 57.6 Kbps) with greater than 512 bytes maximum segment size (MSS) ■ The interval for receiving more than one TCP segment is short
Commitment Level	Unstable

Change History	For information, see "tcp_local_dack_interval (Solaris 10 Releases)" on page 183.

tcp_deferred_acks_max

Description	Specifies the maximum number of TCP segments received from remote destinations (not directly connected) before an acknowledgment (ACK) is generated. TCP segments are measured in units of maximum segment size (MSS) for individual connections. If set to 0 or 1, no ACKs are delayed, assuming all segments are 1 MSS long. The actual number is dynamically calculated for each connection. The value is the default maximum.
Default	2
Range	0 to 16
Dynamic?	Yes
When to Change	Do not change the value. In some circumstances, when the network traffic becomes very bursty because of the delayed ACK effect, decrease the value. Do not decrease this value below 2.
Commitment Level	Unstable

tcp_local_dacks_max

Description	Specifies the maximum number of TCP segments received from directly connected destinations before an acknowledgment (ACK) is generated. TCP segments are measured in units of maximum segment size (MSS) for individual connections. If set to 0 or 1, it means no ACKs are delayed, assuming all segments are 1 MSS long. The actual number is dynamically calculated for each connection. The value is the default maximum.
Default	8
Range	0 to 16
Dynamic?	Yes
When to Change	Do not change the value. In some circumstances, when the network traffic becomes very bursty because of the delayed ACK effect, decrease the value. Do not decrease this value below 2.
Commitment Level	Unstable

tcp_wscale_always

Description	When this parameter is enabled, which is the default setting, TCP always sends a SYN segment with the window scale option, even if the window scale option value is 0. Note that if TCP receives a SYN segment with the window scale option, even if the parameter is disabled, TCP responds with a SYN segment with the window scale option. In addition, the option value is set according to the receive window size. Refer to RFC 1323 for the window scale option.
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes
When to Change	If there is an interoperability problem with an old TCP stack that does not support the window scale option, disable this parameter.
Commitment Level	Unstable

tcp_tstamp_always

Description	If set to 1, TCP always sends a SYN segment with the timestamp option. Note that if TCP receives a SYN segment with the timestamp option, TCP responds with a SYN segment with the timestamp option even if the parameter is set to 0.
Default	0 (disabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes
When to Change	If getting an accurate measurement of round-trip time (RTT) and TCP sequence number wraparound is a problem, enable this parameter. Refer to RFC 1323 for more reasons to enable this option.
Commitment Level	Unstable

tcp_xmit_hiwat

Description	Defines the default send window size in bytes. Refer to "Per-Route Metrics" on page 166 for a discussion of setting a different value on a per-route basis. See also "tcp_max_buf" on page 139.
Default	49,152
Range	4096 to 1,073,741,824
Dynamic?	Yes
When to Change	An application can use setsockopt(3XNET) SO_SNDBUF to change the individual connection's send buffer.
Commitment Level	Unstable

tcp_recv_hiwat

Description	Defines the default receive window size in bytes. Refer to "Per-Route Metrics" on page 166 for a discussion of setting a different value on a per-route basis. See also "tcp_max_buf" on page 139 and "tcp_recv_hiwat_minmss" on page 152.
Default	49,152
Range	2048 to 1,073,741,824
Dynamic?	Yes
When to Change	An application can use setsockopt(3XNET) SO_RCVBUF to change the individual connection's receive buffer.
Commitment Level	Unstable

tcp_max_buf

Description	Defines the maximum buffer size in bytes. This parameter controls how large the send and receive buffers are set to by an application that uses setsockopt(3XNET).
Default	1,048,576
Range	8192 to 1,073,741,824
Dynamic?	Yes

When to Change	If TCP connections are being made in a high-speed network environment, increase the value to match the network link speed.
Commitment Level	Unstable

tcp_cwnd_max

Description	Defines the maximum value of the TCP congestion window (cwnd) in bytes. For more information on the TCP congestion window, refer to RFC 1122 and RFC 2581.
Default	1,048,576
Range	128 to 1,073,741,824
Dynamic?	Yes
When to Change	Even if an application uses `setsockopt(3XNET)` to change the window size to a value higher than `tcp_cwnd_max`, the actual window used can never grow beyond `tcp_cwnd_max`. Thus, `tcp_max_buf` should be greater than `tcp_cwnd_max`.
Commitment Level	Unstable

tcp_slow_start_initial

Description	Defines the maximum initial congestion window (cwnd) size in the maximum segment size (MSS) of a TCP connection. Refer to RFC 2414 on how the initial congestion window size is calculated.
Default	4
Range	1 to 4
Dynamic?	Yes
When to Change	Do not change the value. If the initial cwnd size causes network congestion under special circumstances, decrease the value.
Commitment Level	Unstable

tcp_slow_start_after_idle

Description	The congestion window size in the maximum segment size (MSS) of a TCP connection after it has been idled (no segment received) for a period of one retransmission timeout (RTO). Refer to RFC 2414 on how the initial congestion window size is calculated.
Default	4
Range	1 to 16,384
Dynamic?	Yes
When to Change	For more information, see "tcp_slow_start_initial" on page 140.
Commitment Level	Unstable

tcp_sack_permitted

Description	If set to 2, TCP always sends a SYN segment with the selective acknowledgment (SACK) permitted option. If TCP receives a SYN segment with a SACK-permitted option and this parameter is set to 1, TCP responds with a SACK-permitted option. If the parameter is set to 0, TCP does not send a SACK-permitted option, regardless of whether the incoming segment contains the SACK permitted option. Refer to RFC 2018 for information on the SACK option.
Default	2 (active enabled)
Range	0 (disabled), 1 (passive enabled), or 2 (active enabled)
Dynamic?	Yes
When to Change	SACK processing can improve TCP retransmission performance so it should be actively enabled. Sometimes, the other side can be confused with the SACK option actively enabled. If this confusion occurs, set the value to 1 so that SACK processing is enabled only when incoming connections allow SACK processing.
Commitment Level	Unstable

tcp_rev_src_routes

Description	If set to 0, TCP does not reverse the IP source routing option for incoming connections for security reasons. If set to 1, TCP does the normal reverse source routing.
Default	0 (disabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes
When to Change	If IP source routing is needed for diagnostic purposes, enable it.
Commitment Level	Unstable

tcp_time_wait_interval

Description	Specifies the time in milliseconds that a TCP connection stays in TIME-WAIT state. For more information, refer to RFC 1122, 4.2.2.13.
Default	60,000 (60 seconds)
Range	1 second to 10 minutes
Dynamic?	Yes
When to Change	Do not set the value lower than 60 seconds. For information on changing this parameter, refer to RFC 1122, 4.2.2.13.
Commitment Level	Unstable

tcp_ecn_permitted

Description	Controls Explicit Congestion Notification (ECN) support. If this parameter is set to 0, TCP does not negotiate with a peer that supports the ECN mechanism. If this parameter is set to 1 when initiating a connection, TCP does not tell a peer that it supports ECN mechanism.

However, TCP tells a peer that it supports ECN mechanism when accepting a new incoming connection request if the peer indicates that it supports ECN mechanism in the SYN segment.

If this parameter is set to 2, in addition to negotiating with a peer on the ECN mechanism when accepting connections, TCP indicates in the outgoing SYN segment that it supports the ECN mechanism when TCP makes active outgoing connections.

Refer to RFC 3168 for information on ECN.

Default	1 (passive enabled)
Range	0 (disabled), 1 (passive enabled), or 2 (active enabled)
Dynamic?	Yes
When to Change	ECN can help TCP better handle congestion control. However, there are existing TCP implementations, firewalls, NATs, and other network devices that are confused by this mechanism. These devices do not comply to the IETF standard. Because of these devices, the default value of this parameter is set to 1. In rare cases, passive enabling can still cause problems. Set the parameter to 0 only if absolutely necessary.
Commitment Level	Unstable

tcp_conn_req_max_q

Description	Specifies the default maximum number of pending TCP connections for a TCP listener waiting to be accepted by `accept(3SOCKET)`. See also "`tcp_conn_req_max_q0`" on page 144.
Default	128
Range	1 to 4,294,967,295
Dynamic?	Yes
When to Change	For applications such as web servers that might receive several connection requests, the default value might be increased to match the incoming rate. Do not increase the parameter to a very large value. The pending TCP connections can consume excessive memory. Also, if an application

cannot handle that many connection requests fast enough because the number of pending TCP connections is too large, new incoming requests might be denied.

