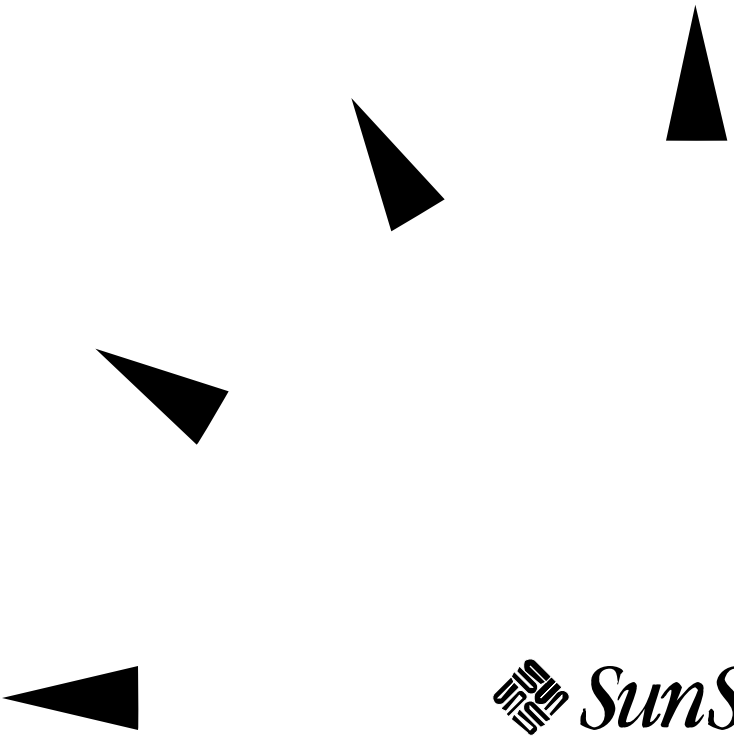
SunOS Reference Manual



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Preface

OVERVIEW

A man page is provided for both the naive user, and sophisticated user who is familiar with the SunOS operating system and is in need of on-line information. A man page is intended to answer concisely the question “What does it do?” The man pages in general comprise a reference manual. They are not intended to be a tutorial.

The following contains a brief description of each section in the man pages and the information it references:

- Section 1 describes, in alphabetical order, commands available with the operating system.

- Section 1M describes, in alphabetical order, commands that are used chiefly for system maintenance and administration purposes.

- Section 2 describes all of the system calls. Most of these calls have one or more error returns. An error condition is indicated by an otherwise impossible returned value.

- Section 3 describes functions found in various libraries, other than those functions that directly invoke UNIX system primitives, which are described in Section 2 of this volume.

- Section 4 outlines the formats of various files. The C structure declarations for the file formats are given where applicable.

- Section 5 contains miscellaneous documentation such as character set tables, etc.

- Section 6 contains available games and demos.

- Section 7 describes various special files that refer to specific hardware peripherals, and device drivers. STREAMS software drivers, modules and the STREAMS-generic set of system calls are also described.

- Section 9 provides reference information needed to write device drivers in the kernel operating systems environment. It describes two device driver interface specifications: the Device Driver Interface (DDI) and the Driver–Kernel Interface (DKI).

- Section 9E describes the DDI/DKI, DDI-only, and DKI-only entry-point routines a developer may include in a device driver.

- Section 9F describes the kernel functions available for use by device drivers.

- Section 9S describes the data structures used by drivers to share information between the driver and the kernel.

Below is a generic format for man pages. The man pages of each manual section generally follow this order, but include only needed headings. For example, if there are no bugs to report, there is no BUGS section. See the intro pages for more information and detail about each section, and **man**(1) for more information about man pages in general.

NAME

This section gives the names of the commands or functions documented, followed by a brief description of what they do.

SYNOPSIS

This section shows the syntax of commands or functions. When a command or file does not exist in the standard path, its full pathname is shown. Literal characters (commands and options) are in **bold** font and variables (arguments, parameters and substitution characters) are in *italic* font. Options and

arguments are alphabetized, with single letter arguments first, and options with arguments next, unless a different argument order is required.

The following special characters are used in this section:

[ ] The option or argument enclosed in these brackets is optional. If the brackets are omitted, the argument *must* be specified.

... Ellipses. Several values may be provided for the previous argument, or the previous argument can be specified multiple times, for example, '*filename* ...'.

| Separator. Only one of the arguments separated by this character can be specified at time.

{ } Braces. The options and/or arguments enclosed within braces are interdependent, such that everything enclosed must be treated as a unit.

PROTOCOL

This section occurs only in subsection 3R to indicate the protocol description file. The protocol specification pathname is always listed in **bold** font.

AVAILABILITY

This section briefly states any limitations on the availabilty of the command. These limitations could be hardware or software specific.

A specification of a class of hardware platform, such as **x86** or **SPARC**, denotes that the command or interface is applicable for the hardware platform specified.

In Section 1 and Section 1M, **AVAILABILITY** indicates which package contains the command being described on the manual page. In order to use the command, the specified package must have been installed with the operating system. If the package was not installed, see **pkgadd**(1) for information on how to upgrade.

MT-LEVEL

This section lists the **MT-LEVEL** of the library functions described in the Section 3 manual pages. The **MT-LEVEL** defines the libraries' ability to support threads. See **Intro**(3) for more information.

DESCRIPTION

This section defines the functionality and behavior of the service. Thus it describes concisely what the command does. It does not discuss OPTIONS or cite EXAMPLES. Interactive commands, subcommands, requests, macros, functions and such, are described under USAGE.

IOCTL

This section appears on pages in Section 7 only. Only the device class which supplies appropriate parameters to the **ioctl**(2) system call is called **ioctl** and generates its own heading. **ioctl** calls for a specific device are listed alphabetically (on the man page for that specific device). **ioctl** calls are used for a particular class of devices all of which have an **io** ending, such as **mtio**(7).

OPTIONS

This lists the command options with a concise summary of what each option does. The options are listed literally and in the order they appear in the SYNOPSIS section. Possible arguments to options are discussed under the option, and where appropriate, default values are supplied.

OPERANDS

This section lists the command operands and describes how they affect the actions of the command.

OUTPUT

This section describes the output - standard output, standard error, or output files - generated by the command.

RETURN VALUES

If the man page documents functions that return values, this section lists these values and describes the conditions under which they are returned. If a function can return only constant values, such as 0 or –1, these values are listed in tagged paragraphs. Otherwise, a single paragraph describes the return values of each function. Functions declared as **void** do not return values, so they are not discussed in RETURN VALUES.

ERRORS

On failure, most functions place an error code in the global variable **errno** indicating why they failed. This section lists alphabetically all error codes a function can generate and describes the conditions that cause each error. When more than one condition can cause the same error, each condition is described in a separate paragraph under the error code.

USAGE

This section is provided as a *guidance* on use. This section lists special rules, features and commands that require in-depth explanations. The subsections listed below are used to explain built-in functionality:

Commands
Modifiers
Variables
Expressions
Input Grammar

EXAMPLES

This section provides examples of usage or of how to use a command or function. Wherever possible a complete example including command line entry and machine response is shown. Whenever an example is given, the prompt is shown as

example%

or if the user must be super-user,

example#

Examples are followed by explanations, variable substitution rules, or returned values. Most examples illustrate concepts from the SYNOPSIS, DESCRIPTION, OPTIONS and USAGE sections.

ENVIRONMENT

This section lists any environment variables that the command or function affects, followed by a brief description of the effect.

EXIT STATUS

This section lists the values the command returns to the calling program or shell and the conditions that cause these values to be returned. Usually, zero is returned for successful completion and values greater than zero for various error conditions.

FILES

This section lists all filenames referred to by the man page, files of interest, and files created or required by commands. Each is followed by a descriptive summary or explanation.

SEE ALSO

This section lists references to other man pages, in-house documentation and outside publications.

DIAGNOSTICS

This section lists diagnostic messages with a brief explanation of the condition causing the error. Messages appear in **bold** font with the exception of variables, which are in *italic* font.

WARNINGS

This section lists warnings about special conditions which could seriously affect your working conditions — this is not a list of diagnostics.

NOTES

This section lists additional information that does not belong anywhere else on the page. It takes the form of an *aside* to the user, covering points of special interest. Critical information is never covered here.

BUGS

This section describes known bugs and wherever possible suggests workarounds.

NAME

FN_attribute_t, fn_attribute_create, fn_attribute_destroy, fn_attribute_copy, fn_attribute_assign, fn_attribute_identifier, fn_attribute_syntax, fn_attribute_valuecount, fn_attribute_first, fn_attribute_next, fn_attribute_add, fn_attribute_remove – an XFN attribute

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_attribute_t ***fn_attribute_create(const FN_identifier_t** ***attribute_id*,
const FN_identifier_t ***attribute_syntax*);

void fn_attribute_destroy(FN_attribute_t ***attr*);

FN_attribute_t ***fn_attribute_copy(const FN_attribute_t** ***attr*);

FN_attribute_t ***fn_attribute_assign(FN_attribute_t** ***dst*,
const FN_attribute_t ***src*);

const FN_identifier_t ***fn_attribute_identifier(const FN_attribute_t***attr*);

const FN_identifier_t ***fn_attribute_syntax(const FN_attribute_t***attr*);

unsigned int fn_attribute_valuecount(const FN_attribute_t*attr*);

const FN_attrvalue_t ***fn_attribute_first(const FN_attribute_t** ***attr*,
void ****iter_pos*);

const FN_attrvalue_t ***fn_attribute_next(const FN_attribute_t** ***attr*,
void ****iter_pos*);

int fn_attribute_add(FN_attribute_t ***attr*, **const FN_attrvalue_t** ***attribute_value*,
unsigned int *exclusive*);

int fn_attribute_remove(FN_attribute_t ***attr*, **const FN_attrvalue_t** ***attribute_value*);

MT-LEVEL

Safe.

DESCRIPTION

An attribute has an attribute identifier, a syntax, and a set of distinct values. Each value is a sequence of octets. The operations associated with objects of type **FN_attribute_t** allow the construction, destruction, and manipulation of an attribute and its value set.

The attribute identifier and its syntax are specified using an **FN_identifier_t**. **fn_attribute_create( )** creates a new attribute object with the given identifier and syntax, and an empty set of values. **fn_attribute_destroy( )** releases the storage associated with *attr*. **fn_attribute_copy( )** returns a copy of the object pointed to by *attr*. **fn_attribute_assign( )** makes a copy of the attribute object pointed to by *src* and assigns it to *dst*, releasing any old contents of *dst*. A pointer to the same object as *dst* is returned.

fn_attribute_identifier( ) returns the attribute identifer of *attr*. **fn_attribute_syntax( )** returns the attribute syntax of *attr*. **fn_attribute_valuecount( )** returns the number of attribute values in *attr*.

fn_attribute_first( ) and **fn_attribute_next( )** are used to enumerate the values of an attribute. Enumeration of the values of an attribute may return the values in any order. **fn_attribute_first( )** returns an attribute value from *attr* and sets the iteration marker *iter_pos*. Subsequent calls to **fn_attribute_next( )** returns the next attribute value identified by *iter_pos* and advances *iter_pos*. Adding or removing values from an attribute invalidates any iteration markers that the caller holds.

fn_attribute_add( ) adds a new value *attribute_value* to *attr*. The operation succeeds (but no change is made) if *attribute_value* is already in *attr* and *exclusive* is zero; the operation fails if *attribute_value* is already in *attr* and *exclusive* is nonzero. **fn_attribute_remove( )** removes *attribute_value* from *attr*. The operation succeeds even if *attribute_value* is not amongst *attr*'s values.

RETURN VALUE

fn_attribute_first( ) returns **0** if the attribute contains no values. **fn_attribute_next( )** returns **0** if there are no more values to be returned in the attribute (as identified by the iteration marker) or if the iteration marker is invalid.

fn_attribute_add( ) and **fn_attribute_remove( )** return **1** if the operation succeeds, **0** if it fails.

APPLICATION USAGE

Manipulation of attributes using the operations described in this manual page does not affect their representation in the underlying naming system. Changes to attributes in the underlying naming system can only be effected through the use of the interfaces described in **xfn_attributes**(3N).

SEE ALSO

FN_attrvalue_t(3N), **FN_attrset_t**(3N), **FN_identifier_t**(3N), **fn_attr_get**(3N), **fn_attr_modify**(3N), **xfn_attributes**(3N), **xfn**(3N)

NAME

FN_attrmodlist_t, fn_attrmodlist_create, fn_attrmodlist_destroy, fn_attrmodlist_copy, fn_attrmodlist_assign, fn_attrmodlist_count, fn_attrmodlist_first, fn_attrmodlist_next, fn_attrmodlist_add – a list of attribute modifications

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_attrmodlist_t ***fn_attrmodlist_create(void);**

void fn_attrmodlist_destroy(FN_attrmodlist_t ****modlist*);**

FN_attrmodlist_t ***fn_attrmodlist_copy(const FN_attrmodlist_t** ****modlist*);**

FN_attrmodlist_t ***fn_attrmodlist_assign(FN_attrmodlist_t** ****dst*,
const FN_attrmodlist_t ****src*);**

unsigned int fn_attrmodlist_count(const FN_attrmodlist_t ****modlist*);**

const FN_attribute_t ***fn_attrmodlist_first(const FN_attrmodlist_t** ****modlist*,
void ***iter_pos*, **unsigned int** ****first_mod_op*);**

const FN_attribute_t ***fn_attrmodlist_next(const FN_attrmodlist_t** ****modlist*,
void ***iter_pos*, **unsigned int** ****mod_op*);**

int fn_attrmodlist_add(FN_attrmodlist_t ****modlist*, **unsigned int** *mod_op*,
const FN_attribute_t ****attr*);**

MT-LEVEL

Safe.

DESCRIPTION

An attribute modification list allows for multiple modification operations to be made on the attributes associated with a single named object. It is used in the **fn_attr_multi_modify( )** operation.

An attribute modification list is a list of attribute modification specifiers. An attribute modification specifier consists of an attribute object and an operation specifier. The attribute's identifier indicates the attribute that is to be operated upon. The attribute's values are used in a manner depending on the operation. The operation specifier is an **unsigned int** that must have one of the values: **FN_ATTR_OP_ADD, FN_ATTR_OP_ADD_EXCLUSIVE, FN_ATTR_OP_REMOVE, FN_ATTR_OP_ADD_VALUES,** or **FN_ATTR_OP_REMOVE_VALUES.** (See **fn_attr_modify( )** for detailed descriptions of these specifiers.) The operations are to be performed in the order in which they appear in the modification list.

fn_attrmodlist_create( ) creates an empty attribute modification list.
fn_attrmodlist_destroy( ) releases the storage associated with *modlist*.
fn_attrmodlist_copy( ) returns a copy of the attribute modification list *modlist*.
fn_attrmodlist_assign( ) makes a copy of *src* and assigns it to *dst*, releasing any old contents of *dst*. It returns a pointer to the same object as *dst*.

fn_attrmodlist_count( ) returns the number attribute modification items in the attribute modification list.

The iterators **fn_attrmodlist_first( )** and **fn_attrmodlist_next( )** return a handle to the attribute part of the modification and return the operation specifier part through an **unsigned int** * parameter. **fn_attrmodlist_first( )** returns the attribute of the first modification item from *modlist* and sets *mod_op* to be the code of the modification operation of that item; *iter_pos* is set after the first modification item. **fn_attrmodlist_next( )** returns the attribute of the next modification item from *modlist* after *iter_pos* and advances *iter_pos*; *mod_op* is set to the code of the modification operation of that item. The order of the items returned during an enumeration is the same as the order by which the items were added to the modification list.

fn_attrmodlist_add( ) adds a new item consisting of the given modification operation code *mod_op* and attribute *attr* to the end of the modification list *modlist*. *attr*'s identifier indicates the attribute that is to be operated upon. *attr*'s values are used in a manner depending on the operation.

RETURN VALUE

fn_attrmodlist_first( ) returns **0** if the modification list is empty. **fn_attrmodlist_next( )** returns **0** if there are no more items on the modification list to be enumerated or if the iteration marker is invalid.

fn_attrmodlist_add( ) returns **1** if the operation succeeds, **0** if the operation fails.

APPLICATION USAGE

Manipulation of attributes using the operations described in this manual page does not affect their representation in the underlying naming system. Changes to attributes in the underlying naming system can only be effected through the use of the interfaces described in **xfn_attributes**(3N).

SEE ALSO

FN_attribute_t(3N), **FN_attrset_t**(3N), **FN_identifier_t**(3N), **fn_attr_multi_modify**(3N), **fn_attr_modify**(3N), **xfn_attributes**(3N), **xfn**(3N)

NAME

FN_attrset_t, fn_attrset_create, fn_attrset_destroy, fn_attrset_copy, fn_attrset_assign, fn_attrset_get, fn_attrset_count, fn_attrset_first, fn_attrset_next, fn_attrset_add, fn_attrset_remove – a set of XFN attributes

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_attrset_t ***fn_attrset_create(void);**

void fn_attrset_destroy(FN_attrset_t ***aset**);

FN_attrset_t ***fn_attrset_copy(const FN_attrset_t** ***aset**);

FN_attrset_t ***fn_attrset_assign(FN_attrset_t** ***dst**, **const FN_attrset_t** ***src**);

const FN_attribute_t ***fn_attrset_get(const FN_attrset_t** ***aset**,
const FN_identifier_t ***attr_id**);

unsigned int fn_attrset_count(const FN_attrset_t ***aset**);

const FN_attribute_t ***fn_attrset_first(const FN_attrset_t** ***aset**,. BI "void " **iter_pos);

const FN_attribute_t ***fn_attrset_next(const FN_attrset_t** ***aset**, **void** ****iter_pos**);

int fn_attrset_add(FN_attrset_t ***aset**, **const FN_attribute_t** ***attr**,
unsigned int *exclusive*);

int fn_attrset_remove(FN_attrset_t ***aset**, **const FN_identifier_t** ***attr_id**);

MT-LEVEL

Safe.

DESCRIPTION

An attribute set is a set of attribute objects with distinct identifiers. The **fn_attr_multi_get( )** operation takes an attribute set as parameter and returns an attribute set. The **fn_attr_get_ids( )** operation returns an attribute set containing the identifiers of the attributes.

Attribute sets are represented by the type **FN_attrset_t**. The following operations are defined for manipulating attribute sets.

fn_attrset_create( ) creates an empty attribute set. **fn_attrset_destroy( )** releases the storage associated with the attribute set *aset*. **fn_attrset_copy( )** returns a copy of the attribute set *aset*. **fn_attrset_assign( )** makes a copy of the attribute set *src* and assigns it to *dst*, releasing any old contents of *dst*. A pointer to the same object as *dst* is returned.

fn_attrset_get( ) returns the attribute with the given identifier *attr_id* from *aset*. **fn_attrset_count( )** returns the number attributes found in the attribute set *aset*.

fn_attrset_first( ) and **fn_attrset_next( )** are functions that can be used to return an enumeration of all the attributes in an attribute set. The attributes are not ordered in any way. There is no guaranteed relation between the order in which items are added to an attribute set and the order of the enumeration. The specification does guarantee that any two enumeration will return the members in the same order, provided that no **fn_attrset_add( )** or **fn_attrset_remove( )** operation was performed on the object in between or during the two enumerations. **fn_attrset_first( )** returns the first attribute

from the set and sets *iter_pos* after the first attribute. **fn_attrset_next( )** returns the attribute following *iter_pos* and advances *iter_pos*.

fn_attrset_add( ) adds the attribute *attr* to the attribute set *aset*, replacing the attribute's values if the identifier of *attr* is not distinct in *aset* and *exclusive* is zero. If *exclusive* is nonzero and the identifier of *attr* is not distinct in *aset*, the operation fails. **fn_attrset_remove( )** removes the attribute with the identifier *attr_id* from *aset*. The operation succeeds even if no such attribute occurs in *aset*.

RETURN VALUE

fn_attrset_first( ) returns **0** if the attribute set is empty. **fn_attrset_next( )** returns **0** if there are no more attributes in the set.

fn_attrset_add( ) and **fn_attrset_remove( )** return **1** if the operation succeeds, and **0** if the operation fails.

APPLICATION USAGE

Manipulation of attributes using the operations described in this manual page does not affect their representation in the underlying naming system. Changes to attributes in the underlying naming system can only be effected through the use of the interfaces described in **xfn_attributes**(3N).

SEE ALSO

FN_attribute_t(3N), **FN_attrvalue_t**(3N), **FN_identifier_t**(3N), **fn_attr_multi_get**(3N), **fn_attr_get_ids**(3N), **xfn_attributes**(3N), **xfn**(3N)

NAME

FN_attrvalue_t – an XFN attribute value

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

DESCRIPTION

The type **FN_attrvalue_t** is used to represent the contents of a single attribute value, within an attribute of type **FN_attribute_t**.

The representation of this structure is defined by XFN as follows:

```
typedef struct {
        size_t length;
        void *contents;
} FN_attrvalue_t;
```

SEE ALSO

FN_attribute_t(3N), **fn_attr_get_values**(3N), **xfn**(3N)

NAME

FN_composite_name_t, fn_composite_name_create, fn_composite_name_destroy, fn_composite_name_from_string, fn_string_from_composite_name, fn_composite_name_copy, fn_composite_name_assign, fn_composite_name_is_empty, fn_composite_name_count, fn_composite_name_first, fn_composite_name_next, fn_composite_name_prev, fn_composite_name_last, fn_composite_name_prefix, fn_composite_name_suffix, fn_composite_name_is_equal, fn_composite_name_is_prefix, fn_composite_name_is_suffix, fn_composite_name_prepend_comp, fn_composite_name_append_comp, fn_composite_name_insert_comp, fn_composite_name_delete_comp, fn_composite_name_prepend_name, fn_composite_name_append_name, fn_composite_name_insert_name – a sequence of component names spanning multiple naming systems

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_composite_name_t ***fn_composite_name_create(void);**

void fn_composite_name_destroy(FN_composite_name_t ***name*);**

FN_composite_name_t ***fn_composite_name_from_string(const FN_string_t** ***str*);**

FN_string_t ***fn_string_from_composite_name(const FN_composite_name_t** ***name*,
unsigned int ***status*);**

FN_composite_name_t ***fn_composite_name_copy(
const FN_composite_name_t** ***name*);**

FN_composite_name_t ***fn_composite_name_assign(FN_composite_name_t** ***dst*,
const FN_composite_name_t ***src*);**

int fn_composite_name_is_empty(const FN_composite_name_t ***name*);**

unsigned int fn_composite_name_count(const FN_composite_name_t ***name*);**

const FN_string_t ***fn_composite_name_first(const FN_composite_name_t** ***name*,
void ***iter_pos*);**

const FN_string_t ***fn_composite_name_next(const FN_composite_name_t** ***name**
void ***iter_pos*);**

const FN_string_t ***fn_composite_name_prev(const FN_composite_name_t** ***name*,
void ***iter_pos*);**

const FN_string_t ***fn_composite_name_last(const FN_composite_name_t** ***name*,
void ***iter_pos*);**

FN_composite_name_t ***fn_composite_name_prefix(
const FN_composite_name_t** ***name*, **const void** ***iter_pos*);**

FN_composite_name_t ***fn_composite_name_suffix(
const FN_composite_name_t** ***name*, **const void** ***iter_pos*);**

int fn_composite_name_is_equal(const FN_composite_name_t ***name*,
const FN_composite_name_t ***name2*, **unsigned int** ***status*);**

int fn_composite_name_is_prefix(const FN_composite_name_t **name*, **const FN_composite_name_t** **prefix*, **void** ***iter_pos*, **unsigned int** **status*);

int fn_composite_name_is_suffix(const FN_composite_name_t **name*, **const FN_composite_name_t** **suffix*, **void** ***iter_pos*, **unsigned int** **status*);

int fn_composite_name_prepend_comp(FN_composite_name_t **name*, **const FN_string_t** **newcomp*);

int fn_composite_name_append_comp(FN_composite_name_t **name*, **const FN_string_t** **newcomp*);

int fn_composite_name_insert_comp(FN_composite_name_t **name*, **void** ***iter_pos*, **const FN_string_t** **newcomp*);

int fn_composite_name_delete_comp(FN_composite_name_t **name*, **void** ***iter_pos*);

int fn_composite_name_prepend_name(FN_composite_name_t **name*, **const FN_composite_name_t** **newcomps*);

int fn_composite_name_append_name(FN_composite_name_t **name*, **const FN_composite_name_t** **newcomps*);

int fn_composite_name_insert_name(FN_composite_name_t **name*, **void** ***iter_pos*, **const FN_composite_name_t** **newcomps*);

MT-LEVEL

Safe.

DESCRIPTION

A composite name is represented by an object of type **FN_composite_name_t**. Each component is a string name, of type **FN_string_t**, from the namespace of a single naming system. It may be an atomic name or a compound name in that namespace.

fn_composite_name_create creates an **FN_composite_name_t** object with zero components. Components may be subsequently added to the composite name using the modify operations described below. **fn_composite_name_destroy** releases any storage associated with the given **FN_composite_name_t** handle.

fn_composite_name_from_string() creates an **FN_composite_name_t** from the string *str* using the XFN composite name syntax. **fn_string_from_composite_name()** returns the standard string form of the given composite name, by concatenating the components of the composite name in a left to right order, each separated by the XFN component separator.

fn_composite_name_copy() returns a copy of the given composite name object. **fn_composite_name_assign()** makes a copy of the composite name object pointed to by *src* and assigns it to *dst*, releasing any old contents of *dst*. A pointer to the same object as *dst* is returned.

fn_composite_name_is_empty() returns **1** if the given composite name is an empty composite name (i.e. consists of a single, empty component name); otherwise, it returns **0**. **fn_composite_name_count()** returns the number of components in the given composite name.

The iteration scheme is based on the exchange of an opaque **void** * argument, *iter_pos*, that serves to record the position of the iteration in the sequence. Conceptually, *iter_pos* records a position between two successive components (or at one of the extreme ends of the sequence).

The function **fn_composite_name_first( )** returns a handle to the **FN_string_t** that is the first component in the name, and sets *iter_pos* to indicate the position immediately following the first component. It returns **0** if the name has no components. Thereafter, successive calls of the **fn_composite_name_next( )** function return pointers to the component following the iteration marker, and advance the iteration marker. If the iteration marker is at the end of the sequence, **fn_composite_name_next( )** returns **0**. Similarly, **fn_composite_name_prev( )** returns the component preceding the iteration pointer and moves the marker back one component. If the marker is already at the beginning of the sequence, **fn_composite_name_prev( )** returns **0**. The function **fn_composite_name_last( )** returns a pointer to the last component of the name and sets the iteration marker immediately preceding this component (so that subsequent calls to **fn_composite_name_prev( )** can be used to step through leading components of the name).

The **fn_composite_name_suffix( )** function returns a composite name consisting of a copy of those components following the supplied iteration marker. The method **fn_composite_name_prefix( )** returns a composite name consisting of those components that precede the iteration marker. Using these functions with an iteration marker that was not initialized using **fn_composite_name_first( ), fn_composite_name_last( ), fn_composite_name_is_prefix( ),** or **fn_composite_name_is_suffix( )** yields undefined and generally undesirable behavior.

The functions **fn_composite_name_is_equal( ), fn_composite_name_is_prefix( ), fn_composite_name_is_suffix( ),** test for equality between composite names or between parts of composite names. For these functions, equality is defined as exact string equality, not name equivalence. A name's syntactic property, such as case-insensitivity, is not taken into account by these functions.

The function **fn_composite_name_is_prefix( )** tests if one composite name is a prefix of another. If so, it returns **1** and sets the iteration marker immediately following the prefix. (So for example, a subsequent call to **fn_composite_name_suffix( )** will return the remainder of the name.) Otherwise it returns **0** and value of the iteration marker is undefined. The function **fn_composite_name_is_suffix( )** is similar. It tests if a one composite name is a suffix of another. If so it returns **1** and sets the iteration marker immediately preceding the suffix.

fn_composite_name_prepend_comp( ) and **fn_composite_name_append_comp( )** prepends and appends a single component to the given composite name, respectively. These operations invalidate any iteration marker the client holds for that object. **fn_composite_name_insert_comp( )** inserts a single component before *iter_pos* to the given composite name and sets *iter_pos* to be immediately after the component just inserted. **fn_composite_name_delete_comp( )** deletes the component located before *iter_pos* from the given composite name and sets *iter_pos* back one component.

The functions **fn_composite_name_prepend_name( ), fn_composite_name_append_name( )** and **fn_composite_name_insert_name( )** perform the same update functions as their *_comp* counterparts, respectively, except that multiple components are being added, rather than single components. For **fn_composite_name_insert_name( )** sets *iter_pos* to be immediately after the name just added.

RETURN VALUE

The functions **fn_composite_name_is_empty( ), fn_composite_name_is_equal( ), fn_composite_name_is_suffix( )** and **fn_composite_name_is_prefix( )** return **1** if the test indicated is true; **0** otherwise.

The update functions **fn_composite_name_prepend_comp( ), fn_composite_name_append_comp( ), fn_composite_name_insert_comp( ), fn_composite_name_delete_comp( )** and their *_name* counterparts return **1** if the update was successful; **0** otherwise.

If a function is expected to return a pointer to an object, a null pointer (**0**) is returned if the function fails.

ERRORS

Code set mismatches that occur during the composition of the string form or during comparisons of composite names are resolved in an implementation-dependent way. **fn_string_from_composite_name( ), fn_composite_name_is_equal( ), fn_composite_name_is_suffix( )** and **fn_composite_name_is_prefix( )** sets *status* to **FN_E_INCOMPATIBLE_CODE_SETS** for composite names whose components have code sets that are determined by the implementation to be incompatible.

SEE ALSO

FN_string_t(3N), **xfn**(3N)

NAME

FN_compound_name_t, fn_compound_name_from_syntax_attrs, fn_compound_name_get_syntax_attrs, fn_compound_name_destroy, fn_string_from_compound_name, fn_compound_name_copy, fn_compound_name_assign, fn_compound_name_count, fn_compound_name_first, fn_compound_name_next, fn_compound_name_prev, fn_compound_name_last, fn_compound_name_prefix, fn_compound_name_suffix, fn_compound_name_is_empty, fn_compound_name_is_equal, fn_compound_name_is_prefix, fn_compound_name_is_suffix, fn_compound_name_prepend_comp, fn_compound_name_append_comp, fn_compound_name_insert_comp, fn_compound_name_delete_comp, fn_compound_name_delete_all – an XFN compound name

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_compound_name_t ***fn_compound_name_from_syntax_attrs(**
const FN_attrset_t ***aset*, **const FN_string_t** ***name*, **FN_status_t** ***status*);

FN_attrset_t ***fn_compound_name_get_syntax_attrs(**
const FN_compound_name_t ***name*);

void fn_compound_name_destroy(FN_compound_name_t ***name*);

FN_string_t ***fn_string_from_compound_name(const FN_compound_name_t** ***name*);

FN_compound_name_t ***fn_compound_name_copy(**
const FN_compound_name_t ***name*);

FN_compound_name_t ***fn_compound_name_assign(FN_compound_name_t** ***dst*,
const FN_compound_name_t ***src*);

unsigned int fn_compound_name_count(const FN_compound_name_t ***name*);

const FN_string_t ***fn_compound_name_first(const FN_compound_name_t** ***name*,
void ****iter_pos*);

const FN_string_t ***fn_compound_name_next(const FN_compound_name_t** ***name*,
void ****iter_pos*);

const FN_string_t ***fn_compound_name_prev(const FN_compound_name_t** ***name*,
void ****iter_pos*);

const FN_string_t ***fn_compound_name_last(const FN_compound_name_t** ***name*,
void ****iter_pos*);

FN_compound_name_t ***fn_compound_name_prefix(**
const FN_compound_name_t ***name*, **const void** ****iter_pos*);

FN_compound_name_t ***fn_compound_name_suffix(**
const FN_compound_name_t ***name*, **const void** ****iter_pos*);

int fn_compound_name_is_empty(const FN_compound_name_t ***name*);

int fn_compound_name_is_equal(" const FN_compound_name_t **name1*, **const FN_compound_name_t** **name2*, **unsigned int** **status*);

int fn_compound_name_is_prefix(const FN_compound_name_t **name*, **const FN_compound_name_t** **pre*, **void** ***iter_pos*, **unsigned int** **status*);

int fn_compound_name_is_suffix(const FN_compound_name_t **name*, **const FN_compound_name_t** **suffix*, **void** ***iter_pos*, **unsigned int** **status*);

int fn_compound_name_prepend_comp(FN_compound_name_t **name*, **const FN_string_t** **atomic_comp*, **unsigned int** **status*);

int fn_compound_name_append_comp(FN_compound_name_t **name*, **const FN_string_t** **atomic_comp*, **unsigned int** **status*);

int fn_compound_name_insert_comp(FN_compound_name_t **name*, **void** ***iter_pos*, **const FN_string_t** **atomic_comp*, **unsigned int** **status*);

int fn_compound_name_delete_comp(FN_compound_name_t **name*, **void** ***iter_pos*);

int fn_compound_name_delete_all(FN_compound_name_t **name*);

MT-LEVEL

Safe.

DESCRIPTION

Most applications treat names as opaque data and hence, the majority of clients of the XFN interface will not need to parse names. Some applications, however, such as browsers, need to parse names. For these applications, XFN provides support in the form of the **FN_compound_name_t** object.

Each naming system in an XFN federation potentially has its own naming conventions. The **FN_compound_name_t** object has associated operations for applications to process compound names that conform to the XFN model of expressing compound name syntax. The XFN syntax model for compound names covers a large number of specific name syntaxes and is expressed in terms of syntax properties of the naming convention. See XFN_compound_syntax.

An **FN_compound_name_t** object is constructed by the operation **fn_compound_name_from_syntax_attrs ,** using a string name and an attribute set containing the "fn_syntax_type" (FN_ID_STRING syntax) attribute identifying the namespace syntax of the string name. The value "standard" (FN_ID_STRING syntax) in the "fn_syntax_type" specifies a syntax model that is by default supported by the **FN_compound_name_t** object. An implementation may support other syntax types instead of the XFN standard syntax model, in which case, the value of the "fn_syntax_type" attribute would be set to an implementation specific string. **fn_compound_name_get_syntax_attrs( )** returns an attribute set containing the syntax attributes that describes the given compound name. **fn_compound_name_destroy( )** releases the storage associated with the given compound name.

fn_string_from_compound_name( ) returns the string form of the given compound name. **fn_compound_name_copy( )** returns a copy of the given compound name. **fn_compound_name_assign( )** makes a copy of the compound name *src* and assigns it to *dst*, releasing any old contents of *dst*. A pointer to the object pointed to by *dst* is returned.

fn_compound_name_count() returns the number of atomic components in the given compound name.

The function **fn_compound_name_first()** returns a handle to the **FN_string_t** that is the first atomic component in the compound name, and sets *iter_pos* to indicate the position immediately following the first component. It returns **0** if the name has no components. Thereafter, successive calls of the **fn_compound_name_next()** function return pointers to the component following the iteration marker, and advance the iteration marker. If the iteration marker is at the end of the sequence, **fn_compound_name_next()** returns **0**. Similarly, **fn_compound_name_prev()** returns the component preceding the iteration pointer and moves the marker back one component. If the marker is already at the beginning of the sequence, **fn_compound_name_prev()** returns **0**. The function **fn_compound_name_last()** returns a pointer to the last component of the name and sets the iteration marker immediately preceding this component (so that subsequent calls to **fn_compound_name_prev()** can be used to step through trailing components of the name).

The **fn_compound_name_suffix()** function returns a compound name consisting of a copy of those components following the supplied iteration marker. The function **fn_compound_name_prefix()** returns a compound name consisting of those components that precede the iteration marker. Using these functions with an iteration marker that was not initialized using **fn_compound_name_first(), fn_compound_name_last(), fn_compound_name_is_prefix(),** or **fn_compound_name_is_suffix()** yields undefined and generally undesirable behavior.

The function **fn_compound_name_is_equal(), fn_compound_name_is_prefix(), fn_compound_name_is_suffix(),** tests for equality between compound names or between parts of compound names. For these functions, equality is defined as name equivalence. A name's syntactic property, such as case-insensitivity, is taken into account by these functions.

The function **fn_compound_name_is_prefix()** tests if one compound name is a prefix of another. If so, it returns **1** and sets the iteration marker immediately following the prefix. (So for example, a subsequent call to **fn_compound_name_suffix()** will return the remainder of the name.) Otherwise it returns **0** and value of the iteration marker is undefined. The function **fn_compound_name_is_suffix()** is similar. It tests if a one compound name is a suffix of another. If so it returns **1** and sets the iteration marker immediately preceding the suffix.

fn_compound_name_prepend_comp() and **fn_compound_name_append_comp()** prepends and appends a single atomic component to the given compound name, respectively. These operations invalidates any interation marker the client holds for that object. **fn_compound_name_insert_comp()** inserts an atomic component before *iter_pos* to the given compound name and sets *iter_pos* to be immediately after the component just inserted. **fn_compound_name_delete_comp()** deletes the atomic component located before *iter_pos* from the given compound name and sets *iter_pos* back one component. **fn_compound_name_delete_all()** deletes all the atomic components from *name.*

RETURN VALUE

The functions **fn_compound_name_is_empty(), fn_compound_name_is_equal(), fn_compound_name_is_suffix()** and **fn_compound_name_is_prefix()** return **1** if the test indicated is true; **0** otherwise.

The update functions **fn_compound_name_prepend_comp(), fn_compound_name_append_comp(), fn_compound_name_insert_comp(), fn_compound_name_delete_comp()** and **fn_compound_name_delete_all()** return **1** if the update was successful; **0** otherwise.

If a function is expected to return a pointer to an object, a null pointer (**0**) is returned if the function fails.

ERRORS

When the function **fn_compound_name_from_syntax_attrs()** fails, it returns in *status* a status code. The possible status codes are:

FN_E_ILLEGAL_NAME
The name supplied to the operation was not a well- formed XFN compound name, or one of the component names was not well-formed according to the syntax of the naming system(s) involved in its resolution.

FN_E_INCOMPATIBLE_CODE_SETS
The code set of the given string is incompatible with that supported by the compound name.

FN_E_INVALID_SYNTAX_ATTRS
The syntax attributes supplied are invalid or insufficient to fully specify the syntax.

FN_E_SYNTAX_NOT_SUPPORTED
The syntax type specified is not supported.

fn_compound_name_is_equal(), fn_compound_name_is_suffix(), fn_compound_name_is_prefix(), fn_compound_name_prepend_comp(), fn_compound_name_append_comp() and **fn_compound_name_insert_comp ()** may **FN_E_INCOMPATIBLE_CODE_SETS** return in *status* when the code set of the given string is incompatible with that of the compound name.

SEE ALSO

FN_attribute_t(3N), **FN_attrset_t**(3N), **FN_composite_name_t**(3N), **FN_status_t**(3N), **FN_string_t**(3N), **fn_ctx_get_syntax_attrs**(3N), **xfn_compound_names**(3N), **xfn**(3N)

NAME

FN_ctx_t – an XFN context

MT-LEVEL

Safe.

DESCRIPTION

An XFN context consists of a set of name to reference bindings. An XFN context is represented by the type **FN_ctx_t** in the client interface. The operations for manipulating an **FN_ctx_t** object are described in detail in separate reference manual pages.

These are the interfaces:

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_ctx_t ∗**fn_ctx_handle_from_initial(FN_status_t** ∗**status);**

FN_ctx_t ∗**fn_ctx_handle_from_ref(const FN_ref_t** ∗*ref*, **FN_status_t** ∗*status*);

FN_ref_t ∗**fn_ctx_get_ref(const FN_ctx_t** ∗*ctx*, **FN_status_t** ∗*status*);

void fn_ctx_handle_destroy(FN_ctx_t ∗*ctx*);

FN_ref_t ∗**fn_ctx_lookup(FN_ctx_t** ∗*ctx*, **const FN_composite_name_t** ∗*name*,
FN_status_t ∗*status*);

FN_namelist_t ∗**fn_ctx_list_names(FN_ctx_t** ∗*ctx*, **const FN_composite_name_t** ∗*name*,
FN_status_t ∗*status*);

FN_string_t ∗**fn_namelist_next(FN_namelist_t** ∗*nl*, **FN_status_t** ∗*status*);

void fn_namelist_destroy(FN_namelist_t ∗*nl*, **FN_status_t** ∗*status*);

FN_bindinglist_t ∗**fn_ctx_list_bindings(FN_ctx_t** ∗*ctx*,
const FN_composite_name_t ∗*name*, **FN_status_t** ∗*status*);

FN_string_t ∗**fn_bindinglist_next(FN_bindinglist_t** ∗*iter*, **FN_ref_t** ∗∗*ref*,
FN_status_t ∗*status*);

void fn_bindinglist_destroy(FN_bindinglist_t ∗*iter_pos*, **FN_status_t** ∗*status*);

int fn_ctx_bind(FN_ctx_t ∗*ctx*, **const FN_composite_name_t** ∗*name*,
const FN_ref_t ∗*ref*, **unsigned int** *exclusive*, **FN_status_t** ∗*status*);

int fn_ctx_unbind(FN_ctx_t ∗*ctx*, **const FN_composite_name_t** ∗*name*,
FN_status_t ∗*status*);

int fn_ctx_rename(FN_ctx_t ∗*ctx*, **const FN_composite_name_t** ∗*oldname*,
const FN_composite_name_t ∗*newname*, **unsigned int** *exclusive*,
FN_status_t ∗*status*);

FN_ref_t ∗**fn_ctx_create_subcontext(FN_ctx_t** ∗*ctx*,
const FN_composite_name_t ∗*name*, **FN_status_t** ∗*status*);

int fn_ctx_destroy_subcontext(FN_ctx_t ∗*ctx*, **const FN_composite_name_t** ∗*name*,
FN_status_t ∗*status*);

FN_ref_t ∗**fn_ctx_lookup_link(FN_ctx_t** ∗*ctx*, **const FN_composite_name_t** ∗*name*,
FN_status_t ∗*status*);

FN_attrset_t ***fn_ctx_get_syntax_attrs(FN_ctx_t** ****ctx*****,**
const FN_composite_name_t ****name*****, FN_status_t** ****status*****);**
The following contains a brief summary of these operations.

fn_ctx_handle_from_initial( ) returns a pointer to an Initial Context that provides a starting point for resolution of composite names. **fn_ctx_handle_from_ref( )** returns a handle to an **FN_ctx_t** object using the given reference *ref*. **fn_ctx_get_ref( )** returns the reference of the context *ctx*. **fn_ctx_handle_destroy( )** releases the resources associated with the **FN_ctx_t** object *ctx;* it does not affect the state of the context itself.

fn_ctx_lookup( ) returns the reference bound to *name* resolved relative to *ctx*. **fn_ctx_list_names( )** is used to enumerate the atomic names bound in the context named by *name* resolved relative to *ctx*. **fn_ctx_list_bindings( )** is used to enumerate the atomic names and their references in the context named by *name* resolved relative to *ctx*.

fn_ctx_bind( ) binds the composite name *name* to a reference *ref* resolved relative to *ctx*. **fn_ctx_unbind( )** unbinds *name* resolved relative to *ctx*. **fn_ctx_rename( )** binds *newname* to the reference bound to *oldname* and unbinds *oldname. oldname* is resolved relative to *ctx; newname* is resolved relative to the target context.

fn_ctx_create_subcontext( ) creates a new context with the given composite name *name* resolved relative to *ctx*. **fn_ctx_destroy_subcontext( )** destroys the context named by *name* resolved relative to *ctx*.

Normal resolution always follow links. **fn_ctx_lookup_link( )** lookups up *name* relative to *ctx* , following links except for the last atomic part of *name* , which must be bound to an XFN link.

fn_ctx_get_syntax_attrs( ) returns an attribute set containing attributes that describes a context's syntax. *name* must name a context.

ERRORS

In each context operation, the caller supplies an **FN_status_t** object as a parameter. The called function sets this status object as described in **FN_status_t**(3N) and **xfn_status_codes**(3N).

APPLICATION USAGE

In most of the operations of the base context interface, the caller supplies a context and a composite name. The supplied name is always interpreted relative to the supplied context.

The operation may eventually be effected on a different context called the operation's *target context*. Each operation has an initial resolution phase that conveys the operation to its target context, and the operation is then applied. The effect (but not necessarily the implementation) is that of doing a lookup on that portion of the name that represents the target context, and then invoking the operation on the target context. The contexts involved only in the resolution phase are called *intermediate contexts.*

Normal resolution of names in context operations always follows XFN links.

SEE ALSO

FN_attrset_t(3N), **FN_composite_name_t**(3N), **FN_ref_t**(3N), **FN_status_t**(3N), **fn_ctx_create_subcontext**(3N), **fn_ctx_bind**(3N), **fn_ctx_destroy_subcontext**(3N), **fn_ctx_handle_destroy**(3N), **fn_ctx_handle_from_initial**(3N),

fn_ctx_handle_from_ref(3N), **fn_ctx_get_ref**(3N), **fn_ctx_get_syntax_attrs**(3N), **fn_ctx_list_names**(3N), **fn_ctx_list_bindings**(3N), **fn_ctx_lookup**(3N), **fn_ctx_lookup_link**(3N), **fn_ctx_rename**(3N), **fn_ctx_unbind**(3N), **xfn_links**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

FN_identifier_t – an XFN identifier

DESCRIPTION

Identifiers are used to identify reference types and address types in an XFN reference, and to identify attributes and their syntax in the attribute operations.

An XFN identifier consists of an **unsigned int**, which determines the format of identifier, and the actual identifier, which is expressed as a sequence of octets.

The representation of this structure is defined by XFN as follows:

```
typedef struct {
        unsigned int format;
        size_t length;
        void *contents;
} FN_identifier_t;
```

XFN defines a small number of standard forms for identifiers.

FN_ID_STRING	The identifier is an ASCII string (**ISO 646**).
FN_ID_DCE_UUID	The identifier is an OSF DCE UUID in string representation. (See the **X/Open DCE RPC**.)
FN_ID_ISO_OID_STRING	The identifier is an ISO OID in ASN.1 dot-separated integer list string format. (See the **ISO ASN.1**.)
FN_ID_ISO_OID_BER	The identifier is an ISO OID in ASN.1 Basic Encoding Rules (BER) format. (See the **ISO BER**.)

FILES

#include <xfn/xfn.h>

SEE ALSO

FN_ref_t(3N), **FN_ref_addr_t**(3N), **FN_attribute_t**(3N), **xfn**(3N)

NAME

FN_ref_addr_t, fn_ref_addr_create, fn_ref_addr_destroy, fn_ref_addr_copy, fn_ref_addr_assign, fn_ref_addr_type, fn_ref_addr_length, fn_ref_addr_data, fn_ref_addr_description – an address in an XFN reference

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_ref_addr_t ***fn_ref_addr_create(const FN_identifier_t** ****type*****, size_t** ***length*****,**
const void ****data*****);**

void fn_ref_addr_destroy(FN_ref_addr_t ****addr*****);**

FN_ref_addr_t ***fn_ref_addr_copy(const FN_ref_addr_t** ****addr*****);**

FN_ref_addr_t ***fn_ref_addr_assign(FN_ref_addr_t** ****dst*****, const FN_ref_addr_t** ****src*****);**

const FN_identifier_t ***fn_ref_addr_type(const FN_ref_addr_t** ****addr*****);**

size_t fn_ref_addr_length(const FN_ref_addr_t ****addr*****);**

const void* **fn_ref_addr_data(const FN_ref_addr_t** ****addr*****);**

FN_string_t ***fn_ref_addr_description(const FN_ref_addr_t** ****addr*****, unsigned int** ***detail*****,**
unsigned int ****more_detail*****);**

MT-LEVEL

Safe.

DESCRIPTION

An XFN reference is represented by the type **FN_ref_t**. An object of this type contains a reference type and a list of addresses. Each address in the list is represented by an object of type **FN_ref_addr_t**. An address consists of an opaque data buffer and a type field, of type **FN_identifier_t**.

fn_ref_addr_create() creates and returns an address with the given type and data. *length* indicates the size of the data. **fn_ref_addr_destroy()** releases the storage associated with the given address. **fn_ref_addr_copy()** returns a copy of the given address object. **fn_ref_addr_assign()** makes a copy of the address pointed to by *src* and assigns it to *dst*, releasing any old contents of *dst*. A pointer to the same object as *dst* is returned.

fn_ref_addr_type() returns the type of the given address. **fn_ref_addr_length()** returns the size of the address in bytes. **fn_ref_addr_data()** returns the contents of the address.

fn_ref_addr_description() returns the implementation-defined textual description of the address. It takes as arguments a number, *detail*, and a pointer to a number *more_detail*. *detail* specifies the level of detail for which the description should be generated; the higher the number, the more detail is to be provided. If *more_detail* is zero, it is ignored. If *more_detail* is non-zero, it is set by the description operation to indicate the next level of detail available, beyond that specified by *detail*. If no higher level of detail is available, *more_detail* is set to *detail*.

APPLICATION USAGE

The address type of an **FN_ref_addr_t** object is intended to identify the mechanism that should be used to reach the object using that address. The client must interpret the contents of the opaque data buffer of the address based on the type of the address, and on the type of the reference that the address is in. However, this interpretation is intended

to occur below the application layer. Most applications developers should not have to manipulate the contents of either address or reference objects themselves. These interfaces would generally be used within service libraries.

Multiple addresses in a single reference are intended to identify multiple communication endpoints for the same conceptual object. Multiple addresses may arise for various reasons, such as the object offering interfaces over more than one communication mechanism.

Manipulation of addresses using the operations described in this manual page does not affect their representation in the underlying naming system. Changes to addresses in the underlying naming system can only be effected through the use of the interfaces described in **FN_ctx_t**(3N).

SEE ALSO

FN_ctx_t(3N), **FN_identifier_t**(3N), **FN_ref_t**(3N), **FN_string_t**(3N), **xfn**(3N)

NAME

FN_ref_t, fn_ref_create, fn_ref_destroy, fn_ref_copy, fn_ref_assign, fn_ref_type, fn_ref_addrcount, fn_ref_first, fn_ref_next, fn_ref_prepend_addr, fn_ref_append_addr, fn_ref_insert_addr, fn_ref_delete_addr, fn_ref_delete_all, fn_ref_create_link, fn_ref_is_link, fn_ref_link_name, fn_ref_description – an XFN reference

SYNOPSIS

cc [*flag* ...] *file* ... –**lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_ref_t ***fn_ref_create(const FN_identifier_t** ***ref_type*);

void fn_ref_destroy(FN_ref_t ***ref*);

FN_ref_t ***fn_ref_copy(const FN_ref_t** ***ref*);

FN_ref_t ***fn_ref_assign(FN_ref_t** ***dst*, **const FN_ref_t** ***src*);

const FN_identifier_t ***fn_ref_type(const FN_ref_t** ***ref*);

unsigned int fn_ref_addrcount(const FN_ref_t ***ref*);

const FN_ref_addr_t ***fn_ref_first(const FN_ref_t** ***ref*, **void** ****iter_pos*);

const FN_ref_addr_t ***fn_ref_next(const FN_ref_t** ***ref*, **void** ****iter_pos*);

int fn_ref_prepend_addr(FN_ref_t ***ref*, **const FN_ref_addr_t** ***addr*);

int fn_ref_append_addr(FN_ref_t ***ref*, **const FN_ref_addr_t** ***addr*);

int fn_ref_insert_addr(FN_ref_t ***ref*, **void** ****iter_pos*, **const FN_ref_addr_t** ***addr*);

int fn_ref_delete_addr(FN_ref_t ***ref*, **void** ****iter_pos*);

int fn_ref_delete_all(FN_ref_t ***ref*);

FN_ref_t ***fn_ref_create_link(const FN_composite_name_t** ***link_name*);

int fn_ref_is_link(const FN_ref_t ***ref*);

FN_composite_name_t ***fn_ref_link_name(const FN_ref_t** ***link_ref*);

FN_string_t ***fn_ref_description(const FN_ref_t** ***ref*, **unsigned int** *detail*, **unsigned int** ****more_detail*);

MT-LEVEL

Safe.

DESCRIPTION

An XFN reference is represented by the type **FN_ref_t**. An object of this type contains a reference type and a list of addresses. The ordering in this list at the time of binding might not be preserved when the reference is returned upon lookup.

The reference type is represented by an object of type **FN_identifier_t**. The reference type is intended to identify the class of object referenced. XFN does not dictate the precise use of this.

Each address is represented by an object of type **FN_ref_addr_t**.

fn_ref_create( ) creates a reference with no address, using *ref_type* as its reference type. Addresses can be later added to the reference using the functions described below. **fn_ref_destroy( )** releases the storage associated with *ref*. **fn_ref_copy( )** creates a copy of *ref* and returns it. **fn_ref_assign( )** creates a copy of *src* and assigns it to *dst*, releasing any

old contents of *dst*. A pointer to the same object as *dst* is returned.

fn_ref_addrcount() returns the number of addresses in the reference *ref*.

fn_ref_first() returns the first address in *ref* and sets *iter_pos* to be after the address. It returns **0** if there is no address in the list. **fn_ref_next()** returns the address following *iter_pos* in *ref* and sets *iter_pos* to be after the address. If the iteration marker *iter_pos* is at the end of the sequence, **fn_ref_next()** returns **0**.

fn_ref_prepend_addr() adds *addr* to the front of the list of addresses in *ref*.
fn_ref_append_addr() adds *addr* to the end of the list of addresses in *ref*.
fn_ref_insert_addr() adds *addr* to *ref* before *iter_pos* and sets *iter_pos* to be immediately after the new reference added. **fn_ref_delete_addr()** deletes the address located before *iter_pos* in the list of addresses in *ref* and sets *iter_pos* back one address.
fn_ref_delete_all() deletes all addresses in *ref*.

fn_ref_create_link() creates a reference using the given composite name *link_name* as an address. The reference type of this reference is defined in **fn_ref_is_link()** tests if *ref* is a link. It returns **1** if it is; **0** if it is not. **fn_ref_link_name()** returns the composite name stored in a link reference. It returns **0** if *link_ref* is not a link.

fn_ref_description() returns a string description of the given reference. It takes as argument an integer, *detail*, and a pointer to an integer *more_detail*. *detail* specifies the level of detail for which the description should be generated; the higher the number, the more detail is to be provided. If *more_detail* is zero, it is ignored. If *more_detail* is non-zero, it is set by the description operation to indicate the next level of detail available, beyond that specified by *detail*. If no higher level of detail is available, *more_detail* is set to *detail*.

RETURN VALUE

The operations **fn_ref_prepend_addr(), fn_ref_append_addr(), fn_ref_insert_addr(), fn_ref_delete_addr()** and **fn_ref_delete_all()** return **1** if the operation succeeds, **0** if the operation fails.

APPLICATION USAGE

The reference type is intended to identify the class of object referenced. XFN does not dictate the precise use of this.

Multiple addresses in a single reference are intended to identify multiple communication endpoints for the same conceptual object. Multiple addresses may arise for various reasons, such as the object offering interfaces over more than one communication mechanism.

The client must interpret the contents of a reference based on the type of the addresses and the type of the reference. However, this interpretation is intended to occur below the application layer. Most applications developers should not have to manipulate the contents of either address or reference objects themselves. These interfaces would generally be used within service libraries.

Manipulation of references using the operations described in this manual page does not affect their representation in the underlying naming system. Changes to references in the underlying naming system can only be effected through the use of the interfaces described in **FN_ctx_t**(3N).

SEE ALSO **FN_composite_name_t**(3N), **FN_ctx_t**(3N), **FN_identifier_t**(3N), **FN_string_t**(3N), **FN_ref_addr_t**(3N), **FN_string_t**(3N), **fn_ctx_lookup**(3N), **fn_ctx_lookup_link**(3N), **xfn_links**(3N), **xfn**(3N)

NAME

FN_status_t, fn_status_create, fn_status_destroy, fn_status_copy, fn_status_assign, fn_status_code, fn_status_remaining_name, fn_status_resolved_name, fn_status_resolved_ref, fn_status_diagnostic_message, fn_status_link_code, fn_status_link_remaining_name, fn_status_link_resolved_name, fn_status_link_resolved_ref, fn_status_link_diagnostic_message, fn_status_is_success, fn_status_set_success, fn_status_set, fn_status_set_code, fn_status_set_remaining_name, fn_status_set_resolved_name, fn_status_set_resolved_ref, fn_status_set_diagnostic_message, fn_status_set_link_code, fn_status_set_link_remaining_name, fn_status_set_link_resolved_name, fn_status_set_link_resolved_ref, fn_status_set_link_diagnostic_message, fn_status_append_resolved_name, fn_status_append_remaining_name, fn_status_advance_by_name, fn_status_description – an XFN status object

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_status_t ***fn_status_create(void);**

void fn_status_destroy(FN_status_t ***stat**);

FN_status_t ***fn_status_copy(const FN_status_t** ***stat**);

FN_status_t ***fn_status_assign(FN_status_t** ***dst**, **const FN_status_t** ***src**);

unsigned int fn_status_code(const FN_status_t ***stat**);

const FN_composite_name_t ***fn_status_remaining_name(const FN_status_t** ***stat**);

const FN_composite_name_t ***fn_status_resolved_name(const FN_status_t** ***stat**);

const FN_ref_t ***fn_status_resolved_ref(const FN_status_t** ***stat**);

const FN_string_t ***fn_status_diagnostic_message(const FN_status_t** ***stat**);

unsigned int fn_status_link_code(const FN_status_t ***stat**);

const FN_composite_name_t ***fn_status_link_remaining_name(**
const FN_status_t ***stat**);

const FN_composite_name_t ***fn_status_link_resolved_name(const FN_status_t** ***stat**);

const FN_ref_t ***fn_status_link_resolved_ref(const FN_status_t** ***stat**);

const FN_string_t ***fn_status_link_diagnostic_message(const FN_status_t** ***stat**);

int fn_status_is_success(const FN_status_t ***stat**);

int fn_status_set_success(FN_status_t ***stat**);

int fn_status_set(FN_status_t ***stat**, **unsigned int** *code*, **const FN_ref_t** ***resolved_ref**,
const FN_composite_name_t ***resolved_name**,
const FN_composite_name_t ***remaining_name**);

int fn_status_set_code(FN_status_t ***stat**, **unsigned int** *code*);

int fn_status_set_remaining_name(FN_status_t ***stat**,
const FN_composite_name_t ***name**);

int fn_status_set_resolved_name(FN_status_t **stat*, **const FN_composite_name_t** **name*);

int fn_status_set_resolved_ref(FN_status_t **stat*, **const FN_ref_t** **ref*);

int fn_status_set_diagnostic_message(FN_status_t **stat*, **const FN_string_t** **msg*);

int fn_status_set_link_code(FN_status_t **stat*, **unsigned int** *code*);

int fn_status_set_link_remaining_name(FN_status_t **stat*, **const FN_composite_name_t** **name*);

int fn_status_set_link_resolved_name(FN_status_t **stat*, **const FN_composite_name_t** **name*);

int fn_status_set_link_resolved_ref(FN_status_t **stat*, **const FN_ref_t** **ref*);

int fn_status_set_link_diagnostic_message(FN_status_t **stat*, **const FN_string_t** **msg*);

int fn_status_append_resolved_name(FN_status_t **stat*, **const FN_composite_name_t** **name*);

int fn_status_append_remaining_name(FN_status_t **stat*, **const FN_composite_name_t** **name*);

int fn_status_advance_by_name(FN_status_t **stat*, **const FN_composite_name_t** **prefix*, **const FN_ref_t** **resolved_ref*);

FN_string_t ***fn_status_description(**const FN_status_t** **stat*, **unsigned int** *detail*, **unsigned int** **more_detail*);

MT-LEVEL

Safe.

