Sun Java System Web Server 7.0 NSAPI Developer's Guide

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Contents

Examples

Preface

This guide discusses how to use Netscape Server Application Programmer's Interface (NSAPI) to build plug-ins that define Server Application Functions (SAFs) to extend and modify Sun™ Java System Web Server 7.0. The guide also provides a reference of the NSAPI functions you can use to define new plug-ins.

Who Should Use This Book

The intended audience for this guide is the person who develops, assembles, and deploys NSAPI plug-ins in a corporate enterprise. This guide assumes you are familiar with the following topics:

- HTTP
- HTML
- NSAPI
- C programming
- Software development processes, including debugging and source code control

Before You Read This Book

Web Server 7.0 can be installed as a stand-alone product or as a component of Sun Java™ Enterprise System (Java ES), a software infrastructure that supports enterprise applications distributed across a network or Internet environment. If you are installing Web Server 7.0 as a component of Java ES, you should be familiar with the system documentation at `http://docs.sun.com/coll/1286.2`.

Web Server 7.0 Documentation Set

The Web Server 7.0 documentation set describes how to install and administer the Web Server. The URL for Web Server 7.0 documentation is http://docs.sun.com/coll/1308.3. For an introduction to Web Server 7.0, refer to the books in the order in which they are listed in the following table.

TABLE P–1 Books in the Web Server 7.0 Documentation Set

Documentation Title	Contents
Sun Java System Web Server 7.0 Documentation Center	Web Server documentation topics organized by tasks and subject
Sun Java System Web Server 7.0 Release Notes	■ Late-breaking information about the software and documentation ■ Supported platforms and patch requirements for installing Web Server
Sun Java System Web Server 7.0 Installation and Migration Guide	Performing installation and migration tasks: ■ Installing Web Server and its various components, ■ Migrating data from Sun ONE Web Server 6.0 or 6.1 to Sun Java System Web Server 7.0
Sun Java System Web Server 7.0 Administrator's Guide	Performing the following administration tasks: ■ Using the Administration and command-line interfaces ■ Configuring server preferences ■ Using server instances ■ Monitoring and logging server activity ■ Using certificates and public key cryptography to secure the server ■ Configuring access control to secure the server ■ Using Java Platform Enterprise Edition (Java EE) security features ■ Deploying applications ■ Managing virtual servers ■ Defining server workload and sizing the system to meet performance needs ■ Searching the contents and attributes of server documents, and creating a text search interface ■ Configuring the server for content compression ■ Configuring the server for web publishing and content authoring using WebDAV
Sun Java System Web Server 7.0 Developer's Guide	Using programming technologies and APIs to do the following: ■ Extend and modify Sun Java System Web Server ■ Dynamically generate content in response to client requests and modify the content of the server
Sun Java System Web Server 7.0 NSAPI Developer's Guide	Creating custom Netscape Server Application Programmer's Interface (NSAPI) plug-ins
Sun Java System Web Server 7.0 Developer's Guide to Java Web Applications	Implementing Java Servlets and JavaServer Pages™ (JSP™) technology in Sun Java System Web Server

TABLE P–1 Books in the Web Server 7.0 Documentation Set *(Continued)*

Documentation Title	Contents
Sun Java System Web Server 7.0 Administrator's Configuration File Reference	Editing configuration files
Sun Java System Web Server 7.0 Performance Tuning, Sizing, and Scaling Guide	Tuning Sun Java System Web Server to optimize performance
Sun Java System Web Server 7.0 Troubleshooting Guide	Troubleshooting Web Server

Related Books

The URL for all documentation about Sun Java Enterprise System (Java ES) and its components is `http://docs.sun.com/app/docs/prod/entsys.06q4`.

Default Paths and File Names

The following table describes the default paths and file names that are used in this book.

TABLE P–2 Default Paths and File Names

Placeholder	Description	Default Value

install_dir	Represents the base installation directory for Web Server 7.0.	Sun Java Enterprise System (Java ES) installations on the Solaris™ platform:
		`/opt/SUNWwbsvr7`
		Java ES installations on the Linux and HP-UX platform:
		`/opt/sun/webserver/`
		Java ES installations on the Windows platform:
		System Drive`:\Program Files\Sun\JavaES5\WebServer7`
		Other Solaris, Linux, and HP-UX installations, non-root user:
		user's home directory`/sun/webserver7`
		Other Solaris, Linux, and HP-UX installations, root user:
		`/sun/webserver7`
		Windows, all installations:
		System Drive`:\Program Files\Sun\WebServer7`

TABLE P–2 Default Paths and File Names *(Continued)*

Placeholder	Description	Default Value

instance_dir	Directory that contains the instance-specific subdirectories.	For Java ES installations, the default location for instances on Solaris:
		`/var/opt/SUNWwbsvr7`
		For Java ES installations, the default location for instances on Linux and HP-UX:
		`/var/opt/sun/webserver7`
		For Java ES installations, the default location for instance on Windows:
		System Drive`:\Program Files\Sun\JavaES5\WebServer7`
		For stand-alone installations, the default location for instance on Solaris, Linux, and HP-UX:
		<*install_dir*>
		For stand-alone installations, the default location for instance on Windows:
		System Drive`:\Program Files\sun\WebServer7`

Typographic Conventions

The following table describes the typographic changes that are used in this book.

TABLE P–3 Typographic Conventions

Typeface	Meaning	Example
`AaBbCc123`	The names of commands, files, and directories, and onscreen computer output	Edit your `.login` file. Use `ls -a` to list all files. `machine_name% you have mail.`
`AaBbCc123`	What you type, contrasted with onscreen computer output	`machine_name%` **`su`** `Password:`
AaBbCc123	A placeholder to be replaced with a real name or value	The command to remove a file is `rm` *filename*.
AaBbCc123	Book titles, new terms, and terms to be emphasized (note that some emphasized items appear bold online)	Read Chapter 6 in the *User's Guide*. A *cache* is a copy that is stored locally. Do *not* save the file.

Symbol Conventions

The following table explains symbols that might be used in this book.

TABLE P–4 Symbol Conventions

Symbol	Description	Example	Meaning
`[ ]`	Contains optional arguments and command options.	`ls [-l]`	The `-l` option is not required.
`{ | }`	Contains a set of choices for a required command option.	`-d {y|n}`	The `-d` option requires that you use either the `y` argument or the `n` argument.
`${ }`	Indicates a variable reference.	`${com.sun.javaRoot}`	References the value of the `com.sun.javaRoot` variable.
-	Joins simultaneous multiple keystrokes.	Control-A	Press the Control key while you press the A key.
+	Joins consecutive multiple keystrokes.	Ctrl+A+N	Press the Control key, release it, and then press the subsequent keys.
→	Indicates menu item selection in a graphical user interface.	File → New → Templates	From the File menu, choose New. From the New submenu, choose Templates.

Accessing Sun Resources Online

The `http://docs.sun.com` (docs.sun.com[SM]) web site enables you to access Sun technical documentation online. You can browse the docs.sun.com archive or search for a specific book title or subject. Books are available as online files in PDF and HTML formats. Both formats are readable by assistive technologies for users with disabilities.

To access the following Sun resources, go to `http://www.sun.com`:

- Downloads of Sun products
- Services and solutions
- Support (including patches and updates)
- Training
- Research
- Communities (for example, Sun Developer Network)

Searching Sun Product Documentation

Besides searching Sun product documentation from the docs.sun.com web site, you can use a search engine by typing the following syntax in the search field:

search-term `site:docs.sun.com`

For example, to search for "Web Server," type the following:

```
Web Server site:docs.sun.com
```

To include other Sun web sites in your search (for example, java.sun.com, www.sun.com, and developers.sun.com), use "`sun.com`" in place of "`docs.sun.com`" in the search field.

Third-Party Web Site References

Third-party URLs are referenced in this document and provide additional, related information.

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◆ ◆ ◆ CHAPTER 1

Creating Custom Server Application Functions

This chapter describes how to write your own NSAPI plug-ins that define custom Server Application Functions (SAFs). Creation of plug-ins allows you to modify or extend the Sun Java System Web Server's built-in functionality. For example, you can modify the server to handle user authorization in a special way or generate dynamic HTML pages based on information in a database.

This chapter has the following sections:

Before writing custom SAFs, you must familiarize yourself with the request-handling process, as described in detail in the *Sun Java System Web Server 7.0 Administrator's Configuration File Reference*. Also, before writing a custom SAF, check to see if a built-in SAF already accomplishes the tasks you have in mind.

See Appendix B for a list of the predefined `Init` SAFs. For information about predefined SAFs used in the `obj.conf` file, see the *Sun Java System Web Server 7.0 Administrator's Configuration File Reference*.

For a complete list of the NSAPI routines for implementing custom SAFs, see Chapter 5.

Future Compatibility Issues

The NSAPI interface might change in a future version of Sun Java System Web Server.

To keep your custom plug-ins upgradable, do the following:

- Make sure plug-in users know how to edit the configuration files (such as `magnus.conf` and `obj.conf`) manually. The plug-in installation software should not be used to edit these configuration files.
- Keep the source code so you can recompile the plug-in.

The SAF Interface

All SAFs (custom and built-in) have the same C interface regardless of the request-handling step for which they are written. SAFs are small functions that are designed for a specific purpose within a specific request-response step. SAFs receive parameters from the directive that invokes them in the `obj.conf` file, from the server, and from previous SAFs.

Here is the C interface for a SAF:

```
int function(pblock *pb, Session *sn, Request *rq);
```

The next section discusses the parameters in detail.

The SAF returns a result code that indicates whether and how it succeeded. The server uses the result code from each function to determine how to proceed with processing the request. For more information on the result codes, see "Result Codes" on page 19.

SAF Parameters

This section discusses the SAF parameters in detail.

The parameters are as follows:

- "pb (parameter block)" on page 18- contains the parameters from the directive that invokes the SAF in the `obj.conf` file.
- "sn (Session)" on page 19- contains information relating to a single TCP/IP session.
- "rq (Request)" on page 19- contains information relating to the current request.

pb (parameter block)

The `pb` parameter is a pointer to a `pblock` data structure that contains values specified by the directive that invokes the SAF. A `pblock` data structure contains a series of name-value pairs.

For example, a directive that invokes the `basic-nsca` function might look like the following:

```
AuthTrans fn=basic-ncsa auth-type=basic dbm=users.db
```

In this case, the `pb` parameter passed to `basic-ncsa` contains name-value pairs that correspond to `auth-type=basic` and `dbm=users.db`.

NSAPI provides a set of functions for working with `pblock` data structures. For example, `pblock_findval()` returns the value for a given name in a `pblock`. For information on working with parameter blocks, see "Parameter Block Manipulation Routines" on page 27.

sn (Session)

The `sn` parameter is a pointer to a `Session` data structure. This parameter contains variables related to an entire session. That is, the time between the opening and closing of the TCP/IP connection between the client and the server. The same `sn` pointer is passed to each SAF called within each request for an entire session. For a list of important fields, see "Session" on page 162.

rq (Request)

The `rq` parameter is a pointer to a `Request` data structure. This parameter contains variables related to the current request, such as the request headers, URI, and local file system path. The same `Request` pointer is passed to each SAF called in the request-response process for an HTTP request. For a list of important fields, see "Request" on page 164.

Result Codes

Upon completion, a SAF returns a result code. The result code indicates what the server should do next.

The result codes are:

- `REQ_PROCEED`

  Indicates that the SAF achieved its objective. For some request-response steps (`AuthTrans`, `NameTrans`, `Service`, and `Error`), this tells the server to proceed to the next request-response step, skipping any other SAFs in the current step. For the other request-response steps (`Input`, `Output`, `Route`, `PathCheck`, `ObjectType`, and `AddLog`), the server proceeds to the next SAF in the current step.

- `REQ_NOACTION`

  Indicates that the SAF took no action. The server continues with the next SAF in the current server step.

- `REQ_ABORTED`

  Indicates that an error occurred and an HTTP response should be sent to the client to indicate the cause of the error. A SAF returning `REQ_ABORTED` should also set the HTTP response status code. If the server finds an `Error` directive matching the status code or reason phrase, the server executes the SAF specified. If not, the server sends a default HTTP response with the status code and reason phrase, in addition to a short HTML page reflecting the status code and reason phrase for the user. The server then goes to the first `AddLog` directive.

- `REQ_EXIT`

  Indicates the connection to the client was lost. This should be returned when the SAF fails in reading or writing to the client. The server then goes to the first `AddLog` directive.

Creating and Using Custom SAFs

Custom SAFs are functions in shared libraries that are loaded and called by the server. Follow these steps to create a custom SAF:

▼ To Create a Custom SAF

1. **Write the Source Code using the NSAPI functions. Each SAF is written for a specific directive.**

2. **Compile and Link the source code to create a shared library (`.so`, `.sl`, or `.dll`) file.**

3. **Load and Initialize the SAF by editing the `magnus.conf` file to:**

   - Load the shared library file containing your custom SAF(s)
   - Initialize the SAF if necessary

4. **Instruct the Server to Call the SAFs by editing `obj.conf` to call your custom SAF(s) at the appropriate time.**

5 **Restart the Server.**

6 **Test the SAF by accessing your server from a browser with a URL that triggers your function.**

The following sections describe these steps in greater detail.

Write the Source Code

Write your custom SAFs using NSAPI functions. For a summary of some of the most commonly used NSAPI functions, see "Overview of NSAPI C Functions" on page 27 and for available routines, see Chapter 5.

For examples of custom SAFs, see Chapter 3.

The signature for all SAFs is as follows:

```
int function(pblock *pb, Session *sn, Request *rq);
```

For more details on the parameters, see "SAF Parameters" on page 18.

You must register your SAFs with the server. SAFs may be registered using the `funcs` parameter of the `load-modules` `Init` SAF or by a call to `func_insert`. A plug-in may define a `nspai_module_init` function that is used to call `func_insert` and perform any other initialization tasks. For more information, see "nsapi_module_init" on page 92 and "func_insert" on page 80.

The server runs as a multi-threaded single process. On UNIX platforms, there are two processes , a parent and a child, for historical reasons. The parent process performs some initialization and forks the child process. The child process performs further initialization and handles all of the HTTP requests.

Keep the following in mind when writing your SAF:

- Write thread-safe code
- Blocking can affect performance
- Write small functions with parameters and configure the parameters in `obj.conf`
- Carefully check and handle all errors (and log the errors so you can determine the source of problems and fix them)

If necessary, write an initialization function that performs initialization tasks required by your new SAFs. The initialization function must be named `nsapi_module_init` and has the same signature as other SAFs:

```
int nsapi_module_init(pblock *pb, Session *sn, Request *rq);
```

SAFs expect to be able to obtain certain types of information from their parameters. In most cases, parameter block (`pblock`) data structures provide the fundamental storage mechanism for these parameters. `pblock` maintains its data as a collection of name-value pairs. For a summary of the most commonly used functions for working with `pblock` structures, see "Parameter Block Manipulation Routines" on page 27.

When defining a SAF, you do not specifically state which directive it is written for. However, each SAF must be written for a specific directive, such as `AuthTrans`, `Service`, and so on. Each directive expects its SAFs to behave in particular ways, and your SAF must conform to the expectations of the directive for which it was written. For details on what each directive expects of its SAFs, see "Required Behavior of SAFs for Each Directive" on page 31.

Compile and Link

Compile and link your code with the native compiler for the target platform. For UNIX, use the `gmake` command. For Windows, use the `nmake` command. For Windows, use Microsoft Visual C++ 6.0 or newer. You must have an import list that specifies all global variables and functions to access from the server binary. Use the correct compiler and linker flags for your platform. Refer to the example Makefile in the *install_dir*`/samples/nsapi` directory.

Adhere to the following guidelines for compiling and linking.

Include Directory and nsapi.h File

Add the *install_dir*`/include` (UNIX) or *install_dir*`\include` (Windows) directory to your makefile to include the `nsapi.h` file.

Libraries

Add the *install_dir*`/bin/https/lib` (UNIX) or *install_dir*`\bin\https\bin` (Windows) library directory to your linker command.

The following table lists the library that you need to link to.

TABLE 1–1 Libraries

Platform	Library
Windows	`ns-httpd40.dll` (in addition to the standard Windows libraries)
HP-UX	`libns-httpd40.sl`
All other UNIX platforms	`libns-httpd40.so`

Linker Commands and Options for Generating a Shared Object

To generate a shared library, use the commands and options listed in the following table.

TABLE 1–2 Linker Commands and Options

Platform	Options
Solaris™ Operating System (SPARC® Platform Edition)	`ld -G` or `cc -G`
Windows	`link -LD`
HP-UX	`cc +Z -b -Wl,+s -Wl,-B,symbolic`
AIX	`cc -p 0 -berok -blibpath:$(LD_RPATH)`
Compaq	`cc -shared`
Linux	`gcc -shared`
IRIX	`cc -shared`

Additional Linker Flags

Use the linker flags in the following table to specify which directories should be searched for shared objects during runtime to resolve symbols.

TABLE 1–3 Linker Flags

Platform	Flags
Solaris SPARC	`-R` *dir*:*dir*
Windows	(no flags, but the `ns-httpd40.dll` file must be in the system PATH variable)
HP-UX	`-Wl,+b,`*dir*`,`*dir*
AIX	`-blibpath:`*dir*:*dir*
Compaq	`-rpath` *dir*:*dir*
Linux	`-Wl,-rpath,`*dir*:*dir*
IRIX	`-Wl,-rpath,`*dir*:*dir*

On UNIX, you can also set the library search path using the `LD_LIBRARY_PATH` environment variable, which must be set when you start the server.

Compiler Flags

The following table lists the flags and defines you need to use for compilation of your source code.

TABLE 1–4 Compiler Flags and Defines

Parameter	Description
Solaris SPARC	`-DXP_UNIX -D_REENTRANT -KPIC -DSOLARIS`
Windows	`-DXP_WIN32 -DWIN32 /MD`
HP-UX	`-DXP_UNIX -D_REENTRANT -DHPUX`
AIX	`-DXP_UNIX -D_REENTRANT -DAIX $(DEBUG)`
Compaq	`-DXP_UNIX -KPIC`
Linux	`-DXP_UNIX -D_REENTRANT -fPIC`
IRIX	`-o32 -exceptions -DXP_UNIX -KPIC`

Compiling and Linking in 64–bit Mode

On Solaris, the server can run in either 32–bit or 64–bit mode. Because a 32–bit shared library cannot be used in a 64–bit process and conversely, you may wish to compile and link two separate shared libraries. By default, the Sun compiler and linker produce 32–bit binaries. To compile and link your plug-in for 64–bit mode on Solaris SPARC, you must use Sun Workshop 5.0 or higher with the `-xarch=v9` flag. To compile and link your plug-in for 64–bit mode on Solaris x86, you must use Sun Java Studio 11 or higher with `-xarch=amd64` flag.

Issues with Using C++ in a NSAPI Plug-in

NSAPI plug-ins are typically written using the C programming language. Using the C++ programming language in an NSAPI plug-in raises special compatibility issues.

On Solaris, the server is built using the new C++ 5 ABI. If your shared library uses C++, it must be compiled with Sun Workshop 5.0 or higher. Sun Java Studio 11 or higher is recommended. Do not use the `-compat=4` option when compiling and linking a shared library that uses C++. When running in 32–bit mode on Solaris SPARC, the server provides some backward compatibility for the old C++ 4 ABI (Sun Workshop 4.2). This backward compatibility may be removed at some future date. For all new NSAPI plug-ins, use the new C++ 5 ABI (Sun Workshop 5.0 or higher).

On Linux, Web Server is built using the gcc 3.2 C++ ABI. If your shared library uses C++, compile with gcc 3.2.x. Because of the volatility of the gcc C++ ABI, it is advised to avoid using C++ in NSAPI plug-ins on Linux.

Load and Initialize the SAF

For each shared library (plug-in) containing custom SAFs to be loaded into the server, add an `Init` directive that invokes the `load-modules` SAF to `magnus.conf`. The `load-modules` SAF loads the shared library and calls the shared library's `nsapi_module_init` function. For more information, see "nsapi_module_init" on page 92.

The syntax for a directive that calls `load-modules` is:

```
Init fn=load-modules
     [shlib=path]
     [funcs="SAF1,...,SAFn"]
     [name1="value1"]...[nameN="valueN"]
```

- `shlib` is the local file system path to the shared library (plug-in).
- `funcs` is an optional comma-separated list of function names to be loaded from the shared library. Function names are case-sensitive. You may use dash a (-) in place of an underscore (_) in function names. There should be no spaces in the function name list.

  If the new SAFs require initialization, you must omit the `funcs` parameter and instead define an `nsapi_module_init` function in your shared library. Any custom parameters on the Init directive will be passed to `nsapi_module_init` in the `pb` parameter block.
- `nameN="`*valueN*`"` are the optional names and values of parameters passed to the shared library's `nsapi_module_init` function in the `pb` parameter block.

Instruct the Server to Call the SAFs

Add directives to `obj.conf` to instruct the server to call each custom SAF at the appropriate time. The syntax for directives is:

Directive `fn=`*function-name* `[`*name1*`="`*value1*`"]...[`*nameN*`="`*valueN*`"]`

- *Directive* is one of the server directives, such as `AuthTrans`, `Service`, and so on.
- *function-name* is the name of the SAF to execute.
- *nameN*`="`*valueN*`"` are the names and values of parameters which are passed to the SAF.

  Depending on what your new SAF does, you might need to add just one directive to `obj.conf`, or you might need to add more than one directive to provide complete instructions for invoking the new SAF.

  For example, if you define a new `AuthTrans` or `PathCheck` SAF, you could just add an appropriate directive in the default object. However, if you define a new `Service` SAF to be invoked only when the requested resource is in a particular directory or has a new kind of file extension, you would need to take extra steps.

If your new `Service` SAF is to be invoked only when the requested resource has a new kind of file extension, you might need to add an entry to the MIME types file so that the `type` value gets set properly during the `ObjectType` stage. Then you could add a `Service` directive to the default object that specifies the desired `type` value.

If your new `Service` SAF is to be invoked only when the requested resource is in a particular directory, you might need to define a `NameTrans` directive that generates a `name` or `ppath` value that matches another object, and then in the new object you could invoke the new `Service` function.

For example, suppose your plug-in defines two new SAFs, `do_small_anim` and `do_big_anim`, which both take `speed` parameters. These functions run animations. All files to be treated as small animations reside in the directory `D:/docs/animations/small`, while all files to be treated as full-screen animations reside in the directory `D:/docs/animations/fullscreen`.

To ensure that the new animation functions are invoked whenever a client sends a request for either a small or full-screen animation, you would add `NameTrans` directives to the default object to translate the appropriate URLs to the corresponding path names and also assign a name to the request.

```
NameTrans fn=pfx2dir
          from="/animations/small"
          dir="D:/docs/animations/small"
          name="small_anim"
NameTrans fn=pfx2dir
          from="/animations/fullscreen"
          dir="D:/docs/animations/fullscreen"
          name="fullscreen_anim"
```

You also need to define objects that contain the `Service` directives that run the animations and specify the `speed` parameter.

```
<Object name="small_anim">
Service fn=do_small_anim speed=40
</Object>
<Object name="fullscreen_anim">
Service fn=do_big_anim speed=20
</Object>
```

Restart the Server

After modifying `obj.conf`, you need to restart the server. A restart is required for all plug-ins that implement SAFs and/or filters.

Test the SAF

Test your SAF by accessing your server from a browser with a URL that triggers your function. For example, if your new SAF is triggered by requests to resources in `http://`*server-name*`/animations/small`, try requesting a valid resource that starts with that URI.

You should disable caching in your browser so that the server is sure to be accessed. In Mozilla Firefox, you may hold the shift key while clicking the Reload button to ensure that the cache is not used.

Examine the access log and error log to help with debugging.

Overview of NSAPI C Functions

NSAPI provides a set of C functions that are used to implement SAFs. They serve several purposes. They provide platform independence across operating system and hardware platforms. They provide improved performance. They are thread-safe which is a requirement for SAFs. They prevent memory leaks. And they provide functionality necessary for implementing SAFs. You should always use these NSAPI routines when defining new SAFs.

This section provides an overview of the function categories available and some of the more commonly used routines. All of the public routines are detailed in Chapter 5.

The main categories of NSAPI functions are:

Parameter Block Manipulation Routines

The parameter block manipulation functions provide routines for locating, adding, and removing entries in a `pblock` data structure:

- "pblock_findval" on page 97 returns the value for a given name in a `pblock`.
- "pblock_nvinsert" on page 99 adds a new name-value pair entry to a `pblock`.
- "pblock_remove" on page 102 removes a `pblock` entry by name from a `pblock`. The entry is not disposed. Use "param_free" on page 95 to free the memory used by the entry.

- "param_free" on page 95 frees the memory for the given `pblock` entry.
- "pblock_pblock2str" on page 101 creates a new string containing all of the name-value pairs from a `pblock` in the form "*name*=*value name*=*value*." This can be a useful function for debugging.

Protocol Utilities for Service SAFs

Protocol utilities provide functionality necessary to implement `Service` SAFs:

- "protocol_status" on page 110 sets the HTTP response status code and reason phrase.
- "protocol_start_response" on page 109 sends the HTTP response and all HTTP headers to the browser.

Memory Management

Memory management routines provide fast, platform-independent versions of the standard memory management routines. They also prevent memory leaks by allocating from a temporary memory (called "pooled" memory) for each request, and then disposing the entire pool after each request. There are wrappers for standard memory routines for using permanent memory. To disable the server's pooled memory allocator for debugging, see the built-in SAF `pool-init` in the *Sun Java System Web Server 7.0 Administrator's Configuration File Reference*.

- "MALLOC" on page 83
- "FREE" on page 78
- "PERM_STRDUP" on page 106
- "REALLOC" on page 113
- "CALLOC" on page 63
- "PERM_MALLOC" on page 104
- "PERM_FREE" on page 104
- "PERM_STRDUP" on page 106
- "PERM_REALLOC" on page 105
- "PERM_CALLOC" on page 103

File I/O

The file I/O functions provide platform-independent, thread-safe file I/O routines.

- "system_fopenRO" on page 124 opens a file for read-only access.
- "system_fopenRW" on page 124 opens a file for read-write access, creating the file if necessary.
- "system_fopenWA" on page 125 opens a file for write-append access, creating the file if necessary.

- "system_fclose" on page 123 closes a file.
- "system_fread" on page 126 reads from a file.
- "system_fwrite" on page 126 writes to a file.
- "system_fwrite_atomic" on page 127 locks the given file before writing to it. This avoids interference between simultaneous writes by multiple processes or threads.

Network I/O

Network I/O functions provide platform-independent, thread-safe network I/O routines. These routines work with SSL when it is enabled.

- "netbuf_grab" on page 91 reads from a network buffer's socket into the network buffer.
- "netbuf_getbytes" on page 89 gets a character from a network buffer.
- "net_flush" on page 84 flushes buffered data.
- "net_read" on page 85 reads bytes from a specified socket into a specified buffer.
- "net_sendfile" on page 86 sends the contents of a specified file to a specified a socket.
- "net_write" on page 87 writes to the network socket.