Note that increasing `tcp_conn_req_max_q` does not mean that applications can have that many pending TCP connections. Applications can use `listen(3SOCKET)` to change the maximum number of pending TCP connections for each socket. This parameter is the maximum an application can use `listen()` to set the number to. Thus, even if this parameter is set to a very large value, the actual maximum number for a socket might be much less than `tcp_conn_req_max_q`, depending on the value used in `listen()`.

Commitment Level: Unstable

tcp_conn_req_max_q0

Description: Specifies the default maximum number of incomplete (three-way handshake not yet finished) pending TCP connections for a TCP listener.

For more information on TCP three-way handshake, refer to RFC 793. See also "`tcp_conn_req_max_q`" on page 143.

Default: 1024

Range: 0 to 4,294,967,295

Dynamic?: Yes

When to Change: For applications such as web servers that might receive excessive connection requests, you can increase the default value to match the incoming rate.

The following explains the relationship between `tcp_conn_req_max_q0` and the maximum number of pending connections for each socket.

When a connection request is received, TCP first checks if the number of pending TCP connections (three-way handshake is done) waiting to be accepted exceeds the maximum (*N*) for the listener. If the connections are excessive, the request is denied. If the number of connections is allowable, then TCP checks if the number of incomplete pending TCP connections exceeds the sum of *N* and `tcp_conn_req_max_q0`. If it does not, the request is accepted. Otherwise, the oldest incomplete pending TCP request is dropped.

Commitment Level	Unstable

tcp_conn_req_min

Description	Specifies the default minimum value for the maximum number of pending TCP connection requests for a listener waiting to be accepted. This is the lowest maximum value of `listen(3SOCKET)` that an application can use.
Default	1
Range	1 to 1024
Dynamic?	Yes
When to Change	This parameter can be a solution for applications that use `listen(3SOCKET)` to set the maximum number of pending TCP connections to a value too low. Increase the value to match the incoming connection request rate.
Commitment Level	Unstable

tcp_rst_sent_rate_enabled

Description	If this parameter is set to 1, the maximum rate of sending a RST segment is controlled by the `ndd` parameter, `tcp_rst_sent_rate`. If this parameter is set to 0, no rate control when sending a RST segment is available.
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes
When to Change	This tunable helps defend against denial of service attacks on TCP by limiting the rate by which a RST segment is sent out. The only time this rate control should be disabled is when strict conformance to RFC 793 is required.
Commitment Level	Unstable

tcp_rst_sent_rate

Description	Sets the maximum number of RST segments that TCP can send out per second.
Default	40
Range	0 to 4,294,967,295
Dynamic?	Yes
When to Change	In a TCP environment, there might be a legitimate reason to generate more RSTs than the default value allows. In this case, increase the default value of this parameter.
Commitment Level	Unstable

tcp_mdt_max_pbufs

Description	Specifies the number of payload buffers that can be carried by a single `M_MULTIDATA` message that is generated by TCP. See also "`ip_multidata_outbound`" on page 133.
Default	16
Range	1 to 16
Dynamic?	Yes
When to Change	Decreasing this parameter might aid in debugging device driver development by limiting the amount of payload buffers per `M_MULTIDATA` message that is generated by TCP.
Commitment Level	Unstable

tcp_naglim_def

Description	This parameter controls the Nagle algorithm threshold. TCP uses the minimum of this parameter and the MSS of a connection to determine when the Nagle algorithm should kick in. For example, if the amount of new data is more than 1 MSS, the data is sent out regardless of the value of this parameter. If this parameter is set to 1, the Nagle is disabled for all TCP connections.
Default	4,096
Range	1 to 65,535

Dynamic?	Yes
When to Change	Real-time applications that need to send data without delay should use `setsockopt()` to set `TCP_NODELAY` to 1 for the sockets needing fast transmission rather than setting the `tcp_naglim_def` parameter.
Commitment Level	Unstable
Change History	For information, see "`tcp_naglim_def` (Solaris 10 Releases)" on page 183.

tcp_smallest_anon_port

Description	This parameter controls the smallest port number TCP can select as an ephemeral port. An application can use an ephemeral port when it creates a connection with a specified protocol and it does not specify a port number. Ephemeral ports are not associated with a specific application. When the connection is closed, the port number can be reused by a different application.
Unit	Port number
Default	32,768
Range	1,024 to 65,534
Dynamic?	Yes
When to Change	When a larger ephemeral port range is required.
Commitment Level	Unstable
Change History	For information, see "`[tcp,sctp,udp]_smallest_anon_port` and `[tcp,sctp,udp]_largest_anon_port` (Solaris 10 Releases)" on page 183.

tcp_largest_anon_port

Description	This parameter controls the largest port number TCP can select as an ephemeral port. An application can use an ephemeral port when it creates a connection with a specified protocol and it does not specify a port number. Ephemeral ports are not associated with a specific application. When the connection is closed, the port number can be reused by a different application.
Unit	Port number

Default	65,535
Range	1,024 to 65,535
Dynamic?	Yes
When to Change	When a larger ephemeral port range is required.
Commitment Level	Unstable
Change History	For information, see "`[tcp,sctp,udp]_smallest_anon_port` and `[tcp,sctp,udp]_largest_anon_port` (Solaris 10 Releases)" on page 183.

TCP/IP Parameters Set in the /etc/system File

The following parameters can be set only in the `/etc/system` file. After the file is modified, reboot the system.

For example, the following entry sets the `ipcl_conn_hash_size` parameter:

```
set ip:ipcl_conn_hash_sizes=value
```

ipcl_conn_hash_size

Description	Controls the size of the connection hash table used by IP. The default value of 0 means that the system automatically sizes an appropriate value for this parameter at boot time, depending on the available memory.
Data Type	Unsigned integer
Default	0
Range	0 to 82,500
Dynamic?	No. The parameter can only be changed at boot time.
When to Change	If the system consistently has tens of thousands of TCP connections, the value can be increased accordingly. Increasing the hash table size means that more memory is wired down, thereby reducing available memory to user applications.
Commitment Level	Unstable

ip_squeue_worker_wait

Description	Governs the maximum delay in waking up a worker thread to process TCP/IP packets that are enqueued on an squeue. An *squeue* is a serialization queue that is used by the TCP/IP kernel code to process TCP/IP packets.
Default	10 milliseconds
Range	0 – 50 milliseconds
Dynamic?	Yes
When to Change	Consider tuning this parameter if latency is an issue, and network traffic is light. For example, if the machine serves mostly interactive network traffic. The default value usually works best on a network file server, a web server, or any server that has substantial network traffic.
Zone Configuration	This parameter can only be set in the global zone.
Commitment Level	Unstable
Change History	For information, see "`ip_squeue_worker_wait` (Solaris 10 11/06 Release)" on page 182.