DESCRIPTION

The result status of operations in the context interface and the attribute interface is encapsulated in an **FN_status_t** object. This object contains information about how the operation completed: whether an error occurred in performing the operation, the nature of the error, and information that helps locate where the error occurred. In the case that the error occurred while resolving an XFN link, the status object contains additional information about that error.

The context status object consists of several items of information.

primary status code	An **unsigned int** code describing the disposition of the operation.
resolved name	In the case of a failure during the resolution phase of the operation, this is the leading portion of the name that was resolved successfully. Resolution may have been successful beyond this point, but the error might not be pinpointed further.
resolved reference	The reference to which resolution was successful (in other words, the reference to which the resolved name is bound).
remaining name	The remaining unresolved portion of the name.
diagnostic message	This contains any diagnostic message returned by the context implementation. This message provides the context

implementation a way of notifying the end-user or administrator of any implementation-specific information related to the returned error status. The diagnostic message could then be used by the end-user or administrator to take appropriate out-of-band action to rectify the problem.

link status code — In the case that an error occurred while resolving an XFN link, the primary status code has the value **FN_E_LINK_ERROR** and the link status code describes the error that occurred while resolving the XFN link.

resolved link name — In the case of a link error, this contains the resolved portion of the name in the XFN link.

resolved link reference — In the case of a link error, this contains the reference to which the resolved link name is bound.

remaining link name — In the case of a link error, this contains the remaining unresolved portion of the name in the XFN link.

link diagnostic message — In the case of a link error, this contains any diagnostic message related to the resolution of the link.

Both the primary status code and the link status code are values of type **unsigned int** that are drawn from the same set of meaningful values. XFN reserves the values 0 through 127 for standard meanings. The values and interpretations for the codes are determined by XFN. See **xfn_status_codes**(3N).

fn_status_create( ) creates a status object with status **FN_SUCCESS. fn_status_destroy( )** releases the storage associated with *stat*. **fn_status_copy( )** returns a copy of the status object *stat*. **fn_status_assign( )** makes a copy of the status object *src* and assigns it to *dst*, releasing any old contents of *dst*. A pointer to the same object as *dst* is returned.

fn_status_code( ) returns the status code. **fn_status_remaining_name( )** returns the remaining part of name to be resolved. **fn_status_resolved_name( )** returns the part of the composite name that has been resolved. **fn_status_resolved_ref( )** returns the reference to which resolution was successful. **fn_status_diagnostic_message** returns any diagnostic message set by the context implementation.

fn_status_link_code( ) returns the link status code. **fn_status_link_remaining_name( )** returns the remaining part of the link name that has not been resolved. **fn_status_link_resolved_name( )** returns the part of the link name that has been resolved. **fn_status_link_resolved_ref( )** returns the reference to which resolution of the link was successful. **fn_status_link_diagnostic_message( )** returns any diagnostic message set by the context implementation during resolution of the link.

fn_status_is_success( ) returns **1** if the status indicates success, **0** otherwise.

fn_status_set_success( ) sets the status code to **FN_SUCCESS** and clears all other parts of *stat*. **fn_status_set( )** sets the non-link contents of the status object *stat*. **fn_status_set_code( )** sets the primary status code field of the status object *stat*.

fn_status_set_remaining_name() sets the remaining name part of the status object *stat* to *name*. **fn_status_set_resolved_name()** sets the resolved name part of the status object *stat* to *name*. **fn_status_set_resolved_ref()** sets the resolved reference part of the status object*stat* to *ref*. **fn_status_set_diagnostic_message()** sets the diagnostic message part of the status object to *msg*.

fn_status_set_link_code() sets the link status code field of the status object *stat* to indicate why resolution of the link failed. **fn_status_set_link_remaining_name()** sets the remaining link name part of the status object *stat* to *name*.
fn_status_set_link_resolved_name() sets the resolved link name part of the status object *stat* to *name*. **fn_status_set_link_resolved_ref()** sets the resolved link reference part of the status object *stat* to *ref*. **fn_status_set_link_diagnostic_message()** sets the link diagnostic message part of the status object to *msg*.

fn_status_append_resolved_name() appends as additional components *name* to the resolved name part of the status object *stat*. **fn_status_append_remaining_name()** appends as additional components *name* to the remaining name part of the status object *stat*. **fn_status_advance_by_name()** removes *prefix* from the remaining name, and appends it to the resolved name. The resolved reference part is set to *resolved_ref*. This operation returns **1** on success, **0** if the *prefix* is not a prefix of the remaining name.

RETURN VALUE

The **fn_status_set_*()** operations return **1** if the operation succeeds, **0** if the operation fails.

SEE ALSO

FN_composite_name_t(3N), **FN_ref_t**(3N), **FN_string_t**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

FN_string_t, fn_string_create, fn_string_destroy, fn_string_from_str, fn_string_from_str_n, fn_string_str, fn_string_from_contents, fn_string_code_set, fn_string_charcount, fn_string_bytecount, fn_string_contents, fn_string_copy, fn_string_assign, fn_string_from_strings, fn_string_from_substring, fn_string_is_empty, fn_string_compare, fn_string_compare_substring, fn_string_next_substring, fn_string_prev_substring – a character string

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_string_t ***fn_string_create(void);**

void fn_string_destroy(FN_string_t ***str**);

FN_string_t ***fn_string_from_str(const unsigned char** ***cstr**);

FN_string_t ***fn_string_from_str_n(const unsigned char** ***cstr**, **size_t** *n*);

const unsigned char ***fn_string_str(const FN_string_t** ***str**, **unsigned int** ***status**);

FN_string_t ***fn_string_from_contents(unsigned long** *code_set*, **const void** ***locale_info**, **size_t** *locale_info_len*, **size_t** *charcount*, **size_t** *bytecount*, **const void** ***contents**, **unsigned int** ***status**);

unsigned long fn_string_code_set(const FN_string_t ***str**, **const void** ****locale_info**, **size_t** ***locale_info_len**);

size_t fn_string_charcount(const FN_string_t ***str**);

size_t fn_string_bytecount(const FN_string_t ***str**);

const void ***fn_string_contents(const FN_string_t** ***str**);

FN_string_t ***fn_string_copy(const FN_string_t** ***str**);

FN_string_t ***fn_string_assign(FN_string_t** ***dst**, **const FN_string_t** ***src**);

FN_string_t ***fn_string_from_strings(unsigned int** ***status**, **const FN_string_t** ***s1**, **const FN_string_t** ***s2**, **...**

FN_string_t ***fn_string_from_substring(const FN_string_t** ***str**, **int** *first*, **int** *last*);

int fn_string_is_empty(const FN_string_t ***str**);

int fn_string_compare(const FN_string_t ***str1**, **const FN_string_t** ***str2**, **unsigned int** *string_case*, **unsigned int** ***status**);

int fn_string_compare_substring(const FN_string_t ***str1**, **int** *first*, **int** *last*, **const FN_string_t** ***str2**, **unsigned int** *string_case*, **unsigned int** ***status**);

int fn_string_next_substring(const FN_string_t ***str**, **const FN_string_t** ***sub**, **int** *index*, **unsigned int** *string_case*, **unsigned int** ***status**);

int fn_string_prev_substring(const FN_string_t ***str**, **const FN_string_t** ***sub**, **int** *index*, **unsigned int** *string_case*, **unsigned int** ***status**);

MT-LEVEL Safe.

DESCRIPTION

The **FN_string_t** type is used to represent character strings in the XFN interface. It provides insulation from specific string representations.

The **FN_string_t** supports multiple code sets. It provides creation functions for character strings of the code set of the current locale setting and a generic creation function for arbitrary code sets. The degree of support for the functions that manipulate **FN_string_t** for arbitrary code sets is implementation-dependent. An XFN implementation is required to support the **ISO 646** code set; all other code sets are optional.

fn_string_destroy( ) releases the storage associated with the given string.

fn_string_create( ) creates an empty string.

fn_string_from_str( ) creates an **FN_string_t** object from the given null terminated string based on the code set of the current locale setting. The number of characters in the string is determined by the code set of the current locale setting. **fn_string_from_str_n( )** is like **fn_string_from_str( )** except only *n* characters from the given string are used. **fn_string_str( )** returns the contents of the given string *str* in the form of a null terminated string in the code set and current locale setting.

fn_string_from_contents( ) creates an **FN_string_t** object using the specified code set *code_set*, locale information *locale_info* and data in the given buffer *contents*. *bytecount* specifies the number of bytes in *contents* and *charcount* specifies the number of characters represented by *contents*.

fn_string_code_set( ) returns the code set associated with the given string object, and if present, the locale information in *locale_info*. **fn_string_charcount( )** returns the number of characters in the given string object. **fn_string_bytecount( )** returns the number of bytes used to represent the given string object. **fn_string_contents( )** returns a pointer to the contents of the given string object.

fn_string_copy( ) returns a copy of the given string object. **fn_string_assign( )** makes a copy of the string object *src* and assigns it to *dst*, releasing any old contents of *dst*. A pointer to the same object as *dst* is returned. **fn_string_from_strings( )** is a function that takes a variable number of arguments (minimum of 2), the last of which must be **NULL** (**0**); it returns a new string object composed of the left to right concatenation of the given strings, in the given order. The support for strings with different code sets and/or locales as arguments to a single invocation of **fn_string_from_strings( )** is implementation-dependent. **fn_string_from_substring( )** returns a new string object consisting of the characters located between *first* and *last* inclusive from *str*. Indexing begins with **0**. If *last* is **FN_STRING_INDEX_LAST** or exceeds the length of the string, the index of the last character of the string is used.

fn_string_is_empty( ) returns whether *str* is an empty string.

Comparison of two strings must take into account code set and locale information. If strings are in the same code set and same locale, case sensitivity is applied according to the case sensitivity rules applicable for the code set and locale; case sensitivity may not necessarily be relevant for all string encodings. If *string_case* is non-zero, case is significant and equality for strings of the same code set is defined as equality between

byte-wise encoded values of the strings. If *string_case* is zero, case is ignored and equality for strings of the same code set is defined using the definition of case-insensitive equality for the specific code set. Support for comparison between strings of different code sets, or lack thereof, is implementation-dependent.

fn_string_compare( ) compares strings *str1* and *str2* and returns **0** if they are equal, non-zero if they are not equal. If two strings are not equal, **fn_string_compare( )** returns a positive value if the difference of *str2* precedes that of *str1* in terms of byte-wise encoded value (with case-sensitivity taken into account when *string_case* is non-zero), and a negative value if the difference of *str1* precedes that of *str2*, in terms of byte-wise encoded value (with case-sensitivity taken into account when *string_case* is non-zero). Such information (positive versus negative return value) may be used by applications that use strings of code sets in which ordering is meaningful; this information is not of general use in internationalized environments. **fn_string_compare_substring( )** is similar to **fn_string_compare( )** except **fn_string_compare_substring( )** compares characters between *first* and *last* inclusive of *str2* with *str1*. Comparison of strings with incompatible code sets returns a negative or positive value (never **0**) depending on the implementation.

fn_string_next_substring( ) returns the index of the next occurrence of *sub* at or after *index* in the string *str*. **FN_STRING_INDEX_NONE** is returned if *sub* does not occur. **fn_string_prev_substring( )** returns the index of the previous occurrence of *sub* at or before *index* in the string *str*. **FN_STRING_INDEX_NONE** is returned if *sub* does not occur. In both of these functions, *string_case* specifies whether the search should take case-sensitivity into account.

ERRORS

fn_string_str( ) returns **0** and sets *status* to **FN_E_INCOMPATIBLE_CODE_SETS** if the given string's representation cannot be converted into the code set of the current locale setting. It is implementation-dependent which code sets can be converted into the code set of the current locale.

Code set mismatches that occur during concatenation, searches or comparisons are resolved in an implementation-dependent way. When an implementation discovers that arguments to substring searches and comparison operations have incompatible code sets, it sets *status* to **FN_E_INCOMPATIBLE_CODE_SETS.** In such cases, **fn_string_from_strings( )** returns **0**. The returned value for comparison operations when there is code set or locale incompatibility is either negative or positive (greater than **0**); it is never **0**;

fn_string_from_contents( ) returns **0** and *status* is set to **FN_E_INCOMPATIBLE_CODE_SETS** if the supplied code set and/or locale information are not supported by the XFN implementation.

SEE ALSO

xfn(3N)

NAME

Intro, intro – introduction to functions and libraries

DESCRIPTION

This section describes functions found in various libraries, other than those functions that directly invoke UNIX system primitives, which are described in Section 2 of this volume. Function declarations can be obtained from the **#include** files indicated on each page. Certain major collections are identified by a letter after the section number:

(3B) These functions constitute the Source Compatibility (with BSD functions) library. It is implemented as a shared object, **libucb.so**, and as an archive, **libucb.a**, but is not automatically linked by the C compilation system. Specify –**lucb** on the **cc** command line to link with this library, which is located in the **/usr/ucb** subdirectory. Header files for this library are located within **/usr/ucbinclude**.

(3C) These functions, together with those of Section 2 and those marked (3S), constitute the standard C library, **libc**, which is automatically linked by the C compilation system. The standard C library is implemented as a shared object, **libc.so**, and as an archive, **libc.a**. C programs are linked with the shared object version of the standard C library by default. Specify –**dn** on the **cc** command line to link with the archive version. (See **cc**(1B) for other overrides, and the ''C Compilation System'' chapter of the *ANSI C Programmer's Guide* for a discussion.) Some functions behave differently in the XPG4 environment. This behavior is noted on the individual manual pages. See **xpg4**(5).

(3E) These functions constitute the ELF access library, **libelf**, (Extensible Linking Formats). This library provides the interface for the creation and analyses of "elf" files; executables, objects, and shared objects. **libelf** is implemented as a shared object, **libelf.so**, and as an archive, **libelf.a**, but is not automatically linked by the C compilation system. Specify –**lelf** on the **cc** command line to link with this library.

(3G) These functions constitute the string pattern-matching & pathname manipulation library, **libgen**. This library is implemented as an archive, **libgen.a**, but not as a shared object, and is not automatically linked by the C compilation system. Specify –**lgen** on the **cc** command line to link with this library.

(3I) These functions constitute the wide character libraries for multi-byte character support, and the international library for messaging. These libraries, **libintl**, and **libw**, are both implemented as shared objects, **libintl.so** and **libw.so**, and as archives, **libintl.a** and **libw.a**. However, they are not automatically linked by the C compilation system; specify –**lintl** or –**lw** on the **cc** command line, as needed to link the appropriate library.

(3K) These functions allow access to the kernel's virtual memory library, which is implemented as a shared object, **libkvm.so**, and as an archive, **libkvm.a**, but is not automatically linked by the C compilation system. Specify –**lkvm** on the **cc** command line to link with this library.

(3M) These functions constitute the math library, **libm**. This library is implemented as a shared object, **libm.so**, and as an archive, **libm.a**, but is not automatically linked by the C compilation system. Specify –**lm** on the **cc** command line to link with this library.

(3N) These functions constitute the Network Service Library, **libnsl**. It is implemented as a shared object, **libnsl.so**, and as an archive, **libnsl.a**, but is not automatically linked by the C compilation system. Specify –**lnsl** on the **cc** command line to link with this library.

Some of the functions documented in man3n incorporate other network libraries, including:
- **libsocket**,
- **libresolv**,
- **librpcsrv**,
- **libnisdb**,
- **librac**,
- **libxfn**, and
- **libkrb**.

(3R) These functions constitute the POSIX.4 Realtime library, **libposix4**. It is implemented only as a shared object, **libposix4.so**, and is not automatically linked by the C compilation system. Specify –**lposix4** on the **cc** command line to link with this library.

(3S) These functions constitute the ''standard I/O package'' (see **stdio**(3S)). They can be compiled using the the standard C library, **libc**, which is automatically linked by the C compilation system. The standard C library is implemented as a shared object, **libc.so**, and as an archive, **libc.a**.

(3T) These functions constitute the threads libraries, **libpthread** and **libthread**. These libraries are used for building multithreaded applications. **libpthread** implements the POSIX threads interface, whereas, **libthread** implements the Solaris threads interface.

Both POSIX threads and Solaris threads can be used within the same application. Their implementations are completely compatible with each other; however, only POSIX threads guarantee portability to other POSIX-compliant environments.

When POSIX and Solaris threads are used in the same application, if there are calls with the same name but different semantics, the POSIX semantic supersedes the Solaris semantic. For example, the call to **fork( )** will imply the **fork1( )** semantic in a program linked with the POSIX threads library, whether or not it is also linked with **-lthread** (Solaris threads).

libpthread and **libthread** are implemented as shared objects, **libpthread.so** and **libthread.so**, but not as archived libraries. **libpthread** and **libthread** are not automatically linked by the C compilation system. Specify –**lpthread** or –**lthread** on the **cc** command line to link with these libraries.

The following functions are optional under POSIX and are not supported in the

current Solaris release.

int pthread_mutexattr_setprotocol(pthread_mutexattr_t **attr,* **int** *protocol***);**
int pthread_mutexattr_getprotocol(const pthread_mutexattr_t **attr,* **int** **protocol***);**
int pthread_mutexattr_setprioceiling(pthread_mutexattr_t **attr,* **int** *prioceiling***);**
int pthread_mutexattr_getprioceiling(const pthread_mutexattr_t **attr,*
int **prioceiling***);**

(3X) Specialized libraries. These functions are contained in libraries including, but not limited to,

- **libadm**,
- **libbsdmalloc**,
- **libcrypt**,
- **libcurses**,
- **libdl**,
- **libform**,
- **libmail**,
- **libmalloc**,
- **libmapmalloc**,
- **libmenu**, and
- **libpanel**.

DEFINITIONS

A character is any bit pattern able to fit into a byte on the machine.

Exception: in some international languages, a "character" may require more than one byte, and is represented in multi-bytes.

The null character is a character with value 0, conventionally represented in the C language as **\ 0**. A character array is a sequence of characters. A null-terminated character array (a *string*) is a sequence of characters, the last of which is the null character. The null string is a character array containing only the terminating null character. A **NULL** pointer is the value that is obtained by casting **0** into a pointer. C guarantees that this value will not match that of any legitimate pointer, so many functions that return pointers return **NULL** to indicate an error. The macro **NULL** is defined in **<stdio.h>**. Types of the form **size_t** are defined in the appropriate headers.

MT-Level of Libraries

Libraries are classified into four categories which define the level of the libraries' ability to support threads.

The MT-Level category of the libraries in this section are shown on each man page under **MT-Level**. Pages containing routines that are of multiple or differing MT-Levels show this under the **NOTES** section.

Safe Safe is simply an attribute of code that can be called from a multithreaded application. It is a generic term used to differentiate between code that is unsafe.

Unsafe An unsafe library contains global and static data that is not protected. It is not safe to use unless the application arranges for only one thread at time to execute within the library. Unsafe libraries may contain routines that are safe; however, most of the library's routines are unsafe to call.

MT-Safe An MT-Safe library is fully prepared for multithreaded access. It protects its global and static data with locks, and can provide a reasonable amount of concurrency. Note that a library can be safe to use, but not MT-Safe. For example, surrounding an entire library with a monitor makes the library safe, but it supports no concurrency so it is not considered MT-Safe. An MT-Safe library must permit a reasonable amount of concurrency. (This definition's purpose is to give precision to what is meant when a library is described as safe. The definition of a "safe" library does not specify if the library supports concurrency. The MT-Safe definition makes it clear that the library is safe, and supports some concurrency. This clarifies the safe definition, which can mean anything from being single threaded to being any degree of multithreaded.)

Async-Signal-Safe

Async-Signal-Safe refers to particular library routines that can be safely called from a signal handler. A thread that is executing an Async-Signal-Safe routine will not deadlock with itself if interrupted by a signal. Signals are only a problem for MT-Safe routines that acquire locks.

Signals are disabled when locks are acquired in Async-Signal-Safe routines. This prevents a signal handler that might acquire the same lock from being called.

The designation "Async-Safe" also indicates "Async-Signal-Safe."

The list of "Async-Signal-Safe" functions includes:

_exit	**kill**	**tcflow**
access	**link**	**tcflush**
aio_error	**lseek**	**tcgetattr**
aio_return	**mkdir**	**tcgetpgrp**
aio_suspend	**mkfifo**	**tcsendbreak**
alarm	**open**	**tcsetattr**
cfgetispeed	**pathconf**	**tcsetpgrp**
cfgetospeed	**pause**	**thr_kill**
cfsetispeed	**pipe**	**thr_sigsetmask**
cfsetospeed	**read**	**time**
chdir	**rename**	**timer_getoverrun**
chmod	**rmdir**	**timer_gettime**
chown	**sem_post**	**timer_settime**
clock_gettime	**sema_post**	**times**
close	**setgid**	**umask**
creat	**setpgid**	**uname**
dup	**setsid**	**unlink**

dup2	**setuid**	**utime**
execle	**sigaction**	**wait**
execve	**sigaddset**	**waitpid**
fcntl	**sigdelset**	**write**
fdatasync	**sigemptyset**	
fork	**sigfillset**	
fstat	**sigismember**	
fsync	**sigpending**	
getegid	**sigprocmask**	
geteuid	**sigqueue**	
getgid	**sigsuspend**	
getgroups	**sleep**	
getpgrp	**stat**	
getpid	**sysconf**	
getppid	**tcdrain**	
getuid	**tcdrain**	

MT-Safe with exceptions

See the **NOTES** sections of these pages for a description of the exceptions.

Safe with exceptions

See the **NOTES** sections of these pages for a description of the exceptions.

Fork1-Safe A fork1-safe library releases the locks it had held whenever **fork1**(2) is called in a Solaris thread program, or **fork**(2) in a POSIX thread program. Calling **fork**(2) in a POSIX thread program has the same semantic as calling **fork1**(2) in a Solaris thread program. All system calls, **libpthread**, and **libthread** are **Fork1-Safe**. Otherwise, you should handle the locking clean-up yourself (see **pthread_atfork**(3T)).

The following table contains reentrant counterparts for unsafe functions. This table is subject to change by SunSoft.

Reentrant functions for libc

Unsafe Function	**Reentrant counterpart**
ctime	**ctime_r**
localtime	**localtime_r**
asctime	**asctime_r**
gmtime	**gmtime_r**
ctermid	**ctermid_r**
getlogin	**getlogin_r**
rand	**rand_r**
readdir	**readdir_r**
strtok	**strtok_r**
tmpnam	**tmpnam_r**

Cancel-Safety

If a multi-threaded application uses **pthread_cancel**(3T) to cancel (i.e., kill a

thread), it is possible that the target thread is killed while holding a resource, such as a lock or allocated memory. If the thread has not installed the appropriate cancellation cleanup handlers to release the resources appropriately (see **pthread_cancel**(3T)), the application is "cancel-unsafe", i.e., it is not safe with respect to cancellation. This unsafety could result in deadlocks due to locks not released by a thread that gets cancelled, or resource leaks; for example, memory not being freed on thread cancellation. All applications that use **pthread_cancel(3T)**should**ensure**that**they**operate cancel-safe environment.

Libraries that have cancellation points and which acquire resources such as locks or allocate memory dynamically, also contribute to the cancel-unsafety of applications that are linked with these libraries. This introduces another level of safety for libraries in a multi-threaded program: cancel-safety.

There are two sub-categories of cancel-safety:

Deferred-cancel-safety and Asynchronous-cancel-safety.

An application is considered to be Deferred-cancel-safe when it is cancel-safe for threads whose cancellation type is PTHREAD_CANCEL_DEFERRED.

An application is considered to be Asynchronous-cancel-safe when it is cancel-safe for threads whose cancellation type is PTHREAD_CANCEL_ASYNCHRONOUS.

Deferred-cancel-safety is easier to achieve than Asynchronous-cancel-safety, since a thread with the deferred cancellation type can be cancelled only at well-defined "cancellation points", whereas a thread with the asynchronous cancellation type can be cancelled anywhere. Since all threads are created by default to have the deferred cancellation type, it may never be necessary to worry about asynchronous cancel safety. Indeed, most applications and libraries are expected to always be Asynchronous-cancel-unsafe.

An application which is Asynchronous-cancel-safe is also, by definition, Deferred-cancel-safe.

FILES

INCDIR	usually **/usr/include**
LIBDIR	usually **/usr/ccs/lib**
LIBDIR/**libc.so**	
LIBDIR/**libc.a**	
LIBDIR/**libgen.a**	
LIBDIR/**libm.a**	
LIBDIR/**libsfm.sa**	
/usr/lib/libc.so.1	

SEE ALSO

ar(1), **cc**(1B), **ld**(1), **nm**(1), **fork**(2), **intro**(2), **stdio**(3S), **pthread_atfork**(3T), **xpg4**(5)

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DIAGNOSTICS

For functions that return floating-point values, error handling varies according to compilation mode. Under the **–Xt** (default) option to **cc**, these functions return the conventional values **0**, ±**HUGE**, or **NaN** when the function is undefined for the given arguments or when the value is not representable. In the **–Xa** and **–Xc** compilation modes, ±**HUGE_VAL** is returned instead of ±**HUGE**. (**HUGE_VAL** and **HUGE** are defined in **math.h** to be infinity and the largest-magnitude single-precision number, respectively.)

NOTES ON MULTITHREAD APPLICATIONS

When compiling a multithreaded application, the **_REENTRANT** flag must be defined on the compile line (**–D_REENTRANT**). This enables new definitions for functions only applicable to multithreaded applications.

When building a singlethreaded application, the _**REENTRANT** flag should be undefined. This generates a binary that is executable on previous Solaris releases, which do not support multithreading.

When linking, it is a requirement that a multithreaded application be constructed to ensure that libthread physically interposes upon the C library. For example, the command line for linking the application using **ld** should be ordered as follows:

example% ld [options] .o's ... –lthread

Note that the behavior of the C library is undefined if **–lc** precedes **–lthread**. And, when linking with **cc**, **–lthread** must be last on the command line.

Unsafe interfaces should be called only from the main thread to ensure the application's safety.

MT-Safe interfaces are denoted in the NOTES section of the functions and libraries man pages. If a man page does not state explicitly that an interface is MT-Safe, the user should assume that the interface is unsafe.

REALTIME APPLICATIONS

Be sure to have set the environment variable **LD_BIND_NOW** to a non-**NULL** value to enable early binding.

NOTES

None of the functions, external variables, or macros should be redefined in the user's programs. Any other name may be redefined without affecting the behavior of other library functions, but such redefinition may conflict with a declaration in an included header.

The headers in *INCDIR* provide function prototypes (function declarations including the types of arguments) for most of the functions listed in this manual. Function prototypes allow the compiler to check for correct usage of these functions in the user's program. The **lint** program checker may also be used and will report discrepancies even if the headers are not included with **#include** statements. Definitions for Sections 2, 3C, and 3S are checked automatically. Other definitions can be included by using the **–l** option to **lint**. (For example, **–lm** includes definitions for **libm**.) Use of **lint** is highly recommended. See the **lint** chapter in *Profiling Tools*.

Users should carefully note the difference between STREAMS and *stream*. STREAMS is a set of kernel mechanisms that support the development of network services and data communication drivers. It is composed of utility routines, kernel facilities, and a set of data structures. A *stream* is a file with its associated buffering. It is declared to be a pointer to a type **FILE** defined in **<stdio.h>**.

In detailed definitions of components, it is sometimes necessary to refer to symbolic names that are implementation-specific, but which are not necessarily expected to be accessible to an application program. Many of these symbolic names describe boundary conditions and system limits.

In this section, for readability, these implementation-specific values are given symbolic names. These names always appear enclosed in curly brackets to distinguish them from symbolic names of other implementation-specific constants that are accessible to application programs by headers. These names are not necessarily accessible to an application program through a header, although they may be defined in the documentation for a particular system.

In general, a portable application program should not refer to these symbolic names in its code. For example, an application program would not be expected to test the length of an argument list given to a routine to determine if it was greater than **{ARG_MAX}**.

Name	**Description**
a64l(3C)	convert between long integer and base-64 ASCII string
abort(3C)	terminate the process abnormally
abs(3C)	return absolute value of integer
accept(3N)	accept a connection on a socket
aclcheck(3)	check the validity of an ACL
aclfrommode(3)	See **acltomode**(3)
aclfrompbits(3)	See **acltopbits**(3)
aclfromtext(3)	See **acltotext**(3)
aclsort(3)	sort an ACL
acltomode(3)	convert an ACL to / from permission bits
acltopbits(3)	convert an ACL to / from permission bits
acltotext(3)	convert an internal representation to / from external representation
acos(3M)	See **trig**(3M)
acosh(3M)	See **hyperbolic**(3M)
addch(3X)	See **curs_addch**(3X)

addchnstr(3X)	See **curs_addchstr**(3X)
addchstr(3X)	See **curs_addchstr**(3X)
addnstr(3X)	See **curs_addstr**(3X)
addnwstr(3X)	See **curs_addwstr**(3X)
addsev(3C)	define additional severities
addseverity(3C)	build a list of severity levels for an application for use with fmtmsg
addstr(3X)	See **curs_addstr**(3X)
addwch(3X)	See **curs_addwch**(3X)
addwchnstr(3X)	See **curs_addwchstr**(3X)
addwchstr(3X)	See **curs_addwchstr**(3X)
addwstr(3X)	See **curs_addwstr**(3X)
adjcurspos(3X)	See **curs_alecompat**(3X)
advance(3G)	See **regexpr**(3G)
aiocancel(3)	cancel an asynchronous operation
aio_cancel(3R)	cancel asynchronous I/O request
aio_error(3R)	See **aio_return**(3R)
aio_fsync(3R)	asynchronous file synchronization
aioread(3)	asynchronous I/O operations
aio_read(3R)	asynchronous read and write operations
aio_return(3R)	retrieve return or error status of asynchronous I/O operation
aio_suspend(3R)	wait for asynchronous I/O request
aiowait(3)	wait for completion of asynchronous I/O operation
aiowrite(3)	See **aioread**(3)
aio_write(3R)	See **aio_read**(3R)
alloca(3C)	See **malloc**(3C)
alphasort(3B)	See **scandir**(3B)
arc(3)	See **plot**(3)
ascftime(3C)	See **strftime**(3C)
asctime(3C)	See **ctime**(3C)
asctime_r(3C)	See **ctime**(3C)
asin(3M)	See **trig**(3M)
asinh(3M)	See **hyperbolic**(3M)

assert(3C)	verify program assertion
asysmem(3)	See **sysmem**(3)
atan(3M)	See **trig**(3M)
atan2(3M)	See **trig**(3M)
atanh(3M)	See **hyperbolic**(3M)
atexit(3C)	add program termination routine
atof(3C)	See **strtod**(3C)
atoi(3C)	See **strtol**(3C)
atol(3C)	See **strtol**(3C)
atoll(3C)	See **strtol**(3C)
attroff(3X)	See **curs_attr**(3X)
attron(3X)	See **curs_attr**(3X)
attrset(3X)	See **curs_attr**(3X)
au_close(3)	See **au_open**(3)
au_open(3)	construct and write audit records
au_preselect(3)	preselect an audit event
authdes_create(3N)	See **rpc_soc**(3N)
authdes_getucred(3N)	See **secure_rpc**(3N)
authdes_seccreate(3N)	See **secure_rpc**(3N)
auth_destroy(3N)	See **rpc_clnt_auth**(3N)
authkerb_getucred(3N)	See **kerberos_rpc**(3N)
authkerb_seccreate(3N)	See **kerberos_rpc**(3N)
authnone_create(3N)	See **rpc_clnt_auth**(3N)
authsys_create(3N)	See **rpc_clnt_auth**(3N)
authsys_create_default(3N)	See **rpc_clnt_auth**(3N)
authunix_create(3N)	See **rpc_soc**(3N)
authunix_create_default(3N)	See **rpc_soc**(3N)
au_to(3)	create audit record tokens
au_to_arg(3)	See **au_to**(3)
au_to_attr(3)	See **au_to**(3)
au_to_data(3)	See **au_to**(3)
au_to_groups(3)	See **au_to**(3)
au_to_in_addr(3)	See **au_to**(3)
au_to_ipc(3)	See **au_to**(3)
au_to_ipc_perm(3)	See **au_to**(3)

au_to_iport(3)	See **au_to**(3)
au_to_me(3)	See **au_to**(3)
au_to_opaque(3)	See **au_to**(3)
au_to_path(3)	See **au_to**(3)
au_to_process(3)	See **au_to**(3)
au_to_return(3)	See **au_to**(3)
au_to_socket(3)	See **au_to**(3)
au_to_text(3)	See **au_to**(3)
au_user_mask(3)	get user's binary preselection mask
au_write(3)	See **au_open**(3)
basename(3G)	return the last element of a path name
baudrate(3X)	See **curs_termattrs**(3X)
bcmp(3C)	See **bstring**(3C)
bcopy(3C)	See **bstring**(3C)
beep(3X)	See **curs_beep**(3X)
bessel(3M)	Bessel functions
bgets(3G)	read stream up to next delimiter
bind(3N)	bind a name to a socket
bindtextdomain(3I)	See **gettext**(3I)
bkgd(3X)	See **curs_bkgd**(3X)
bkgdset(3X)	See **curs_bkgd**(3X)
border(3X)	See **curs_border**(3X)
bottom_panel(3X)	See **panel_top**(3X)
box(3)	See **plot**(3)
box(3X)	See **curs_border**(3X)
bsdmalloc(3X)	memory allocator
bsearch(3C)	binary search a sorted table
bstring(3C)	bit and byte string operations
bufsplit(3G)	split buffer into fields
byteorder(3N)	convert values between host and network byte order
bzero(3C)	See **bstring**(3C)
calloc(3C)	See **malloc**(3C)
calloc(3X)	See **malloc**(3X)
calloc(3X)	See **mapmalloc**(3X)

callrpc(3N)	See **rpc_soc**(3N)
cancellation(3T)	canceling execution of a thread
can_change_color(3X)	See **curs_color**(3X)
catclose(3C)	See **catopen**(3C)
catgets(3C)	read a program message
catopen(3C)	open / close a message catalog
cbc_crypt(3)	See **des_crypt**(3)
cbreak(3X)	See **curs_inopts**(3X)
cbrt(3M)	See **sqrt**(3M)
ceil(3M)	See **floor**(3M)
cfgetispeed(3)	See **termios**(3)
cfgetospeed(3)	See **termios**(3)
cfree(3X)	See **mapmalloc**(3X)
cfsetispeed(3)	See **termios**(3)
cfsetospeed(3)	See **termios**(3)
cftime(3C)	See **strftime**(3C)
circle(3)	See **plot**(3)
clear(3X)	See **curs_clear**(3X)
clearerr(3S)	See **ferror**(3S)
clearok(3X)	See **curs_outopts**(3X)
clnt_broadcast(3N)	See **rpc_soc**(3N)
clnt_call(3N)	See **rpc_clnt_calls**(3N)
clnt_control(3N)	See **rpc_clnt_create**(3N)
clnt_create(3N)	See **rpc_clnt_create**(3N)
clnt_create_timed(3N)	See **rpc_clnt_create**(3N)
clnt_create_vers(3N)	See **rpc_clnt_create**(3N)
clnt_destroy(3N)	See **rpc_clnt_create**(3N)
clnt_dg_create(3N)	See **rpc_clnt_create**(3N)
clnt_freeres(3N)	See **rpc_clnt_calls**(3N)
clnt_geterr(3N)	See **rpc_clnt_calls**(3N)
clnt_pcreateerror(3N)	See **rpc_clnt_create**(3N)
clnt_perrno(3N)	See **rpc_clnt_calls**(3N)
clnt_perror(3N)	See **rpc_clnt_calls**(3N)
clnt_raw_create(3N)	See **rpc_clnt_create**(3N)
clntraw_create(3N)	See **rpc_soc**(3N)

clnt_spcreateerror(3N)	See **rpc_clnt_create**(3N)
clnt_sperrno(3N)	See **rpc_clnt_calls**(3N)
clnt_sperror(3N)	See **rpc_clnt_calls**(3N)
clnttcp_create(3N)	See **rpc_soc**(3N)
clnt_tli_create(3N)	See **rpc_clnt_create**(3N)
clnt_tp_create(3N)	See **rpc_clnt_create**(3N)
clnt_tp_create_timed(3N)	See **rpc_clnt_create**(3N)
clntudp_bufcreate(3N)	See **rpc_soc**(3N)
clntudp_create(3N)	See **rpc_soc**(3N)
clnt_vc_create(3N)	See **rpc_clnt_create**(3N)
clock(3C)	report CPU time used
clock_getres(3R)	See **clock_settime**(3R)
clock_gettime(3R)	See **clock_settime**(3R)
clock_settime(3R)	high-resolution clock operations
closedir(3C)	See **directory**(3C)
closelog(3)	See **syslog**(3)
closepl(3)	See **plot**(3)
closevt(3)	See **plot**(3)
clrtobot(3X)	See **curs_clear**(3X)
clrtoeol(3X)	See **curs_clear**(3X)
color_content(3X)	See **curs_color**(3X)
compile(3G)	See **regexpr**(3G)
cond_broadcast(3T)	See **condition**(3T)
cond_destroy(3T)	See **condition**(3T)
cond_init(3T)	See **condition**(3T)
condition(3T)	condition variables
cond_signal(3T)	See **condition**(3T)
cond_timedwait(3T)	See **condition**(3T)
cond_wait(3T)	See **condition**(3T)
confstr(3C)	get configurable variables
connect(3N)	initiate a connection on a socket
cont(3)	See **plot**(3)
conv(3C)	translate characters
copylist(3G)	copy a file into memory
copysign(3M)	See **ieee_functions**(3M)

copywin(3X)	See **curs_overlay**(3X)
cos(3M)	See **trig**(3M)
cosh(3M)	See **hyperbolic**(3M)
crypt(3C)	generate encryption
cset(3I)	get information on EUC codesets
csetcol(3I)	See **cset**(3I)
csetlen(3I)	See **cset**(3I)
csetno(3I)	See **cset**(3I)
ctermid(3S)	generate path name for controlling terminal
ctermid_r(3S)	See **ctermid**(3S)
ctime(3C)	convert date and time to string
ctime_r(3C)	See **ctime**(3C)
ctype(3C)	character handling
current_field(3X)	See **form_page**(3X)
current_item(3X)	See **menu_item_current**(3X)
curs_addch(3X)	add a character (with attributes) to a curses window and advance cursor
curs_addchstr(3X)	add string of characters (and attributes) to a curses window
curs_addstr(3X)	add a string of characters to a curses window and advance cursor
curs_addwch(3X)	add a wchar_t character (with attributes) to a curses window and advance cursor
curs_addwchstr(3X)	add string of wchar_t characters (and attributes) to a curses window
curs_addwstr(3X)	add a string of wchar_t characters to a curses window and advance cursor
curs_alecompat(3X)	these functions are added to ALE curses library for moving the cursor by character.
curs_attr(3X)	curses character and window attribute control routines
curs_beep(3X)	curses bell and screen flash routines
curs_bkgd(3X)	curses window background manipulation routines
curs_border(3X)	create curses borders, horizontal and

	vertical lines
curs_clear(3X)	clear all or part of a curses window
curs_color(3X)	curses color manipulation routines
curs_delch(3X)	delete character under cursor in a curses window
curs_deleteln(3X)	delete and insert lines in a curses window
curses(3X)	CRT screen handling and optimization package
curs_getch(3X)	get (or push back) characters from curses terminal keyboard
curs_getstr(3X)	get character strings from curses terminal keyboard
curs_getwch(3X)	get (or push back) wchar_t characters from curses terminal keyboard
curs_getwstr(3X)	get wchar_t character strings from curses terminal keyboard
curs_getyx(3X)	get curses cursor and window coordinates
curs_inch(3X)	get a character and its attributes from a curses window
curs_inchstr(3X)	get a string of characters (and attributes) from a curses window
curs_initscr(3X)	curses screen initialization and manipulation routines
curs_inopts(3X)	curses terminal input option control routines
curs_insch(3X)	insert a character before the character under the cursor in a curses window
curs_insstr(3X)	insert string before character under the cursor in a curses window
curs_instr(3X)	get a string of characters from a curses window
curs_inswch(3X)	insert a wchar_t character before the character under the cursor in a curses window
curs_inswstr(3X)	insert wchar_t string before character under the cursor in a curses window
curs_inwch(3X)	get a wchar_t character and its

	attributes from a curses window
curs_inwchstr(3X)	get a string of wchar_t characters (and attributes) from a curses window
curs_inwstr(3X)	get a string of wchar_t characters from a curses window
curs_kernel(3X)	low-level curses routines
curs_move(3X)	move curses window cursor
curs_outopts(3X)	curses terminal output option control routines
curs_overlay(3X)	overlap and manipulate overlapped curses windows
curs_pad(3X)	create and display curses pads
curs_printw(3X)	print formatted output in curses windows
curs_refresh(3X)	refresh curses windows and lines
curs_scanw(3X)	convert formatted input from a curses widow
curs_scr_dump(3X)	read (write) a curses screen from (to) a file
curs_scroll(3X)	scroll a curses window
curs_set(3X)	See **curs_kernel**(3X)
curs_slk(3X)	curses soft label routines
curs_termattrs(3X)	curses environment query routines
curs_termcap(3X)	curses interfaces (emulated) to the termcap library
curs_terminfo(3X)	curses interfaces to terminfo database
curs_touch(3X)	curses refresh control routines
curs_util(3X)	curses miscellaneous utility routines
curs_window(3X)	create curses windows
cuserid(3S)	get character login name of the user
data_ahead(3X)	See **form_data**(3X)
data_behind(3X)	See **form_data**(3X)
db_add_entry(3N)	See **nis_db**(3N)
db_checkpoint(3N)	See **nis_db**(3N)
db_create_table(3N)	See **nis_db**(3N)

db_destroy_table(3N)	See **nis_db**(3N)
db_first_entry(3N)	See **nis_db**(3N)
db_free_result(3N)	See **nis_db**(3N)
db_initialize(3N)	See **nis_db**(3N)
db_list_entries(3N)	See **nis_db**(3N)
dbm(3B)	data base subroutines
dbm_clearerr(3)	See **ndbm**(3)
dbm_close(3)	See **ndbm**(3)
dbmclose(3B)	See **dbm**(3B)
dbm_delete(3)	See **ndbm**(3)
dbm_error(3)	See **ndbm**(3)
dbm_fetch(3)	See **ndbm**(3)
dbm_firstkey(3)	See **ndbm**(3)
dbminit(3B)	See **dbm**(3B)
dbm_nextkey(3)	See **ndbm**(3)
dbm_open(3)	See **ndbm**(3)
dbm_store(3)	See **ndbm**(3)
db_next_entry(3N)	See **nis_db**(3N)
db_remove_entry(3N)	See **nis_db**(3N)
db_reset_next_entry(3N)	See **nis_db**(3N)
db_standby(3N)	See **nis_db**(3N)
db_table_exists(3N)	See **nis_db**(3N)
db_unload_table(3N)	See **nis_db**(3N)
dcgettext(3I)	See **gettext**(3I)
decimal_to_double(3)	See **decimal_to_floating**(3)
decimal_to_extended(3)	See **decimal_to_floating**(3)
decimal_to_floating(3)	convert decimal record to floating-point value
decimal_to_quadruple(3)	See **decimal_to_floating**(3)
decimal_to_single(3)	See **decimal_to_floating**(3)
def_prog_mode(3X)	See **curs_kernel**(3X)
def_shell_mode(3X)	See **curs_kernel**(3X)
delay_output(3X)	See **curs_util**(3X)
delch(3X)	See **curs_delch**(3X)
del_curterm(3X)	See **curs_terminfo**(3X)

delete(3B)	See **dbm**(3B)
deleteln(3X)	See **curs_deleteln**(3X)
del_panel(3X)	See **panel_new**(3X)
delscreen(3X)	See **curs_initscr**(3X)
delwin(3X)	See **curs_window**(3X)
demangle(3)	decode a C++ encoded symbol name
derwin(3X)	See **curs_window**(3X)
des_crypt(3)	fast DES encryption
DES_FAILED(3)	See **des_crypt**(3)
des_failed(3)	See **des_crypt**(3)
des_setparity(3)	See **des_crypt**(3)
dgettext(3I)	See **gettext**(3I)
dial(3N)	establish an outgoing terminal line connection
difftime(3C)	computes the difference between two calendar times
directory(3C)	directory operations
dirname(3G)	report the parent directory name of a file path name
div(3C)	compute the quotient and remainder
dladdr(3X)	translate address to symbolic information
dlclose(3X)	close a shared object
dlerror(3X)	get diagnostic information
dlopen(3X)	open a shared object
dlsym(3X)	get the address of a symbol in a shared object
dn_comp(3N)	See **resolver**(3N)
dn_expand(3N)	See **resolver**(3N)
doconfig(3N)	execute a configuration script
double_to_decimal(3)	See **floating_to_decimal**(3)
doupdate(3X)	See **curs_refresh**(3X)
drand48(3C)	generate uniformly distributed pseudo-random numbers
dup2(3C)	duplicate an open file descriptor
dup_field(3X)	See **form_field_new**(3X)

dupwin(3X)	See **curs_window**(3X)
dynamic_field_info(3X)	See **form_field_info**(3X)
ecb_crypt(3)	See **des_crypt**(3)
echo(3X)	See **curs_inopts**(3X)
echochar(3X)	See **curs_addch**(3X)
echowchar(3X)	See **curs_addwch**(3X)
econvert(3)	output conversion
ecvt(3)	See **econvert**(3)
ecvt(3C)	convert floating-point number to string
_edata(3C)	See **end**(3C)
edata(3C)	See **end**(3C)
elf(3E)	object file access library
elf32_fsize(3E)	return the size of an object file type
elf32_getehdr(3E)	retrieve class-dependent object file header
elf32_getphdr(3E)	retrieve class-dependent program header table
elf32_getshdr(3E)	retrieve class-dependent section header
elf32_newehdr(3E)	See **elf32_getehdr**(3E)
elf32_newphdr(3E)	See **elf32_getphdr**(3E)
elf32_xlatetof(3E)	class-dependent data translation
elf32_xlatetom(3E)	See **elf32_xlatetof**(3E)
elf_begin(3E)	process ELF object files
elf_cntl(3E)	control an elf file descriptor
elf_end(3E)	See **elf_begin**(3E)
elf_errmsg(3E)	error handling
elf_errno(3E)	See **elf_errmsg**(3E)
elf_fill(3E)	set fill byte
elf_flagdata(3E)	manipulate flags
elf_flagehdr(3E)	See **elf_flagdata**(3E)
elf_flagelf(3E)	See **elf_flagdata**(3E)
elf_flagphdr(3E)	See **elf_flagdata**(3E)
elf_flagscn(3E)	See **elf_flagdata**(3E)
elf_flagshdr(3E)	See **elf_flagdata**(3E)

elf_getarhdr(3E)	retrieve archive member header
elf_getarsym(3E)	retrieve archive symbol table
elf_getbase(3E)	get the base offset for an object file
elf_getdata(3E)	get section data
elf_getident(3E)	retrieve file identification data
elf_getscn(3E)	get section information
elf_hash(3E)	compute hash value
elf_kind(3E)	determine file type
elf_memory(3E)	See **elf_begin**(3E)
elf_ndxscn(3E)	See **elf_getscn**(3E)
elf_newdata(3E)	See **elf_getdata**(3E)
elf_newscn(3E)	See **elf_getscn**(3E)
elf_next(3E)	See **elf_begin**(3E)
elf_nextscn(3E)	See **elf_getscn**(3E)
elf_rand(3E)	See **elf_begin**(3E)
elf_rawdata(3E)	See **elf_getdata**(3E)
elf_rawfile(3E)	retrieve uninterpreted file contents
elf_strptr(3E)	make a string pointer
elf_update(3E)	update an ELF descriptor
elf_version(3E)	coordinate ELF library and application versions
encrypt(3C)	See **crypt**(3C)
end(3C)	last locations in program
_end(3C)	See **end**(3C)
endac(3)	See **getacinfo**(3)
endauclass(3)	See **getauclassent**(3)
endauevent(3)	See **getauevent**(3)
endauuser(3)	See **getauusernam**(3)
endgrent(3C)	See **getgrnam**(3C)
endhostent(3N)	See **gethostbyname**(3N)
endnetconfig(3N)	See **getnetconfig**(3N)
endnetent(3N)	See **getnetbyname**(3N)
endnetgrent(3N)	See **getnetgrent**(3N)
endnetpath(3N)	See **getnetpath**(3N)
endprotoent(3N)	See **getprotobyname**(3N)

endpwent(3C)	See **getpwnam**(3C)
endrpcent(3N)	See **getrpcbyname**(3N)
endservent(3N)	See **getservbyname**(3N)
endspent(3C)	See **getspnam**(3C)
endusershell(3C)	See **getusershell**(3C)
endutent(3C)	See **getutent**(3C)
endutxent(3C)	See **getutxent**(3C)
endwin(3X)	See **curs_initscr**(3X)
erand48(3C)	See **drand48**(3C)
erase(3)	See **plot**(3)
erase(3X)	See **curs_clear**(3X)
erasechar(3X)	See **curs_termattrs**(3X)
erf(3M)	error functions
erfc(3M)	See **erf**(3M)
errno(3C)	See **perror**(3C)
_etext(3C)	See **end**(3C)
etext(3C)	See **end**(3C)
ether_aton(3N)	See **ethers**(3N)
ether_hostton(3N)	See **ethers**(3N)
ether_line(3N)	See **ethers**(3N)
ether_ntoa(3N)	See **ethers**(3N)
ether_ntohost(3N)	See **ethers**(3N)
ethers(3N)	Ethernet address mapping operations
euccol(3I)	See **euclen**(3I)
euclen(3I)	get byte length and display width of EUC characters
eucscol(3I)	See **euclen**(3I)
exit(3C)	terminate process
exp(3M)	exponential, logarithm, power
expm1(3M)	See **exp**(3M)
extended_to_decimal(3)	See **floating_to_decimal**(3)
fabs(3M)	See **ieee_functions**(3M)
fattach(3C)	attach a STREAMS-based file descriptor to an object in the file system name space

fclose(3S)	close or flush a stream
fconvert(3)	See **econvert**(3)
fcvt(3)	See **econvert**(3)
fcvt(3C)	See **ecvt**(3C)
fdatasync(3R)	synchronize a file's data
fdetach(3C)	detach a name from a STREAMS-based file descriptor
fdopen(3S)	See **fopen**(3S)
feof(3S)	See **ferror**(3S)
ferror(3S)	stream status inquiries
fetch(3B)	See **dbm**(3B)
fflush(3S)	See **fclose**(3S)
ffs(3C)	find first set bit
fgetc(3S)	See **getc**(3S)
fgetgrent(3C)	See **getgrnam**(3C)
fgetgrent_r(3C)	See **getgrnam**(3C)
fgetpos(3C)	See **fsetpos**(3C)
fgetpwent(3C)	See **getpwnam**(3C)
fgetpwent_r(3C)	See **getpwnam**(3C)
fgets(3S)	See **gets**(3S)
fgetspent(3C)	See **getspnam**(3C)
fgetspent_r(3C)	See **getspnam**(3C)
fgetwc(3I)	See **getwc**(3I)
fgetws(3I)	See **getws**(3I)
field_arg(3X)	See **form_field_validation**(3X)
field_back(3X)	See **form_field_attributes**(3X)
field_buffer(3X)	See **form_field_buffer**(3X)
field_count(3X)	See **form_field**(3X)
field_fore(3X)	See **form_field_attributes**(3X)
field_index(3X)	See **form_page**(3X)
field_info(3X)	See **form_field_info**(3X)
field_init(3X)	See **form_hook**(3X)
field_just(3X)	See **form_field_just**(3X)
field_opts(3X)	See **form_field_opts**(3X)
field_opts_off(3X)	See **form_field_opts**(3X)

field_opts_on(3X)	See **form_field_opts**(3X)
field_pad(3X)	See **form_field_attributes**(3X)
field_status(3X)	See **form_field_buffer**(3X)
field_term(3X)	See **form_hook**(3X)
field_type(3X)	See **form_field_validation**(3X)
field_userptr(3X)	See **form_field_userptr**(3X)
fileno(3S)	See **ferror**(3S)
file_to_decimal(3)	See **string_to_decimal**(3)
filter(3X)	See **curs_util**(3X)
finite(3C)	See **isnan**(3C)
firstkey(3B)	See **dbm**(3B)
flash(3X)	See **curs_beep**(3X)
floating_to_decimal(3)	convert floating-point value to decimal record
flock(3B)	apply or remove an advisory lock on an open file
flockfile(3S)	acquire and release stream lock
floor(3M)	round to integral value in floating-point format
flushinp(3X)	See **curs_util**(3X)
fmod(3M)	See **ieee_functions**(3M)
fmtmsg(3C)	display a message on stderr or system console
fn_attr_get(3N)	return specified attribute associated with name
fn_attr_get_ids(3N)	get a list of the identifiers of all attributes associated with named object
fn_attr_get_values(3N)	return values of an attribute
fn_attribute_add(3N)	See **FN_attribute_t**(3N)
fn_attribute_assign(3N)	See **FN_attribute_t**(3N)
fn_attribute_copy(3N)	See **FN_attribute_t**(3N)
fn_attribute_create(3N)	See **FN_attribute_t**(3N)
fn_attribute_destroy(3N)	See **FN_attribute_t**(3N)
fn_attribute_first(3N)	See **FN_attribute_t**(3N)
fn_attribute_identifier(3N)	See **FN_attribute_t**(3N)
fn_attribute_next(3N)	See **FN_attribute_t**(3N)

fn_attribute_remove(3N)	See **FN_attribute_t**(3N)
fn_attribute_syntax(3N)	See **FN_attribute_t**(3N)
FN_attribute_t(3N)	an XFN attribute
fn_attribute_valuecount(3N)	See **FN_attribute_t**(3N)
fn_attr_modify(3N)	modify specified attribute associated with name
fn_attrmodlist_add(3N)	See **FN_attrmodlist_t**(3N)
fn_attrmodlist_assign(3N)	See **FN_attrmodlist_t**(3N)
fn_attrmodlist_copy(3N)	See **FN_attrmodlist_t**(3N)
fn_attrmodlist_count(3N)	See **FN_attrmodlist_t**(3N)
fn_attrmodlist_create(3N)	See **FN_attrmodlist_t**(3N)
fn_attrmodlist_destroy(3N)	See **FN_attrmodlist_t**(3N)
fn_attrmodlist_first(3N)	See **FN_attrmodlist_t**(3N)
fn_attrmodlist_next(3N)	See **FN_attrmodlist_t**(3N)
FN_attrmodlist_t(3N)	a list of attribute modifications
fn_attr_multi_get(3N)	return multiple attributes associated with named object
fn_attr_multi_modify(3N)	modify multiple attributes associated with named object
fn_attrset_add(3N)	See **FN_attrset_t**(3N)
fn_attrset_assign(3N)	See **FN_attrset_t**(3N)
fn_attrset_copy(3N)	See **FN_attrset_t**(3N)
fn_attrset_count(3N)	See **FN_attrset_t**(3N)
fn_attrset_create(3N)	See **FN_attrset_t**(3N)
fn_attrset_destroy(3N)	See **FN_attrset_t**(3N)
fn_attrset_first(3N)	See **FN_attrset_t**(3N)
fn_attrset_get(3N)	See **FN_attrset_t**(3N)
fn_attrset_next(3N)	See **FN_attrset_t**(3N)
fn_attrset_remove(3N)	See **FN_attrset_t**(3N)
FN_attrset_t(3N)	a set of XFN attributes
FN_attrvalue_t(3N)	an XFN attribute value
fn_bindinglist_destroy(3N)	See **fn_ctx_list_bindings**(3N)
fn_bindinglist_next(3N)	See **fn_ctx_list_bindings**(3N)
FN_bindinglist_t(3N)	See **fn_ctx_list_bindings**(3N)
fn_composite_name_append_comp(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_append_name(3N)	See **FN_composite_name_t**(3N)

fn_composite_name_assign(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_copy(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_count(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_create(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_delete_comp(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_destroy(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_first(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_from_string(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_insert_comp(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_insert_name(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_is_empty(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_is_equal(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_is_prefix(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_is_suffix(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_last(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_next(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_prefix(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_prepend_comp(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_prepend_name(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_prev(3N)	See **FN_composite_name_t**(3N)
fn_composite_name_suffix(3N)	See **FN_composite_name_t**(3N)
FN_composite_name_t(3N)	a sequence of component names spanning multiple naming systems
fn_compound_name_append_comp(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_assign(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_copy(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_count(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_delete_all(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_delete_comp(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_destroy(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_first(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_from_syntax_attrs(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_get_syntax_attrs(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_insert_comp(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_is_empty(3N)	See **FN_compound_name_t**(3N)

fn_compound_name_is_equal(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_is_prefix(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_is_suffix(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_last(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_next(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_prefix(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_prepend_comp(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_prev(3N)	See **FN_compound_name_t**(3N)
fn_compound_name_suffix(3N)	See **FN_compound_name_t**(3N)
FN_compound_name_t(3N)	an XFN compound name
fn_ctx_bind(3N)	bind a reference to a name
fn_ctx_create_subcontext(3N)	create a subcontext in a context
fn_ctx_destroy_subcontext(3N)	destroy the named context and remove its binding from the parent context
fn_ctx_get_ref(3N)	return a context's reference
fn_ctx_get_syntax_attrs(3N)	return syntax attributes associated with named context
fn_ctx_handle_destroy(3N)	release storage associated with context handle
fn_ctx_handle_from_initial(3N)	return a handle to the Initial Context
fn_ctx_handle_from_ref(3N)	construct a handle to a context object using the given reference
fn_ctx_list_bindings(3N)	list the atomic names and references bound in a context
fn_ctx_list_names(3N)	list the atomic names bound in a context
fn_ctx_lookup(3N)	look up name in context
fn_ctx_lookup_link(3N)	look up the link reference bound to a name
fn_ctx_rename(3N)	rename the name of a binding
FN_ctx_t(3N)	an XFN context
fn_ctx_unbind(3N)	unbind a name from a context
FN_identifier_t(3N)	an XFN identifier
fnmatch(3C)	match filename or path name
fn_multigetlist_destroy(3N)	See **fn_attr_multi_get**(3N)
fn_multigetlist_next(3N)	See **fn_attr_multi_get**(3N)