Threads

Thread functions include functions for creating your own threads that are compatible with the server's threads. There are also routines for critical sections and condition variables.

- "systhread_start" on page 134 creates a new thread.
- "systhread_sleep" on page 134 puts a thread to sleep for a given time.
- "crit_init" on page 68 creates a new critical section variable.
- "crit_enter" on page 67 gains ownership of a critical section.
- "crit_exit" on page 67 surrenders ownership of a critical section.
- "crit_terminate" on page 68 disposes of a critical section variable.
- "condvar_init" on page 65 creates a new condition variable.
- "condvar_notify" on page 65 awakens any threads blocked on a condition variable.
- "condvar_wait" on page 66 blocks on a condition variable.
- "condvar_terminate" on page 66 disposes of a condition variable.
- "prepare_nsapi_thread" on page 107 allows threads that are not created by the server to act like server-created threads.

Utilities

Utility functions include platform-independent, thread-safe versions of many standard library functions (such as string manipulation), as well as new utilities useful for NSAPI.

- "daemon_atrestart" on page 69 registers a user function to be called when the server is sent a restart signal (`HUP`) or at shutdown.
- "util_hostname" on page 142 gets the local host name as a fully qualified domain name.
- "util_later_than" on page 143 compares two dates.
- "util_sprintf" on page 145 is the same as the standard library routine `sprintf()`.
- "util_strftime" on page 147 is the same as the standard library routine `strftime()`.
- "util_uri_escape" on page 148 converts the special characters in a string into URI-escaped format.
- "util_uri_unescape" on page 150 converts the URI-escaped characters in a string back into special characters.

Note – You cannot use an embedded null in a string, because NSAPI functions assume that a null is the end of the string. Therefore, passing unicode-encoded content through an NSAPI plug-in does not work.

Virtual Server

The virtual server functions provide routines for retrieving information about virtual servers.

- "request_get_vs" on page 114 finds the virtual server to which a request is directed.
- "vs_alloc_slot" on page 151 allocates a new slot for storing a pointer to data specific to a certain virtual server.
- "vs_get_data" on page 152 finds the value of a pointer to data for a given virtual server and slot.
- "vs_get_default_httpd_object" on page 153 obtains a pointer to the default (or root) object from the virtual server's virtual server class configuration.
- "vs_get_doc_root" on page 153 finds the document root for a virtual server.
- "vs_get_httpd_objset" on page 154 obtains a pointer to the virtual server class configuration for a given virtual server.
- "vs_get_id" on page 154 finds the ID of a virtual server.
- "vs_get_mime_type" on page 155 determines the MIME type that would be returned in the `content-type:` header for the given URI.
- "vs_lookup_config_var" on page 155 finds the value of a configuration variable for a given virtual server.

- "vs_register_cb" on page 156 allows a plug-in to register functions that will receive notifications of virtual server initialization and destruction events.
- "vs_set_data" on page 156 sets the value of a pointer to data for a given virtual server and slot.
- "vs_translate_uri" on page 157 translates a URI as though it were part of a request for a specific virtual server.

Required Behavior of SAFs for Each Directive

When writing a new SAF, you should define it to do certain things, depending on which stage of the request-handling process will invoke it. For example, SAFs to be invoked during the `Init` stage must conform to different requirements than SAFs to be invoked during the `Service` stage.

The `rq` parameter is the primary mechanism for passing along information throughout the request-response process. On input to a SAF, `rq` contains whatever values were inserted or modified by previously executed SAFs. On output, `rq` contains any modifications or additional information inserted by the SAF. Some SAFs depend on the existence of specific information provided at an earlier step in the process. For example, a `PathCheck` SAF retrieves values in `rq->vars` that were previously inserted by an `AuthTrans` SAF.

This section outlines the expected behavior of SAFs used at each stage in the request-handling process.

- "Init SAFs" on page 32
- "AuthTrans SAFs" on page 32
- "NameTrans SAFs" on page 32
- "PathCheck SAFs" on page 32
- "ObjectType SAFs" on page 33
- "Input SAFs" on page 33
- "Output SAFs" on page 33
- "Service SAFs" on page 33
- "Error SAFs" on page 34
- "AddLog SAFs" on page 34

  For more detailed information about these SAFs, see the *Sun Java System Web Server 7.0 Administrator's Configuration File Reference*.

Init SAFs

- Purpose: Initialize at startup.
- Called at server startup and restart.
- `rq` and `sn` are NULL.
- Initialize any shared resources such as files and global variables.
- Can register callback function with `daemon_atrestart()` to clean up.
- On error, insert `error` parameter into `pb` describing the error and return `REQ_ABORTED`.
- If successful, return `REQ_PROCEED`.

AuthTrans SAFs

- Purpose: Verify any authorization information.
- Return `REQ_PROCEED` if the user was successfully and completely authenticated, `REQ_NOACTION` otherwise.

NameTrans SAFs

- Purpose: Convert logical URI to physical path.
- Perform operations on logical path (`ppath` in `rq->vars`) to convert it into a full local file system path.
- Return `REQ_PROCEED` if `ppath` in `rq->vars` contains the full local file system path, or `REQ_NOACTION` if not.
- To redirect the client to another site, add `url` to `rq->vars` with full URL (for example, `http://www.sun.com/`). Call `protocol_status()` to set HTTP response status to `PROTOCOL_REDIRECT`, `NULL`. Return `REQ_ABORTED`.

PathCheck SAFs

- Purpose: Check path validity and user's access rights.
- Check `auth-type`, `auth-user`, and/or `auth-group` in `rq->vars`.
- Return `REQ_PROCEED` if user (and group) is authorized for this area (`ppath` in `rq->vars`).
- If not authorized, insert `WWW-Authenticate` to `rq->srvhdrs` with a value such as: `Basic; Realm=\"Our private area\"`. Call `protocol_status()` to set HTTP response status to `PROTOCOL_UNAUTHORIZED`. Return `REQ_ABORTED`.

ObjectType SAFs

- Purpose: Determine `content-type` of data.
- If `content-type` in `rq->srvhdrs` already exists, return `REQ_NOACTION`.
- Determine the MIME type and create `content-type` in `rq->srvhdrs`
- Return `REQ_PROCEED` if `content-type` is created, `REQ_NOACTION` otherwise.

Input SAFs

- Purpose: Insert filters that process incoming (client-to-server) data.
- `Input` SAFs are executed when a plug-in or the server first attempts to read entity body data from the client.
- `Input` SAFs are executed at most once per request.
- Return `REQ_PROCEED` to indicate success, or `REQ_NOACTION` to indicate it performed no action.

Output SAFs

- Purpose: Insert filters that process outgoing (server-to-client) data.
- `Output` SAFs are executed when a plug-in or the server first attempts to write entity body data from the client.
- `Output` SAFs are executed at most once per request.
- Return `REQ_PROCEED` to indicate success, or `REQ_NOACTION` to indicate it performed no action.

Service SAFs

- Purpose: Generate and send the response to the client.
- A `Service` SAF is only called if each of the optional parameters `type`, `method`, and `query` specified in the directive in `obj.conf` match the request.
- Remove existing `content-type` from `rq->srvhdrs`. Insert correct `content-type` in `rq->srvhdrs`.
- Create any other headers in `rq->srvhdrs`.
- Call "protocol_status" on page 110 to set HTTP response status.
- Call "protocol_start_response" on page 109 to send HTTP response and headers.
- Generate and send data to the client using "net_write" on page 87.
- Return `REQ_PROCEED` if successful, `REQ_EXIT` on write error, `REQ_ABORTED` on other failures.

Error SAFs

- Purpose: Respond to an HTTP status error condition.
- The `Error` SAF is only called if each of the optional parameters `code` and `reason` specified in the directive in `obj.conf` match the current error.
- `Error` SAFs do the same as `Service` SAFs, but only in response to an HTTP status error condition.

AddLog SAFs

- Purpose: Log the transaction to a log file.
- `AddLog` SAFs can use any data available in `pb`, `sn`, or `rq` to log this transaction.
- Return `REQ_PROCEED`.

CGI to NSAPI Conversion

You may have a need to convert a CGI variable into an SAF using NSAPI. Since the CGI environment variables are not available to NSAPI, you retrieve them from the NSAPI parameter blocks. The table below indicates how each CGI environment variable can be obtained in NSAPI.

Keep in mind that your code must be thread-safe under NSAPI. You should use NSAPI functions that are thread-safe. Also, you should use the NSAPI memory management and other routines for speed and platform independence.

TABLE 1–5 Parameter Blocks for CGI Variables

CGI getenv()	NSAPI
`AUTH_TYPE`	`pblock_findval("auth-type", rq->vars);`
`AUTH_USER`	`pblock_findval("auth-user", rq->vars);`
`CONTENT_LENGTH`	`pblock_findval("content-length", rq->headers);`
`CONTENT_TYPE`	`pblock_findval("content-type", rq->headers);`
`GATEWAY_INTERFACE`	`"CGI/1.1"`
`HTTP_*`	`pblock_findval( "*", rq->headers);` (* is lowercase; dash replaces underscore)
`PATH_INFO`	`pblock_findval("path-info", rq->vars);`

TABLE 1–5 Parameter Blocks for CGI Variables *(Continued)*

CGI getenv()	NSAPI
PATH_TRANSLATED	`pblock_findval("path-translated", rq->vars);`
QUERY_STRING	`pblock_findval("query", rq->reqpb);`
REMOTE_ADDR	`pblock_findval("ip", sn->client);`
REMOTE_HOST	`session_dns(sn) ? session_dns(sn) : pblock_findval("ip", sn->client);`
REMOTE_IDENT	`pblock_findval( "from", rq->headers);` (not usually available)
REMOTE_USER	`pblock_findval("auth-user", rq->vars);`
REQUEST_METHOD	`pblock_findval("method", req->reqpb);`
SCRIPT_NAME	`pblock_findval("uri", rq->reqpb);`
SERVER_NAME	`char *util_hostname();`
SERVER_PORT	`conf_getglobals()->Vport;` (as a string)
SERVER_PROTOCOL	`pblock_findval("protocol", rq->reqpb);`
SERVER_SOFTWARE	`system_version()`
Sun Java System-specific:	
CLIENT_CERT	`pblock_findval("auth-cert", rq->vars) ;`
HOST	`char *session_maxdns(sn);` (may be null)
HTTPS	`security_active ? "ON" : "OFF";`
HTTPS_KEYSIZE	`pblock_findval("keysize", sn->client);`
HTTPS_SECRETKEYSIZE	`pblock_findval("secret-keysize", sn->client);`
SERVER_URL	`protocol_uri2url_dynamic("","", sn, rq);`

◆ ◆ ◆ CHAPTER 2

Creating Custom Filters

This chapter describes how to create custom filters that can be used to intercept and possibly modify the content presented to or generated by another function.

This chapter has the following sections:

Future Compatibility Issues

The NSAPI interface may change in a future version of Sun Java System Web Server.

To keep your custom plug-ins upgradable, do the following:

- Make sure plug-in users know how to edit the configuration files (such as `magnus.conf` and `obj.conf`) manually. The plug-in installation software should not be used to edit these configuration files.
- Keep the source code so you can recompile the plug-in.

The NSAPI Filter Interface

The NSAPI filter interface complements the NSAPI Server Application Function (SAF) interface. Filters make it possible to intercept and possibly modify data sent to and from the server. The server communicates with a filter by calling the filter's filter methods. Each filter implements one or more filter methods. A filter method is a C function that performs a specific operation, such as processing data sent by the server.

Filter Methods

This section describes the filter methods that a filter can implement. To create a filter, a filter developer implements one or more of these methods.

This section describes the following filter methods:

- insert
- remove
- flush
- read
- write
- writev
- sendfile

For more information about these methods, see Chapter 5.

C Prototypes for Filter Methods

Following is a list of C prototypes for the filter methods:

```
int insert(FilterLayer *layer, pblock *pb);
void remove(FilterLayer *layer);
int flush(FilterLayer *layer);
int read(FilterLayer *layer, void *buf, int amount, int timeout);
int write(FilterLayer *layer, const void *buf, int amount);
int writev(FilterLayer *layer, const struct iovec *iov, int iov_size);
int sendfile(FilterLayer *layer, sendfiledata *sfd);
```

The `layer` parameter is a pointer to a `FilterLayer` data structure, which contains variables related to a particular instance of a filter.

Following is a list of the most important fields in the `FilterLayer` data structure:

- `context->sn`: Contains information relating to a single TCP/IP session (the same `sn` pointer that's passed to SAFs).
- `context->rq`: Contains information relating to the current request (the same `rq` pointer that's passed to SAFs).
- `context->data`: Pointer to filter-specific data.
- `lower`: A platform-independent socket descriptor used to communicate with the next filter in the stack.

  The meaning of the `context->data` field is defined by the filter developer. Filters that must maintain state information across filter method calls can use `context->data` to store that information.

  For more information about `FilterLayer`, see "FilterLayer" on page 167.

insert

The `insert` filter method is called when an SAF such as `insert-filter` calls the `filter_insert` function to request that a specific filter be inserted into the filter stack. Each filter must implement the `insert` filter method.

When `insert` is called, the filter can determine whether it should be inserted into the filter stack. For example, the filter could inspect the `content-type` header in the `rq->srvhdrs` pblock to determine whether it is interested in the type of data that will be transmitted. If the filter should not be inserted, the `insert` filter method should indicate this by returning `REQ_NOACTION`.

If the filter should be inserted, the `insert` filter method provides an opportunity to initialize this particular instance of the filter. For example, the `insert` method could allocate a buffer with `MALLOC` and store a pointer to that buffer in `layer->context->data`.

The filter is not part of the filter stack until after `insert` returns. As a result, the `insert` method should not attempt to read from, write to, or otherwise interact with the filter stack.

See Also

insert in Chapter 5

remove

The `remove` filter method is called when a filter stack is destroyed (that is, when the corresponding socket descriptor is closed), when the server finishes processing the request the filter was associated with, or when an SAF such as `remove-filter` calls the `filter_remove` function. The `remove` filter method is optional.

The `remove` method can be used to clean up any data the filter allocated in `insert` and to pass any buffered data to the next filter by calling `net_write(layer->lower, ...)`.

See Also

remove in Chapter 5

flush

The `flush` filter method is called when a filter or SAF calls the `net_flush` function. The `flush` method should pass any buffered data to the next filter by calling `net_write(layer->lower, ...)`. The `flush` method is optional, but it should be implemented by any filter that buffers outgoing data.

See Also

flush in Chapter 5

read

The `read` filter method is called when a filter or SAF calls the `net_read` function. Filters that are interested in incoming data (data sent from a client to the server) implement the `read` filter method.

Typically, the `read` method will attempt to obtain data from the next filter by calling `net_read(layer->lower, ...)`. The `read` method may then modify the received data before returning it to its caller.

See Also

read in Chapter 5

write

The `write` filter method is called when a filter or SAF calls the `net_write` function. Filters that are interested in outgoing data (data sent from the server to a client) implement the `write` filter method.

Typically, the `write` method will pass data to the next filter by calling `net_write(layer->lower, ...)`. The `write` method may modify the data before calling `net_write`. For example, the `http-compression` filter compresses data before passing it on to the next filter.

If a filter implements the `write` filter method but does not pass the data to the next layer before returning to its caller (that is, if the filter buffers outgoing data), the filter should also implement the `flush` method.

See Also

write in Chapter 5

sendfile

The `sendfile` filter method performs a function similar to the `writev` filter method, but it sends a file directly instead of first copying the contents of the file into a buffer. It is not necessary to implement the `sendfile` filter method; if a filter implements the `write` filter method but not the `sendfile` filter method, the server will use the `write` method instead of the `sendfile` method. A filter should not implement the `sendfile` method unless it also implements the `write` method.

Under some circumstances, the server may run slightly faster when filters that implement the `write` filter method also implement the `sendfile` filter method.

See Also

sendfile in Chapter 5

writev

The `writev` filter method performs the same function as the `write` filter method, but the format of its parameters is different. It is not necessary to implement the `writev` filter method; if a filter implements the `write` filter method but not the `writev` filter method, the server uses the `write` method instead of the `writev` method. A filter should not implement the `writev` method unless it also implements the `write` method.

Under some circumstances, the server may run slightly faster when filters that implement the `write` filter method also implement the `writev` filter method.

See Also

writev in Chapter 5

Position of Filters in the Filter Stack

All data sent to the server (such as the result of an HTML form) or sent from the server (such as the output of a JSP page) is passed through a set of filters known as a filter stack. The server creates a separate filter stack for each connection. While processing a request, individual filters can be inserted into and removed from the stack.

Different types of filters occupy different positions within a filter stack. Filters that deal with application-level content (such filters that translates a page from XHTML to HTML) occupy a higher position than filters that deal with protocol-level issues (such as filters that format HTTP responses). When two or more filters are defined to occupy the same position in the filter stack, filters that were inserted later will appear higher than filters that were inserted earlier.

Filters positioned higher in the filter stack are given an earlier opportunity to process outgoing data, while filters positioned lower in the stack are given an earlier opportunity to process incoming data. For example, in the following figure, the `xml-to-xhtml` filter is given an earlier opportunity to process outgoing data than the `xhtml-to-html` filter.

When you create a filter with the `filter_create` function, you specify what position your filter should occupy in the stack. You can also use the `init-filter-order` `Init` SAF to control the position of specific filters within filter stacks. For example, `init-filter-order` can be used to ensure that a filter that converts outgoing XML to XHTML is inserted above a filter that converts outgoing XHTML to HTML.

For more information, see "filter_create" on page 73 and `init-filter-order` in the *Sun Java System Web Server 7.0 Administrator's Configuration File Reference*.

Filters that Alter Content-Length

Filters that can alter the length of an incoming request body or outgoing response body must take special steps to ensure interoperability with other filters and SAFs.

Filters that process incoming data are referred to as input filters. If an input filter can alter the length of the incoming request body (for example, if a filter decompresses incoming data) and there is a `Content-Length` header in the `rq->headers` pblock, the filter's `insert` filter method should remove the `Content-Length` header and replace it with a `Transfer-encoding: identity` header as follows:

```
pb_param *pp;

pp = pblock_remove("content-length", layer->context->rq->headers);
if (pp != NULL) {
   param_free(pp);
   pblock_nvinsert("transfer-encoding", "identity", layer->context->rq->headers);
}
```

Because some SAFs expect a `content-length` header when a request body is present, before calling the first `Service` SAF the server will insert all relevant filters, read the entire request body, and compute the length of the request body after it has been passed through all input filters. However, by default, the server will read at most 8192 bytes of request body data. If the request body exceeds 8192 bytes after being passed through the relevant input filters, the request will be cancelled. For more information, see the description of `ChunkedRequestBufferSize` in the "Syntax and Use of `obj.conf`" chapter in the *Sun Java System Web Server 7.0 Administrator's Configuration File Reference*.

Filters that process outgoing data are referred to as output filters. If an output filter can alter the length of the outgoing response body (for example, if the filter compresses outgoing data), the filter's `insert` filter method should remove the `Content-Length` header from `rq->srvhdrs` as follows:

```
pb_param *pp;

pp = pblock_remove("content-length", layer->context->rq->srvhdrs);
if (pp != NULL)
   param_free(pp);
```

Creating and Using Custom Filters

Custom filters are defined in shared libraries that are loaded and called by the server. The general steps for creating a custom filter are as follows:

▼ To create a custom filter

1. **"Write the Source Code" on page 44 using the NSAPI functions.**
2. **"Compile and Link" on page 44 the source code to create a shared library (`.so`, `.sl`, or `.dll`) file.**
3. **"Load and Initialize the Filter" on page 45 by editing the `magnus.conf` file.**
4. **"Instruct the Server to Insert the Filter" on page 45 by editing the `obj.conf` file to insert your custom filter(s) at the appropriate time.**
5. **"Restart the Server" on page 46.**

6 **"Test the Filter" on page 46 by accessing your server from a browser with a URL that triggers your filter.**

These steps are described in greater detail in the following sections.

Write the Source Code

Write your custom filter methods using NSAPI functions. For a summary of the NSAPI functions specific to filter development, see"Overview of NSAPI Functions for Filter Development" on page 46 and "Filter Methods" on page 38 for the filter method prototypes.

The filter must be created by a call to `filter_create`. Typically, each plug-in defines an `nsapi_module_init` function that is used to call `filter_create` and perform any other initialization tasks. For more information, see "nsapi_module_init" on page 92 and "filter_create" on page 73.

Filter methods are invoked whenever the server or an SAF calls certain NSAPI functions such as `net_write` or `filter_insert`. As a result, filter methods can be invoked from any thread and should only block using NSAPI functions (for example, `crit_enter` and `net_read`). If a filter method blocks using other functions (for example, the Windows `WaitForMultipleObjects` and `ReadFile` functions), the server may hang. Also, shared objects that define filters should be loaded with the `NativeThread="no"` flag, as described in "Load and Initialize the Filter" on page 45.

If a `filter` method must block using a non-NSAPI function, `KernelThreads 1` should be set in `magnus.conf`. For more information about `KernelThreads`, see the description in the chapter *Syntax and Use of* `magnus.conf` in the *Sun Java System Web Server 7.0 Administrator's Configuration File Reference*.

Keep the following in mind when writing your filter:

- Write thread-safe code
- IO should only be performed using the NSAPI functions documented in "File I/O" on page 28
- Thread synchronization should only be performed using NSAPI functions documented in "Threads" on page 29
- Blocking may affect performance
- Carefully check and handle all errors

For examples of custom filters, see Chapter 3.

Compile and Link

Filters are compiled and linked in the same way as SAFs. For more information, see "Compile and Link" on page 22.

Load and Initialize the Filter

For each shared library (plug-in) containing custom filters to be loaded into the server, add an `Init` directive that invokes the `load-modules` SAF to `magnus.conf`. The syntax for a directive that loads a filter plug-in is:

```
Init fn=load-modules shlib=path NativeThread="no"
```

- `shlib` is the local file system path to the shared library (plug-in).
- `NativeThread` indicates whether the plug-in requires native threads. Filters should be written to run on any type of thread (see "Write the Source Code" on page 44).

  When the server encounters such a directive, it calls the plug-in's `nsapi_module_init` function to initialize the filter.

Instruct the Server to Insert the Filter

Add an `Input` or `Output` directive to `obj.conf` to instruct the server to insert your filter into the filter stack. The format of the directive is as follows:

Directive `fn=insert-filter filter="`*filter-name*`" [`*name1*`="`*value1*`"]...[`*nameN*`="`*valueN*`"]`

- *Directive* is `Input` or `Output`.
- *filter-name* is the name of the filter, as passed to `filter_create`, to insert.
- *nameN*`="`*valueN*`"` are the names and values of parameters that are passed to the filter's `insert` filter method.

  Filters that process incoming data should be inserted using an `Input` directive. Filters that process outgoing data should be inserted using an `Output` directive.

  To ensure that your filter is inserted whenever a client sends a request, add the `Input` or `Output` directive to the default object. For example, the following portion of `obj.conf` instructs the server to insert a filter named `example-replace` and pass it two parameters, `from` and `to`:

```
<Object name="default">
Output fn=insert-filter
       filter="example-replace"
       from="Old String"
       to="New String"
...
</Object>
```

Restart the Server

For the server to load your plug-in, you must restart the server. A restart is required for all plug-ins that implement SAFs and/or filters.

Test the Filter

Test your filter by accessing your server from a web browser. You should disable caching in your web browser so that the server is sure to be accessed. In Mozilla Firefox, you may hold the shift key while clicking the Reload button to ensure that the cache is not used. Examine the access and error logs to help with debugging.

Overview of NSAPI Functions for Filter Development

NSAPI provides a set of C functions that are used to implement SAFs and filters. This section lists the functions that are specific to the development of filters. All of the public routines are described in detail in Chapter 5.

The NSAPI functions specific to the development of filters are:

- "filter_create" on page 73 creates a new filter
- "filter_insert" on page 75 inserts the specified filter into a filter stack
- "filter_remove" on page 77 removes the specified filter from a filter stack
- "filter_name" on page 77 returns the name of the specified filter
- "filter_find" on page 75 finds an existing filter given a filter name
- "filter_layer" on page 76 returns the layer in a filter stack that corresponds to the specified filter

◆ ◆ ◆ CHAPTER 3

Examples of Custom SAFs and Filters

This chapter provides examples of custom Sever Application Functions (SAFs) and filters for each directive in the request-response process. You may wish to use these examples as the basis for implementing your own custom SAFs and filters. For more information about creating your own custom SAFs, see Chapter 1 and for information about creating your own filters, see Chapter 2.

Before writing custom SAFs, you should be familiar with the request-response process and the role of the configuration file `obj.conf`. See the *Sun Java System Web Server 7.0 Administrator's Configuration File Reference* for information on the `obj.conf` file.

Before writing your own SAF, check to see if an existing SAF serves your purpose. The predefined SAFs are discussed in the *Sun Java System Web Server 7.0 Administrator's Configuration File Reference*.

For a list of the NSAPI functions for creating new SAFs, see Chapter 5.

This chapter has the following sections:

Examples Bundled With the Server

The *install_dir*/samples/nsapi directory contains examples of source code for SAFs.

You can use the example.mak (Windows) or Makefile (UNIX) makefile in the same directory to compile the examples and create shared libraries containing the functions in all of the example files.

To test an example, load the examples shared library into the server by adding the following directive in the Init section of magnus.conf:

```
Init fn=load-modules
     shlib=examples.so/dll
     funcs=function1,...,functionN
```

The shlib parameter specifies the path to the shared library (for example, ../../samples/nsapi/examples.so), and the funcs parameter specifies the functions to load from the shared library.

If the example uses an initialization function, be sure to specify the initialization function in the funcs argument to load-modules, and also add an Init directive to call the initialization function.

For example, the PathCheck example implements the restrict-by-acf function, which is initialized by the acf-init function. The following directive loads both these functions:

```
Init fn=load-modules
     shlib="path"
     funcs=acf-init,restrict-by-acf
```

The following directive calls the acf-init function during server initialization:

```
Init fn=acf-init file=extra-arg
```

To invoke the new SAF at the appropriate step in the response handling process, add an appropriate directive in the object to which it applies, for example:

```
PathCheck fn=restrict-by-acf
```

After adding new Init directives to magnus.conf, you always need to restart the Web Server to load the changes, since Init directives are only applied during server initialization.

AuthTrans Example

This simple example of an `AuthTrans` function demonstrates how to use your own custom ways of verifying the user name and password that a remote client provided is accurate. This program uses a hard-coded table of user names and passwords and checks a given user's password against the one in the static data array. The *userdb* parameter is not used in this function.

`AuthTrans` directives work in conjunction with `PathCheck` directives. Generally, an `AuthTrans` function checks if the user name and password associated with the request are acceptable, but it does not allow or deny access to the request; it leaves that to a `PathCheck` function.

`AuthTrans` functions get the user name and password from the headers associated with the request. When a client initially makes a request, the user name and password are unknown so the `AuthTrans` function and `PathCheck` function work together to reject the request, since they can't validate the user name and password. When the client receives the rejection, the usual response is for it to present a dialog box asking the user for their user name and password, and then the client submits the request again, this time including the user name and password in the headers.