TCP Parameters With Additional Cautions

Changing the following parameters is not recommended.

tcp_keepalive_interval

Description	This `ndd` parameter sets a probe interval that is first sent out after a TCP connection is idle on a system-wide basis. Oracle Solaris supports the TCP keep-alive mechanism as described in RFC 1122. This mechanism is enabled by setting the `SO_KEEPALIVE` socket option on a TCP socket. If `SO_KEEPALIVE` is enabled for a socket, the first keep-alive probe is sent out after a TCP connection is idle for two hours, the default value of the `tcp_keepalive_interval` parameter. If the peer does not respond to the probe after eight minutes, the TCP connection is aborted. You can also use the `TCP_KEEPALIVE_THRESHOLD` socket option on individual applications to override the default interval so that each

	application can have its own interval on each socket. The option value is an unsigned integer in milliseconds. See also `tcp(7P)`.
Default	2 hours
Range	10 seconds to 10 days
Units	Unsigned integer (milliseconds)
Dynamic?	Yes
When to Change	Do not change the value. Lowering it may cause unnecessary network traffic and might also increase the chance of premature termination of the connection because of a transient network problem.
Commitment Level	Unstable

tcp_ip_abort_interval

Description	Specifies the default total retransmission timeout value for a TCP connection. For a given TCP connection, if TCP has been retransmitting for `tcp_ip_abort_interval` period of time and it has not received any acknowledgment from the other endpoint during this period, TCP closes this connection. For TCP retransmission timeout (RTO) calculation, refer to RFC 1122, 4.2.3. See also "`tcp_rexmit_interval_max`" on page 151.
Default	5 minutes
Range	500 milliseconds to 1193 hours
Dynamic?	Yes
When to Change	Do not change this value. See "`tcp_rexmit_interval_max`" on page 151 for exceptions.
Commitment Level	Unstable

tcp_rexmit_interval_initial

Description	Specifies the default initial retransmission timeout (RTO) value for a TCP connection. Refer to "Per-Route Metrics" on page 166 for a discussion of setting a different value on a per-route basis.
Default	3 seconds
Range	1 millisecond to 20 seconds
Dynamic?	Yes

When to Change: Do not change this value. Lowering the value can result in unnecessary retransmissions.

Commitment Level: Unstable

tcp_rexmit_interval_max

Description: Defines the default maximum retransmission timeout value (RTO). The calculated RTO for all TCP connections cannot exceed this value. See also "tcp_ip_abort_interval" on page 150.

Default: 60 seconds

Range: 1 millisecond to 2 hours

Dynamic?: Yes

When to Change: Do not change the value in a normal network environment.

If, in some special circumstances, the round-trip time (RTT) for a connection is about 10 seconds, you can increase this value. If you change this value, you should also change the `tcp_ip_abort_interval` parameter. Change the value of `tcp_ip_abort_interval` to at least four times the value of `tcp_rexmit_interval_max`.

Commitment Level: Unstable

tcp_rexmit_interval_min

Description: Specifies the default minimum retransmission time out (RTO) value. The calculated RTO for all TCP connections cannot be lower than this value. See also "tcp_rexmit_interval_max" on page 151.

Default: 400 milliseconds

Range: 1 millisecond to 20 seconds

Dynamic?: Yes

When to Change: Do not change the value in a normal network environment.

TCP's RTO calculation should cope with most RTT fluctuations. If, in some very special circumstances, the round-trip time (RTT) for a connection is about 10 seconds, increase this value. If you change this value, you should change the `tcp_rexmit_interval_max` parameter. Change the value of `tcp_rexmit_interval_max` to at least eight times the value of `tcp_rexmit_interval_min`.

Commitment Level: Unstable

tcp_rexmit_interval_extra

Description	Specifies a constant added to the calculated retransmission time out value (RTO).
Default	0 milliseconds
Range	0 to 2 hours
Dynamic?	Yes
When to Change	Do not change the value. When the RTO calculation fails to obtain a good value for a connection, you can change this value to avoid unnecessary retransmissions.
Commitment Level	Unstable

tcp_tstamp_if_wscale

Description	If this parameter is set to 1, and the window scale option is enabled for a connection, TCP also enables the `timestamp` option for that connection.
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes
When to Change	Do not change this value. In general, when TCP is used in high-speed network, protection against sequence number wraparound is essential. Thus, you need the `timestamp` option.
Commitment Level	Unstable

tcp_recv_hiwat_minmss

Description	Controls the default minimum receive window size. The minimum is `tcp_recv_hiwat_minmss` times the size of maximum segment size (MSS) of a connection.
Default	4
Range	1 to 65,536
Dynamic?	Yes
When to Change	Do not change the value. If changing it is necessary, do not change the value lower than 4.

Commitment Level	Unstable

UDP Tunable Parameters

udp_xmit_hiwat

Description	Defines the default maximum UDP socket datagram size. For more information, see "udp_max_buf" on page 155.
Default	57,344 bytes
Range	1,024 to 1,073,741,824 bytes
Dynamic?	Yes
When to Change	Note that an application can use setsockopt(3XNET) SO_SNDBUF to change the size for an individual socket. In general, you do not need to change the default value.
Commitment Level	Unstable

udp_recv_hiwat

Description	Defines the default maximum UDP socket receive buffer size. For more information, see "udp_max_buf" on page 155.
Default	57,344 bytes
Range	128 to 1,073,741,824 bytes
Dynamic?	Yes
When to Change	Note that an application can use setsockopt(3XNET) SO_RCVBUF to change the size for an individual socket. In general, you do not need to change the default value.
Commitment Level	Unstable

udp_smallest_anon_port

Description	This parameter controls the smallest port number UDP can select as an ephemeral port. An application can use an ephemeral port when it creates a connection with a specified protocol and it does not specify a

	port number. Ephemeral ports are not associated with a specific application. When the connection is closed, the port number can be reused by a different application.
Unit	Port number
Default	32,768
Range	1,024 to 65,534
Dynamic?	Yes
When to Change	When a larger ephemeral port range is required.
Commitment Level	Unstable
Change History	For information, see "`[tcp,sctp,udp]_smallest_anon_port` and `[tcp,sctp,udp]_largest_anon_port` (Solaris 10 Releases)" on page 183.

udp_largest_anon_port

Description	This parameter controls the largest port number UDP can select as an ephemeral port. An application can use an ephemeral port when it creates a connection with a specified protocol and it does not specify a port number. Ephemeral ports are not associated with a specific application. When the connection is closed, the port number can be reused by a different application.
Unit	Port number
Default	65,535
Range	1,024 to 65,535
Dynamic?	Yes
When to Change	When a larger ephemeral port range is required.
Commitment Level	Unstable
Change History	For information, see "`[tcp,sctp,udp]_smallest_anon_port` and `[tcp,sctp,udp]_largest_anon_port` (Solaris 10 Releases)" on page 183.

udp_do_checksum

Description	This parameter controls whether UDP calculates the checksum on outgoing UDP/IPv4 packets.
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes
When to Change	Do not change this parameter.
Commitment Level	Unstable
Change History	For information, see "udp_do_checksum (Solaris 10 Releases)" on page 183.

UDP Parameter With Additional Caution

Changing the following parameter is not recommended.

udp_max_buf

Description	Controls how large send and receive buffers can be for a UDP socket.
Default	2,097,152 bytes
Range	65,536 to 1,073,741,824 bytes
Dynamic?	Yes
When to Change	Do not change the value. If this parameter is set to a very large value, UDP socket applications can consume too much memory.
Commitment Level	Unstable

IPQoS Tunable Parameter

ip_policy_mask

Description	Enables or disables IPQoS processing in any of the following callout positions: forward outbound, forward inbound, local outbound, and local inbound. This parameter is a bitmask as follows:

Not Used	Not Used	Not Used	Not Used	Forward Outbound	Forward Inbound	Local Outbound	Local Inbound
X	X	X	X	0	0	0	0

A `1` in any of the position masks or disables IPQoS processing in that particular callout position. For example, a value of `0x01` disables IPQoS processing for all the local inbound packets.

Default — The default value is 0, meaning that IPQoS processing is enabled in all the callout positions.

Range — 0 (0x00) to 15 (0x0F). A value of 15 indicates that IPQoS processing is disabled in all the callout positions.

Dynamic? — Yes

When to Change — If you want to enable or disable IPQoS processing in any of the callout positions.

Commitment Level — Unstable

SCTP Tunable Parameters

sctp_max_init_retr

Description — Controls the maximum number of attempts an SCTP endpoint should make at resending an INIT chunk. The SCTP endpoint can use the SCTP initiation structure to override this value.

Default — 8

Range — 0 to 128

Dynamic? — Yes

When to Change — The number of INIT retransmissions depend on "`sctp_pa_max_retr`" on page 157. Ideally, `sctp_max_init_retr` should be less than or equal to `sctp_pa_max_retr`.