FN_multigetlist_t(3N)	See **fn_attr_multi_get**(3N)
fn_namelist_destroy(3N)	See **fn_ctx_list_names**(3N)
fn_namelist_next(3N)	See **fn_ctx_list_names**(3N)
FN_namelist_t(3N)	See **fn_ctx_list_names**(3N)
fn_ref_addr_assign(3N)	See **FN_ref_addr_t**(3N)
fn_ref_addr_copy(3N)	See **FN_ref_addr_t**(3N)
fn_ref_addrcount(3N)	See **FN_ref_t**(3N)
fn_ref_addr_create(3N)	See **FN_ref_addr_t**(3N)
fn_ref_addr_data(3N)	See **FN_ref_addr_t**(3N)
fn_ref_addr_description(3N)	See **FN_ref_addr_t**(3N)
fn_ref_addr_destroy(3N)	See **FN_ref_addr_t**(3N)
fn_ref_addr_length(3N)	See **FN_ref_addr_t**(3N)
FN_ref_addr_t(3N)	an address in an XFN reference
fn_ref_addr_type(3N)	See **FN_ref_addr_t**(3N)
fn_ref_append_addr(3N)	See **FN_ref_t**(3N)
fn_ref_assign(3N)	See **FN_ref_t**(3N)
fn_ref_copy(3N)	See **FN_ref_t**(3N)
fn_ref_create(3N)	See **FN_ref_t**(3N)
fn_ref_create_link(3N)	See **FN_ref_t**(3N)
fn_ref_delete_addr(3N)	See **FN_ref_t**(3N)
fn_ref_delete_all(3N)	See **FN_ref_t**(3N)
fn_ref_description(3N)	See **FN_ref_t**(3N)
fn_ref_destroy(3N)	See **FN_ref_t**(3N)
fn_ref_first(3N)	See **FN_ref_t**(3N)
fn_ref_insert_addr(3N)	See **FN_ref_t**(3N)
fn_ref_is_link(3N)	See **FN_ref_t**(3N)
fn_ref_link_name(3N)	See **FN_ref_t**(3N)
fn_ref_next(3N)	See **FN_ref_t**(3N)
fn_ref_prepend_addr(3N)	See **FN_ref_t**(3N)
FN_ref_t(3N)	an XFN reference
fn_ref_type(3N)	See **FN_ref_t**(3N)
fn_status_advance_by_name(3N)	See **FN_status_t**(3N)
fn_status_append_remaining_name(3N)	See **FN_status_t**(3N)
fn_status_append_resolved_name(3N)	See **FN_status_t**(3N)
fn_status_assign(3N)	See **FN_status_t**(3N)

fn_status_code(3N)	See **FN_status_t**(3N)
fn_status_copy(3N)	See **FN_status_t**(3N)
fn_status_create(3N)	See **FN_status_t**(3N)
fn_status_description(3N)	See **FN_status_t**(3N)
fn_status_destroy(3N)	See **FN_status_t**(3N)
fn_status_diagnostic_message(3N)	See **FN_status_t**(3N)
fn_status_is_success(3N)	See **FN_status_t**(3N)
fn_status_link_code(3N)	See **FN_status_t**(3N)
fn_status_link_diagnostic_message(3N)	See **FN_status_t**(3N)
fn_status_link_remaining_name(3N)	See **FN_status_t**(3N)
fn_status_link_resolved_name(3N)	See **FN_status_t**(3N)
fn_status_link_resolved_ref(3N)	See **FN_status_t**(3N)
fn_status_remaining_name(3N)	See **FN_status_t**(3N)
fn_status_resolved_name(3N)	See **FN_status_t**(3N)
fn_status_resolved_ref(3N)	See **FN_status_t**(3N)
fn_status_set(3N)	See **FN_status_t**(3N)
fn_status_set_code(3N)	See **FN_status_t**(3N)
fn_status_set_diagnostic_message(3N)	See **FN_status_t**(3N)
fn_status_set_link_code(3N)	See **FN_status_t**(3N)
fn_status_set_link_diagnostic_message(3N)	See **FN_status_t**(3N)
fn_status_set_link_remaining_name(3N)	See **FN_status_t**(3N)
fn_status_set_link_resolved_name(3N)	See **FN_status_t**(3N)
fn_status_set_link_resolved_ref(3N)	See **FN_status_t**(3N)
fn_status_set_remaining_name(3N)	See **FN_status_t**(3N)
fn_status_set_resolved_name(3N)	See **FN_status_t**(3N)
fn_status_set_resolved_ref(3N)	See **FN_status_t**(3N)
fn_status_set_success(3N)	See **FN_status_t**(3N)
FN_status_t(3N)	an XFN status object
fn_string_assign(3N)	See **FN_string_t**(3N)
fn_string_bytecount(3N)	See **FN_string_t**(3N)
fn_string_charcount(3N)	See **FN_string_t**(3N)
fn_string_code_set(3N)	See **FN_string_t**(3N)
fn_string_compare(3N)	See **FN_string_t**(3N)
fn_string_compare_substring(3N)	See **FN_string_t**(3N)
fn_string_contents(3N)	See **FN_string_t**(3N)

fn_string_copy(3N)	See **FN_string_t**(3N)
fn_string_create(3N)	See **FN_string_t**(3N)
fn_string_destroy(3N)	See **FN_string_t**(3N)
fn_string_from_composite_name(3N)	See **FN_composite_name_t**(3N)
fn_string_from_compound_name(3N)	See **FN_compound_name_t**(3N)
fn_string_from_contents(3N)	See **FN_string_t**(3N)
fn_string_from_str(3N)	See **FN_string_t**(3N)
fn_string_from_strings(3N)	See **FN_string_t**(3N)
fn_string_from_str_n(3N)	See **FN_string_t**(3N)
fn_string_from_substring(3N)	See **FN_string_t**(3N)
fn_string_is_empty(3N)	See **FN_string_t**(3N)
fn_string_next_substring(3N)	See **FN_string_t**(3N)
fn_string_prev_substring(3N)	See **FN_string_t**(3N)
fn_string_str(3N)	See **FN_string_t**(3N)
FN_string_t(3N)	a character string
fn_valuelist_destroy(3N)	See **fn_attr_get_values**(3N)
fn_valuelist_next(3N)	See **fn_attr_get_values**(3N)
FN_valuelist_t(3N)	See **fn_attr_get_values**(3N)
fopen(3B)	open a stream
fopen(3S)	open a stream
form_cursor(3X)	position forms window cursor
form_data(3X)	tell if forms field has off-screen data ahead or behind
form_driver(3X)	command processor for the forms subsystem
form_field(3X)	connect fields to forms
form_field_attributes(3X)	format the general display attributes of forms
form_field_buffer(3X)	set and get forms field attributes
form_field_info(3X)	get forms field characteristics
form_field_just(3X)	format the general appearance of forms
form_field_new(3X)	create and destroy forms fields
form_field_opts(3X)	forms field option routines
form_fields(3X)	See **form_field**(3X)
form_fieldtype(3X)	forms fieldtype routines

form_field_userptr(3X)	associate application data with forms
form_field_validation(3X)	forms field data type validation
form_hook(3X)	assign application-specific routines for invocation by forms
form_init(3X)	See **form_hook**(3X)
form_new(3X)	create and destroy forms
form_new_page(3X)	forms pagination
form_opts(3X)	forms option routines
form_opts_off(3X)	See **form_opts**(3X)
form_opts_on(3X)	See **form_opts**(3X)
form_page(3X)	set forms current page and field
form_post(3X)	write or erase forms from associated subwindows
forms(3X)	character based forms package
form_sub(3X)	See **form_win**(3X)
form_term(3X)	See **form_hook**(3X)
form_userptr(3X)	associate application data with forms
form_win(3X)	forms window and subwindow association routines
fpclass(3C)	See **isnan**(3C)
fpgetmask(3C)	See **fpgetround**(3C)
fpgetround(3C)	IEEE floating-point environment control
fpgetsticky(3C)	See **fpgetround**(3C)
fprintf(3B)	See **printf**(3B)
fprintf(3S)	See **printf**(3S)
fpsetmask(3C)	See **fpgetround**(3C)
fpsetround(3C)	See **fpgetround**(3C)
fpsetsticky(3C)	See **fpgetround**(3C)
fputc(3S)	See **putc**(3S)
fputs(3S)	See **puts**(3S)
fputwc(3I)	See **putwc**(3I)
fputws(3I)	See **putws**(3I)
fread(3S)	buffered binary input/output
free(3C)	See **malloc**(3C)
free(3X)	See **bsdmalloc**(3X)

free(3X)	See **malloc**(3X)
free(3X)	See **mapmalloc**(3X)
free_field(3X)	See **form_field_new**(3X)
free_fieldtype(3X)	See **form_fieldtype**(3X)
free_form(3X)	See **form_new**(3X)
free_item(3X)	See **menu_item_new**(3X)
free_menu(3X)	See **menu_new**(3X)
freenetconfigent(3N)	See **getnetconfig**(3N)
freopen(3B)	See **fopen**(3B)
freopen(3S)	See **fopen**(3S)
frexp(3C)	manipulate parts of floating-point numbers
fscanf(3S)	See **scanf**(3S)
fseek(3S)	reposition a file pointer in a stream
fsetpos(3C)	reposition a file pointer in a stream
fsync(3C)	synchronize a file's in-memory state with that on the physical medium
ftell(3S)	See **fseek**(3S)
ftime(3C)	get date and time
ftok(3C)	See **stdipc**(3C)
ftrylockfile(3S)	See **flockfile**(3S)
ftruncate(3C)	See **truncate**(3C)
ftw(3C)	walk a file tree
func_to_decimal(3)	See **string_to_decimal**(3)
funlockfile(3S)	See **flockfile**(3S)
fwrite(3S)	See **fread**(3S)
gamma(3M)	See **lgamma**(3M)
gamma_r(3M)	See **lgamma**(3M)
gconvert(3)	See **econvert**(3)
gcvt(3)	See **econvert**(3)
gcvt(3C)	See **ecvt**(3C)
getacdir(3)	See **getacinfo**(3)
getacflg(3)	See **getacinfo**(3)
getacinfo(3)	get audit control file information
getacmin(3)	See **getacinfo**(3)
getacna(3)	See **getacinfo**(3)

getauclassent(3)	get audit_class entry
getauclassent_r(3)	See **getauclassent**(3)
getauclassnam(3)	See **getauclassent**(3)
getauclassnam_r(3)	See **getauclassent**(3)
getauditflags(3)	convert audit flag specifications
getauditflagsbin(3)	See **getauditflags**(3)
getauditflagschar(3)	See **getauditflags**(3)
getauevent(3)	get audit_user entry
getauevent_r(3)	See **getauevent**(3)
getauevnam(3)	See **getauevent**(3)
getauevnam_r(3)	See **getauevent**(3)
getauevnonam(3)	See **getauevent**(3)
getauevnum(3)	See **getauevent**(3)
getauevnum_r(3)	See **getauevent**(3)
getauuserent(3)	See **getauusernam**(3)
getauusernam(3)	get audit_user entry
getbegyx(3X)	See **curs_getyx**(3X)
getc(3S)	get character or word from a stream
getch(3X)	See **curs_getch**(3X)
getchar(3S)	See **getc**(3S)
getchar_unlocked(3S)	See **getc**(3S)
getc_unlocked(3S)	See **getc**(3S)
getcwd(3C)	get pathname of current working directory
getdate(3C)	convert user format date and time
getdtablesize(3C)	get file descriptor table size
getenv(3C)	return value for environment name
getfauditflags(3)	generates the process audit state
getgrent(3C)	See **getgrnam**(3C)
getgrent_r(3C)	See **getgrnam**(3C)
getgrgid(3C)	See **getgrnam**(3C)
getgrgid_r(3C)	See **getgrnam**(3C)
getgrnam(3C)	get group entry
getgrnam_r(3C)	See **getgrnam**(3C)
gethostbyaddr(3N)	See **gethostbyname**(3N)

gethostbyaddr_r(3N)	See **gethostbyname**(3N)
gethostbyname(3N)	get network host entry
gethostbyname_r(3N)	See **gethostbyname**(3N)
gethostent(3N)	See **gethostbyname**(3N)
gethostent_r(3N)	See **gethostbyname**(3N)
gethostid(3C)	get unique identifier of current host
gethostname(3C)	get / set name of current host
gethrtime(3C)	get high resolution time
gethrvtime(3C)	See **gethrtime**(3C)
getlogin(3C)	get login name
getlogin_r(3C)	See **getlogin**(3C)
getmaxyx(3X)	See **curs_getyx**(3X)
getmntany(3C)	See **getmntent**(3C)
getmntent(3C)	get mnttab file information
get_myaddress(3N)	See **rpc_soc**(3N)
getnetbyaddr(3N)	See **getnetbyname**(3N)
getnetbyaddr_r(3N)	See **getnetbyname**(3N)
getnetbyname(3N)	get network entry
getnetbyname_r(3N)	See **getnetbyname**(3N)
getnetconfig(3N)	get network configuration database entry
getnetconfigent(3N)	See **getnetconfig**(3N)
getnetent(3N)	See **getnetbyname**(3N)
getnetent_r(3N)	See **getnetbyname**(3N)
getnetgrent(3N)	get network group entry
getnetgrent_r(3N)	See **getnetgrent**(3N)
getnetname(3N)	See **secure_rpc**(3N)
getnetpath(3N)	get / etc / netconfig entry corresponding to NETPATH component
getnwstr(3X)	See **curs_getwstr**(3X)
getopt(3C)	get option letter from argument vector
getpagesize(3C)	get system page size
getparyx(3X)	See **curs_getyx**(3X)
getpass(3C)	read a password
getpeername(3N)	get name of connected peer

getpriority(3C)	get / set scheduling priority for process, process group or user
getprotobyname(3N)	get protocol entry
getprotobyname_r(3N)	See **getprotobyname**(3N)
getprotobynumber(3N)	See **getprotobyname**(3N)
getprotobynumber_r(3N)	See **getprotobyname**(3N)
getprotoent(3N)	See **getprotobyname**(3N)
getprotoent_r(3N)	See **getprotobyname**(3N)
getpublickey(3N)	retrieve public or secret key
getpw(3C)	get passwd entry from UID
getpwent(3C)	See **getpwnam**(3C)
getpwent_r(3C)	See **getpwnam**(3C)
getpwnam(3C)	get password entry
getpwnam_r(3C)	See **getpwnam**(3C)
getpwuid(3C)	See **getpwnam**(3C)
getpwuid_r(3C)	See **getpwnam**(3C)
getrpcbyname(3N)	get RPC entry
getrpcbyname_r(3N)	See **getrpcbyname**(3N)
getrpcbynumber(3N)	See **getrpcbyname**(3N)
getrpcbynumber_r(3N)	See **getrpcbyname**(3N)
getrpcent(3N)	See **getrpcbyname**(3N)
getrpcent_r(3N)	See **getrpcbyname**(3N)
getrpcport(3N)	See **rpc_soc**(3N)
getrusage(3C)	get information about resource utilization
gets(3S)	get a string from a stream
getsecretkey(3N)	See **getpublickey**(3N)
getservbyname(3N)	get service entry
getservbyname_r(3N)	See **getservbyname**(3N)
getservbyport(3N)	See **getservbyname**(3N)
getservbyport_r(3N)	See **getservbyname**(3N)
getservent(3N)	See **getservbyname**(3N)
getservent_r(3N)	See **getservbyname**(3N)
getsockname(3N)	get socket name
getsockopt(3N)	get and set options on sockets
getspent(3C)	See **getspnam**(3C)

getspent_r(3C)	See **getspnam**(3C)
getspnam(3C)	get password entry
getspnam_r(3C)	See **getspnam**(3C)
getstr(3X)	See **curs_getstr**(3X)
getsubopt(3C)	parse suboptions from a string
getsyx(3X)	See **curs_kernel**(3X)
gettext(3I)	message handling functions
gettimeofday(3B)	get or set the date and time
gettimeofday(3C)	get or set the date and time
gettxt(3C)	retrieve a text string
getusershell(3C)	get legal user shells
getutent(3C)	access utmp file entry
getutid(3C)	See **getutent**(3C)
getutline(3C)	See **getutent**(3C)
getutmp(3C)	See **getutxent**(3C)
getutmpx(3C)	See **getutxent**(3C)
getutxent(3C)	access utmpx file entry
getutxid(3C)	See **getutxent**(3C)
getutxline(3C)	See **getutxent**(3C)
getvfsany(3C)	See **getvfsent**(3C)
getvfsent(3C)	get vfstab file entry
getvfsfile(3C)	See **getvfsent**(3C)
getvfsspec(3C)	See **getvfsent**(3C)
getw(3S)	See **getc**(3S)
getwc(3I)	convert EUC character from the stream to Process Code
getwch(3X)	See **curs_getwch**(3X)
getwchar(3I)	See **getwc**(3I)
getwd(3C)	get current working directory path-name
getwidth(3I)	get codeset information
getwin(3X)	See **curs_util**(3X)
getws(3I)	convert a string of EUC characters from the stream to Process Code
getwstr(3X)	See **curs_getwstr**(3X)
getyx(3X)	See **curs_getyx**(3X)

glob(3C)	generate path names matching a pattern
globfree(3C)	See **glob**(3C)
gmatch(3G)	shell global pattern matching
gmtime(3C)	See **ctime**(3C)
gmtime_r(3C)	See **ctime**(3C)
grantpt(3C)	grant access to the slave pseudo-terminal device
gsignal(3C)	See **ssignal**(3C)
halfdelay(3X)	See **curs_inopts**(3X)
has_colors(3X)	See **curs_color**(3X)
has_ic(3X)	See **curs_termattrs**(3X)
has_il(3X)	See **curs_termattrs**(3X)
hasmntopt(3C)	See **getmntent**(3C)
havedisk(3N)	See **rstat**(3N)
hcreate(3C)	See **hsearch**(3C)
hdestroy(3C)	See **hsearch**(3C)
hide_panel(3X)	See **panel_show**(3X)
host2netname(3N)	See **secure_rpc**(3N)
hsearch(3C)	manage hash search tables
htonl(3N)	See **byteorder**(3N)
htons(3N)	See **byteorder**(3N)
hyperbolic(3M)	hyperbolic functions
hypot(3M)	Euclidean distance
iconv(3)	code conversion function
iconv_close(3)	code conversion deallocation function
iconv_open(3)	code conversion allocation function
idcok(3X)	See **curs_outopts**(3X)
idlok(3X)	See **curs_outopts**(3X)
ieee_functions(3M)	appendix and related miscellaneous functions for IEEE arithmetic
ieee_test(3M)	IEEE test functions for verifying standard compliance
ilogb(3M)	See **ieee_functions**(3M)
immedok(3X)	See **curs_outopts**(3X)

inch(3X)	See **curs_inch**(3X)
inchnstr(3X)	See **curs_inchstr**(3X)
inchstr(3X)	See **curs_inchstr**(3X)
index(3C)	string operations
inet(3N)	Internet address manipulation
inet_addr(3N)	See **inet**(3N)
inet_lnaof(3N)	See **inet**(3N)
inet_makeaddr(3N)	See **inet**(3N)
inet_netof(3N)	See **inet**(3N)
inet_network(3N)	See **inet**(3N)
inet_ntoa(3N)	See **inet**(3N)
init_color(3X)	See **curs_color**(3X)
initgroups(3C)	initialize the supplementary group access list
init_pair(3X)	See **curs_color**(3X)
initscr(3X)	See **curs_initscr**(3X)
initstate(3C)	See **random**(3C)
innetgr(3N)	See **getnetgrent**(3N)
innstr(3X)	See **curs_instr**(3X)
innwstr(3X)	See **curs_inwstr**(3X)
insch(3X)	See **curs_insch**(3X)
insdelln(3X)	See **curs_deleteln**(3X)
insertln(3X)	See **curs_deleteln**(3X)
insnstr(3X)	See **curs_insstr**(3X)
insnwstr(3X)	See **curs_inswstr**(3X)
insque(3C)	insert / remove element from a queue
insstr(3X)	See **curs_insstr**(3X)
instr(3X)	See **curs_instr**(3X)
inswch(3X)	See **curs_inswch**(3X)
inswstr(3X)	See **curs_inswstr**(3X)
intrflush(3X)	See **curs_inopts**(3X)
inwch(3X)	See **curs_inwch**(3X)
inwchnstr(3X)	See **curs_inwchstr**(3X)
inwchstr(3X)	See **curs_inwchstr**(3X)
inwstr(3X)	See **curs_inwstr**(3X)

isalnum(3C)	See **ctype**(3C)
isalpha(3C)	See **ctype**(3C)
isascii(3C)	See **ctype**(3C)
isastream(3C)	test a file descriptor
isatty(3C)	See **ttyname**(3C)
iscntrl(3C)	See **ctype**(3C)
isdigit(3C)	See **ctype**(3C)
isencrypt(3G)	determine whether a buffer of characters is encrypted
isendwin(3X)	See **curs_initscr**(3X)
isenglish(3I)	See **iswalpha**(3I)
isgraph(3C)	See **ctype**(3C)
isideogram(3I)	See **iswalpha**(3I)
is_linetouched(3X)	See **curs_touch**(3X)
islower(3C)	See **ctype**(3C)
isnan(3C)	determine type of floating-point number
isnan(3M)	See **ieee_functions**(3M)
isnand(3C)	See **isnan**(3C)
isnanf(3C)	See **isnan**(3C)
isnumber(3I)	See **iswalpha**(3I)
isphonogram(3I)	See **iswalpha**(3I)
isprint(3C)	See **ctype**(3C)
ispunct(3C)	See **ctype**(3C)
isspace(3C)	See **ctype**(3C)
isspecial(3I)	See **iswalpha**(3I)
isupper(3C)	See **ctype**(3C)
iswalnum(3I)	See **iswalpha**(3I)
iswalpha(3I)	Process Code character classification macros and functions
iswascii(3I)	See **iswalpha**(3I)
iswcntrl(3I)	See **iswalpha**(3I)
iswctype(3I)	test character for specified class
iswdigit(3I)	See **iswalpha**(3I)
iswgraph(3I)	See **iswalpha**(3I)
is_wintouched(3X)	See **curs_touch**(3X)

iswlower(3I)	See **iswalpha**(3I)
iswprint(3I)	See **iswalpha**(3I)
iswpunct(3I)	See **iswalpha**(3I)
iswspace(3I)	See **iswalpha**(3I)
iswupper(3I)	See **iswalpha**(3I)
iswxdigit(3I)	See **iswalpha**(3I)
isxdigit(3C)	See **ctype**(3C)
item_count(3X)	See **menu_items**(3X)
item_description(3X)	See **menu_item_name**(3X)
item_index(3X)	See **menu_item_current**(3X)
item_init(3X)	See **menu_hook**(3X)
item_name(3X)	See **menu_item_name**(3X)
item_opts(3X)	See **menu_item_opts**(3X)
item_opts_off(3X)	See **menu_item_opts**(3X)
item_opts_on(3X)	See **menu_item_opts**(3X)
item_term(3X)	See **menu_hook**(3X)
item_userptr(3X)	See **menu_item_userptr**(3X)
item_value(3X)	See **menu_item_value**(3X)
item_visible(3X)	See **menu_item_visible**(3X)
j0(3M)	See **bessel**(3M)
j1(3M)	See **bessel**(3M)
jn(3M)	See **bessel**(3M)
jrand48(3C)	See **drand48**(3C)
kerberos(3N)	Kerberos authentication library
kerberos_rpc(3N)	library routines for remote procedure calls using Kerberos authentication
key_decryptsession(3N)	See **secure_rpc**(3N)
key_encryptsession(3N)	See **secure_rpc**(3N)
key_gendes(3N)	See **secure_rpc**(3N)
keyname(3X)	See **curs_util**(3X)
keypad(3X)	See **curs_inopts**(3X)
key_secretkey_is_set(3N)	See **secure_rpc**(3N)
key_setsecret(3N)	See **secure_rpc**(3N)
killchar(3X)	See **curs_termattrs**(3X)
killpg(3C)	send signal to a process group

krb_get_admhst(3N)	See **krb_realmofhost**(3N)
krb_get_cred(3N)	See **kerberos**(3N)
krb_get_krbhst(3N)	See **krb_realmofhost**(3N)
krb_get_lrealm(3N)	See **krb_realmofhost**(3N)
krb_get_phost(3N)	See **krb_realmofhost**(3N)
krb_kntoln(3N)	See **kerberos**(3N)
krb_mk_err(3N)	See **kerberos**(3N)
krb_mk_req(3N)	See **kerberos**(3N)
krb_mk_safe(3N)	See **kerberos**(3N)
krb_net_read(3N)	See **krb_sendauth**(3N)
krb_net_write(3N)	See **krb_sendauth**(3N)
krb_rd_err(3N)	See **kerberos**(3N)
krb_rd_req(3N)	See **kerberos**(3N)
krb_rd_safe(3N)	See **kerberos**(3N)
krb_realmofhost(3N)	additional Kerberos utility routines
krb_recvauth(3N)	See **krb_sendauth**(3N)
krb_sendauth(3N)	Kerberos routines for sending authentication via network stream sockets
krb_set_key(3N)	See **kerberos**(3N)
krb_set_tkt_string(3N)	set Kerberos ticket cache file name
kstat(3K)	kernel statistics facility
kstat_chain_update(3K)	update the kstat header chain
kstat_close(3K)	See **kstat_open**(3K)
kstat_data_lookup(3K)	See **kstat_lookup**(3K)
kstat_lookup(3K)	find a kstat by name
kstat_open(3K)	initialize kernel statistics facility
kstat_read(3K)	read or write kstat data
kstat_write(3K)	See **kstat_read**(3K)
kvm_close(3K)	See **kvm_open**(3K)
kvm_getcmd(3K)	See **kvm_getu**(3K)
kvm_getproc(3K)	See **kvm_nextproc**(3K)
kvm_getu(3K)	get the u-area or invocation arguments for a process
kvm_kread(3K)	See **kvm_read**(3K)
kvm_kwrite(3K)	See **kvm_read**(3K)

kvm_nextproc(3K)	read system process structures
kvm_nlist(3K)	get entries from kernel symbol table
kvm_open(3K)	specify a kernel to examine
kvm_read(3K)	copy data to or from a kernel image or running system
kvm_setproc(3K)	See **kvm_nextproc**(3K)
kvm_uread(3K)	See **kvm_read**(3K)
kvm_uwrite(3K)	See **kvm_read**(3K)
kvm_write(3K)	See **kvm_read**(3K)
l64a(3C)	See **a64l**(3C)
label(3)	See **plot**(3)
labs(3C)	See **abs**(3C)
lckpwdf(3C)	manipulate shadow password database lock file
lcong48(3C)	See **drand48**(3C)
ldexp(3C)	See **frexp**(3C)
ldiv(3C)	See **div**(3C)
leaveok(3X)	See **curs_outopts**(3X)
lfind(3C)	See **lsearch**(3C)
lfmt(3C)	display error message in standard format and pass to logging and monitoring services
lgamma(3M)	log gamma function
lgamma_r(3M)	See **lgamma**(3M)
libpthread(3T)	See **threads**(3T)
libthread(3T)	See **threads**(3T)
libthread_db(3T)	interface to libthread threads information
line(3)	See **plot**(3)
link_field(3X)	See **form_field_new**(3X)
link_fieldtype(3X)	See **form_fieldtype**(3X)
linmod(3)	See **plot**(3)
lio_listio(3R)	list directed I/O
listen(3N)	listen for connections on a socket
llabs(3C)	See **abs**(3C)
lldiv(3C)	See **div**(3C)

lltostr(3C)	See **strtol**(3C)
localeconv(3C)	get numeric formatting information
localtime(3C)	See **ctime**(3C)
localtime_r(3C)	See **ctime**(3C)
lockf(3C)	record locking on files
log(3M)	See **exp**(3M)
log10(3M)	See **exp**(3M)
log1p(3M)	See **exp**(3M)
logb(3C)	See **frexp**(3C)
logb(3M)	See **ieee_test**(3M)
_longjmp(3B)	See **setjmp**(3B)
longjmp(3B)	See **setjmp**(3B)
longjmp(3C)	See **setjmp**(3C)
longname(3X)	See **curs_termattrs**(3X)
lrand48(3C)	See **drand48**(3C)
lsearch(3C)	linear search and update
madvise(3)	provide advice to VM system
maillock(3X)	manage lockfile for user's mailbox
major(3C)	See **makedev**(3C)
makecontext(3C)	manipulate user contexts
makedev(3C)	manage a device number
mallinfo(3X)	See **malloc**(3X)
malloc(3C)	memory allocator
malloc(3X)	memory allocator
malloc(3X)	See **bsdmalloc**(3X)
mallopt(3X)	See **malloc**(3X)
mapmalloc(3X)	memory allocator
matherr(3M)	math library exception-handling function
mbchar(3C)	multibyte character handling
mblen(3C)	See **mbchar**(3C)
mbstowcs(3C)	See **mbstring**(3C)
mbstring(3C)	multibyte string functions
mbtowc(3C)	See **mbchar**(3C)
mctl(3B)	memory management control

media_findname(3X)	convert a supplied name into an absolute pathname that can be used to access removable media
media_getattr(3X)	get and set media attributes
media_setattr(3X)	See **media_getattr**(3X)
memalign(3C)	See **malloc**(3C)
memccpy(3C)	See **memory**(3C)
memchr(3C)	See **memory**(3C)
memcmp(3C)	See **memory**(3C)
memcpy(3C)	See **memory**(3C)
memmove(3C)	See **memory**(3C)
memory(3C)	memory operations
memset(3C)	See **memory**(3C)
menu_attributes(3X)	control menus display attributes
menu_back(3X)	See **menu_attributes**(3X)
menu_cursor(3X)	correctly position a menus cursor
menu_driver(3X)	command processor for the menus subsystem
menu_fore(3X)	See **menu_attributes**(3X)
menu_format(3X)	set and get maximum numbers of rows and columns in menus
menu_grey(3X)	See **menu_attributes**(3X)
menu_hook(3X)	assign application-specific routines for automatic invocation by menus
menu_init(3X)	See **menu_hook**(3X)
menu_item_current(3X)	set and get current menus items
menu_item_name(3X)	get menus item name and description
menu_item_new(3X)	create and destroy menus items
menu_item_opts(3X)	menus item option routines
menu_items(3X)	connect and disconnect items to and from menus
menu_item_userptr(3X)	associate application data with menus items
menu_item_value(3X)	set and get menus item values
menu_item_visible(3X)	tell if menus item is visible
menu_mark(3X)	menus mark string routines

menu_new(3X)	create and destroy menus
menu_opts(3X)	menus option routines
menu_opts_off(3X)	See **menu_opts**(3X)
menu_opts_on(3X)	See **menu_opts**(3X)
menu_pad(3X)	See **menu_attributes**(3X)
menu_pattern(3X)	set and get menus pattern match buffer
menu_post(3X)	write or erase menus from associated subwindows
menus(3X)	character based menus package
menu_sub(3X)	See **menu_win**(3X)
menu_term(3X)	See **menu_hook**(3X)
menu_userptr(3X)	associate application data with menus
menu_win(3X)	menus window and subwindow association routines
meta(3X)	See **curs_inopts**(3X)
minor(3C)	See **makedev**(3C)
mkdirp(3G)	create, remove directories in a path
mkfifo(3C)	create a new FIFO
mktemp(3C)	make a unique file name
mktime(3C)	converts a tm structure to a calendar time
mlock(3C)	lock (or unlock) pages in memory
mlockall(3C)	lock or unlock address space
modf(3C)	See **frexp**(3C)
modff(3C)	See **frexp**(3C)
monitor(3C)	prepare process execution profile
move(3)	See **plot**(3)
move(3X)	See **curs_move**(3X)
move_field(3X)	See **form_field**(3X)
movenextch(3X)	See **curs_alecompat**(3X)
move_panel(3X)	See **panel_move**(3X)
moveprevch(3X)	See **curs_alecompat**(3X)
mq_close(3R)	close a message queue
mq_getattr(3R)	See **mq_setattr**(3R)

mq_notify(3R)	notify process (or thread) that a message is available on a queue
mq_open(3R)	open a message queue
mq_receive(3R)	receive a message from a message queue
mq_send(3R)	send a message to a message queue
mq_setattr(3R)	set / get message queue attributes
mq_unlink(3R)	remove a message queue
mrand48(3C)	See **drand48**(3C)
msync(3C)	synchronize memory with physical storage
munlock(3C)	See **mlock**(3C)
munlockall(3C)	See **mlockall**(3C)
mutex(3T)	mutual exclusion locks
mutex_destroy(3T)	See **mutex**(3T)
mutex_init(3T)	See **mutex**(3T)
mutex_lock(3T)	See **mutex**(3T)
mutex_trylock(3T)	See **mutex**(3T)
mutex_unlock(3T)	See **mutex**(3T)
mvaddch(3X)	See **curs_addch**(3X)
mvaddchnstr(3X)	See **curs_addchstr**(3X)
mvaddchstr(3X)	See **curs_addchstr**(3X)
mvaddnstr(3X)	See **curs_addstr**(3X)
mvaddnwstr(3X)	See **curs_addwstr**(3X)
mvaddstr(3X)	See **curs_addstr**(3X)
mvaddwch(3X)	See **curs_addwch**(3X)
mvaddwchnstr(3X)	See **curs_addwchstr**(3X)
mvaddwchstr(3X)	See **curs_addwchstr**(3X)
mvaddwstr(3X)	See **curs_addwstr**(3X)
mvcur(3X)	See **curs_terminfo**(3X)
mvdelch(3X)	See **curs_delch**(3X)
mvderwin(3X)	See **curs_window**(3X)
mvgetch(3X)	See **curs_getch**(3X)
mvgetnwstr(3X)	See **curs_getwstr**(3X)
mvgetstr(3X)	See **curs_getstr**(3X)
mvgetwch(3X)	See **curs_getwch**(3X)

mvgetwstr(3X)	See **curs_getwstr**(3X)
mvinch(3X)	See **curs_inch**(3X)
mvinchnstr(3X)	See **curs_inchstr**(3X)
mvinchstr(3X)	See **curs_inchstr**(3X)
mvinnstr(3X)	See **curs_instr**(3X)
mvinnwstr(3X)	See **curs_inwstr**(3X)
mvinsch(3X)	See **curs_insch**(3X)
mvinsnstr(3X)	See **curs_insstr**(3X)
mvinsnwstr(3X)	See **curs_inswstr**(3X)
mvinsstr(3X)	See **curs_insstr**(3X)
mvinstr(3X)	See **curs_instr**(3X)
mvinswch(3X)	See **curs_inswch**(3X)
mvinswstr(3X)	See **curs_inswstr**(3X)
mvinwch(3X)	See **curs_inwch**(3X)
mvinwchnstr(3X)	See **curs_inwchstr**(3X)
mvinwchstr(3X)	See **curs_inwchstr**(3X)
mvinwstr(3X)	See **curs_inwstr**(3X)
mvprintw(3X)	See **curs_printw**(3X)
mvscanw(3X)	See **curs_scanw**(3X)
mvwaddch(3X)	See **curs_addch**(3X)
mvwaddchnstr(3X)	See **curs_addchstr**(3X)
mvwaddchstr(3X)	See **curs_addchstr**(3X)
mvwaddnstr(3X)	See **curs_addstr**(3X)
mvwaddnwstr(3X)	See **curs_addwstr**(3X)
mvwaddstr(3X)	See **curs_addstr**(3X)
mvwaddwch(3X)	See **curs_addwch**(3X)
mvwaddwchnstr(3X)	See **curs_addwchstr**(3X)
mvwaddwchstr(3X)	See **curs_addwchstr**(3X)
mvwaddwstr(3X)	See **curs_addwstr**(3X)
mvwdelch(3X)	See **curs_delch**(3X)
mvwgetch(3X)	See **curs_getch**(3X)
mvwgetnwstr(3X)	See **curs_getwstr**(3X)
mvwgetstr(3X)	See **curs_getstr**(3X)
mvwgetwch(3X)	See **curs_getwch**(3X)
mvwgetwstr(3X)	See **curs_getwstr**(3X)

mvwin(3X)	See **curs_window**(3X)
mvwinch(3X)	See **curs_inch**(3X)
mvwinchnstr(3X)	See **curs_inchstr**(3X)
mvwinchstr(3X)	See **curs_inchstr**(3X)
mvwinnstr(3X)	See **curs_instr**(3X)
mvwinnwstr(3X)	See **curs_inwstr**(3X)
mvwinsch(3X)	See **curs_insch**(3X)
mvwinsnstr(3X)	See **curs_insstr**(3X)
mvwinsnwstr(3X)	See **curs_inswstr**(3X)
mvwinsstr(3X)	See **curs_insstr**(3X)
mvwinstr(3X)	See **curs_instr**(3X)
mvwinswch(3X)	See **curs_inswch**(3X)
mvwinswstr(3X)	See **curs_inswstr**(3X)
mvwinwch(3X)	See **curs_inwch**(3X)
mvwinwchnstr(3X)	See **curs_inwchstr**(3X)
mvwinwchstr(3X)	See **curs_inwchstr**(3X)
mvwinwstr(3X)	See **curs_inwstr**(3X)
mvwprintw(3X)	See **curs_printw**(3X)
mvwscanw(3X)	See **curs_scanw**(3X)
nanosleep(3R)	high resolution sleep
napms(3X)	See **curs_kernel**(3X)
nc_perror(3N)	See **getnetconfig**(3N)
nc_sperror(3N)	See **getnetconfig**(3N)
ndbm(3)	data base subroutines
netdir(3N)	generic transport name-to-address translation
netdir_free(3N)	See **netdir**(3N)
netdir_getbyaddr(3N)	See **netdir**(3N)
netdir_getbyname(3N)	See **netdir**(3N)
netdir_mergeaddr(3N)	See **netdir**(3N)
netdir_options(3N)	See **netdir**(3N)
netdir_perror(3N)	See **netdir**(3N)
netdir_sperror(3N)	See **netdir**(3N)
netname2host(3N)	See **secure_rpc**(3N)
netname2user(3N)	See **secure_rpc**(3N)

new_field(3X)	See **form_field_new**(3X)
new_fieldtype(3X)	See **form_fieldtype**(3X)
new_form(3X)	See **form_new**(3X)
new_item(3X)	See **menu_item_new**(3X)
new_menu(3X)	See **menu_new**(3X)
newpad(3X)	See **curs_pad**(3X)
new_page(3X)	See **form_new_page**(3X)
new_panel(3X)	See **panel_new**(3X)
newterm(3X)	See **curs_initscr**(3X)
newwin(3X)	See **curs_window**(3X)
nextafter(3C)	See **frexp**(3C)
nextafter(3M)	See **ieee_functions**(3M)
nextkey(3B)	See **dbm**(3B)
nftw(3C)	See **ftw**(3C)
nice(3B)	change priority of a process
nis_add(3N)	See **nis_names**(3N)
nis_add_entry(3N)	See **nis_tables**(3N)
nis_addmember(3N)	See **nis_groups**(3N)
nis_checkpoint(3N)	See **nis_ping**(3N)
nis_clone_object(3N)	See **nis_subr**(3N)
nis_creategroup(3N)	See **nis_groups**(3N)
nis_db(3N)	NIS+ Database access functions
nis_destroygroup(3N)	See **nis_groups**(3N)
nis_destroy_object(3N)	See **nis_subr**(3N)
nis_dir_cmp(3N)	See **nis_subr**(3N)
nis_domain_of(3N)	See **nis_subr**(3N)
nis_error(3N)	display NIS+ error messages
nis_first_entry(3N)	See **nis_tables**(3N)
nis_freenames(3N)	See **nis_subr**(3N)
nis_freeresult(3N)	See **nis_names**(3N)
nis_freeservlist(3N)	See **nis_server**(3N)
nis_freetags(3N)	See **nis_server**(3N)
nis_getnames(3N)	See **nis_subr**(3N)
nis_getservlist(3N)	See **nis_server**(3N)
nis_groups(3N)	NIS+ group manipulation functions

nis_ismember(3N)	See **nis_groups**(3N)
nis_leaf_of(3N)	See **nis_subr**(3N)
nis_lerror(3N)	See **nis_error**(3N)
nis_list(3N)	See **nis_tables**(3N)
nis_local_directory(3N)	See **nis_local_names**(3N)
nis_local_group(3N)	See **nis_local_names**(3N)
nis_local_host(3N)	See **nis_local_names**(3N)
nis_local_names(3N)	NIS+ local names
nis_local_principal(3N)	See **nis_local_names**(3N)
nis_lookup(3N)	See **nis_names**(3N)
__nis_map_group(3N)	See **nis_groups**(3N)
nis_map_group(3N)	See **nis_groups**(3N)
nis_mkdir(3N)	See **nis_server**(3N)
nis_modify(3N)	See **nis_names**(3N)
nis_modify_entry(3N)	See **nis_tables**(3N)
nis_name_of(3N)	See **nis_subr**(3N)
nis_names(3N)	NIS+ namespace functions
nis_next_entry(3N)	See **nis_tables**(3N)
nis_objects(3N)	NIS+ object formats
nis_perror(3N)	See **nis_error**(3N)
nis_ping(3N)	misc NIS+ log administration functions
nis_print_group_entry(3N)	See **nis_groups**(3N)
nis_print_object(3N)	See **nis_subr**(3N)
nis_remove(3N)	See **nis_names**(3N)
nis_remove_entry(3N)	See **nis_tables**(3N)
nis_removemember(3N)	See **nis_groups**(3N)
nis_rmdir(3N)	See **nis_server**(3N)
nis_server(3N)	miscellaneous NIS+ functions
nis_servstate(3N)	See **nis_server**(3N)
nis_sperrno(3N)	See **nis_error**(3N)
nis_sperror(3N)	See **nis_error**(3N)
nis_sperror_r(3N)	See **nis_error**(3N)
nis_stats(3N)	See **nis_server**(3N)
nis_subr(3N)	NIS+ subroutines

nis_tables(3N)	NIS+ table functions
nis_verifygroup(3N)	See **nis_groups**(3N)
nl(3X)	See **curs_outopts**(3X)
nlist(3B)	get entries from symbol table
nlist(3E)	get entries from name list
nl_langinfo(3C)	language information
nlsgetcall(3N)	get client's data passed via the listener
nlsprovider(3N)	get name of transport provider
nlsrequest(3N)	format and send listener service request message
nocbreak(3X)	See **curs_inopts**(3X)
nodelay(3X)	See **curs_inopts**(3X)
noecho(3X)	See **curs_inopts**(3X)
nonl(3X)	See **curs_outopts**(3X)
noqiflush(3X)	See **curs_inopts**(3X)
noraw(3X)	See **curs_inopts**(3X)
NOTE(3X)	annotate source code with info for tools
_NOTE(3X)	See **NOTE**(3X)
notimeout(3X)	See **curs_inopts**(3X)
nrand48(3C)	See **drand48**(3C)
ntohl(3N)	See **byteorder**(3N)
ntohs(3N)	See **byteorder**(3N)
offsetof(3C)	offset of structure member
opendir(3C)	See **directory**(3C)
openlog(3)	See **syslog**(3)
openpl(3)	See **plot**(3)
openvt(3)	See **plot**(3)
overlay(3X)	See **curs_overlay**(3X)
overwrite(3X)	See **curs_overlay**(3X)
p2close(3G)	See **p2open**(3G)
p2open(3G)	open, close pipes to and from a command
pair_content(3X)	See **curs_color**(3X)
panel_above(3X)	panels deck traversal primitives

panel_below(3X)	See **panel_above**(3X)
panel_hidden(3X)	See **panel_show**(3X)
panel_move(3X)	move a panels window on the virtual screen
panel_new(3X)	create and destroy panels
panels(3X)	character based panels package
panel_show(3X)	panels deck manipulation routines
panel_top(3X)	panels deck manipulation routines
panel_update(3X)	panels virtual screen refresh routine
panel_userptr(3X)	associate application data with a panels panel
panel_window(3X)	get or set the current window of a panels panel
pathfind(3G)	search for named file in named directories
pclose(3S)	See **popen**(3S)
pechochar(3X)	See **curs_pad**(3X)
pechowchar(3X)	See **curs_pad**(3X)
perror(3C)	print system error messages
pfmt(3C)	display error message in standard format
plock(3C)	lock or unlock into memory process, text, or data
plot(3)	graphics interface
pmap_getmaps(3N)	See **rpc_soc**(3N)
pmap_getport(3N)	See **rpc_soc**(3N)
pmap_rmtcall(3N)	See **rpc_soc**(3N)
pmap_set(3N)	See **rpc_soc**(3N)
pmap_unset(3N)	See **rpc_soc**(3N)
pnoutrefresh(3X)	See **curs_pad**(3X)
point(3)	See **plot**(3)
popen(3S)	initiate pipe to / from a process
pos_form_cursor(3X)	See **form_cursor**(3X)
pos_menu_cursor(3X)	See **menu_cursor**(3X)
post_form(3X)	See **form_post**(3X)
post_menu(3X)	See **menu_post**(3X)

pow(3M)	See **exp**(3M)
prefresh(3X)	See **curs_pad**(3X)
printf(3B)	formatted output conversion
printf(3S)	print formatted output
printw(3X)	See **curs_printw**(3X)
proc_service(3T)	process service interface
psiginfo(3C)	See **psignal**(3C)
psignal(3B)	system signal messages
psignal(3C)	system signal messages
ps_lcontinue(3T)	See **proc_service**(3T)
ps_lgetfpregs(3T)	See **proc_service**(3T)
ps_lgetregs(3T)	See **proc_service**(3T)
ps_lgetxregs(3T)	See **proc_service**(3T)
ps_lgetxregsize(3T)	See **proc_service**(3T)
ps_lsetfpregs(3T)	See **proc_service**(3T)
ps_lsetregs(3T)	See **proc_service**(3T)
ps_lsetxregs(3T)	See **proc_service**(3T)
ps_lstop(3T)	See **proc_service**(3T)
ps_pcontinue(3T)	See **proc_service**(3T)
ps_pdread(3T)	See **proc_service**(3T)
ps_pdwrite(3T)	See **proc_service**(3T)
ps_pglobal_lookup(3T)	See **proc_service**(3T)
ps_plog(3T)	See **proc_service**(3T)
ps_pstop(3T)	See **proc_service**(3T)
ps_ptread(3T)	See **proc_service**(3T)
ps_ptwrite(3T)	See **proc_service**(3T)
pthread_atfork(3T)	register fork handlers
pthread_attr_destroy(3T)	See **pthread_attr_init**(3T)
pthread_attr_getdetachstate(3T)	See **pthread_attr_init**(3T)
pthread_attr_getinheritsched(3T)	See **pthread_attr_init**(3T)
pthread_attr_getschedparam(3T)	See **pthread_attr_init**(3T)
pthread_attr_getschedpolicy(3T)	See **pthread_attr_init**(3T)
pthread_attr_getscope(3T)	See **pthread_attr_init**(3T)
pthread_attr_getstackaddr(3T)	See **pthread_attr_init**(3T)
pthread_attr_getstacksize(3T)	See **pthread_attr_init**(3T)

pthread_attr_init(3T)	thread creation attributes
pthread_attr_setdetachstate(3T)	See **pthread_attr_init**(3T)
pthread_attr_setinheritsched(3T)	See **pthread_attr_init**(3T)
pthread_attr_setschedparam(3T)	See **pthread_attr_init**(3T)
pthread_attr_setschedpolicy(3T)	See **pthread_attr_init**(3T)
pthread_attr_setscope(3T)	See **pthread_attr_init**(3T)
pthread_attr_setstackaddr(3T)	See **pthread_attr_init**(3T)
pthread_attr_setstacksize(3T)	See **pthread_attr_init**(3T)
pthread_cancel(3T)	See **cancellation**(3T)
pthread_cleanup_pop(3T)	See **cancellation**(3T)
pthread_cleanup_push(3T)	See **cancellation**(3T)
pthread_condattr_destroy(3T)	See **pthread_condattr_init**(3T)
pthread_condattr_getpshared(3T)	See **pthread_condattr_init**(3T)
pthread_condattr_init(3T)	condition variable initialization attributes
pthread_condattr_setpshared(3T)	See **pthread_condattr_init**(3T)
pthread_cond_broadcast(3T)	See **condition**(3T)
pthread_cond_destroy(3T)	See **condition**(3T)
pthread_cond_init(3T)	See **condition**(3T)
pthread_cond_signal(3T)	See **condition**(3T)
pthread_cond_timedwait(3T)	See **condition**(3T)
pthread_cond_wait(3T)	See **condition**(3T)
pthread_create(3T)	thread creation
pthread_detach(3T)	dynamically detaching a thread
pthread_equal(3T)	compare thread IDs
pthread_exit(3T)	thread termination
pthread_getschedparam(3T)	See **pthread_setschedparam**(3T)
pthread_getspecific(3T)	See **pthread_key_create**(3T)
pthread_join(3T)	wait for thread termination
pthread_key_create(3T)	thread-specific-data functions
pthread_key_delete(3T)	See **pthread_key_create**(3T)
pthread_kill(3T)	send a signal to a thread
pthread_mutexattr_destroy(3T)	See **pthread_mutexattr_init**(3T)
pthread_mutexattr_getprioceiling(3T)	See **pthread_mutexattr_init**(3T)
pthread_mutexattr_getprotocol(3T)	See **pthread_mutexattr_init**(3T)

pthread_mutexattr_getpshared(3T) See **pthread_mutexattr_init**(3T)
pthread_mutexattr_init(3T) mutex initialization attributes
pthread_mutexattr_setprioceiling(3T) See **pthread_mutexattr_init**(3T)
pthread_mutexattr_setprotocol(3T) See **pthread_mutexattr_init**(3T)
pthread_mutexattr_setpshared(3T) See **pthread_mutexattr_init**(3T)
pthread_mutex_destroy(3T) See **mutex**(3T)
pthread_mutex_getprioceiling(3T) See **pthread_mutex_setprioceiling**(3T)
pthread_mutex_init(3T) See **mutex**(3T)
pthread_mutex_lock(3T) See **mutex**(3T)
pthread_mutex_setprioceiling(3T) change the priority ceiling of a mutex
pthread_mutex_trylock(3T) See **mutex**(3T)
pthread_mutex_unlock(3T) See **mutex**(3T)
pthread_once(3T) dynamic package initialization
pthreads(3T) See **threads**(3T)
pthread_self(3T) get calling thread's ID
pthread_setcancelstate(3T) See **cancellation**(3T)
pthread_setcanceltype(3T) See **cancellation**(3T)
pthread_setschedparam(3T) dynamic access to thread scheduling
pthread_setspecific(3T) See **pthread_key_create**(3T)
pthread_sigmask(3T) change and / or examine calling thread's signal mask
pthread_testcancel(3T) See **cancellation**(3T)
ptsname(3C) get name of the slave pseudo-terminal device
publickey(3N) See **getpublickey**(3N)
putc(3S) put character or word on a stream
putchar(3S) See **putc**(3S)
putchar_unlocked(3S) See **putc**(3S)
putc_unlocked(3S) See **putc**(3S)
putenv(3C) change or add value to environment
putmntent(3C) See **getmntent**(3C)
putp(3X) See **curs_terminfo**(3X)
putpwent(3C) write password file entry
puts(3S) put a string on a stream
putspent(3C) write shadow password file entry

pututline(3C)	See **getutent**(3C)
pututxline(3C)	See **getutxent**(3C)
putw(3S)	See **putc**(3S)
putwc(3I)	convert Process Code character to EUC and put on a stream
putwchar(3I)	See **putwc**(3I)
putwin(3X)	See **curs_util**(3X)
putws(3I)	convert a string of Process Code characters to EUC characters and put it on a stream
qeconvert(3)	See **econvert**(3)
qfconvert(3)	See **econvert**(3)
qgconvert(3)	See **econvert**(3)
qiflush(3X)	See **curs_inopts**(3X)
qsort(3C)	quick sort
quadruple_to_decimal(3)	See **floating_to_decimal**(3)
rac_drop(3N)	See **rpc_rac**(3N)
rac_poll(3N)	See **rpc_rac**(3N)
rac_recv(3N)	See **rpc_rac**(3N)
rac_send(3N)	See **rpc_rac**(3N)
raise(3C)	send signal to program
rand(3B)	simple random number generator
rand(3C)	simple random-number generator
random(3C)	better random number generator; routines for changing generators
rand_r(3C)	See **rand**(3C)
raw(3X)	See **curs_inopts**(3X)
rcmd(3N)	routines for returning a stream to a remote command
readdir(3B)	read a directory entry
readdir(3C)	See **directory**(3C)
readdir_r(3C)	See **directory**(3C)
read_vtoc(3X)	read and write a disk's VTOC
realloc(3C)	See **malloc**(3C)
realloc(3X)	See **bsdmalloc**(3X)
realloc(3X)	See **malloc**(3X)

realloc(3X)	See **mapmalloc**(3X)
realpath(3C)	returns the real file name
reboot(3C)	reboot system or halt processor
re_comp(3C)	See **regex**(3C)
recv(3N)	receive a message from a socket
recvfrom(3N)	See **recv**(3N)
recvmsg(3N)	See **recv**(3N)
redrawwin(3X)	See **curs_refresh**(3X)
re_exec(3C)	See **regex**(3C)
refresh(3X)	See **curs_refresh**(3X)
regcmp(3G)	compile and execute regular expression
regcomp(3C)	regular expression matching
regerror(3C)	See **regcomp**(3C)
regex(3C)	regular expression handler
regex(3G)	See **regcmp**(3G)
regexec(3C)	See **regcomp**(3C)
regexpr(3G)	regular expression compile and match routines
regfree(3C)	See **regcomp**(3C)
registerrpc(3N)	See **rpc_soc**(3N)
remainder(3M)	See **ieee_functions**(3M)
remove(3C)	remove file
remque(3C)	See **insque**(3C)
replace_panel(3X)	See **panel_window**(3X)
reset_prog_mode(3X)	See **curs_kernel**(3X)
reset_shell_mode(3X)	See **curs_kernel**(3X)
resetty(3X)	See **curs_kernel**(3X)
res_init(3N)	See **resolver**(3N)
res_mkquery(3N)	See **resolver**(3N)
resolver(3N)	resolver routines
res_search(3N)	See **resolver**(3N)
res_send(3N)	See **resolver**(3N)
restartterm(3X)	See **curs_terminfo**(3X)
rewind(3S)	See **fseek**(3S)
rewinddir(3C)	See **directory**(3C)

rexec(3N)	return stream to a remote command
rindex(3C)	See **index**(3C)
rint(3M)	See **floor**(3M)
ripoffline(3X)	See **curs_kernel**(3X)
rmdirp(3G)	See **mkdirp**(3G)
rnusers(3N)	See **rusers**(3N)
rpc(3N)	library routines for remote procedure calls
rpcb_getaddr(3N)	See **rpcbind**(3N)
rpcb_getmaps(3N)	See **rpcbind**(3N)
rpcb_gettime(3N)	See **rpcbind**(3N)
rpcbind(3N)	library routines for RPC bind service
rpcb_rmtcall(3N)	See **rpcbind**(3N)
rpc_broadcast(3N)	See **rpc_clnt_calls**(3N)
rpc_broadcast_exp(3N)	See **rpc_clnt_calls**(3N)
rpcb_set(3N)	See **rpcbind**(3N)
rpcb_unset(3N)	See **rpcbind**(3N)
rpc_call(3N)	See **rpc_clnt_calls**(3N)
rpc_clnt_auth(3N)	library routines for client side remote procedure call authentication
rpc_clnt_calls(3N)	library routines for client side calls
rpc_clnt_create(3N)	library routines for dealing with creation and manipulation of CLIENT handles
rpc_control(3N)	library routine for manipulating global RPC attributes for client and server applications
rpc_createerr(3N)	See **rpc_clnt_create**(3N)
rpc_rac(3N)	remote asynchronous calls
rpc_reg(3N)	See **rpc_svc_reg**(3N)
rpc_soc(3N)	obsolete library routines for RPC
rpc_svc_calls(3N)	library routines for RPC servers
rpc_svc_create(3N)	library routines for the creation of server handles
rpc_svc_err(3N)	library routines for server side remote procedure call errors
rpc_svc_reg(3N)	library routines for registering