In this example, the `hardcoded-auth` function, which is invoked during the `AuthTrans` step, checks if the user name and password correspond to an entry in the hard-coded table of users and passwords.

Installing the AuthTrans Example

To install the function on the Web Server, add the following `Init` directive to `magnus.conf` to load the compiled function:

```
Init fn=load-modules
     shlib="path"
     funcs=hardcoded-auth
```

Inside the default object in `obj.conf`, add the following `AuthTrans` directive:

```
AuthTrans fn=basic-auth
          auth-type="basic"
          userfn=hardcoded-auth
          userdb=unused
```

Note that this function does not actually enforce authorization requirements, it only takes given information and tells the server if it is correct or not. The `PathCheck` function `require-auth` performs the enforcement, so add the following `PathCheck` directive as well:

```
PathCheck fn=require-auth
          realm="test realm"
          auth-type="basic"
```

The source code for this example is in the `auth.c` file in the *install_dir*`/samples/nsapi/`directory.

NameTrans Example

The `ntrans.c` file in the `samples/nsapi` subdirectory of the server root directory contains source code for two example `NameTrans` functions:

- `explicit_pathinfo`

  This example allows the use of explicit extra path information in a URL.

- `https_redirect`

  This example redirects the URL if the client is a particular version of Netscape Navigator.

  This section discusses the first example. Look at the source code in `ntrans.c` for the second example.

Note – A `NameTrans` function is used primarily to convert the logical URL in `ppath` in `rq->vars` to a physical path name. However, the example discussed here, `explicit_pathinfo`, does not translate the URL into a physical path name; it changes the value of the requested URL. See the second example, `https_redirect`, in `ntrans.c` for an example of a `NameTrans` function that converts the value of `ppath` in `rq->vars` from a URL to a physical path name.

The `explicit_pathinfo` example allows URLs to explicitly include extra path information for use by a CGI program. The extra path information is delimited from the main URL by a specified separator, such as a comma.

For example:

```
http://server-name/cgi/marketing,/jan/releases/hardware
```

In this case, the URL of the requested resource (which would be a CGI program) is `http://`*server-name*`/cgi/marketing`, and the extra path information to give to the CGI program is `/jan/releases/hardware`.

When choosing a separator, be sure to pick a character that is never used as part of the real URL.

The `explicit_pathinfo` function reads the URL, strips out everything following the comma, and puts it in the `path-info` field of the `vars` field in the `request` object (`rq->vars`). CGI programs can access this information through the `PATH_INFO` environment variable.

One side effect of `explicit_pathinfo` is that the `SCRIPT_NAME` CGI environment variable has the separator character tacked onto the end.

`NameTrans` directives usually return `REQ_PROCEED` when they change the path, so that the server does not process any more `NameTrans` directives. However, in this case we want name translation to continue after we have extracted the path info, since we have not yet translated the URL to a physical path name.

Installing the NameTrans Example

To install the function on the Web Server, add the following `Init` directive to `magnus.conf` to load the compiled function:

```
Init fn=load-modules
     shlib="path"
     funcs=explicit-pathinfo
```

Inside the default object in `obj.conf`, add the following `NameTrans` directive:

```
NameTrans fn=explicit-pathinfo
          separator=","
```

This `NameTrans` directive should appear before other `NameTrans` directives in the default object.

The source code for this example is in the `ntrans.c` file in the *install_dir*`/samples/nsapi/`directory.

PathCheck Example

The example in this section demonstrates how to implement a custom SAF for performing path checks. This example simply checks if the requesting host is on a list of allowed hosts.

The `Init` function `acf-init` loads a file containing a list of allowable IP addresses with one IP address per line. The `PathCheck` function `restrict_by_acf` gets the IP address of the host that is making the request and checks if it is on the list. If the host is on the list, it is allowed access; otherwise, access is denied.

For simplicity, the `stdio` library is used to scan the IP addresses from the file.

Installing the PathCheck Example

To load the shared object containing your functions, add the following line in the `Init` section of the `magnus.conf` file:

```
Init fn=load-modules
     shlib="path"
     funcs=acf-init, restrict-by-acf
```

To call `acf-init` to read the list of allowable hosts, add the following line to the `Init` section in `magnus.conf`. (This line must come after the one that loads the library containing `acf-init`).

```
Init fn=acf-init
     file=fileContainingHostsList
```

To execute your custom SAF during the request-response process for some object, add the following line to that object in the `obj.conf` file:

```
PathCheck fn=restrict-by-acf
```

The source code for this example is in `pcheck.c` in the *install_dir*`/samples/nsapi/`directory.

ObjectType Example

The example in this section demonstrates how to implement `html2shtml`, a custom SAF that instructs the server to treat a `.html` file as a `.shtml` file if a `.shtml` version of the requested file exists.

A well-behaved `ObjectType` function checks if the content type is already set, and if so, does nothing except return `REQ_NOACTION`.

```
if(pblock_findval("content-type", rq->srvhdrs))
    return REQ_NOACTION;
```

The primary task an `ObjectType` directive needs to perform is to set the content type (if it is not already set). This example sets it to `magnus-internal/parsed-html` in the following lines:

```
/* Set the content-type to magnus-internal/parsed-html */
pblock_nvinsert("content-type", "magnus-internal/parsed-html",
                rq->srvhdrs);
```

The `html2shtml` function looks at the requested file name. If it ends with `.html`, the function looks for a file with the same base name, but with the extension `.shtml` instead. If it finds one, it uses that path and informs the server that the file is parsed HTML instead of regular HTML. Note that this requires an extra `stat` call for every HTML file accessed.

Installing the ObjectType Example

To load the shared object containing your function, add the following line in the `Init` section of the `magnus.conf` file:

```
Init fn=load-modules
     shlib="path"
     funcs=html2shtml
```

To execute the custom SAF during the request-response process for some object, add the following line to that object in the `obj.conf` file:

```
ObjectType fn=html2shtml
```

The source code for this example is in `otype.c` in the *install_dir*`/samples/nsapi/` directory.

Output Example

This section describes an example NSAPI filter named `example-replace`, which examines outgoing data and substitutes one string for another. It shows how you can create a filter that intercepts and modifies outgoing data.

Installing the Output Example

To load the filter, add the following line in the `Init` section of the `magnus.conf` file:

```
Init fn="load-modules"
     shlib=yourlibrary
     NativeThread="no"
```

To execute the filter during the request-response process for some object, add the following line to that object in the `obj.conf` file:

```
Output fn="insert-filter"
       type="text/*"
       filter="example-replace"
       from="iPlanet" to="Sun ONE"
```

The source code for this example is in the `replace.c` file in the *install_dir*`/samples/nsapi/` directory.

Service Example

This section discusses a very simple `Service` function called `simple_service.` All this function does is send a message in response to a client request. The message is initialized by the `init_simple_service` function during server initialization.

For a more complex example, see the file `service.c` in the `examples` directory, which is discussed in "More Complex Service Example" on page 55.

Installing the Service Example

To load the shared object containing your functions, add the following line in the `Init` section of the `magnus.conf` file:

```
Init fn=load-modules
     shlib=yourlibrary
     funcs=simple-service-init,simple-service
```

To call the `simple-service-init` function to initialize the message representing the generated output, add the following line to the `Init` section in `magnus.conf`. (This line must come after the one that loads the library containing `simple-service-init`.)

```
Init fn=simple-service-init
     generated-output="<H1>Generated output msg</H1>"
```

To execute the custom SAF during the request-response process for some object, add the following line to that object in the `obj.conf` file:

```
Service type="text/html"
        fn=simple-service
```

The `type="text/html"` argument indicates that this function is invoked during the `Service` stage only if the `content-type` has been set to `text/html`.

The source code for this example is in the `service.c` file in the *install_dir*`/samples/nsapi` directory.

More Complex Service Example

The `send-images` function is a custom SAF that replaces the `doit.cgi` demonstration available on the iPlanet home pages. When a file is accessed as `/dir1/dir2/something.picgroup`, the `send-images` function checks if the file is being accessed by a Mozilla/1.1 browser. If not, it sends a short error message. The file `something.picgroup` contains a list of lines, each of which specifies a file name followed by a `content-type` (for example, `one.gif image/gif`).

To load the shared object containing your function, add the following line at the beginning of the `magnus.conf` file:

```
Init fn=load-modules shlib=yourlibrary funcs=send-images
```

Also, add the following line to the `mime.types` file:

```
type=magnus-internal/picgroup exts=picgroup
```

To execute the custom SAF during the request-response process for some object, add the following line to that object in the `obj.conf` file (`send-images` takes an optional parameter, `delay`, which is not used for this example):

```
Service method=(GET|HEAD) type=magnus-internal/picgroup fn=send-images
```

The source code for this example is in the `service.c` file in the *install_dir*`/samples/nsapi` directory.

AddLog Example

The example in this section demonstrates how to implement `brief-log`, a custom SAF for logging only three items of information about a request: the IP address, the method, and the URI (for example, `198.93.95.99 GET /jocelyn/dogs/homesneeded.html`).

Installing the AddLog Example

To load the shared object containing your functions, add the following line in the `Init` section of the `magnus.conf` file:

```
Init fn=load-modules
     shlib=yourlibrary
     funcs=brief-init,brief-log
```

To call `brief-init` to open the log file, add the following line to the `Init` section in `magnus.conf`. (This line must come after the one that loads the library containing `brief-init`.)

```
Init fn=brief-init
     file=/tmp/brief.log
```

To execute your custom SAF during the `AddLog` stage for some object, add the following line to that object in the `obj.conf` file:

```
AddLog fn=brief-log
```

The source code for this example is in `addlog.c` file in the *install_dir*`/samples/nsapi` directory.

Quality of Service Example

The code for the `qos-handler` (`AuthTrans`) and `qos-error` (`Error`) SAFs is provided as an example in case you want to define your own SAFs for quality of service handling.

For more information about predefined SAFs, see the *Sun Java System Web Server 7.0 Administrator's Configuration File Reference*.

Installing the Quality of Service Example

Inside the default object in `obj.conf`, add the following `AuthTrans` and `Error` directives:

```
AuthTrans fn=qos-handler
...
Error fn=qos-error code=503
```

The source code for this example is in the `qos.c` file in the `samples/nsapi` subdirectory of the server root directory.

◆ ◆ ◆ CHAPTER 4

Creating Custom Server-parsed HTML Tags

This chapter describes the procedure to create customer server-parsed HTML tags. This chapter contains the following sections:

Defining Custom Server-parsed HTML Tags

HTML files can contain tags that are executed on the server. For general information about server-parsed HTML tags, see the *Sun Java System Web Server 7.0 Developer's Guide*.

In Web Server 7.0, you can define your own server-side tags. For example, you could define the tag `HELLO` to invoke a function that prints `Hello World!` You could have the following code in your `hello.shtml` file:

```
<html>
    <head>
            <title>shtml custom tag example</title>
    </head>
    <body>
            <!--#HELLO-->
    </body>
</html>
```

When the browser displays this code, each occurrence of the `HELLO` tag calls the function.

The steps for defining a customized server-parsed tag are listed below, and described in this chapter.

▼ To Define Customer Server-parsed HTML Tags

1 **"Define the Functions that Implement the Tag" on page 58.**

You must define the tag execution function. You must also define other functions that are called on tag loading and unloading, and on page loading and unloading.

2 **"Write an Initialization Function" on page 61.**

Write an initialization function that registers the tag using the `shtml_add_tag` function.

3 **"Load the New Tag into the Server" on page 62.**

Define the Functions that Implement the Tag

Define the functions that implement the tags in C, using NSAPI.

- Include the header `shtml_public.h`, which is in the directory *install_dir*`/include/shtml`.
- Link against the SHTML shared library in the *install_dir*`/lib` directory. On Windows, the SHTML shared library is named `sshtml.dll`. On UNIX platforms, it is named `libShtml.so` or `libShtml.sl`.

  `ShtmlTagExecuteFunc` is the actual tag handler. It gets called with the usual NSAPI `pblock`, `Session`, and `Request` variables. In addition, it also gets passed the `TagUserData` created from the result of executing the tag loading and page loading functions (if defined) for that tag.

  The signature for the tag execution function is:

```
typedef int (*ShtmlTagExecuteFunc)
            (pblock*, Session*, Request*, TagUserData, TagUserData);
```

  Write the body of the tag execution function to generate the output to replace the tag in the `.shtml` page. Do this in the usual NSAPI way, using the `net_write` NSAPI function, which writes a specified number of bytes to a specified socket from a specified buffer.

  For more information about writing NSAPI plug-ins, see Chapter 1.

  For more information about `net_write` and other NSAPI functions, see Chapter 5.

  The tag execution function must return an `int` that indicates whether the server should proceed to the next instruction in `obj.conf`, which is one of:

- `REQ_PROCEED` -- the execution was successful
- `REQ_NOACTION` -- nothing happened
- `REQ_ABORTED` - - an error occurred
- `REQ_EXIT` - - the connection was lost

The other functions you must define for your tag are:

- `ShtmlTagInstanceLoad`

  This is called when a page containing the tag is parsed. It is not called if the page is retrieved from the browser's cache. It serves as a constructor, the result of which is cached and is passed into `ShtmlTagExecuteFunc` whenever the execution function is called.

- `ShtmlTagInstanceUnload`

  This is a destructor for cleaning up whatever was created in the `ShtmlTagInstanceLoad` function. It gets passed the result that was originally returned from the `ShtmlTagInstanceLoad` function.

- `ShtmlTagPageLoadFunc`

  This is called when a page containing the tag is executed, regardless of whether the page is still in the browser's cache. This provides a way to make information persistent between occurrences of the same tag on the same page.

- `ShtmlTagPageUnLoadFn`

  This is called after a page containing the tag has executed. It provides a way to clean up any allocations done in a `ShtmlTagPageLoadFunc` and hence gets passed the result returned from the `ShtmlTagPageLoadFunc`.

  The signatures for these functions are:

```
#define TagUserData void*
typedef TagUserData (*ShtmlTagInstanceLoad)
                    (const char* tag, pblock*, const char*, size_t);
typedef void (*ShtmlTagInstanceUnload)(TagUserData);
typedef int (*ShtmlTagExecuteFunc)
            (pblock*, Session*, Request*, TagUserData, TagUserData);
typedef TagUserData (*ShtmlTagPageLoadFunc)
                    (block* pb, Session*, Request*);
typedef void (*ShtmlTagPageUnLoadFunc)(TagUserData);
```

Following is the code that implements the `HELLO` tag:

```
/*
 * mytag.c: NSAPI functions to implement #HELLO SSI calls
*/
#include "nsapi.h"
#include "shtml/shtml_public.h"
/* FUNCTION : mytag_con
 *
 * DESCRIPTION: ShtmlTagInstanceLoad function
 */
#ifdef __cplusplus
```

```
extern "C"
#endif
TagUserData
mytag_con(const char* tag, pblock* pb, const char* c1, size_t t1)
{
    return NULL;
}
/* FUNCTION : mytag_des
 *
 * DESCRIPTION: ShtmlTagInstanceUnload
 */
#ifdef __cplusplus
extern "C"
#endif
void
mytag_des(TagUserData v1)
{
}
/* FUNCTION : mytag_load
 * DESCRIPTION: ShtmlTagPageLoadFunc
 */
#ifdef __cplusplus
extern "C"
#endif
TagUserData
mytag_load(pblock *pb, Session *sn, Request *rq)
{
    return NULL;
}
/* FUNCTION : mytag_unload
 *
 * DESCRIPTION: ShtmlTagPageUnloadFunc
 */
#
#ifdef __cplusplus
extern "C"
#endif
void
mytag_unload(TagUserData v2)
{
}
/* FUNCTION : mytag
    * DESCRIPTION: ShtmlTagExecuteFunc
 */
#ifdef __cplusplus
extern "C"
#endif
int
```

```
mytag(pblock* pb, Session* sn, Request* rq, TagUserData t1, TagUserData t2)
{
    char* buf;
    int length;
    char* client;
    buf = (char *) MALLOC(100*sizeof(char));
    length = util_sprintf(buf, "<h1>Hello World! </h1>", client);
    if (net_write(sn->csd, buf, length) == IO_ERROR)
    {
        FREE(buf);
        return REQ_ABORTED;
    }
    FREE(buf);
    return REQ_PROCEED;
}
/* FUNCTION : mytag_init
    * DESCRIPTION: initialization function, calls shtml_add_tag() to
* load new tag
*/
#
#ifdef __cplusplus
extern "C"
#endif
int
mytag_init(pblock* pb, Session* sn, Request* rq)
{
    int retVal = 0;
// NOTE: ALL arguments are required in the shtml_add_tag() function
    retVal = shtml_add_tag("HELLO", mytag_con, mytag_des, mytag, mytag_load, mytag_unload);
return retVal;
}
/* end mytag.c */
```

Write an Initialization Function

In the initialization function for the shared library that defines the new tag, register the tag using the function `shtml_add_tag`. The signature is:

```
NSAPI_PUBLIC int shtml_add_tag (
             const char* tag,
             ShtmlTagInstanceLoad ctor,
             ShtmlTagInstanceUnload dtor,
             ShtmlTagExecuteFunc execFn,
             ShtmlTagPageLoadFunc pageLoadFn,
             ShtmlTagPageUnLoadFunc pageUnLoadFn);
```

Any of these arguments can return NULL except for the `tag` and `execFn`.

Load the New Tag into the Server

After creating the shared library that defines the new tag, you load the library into the Web Server in the usual way for NSAPI plug-ins. That is, add the following directives to the configuration file `magnus.conf`:

Add an `Init` directive whose `fn` parameter is `load-modules` and whose `shlib` parameter is the shared library to load. For example, if you compiled your tag into the shared object *install_dir*`/hello.so`, it would be:

```
Init funcs="mytag,mytag_init" shlib="install_dir/hello.so" fn="load-modules"
```

Add another `Init` directive whose `fn` parameter is the initialization function in the shared library that uses `shtml_add_tag` to register the tag. For example:

```
Init fn="mytag_init"
```

◆ ◆ ◆ CHAPTER 5

NSAPI Function and Macro Reference

This chapter lists all the public C functions and macros of the Netscape Server Applications Programming Interface (NSAPI). These are the functions you use when writing your own Server Application Functions (SAFs) and filters.

Each function provides the name, syntax, parameters, return value, a description of what the function does, and sometimes an example of its use and a list of related functions.

For more information on data structures, see Chapter 6.

NSAPI Functions and Macros

For an alphabetical list of function names, see Appendix B.

C	D	F	I	L	M	N	P	R	S	U	V	W

C

CALLOC

The `CALLOC` macro is a platform-independent substitute for the C library routine `calloc`. It allocates `size` bytes from the request's memory pool and initializes the memory to zeros. The memory can be explicitly freed by a call to `FREE`. If the memory is not explicitly freed, it is automatically freed after processing of the current request has been completed. If pooled memory has been disabled in the configuration file (with the `pool-init` built-in SAF), `PERM-CALLOC` and `CALLOC` both obtain their memory from the system heap. However, since memory allocated by `CALLOC`is automatically freed, it should not be shared with threads.

Syntax

```
void *CALLOC(int size)
```

Returns

A void pointer to a block of memory.

Parameters

`int size` is the number of bytes to allocate.

Example

```
char *name;
name = (char *) CALLOC(100);
```

See Also

"FREE" on page 78, "MALLOC" on page 83, "REALLOC" on page 113, "STRDUP" on page 121, "PERM_CALLOC" on page 103

cinfo_find

The `cinfo_find()` function uses the MIME types information to find the type, encoding, and/or language based on the extension(s) of the Universal Resource Identifier (URI) or local file name. Use this information to send headers (`rq->srvhdrs`) to the client indicating the `content-type`, `content-encoding`, and `content-language` of the data it will be receiving from the server.

The name used is everything after the last slash (/) or the whole string if no slash is found. File name extensions are not case-sensitive. The name can contain multiple extensions separated by period (.) to indicate type, encoding, or language. For example, the URI `a/b/filename.jp.txt.zip` represents a Japanese language, text/plain type, zip encoded file.

Syntax

```
cinfo *cinfo_find(char *uri);
```

Returns

A pointer to a newly allocated `cinfo` structure if the find succeeds, or NULL if the find fails.

The `cinfo` structure that is allocated and returned contains pointers to the `content-type`, `content-encoding`, and `content-language`, if found. Each structure points to static data in the types database, or NULL if not found. Do not free these pointers. You should free the `cinfo` structure after using it.

Parameters

`char *uri` is a Universal Resource Identifier (URI) or local file name. Multiple file name extensions should be separated by periods (.).

condvar_init

The `condvar_init` function is a critical-section function that initializes and returns a new condition variable associated with a specified critical-section variable. You can use the condition variable to prevent interference between two threads of execution.

Syntax

```
CONDVAR condvar_init(CRITICAL id);
```

Returns

A newly allocated condition variable (`CONDVAR`).

Parameters

`CRITICAL id` is a critical-section variable.

See Also

"condvar_notify" on page 65, "condvar_terminate" on page 66, "condvar_wait" on page 66, "crit_init" on page 68, "crit_enter" on page 67, "crit_exit" on page 67, "crit_terminate" on page 68

condvar_notify

The `condvar_notify` function is a critical-section function that awakens any threads that are blocked on the given critical-section variable. Use this function to awaken threads of execution of a given critical section. First, use `crit_enter` to gain ownership of the critical section. Then use the returned critical-section variable to call `condvar_notify` to awaken the threads. Finally, when `condvar_notify` returns, call `crit_exit` to surrender ownership of the critical section.

Syntax

```
void condvar_notify(CONDVAR cv);
```

Returns

`void`

Parameters

`CONDVAR cv` is a condition variable.

See Also

"condvar_init" on page 65, "condvar_terminate" on page 66, "condvar_wait" on page 66, "crit_init" on page 68, "crit_enter" on page 67, "crit_exit" on page 67, "crit_terminate" on page 68

condvar_terminate

The `condvar_terminate` function is a critical-section function that frees a condition variable. Use this function to free a previously allocated condition variable.

Caution – Terminating a condition variable that is in use can lead to unpredictable results.

Syntax

```
void condvar_terminate(CONDVAR cv);
```

Returns

`void`

Parameters

`CONDVAR cv` is a condition variable.

See Also

"condvar_init" on page 65, "condvar_notify" on page 65, "condvar_wait" on page 66, "crit_init" on page 68, "crit_enter" on page 67, "crit_exit" on page 67, "crit_terminate" on page 68

condvar_wait

The `condvar_wait` function is a critical-section function that blocks on a given condition variable. Use this function to wait for a critical section (specified by a condition variable argument) to become available. The calling thread is blocked until another thread calls `condvar_notify` with the same condition variable argument. The caller must have entered the critical section associated with this condition variable before calling `condvar_wait`.

Syntax

```
void condvar_wait(CONDVAR cv);
```

Returns

`void`

Parameters

`CONDVAR cv` is a condition variable.

See Also

"condvar_init" on page 65, "condvar_terminate" on page 66, "condvar_notify" on page 65, "crit_init" on page 68, "crit_enter" on page 67, "crit_exit" on page 67, "crit_terminate" on page 68

crit_enter

The `crit_enter` function is a critical-section function that attempts to enter a critical section. Use this function to gain ownership of a critical section. If another thread already owns the section, the calling thread is blocked until the first thread surrenders ownership by calling `crit_exit`.

Syntax

```
void crit_enter(CRITICAL crvar);
```

Returns

`void`

Parameters

`CRITICAL crvar` is a critical-section variable.

See Also

"crit_init" on page 68, "crit_exit" on page 67, "crit_terminate" on page 68

crit_exit

The `crit_exit` function is a critical-section function that surrenders ownership of a critical section. Use this function to surrender ownership of a critical section. If another thread is blocked waiting for the section, the block is removed and the waiting thread is given ownership of the section.

Syntax

```
void crit_exit(CRITICAL crvar);
```

Returns

`void`

Parameters

`CRITICAL crvar` is a critical-section variable.

See Also

"crit_init" on page 68, "crit_enter" on page 67, "crit_terminate" on page 68

crit_init

The `crit_init` function is a critical-section function that creates and returns a new critical-section variable (a variable of type `CRITICAL`). Use this function to obtain a new instance of a variable of type `CRITICAL` (a critical-section variable). Use this variable to prevent interference between two threads of execution. At the time of creating this variable, no thread owns the critical section.

Caution – Threads must not own or be waiting for the critical section when `crit_terminate` is called.

Syntax

```
CRITICAL crit_init(void);
```

Returns

A newly allocated critical-section variable (`CRITICAL`).

Parameters

none

See Also

"crit_enter" on page 67, "crit_exit" on page 67, "crit_terminate" on page 68

crit_terminate

The `crit_terminate` function is a critical-section function that removes a previously allocated critical-section variable (a variable of type `CRITICAL`). Use this function to release a critical-section variable previously obtained by a call to `crit_init`.

Syntax

```
void crit_terminate(CRITICAL crvar);
```

Returns

`void`

Parameters

`CRITICAL crvar` is a critical-section variable.

See Also

"crit_init" on page 68, "crit_enter" on page 67, "crit_exit" on page 67

D

daemon_atrestart

The `daemon_atrestart` function lets you register a callback function named `fn` to be used when the server terminates. Use this function when you need a callback function to deallocate resources allocated by an initialization function. The `daemon_atrestart` function is a generalization of the `magnus_atrestart` function.

The `magnus.conf` directives `TerminateTimeout` and `ChildRestartCallback` also affect the callback of NSAPI functions.

Syntax

```
void daemon_atrestart(void (*fn)(void *), void *data);
```

Returns

`void`

Parameters

`void (* fn) (void *)` is the callback function.

`void *data` is the parameter passed to the callback function when the server is restarted.

Example

```
/* Register the log_close function, passing it NULL */
/* to close *a log file when the server is */
/* restarted or shutdown. */
daemon_atrestart(log_close, NULL);
NSAPI_PUBLIC void log_close(void *parameter)
                  {system_fclose(global_logfd);}
```

F

filebuf_buf2sd

The `filebuf_buf2sd` function sends a file buffer to a socket (descriptor) and returns the number of bytes sent.

Use this function to send the contents of an entire file to the client.

Syntax

```
int filebuf_buf2sd(filebuf *buf, SYS_NETFD sd);
```

Returns

The number of bytes sent to the socket if successful, or the constant `IO_ERROR` if the file buffer cannot be sent.

Parameters

`filebuf *buf` is the file buffer that must already have been opened.

`SYS_NETFD sd` is the platform-independent socket descriptor. Normally this is obtained from the `csd` (client socket descriptor) field of the `sn` (`session`) structure.

Example

```
if (filebuf_buf2sd(buf, sn->csd) == IO_ERROR)
   return(REQ_EXIT);
```

See Also

"filebuf_close" on page 71, "filebuf_open" on page 72, "filebuf_open_nostat" on page 72, "filebuf_getc" on page 71

filebuf_close

The `filebuf_close` function deallocates a file buffer and closes its associated file.

Generally, use `filebuf_open` first to open a file buffer, and then `filebuf_getc` to access the information in the file. After you have finished using the file buffer, use `filebuf_close` to close it.