Commitment Level — Unstable

sctp_pa_max_retr

Description	Controls the maximum number of retransmissions (over all paths) for an SCTP association. The SCTP association is aborted when this number is exceeded.
Default	10
Range	1 to 128
Dynamic?	Yes
When to Change	The maximum number of retransmissions over all paths depend on the number of paths and the maximum number of retransmission over each path. Ideally, `sctp_pa_max_retr` should be set to the sum of "`sctp_pp_max_retr`" on page 157 over all available paths. For example, if there are 3 paths to the destination and the maximum number of retransmissions over each of the 3 paths is 5, then `sctp_pa_max_retr` should be set to less than or equal to 15. (See the Note in Section 8.2, RFC 2960.)
Commitment Level	Unstable

sctp_pp_max_retr

Description	Controls the maximum number of retransmissions over a specific path. When this number is exceeded for a path, the path (destination) is considered unreachable.
Default	5
Range	1 to 128
Dynamic?	Yes
When to Change	Do not change this value to less than 5.
Commitment Level	Unstable

sctp_cwnd_max

Description	Controls the maximum value of the congestion window for an SCTP association.
Default	1,048,576
Range	128 to 1,073,741,824
Dynamic?	Yes

When to Change: Even if an application uses `setsockopt(3XNET)` to change the window size to a value higher than `sctp_cwnd_max`, the actual window used can never grow beyond `sctp_cwnd_max`. Thus, "`sctp_max_buf`" on page 161 should be greater than `sctp_cwnd_max`.

Commitment Level: Unstable

sctp_ipv4_ttl

Description: Controls the time to live (TTL) value in the IP version 4 header for the outbound IP version 4 packets on an SCTP association.

Default: 64
Range: 1 to 255

Dynamic?: Yes

When to Change: Generally, you do not need to change this value. Consider increasing this parameter if the path to the destination is likely to span more than 64 hops.

Commitment Level: Unstable

sctp_heartbeat_interval

Description: Computes the interval between HEARTBEAT chunks to an idle destination, that is allowed to heartbeat.

An SCTP endpoint periodically sends an HEARTBEAT chunk to monitor the reachability of the idle destinations transport addresses of its peer.

Default: 30 seconds
Range: 0 to 86,400 seconds

Dynamic?: Yes

When to Change: Refer to RFC 2960, section 8.3.

Commitment Level: Unstable

sctp_new_secret_interval

Description: Determines when a new secret needs to be generated. The generated secret is used to compute the MAC for a cookie.

Default	2 minutes
Range	0 to 1,440 minutes
Dynamic?	Yes
When to Change	Refer to RFC 2960, section 5.1.3.
Commitment Level	Unstable

sctp_initial_mtu

Description	Determines the initial maximum send size for an SCTP packet including the length of the IP header.
Default	1500 bytes
Range	68 to 65,535
Dynamic?	Yes
When to Change	Increase this parameter if the underlying link supports frame sizes that are greater than 1500 bytes.
Commitment Level	Unstable

sctp_deferred_ack_interval

Description	Sets the time-out value for SCTP delayed acknowledgment (ACK) timer in milliseconds.
Default	100 milliseconds
Range	1 to 60,000 milliseconds
Dynamic?	Yes
When to Change	Refer to RFC 2960, section 6.2.
Commitment Level	Unstable

sctp_ignore_path_mtu

Description	Enables or disables path MTU discovery.
Default	0 (disabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes

When to Change	Enable this parameter if you want to ignore MTU changes along the path. However, doing so might result in IP fragmentation if the path MTU decreases.
Commitment Level	Unstable

sctp_initial_ssthresh

Description	Sets the initial slow start threshold for a destination address of the peer.
Default	102,400
Range	1024 to 4,294,967,295
Dynamic?	Yes
When to Change	Refer to RFC 2960, section 7.2.1.
Commitment Level	Unstable

sctp_xmit_hiwat

Description	Sets the default send window size in bytes. See also "sctp_max_buf" on page 161.
Default	102,400
Range	8,192 to 1,073,741,824
Dynamic?	Yes
When to Change	An application can use getsockopt(3SOCKET) SO_SNDBUF to change the individual association's send buffer.
Commitment Level	Unstable

sctp_xmit_lowat

Description	Controls the lower limit on the send window size.
Default	8,192
Range	8,192 to 1,073,741,824
Dynamic?	Yes
When to Change	Generally, you do not need to change this value. This parameter sets the minimum size required in the send buffer for the socket to be

marked writable. If required, consider changing this parameter in accordance with "sctp_xmit_hiwat" on page 160.

Commitment Level: Unstable

sctp_recv_hiwat

Description: Controls the default receive window size in bytes. See also "sctp_max_buf" on page 161.

Default: 102,400
Range: 8,192 to 1,073,741,824

Dynamic?: Yes

When to Change: An application can use getsockopt(3SOCKET) SO_RCVBUF to change the individual association's receive buffer.

Commitment Level: Unstable

sctp_max_buf

Description: Controls the maximum buffer size in bytes. It controls how large the send and receive buffers are set to by an application that uses getsockopt(3SOCKET).

Default: 1,048,576
Range: 8,192 to 1,073,741,824

Dynamic?: Yes

When to Change: Increase the value of this parameter to match the network link speed if associations are being made in a high-speed network environment.

Commitment Level: Unstable

sctp_ipv6_hoplimit

Description: Sets the value of the hop limit in the IP version 6 header for the outbound IP version 6 packets on an SCTP association.

Default: 60
Range: 0 to 255

Dynamic?: Yes

When to Change	Generally, you do not need to change this value. Consider increasing this parameter if the path to the destination is likely to span more than 60 hops.
Commitment Level	Unstable

sctp_rto_min

Description	Sets the lower bound for the retransmission timeout (RTO) in milliseconds for all the destination addresses of the peer.
Default	1,000
Range	500 to 60,000
Dynamic?	Yes
When to Change	Refer to RFC 2960, section 6.3.1.
Commitment Level	Unstable

sctp_rto_max

Description	Controls the upper bound for the retransmission timeout (RTO) in milliseconds for all the destination addresses of the peer.
Default	60,000
Range	1,000 to 60,000,000
Dynamic?	Yes
When to Change	Refer to RFC 2960, section 6.3.1.
Commitment Level	Unstable

sctp_rto_initial

Description	Controls the initial retransmission timeout (RTO) in milliseconds for all the destination addresses of the peer.
Default	3,000
Range	1,000 to 60,000,000
Dynamic?	Yes
When to Change	Refer to RFC 2960, section 6.3.1.
Commitment Level	Unstable

sctp_cookie_life

Description	Sets the lifespan of a cookie in milliseconds.
Default	60,000
Range	10 to 60,000,000
Dynamic?	Yes
When to Change	Generally, you do not need to change this value. This parameter might be changed in accordance with "sctp_rto_max" on page 162.
Commitment Level	Unstable

sctp_max_in_streams

Description	Controls the maximum number of inbound streams permitted for an SCTP association.
Default	32
Range	1 to 65,535
Dynamic?	Yes
When to Change	Refer to RFC 2960, section 5.1.1.
Commitment Level	Unstable

sctp_initial_out_streams

Description	Controls the maximum number of outbound streams permitted for an SCTP association.
Default	32
Range	1 to 65,535
Dynamic?	Yes
When to Change	Refer to RFC 2960, section 5.1.1.
Commitment Level	Unstable

sctp_shutack_wait_bound

Description	Controls the maximum time, in milliseconds, to wait for a SHUTDOWN ACK after having sent a SHUTDOWN chunk.