	servers
rpc_xdr(3N)	XDR library routines for remote procedure calls
rresvport(3N)	See **rcmd**(3N)
rstat(3N)	get performance data from remote kernel
ruserok(3N)	See **rcmd**(3N)
rusers(3N)	return information about users on remote machines
rwall(3N)	write to specified remote machines
rwlock(3T)	multiple readers, single writer locks
rwlock_destroy(3T)	See **rwlock**(3T)
rwlock_init(3T)	See **rwlock**(3T)
rw_rdlock(3T)	See **rwlock**(3T)
rw_tryrdlock(3T)	See **rwlock**(3T)
rw_trywrlock(3T)	See **rwlock**(3T)
rw_unlock(3T)	See **rwlock**(3T)
rw_wrlock(3T)	See **rwlock**(3T)
savetty(3X)	See **curs_kernel**(3X)
scalb(3C)	See **frexp**(3C)
scalb(3M)	See **ieee_test**(3M)
scalbn(3M)	See **ieee_functions**(3M)
scale_form(3X)	See **form_win**(3X)
scale_menu(3X)	See **menu_win**(3X)
scandir(3B)	scan a directory
scanf(3S)	convert formatted input
scanw(3X)	See **curs_scanw**(3X)
sched_getparam(3R)	See **sched_setparam**(3R)
sched_get_priority_max(3R)	get scheduling parameter limits
sched_get_priority_min(3R)	See **sched_get_priority_max**(3R)
sched_getscheduler(3R)	See **sched_setscheduler**(3R)
sched_rr_get_interval(3R)	See **sched_get_priority_max**(3R)
sched_setparam(3R)	set / get scheduling parameters
sched_setscheduler(3R)	set / get scheduling policy and scheduling parameters
sched_yield(3R)	yield processor

scr_dump(3X)	See **curs_scr_dump**(3X)
scr_init(3X)	See **curs_scr_dump**(3X)
scrl(3X)	See **curs_scroll**(3X)
scroll(3X)	See **curs_scroll**(3X)
scrollok(3X)	See **curs_outopts**(3X)
scr_restore(3X)	See **curs_scr_dump**(3X)
scr_set(3X)	See **curs_scr_dump**(3X)
seconvert(3)	See **econvert**(3)
secure_rpc(3N)	library routines for secure remote procedure calls
seed48(3C)	See **drand48**(3C)
seekdir(3C)	See **directory**(3C)
select(3C)	synchronous I / O multiplexing
sema_destroy(3T)	See **semaphore**(3T)
sema_init(3T)	See **semaphore**(3T)
semaphore(3T)	semaphores
sema_post(3T)	See **semaphore**(3T)
sema_trywait(3T)	See **semaphore**(3T)
sema_wait(3T)	See **semaphore**(3T)
sem_close(3R)	close a named semaphore
sem_destroy(3R)	destroy an unnamed semaphore
sem_getvalue(3R)	get the value of a semaphore
sem_init(3R)	initialize an unnamed semaphore
sem_open(3R)	initialize / open a named semaphore
sem_post(3R)	increment the count of a semaphore
sem_trywait(3R)	See **sem_wait**(3R)
sem_unlink(3R)	remove a named semaphore
sem_wait(3R)	acquire or wait for a semaphore
send(3N)	send a message from a socket
sendmsg(3N)	See **send**(3N)
sendto(3N)	See **send**(3N)
setac(3)	See **getacinfo**(3)
setauclass(3)	See **getauclassent**(3)
setauevent(3)	See **getauevent**(3)
setauuser(3)	See **getauusernam**(3)

setbuf(3S)	assign buffering to a stream
setbuffer(3C)	assign buffering to a stream
setcat(3C)	define default catalog
set_current_field(3X)	See **form_page**(3X)
set_current_item(3X)	See **menu_item_current**(3X)
set_curterm(3X)	See **curs_terminfo**(3X)
set_field_back(3X)	See **form_field_attributes**(3X)
set_field_buffer(3X)	See **form_field_buffer**(3X)
set_field_fore(3X)	See **form_field_attributes**(3X)
set_field_init(3X)	See **form_hook**(3X)
set_field_just(3X)	See **form_field_just**(3X)
set_field_opts(3X)	See **form_field_opts**(3X)
set_field_pad(3X)	See **form_field_attributes**(3X)
set_field_status(3X)	See **form_field_buffer**(3X)
set_field_term(3X)	See **form_hook**(3X)
set_field_type(3X)	See **form_field_validation**(3X)
set_fieldtype_arg(3X)	See **form_fieldtype**(3X)
set_fieldtype_choice(3X)	See **form_fieldtype**(3X)
set_field_userptr(3X)	See **form_field_userptr**(3X)
set_form_fields(3X)	See **form_field**(3X)
set_form_init(3X)	See **form_hook**(3X)
set_form_opts(3X)	See **form_opts**(3X)
set_form_page(3X)	See **form_page**(3X)
set_form_sub(3X)	See **form_win**(3X)
set_form_term(3X)	See **form_hook**(3X)
set_form_userptr(3X)	See **form_userptr**(3X)
set_form_win(3X)	See **form_win**(3X)
setgrent(3C)	See **getgrnam**(3C)
sethostent(3N)	See **gethostbyname**(3N)
sethostname(3C)	See **gethostname**(3C)
set_item_init(3X)	See **menu_hook**(3X)
set_item_opts(3X)	See **menu_item_opts**(3X)
set_item_term(3X)	See **menu_hook**(3X)
set_item_userptr(3X)	See **menu_item_userptr**(3X)
set_item_value(3X)	See **menu_item_value**(3X)

setjmp(3B)	non-local goto
_setjmp(3B)	See **setjmp**(3B)
setjmp(3C)	non-local goto
setkey(3C)	See **crypt**(3C)
setlabel(3C)	define the label for pfmt() and lfmt().
setlinebuf(3C)	See **setbuffer**(3C)
setlocale(3C)	modify and query a program's locale
setlogmask(3)	See **syslog**(3)
set_max_field(3X)	See **form_field_buffer**(3X)
set_menu_back(3X)	See **menu_attributes**(3X)
set_menu_fore(3X)	See **menu_attributes**(3X)
set_menu_format(3X)	See **menu_format**(3X)
set_menu_grey(3X)	See **menu_attributes**(3X)
set_menu_init(3X)	See **menu_hook**(3X)
set_menu_items(3X)	See **menu_items**(3X)
set_menu_mark(3X)	See **menu_mark**(3X)
set_menu_opts(3X)	See **menu_opts**(3X)
set_menu_pad(3X)	See **menu_attributes**(3X)
set_menu_pattern(3X)	See **menu_pattern**(3X)
set_menu_sub(3X)	See **menu_win**(3X)
set_menu_term(3X)	See **menu_hook**(3X)
set_menu_userptr(3X)	See **menu_userptr**(3X)
set_menu_win(3X)	See **menu_win**(3X)
setnetconfig(3N)	See **getnetconfig**(3N)
setnetent(3N)	See **getnetbyname**(3N)
setnetgrent(3N)	See **getnetgrent**(3N)
setnetpath(3N)	See **getnetpath**(3N)
set_new_page(3X)	See **form_new_page**(3X)
set_panel_userptr(3X)	See **panel_userptr**(3X)
setpriority(3C)	See **getpriority**(3C)
setprotoent(3N)	See **getprotobyname**(3N)
setpwent(3C)	See **getpwnam**(3C)
setrpcent(3N)	See **getrpcbyname**(3N)
setscrreg(3X)	See **curs_outopts**(3X)
setservent(3N)	See **getservbyname**(3N)

setsockopt(3N)	See **getsockopt**(3N)
setspent(3C)	See **getspnam**(3C)
setstate(3C)	See **random**(3C)
setsyx(3X)	See **curs_kernel**(3X)
set_term(3X)	See **curs_initscr**(3X)
setterm(3X)	See **curs_terminfo**(3X)
settimeofday(3B)	See **gettimeofday**(3B)
settimeofday(3C)	See **gettimeofday**(3C)
set_top_row(3X)	See **menu_item_current**(3X)
setupterm(3X)	See **curs_terminfo**(3X)
setusershell(3C)	See **getusershell**(3C)
setutent(3C)	See **getutent**(3C)
setutxent(3C)	See **getutxent**(3C)
setvbuf(3S)	See **setbuf**(3S)
sfconvert(3)	See **econvert**(3)
sgconvert(3)	See **econvert**(3)
shm_open(3R)	open a shared memory object
shm_unlink(3R)	remove a shared memory object
show_panel(3X)	See **panel_show**(3X)
shutdown(3N)	shut down part of a full-duplex connection
sig2str(3C)	See **str2sig**(3C)
sigaddset(3C)	See **sigsetops**(3C)
sigblock(3B)	block signals
sigdelset(3C)	See **sigsetops**(3C)
sigemptyset(3C)	See **sigsetops**(3C)
sigfillset(3C)	See **sigsetops**(3C)
sigfpe(3)	signal handling for specific SIGFPE codes
sighold(3C)	See **signal**(3C)
sigignore(3C)	See **signal**(3C)
siginterrupt(3B)	allow signals to interrupt functions
sigismember(3C)	See **sigsetops**(3C)
siglongjmp(3C)	See **setjmp**(3C)
sigmask(3B)	See **sigblock**(3B)
signal(3B)	simplified software signal facilities

signal(3C)	simplified signal management for application processes
significand(3M)	See **ieee_test**(3M)
sigpause(3B)	See **sigblock**(3B)
sigpause(3C)	See **signal**(3C)
sigqueue(3R)	queue a signal to a process
sigrelse(3C)	See **signal**(3C)
sigset(3C)	See **signal**(3C)
sigsetjmp(3C)	See **setjmp**(3C)
sigsetmask(3B)	See **sigblock**(3B)
sigsetops(3C)	manipulate sets of signals
sigstack(3B)	set and / or get signal stack context
sigtimedwait(3R)	See **sigwaitinfo**(3R)
sigvec(3B)	software signal facilities
sigwaitinfo(3R)	wait for queued signals
sin(3M)	See **trig**(3M)
single_to_decimal(3)	See **floating_to_decimal**(3)
sinh(3M)	See **hyperbolic**(3M)
sleep(3B)	suspend execution for interval
sleep(3C)	suspend execution for interval
slk_attroff(3X)	See **curs_slk**(3X)
slk_attron(3X)	See **curs_slk**(3X)
slk_attrset(3X)	See **curs_slk**(3X)
slk_clear(3X)	See **curs_slk**(3X)
slk_init(3X)	See **curs_slk**(3X)
slk_label(3X)	See **curs_slk**(3X)
slk_noutrefresh(3X)	See **curs_slk**(3X)
slk_refresh(3X)	See **curs_slk**(3X)
slk_restore(3X)	See **curs_slk**(3X)
slk_set(3X)	See **curs_slk**(3X)
slk_touch(3X)	See **curs_slk**(3X)
socket(3N)	create an endpoint for communication
socketpair(3N)	create a pair of connected sockets
space(3)	See **plot**(3)
spray(3N)	scatter data in order to test the

	network
sprintf(3B)	See **printf**(3B)
sprintf(3S)	See **printf**(3S)
sqrt(3M)	square root, cube root
srand(3B)	See **rand**(3B)
srand(3C)	See **rand**(3C)
srand48(3C)	See **drand48**(3C)
srandom(3C)	See **random**(3C)
sscanf(3S)	See **scanf**(3S)
ssignal(3C)	software signals
standend(3X)	See **curs_attr**(3X)
standout(3X)	See **curs_attr**(3X)
start_color(3X)	See **curs_color**(3X)
stdio(3S)	standard buffered input/output package
stdipc(3C)	standard interprocess communication package
step(3G)	See **regexpr**(3G)
store(3B)	See **dbm**(3B)
str(3G)	See **strfind**(3G)
str2sig(3C)	translation between signal name and signal number
strcadd(3G)	See **strccpy**(3G)
strcasecmp(3C)	See **string**(3C)
strcat(3C)	See **string**(3C)
strccpy(3G)	copy strings, compressing or expanding escape codes
strchr(3C)	See **string**(3C)
strcmp(3C)	See **string**(3C)
strcoll(3C)	string collation
strcpy(3C)	See **string**(3C)
strcspn(3C)	See **string**(3C)
strdup(3C)	See **string**(3C)
streadd(3G)	See **strccpy**(3G)
strecpy(3G)	See **strccpy**(3G)
strerror(3C)	get error message string

strfind(3G)	string manipulations
strfmon(3C)	convert monetary value to string
strftime(3C)	convert date and time to string
string(3C)	string operations
string_to_decimal(3)	parse characters into decimal record
strlen(3C)	See **string**(3C)
strncasecmp(3C)	See **string**(3C)
strncat(3C)	See **string**(3C)
strncmp(3C)	See **string**(3C)
strncpy(3C)	See **string**(3C)
strpbrk(3C)	See **string**(3C)
strptime(3C)	date and time conversion
strrchr(3C)	See **string**(3C)
strrspn(3G)	See **strfind**(3G)
strsignal(3C)	get error message string
strspn(3C)	See **string**(3C)
strstr(3C)	See **string**(3C)
strtod(3C)	convert string to double-precision number
strtok(3C)	See **string**(3C)
strtok_r(3C)	See **string**(3C)
strtol(3C)	conversion routines
strtoll(3C)	See **strtol**(3C)
strtoul(3C)	See **strtol**(3C)
strtoull(3C)	See **strtol**(3C)
strtrns(3G)	See **strfind**(3G)
strxfrm(3C)	string transformation
subpad(3X)	See **curs_pad**(3X)
subwin(3X)	See **curs_window**(3X)
svc_auth_reg(3N)	See **rpc_svc_reg**(3N)
svc_control(3N)	See **rpc_svc_create**(3N)
svc_create(3N)	See **rpc_svc_create**(3N)
svc_destroy(3N)	See **rpc_svc_create**(3N)
svc_dg_create(3N)	See **rpc_svc_create**(3N)
svc_dg_enablecache(3N)	See **rpc_svc_calls**(3N)

svc_done(3N)	See **rpc_svc_calls**(3N)
svcerr_auth(3N)	See **rpc_svc_err**(3N)
svcerr_decode(3N)	See **rpc_svc_err**(3N)
svcerr_noproc(3N)	See **rpc_svc_err**(3N)
svcerr_noprog(3N)	See **rpc_svc_err**(3N)
svcerr_progvers(3N)	See **rpc_svc_err**(3N)
svcerr_systemerr(3N)	See **rpc_svc_err**(3N)
svcerr_weakauth(3N)	See **rpc_svc_err**(3N)
svc_exit(3N)	See **rpc_svc_calls**(3N)
svcfd_create(3N)	See **rpc_soc**(3N)
svc_fd_create(3N)	See **rpc_svc_create**(3N)
svc_fds(3N)	See **rpc_soc**(3N)
svc_fdset(3N)	See **rpc_svc_calls**(3N)
svc_freeargs(3N)	See **rpc_svc_calls**(3N)
svc_getargs(3N)	See **rpc_svc_calls**(3N)
svc_getcaller(3N)	See **rpc_soc**(3N)
svc_getreq(3N)	See **rpc_soc**(3N)
svc_getreq_common(3N)	See **rpc_svc_calls**(3N)
svc_getreq_poll(3N)	See **rpc_svc_calls**(3N)
svc_getreqset(3N)	See **rpc_svc_calls**(3N)
svc_getrpccaller(3N)	See **rpc_svc_calls**(3N)
svc_kerb_reg(3N)	See **kerberos_rpc**(3N)
svc_pollset(3N)	See **rpc_svc_calls**(3N)
svcraw_create(3N)	See **rpc_soc**(3N)
svc_raw_create(3N)	See **rpc_svc_create**(3N)
svc_reg(3N)	See **rpc_svc_reg**(3N)
svc_register(3N)	See **rpc_soc**(3N)
svc_run(3N)	See **rpc_svc_calls**(3N)
svc_sendreply(3N)	See **rpc_svc_calls**(3N)
svctcp_create(3N)	See **rpc_soc**(3N)
svc_tli_create(3N)	See **rpc_svc_create**(3N)
svc_tp_create(3N)	See **rpc_svc_create**(3N)
svcudp_bufcreate(3N)	See **rpc_soc**(3N)
svcudp_create(3N)	See **rpc_soc**(3N)
svc_unreg(3N)	See **rpc_svc_reg**(3N)

svc_unregister(3N)	See **rpc_soc**(3N)
svc_vc_create(3N)	See **rpc_svc_create**(3N)
swab(3C)	swap bytes
swapcontext(3C)	See **makecontext**(3C)
syncok(3X)	See **curs_window**(3X)
syscall(3B)	indirect system call
sysconf(3C)	get configurable system variables
syslog(3)	control system log
sysmem(3)	return physical memory information
sys_siglist(3B)	See **psignal**(3B)
system(3S)	issue a shell command
t_accept(3N)	accept a connect request
taddr2uaddr(3N)	See **netdir**(3N)
t_alloc(3N)	allocate a library structure
tan(3M)	See **trig**(3M)
tanh(3M)	See **hyperbolic**(3M)
t_bind(3N)	bind an address to a transport end-point
tcdrain(3)	See **termios**(3)
tcflow(3)	See **termios**(3)
tcflush(3)	See **termios**(3)
tcgetattr(3)	See **termios**(3)
tcgetpgrp(3)	See **termios**(3)
tcgetsid(3)	See **termios**(3)
t_close(3N)	close a transport endpoint
t_connect(3N)	establish a connection with another transport user
tcsendbreak(3)	See **termios**(3)
tcsetattr(3)	See **termios**(3)
tcsetpgrp(3)	See **termios**(3)
tcsetpgrp(3C)	set foreground process group id of terminal
tdelete(3C)	See **tsearch**(3C)
td_init(3T)	See **libthread_db**(3T)
td_log(3T)	See **libthread_db**(3T)
td_ta_delete(3T)	See **libthread_db**(3T)

td_ta_get_nthreads(3T)	See **libthread_db**(3T)
td_ta_get_ph(3T)	See **libthread_db**(3T)
td_ta_map_id2thr(3T)	See **libthread_db**(3T)
td_ta_map_lwp2thr(3T)	See **libthread_db**(3T)
td_ta_new(3T)	See **libthread_db**(3T)
td_ta_thr_iter(3T)	See **libthread_db**(3T)
td_ta_tsd_iter(3T)	See **libthread_db**(3T)
td_thr_getfpregs(3T)	See **libthread_db**(3T)
td_thr_getgregs(3T)	See **libthread_db**(3T)
td_thr_get_info(3T)	See **libthread_db**(3T)
td_thr_getxregs(3T)	See **libthread_db**(3T)
td_thr_getxregsize(3T)	See **libthread_db**(3T)
td_thr_setfpregs(3T)	See **libthread_db**(3T)
td_thr_setgregs(3T)	See **libthread_db**(3T)
td_thr_setprio(3T)	See **libthread_db**(3T)
td_thr_setsigpending(3T)	See **libthread_db**(3T)
td_thr_setxregs(3T)	See **libthread_db**(3T)
td_thr_sigsetmask(3T)	See **libthread_db**(3T)
td_thr_tsd(3T)	See **libthread_db**(3T)
td_thr_validate(3T)	See **libthread_db**(3T)
telldir(3C)	See **directory**(3C)
tempnam(3S)	See **tmpnam**(3S)
termattrs(3X)	See **curs_termattrs**(3X)
termios(3)	general terminal interface
termname(3X)	See **curs_termattrs**(3X)
t_error(3N)	produce error message
textdomain(3I)	See **gettext**(3I)
tfind(3C)	See **tsearch**(3C)
t_free(3N)	free a library structure
tgetent(3X)	See **curs_termcap**(3X)
tgetflag(3X)	See **curs_termcap**(3X)
t_getinfo(3N)	get protocol-specific service information
tgetnum(3X)	See **curs_termcap**(3X)
t_getstate(3N)	get the current state

tgetstr(3X)	See **curs_termcap**(3X)
tgoto(3X)	See **curs_termcap**(3X)
thr_continue(3T)	See **thr_suspend**(3T)
thr_create(3T)	See **pthread_create**(3T)
threads(3T)	thread libraries: libpthread and libthread
thr_exit(3T)	See **pthread_exit**(3T)
thr_getconcurrency(3T)	See **thr_setconcurrency**(3T)
thr_getprio(3T)	See **pthread_setschedparam**(3T)
thr_getspecific(3T)	See **pthread_key_create**(3T)
thr_join(3T)	See **pthread_join**(3T)
thr_keycreate(3T)	See **pthread_key_create**(3T)
thr_kill(3T)	See **pthread_kill**(3T)
thr_main(3T)	identify the main thread
thr_min_stack(3T)	returns the minimum-allowable size for a thread's stack
thr_self(3T)	See **pthread_self**(3T)
thr_setconcurrency(3T)	get / set thread concurrency level
thr_setprio(3T)	See **pthread_setschedparam**(3T)
thr_setspecific(3T)	See **pthread_key_create**(3T)
thr_sigsetmask(3T)	See **pthread_sigmask**(3T)
thr_stksegment(3T)	get thread stack bottom and stack size
thr_suspend(3T)	suspend or continue thread execution
thr_yield(3T)	thread yield to another thread
tigetflag(3X)	See **curs_terminfo**(3X)
tigetnum(3X)	See **curs_terminfo**(3X)
tigetstr(3X)	See **curs_terminfo**(3X)
timeout(3X)	See **curs_inopts**(3X)
timer_create(3R)	create a timer
timer_delete(3R)	delete a per-LWP timer
timer_getoverrun(3R)	See **timer_settime**(3R)
timer_gettime(3R)	See **timer_settime**(3R)
timer_settime(3R)	high-resolution timer operations
times(3B)	get process times

t_listen(3N)	listen for a connect request
t_look(3N)	look at the current event on a transport endpoint
tmpfile(3S)	create a temporary file
tmpnam(3S)	create a name for a temporary file
tmpnam_r(3S)	See **tmpnam**(3S)
TNF_DECLARE_RECORD(3X)	TNF type extension interface for probes
TNF_DEFINE_RECORD_1(3X)	See **TNF_DECLARE_RECORD**(3X)
TNF_DEFINE_RECORD_2(3X)	See **TNF_DECLARE_RECORD**(3X)
TNF_DEFINE_RECORD_3(3X)	See **TNF_DECLARE_RECORD**(3X)
TNF_DEFINE_RECORD_4(3X)	See **TNF_DECLARE_RECORD**(3X)
TNF_DEFINE_RECORD_5(3X)	See **TNF_DECLARE_RECORD**(3X)
TNF_PROBE(3X)	probe insertion interface
TNF_PROBE_0(3X)	See **TNF_PROBE**(3X)
TNF_PROBE_1(3X)	See **TNF_PROBE**(3X)
TNF_PROBE_2(3X)	See **TNF_PROBE**(3X)
TNF_PROBE_3(3X)	See **TNF_PROBE**(3X)
TNF_PROBE_4(3X)	See **TNF_PROBE**(3X)
TNF_PROBE_5(3X)	See **TNF_PROBE**(3X)
tnf_process_disable(3X)	probe control internal interface
tnf_process_enable(3X)	See **tnf_process_disable**(3X)
tnf_thread_disable(3X)	See **tnf_process_disable**(3X)
tnf_thread_enable(3X)	See **tnf_process_disable**(3X)
toascii(3C)	See **conv**(3C)
_tolower(3C)	See **conv**(3C)
tolower(3C)	See **conv**(3C)
t_open(3N)	establish a transport endpoint
top_panel(3X)	See **panel_top**(3X)
top_row(3X)	See **menu_item_current**(3X)
t_optmgmt(3N)	manage options for a transport endpoint
touchline(3X)	See **curs_touch**(3X)
touchwin(3X)	See **curs_touch**(3X)
_toupper(3C)	See **conv**(3C)
toupper(3C)	See **conv**(3C)

towlower(3I)	See **wconv**(3I)
towupper(3I)	See **wconv**(3I)
tparm(3X)	See **curs_terminfo**(3X)
tputs(3X)	See **curs_termcap**(3X)
tputs(3X)	See **curs_terminfo**(3X)
t_rcv(3N)	receive data or expedited data sent over a connection
t_rcvconnect(3N)	receive the confirmation from a connect request
t_rcvdis(3N)	retrieve information from disconnect
t_rcvrel(3N)	acknowledge receipt of an orderly release indication
t_rcvudata(3N)	receive a data unit
t_rcvuderr(3N)	receive a unit data error indication
trig(3M)	trigonometric functions
truncate(3C)	set a file to a specified length
tsearch(3C)	manage binary search trees
t_snd(3N)	send data or expedited data over a connection
t_snddis(3N)	send user-initiated disconnect request
t_sndrel(3N)	initiate an orderly release
t_sndudata(3N)	send a data unit
t_strerror(3N)	get error message string
t_sync(3N)	synchronize transport library
ttyname(3C)	find name of a terminal
ttyname_r(3C)	See **ttyname**(3C)
ttyslot(3C)	find the slot in the utmp file of the current user
t_unbind(3N)	disable a transport endpoint
twalk(3C)	See **tsearch**(3C)
typeahead(3X)	See **curs_inopts**(3X)
tzset(3C)	See **ctime**(3C)
tzsetwall(3C)	See **ctime**(3C)
uaddr2taddr(3N)	See **netdir**(3N)
ualarm(3C)	schedule signal after interval in microseconds

ulckpwdf(3C)	See **lckpwdf**(3C)
ulltostr(3C)	See **strtol**(3C)
unctrl(3X)	See **curs_util**(3X)
ungetc(3S)	push character back onto input stream
ungetch(3X)	See **curs_getch**(3X)
ungetwc(3I)	push a Process Code character back into input stream
ungetwch(3X)	See **curs_getwch**(3X)
unlockpt(3C)	unlock a pseudo-terminal master/slave pair
unordered(3C)	See **isnan**(3C)
unpost_form(3X)	See **form_post**(3X)
unpost_menu(3X)	See **menu_post**(3X)
untouchwin(3X)	See **curs_touch**(3X)
update_panels(3X)	See **panel_update**(3X)
updwtmp(3C)	See **getutxent**(3C)
updwtmpx(3C)	See **getutxent**(3C)
use_env(3X)	See **curs_util**(3X)
user2netname(3N)	See **secure_rpc**(3N)
usleep(3C)	suspend execution for interval in microseconds
utmpname(3C)	See **getutent**(3C)
utmpxname(3C)	See **getutxent**(3C)
valloc(3C)	See **malloc**(3C)
vfprintf(3B)	See **printf**(3B)
vfprintf(3S)	See **vprintf**(3S)
vidattr(3X)	See **curs_terminfo**(3X)
vidputs(3X)	See **curs_terminfo**(3X)
vlfmt(3C)	display error message in standard format and pass to logging and monitoring services
volmgt_check(3X)	have Volume Management check for media
volmgt_inuse(3X)	check whether or not Volume Management is managing a pathname

volmgt_root(3X)	return the Volume Management root directory
volmgt_running(3X)	return whether or not Volume Management is running
volmgt_symdev(3X)	See **volmgt_symname**(3X)
volmgt_symname(3X)	convert between Volume Management symbolic names, and the devices that correspond to them
vpfmt(3C)	display error message in standard format and pass to logging and monitoring services
vprintf(3B)	See **printf**(3B)
vprintf(3S)	print formatted output of a variable argument list
vsprintf(3B)	See **printf**(3B)
vsprintf(3S)	See **vprintf**(3S)
vsyslog(3)	log message with a varargs argument list
vwprintw(3X)	See **curs_printw**(3X)
vwscanw(3X)	See **curs_scanw**(3X)
waddch(3X)	See **curs_addch**(3X)
waddchnstr(3X)	See **curs_addchstr**(3X)
waddchstr(3X)	See **curs_addchstr**(3X)
waddnstr(3X)	See **curs_addstr**(3X)
waddnwstr(3X)	See **curs_addwstr**(3X)
waddstr(3X)	See **curs_addstr**(3X)
waddwch(3X)	See **curs_addwch**(3X)
waddwchnstr(3X)	See **curs_addwchstr**(3X)
waddwchstr(3X)	See **curs_addwchstr**(3X)
waddwstr(3X)	See **curs_addwstr**(3X)
wadjcurspos(3X)	See **curs_alecompat**(3X)
wait(3B)	wait for process to terminate or stop
wait3(3B)	See **wait**(3B)
wait3(3C)	wait for process to terminate or stop
wait4(3B)	See **wait**(3B)
wait4(3C)	See **wait3**(3C)
waitpid(3B)	See **wait**(3B)

watof(3I)	See **wcstod**(3I)
watoi(3I)	See **wcstol**(3I)
watol(3I)	See **wcstol**(3I)
watoll(3I)	See **wcstol**(3I)
wattroff(3X)	See **curs_attr**(3X)
wattron(3X)	See **curs_attr**(3X)
wattrset(3X)	See **curs_attr**(3X)
wbkgd(3X)	See **curs_bkgd**(3X)
wbkgdset(3X)	See **curs_bkgd**(3X)
wborder(3X)	See **curs_border**(3X)
wclear(3X)	See **curs_clear**(3X)
wclrtobot(3X)	See **curs_clear**(3X)
wclrtoeol(3X)	See **curs_clear**(3X)
wconv(3I)	Process Code character conversion macros
wcscat(3I)	See **wcstring**(3I)
wcschr(3I)	See **wcstring**(3I)
wcscmp(3I)	See **wcstring**(3I)
wcscoll(3I)	wide character string comparison using collating information
wcscpy(3I)	See **wcstring**(3I)
wcscspn(3I)	See **wcstring**(3I)
wcsetno(3I)	See **cset**(3I)
wcsftime(3I)	convert date and time to wide character string
wcslen(3I)	See **wcstring**(3I)
wcsncat(3I)	See **wcstring**(3I)
wcsncmp(3I)	See **wcstring**(3I)
wcsncpy(3I)	See **wcstring**(3I)
wcspbrk(3I)	See **wcstring**(3I)
wcsrchr(3I)	See **wcstring**(3I)
wcsspn(3I)	See **wcstring**(3I)
wcstod(3I)	convert wide character string to double-precision number
wcstok(3I)	See **wcstring**(3I)
wcstol(3I)	convert wide character string to long

	integer
wcstombs(3C)	See **mbstring**(3C)
wcstoul(3I)	convert wide character string to unsigned long
wcstring(3I)	wide character string operations
wcswcs(3I)	See **wcstring**(3I)
wcswidth(3I)	See **wcstring**(3I)
wcsxfrm(3I)	wide character string transformation
wctomb(3C)	See **mbchar**(3C)
wctype(3I)	define character class
wcursyncup(3X)	See **curs_window**(3X)
wcwidth(3I)	See **wcstring**(3I)
wdelch(3X)	See **curs_delch**(3X)
wdeleteln(3X)	See **curs_deleteln**(3X)
wechochar(3X)	See **curs_addch**(3X)
wechowchar(3X)	See **curs_addwch**(3X)
werase(3X)	See **curs_clear**(3X)
wgetch(3X)	See **curs_getch**(3X)
wgetnstr(3X)	See **curs_getstr**(3X)
wgetnwstr(3X)	See **curs_getwstr**(3X)
wgetstr(3X)	See **curs_getstr**(3X)
wgetwch(3X)	See **curs_getwch**(3X)
wgetwstr(3X)	See **curs_getwstr**(3X)
whline(3X)	See **curs_border**(3X)
WIFEXITED(3B)	See **wait**(3B)
WIFSIGNALED(3B)	See **wait**(3B)
WIFSTOPPED(3B)	See **wait**(3B)
winch(3X)	See **curs_inch**(3X)
winchnstr(3X)	See **curs_inchstr**(3X)
winchstr(3X)	See **curs_inchstr**(3X)
windex(3I)	See **wcstring**(3I)
winnstr(3X)	See **curs_instr**(3X)
winnwstr(3X)	See **curs_inwstr**(3X)
winsch(3X)	See **curs_insch**(3X)
winsdelln(3X)	See **curs_deleteln**(3X)

winsertln(3X)	See **curs_deleteln**(3X)
winsnstr(3X)	See **curs_insstr**(3X)
winsnwstr(3X)	See **curs_inswstr**(3X)
winsstr(3X)	See **curs_insstr**(3X)
winstr(3X)	See **curs_instr**(3X)
winswch(3X)	See **curs_inswch**(3X)
winswstr(3X)	See **curs_inswstr**(3X)
winwch(3X)	See **curs_inwch**(3X)
winwchnstr(3X)	See **curs_inwchstr**(3X)
winwchstr(3X)	See **curs_inwchstr**(3X)
winwstr(3X)	See **curs_inwstr**(3X)
wmove(3X)	See **curs_move**(3X)
wmovenextch(3X)	See **curs_alecompat**(3X)
wmoveprevch(3X)	See **curs_alecompat**(3X)
wnoutrefresh(3X)	See **curs_refresh**(3X)
wordexp(3C)	perform word expansions
wordfree(3C)	See **wordexp**(3C)
wprintw(3X)	See **curs_printw**(3X)
wredrawln(3X)	See **curs_refresh**(3X)
wrefresh(3X)	See **curs_refresh**(3X)
wrindex(3I)	See **wcstring**(3I)
write_vtoc(3X)	See **read_vtoc**(3X)
wscanw(3X)	See **curs_scanw**(3X)
wscasecmp(3I)	See **wstring**(3I)
wscat(3I)	See **wcstring**(3I)
wschr(3I)	See **wcstring**(3I)
wscmp(3I)	See **wcstring**(3I)
wscol(3I)	See **wstring**(3I)
wscoll(3I)	See **wcscoll**(3I)
wscpy(3I)	See **wcstring**(3I)
wscrl(3X)	See **curs_scroll**(3X)
wscspn(3I)	See **wcstring**(3I)
wsdup(3I)	See **wstring**(3I)
wsetscrreg(3X)	See **curs_outopts**(3X)
wslen(3I)	See **wcstring**(3I)

wsncasecmp(3I)	See **wstring**(3I)
wsncat(3I)	See **wcstring**(3I)
wsncmp(3I)	See **wcstring**(3I)
wsncpy(3I)	See **wcstring**(3I)
wspbrk(3I)	See **wcstring**(3I)
wsprintf(3I)	formatted output conversion
wsrchr(3I)	See **wcstring**(3I)
wsscanf(3I)	formatted input conversion
wsspn(3I)	See **wcstring**(3I)
wstandend(3X)	See **curs_attr**(3X)
wstandout(3X)	See **curs_attr**(3X)
wstod(3I)	See **wcstod**(3I)
wstok(3I)	See **wcstring**(3I)
wstol(3I)	See **wcstol**(3I)
wstring(3I)	Process Code string operations
wsxfrm(3I)	See **wcsxfrm**(3I)
wsyncdown(3X)	See **curs_window**(3X)
wsyncup(3X)	See **curs_window**(3X)
wtimeout(3X)	See **curs_inopts**(3X)
wtouchln(3X)	See **curs_touch**(3X)
wvline(3X)	See **curs_border**(3X)
xdr(3N)	library routines for external data representation
xdr_accepted_reply(3N)	See **rpc_xdr**(3N)
xdr_admin(3N)	library routines for external data representation
xdr_array(3N)	See **xdr_complex**(3N)
xdr_authsys_parms(3N)	See **rpc_xdr**(3N)
xdr_authunix_parms(3N)	See **rpc_soc**(3N)
xdr_bool(3N)	See **xdr_simple**(3N)
xdr_bytes(3N)	See **xdr_complex**(3N)
xdr_callhdr(3N)	See **rpc_xdr**(3N)
xdr_callmsg(3N)	See **rpc_xdr**(3N)
xdr_char(3N)	See **xdr_simple**(3N)
xdr_complex(3N)	library routines for external data representation

xdr_control(3N)	See **xdr_admin**(3N)
xdr_create(3N)	library routines for external data representation stream creation
xdr_destroy(3N)	See **xdr_create**(3N)
xdr_double(3N)	See **xdr_simple**(3N)
xdr_enum(3N)	See **xdr_simple**(3N)
xdr_float(3N)	See **xdr_simple**(3N)
xdr_free(3N)	See **xdr_simple**(3N)
xdr_getpos(3N)	See **xdr_admin**(3N)
xdr_hyper(3N)	See **xdr_simple**(3N)
xdr_inline(3N)	See **xdr_admin**(3N)
xdr_int(3N)	See **xdr_simple**(3N)
xdr_long(3N)	See **xdr_simple**(3N)
xdr_longlong_t(3N)	See **xdr_simple**(3N)
xdrmem_create(3N)	See **xdr_create**(3N)
xdr_opaque(3N)	See **xdr_complex**(3N)
xdr_opaque_auth(3N)	See **rpc_xdr**(3N)
xdr_pointer(3N)	See **xdr_complex**(3N)
xdr_quadruple(3N)	See **xdr_simple**(3N)
xdrrec_create(3N)	See **xdr_create**(3N)
xdrrec_endofrecord(3N)	See **xdr_admin**(3N)
xdrrec_eof(3N)	See **xdr_admin**(3N)
xdrrec_readbytes(3N)	See **xdr_admin**(3N)
xdrrec_skiprecord(3N)	See **xdr_admin**(3N)
xdr_reference(3N)	See **xdr_complex**(3N)
xdr_rejected_reply(3N)	See **rpc_xdr**(3N)
xdr_replymsg(3N)	See **rpc_xdr**(3N)
xdr_setpos(3N)	See **xdr_admin**(3N)
xdr_short(3N)	See **xdr_simple**(3N)
xdr_simple(3N)	library routines for external data representation
xdr_sizeof(3N)	See **xdr_admin**(3N)
xdrstdio_create(3N)	See **xdr_create**(3N)
xdr_string(3N)	See **xdr_complex**(3N)
xdr_u_char(3N)	See **xdr_simple**(3N)
xdr_u_hyper(3N)	See **xdr_simple**(3N)

xdr_u_int(3N)	See **xdr_simple**(3N)
xdr_u_long(3N)	See **xdr_simple**(3N)
xdr_u_longlong_t(3N)	See **xdr_simple**(3N)
xdr_union(3N)	See **xdr_complex**(3N)
xdr_u_short(3N)	See **xdr_simple**(3N)
xdr_vector(3N)	See **xdr_complex**(3N)
xdr_void(3N)	See **xdr_simple**(3N)
xdr_wrapstring(3N)	See **xdr_complex**(3N)
xfn(3N)	overview of the XFN interface
xfn_attributes(3N)	an overview of XFN attribute operations
xfn_composite_names(3N)	XFN composite syntax: an overview of the syntax for XFN composite name
xfn_compound_names(3N)	XFN compound syntax: an overview of XFN model for compound name parsing
xfn_links(3N)	XFN links: an overview of XFN links
xfn_status_codes(3N)	descriptions of XFN status codes
xprt_register(3N)	See **rpc_svc_reg**(3N)
xprt_unregister(3N)	See **rpc_svc_reg**(3N)
y0(3M)	See **bessel**(3M)
y1(3M)	See **bessel**(3M)
yn(3M)	See **bessel**(3M)
yp_all(3N)	See **ypclnt**(3N)
yp_bind(3N)	See **ypclnt**(3N)
ypclnt(3N)	NIS Version 2 client interface
yperr_string(3N)	See **ypclnt**(3N)
yp_first(3N)	See **ypclnt**(3N)
yp_get_default_domain(3N)	See **ypclnt**(3N)
yp_master(3N)	See **ypclnt**(3N)
yp_match(3N)	See **ypclnt**(3N)
yp_next(3N)	See **ypclnt**(3N)
yp_order(3N)	See **ypclnt**(3N)
ypprot_err(3N)	See **ypclnt**(3N)
yp_unbind(3N)	See **ypclnt**(3N)

yp_update(3N) change NIS information

NAME

NOTE, _NOTE – annotate source code with info for tools

SYNOPSIS

#include <note.h>
NOTE(*NoteInfo***)**

or

#include <sys/note.h>
_NOTE(*NoteInfo***)**

MT-LEVEL

Safe

DESCRIPTION

These macros are used to embed information for tools in program source. A use of one of these macros is called an “annotation”. A tool may define a set of such annotations which can then be used to provide the tool with information that would otherwise be unavailable from the source code.

Annotations should, in general, provide documentation useful to the human reader. If information is of no use to a human trying to understand the code but is necessary for proper operation of a tool, use another mechanism for conveying that information to the tool (one which does not involve adding to the source code), so as not to detract from the readability of the source. The following is an example of an annotation which provides information of use to a tool and to the human reader (in this case, which data are protected by a particular lock, an annotation defined by the static lock analysis tool **lock_lint**).

NOTE(MUTEX_PROTECTS_DATA(foo_lock, foo_list Foo))

Such annotations do not represent executable code; they are neither statements nor declarations. They should not be followed by a semicolon. If a compiler or tool that analyzes C source does not understand this annotation scheme, then the tool will ignore the annotations. (For such tools, **NOTE(***x***)** expands to nothing.)

Annotations may only be placed at particular places in the source. These places are where the following C constructs would be allowed:

- a top-level declaration (that is, a declaration not within a function or other construct)
- a declaration or statement within a block (including the block which defines a function)
- a member of a **struct** or **union**.

Annotations are not allowed in any other place. For example, the following are illegal:

x = y + NOTE(...) z ;
typedef NOTE(...) unsigned int uint ;

While **NOTE** and **_NOTE** may be used in the places described above, a particular type of annotation may only be allowed in a subset of those places. For example, a particular annotation may not be allowed inside a **struct** or **union** definition.

NOTE vs _NOTE

Ordinarily, **NOTE** should be used rather than **_NOTE**, since use of **_NOTE** technically makes a program non-portable. However, it may be inconvenient to use **NOTE** for this purpose in existing code if **NOTE** is already heavily used for another purpose. In this case one should use a different macro and write a header file similar to **/usr/include/note.h** which maps that macro to **_NOTE** in the same manner. For example, the following makes **FOO** such a macro:

```
#ifndef _FOO_H
#define _FOO_H
#define FOO _NOTE
#include <sys/note.h>
#endif
```

Public header files which span projects should use **_NOTE** rather than **NOTE**, since **NOTE** may already be used by a program which needs to include such a header file.

NoteInfo **Argument**

The actual *NoteInfo* used in an annotation should be specified by a tool that deals with program source (see the documentation for the tool to determine which annotations, if any, it understands).

NoteInfo must have one of the following forms:

NoteName
*NoteName***(***Args***)**

where *NoteName* is simply an identifier which indicates the type of annotation, and *Args* is something defined by the tool that specifies the particular *NoteName*. The general restrictions on *Args* are that it be compatible with an ANSI C tokenizer and that unquoted parentheses be balanced (so that the end of the annotation can be determined without intimate knowledge of any particular annotation).

SEE ALSO

note(4)

NAME

TNF_DECLARE_RECORD, TNF_DEFINE_RECORD_1, TNF_DEFINE_RECORD_2, TNF_DEFINE_RECORD_3, TNF_DEFINE_RECORD_4, TNF_DEFINE_RECORD_5 – TNF type extension interface for probes

SYNOPSIS

cc [*flag* ...] *file* ... [–**ltnfprobe**] [*library* ...]

#include <tnf/probe.h>

TNF_DECLARE_RECORD(*c_type***,** *tnf_type***);**

TNF_DEFINE_RECORD_1(*c_type***,** *tnf_type***,**
*tnf_member_type_1***,** *c_member_name_1***)**

TNF_DEFINE_RECORD_2(*c_type***,** *tnf_type***,**
*tnf_member_type_1***,** *c_member_name_1***,**
*tnf_member_type_2***,** *c_member_name_2***)**

TNF_DEFINE_RECORD_3(*c_type***,** *tnf_type***,**
*tnf_member_type_1***,** *c_member_name_1***,**
*tnf_member_type_2***,** *c_member_name_2***,**
*tnf_member_type_3***,** *c_member_name_3***)**

TNF_DEFINE_RECORD_4(*c_type***,** *tnf_type***,**
*tnf_member_type_1***,** *c_member_name_1***,**
*tnf_member_type_2***,** *c_member_name_2***,**
*tnf_member_type_3***,** *c_member_name_3***,**
*tnf_member_type_4***,** *c_member_name_4***)**

TNF_DEFINE_RECORD_5(*c_type***,** *tnf_type***,**
*tnf_member_type_1***,** *c_member_name_1***,**
*tnf_member_type_2***,** *c_member_name_2***,**
*tnf_member_type_3***,** *c_member_name_3***,**
*tnf_member_type_4***,** *c_member_name_4***,**
*tnf_member_type_5***,** *c_member_name_5***)**

AVAILABILITY

SUNWtnfd

MT-LEVEL

MT-Safe.

DESCRIPTION

This macro interface is used to extend the TNF (Trace Normal Form) types that can be used in **TNF_PROBE**(3X).

There should be only one **TNF_DECLARE_RECORD** and one **TNF_DEFINE_RECORD** per new type being defined. The **TNF_DECLARE_RECORD** should precede the **TNF_DEFINE_RECORD**. It can be in a header file that multiple source files share if those source files need to use the *tnf_type* being defined. The **TNF_DEFINE_RECORD** should only appear in one of the source files.

The **TNF_DEFINE_RECORD** macro interface defines a function as well as a couple of data structures. Hence, this interface has to be used in a source file (.c or .cc file) at file scope and not inside a function.

Note that there is no semicolon after the **TNF_DEFINE_RECORD** interface. Having one will generate a compiler warning.

Compiling with the preprocessor option –**DNPROBE** (see **cc**(1B)), or with the preprocessor control statement **#define NPROBE** ahead of the **#include <tnf/probe.h>** statement, will stop the TNF type extension code from being compiled into the program.

c_type *c_type* must be a C struct type. It is the template from which the new *tnf_type* is being created. Not all elements of the C struct need be provided in the TNF type being defined.

tnf_type *tnf_type* is the name being given to the newly created type. Use of this interface uses the name space prefixed by *tnf_type*. So, if a new type called "xxx_type" is defined by a library, then the library should not use "xxx_type" as a prefix in any other symbols it defines. The policy on managing the type name space is the same as managing any other name space in a library i.e., prefix any new TNF types by the unique prefix that the rest of the symbols in the library use. This would prevent name space collisions when linking multiple libraries that define new TNF types. For example, if a library libpalloc.so uses the prefix "pal" for all symbols it defines, then it should also use the prefix "pal" for all new TNF types being defined.

tnf_member_type_n *tnf_member_type_n* is the TNF type of the *n*th provided member of the C structure.

tnf_member_name_n *tnf_member_name_n* is the name of the *n*th provided member of the C structure.

EXAMPLES This example shows how a new TNF type is defined and used in a probe. This code is assumed to be part of a fictitious library called "libpalloc.so" which uses the prefix "pal" for all it's symbols.

```
#include <tnf/probe.h>

typedef struct pal_header {
        long    size;
        char *  descriptor;
        struct pal_header *next;
} pal_header_t;

TNF_DECLARE_RECORD(pal_header_t, pal_tnf_header);
TNF_DEFINE_RECORD_2(pal_header_t, pal_tnf_header,
                    tnf_long,   size,
                    tnf_string, descriptor)

/*
 * Note: name space prefixed by pal_tnf_header should not be used by this
 *       client anymore.
 */

void
```

```
pal_free(pal_header_t *header_p)
{
        int state;

        TNF_PROBE_2(pal_free_start, "palloc pal_free",
                "sunw%debug entering pal_free",
                tnf_long,       state_var,  state,
                pal_tnf_header, header_var, header_p);
        . . .
}
```

SEE ALSO **prex**(1), **tnfdump**(1), **TNF_PROBE**(3X), **tnf_process_disable**(3X)

NOTES It is possible to make a *tnf_type* definition be recursive or mutually recursive e.g. a structure that uses the "next" field to point to itself (a linked list). If such a structure is sent in to a **TNF_PROBE**(3X), then the entire linked list will be logged to the trace file (until the "next" field is **NULL).** But, if the list is circular, it will result in an infinite loop. To break the recursion, either don't include the "next" field in the *tnf_type*, or define the type of the "next" member as **tnf_opaque.**

NAME

TNF_PROBE, TNF_PROBE_0, TNF_PROBE_1, TNF_PROBE_2, TNF_PROBE_3, TNF_PROBE_4, TNF_PROBE_5 – probe insertion interface

SYNOPSIS

cc [*flag* ...] *file* ... [–**ltnfprobe**] [*library* ...]

#include <tnf/probe.h>

TNF_PROBE_0(*name***,** *keys***,** *detail***);**

TNF_PROBE_1(*name***,** *keys***,** *detail***,**
*arg_type_1***,** *arg_name_1***,** *arg_value_1***);**

TNF_PROBE_2(*name***,** *keys***,** *detail***,**
*arg_type_1***,** *arg_name_1***,** *arg_value_1***,**
*arg_type_2***,** *arg_name_2***,** *arg_value_2***);**

TNF_PROBE_3(*name***,** *keys***,** *detail***,**
*arg_type_1***,** *arg_name_1***,** *arg_value_1***,**
*arg_type_2***,** *arg_name_2***,** *arg_value_2***,**
*arg_type_3***,** *arg_name_3***,** *arg_value_3***);**

TNF_PROBE_4(*name***,** *keys***,** *detail***,**
*arg_type_1***,** *arg_name_1***,** *arg_value_1***,**
*arg_type_2***,** *arg_name_2***,** *arg_value_2***,**
*arg_type_3***,** *arg_name_3***,** *arg_value_3***,**
*arg_type_4***,** *arg_name_4***,** *arg_value_4***);**

TNF_PROBE_5(*name***,** *keys***,** *detail***,**
*arg_type_1***,** *arg_name_1***,** *arg_value_1***,**
*arg_type_2***,** *arg_name_2***,** *arg_value_2***,**
*arg_type_3***,** *arg_name_3***,** *arg_value_3***,**
*arg_type_4***,** *arg_name_4***,** *arg_value_4***,**
*arg_type_5***,** *arg_name_5***,** *arg_value_5***);**

AVAILABILITY

SUNWtnfd

MT-LEVEL

MT-Safe.

DESCRIPTION

This macro interface is used to insert probes into C or C++ code. Probes can be placed anywhere in these programs including .init sections, .fini sections, multi-threaded code, shared objects, and shared objects opened by **dlopen**(3X). Probes can be used to generate trace data for performance analysis or to write debugging output to stderr. Probes are controlled by **prex**(1).

The trace data is logged to a trace file in Trace Normal Form (TNF) — the interface for the user to specify the name and size of the trace file is describe in **prex**(1). The trace file can be thought of as a least recently used circular buffer. Once the file has been filled, newer events will overwrite the older ones.

Compiling with the preprocessor option –**DNPROBE** (see **cc**(1B)), or with the preprocessor control statement **#define NPROBE** ahead of the **#include <tnf/probe.h>** statement, will stop probes from being compiled into the program.

name *name* of the probe and should follow the syntax guidelines for identifiers in ANSI C. The use of *name* declares it—hence no separate declaration is necessary. This is a block scope declaration, so it does not affect the name space of the program.

keys *keys* is a string of space separated keywords that specify the groups that the probe belongs to. A semicolon or single quotation mark is not allowed in this string. If any of the groups are enabled, the probe is enabled. *keys* cannot be a variable—it has to be an inlined string.

detail *detail* is a string that consists of <attribute> <value> pairs that are each separated by a semicolon. The first word (up to a space) is considered to be the attribute and the rest of the string (up to the semicolon) is considered the value. The value is optional. Semicolons or single quotation marks are not allowed in either the attribute or the value. *detail* is used for two reasons—first, it can be used to supply an attribute that a user can type into **prex**(1) to select probes. For example, if a user defines an attribute called color, then **prex**(1) can select probes based on the value of color. Secondly, *detail* can be used to annotate a probe with a string that is written out to a trace file only once. **prex**(1) uses spaces to tokenize the value when searching for a match. Spaces around the semicolon delimiter are allowed. *detail* cannot be a variable—it has to be an inlined string. For example, the *detail* string:

"**sunw%debug entering function A; comX%exception no file; comY%func_entry; comY%color red blue**"

consists of 4 units:

attribute	value	values that prex matches on
sunw%debug	entering function A	entering <or> function <or> A
comX%exception	no file	no <or> file
comY%func_entry		/.*/ (regular expression)
comY%color	red blue	red <or> blue

Attribute names have to be prefixed by the vendor stock symbol followed by the '%' character. This is to avoid collisions in the attribute name space. All attributes that do not have a '%' character are reserved. These are the predefined attributes:

attribute	semantics
name	name of probe
keys	keys of the probe (value is space separated tokens)
file	file name of the probe
line	line number of the probe
slots	slot names of the probe event (i.e., *arg_name_n*)
enable	off => probe is disabled. on => probe is enabled.
trace	off => tracing is off. on => tracing is on.
object	the executable or shared object that this probe is in.

funcs list of probe functions connected to this probe.

arg_type_n This is the type of the *n*th argument. These are the predefined TNF types:

tnf type	associated C type (and semantics)
tnf_long	int, long
tnf_ulong	unsigned int, unsigned long
tnf_longlong	long long (if implemented in compilation system)
tnf_ulonglong	unsigned long long (if implemented in compilation system)
tnf_float	float
tnf_double	double
tnf_string	char *
tnf_opaque	void *

To define new TNF types that are records consisting of the predefined TNF types or references to other user defined types, use the interface specified in **TNF_DECLARE_RECORD**(3X).

arg_name_n *arg_name_n* is the name that the user wants associated with the *n*th argument. It should not have any quotes around it and should follow the syntax guidelines for identifiers in ANSI C. The string version of *arg_name_n* is stored for every probe and can be accessed as the attribute 'slots' as mentioned above.

arg_value_n *arg_value_n* is evaluated to yield a value to be included in the trace file. A read access is done on any variables that are in mentioned in *arg_value_n*. In a multi-threaded program, it is the user's responsibility to place locks around the **TNF_PROBE** macro if *arg_value_n* contains a variable that needs to be read protected.

EXAMPLES In this example, probes are placed at the entry and exit of a function to determine how much time is spent in the function. The function entry probe also logs the arguments to the function. When a probe is encountered at run time and if it is enabled for tracing, a record is generated to a trace file. All probes log the time when it was encountered and also have a reference to a tag record which has information like the file name, line number, *name*, *keys*, and *detail* of the probe. The first probe *work_args* will also log the value of the two arguments of the probe (*state* and *message*).

```
#include <tnf/probe.h>
int
work(int state, char *message)
{
        TNF_PROBE_2(work_start, "work_module work",
                "sunw%debug in function work",
                tnf_long,   int_input,    state,
                tnf_string, string_input, message);
              . . .
        TNF_PROBE_0(work_end, "work_module work", "");
}
```

SEE ALSO

cc(1B), **ld**(1), **prex**(1), **tnfdump**(1), **dlopen**(3X), **TNF_DECLARE_RECORD**(3X), **tnf_process_disable**(3X)

NOTES

If attaching to a running program with **prex**(1) to control the probes, compile the program with –**ltnfprobe** or start the program with the environment variable **LD_PRELOAD** set to libtnfprobe.so.1 (see **ld**(1)). If **libtnfprobe** is explicitly linked in to the program, it has to be before **libthread** on the link line.

NAME

a64l, l64a – convert between long integer and base-64 ASCII string

SYNOPSIS

```
#include <stdlib.h>
long a64l(const char *s);
char *l64a(long l);
```

MT-LEVEL

MT-Safe

DESCRIPTION

These functions are used to maintain numbers stored in base-64 ASCII characters. These characters define a notation by which long integers can be represented by up to six characters; each character represents a ''digit'' in a radix-64 notation.

The characters used to represent ''digits'' are **.** for 0, **/** for 1, **0** through **9** for 2–11, **A** through **Z** for 12–37, and **a** through **z** for 38–63.

a64l( ) takes a pointer to a null-terminated base-64 representation and returns a corresponding **long** value. If the string pointed to by *s* contains more than six characters, **a64l( )** will use the first six.

a64l( ) scans the character string from left to right with the least significant digit on the left, decoding each character as a 6-bit radix-64 number.

l64a( ) takes a **long** argument and returns a pointer to the corresponding base-64 representation. If the argument is 0, **l64a( )** returns a pointer to a null string.

NOTES

The value returned by **l64a( )** is a pointer into a static buffer, the contents of which are overwritten by each call. In the case of multithreaded applications, the return value is a pointer to thread specific data.

NAME

abort – terminate the process abnormally

SYNOPSIS

```
#include <stdlib.h>
void abort(void);
```

MT-LEVEL

Safe

DESCRIPTION

abort( ) causes abnormal process termination to occur, unless the signal **SIGABRT** is being caught and the signal handler does not return. The abnormal termination processing includes at least the effect of **fclose**(3S) on all open streams and message catalogue descriptors, and the default actions defined for **SIGABRT**. The **SIGABRT** signal is sent to the calling process as if by means of the **raise**(3C) function with the argument **SIGABRT**.

The status made available to **wait**(2) or **waitpid**(2) by **abort** will be that of a process terminated by the **SIGABRT** signal. **abort** will override blocking or ignoring the **SIGABRT** signal.

SEE ALSO

exit(2), **getrlimit**(2), **kill**(2), **wait**(2), **waitpid**(2), **fclose**(3S), **raise**(3C), **signal**(3C)

NOTES

Catching the signal is intended to provide the application writer with a portable means to abort processing, free from possible interference from any implementation-provided library functions. If **SIGABRT** is neither caught nor ignored, and the current directory is writable, a core dump may be produced.

NAME

abs, labs, llabs – return absolute value of integer

SYNOPSIS

```
#include <stdlib.h>
int abs(int val);
long labs(long lval);
long long llabs(long long llval);
```

MT-LEVEL

MT-Safe

DESCRIPTION

abs( ) returns the absolute value of its **int** operand. **labs( )** returns the absolute value of its **long** operand. **llabs( )** returns the absolute value of its **long long** operand.

NOTES

In 2's-complement representation, the absolute value of the largest magnitude negative integral value is undefined.