Syntax

```
void filebuf_close(filebuf *buf);
```

Returns

```
void
```

Parameters

`filebuf *buf` is the file buffer previously opened with `filebuf_open`.

Example

```
filebuf_close(buf);
```

See Also

"filebuf_open" on page 72, "filebuf_open_nostat" on page 72, "filebuf_buf2sd" on page 70, "filebuf_getc" on page 71

filebuf_getc

The `filebuf_getc` function retrieves a character from the current file position and returns it as an integer. It then increments the current file position.

Use `filebuf_getc` to sequentially read characters from a buffered file.

Syntax

```
filebuf_getc(filebuf b);
```

Returns

An integer containing the character retrieved, or the constant `IO_EOF` or `IO_ERROR` upon an end of file or error.

Parameters

`filebuf b` is the name of the file buffer.

See Also

"filebuf_close" on page 71, "filebuf_buf2sd" on page 70, "filebuf_open" on page 72, "filter_create" on page 73

filebuf_open

The `filebuf_open` function opens a new file buffer for a previously opened file. It returns a new buffer structure. Buffered files provide more efficient file access by guaranteeing the use of buffered file I/O in environments where it is not supported by the operating system.

Syntax

```
filebuf *filebuf_open(SYS_FILE fd, int sz);
```

Returns

A pointer to a new buffer structure to hold the data if successful, or NULL if no buffer can be opened.

Parameters

`SYS_FILE fd` is the platform-independent file descriptor of the file which has already been opened.

`int sz` is the size, in bytes, to be used for the buffer.

Example

```
filebuf *buf = filebuf_open(fd, FILE_BUFFERSIZE);
if (!buf)
{
   system_fclose(fd);
}
```

See Also

"filebuf_getc" on page 71, "filebuf_buf2sd" on page 70, "filebuf_close" on page 71, "filebuf_open_nostat" on page 72

filebuf_open_nostat

The `filebuf_open_nostat` function opens a new file buffer for a previously opened file. It returns a new buffer structure. Buffered files provide more efficient file access by guaranteeing the use of buffered file I/O in environments where it is not supported by the operating system.

This function is the same as `filebuf_open`, but is more efficient, since it does not need to call the `request_stat_path` function. It requires that the stat information be passed in.

Syntax

```
filebuf* filebuf_open_nostat(SYS_FILE fd, int sz, struct stat *finfo);
```

Returns

A pointer to a new buffer structure to hold the data if successful, or NULL if no buffer can be opened.

Parameters

`SYS_FILE fd` is the platform-independent file descriptor of the file that has already been opened.

`int sz` is the size, in bytes, to be used for the buffer.

`struct stat *finfo` is the file information of the file. Before calling the `filebuf_open_nostat` function, you must call the `request_stat_path` function to retrieve the file information.

Example

```
filebuf *buf = filebuf_open_nostat(fd, FILE_BUFFERSIZE, &finfo);
if (!buf)
{
  system_fclose(fd);
}
```

See Also

"filebuf_close" on page 71, "filebuf_open" on page 72, "filebuf_getc" on page 71, "filebuf_buf2sd" on page 70

filter_create

The `filter_create` function defines a new filter.

The `name` parameter specifies a unique name for the filter. If a filter with the specified name already exists, it will be replaced.

Names beginning with `magnus-` or `server-` are reserved by the server.

The `order` parameter indicates the position of the filter in the filter stack by specifying what class of functionality the filter implements.

The following table describes parameters allowed constants and their associated meanings for the `filter_create` function. The left column lists the name of the constant, the middle column describes the functionality the filter implements, and the right column lists the position the filter occupies in the filter stack.

TABLE 5–1 filter-create constants

Constant	Functionality Filter Implements	Position in Filter Stack
`FILTER_CONTENT_TRANSLATION`	Translates content from one form to another (for example, XSLT)	Top
`FILTER_CONTENT_CODING`	Encodes content (for example, HTTP gzip compression)	Middle
`FILTER_TRANSFER_CODING`	Encodes entity bodies for transmission (for example, HTTP chunking)	Bottom

The `methods` parameter specifies a pointer to a `FilterMethods` structure. Before calling `filter_create`, you must initialize the `FilterMethods` structure using the `FILTER_METHODS_INITIALIZER` macro, and then assign function pointers to the individual `FilterMethods` members (for example, `insert`, `read`, `write`, and so on) that correspond to the filter methods the filter supports.

`filter_create` returns `const Filter *`, a pointer to an opaque representation of the filter. This value can be passed to `filter_insert` to insert the filter in a particular filter stack.

Syntax

```
const Filter *filter_create(const char *name, int order,
                            const FilterMethods *methods);
```

Returns

The `const Filter *` that identifies the filter or NULL if an error occurs.

Parameters

`const char *name` is the name of the filter.

`int order` is one of the order constants above.

`const FilterMethods *methods` contains pointers to the filter methods the filter supports.

Example

```
FilterMethods methods = FILTER_METHODS_INITIALIZER;
const Filter *filter;
/* This filter will only support the "read" filter method */
methods.read = my_input_filter_read;
/* Create the filter */
filter = filter_create("my-input-filter", FILTER_CONTENT_TRANSLATION,
                       &methods);
```

See Also

"filter_insert" on page 75, "insert" on page 81, "flush" on page 78, "read" on page 112, "sendfile" on page 117, "write" on page 158, "writev" on page 159, "FilterMethods" on page 167

filter_find

The `filter_find` function finds the filter with the specified name.

Syntax

```
const Filter *filter_find(const char *name);
```

Returns

The `const Filter *` that identifies the filter, or NULL if the specified filter does not exist.

Parameters

`const char *name` is the name of the filter of interest.

filter_insert

The `filter_insert` function inserts a filter into a filter stack, creating a new filter layer and installing the filter at that layer. The filter layer's position in the stack is determined by the order value specified when `filter_create` was called, and any explicit ordering configured by `init-filter-order`. If a filter layer with the same order value already exists in the stack, the new layer is inserted above that layer.

Parameters are passed to the filter using the `pb` and `data` parameters. The semantics of the data parameter are defined by individual filters. However, all filters must be able to handle a data parameter of NULL.

Note – When possible, plug-in developers should avoid calling `filter_insert` directly, and instead use the `insert-filter` SAF.

Syntax

```
int filter_insert(SYS_NETFD sd, pblock *pb, Session *sn, Request *rq,
                  void *data, const Filter *filter);
```

Returns

`REQ_PROCEED` if the specified filter was inserted successfully, or `REQ_NOACTION` if the specified filter was not inserted because it was not required. Any other return value indicates an error.

Parameters

`SYS_NETFD sd` is NULL (reserved for future use).

`pblock *pb` is a set of parameters to pass to the specified filter\qs init method.

`Session *sn` is the Session.

`Request *rq` is the Request.

`void *data` is filter-defined private data.

`const Filter *filter` is the filter to insert.

See Also

"filter_create" on page 73

filter_layer

The `filter_layer` function returns the layer in a filter stack that corresponds to the specified filter.

Syntax

```
FilterLayer *filter_layer(SYS_NETFD sd, const Filter *filter);
```

Returns

The topmost `FilterLayer *` associated with the specified filter, or NULL if the specified filter is not part of the specified filter stack.

Parameters

`SYS_NETFD sd` is the filter stack to inspect.

`const Filter *filter` is the filter of interest.

filter_name

The `filter_name` function returns the name of the specified filter. The caller should not free the returned string.

Syntax

```
const char *filter_name(const Filter *filter);
```

Returns

The name of the specified filter, or NULL if an error occurred.

Parameters

`const Filter *filter` is the filter of interest.

filter_remove

The `filter_remove` function removes the specified filter from the specified filter stack, destroying a filter layer. If the specified filter was inserted into the filter stack multiple times, only the topmost filter layer of the filter is destroyed.

Note – When possible, plug-in developers should avoid calling `filter_remove` directly, and instead use the remove-filter SAF (applicable in `Input`-, `Output`-, `Service`-, and `Error`-class directives).

Syntax

```
int filter_remove(SYS_NETFD sd, const Filter *filter);
```

Returns

`REQ_PROCEED` if the specified filter was removed successfully or `REQ_NOACTION` if the specified filter was not part of the filter stack. Any other return value indicates an error.

Parameters

`SYS_NETFD sd` is the filter stack, `sn->csd`.

`const Filter *filter` is the filter to remove.

flush

The `flush` filter method is called when buffered data should be sent. Filters that buffer outgoing data should implement the `flush` filter method.

Upon receiving control, a `flush` implementation must write any buffered data to the filter layer immediately below it. Before returning success, a `flush` implementation must successfully call the `net_flush` function:

```
net_flush(layer->lower).
```

Syntax

```
int flush(FilterLayer *layer);
```

Returns

`0` on success or `-1` if an error occurs.

Parameters

`FilterLayer *layer` is the filter layer the filter is installed in.

Example

```
int myfilter_flush(FilterLayer *layer)
{
    MyFilterContext context = (MyFilterContext *)layer->context->data;
    if (context->buf.count) {
       int rv;
       rv = net_write(layer->lower, context->buf.data, context->buf.count);
       if (rv != context->buf.count)
           return -1; /* failed to flush data */
       context->buf.count = 0;
    }
    return net_flush(layer->lower);
}
```

See Also

"net_flush" on page 84, "filter_create" on page 73

FREE

The `FREE` macro is a platform-independent substitute for the C library routine `free`. It deallocates the space previously allocated by `MALLOC`, `CALLOC`, or `STRDUP` from the request's memory pool.

Note – Calling FREE for a block that was allocated with PERM_MALLOC, PERM_CALLOC, or PERM_STRDUP will not work.

Syntax

```
FREE(void *ptr);
```

Returns

```
void
```

Parameters

void *ptr is a (void *) pointer to a block of memory. If the pointer is not the one created by MALLOC, CALLOC, or STRDUP, the behavior is undefined.

Example

```
char *name;
name = (char *) MALLOC(256);
...
...
FREE(name);
```

See Also

"CALLOC" on page 63, "MALLOC" on page 83, "REALLOC" on page 113, "STRDUP" on page 121, "PERM_FREE" on page 104

func_exec

The func_exec function executes the function named by the fn entry in a specified pblock. If the function name is not found, it logs the error and returns REQ_ABORTED.

You can use this function to execute a built-in Server Application Function (SAF) by identifying it in the pblock.

Syntax

```
int func_exec(pblock *pb, Session *sn, Request *rq);
```

Returns

The value returned by the executed function, or the constant if successful. REQ_ABORTED, if no function is executed.

Parameters

`pblock pb` is the `pblock` containing the function name (`fn`) and parameters.

`Session *sn` is the Session.

`Request *rq` is the Request.

The `Session` and `Request` parameters are the same as the ones passed into your SAF.

See Also

"log_error" on page 82

func_find

The `func_find` function returns a pointer to the function specified by `name`. If the function does not exist, it returns NULL.

Syntax

```
FuncPtr func_find(char *name);
```

Returns

A pointer to the chosen function, suitable for de-referencing, or NULL if the function is not found.

Parameters

`char *name` is the name of the function.

Example

```
/* this block of code does the same thing as func_exec */
char *afunc = pblock_findval("afunction", pb);
FuncPtr afnptr = func_find(afunc);
 if (afnptr)
     return (afnptr)(pb, sn, rq);
```

See Also

"func_exec" on page 79

func_insert

The `func_insert` function dynamically inserts a named function into the server's table of functions. This function should only be called during the `Init` stage.

Syntax

```
FuncStruct *func_insert(char *name, FuncPtr fn);
```

Returns

The `FuncStruct` structure that identifies the newly inserted function. The caller should not modify the contents of the `FuncStruct` structure.

Parameters

`char *name` is the name of the function.

`FuncPtr fn` is the pointer to the function.

Example

```
func_insert("my-service-saf", &my_service_saf);
```

See Also

"func_exec" on page 79, "func_find" on page 80

I

insert

The `insert` filter method is called when a filter is inserted into a filter stack by the `filter_insert` function or `insert-filter` SAF.

Syntax

```
int insert(FilterLayer *layer, pblock *pb);
```

Returns

`REQ_PROCEED` if the filter should be inserted into the filter stack, `REQ_NOACTION` if the filter should not be inserted because it is not required, or `REQ_ABORTED` if the filter should not be inserted because of an error.

Parameters

`FilterLayer *layer` is the filter layer at which the filter is being inserted.

`pblock *pb` is the set of parameters passed to `filter_insert` or specified by the `fn="insert-filter"` directive.

Example

```
int myfilter_insert(FilterLayer *layer, pblock *pb)
{
        if (pblock_findval("dont-insert-filter", pb))
        return REQ_NOACTION;
        return REQ_PROCEED;
}
...

FilterMethods myfilter_methods = FILTER_METHODS_INITIALIZER;
const Filter *myfilter;

myfilter_methods.insert = &myfilter_insert;
myfilter = filter_create("myfilter", &myfilter_methods);
...
```

See Also

"filter_insert" on page 75, "filter_create" on page 73

L

log_error

The `log_error` function creates an entry in an error log, recording the date, the severity, and a description of the error.

Syntax

```
int log_error(int degree, char *func, Session *sn, Request *rq, char *fmt, ...);
```

Returns

`0` if the log entry is created, or `-1` if the log entry is not created.

Parameters

`int degree` specifies the severity of the error. It must be one of the following constants:

- `LOG_VERBOSE` - debug message
- `LOG_VERBOSE` - debug message
- `LOG_INFORM` - information message
- `LOG_WARN`- warning
- `LOG_FAILURE` - operation failed

- LOG_MISCONFIG- misconfiguration
- LOG_SECURITY - authentication or authorization failure
- LOG_CATASTROPHE - nonrecoverable server error

char *func is the name of the function where the error has occurred.

Session *sn is the Session.

Request *rq is the Request.

char *fmt specifies the format for the printf function that delivers the message.

Example

```
log_error(LOG_WARN, "send-file", sn, rq, "error opening buffer from %s (%s)"),
          path, system_errmsg(fd));
```

See Also

"func_exec" on page 79

M

MALLOC

The MALLOC macro is a platform-independent substitute for the C library routine malloc. It allocates size bytes from the requests's memory pool. The memory can be explicitly freed by a call to FREE. If the memory is not explicitly freed, it is automatically freed after processing of the current request has been completed. If pooled memory has been disabled in the configuration file (with the pool-initbuilt-in SAF), PERM_MALLOC and MALLOC both obtain their memory from the system heap. However, since memory allocated by MALLOC is automatically freed, it should not be shared between threads.

If pooled memory has been disabled in the configuration file (with the pool-init built-in SAF), PERM_MALLOC and MALLOC both obtain their memory from the system heap.

Syntax

```
void *MALLOC(int size)
```

Returns

A void pointer to a block of memory.

Parameters

int size is the number of bytes to allocate.

Example

```
/* Allocate 256 bytes for a name */
char *name;
name = (char *) MALLOC(256);
```

See Also

"FREE" on page 78, "CALLOC" on page 63, "REALLOC" on page 113, "STRDUP" on page 121, "PERM_MALLOC" on page 104

N

net_flush

The `net_flush` function flushes any buffered data. If you require that data be sent immediately, call `net_flush` after calling the network output functions such as `net_write` or `net_sendfile`.

Syntax

```
int net_flush(SYS_NETFD sd);
```

Returns

`0` on success, or a negative value if an error occurs.

Parameters

`SYS_NETFD sd` is the socket to flush.

Example

```
net_write(sn->csd, "Please wait... ", 15);
net_flush(sn->csd);
/* Perform some time-intensive operation */
...
net_write(sn->csd, "Thank you.\n", 11);
```

See Also

"net_write" on page 87, "net_sendfile" on page 86

net_ip2host

The `net_ip2host` function transforms a textual IP address into a fully-qualified domain name and returns it.

Note – This function works only if the `DNS` directive is enabled in the `magnus.conf` file.

Syntax

```
char *net_ip2host(char *ip, int verify);
```

Returns

A new string containing the fully-qualified domain name if the transformation is accomplished, or NULL if the transformation is not accomplished.

Parameters

`char *ip` is the IP address as a character string in dotted-decimal notation: `nnn.nnn.nnn.nnn`

`int verify`, if nonzero, specifies that the function should verify the fully-qualified domain name. Though this requires an extra query, you should use it when checking the access control.

net_read

The `net_read` function reads bytes from a specified socket into a specified buffer. The function waits to receive data from the socket until either at least one byte is available in the socket or the specified time has elapsed.

Syntax

```
int net_read (SYS_NETFD sd, char *buf, int sz, int timeout);
```

Returns

The number of bytes read, which will not exceed the maximum size, `sz`. A negative value is returned if an error has occurred, in which case `errno` is set to the constant `ETIMEDOUT` if the operation did not complete before `timeout` seconds elapsed.

Parameters

`SYS_NETFD sd` is the platform-independent socket descriptor.

`char *buf` is the buffer to receive the bytes.

`int sz` is the maximum number of bytes to read.

`int timeout` is the number of seconds to allow for the read operation before returning. The purpose of `timeout` is not to return because not enough bytes were read in the given time, but to limit the amount of time devoted to waiting until some data arrives.

See Also

"net_write" on page 87

net_sendfile

The `net_sendfile` function sends the contents of a specified file to a specified a socket. Either the whole file or a fraction might be sent, and the contents of the file might optionally be preceded and/or followed by caller-specified data.

Parameters are passed to `net_sendfile` in the `sendfiledata` structure. Before invoking `net_sendfile`, the caller must initialize every `sendfiledata` structure member.

Syntax

```
int net_sendfile(SYS_NETFD sd, const sendfiledata *sfd);
```

Returns

A positive number indicating the number of bytes successfully written, including the headers, file contents, and trailers. A negative value indicating an error.

Parameters

`SYS_NETFD sd` is the socket to write to.

`const sendfiledata *sfd` identifies the data to send.

Example

The following `Service` SAF sends a file bracketed by the strings "begin" and "end."

```
#include <string.h>
#include "nsapi.h"

NSAPI_PUBLIC int service_net_sendfile(pblock *pb, Session *sn, Request *rq)
{
    char *path;
    SYS_FILE fd;
    struct sendfiledata sfd;
    int rv;

    path = pblock_findval("path", rq->vars);
    fd = system_fopenRO(path);
    if (!fd) {
        log_error(LOG_MISCONFIG, "service-net-sendfile", sn, rq,
                  "Error opening %s (%s)", path, system_errmsg());
```

```
        return REQ_ABORTED;
    }

    sfd.fd = fd;                    /* file to send */
    sfd.offset = 0;                 /* start sending from the beginning */
    sfd.len = 0;                    /* send the whole file */
    sfd.header = "begin";           /* header data to send before the file */
    sfd.hlen = strlen(sfd.header);  /* length of header data */
    sfd.trailer = "end";            /* trailer data to send after the file */
    sfd.tlen = strlen(sfd.trailer); /* length of trailer data */

    /* send the headers, file, and trailers to the client */
    rv = net_sendfile(sn->csd, &sfd);

    system_fclose(fd);

    if (rv < 0) {
        log_error(LOG_INFORM, "service-net-sendfile", sn, rq,
                  "Error sending %s (%s)", path,
                  system_errmsg());
        return REQ_ABORTED;
    }

    return REQ_PROCEED;
}
```

See Also

"net_flush" on page 84

net_write

The `net_write` function writes a specified number of bytes to a specified socket from a specified buffer.

Syntax

```
int net_write(SYS_NETFD sd, char *buf, int sz);
```

Returns

The number of bytes written, which may be less than the requested size if an error occurs.

Parameters

`SYS_NETFD sd` is the platform-independent socket descriptor.

`char *buf` is the buffer containing the bytes.

`int sz` is the number of bytes to write.

Example

```
if (net_write(sn->csd, FIRSTMSG, strlen(FIRSTMSG)) == IO_ERROR)
    return REQ_EXIT;
```

See Also

"net_read" on page 85

netbuf_buf2sd

The `netbuf_buf2sd` function sends a buffer to a socket. You can use this function to send data from IPC pipes to the client.

Syntax

```
int netbuf_buf2sd(netbuf *buf, SYS_NETFD sd, int len);
```

Returns

The number of bytes transferred to the socket, if successful, or the constant `IO_ERROR` if unsuccessful.

Parameters

`netbuf *buf` is the buffer to send.

`SYS_NETFD sd` is the platform-independent identifier of the socket.

`int len` is the length of the buffer.

See Also

"netbuf_close" on page 88, "netbuf_getc" on page 90, "netbuf_getbytes" on page 89, "netbuf_grab" on page 91, "netbuf_open" on page 91

netbuf_close

The `netbuf_close` function deallocates a network buffer and closes its associated files. Use this function when you need to deallocate the network buffer and close the socket.

You should never close the `netbuf` parameter in a `session` structure.

Syntax

```
void netbuf_close(netbuf *buf);
```

Returns

```
void
```

Parameters

`netbuf *buf` is the buffer to close.

See Also

"netbuf_buf2sd" on page 88, "netbuf_getc" on page 90, "netbuf_getbytes" on page 89, "netbuf_grab" on page 91, "netbuf_open" on page 91

netbuf_getbytes

The `netbuf_getbytes` function reads bytes from a network buffer into a caller-supplied buffer. If the network buffer is empty, the function waits to receive data from the network buffer's socket until either at least one byte is available from the socket or the network buffer's timeout has elapsed.

Syntax

```
int netbuf_getbytes(netbuf *buf, char *buffer, int sz);
```

Returns

The number of bytes placed into buffer (between 1 and sz) if the operation is successful, the constant `NETBUF_EOF` on end of file, or the constant `NETBUF_ERROR` if an error occurred.

Parameters

`netbuf *buf` is the buffer from which to retrieve bytes.

`char *buffer` is the caller-supplied buffer that receives the bytes.

`int sz` is the maximum number of bytes to read.

Example

```
int cl = 0;

* Read the entire request body */
```

```
for (;;) {
     char mybuf[1024];
     int rv;

     rv = netbuf_getbytes(sn->inbuf, mybuf, sizeof(mybuf));
     if (rv == NETBUF_EOF) {
         log_error(LOG_INFORM, "mysaf", sn, rq,
                   "Received %d byte(s)",
                   cl);
         break;
     }
     if (rv == NETBUF_ERROR) {
         log_error(LOG_FAILURE, "mysaf", sn, rq,
                   "Error reading request body (%s)",
                   cl, system_errmsg());
         break;     }

     cl += rv;
}
```

See Also

"netbuf_buf2sd" on page 88, "netbuf_close" on page 88, "netbuf_getc" on page 90, "netbuf_grab" on page 91, "netbuf_open" on page 91

netbuf_getc

The `netbuf_getc` function retrieves a character from the cursor position of the network buffer specified by `b`.

Note – Because the constant `IO_EOF` has a value of 0, `netbuf_getc` cannot be used to read data that may contain a null character. To read binary data, use "netbuf_getbytes" on page 89 or "netbuf_grab" on page 91.

Syntax

```
netbuf_getc(netbuf b);
```

Returns

The integer representing the character if a character is retrieved, or the constant `IO_EOF` or `IO_ERROR` for end of file or error.

Parameters

`netbuf b` is the buffer from which to retrieve one character.

See Also

"netbuf_buf2sd" on page 88, "netbuf_close" on page 88, "netbuf_getbytes" on page 89, "netbuf_grab" on page 91, "netbuf_open" on page 91

netbuf_grab

The `netbuf_grab` function reads `sz` number of bytes from the network buffer's (`buf`) socket into the network buffer. If the buffer is not large enough it is resized. The data can be retrieved from `buf->inbuf` on success.

This function is used by the function `netbuf_buf2sd`.

Syntax

```
int netbuf_grab(netbuf *buf, int sz);
```

Returns

The number of bytes actually read (between `1` and `sz`) if the operation is successful, or the constant `IO_EOF` or `IO_ERROR` for end of file or error.

Parameters

`netbuf *buf` is the buffer to read into.

`int sz` is the number of bytes to read.

See Also

"netbuf_buf2sd" on page 88, "netbuf_close" on page 88, "netbuf_getbytes" on page 89, "netbuf_getc" on page 90, "netbuf_open" on page 91

netbuf_open

The `netbuf_open` function opens a new network buffer and returns it. You can use `netbuf_open` to create a `netbuf` structure and start using buffered I/O on a socket.

Syntax

```
netbuf* netbuf_open(SYS_NETFD sd, int sz);
```

Returns

A pointer to a new `netbuf` structure (network buffer).

Parameters

`SYS_NETFD sd` is the platform-independent identifier of the socket.

`int sz` is the number of characters to allocate for the network buffer.

See Also

"netbuf_buf2sd" on page 88, "netbuf_close" on page 88, "netbuf_getc" on page 90, "netbuf_getbytes" on page 89, "netbuf_grab" on page 91

nsapi_module_init

Define the `nsapi_module_init` function, which is a module initialization entry point that enables a plug-in to create filters when it is loaded. When an NSAPI module contains an `nsapi_module_init` function, the server will call that function immediately after loading the module. The `nsapi_module_init` presents the same interface as an `Init` SAF, and it must follow the same rules.

The `nsapi_module_init` function is used to register SAFs with `func_insert`, create filters with `filter_create`, register virtual server initialization/destruction callbacks with `vs_register_cb`, and perform other initialization tasks.

Syntax

```
int nsapi_module_init(pblock *pb, Session *sn, Request *rq);
```

Returns

`REQ_PROCEED` on success, or `REQ_ABORTED` on error.

Parameters

`pblock *pb` is a set of parameters specified by the `fn="load-modules"` directive.

`Session *sn` (the Session) is NULL.

`Request *rq` (the Request) is NULL.

See Also

"filter_create" on page 73, "func_insert" on page 80, "vs_register_cb" on page 156

NSAPI_RUNTIME_VERSION

The `NSAPI_RUNTIME_VERSION` macro defines the NSAPI version available at runtime. This is the same as the highest NSAPI version supported by the server the plug-in is running in. The NSAPI version is encoded as in `USE_NSAPI_VERSION`.

The value returned by the NSAPI_RUNTIME_VERSION macro is valid only in iPlanet™ Web Server 6.0, Netscape Enterprise Server 6.0, Sun ONE Web Server 6.1 and Sun Java System Web Server 7.0. That is, the server must support NSAPI 3.1 for this macro to return a valid value. Additionally, to use NSAPI_RUNTIME_VERSION, you must compile against an nsapi.h header file that supports NSAPI 3.2 or higher.

You must not attempt to set the value of the NSAPI_RUNTIME_VERSION macro directly. Instead, see the USE_NSAPI_VERSION macro.

Syntax

```
int NSAPI_RUNTIME_VERSION
```

Example

```
NSAPI_PUBLIC int log_nsapi_runtime_version(pblock *pb, Session *sn, Request *rq)
{
    log_error(LOG_INFORM, "log-nsapi-runtime-version", sn, rq,
                "Server supports NSAPI version %d.%d\n",
                NSAPI_RUNTIME_VERSION / 100,
                NSAPI_RUNTIME_VERSION % 100);
     return REQ_PROCEED;
}
```

See Also

"filter_create" on page 73, "func_insert" on page 80, "vs_register_cb" on page 156

NSAPI_VERSION

The NSAPI_VERSION macro defines the NSAPI version used at compile time. This value is determined by the value of the USE_NSAPI_VERSION macro or by the highest NSAPI version supported by the nsapi.h header the plug-in was compiled against. The NSAPI version is encoded as in USE_NSAPI_VERSION.