Default	60,000
Range	0 to 300,000
Dynamic?	Yes
When to Change	Generally, you do not need to change this value. This parameter might be changed in accordance with "sctp_rto_max" on page 162.
Commitment Level	Unstable

sctp_maxburst

Description	Sets the limit on the number of segments to be sent in a burst.
Default	4
Range	2 to 8
Dynamic?	Yes
When to Change	You do not need to change this parameter. You might change it for testing purposes.
Commitment Level	Unstable

sctp_addip_enabled

Description	Enables or disables SCTP dynamic address reconfiguration.
Default	0 (disabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes
When to Change	The parameter can be enabled if dynamic address reconfiguration is needed. Due to security implications, enable this parameter only for testing purposes.
Commitment Level	Unstable

sctp_prsctp_enabled

Description	Enables or disables the partial reliability extension (RFC 3758) to SCTP.
Default	1 (enabled)
Range	0 (disabled) or 1 (enabled)
Dynamic?	Yes

When to Change	Disable this parameter if partial reliability is not supported in your SCTP environment.
Commitment Level	Unstable

sctp_smallest_anon_port

Description	This parameter controls the smallest port number SCTP can select as an ephemeral port. An application can use an ephemeral port when it creates a connection with a specified protocol and it does not specify a port number. Ephemeral ports are not associated with a specific application. When the connection is closed, the port number can be reused by a different application.
Unit	Port number
Default	32,768
Range	1,024 to 65,534
Dynamic?	Yes
When to Change	When a larger ephemeral port range is required.
Commitment Level	Unstable

sctp_largest_anon_port

Description	This parameter controls the largest port number SCTP can select as an ephemeral port. An application can use an ephemeral port when it creates a connection with a specified protocol and it does not specify a port number. Ephemeral ports are not associated with a specific application. When the connection is closed, the port number can be reused by a different application.
Unit	Port number
Default	32,768
Range	1,024 to 65,534
Dynamic?	Yes
When to Change	When a larger ephemeral port range is required.
Commitment Level	Unstable

Per-Route Metrics

You can use per-route metrics to associate some properties with IPv4 and IPv6 routing table entries.

For example, a system has two different network interfaces, a fast Ethernet interface and a gigabit Ethernet interface. The system default `tcp_recv_hiwat` is 49,152 bytes. This default is sufficient for the fast Ethernet interface, but may not be sufficient for the gigabit Ethernet interface.

Instead of increasing the system's default for `tcp_recv_hiwat`, you can associate a different default TCP receive window size to the gigabit Ethernet interface routing entry. By making this association, all TCP connections going through the route will have the increased receive window size.

For example, the following is in the routing table (`netstat -rn`), assuming IPv4:

```
192.123.123.0        192.123.123.4        U       1      4  hme0
192.123.124.0        192.123.124.4        U       1      4  ge0
default              192.123.123.1        UG      1      8
```

In this example, do the following:

```
# route change -net 192.123.124.0 -recvpipe x
```

Then, all connections going to the `192.123.124.0` network, which is on the `ge0` link, use the receive buffer size *x*, instead of the default `49,152` receive window size.

If the destination is in the `a.b.c.d` network, and no specific routing entry exists for that network, you can add a prefix route to that network and change the metric. For example:

```
# route add -net a.b.c.d 192.123.123.1 -netmask w.x.y.z
# route change -net a.b.c.d -recvpipe y
```

Note that the prefix route's gateway is the default router. Then, all connections going to that network use the receive buffer size *y*. If you have more than one interface, use the `-ifp` argument to specify which interface to use. This way, you can control which interface to use for specific destinations. To verify the metric, use the `route`(1M) `get` command.

◆ ◆ ◆ CHAPTER 5

Network Cache and Accelerator Tunable Parameters

This chapter describes some of the Network Cache and Accelerator (NCA) tunable parameters.

Where to Find Tunable Parameters Information

Tunable Parameter	For Information
Oracle Solaris kernel tunables	Chapter 2, "Oracle Solaris Kernel Tunable Parameters"
NFS tunable parameters	Chapter 3, "NFS Tunable Parameters"
Internet Protocol Suite tunable parameters	Chapter 4, "Internet Protocol Suite Tunable Parameters"

Tuning NCA Parameters

Setting these parameters is appropriate on a system that is a dedicated web server. These parameters allocate more memory for caching pages. You can set all of the tuning parameters described in this chapter in the `/etc/system` file.

For information on adding tunable parameters to the `/etc/system` file, see "Tuning the Oracle Solaris Kernel" on page 26.

nca:nca_conn_hash_size

Description	Controls the hash table size in the NCA module for all TCP connections, adjusted to the nearest prime number.
Default	383 hash table entries
Range	0 to 201,326,557
Dynamic?	No
When to Change	When the NCA's TCP hash table is too small to keep track of the incoming TCP connections. This situation causes many TCP connections to be grouped together in the same hashtable entry. This situation is indicated when NCA is receiving many TCP connections, and system performance decreases.
Commitment Level	Unstable

nca:nca_conn_req_max_q

Description	Defines the maximum number of pending TCP connections for NCA to listen on.
Default	256 connections
Range	0 to 4,294,967,295
Dynamic?	No
When to Change	When NCA closes a connection immediately after it is established because it already has too many established TCP connections. If NCA is receiving many TCP connections and can handle a larger load, but is refusing any more connections, increase this parameter. Doing so allows NCA to handle more simultaneous TCP connections.
Commitment Level	Unstable

nca:nca_conn_req_max_q0

Description	Defines the maximum number of incomplete (three-way handshake not yet finished) pending TCP connections for NCA to listen on.
Default	1024 connections
Range	0 to 4,294,967,295
Dynamic?	No

When to Change	When NCA refuses to accept any more TCP connections because it already has too many pending TCP connections. If NCA is receiving many TCP connections and can handle a larger load, but is refusing any more connections, increase this parameter. Doing so allows NCA to handle more simultaneous TCP connections.
Commitment Level	Unstable

nca:nca_ppmax

Description	Specifies the maximum amount of physical memory (in pages) used by NCA for caching the pages. This value should not be more than 75 percent of total memory.
Default	25 percent of physical memory
Range	1 percent to maximum amount of physical memory
Dynamic?	No
When to Change	When using NCA on a system with more than 512 MB of memory. If a system has a lot of physical memory that is not being used, increase this parameter. Then, NCA will efficiently use this memory to cache new objects. As a result, system performance will increase. This parameter should be increased in conjunction with `nca_vpmax`, unless you have a system with more physical memory than virtual memory (a 32-bit kernel that has greater than 4 GB memory). Use `pagesize(1)` to determine your system's page size.
Commitment Level	Unstable

nca:nca_vpmax

Description	Specifies the maximum amount of virtual memory (in pages) used by NCA for caching pages. This value should not be more than 75 percent of the total memory.
Default	25 percent of virtual memory
Range	1 percent to maximum amount of virtual memory
Dynamic?	No
When to Change	When using NCA on a system with more than 512 MB of memory. If a system has a lot of virtual memory that is not being used, increase this

parameter. Then, NCA will efficiently use this memory to cache new objects. As a result, system performance will increase.

This parameter should be increased in conjunction with `nca_ppmax`. Set this parameter about the same value as `nca_vpmax`, unless you have a system with more physical memory than virtual memory.

Commitment Level: Unstable

General System Tuning for the NCA

In addition to setting the NCA parameters, you can do some general system tuning to benefit NCA performance. If you are using gigabit Ethernet (`ge` driver), you should set the interface in interrupt mode for better results.

For example, a system with 4 GB of memory that is booted under 64-bit kernel should have the following parameters set in the `/etc/system` file. Use `pagesize` to determine your system's page size.

```
set sq_max_size=0
set ge:ge_intr_mode=1
set nca:nca_conn_hash_size=82500
set nca:nca_conn_req_max_q=100000
set nca:nca_conn_req_max_q0=100000
set nca:nca_ppmax=393216
set nca:nca_vpmax=393216
```

sq_max_size

Description: Sets the depth of the syncq (number of messages) before a destination STREAMS queue generates a `QFULL` message.

Default: 10000 messages

Range: 0 (unlimited) to MAXINT

Dynamic?: No

When to Change: When NCA is running on a system with a lot of memory, increase this parameter to allow drivers to queue more packets of data. If a server is under heavy load, increase this parameter so that modules and drivers can process more data without dropping packets or getting backlogged.