NAME

accept – accept a connection on a socket

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>

int accept(int s, struct sockaddr *addr, int *addrlen);
```

MT-LEVEL

Safe

DESCRIPTION

The argument *s* is a socket that has been created with **socket**(3N) and bound to an address with **bind**(3N), and that is listening for connections after a call to **listen**(3N). **accept( )** extracts the first connection on the queue of pending connections, creates a new socket with the properties of *s*, and allocates a new file descriptor, *ns*, for the socket. If no pending connections are present on the queue and the socket is not marked as non-blocking, **accept( )** blocks the caller until a connection is present. If the socket is marked as non-blocking and no pending connections are present on the queue, **accept( )** returns an error as described below. **accept( )** uses the **netconfig**(4) file to determine the STREAMS device file name associated with *s*. This is the device on which the connect indication will be accepted. The accepted socket, *ns*, is used to read and write data to and from the socket that connected to *ns*; it is not used to accept more connections. The original socket (*s*) remains open for accepting further connections.

The argument *addr* is a result parameter that is filled in with the address of the connecting entity as it is known to the communications layer. The exact format of the *addr* parameter is determined by the domain in which the communication occurs.

addrlen is a value-result parameter. Initially, it contains the amount of space pointed to by *addr*; on return it contains the length in bytes of the address returned.

accept( ) is used with connection-based socket types, currently with **SOCK_STREAM**.

It is possible to **select**(3C) or **poll**(2) a socket for the purpose of an **accept( )** by selecting or polling it for a read. However, this will only indicate when a connect indication is pending; it is still necessary to call **accept( )**.

RETURN VALUES

accept( ) returns **–1** on error. If it succeeds, it returns a non-negative integer that is a descriptor for the accepted socket.

ERRORS

accept( ) will fail if:

EBADF	The descriptor is invalid.
EINTR	The accept attempt was interrupted by the delivery of a signal.
ENODEV	The protocol family and type corresponding to *s* could not be found in the **netconfig** file.
ENOMEM	There was insufficient user memory available to complete the operation.

ENOSR	There were insufficient STREAMS resources available to complete the operation.
ENOTSOCK	The descriptor does not reference a socket.
EOPNOTSUPP	The referenced socket is not of type **SOCK_STREAM**.
EPROTO	A protocol error has occurred; for example, the STREAMS protocol stack has not been initialized.
EWOULDBLOCK	The socket is marked as non-blocking and no connections are present to be accepted.

SEE ALSO **poll**(2), **bind**(3N), **connect**(3N), **listen**(3N), **select**(3C), **socket**(3N), **netconfig**(4)

NAME

aclcheck – check the validity of an ACL

SYNOPSIS

#include <sys/acl.h>

int aclcheck(aclent_t **aclbufp*, **int** *nentries*, **int** **which*);

DESCRIPTION

aclcheck() checks the validity of an ACL pointed to by *aclbufp*. *nentries* is the number of entries contained in the buffer. *which* returns the index of the first entry that is invalid.

The function verifies that an ACL pointed to by *aclbufp* is valid according to the following rules:

There must be exactly one **group_obj** ACL entry.

There must be exactly one **user_obj** ACL entry.

There must be exactly one **other_obj** ACL entry.

If there are any **group** ACL entries, then the group ID in each group ACL entry must be unique.

If there are any **user** ACL entries, then the user ID in each user ACL entry must be unique.

If there are any **group** or **user** ACL entries, then there must be exactly one **class_obj** ACL entry.

If there are any default ACL entries, then the following apply:

There must be exactly one default **group_obj** ACL entry.

There must be exactly one default **other_obj** ACL entry.

There must be exactly one default **user_obj** ACL entry.

If there are any default **group** entries, then the group ID in each default **group** ACL entry must be unique.

If there are any default **user** entries, then the user ID in each default **user** ACL entry must be unique.

If there are any default **group** or **user** entries, then there must be exactly one default **class_obj** ACL entry.

If any of the above rules are violated, then the function fails with *errno* set to **EINVAL.**

RETURN VALUES

If the ACL is valid, **alcheck( )** will return **0**. Otherwise **errno** is set to **EINVAL** and return code is set to one of the following.

GRP_ERROR	There is more than one (default) **group_obj** ACL entry.
USER_ERROR	There is more than one (default) **user_obj** ACL entry.
CLASS_ERROR	There is more than one (default) **class_obj** ACL entry.
OTHER_ERROR	There is more than one (default) **other_obj** ACL entry.
DUPLICATE_ERROR	Duplicate (default) entries of **user** or **group**.
ENTRY_ERROR	The entry type is invalid.

MISS_ERROR	Missing (default) **group_obj, user_obj, class_obj**, or **other_obj** entries. *which* returns **-1** in this case.
MEM_ERROR	The system can't allocate any memory. *which* returns **-1** in this case.

SEE ALSO **acl**(2), **aclsort**(3)

NAME

aclsort – sort an ACL

SYNOPSIS

```
#include <sys/acl.h>
int aclsort(int nentries, int calclass, aclent_t *aclbufp);
```

DESCRIPTION

aclbufp points to a buffer containing ACL entries. *nentries* specifies the number of ACL entries in the buffer. *calclass*, if non-zero, indicates that the **CLASS_OBJ** permissions should be recalculated. The union of the permission bits associated with all ACL entries in the buffer other than **CLASS_OBJ, OTHER_OBJ**, and **USER_OBJ** is calculated. The result is copied to the permission bits associated with the **CLASS_OBJ** entry.

aclsort() sorts the contents of the ACL buffer as follows:

Entries will be in the order **USER_OBJ, USER, GROUP_OBJ, GROUP, CLASS_OBJ, OTHER_OBJ, DEF_USER_OBJ, DEF_USER, DEF_GROUP_OBJ, DEF_GROUP, DEF_CLASS_OBJ,** and **DEF_OTHER_OBJ.**

Entries of type **USER, GROUP, DEF_USER,** and **DEF_GROUP** will sorted in increasing order by id.

aclsort( ) will succeed if all of the following are true:

There is exactly one entry each of type **USER_OBJ, GROUP_OBJ, CLASS_OBJ,** and **OTHER_OBJ**.

There is exactly one entry each of type **DEF_USER_OBJ, DEF_GROUP_OBJ, DEF_CLASS_OBJ**, and **DEF_OTHER_OBJ** if there are any default entries.

Entries of type **USER, GROUP, DEF_USER**, or **DEF_GROUP** may not contain duplicate entries. A duplicate entry is one of the same type containing the same numeric id.

RETURN VALUES

Upon successful completion, the return value is **0**. Otherwise, the return value is **-1**.

SEE ALSO

acl(2), **aclcheck**(3)

NAME

acltomode, aclfrommode – convert an ACL to/from permission bits

SYNOPSIS

#include <sys/types.h>

#include <sys/acl.h>

int acltomode(aclent_t **aclbufp***, int** *nentries*, **mode_t** **modep***);**

int aclfrommode(aclent_t **aclbufp***, int** *nentries*, **mode_t** **modep***);**

DESCRIPTION

acltomode() converts an ACL pointed to by *aclbufp* into permission bits. If the **USER_OBJ** ACL entry, **GROUP_OBJ** ACL entry, or the **OTHER_OBJ** ACL entry cannot be found in the ACL buffer, then the function fails with **errno** set to **EINVAL.**

The **USER_OBJ** ACL entry permission bits are copied to the file owner class bits in the permission bits buffer. The **OTHER_OBJ** ACL entry permission bits are copied to the file other class bits in the permission bits buffer. If there is a **CLASS_OBJ** ACL entry, then the **CLASS_OBJ** ACL entry permission bits are copied to the file group class bits in the permission bits buffer. Otherwise, the **GROUP_OBJ** ACL entry permission bits are copied to the file group class bits in the permission bits buffer.

aclfrommode() converts permission bits into an ACL pointed to by *aclbufp*. If the **USER_OBJ** ACL entry, **GROUP_OBJ** ACL entry, or the **OTHER_OBJ** ACL entry cannot be found in the ACL buffer, then the function fails with **errno** set to *EINVAL*.

The file owner class bits from permission bits buffer are copied to the **USER_OBJ** ACL entry. The file other class bits from permission bits buffer are copied to the **OTHER_OBJ** ACL entry. If there is a **CLASS_OBJ** ACL entry, then the file group class bits from permission bits buffer are copied to the **CLASS_OBJ** ACL entry, and the **GROUP_OBJ** ACL entry is not modified. Otherwise, the file group class bits from permission bits buffer are copied to the **GROUP_OBJ** ACL entry.

nentries is the number of ACL entries in the buffer pointed to by *aclbufp*.

RETURN VALUES

Upon successful completion, the function returns **0**. Otherwise, a value of **-**1 is returned and **errno** is set to indicate the error.

SEE ALSO

acl(2)

NAME

acltopbits, aclfrompbits – convert an ACL to/from permission bits

SYNOPSIS

```
#include <sys/types.h>
#include <sys/acl.h>
int acltopbits(aclent_t *aclbufp, int nentries, mode_t *pbitsp );
int aclfrompbits(aclent_t *aclbufp, int nentries, mode_t *pbitsp );
```

DESCRIPTION

acltopbits( ) converts an ACL pointed to by *aclbufp* into permission bits. If the **USER_OWNER** ACL entry, **GROUP_OWNER** ACL entry, or the **OTHER** ACL entry cannot be found in the ACL buffer, then the function fails with **errno** set to **EINVAL.**

The **USER_OWNER** ACL entry permission bits are copied to the file owner class bits in the permission bits buffer. The **OTHER** ACL entry permission bits are copied to the file other class bits in the permission bits buffer. If there is a **MASK** ACL entry, then the **MASK** ACL entry permission bits are copied to the file group class bits in the permission bits buffer. Otherwise, the **GROUP_OWNER** ACL entry permission bits are copied to the file group class bits in the permission bits buffer.

aclfrompbits( ) converts permission bits into an ACL pointed to by *aclbufp*. If the **USER_OWNER** ACL entry, **GROUP_OWNER** ACL entry, or the **OTHER** ACL entry cannot be found in the ACL buffer, then the function fails with **errno** set to **EINVAL**.

The file owner class bits from permission bits buffer are copied to the **USER_OWNER** ACL entry. The file other class bits from permission bits buffer are copied to the **OTHER** ACL entry. If there is a **MASK** ACL entry, then the file group class bits from permission bits buffer are copied to the **MASK** ACL entry, and the **GROUP_OWNER** ACL entry is not modified. Otherwise, the file group class bits from permission bits buffer are copied to the **GROUP_OWNER** ACL entry.

nentries is the number of ACL entries in the buffer pointed to by *aclbufp*.

RETURN VALUES

Upon successful completion, the function returns **0**. Otherwise, a value of **-1** is returned and **errno** is set to indicate the error.

SEE ALSO

acl(2)

NAME

acltotext, aclfromtext – convert an internal representation to/from external representation

SYNOPSIS

```
#include <sys/acl.h>

char *acltotext(aclent_t *aclbufp, int aclcnt);

aclent_t *aclfromtext(char *acltextp, int *aclcnt);
```

DESCRIPTION

acltotext() converts an internal ACL representation pointed to by *aclbufp* into an external ACL representation. The space for the external text string is obtained using **malloc**(3C). The caller is responsible for freeing the space when it's done.

aclfromtext() converts an external ACL representation pointed to by *acltextp* into an internal ACL representation. The space for the list of ACL entries is obtained using **malloc**(3C). The caller is responsible for freeing the space when it's done. *aclcnt* is returned to indicate the number of acl entries found.

An external ACL representation is defined as follows:

<acl_entry>[,<acl_entry>]...

Each <acl_entry> contains one ACL entry. The external representation of an ACL entry contains three colon-separated fields. The first field contains the ACL entry tag type. The entry type keywords are defined as:

user	This ACL entry with no uid specified in the ACL entry id field specifies the access granted to the owner of the object. Otherwise, this ACL entry specifies the access granted to a specific user-name or user-id number.
group	This ACL entry with no gid specified in the ACL entry id field specifies the access granted to the owning group of the object. Otherwise, this ACL entry specifies the access granted to a specific group-name or group-id number.
other	This ACL entry specifies the access granted to any user or group that does not match any other ACL entry.
mask	This ACL entry specifies the maximum access granted to user or group entries.
defaultuser	This ACL entry with no uid specified in the ACL entry id field specifies the default access granted to the owner of the object. Otherwise, this ACL entry specifies the default access granted to a specific user-name or user-id number.
defaultgroup	This ACL entry with no gid specified in the ACL entry id field specifies the default access granted to the owning group of the object. Otherwise, this ACL entry specifies the default access granted to a specific group-name or group-id number.
defaultother	This ACL entry specifies the default access for other entry.
defaultmask	This ACL entry specifies the default access for mask entry.

The second field contains the ACL entry id. It is as follows:

uid This field specifies a user-name, or user-id if there is no user-name associated with the user-id number.

gid This field specifies a group-name, or group-id if there is no group-name associated with the group-id number.

empty It is used by user, group, other, and mask ACL entry types.

The third field contains the following symbolic discretionary access permissions:

r read permission

w write permission

x execute/search permission

- no access

RETURN VALUES

Upon successful completion, the function returns a pointer to a text string (**acltotext()**) or to a list of ACL entries (**aclfromtext()**). Otherwise, it returns **NULL.**

SEE ALSO

acl(2), **malloc**(3C)

NAME

addsev – define additional severities

SYNOPSIS

int addsev(int *int_val***, const char** ∗*string***);**

MT-LEVEL

MT-safe

DESCRIPTION

The function **addsev()** defines additional severities for use in subsequent calls to **pfmt()** or **lfmt()**. **addsev()** associates an integer value *int_val* in the range [5-255] with a character *string*. It overwrites any previous string association with *int_val* and *string*.

If *int_val* is ORed with the *flags* passed to subsequent calls **pfmt()** or **lfmt()**, *string* will be used as severity.

Passing a NULL *string* removes the severity.

Add-on severities are only effective within the applications defining them.

RETURN VALUE

addsev() returns **0** in case of success, **–1** otherwise.

USAGE

Only the standard severities are automatically displayed per the locale in effect at run-time. An application must provide the means for displaying locale-specific versions of add-on severities.

EXAMPLE

```
#define Panic    5
setlabel("APPL");
setcat("my_appl");
addsev(Panic, gettxt(":26", "PANIC"));
/* ... */
lfmt(stderr, MM_SOFT|MM_APPL|PANIC, ":12:Cannot locate database\n");
```

will display the message to *stderr* and forward to the logging service:

```
APPL: PANIC: Cannot locate database
```

SEE ALSO

gettxt(3C), **lfmt**(3C), **pfmt**(3C).

NAME

addseverity – build a list of severity levels for an application for use with fmtmsg

SYNOPSIS

```
#include <fmtmsg.h>
int addseverity(int severity, const char *string);
```

MT-LEVEL

Safe

DESCRIPTION

The **addseverity( )** function builds a list of severity levels for an application to be used with the message formatting facility, **fmtmsg( )**. *severity* is an integer value indicating the seriousness of the condition, and *string* is a pointer to a string describing the condition (string is not limited to a specific size).

If **addseverity( )** is called with an integer value that has not been previously defined, the function adds that new severity value and print string to the existing set of standard severity levels.

If **addseverity( )** is called with an integer value that has been previously defined, the function redefines that value with the new print string. Previously defined severity levels may be removed by supplying the **NULL** string. If **addseverity( )** is called with a negative number or an integer value of 0, 1, 2, 3, or 4, the function fails and returns –1. The values 0–4 are reserved for the standard severity levels and cannot be modified. Identifiers for the standard levels of severity are:

MM_HALT	Indicates that the application has encountered a severe fault and is halting. Produces the print string **HALT**.
MM_ERROR	Indicates that the application has detected a fault. Produces the print string **ERROR**.
MM_WARNING	Indicates a condition that is out of the ordinary, that might be a problem, and should be watched. Produces the print string **WARNING**.
MM_INFO	Provides information about a condition that is not in error. Produces the print string **INFO**.
MM_NOSEV	Indicates that no severity level is supplied for the message.

Severity levels may also be defined at run time using the **SEV_LEVEL** environment variable (see **fmtmsg**(3C)).

EXAMPLES

When the function **addseverity( )** is used as follows:

```
addseverity(7,"ALERT")
```

the following call to **fmtmsg( )**:

```
fmtmsg(MM_PRINT, "UX:cat", 7, "invalid syntax", "refer to manual",
"UX:cat:001")
```

produces:

```
UX:cat: ALERT: invalid syntax
TO FIX: refer to manual   UX:cat:001
```

RETURN VALUES

addseverity( ) returns **MM_OK** on success or **MM_NOTOK** on failure.

SEE ALSO

fmtmsg(1), **fmtmsg**(3C), **gettxt**(3C), **printf**(3S)

NAME

aio_cancel – cancel asynchronous I / O request

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <aio.h>
int aio_cancel(int fildes, struct aiocb *aiocbp );
```

MT-LEVEL

MT-Safe

DESCRIPTION

aio_cancel( ) attempts to cancel either one or all outstanding asynchronous I / O requests pending on the file descriptor specified by *fildes*. If *aiocbp* is **NULL**, then all such outstanding cancelable requests are canceled; otherwise, the individual request referenced by *aiocbp* references will be canceled.

Normal completion notification occurs even for asynchronous I / O operations that are successfully canceled. If there are requests which cannot be canceled, then the normal asynchronous completion process takes place for those requests, and their associated **aiocb** structures are not modified.

```
struct aiocb {
  int              aio_fildes;      /* file descriptor */
  volatile void    *aio_buf;        /* buffer location */
  size_t           aio_nbytes;      /* length of transfer */
  off_t            aio_offset;      /* file offset */
  int              aio_reqprio;     /* request priority offset */
  struct sigevent  aio_sigevent;    /* signal number and offset */
  int              aio_lio_opcode;  /* listio operation */
};
struct sigevent {
  int              sigev_notify;    /* notification mode */
  int              sigev_signo;     /* signal number */
  union sigval     sigev_value;     /* signal value */
};
union sigval {
  int              sival_int;       /* integer value */
  void             *sival_ptr;      /* pointer value */
};
```

RETURN VALUES

If the requested operation(s) were canceled, **aio_cancel( )** returns **AIO_CANCELED**. But if at least one of the requested operation(s) cannot be canceled because it is in progress, then **AIO_NOTCANCELED** is returned, and the application may determine the state of affairs for these operation(s) by using **aio_error**(3R). If all of the operation(s) had already completed, **AIO_ALLDONE** is returned. Otherwise, **aio_cancel( )** returns **-1,** and sets **errno** to indicate the error condition.

ERRORS

EBADF *fildes* is not a valid file descriptor.

ENOSYS **aio_cancel( )** is not supported by this implementation.

SEE ALSO

aio_return(3R), **aio_read**(3R)

NOTES

Applications compiled under Solaris 2.3 and 2.4 and using POSIX aio must be recompiled to work correctly when Solaris supports the Asynchronous Input and Output option.

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Asynchronous Input and Output option. It is our intention to provide support for these interfaces in future releases.

NAME

aio_fsync – asynchronous file synchronization

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

#include <aio.h>

int aio_fsync(int *op***, aiocb** **aiocbp***);**

MT-LEVEL

MT-Safe

DESCRIPTION

aio_fsync( ) queues an asynchronous **fsync**(3C) or **fdatasync**(3R) request for all the currently queued I / O operations on the file referenced by *aiocbp->***aio_fildes,** and returns control immediately. This request is serviced concurrently with other activity of the process. If *op* is **O_DSYNC**, all I / O operations are completed by a call to **fdatasync**(3R) (synchronized I / O data integrity completion). If *op* is **O_SYNC**, all I / O operations are completed by a call to **fsync**(3C) (synchronized I / O file integrity completion). (see **fcntl**(5) definitions of **O_DSYNC** and **O_SYNC**.)

When the request is queued, the error status for the operation is **EINPROGRESS**. When all data has been successfully transferred, the error status is reset to reflect the success or failure of the operation. **aio_return**(3R) and **aio_error**(3R) may be used with this *aiocbp* value to monitor both the return and the error status of the asynchronous operation while it is proceeding.

*aiocbp->***aio_sigevent** defines the signal to be generated upon I / O completion.
If *aiocbp->***aio_sigevent.sigev_signo** is non-zero, then a signal will be generated when all I / O operations have achieved syncronized I / O completion.

```
struct aiocb {
  int              aio_fildes;      /* file descriptor */
  volatile void    *aio_buf;        /* buffer location */
  size_t           aio_nbytes;      /* length of transfer */
  off_t            aio_offset;      /* file offset */
  int              aio_reqprio;     /* request priority offset */
  struct sigevent  aio_sigevent;    /* signal number and offset */
  int              aio_lio_opcode;  /* listio operation */
};
struct sigevent {
  int              sigev_notify;    /* notification mode */
  int              sigev_signo;     /* signal number */
  union sigval     sigev_value;     /* signal value */
};
union sigval {
  int              sival_int;       /* integer value */
  void             *sival_ptr;      /* pointer value */
};
```

RETURN VALUES

If the I / O operation is successfully queued, **aio_fsync( )** returns **0.** Otherwise, it returns **- 1,** and sets **errno** to indicate the error condition.

ERRORS

EAGAIN	The requested asynchronous operation was not queued due to temporary resource limitations.
EBADF	*aiocbp->* **aio_fildes** is not a valid file descriptor open for writing.
EINVAL	This implementation does not support synchronized I / O for this file.
	A value of *op* other than **O_DSYNC** or **O_SYNC** was specified.
ENOSYS	**aio_fsync( )** is not supported by this implementation.

SEE ALSO

fcntl(2), **open**(2), **read**(2), **write**(2), **fsync**(3C), **aio_error**(3R), **aio_return**(3R), **fdatasync**(3R), **fcntl**(5)

NOTES

If **aio_fsync( )** fails, outstanding I / O operations are not guaranteed to have been completed.

Applications compiled under Solaris 2.3 and 2.4 and using POSIX aio must be recompiled to work correctly when Solaris supports the Asynchronous Input and Output option.

BUGS

In Solaris 2.5, these functions always return –**1** and set **errno** to **ENOSYS,** because this release does not support the Asynchronous Input and Output option. It is our intention to provide support for these interfaces in future releases.

NAME

aio_read, aio_write – asynchronous read and write operations

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <aio.h>

int aio_read(struct aiocb *aiocbp);

int aio_write(struct aiocb *aiocbp);

struct aiocb {
  int              aio_fildes;     /* file descriptor */
  volatile void    *aio_buf;       /* buffer location */
  size_t           aio_nbytes;     /* length of transfer */
  off_t            aio_offset;     /* file offset */
  int              aio_reqprio;    /* request priority offset */
  struct sigevent  aio_sigevent;   /* signal number and offset */
  int              aio_lio_opcode; /* listio operation */
};

struct sigevent {
  int              sigev_notify;   /* notification mode */
  int              sigev_signo;    /* signal number */
  union sigval     sigev_value;    /* signal value */
};

union sigval {
  int              sival_int;      /* integer value */
  void             *sival_ptr;     /* pointer value */
};
```

MT-LEVEL

MT-Safe

DESCRIPTION

aio_read( ) queues an asynchronous read request, and returns control immediately. Rather than blocking until completion, the read operation continues concurrently with other activity of the process.

Upon enqueuing the request, the calling process reads *aiocbp*->**nbytes** from the file referred to by *aiocbp*->**fildes** into the buffer pointed to by *aiocbp*->**aio_buf**. *aiocbp*->**offset** marks the absolute position from the beginning of the file (in bytes) at which the read begins.

aio_write( ) queues an asynchronous write request, and returns control immediately. Rather than blocking until completion, the write operation continues concurrently with other activity of the process.

Upon enqueuing the request, the calling process writes *aiocbp*->**nbytes** from the buffer pointed to by *aiocbp*->**aio_buf** into the file referred to by *aiocbp*->**fildes**. If **O_APPEND** is set for *aiocbp*->**fildes**, **aio_write( )** operations append to the file in the same order as the calls were made.

If **O_APPEND** is not set for the file descriptor, then the write operation will occur at the absolute position from the beginning of the file plus *aiocbp*->**offset** (in bytes).

These asynchronous operations are submitted at a priority equal to the calling process' scheduling priority minus *aiocbp*->**aio_reqprio**.

aiocb->**aio_sigevent** defines both the signal to be generated and how the calling process will be notified upon I/O completion. If **aio_sigevent.sigev_notify** is **SIGEV_NONE**, then no signal will be posted upon I/O completion, but the error status and the return status for the operation will be set appropriately. If **aio_sigevent.sigev_notify** is **SIGEV_SIGNAL**, then the signal specified in **aio_sigevent.sigev_signo** will be sent to the process. If the **SA_SIGINFO** flag is set for that signal number, then the signal will be queued to the process and the value specified in **aio_sigevent.sigev_value** will be the **si_value** component of the generated signal (see **siginfo**(5)).

RETURN VALUES

If the I/O operation is successfully queued, **aio_read( )** and **aio_write( )** return **0,** otherwise, they return **-1,** and set **errno** to indicate the error condition. *aiocbp* may be used as an argument to **aio_error**(3R) and **aio_return**(3R) in order to determine the error status and the return status of the asynchronous operation while it is proceeding.

ERRORS

EAGAIN The requested asynchronous I/O operation was not queued due to system resource limitations.

ENOSYS **aio_read( )** or is not supported by this implementation.

EBADF If the calling function is **aio_read( ),** and *aiocbp*->**fildes** is not a valid file descriptor open for reading. If the calling function is **aio_write( ),** and *aiocbp*->**fildes** is not a valid file descriptor open for writing.

EINVAL
- The file offset value implied by *aiocbp*->**aio_offset** would be invalid,
- *aiocbp*->**aio_reqprio** is not a valid value, or
- *aiocbp*->**aio_nbytes** is an invalid value.

ECANCELED The requested I/O was canceled before the I/O completed due to an explicit **aio_cancel**(3R) request.

EINVAL The file offset value implied by *aiocbp*->**aio_offset** would be invalid.

SEE ALSO

close(2), **exec**(2), **exit**(2), **fork**(2), **lseek**(2), **read**(2), **write**(2), **aio_cancel**(3R), **aio_return**(3R), **lio_listio**(3R), **siginfo**(5)

NOTES

For portability, the application should set *aiocb*->**aio_reqprio** to **0.**

Applications compiled under Solaris 2.3 and 2.4 and using POSIX aio must be recompiled to work correctly when Solaris supports the Asynchronous Input and Output option.

BUGS

In Solaris 2.5, these functions always return –**1** and set **errno** to **ENOSYS,** because this release does not support the Asynchronous Input and Output option. It is our intention to provide support for these interfaces in future releases.

NAME

aio_return, aio_error – retrieve return or error status of asynchronous I / O operation

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <aio.h>

ssize_t aio_return(struct aiocb * aiocbp);

int aio_error(const struct aiocb * aiocbp);

struct aiocb {
  int               aio_fildes;     /* file descriptor */
  volatile void     *aio_buf;       /* buffer location */
  size_t            aio_nbytes;     /* length of transfer */
  off_t             aio_offset;     /* file offset */
  int               aio_reqprio;    /* request priority offset */
  struct sigevent   aio_sigevent;   /* signal number and offset */
  int               aio_lio_opcode; /* listio operation */
};

struct sigevent {
  int               sigev_notify;   /* notification mode */
  int               sigev_signo;    /* signal number */
  union sigval      sigev_value;    /* signal value */
};

union sigval {
  int               sival_int;      /* integer value */
  void              *sival_ptr;     /* pointer value */
};
```

MT-LEVEL

Async-Signal-Safe

DESCRIPTION

aio_return( ) returns the return status of the asynchronous I / O request associated with the **aiocb** structure pointed to by *aiocbp.*

aio_error( ) returns the error status of the asynchronous I / O request associated with the **aiocb** structure pointed to by *aiocbp.*

aio_return( ) should be called only once to retrieve the valid return status of a given asynchronous operation, after **aio_error( )** has returned a value other than **EINPROGRESS.**

RETURN VALUES

If the asynchronous I / O operation has completed successfully, **aio_return( )** returns the return status, as described for **read**(2), **write**(2), and **fsync**(3C).

If the asynchronous I / O operation has completed successfully, **aio_error( )** returns **0.** If the operation has not yet completed, then **EINPROGRESS** is returned. If the asynchronous I / O operation has completed unsuccessfully, then the error status, as described for **read**(2), **write**(2), and **fsync**(3C) is returned.

If unsuccessful, **aio_return( )** or **aio_error( )** return **-1,** and set **errno** to indicate the error condition.

ERRORS

EINVAL *aiocbp* does not reference an asynchronous operation which has completed or failed.

ENOSYS **aio_return( )** or **aio_error( )** is not supported by this implementation.

EXAMPLES

```
#include <aio.h>
#include <errno.h>
#include <signal.h>
struct aiocb        my_aiocb;
struct sigaction    my_sigaction;
void                my_aio_handler(int, siginfo_t *, void *);
...
my_sigaction.sa_flags = SA_SIGINFO;
my_sigaction.sa_sigaction = my_aio_handler;
sigsetempty(&my_sigaction.sa_mask);
(void) sigaction(SIGRTMIN, &my_sigaction, NULL);
...
my_aiocb.aio_sigevent.sigev_notify = SIGEV_SIGNAL;
my_aiocb.aio_sigevent.sigev_signo = SIGRTMIN;
my_aiocb.aio_sigevent.sigev_value.sival_ptr = &myaiocb;
...
(void) aio_read(&my_aiocb);
...
void
my_aio_handler(int signo, siginfo_t *siginfo, void *context) {
int    my_errno;
struct aiocb   *my_aiocbp;
my_aiocbp = siginfo.si_value.sival_ptr;
      if ((my_errno = aio_error(my_aiocb)) != EINPROGRESS) {
            int      my_status = aio_return(my_aiocb);
            if (my_status >= 0){/* start another operation */
                  ...
            } else   {                    /* handle I/O error */
                  ...
            }
      }
}
```

SEE ALSO

close(2), **exec**(2), **exit**(2), **fork**(2), **lseek**(2), **read**(2), **write**(2), **fsync**(3C), **aio_cancel**(3R), **aio_fsync**(3R), **aio_read**(3R), **lio_listio**(3R)

NOTES

Applications compiled under Solaris 2.3 and 2.4 and using POSIX aio must be recompiled to work correctly when Solaris supports the Asynchronous Input and Output option.

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Asynchronous Input and Output option. It is our intention to provide support for these interfaces in future releases.

NAME

aio_suspend – wait for asynchronous I / O request

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

#include <aio.h>

int aio_suspend(const struct aiocb ∗ **const** *list***[ ], int** *nent*,
const struct timespec ∗*timeout***);**

MT-LEVEL

Async-Signal-Safe

DESCRIPTION

aio_suspend( ) suspends the caller until at least one of the asynchronous I / O operations referenced by *list* has completed, until a signal interrupts the function, or, if *timeout* is not NULL, until the time interval specified by *timeout* has passed. If any of the **aiocb** structures in the list corresponds to a completed asynchronous I / O operation (i.e., the error status for the operation is not equal to **EINPROGRESS**), at the time of the call, the function returns without suspending the caller.

If the time interval indicated in the **timespec** structure pointed to by *timeout* passes before any of the I / O operations referenced by *list* are completed, then **aio_suspend( )** returns with an error.

list is an array of pointers to asynchronous I / O control blocks. *nent* indicates the number of elements in this array. Each **aiocb** structure pointed to must have been used in initiating an asynchronous I / O request via **aio_read**(3R), **aio_write**(3R), **aio_fsync**(3R), or **lio_listio**(3R). This array may contain NULL pointers which will be ignored.

```
struct aiocb {
  int              aio_fildes;      /* file descriptor */
  volatile void    *aio_buf;        /* buffer location */
  size_t           aio_nbytes;      /* length of transfer */
  off_t            aio_offset;      /* file offset */
  int              aio_reqprio;     /* request priority offset */
  struct sigevent  aio_sigevent;    /* signal number and offset */
  int              aio_lio_opcode;  /* listio operation */
};
struct sigevent {
  int              sigev_notify;    /* notification mode */
  int              sigev_signo;     /* signal number */
  union sigval     sigev_value;     /* signal value */
};
union sigval {
  int              sival_int;       /* integer value */
  void             *sival_ptr;      /* pointer value */
};
```

```
struct timespec {
  time_t        tv_sec;        /* seconds */
  long          tv_nsec;       /* and nanoseconds */
};
```

RETURN VALUES

If **aio_suspend( )** returns after one or more asynchronous I / O operations have completed, it returns **0.** Otherwise, it returns **-1,** and sets **errno** to indicate the error condition.

The application may determine which asynchronous I / O had completed with both the associated error and return status of **aio_return**(3R), and **aio_error**(3R).

ERRORS

EAGAIN — No asynchronous I / O indicated in the list referenced by *list* completed in the time interval indicated by *timeout.*

EINTR — A signal interrupted the **aio_suspen( )** function. Note that, since each asynchronous I / O operation may possibly provoke a signal when it completes, this error return may be caused by the completion of one (or more) of the very I / O operations being awaited.

ENOSYS — **aio_suspend( )** is not supported by this implementation.

SEE ALSO

aio_fsync(3R), **aio_read**(3R), **aio_return**(3R), **aio_write**(3R), **lio_listio**(3R)

NOTES

Applications compiled under Solaris 2.3 and 2.4 and using POSIX aio must be recompiled to work correctly when Solaris supports the Asynchronous Input and Output option.

BUGS

In Solaris 2.5, these functions always return –**1** and set **errno** to **ENOSYS,** because this release does not support the Asynchronous Input and Output option. It is our intention to provide support for these interfaces in future releases.

NAME

aiocancel – cancel an asynchronous operation

SYNOPSIS

cc [*flag* ...] *file* ... **–laio** [*library* ...]

#include <sys/asynch.h>

int aiocancel(aio_result_t **resultp***);**

MT-LEVEL

Unsafe

DESCRIPTION

aiocancel( ) cancels the asynchronous operation associated with the result buffer pointed to by *resultp*. It may not be possible to immediately cancel an operation which is in progress and in this case, **aiocancel( )** will not wait to cancel it.

Upon successful completion, **aiocancel( )** returns 0 and the requested operation is cancelled. The application will not receive the **SIGIO** completion signal for an asynchronous operation that is successfully cancelled.

RETURN VALUES

aiocancel( ) returns:

0 on success.

–1 on failure and sets **errno** to indicate the error.

ERRORS

aiocancel( ) will fail if any of the following are true:

EACCES The parameter *resultp* does not correspond to any outstanding asynchronous operation, although there is at least one currently outstanding.

EFAULT *resultp* points to an address outside the address space of the requesting process. See **NOTES** below.

EINVAL There are not any outstanding requests to cancel.

SEE ALSO

aioread(3), **aiowait**(3)

NOTES

Passing an illegal address as *resultp* will result in setting **errno** to EFAULT *only* if it is detected by the application process.

NAME

aioread, aiowrite – asynchronous I / O operations

SYNOPSIS

cc [*flag* ...] *file* ... **–laio** [*library* ...]

#include <sys/asynch.h>

int aioread(int *fildes***, char** **bufp***, int** *bufs***, off_t** *offset***, int** *whence***, aio_result_t** **resultp***);**

int aiowrite(int *fildes***, const char** **bufp***, int** *bufs***, off_t** *offset***, int** *whence***, aio_result_t** **resultp***);**

MT-LEVEL

Unsafe

DESCRIPTION

aioread( ) initiates one asynchronous **read**(2) and returns control to the calling program. The **read( )** continues concurrently with other activity of the process. An attempt is made to read *bufs* bytes of data from the object referenced by the descriptor *fildes* into the buffer pointed to by *bufp*.

aiowrite( ) initiates one asynchronous **write**(2) and returns control to the calling program. The **write( )** continues concurrently with other activity of the process. An attempt is made to write *bufs* bytes of data from the buffer pointed to by *bufp* to the object referenced by the descriptor *fildes*.

On objects capable of seeking, the I / O operation starts at the position specified by *whence* and *offset*. These parameters have the same meaning as the corresponding parameters to the **lseek**(2) function. On objects not capable of seeking the I / O operation always start from the current position and the parameters *whence* and *offset* are ignored. The seek pointer for objects capable of seeking is not updated by **aioread( )** or **aiowrite( )**. Sequential asynchronous operations on these devices must be managed by the application using the *whence* and *offset* parameters.

The result of the asynchronous operation is stored in the structure pointed to by *resultp*:

```
int aio_return;        /* return value of read( ) or write( ) */
int aio_errno;         /* value of errno for read( ) or write( ) */
```

Upon completion of the operation both *aio_return* and *aio_errno* are set to reflect the result of the operation. AIO_INPROGRESS is not a value used by the system so the client may detect a change in state by initializing *aio_return* to this value.

The application supplied buffer *bufp* should not be referenced by the application until after the operation has completed. While the operation is *in progress*, this buffer is in use by the operating system.

Notification of the completion of an asynchronous I / O operation may be obtained synchronously through the **aiowait**(3) function, or asynchronously by installing a signal handler for the **SIGIO** signal. Asynchronous notification is accomplished by sending the process a **SIGIO** signal. If a signal handler is not installed for the **SIGIO** signal, asynchronous notification is disabled. The delivery of this instance of the **SIGIO** signal is reliable in that a signal delivered while the handler is executing is not lost. If the client ensures that **aiowait**(3) returns nothing (using a polling timeout) before returning from the signal handler, no asynchronous I / O notifications are lost. The **aiowait**(3) function is the only

way to dequeue an asynchronous notification. Note: **SIGIO** may have several meanings simultaneously: for example, that a descriptor generated **SIGIO** and an asynchronous operation completed. Further, issuing an asynchronous request successfully guarantees that space exists to queue the completion notification.

close(2), **exit**(2) and **execve()** (see **exec**(2)) will block until all pending asynchronous I/O operations can be canceled by the system.

It is an error to use the same result buffer in more than one outstanding request. These structures may only be reused after the system has completed the operation.

RETURN VALUES

aioread() and **aiowrite()** return:

0 on success.

–1 on failure and set **errno** to indicate the error.

ERRORS

EAGAIN	The number of asynchronous requests that the system can handle at any one time has been exceeded
EBADF	*fildes* is not a valid file descriptor open for reading.
EFAULT	At least one of *bufp* points to an address outside the address space of the requesting process. See **NOTES** below.
EINVAL	The parameter *resultp* is currently being used by an outstanding asynchronous request.
ENOMEM	Memory resources are unavailable to initiate request.

SEE ALSO

close(2), **exec**(2), **exit**(2), **lseek**(2), **open**(2), **read**(2), **write**(2), **aiocancel**(3), **aiowait**(3), **sigvec**(3B)

NOTES

Passing an illegal address to *bufp* will result in setting **errno** to EFAULT *only* if it is detected by the application process.

NAME

aiowait – wait for completion of asynchronous I/O operation

SYNOPSIS

cc [*flag* ...] *file* ... **–laio** [*library* ...]

#include <sys/asynch.h>
#include <sys/time.h>

aio_result_t ***aiowait(const struct timeval** ***timeout*);**

MT-LEVEL

Unsafe

DESCRIPTION

aiowait() suspends the calling process until one of its outstanding asynchronous I/O operations completes. This provides a synchronous method of notification.

If *timeout* is a non-zero pointer, it specifies a maximum interval to wait for the completion of an asynchronous I/O operation. If *timeout* is a zero pointer, then **aiowait()** blocks indefinitely. To effect a poll, the *timeout* parameter should be non-zero, pointing to a zero-valued *timeval* structure.

The *timeval* structure is defined in **<sys/time.h>** and contains the following members:

```
long  tv_sec;        /* seconds */
long  tv_usec;       /* and microseconds */
```

RETURN VALUES

On success, **aiowait()** returns a pointer to the result structure used when the completed asynchronous I/O operation was requested. On failure, it returns –1 and sets **errno** to indicate the error. **aiowait() returns** 0 if the time limit expires.

ERRORS

EFAULT *timeout* points to an address outside the address space of the requesting process. See **NOTES** below.

EINTR A signal was delivered before an asynchronous I/O operation completed.

The time limit expired.

EINVAL There are no outstanding asynchronous I/O requests.

SEE ALSO

aiocancel(3), **aioread**(3)

NOTES

aiowait() is the only way to dequeue an asynchronous notification. It may be used either inside a **SIGIO** signal handler or in the main program. One **SIGIO** signal may represent several queued events.

Passing an illegal address as *timeout* will result in setting **errno** to EFAULT *only* if it is detected by the application process.

NAME

assert – verify program assertion

SYNOPSIS

#include <assert.h>

void assert(int *expression***);**

MT-LEVEL

Safe

DESCRIPTION

This macro is useful for putting diagnostics into programs. When it is executed, if *expression* is false (zero), **assert( )** prints

Assertion failed: *expression***, file** *xyz***, line** *nnn*

on the standard error output and aborts. In the error message, *xyz* is the name of the source file and *nnn* the source line number of the **assert( )** statement. The latter are respectively the values of the preprocessor macros **__FILE__** and **__LINE__**.

Compiling with the preprocessor option –**DNDEBUG** (see **cc**(1B)), or with the preprocessor control statement **#define NDEBUG** ahead of the **#include <assert.h>** statement, will stop assertions from being compiled into the program.

SEE ALSO

cc(1B), **abort**(3C), **gettext**(3I), **setlocale**(3C)

NOTES

If the application is linked with –**lintl**, then messages printed from this function are in the native language specified by the **LC_MESSAGES** locale category; see **setlocale**(3C).

Since **assert( )** is implemented as a macro, the *expression* may not contain any string literals.

NAME

atexit – add program termination routine

SYNOPSIS

```
#include <stdlib.h>
int atexit(void (*func)(void));
```

MT-LEVEL

Safe

DESCRIPTION

atexit( ) adds the function *func( )* to a list of functions to be called without arguments on normal termination of the program. Normal termination occurs by either a call to the **exit( )** function or a return from **main( )**. At most 32 functions may be registered by **atexit( )**; the functions will be called in the reverse order of their registration.

RETURN VALUES

atexit( ) returns 0 if the registration succeeds, nonzero if it fails.

SEE ALSO

exit(3C)

NAME

au_open, au_close, au_write – construct and write audit records

SYNOPSIS

cc [*flag* ...] *file* ... **–lbsm –lsocket –lnsl –lintl** [*library* ...]

#include <bsm/libbsm.h>

int au_close(int *d***, int** *keep***, short** *event***);**

int au_open(*void***);**

int au_write(int *d***, token_t** ∗*m***);**

MT-LEVEL

Safe.

AVAILABILITY

The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See **bsmconv**(1M) for more information.

DESCRIPTION

au_open( ) allocates an audit record descriptor to which audit tokens can be written using **au_write( )**.

au_close( ) terminates the life of an audit record *d* of type *event* started by **au_open( )**. If the *keep* parameter is zero, the data contained therein is discarded and the memory used is given up by calling **free**(3C). Otherwise, the additional parameters are used to create a header token. Depending on the audit policy information obtained by **auditon**(2), additional tokens such as *sequence* and *trailer* tokens may be added to the record. **au_close( )** finally writes the record to the audit trail by calling **audit**(2).

au_write( ) adds the audit token pointed to by *m* to the audit record identified by the descriptor *d*. After this call is made the audit token is no longer available to the caller.

RETURN VALUES

A successful invocation for all calls will return a **0**.

au_open( ) returns **–1** if a descriptor could not be allocated. **au_close( )** and **au_write( )** return **–1** if *d* is not a valid descriptor or if **audit**(2) experienced an error. In the latter case, **errno** is set to indicate the error.

SEE ALSO

bsmconv(1M), **audit**(2), **auditon**(2), **au_preselect**(3), **au_to**(3), **free**(3C)

NAME

au_preselect – preselect an audit event

SYNOPSIS

cc [*flag* ...] *file* ... **–lbsm –lsocket –lnsl –lintl** [*library* ...]

#include <bsm/libbsm.h>

int au_preselect(au_event_t *event***, au_mask_t** **mask_p***, int** *sorf***, int** *flag***);**

MT_LEVEL

MT-Safe.

AVAILABILITY

The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See **bsmconv**(1M) for more information.

DESCRIPTION

au_preselect( ) determines whether or not the audit event *event* is preselected against the binary preselection mask pointed to by *mask_p* (usually obtained by a call to **getaudit**(2)). **au_preselect( )** looks up the classes associated with *event* in **audit_event**(4) and compares them with the classes in *mask_p*. If the classes associated with *event* match the classes in the specified portions of the binary preselection mask pointed to by *mask_p*, the event is said to be preselected.

sorf indicates whether the comparison is made with the success portion, the failure portion or both portions of the mask pointed to by *mask_p*.

The following are the valid values of *sorf*:

AU_PRS_SUCCESS
Compare the event class with the success portion of the preselection mask.

AU_PRS_FAILURE
Compare the event class with the failure portion of the preselection mask.

AU_PRS_BOTH
Compare the event class with both the success and failure portions of the preselection mask.

flag tells **au_preselect( )** how to read the **audit_event**(4) database. Upon initial invocation, **au_preselect( )** reads the **audit_event**(4) database and allocates space in an internal cache for each entry with **malloc**(3C). In subsequent invocations, the value of *flag* determines where **au_preselect()** obtains audit event information. The following are the valid values of *flag*:

AU_PRS_REREAD
Get audit event information by searching the **audit_event**(4) database.

AU_PRS_USECACHE
Get audit event information from internal cache created upon the initial invocation. This option is much faster.

RETURN VALUES

au_preselect( ) returns:

- **0** *event* is not preselected.
- **1** *event* is preselected.
- **-1** An error occurred. **au_preselect( )** couldn't allocate memory or couldn't find *event* in the **audit_event**(4) database.

FILES

/etc/security/audit_class	maps audit class number to audit class names and descriptions
/etc/security/audit_event	maps audit even number to audit event names and associates

SEE ALSO

bsmconv(1M), **getaudit**(2), **au_open**(3), **getauclassent**(3), **getauevent**(3), **malloc**(3C), **audit_class**(4), **audit_event**(4)

NOTES

au_preselect( ) is normally called prior to constructing and writing an audit record. If the event is not preselected, the overhead of constructing and writing the record can be saved.

NAME

au_to, au_to_arg, au_to_attr, au_to_data, au_to_groups, au_to_in_addr, au_to_ipc, au_to_ipc_perm, au_to_iport, au_to_me, au_to_opaque, au_to_path, au_to_process, au_to_return, au_to_socket, au_to_text – create audit record tokens

SYNOPSIS

cc [*flag* ...] *file* ... **–lbsm –lsocket –lnsl –lintl** [*library* ...]

```
#include <sys/types.h>
#include <sys/vnode.h>
#include <netinet/in.h>
#include <bsm/libbsm.h>

token_t *au_to_arg( char n, char *text, u_long v);

token_t *au_to_attr( struct vattr *attr);

token_t *au_to_cmd( u_long argc, char **argv, char **envp);

token_t *au_to_data(char unit_print, char unit_type, char unit_count, char *p);

token_t *au_to_groups( int *groups);

token_t *au_to_in_addr( struct inaddr *internet_addr);

token_t *au_to_iport( u_short iport);

token_t *au_to_ipc( int id);

token_t *au_to_ipc_perm( struct ipc_perm *perm);

token_t *au_to_iport( u_short iport);

token_t *au_to_me( void);

token_t *au_to_newgroups(int n, int *groups);

token_t *au_to_opaque( char *data, short bytes);

token_t *au_to_path( char *path);

token_t *au_to_process ( au_id_t auid, uid_t euid, gid_t egid, uid_t ruid, gid_t rgid,
        au_asid_t sid, au_tid_t *tid);

token_t *au_to_return( char number, u_int value);

token_t *au_to_socket( struct socket *so);

token_t *au_to_subject( au_id_t auid, uid_t euid, gid_t egid, uid_t ruid, gid_t rgid,
        pid_t pid, au_asid_t sid, au_tid_t *tid );

token_t *au_to_text( char *text);
```

MT_LEVEL

MT-Safe.

AVAILABILITY

The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See **bsmconv**(1M) for more information.

DESCRIPTION

au_to_arg() formats the data in *v* into an ''argument token.'' The *n* argument indicates the argument number. The *text* argument is a null terminated string describing the argument.

au_to_attr( ) formats the data pointed to by *attr* into a ''vnode attribute token.''

au_to_data( ) formats the data pointed to by *p* into an ''arbitrary data token.'' The *unit_print* parameter determines the preferred display base of the data and is one of **AUP_BINARY**, **AUP_OCTAL**, **AUP_DECIMAL**, **AUP_HEX**, or **AUP_STRING**. The *unit_type* parameter defines the basic unit of data and is one of **AUR_BYTE**, **AUR_CHAR**, **AUR_SHORT**, **AUR_INT**, or **AUR_LONG**. The *unit_count* parameter specifies the number of basic data units to be used and must be positive.

au_to_groups( ) formats the array of 16 integers pointed to by *groups* into a ''groups token.''

au_to_in_addr( ) formats the data pointed to by *internet_addr* into an ''internet address token.''

au_to_ipc( ) formats the data in the *id* parameter into an ''interprocess communications id token.''

au_to_ipc_perm( ) formats the data pointed to by *perm* into an ''interprocess communications permission token.''

au_to_iport( ) formats the data pointed to by *iport* into an ''ip port address token.''

au_to_me( ) collects audit information from the current process and creates a ''subject token'' by calling **au_to_subject( )**.

au_to_newgroups( ) formats the array of *n* integers pointed to by *groups* into a ''newgroups token.''

au_to_subject( ) formats an *auid* (audit user ID), an *euid* (effective user ID), an *egid* (effective group ID), a *ruid* (real user ID), an *rgid* (real group ID), a *pid* (process ID), an *sid* (audit session ID), an *tid* (audit terminal ID), into a ''subject token.''

au_to_opaque( ) formats the *bytes* bytes pointed to by *data* into an ''opaque token.'' The value of *size* must be positive.

au_to_path( ) formats the path name pointed to by *path* into a ''path token.''

au_to_process( ) formats an *auid* (audit user ID), an *euid* (effective user ID), an *egid* (effective group ID), a *ruid* (real user ID), a *rgid* (real group ID), a *pid* (process ID), an *sid* (audit session ID), and a *tid* (audit terminal ID), into a ''process token.'' A process token should be used when the process is the object of an action (ie. when the process is the receiver of a signal).

au_to_return( ) formats an error number *number* and a return value *value* into a ''return value token.''

au_to_socket( ) format the data pointed to by *so* into a ''socket token.''

au_to_text( ) formats the **NULL** terminated string pointed to by *text* into a ''text token.''

RETURN VALUES

These functions return **NULL** if memory cannot be allocated to put the resultant token into, or if an error in the input is detected.

SEE ALSO **bsmconv**(1M), **au_open**(3)

NAME

au_user_mask – get user's binary preselection mask

SYNOPSIS

cc [*flag* ...] *file* ... **–lbsm –lsocket –lnsl –lintl** [*library* ...]

#include <bsm/libbsm.h>

int au_user_mask(char ∗*username,* **au_mask_t** ∗*mask_p***);**

MT-LEVEL

MT-Safe.

AVAILABILITY

The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See **bsmconv**(1M) for more information.

DESCRIPTION

au_user_mask( ) reads the default, system wide audit classes from **audit_control**(4), combines them with the per-user audit classes from the **audit_user**(4) database, and updates the binary preselection mask pointed to by *mask_p* with the combined value.

The audit flags in the *flags* field of the **audit_control**(4) database and the *always-audit-flags* and *never-audit-flags* from the **audit_user**(4) database represent binary audit classes. These fields are combined by **au_preselect**(3) as follows:

mask = (*flags* + *always-audit-flags*) – *never-audit-flags*

au_user_mask( ) only fails if both the both the **audit_control**(4) and the **audit_user**(4) database entries could not be retrieved. This allows for flexible configurations.

RETURN VALUES

au_user_mask( ) returns:

0	Success.
-1	Failure. Both the **audit_control**(4) and the **audit_user**(4) database entries could not be retrieved.

FILES

/etc/security/audit_control	contains default parameters read by the audit daemon, **auditd**(1M)
/etc/security/audit_user	stores per-user audit event mask

SEE ALSO

login(1), **bsmconv**(1M), **getaudit**(2), **setaudit**(2), **au_preselect**(3), **getacinfo**(3), **getauusernam**(3), **audit_control**(4), **audit_user**(4)

NOTES

au_user_mask( ) should be called by programs like **login**(1) which set a process's preselection mask with **setaudit**(2). **getaudit**(2) should be used to obtain audit characteristics for the current process.

NAME

basename – return the last element of a path name

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

char ∗**basename(char** ∗*path***);**

MT-LEVEL

MT-Safe

DESCRIPTION

Given a pointer to a null-terminated character string that contains a path name, **basename( )** returns a pointer to the last element of *path*. Trailing ‘‘/’’ characters are deleted.

If *path* or ∗*path* is zero, pointer to a static constant ‘‘**.**’’ is returned.

EXAMPLES

Input string	Output pointer
/usr/lib	**lib**
/usr/	**usr**
/	**/**

SEE ALSO

basename(1), **dirname**(3G)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

bessel, j0, j1, jn, y0, y1, yn – Bessel functions

SYNOPSIS

cc [*flag* ...] *file* ... **–lm** [*library* ...]

#include <math.h>

double j0(double *x***);**

double j1(double *x***);**

double jn(int *n***, double** *x***);**

double y0(double *x***);**

double y1(double *x***);**

double yn(int *n***, double** *x***);**

MT-LEVEL

MT-Safe

DESCRIPTION

These functions calculate Bessel functions of the first and second kinds for real arguments and integer orders.

RETURN VALUES

For exceptional cases, **matherr**(3M) tabulates the values to be returned as dictated by various Standards.

SEE ALSO

exp(3M), **matherr**(3M)

DIAGNOSTICS

In IEEE 754 mode (i.e. the **–xlibmieee cc** compilation option), the functions **y0( )**, **y1( )**, and **yn( )** have logarithmic singularities at the origin, so they treat zero and negative arguments the way **log( )** does, as described in **exp**(3M). Such arguments are unexceptional for **j0( )**, **j1( )**, and **jn( )**.

NAME

bgets – read stream up to next delimiter

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

char ∗**bgets(char** ∗*buffer*, **size_t** ∗*count*, **FILE** ∗*stream*, **const char** ∗*breakstring*);

MT-LEVEL

MT-Safe

DESCRIPTION

bgets( ) reads characters from *stream* into *buffer* until either *count* is exhausted or one of the characters in *breakstring* is encountered in the stream. The read data is terminated with a null byte ('**\0**') and a pointer to the trailing null is returned. If a *breakstring* character is encountered, the last non-null is the delimiter character that terminated the scan.

Note that, except for the fact that the returned value points to the **end** of the read string rather than to the beginning, the call

bgets(buffer, sizeof buffer, stream, "\n");

is identical to

fgets (buffer, sizeof buffer, stream);

There is always enough room reserved in the buffer for the trailing null.

If *breakstring* is a null pointer, the value of *breakstring* from the previous call is used. If *breakstring* is null at the first call, no characters will be used to delimit the string.

RETURN VALUES

NULL is returned on error or end-of-file. Reporting the condition is delayed to the next call if any characters were read but not yet returned.