You must not attempt to set the value of the NSAPI_VERSION macro directly. Instead, see the USE_NSAPI_VERSION macro.

Syntax

```
int NSAPI_VERSION
```

Example

```
NSAPI_PUBLIC int log_nsapi_compile_time_version(pblock *pb, Session *sn, Request *rq)
{
    log_error(LOG_INFORM, "log-nsapi-compile-time-version", sn, rq,
            "Plugin compiled against NSAPI version %d.%d\n",
            NSAPI_VERSION / 100,
            NSAPI_VERSION % 100);
    return REQ_PROCEED;
}
```

See Also

"NSAPI_RUNTIME_VERSION" on page 92, "USE_NSAPI_VERSION" on page 136

P

param_create

The `param_create` function creates a `pb_param` structure containing a specified name and value. The name and value are copied. Use this function to prepare a `pb_param` structure to be used in calls to `pblock` routines such as `pblock_pinsert`.

Syntax

```
pb_param *param_create(char *name, char *value);
```

Returns

A pointer to a new `pb_param` structure.

Parameters

`char *name` is the string containing the name.

`char *value` is the string containing the value.

Example

```
pb_param *newpp = param_create("content-type","text/plain");
pblock_pinsert(newpp, rq->srvhdrs);
```

See Also

"param_free" on page 95, "pblock_pinsert" on page 101, "pblock_remove" on page 102

param_free

The `param_free` function frees the `pb_param` structure specified by `pp` and its associated structures. Use the `param_free` function to dispose a `pb_param` after removing it from a `pblock` with `pblock_remove`.

Syntax

```
int param_free(pb_param *pp);
```

Returns

`1` if the parameter is freed or `0` if the parameter is NULL.

Parameters

`pb_param *pp` is the name-value pair stored in a pblock.

Example

```
if (param_free(pblock_remove("content-type", rq-srvhdrs)))
return; /* we removed it */
```

See Also

"param_create" on page 94, "pblock_pinsert" on page 101, "pblock_remove" on page 102

pblock_copy

The `pblock_copy` function copies the entries of the source `pblock` and adds them into the destination `pblock`. Any previous entries in the destination `pblock` are left intact.

Syntax

```
void pblock_copy(pblock *src, pblock *dst);
```

Returns

`void`

Parameters

`pblock *src` is the source pblock.

`pblock *dst` is the destination pblock.

Names and values are newly allocated so that the original `pblock` may be freed, or the new `pblock` changed without affecting the original `pblock`.

See Also

"pblock_create" on page 96, "pblock_dup" on page 96, "pblock_free" on page 98, "pblock_find" on page 97, "pblock_findval" on page 97, "pblock_remove" on page 102, "pblock_nvinsert" on page 99

pblock_create

The `pblock_create` function creates a new `pblock`. The `pblock` maintains an internal hash table for fast name-value pair lookups. Because the `pblock` is allocated from the request's memory pool, it should not be shared between threads.

Syntax

```
pblock *pblock_create(int n);
```

Returns

A pointer to a newly allocated `pblock`.

Parameters

`int n` is the size of the hash table (number of name-value pairs) for the `pblock`.

See Also

"pblock_copy" on page 95, "pblock_dup" on page 96, "pblock_find" on page 97, "pblock_findval" on page 97, "pblock_free" on page 98, "pblock_nvinsert" on page 99, "pblock_remove" on page 102

pblock_dup

The `pblock_dup` function duplicates a `pblock`. It is equivalent to a sequence of `pblock_create` and `pblock_copy`.

Syntax

```
pblock *pblock_dup(pblock *src);
```

Returns

A pointer to a newly allocated `pblock`.

Parameters

`pblock *src` is the source `pblock`.

See Also

"pblock_create" on page 96, "pblock_find" on page 97, "pblock_findval" on page 97, "pblock_free" on page 98, "pblock_nvinsert" on page 99, "pblock_remove" on page 102

pblock_find

The `pblock_find` macro finds a specified name-value pair entry in a `pblock`, and returns the `pb_param` structure. If you only want the value associated with the name, use the `pblock_findval` function.

Note – Parameter names are case-sensitive. By convention, lowercase names are used for parameters that correspond to HTTP header fields.

Syntax

```
pb_param *pblock_find(char *name, pblock *pb);
```

Returns

A pointer to the `pb_param` structure if found, or NULL if name is not found.

Parameters

`char *name` is the name of a name-value pair.

`pblock *pb` is the `pblock` to be searched.

See Also

"pblock_copy" on page 95, "pblock_dup" on page 96, "pblock_findval" on page 97, "pblock_free" on page 98, "pblock_nvinsert" on page 99, "pblock_remove" on page 102

pblock_findval

The `pblock_findval` function finds the value associated with a specified name in a `pblock`. If you want the `pb_param` structure of the `pblock`, use the `pblock_find` function.

The pointer returned is a pointer into the `pblock`. Do not FREE it. If you want to modify it, do a `STRDUP` and modify the copy.

Note – Parameter names are case-sensitive. By convention, lowercase names are used for parameters that correspond to HTTP header fields.

Syntax

```
char *pblock_findval(char *name, pblock *pb);
```

Returns

A string containing the value associated with the name if found, or NULL if no match is found.

Parameters

`char *name` is the name of a name-value pair.

`pblock *pb` is the `pblock` to be searched.

Example

See "pblock_nvinsert" on page 99.

See Also

"pblock_create" on page 96, "pblock_copy" on page 95, "pblock_find" on page 97, "pblock_free" on page 98, "pblock_nvinsert" on page 99, "pblock_remove" on page 102, "request_header" on page 115

pblock_free

The `pblock_free` function frees a specified `pblock` and any entries inside it. If you want to save a variable in the `pblock`, remove the variable using the function `pblock_remove` and save the resulting pointer.

Syntax

```
void pblock_free(pblock *pb);
```

Returns

`void`

Parameters

`pblock *pb` is the `pblock` to be freed.

See Also

"pblock_copy" on page 95, "pblock_create" on page 96, "pblock_dup" on page 96, "pblock_find" on page 97, "pblock_findval" on page 97, "pblock_nvinsert" on page 99, "pblock_remove" on page 102

pblock_nninsert

The `pblock_nninsert` function creates a new entry with a given name and a numeric value in the specified `pblock`. The numeric value is first converted into a string. The name and value parameters are copied.

Note – Parameter names are case-sensitive. By convention, lowercase names are used for parameters that correspond to HTTP header fields.

Syntax

```
pb_param *pblock_nninsert(char *name, int value, pblock *pb);
```

Returns

A pointer to the new `pb_param` structure.

Parameters

`char *name` is the name of the new entry.

`int value` is the numeric value being inserted into the `pblock`. This parameter must be an integer. If the value you assign is not a number, then instead use the function `pblock_nvinsert` to create the parameter.

`pblock *pb` is the `pblock` into which the insertion occurs.

See Also

"pblock_copy" on page 95, "pblock_create" on page 96, "pblock_find" on page 97, "pblock_free" on page 98, "pblock_nvinsert" on page 99, "pblock_remove" on page 102, "pblock_str2pblock" on page 102

pblock_nvinsert

The `pblock_nvinsert` function creates a new entry with a given name and character value in the specified `pblock`. The name and value parameters are copied.

Note – Parameter names are case-sensitive. By convention, lowercase names are used for parameters that correspond to HTTP header fields.

Syntax

```
pb_param *pblock_nvinsert(char *name, char *value, pblock *pb);
```

Returns

A pointer to the newly allocated `pb_param` structure.

Parameters

`char *name` is the name of the new entry.

`char *value` is the string value of the new entry.

`pblock *pb` is the `pblock` into which the insertion occurs.

Example

```
pblock_nvinsert("content-type", "text/html", rq->srvhdrs);
```

See Also

"pblock_copy" on page 95, "pblock_create" on page 96, "pblock_find" on page 97, "pblock_free" on page 98, "pblock_nninsert" on page 99, "pblock_remove" on page 102, "pblock_str2pblock" on page 102

pblock_pb2env

The `pblock_pb2env` function copies a specified `pblock` into a specified environment. The function creates one new environment entry for each name-value pair in the `pblock`. Use this function to send `pblock` entries to a program that you are going to execute.

Syntax

```
char **pblock_pb2env(pblock *pb, char **env);
```

Returns

A pointer to the environment.

Parameters

`pblock *pb` is the `pblock` to be copied.

`char **env` is the environment into which the `pblock` is to be copied.

See Also

"pblock_copy" on page 95, "pblock_create" on page 96, "pblock_find" on page 97, "pblock_free" on page 98, "pblock_nvinsert" on page 99, "pblock_remove" on page 102, "pblock_str2pblock" on page 102

pblock_pblock2str

The `pblock_pblock2str` function copies all parameters of a specified `pblock` into a specified string. The function allocates additional non-heap space for the string if needed.

Use this function to stream the `pblock` for archival and other purposes.

Syntax

```
char *pblock_pblock2str(pblock *pb, char *str);
```

Returns

The new version of the `str` parameter. If `str` is NULL, this is a new string; otherwise, it is a reallocated string. In either case, it is allocated from the request's memory pool.

Parameters

`pblock *pb` is the `pblock` to be copied.

`char *str` is the string into which the `pblock` is to be copied. It must have been allocated by `MALLOC` or `REALLOC`, not by `PERM_MALLOC` or `PERM_REALLOC` (which allocate from the system heap).

Each name-value pair in the string is separated from its neighbor pair by a space, and is in the format *name*`="`*value*`."`

See Also

"pblock_copy" on page 95, "pblock_create" on page 96, "pblock_find" on page 97, "pblock_free" on page 98, "pblock_nvinsert" on page 99, "pblock_remove" on page 102, "pblock_str2pblock" on page 102

pblock_pinsert

The function `pblock_pinsert` inserts a `pb_param` structure into a `pblock`.

Note – Parameter names are case-sensitive. By convention, lowercase names are used for parameters that correspond to HTTP header fields.

Syntax

```
void pblock_pinsert(pb_param *pp, pblock *pb);
```

Returns

`void`

Parameters

`pb_param *pp` is the `pb_param` structure to insert.

`pblock *pb` is the `pblock`.

See Also

"pblock_copy" on page 95, "pblock_create" on page 96, "pblock_find" on page 97, "pblock_free" on page 98, "pblock_nvinsert" on page 99, "pblock_remove" on page 102, "pblock_str2pblock" on page 102

pblock_remove

The `pblock_remove` macro removes a specified name-value entry from a specified `pblock`. If you use this function, you must call `param_free` to deallocate the memory used by the `pb_param` structure.

Syntax

```
pb_param *pblock_remove(char *name, pblock *pb);
```

Returns

A pointer to the named `pb_param` structure if it is found, or NULL if the named `pb_param` is not found.

Parameters

`char *name` is the name of the `pb_param` to be removed.

`pblock *pb` is the `pblock` from which the name-value entry is to be removed.

See Also

"pblock_copy" on page 95, "pblock_create" on page 96, "pblock_find" on page 97, "pblock_free" on page 98, "pblock_nvinsert" on page 99, "param_create" on page 94, "param_free" on page 95

pblock_str2pblock

The `pblock_str2pblock` function scans a string for parameter pairs, adds them to a `pblock`, and returns the number of parameters added.

Syntax

```
int pblock_str2pblock(char *str, pblock *pb);
```

Returns

The number of parameter pairs added to the `pblock`, if any, or `-1` if an error occurs.

Parameters

`char *str` is the string to be scanned.

The name-value pairs in the string can have the format *name*=*value* or *name*="*value*."

All backslashes (\) must be followed by a literal character. If string values are found with no unescaped = signs (no `name=`), it assumes the names `1`, `2`, `3`, and so on, depending on the string position. For example, if `pblock_str2pblock` finds `"some strings together,"` the function treats the strings as if they appeared in name-value pairs as `1="some" 2="strings" 3="together."`

`pblock *pb` is the `pblock` into which the name-value pairs are stored.

See Also

"pblock_copy" on page 95, "pblock_create" on page 96, "pblock_find" on page 97, "pblock_free" on page 98, "pblock_nvinsert" on page 99, "pblock_remove" on page 102, "pblock_pblock2str" on page 101

PERM_CALLOC

The `PERM_CALLOC` macro is a platform-independent substitute for the C library routine `calloc`. It allocates `size` bytes of memory and initializes the memory to zeros. The memory persists after processing the current request has been completed. The memory should be explicitly freed by a call to `PERM_FREE`.

Syntax

```
void *PERM_CALLOC(int size)
```

Returns

A void pointer to a block of memory.

Parameters

`int size` is the number of bytes to allocate.

Example

```
char **name;
name = (char **) PERM_CALLOC(100 * sizeof(char *));
```

See Also

"CALLOC" on page 63, "PERM_FREE" on page 104, "PERM_STRDUP" on page 106, "PERM_MALLOC" on page 104, "PERM_REALLOC" on page 105

PERM_FREE

The `PERM_FREE` macro is a platform-independent substitute for the C library routine `free`. It deallocates the persistent space previously allocated by `PERM_MALLOC`, `PERM_CALLOC`, or `PERM_STRDUP`.

Note – Calling `PERM_FREE` for a block that was allocated with `MALLOC`, `CALLOC`, or `STRTUP` will not work.

Syntax

```
PERM_FREE(void *ptr);
```

Returns

`void`

Parameters

`void *ptr` is a `(void *)` pointer to block of memory. If the pointer is not the one created by `PERM_MALLOC`, `PERM_CALLOC`, or `PERM_STRDUP`, the behavior is undefined.

Example

```
char *name;
name = (char *) PERM_MALLOC(256);
...
PERM_FREE(name);
```

See Also

"FREE" on page 78, "PERM_MALLOC" on page 104, "PERM_CALLOC" on page 103, "PERM_REALLOC" on page 105, "PERM_STRDUP" on page 106

PERM_MALLOC

The `PERM_MALLOC` macro is a platform-independent substitute for the C library routine `malloc`. It provides allocation of memory that persists after the request that is being processed has been completed.

Syntax

```
void *PERM_MALLOC(int size)
```

Returns

A void pointer to a block of memory.

Parameters

`int size` is the number of bytes to allocate.

Example

```
/* Allocate 256 bytes for a name */
char *name;
name = (char *) PERM_MALLOC(256);
```

See Also

"MALLOC" on page 83, "PERM_FREE" on page 104, "PERM_STRDUP" on page 106, "PERM_CALLOC" on page 103, "PERM_REALLOC" on page 105

PERM_REALLOC

The `PERM_REALLOC` macro is a platform-independent substitute for the C library routine `realloc`. It changes the size of a specified memory block that was originally created by `PERM_MALLOC`, `PERM_CALLOC`, or `PERM_STRDUP`. The contents of the object remains unchanged up to the lesser of the old and new sizes. If the new size is larger, the new space is uninitialized.

Caution – Calling `PERM_REALLOC` for a block that was allocated with `MALLOC`, `CALLOC`, or `STRDUP` does not work.

Syntax

```
void *PERM_REALLOC(vod *ptr, int size)
```

Returns

A void pointer to a block of memory.

Parameters

`void *ptr` a void pointer to a block of memory created by `PERM_MALLOC`, `PERM_CALLOC`, or `PERM_STRDUP`.

`int size` is the number of bytes to which the memory block should be resized.

Example

```
char *name;
name = (char *) PERM_MALLOC(256);
if (NotBigEnough())
   name = (char *) PERM_REALLOC(name, 512);
```

See Also

"REALLOC" on page 113, "PERM_CALLOC" on page 103, "PERM_MALLOC" on page 104, "PERM_FREE" on page 104, "PERM_STRDUP" on page 106

PERM_STRDUP

The `PERM_STRDUP` macro is a platform-independent substitute for the C library routine `strdup`. It creates a new copy of a string in memory that persists after the request that is being processed has been completed. If pooled memory has been disabled in the configuration file (with the `pool-init` built-in SAF), `PERM_STRDUP` and `STRDUP` both obtain their memory from the system heap.

The `PERM_STRDUP` routine is functionally equivalent to:

```
newstr = (char *) PERM_MALLOC(strlen(str) + 1);strcpy(newstr, str);
```

A string created with `PERM_STRDUP` should be disposed with `PERM_FREE`.

Syntax

```
char *PERM_STRDUP(char *ptr);
```

Returns

A pointer to the new string.

Parameters

`char *ptr` is a pointer to a string.

See Also

"PERM_MALLOC" on page 104, "PERM_FREE" on page 104, "PERM_CALLOC" on page 103, "PERM_REALLOC" on page 105, "MALLOC" on page 83, "FREE" on page 78, "STRDUP" on page 121, "CALLOC" on page 63, "REALLOC" on page 113

prepare_nsapi_thread

The `prepare_nsapi_thread` function allows threads that are not created by the server to act like server-created threads. This function must be called before any NSAPI functions are called from a thread that is not server-created.

Syntax

```
void prepare_nsapi_thread(Request *rq, Session *sn);
```

Returns

`void`

Parameters

`Request *rq` is the Request.

`Session *sn` is the Session.

The `Request` and `Session` parameters are the same as the ones passed into your SAF.

See Also

"protocol_start_response" on page 109

protocol_dump822

The `protocol_dump822` function prints headers from a specified `pblock` into a specific buffer, with a specified size and position. Use this function to serialize the headers so that they can be sent, for example, in a mail message.

Syntax

```
char *protocol_dump822(pblock *pb, char *t, int *pos, int tsz);
```

Returns

A pointer to the buffer, which will be reallocated if necessary.

The function also modifies `*pos` to the end of the headers in the buffer.

Parameters

`pblock *pb` is the `pblock` structure.

`char *t` is the buffer, allocated with `MALLOC`, `CALLOC`, or `STRDUP`.

`int *pos` is the position within the buffer at which the headers are to be dumped.

`int tsz` is the size of the buffer.

See Also

"protocol_start_response" on page 109, "protocol_status" on page 110

protocol_set_finfo

The `protocol_set_finfo` function retrieves the `content-length` and `last-modified` date from a specified `stat` structure and adds them to the response headers (`rq->srvhdrs`). Call `protocol_set_finfo` before calling `protocol_start_response`.

Syntax

```
int protocol_set_finfo(Session *sn, Request *rq, struct stat *finfo);
```

Returns

The constant `REQ_PROCEED` if the request can proceed normally, or the constant `REQ_ABORTED` if the function should treat the request normally but not send any output to the client.

Parameters

`Session *sn` is the Session.

`Request *rq` is the Request.

The `Session` and `Request` parameters are the same as the ones passed into your SAF.

`stat *finfo` is the `stat` structure for the file.

The `stat` structure contains the information about the file from the file system. You can get the `stat` structure info using `request_stat_path`.

See Also

"protocol_start_response" on page 109, "protocol_status" on page 110

protocol_start_response

The `protocol_start_response` function initiates the HTTP response for a specified session and request. If the protocol version is HTTP/0.9, the function does nothing, because that version has no concept of status. If the protocol version is HTTP/1.0 or higher, the function sends a status line followed by the response headers. Because of buffering, the status line and response headers might not be sent immediately. To flush the status line and response headers, use the `net_flush` function. Use this function to set up HTTP and prepare the client and server to receive the body (or data) of the response.

Note – If you do not want the server to send the status line and response headers, set `rq->senthdrs = 1` before calling `protocol_start_response` or sending any data to the client.

Syntax

```
int protocol_start_response(Session *sn, Request *rq);
```

Returns

The constant `REQ_PROCEED` if the operation succeeds, in which case you should send the data you were preparing to send.

The constant `REQ_NOACTION` if the operation succeeds but the request method is `HEAD`, in which case no data should be sent to the client.

The constant `REQ_ABORTED` if the operation fails.

Parameters

`Session *sn` is the Session.

`Request *rq` is the Request.

The `Session` and `Request` parameters are the same as the ones passed into your SAF.

Example

```
/* REQ_NOACTION means the request was HEAD */
if (protocol_start_response(sn, rq) == REQ_NOACTION)
{
  filebuf_close(groupbuf); /* close our file*/
  return REQ_PROCEED;
}
```

See Also

"protocol_status" on page 110, "net_flush" on page 84

protocol_status

The `protocol_status` function sets the session status to indicate whether an error condition occurred. If the reason string is NULL, the server attempts to find a reason string for the given status code. If it finds none, it returns `Unknown reason`. The reason string is sent to the client in the HTTP response line. Use this function to set the status of the response before calling the function `protocol_start_response` or returning `REQ_ABORTED`.

For the complete list of valid status code constants, refer to the `nsapi.h` file.

Syntax

```
void protocol_status(Session *sn, Request *rq, int n, char *r);
```

Returns

`void`

Parameters

`Session *sn` is the Session.

`Request *rq` is the Request.

The `Session` and `Request` parameters are the same as the ones passed into your SAF.

`int n` is an HTTP status code constants above.

`char *r` is the reason string.

Example

```
/* if we find extra path-info, the URL was bad so tell the */
/* browser it was not found */
if (t = pblock_findval("path-info", rq->vars))
{
         protocol_status(sn, rq, PROTOCOL_NOT_FOUND, NULL);
      log_error(LOG_WARN, "function-name", sn, rq, "%s not found",path);
      return REQ_ABORTED;
}
```

See Also

"protocol_start_response" on page 109

protocol_uri2url

The `protocol_uri2url` function takes strings containing the given URI prefix and URI suffix, and creates a newly allocated, fully qualified URL in the form `http://(server):(port)(prefix)(suffix)`. See `protocol_uri2url_dynamic`.

If you want to omit either the URI prefix or suffix, use `""` instead of NULL as the value for either parameter.

Syntax

```
char *protocol_uri2url(char *prefix, char *suffix);
```

Returns

A new string containing the URL.

Parameters

`char *prefix` is the prefix.

`char *suffix` is the suffix.

See Also

"pblock_nvinsert" on page 99, "protocol_start_response" on page 109, "protocol_status" on page 110, "protocol_uri2url_dynamic" on page 111

protocol_uri2url_dynamic

The `protocol_uri2url` function takes strings containing the given URI prefix and URI suffix, and creates a newly allocated, fully qualified URL in the form `http://(server):(port)(prefix)(suffix)`.

If you want to omit either the URI prefix or suffix, use `""` instead of NULL as the value for either parameter.

The `protocol_uri2url_dynamic` function is similar to the `protocol_uri2url` function, but should be used whenever the `Session` and `Request` structures are available. This ensures that the URL it constructs refers to the host that the client specified.

Syntax

```
char *protocol_uri2url(char *prefix, char *suffix, Session *sn, Request *rq);
```

Returns

A new string containing the URL.

Parameters

`char *prefix` is the prefix.

`char *suffix` is the suffix.

`Session *sn` is the Session.

`Request *rq` is the Request.

The `Session` and `Request` parameters are the same as the ones passed into your SAF.

See Also

"protocol_start_response" on page 109, "protocol_status" on page 110

R

read

The `read` filter method is called when input data is required. Filters that modify or consume incoming data should implement the `read` filter method.

Upon receiving control, a read implementation should fill `buf` with up to `amount` bytes of input data. This data can be obtained by calling the "net_read" on page 85 function, as shown in the example below.

Syntax

```
int read(FilterLayer *layer, void *buf, int amount, int timeout);
```

Returns

The number of bytes placed in `buf` on success. `0` if no data is available, or a negative value if an error occurs.

Parameters

`FilterLayer *layer` is the filter layer in which the filter is installed.

`void *buf` is the buffer in which data should be placed.

`int amount` is the maximum number of bytes that should be placed in the buffer.

`int timeout` is the number of seconds to allow the `read` operation to return. The purpose of `timeout` is not to return because not enough bytes were read in the given time, but to limit the amount of time devoted to waiting until some data arrives.

Example

```
int myfilter_read(FilterLayer *layer, void *buf, int amount, int timeout)
{
    return net_read(layer->lower, buf, amount, timeout);
}
```

See Also

"net_read" on page 85, "filter_create" on page 73

REALLOC

The REALLOC macro is a platform-independent substitute for the C library routine realloc. It changes the size of a specified memory block that was originally created by MALLOC, CALLOC, or STRDUP. The contents of the object remains unchanged up to the lesser of the old and new sizes. If the new size is larger, the new space is uninitialized.

Caution – Calling REALLOC for a block that was allocated with PERM_MALLOC, PERM_CALLOC, or PERM_STRDUP will not work.

Syntax

```
void *REALLOC(void *ptr, int size);
```

Returns

A pointer to the new space if the request is satisfied.

Parameters

void *ptr is a (void *) pointer to a block of memory. If the pointer is not the one created by MALLOC, CALLOC, or STRDUP, the behavior is undefined.

int size is the number of bytes to allocate.

Example

```
char *name;
name = (char *) MALLOC(256);
if (NotBigEnough())
    name = (char *) REALLOC(name, 512);
```

See Also

"CALLOC" on page 63, "MALLOC" on page 83, "FREE" on page 78, "STRDUP" on page 121, "PERM_REALLOC" on page 105

remove

The `remove` filter method is called when the filter stack is destroyed, or when a filter is removed from a filter stack by the `filter_remove` function or `remove-filter` SAF.

Note – It may be too late to flush buffered data when the `remove` method is invoked. For this reason, filters that buffer outgoing data should implement the `flush` filter method.

Syntax

```
void remove(FilterLayer *layer);
```

Returns

`void`

Parameters

`FilterLayer *layer` is the filter layer in which the filter is installed.

See Also

"flush" on page 78, "filter_remove" on page 77, "filter_create" on page 73

request_get_vs

The `request_get_vs` function finds the `VirtualServer*` to which a request is directed.

The returned `VirtualServer*` is valid only for the current request. To retrieve a virtual server ID that is valid across requests, use "vs_get_id" on page 154.

Syntax

```
const VirtualServer* request_get_vs(Request* rq);
```

Returns

The `VirtualServer*` to which the request is directed.

Parameters

`Request *rq` is the request for which the `VirtualServer*` is returned.

See Also

"vs_get_id" on page 154

request_header

The request_header function finds an entry in the pblock containing the client's HTTP request headers (rq->headers). You must use this function rather than pblock_findval when accessing the client headers, since the server might begin processing the request before the headers have been completely read.

Syntax

```
int request_header(char *name, char **value, Session *sn, Request *rq);
```

Returns

A result code, REQ_PROCEED if the header was found, REQ_ABORTED if the header was not found, REQ_EXIT if there was an error reading from the client.

Parameters

char *name is the name of the header.

char **value is the address where the function will place the value of the specified header. If none is found, the function stores a NULL.

Session *sn is the Session.

Request *rq is the Request.

The Session and Request parameters are the same as the ones passed into your SAF.

See Also

request_create, request_free

request_stat_path

The request_stat_path function returns the file information structure for a specified path or, if none is specified, the path entry in the vars pblock in the specified request structure. If the resulting file name points to a file that the server can read, request_stat_path returns a new file information structure. This structure contains information on the size of the file, its owner, when it was created, and when it was last modified.