Commitment Level: Unstable

ge:ge_intr_mode

Description	Enables the ge driver to send packets directly to the upper communication layers rather than queue the packets
Default	0 (queue packets to upper layers)
Range	0 (enable) or 1 (disable)
Dynamic?	No
When to Change	When NCA is enabled, set this parameter to 1 so that the packet is delivered to NCA in interrupt mode for faster processing.
Commitment Level	Unstable

◆ ◆ ◆ CHAPTER 6

System Facility Parameters

This chapter describes most of the parameters default values for various system facilities.

System Default Parameters

The functioning of various system facilities is governed by a set of values that are read by each facility on startup. The values stored in a file for each facility are located in the `/etc/default` directory. Not every system facility has a file located in this directory.

autofs

This facility enables you to configure `autofs` parameters such as automatic timeout, displaying or logging status messages, browsing autofs mount points, and tracing. For details, see `autofs(4)`.

cron

This facility enables you to disable or enable `cron` logging.

devfsadm

This file is not currently used.

dhcpagent

Client usage of DHCP is provided by the `dhcpagent` daemon. When `ifconfig` identifies an interface that has been configured to receive its network configuration from DHCP, it starts the client daemon to manage that interface.

For more information, see the `/etc/default/dhcpagent` information in the FILES section of `dhcpagent(1M)`.

fs

File system administrative commands have a generic and file system-specific portion. If the file system type is not explicitly specified with the `-F` option, a default is applied. The value is specified in this file. For more information, see the Description section of `default_fs(4)`.

ftp

This facility enables you to set the `ls` command behavior to the RFC 959 `NLST` command. The default `ls` behavior is the same as in the previous Solaris release.

For details, see `ftp(4)`.

inetinit

This facility enables you to configure TCP sequence numbers and to enable or disable support for 6to4 relay routers.

init

For details, see the `/etc/default/init` information in the FILES section of `init(1M)`.

All values in the file are placed in the environment of the shell that `init` invokes in response to a single user boot request. The `init` process also passes these values to any commands that it starts or restarts from the `/etc/inittab` file.

ipsec

This facility enables you to configure parameters, such as IKE daemon debugging information and the `ikeadm` privilege level.

kbd

For details, see the Extended Description section of `kbd(1)`.

keyserv

For details, see the `/etc/default/keyserv` information in the FILES section of `keyserv(1M)`.

login

For details, see the `/etc/default/login` information in the FILES section of `login(1)`.

lu

This file contains default settings for the Oracle Solaris Live Upgrade feature.

mpathd

This facility enables you to set `in.mpathd` configuration parameters.

For details, see `in.mpathd(1M)`.

nfs

This facility enables you to set NFS daemon configuration parameters.

For details, see nfs(4).

nfslogd

For details, see the Description section of nfslogd(1M).

nss

This facility enables you to configure initgroups(3C) lookup parameters.

For details, see nss(4).

passwd

For details, see the /etc/default/passwd information in the FILES section of passwd(1).

power

For details, see the /etc/default/power information in the FILES section of pmconfig(1M).

rpc.nisd

For details, see the /etc/default/rpc.nisd information in the FILES section of rpc.nisd(1M).

su

For details, see the /etc/default/su information in the FILES section of su(1M).

syslog

For details, see the /etc/default/syslogd information in the FILES section of syslogd(1M).

sys-suspend

For details, see the `/etc/default/sys-suspend` information in the FILES section of `sys-suspend(1M)`.

tar

For a description of the `-f` function modifier, see `tar(1)`.

If the `TAPE` environment variable is not present and the value of one of the arguments is a number and `-f` is not specified, the number matching the archive*N* string is looked up in the `/etc/default/tar` file. The value of the archive*N* string is used as the output device with the blocking and size specifications from the file.

For example:

```
% tar -c 2 /tmp/*
```

This command writes the output to the device specified as `archive2` in the `/etc/default/tar` file.

telnetd

This file identifies the default `BANNER` that is displayed upon a telnet connection.

utmpd

The `utmpd` daemon monitors `/var/adm/utmpx` (and `/var/adm/utmp` in earlier Solaris versions) to ensure that `utmp` entries inserted by non-root processes by `pututxline(3C)` are cleaned up on process termination.

Two entries in `/etc/default/utmpd` are supported:

- `SCAN_PERIOD` – The number of seconds that `utmpd` sleeps between checks of `/proc` to see if monitored processes are still alive. The default is 300.
- `MAX_FDS` – The maximum number of processes that `utmpd` attempts to monitor. The default value is 4096 and should never need to be changed.

yppasswdd

This facility enables you to configure whether a user can successfully set a login shell to a restricted shell when using the `passwd -r nis -e` command.

For details, see `rpc.yppasswdd(1M)`.

◆ ◆ ◆ APPENDIX A

Tunable Parameters Change History

This chapter describes the change history of specific tunable parameters. If a parameter is in this section, it has changed from a previous release. Parameters whose functionality has been removed are listed also.

Kernel Parameters

Process-Sizing Tunables

max_nprocs (Solaris 10 Releases)

The Solaris 10 description section was updated by removing the text "sun4m."

General Driver Parameter

ddi_msix_alloc_limit (Solaris 10 Releases)

This parameter is new in the Solaris 10 10/09 release. For more information, see "ddi_msix_alloc_limit" on page 62.

General I/O Tunable Parameters

maxphys (Solaris 10 Releases)

The default value is updated to include sun4v systems. For more information, see "maxphys" on page 63.

General Kernel and Memory Parameters

zfs_arc_min (Solaris 10 Releases)

This parameter description is newly documented in the Solaris 10 10/09 release. For more information, see "zfs_arc_min" on page 33.

zfs_arc_max (Solaris 10 Releases)

This parameter description is newly documented in the Solaris 10 10/09 release. For more information, see "zfs_arc_max" on page 33.

noexec_user_stack (Solaris 10 Releases)

The Solaris 10 description section was updated by removing the text "and sun4m" and adding the text "64–bit SPARC and AMD64."

lwp_default_stksize (Solaris 10 Releases)

The Solaris 10 description section was updated by adding default and maximum values for AMD64.

The Solaris 10 default value for SPARC platforms was changed to 24,576.

fsflush and Related Parameters

dopageflush (Solaris 10 Releases)

The description was clarified by including that number of *physical* memory pages are examined.

Paging-Related Tunable Parameters

maxpgio (Solaris 10 Releases)

In the Solaris 10 versions, the range value was incorrectly documented as 1 to 1024. The actual range depends on system architecture and I/O subsystems. For more information, see "maxpgio" on page 57.

General File System Parameters

ncsize (Solaris 10 Releases)

In previous Solaris 10 releases, the default value of the `ncsize` parameter was incorrectly described as follows:

4 x (`v.v_proc` + `maxusers`) + 320 / 100

The correct default value is as follows:

(4 x (`v.v_proc` + `maxusers`) + 320) + (4 x (`v.v_proc` + `maxusers`) + 320 / 100

For more information, see "`ncsize`" on page 65.

TMPFS Parameters

tmpfs:tmpfs_maxkmem (Solaris 10 Releases)

The range description is updated to include sun4v systems. For more information, see "`tmpfs:tmpfs_maxkmem`" on page 75.

SPARC System Specific Parameters (Solaris 10 Releases)

The title of the SPARC System Specific Parameters section was revised in the Solaris 10 8/07 release to include sun4v systems.

default_tsb_size (Solaris 10 Releases)

The default description has changed. For more information, see "`default_tsb_size`" on page 86.

enable_tsb_rss_sizing (Solaris 10 Releases)

The description and default and range values have changed. For more information, see "`enable_tsb_rss_sizing`" on page 87.

tsb_rss_factor (Solaris 10 Releases)

The when to change example text was changed to this:

For example, changing `tsb_rss_factor` to 256 (effectively, 50%) instead of 384 (effectively, 75%) might help eliminate virtual address conflicts in the TSB in some cases, but will use more kernel memory, particularly on a heavily loaded system.

NFS Tunable Parameters

nfs:nfs3_nra (Solaris 10 Releases)

The default value was incorrectly documented in previous Solaris 10 releases. The default value is 4.