EXAMPLES

```
#include        <libgen.h>

char buffer[8];
/* read in first user name from /etc/passwd */
fp = fopen("/etc/passwd","r");
bgets(buffer, 8, fp, ":");
```

SEE ALSO

gets(3S)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

bind – bind a name to a socket

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>

int bind(int s, const struct sockaddr *name, int namelen);
```

MT-LEVEL

Safe

DESCRIPTION

bind( ) assigns a name to an unnamed socket. When a socket is created with **socket**(3N), it exists in a name space (address family) but has no name assigned. **bind( )** requests that the name pointed to by *name* be assigned to the socket.

RETURN VALUES

If the bind is successful, **0** is returned. A return value of **–1** indicates an error, which is further specified in the global **errno**.

ERRORS

The **bind( )** call will fail if:

EACCES	The requested address is protected and the current user has inadequate permission to access it.
EADDRINUSE	The specified address is already in use.
EADDRNOTAVAIL	The specified address is not available on the local machine.
EBADF	*s* is not a valid descriptor.
EINVAL	*namelen* is not the size of a valid address for the specified address family.
EINVAL	The socket is already bound to an address.
ENOSR	There were insufficient STREAMS resources for the operation to complete.
ENOTSOCK	*s* is a descriptor for a file, not a socket.

The following errors are specific to binding names in the UNIX domain:

EACCES	Search permission is denied for a component of the path prefix of the pathname in *name*.
EIO	An I/O error occurred while making the directory entry or allocating the inode.
EISDIR	A null pathname was specified.
ELOOP	Too many symbolic links were encountered in translating the pathname in *name*.
ENOENT	A component of the path prefix of the pathname in *name* does not exist.

ENOTDIR	A component of the path prefix of the pathname in *name* is not a directory.
EROFS	The inode would reside on a read-only file system.

SEE ALSO

unlink(2), **socket**(3N)

NOTES

Binding a name in the UNIX domain creates a socket in the file system that must be deleted by the caller when it is no longer needed (using **unlink**(2)).

The rules used in name binding vary between communication domains.

NAME

bsdmalloc, malloc, free, realloc – memory allocator

SYNOPSIS

cc [*flag* ...] *file* ... **–lbsdmalloc** [*library* ...]

char ∗**malloc(** *size***)**
unsigned *size***;**

int free(*ptr***)**
char ∗ *ptr***;**

char ∗**realloc(** *ptr, size***)**
char ∗*ptr***;**
unsigned size**;**

DESCRIPTION

These routines provide a general-purpose memory allocation package. They maintain a table of free blocks for efficient allocation and coalescing of free storage. When there is no suitable space already free, the allocation routines call **sbrk**(2) to get more memory from the system. Each of the allocation routines returns a pointer to space suitably aligned for storage of any type of object. Each returns a NULL pointer if the request cannot be completed (see **DIAGNOSTICS**).

malloc( ) returns a pointer to a block of at least *size* bytes, which is appropriately aligned.

free( ) releases a previously allocated block. Its argument is a pointer to a block previously allocated by **malloc( )** or **realloc( )**.

realloc( ) changes the size of the block referenced by *ptr* to *size* bytes and returns a pointer to the (possibly moved) block. The contents will be unchanged up to the lesser of the new and old sizes. If unable to honor a reallocation request, **realloc( )** leaves its first argument unaltered. For backwards compatibility, **realloc( )** accepts a pointer to a block freed since the most recent call to **malloc( )** or **realloc( )**.

RETURN VALUES

malloc( ) and **realloc( )** return a **NULL** pointer if there is not enough available memory. When **realloc( )** returns **NULL**, the block pointed to by *ptr* is left intact.

ERRORS

If **malloc( )** or **realloc( )** returns unsuccessfully, **errno** will be set to indicate the following:

ENOMEM *size* bytes of memory exceeds the physical limits of your system, and cannot be allocated.

EAGAIN There is not enough memory available AT THIS POINT IN TIME to allocate *size* bytes of memory; but the application could try again later.

SEE ALSO

brk(2), **malloc**(3C), **malloc**(3X), **mapmalloc**(3X)

WARNINGS

Use of libbsdmalloc renders an application non-SCD compliant.

libbsdmalloc routines are incompatible with the memory allocation routines in the standard C-library (libc): **malloc**(3C), **alloca**(3C), **calloc**(3C), **free**(3C), **memalign**(3C), **realloc**(3C), and **valloc**(3C).

NOTES

Using **realloc( )** with a block freed before the most recent call to **malloc( )** or **realloc( )** will result in an error.

malloc( ) and **realloc( )** return a non-NULL pointer if *size* is 0. These pointers should not be dereferenced.

Always cast the value returned by **malloc( )** and **realloc( )**.

Comparative Features of **bsdmalloc( ), malloc**(3X), and **malloc**(3C)**:**

- The **bsdmalloc( )** routines afford better performance, but are space-inefficient.
- The **malloc**(3X) routines are space-efficient, but have slower performance.
- The standard, fully SCD-compliant **malloc**(3C) routines are a trade-off between performance and space-efficiency.

free( ) does not set **errno.**

NAME

bsearch – binary search a sorted table

SYNOPSIS

```
#include <stdlib.h>

void *bsearch(const void *key, const void *base, size_t nel, size_t size,
        int (*compar)(const void *, const void *));
```

MT-LEVEL

Safe

DESCRIPTION

bsearch( ) is a binary search routine generalized from Knuth (6.2.1) Algorithm B. It returns a pointer into a table (an array) indicating where a datum may be found or a null pointer if the datum cannot be found. The table must be previously sorted in increasing order according to a comparison function pointed to by *compar*. *key* points to a datum instance to be sought in the table. *base* points to the element at the base of the table. *nel* is the number of elements in the table. *size* is the number of bytes in each element. The function pointed to by *compar* is called with two arguments that point to the elements being compared. The function must return an integer less than, equal to, or greater than 0 as accordingly the first argument is to be considered less than, equal to, or greater than the second.

RETURN VALUES

A null pointer is returned if the key cannot be found in the table.

EXAMPLES

The example below searches a table containing pointers to nodes consisting of a string and its length. The table is ordered alphabetically on the string in the node pointed to by each entry.

This program reads in strings and either finds the corresponding node and prints out the string and its length, or prints an error message.

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>

struct node {                   /* these are stored in the table */
        char *string;
        int length;
};
static struct node table[] = {  /* table to be searched */
        { "asparagus", 10 },
        { "beans", 6 },
        { "tomato", 7 },
        { "watermelon", 11 },
};

main( )
{
        struct node *node_ptr, node;
```

```
        /* routine to compare 2 nodes */
        static int node_compare(const void *, const void *);
        char str_space[20];   /* space to read string into */

        node.string = str_space;
        while (scanf("%20s", node.string) != EOF) {
                node_ptr = bsearch( &node,
                        table, sizeof(table)/sizeof(struct node),
                        sizeof(struct node), node_compare);
                if (node_ptr != NULL) {
                        (void) printf("string = %20s, length = %d\n",
                                node_ptr->string, node_ptr->length);
                } else {
                        (void)printf("not found: %20s\n", node.string);
                }
        }
        return(0);
}

/* routine to compare two nodes based on an  */
/* alphabetical ordering of the string field */
static int
node_compare(const void *node1, const void *node2) {
        return (strcmp(
                        ((const struct node *)node1)->string,
                        ((const struct node *)node2)->string));
}
```

SEE ALSO

hsearch(3C), **lsearch**(3C), **qsort**(3C), **tsearch**(3C)

NOTES

The pointers to the key and the element at the base of the table should be of type pointer-to-*element*.

The comparison function need not compare every byte, so arbitrary data may be contained in the elements in addition to the values being compared.

If the number of elements in the table is less than the size reserved for the table, *nel* should be the lower number.

NAME

bstring, bcopy, bcmp, bzero – bit and byte string operations

SYNOPSIS

```
#include <strings.h>
void bcopy(const void *s1, void *s2, size_t n);
int bcmp(const void *s1, const void *s2, size_t n);
void bzero(void *s, size_t n);
```

DESCRIPTION

The functions **bcopy( )**, **bcmp( )**, and **bzero( )** operate on variable length strings of bytes. They do not check for null bytes as the routines in **string**(3C) do.

bcopy( ) copies *n* bytes from string *s1* to the string *s2*. Overlapping strings are handled correctly.

bcmp( ) compares byte string *s1* against byte string *s2*, returning zero if they are identical, 1 otherwise. Both strings are assumed to be *n* bytes long. **bcmp( )** using *n* zero bytes always returns zero.

bzero( ) places *n* 0 bytes in the string *s*.

WARNINGS

The **bcmp( )** and **bcopy( )** routines take parameters backwards from **strcmp** and **strcpy**, respectively. See **string**(3C).

SEE ALSO

memory(3C), **string**(3C)

NAME

bufsplit – split buffer into fields

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

size_t bufsplit(char ∗*buf***, size_t** *n***, char** ∗∗*a***);**

MT-LEVEL

MT-Safe

DESCRIPTION

bufsplit( ) examines the buffer, *buf*, and assigns values to the pointer array, *a*, so that the pointers point to the first *n* fields in *buf* that are delimited by tabs or new-lines.

To change the characters used to separate fields, call **bufsplit( )** with *buf* pointing to the string of characters, and *n* and *a* set to zero. For example, to use '**:**', '**.**', and '**,**' as separators along with tab and new-line:

```
bufsplit (":.,\t\n", 0, (char**)0 );
```

RETURN VALUES

The number of fields assigned in the array *a*. If *buf* is zero, the return value is zero and the array is unchanged. Otherwise the value is at least one. The remainder of the elements in the array are assigned the address of the null byte at the end of the buffer.

EXAMPLES

```
/*
 * set a[0] = "This", a[1] = "is", a[2] = "a",
 * a[3] = "test"
 */
bufsplit("This\tis\ta\ttest\n", 4, a);
```

NOTES

bufsplit( ) changes the delimiters to null bytes in *buf*.

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

byteorder, htonl, htons, ntohl, ntohs – convert values between host and network byte order

SYNOPSIS

```
#include <sys/types.h>
#include <netinet/in.h>

ulong htonl(u_long hostlong);

u_short htons(u_short hostshort);

u_long ntohl(u_long netlong);

u_short ntohs(u_short netshort);
```

MT-LEVEL

Safe

DESCRIPTION

These routines convert 16 and 32 bit quantities between network byte order and host byte order. On some architectures these routines are defined as NULL macros in the include file **<netinet/in.h>**. On other architectures, if their host byte order is different from network byte order, these routines are functional.

These routines are most often used in conjunction with Internet addresses and ports as returned by **gethostent( )** and **getservent( )**. (See **gethostbyname**(3N) and **getservbyname**(3N) respectively.)

SEE ALSO

gethostbyname(3N), **getservbyname**(3N)

NAME

cancellation, pthread_cancel, pthread_setcancelstate, pthread_setcanceltype, pthread_testcancel, pthread_cleanup_push, pthread_cleanup_pop – canceling execution of a thread

SYNOPSIS

```
#include <pthread.h>

int pthread_cancel(pthread_t target_thread);

int pthread_setcancelstate(int state, int *oldstate);
int pthread_setcanceltype(int type, int *oldtype);
void pthread_testcancel( );

void pthread_cleanup_push(void (*handler)(void *) void *arg);
void pthread_cleanup_pop(int execute);
```

MT-LEVEL

MT-Safe

DESCRIPTION

Thread cancellation enables a thread to terminate the execution of any other thread in the process. When the notice of cancellation is acted upon, the target thread (the thread being cancelled) is allowed to hold pending cancellation requests in several ways and to perform application-specific cleanup processing.

As a thread acquires resources around areas where it may get cancelled (i.e., before a cancellation point), it needs to push cancellation cleanup handlers along with the acquisition of these resources. The cleanup handlers release these resources and are invoked only if the thread were to be cancelled. As the thread leaves the last cancellation point before releasing a resource, it needs to pop the cleanup handler it had pushed earlier for this resource.

When a thread is cancelled, all the currently stacked cleanup handlers are executed and thread execution is terminated when the last cancellation cleanup handler returns. Its exit status of **PTHREAD_CANCELED** is then available to any threads joining with the cancelled target thread.

The thread's cancellation state and type determine when a thread could get cancelled.

State

PTHREAD_CANCEL_DISABLE
All cancellation requests to the *target_thread* are held pending.

PTHREAD_CANCEL_ENABLE
Cancellation requests are acted upon, depending upon the thread's cancellation type:

PTHREAD_CANCEL_ASYNCHRONOUS
If the cancellation state is enabled, new or pending cancellation requests may be acted upon at any time.

PTHREAD_CANCEL_DEFERRED
Cancellation requests are held pending until a cancellation point (see below) is reached.

Disabling cancellation will cause the setting of the cancellation type to be ineffective because all cancellation requests are held pending; however, when cancellation is enabled again, the new type will be in effect. The cancellation state is set to enabled, by default.

Type

When the cancellation state is disabled, a thread's cancellation *type* is meaningless. The following cancellation types behave as follows when enabled:

PTHREAD_CANCEL_ASYNCHRONOUS
Receipt of a **pthread_cancel( )** call will result in an immediate cancellation.

PTHREAD_CANCEL_DEFERRED
Cancellation will not occur until the target thread reaches a cancellation point (see below). Receipt of a **pthread_cancel( )** call will result in an immediate cancellation at this cancellation point.

The cancellation type is set to **PTHREAD_CANCEL_DEFERRED,** by default.

Cancellation Points

Cancellation begins at a point in a thread's execution when pending cancellation requests are tested and the cancellation state is found to be enabled. This is called the cancellation point.

A cancellation point can be explicitly set by inserting a call to the **pthread_testcancel( )** function.

In addition to explicit **pthread_testcancel( )** cancellation points, implicit cancellation points occur at a defined list of system entry points. Typically, any call that might require a long term wait should be a cancellation point. Operations need only check for pending cancellation requests when the operation is about to block indefinitely. This includes threads waiting in **pthread_cond_wait**(3T) and **pthread_cond_timedwait**(3T), threads waiting for the termination of another thread in **pthread_join**(3T), and threads blocked on **sigwait**(2).

POSIX has also defined several other functions (in **libc** and **libposix4**), as implicit cancellation points. In general, these are functions in which threads may block:

aio_suspend(3R), **close**(2), **creat**(2), **fcntl**(2), **fsync**(3C), **mq_receive**(3R), **mq_send**(3R), **msync**(3C), **nanosleep**(3R), **open**(2), **pause**(2), **pthread_cond_timedwait**(3T), **pthread_cond_wait**(3T), **pthread_join**(3T), **pthread_testcancel**, **read**(2), **sem_wait**(3R), **sigwaitinfo**(3R), **sigsuspend**(2), **sigtimedwait**(3R), **sigwait**(2), **sleep**(3C), **system**(3S), **tcdrain**(3), **wait**(2), **waitpid**(2), and **write**(2).

A cancellation point may also occur when a thread is executing the following functions:

closedir(3C), **ctermid**(3S), **fclose**(3S), **fcntl**(2), **fflush**(3S), **fgetc**(3S), **fgets**(3S), **fopen**(3S), **fprintf**(3S), **fputc**(3S), **fputs**(3S), **fread**(3S), **freopen**(3S), **fscanf**(3S), **fseek**(3S), **ftell**(3S), **fwrite**(3S), **getc**(3S), **getc_unlocked**(3S), **getchar**(3S), **getchar_unlocked**(3S), **getcwd**(3C), **getgrgid**(3C), **getgrgid_r**(3C), **getgrnam**(3C), **getgrnam_r**(3C), **getlogin**(3C), **getlogin_r**(3C), **getpwnam**(3C), **getpwnam_r**(3C), **getpwuid**(3C), **getpwuid_r**(3C), **gets**(3S), **lseek**(2), **rename**(2), **opendir**(3C), **perror**(3C), **printf**(3S), **putc**(3S), **putc_unlocked**(3S), **putchar**(3S), **putchar_unlocked**(3S), **puts**(3S), **readdir**(3C),

remove(3C), **rewind**(3S), **rewinddir**(3C), **scanf**(3S), **tmpfile**(3S), **ttyname**(3C), **ttyname_r**(3C), **ungetc**(3S), and **unlink**(2).

Cleanup Handling

An application should set up a cancellation cleanup handling function to restore any resources before a thread reaches a cancellation point. Specified cancellation points allow programmers to easily keep track of actions needed in a cancellation cleanup handler. A thread should only be made asynchronously cancelable when it is not in the process of acquiring or releasing resources (or locks), or otherwise, not in a difficult or impossible recover state.

When a cancellation request is acted upon, the routines in the list are invoked one-by-one in LIFO (last-in, first-out) order. When a scope's cancellation cleanup handler is invoked, the storage for that scope remains valid.

pthread_cancel

pthread_cancel( ) requests that *target_thread* be canceled. If the *target_thread*'s cancellation state is enabled, the **pthread_cancel( )** call will result in an immediate cancellation, if the target thread has the **PTHREAD_CANCEL_ASYNCHRONOUS** type set. Cancellation cleanup handlers for *target_thread* are called when the cancellation is acted on. Upon return of the last cancellation cleanup handler, the thread-specific data destructor functions are called for *target_thread*. *target_thread* is terminated when the last destructor function returns.

pthread_setcancelstate

pthread_setcancelstate( ) atomically sets the calling thread's cancellation state to the specified *state* and, if *oldstate* is not **NULL,** stores the previous cancellation state in *oldstate*.

A cancellation point occurs in the calling thread once the state is set if **pthread_setcancelstate( )** is called with **PTHREAD_CANCEL_ENABLE**, and *type* is **PTHREAD_CANCEL_ASYNCHRONOUS**.

pthread_setcanceltype

pthread_setcanceltype( ) atomically sets the calling thread's cancellation type to the specified *type* and, if *oldtype* is not **NULL,** stores the previous cancellation type in *oldtype*.

If **pthread_setcanceltype( )** is called with **PTHREAD_CANCEL_ASYNCHRONOUS**, and if *state* is **PTHREAD_CANCEL_ENABLE**, a cancellation point is set in the calling thread after *type* is specified.

Legal values for *state* are **PTHREAD_CANCEL_ENABLE** and **PTHREAD_CANCEL_DISABLE**. Legal values for *type* are **PTHREAD_CANCEL_DEFERRED** and **PTHREAD_CANCEL_ASYNCHRONOUS**. The cancellation state and type for newly created threads, including the thread in which **main( )** was first invoked, are **PTHREAD_CANCEL_ENABLE** and **PTHREAD_CANCEL_DEFERRED**, respectively.

pthread_testcancel

pthread_testcancel( ) creates a cancellation point in the calling thread; it has no effect if cancellation is disabled.

pthread_cleanup_push

pthread_cleanup_push( ) pushes the specified cancellation cleanup handler routine, *handler*, onto the cancellation cleanup stack of the calling thread. When the thread exits, implicitly or explicitly, or is cancelled, its cancellation cleanup handler is popped from the cancellation cleanup stack and invoked with the argument *arg*. The thread acts upon

a cancellation request, or the thread calls **pthread_cleanup_pop( )** with a non-zero *execute* argument.

pthread_cleanup_pop

pthread_cleanup_pop( ) removes the cleanup handler routine at the top of the cancellation cleanup stack of the calling thread and executes it if *execute* is non-zero.

If there are any calls to **pthread_cleanup_push( )** or **pthread_cleanup_pop( )** made without the matching call after the jump buffer is full, the effect of calling **longjmp**(3C) or **siglongjmp**(3C) is undefined.

Calls to **longjmp( )** or **siglongjmp( )** from within a cancellation cleanup handler is also undefined unless the jump buffer was also filled in the cancellation cleanup handler.

RETURN VALUES

If successful, **pthread_cancel( ) pthread_setcancelstate( )** and **pthread_setcanceltype( )** returns **0**; otherwise, an error number is returned.

pthread_testcancel( ), **pthread_cleanup_push( )**. and **pthread_cleanup_pop( )** are statements and do not return anything.

ERRORS

For each of the following conditions, **pthread_cancel( )** returns the corresponding error number if the condition is detected:

ESRCH No thread was found corresponding to that specified by the *target_thread* ID.

For each of the following conditions, **pthread_setcancelstate( )** returns the corresponding error if the condition is detected:

EINVAL The specified state is not **PTHREAD_CANCEL_ENABLE** or **PTHREAD_CANCEL_DISABLE**.

For each of the following conditions, **pthread_setcanceltype( )** returns the corresponding error if the condition is detected:

EINVAL The specified type is not **PTHREAD_CANCEL_DEFERRED** or **PTHREAD_CANCEL_ASYNCHRONOUS**.

SEE ALSO

condition(3T), **pthread_exit**(3T), **pthread_join**(3T), **setjmp**(3C)

NOTES

Please see **Intro**(3) for the notion of cancel-safety, Deferred-cancel-safety, and Asynchronous-cancel-safety. All libraries that have cancellation points but do not push / pop cancellation cleanup handlers are cancel-unsafe. If they push / pop cancellation handlers around cancellation points, they would become Deferred-cancel-safe, but could still be Asynchronous-cancel-unsafe.

In general, on Solaris, unless stated otherwise, all libraries are Asynchronous-cancel-unsafe and they may always remain so, because it may be too expensive for the common case (which is deferred cancellation) to make them Asynchronous-cancel-safe.

Libraries that do not have cancellation points are, by definition, Deferred-cancel-safe. Libraries that do have cancellation points but do not acquire any resources, such as locks or memory around these cancellation points, are also Deferred-cancel-safe. Those libraries which acquire locks and / or other resources before cancellation points are

Deferred-cancel-unsafe. Currently, there does not exist any labeling of libraries on Solaris about their cancel-safety status.

Applications can ensure cancel-safety of libraries by disabling cancellation before entering the library and restoring the old cancellation state on exit from the library.

Solaris threads do not offer this functionality.

Use of asynchronous cancellation while holding resources that need to be released may result in resource loss. Similarly, cancellation scopes may be safely manipulated (pushed and popped) only when the thread is in the deferred or disabled cancellation states.

For every **push( )** there must be the same number of **pop( )**s to compile the application.

EXAMPLES

The following is a short C++ example that shows the pushing / popping of cancellation handlers, the disabling / enabling of cancellation, the usage of **pthread_testcancel( )**, etc. The cancellation handler is "**free_res( ),**" which is a dummy function that simply prints a message in this example, but in a real application would actually free resources. The function "**f2( )**" is called from the main thread, and it goes deep into its call stack by calling itself recursively.

Before **f2( )** starts running, the newly created thread has most likely posted a cancellation on the main thread since the main thread calls **thr_yield( )** right after creating thread2. Since cancellation has been disabled in the main thread initially, via the call to **pthread_setcancelstate( )**, the call to **f2( )** from **main( )** proceeds fine with "**X**" being constructed at each recursive call, although the main thread has a pending cancellation. Now, when **f2( )** is called for the fifty-first time (i.e., when "i == 50"), **f2( )** enables cancellation by calling **pthread_setcancelstate( )** and then establishes a cancellation point for itself by calling **pthread_testcancel( )**.

Instead of **pthread_testcancel( )**, there could have been a call to a cancellation point such as **read**(2) or **write**(2), which would have a similar effect (i.e., cause the caller to get cancelled at this point since there is a pending cancellation). Hence, the **main( )** thread gets cancelled at the fifty-first iteration and then all the cleanup handlers that were pushed, are called in sequence; this is indicated by the calls to **free_res( )** and the calls to the destructor for "**X**". At each level, the C++ runtime calls the destructor for **X** and then the cancellation handler, **free_res( )**. The print messages from **free_res( )** and **X**'s destructor show the sequence of calls.

At the end, the main thread is joined by thread2, and since the main thread has been cancelled, its return status is **PTHREAD_CANCELED,** which is obtained from the **pthread_join( )**. This status is printed out and then thread2 returns, killing the process, since it is the last thread in the process.

```
#include <pthread.h>
#include <string.h>
extern "C" void thr_yield(void);

extern "C" void printf(...);

struct X {
```

```
        int x;
        X(int i){x = i; printf("X(%d) constructed.0, i);}
        ~X(){ printf("X(%d) destroyed.0, x);}
};

void
free_res(void *i)
{
        printf("Freeing `%d`0,i);
}

char* f2(int i)
{
   try {
        X dummy(i);
        pthread_cleanup_push(free_res, (void *)i);
        if (i == 50){
           pthread_setcancelstate(PTHREAD_CANCEL_ENABLE, NULL);
           pthread_testcancel();
        }
        f2(i+1);
        pthread_cleanup_pop(0);
   }
   catch (int) {
        printf("Error: In handler.0);
   }
   return "f2";
}

void *
thread2(void *tid)
{
        void *sts;

        printf("I am new thread :%d0, pthread_self());

        pthread_cancel((pthread_t)tid);

        pthread_join((pthread_t)tid, &sts);

        printf("main thread cancelled due to %d0, sts);

        return (sts);

}
```

```
main()
{
   pthread_setcancelstate(PTHREAD_CANCEL_DISABLE, NULL);
   pthread_create(NULL, NULL, thread2, (void *)pthread_self());
   thr_yield();
   printf("Returned from %s0, f2(0));
}
```

NAME

catgets – read a program message

SYNOPSIS

```
#include <nl_types.h>
char *catgets(nl_catd catd, int set_num, int msg_num, char *s);
```

MT-LEVEL

MT-Safe

DESCRIPTION

catgets( ) attempts to read message *msg_num*, in set *set_num*, from the message catalog identified by *catd*. *catd* is a catalog descriptor returned from an earlier call to **catopen( )**. *s* points to a default message string which will be returned by **catgets( )** if the identified message catalog is not currently available.

RETURN VALUES

If the identified message is retrieved successfully, **catgets( )** returns a pointer to an internal buffer area containing the null terminated message string. If the call is unsuccessful for any reason, **catgets( )** returns a pointer to *s*.

SEE ALSO

gencat(1), **catopen**(3C), **setlocale**(3C), **gettext**(3I)

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NOTES

catgets can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

catopen, catclose – open / close a message catalog

SYNOPSIS

```
#include <nl_types.h>
nl_catd catopen(char *name, int oflag);
int catclose(nl_catd catd);
```

MT-LEVEL

MT-Safe

DESCRIPTION

catopen( ) opens a message catalog and returns a message catalog descriptor. *name* specifies the name of the message catalog to be opened. If *name* contains a " / ", then *name* specifies a complete pathname for the message catalog; otherwise, the environment variable **NLSPATH** is used and **/usr/lib/locale/***locale***/LC_MESSAGES** must exist. If **NLSPATH** does not exist in the environment, or if a message catalog cannot be opened in any of the paths specified by **NLSPATH**, then the default path **/usr/lib/locale/***locale***/LC_MESSAGES** is used. In C locale, **catopen( )** will always succeed without even looking at the implementation defined default search path.

The names of message catalogs, and their location in the filesystem, can vary from one system to another. Individual applications can choose to name or locate message catalogs according to their own special needs. A mechanism is therefore required to specify where the catalog resides.

The **NLSPATH** variable provides both the location of message catalogs, in the form of a search path, and the naming conventions associated with message catalog files. For example:

NLSPATH=/nlslib/%L/%N.cat:/nlslib/%N/%L

The metacharacter **%** introduces a substitution field, where **%L** substitutes the current setting of either the **LANG** environment variable, if the value of *oflag* is **0**, or the **LC_MESSAGES** category, if the value of *oflag* is **NL_CAT_LOCALE**, and **%N** substitutes the value of the *name* parameter passed to **catopen( )**. Thus, in the above example, **catopen( )** will search in **/nlslib/$LANG/***name***.cat**, if *oflag* is **0**, or in **/nlslib/{LC_MESSAGES}/***name***.cat**, if *oflag* is **NL_CAT_LOCALE**.

NLSPATH will normally be set up on a system wide basis (in **/etc/profile**) and thus makes the location and naming conventions associated with message catalogs transparent to both programs and users.

The full set of metacharacters is:

%N The value of the name parameter passed to **catopen( )**.

%L The value of **LANG** or **LC_MESSAGES**.

%l The value of the *language* element of **LANG** or **LC_MESSAGES**.

%t The value of the *territory* element of **LANG** or **LC_MESSAGES**.

%c The value of the *codeset* element of **LANG** or **LC_MESSAGES**.

%% A single %.

The **LANG** environment variable provides the ability to specify the user's requirements for native languages, local customs and character set, as an ASCII string in the form

LANG=language[_territory[.codeset]]

A user who speaks German as it is spoken in Austria and has a terminal which operates in ISO 8859/1 codeset, would want the setting of the **LANG** variable to be

LANG=De_A.88591

With this setting it should be possible for that user to find any relevant catalogs should they exist.

Should the **LANG** variable not be set then the value of **LC_MESSAGES** as returned by **setlocale( )** is used. If this is **NULL** then the default path as defined in **nl_types( )** is used.

If the value of *oflag* argument is **0**, the **LANG** environment variable is used to locate the catalogue without regard to the **LC_MESSAGES** category. If the *oflag* argument is **NL_CAT_LOCALE**, the **LC_MESSAGES** category is used to locate the message catalogue.

catclose( ) closes the message catalog identified by *catd*.

RETURN VALUES

If successful, **catopen( )** returns a message catalog descriptor for use on subsequent calls to **catgets( )** and **catclose( )**; otherwise **catopen( )** returns **(nl_catd) –1**.

catclose( ) returns **0** if successful, otherwise **–1**.

SEE ALSO

gencat(1), **gettext**(3I), **catgets**(3C), **setlocale**(3C), **environ**(5), **nl_types**(5)

NOTES

catopen and **catclose** can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

clock – report CPU time used

SYNOPSIS

```
#include <time.h>
clock_t clock(void);
```

MT-LEVEL

MT-Safe

DESCRIPTION

clock( ) returns the amount of CPU time (in microseconds) used since the first call to **clock( )** in the calling process. The time reported is the sum of the user and system times of the calling process and its terminated child processes for which it has executed the **wait( )** function, the **pclose( )** function, or the **system( )** function.

Dividing the value returned by **clock( )** by the constant **CLOCKS_PER_SEC**, defined in the <**time.h**> header, will give the time in seconds.

SEE ALSO

times(2), **wait**(2), **popen**(3S), **system**(3S)

NOTES

The value returned by **clock( )** is defined in microseconds for compatibility with systems that have CPU clocks with much higher resolution. Because of this, the value returned will wrap around after accumulating only 2147 seconds of CPU time (about 36 minutes). If the process time used is not available or cannot be represented, clock returns the value **(clock_t) –1**.

NAME

clock_settime, clock_gettime, clock_getres – high-resolution clock operations

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <time.h>
int clock_settime(clockid_t clock_id, const struct timespec *tp);
int clock_gettime(clockid_t clock_id, struct timespec *tp);
int clock_getres(clockid_t clock_id, struct timespec *res);
struct  timespec {
        time_t          tv_sec;         /* seconds */
        long            tv_nsec;        /* and nanoseconds */
};
```

MT-LEVEL

clock_gettime( ) is Async-Signal-Safe

DESCRIPTION

clock_settime( ) sets the specified clock, *clock_id,* to the value specified by *tp.* The calling process must have an effective user ID of **0.**

clock_gettime( ) returns the current value *tp* for the specified clock, *clock_id.*

The resolution of any clock can be obtained by calling **clock_getres( ).** If *res* is not NULL, the resolution of the specified clock is stored in *res.*

The *clock_id* for the real-time clock for the system is **CLOCK_REALTIME**. The values returned by **clock_gettime( )** and specified by **clock_settime( )** represent the amount of time (in seconds and nanoseconds) since 00:00 Universal Coordinated Time, January 1, 1970.

RETURN VALUES

clock_settime( ), clock_gettime( ), and **clock_getres( )** return **0** upon success, otherwise they return **-1** and set **errno** to indicate the error condition.

ERRORS

EINVAL *clock_id* does not specify a known clock.

The *tp* argument to **clock_settime( )** is outside the range for the given clock id.

The *tp* argument to **clock_settime( )** specified a nanosecond value less than zero or greater than or equal to 1,000,000,000.

ENOSYS **clock_settime( ), clock_gettime( ),** or **clock_getres( )** is not supported by this implementation.

EPERM The requesting process does not have the appropriate privilege to set the specified clock.

SEE ALSO

time(2), **ctime**(3C), **timer_gettime**(3R)

NOTES

Clock resolutions are implementation defined and are not settable by a process. Time values that are between two consecutive non-negative integer multiples of the resolution of the specified clock are truncated down to the smaller multiple of the resolution.

NAME

condition, pthread_cond_init, pthread_cond_wait, pthread_cond_timedwait, pthread_cond_signal, pthread_cond_broadcast, pthread_cond_destroy, cond_init, cond_wait, cond_timedwait, cond_signal, cond_broadcast, cond_destroy – condition variables

SYNOPSIS

POSIX

cc [*flag* ...] *file* ... **–lpthread** [*library* ...]

```
#include <pthread.h>
int pthread_cond_init(pthread_cond_t *cond, const pthread_condattr_t *attr);
int pthread_cond_wait(pthread_cond_t *cond, pthread_mutex_t *mutex);
int pthread_cond_timedwait(pthread_cond_t *cond, pthread_mutex_t *mutex,
    const struct timespec *abstime);
int pthread_cond_signal(pthread_cond_t *cond);
int pthread_cond_broadcast(pthread_cond_t *cond);
int pthread_cond_destroy(pthread_cond_t *cond);
```

Solaris

cc [*flag* ...] *file* ... **–lthread** [*library* ...]

```
#include <thread.h>
#include <synch.h>
int cond_init(cond_t *cvp, int type, int arg);
int cond_wait(cond_t *cvp, mutex_t *mp);
int cond_timedwait(cond_t *cvp, mutex_t *mp, timestruc_t *abstime);
int cond_signal(cond_t *cvp);
int cond_broadcast(cond_t *cvp);
int cond_destroy(cond_t *cvp);
```

MT-LEVEL

MT-Safe

DESCRIPTION

Occasionally, a thread running within a mutex needs to wait for an event, in which case, it blocks or sleeps. When a thread is waiting for another thread to communicate its disposition, it uses a condition variable in conjunction with a mutex. Although a mutex is exclusive and the code it protects is sharable (at certain moments), condition variables enable the synchronization of differing events that share a mutex, but not necessarily data. Several condition variables may be used by threads to signal each other when a task is complete, which then allows the next waiting thread to take ownership of the mutex.

A condition variable enables threads to atomically block and test the condition under the protection of a mutual exclusion lock (mutex) until the condition is satisfied. If the condition is false, a thread blocks on a condition variable and atomically releases the mutex that is waiting for the condition to change. If another thread changes the condition, it may wake up waiting threads by signaling the associated condition variable. The waiting threads, upon awakening, reacquire the mutex and re-evaluate the condition.

Initialize

Condition variables and mutexes should be global. Condition variables that are allocated in writable memory can synchronize threads among processes if they are shared by the cooperating processes (see **mmap**(2)) and are initialized for this purpose.

The scope of a condition variable is either intra-process or inter-process. This is dependent upon whether the argument is passed implicitly or explicitly to the initialization of that condition variable. A condition variable does not need to be explicitly initialized. A condition variable is initialized with all zeros, by default, and its scope is set to within the calling process. For inter-process synchronization, a condition variable must be initialized once, and only once, before use.

A condition variable must not be simultaneously initialized by multiple threads or reinitialized while in use by other threads.

Condition variables attributes may be set to the default or customized at initialization. POSIX threads even allow the default values to be customized. Establishing these attributes varies depending upon whether POSIX or Solaris threads are used. Similar to the distinctions between POSIX and Solaris thread creation, POSIX condition variables implement the default, intra-process, unless an attribute object is modified for inter-process prior to the initialization of the condition variable. Solaris condition variables also implement as the default, intra-process; however, they set this attribute according to the argument, *type*, passed to their initialization function.

POSIX Initialize

POSIX condition variables mutexes, and threads use attributes objects in the same manner; they are initialized with the configuration of an attributes object (see **pthread_condattr_init**(3T)). The **pthread_cond_init( )** function initializes the condition variable referenced by *cond* with attributes referenced by *attr*. If *attr* is **NULL**, the default condition variable attributes are used, which is the same as passing the address of a default condition variable attributes object. When the initialization is complete, the state of the condition variable is then initialized. If a default condition variable is used, then only threads created within the same process can operate on the initialized condition variable.

A condition variable can possess two different types of shared-scope behavior, which is determined by the second argument to **pthread_condattr_setpshared**(3T). This argument can be set to either of the following:

PTHREAD_PROCESS_PRIVATE The condition variable can synchronize threads only in this process. The **PTHREAD_PROCESS_PRIVATE** POSIX setting for process scope is equivalent to the **USYNC_THREAD** flag to **cond_init( )** in the Solaris API. This is the default.

PTHREAD_PROCESS_SHARED The condition variable can synchronize threads in this process and other processes. Only one process should initialize the condition variable. The **PTHREAD_PROCESS_SHARED** POSIX setting for system-wide scope is equivalent to the **USYNC_PROCESS** flag to **cond_init( )** in the Solaris API.

Initializing condition variables can also be accomplished by allocating-in zeroed memory (default), in which case, **PTHREAD_PROCESS_PRIVATE** is assumed. The same condition variable must not be simultaneously initialized by multiple threads nor re-initialized while in use by other threads.

If default condition variable attributes are used, statically allocated condition variables can be initialized by the macro **PTHREAD_COND_INITIALIZER.** The effect is the same as a dynamic initialization by a call to **pthread_cond_init( )** with parameter *attr* specified as **NULL**, except error checks are not performed.

Default condition variable initialization (intra-process):

```
pthread_cond_t      cvp;
pthread_condattr_t  cv_attr;

pthread_cond_init(&cvp, NULL); /* initialize cv with defaults */
  OR
pthread_condattr_init(&cv_attr); /* initialize cv_attr with defaults */
pthread_cond_init(&cvp, &cv_attr);      /* initialize cv with default cv_attr */
  OR
pthread_condattr_setpshared(&cv_attr, PTHREAD_PROCESS_PRIVATE);
pthread_cond_init(&cvp, &cv_attr);      /* initialize cv with defaults */
  OR
pthread_cond_t  cond  = PTHREAD_COND_INITIALIZER;
  OR
pthread_cond_t cond;
cond = calloc(1, sizeof (pthread_cond_t));
```

Customized condition variable initialization (inter-process):

```
pthread_condattr_init(&cv_attr); /* initialize cv_attr with defaults */
pthread_condattr_setpshared(&cv_attr, PTHREAD_PROCESS_SHARED);
pthread_cond_init(&cvp, &cv_attr);      /* initialize cv with inter-process scope */
```

Solaris Initialize

cond_init( ) initializes the condition variable pointed to by *cvp*. A condition variable can have several different types of behavior, specified by *type*. No current type uses *arg* although a future type may specify additional behavior parameters via *arg*. *type* may be one of the following:

USYNC_THREAD The condition variable can synchronize threads only in this process. The **USYNC_THREAD** Solaris condition variable type for process scope is equivalent to the POSIX condition variable attribute setting **PTHREAD_PROCESS_PRIVATE.** *arg* is ignored.

USYNC_PROCESS The condition variable can synchronize threads in this process and other processes. Only one process should initialize the condition variable. The **USYNC_PROCESS** Solaris condition variable type for system-wide scope is equivalent to the POSIX condition variable

attribute setting **PTHREAD_PROCESS_SHARED.** *arg* is ignored.

Initializing condition variables can also be accomplished by allocating in zeroed memory, in which case, a *type* of **USYNC_THREAD** is assumed.

If default condition variable attributes are used, statically allocated condition variables can be initialized by the macro **DEFAULTCV.**

Default condition variable initialization (intra-process):

```
cond_t    cvp;

cond_init(&cvp, NULL, NULL); /* initialize condition variable with default */
 OR
cond_init(&cvp, USYNC_THREAD, NULL);
 OR
cond_t cond  = DEFAULTCV;
```

Customized condition variable initialization (inter-process):

```
cond_init(&cvp, USYNC_PROCESS, NULL); /* initialize cv with inter-process scope */
```

Condition Wait

The condition wait interface allows a thread to wait for a condition and atomically release the associated mutex that it needs to hold to check the condition. The thread waits for another thread to make the condition true and that thread's resulting call to signal and wakeup the waiting thread.

POSIX Wait

pthread_cond_wait( ) and **pthread_cond_timedwait( )** block on a condition variable, which atomically release the mutex pointed to by *mp* and cause the calling thread to block on the condition variable pointed to by *cond*. The blocked thread may be awakened by **pthread_cond_signal( )**, **pthread_cond_broadcast( )**, or interrupted by a UNIX signal.

These functions atomically release the *mutex*, causing the calling thread to block on the condition variable *cond*.

Upon successful completion, the mutex is locked and owned by the calling thread.

pthread_cond_timedwait( ) is the same as **pthread_cond_wait( )**, except an error is returned if the system time equals or exceeds the time specified by *abstime* before the condition *cond* is signaled or broadcasted, or if the absolute time specified by *abstime* has already passed at the time of the call. When timeouts occur, **pthread_cond_timedwait( )** releases and reacquires the mutex referenced by *mutex*.

When using condition variables, there is always a boolean predicate involving shared variables related to each condition wait that is true, if the thread should proceed. Since the return from **pthread_cond_wait( )** or **pthread_cond_timedwait( )** does not indicate anything about the value of this predicate, the predicate should be reevaluated on return. Unwanted wakeups from **pthread_cond_wait( )** or **pthread_cond_timedwait( )** may occur.

The functions **pthread_cond_wait( )** and **pthread_cond_timedwait( )** are cancellation points. If a cancellation request is acted upon while in a condition wait when the cancellation enable state of a thread is set to **PTHREAD_CANCEL_DEFERRED**, the mutex will be reacquired before calling the first cancellation cleanup handler. In other words, the thread is unblocked, allowed to execute up to the point of returning from the call to **pthread_cond_wait( )** or **pthread_cond_timedwait( )**, but then notices the cancellation request and, instead of returning to the caller of **pthread_cond_wait( )** or **pthread_cond_timedwait( )**, it starts the thread cancellation activities including cancellation cleanup handlers.

A thread that is unblocked because it was canceled while blocked in a call to **pthread_cond_wait( )** or **pthread_cond_timedwait( )** does not awaken anyone else asleep on the condition.

Solaris Wait

cond_wait( ) atomically releases the mutex pointed to by *mp* and causes the calling thread to block on the condition variable pointed to by *cvp*. The blocked thread may be awakened by **cond_signal( )**, **cond_broadcast( )**, or when interrupted by delivery of a UNIX signal or a **fork( )**.

cond_wait( ) and **cond_timedwait( )** always return with the mutex locked and owned by the calling thread even when returning an error.

Condition Signaling

A condition signal allows a thread to unblock the next thread waiting on the condition variable, whereas, a condition broadcast allows a thread to unblock all threads waiting on the condition variable.

POSIX Signal and Broadcast

pthread_cond_signal( ) and **pthread_cond_broadcast( )** unblock threads blocked on a condition variable.

pthread_cond_signal( ) unblocks at least one thread blocked on the specified condition variable *cond*, if any threads are blocked on *cond*.

pthread_cond_broadcast( ) unblocks all threads blocked on the condition variable *cond*.

pthread_cond_signal( ) and **pthread_cond_broadcast( )** have no effect if there are no threads blocked on *cond*.

pthread_cond_signal( ) or **pthread_cond_broadcast( )** may be called by a thread regardless of whether it owns the mutex which threads calling **pthread_cond_wait( )** or **pthread_cond_timedwait( )** have associated with the condition variable during their waits. However, if predictable scheduling behavior is required, then that mutex should be locked by the thread calling **pthread_cond_signal( )** or **pthread_cond_broadcast( )**.

Solaris Signal and Broadcast

cond_signal( ) unblocks one thread that is blocked on the condition variable pointed to by *cvp*.

cond_broadcast( ) unblocks all threads that are blocked on the condition variable pointed to by *cvp*.

If no threads are blocked on the condition variable, then **cond_signal( )** and **cond_broadcast( )** have no effect.

Both functions should be called under the protection of the same mutex that is used with the condition variable being signaled. Otherwise, the condition variable may be signaled between the test of the associated condition and blocking in **cond_wait( )**. This can cause an infinite wait.

Destroy

The condition destroy functions destroy any state, but not the space, associated with the condition variable.

POSIX Destroy

pthread_cond_destroy( ) destroys the condition variable specified by *cond*. The space for destroying the condition variable is not freed.

Solaris Destroy

cond_destroy( ) destroys any state associated with the condition variable pointed to by *cvp*. The space for storing the condition variable is not freed.

RETURN VALUES

0 is returned when any of these functions are successful. A non-zero value indicates an error, except **pthread_timedwait( )**, which returns **ETIME**.

ERRORS

These functions fail and return the corresponding value if any of the following conditions are detected:

EFAULT *cond*, *attr*, *cvp*, *arg*, *abstime*, or *mutex* point to an illegal address.

EINVAL Invalid argument.
For **pthread_cond_init( )**, the value specified for *attr* is invalid.

For **cond_init( )**, *type* is not a recognized type.

For **pthread_cond_timedwait( )** or **cond_timedwait( )**, the specified number of seconds, *abstime*, is greater than pgm_start_time + 50,000,000, where pgm_start_time is the start time of the application, or the number of nanoseconds is greater than or equal to 1,000,000,000.

cond_wait( ) or **cond_timedwait( )** fails and returns the corresponding value if any of the following conditions are detected:

EINTR The wait was interrupted by a signal or **fork( )**.

pthread_cond_timedwait( ) or **cond_timedwait( )** fails and returns the corresponding value if any of the following conditions are detected:

ETIME The time specified by *abstime* has passed.

SEE ALSO

mmap(2), **fork**(2), **signal**(3C), **mutex**(3T), **pthread_condattr_init**(3T)

NOTES

The only policy currently supported is **SCHED_OTHER**. In Solaris, under the **SCHED_OTHER** policy, there is no established order in which threads are unblocked.

If more than one thread is blocked on a condition variable, the order in which threads are unblocked is determined by the scheduling policy. When each thread, unblocked as a result of a **pthread_cond_signal( )** or **pthread_cond_broadcast( )**, returns from its call to

pthread_cond_wait( ) or **pthread_cond_timedwait( )**, the thread owns the mutex with which it called **pthread_cond_wait( )** or **pthread_cond_timedwait( )**. The thread(s) that are unblocked compete for the mutex according to the scheduling policy, and as if each had called **pthread_mutex_lock**(3T).

When **cond_wait( )** returns the value of the condition is indeterminate and must be reevaluated.

cond_timedwait( ) is similar to **cond_wait( )**, except that the calling thread will not wait for the condition to become true past the absolute time specified by *abstime*. Note that **cond_timedwait( )** may continue to block as it trys to reacquire the mutex pointed to by *mp*, which may be locked by another thread. If *abstime* then **cond_timedwait( )** returns because of a timeout, it returns the error code **ETIME .**

NAME

confstr – get configurable variables

SYNOPSIS

```
#include <unistd.h>
size_t confstr(int name, char *buf, size_t len);
```

MT-LEVEL

Mt-Safe

DESCRIPTION

The **confstr( )** function provides a method for applications to get configuration-defined string values. Its use and purpose are similar to the **sysconf**(3C) function, but it is used where string values rather than numeric values are returned.

The *name* argument represents the system variable to be queried. The implementation supports the *name* value of **_CS_PATH**, defined in **<unistd.h>**. It may support others.

If *len* is not **0**, and if *name* has a configuration-defined value, **confstr( )** copies that value into the *len*-byte buffer pointed to by *buf*. If the string to be returned is longer than *len* bytes, including the terminating null, then **confstr( )** truncates the string to *len*–1 bytes and null-terminates the result. The application can detect that the string was truncated by comparing the value returned by **confstr( )** with *len*.

If *len* is **0**, and *buf* is a null pointer, then **confstr( )** still returns the integer value as defined below, but does not return the string. If *len* is **0** but *buf* is not a null pointer, the result is unspecified.

RETURN VALUES

If *name* has a configuration-defined value, the **confstr( )** function returns the size of buffer that would be needed to hold the entire configuration-defined value. If this return value is greater than *len*, the string returned in *buf* is truncated.

If *name* is invalid, **confstr( )** returns **0** and sets **errno** to indicate the error.

If *name* does not have a configuration-defined value, **confstr( )** returns **0** and leaves **errno** unchanged.

ERRORS

The **confstr( )** funciton will fail if:

EINVAL The value of the *name* argument is invalid.

SEE ALSO

pathconf(2), **sysconf**(3C)

NAME

connect – initiate a connection on a socket

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>

int connect(int s, struct sockaddr *name, int namelen);
```

MT-LEVEL

Safe

DESCRIPTION

The parameter *s* is a socket. If it is of type **SOCK_DGRAM**, **connect( )** specifies the peer with which the socket is to be associated; this address is the address to which datagrams are to be sent if a receiver is not explicitly designated; it is the only address from which datagrams are to be received. If the socket *s* is of type **SOCK_STREAM**, **connect( )** attempts to make a connection to another socket. The other socket is specified by *name*. *name* is an address in the communication space of the socket. Each communication space interprets the *name* parameter in its own way. If *s* is not bound, then it will be bound to an address selected by the underlying transport provider. Generally, stream sockets may successfully **connect( )** only once; datagram sockets may use **connect( )** multiple times to change their association. Datagram sockets may dissolve the association by connecting to a null address.

RETURN VALUES

If the connection or binding succeeds, **0** is returned. Otherwise, **–1** is returned and sets **errno** to indicate the error.

ERRORS

The call fails if:

EACCES	Search permission is denied for a component of the path prefix of the pathname in *name*.
EADDRINUSE	The address is already in use.
EADDRNOTAVAIL	The specified address is not available on the remote machine.
EAFNOSUPPORT	Addresses in the specified address family cannot be used with this socket.
EALREADY	The socket is non-blocking and a previous connection attempt has not yet been completed.
EBADF	*s* is not a valid descriptor.
ECONNREFUSED	The attempt to connect was forcefully rejected. The calling program should **close**(2) the socket descriptor, and issue another **socket**(3N) call to obtain a new descriptor before attempting another **connect( )** call.

EINPROGRESS	The socket is non-blocking and the connection cannot be completed immediately. It is possible to **select**(3C) for completion by selecting the socket for writing. However, this is only possible if the socket STREAMS module is the topmost module on the protocol stack with a write service procedure. This will be the normal case.
EINTR	The connection attempt was interrupted before any data arrived by the delivery of a signal.
EINVAL	*namelen* is not the size of a valid address for the specified address family.
EIO	An I/O error occurred while reading from or writing to the file system.
EISCONN	The socket is already connected.
ELOOP	Too many symbolic links were encountered in translating the pathname in *name*.
ENETUNREACH	The network is not reachable from this host.
ENOENT	A component of the path prefix of the pathname in *name* does not exist.
ENOENT	The socket referred to by the pathname in *name* does not exist.
ENOSR	There were insufficient **STREAMS** resources available to complete the operation.
ENXIO	The server exited before the connection was complete.
ETIMEDOUT	Connection establishment timed out without establishing a connection.

The following errors are specific to connecting names in the UNIX domain. These errors may not apply in future versions of the UNIX IPC domain.

ENOTDIR	A component of the path prefix of the pathname in *name* is not a directory.
ENOTSOCK	*s* is not a socket.
ENOTSOCK	*name* is not a socket.
EPROTOTYPE	The file referred to by *name* is a socket of a type other than type *s* (for example, *s* is a **SOCK_DGRAM** socket, while *name* refers to a **SOCK_STREAM** socket).

SEE ALSO **close**(2), **accept**(3N), **getsockname**(3N), **select**(3C), **socket**(3N)

NAME

conv, toupper, tolower, _toupper, _tolower, toascii – translate characters

SYNOPSIS

```
#include <ctype.h>
int toupper(int c);
int tolower(int c);
int _toupper(int c);
int _tolower(int c);
int toascii(int c);
```

MT-LEVEL

MT-Safe with exceptions

DESCRIPTION

toupper( ) and **tolower( )** have as their domain the range of the function **getc( )**: all values represented in an **unsigned char** and the value of the macro **EOF** as defined in **stdio.h**. If the argument of **toupper( )** represents a lower-case letter, the result is the corresponding upper-case letter. If the argument of **tolower( )** represents an upper-case letter, the result is the corresponding lower-case letter. All other arguments in the domain are returned unchanged.

The macros **_toupper( )** and **_tolower( )** accomplish the same things as **toupper( )** and **tolower( )**, respectively, but have restricted domains and are faster. **_toupper( )** requires a lower-case letter as its argument; its result is the corresponding upper-case letter. **_tolower( )** requires an upper-case letter as its argument; its result is the corresponding lower-case letter. Arguments outside the domain cause undefined results.

toascii( ) yields its argument with all bits turned off that are not part of a standard 7-bit ASCII character; it is intended for compatibility with other systems.

toupper( ), **tolower( )**, **_toupper( )**, and **_tolower( )** are affected by **LC_CTYPE**. In the "C" locale, or in a locale where shift information is not defined, these functions determine the case of characters according to the rules of the ASCII-coded character set. Characters outside the ASCII range of characters are returned unchanged.

SEE ALSO

ctype(3C), **setlocale**(3C), **getc**(3S), **environ**(5)

NOTES

toupper, **tolower**, **_toupper**, **_tolower** and **toascii** can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

copylist – copy a file into memory

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

char ∗**copylist(const char** ∗*filenm***, off_t** ∗*szptr***);**

MT-LEVEL

MT-Safe

DESCRIPTION

copylist( ) copies a list of items from a file into freshly allocated memory, replacing newlines with null characters. It expects two arguments: a pointer *filenm* to the name of the file to be copied, and a pointer *szptr* to a variable where the size of the file will be stored.

Upon success, **copylist( )** returns a pointer to the memory allocated. Otherwise it returns NULL if it has trouble finding the file, calling **malloc( )**, or reading the file.

EXAMPLES

```
/* read "file" into buf */
off_t size;
char *buf;
buf = copylist("file", &size);
if (buf) {
        for (i=0; i<size; i++)
                if (buf[i])
                        putchar(buf[i]);
                else
                        putchar{'\n');
        }
} else {
        fprintf(stderr, "%s: Copy failed for "file".\n", argv[0]);
        exit (1);
}
```

SEE ALSO

malloc(3C)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

crypt, setkey, encrypt – generate encryption

SYNOPSIS

```
#include <crypt.h>
char *crypt(const char *key, const char *salt);
void setkey(const char *key);
void encrypt(char *block, int edflag);
```

MT-LEVEL

Safe

DESCRIPTION

crypt( ) is the password encryption function. It is based on a one-way encryption algorithm with variations intended (among other things) to frustrate use of hardware implementations of a key search.

key is the input string to encrypt, for instance, a user's typed password. Only the first eight characters are used; the rest are ignored. *salt* is a two-character string chosen from the set [**a-zA-Z0-9./**]; this string is used to perturb the hashing algorithm in one of 4096 different ways, after which the input string is used as the key to repeatedly encrypt a constant string. The returned value points to the encrypted input string. The first two characters of the return value are the *salt* itself.

The **setkey( )** and **encrypt( )** functions provide (rather primitive) access to the actual hashing algorithm. The argument of **setkey( )** is a character array of length 64 containing only the characters with numerical value 0 and 1. This string is divided into groups of 8, the low-order bit in each group is ignored; this gives a 56-bit key that is set into the machine. This is the key that will be used with the hashing algorithm to encrypt the string *block* with the **encrypt( )** function.

The *block* argument of **encrypt( )** is a character array of length 64 containing only the characters with numerical value 0 and 1. The argument array is modified in place to a similar array representing the bits of the argument after having been subjected to the hashing algorithm using the key set by **setkey( )**. The argument *edflag*, indicating decryption rather than encryption, is ignored; use **encrypt( )** in **libcrypt( )** (see **crypt**(3X)) for decryption.

RETURN VALUES

If *edflag* is set to anything other than zero, **errno** will be set to **ENOSYS**.

SEE ALSO

login(1), **passwd**(1), **crypt**(3X), **getpass**(3C), **passwd**(4)

NOTES

The return value for **crypt( )** points to static data that are overwritten by each call. In the case of multithreaded applications, the return value is a pointer to thread specific data.

NAME

cset, csetlen, csetcol, csetno, wcsetno – get information on EUC codesets

SYNOPSIS

```
#include <euc.h>
int csetlen(int codeset);
int csetcol(int codeset);
int csetno(unsigned char c);
#include <widec.h>
int wcsetno(wchar_t pc);
```

MT-LEVEL

MT-Safe with exceptions

DESCRIPTION

Both **csetlen( )** and **csetcol( )** take a code set number *codeset*, which must be 0, 1, 2, or 3. **csetlen( )** returns the number of bytes needed to represent a character of the given Extended Unix Code (EUC) code set, excluding the single-shift characters SS2 and SS3 for codesets 2 and 3. **csetcol( )** returns the number of columns a character in the given EUC code set would take on the display.

csetno( ) is a macro that returns a codeset number (0, 1, 2, or 3) for the EUC character whose first byte is *c*. For example,

```
#include<euc.h>
...
x+=csetcol(csetno(c));
```

increments a counter ''x'' (such as the cursor position) by the width of the character whose first byte is *c*.

wcsetno( ) is a macro that returns a codeset number (0, 1, 2, or 3) for the given process code character *pc*. For example,

```
#include<euc.h>
#include<widec.h>
...
x+=csetcol(wcsetno(pc));
```

increments a counter ''x'' (such as the cursor position) by the width of the Process Code character *pc*.

SEE ALSO

setlocale(3C)**, euclen**(3I)

NOTES

cset, **csetlen**, **csetcol**, **csetno** and **wcsetno** can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

ctermid, ctermid_r – generate path name for controlling terminal

SYNOPSIS

```
#include <stdio.h>
char *ctermid(char *s);
char *ctermid_r(char *s);
```

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

ctermid( ) generates the path name of the controlling terminal for the current process, and stores it in a string.

If *s* is a **NULL** pointer, the string is stored in an internal static area, the contents of which are overwritten at the next call to **ctermid( )**, and the address of which is returned. Otherwise, *s* is assumed to point to a character array of at least **L_ctermid** elements; the path name is placed in this array and the value of *s* is returned. The constant **L_ctermid** is defined in the header **<stdio.h>**.

ctermid_r( ) has the same functionality as **ctermid( )** except that if *s* is a **NULL** pointer, the function returns **NULL.**

SEE ALSO

ttyname(3C)

NOTES

The **ctermid_r( )** interface is as proposed in the POSIX.4a Draft #6 document, and is subject to change to be compliant to the standard when it is accepted.

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

The difference between **ctermid( )** and **ttyname**(3C) is that **ttyname( )** must be handed a file descriptor and returns the actual name of the terminal associated with that file descriptor, while **ctermid( )** returns a string (**/dev/tty**) that will refer to the terminal if used as a file name. Thus **ttyname( )** is useful only if the process already has at least one file open to a terminal.

ctermid( ) is unsafe in multithreaded applications. **ctermid_r( )** is MT-Safe, and should be used instead.

NAME

ctime, ctime_r, localtime, localtime_r, gmtime, gmtime_r, asctime, asctime_r, tzset, tzsetwall – convert date and time to string

SYNOPSIS

```
#include <time.h>
char *ctime(const time_t *clock);
char *ctime_r(const time_t *clock, char *buf, int buflen);
struct tm *localtime(const time_t *clock);
struct tm *localtime_r(const time_t *clock, struct tm *res);
struct tm *gmtime(const time_t *clock);
struct tm *gmtime_r(const time_t *clock, struct tm *res);
char *asctime(const struct tm *tm);
char *asctime_r(const struct tm *tm,char *buf, int buflen);
extern time_t timezone, altzone;
extern int daylight;
extern char *tzname[2];
void tzset(void);
void tzsetwall(void);
```

POSIX

```
cc [ flag... ] file ... –D_POSIX_PTHREAD_SEMANTICS [ library... ]
char *ctime_r(const time_t *clock, char *buf,);
char *asctime_r(const struct tm *tm,char *buf,);
```

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

ctime( ), **localtime( )**, and **gmtime( )** accept arguments of type **time_t**, pointed to by **clock( )**, representing the time in seconds since 00:00:00 UTC, January 1, 1970. **ctime( )** returns a pointer to a 26-character string as shown below. Time zone and daylight savings corrections are made before string generation. The fields are constant width:

```
Fri Sep 13 00:00:00 1986\n\0
```

ctime_r( ) has the same functionality as **ctime( )** except that the caller must supply a buffer *buf* with length *buflen* to store the result; *buf* must be at least 26 bytes. The POSIX **ctime_r( )** routine does not take a *buflen* parameter.

localtime( ) and **gmtime( )** return pointers to **tm** structures (see below). **localtime( )** corrects for the main time zone and possible alternate (''daylight savings'') time zone; **gmtime( )** converts directly to Coordinated Universal Time (UTC), which is what the UNIX system uses internally.

localtime_r( ) and **gmtime_r( )** have the same functionality as **localtime( )** and **gmtime( )** respectively, except that the caller must supply a buffer *res* to store the result.

asctime( ) converts a **tm** structure to a 26-character string, as shown in the above example, and returns a pointer to the string.

asctime_r( ) has the same functionality as **asctime( )** except that the caller must supply a buffer *buf* with length *buflen* for the result to be stored. *buf* must be at least 26 bytes. The POSIX **asctime_r( )** routine does not take a *buflen* parameter. **asctime_r( )** returns a pointer to *buf* upon success. In case of failure, **NULL** is returned and **errno** is set.