You should use request_stat_path to retrieve information on the file you are currently accessing (instead of calling stat directly), because this function keeps track of previous calls for the same path and returns its cached information.

Syntax

```
struct stat *request_stat_path(char *path, Request *rq);
```

Returns

Returns a pointer to the file information structure for the file named by the `path` parameter. Do not free this structure. Returns NULL if the file is not valid or the server cannot read it. In this case, it also leaves an error message describing the problem in `rq->staterr`.

Parameters

`char *path` is the string containing the name of the path. If the value of `path` is NULL, the function uses the `path` entry in the `vars` pblock in the `request` structure denoted by `rq`.

`Request *rq` is the request identifier for a Server Application Function call.

Example

```
fi = request_stat_path(path, rq);
```

See Also

`request_create`, `request_free`, "request_header" on page 115

request_translate_uri

The `request_translate_uri` function performs virtual to physical mapping on a specified URI during a specified session. Use this function to determine the file to be sent back if a given URI is accessed.

Syntax

```
char *request_translate_uri(char *uri, Session *sn);
```

Returns

A path string if it performed the mapping, or NULL if it could not perform the mapping.

Parameters

`char *uri` is the name of the URI.

`Session *sn` is the `Session` parameter that is passed into your SAF.

See Also

`request_create`, `request_free`, "request_header" on page 115

S

sendfile

The `sendfile` filter method is called when the contents of a file are to be sent. Filters that modify or consume outgoing data can choose to implement the `sendfile` filter method.

If a filter implements the `write` filter method but not the `sendfile` filter method, the server will automatically translate "net_sendfile" on page 86 calls to "net_write" on page 87 calls. As a result, filters interested in the outgoing data stream do not need to implement the `sendfile` filter method. However, for performance reasons, it is beneficial for filters that implement the `write` filter method to also implement the `sendfile` filter method.

Syntax

```
int sendfile(FilterLayer *layer, const sendfiledata *data);
```

Returns

The number of bytes consumed, which may be less than the requested amount if an error occurred.

Parameters

`FilterLayer *layer` is the filter layer in which the filter is installed.

`const sendfiledata *sfd` identifies the data to send.

Example

```
int myfilter_sendfile(FilterLayer *layer, const sendfiledata *sfd)
{
    return net_sendfile(layer->lower, sfd);
}
```

See Also

"net_sendfile" on page 86, "filter_create" on page 73

session_dns

The `session_dns` function resolves the IP address of the client associated with a specified session into its DNS name. It returns a newly allocated string. You can use `session_dns` to change the numeric IP address into something more readable.

The `session_maxdns` function verifies that the client is who it claims to be; the `session_dns` function does not perform this verification.

Note – This function works only if the `DNS` directive is enabled in the `magnus.conf` file. For more information, see Appendix B.

Syntax

```
char *session_dns(Session *sn);
```

Returns

A string containing the host name, or NULL if the DNS name cannot be found for the IP address.

Parameters

`Session *sn` is the Session.

The `Session` is the same as the one passed to your SAF.

session_maxdns

The `session_maxdns` function resolves the IP address of the client associated with a specified session into its DNS name. It returns a newly allocated string. You can use `session_maxdns` to change the numeric IP address into something more readable.

Note – This function works only if the `DNS` directive is enabled in the `magnus.conf` file. For more information, see Appendix B.

Syntax

```
char *session_maxdns(Session *sn);
```

Returns

A string containing the host name, or NULL if the DNS name cannot be found for the IP address.

Parameters

`Session *sn` is the `Session`.

The `Session` is the same as the one passed to your SAF.

shexp_casecmp

The `shexp_casecmp` function validates a specified shell expression and compares it with a specified string. It returns one of three possible values representing match, no match, and invalid comparison. The comparison (in contrast to that of the `shexp_cmp` function) is not case-sensitive.

Use this function if you have a shell expression like `*.netscape.com` and make sure that a string matches it, such as `foo.netscape.com`.

Syntax

```
int shexp_casecmp(char *str, char *exp);
```

Returns

`0` if a match was found.

`1` if no match was found.

`-1` if the comparison resulted in an invalid expression.

Parameters

`char *str` is the string to be compared.

`char *exp` is the shell expression (wildcard pattern) to compare against.

See Also

"shexp_cmp" on page 119, "shexp_match" on page 120, "shexp_valid" on page 121

shexp_cmp

The `shexp_cmp` function validates a specified shell expression and compares it with a specified string. It returns one of three possible values representing match, no match, and invalid comparison. The comparison (in contrast to that of the `shexp_casecmp` function) is case-sensitive.

Use this function for a shell expression like `*.netscape.com` and make sure that a string matches it, such as `foo.netscape.com`.

Syntax

```
int shexp_cmp(char *str, char *exp);
```

Returns

`0` if a match was found.

`1` if no match was found.

`-1` if the comparison resulted in an invalid expression.

Parameters

`char *str` is the string to be compared.

`char *exp` is the shell expression (wildcard pattern) to compare against.

Example

```
/* Use wildcard match to see if this path is one we want */
char *path;
char *match = "/usr/netscape/*";
if (shexp_cmp(path, match) != 0)
    return REQ_NOACTION;   /* no match */
```

See Also

"shexp_casecmp" on page 119, "shexp_match" on page 120, "shexp_valid" on page 121

shexp_match

The `shexp_match` function compares a specified pre-validated shell expression against a specified string. It returns one of three possible values representing match, no match, and invalid comparison. The comparison (in contrast to that of the `shexp_casecmp` function) is case-sensitive.

The `shexp_match` function does not perform validation of the shell expression; instead the function assumes that you have already called `shexp_valid`.

Use this function for a shell expression such as `*.netscape.com`, and make sure that a string matches it, such as `foo.netscape.com`.

Syntax

```
int shexp_match(char *str, char *exp);
```

Returns

`0` if a match was found.

`1` if no match was found.

`-1` if the comparison resulted in an invalid expression.

Parameters

`char *str` is the string to be compared.

`char *exp` is the prevalidated shell expression (wildcard pattern) to compare against.

See Also

"shexp_casecmp" on page 119, "shexp_cmp" on page 119, "shexp_valid" on page 121

shexp_valid

The `shexp_valid` function validates a specified shell expression named by `exp`. Use this function to validate a shell expression before using the function `shexp_match` to compare the expression with a string.

Syntax

```
int shexp_valid(char *exp);
```

Returns

The constant `NON_SXP` if `exp` is a standard string.

The constant `INVALID_SXP` if `exp` is a shell expression, but invalid.

The constant `VALID_SXP` if `exp` is a valid shell expression.

Parameters

`char *exp` is the shell expression (wildcard pattern) to be validated.

See Also

"shexp_casecmp" on page 119, "shexp_match" on page 120, "shexp_cmp" on page 119

STRDUP

The `STRDUP` macro is a platform-independent substitute for the C library routine `strdup`. It creates a new copy of a string in the request's memory pool. The memory can be explicitly freed by a call to `FREE`. If the memory is not explicitly freed, it is automatically freed after processing the current request. If pooled memory has been disabled in the configuration file (with the `pool-init` built-in SAF), `PERM_STRDUP` and `STRDUP` both obtain their memory from the system heap. However, since the memory allocated by `STRDUP` is automatically freed, it should not be shared between threads.

The STRDUP routine is functionally equivalent to:

```
newstr = (char *) MALLOC(strlen(str) + 1);
strcpy(newstr, str);
```

Syntax

```
char *STRDUP(char *ptr);
```

Returns

A pointer to the new string.

Parameters

`char *ptr` is a pointer to a string.

Example

```
char *name1 = "MyName";
char *name2 = STRDUP(name1);
```

See Also

"CALLOC" on page 63,"MALLOC" on page 83, "FREE" on page 78, "REALLOC" on page 113, "PERM_STRDUP" on page 106

system_errmsg

The `system_errmsg` function returns the last error that occurred from the most recent system call. This function is implemented as a macro that returns an entry from the global array `sys_errlist`. Use this macro to help with I/O error diagnostics.

Syntax

```
char *system_errmsg(int param1);
```

Returns

A string containing the text of the latest error message that resulted from a system call. Do not FREE this string.

Parameters

`int param1` is reserved, and should always have the value 0.

See Also

"system_fopenRO" on page 124, "system_fopenRW" on page 124, "system_fopenWA" on page 125, "system_lseek" on page 129, "system_fread" on page 126, "system_fwrite" on page 126, "system_fwrite_atomic" on page 127, "system_flock" on page 123, "system_ulock" on page 130, "system_fclose" on page 123

system_fclose

The `system_fclose` function closes a specified file descriptor. The `system_fclose` function must be called for every file descriptor opened by any of the `system_fopen` functions.

Syntax

```
int system_fclose(SYS_FILE fd);
```

Returns

`0` if the close succeeds, or the constant `IO_ERROR` if the close fails.

Parameters

`SYS_FILE fd` is the platform-independent file descriptor.

Example

```
SYS_FILE logfd;
system_fclose(logfd);
```

See Also

"system_errmsg" on page 122, "system_fopenRO" on page 124, "system_fopenRW" on page 124, "system_fopenWA" on page 125, "system_lseek" on page 129, "system_fread" on page 126, "system_fwrite" on page 126, "system_fwrite_atomic" on page 127, "system_flock" on page 123, "system_ulock" on page 130

system_flock

The `system_flock` function locks the specified file against interference from other processes. Use `system_flock` if you do not want other processes to use the file you currently have open. Overusing file locking can cause performance degradation and possibly lead to deadlocks.

Syntax

```
int system_flock(SYS_FILE fd);
```

Returns

The constant `IO_OKAY` if the lock succeeds, or the constant `IO_ERROR` if the lock fails.

Parameters

`SYS_FILE fd` is the platform-independent file descriptor.

See Also

"system_errmsg" on page 122, "system_fopenRO" on page 124, "system_fopenRW" on page 124, "system_fopenWA" on page 125, "system_lseek" on page 129, "system_fread" on page 126, "system_fwrite" on page 126, "system_fwrite_atomic" on page 127, "system_ulock" on page 130, "system_fclose" on page 123

system_fopenRO

The `system_fopenRO` function opens the file identified by `path` in read-only mode and returns a valid file descriptor. Use this function to open files that will not be modified by your program. In addition, you can use `system_fopenRO` to open a new file buffer structure using `filebuf_open`.

Syntax

```
SYS_FILE system_fopenRO(char *path);
```

Returns

The system-independent file descriptor (`SYS_FILE`) if the open succeeds, or `0` if the open fails.

Parameters

`char *path` is the file name.

See Also

"system_errmsg" on page 122, "system_fopenRW" on page 124, "system_fopenWA" on page 125, "system_lseek" on page 129, "system_fread" on page 126, "system_fwrite" on page 126, "system_fwrite_atomic" on page 127, "system_flock" on page 123, "system_ulock" on page 130, "system_fclose" on page 123

system_fopenRW

The `system_fopenRW` function opens the file identified by `path` in read-write mode and returns a valid file descriptor. If the file already exists, `system_fopenRW` does not truncate it. Use this function to open files that can be read and written by your program.

Syntax

```
SYS_FILE system_fopenRW(char *path);
```

Returns

The system-independent file descriptor (`SYS_FILE`) if the open succeeds, or `0` if the open fails.

Parameters

`char *path` is the file name.

Example

```
SYS_FILE fd;
fd = system_fopenRO(pathname);
if (fd == SYS_ERROR_FD)
break;
```

See Also

"system_errmsg" on page 122, "system_fopenRO" on page 124, "system_fopenWA" on page 125, "system_lseek" on page 129, "system_fread" on page 126, "system_fwrite" on page 126, "system_fwrite_atomic" on page 127, "system_flock" on page 123, "system_ulock" on page 130, "system_fclose" on page 123

system_fopenWA

The `system_fopenWA` function opens the file identified by `path` in write-append mode and returns a valid file descriptor. Use this function to open those files to which your program will append data.

Syntax

```
SYS_FILE system_fopenWA(char *path);
```

Returns

The system-independent file descriptor (`SYS_FILE`) if the open succeeds, or `0` if the open fails.

Parameters

`char *path` is the file name.

See Also

"system_errmsg" on page 122, "system_fopenRO" on page 124, "system_fopenRW" on page 124, "system_lseek" on page 129, "system_fread" on page 126, "system_fwrite" on page 126, "system_fwrite_atomic" on page 127, "system_flock" on page 123, "system_ulock" on page 130, "system_fclose" on page 123

system_fread

The `system_fread` function reads a specified number of bytes from a specified file into a specified buffer. It returns the number of bytes read. Before `system_fread` can be used, you must open the file using any of the `system_fopen` functions (except `system_fopenWA`).

Syntax

```
int system_fread(SYS_FILE fd, char *buf, int sz);
```

Returns

The number of bytes read, which may be less than the requested size if an error occurs, or the end of the file was reached before that number of characters were obtained.

Parameters

`SYS_FILE fd` is the platform-independent file descriptor.

`char *buf` is the buffer to receive the bytes.

`int sz` is the number of bytes to read.

See Also

"system_errmsg" on page 122, "system_fopenRO" on page 124, "system_fopenRW" on page 124, "system_fopenWA" on page 125, "system_lseek" on page 129, "system_fwrite" on page 126, "system_fwrite_atomic" on page 127, "system_flock" on page 123, "system_ulock" on page 130, "system_fclose" on page 123

system_fwrite

The `system_fwrite` function writes a specified number of bytes from a specified buffer into a specified file.

Before `system_fwrite` can be used, you must open the file using any of the `system_fopen` functions (except `system_fopenRO`).

Syntax

```
int system_fwrite(SYS_FILE fd, char *buf, int sz);
```

Returns

The constant `IO_OKAY` if the write succeeds, or the constant `IO_ERROR` if the write fails.

Parameters

`SYS_FILE fd` is the platform-independent file descriptor.

`char *buf` is the buffer containing the bytes to be written.

`int sz` is the number of bytes to write to the file.

See Also

"system_errmsg" on page 122, "system_fopenRO" on page 124, "system_fopenRW" on page 124, "system_fopenWA" on page 125, "system_lseek" on page 129, "system_fread" on page 126, "system_fwrite" on page 126, "system_fwrite_atomic" on page 127, "system_flock" on page 123, "system_ulock" on page 130, "system_fclose" on page 123

system_fwrite_atomic

The `system_fwrite_atomic` function writes a specified number of bytes from a specified buffer into a specified file. This function also locks the file prior to performing the write, and then unlocks it when done, thereby avoiding interference between simultaneous write actions. Before `system_fwrite_atomic` can be used, you must open the file using any of the `system_fopen` functions, except `system_fopenRO`.

Syntax

```
int system_fwrite_atomic(SYS_FILE fd, char *buf, int sz);
```

Returns

The constant `IO_OKAY` if the write/lock succeeds, or the constant `IO_ERROR` if the write/lock fails.

Parameters

`SYS_FILE fd` is the platform-independent file descriptor.

`char *buf` is the buffer containing the bytes to be written.

`int sz` is the number of bytes to write to the file.

Example

```
SYS_FILE logfd;
char *logmsg = "An error occurred.";
system_fwrite_atomic(logfd, logmsg, strlen(logmsg));
```

See Also

"system_errmsg" on page 122, "system_fopenRO" on page 124, "system_fopenRW" on page 124, "system_fopenWA" on page 125, "system_lseek" on page 129, "system_fread" on page 126, "system_fwrite" on page 126, "system_flock" on page 123, "system_ulock" on page 130, "system_fclose" on page 123

system_gmtime

The `system_gmtime` function is a thread-safe version of the standard `gmtime` function. It returns the current time adjusted to Greenwich Mean Time.

Syntax

```
struct tm *system_gmtime(const time_t *tp, const struct tm *res);
```

Returns

A pointer to a calendar time (`tm`) structure containing the GMT time. Depending on your system, the pointer may point to the data item represented by the second parameter, or it may point to a statically-allocated item. For portability, do not assume either situation.

Parameters

`time_t *tp` is an arithmetic time.

`tm *res` is a pointer to a calendar time (`tm`) structure.

Example

```
time_t tp;
struct tm res, *resp;
tp = time(NULL);
resp = system_gmtime(&tp, &res);
```

See Also

"system_localtime" on page 128, "util_strftime" on page 147

system_localtime

The `system_localtime` function is a thread-safe version of the standard `localtime` function. It returns the current time in the local time zone.

Syntax

```
struct tm *system_localtime(const time_t *tp, const struct tm *res);
```

Returns

A pointer to a calendar time (`tm`) structure containing the local time. Depending on your system, the pointer may point to the data item represented by the second parameter, or it may point to a statically-allocated item. For portability, do not assume either situation.

Parameters

`time_t *tp` is an arithmetic time.

`tm *res` is a pointer to a calendar time (`tm`) structure.

See Also

"system_gmtime" on page 128, "util_strftime" on page 147

system_lseek

The `system_lseek` function sets the file position of a file. This affects where data from `system_fread` or `system_fwrite` is read or written.

Syntax

```
int system_lseek(SYS_FILE fd, int offset, int whence);
```

Returns

The offset, in bytes, of the new position from the beginning of the file if the operation succeeds, or `-1` if the operation fails.

Parameters

`SYS_FILE fd` is the platform-independent file descriptor.

`int offset` is a number of bytes relative to `whence`. It may be negative.

`int whence` is one of the following constants:

`SEEK_SET`, from the beginning of the file.

`SEEK_CUR`, from the current file position.

`SEEK_END`, from the end of the file.

See Also

"system_errmsg" on page 122, "system_fopenRO" on page 124, "system_fopenRW" on page 124, "system_fopenWA" on page 125, "system_fread" on page 126, "system_fwrite" on page 126, "system_fwrite_atomic" on page 127, "system_flock" on page 123, "system_ulock" on page 130, "system_fclose" on page 123

system_rename

The `system_rename` function renames a file. It does not work on directories if the old and new directories are on different file systems.

Syntax

```
int system_rename(char *old, char *new);
```

Returns

`0` if the operation succeeds, or `-1` if the operation fails.

Parameters

`char *old` is the old name of the file.

`char *new` is the new name for the file.

system_ulock

The `system_ulock` function unlocks the specified file that has been locked by the function `system_lock`. For more information about locking, see `system_flock`.

Syntax

```
int system_ulock(SYS_FILE fd);
```

Returns

The constant `IO_OKAY` if the operation succeeds, or the constant `IO_ERROR` if the operation fails.

Parameters

`SYS_FILE fd` is the platform-independent file descriptor.

See Also

"system_errmsg" on page 122, "system_fopenRO" on page 124, "system_fopenRW" on page 124, "system_fopenWA" on page 125, "system_fread" on page 126, "system_fwrite" on page 126, "system_fwrite_atomic" on page 127, "system_flock" on page 123, "system_fclose" on page 123

system_unix2local

The `system_unix2local` function converts a specified UNIX-style path name to a local file system path name. Use this function when you have a file name in the UNIX format (such as one containing forward slashes), and you need to access a file on another system such as Windows. You can use `system_unix2local` to convert the UNIX file name into the format that Windows accepts. In the UNIX environment this function does nothing, but can be called for portability.

Syntax

```
char *system_unix2local(char *path, char *lp);
```

Returns

A pointer to the local file system path string.

Parameters

`char *path` is the UNIX-style path name to be converted.

`char *lp` is the local path name.

You must allocate the parameter `lp`, and it must contain enough space to hold the local path name.

See Also

"system_fclose" on page 123, "system_flock" on page 123, "system_fopenRO" on page 124, "system_fopenRW" on page 124, "system_fopenWA" on page 125, "system_fwrite" on page 126

systhread_attach

The `systhread_attach` function makes an existing thread into a platform-independent thread.

Syntax

```
SYS_THREAD systhread_attach(void);
```

Returns

A SYS_THREAD pointer to the platform-independent thread.

Parameters

none

See Also

"systhread_current" on page 132, "systhread_getdata" on page 132, "systhread_newkey" on page 133, "systhread_setdata" on page 133, "systhread_sleep" on page 134, "systhread_start" on page 134, "systhread_timerset" on page 135

systhread_current

The systhread_current function returns a pointer to the current thread.

Syntax

```
SYS_THREAD systhread_current(void);
```

Returns

A SYS_THREAD pointer to the current thread.

Parameters

none

See Also

"systhread_getdata" on page 132, "systhread_newkey" on page 133, "systhread_setdata" on page 133, "systhread_sleep" on page 134, "systhread_start" on page 134, "systhread_timerset" on page 135

systhread_getdata

The systhread_getdata function gets data that is associated with a specified key in the current thread.

Syntax

```
void *systhread_getdata(int key);
```

Returns

A pointer to the data that was earlier used with the `systhread_setkey` function from the current thread, using the same value of `key` if the call succeeds. Returns NULL if the call does not succeed; for example, if the `systhread_setkey` function was never called with the specified key during this session.

Parameters

`int key` is the value associated with the stored data by a `systhread_setdata` function. Keys are assigned by the `systhread_newkey` function.

See Also

"systhread_current" on page 132, "systhread_newkey" on page 133, "systhread_setdata" on page 133, "systhread_sleep" on page 134, "systhread_start" on page 134, "systhread_timerset" on page 135

systhread_newkey

The `systhread_newkey` function allocates a new integer key (identifier) for thread-private data. Use this key to identify a variable that you want to localize to the current thread, then use the `systhread_setdata` function to associate a value with the key.

Syntax

```
int systhread_newkey(void);
```

Returns

An integer key.

Parameters

none

See Also

"systhread_current" on page 132, "systhread_getdata" on page 132, "systhread_setdata" on page 133, "systhread_sleep" on page 134, "systhread_start" on page 134, "systhread_timerset" on page 135

systhread_setdata

The `systhread_setdata` function associates data with a specified key number for the current thread. Keys are assigned by the `systhread_newkey` function.

Syntax

```
void systhread_setdata(int key, void *data);
```

Returns

`void`

Parameters

`int key` is the priority of the thread.

`void *data` is the pointer to the string of data to be associated with the value of `key`.

See Also

"systhread_current" on page 132, "systhread_getdata" on page 132, "systhread_newkey" on page 133, "systhread_sleep" on page 134, "systhread_start" on page 134, "systhread_timerset" on page 135

systhread_sleep

The `systhread_sleep` function puts the calling thread to sleep for a given time.

Syntax

```
void systhread_sleep(int milliseconds);
```

Returns

`void`

Parameters

`int milliseconds` is the number of milliseconds the thread is to sleep.

See Also

"systhread_current" on page 132, "systhread_getdata" on page 132, "systhread_newkey" on page 133, "systhread_setdata" on page 133, "systhread_start" on page 134, "systhread_timerset" on page 135

systhread_start

The `systhread_start` function creates a thread with the given priority, allocates a stack of a specified number of bytes, and calls a specified function with a specified argument.

Syntax

```
SYS_THREAD systhread_start(int prio, int stksz, void (*fn)(void *), void *arg);
```

Returns

A new `SYS_THREAD` pointer if the call succeeds, or the constant `SYS_THREAD_ERROR` if the call does not succeed.

Parameters

`int prio` is the priority of the thread. Priorities are system-dependent.

`int stksz` is the stack size in bytes. If `stksz` is zero (`0`), the function allocates a default size.

`void (*fn)(void *)` is the function to call.

`void *arg` is the argument for the `fn` function.

See Also

"systhread_current" on page 132, "systhread_getdata" on page 132, "systhread_newkey" on page 133, "systhread_setdata" on page 133, "systhread_sleep" on page 134, "systhread_timerset" on page 135

systhread_timerset

The `systhread_timerset` function starts or resets the interrupt timer interval for a thread system.

Because most systems do not allow the timer interval to be changed, this should be considered a suggestion, rather than a command.

Syntax

```
void systhread_timerset(int usec);
```

Returns

`void`

Parameters

`int usec` is the time, in microseconds

See Also

"systhread_current" on page 132, "systhread_getdata" on page 132, "systhread_newkey" on page 133, "systhread_setdata" on page 133, "systhread_sleep" on page 134, "systhread_start" on page 134

U

USE_NSAPI_VERSION

Plug-in developers can define the `USE_NSAPI_VERSION` macro before including the `nsapi.h` header file to request a particular version of NSAPI. The requested NSAPI version is encoded by multiplying the major version number by 100 and then adding this to the minor version number. For example, the following code requests NSAPI 3.2 features:

```
#define USE_NSAPI_VERSION 302 /* We want NSAPI 3.2 (Web Server 6.1) */
#include "nsapi.h"
```

To develop a plug-in that is compatible across multiple server versions, define `USE_NSAPI_VERSION` to the highest NSAPI version supported by all of the target server versions.

The following table lists server versions and the highest NSAPI version supported by each:

TABLE 5–2 NSAPI Versions Supported by Different Servers

Server Version	NSAPI Version
iPlanet Web Server 4.1	3.0
iPlanet Web Server 6.0	3.1
Netscape Enterprise Server 6.0	3.1
Netscape Enterprise Server 6.1	3.1
Sun ONE Application Server 7.0	3.1
Sun ONE Web Server 6.1	3.2
Sun Java System Web Proxy Server 4.0	3.3
Sun Java System Web Server 7.0	3.3

It is an error to request a version of NSAPI higher than the highest version supported by the `nsapi.h` header that the plug-in is being compiled against. Additionally, to use `USE_NSAPI_VERSION`, you must compile against an `nsapi.h` header file that supports NSAPI 3.2 or higher.

Syntax

```
int USE_NSAPI_VERSION
```

Example

The following code can be used when building a plug-in designed to work with iPlanet Web Server 4.1 and Sun Java System Web Server 7.0:

```
#define USE_NSAPI_VERSION 300 /* We want NSAPI 3.0 (Web Server 4.1) */
#include "nsapi.h"
```

See Also

"NSAPI_RUNTIME_VERSION" on page 92, "NSAPI_VERSION" on page 93

util_can_exec

UNIX Only

The `util_can_exec` function checks that a specified file can be executed, returning either a `1` (executable) or a `0`. The function checks if the file can be executed by the user with the given user and group ID.

Use this function before executing a program using the `exec` system call.

Syntax

```
int util_can_exec(struct stat *finfo, uid_t uid, gid_t gid);
```

Returns

`1` if the file is executable, or `0` if the file is not executable.

Parameters

`stat *finfo` is the `stat` structure associated with a file.

`uid_t uid` is the UNIX user id.

`gid_t gid` is the UNIX group id. Together with `uid`, this determines the permissions of the UNIX user.

See Also

"util_env_create" on page 139, "util_getline" on page 141, "util_hostname" on page 142

util_chdir2path

The `util_chdir2path` function changes the current working directory. Because a server process can service multiple requests concurrently but has only a single current working directory, this function should not be used.

Syntax

```
int util_chdir2path(char *path);
```

Returns

`0` if the directory change succeeds, or `-1` if the directory can not bee changed.

Parameters

`char *path` is the name of a directory.

The parameter must be a writable string.

util_cookie_find

The `util_cookie_find` function finds a specific cookie in a cookie string and returns its value.

Syntax

```
char *util_cookie_find(char *cookie, char *name);
```

Returns

If successful, returns a pointer to the NULL-terminated value of the cookie. Otherwise, returns NULL. This function modifies the cookie string parameter by null-terminating the name and value.