TCP/IP Tunable Parameters

ip_forward_src_routed and ip6_forward_src_routed (Solaris 10 Releases)

The default value of these parameters was incorrectly documented in previous Solaris 10 releases. The correct default value is disabled. For more information, see "ip_forward_src_routed and ip6_forward_src_routed" on page 131.

ip_multidata_outbound (Solaris 10 Releases)

This parameter was enhanced in the Solaris 10 release to deliver IP fragments in batches to the network driver. For more information, see "ip_multidata_outbound" on page 133.

ip_squeue_fanout (Solaris 10 11/06 Release)

Zone configuration information was added in the Solaris 10 8/07 release. For more information, see "ip_squeue_fanout" on page 133.

ip_squeue_worker_wait (Solaris 10 11/06 Release)

Zone configuration information was added in the Solaris 10 8/07 release. For more information, see "ip_squeue_worker_wait" on page 149 In addition, this parameter was moved to "TCP/IP Parameters Set in the /etc/system File" on page 148.

ip_soft_rings_cnt (Solaris 10 11/06 Release)

Zone configuration information was added in the Solaris 10 8/07 release. For more information, see "ip_soft_rings_cnt" on page 134.

ip_squeue_write (Solaris 10 Releases)

This parameter was incorrectly documented in the Solaris 10 release. It has been removed.

tcp_local_dack_interval (Solaris 10 Releases)

The range of this parameter was incorrectly documented in previous Solaris releases. The correct range is 10 milliseconds to 1 minute.

[tcp,sctp,udp]_smallest_anon_port and [tcp,sctp,udp]_largest_anon_port (Solaris 10 Releases)

These parameters are newly documented in the Solaris 10 8/11 release.

- "sctp_smallest_anon_port" on page 165
- "sctp_largest_anon_port" on page 165
- "tcp_smallest_anon_port" on page 147
- "tcp_largest_anon_port" on page 147
- "udp_smallest_anon_port" on page 153
- "udp_largest_anon_port" on page 154

tcp_naglim_def (Solaris 10 Releases)

The "tcp_naglim_def" on page 146 parameter is newly documented in the Solaris 10 8/11 release.

udp_do_checksum (Solaris 10 Releases)

The "udp_do_checksum" on page 155 parameter is newly documented in the Solaris 10 8/11 release.

Parameters That Are Obsolete or Have Been Removed

The following section describes parameters that are obsolete or have been removed from more recent Solaris releases.

rstchown

This parameter is obsolete starting in the Oracle Solaris 10 8/11 release.

Description	Indicates whether the POSIX semantics for the `chown` system call are in effect. POSIX semantics are as follows: ■ A process cannot change the owner of a file, unless it is running with UID 0. ■ A process cannot change the group ownership of a file to a group in which it is not currently a member, unless it is running as UID 0. For more information, see `chown(2)`.
Data Type	Signed integer
Default	1, indicating that POSIX semantics are used
Range	0 = POSIX semantics not in force or 1 = POSIX semantics used
Units	Toggle (on/off)
Dynamic?	Yes
Validation	None
When to Change	When POSIX semantics are not wanted. Note that turning off POSIX semantics opens the potential for various security holes. Doing so also opens the possibility of a user changing ownership of a file to another user and being unable to retrieve the file without intervention from the user or the system administrator.
Commitment Level	Obsolete

System V Message Queue Parameters

msgsys:msginfo_msgmni

Obsolete in the Solaris 10 release.

Description	Maximum number of message queues that can be created.
Data Type	Signed integer

Default	50
Range	0 to MAXINT
Dynamic?	No. Loaded into `msgmni` field of `msginfo` structure.
Validation	None
When to Change	When `msgget(2)` calls return with an error of `ENOSPC` or at the recommendation of a software vendor.
Commitment Level	Unstable

msgsys:msginfo_msgtql

Obsolete in the Solaris 10 release.

Description	Maximum number of messages that can be created. If a `msgsnd` call attempts to exceed this limit, the request is deferred until a message header is available. Or, if the request has set the `IPC_NOWAIT` flag, the request fails with the error `EAGAIN`.
Data Type	Signed integer
Default	40
Range	0 to MAXINT
Dynamic?	No. Loaded into `msgtql` field of `msginfo` structure.
Validation	None
When to Change	When `msgsnd()` calls block or return with error of `EGAIN`, or at the recommendation of a software vendor.
Commitment Level	Unstable

msgsys:msginfo_msgmnb

Obsolete in the Solaris 10 release.

Description	Maximum number of bytes that can be on any one message queue.
Data Type	Unsigned long
Default	4096
Range	0 to amount of physical memory
Units	Bytes
Dynamic?	No. Loaded into `msgmnb` field of `msginfo` structure.
Validation	None

When to Change	When `msgsnd()` calls block or return with an error of `EGAIN`, or at the recommendation of a software vendor.
Commitment Level	Unstable

msgsys:msginfo_msgssz

Removed in the Solaris 10 release.

Description	Specifies size of chunks system uses to manage space for message buffers.
Data Type	Signed integer
Default	40
Range	0 to MAXINT
Dynamic?	No. Loaded into `msgtql` field of `msginfo`structure.
Validation	The space consumed by the maximum number of data structures that would be created to support the messages and queues is compared to 25% of the available kernel memory at the time the module is loaded. If the number is too big, the message queue module refuses to load and the facility is unavailable. This computation does include the space that might be consumed by the messages. This situation occurs only when the module is first loaded.
When to Change	When the default value is not enough. Generally changed at the recommendation of software vendors.
Commitment Level	Obsolete

msgsys:msginfo_msgmap

Removed in the Solaris 10 release.

Description	Number of messages the system supports.
Data Type	Signed integer
Default	100
Range	0 to MAXINT
Dynamic?	No
Validation	The space consumed by the maximum number of data structures that would be created to support the messages and queues is compared to 25% of the available kernel memory at the time the module is loaded. If the number is too big, the message queue module refuses to load and

	the facility is unavailable. This computation does include the space that might be consumed by the messages. This situation occurs only when the module is first loaded.
When to Change	When the default value is not enough. Generally changed at the recommendation of software vendors.
Commitment Level	Obsolete

msgsys:msginfo_msgseg

Removed in the Solaris 10 release.

Description	Number of `msginfo_msgssz` segments the system uses as a pool for available message memory. Total memory available for messages is `msginfo_msgseg * msginfo_msgssz`.
Data Type	Signed short
Default	1024
Range	0 to 32,767
Dynamic?	No
Validation	The space consumed by the maximum number of data structures that would be created to support the messages and queues is compared to 25% of the available kernel memory at the time the module is loaded. If the number is too big, the message queue module refuses to load and the facility is unavailable. This computation does not include the space that might be consumed by the messages. This situation occurs only when the module is first loaded.
When to Change	When the default value is not enough. Generally changed at the recommendation of software vendors.
Commitment Level	Obsolete

msgsys:msginfo_msgmax

Removed in the Solaris 10 release.

Description	Maximum size of System V message.
Data Type	Unsigned long
Default	2048
Range	0 to amount of physical memory
Units	Bytes

Dynamic?	No. Loaded into `msgmax` field of `msginfo` structure.
Validation	None
When to Change	When `msgsnd(2)` calls return with error of `EINVAL` or at the recommendation of a software vendor.
Commitment Level	Unstable

System V Semaphore Parameters

semsys:seminfo_semmni

Obsolete in the Solaris 10 release.

Description	Specifies the maximum number of semaphore identifiers.
Data Type	Signed integer
Default	10
Range	1 to 65,535
Dynamic?	No
Validation	Compared to `SEMA_INDEX_MAX` (currently 65,535) and reset to that value if larger. A warning message is written to the console, messages file, or both.
When to Change	When the default number of sets is not enough. Generally changed at the recommendation of software vendors. No error messages are displayed when an attempt is made to create more sets than are currently configured. Instead, the application receives a return code of `ENOSPC` from a `semget` call. For more information, see `semget(2)`.
Commitment Level	Unstable

semsys:seminfo_semmsl

Obsolete in the Solaris 10 release.