Declarations of all the functions and externals, and the **tm** structure, are in the **time.h** header. The members of the **tm** structure are:

```
int     tm_sec;         /* seconds after the minute — [0, 61] */
                        /* for leap seconds */
int     tm_min;         /* minutes after the hour — [0, 59] */
int     tm_hour;        /* hour since midnight — [0, 23] */
int     tm_mday;        /* day of the month — [1, 31] */
int     tm_mon;         /* months since January — [0, 11] */
int     tm_year;        /* years since 1900 */
int     tm_wday;        /* days since Sunday — [0, 6] */
int     tm_yday;        /* days since January 1 — [0, 365] */
int     tm_isdst;       /* flag for alternate daylight savings time */
```

The value of **tm_isdst** is positive if daylight savings time is in effect, zero if daylight savings time is not in effect, and negative if the information is not available. (Previously, the value of **tm_isdst** was defined as non-zero if daylight savings was in effect.)

The external **time_t** variable **altzone** contains the difference, in seconds, between Coordinated Universal Time and the alternate time zone. The external variable **timezone** contains the difference, in seconds, between UTC and local standard time. The external variable **daylight** indicates whether time should reflect daylight savings time. Both **timezone** and **altzone** default to 0 (UTC). The external variable **daylight** is non-zero if an alternate time zone exists. The time zone names are contained in the external variable **tzname**, which by default is set to:

```
char *tzname[2] = { "GMT", " " };
```

These functions know about the peculiarities of this conversion for various time periods for the U.S. (specifically, the years 1974, 1975, and 1987). They will handle the new daylight savings time starting with the first Sunday in April, 1987.

tzset( ) uses the contents of the environment variable **TZ** to override the value of the different external variables. The function **tzset( )** is called by **asctime( )** and may also be called by the user. See **environ**(5) for a description of the **TZ** environment variable.

Starting and ending times are relative to the current local time zone. If the alternate time zone start and end dates and the time are not provided, the days for the United States that year will be used and the time will be 2 AM. If the start and end dates are provided but the time is not provided, the time will be 2 AM. The effects of **tzset( )** change the values of the external variables **timezone**, **altzone**, **daylight**, and **tzname**.

Note that in most installations, **TZ** is set to the correct value by default when the user logs on, using the local **/etc/default/init** file (see **TIMEZONE**(4)).

tzsetwall( ) sets things up so that **localtime( )** returns the best available approximation of local wall clock time.

LC_TIME determines how these functions handle date and time formats. In the "C" locale, date and time handling follow the U.S. rules.

ERRORS

ctime_r( ) and **asctime_r( )** will fail if the following is true:

ERANGE length of the buffer supplied by caller is not large enough to store the result.

EXAMPLES

tzset( ) scans the contents of the environment variable and assigns the different fields to the respective variable. For example, the most complete setting for New Jersey in 1986 could be

EST5EDT4,116/2:00:00,298/2:00:00

or simply

EST5EDT

An example of a southern hemisphere setting such as the Cook Islands could be

KDT9:30KST10:00,63/5:00,302/20:00

In the longer version of the New Jersey example of **TZ, tzname**[*0*] is EST, **timezone** will be set to 5∗60∗60, **tzname**[*1*] is EDT, **altzone** will be set to 4∗60∗60, the starting date of the alternate time zone is the 117th day at 2 AM, the ending date of the alternate time zone is the 299th day at 2 AM (using zero-based Julian days), and **daylight** will be set positive. Starting and ending times are relative to the current local time zone. If the alternate time zone start and end dates and the time are not provided, the days for the United States that year will be used and the time will be 2 AM. If the start and end dates are provided but the time is not provided, the time will be 2 AM. The effects of **tzset( )** are thus to change the values of the external variables **timezone**, **altzone**, **daylight**, and **tzname**. **ctime( )**, **localtime( )**, **mktime( )**, and **strftime( )** will also update these external variables as if they had called **tzset( )** at the time specified by the **time_t** or **struct tm** value that they are converting.

FILES

/usr/lib/locale/*locale***/LC_TIME/time**
file containing locale specific date and time information

SEE ALSO

time(2), **getenv**(3C), **mktime**(3C), **printf**(3S), **putenv**(3C), **setlocale**(3C), **strftime**(3C), **TIMEZONE**(4), **environ**(5)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

The return values for **ctime( )**, **localtime( )**, and **gmtime( )** point to static data whose content is overwritten by each call.

Setting the time during the interval of change from **timezone** to **altzone** or vice versa can produce unpredictable results. The system administrator must change the Julian start and end days annually.

asctime( ), **ctime( )**, **gmtime( )** and **localtime( )** are unsafe in multi-thread applications. **asctime_r( )**, **ctime_r( )**, **gmtime_r( )** and **localtime_r( )** are MT-Safe, and should be used instead. **tzset( )** and **tzsetwall( )** are unsafe in multi-thread applications.

All interfaces on this page with an _r suffix are as specified in POSIX 1003.1c Draft #10.

NAME

ctype, isdigit, isxdigit, islower, isupper, isalpha, isalnum, isspace, iscntrl, ispunct, isprint, isgraph, isascii – character handling

SYNOPSIS

```
#include <ctype.h>
int isalpha(int c);
int isupper(int c);
int islower(int c);
int isdigit(int c);
int isxdigit(int c);
int isalnum(int c);
int isspace(int c);
int ispunct(int c);
int isprint(int c);
int isgraph(int c);
int iscntrl(int c);
int isascii(int c);
```

MT-LEVEL

MT-Safe with exceptions

DESCRIPTION

These macros classify character-coded integer values. Each is a predicate returning non-zero for true, zero for false. The behavior of these macros, except **isascii( )**, is affected by the current locale (see **setlocale**(3C)). To modify the behavior, change the **LC_TYPE** category in **setlocale( )**, that is, **setlocale(LC_CTYPE,** *newlocale***)**. In the **C** locale, or in a locale where character type information is not defined, characters are classified according to the rules of the US-ASCII 7-bit coded character set.

The macro **isascii( )** is defined on all integer values; the rest are defined only where the argument is an **int**, the value of which is representable as an **unsigned char**, or **EOF**, which is defined by the **stdio.h** header and represents end-of-file.

isalpha( ) tests for any character for which **isupper( )** or **islower( )** is true, or any character that is one of an implementation-defined set of characters for which none of **iscntrl( )**, **isdigit( )**, **ispunct( )**, or **isspace( )** is true. In the **C** locale, **isalpha( )** returns true only for the characters for which **isupper( )** or **islower( )** is true.

isupper( ) tests for any character that is an upper-case letter or is one of an implementation-defined set of characters for which none of **iscntrl( )**, **isdigit( )**, **ispunct( )**, **isspace( )**, or **islower( )** is true. In the **C** locale, **isupper( )** returns true only for the characters defined as upper-case ASCII characters.

islower( ) tests for any character that is a lower-case letter or is one of an implementation-defined set of characters for which none of **iscntrl( )**, **isdigit( )**, **ispunct( )**, **isspace( )**, or **isupper( )** is true. In the **C** locale, **islower( )** returns true only for the characters defined as lower-case ASCII characters.

isdigit( ) tests for any decimal-digit character.

isxdigit( ) tests for any hexadecimal-digit character (**[0–9]**, **[A–F]** or **[a–f]**).

isalnum( ) tests for any character for which **isalpha( )** or **isdigit( )** is true (letter or digit).

isspace( ) tests for any space, tab, carriage-return, newline, vertical-tab or form-feed (standard white-space characters) or for one of an implementation-defined set of characters for which **isalnum( )** is false. In the **C** locale, **isspace( )** returns true only for the standard white-space characters.

ispunct( ) tests for any printing character which is neither a space nor a character for which **isalnum( )** is true.

isprint( ) tests for any printing character, including space (" ").

isgraph( ) tests for any printing character, except space.

iscntrl( ) tests for any ‘‘control character’’ as defined by the character set.

isascii( ) tests for any ASCII character, code between **0** and **0177** inclusive.

All the character classification macros and the conversion functions and macros use a table lookup.

Functions exist for all the above-defined macros. To get the function form, the macro name must be undefined (for example, **#undef isdigit**).

RETURN VALUES

If the argument to any of the character handling macros is not in the domain of the function, the result is undefined.

FILES

/usr/lib/locale/*locale***/LC_CTYPE**

SEE ALSO

chrtbl(1M), **setlocale**(3C), **stdio**(3S), **ascii**(5), **environ**(5)

NOTES

isdigit( ), **isxdigit( )**, **islower( )**, **isupper( )**, **isalpha( )**, **isalnum( )**, **isspace( )**, **iscntrl( )**, **ispunct( )**, **isprint( )**, **isgraph( )** and **isascii( )** can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

curs_addch, addch, waddch, mvaddch, mvwaddch, echochar, wechochar – add a character (with attributes) to a curses window and advance cursor

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int addch(chtype *ch***);**

int waddch(WINDOW ∗*win***, chtype** *ch***);**

int mvaddch(int *y***, int** *x***, chtype** *ch***);**

int mvwaddch(WINDOW ∗*win***, int** *y***, int** *x***, chtype** *ch***);**

int echochar(chtype *ch***);**

int wechochar(WINDOW ∗*win*, **chtype** *ch);*

MT-LEVEL

Unsafe

DESCRIPTION

With the **addch( )**, **waddch( )**, **mvaddch( )**, and **mvwaddch( )** routines, the character *ch* is put into the window at the current cursor position of the window and the position of the window cursor is advanced. Its function is similar to that of **putchar( )**. At the right margin, an automatic newline is performed. At the bottom of the scrolling region, if **scrollok( )** is enabled, the scrolling region is scrolled up one line.

If *ch* is a tab, newline, or backspace, the cursor is moved appropriately within the window. A newline also does a **clrtoeol( )** before moving. Tabs are considered to be at every eighth column. If *ch* is another control character, it is drawn in the ^*X* notation. Calling **winch( )** after adding a control character does not return the control character, but instead returns the representation of the control character. See **curs_inch**(3X).

Video attributes can be combined with a character by OR-ing them into the parameter. This results in these attributes also being set. (The intent here is that text, including attributes, can be copied from one place to another using **inch( )** and **addch( )**.) (see **standout( )**, predefined video attribute constants, on the **curs_attr**(3X) page).

The **echochar( )** and **wechochar( )** routines are functionally equivalent to a call to **addch( )** followed by a call to **refresh( )**, or a call to **waddch** followed by a call to **wrefresh( )**. The knowledge that only a single character is being output is taken into consideration and, for non-control characters, a considerable performance gain might be seen by using these routines instead of their equivalents.

Line Graphics

The following variables may be used to add line drawing characters to the screen with routines of the **addch( )** family. When variables are defined for the terminal, the **A_ALTCHARSET** bit is turned on (see **curs_attr**(3X)). Otherwise, the default character listed below is stored in the variable. The names chosen are consistent with the VT100 nomenclature.

Name	*Default*	*Glyph Description*
ACS_ULCORNER	+	upper left-hand corner
ACS_LLCORNER	+	lower left-hand corner
ACS_URCORNER	+	upper right-hand corner
ACS_LRCORNER	+	lower right-hand corner
ACS_RTEE	+	right tee (⊣)
ACS_LTEE	+	left tee (⊢)
ACS_BTEE	+	bottom tee (⊥)
ACS_TTEE	+	top tee (⊤)
ACS_HLINE	–	horizontal line
ACS_VLINE	|	vertical line
ACS_PLUS	+	plus
ACS_S1	–	scan line 1
ACS_S9	_	scan line 9
ACS_DIAMOND	+	diamond
ACS_CKBOARD	:	checker board (stipple)
ACS_DEGREE	'	degree symbol
ACS_PLMINUS	#	plus / minus
ACS_BULLET	o	bullet
ACS_LARROW	<	arrow pointing left
ACS_RARROW	>	arrow pointing right
ACS_DARROW	v	arrow pointing down
ACS_UARROW	^	arrow pointing up
ACS_BOARD	#	board of squares
ACS_LANTERN	#	lantern symbol
ACS_BLOCK	#	solid square block

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curs_attr(3X), **curs_clear**(3X), **curs_inch**(3X), **curs_outopts**(3X), **curs_refresh**(3X), **curses**(3X), **putc**(3S)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **addch( )**, **mvaddch( )**, **mvwaddch( )**, and **echochar( )** may be macros.

NAME

curs_addchstr, addchstr, addchnstr, waddchstr, waddchnstr, mvaddchstr, mvaddchnstr, mvwaddchstr, mvwaddchnstr – add string of characters (and attributes) to a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int addchstr(chtype ∗*chstr***);**

int addchnstr(chtype ∗*chstr***, int** *n***);**

int waddchstr(WINDOW ∗*win***, chtype** ∗*chstr***);**

int waddchnstr(WINDOW ∗*win***, chtype** ∗*chstr***, int** *n***);**

int mvaddchstr(int *y***, int** *x***, chtype** ∗*chstr***);**

int mvaddchnstr(int *y***, int** *x***, chtype** ∗*chstr***, int** *n***);**

int mvwaddchstr(WINDOW ∗**win, int** *y***, int** *x***, chtype** ∗*chstr***);**

int mvwaddchnstr(WINDOW ∗**win, int** *y***, int** *x***, chtype** ∗*chstr***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

All of these routines copy *chstr* directly into the window image structure starting at the current cursor position. The four routines with *n* as the last argument copy at most *n* elements, but no more than will fit on the line. If **n**=**-1** then the whole string is copied, to the maximum number that fit on the line.

The position of the window cursor is **not** advanced. These routines works faster than **waddnstr( )** (see **curs_addstr**(3X)) because they merely copy *chstr* into the window image structure. On the other hand, care must be taken when using these functions because they do not perform any kind of checking (such as for the newline character), they do not advance the current cursor position, and they truncate the string, rather then wrapping it around to the next line.

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curs_addstr(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that all routines except **waddchnstr( )** and **waddchstr( )** may be macros.

NAME

curs_addstr, addstr, addnstr, waddstr, waddnstr, mvaddstr, mvaddnstr, mvwaddstr, mvwaddnstr – add a string of characters to a curses window and advance cursor

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int addstr(char ∗*str***);**

int addnstr(char ∗*str***, int** *n***);**

int waddstr(WINDOW ∗*win***, char** ∗*str***);**

int waddnstr(WINDOW ∗*win***, char** ∗*str***, int** *n***);**

int mvaddstr(int *y***, int** *x***, char** ∗*str***);**

int mvaddnstr(int *y***, int** *x***, char** ∗*str***, int** *n***);**

int mvwaddstr(WINDOW ∗*win***, int** *y***, int** *x***, char** ∗*str***);**

int mvwaddnstr(WINDOW ∗*win***, int** *y***, int** *x***, char** ∗*str***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

All of these routines write all the characters of the null terminated character string *str* on the given window. It is similar to calling **waddch( )** once for each character in the string. The four routines with *n* as the last argument write at most *n* characters. If *n* is negative, then the entire string will be added.

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curs_addch(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that all routines except **waddstr( )** and **waddnstr( )** may not be macros.

NAME

curs_addwch, addwch, waddwch, mvaddwch, mvwaddwch, echowchar, wechowchar – add a wchar_t character (with attributes) to a curses window and advance cursor

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]
#include <curses.h>

int addwch(chtype *wch***);**

int waddwch(WINDOW ∗*win***, chtype** *wch***);**

int mvaddwch(int *y***, int** *x***, chtype** *wch* **);**

int mvwaddwch(WINDOW ∗*win***, int** *y***, int** *x* **, chtype** *wch***);**

int echowchar(chtype *wch***);**

int wechowchar(WINDOW ∗*win***, chtype** *wch***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **addwch( )**,**waddwch( )**,**mvaddwch( )**, and **mvwaddwch( )** routines put the character *wch*, holding a **wchar_t** character, into the window at the current cursor position of the window and advance the position of the window cursor. Their function is similar to that of **putwchar**(3I) in the C multibyte library. At the right margin, an automatic newline is performed. At the bottom of the scrolling region, if **scrollok** is enabled, the scrolling region is scrolled up one line.

If *wch* is a tab, newline, or backspace, the cursor is moved appropriately within the window. A newline also does a **clrtoeol**(3X) before moving. Tabs are considered to be at every eighth column. If *wch* is another control character, it is drawn in the ^*X* notation. Calling **winwch**(3X) after adding a control character does not return the control character, but instead returns the representation of the control character.

Video attributes can be combined with a **wchar_t** character by OR-ing them into the parameter. This results in these attributes also being set. (The intent here is that text, including attributes, can be copied from one place to another using **inwch( )** and **addwch( )**.) See **standout**(3X), predefined video attribute constants.

The **echowchar( )** and **wechowchar( )** routines are functionally equivalent to a call to **addwch( )** followed by a call to **refresh**(3X), or a call to **waddwch( )** followed by a call to **wrefresh**(3X). The knowledge that only a single character is being output is taken into consideration and, for non-control characters, a considerable performance gain might be seen by using these routines instead of their equivalents.

Line Graphics

The following variables may be used to add line drawing characters to the screen with routines of the **addwch( )** family. When variables are defined for the terminal, the **A_ALTCHARSET** bit is turned on. (See **curs_attr**(3X)). Otherwise, the default character listed below is stored in the variable. The names chosen are consistent with the VT100 nomenclature.

Name	*Default*	*Glyph Description*
ACS_ULCORNER	**+**	upper left-hand corner
ACS_LLCORNER	**+**	lower left-hand corner
ACS_URCORNER	**+**	upper right-hand corner
ACS_LRCORNER	**+**	lower right-hand corner
ACS_RTEE	**+**	right tee (┤)
ACS_LTEE	**+**	left tee (├)
ACS_BTEE	**+**	bottom tee (┴)
ACS_TTEE	**+**	top tee (┬)
ACS_HLINE	–	horizontal line
ACS_VLINE	**|**	vertical line
ACS_PLUS	**+**	plus
ACS_S1	–	scan line 1
ACS_S9	_	scan line 9
ACS_DIAMOND	**+**	diamond
ACS_CKBOARD	**:**	checker board (stipple)
ACS_DEGREE	**'**	degree symbol
ACS_PLMINUS	**#**	plus/minus
ACS_BULLET	**o**	bullet
ACS_LARROW	**<**	arrow pointing left
ACS_RARROW	**>**	arrow pointing right
ACS_DARROW	**v**	arrow pointing down
ACS_UARROW	**^**	arrow pointing up
ACS_BOARD	**#**	board of squares
ACS_LANTERN	**#**	lantern symbol
ACS_BLOCK	**#**	solid square block

RETURN VALUE

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion, unless otherwise noted in the preceding routine descriptions.

SEE ALSO

putwchar(3I), **clrtoeol**(3X), **curses**(3X), **curs_attr**(3X), **curs_inwch**(3X), **curs_outopts**(3X), **refresh**(3X), **standout**(3X), **winwch**(3X), **wrefresh**(3X)

NOTES

The header file **<curses.h>** automatically includes the header files **<stdio.h>**, **<unctrl.h>** and **<widec.h>**.

Note that **addwch()**, **mvaddwch()**, **mvwaddwch()**, and **echowchar()** may be macros.

None of these routines can use the color attribute in **chtype**.

NAME

curs_addwchstr, addwchstr, addwchnstr, waddwchstr, waddwchnstr, mvaddwchstr, mvaddwchnstr, mvwaddwchstr, mvwaddwchnstr – add string of wchar_t characters (and attributes) to a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lcurses** [*library* . .]
#include <curses.h>

int addwchstr(chtype ∗*wchstr***);**

int addwchnstr(chtype ∗*wchstr***, int** *n***);**

int waddwchstr(WINDOW ∗*win***, chtype** ∗*wchstr***);**

int waddwchnstr(WINDOW ∗*win***, chtype** ∗*wchstr***, int** *n* **);**

int mvaddwchstr(int *y***, int** *x***, chtype** ∗*wchstr* **);**

int mvaddwchnstr(int *y***, int** *x***, chtype** ∗ *wchstr***, int** *n***);**

int mvwaddwchstr(WINDOW ∗*win***, int** *y***, int** *x***, chtype** ∗*wchstr***);**

int mvwaddwchnstr(WINDOW ∗*win***, int** *y***, int** *x***, chtype** ∗*wchstr***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

All of these routines copy *wchstr*, which points to a string of **wchar_t** characters, directly into the window image structure starting at the current cursor position. The four routines with *n* as the last argument copy at most *n* elements, but no more than will fit on the line. If **n**=**–1** then the whole string is copied, to the maximum number that fit on the line.

The position of the window cursor is not advanced. These routines work faster than **waddnwstr**(3X) because they merely copy *wchstr* into the window image structure. On the other hand, care must be taken when using these functions because they don't perform any kind of checking (such as for the newline character), they do not advance the current cursor position, and they truncate the string, rather than wrapping it around to the new line.

RETURN VALUE

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion, unless otherwise noted in the preceding routine descriptions.

SEE ALSO

curses(3X), **waddnwstr**(3X)

NOTES

The header file **<curses.h>** automatically includes the header files **<stdio.h>**, **<unctrl.h>** and **<widec.h>**.

Note that all routines except **waddwchnstr( )** may be macros.

None of these routines can use the color attribute in **chtype**.

NAME

curs_addwstr, addwstr, addnwstr, waddwstr, waddnwstr, mvaddwstr, mvaddnwstr, mvwaddwstr, mvwaddnwstr – add a string of wchar_t characters to a curses window and advance cursor

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]
#include <curses.h>

int addwstr(wchar_t ∗*wstr***);**

int addnwstr(wchar_t ∗*wstr***, int** *n***);**

int waddwstr(WINDOW ∗*win***, wchar_t** ∗*wstr***);**

int waddnwstr(WINDOW ∗*win***, wchar_t** ∗*wstr***, int** *n***);**

int mvaddwstr(int *y***, int** *x***, wchar_t** ∗ *wstr***);**

int mvaddnwstr(int *y* **, int** *x***, wchar_t** ∗ *wstr***, int** *n***);**

int mvwaddwstr(WINDOW ∗*win***, int** *y***, int** *x***, wchar_t** ∗*wstr***);**

int mvwaddnwstr(WINDOW ∗*win***, int** *y***, int** *x* **, wchar_t** ∗*wstr***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

All of these routines write all the characters of the null-terminated **wchar_t** character string **wstr** on the given window. The effect is similar to calling **waddwch**(3X) once for each **wchar_t** character in the string. The four routines with *n* as the last argument write at most *n* **wchar_t** characters. If *n* is negative, then the entire string will be added.

RETURN VALUE

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curses(3X), **waddwch**(3X)

NOTES

The header file **<curses.h>** automatically includes the header files **<stdio.h>**, **<unctrl.h>** and **<widec.h**.

Note that all of these routines except **waddwstr( )** and **waddnwstr( )** may be macros.

NAME

curs_alecompat, movenextch, wmovenextch, moveprevch, wmoveprevch, adjcurspos, wadjcurspos – these functions are added to ALE curses library for moving the cursor by character.

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]
#include <curses.h>

int movenextch(void);

int wmovenextch(WINDOW ∗*win***);**

int moveprevch(void);

int wmoveprevch(WINDOW ∗*win***);**

int adjcurspos(void);

int wadjcurspos(WINDOW ∗*win***);**

MT-LEVEL

Unsafe

DESCRIPTION

movenextch( ) and **wmovenextch( )** move the cursor to the next character to the right. If the next character is a multicolumn character, the cursor is positioned on the first (left-most) column of that character. The new cursor position will be on the next character, even if the cursor was originally positioned on the left-most column of a multicolumn character. Note that the simple cursor increment (**++x**) does not guarantee movement to the next character, if the cursor was originally positioned on a multicolumn character. **getyx**(3X) can be used to find the new position.

moveprevc( ) and **wmoveprevch( )** routines are the opposite of **movenextch( )** and **wmovenextch( )**, moving the cursor to the left-most column of the previous character.

adjcurspos( ) and **wadjcurspos( )** move the cursor to the first(left-most) column of the multicolumn character that the cursor is presently on. If the cursor is already on the first column, or if the cursor is on a single-column character, these routines will have no effect.

RETURN VALUE

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curses(3X), **getyx**(3X)

NOTES

The header file **<curses.h>** automatically includes the header files **<stdio.h>** , **<unctrl.h>** and **<widec.h>**.

Note that **movenextch()**, **moveprevch()**, and **adjcurspos()** may be macros.

NAME

curs_attr, attroff, wattroff, attron, wattron, attrset, wattrset, standend, wstandend, standout, wstandout – curses character and window attribute control routines

SYNOPSIS

cc [*flag* ...] *file* ... –**lcurses** [*library* ..]

#include <curses.h>

int attroff(int *attrs***);**

int wattroff(WINDOW ∗*win***, int** *attrs***);**

int attron(int *attrs***);**

int wattron(WINDOW ∗*win***, int** *attrs***);**

int attrset(int *attrs***);**

int wattrset(WINDOW ∗*win***, int** *attrs***);**

int standend(void);

int wstandend(WINDOW ∗*win***);**

int standout (void);

int wstandout(WINDOW ∗*win***);**

MT-LEVEL

Unsafe

DESCRIPTION

All of these routines manipulate the current attributes of the named window. The current attributes of a window are applied to all characters that are written into the window with **waddch( )**, **waddstr( )**, and **wprintw( )**. Attributes are a property of the character, and move with the character through any scrolling and insert / delete line / character operations. To the extent possible on the particular terminal, they are displayed as the graphic rendition of characters put on the screen.

The routine **attrset( )** sets the current attributes of the given window to *attrs*. The routine **attroff( )** turns off the named attributes without turning any other attributes on or off. The routine **attron( )** turns on the named attributes without affecting any others. The routine **standout( )** is the same as **attron(A_STANDOUT)**. The routine **standend( )** is the same as **attrset**( )**,** that is, it turns off all attributes.

Attributes

The following video attributes, defined in **<curses.h>**, can be passed to the routines **attron( )**, **attroff( )**, and **attrset( )**, or OR-ed with the characters passed to **addch( )**.

A_STANDOUT	**Best highlighting mode of the terminal.**
A_UNDERLINE	**Underlining**
A_REVERSE	**Reverse video**
A_BLINK	**Blinking**
A_DIM	**Half bright**
A_BOLD	**Extra bright or bold**
A_ALTCHARSET	**Alternate character set**
A_CHARTEXT	**Bit-mask to extract a character**
COLOR_PAIR(*n***)**	**Color-pair number** *n*

The following macro is the reverse of **COLOR_PAIR(***n***)**:

PAIR_NUMBER(*attrs***)** Returns the pair number associated with the **COLOR_PAIR(***n***)** attribute.

RETURN VALUES

These routines always return 1.

SEE ALSO

curs_addch(3X), **curs_addstr**(3X), **curs_printw**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **attroff( )**, **wattroff( )**, **attron( )**, **wattron( )**, **wattrset( )**, **standend( )**, and **standout( )** may be macros.

NAME

curs_beep, beep, flash – curses bell and screen flash routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int beep(void);

int flash(void);

MT-LEVEL

Unsafe

DESCRIPTION

The **beep( )** and **flash( )** routines are used to signal the terminal user. The routine **beep( )** sounds the audible alarm on the terminal, if possible; if that is not possible, it flashes the screen (visible bell), if that is possible. The routine **flash( )** flashes the screen, and if that is not possible, sounds the audible signal. If neither signal is possible, nothing happens. Nearly all terminals have an audible signal (bell or beep), but only some can flash the screen.

RETURN VALUES

These routines always return **OK**.

SEE ALSO

curses(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

NAME

curs_bkgd, bkgd, bkgdset, wbkgdset, wbkgd – curses window background manipulation routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int bkgd(chtype *ch***);**

void bkgdset(chtype *ch***);**

void wbkgdset(WINDOW ∗*win***, chtype** *ch***);**

int wbkgd(WINDOW ∗*win***, chtype** *ch***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **bkgdsets( )** and **wbkgdset( )** routines manipulate the background of the named window. Background is a **chtype** consisting of any combination of attributes and a character. The attribute part of the background is combined (ORed) with all non-blank characters that are written into the window with **waddch( )**. Both the character and attribute parts of the background are combined with the blank characters. The background becomes a property of the character and moves with the character through any scrolling and insert / delete line / character operations. To the extent possible on a particular terminal, the attribute part of the background is displayed as the graphic rendition of the character put on the screen.

The **bkgd( )** and **wbkgd( )** routines combine the new background with every position in the window. Background is any combination of attributes and a character. Only the attribute part is used to set the background of non-blank characters, while both character and attributes are used for blank positions. To the extent possible on a particular terminal, the attribute part of the background is displayed as the graphic rendition of the character put on the screen.

RETURN VALUES

bkgd( ) and **wbkgd( )** return the integer **OK**, or a non-negative integer, if **immedok( )** is set. See **curs_outopts**(3X).

SEE ALSO

curs_addch(3X), **curs_outopts**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **bkgdset( )** and **bkgd( )** may be macros.

NAME

curs_border, border, wborder, box, whline, wvline – create curses borders, horizontal and vertical lines

SYNOPSIS

cc [*flag* ...] *file* ... **–lcurses** [*library* ..]

#include <curses.h>

int border(chtype *ls***, chtype** *rs***, chtype** *ts***, chtype** *bs***, chtype** *tl***, chtype** *tr***, chtype** *bl***, chtype** *br***);**

int wborder(WINDOW ∗*win***, chtype** *ls***, chtype** *rs***, chtype** *ts***, chtype** *bs***, chtype** *tl***, chtype** *tr***, chtype** *bl***, chtype** *br***);**

int box(WINDOW ∗*win***, chtype** *verch***, chtype** *horch***);**

int hline(chtype *ch***, int** *n***);**

int whline(WINDOW ∗*win***, chtype** *ch***, int** *n***);**

int vline(chtype *ch***, int** *n***);**

int wvline(WINDOW ∗*win***, chtype** *ch***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

With the **border( )**, **wborder( )**, and **box( )** routines, a border is drawn around the edges of the window. The arguments and attributes are:

ls	left side of the border
rs	right side of the border
ts	top side of the border
bs	bottom side of the border
tl	top left-hand corner
tr	top right-hand corner
bl	bottom left-hand corner
br	bottom right-hand corner

If any of these arguments is zero, then the following default values (defined in **<curses.h>**) are used respectively instead: **ACS_VLINE**, **ACS_VLINE**, **ACS_HLINE**, **ACS_HLINE**, **ACS_ULCORNER**, **ACS_URCORNER**, **ACS_BLCORNER**, **ACS_BRCORNER**.

box(*win***,** *verch***,** *horch***)** is a shorthand for the following call:

wborder(*win***,** *verch***,** *verch***,** *horch***, horch , 0, 0, 0, 0)**

hline( ) and **whline( )** draw a horizontal (left to right) line using *ch* starting at the current cursor position in the window. The current cursor position is not changed. The line is at most *n* characters long, or as many as fit into the window.

vline( ) and **wvline( )** draw a vertical (top to bottom) line using *ch* starting at the current cursor position in the window. The current cursor position is not changed. The line is at most *n* characters long, or as many as fit into the window.

RETURN VALUES

All routines return the integer **OK**, or a non-negative integer if **immedok( )** is set. See **curs_outopts**(3X).

SEE ALSO

curs_outopts(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **border( )** and **box( )** may be macros.

NAME

curs_clear, erase, werase, clear, wclear, clrtobot, wclrtobot, clrtoeol, wclrtoeol – clear all or part of a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int erase(void);

int werase(WINDOW ∗*win***);**

int clear(void);

int wclear(WINDOW ∗*win***);**

int clrtobot(void);

int wclrtobot (WINDOW ∗*win***);**

int clrtoeol(void);

int wclrtoeol(WINDOW ∗*win***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **erase( )** and **werase( )** routines copy blanks to every position in the window.

The **clear( )** and **wclear( )** routines are like **erase( )** and **werase( )**, but they also call **clearok( ) ,** so that the screen is cleared completely on the next call to **wrefresh( )** for that window and repainted from scratch.

The **clrtobot( )** and **wclrtobot( )** routines erase all lines below the cursor in the window. Also, the current line to the right of the cursor, inclusive, is erased.

The **clrtoeol( )** and **wclrtoeol( )** routines erase the current line to the right of the cursor, inclusive.

RETURN VALUES

All routines return the integer **OK**, or a non-negative integer if **immedok( )** is set. See **curs_outopts**(3X).

SEE ALSO

curs_outopts(3X), **curs_refresh**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **erase( )**, **werase( )**, **clear( )**, **wclear( )**, **clrtobot( )**, and **clrtoeol( )** may be macros.

NAME

curs_color, start_color, init_pair, init_color, has_colors, can_change_color, color_content, pair_content – curses color manipulation routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int start_color(void);

int init_pair(short *pair***, short** *fg***, short** *bg***);**

int init_color(short *color***, short** *red***, short** *green***, short** *blue***);**

bool has_colors(void);

bool can_change_color(void);

int color_content(short *color***, short** **redp***, short** **greenp***, short** **bluep***);**

int pair_content(short *pair***, short** **fgp***, short** **bgp***);**

MT-LEVEL

Unsafe

DESCRIPTION

Overview

curses provides routines that manipulate color on color alphanumeric terminals. To use these routines **start_color( )** must be called, usually right after **initscr( )**. See **curs_initscr**(3X). Colors are always used in pairs (referred to as color-pairs). A color-pair consists of a foreground color (for characters) and a background color (for the field on which the characters are displayed). A programmer initializes a color-pair with the routine **init_pair**. After it has been initialized, **COLOR_PAIR(***n***)**, a macro defined in **<curses.h>**, can be used in the same ways other video attributes can be used. If a terminal is capable of redefining colors, the programmer can use the routine **init_color( )** to change the definition of a color. The routines **has_colors( )** and **can_change_color( )** return **TRUE** or **FALSE**, depending on whether the terminal has color capabilities and whether the programmer can change the colors. The routine **color_content( )** allows a programmer to identify the amounts of red, green, and blue components in an initialized color. The routine **pair_content( )** allows a programmer to find out how a given color-pair is currently defined.

Routine Descriptions

The **start_color( )** routine requires no arguments. It must be called if the programmer wants to use colors, and before any other color manipulation routine is called. It is good practice to call this routine right after **initscr( )**. **start_color( )** initializes eight basic colors (black, red, green, yellow, blue, magenta, cyan, and white), and two global variables, **COLORS** and **COLOR_PAIRS** (respectively defining the maximum number of colors and color-pairs the terminal can support). It also restores the colors on the terminal to the values they had when the terminal was just turned on.

The **init_pair( )** routine changes the definition of a color-pair. It takes three arguments: the number of the color-pair to be changed, the foreground color number, and the background color number. The value of the first argument must be between **1** and **COLOR_PAIRS**–1. The value of the second and third arguments must be between 0 and

COLORS. If the color-pair was previously initialized, the screen is refreshed and all occurrences of that color-pair is changed to the new definition.

The **init_color( )** routine changes the definition of a color. It takes four arguments: the number of the color to be changed followed by three RGB values (for the amounts of red, green, and blue components). The value of the first argument must be between **0** and **COLORS**. (See the section **Colors** for the default color index.) Each of the last three arguments must be a value between 0 and 1000. When **init_color( )** is used, all occurrences of that color on the screen immediately change to the new definition.

The **has_colors( )** routine requires no arguments. It returns **TRUE** if the terminal can manipulate colors; otherwise, it returns **FALSE**. This routine facilitates writing terminal-independent programs. For example, a programmer can use it to decide whether to use color or some other video attribute.

The **can_change_color( )** routine requires no arguments. It returns **TRUE** if the terminal supports colors and can change their definitions; other, it returns **FALSE**. This routine facilitates writing terminal-independent programs.

The **color_content( )** routine gives users a way to find the intensity of the red, green, and blue (RGB) components in a color. It requires four arguments: the color number, and three addresses of **short**s for storing the information about the amounts of red, green, and blue components in the given color. The value of the first argument must be between 0 and **COLORS**. The values that are stored at the addresses pointed to by the last three arguments are between 0 (no component) and 1000 (maximum amount of component).

The **pair_content( )** routine allows users to find out what colors a given color-pair consists of. It requires three arguments: the color-pair number, and two addresses of **short**s for storing the foreground and the background color numbers. The value of the first argument must be between 1 and **COLOR_PAIRS**–1. The values that are stored at the addresses pointed to by the second and third arguments are between 0 and **COLORS**.

Colors

In **<curses.h>** the following macros are defined. These are the default colors. **curses** also assumes that **COLOR_BLACK** is the default background color for all terminals.

COLOR_BLACK
COLOR_RED
COLOR_GREEN
COLOR_YELLOW
COLOR_BLUE
COLOR_MAGENTA
COLOR_CYAN
COLOR_WHITE

RETURN VALUES

All routines that return an integer return **ERR** upon failure and **OK** upon successful completion.

SEE ALSO

curs_attr(3X), **curs_initscr**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

NAME

curs_delch, delch, wdelch, mvdelch, mvwdelch – delete character under cursor in a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int delch(void);

int wdelch(WINDOW ∗*win***);**

int mvdelch(int *y***, int** *x***);**

int mvwdelch(WINDOW ∗*win***, int** *y***, int** *x***);**

MT-LEVEL

Unsafe

DESCRIPTION

With these routines the character under the cursor in the window is deleted; all characters to the right of the cursor on the same line are moved to the left one position and the last character on the line is filled with a blank. The cursor position does not change (after moving to *y*, *x*, if specified). This does not imply use of the hardware delete character feature.

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curses(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **delch( )**, **mvdelch( )**, and **mvwdelch( )** may be macros.

NAME

curs_deleteln, deleteln, wdeleteln, insdelln, winsdelln, insertln, winsertln – delete and insert lines in a curses window

SYNOPSIS

cc [*flag* ...] *file* ... –**lcurses** [*library* ..]

#include <curses.h>

int deleteln(void);

int wdeleteln(WINDOW ∗*win***);**

int insdelln(int *n***);**

int winsdelln(WINDOW ∗*win***, int** *n***);**

int insertln(void);

int winsertln(WINDOW ∗*win***);**

MT-LEVEL

Unsafe

DESCRIPTION

With the **deleteln( )** and **wdeleteln( )** routines, the line under the cursor in the window is deleted; all lines below the current line are moved up one line. The bottom line of the window is cleared. The cursor position does not change. This does not imply use of a hardware delete line feature.

With the **insdelln( )** and **winsdelln( )** routines, for positive *n*, insert *n* lines into the specified window above the current line. The *n* bottom lines are lost. For negative *n*, delete *n* lines (starting with the one under the cursor), and move the remaining lines up. The bottom *n* lines are cleared. The current cursor position remains the same.

With the **insertln( )** and **insertln( )** routines, a blank line is inserted above the current line and the bottom line is lost. This does not imply use of a hardware insert line feature.

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curses(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that all but **winsdelln( )** may be macros.

NAME

curs_getch, getch, wgetch, mvgetch, mvwgetch, ungetch – get (or push back) characters from curses terminal keyboard

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lcurses** [*library* . .]

#include <curses.h>

int getch(void);

int wgetch(WINDOW ∗*win***);**

int mvgetch(int *y***, int** *x***);**

int mvwgetch(WINDOW ∗*win***, int** *y***, int** *x***);**

int ungetch(int *ch***);**

MT-LEVEL

Unsafe

DESCRIPTION

With the **getch( )**, **wgetch( )**, **mvgetch( )**, and **mvwgetch( )** routines a character is read from the terminal associated with the window. In no-delay mode, if no input is waiting, the value **ERR** is returned. In delay mode, the program waits until the system passes text through to the program. Depending on the setting of **cbreak( )**, this is after one character (cbreak mode), or after the first newline (nocbreak mode). In half-delay mode, the program waits until a character is typed or the specified timeout has been reached. Unless **noecho( )** has been set, the character will also be echoed into the designated window.

If the window is not a pad, and it has been moved or modified since the last call to **wrefresh( )**, **wrefresh( )** will be called before another character is read.

If **keypad( )** is **TRUE**, and a function key is pressed, the token for that function key is returned instead of the raw characters. Possible function keys are defined in **<curses.h>** with integers beginning with **0401**, whose names begin with **KEY_**. If a character that could be the beginning of a function key (such as escape) is received, **curses** sets a timer. If the remainder of the sequence does not come in within the designated time, the character is passed through; otherwise, the function key value is returned. For this reason, many terminals experience a delay between the time a user presses the escape key and the escape is returned to the program. Since tokens returned by these routines are outside the ASCII range, they are not printable.

The **ungetch( )** routine places *ch* back onto the input queue to be returned by the next call to **wgetch( )**.

Function Keys

The following function keys, defined in **<curses.h>**, might be returned by **getch( )** if **keypad( )** has been enabled. Note that not all of these may be supported on a particular terminal if the terminal does not transmit a unique code when the key is pressed or if the definition for the key is not present in the *terminfo* database.

Name	*Key name*
KEY_BREAK	Break key
KEY_DOWN	The four arrow keys ...
KEY_UP	
KEY_LEFT	
KEY_RIGHT	
KEY_HOME	Home key (upward+left arrow)
KEY_BACKSPACE	Backspace
KEY_F0	Function keys; space for 64 keys is reserved.
KEY_F(*n*)	For $0 \le n \le 63$
KEY_DL	Delete line
KEY_IL	Insert line
KEY_DC	Delete character
KEY_IC	Insert char or enter insert mode
KEY_EIC	Exit insert char mode
KEY_CLEAR	Clear screen
KEY_EOS	Clear to end of screen
KEY_EOL	Clear to end of line
KEY_SF	Scroll 1 line forward
KEY_SR	Scroll 1 line backward (reverse)
KEY_NPAGE	Next page
KEY_PPAGE	Previous page
KEY_STAB	Set tab
KEY_CTAB	Clear tab
KEY_CATAB	Clear all tabs
KEY_ENTER	Enter or send
KEY_SRESET	Soft (partial) reset
KEY_RESET	Reset or hard reset
KEY_PRINT	Print or copy
KEY_LL	Home down or bottom (lower left). Keypad is arranged like this: **A1 up A3** **left B2 right** **C1 down C3**
KEY_A1	Upper left of keypad
KEY_A3	Upper right of keypad
KEY_B2	Center of keypad
KEY_C1	Lower left of keypad
KEY_C3	Lower right of keypad
KEY_BTAB	Back tab key
KEY_BEG	Beg(inning) key
KEY_CANCEL	Cancel key
KEY_CLOSE	Close key
KEY_COMMAND	Cmd (command) key
KEY_COPY	Copy key
KEY_CREATE	Create key
KEY_END	End key
KEY_EXIT	Exit key
KEY_FIND	Find key

KEY_HELP	Help key
KEY_MARK	Mark key
KEY_MESSAGE	Message key
KEY_MOVE	Move key
KEY_NEXT	Next object key
KEY_OPEN	Open key
KEY_OPTIONS	Options key
KEY_PREVIOUS	Previous object key
KEY_REDO	Redo key
KEY_REFERENCE	Reference key
KEY_REFRESH	Refresh key
KEY_REPLACE	Replace key
KEY_RESTART	Restart key
KEY_RESUME	Resume key
KEY_SAVE	Save key
KEY_SBEG	Shifted beginning key
KEY_SCANCEL	Shifted cancel key
KEY_SCOMMAND	Shifted command key
KEY_SCOPY	Shifted copy key
KEY_SCREATE	Shifted create key
KEY_SDC	Shifted delete char key
KEY_SDL	Shifted delete line key
KEY_SELECT	Select key
KEY_SEND	Shifted end key
KEY_SEOL	Shifted clear line key
KEY_SEXIT	Shifted exit key
KEY_SFIND	Shifted find key
KEY_SHELP	Shifted help key
KEY_SHOME	Shifted home key
KEY_SIC	Shifted input key
KEY_SLEFT	Shifted left arrow key
KEY_SMESSAGE	Shifted message key
KEY_SMOVE	Shifted move key
KEY_SNEXT	Shifted next key
KEY_SOPTIONS	Shifted options key
KEY_SPREVIOUS	Shifted prev key
KEY_SPRINT	Shifted print key
KEY_SREDO	Shifted redo key
KEY_SREPLACE	Shifted replace key
KEY_SRIGHT	Shifted right arrow
KEY_SRSUME	Shifted resume key
KEY_SSAVE	Shifted save key
KEY_SSUSPEND	Shifted suspend key
KEY_SUNDO	Shifted undo key
KEY_SUSPEND	Suspend key
KEY_UNDO	Undo key

RETURN VALUES

All routines return the integer **ERR** upon failure. The **ungetch( )** routine returns an integer value other than **ERR** upon successful completion. The other routines return the next input character or function key code upon successful completion.

SEE ALSO

curs_inopts(3X), **curs_move**(3X), **curs_refresh**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Use of the escape key for a single character function is discouraged.

When using **getch( )**, **wgetch( )**, **mvgetch( )**, or **mvwgetch( )**, **nocbreak** mode (**nocbreak( )**) and **echo** mode (**echo( )**) should not be used at the same time. Depending on the state of the tty driver when each character is typed, the program may produce undesirable results.

Note that **getch( )**, **mvgetch( )**, and **mvwgetch( )** may be macros.

NAME

curs_getstr, getstr, wgetstr, mvgetstr, mvwgetstr, wgetnstr – get character strings from curses terminal keyboard

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lcurses** [*library* . .]

#include <curses.h>

int getstr(char ∗**str);**

int wgetstr(WINDOW ∗*win***, char** ∗*str***);**

int mvgetstr(int *y***, int** *x***, char** ∗*str***);**

int mvwgetstr(WINDOW ∗*win***, int** *y***, int** *x***, char** ∗*str***);**

int wgetnstr(WINDOW ∗*win***, char** ∗*str***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

The effect of **getstr( )** is as though a series of calls to **getch( )** were made, until a newline or carriage return is received. The resulting value is placed in the area pointed to by the character pointer *str*. **wgetnstr( )** reads at most *n* characters, thus preventing a possible overflow of the input buffer. The user's erase and kill characters are interpreted, as well as any special keys (such as function keys, HOME key, CLEAR key, etc.).

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curs_getch(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **getstr( )**, **mvgetstr( )**, and **mvwgetstr( )** may be macros.

NAME

curs_getwch, getwch, wgetwch, mvgetwch, mvwgetwch, ungetwch – get (or push back) wchar_t characters from curses terminal keyboard

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]
#include <curses.h>

int getwch(void);

int wgetwch(WINDOW ∗*win***);**

int mvgetwch(int *y***, int** *x***);**

int mvwgetwch(WINDOW ∗*win***, int** *y***, int** *x* **);**

int ungetwch(int *wch***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **getwch( )**, **wgetwch( )**, **mvgetwch( )**, and **mvwgetwch( )** routines read an EUC character from the terminal associated with the window, transform it into a **wchar_t** character, and return a **wchar_t** character. In no-delay mode, if no input is waiting, the value **ERR** is returned. In delay mode, the program waits until the system passes text through to the program. Depending on the setting of **cbreak**, this is after one character (**cbreak** mode), or after the first newline (**nocbreak** mode). In **half-delay** mode, the program waits until a character is typed or the specified timeout has been reached. Unless **noecho** has been set, the character will also be echoed into the designated window.

If the window is not a pad, and it has been moved or modified since the last call to **wrefresh**(3X), **wrefresh** will be called before another character is read.

If **keypad** is **TRUE**, and a function key is pressed, the token for that function key is returned instead of the raw characters. Possible function keys are defined in **<curses.h>** with integers beginning with **0401**, whose names begin with **KEY_.** If a character that could be the beginning of a function key (such as escape) is received, **curses**(3X) sets a timer. If the remainder of the sequence does not come in within the designated time, the character is passed through; otherwise, the function key value is returned. For this reason, many terminals experience a delay between the time a user presses the escape key and the escape is returned to the program.

The **ungetwch( )** routine places **wch** back onto the input queue to be returned by the next call to **wgetwch( )**.

Function Keys

The following function keys, defined in **<curses.h>**, might be returned by **getwch( )** if **keypad** has been enabled. Note that not all of these may be supported on a particular terminal if the terminal does not transmit a unique code when the key is pressed or if the definition for the key is not present in the **terminfo**(4) database.

Name	*Key name*
KEY_BREAK	Break key
KEY_DOWN	The four arrow keys ...
KEY_UP	
KEY_LEFT	
KEY_RIGHT	
KEY_HOME	Home key (upward+left arrow)
KEY_BACKSPACE	Backspace
KEY_F0	Function keys; space for 64 keys is reserved.
KEY_F(*n*)	For $0 \le n \le 63$
KEY_DL	Delete line
KEY_IL	Insert line
KEY_DC	Delete character
KEY_IC	Insert char or enter insert mode
KEY_EIC	Exit insert char mode
KEY_CLEAR	Clear screen
KEY_EOS	Clear to end of screen
KEY_EOL	Clear to end of line
KEY_SF	Scroll 1 line forward
KEY_SR	Scroll 1 line backward (reverse)
KEY_NPAGE	Next page
KEY_PPAGE	Previous page
KEY_STAB	Set tab
KEY_CTAB	Clear tab
KEY_CATAB	Clear all tabs
KEY_ENTER	Enter or send
KEY_SRESET	Soft (partial) reset
KEY_RESET	Reset or hard reset
KEY_PRINT	Print or copy
KEY_LL	Home down or bottom (lower left). Keypad is arranged like this: **A1 up A3** **left B2 right** **C1 down C3**
KEY_A1	Upper left of keypad
KEY_A3	Upper right of keypad
KEY_B2	Center of keypad
KEY_C1	Lower left of keypad
KEY_C3	Lower right of keypad
KEY_BTAB	Back tab key
KEY_BEG	Beg(inning) key
KEY_CANCEL	Cancel key
KEY_CLOSE	Close key
KEY_COMMAND	Cmd (command) key
KEY_COPY	Copy key
KEY_CREATE	Create key
KEY_END	End key
KEY_EXIT	Exit key
KEY_FIND	Find key

KEY_HELP	Help key
KEY_MARK	Mark key
KEY_MESSAGE	Message key
KEY_MOVE	Move key
KEY_NEXT	Next object key
KEY_OPEN	Open key
KEY_OPTIONS	Options key
KEY_PREVIOUS	Previous object key
KEY_REDO	Redo key
KEY_REFERENCE	Reference key
KEY_REFRESH	Refresh key
KEY_REPLACE	Replace key
KEY_RESTART	Restart key
KEY_RESUME	Resume key
KEY_SAVE	Save key
KEY_SBEG	Shifted beginning key
KEY_SCANCEL	Shifted cancel key
KEY_SCOMMAND	Shifted command key
KEY_SCOPY	Shifted copy key
KEY_SCREATE	Shifted create key
KEY_SDC	Shifted delete char key
KEY_SDL	Shifted delete line key
KEY_SELECT	Select key
KEY_SEND	Shifted end key
KEY_SEOL	Shifted clear line key
KEY_SEXIT	Shifted exit key
KEY_SFIND	Shifted find key
KEY_SHELP	Shifted help key
KEY_SHOME	Shifted home key
KEY_SIC	Shifted input key
KEY_SLEFT	Shifted left arrow key
KEY_SMESSAGE	Shifted message key
KEY_SMOVE	Shifted move key
KEY_SNEXT	Shifted next key
KEY_SOPTIONS	Shifted options key
KEY_SPREVIOUS	Shifted prev key
KEY_SPRINT	Shifted print key
KEY_SREDO	Shifted redo key
KEY_SREPLACE	Shifted replace key
KEY_SRIGHT	Shifted right arrow
KEY_SRSUME	Shifted resume key
KEY_SSAVE	Shifted save key
KEY_SSUSPEND	Shifted suspend key
KEY_SUNDO	Shifted undo key
KEY_SUSPEND	Suspend key
KEY_UNDO	Undo key

RETURN VALUE

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curses(3X), **curs_inopts**(3X), **curs_move**(3X), **wrefresh**(3X), **terminfo**(4)

NOTES

The header file **<curses.h>** automatically includes the header files **<stdio.h>** . **<unctrl.h>** and **<widec.h>**.

Use of the escape key by a programmer for a single character function is discouraged.

When using **getwch( )**, **wgetwch( )**, **mvgetwch( )**, or **mvwgetwch( )**, **nocbreak** mode and **echo** mode should not be used at the same time. Depending on the state of the tty driver when each character is typed, the program may produce undesirable results.

Note that **getwch( )**, **mvgetwch( )**, and **mvwgetwch( )** may be macros.

NAME

curs_getwstr, getwstr, getnwstr, wgetwstr, wgetnwstr, mvgetwstr, mvgetnwstr, mvwgetwstr, mvwgetnwstr – get wchar_t character strings from curses terminal keyboard

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]
#include <curses.h>

int getwstr(wchar_t ∗*wstr***);**

int getnwstr(wchar_t ∗*wstr***, int** *n***);**

int wgetwstr(WINDOW ∗*win***, wchar_t** ∗*wstr***);**

int wgetnwstr(WINDOW ∗*win***, wchar_t** ∗ *wstr***, int** *n***);**

int mvgetwstr(int *y***, int** *x***, wchar_t** ∗*wstr***);**

int mvgetnwstr(int *y***, int** *x***, wchar_t** ∗ *wstr***, int** *n***);**

int mvwgetwstr(WINDOW ∗*win***, int** *y***, int** *x***, wchar_t** ∗*wstr***);**

int mvwgetnwstr(WINDOW ∗*win***, int** *y***, int** *x*, **wchar_t** ∗*wstr***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

The effect of **getwstr( )** is as though a series of calls to **getwch**(3X) were made, until a newline and carriage return is received. The resulting value is placed in the area pointed to by the **wchar_t** pointer *wstr*. **getnwstr( )** reads at most *n* **wchar_t** characters, thus preventing a possible overflow of the input buffer. The user's erase and kill characters are interpreted, as well as any special keys (such as function keys, HOME key, CLEAR key, etc.).

RETURN VALUE

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curses(3X), **getwch**(3X)

NOTES

The header file **<curses.h>** automatically includes the header files **<stdio.h>**, **<unctrl.h>**, and **<widec.h>**.

Note that all routines except **wgetnwstr( )** may be macros.

NAME

curs_getyx, getyx, getparyx, getbegyx, getmaxyx – get curses cursor and window coordinates

SYNOPSIS

cc [*flag* ...] *file* ... –**lcurses** [*library* ..]

#include <curses.h>

void getyx(WINDOW ∗*win***, int** *y***, int** *x***);**

void getparyx(WINDOW ∗*win***, int** *y***, int** *x***);**

void getbegyx(WINDOW ∗*win***, int** *y***, int** *x***);**

void getmaxyx(WINDOW ∗*win***, int** *y***, int** *x***);**

MT-LEVEL

Unsafe

DESCRIPTION

With the **getyx( )** macro, the cursor position of the window is placed in the two integer variables *y* and *x*.

With the **getparyx( )** macro, if *win* is a subwindow, the beginning coordinates of the subwindow relative to the parent window are placed into two integer variables, *y* and *x*. Otherwise, **–1** is placed into *y* and *x*.

Like **getyx( ) ,** the **getbegyx( )** and **getmaxyx( )** macros store the current beginning coordinates and size of the specified window.

RETURN VALUES

The return values of these macros are undefined (that is, they should not be used as the right-hand side of assignment statements).

SEE ALSO

curses(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that all of these interfaces are macros and that ‘‘**&**’’ is not necessary before the variables *y* and *x*.

NAME

curs_inch, inch, winch, mvinch, mvwinch – get a character and its attributes from a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

chtype inch(void);

chtype winch(WINDOW ∗*win***);**

chtype mvinch(int *y***, int** *x***);**

chtype mvwinch(WINDOW ∗*win***, int** *y***, int** *x***);**

MT-LEVEL

Unsafe

DESCRIPTION

With these routines, the character, of type **chtype( )**, at the current position in the named window is returned. If any attributes are set for that position, their values are OR-ed into the value returned. Constants defined in **<curses.h>** can be used with the logical AND (**&**) operator to extract the character or attributes alone.

Attributes

The following bit-masks may be AND-ed with characters returned by **winch( )**.

A_CHARTEXT	Bit-mask to extract character
A_ATTRIBUTES	Bit-mask to extract attributes
A_COLOR	Bit-mask to extract color-pair field information

SEE ALSO

curses(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that all of these routines may be macros.

NAME

curs_inchstr, inchstr, inchnstr, winchstr, winchnstr, mvinchstr, mvinchnstr, mvwinchstr, mvwinchnstr – get a string of characters (and attributes) from a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int inchstr(chtype ∗*chstr***);**

int inchnstr(chtype ∗*chstr***, int** *n***);**

int winchstr(WINDOW ∗*win***, chtype** ∗*chstr***);**

int winchnstr(WINDOW ∗*win***, chtype** ∗*chstr***, int** *n***);**

int mvinchstr(int *y***, int** *x***, chtype** ∗*chstr***);**

int mvinchnstr(int *y***, int** *x***, chtype** ∗*chstr***, int** *n***);**

int mvwinchstr(WINDOW ∗*win***, int** *y***, int** *x***, chtype** ∗*chstr***);**

int mvwinchnstr(WINDOW ∗*win***, int** *y***, int** *x***, chtype** ∗*chstr***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

With these routines, a string of type **chtype( )**, starting at the current cursor position in the named window and ending at the right margin of the window, is returned. The four functions with *n* as the last argument, return the string at most *n* characters long. Constants defined in **<curses.h>** can be used with the **&** (logical AND) operator to extract the character or the attribute alone from any position in the *chstr* (see **curs_inch**(3X)).