Parameters

`char *cookie` is the value of the `Cookie:` request header.

`char *name` is the name of the cookie whose value is to be retrieved.

util_env_find

The `util_env_find` function locates the string denoted by a name in a specified environment and returns the associated value. Use this function to find an entry in an environment.

Syntax

```
char *util_env_find(char **env, char *name);
```

Returns

The value of the environment variable if it is found, or NULL if the string was not found.

Parameters

`char **env` is the environment.

`char *name` is the name of an environment variable in `env`.

See Also

"util_env_replace" on page 140, "util_env_str" on page 140, "util_env_free" on page 139, "util_env_create" on page 139

util_env_create

The `util_env_create` function creates and allocates the environment specified by `env`, returns a pointer to the environment. If the parameter `env` is NULL, the function allocates a new environment. Use `util_env_create` to create an environment when executing a new program.

Syntax

```
#include <base/util.h>
char **util_env_create(char **env, int n, int *pos);
```

Returns

A pointer to an environment.

Parameters

`char **env` is the environment or NULL.

`int n` is the maximum number of environment entries that you want in the environment.

`int *pos` is an integer that keeps track of the number of entries used in the environment.

See Also

"util_env_replace" on page 140, "util_env_str" on page 140, "util_env_free" on page 139, "util_env_find" on page 138

util_env_free

The `util_env_free` function frees a specified environment. Use this function to deallocate an environment you created using the function `util_env_create`.

Syntax

```
void util_env_free(char **env);
```

Returns

`void`

Parameters

`char **env` is the environment to be freed.

See Also

"util_env_replace" on page 140, "util_env_str" on page 140, "util_env_create" on page 139, "util_env_create" on page 139

util_env_replace

The `util_env_replace` function replaces the occurrence of the variable denoted by a name in a specified environment with a specified value. Use this function to change the value of a setting in an environment.

Syntax

```
void util_env_replace(char **env, char *name, char *value);
```

Returns

`void`

Parameters

`char **env` is the environment.

`char *name` is the name of a name-value pair.

`char *value` is the new value to be stored.

See Also

"util_env_str" on page 140, "util_env_free" on page 139, "util_env_create" on page 139, "util_env_create" on page 139

util_env_str

The `util_env_str` function creates an environment entry and returns it. This function does not check for non-alphanumeric symbols in the name (such as the equal sign "="). You can use this function to create a new environment entry.

Syntax

```
char *util_env_str(char *name, char *value);
```

Returns

A newly allocated string containing the name-value pair.

Parameters

`char *name` is the name of a name-value pair.

`char *value` is the new value to be stored.

See Also

"util_env_replace" on page 140, "util_env_free" on page 139, "util_env_create" on page 139, "util_env_create" on page 139

util_getline

The `util_getline` function scans the specified file buffer to find a line feed or carriage return/line feed terminated string. The string is copied into the specified buffer, and NULL-terminates it. The function returns a value that indicates whether the operation stored a string in the buffer, encountered an error, or reached the end of the file.

Use this function to scan lines out of a text file, such as a configuration file.

Syntax

```
int util_getline(filebuf *buf, int lineno, int maxlen, char *line);
```

Returns

`0` if successful; `line` contains the string.

`1` if the end of file is reached; `line` contains the string.

`-1` if an error occurs; `line` contains a description of the error.

Parameters

`filebuf *buf` is the file buffer to be scanned.

`int lineno` is used to include the line number in the error message when an error occurs. The caller is responsible for making sure the line number is accurate.

`int maxlen` is the maximum number of characters that can be written into `l`.

`char *l` is the buffer in which to store the string. The user is responsible for allocating and deallocating `line`.

util_hostname

The `util_hostname` function retrieves the local host name and returns it as a string. If the function cannot find a fully-qualified domain name, it returns NULL. You can reallocate or free this string. Use this function to determine the name of the system you are on.

Syntax

```
char *util_hostname(void);
```

Returns

A string containing that name, if a fully-qualified domain name is found; otherwise, returns NULL.

Parameters

none

util_is_mozilla

The `util_is_mozilla` function checks whether a specified user-agent header string is a mozilla browser of at least a specified revision level, returning a `1` if it is, and `0` otherwise. It uses strings to specify the revision level to avoid ambiguities such as 1.56 > 1.5.

Syntax

```
int util_is_mozilla(char *ua, char *major, char *minor);
```

Returns

1 if the user-agent is a mozilla browser, or 0 if the user-agent is not a mozilla browser.

Parameters

`char *ua` is the user-agent string from the request headers.

`char *major` is the major release number (to the left of the decimal point).

`char *minor` is the minor release number (to the right of the decimal point).

See Also

"util_is_url" on page 143, "util_later_than" on page 143

util_is_url

The `util_is_url` function checks whether a string is a URL, returns `1` if it is a URL and `0` otherwise. The string is a URL if it begins with alphabetic characters followed by a colon (:).

Syntax

```
int util_is_url(char *url);
```

Returns

`1` if the string specified by `url` is a URL, or `0` if the string specified by `url` is not a URL.

Parameters

`char *url` is the string to be examined.

See Also

"util_is_mozilla" on page 142, "util_later_than" on page 143

util_itoa

The `util_itoa` function converts a specified integer to a string, and returns the length of the string. Use this function to create a textual representation of a number.

Syntax

```
int util_itoa(int i, char *a);
```

Returns

The length of the string created.

Parameters

`int i` is the integer to be converted.

`char *a` is the ASCII string that represents the value. The user is responsible for the allocation and deallocation of `a`, and it should be at least 32 bytes long.

util_later_than

The `util_later_than` function compares the date specified in a time structure against a date specified in a string. If the date in the string is later than or equal to the one in the time structure, the function returns `1`. Use this function to handle RFC 822, RFC 850, and ctime formats.

Syntax

```
int util_later_than(struct tm *lms, char *ims);
```

Returns

`1` if the date represented by `ims` is the same as or later than that represented by the `lms`, or `0` if the date represented by `ims` is earlier than that represented by the `lms`.

Parameters

`tm *lms` is the time structure containing a date.

`char *ims` is the string containing a date.

See Also

"util_strftime" on page 147

util_sh_escape

The `util_sh_escape` funct ion parses a specified string and places a backslash (\) in front of any shell-special characters, returning the resultant string. Use this function to ensure that strings from clients do not cause a shell to do anything unexpected.

The shell-special characters are the space plus the following characters:

```
&;‘’"|*?~<>^()[]{}$\#!
```

Syntax

```
char *util_sh_escape(char *s);
```

Returns

A newly allocated string.

Parameters

`char *s` is the string to be parsed.

See Also

"util_uri_escape" on page 148

util_snprintf

The `util_snprintf` function formats a specified string, using a specified format, into a specified buffer using the `printf`-style syntax and performs bounds checking. It returns the number of characters in the formatted buffer.

For more information, see the documentation on the `printf` function for the runtime library of your compiler.

Syntax

```
int util_snprintf(char *s, int n, char *fmt, ...);
```

Returns

The number of characters formatted into the buffer.

Parameters

`char *s` is the buffer to receive the formatted string.

`int n` is the maximum number of bytes allowed to be copied.

`char *fmt` is the format string. The function handles only `%d` and `%s` strings; it does not handle any width or precision strings.

`...` represents a sequence of parameters for the `printf` function.

See Also

"util_sprintf" on page 145 , "util_vsnprintf" on page 150, "util_vsprintf" on page 151

util_sprintf

The `util_sprintf` function formats a specified string, using a specified format, into a specified buffer, using the `printf`-style syntax without bounds checking. It returns the number of characters in the formatted buffer.

Because `util_sprintf` does not perform bounds checking, use this function only if you are certain that the string fits the buffer. Otherwise, use the function `util_snprintf`. For more information, see the documentation on the `printf` function for the runtime library of your compiler.

Syntax

```
int util_sprintf(char *s, char *fmt, ...);
```

Returns

The number of characters formatted into the buffer.

Parameters

`char *s` is the buffer to receive the formatted string.

`char *fmt` is the format string. The function handles only `%d` and `%s` strings; it does not handle any width or precision strings.

`...` represents a sequence of parameters for the `printf` function.

Example

```
char *logmsg;
int len;
logmsg = (char *) MALLOC(256);
len = util_sprintf(logmsg, "%s %s %s\n", ip, method, uri);
```

See Also

"util_snprintf" on page 145, "util_vsnprintf" on page 150, "util_vsprintf" on page 151

util_strcasecmp

The `util_strcasecmp` function performs a comparison of two alphanumeric strings and returns a `-1`, `0`, or `1` to signal which is larger or that they are identical.

The comparison is not case-sensitive.

Syntax

```
int util_strcasecmp(const char *s1, const char *s2);
```

Returns

`1` if `s1` is greater than `s2`.

`0` if `s1` is equal to `s2`.

`-1` if `s1` is less than `s2`.

Parameters

`char *s1` is the first string.

`char *s2` is the second string.

See Also

"util_strncasecmp" on page 147

util_strftime

The `util_strftime` function translates a `tm` structure, which is a structure describing a system time, into a textual representation. It is a thread-safe version of the standard `strftime` function.

Syntax

```
int util_strftime(char *s, const char *format, const struct tm *t);
```

Returns

The number of characters placed into `s`, not counting the terminating NULL character.

Parameters

`char *s` is the string buffer to put the text into. There is no bounds checking, so you must make sure that the buffer is large enough for the text of the date.

`const char *format` is a format string, a bit like a `printf` string in that it consists of text with certain `%x` substrings. You can use the constant `HTTP_DATE_FMT` to create date strings in the standard Internet format. For more information, see the documentation on the `printf` function for the runtime library of your compiler. For more information on time formats, see the *Sun Java System Web Server 7.0 Administrator's Configuration File Reference*.

`const struct tm *t` is a pointer to a calendar time (`tm`) structure, usually created by the function `system_localtime` or `system_gmtime`.

See Also

"system_localtime" on page 128, "system_gmtime" on page 128

util_strncasecmp

The `util_strncasecmp` function performs a comparison of the first `n` characters in the alphanumeric strings and returns a `-1`, `0`, or `1` to signal which is larger or that they are identical.

The function's comparison is not case-sensitive.

Syntax

```
int util_strncasecmp(const char *s1, const char *s2, int n);
```

Returns

`1` if `s1` is greater than `s2`.

`0` if `s1` is equal to `s2`.

`-1` if `s1` is less than `s2`.

Parameters

`char *s1` is the first string.

`char *s2` is the second string.

`int n` is the number of initial characters to compare.

See Also

"util_strcasecmp" on page 146

util_uri_escape

The `util_uri_escape` function converts any special characters in the URI into the URI format (`%XX`, where `XX` is the hexadecimal equivalent of the ASCII character), and returns the escaped string. The special characters are `%?#:+&*"<>`, space, carriage return, and line feed.

Use `util_uri_escape` before sending a URI back to the client.

Syntax

```
char *util_uri_escape(char *d, char *s);
```

Returns

The string (possibly newly allocated) with escaped characters replaced.

Parameters

`char *d` is a string. If `d` is not NULL, the function copies the formatted string into `d` and returns it. If `d` is NULL, the function allocates a properly sized string and copies the formatted special characters into the new string, then returns it.

The `util_uri_escape` function does not check bounds for the parameter `d`. Therefore, if `d` is not NULL, it should be at least three times as large as the string `s`.

`char *s` is the string containing the original unescaped URI.

See Also

"util_uri_is_evil" on page 149, "util_uri_parse" on page 149, "util_uri_unescape" on page 150

util_uri_is_evil

The `util_uri_is_evil` function checks a specified URI for insecure path characters. Insecure path characters include `//`, `/./`, `/../` and `/.`, `/..` (also for Windows `./`) at the end of the URI. Use this function to see if a URI requested by the client is insecure.

Syntax

```
int util_uri_is_evil(char *t);
```

Returns

`1` if the URI is insecure, or `0` if the URI is OK.

Parameters

`char *t` is the URI to be checked.

See Also

"util_uri_parse" on page 149, "util_uri_escape" on page 148

util_uri_parse

The `util_uri_parse` function converts `//`, `/./`, and `/*/../` into `/` in the specified URI (where * is any character other than `/`). You can use this function to convert a URI's bad sequences into valid ones. First, use the function `util_uri_is_evil` to determine whether the function has a bad sequence.

Syntax

```
void util_uri_parse(char *uri);
```

Returns

`void`

Parameters

`char *uri` is the URI to be converted.

See Also

"util_uri_is_evil" on page 149, "util_uri_unescape" on page 150

util_uri_unescape

The `util_uri_unescape` function converts the encoded characters of a URI into their ASCII equivalents. Encoded characters appear as `%XX`, where `XX` is a hexadecimal equivalent of the character.

Note – You cannot use an embedded null in a string, because NSAPI functions assume that a null is the end of the string. Therefore, passing unicode-encoded content through an NSAPI plug-in does not work.

Syntax

```
void util_uri_unescape(char *uri);
```

Returns

`void`

Parameters

`char *uri` is the URI to be converted.

See Also

"util_uri_escape" on page 148"util_uri_is_evil" on page 149, "util_uri_parse" on page 149

util_vsnprintf

The `util_vsnprintf` function formats a specified string, using a specified format, into a specified buffer using the `vprintf`-style syntax and performs bounds checking. It returns the number of characters in the formatted buffer.

For more information, see the documentation on the `printf` function for the runtime library of your compiler.

Syntax

```
int util_vsnprintf(char *s, int n, register char *fmt, va_list args);
```

Returns

The number of characters formatted into the buffer.

Parameters

`char *s` is the buffer to receive the formatted string.

`int n` is the maximum number of bytes allowed to be copied.

`register char *fmt` is the format string. The function handles only `%d` and `%s` strings; it does not handle any width or precision strings.

`va_list args` is an STD argument variable obtained from a previous call to `va_start`.

See Also

"util_sprintf" on page 145, "util_vsprintf" on page 151

util_vsprintf

The `util_vsprintf` function formats a specified string, using a specified format, into a specified buffer using the `vprintf`-style syntax without bounds checking. It returns the number of characters in the formatted buffer.

For more information, see the documentation on the `printf` function for the runtime library of your compiler.

Syntax

```
int util_vsprintf(char *s, register char *fmt, va_list args);
```

Returns

The number of characters formatted into the buffer.

Parameters

`char *s` is the buffer to receive the formatted string.

`register char *fmt` is the format string. The function handles only `%d` and `%s` strings; it does not handle any width or precision strings.

`va_list args` is an STD argument variable obtained from a previous call to `va_start`.

See Also

"util_snprintf" on page 145, "util_vsnprintf" on page 150

V

vs_alloc_slot

The `vs_alloc_slot` function allocates a new slot for storing a pointer to data specific to a certain `VirtualServer*`. The returned slot number can be used in subsequent `vs_set_data` and `vs_get_data` calls. The returned slot number is valid for any `VirtualServer*`.

The value of the pointer (which may be returned by a call to "vs_set_data" on page 156) defaults to NULL for every `VirtualServer*`.

Syntax

```
int vs_alloc_slot(void);
```

Returns

A slot number if succeeds, or `-1` if fails.

See Also

"vs_get_data" on page 152, "vs_set_data" on page 156

vs_get_data

The `vs_get_data` function finds the value of a pointer to data for a given `VirtualServer*` and `slot`. The `slot` must be a slot number returned from `vs_alloc_slot` or `vs_set_data`.

Syntax

```
void* vs_get_data(const VirtualServer* vs, int slot);
```

Returns

The value of the pointer previously stored using `vs_set_data` or NULL on failure.

Parameters

`const VirtualServer* vs` represents the virtual server to query the pointer for.

`int slot` is the slot number to retrieve the pointer from.

See Also

"vs_set_data" on page 156, "vs_alloc_slot" on page 151

vs_get_default_httpd_object

The `vs_get_default_httpd_object` function obtains a pointer to the default (or root) `httpd_object` from the virtual server's `httpd_objset` (in the configuration defined by the `obj.conf` file of the virtual server class). The default object is typically named `default`. Plug-ins may only modify the `httpd_object` at `VSInitFunc` time (see "vs_register_cb" on page 156 for an explanation of `VSInitFunc` time).

Do not FREE the returned object.

Syntax

```
httpd_object* vs_get_default_httpd_object(VirtualServer* vs);
```

Returns

A pointer the default `httpd_object`, or NULL on failure. Do not FREE this object.

Parameters

`VirtualServer*` `vs` represents the virtual server for which to find the default object.

See Also

"vs_get_httpd_objset" on page 154, "vs_register_cb" on page 156

vs_get_doc_root

The `vs_get_doc_root` function finds the document root for a virtual server. The returned string is the full operating system path to the document root.

The caller should FREE the returned string when done with it.

Syntax

```
char* vs_get_doc_root(const VirtualServer* vs);
```

Returns

A pointer to a string representing the full operating system path to the document root. It is the caller's responsibility to FREE this string.

Parameters

`const VirtualServer*` vs represents the virtual server for which to find the document root.

vs_get_httpd_objset

The `vs_get_httpd_objset` function obtains a pointer to the `httpd_objset` (the configuration defined by the `obj.conf` file of the virtual server class) for a given virtual server. Plug-ins may only modify the `httpd_objset` at `VSInitFunc` time (see "vs_register_cb" on page 156 for an explanation of `VSInitFunc` time).

Do not FREE the returned `objset`.

Syntax

```
httpd_objset* vs_get_httpd_objset(VirtualServer* vs);
```

Returns

A pointer to the `httpd_objset`, or NULL on failure. Do not FREE this `objset`.

Parameters

`VirtualServer*` `vs` represents the virtual server for which to find the `objset`.

See Also

"vs_get_default_httpd_object" on page 153, "vs_register_cb" on page 156

vs_get_id

The `vs_get_id` function finds the ID of a `VirtualServer*`.

The ID of a virtual server is a unique null-terminated string that remains constant across configurations. Note that while IDs remain constant across configurations, the value of `VirtualServer*` pointers do not.

Do not FREE the virtual server ID string. If called during request processing, the string will remain valid for the duration of the current request. If called during `VSInitFunc` processing, the string will remain valid until after the corresponding `VSDestroyFunc` function has returned (see "vs_register_cb" on page 156).

To retrieve a `VirtualServer*` that is valid only for the current request, use "request_get_vs" on page 114.

Syntax

```
const char* vs_get_id(const VirtualServer* vs);
```

Returns

A pointer to a string representing the virtual server ID. Do not FREE this string.

Parameters

`const VirtualServer*` vs represents the virtual server of interest.

See Also

"vs_register_cb" on page 156, "request_get_vs" on page 114

vs_get_mime_type

The `vs_get_mime_type` function determines the MIME type that would be returned in the `content-type:` header for the given URI.

The caller should FREE the returned string when done with it.

Syntax

```
char* vs_get_mime_type(const VirtualServer* vs, const char* uri);
```

Returns

A pointer to a string representing the MIME type. It is the caller's responsibility to FREE this string.

Parameters

`const VirtualServer* vs` represents the virtual server of interest.

`const char* uri` is the URI whose MIME type is of interest.

vs_lookup_config_var

The `vs_lookup_config_var` function finds the value of a configuration variable for a given virtual server.

Do not FREE the returned string.

Syntax

```
const char* vs_lookup_config_var(const VirtualServer* vs, const char* name);
```

Returns

A pointer to a string representing the value of variable name on success, or NULL if variable name was not found. Do not FREE this string.

Parameters

`const VirtualServer*` vs represents the virtual server of interest.

`const char*` `name` is the name of the configuration variable.

vs_register_cb

The `vs_register_cb` function allows a plug-in to register functions that will receive notifications of virtual server initialization and destruction events. The `vs_register_cb` function would typically be called from an `Init` SAF in `magnus.conf`.

When a new configuration is loaded, all registered `VSInitFunc` (virtual server initialization) callbacks are called for each of the virtual servers before any requests are served from the new configuration. `VSInitFunc` callbacks are called in the same order they were registered; that is, the first callback registered is the first called.

When the last request has been served from an old configuration, all registered `VSDestroyFunc` (virtual server destruction) callbacks are called for each of the virtual servers before any virtual servers are destroyed. `VSDestroyFunc` callbacks are called in reverse order; that is, the first callback registered is the last called.

Either `initfn` or `destroyfn` may be NULL if the caller is not interested in callbacks for initialization or destruction, respectively.

Syntax

```
int vs_register_cb(VSInitFunc* initfn, VSDestroyFunc* destroyfn);
```

Returns

The constant `REQ_PROCEED` if the operation succeeds.

The constant `REQ_ABORTED` if the operation fails.

Parameters

`VSInitFunc*` `initfn` is a pointer to the function to call at virtual server initialization time, or NULL if the caller is not interested in virtual server initialization events.

`VSDestroyFunc*` `destroyfn` is a pointer to the function to call at virtual server destruction time, or NULL if the caller is not interested in virtual server destruction events.

vs_set_data

The `vs_set_data` function sets the value of a pointer to data for a given virtual server and slot. The `*slot` must be `-1` or a slot number returned from `vs_alloc_slot`. If `*slot` is `-1`, `vs_set_data` calls `vs_alloc_slot` implicitly and returns the new slot number in `*slot`.

Note that the stored pointer is maintained on a per-`VirtualServer*` basis, not a per-ID basis. Distinct `VirtualServer*`s from different configurations might exist simultaneously with the same virtual server IDs. However, since these are distinct `VirtualServer*`s, they each have their own `VirtualServer*`-specific data. As a result, `vs_set_data` should generally not be called outside of `VSInitFunc` processing (see "vs_register_cb" on page 156 for an explanation of `VSInitFunc` processing).

Syntax

```
void* vs_set_data(const VirtualServer* vs, int* slot, void* data);
```

Returns

Data on success, or NULL on failure.

Parameters

`const VirtualServer* vs` represents the virtual server to set the pointer for.

`int* slot` is the slot number to store the pointer at.

`void* data` is the pointer to store.

See Also

"vs_get_data" on page 152, "vs_alloc_slot" on page 151, "vs_register_cb" on page 156

vs_translate_uri

The `vs_translate_uri` function translates a URI as though it were part of a request for a specific virtual server. The returned string is the full operating system path.

The caller should FREE the returned string when done with it.

Syntax

```
char* vs_translate_uri(const VirtualServer* vs, const char* uri);
```

Returns

A pointer to a string representing the full operating system path for the given URI. It is the caller\qs responsibility to FREE this string.

Parameters

`const VirtualServer* vs` represents the virtual server for which to translate the URI.

`const char* uri` is the URI to translate to an operating system path.

W

write

The `write` filter method is called when output data is to be sent. Filters that modify or consume outgoing data should implement the `write` filter method.

Upon receiving control, a write implementation should first process the data as necessary, and then pass it on to the next filter layer; for example, by calling `net_write(layer->lower, ...,)`. If the filter buffers outgoing data, it should implement the `flush` filter method.

Syntax

```
int write(FilterLayer *layer, const void *buf, int amount);
```

Returns

The number of bytes consumed, which may be less than the requested amount if an error occurred.

Parameters

`FilterLayer *layer` is the filter layer in which the filter is installed.

`const void *buf` is the buffer that contains the outgoing data.

`int amount` is the number of bytes in the buffer.

Example

```
int myfilter_write(FilterLayer *layer, const void *buf, int amount)
{
    return net_write(layer->lower, buf, amount);
}
```

See Also

"flush" on page 78, "net_write" on page 87, "writev" on page 159, "filter_create" on page 73

writev

The `writev` filter method is called when multiple buffers of output data are to be sent. Filters that modify or consume outgoing data may choose to implement the `writev` filter method.

If a filter implements the `write` filter method but not the `writev` filter method, the server automatically translates `net_writev` calls to `net_write` calls. As a result, filters interested in the outgoing data stream do not need to implement the `writev` filter method. However, for performance reasons, it is beneficial for filters that implement the `write` filter method to also implement the `writev` filter method.

Syntax

```
int writev(FilterLayer *layer, const struct iovec *iov, int iov_size);
```

Returns

The number of bytes consumed, which may be less than the requested amount if an error occurred.

Parameters

`FilterLayer *layer` is the filter layer the filter is installed in.

`const struct iovec *iov` is an array of `iovec` structures, each of which contains outgoing data.

`int iov_size` is the number of `iovec` structures in the `iov` array.

Example

```
int myfilter_writev(FilterLayer *layer, const struct iovec *iov, int iov_size)
{
    return net_writev(layer->lower, iov, iov_size);
}
```

See Also

"flush" on page 78, "net_write" on page 87, "write" on page 158, "filter_create" on page 73

◆ ◆ ◆ CHAPTER 6

Data Structure Reference

NSAPI uses many data structures that are defined in the `nsapi.h` header file, which is in the *install_dir*`/include` directory.

This chapter describes public data structures in `nsapi.h`.

Note – The data structures in `nsapi.h` that are not described in this chapter are considered private and may change incompatibly in future releases. Some of the data structures described in this chapter may contain additional, undocumented fields. These fields are also considered private and may change incompatibly in future releases. Additional fields may be added in future release, so do not make assumptions regarding the size of data structures.

This chapter has the following sections:

Public Data Structures

This section describes the following data structures in `nsapi.h`.

Session

A session is the time between the opening and closing of the connection between the client and the server.

The following list describes the most important fields in this data structure:

- `sn->client`

  Pointer to a `pblock` containing information about the client such as its IP address, DNS name, or certificate.

- `sn->csd`

  Platform-independent client socket descriptor. This is passed to the routines for reading from and writing to the client.

The `Session` data structure holds variables that apply to a client, regardless of the requests being sent.

```
typedef struct {
    /* Information about the remote client */
    pblock *client;

    /* The socket descriptor to the remote client */
    SYS_NETFD csd;

    /* The input buffer for that socket descriptor */
    netbuf *inbuf;
} Session;
```

The following list describes the most important fields in the `Session` data structure:

- `client` - Pointer to a `pblock` containing information about the client such as its IP address, DNS name, or SSL certificate. The `ip` parameter contains the client's IP address. Do not modify the contents of this `pblock`.
- `csd` - The platform-independent client socket descriptor used to communicate with the client. This can be passed to routines such as `net_write` to send output to the client.
- `inbuf` - Pointer to the input buffer for the client socket descriptor. This can be passed to routines such as `netbuf_grab` or `netbuf_getc` to receive input from the client.