Description	Specifies the maximum number of System V semaphores per semaphore identifier.
Data Type	Signed integer
Default	25

Range	1 to MAXINT
Dynamic?	No
Validation	The amount of space that could possibly be consumed by the semaphores and their supporting data structures is compared to 25 percent of the kernel memory available at the time the module is first loaded. If the memory threshold is exceeded, the module refuses to load and the semaphore facility is not available.
When to Change	When the default value is not enough. Generally changed at the recommendation of software vendors. No error messages are displayed when an attempt is made to create more semaphores in a set than are currently configured. The application sees a return code of `EINVAL` from a `semget(2)` call.
Commitment Level	Unstable

semsys:seminfo_semopm

Obsolete in the Solaris 10 release.

Description	Specifies the maximum number of System V semaphore operations per `semop` call. This parameter refers to the number of `sembufs` in the `sops` array that is provided to the `semop()` system call. For more information, see `semop(2)`.
Data Type	Signed integer
Default	10
Range	1 to MAXINT
Dynamic?	No
Validation	The amount of space that could possibly be consumed by the semaphores and their supporting data structures is compared to 25 percent of the kernel memory available at the time the module is first loaded. If the memory threshold is exceeded, the module refuses to load and the semaphore facility is not available.
When to Change	When the default value is not enough. Generally changed at the recommendation of software vendors. No error messages are displayed when an attempt is made to perform more semaphore operations in a single `semop` call than are currently allowed. Instead, the application receives a return code of `E2BIG` from a `semop()` call.
Commitment Level	Unstable

semsys:seminfo_semmns

Removed in the Solaris 10 release.

Description	Maximum number of System V semaphores on the system.
Data Type	Signed integer
Default	60
Range	1 to MAXINT
Dynamic?	No
Validation	The amount of space that could possibly be consumed by the semaphores and their supporting data structures is compared to 25% of the kernel memory available at the time the module is first loaded. If the memory threshold is exceeded, the module refuses to load and the semaphore facility is not available.
When to Change	When the default number of semaphores is not enough. Generally changed at the recommendation of software vendors. No error messages are displayed when an attempt is made to create more semaphores than are currently configured. The application sees a return code of `ENOSPC` from a `semget(2)` call.
Commitment Level	Unstable

semsys:seminfo_semmnu

Removed in the Solaris 10 release.

Description	Total number of undo structures supported by the System V semaphore system.
Data Type	Signed integer
Default	30
Range	1 to MAXINT
Dynamic?	No
Validation	The amount of space that could possibly be consumed by the semaphores and their supporting data structures is compared to 25% of the kernel memory available at the time the module is first loaded. If the memory threshold is exceeded, the module refuses to load and the semaphore facility is not available.
When to Change	When the default value is not enough. Generally changed at the recommendation of software vendors. No error message is displayed when an attempt is made to perform more undo operations than are

	currently configured. The application sees a return value of ENOSPC from a semop(2) call when the system runs out of undo structures.
Commitment Level	Unstable

semsys:seminfo_semume

Description	Removed in the Solaris 10 release. Maximum number of System V semaphore undo structures that can be used by any one process.
Data Type	Signed integer
Default	10
Range	1 to MAXINT
Dynamic?	No
Validation	The amount of space that could possibly be consumed by the semaphores and their supporting data structures is compared to 25% of the kernel memory available at the time the module is first loaded. If the memory threshold is exceeded, the module refuses to load and the semaphore facility is not available.
When to Change	When the default value is not enough. Generally changed at the recommendation of software vendors. No error messages are displayed when an attempt is made to perform more undo operations than are currently configured. The application sees a return code of EINVAL from a semop(2) call.
Commitment Level	Unstable

semsys:seminfo_semvmx

Removed in the Solaris 10 release.

Description	Maximum value a semaphore can be set to.
Data Type	Unsigned short
Default	32,767
Range	1 to 65,535
Dynamic?	No
Validation	None
When to Change	When the default value is not enough. Generally changed at the recommendation of software vendors. No error messages are displayed

	when the maximum value is exceeded. The application sees a return code of ERANGE from a semop(2) call.
Commitment Level	Unstable

semsys:seminfo_semaem

Removed in the Solaris 10 release.

Description	Maximum value that a semaphore's value in an undo structure can be set to.
Data Type	Unsigned short
Default	16,384
Range	1 to 65,535
Dynamic?	No
Validation	None
When to Change	When the default value is not enough. Generally changed at the recommendation of software vendors. No error messages are displayed when an attempt is made to perform more undo operations than are currently configured. The application sees a return code of EINVAL from a semop(2) call.
Commitment Level	Unstable

System V Shared Memory Parameters

shmsys:shminfo_shmmni

Obsolete in the Solaris 10 release.

Description	System wide limit on number of shared memory segments that can be created.
Data Type	Signed integer
Default	100
Range	0 to MAXINT
Dynamic?	No. Loaded into shmmni field of shminfo structure.
Validation	The amount of space consumed by the maximum possible number of data structures to support System V shared memory is checked against

	25% of the currently available kernel memory at the time the module is loaded. If the memory consumed is too large, the attempt to load the module fails.
When to Change	When the system limits are too low. Generally changed on the recommendation of software vendors.
Commitment Level	Unstable

shmsys:shminfo_shmmax

Obsolete in the Solaris 10 release.

Description	Maximum size of system V shared memory segment that can be created. This parameter is an upper limit that is checked before the application sees if it actually has the physical resources to create the requested memory segment. Attempts to create a shared memory section whose size is zero or whose size is larger than the specified value will fail with an `EINVAL` error. This parameter specifies only the largest value the operating system can accept for the size of a shared memory segment. Whether the segment can be created depends entirely on the amount of swap space available on the system and, for a 32-bit process, whether there is enough space available in the process's address space for the segment to be attached.
Data Type	Unsigned long
Default	8,388,608
Range	0 - MAXUINT32 on 32-bit systems, 0 – MAXUINT64 on 64-bit systems
Units	Bytes
Dynamic?	No. Loaded into `shmmax` field of `shminfo` structure.
Validation	None
When to Change	When the default value is too low. Generally changed at the recommendation of software vendors, but unless the size of a shared memory segment needs to be constrained, setting this parameter to the maximum possible value has no side effects.
Commitment Level	Unstable

◆ ◆ ◆ APPENDIX B

Revision History for This Manual

This section describes the revision history for this manual.

Current Version: Oracle Solaris 10 8/11 Release

The current version of this manual applies to the Oracle Solaris 10 8/11 release.

New or Changed Parameters in the Oracle Solaris Release

The following sections describe new, changed, or obsolete kernel tunables.

- **Oracle Solaris 10 8/11**: The `rstchown` parameter is obsolete. For more information, see "What's New in Oracle Solaris System Tuning?" on page 17.
- **Oracle Solaris 10 8/11**: This release includes the `ngroups_max` parameter description. For more information, see "`ngroups_max`" on page 46.
- **Solaris 10 10/09**: This release includes the `zfs_arc_min` and `zfs_arc_max` parameter descriptions. For more information, see "`zfs_arc_min`" on page 33 and "`zfs_arc_max`" on page 33.
- **Solaris 10 10/09**: This release includes the `ddi_msix_alloc_limit` parameter that can be used to increase the number of MSI-X interrupts that a device instance can allocate. For more information, see "`ddi_msix_alloc_limit`" on page 62.
- **Solaris 10 10/09**: Memory locality group parameters are provided in this release. For more information about these parameters, see "Locality Group Parameters" on page 88.
- **Solaris 10 5/08**: The translation storage buffers parameters in the "SPARC System Specific Parameters" on page 84 section have been revised to provide better information. In this release, the following parameters have changed:

- "default_tsb_size" on page 86 – The default text has been clarified.
- "enable_tsb_rss_sizing" on page 87 – The default text was incorrect and has been corrected.
- "tsb_rss_factor" on page 87 – The example section referred to percentages rather than the more appropriate parameter units. This issue has been resolved.

- **Solaris 10 8/07**: Parameter information was updated to include sun4v systems. For more information, see the following references:
  - "maxphys" on page 63
  - "tmpfs:tmpfs_maxkmem" on page 75
  - "SPARC System Specific Parameters" on page 84

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