Note – The `Session` NSAPI data structure cannot be used concurrently by multiple threads. It is an error to retain any reference to a `Session` or its contents after processing of the current request is complete.

pblock

The parameter block is the hash table that holds `pb_entry` structures. Its contents are transparent to most code. This data structure is frequently used in NSAPI; it provides the basic mechanism for packaging up parameters and values. There are many functions for creating and managing parameter blocks, and for extracting, adding, and deleting entries. See the functions whose names start with `pblock_` in Chapter 5. You do not need to write code that accesses `pblock` data fields directly.

```
typedef struct {
    int hsize;
    struct pb_entry **ht;
} pblock;
```

Note – The `pblock` NSAPI data structure can not be used concurrently by multiple threads. It is an error to retain any reference to a `pblock` or its contents after processing of the current request is complete.

pb_entry

The `pb_entry` is a single element in the parameter block.

```
struct pb_entry {
    pb_param *param;
    struct pb_entry *next;
};
```

pb_param

The `pb_param` represents a name-value pair, as stored in a `pb_entry`.

```
typedef struct {
    char *name,*value;
} pb_param;
```

Request

The `Request` data structure describes an HTTP transaction (for example, the variables include the client's HTTP request headers).

```
typedef struct{
    */Server working variables */
    pblock *vars;

    /* The method, URI, and protocol revision of this request */
    pblock *reqpb;

    /* Protocol specific headers */
    int loadhdrs;
    pblock *headers;

    /* Server's response headers */
    int senthdrs;
    pblock *srvhdrs;

    /* The object set constructed to fulfill this request */
    httpd_objset *os;
} Request;
```

The following list describes the most important fields in the `Request` data structure:

- `vars`- Pointer to a `pblock` containing information about request-response processing. SAFs may modify the contents of this `pblock` according to the rules established in "Required Behavior of SAFs for Each Directive" on page 31.
- `reqpb` - Pointer to a `pblock` containing information about the client's HTTP request. The `method` parameter contains the HTTP request method, the `uri`parameter contains the path portion of the requested URI, the optional query parameter contains any query string from the requested URI, and the `protocol` parameter contains the HTTP protocol version. Do not modify the contents of this `pblock`.
- `headers` - Pointer to a `pblock` containing the client's HTTP request headers. By convention, all parameter names are lowercase. Do not modify the contents of this `pblock`.
- `senthdrs` - Indicates whether the server has sent HTTP response headers. Service SAFs may set `rq->senthdrs = 1` to prevent the server from sending HTTP response headers.
- `srvhdrs` - Pointer to a `pblock` containing the server's HTTP response headers. By convention, all parameter names are lowercase. SAFs and filters may modify the contents of this `pblock`.

Note – The `Request` NSAPI data structure cannot be used concurrently by multiple threads. It is an error to retain any references to a `Request` or its contents after processing of the current request.

stat

When a program calls the `stat( )` function for a given file, the system returns a structure that provides information about the file. The specific details of the structure should be obtained from the implementation of your platform, but the basic outline of the structure is as follows:

```
struct stat {
    dev_t      st_dev;     /* device of inode */
    inot_t     st_ino;     /* inode number */
    short      st_mode;    /* mode bits */
    short      st_nlink;   /* number of links to file /*
    short      st_uid;     /* owner's user id */
    short      st_gid;     /* owner's group id */
    dev_t      st_rdev;    /* for special files */
    off_t      st_size;    /* file size in characters */
    time_t     st_atime;   /* time last accessed */
    time_t     st_mtime;   /* time last modified */
    time_t     st_ctime;   /* time inode last changed*/
}
```

The elements that are most significant for server plug-in API activities are `st_size`, `st_atime`, `st_mtime`, and `st_ctime`.

shmem_s

```
typedef struct {
    void       *data;   /* the data */
    HANDLE     fdmap;
    int        size;    /* the maximum length of the data */
    char       *name;   /* internal use: filename to unlink if exposed */
    SYS_FILE   fd;      /* internal use: file descriptor for region */
} shmem_s;
```

cinfo

The `cinfo` data structure records the content information for a file.

```
typedef struct {
    char    *type;
            /* Identifies what kind of data is in the file*/
    char    *encoding;
             /* encoding identifies any compression or other /*
            /* content-independent transformation that's been /*
            /* applied to the file, such as uuencode)*/
    char    *language;
            /* Identifies the language a text document is in. */
} cinfo;
```

sendfiledata

The `sendfiledata` data structure is used to pass parameters to the `net_sendfile` function. It is also passed to the `sendfile` method in an installed filter in response to a `net_sendfile` call.

```
typedef struct {
    SYS_FILE fd;         /* file to send */
    size_t offset;       /* offset in file to start sending from */
    size_t len;          /* number of bytes to send from file */
    const void *header;  /* data to send before file */
    int hlen;            /* number of bytes to send before file */
    const void *trailer; /* data to send after file */
    int tlen;            /* number of bytes to send after file */
} sendfiledata;
```

Filter

The `Filter` data structure is an opaque representation of a filter. A `Filter` structure is created by calling "filter_create" on page 73.

```
typedef struct Filter Filter;
```

FilterContext

The `FilterContext` data structure stores the context associated with a particular filter layer. Filter layers are created by calling "filter_insert" on page 75.

Filter developers may use the data member to store filter-specific context information.

```
typedef struct {
    pool_handle_t *pool; /* pool context was allocated from */
    Session *sn;         /* session being processed */
    Request *rq;         /* request being processed */
    void *data;          /* filter-defined private data */
} FilterContext;
```

FilterLayer

The `FilterLayer` data structure represents one layer in a filter stack. The `FilterLayer` structure identifies the filter installed at that layer. It provides pointers to layer-specific context and a filter stack that represents the layer immediately below it in the filter stack.

```
typedef struct {
    Filter *filter; /* the filter at this layer in the filter stack */
    FilterContext *context; /* context for the filter */
    SYS_NETFD lower; /* access to the next filter layer in the stack */
} FilterLayer;
```

FilterMethods

The `FilterMethods` data structure is passed to "filter_create" on page 73 to define the filter methods that a filter supports. Each new `FilterMethods` instance must be initialized with the `FILTER_METHODS_INITIALIZER` macro. For each filter method that a filter supports, the corresponding `FilterMethods` member should point to a function that implements that filter method.

```
typedef struct {
    size_t size;
    FilterInsertFunc *insert;
    FilterRemoveFunc *remove;
    FilterFlushFunc *flush;
    FilterReadFunc *read;
    FilterWriteFunc *write;
    FilterWritevFunc *writev;
    FilterSendfileFunc *sendfile;
} FilterMethods;
```

◆ ◆ ◆ CHAPTER 7

Dynamic Results Caching Functions

The functions described in this chapter allow you to write a results caching plug-in for Sun Java System Web Server. A results caching plug-in, which is a `Service` SAF, caches data, a page, or part of a page in the web server address space, which the Web Server can refresh periodically on demand. An `Init` SAF initializes the callback function that performs the refresh.

A results caching plug-in can generate a page for a request in three parts:

- A header, such as a page banner, which changes for every request
- A body, which changes less frequently
- A footer, which also changes for every request

  Without this feature, a plug-in would have to generate the whole page for every request (unless an `IFRAME` is used, where the header or footer is sent in the first response along with an `IFRAME` pointing to the body; in this case the browser must send another request for the `IFRAME`).

  If the body of a page has not changed, the plug-in needs to generate only the header and footer and to call the `dr_net_write` function (instead of `net_write`) with the following arguments:

- header
- footer
- handle to cache
- key to identify the cached object

  The web server constructs the whole page by fetching the body from the cache. If the cache has expired, it calls the refresh function and sends the refreshed page back to the client.

  An `Init` SAF that is visible to the plug-in creates the handle to the cache. The `Init` SAF must pass the following parameters to the `dr_cache_init` function:

- `RefreshFunctionPointer`
- `FreeFunctionPointer`

- `KeyComparatorFunctionPtr`
- `RefreshInterval`

  The `RefreshInterval` value must be a `PRIntervalTime` type. For more information, see the NSPR reference at:

  ```
  http://www.mozilla.org/projects/nspr/reference/html/index.html
  ```

  As an alternative, if the body is a file that is present in a directory within the web server system machine, the plug-in can generate the header and footer and call the `fc_net_write` function along with the file name.

  This chapter lists the most important functions a results caching plug-in can use. For more information, see the following file:

  install_dir`/include/drnsapi.h`

This chapter has the following sections:

Functions

This section describes the dynamic result cache functions.

dr_cache_destroy

The `dr_cache_destroy` function destroys and frees resources associated with a previously created and used cache handle. This handle cannot be used in subsequent calls to any of the above functions unless another `dr_cache_init` is performed.

Syntax

```
void dr_cache_destroy(DrHdl *hdl);
```

Parameters

`DrHdl *hdl` is a pointer to a previously initialized handle to a cache (see `dr_cache_init`).

Returns

`void`

Example

```
dr_cache_destroy(&myHdl);
```

dr_cache_init

The `dr_cache_init` function creates a persistent handle to the cache, or NULL on failure. It is called by an `Init` SAF.

Syntax

```
PRInt32 dr_cache_init(DrHdl *hdl, RefreshFunc_t ref, FreeFunc_t fre,
                      CompareFunc_t cmp, PRUint32 maxEntries,
                      PRIntervalTime maxAge);
```

Returns

`1` if successful.

`0` if an error occurs.

Parameters

The following table describes parameters for the `dr_cache_init` function.

TABLE 7–1 dr_cache_init parameters

Parameter	Description
`DrHdl hdl`	Pointer to an unallocated handle.
`RefreshFunc_t ref`	pointer to a cache refresh function. This can be NULL; see the `DR_CHECK` flag and `DR_EXPIR` return value for `dr_net_write`.
`FreeFunc_t fre`	Pointer to a function that frees an entry.
`CompareFunc_t cmp`	Pointer to a key comparator function.
`PRUint32 maxEntriesp`	Maximum number of entries possible in the cache for a given `hdl`.
`PRIntervalTime maxAgep`	The maximum amount of time that an entry is valid. If `0`, the cache never expires.

Example

```
if(!dr_cache_init(&hdl, (RefreshFunc_t)FnRefresh, (FreeFunc_t)FnFree,
  (CompareFunc_t)FnCompare, 150000, PR_SecondsToInterval(7200)))
{
    ereport(LOG_FAILURE, "dr_cache_init() failed");
    return(REQ_ABORTED);
}
```

dr_cache_refresh

The `dr_cache_refresh` function provides a way to refresh a cache entry when the plug-in requires it. This can be achieved by passing NULL for the `ref` parameter in `dr_cache_init` and by passing `DR_CHECK` in a `dr_net_write` call. If `DR_CHECK` is passed to `dr_net_write` and it returns with `DR_EXPIR`, the plug-in should generate new content in the entry and call `dr_cache_refresh` with that entry before calling `dr_net_write` again to send the response.

The plug-in may simply decide to replace the cached entry even if it has not expired (based on some other business logic). The `dr_cache_refresh` function is useful in this case. This way the plug-in does the cache refresh management actively by itself.

Syntax

```
PRInt32 dr_cache_refresh(DrHdl hdl, const char *key,
                         PRUint32 klen, PRIntervalTime timeout,
                         Entry *entry, Request *rq, Session *sn);
```

Returns

`1` if successful.

`0` if an error occurs.

Parameters

The following table describes parameters for the `dr_cache_refresh` function.

TABLE 7–2 dr_cache_refresh parameters

Parameter	Description
`DrHdl hdl`	Persistent handle created by the `dr_cache_init` function.
`const char *key`	Key to cache, search, or refresh.
`PRUint32 klen`	Length of the key in bytes.

TABLE 7–2 dr_cache_refresh parameters *(Continued)*

Parameter	Description
`PRIntervalTime timeout`	Expiration time of this entry; if a value of `0` is passed, the `maxAge` value passed to `dr_cache_init` is used.
`Entry *entry`	The not NULL entry to be cached.
`Request *rq`	Pointer to the request.
`Session *sn`	Pointer to the session.

Example

```
Entry entry;
char *key = "MOVIES"
GenNewMovieList(&entry.data, &entry.dataLen);  // Implemented by
                                               // plugin developer
if(!dr_cache_refresh(hdl, key, strlen(key), 0, &entry, rq, sn))
{
    ereport(LOG_FAILURE, "dr_cache_refresh() failed");
    return REQ_ABORTED;
}
```

dr_net_write

The `dr_net_write` function sends a response back to the requestor after constructing the full page with `hdr`, the content of the cached entry as the body (located using the `key`), and `ftr`. The `hdr`, `ftr`, or `hdl` can be NULL, but not all of them can be NULL. If `hdl` is NULL, no cache lookup is done; the caller must pass `DR_NONE` as the flag.

By default, this function refreshes the cache entry if it has expired by making a call to the `ref` function passed to `dr_cache_init`. If no cache entry is found with the specified `key`, this function adds a new cache entry by calling the `ref` function before sending out the response. However, if the `DR_CHECK` flag is passed in the `flags` parameter and if either the cache entry has expired or the cache entry corresponding to the `key` does not exist, `dr_net_write` does not send any data out. Instead, it returns with `DR_EXPIR`.

If `ref` (passed to `dr_cache_init`) is NULL, the `DR_CHECK` flag is not passed in the `flags` parameter, and the cache entry corresponding to the `key` has expired or does not exist, then `dr_net_write` fails with `DR_ERROR`. However, `dr_net_write` refreshes the cache if `ref` is not NULL and `DR_CHECK` is not passed.

If `ref` (passed to `dr_cache_init`) is NULL and the `DR_CHECK` flag is not passed but `DR_IGNORE` is passed and the entry is present in the cache, `dr_net_write` sends out the response even if the entry has expired. However, if the entry is not found, `dr_net_write` returns `DR_ERROR`.

If ref (passed to dr_cache_init) is not NULL and the DR_CHECK flag is not passed but DR_IGNORE is passed and the entry is present in the cache, dr_net_write sends out the response even if the entry has expired. However, if the entry is not found, dr_net_write calls the ref function and stores the new entry returned from ref before sending out the response.

Syntax

```
PRInt32 dr_net_write(DrHdl hdl, const char *key, PRUint32 klen, const char *hdr,
                     const char *ftr, PRUint32 hlen, PRUint32 flen,
                     PRIntervalTime timeout, PRUint32 flags,
                     Request *rq, Session *sn);
```

Returns

IO_OKAY if successful.

IO_ERROR if an error occurs.

DR_ERROR if an error in cache handling occurs.

DR_EXPIR if the cache has expired.

Parameters

The following table describes parameters for the dr_net_write function.

TABLE 7–3 dr_net_write parameters

Parameter	Description
DrHdl hdl	Persistent handle created by the dr_cache_init function.
const char *key	Key to cache, search, or refresh.
PRUint32 klen	Length of the key in bytes.
const char *hdr	Any header data (which can be NULL).
const char *ftr	Any footer data (which can be NULL).
PRUint32 hlen	Length of the header data in bytes (which can be 0).
PRUint32 flen	Length of the footer data in bytes (which can be 0).
PRIntervalTime timeout	Timeout before this function aborts.
PRUint32 flags	ORed directives for this function (see the Flags table, below).
Request *rq	Pointer to the request.
Session *sn	Pointer to the session.

Flags

The following table describes flags for `dr_net_write`.

TABLE 7–4 Flags for dr_net_write

Flag	Description
`DR_NONE`	Specifies that no cache is used, so the function works as `net_write` does; `DrHdl` can be NULL.
`DR_FORCE`	Forces the cache to refresh, even if it has not expired.
`DR_CHECK`	Returns `DR_EXPIR` if the cache has expired; if the calling function has not provided a refresh function and this flag is not used, `DR_ERROR` is returned.
`DR_IGNORE`	Ignores cache expiration and sends out the cache entry even if it has expired.
`DR_CNTLEN`	Supplies the `Content-Length` header and does a `PROTOCOL_START_RESPONSE`.
`DR_PROTO`	Does a `PROTOCOL_START_RESPONSE`.

Example

```
if(dr_net_write(Dr, szFileName, iLenK, NULL, NULL, 0, 0, 0,
                DR_CNTLEN | DR_PROTO, rq, sn) == IO_ERROR)
{
    return(REQ_EXIT);
}
```

fc_open

The `fc_open` function returns a pointer to `PRFileDesc` that refers to an open file (`fileName`). The `fileName` must be the full path name of an existing file. The file is opened in read mode only. The application calling this function should not modify the currency of the file pointed to by the `PRFileDesc *` unless the `DUP_FILE_DESC` is also passed to this function. In other words, the application (at minimum) should not issue a read operation based on this pointer that would modify the currency for the `PRFileDesc *`. If such a read operation is required (that may change the currency for the `PRFileDesc *`), then the application should call this function with the argument `DUP_FILE_DESC`.

On a successful call to this function, a valid pointer to `PRFileDesc` is returned and the handle '`FcHdl`' is properly initialized. The size information for the file is stored in the '`fileSize`' member of the handle.

Syntax

```
PRFileDesc *fc_open(const char *fileName, FcHdl *hDl,PRUint32 flags, Session *sn, Request *rq);
```

Returns

Pointer to `PRFileDesc`, or `NULL` on failure.

Parameters

`const char *fileName` is the full path name of the file to be opened.

`FcHdl*hDl` is a valid pointer to a structure of type `FcHdl`.

PRUint32 flags can be 0 or `DUP_FILE_DESC`.

`Session *sn` is a pointer to the session.

`Request *rq` is a pointer to the request.

fc_close

The `fc_close` function closes a file opened using `fc_open`. This function should only be called with files opened using `fc_open`.

Syntax

```
void fc_close(PRFileDesc *fd, FcHdl *hDl;
```

Returns

`void`

Parameters

`PRFileDesc *fd` is a valid pointer returned from a prior call to `fc_open`.

`FcHdl *hDl` is a valid pointer to a structure of type `FcHdl`. This pointer must have been initialized by a prior call to `fc_open`.

fc_net_write

The `fc_net_write` function is used to send a header and/or footer and a file that exists somewhere in the system. The `fileName` should be the full path name of a file.

Syntax

```
PRInt32 fc_net_write(const char *fileName, const char *hdr,
                     const char *ftr, PRUint32 hlen,
                     PRUint32 flen, PRUint32 flags,
                     PRIntervalTime timeout, Session *sn, Request *rq);
```

Returns

`IO_OKAY` if successful.

`IO_ERROR` if an error occurs.

`FC_ERROR` if an error in file handling occurs.

Parameters

The following table describes parameters for the `fc_net_write` function.

TABLE 7–5 fc_net_write parameters

Parameter	Description
`const char *fileName`	File to be inserted.
`const char *hdr`	Any header data (which can be NULL).
`const char *ftr`	Any footer data (which can be NULL).
`PRUint32 hlen`	Length of the header data in bytes (which can be 0).
`PRUint32 flen`	Length of the footer data in bytes (which can be 0).
`PRUint32 flags`	ORed directives for this function (see the Flags table, below).
`PRIntervalTime timeout`	Timeout before this function aborts.
`Request *rq`	Pointer to the request.
`Session *sn`	Pointer to the session.

Flags

The following table describes flags for `fc_net_write`.

TABLE 7–6 Flags for fc_net_write

Flag	Description
`FC_CNTLEN`	Supplies the `Content-Length` header and does a `PROTOCOL_START_RESPONSE`.
`FC_PROTO`	Does a `PROTOCOL_START_RESPONSE`.

Example

```
const char *fileName = "/docs/myads/file1.ad";
char *hdr = GenHdr(); // Implemented by plugin
char *ftr = GenFtr(); // Implemented by plugin

if(fc_net_write(fileName, hdr, ftr, strlen(hdr), strlen(ftr),
    FC_CNTLEN, PR_INTERVAL_NO_TIMEOUT, sn, rq) != IO_OKEY)
{
    ereport(LOG_FAILURE, "fc_net_write() failed");
    return REQ_ABORTED;
}
```

◆ ◆ ◆ APPENDIX A

Hypertext Transfer Protocol

The Hypertext Transfer Protocol (HTTP) is a protocol (a set of rules that describes how information is exchanged) that allows a client (such as a web browser) and a web server to communicate with each other.

HTTP is based on a request-response model. The browser opens a connection to the server and sends a request to the server. The server processes the request and generates a response, which it sends to the browser. The server then closes the connection.

This chapter provides a short introduction to a few HTTP basics. For more information on HTTP, see the IETF home page at:

`http://www.ietf.org/home.html`

This chapter has the following sections:

Compliance

Sun Java System Web Server supports HTTP/1.1. The server is conditionally compliant with the HTTP/1.1 proposed standard, as approved by the Internet Engineering Steering Group (IESG), and the Internet Engineering Task Force (IETF) HTTP working group.

For more information on the criteria for being conditionally compliant, see the Hypertext Transfer Protocol -- HTTP/1.1 specification (RFC 2616) at:
`http://www.ietf.org/rfc/rfc2616.txt`

Requests

A request from a browser to a server includes the following information:

Request Method, URI, and Protocol Version

A browser can request information using a number of methods. The commonly used methods are:

- `GET` -- Requests the specified resource, such as a document or image
- `HEAD` -- Requests only the header information for the document
- `POST` -- Requests that the server accept some data from the browser, such as form input for a CGI program
- `PUT` -- Replaces the contents of a server's document with data from the browser

Request Headers

The browser can send headers to the server. Most of these request headers are optional.

The following table lists some of the commonly used request headers.

TABLE A–1 Common Request Headers

Request Header	Description
`Accept`	File types the browser can accept.
`Authorization`	Used if the browser wants to authenticate itself with a server; information such as the user name and password are included.
`User-Agent`	Name and version of the browser software.
`Referer`	URL of the document.
`Host`	Internet host and port number of the resource being requested.

Request Data

If the browser has made a `POST` or `PUT` request, it sends data after the blank line following the request headers. If the browser sends a `GET` or `HEAD` request, there is no data to send.

Responses

The server's response includes the following:

HTTP Protocol Version, Status Code, and Reason Phrase

The server sends back a status code, which is a three-digit numeric code. The five categories of status codes are:

- `100-199` a provisional response.
- `200-299` a successful transaction.
- `300-399` the requested resource should be retrieved from a different location.
- `400-499` an error was caused by the browser.
- `500-599` a serious error occurred in the server.

The following table lists some common status codes.

TABLE A–2 Common HTTP Status Codes

Status Code	Meaning
`200`	OK; request has succeeded for the method used (`GET`, `POST`, `HEAD`).
`201`	The request has resulted in the creation of a new resource reference by the returned URI.
`206`	The server has sent a response to byte range requests.
`302`	Found. Redirection to a new URL. The original URL has moved. This is not an error; most browsers will get the new page.
`304`	Use a local copy. If a browser already has a page in its cache, and the page is requested again, some browsers (such as Netscape Navigator) relay to the web server the "last-modified" timestamp on the browser's cached copy. If the copy on the server is not newer than the browser's copy, the server returns a 304 code instead of returning the page, reducing unnecessary network traffic. This is not an error.
`400`	Sent if the request is not a valid HTTP/1.0 or HTTP/1.1 request. For example HTTP/1.1 requires a host to be specified either in the `Host` header or as part of the URI on the request line.

TABLE A–2 Common HTTP Status Codes *(Continued)*

Status Code	Meaning
401	Unauthorized. The user requested a document but did not provide a valid user name or password.
403	Forbidden. Access to this URL is forbidden.
404	Not found. The document requested is not on the server. This code can also be sent if the server is configured to protect the document for unauthorized personnel.
408	If the client starts a request but does not complete it within the keep-alive timeout configured in the server, then this response will be sent and the connection closed. The request can be repeated with another open connection.
411	The client submitted a `POST` request with chunked encoding, which is of variable length. However, the resource or application on the server requires a fixed length - a `Content-Length` header to be present. This code tells the client to resubmit its request with content-length.
413	Some applications (e.g., certain NSAPI plug-ins) cannot handle very large amounts of data, so returns this error code.
414	The URI is longer than the maximum the web server is willing to serve.
416	Data was requested outside the range of a file.
500	Server error. A server-related error occurred. The server administrator must check the error log in the server.
503	Sent if the quality of service mechanism was enabled and bandwidth or connection limits were attained. The server then serves requests with that code.

Response Headers

The response headers contain information about the server and the response data.

The following table lists some common response headers.

TABLE A–3 Common Response Headers

Response Header	Description
`Server`	Name and version of the web server.
`Date`	Current date (in Greenwich Mean Time).
`Last-Modified`	Date when the document was last modified.
`Expires`	Date when the document expires.
`content-length`	Length of the data that follows (in bytes).

TABLE A–3 Common Response Headers *(Continued)*

Response Header	Description
`content-type`	MIME type of the data that follows.
`WWW-Authenticate`	Used during authentication and includes information that tells the browser software what is necessary for authentication (such as user name and password).

Response Data

The server sends a blank line after the last header. It then sends the response data such as an image or an HTML page.

◆ ◆ ◆ APPENDIX B

Alphabetical List of NSAPI Functions and Macros

This appendix provides an alphabetical list for the easy lookup of NSAPI functions and macros.

NSAPI Functions and Macros

C CALLOC

cinfo_find

condvar_init

condvar_notify

condvar_terminate

condvar_wait

crit_enter

crit_exit

crit_init

crit_terminate

D daemon_atrestart

F fc_close

fc_open

filebuf_buf2sd

filebuf_close

filebuf_getc

filebuf_open

filebuf_open_nostat

filter_find

filter_insert

filter_layer

filter_name

filter_remove

filter-create

flush

FREE

func_exec

func_find

func_insert

I insert

L log_error

M MALLOC

N net_flush

net_ip2host

net_read

net_sendfile

net_write

netbuf_buf2sd

netbuf_getbytes

netbuf_close

netbuf_getc

netbuf_grab

netbuf_open

nsapi_module_init

NSAPI_RUNTIME_VERSION

NSAPI_VERSION

P param_create

param_free

pblock_copy

pblock_create

pblock_dup

pblock_find

pblock_findval

pblock_free

pblock_nninsert

pblock_nvinsert

pblock_pb2env

pblock_pblock2str

pblock_pinsert

pblock_remove

pblock_str2pblock

PERM_CALLOC

PERM_FREE

PERM_MALLOC

PERM_REALLOC

PERM_STRDUP

prepare_nsapi_thread

protocol_dump822

protocol_set_finfo

protocol_start_response

protocol_status

protocol_uri2url

protocol_uri2url_dynamic

R read

REALLOC

remove

request_get_vs

request_header

request_stat_path

request_translate_uri

S sendfile

session_dns

session_maxdns

shexp_casecmp

shexp_cmp

shexp_match

shexp_valid

STRDUP

system_errmsg

system_fclose

system_flock

system_fopenRO

system_fopenRW

system_fopenWA

system_fread

system_fwrite

system_fwrite_atomic

system_gmtime

system_localtime

system_lseek

system_rename

system_ulock

system_unix2local

systhread_attach

systhread_current

systhread_getdata

systhread_newkey

systhread_setdata

systhread_sleep

systhread_start

systhread_timerset

U USE_NSAPI_VERSION

util_can_exec

util_chdir2path

util_chdir2path

util_cookie_find

util_env_find

util_env_free

util_env_replace

util_env_str

util_getline

util_hostname

util_is_mozilla

util_is_url

util_itoa

util_later_than

util_sh_escape

util_snprintf

util_sprintf

util_strcasecmp

util_strftime

util_strncasecmp

util_uri_escape

util_uri_is_evil

util_uri_parse

util_uri_unescape

util_vsnprintf

util_vsprintf

V vs_alloc_slot

vs_get_data

vs_get_default_httpd_object

vs_get_doc_root

vs_get_httpd_objset

vs_get_id

vs_get_mime_type

vs_lookup_config_var

vs_register_cb

vs_set_data

vs_translate_uri

W write

writev

Index

C

D

E

F

G

H

P

Q

R

S

T

U

V

W