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curs_inch(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that all routines except **winchnstr( )** may be macros.

NAME

curs_initscr, initscr, newterm, endwin, isendwin, set_term, delscreen – curses screen initialization and manipulation routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

WINDOW ∗**initscr(void);**

int endwin(void);

int isendwin(void);

SCREEN ∗**newterm(char** ∗*type***, FILE** ∗*outfd***, FILE** ∗*infd***);**

SCREEN ∗**set_term(SCREEN** ∗*new***);**

void delscreen(SCREEN∗ *sp***);**

MT-LEVEL

Unsafe

DESCRIPTION

initscr( ) is almost always the first routine that should be called (the exceptions are **slk_init( )**, **filter( )**, **ripoffline( )**, **use_env( )** and, for multiple-terminal applications, **newterm( )**.) This determines the terminal type and initializes all **curses** data structures. **initscr( )** also causes the first call to **refresh( )** to clear the screen. If errors occur, **initscr( )** writes an appropriate error message to standard error and exits; otherwise, a pointer is returned to **stdscr( )**. If the program needs an indication of error conditions, **newterm( )** should be used instead of **initscr( )**; **initscr( )** should only be called once per application.

A program that outputs to more than one terminal should use the **newterm( )** routine for each terminal instead of **initscr( )**. A program that needs an indication of error conditions, so it can continue to run in a line-oriented mode if the terminal cannot support a screen-oriented program, would also use this routine. The routine **newterm( )** should be called once for each terminal. It returns a variable of type **SCREEN** ∗ which should be saved as a reference to that terminal. The arguments are the *type* of the terminal to be used in place of **$TERM**, a file pointer for output to the terminal, and another file pointer for input from the terminal (if *type* is **NULL**, **$TERM** will be used). The program must also call **endwin( )** for each terminal being used before exiting from curses. If **newterm( )** is called more than once for the same terminal, the first terminal referred to must be the last one for which **endwin( )** is called.

A program should always call **endwin( )** before exiting or escaping from **curses** mode temporarily. This routine restores tty modes, moves the cursor to the lower left-hand corner of the screen and resets the terminal into the proper non-visual mode. Calling **refresh( )** or **doupdate( )** after a temporary escape causes the program to resume visual mode.

The **isendwin( )** routine returns **TRUE** if **endwin( )** has been called without any subsequent calls to **wrefresh( )**, and **FALSE** otherwise.

The **set_term( )** routine is used to switch between different terminals. The screen reference **new** becomes the new current terminal. The previous terminal is returned by the routine. This is the only routine which manipulates **SCREEN** pointers; all other routines

affect only the current terminal.

The **delscreen( )** routine frees storage associated with the **SCREEN** data structure. The **endwin( )** routine does not do this, so **delscreen( )** should be called after **endwin( )** if a particular **SCREEN** is no longer needed.

RETURN VALUES

endwin( ) returns the integer **ERR** upon failure and **OK** upon successful completion.

Routines that return pointers always return **NULL** on error.

SEE ALSO

curs_kernel(3X), **curs_refresh**(3X), **curs_slk**(3X), **curs_util**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **initscr( )** and **newterm( )** may be macros.

NAME

curs_inopts, cbreak, nocbreak, echo, noecho, halfdelay, intrflush, keypad, meta, nodelay, notimeout, raw, noraw, noqiflush, qiflush, timeout, wtimeout, typeahead – curses terminal input option control routines

SYNOPSIS

cc [*flag* ...] *file* ... **–lcurses** [*library* ..]

#include <curses.h>

int cbreak(void);

int nocbreak(void);

int echo(void);

int noecho(void);

int halfdelay(int *tenths***);**

int intrflush(WINDOW ∗*win***, bool** *bf***);**

int keypad(WINDOW ∗*win***, bool** *bf***);**

int meta(WINDOW ∗*win***, bool** *bf***);**

int nodelay(WINDOW ∗*win***, bool** *bf***);**

int notimeout(WINDOW ∗*win***, bool** *bf***);**

int raw(void);

int noraw(void);

void noqiflush(void);

void qiflush(void);

void timeout(int *delay***);**

void wtimeout(WINDOW ∗*win***, int** *delay***);**

int typeahead(int *fildes***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **cbreak( )** and **nocbreak( )** routines put the terminal into and out of **cbreak( )** mode, respectively. In this mode, characters typed by the user are immediately available to the program, and erase / kill character-processing is not performed. When out of this mode, the tty driver buffers the typed characters until a newline or carriage return is typed. Interrupt and flow control characters are unaffected by this mode. Initially the terminal may or may not be in **cbreak( )** mode, as the mode is inherited; therefore, a program should call **cbreak( )** or **nocbreak( )** explicitly. Most interactive programs using **curses** set the **cbreak( )** mode.

Note that **cbreak( )** overrides **raw( )**. (See **curs_getch**(3X) for a discussion of how these routines interact with **echo( )** and **noecho( )**.)

The **echo( )** and **noecho( )** routines control whether characters typed by the user are echoed by **getch( )** as they are typed. Echoing by the tty driver is always disabled, but initially **getch( )** is in echo mode, so characters typed are echoed. Authors of most

interactive programs prefer to do their own echoing in a controlled area of the screen, or not to echo at all, so they disable echoing by calling **noecho( )**. (See **curs_getch**(3X) for a discussion of how these routines interact with **cbreak( )** and **nocbreak( )**.)

The **halfdelay( )** routine is used for half-delay mode, which is similar to **cbreak( )** mode in that characters typed by the user are immediately available to the program. However, after blocking for *tenths* tenths of seconds, **ERR** is returned if nothing has been typed. The value of *tenths* must be a number between 1 and 255. Use **nocbreak( )** to leave half-delay mode.

If the **intrflush( )** option is enabled, (*bf* is **TRUE**), when an interrupt key is pressed on the keyboard (interrupt, break, quit) all output in the tty driver queue will be flushed, giving the effect of faster response to the interrupt, but causing **curses** to have the wrong idea of what is on the screen. Disabling (*bf* is **FALSE**), the option prevents the flush. The default for the option is inherited from the tty driver settings. The window argument is ignored.

The **keypad( )** option enables the keypad of the user's terminal. If enabled (*bf* is **TRUE**), the user can press a function key (such as an arrow key) and **wgetch( )** returns a single value representing the function key, as in **KEY_LEFT**. If disabled (*bf* is **FALSE**), **curses** does not treat function keys specially and the program has to interpret the escape sequences itself. If the keypad in the terminal can be turned on (made to transmit) and off (made to work locally), turning on this option causes the terminal keypad to be turned on when **wgetch( )** is called. The default value for keypad is false.

Initially, whether the terminal returns 7 or 8 significant bits on input depends on the control mode of the tty driver (see **termio**(7I)). To force 8 bits to be returned, invoke **meta**(*win*, **TRUE**). To force 7 bits to be returned, invoke **meta**(*win*, **FALSE**). The window argument, *win*, is always ignored. If the terminfo capabilities **smm** (meta_on) and **rmm** (meta_off) are defined for the terminal, **smm** is sent to the terminal when **meta**(*win*, **TRUE)** is called and **rmm** is sent when **meta**(*win*, **FALSE**) is called.

The **nodelay( )** option causes **getch( )** to be a non-blocking call. If no input is ready, **getch( )** returns **ERR**. If disabled (*bf* is **FALSE**), **getch( )** waits until a key is pressed.

While interpreting an input escape sequence, **wgetch( )** sets a timer while waiting for the next character. If **notimeout(***win*, **TRUE**) is called, then **wgetch( )** does not set a timer. The purpose of the timeout is to differentiate between sequences received from a function key and those typed by a user.

With the **raw( )** and **noraw( )** routines, the terminal is placed into or out of raw mode. Raw mode is similar to **cbreak( )** mode, in that characters typed are immediately passed through to the user program. The differences are that in raw mode, the interrupt, quit, suspend, and flow control characters are all passed through uninterpreted, instead of generating a signal. The behavior of the BREAK key depends on other bits in the tty driver that are not set by **curses**.

When the **noqiflush( )** routine is used, normal flush of input and output queues associated with the **INTR**, **QUIT** and **SUSP** characters will not be done (see **termio**(7I)). When **qiflush( )** is called, the queues will be flushed when these control characters are read.

The **timeout( )** and **wtimeout( )** routines set blocking or non-blocking read for a given window. If *delay* is negative, blocking read is used (that is, waits indefinitely for input). If *delay* is zero, then non-blocking read is used (that is, read returns **ERR** if no input is waiting). If *delay* is positive, then read blocks for *delay* milliseconds, and returns **ERR** if there is still no input. Hence, these routines provide the same functionality as **nodelay( )**, plus the additional capability of being able to block for only *delay* milliseconds (where *delay* is positive).

curses does ''line-breakout optimization'' by looking for typeahead periodically while updating the screen. If input is found, and it is coming from a tty, the current update is postponed until **refresh( )** or **doupdate( )** is called again. This allows faster response to commands typed in advance. Normally, the input FILE pointer passed to **newterm( )**, or **stdin** in the case that **initscr( )** was used, will be used to do this typeahead checking. The **typeahead( )** routine specifies that the file descriptor *fildes* is to be used to check for typeahead instead. If *fildes* is –1, then no typeahead checking is done.

RETURN VALUES

All routines that return an integer return **ERR** upon failure and an integer value other than **ERR** upon successful completion, unless otherwise noted in the preceding routine descriptions.

SEE ALSO

curs_getch(3X), **curs_initscr**(3X), **curses**(3X), **termio**(7I)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **echo( )**, **noecho( )**, **halfdelay( )**, **intrflush( )**, **meta( )**, **nodelay( )**, **notimeout( )**, **noqiflush( )**, **qiflush( )**, **timeout( )**, and **wtimeout( )** may be macros.

NAME

curs_insch, insch, winsch, mvinsch, mvwinsch – insert a character before the character under the cursor in a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int insch(chtype *ch***);**

int winsch(WINDOW ∗*win***, chtype** *ch***);**

int mvinsch(int *y***, int** *x***, chtype** *ch***);**

int mvwinsch(WINDOW ∗*win***, int** *y***, int** *x***, chtype** *ch***);**

MT-LEVEL

Unsafe

DESCRIPTION

With these routines, the character *ch* is inserted before the character under the cursor. All characters to the right of the cursor are moved one space to the right, with the possibility of the rightmost character on the line being lost. The cursor position does not change (after moving to *y*, *x*, if specified). (This does not imply use of the hardware insert character feature.)

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curses(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **insch( )**, **mvinsch( )**, and **mvwinsch( )** may be macros.

NAME

curs_insstr, insstr, insnstr, winsstr, winsnstr, mvinsstr, mvinsnstr, mvwinsstr, mvwinsnstr – insert string before character under the cursor in a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int insstr(char **str***);**

int insnstr(char **str***, int** *n***);**

int winsstr(WINDOW **win***, char** **str***);**

int winsnstr(WINDOW **win***, char** **str***, int** *n***);**

int mvinsstr(int *y***, int** *x***, char** **str***);**

int mvinsnstr(int *y***, int** *x***, char** **str***, int** *n***);**

int mvwinsstr(WINDOW **win***, int** *y***, int** *x***, char** **str***);**

int mvwinsnstr(WINDOW **win***, int** *y***, int** *x***, char** **str***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

With these routines, a character string (as many characters as will fit on the line) is inserted before the character under the cursor. All characters to the right of the cursor are moved to the right, with the possibility of the rightmost characters on the line being lost. The cursor position does not change (after moving to *y*, *x*, if specified). (This does not imply use of the hardware insert character feature.) The four routines with *n* as the last argument insert at most *n* characters. If *n*<=0, then the entire string is inserted.

If a character in *str* is a tab, newline, carriage return or backspace, the cursor is moved appropriately within the window. A newline also does a **clrtoeol( )** before moving. Tabs are considered to be at every eighth column. If a character in *str* is another control character, it is drawn in the ^*X* notation. Calling **winch( )** after adding a control character (and moving to it, if necessary) does not return the control character, but instead returns the representation of the control character.

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curs_clear(3X), **curs_inch**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that all but **winsnstr( )** may be macros.

NAME

curs_instr, instr, innstr, winstr, winnstr, mvinstr, mvinnstr, mvwinstr, mvwinnstr – get a string of characters from a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int instr(char **str****);**

int innstr(char **str***, int** *n***);**

int winstr(WINDOW **win***, char** **str***);**

int winnstr(WINDOW **win***, char** **str***, int** *n***);**

int mvinstr(int *y***, int** *x***, char** **str***);**

int mvinnstr(int *y***, int** *x***, char** **str***, int** *n***);**

int mvwinstr(WINDOW **win***, int** *y***, int** *x***, char** **str***);**

int mvwinnstr(WINDOW **win***, int** *y***, int** *x***, char** **str***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

These routines return a string of characters in *str*, starting at the current cursor position in the named window and ending at the right margin of the window. Attributes are stripped from the characters. The four functions with *n* as the last argument return the string at most *n* characters long.

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curses(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that all routines except **winnstr( )** may be macros.

NAME

curs_inswch, inswch, winswch, mvinswch, mvwinswch – insert a wchar_t character before the character under the cursor in a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]
#include <curses.h>

int inswch(chtype *wch***);**

int winswch(WINDOW ∗*win***, chtype** *wch***);**

int mvinswch(int *y***, int** *x***, chtype** *wch***);**

int mvwinswch(WINDOW ∗*win***, int** *y***, int** *x***, chtype** *wch***);**

MT-LEVEL

Unsafe

DESCRIPTION

These routines insert the character *wch*, holding a **wchar_t** character, before the character under the cursor. All characters to the right of the cursor are moved one space to the right, with the possibility of the rightmost character on the line being lost. The cursor position does not change (after moving to *y*, *x*, if specified). (This does not imply use of the hardware insert character feature.)

RETURN VALUE

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curses(3X)

NOTES

The header file **<curses.h>** automatically includes the header files **<stdio.h>**, **<unctrl.h>** and **<widec.h>**.

Note that **inswch( )**, **mvinswch( )**, and **mvwinswch( )** may be macros.

None of these routines can use the color attribute in **chtype**.

NAME

curs_inswstr, inswstr, insnwstr, winswstr, winsnwstr, mvinswstr, mvinsnwstr, mvwinswstr, mvwinsnwstr – insert wchar_t string before character under the cursor in a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]
#include <curses.h>

int inswstr(wchar_t **wstr***);**

int insnwstr(wchar_t **wstr***, int** *n***);**

int winswstr(WINDOW **win***, wchar_t** **wstr***);**

int winsnwstr(WINDOW **win***, wchar_t** **wstr,* **int** *n***);**

int mvinswstr(int *y***, int** *x,* **wchar_t** ** wstr***);**

int mvinsnwstr(int *y***, int** *x***, wchar_t** **wstr,* **int** *n***);**

int mvwinswstr(WINDOW **win***, int** *y***, int** *x,* **wchar_t** **wstr***);**

int mvwinsnwstr(WINDOW **win***, int** *y***, int** *x,* **wchar_t** **wstr***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

These routines insert a **wchar_t** character string (as many **wchar_t** characters as will fit on the line) before the character under the cursor. All characters to the right of the cursor are moved to the right, with the possibility of the rightmost characters on the line being lost. The cursor position does not change (after moving to *y*, *x*, if specified). (This does not imply use of the hardware insert character feature.) The four routines with *n* as the last argument insert at most *n* **wchar_t** characters. If *n*<=0, then the entire string is inserted.

If a character in *wstr* is a tab, newline, carriage return, or backspace, the cursor is moved appropriately within the window. A newline also does a **clrtoeol**(3X) before moving. Tabs are considered to be at every eighth column. If a character in *wstr* is another control character, it is drawn in the ^*X* notation. Calling **winwch**(3X) after adding a control character (and moving to it, if necessary) does not return the control character, but instead returns the representation of the control character.

RETURN VALUE

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

clrtoeol(3X), **curses**(3X), **winwch**(3X)

NOTES

The header file **<curses.h>** automatically includes the header files **<stdio.h>**, **<unctrl.h>** and **<widec.h>**.

Note that all but **winsnwstr( )** may be macros.

NAME

curs_inwch, inwch, winwch, mvinwch, mvwinwch – get a wchar_t character and its attributes from a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]
#include <curses.h>

chtype inwch(void);

chtype winwch(WINDOW ∗*win***);**

chtype mvinwch(int *y***, int** *x***);**

chtype mvwinwch(WINDOW ∗*win***, int** *y* **, int** *x***);**

MT-LEVEL

Unsafe

DESCRIPTION

These routines return the **wchar_t** character, of type **chtype**, at the current position in the named window. If any attributes are set for that position, their values are OR-ed into the value returned. Constants defined in **<curses.h>** can be used with the logical AND (**&**) operator to extract the character or attributes alone.

Attributes

The following bit-masks may be AND-ed with characters returned by **winwch( )**.

A_WCHARTEXT	Bit-mask to extract character
A_WATTRIBUTES	Bit-mask to extract attributes

SEE ALSO

curses(3X)

NOTES

The header file **<curses.h>** automatically includes the header files **<stdio.h>**, **<unctrl.h>**, and **<widec.h>**.

Note that all of these routines may be macros.

None of these routines can use the color attribute in **chtype**.

NAME

curs_inwchstr, inwchstr, inwchnstr, winwchstr, winwchnstr, mvinwchstr, mvinwchnstr, mvwinwchstr, mvwinwchnstr – get a string of wchar_t characters (and attributes) from a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]
#include <curses.h>

int inwchstr(chtype ∗*wchstr***);**

int inwchnstr(chtype ∗*wchstr***, int** *n***);**

int winwchstr(WINDOW ∗*win***, chtype** ∗*wchstr***);**

int winwchnstr(WINDOW ∗*win***, chtype** ∗*wchstr***, int** *n***);**

int mvinwchstr(int *y***, int** *x***, chtype** ∗*wchstr***);**

int mvinwchnstr(int *y***, int** *x***, chtype** ∗*wchstr***, int** *n***);**

int mvwinwchstr(WINDOW ∗*win***, int** *y***, int** *x***, chtype** ∗*wchstr***);**

int mvwinwchnstr(WINDOW ∗*win***, int** *y***, int** *x***, chtype** ∗*wchstr***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

These routines return a string of type **chtype**, holding **wchar_t** characters, starting at the current cursor position in the named window and ending at the right margin of the window. The four functions with *n* as the last argument, return the string at most *n* **wchar_t** characters long. Constants defined in **<curses.h>** can be used with the logical AND (**&**) operator to extract the **wchar_t** character or the attribute alone from any position in the *wchstr* (see **curs_inwch**(3X)).

RETURN VALUE

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curses(3X), **curs_inwch**(3X)

NOTES

The header file **<curses.h>** automatically includes the header files **<stdio.h>**, **<unctrl.h>**, and **<widec.h>**.

Note that all routines except **winwchnstr( )** may be macros.

None of these routines can use the color attribute in **chtype**.

NAME

curs_inwstr, inwstr, innwstr, winwstr, winnwstr, mvinwstr, mvinnwstr, mvwinwstr, mvwinnwstr – get a string of wchar_t characters from a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]
#include <curses.h>

int inwstr(wchar_t ∗*wstr***);**

int innwstr(wchar_t ∗*wstr***, int** *n***);**

int winwstr(WINDOW ∗*win***, wchar_t** ∗*wstr***);**

int winnwstr(WINDOW ∗*win***, wchar_t** ∗*wstr***, int** *n***);**

int mvinwstr(int *y***, int** *x***, wchar_t** ∗*wstr***);**

int mvinnwstr(int *y***, int** *x***, wchar_t** ∗*wstr***, int** *n***);**

int mvwinwstr(WINDOW ∗*win***, int** *y***, int** *x***, wchar_t** ∗*wstr***);**

int mvwinnwstr(WINDOW ∗*win***, int** *y***, int** *x***, wchar_t** ∗*wstr***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

These routines return the string of **wchar_t** characters in *wstr* starting at the current cursor position in the named window and ending at the right margin of the window. Attributes are stripped from the characters. The four functions with *n* as the last argument return the string at most *n* **wchar_t** characters long.

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curses(3X)

NOTES

The header file **<curses.h>** automatically includes the header files **<stdio.h>**, **<unctrl.h>** and **<widec.h>**.

Note that all routines except **winnwstr( )** may be macros.

NAME

curs_kernel, def_prog_mode, def_shell_mode, reset_prog_mode, reset_shell_mode, resetty, savetty, getsyx, setsyx, ripoffline, curs_set, napms – low-level curses routines

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lcurses** [*library* . .]

#include <curses.h>

int def_prog_mode(void);

int def_shell_mode(void);

int reset_prog_mode(void);

int reset_shell_mode(void);

int resetty(void);

int savetty(void);

int getsyx(int *y***, int** *x***);**

int setsyx(int *y***, int** *x***);**

int ripoffline(int *line***, int (*****init)(WINDOW** ***, int));**

int curs_set(int *visibility***);**

int napms(int *ms***);**

MT-LEVEL

Unsafe

DESCRIPTION

The following routines give low-level access to various **curses** functionality. Theses routines typically are used inside library routines.

The **def_prog_mode( )** and **def_shell_mode( )** routines save the current terminal modes as the ‘‘program’’ (in **curses**) or ‘‘shell’’ (not in **curses)** state for use by the **reset_prog_mode( )** and **reset_shell_mode( )** routines. This is done automatically by **initscr( )**.

The **reset_prog_mode( )** and **reset_shell_mode( )** routines restore the terminal to ‘‘program’’ (in **curses**) or ‘‘shell’’ (out of **curses**) state. These are done automatically by **endwin( )** and, after an **endwin( )**, by **doupdate( )**, so they normally are not called.

The **resetty( )** and **savetty( )** routines save and restore the state of the terminal modes. **savetty( )** saves the current state in a buffer and **resetty( )** restores the state to what it was at the last call to **savetty( )**.

With the **getsyx( )** routine, the current coordinates of the virtual screen cursor are returned in *y* and *x*. If **leaveok( )** is currently **TRUE**, then **–1**,**–1** is returned. If lines have been removed from the top of the screen, using **ripoffline( )**, *y* and *x* include these lines; therefore, *y* and *x* should be used only as arguments for **setsyx( )**.

With the **setsyx( )** routine, the virtual screen cursor is set to *y*, *x*. If *y* and *x* are both **–1**, then **leaveok( )** is set. The two routines **getsyx( )** and **setsyx( )** are designed to be used by a library routine, which manipulates **curses** windows but does not want to change the current position of the program’s cursor. The library routine would call **getsyx( )** at the beginning, do its manipulation of its own windows, do a **wnoutrefresh( )** on its windows,

call **setsyx( )**, and then call **doupdate( )**.

The **ripoffline( )** routine provides access to the same facility that **slk_init( )** (see **curs_slk**(3X)) uses to reduce the size of the screen. **ripoffline( )** must be called before **initscr( )** or **newterm( )** is called. If *line* is positive, a line is removed from the top of **stdscr( )**; if *line* is negative, a line is removed from the bottom. When this is done inside **initscr( )**, the routine **init( )** (supplied by the user) is called with two arguments: a window pointer to the one-line window that has been allocated and an integer with the number of columns in the window. Inside this initialization routine, the integer variables **LINES** and **COLS** (defined in **<curses.h>**) are not guaranteed to be accurate and **wrefresh( )** or **doupdate( )** must not be called. It is allowable to call **wnoutrefresh( )** during the initialization routine.

ripoffline( ) can be called up to five times before calling **initscr( )** or **newterm( )**.

With the **curs_set( )** routine, the cursor state is set to invisible, normal, or very visible for *visibility* equal to **0**, **1**, or **2** respectively. If the terminal supports the *visibility* requested, the previous *cursor* state is returned; otherwise, **ERR** is returned.

The **napms( )** routine is used to sleep for *ms* milliseconds.

RETURN VALUES

Except for **curs_set( )**, these routines always return **OK**. **curs_set( )** returns the previous cursor state, or **ERR** if the requested *visibility* is not supported.

SEE ALSO

curs_initscr(3X), **curs_outopts**(3X), **curs_refresh**(3X), **curs_scr_dump**(3X), **curs_slk**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **getsyx( )** is a macro, so an ampersand (**&**) is not necessary before the variables *y* and *x*.

NAME

curs_move, move, wmove – move curses window cursor

SYNOPSIS

```
#include <curses.h>
int move(int y, int x);
int wmove(WINDOW *win, int y, int x);
```

MT-LEVEL

Unsafe

DESCRIPTION

With these routines, the cursor associated with the window is moved to line *y* and column *x*. This routine does not move the physical cursor of the terminal until **refresh()** is called. The position specified is relative to the upper left-hand corner of the window, which is (0,0).

RETURN VALUES

These routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curs_refresh(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **move()** may be a macro.

NAME

curs_outopts, clearok, idlok, idcok, immedok, leaveok, setscrreg, wsetscrreg, scrollok, nl, nonl – curses terminal output option control routines

SYNOPSIS

cc [*flag* ...] *file* ... –**lcurses** [*library* ..]

#include <curses.h>

int clearok(WINDOW ∗*win***, bool** *bf***);**

int idlok(WINDOW ∗*win***, bool** *bf***);**

void idcok(WINDOW ∗*win***, bool** *bf***);**

void immedok(WINDOW ∗*win***, bool** *bf***);**

int leaveok(WINDOW ∗*win***, bool** *bf***);**

int setscrreg(int *top***, int** *bot***);**

int wsetscrreg(WINDOW ∗*win***, int** *top***, int** *bot***);**

int scrollok(WINDOW ∗*win***, bool** *bf***);**

int nl(void);

int nonl(void);

MT-LEVEL

Unsafe

DESCRIPTION

These routines set options that deal with output within **curses**. All options are initially **FALSE**, unless otherwise stated. It is not necessary to turn these options off before calling **endwin( )**.

With the **clearok( )** routine, if enabled (*bf* is **TRUE**), the next call to **wrefresh( )** with this window will clear the screen completely and redraw the entire screen from scratch. This is useful when the contents of the screen are uncertain, or in some cases for a more pleasing visual effect. If the *win* argument to **clearok( )** is the global variable **curscr( )**, the next call to **wrefresh( )** with any window causes the screen to be cleared and repainted from scratch.

With the **idlok( )** routine, if enabled (*bf* is **TRUE**), **curses** considers using the hardware insert / delete line feature of terminals so equipped. If disabled (*bf* is **FALSE) , curses** very seldom uses this feature. (The insert / delete character feature is always considered.) This option should be enabled only if the application needs insert / delete line, for example, for a screen editor. It is disabled by default because insert / delete line tends to be visually annoying when used in applications where it isn't really needed. If insert / delete line cannot be used, **curses** redraws the changed portions of all lines.

With the **idcok( )** routine, if enabled (*bf* is **TRUE**), **curses** considers using the hardware insert / delete character feature of terminals so equipped. This is enabled by default.

With the **immedok( )** routine, if enabled (*bf* is **TRUE**), any change in the window image, such as the ones caused by **waddch( )**, **wclrtobot( )**, **wscrl( )**, etc., automatically cause a call to **wrefresh( )**. However, it may degrade the performance considerably, due to repeated calls to **wrefresh( )**. It is disabled by default. Normally, the hardware cursor is

left at the location of the window cursor being refreshed. The **leaveok( )** option allows the cursor to be left wherever the update happens to leave it. It is useful for applications where the cursor is not used, since it reduces the need for cursor motions. If possible, the cursor is made invisible when this option is enabled.

The **setscrreg( )** and **wsetscrreg( )** routines allow the application programmer to set a software scrolling region in a window. *top* and *bot* are the line numbers of the top and bottom margin of the scrolling region. (Line 0 is the top line of the window.) If this option and **scrollok( )** are enabled, an attempt to move off the bottom margin line causes all lines in the scrolling region to scroll up one line. Only the text of the window is scrolled. (Note that this has nothing to do with the use of a physical scrolling region capability in the terminal, like that in the VT100. If **idlok( )** is enabled and the terminal has either a scrolling region or insert / delete line capability, they will probably be used by the output routines.)

The **scrollok( )** option controls what happens when the cursor of a window is moved off the edge of the window or scrolling region, either as a result of a newline action on the bottom line, or typing the last character of the last line. If disabled, (*bf* is **FALSE**), the cursor is left on the bottom line. If enabled, (*bf* is **TRUE**), **wrefresh( )** is called on the window, and the physical terminal and window are scrolled up one line. (Note that in order to get the physical scrolling effect on the terminal, it is also necessary to call **idlok( )**.)

The **nl( )** and **nonl( )** routines control whether newline is translated into carriage return and linefeed on output, and whether return is translated into newline on input. Initially, the translations do occur. By disabling these translations using **nonl( )**, **curses** is able to make better use of the linefeed capability, resulting in faster cursor motion.

RETURN VALUES

setscrreg( ) and **wsetscrreg( )** return **OK** upon success and **ERR** upon failure. All other routines that return an integer always return **OK**.

SEE ALSO

curs_addch(3X), **curs_clear**(3X), **curs_initscr**(3X), **curs_refresh**(3X), **curs_scroll**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **clearok( )**, **leaveok( )**, **scrollok( )**, **idcok( )**, **nl( )**, **nonl( )**, and **setscrreg( )** may be macros.

The **immedok( )** routine is useful for windows that are used as terminal emulators.

NAME

curs_overlay, overlay, overwrite, copywin – overlap and manipulate overlapped curses windows

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int overlay(WINDOW **srcwin***, WINDOW** **dstwin***);**

int overwrite(WINDOW **srcwin***, WINDOW** **dstwin***);**

int copywin(WINDOW **srcwin***, WINDOW** **dstwin***, int** *sminrow***, int** *smincol***, int** *dminrow***, int** *dmincol***, int** *dmaxrow***, int** *dmaxcol***, int** *overlay***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **overlay( )** and **overwrite( )** routines overlay *srcwin* on top of *dstwin*. *scrwin* and *dstwin* are not required to be the same size; only text where the two windows overlap is copied. The difference is that **overlay( )** is non-destructive (blanks are not copied) whereas **overwrite( )** is destructive.

The **copywin( )** routine provides a finer granularity of control over the **overlay( )** and **overwrite( )** routines. Like in the **prefresh( )** routine, a rectangle is specified in the destination window, (*dminrow*, *dmincol*) and (*dmaxrow*, *dmaxcol*), and the upper-left-corner coordinates of the source window, (*sminrow*, *smincol*). If the argument *overlay* is **true**, then copying is non-destructive, as in **overlay( )**.

RETURN VALUES

Routines that return an integer return **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curs_pad(3X), **curs_refresh**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **overlay( )** and **overwrite** may be macros.

NAME

curs_pad, newpad, subpad, prefresh, pnoutrefresh, pechochar, pechowchar – create and display curses pads

SYNOPSIS

cc [*flag* ...] *file* ... –**lcurses** [*library* ..]
#include <curses.h>

WINDOW ∗**newpad(int** *nlines***, int** *ncols***);**

WINDOW ∗**subpad(WINDOW** ∗*orig***, int** *nlines***, int** *ncols***, int** *begin_y***, int** *begin_x***);**

int prefresh(WINDOW ∗*pad***, int** *pminrow***, int** *pmincol***, int** *sminrow***, int** *smincol***, int** *smaxrow***, int** *smaxcol***);**

int pnoutrefresh(WINDOW ∗*pad***, int** *pminrow***, int** *pmincol***, int** *sminrow***, int** *smincol***, int** *smaxrow***, int** *smaxcol***);**

int pechochar(WINDOW ∗*pad***, chtype** *ch***);**

int pechowchar(WINDOW ∗*pad***, chtype** *wch***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **newpad( )** routine creates and returns a pointer to a new pad data structure with the given number of lines, *nlines*, and columns, *ncols*. A pad is like a window, except that it is not restricted by the screen size, and is not necessarily associated with a particular part of the screen. Pads can be used when a large window is needed, and only a part of the window will be on the screen at one time. Automatic refreshes of pads (for example, from scrolling or echoing of input) do not occur. It is not legal to call **wrefresh**(3X) with a *pad* as an argument; the routines **prefresh( )** or **pnoutrefresh( )** should be called instead. Note that these routines require additional parameters to specify the part of the pad to be displayed and the location on the screen to be used for the display.

The **subpad( )** routine creates and returns a pointer to a subwindow within a pad with the given number of lines, *nlines*, and columns, *ncols*. Unlike **subwin**(3X), which uses screen coordinates, the window is at position (*begin_x*, *begin_y*) on the pad. The window is made in the middle of the window *orig*, so that changes made to one window affect both windows. During the use of this routine, it will often be necessary to call **touchwin**(3X) or **touchline**(3X) on *orig* before calling **prefresh( )**.

The **prefresh( )** and **pnoutrefresh( )** routines are analogous to **wrefresh**(3X) and **wnoutrefresh**(3X) except that they relate to pads instead of windows. The additional parameters are needed to indicate what part of the pad and screen are involved. *pminrow* and *pmincol* specify the upper left-hand corner of the rectangle to be displayed in the pad. *sminrow*, *smincol*, *smaxrow*, and *smaxcol* specify the edges of the rectangle to be displayed on the screen. The lower right-hand corner of the rectangle to be displayed in the pad is calculated from the screen coordinates, since the rectangles must be the same size. Both rectangles must be entirely contained within their respective structures. Negative values of *pminrow*, *pmincol*, *sminrow*, or *smincol* are treated as if they were zero.

The **pechochar( )** routine is functionally equivalent to a call to **addch**(3X) followed by a call to **refresh**(3X), a call to **waddch**(3X) followed by a call to **wrefresh**(3X), or a call to **waddch**(3X) followed by a call to **prefresh( )**. The knowledge that only a single character is being output is taken into consideration and, for non-control characters, a considerable performance gain might be seen by using these routines instead of their equivalents. In the case of **pechochar( )**, the last location of the pad on the screen is reused for the arguments to **prefresh( )**.

RETURN VALUES

Routines that return an integer return **ERR** upon failure and an integer value other than **ERR** upon successful completion.

Routines that return pointers return **NULL** on error.

SEE ALSO

addch(3X), **curses**(3X), **refresh**(3X), **subwin**(3X), **touchline**(3X), **touchwin**(3X), **waddch**(3X), **wnoutrefresh**(3X), **wrefresh**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>**, **<unctrl.h>** and **<widec.h>**.

Note that **pechochar( )** may be a macro.

NAME

curs_printw, printw, wprintw, mvprintw, mvwprintw, vwprintw – print formatted output in curses windows

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int printw(char ∗*fmt*, /∗ *arg* ∗/ **. . .);**

int wprintw(WINDOW ∗*win*, **char** ∗*fmt*, /∗ *arg* ∗/ **. . .);**

int mvprintw(int *y*, **int** *x*, **char** ∗*fmt*, /∗ *arg* ∗/ **. . .);**

int mvwprintw(WINDOW ∗*win*, **int** *y*, **int** *x*, **char** ∗*fmt*, /∗ *arg* ∗/ **. . .);**

#include <varargs.h>

int vwprintw(WINDOW ∗*win*, **char** ∗*fmt*, /∗ *varglist* ∗/ **. . .);**

MT-LEVEL

Unsafe

DESCRIPTION

The **printw( )**, **wprintw( )**, **mvprintw( )**, and **mvwprintw( )** routines are analogous to **printf( )** (see **printf**(3S)). In effect, the string that would be output by **printf( )** is output instead as though **waddstr( )** were used on the given window.

The **vwprintw( )** routine is analogous to **vprintf( )** (see **vprintf**(3S)) and performs a **wprintw( )** using a variable argument list. The third argument is a **va_list**, a pointer to a list of arguments, as defined in **<varargs.h>**.

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curses(3X), **printf**(3S), **vprintf**(3S)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

NAME

curs_refresh, refresh, wrefresh, wnoutrefresh, doupdate, redrawwin, wredrawln – refresh curses windows and lines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int refresh(void);

int wrefresh(WINDOW ∗*win***);**

int wnoutrefresh(WINDOW ∗*win***);**

int doupdate(void);

int redrawwin(WINDOW ∗*win***);**

int wredrawln(WINDOW ∗*win***, int** *beg_line***, int** *num_lines***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **refresh( )** and **wrefresh( )** routines (or **wnoutrefresh( )** and **doupdate( )**) must be called to get any output on the terminal, as other routines merely manipulate data structures. The routine **wrefresh( )** copies the named window to the physical terminal screen, taking into account what is already there in order to do optimizations. The **refresh( )** routine is the same, using **stdscr** as the default window. Unless **leaveok( )** has been enabled, the physical cursor of the terminal is left at the location of the cursor for that window.

The **wnoutrefresh( )** and **doupdate( )** routines allow multiple updates with more efficiency than **wrefresh( )** alone. In addition to all the window structures, **curses** keeps two data structures representing the terminal screen: a physical screen, describing what is actually on the screen, and a virtual screen, describing what the programmer wants to have on the screen.

The routine **wrefresh( )** works by first calling **wnoutrefresh( )**, which copies the named window to the virtual screen, and then calling **doupdate( )**, which compares the virtual screen to the physical screen and does the actual update. If the programmer wishes to output several windows at once, a series of calls to **wrefresh( )** results in alternating calls to **wnoutrefresh( )** and **doupdate( )**, causing several bursts of output to the screen. By first calling **wnoutrefresh( )** for each window, it is then possible to call **doupdate( )** once, resulting in only one burst of output, with fewer total characters transmitted and less CPU time used. If the *win* argument to **wrefresh( )** is the global variable **curscr**, the screen is immediately cleared and repainted from scratch.

The **redrawwin( )** routine indicates to **curses** that some screen lines are corrupted and should be thrown away before anything is written over them. These routines could be used for programs such as editors, which want a command to redraw some part of the screen or the entire screen. The routine **redrawln( )** is preferred over **redrawwin( )** where a noisy communication line exists and redrawing the entire window could be subject to even more communication noise. Just redrawing several lines offers the possibility that they would show up unblemished.

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curs_outopts(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **refresh( )** and **redrawwin( )** may be macros.

NAME

curs_scanw, scanw, wscanw, mvscanw, mvwscanw, vwscanw – convert formatted input from a curses widow

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int scanw(char ∗*fmt*, /∗ *arg* ∗/ **. . .);**

int wscanw(WINDOW ∗*win*, **char** ∗*fmt*, /∗ *arg* ∗/ **. . .);**

int mvscanw(int *y*, **int** *x*, **char** ∗*fmt*, /∗ *arg* ∗/ **. . .);**

int mvwscanw(WINDOW ∗**win, int** *y*, **int** *x*, **char** ∗*fmt*, /∗ *arg* ∗/ **. . .);**

int vwscanw(WINDOW ∗*win*, **char** ∗*fmt*, **va_list** *varglist***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **scanw( )**, **wscanw( )**, and **mvscanw( )** routines correspond to **scanf( )** (see **scanf**(3S)). The effect of these routines is as though **wgetstr( )** were called on the window, and the resulting line used as input for the scan. Fields which do not map to a variable in the *fmt* field are lost.

The **vwscanw( )** routine is similar to **vwprintw( ) in that** it performs a **wscanw( )** using a variable argument list. The third argument is a *va_list*, a pointer to a list of arguments, as defined in **<varargs.h>**.

RETURN VALUES

vwscanw( ) returns **ERR** on failure and an integer equal to the number of fields scanned on success.

Applications may interrogate the return value from the **scanw**, **wscanw( )**, **mvscanw( )**, and **mvwscanw( )** routines to determine the number of fields which were mapped in the call.

SEE ALSO

curs_getstr(3X), **curs_printw**(3X), **curses**(3X), **scanf**(3S)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

NAME

curs_scr_dump, scr_dump, scr_restore, scr_init, scr_set – read (write) a curses screen from (to) a file

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int scr_dump(char ∗*filename***);**

int scr_restore(char ∗*filename***);**

int scr_init(char ∗*filename***);**

int scr_set(char ∗*filename***);**

MT-LEVEL

Unsafe

DESCRIPTION

With the **scr_dump( )** routine, the current contents of the virtual screen are written to the file *filename*.

With the **scr_restore( )** routine, the virtual screen is set to the contents of *filename*, which must have been written using **scr_dump( )**. The next call to **doupdate( )** restores the screen to the way it looked in the dump file.

With the **scr_init( )** routine, the contents of *filename* are read in and used to initialize the **curses** data structures about what the terminal currently has on its screen. If the data is determined to be valid, **curses** bases its next update of the screen on this information rather than clearing the screen and starting from scratch. **scr_init( )** is used after **initscr( )** or a **system**(3S) call to share the screen with another process which has done a **scr_dump( )** after its **endwin( )** call. The data is declared invalid if the time-stamp of the tty is old or the terminfo capabilities **rmcup( )** and **nrrmc( )** exist.

The **scr_set( )** routine is a combination of **scr_restore( )** and **scr_init( )**. It tells the program that the information in *filename* is what is currently on the screen, and also what the program wants on the screen. This can be thought of as a screen inheritance function.

To read (write) a window from (to) a file, use the **getwin( )** and **putwin( )** routines (see **curs_util**(3X)).

RETURN VALUES

All routines return the integer **ERR** upon failure and **OK** upon success.

SEE ALSO

curs_initscr(3X), **curs_refresh**(3X), **curs_util**(3X), **curses**(3X), **system**(3S)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **scr_init( )**, **scr_set( )**, and **scr_restore( )** may be macros.

NAME

curs_scroll, scroll, scrl, wscrl – scroll a curses window

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lcurses** [*library* . .]

#include <curses.h>

int scroll(WINDOW ∗*win***);**

int scrl(int *n***);**

int wscrl(WINDOW ∗*win***, int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

With the **scroll( )** routine, the window is scrolled up one line. This involves moving the lines in the window data structure. As an optimization, if the scrolling region of the window is the entire screen, the physical screen is scrolled at the same time.

With the **scrl( )** and **wscrl( )** routines, for positive *n* scroll the window up *n* lines (line *i*+*n* becomes *i*); otherwise scroll the window down *n* lines. This involves moving the lines in the window character image structure. The current cursor position is not changed.

For these functions to work, scrolling must be enabled via **scrollok( )**.

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curs_outopts(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **scrl( )** and **scroll( )** may be macros.

NAME

curs_slk, slk_init, slk_set, slk_refresh, slk_noutrefresh, slk_label, slk_clear, slk_restore, slk_touch, slk_attron, slk_attrset, slk_attroff – curses soft label routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int slk_init(int *fmt***);**

int slk_set(int *labnum***, char** ∗*label***, int** *fmt***);**

int slk_refresh(void);

int slk_noutrefresh(void);

char ∗**slk_label(int** *labnum***);**

int slk_clear(void);

int slk_restore(void);

int slk_touch(void);

int slk_attron(chtype *attrs***);**

int slk_attrset(chtype *attrs***);**

int slk_attroff(chtype *attrs***);**

MT-LEVEL

Unsafe

DESCRIPTION

curses manipulates the set of soft function-key labels that exist on many terminals. For those terminals that do not have soft labels, **curses** takes over the bottom line of **stdscr**, reducing the size of **stdscr** and the variable **LINES**. **curses** standardizes on eight labels of up to eight characters each.

To use soft labels, the **slk_init( )** routine must be called before **initscr( )** or **newterm( )** is called. If **initscr( )** eventually uses a line from **stdscr** to emulate the soft labels, then *fmt* determines how the labels are arranged on the screen. Setting *fmt* to **0** indicates a 3-2-3 arrangement of the labels; **1** indicates a 4-4 arrangement.

With the **slk_set( )** routine, *labnum* is the label number, from **1** to **8**. *label* is the string to be put on the label, up to eight characters in length. A null string or a null pointer sets up a blank label. *fmt* is either **0**, **1**, or **2**, indicating whether the label is to be left-justified, centered, or right-justified, respectively, within the label.

The **slk_refresh( )** and **slk_noutrefresh( )** routines correspond to the **wrefresh( )** and **wnoutrefresh( )** routines.

With the **slk_label( )** routine, the current label for label number *labnum* is returned with leading and trailing blanks stripped.

With the **slk_clear( )** routine, the soft labels are cleared from the screen.

With the **slk_restore( )** routine, the soft labels are restored to the screen after a **slk_clear( )** is performed.

With the **slk_touch( )** routine, all the soft labels are forced to be output the next time a **slk_noutrefresh( )** is performed.

The **slk_attron( )**, **slk_attrset( )**, and **slk_attroff( )** routines correspond to **attron( )**, **attrset( )**, and **attroff( )**. They have an effect only if soft labels are simulated on the bottom line of the screen.

RETURN VALUES

Routines that return an integer return **ERR** upon failure and an integer value other than **ERR** upon successful completion.

slk_label( ) returns **NULL** on error.

SEE ALSO

curs_attr(3X), **curs_initscr**(3X), **curs_refresh**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Most applications would use **slk_noutrefresh( )** because a **wrefresh( )** is likely to follow soon.

NAME

curs_termattrs, baudrate, erasechar, has_ic, has_il, killchar, longname, termattrs, termname – curses environment query routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int baudrate(void);

char erasechar(void);

int has_ic(void);

int has_il(void);

char killchar(void);

char *longname(void);

chtype termattrs(void);

char *termname(void);

MT-LEVEL

Unsafe

DESCRIPTION

The **baudrate( )** routine returns the output speed of the terminal. The number returned is in bits per second, for example **9600**, and is an integer.

With the **erasechar( )** routine, the user's current erase character is returned.

The **has_ic( )** routine is true if the terminal has insert- and delete-character capabilities.

The **has_il( )** routine is true if the terminal has insert- and delete-line capabilities, or can simulate them using scrolling regions. This might be used to determine if it would be appropriate to turn on physical scrolling using **scrollok( )**.

With the **killchar( )** routine, the user's current line kill character is returned.

The **longname( )** routine returns a pointer to a static area containing a verbose description of the current terminal. The maximum length of a verbose description is 128 characters. It is defined only after the call to **initscr( )** or **newterm( )**. The area is overwritten by each call to **newterm( )** and is not restored by **set_term( )**, so the value should be saved between calls to **newterm( )** if **longname( )** is going to be used with multiple terminals.

If a given terminal doesn't support a video attribute that an application program is trying to use, **curses** may substitute a different video attribute for it. The **termattrs( )** function returns a logical OR of all video attributes supported by the terminal. This information is useful when a **curses** program needs complete control over the appearance of the screen.

The **termname( )** routine returns the value of the environment variable **TERM** (truncated to 14 characters).

RETURN VALUES

longname( ) and **termname( )** return **NULL** on error.

Routines that return an integer return **ERR** upon failure and an integer value other than **ERR** upon successful completion.

SEE ALSO

curs_initscr(3X), **curs_outopts**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **termattrs( )** may be a macro.

NAME

curs_termcap, tgetent, tgetflag, tgetnum, tgetstr, tgoto, tputs – curses interfaces (emulated) to the termcap library

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>
#include <term.h>

int tgetent(char ∗*bp***, char** ∗*name***);**

int tgetflag(char *id***[2]);**

int tgetnum(char *id***[2]);**

char ∗**tgetstr(char** *id***[2], char** ∗∗*area***);**

char ∗**tgoto(char** ∗*cap***, int** *col***, int** *row***);**

int tputs(char ∗*str***, int** *affcnt***, int (**∗**putc)(void));**

MT-LEVEL

Unsafe

DESCRIPTION

These routines are included as a conversion aid for programs that use the *termcap* library. Their parameters are the same and the routines are emulated using the *terminfo* database. These routines are supported at Level 2 and should not be used in new applications.

The **tgetent( )** routine looks up the termcap entry for *name*. The emulation ignores the buffer pointer *bp*.

The **tgetflag( )** routine gets the boolean entry for *id*.

The **tgetnum( )** routine gets the numeric entry for *id*.

The **tgetstr( )** routine returns the string entry for *id*. Use **tputs( )** to output the returned string.

The **tgoto( )** routine instantiates the parameters into the given capability. The output from this routine is to be passed to **tputs( )**.

The **tputs( )** routine is described on the **curs_terminfo**(3X) manual page.

RETURN VALUES

Routines that return an integer return **ERR** upon failure and an integer value other than **ERR** upon successful completion.

Routines that return pointers return **NULL** on error.

SEE ALSO

curs_terminfo(3X), **curses**(3X), **putc**(3S)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

NAME

curs_terminfo, setupterm, setterm, set_curterm, del_curterm, restartterm, tparm, tputs, putp, vidputs, vidattr, mvcur, tigetflag, tigetnum, tigetstr – curses interfaces to terminfo database

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lcurses** [*library* . .]

#include <curses.h>
#include <term.h>

int setupterm(char ∗*term***, int** *fildes***, int** ∗*errret***);**

int setterm(char ∗*term***);**

int set_curterm(TERMINAL ∗*nterm***);**

int del_curterm(TERMINAL ∗*oterm***);**

int restartterm(char ∗*term***, int** *fildes***, int** ∗*errret***);**

char ∗**tparm(char** ∗*str***, long int** *p1***, long int** *p2***, long int** *p3***, long int** *p4***, long int** *p5***, long int** *p6***, long int** *p7***, long int** *p8***, long int** *p9***);**

int tputs(char ∗*str***, int** *affcnt***, int (**∗**putc)(char));**

int putp(char ∗*str***);**

int vidputs(chtype *attrs***, int (**∗**putc)(char));**

int vidattr(chtype *attrs***);**

int mvcur(int *oldrow***, int** *oldcol***, int** *newrow***, int** *newcol***);**

int tigetflag(char ∗*capname***);**

int tigetnum(char ∗*capname***);**

int tigetstr(char ∗*capname***);**

MT-LEVEL

Unsafe

DESCRIPTION

These low-level routines must be called by programs that have to deal directly with the *terminfo* database to handle certain terminal capabilities, such as programming function keys. For all other functionality, **curses** routines are more suitable and their use is recommended.

Initially, **setupterm( )** should be called. Note that **setupterm( )** is automatically called by **initscr( )** and **newterm( )**. This defines the set of terminal-dependent variables (listed in **terminfo**(4)). The *terminfo* variables **lines** and **columns** are initialized by **setupterm( )** as follows: If **use_env(FALSE)** has been called, values for **lines** and **columns** specified in *terminfo* are used. Otherwise, if the environment variables **LINES** and **COLUMNS** exist, their values are used. If these environment variables do not exist and the program is running in a window, the current window size is used. Otherwise, if the environment variables do not exist, the values for **lines** and **columns** specified in the *terminfo* database are used.

The headers **<curses.h>** and **<term.h>** should be included (in this order) to get the definitions for these strings, numbers, and flags. Parameterized strings should be passed through **tparm( )** to instantiate them. All *terminfo* strings (including the output of

tparm( )) should be printed with **tputs( )** or **putp( )**. Call the **reset_shell_mode( )** routine to restore the tty modes before exiting (see **curs_kernel**(3X)). Programs which use cursor addressing should output **enter_ca_mode** upon startup and should output **exit_ca_mode** before exiting. Programs desiring shell escapes should call **reset_shell_mode** and output **exit_ca_mode** before the shell is called and should output **enter_ca_mode** and call **reset_prog_mode** after returning from the shell.

The **setupterm( )** routine reads in the *terminfo* database, initializing the *terminfo* structures, but does not set up the output virtualization structures used by **curses**. The terminal type is the character string *term;* if *term* is null, the environment variable **TERM** is used. All output is to file descriptor *fildes* which is initialized for output. If *errret* is not null, then **setupterm( )** returns **OK** or **ERR** and stores a status value in the integer pointed to by *errret*. A status of **1** in *errret* is normal, **0** means that the terminal could not be found, and **–1** means that the *terminfo* database could not be found. If *errret* is null, **setupterm( )** prints an error message upon finding an error and exits. Thus, the simplest call is:

setupterm((char *)0, 1, (int *)0);,

which uses all the defaults and sends the output to **stdout**.

The **setterm( )** routine is being replaced by **setupterm( )**. The call:

setupterm(*term***, 1, (int *)0)**

provides the same functionality as **setterm(***term***)**. The **setterm( )** routine is included here for compatibility and is supported at Level 2.

The **set_curterm( )** routine sets the variable **cur_term** to *nterm*, and makes all of the *terminfo* boolean, numeric, and string variables use the values from *nterm*.

The **del_curterm( )** routine frees the space pointed to by *oterm* and makes it available for further use. If *oterm* is the same as **cur_term**, references to any of the *terminfo* boolean, numeric, and string variables thereafter may refer to invalid memory locations until another **setupterm( )** has been called.

The **restartterm( )** routine is similar to **setupterm( )** and **initscr( )**, except that it is called after restoring memory to a previous state. It assumes that the windows and the input and output options are the same as when memory was saved, but the terminal type and baud rate may be different.

The **tparm( )** routine instantiates the string *str* with parameters *pi*. A pointer is returned to the result of *str* with the parameters applied.

The **tputs( )** routine applies padding information to the string *str* and outputs it. The *str* must be a terminfo string variable or the return value from **tparm( )**, **tgetstr( )**, or **tgoto( )**. *affcnt* is the number of lines affected, or 1 if not applicable. *putc* is a **putchar( )**-like routine to which the characters are passed, one at a time.

The **putp( )** routine calls **tputs(***str***, 1, putchar)**. Note that the output of **putpA( )** always goes to **stdout**, not to the *fildes* specified in **setupterm( )**.

The **vidputs( )** routine displays the string on the terminal in the video attribute mode *attrs*, which is any combination of the attributes listed in **curses**(3X). The characters are passed to the **putchar( )**-like routine **putc( ) .**

The **vidattr( )** routine is like the **vidputs( )** routine, except that it outputs through **putchar( )**.

The **mvcur( )** routine provides low-level cursor motion.

The **tigetflag( )**, **tigetnum( )** and **tigetstr( )** routines return the value of the capability corresponding to the *terminfo capname* passed to them, such as **xenl**.

With the **tigetflag( )** routine, the value **–1** is returned if *capname* is not a boolean capability.

With the **tigetnum( )** routine, the value **–2** is returned if *capname* is not a numeric capability.

With the **tigetstr( )** routine, the value **(char *)–1** is returned if *capname* is not a string capability.

The *capname* for each capability is given in the table column entitled *capname* code in the capabilities section of **terminfo**(4).

char *boolnames, ***boolcodes**, ***boolfnames**

char *numnames, ***numcodes**, ***numfnames**

char *strnames, ***strcodes**, ***strfnames**

These null-terminated arrays contain the *capnames*, the *termcap* codes, and the full C names, for each of the *terminfo* variables.

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion, unless otherwise noted in the preceding routine descriptions.

Routines that return pointers always return **NULL** on error.

SEE ALSO

curs_initscr(3X), **curs_kernel**(3X), **curs_termcap**(3X), **curses**(3X), **putc**(3S), **terminfo**(4)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

The **setupterm( )** routine should be used in place of **setterm( )**.

Note that **vidattr( )** and **vidputs( )** may be macros.

NAME

curs_touch, touchwin, touchline, untouchwin, wtouchln, is_linetouched, is_wintouched – curses refresh control routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

int touchwin(WINDOW ∗*win***);**

int touchline(WINDOW ∗*win***, int** *start***, int** *count***);**

int untouchwin(WINDOW ∗*win***);**

int wtouchln(WINDOW ∗*win***, int** *y***, int** *n***, int** *changed***);**

int is_linetouched(WINDOW ∗*win***, int** *line***);**

int is_wintouched(WINDOW ∗*win***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **touchwin( )** and **touchline( )** routines throw away all optimization information about which parts of the window have been touched, by pretending that the entire window has been drawn on. This is sometimes necessary when using overlapping windows, since a change to one window affects the other window, but the records of which lines have been changed in the other window do not reflect the change. The routine **touchline( )** only pretends that *count* lines have been changed, beginning with line *start*.

The **untouchwin( )** routine marks all lines in the window as unchanged since the last call to **wrefresh( )**.

The **wtouchln( )** routine makes *n* lines in the window, starting at line *y*, look as if they have (*changed***=1**) or have not (*changed***=0**) been changed since the last call to **wrefresh( )**.

The **is_linetouched( )** and **is_wintouched( )** routines return **TRUE** if the specified line / window was modified since the last call to **wrefresh( )**; otherwise they return **FALSE**. In addition, **is_linetouched( )** returns **ERR** if *line* is not valid for the given window.

RETURN VALUES

All routines return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion, unless otherwise noted in the preceding routine descriptions.

SEE ALSO

curs_refresh(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that all routines except **wtouchln( )** may be macros.

NAME

curs_util, unctrl, keyname, filter, use_env, putwin, getwin, delay_output, flushinp – curses miscellaneous utility routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lcurses** [*library* . .]

#include <curses.h>

char ∗**unctrl(chtype** *c***);**

char ∗**keyname(int** *c***);**

int filter(void);

void use_env(char *bool***);**

int putwin(WINDOW ∗*win***, FILE** ∗*filep***);**

WINDOW ∗**getwin(FILE** ∗*filep***);**

int delay_output(int *ms***);**

int flushinp(void);

MT-LEVEL

Unsafe

DESCRIPTION

The **unctrl( )** macro expands to a character string which is a printable representation of the character *c*. Control characters are displayed in the ^*X* notation. Printing characters are displayed as is.

With the **keyname( )** routine, a character string corresponding to the key *c* is returned.

The **filter( )** routine, if used, is called before **initscr( )** or **newterm( )** are called. It makes **curses** think that there is a one-line screen. **curses** does not use any terminal capabilities that assume that they know on what line of the screen the cursor is positioned.

The **use_env( )** routine, if used, is called before **initscr( )** or **newterm( )** are called. When called with **FALSE** as an argument, the values of **lines** and **columns** specified in the *terminfo* database will be used, even if environment variables **LINES** and **COLUMNS** (used by default) are set, or if **curses** is running in a window (in which case default behavior would be to use the window size if **LINES** and **COLUMNS** are not set).

With the **putwin( )** routine, all data associated with window *win* is written into the file to which *filep* points. This information can be later retrieved using the **getwin( )** function.

The **getwin( )** routine reads window related data stored in the file by **putwin( )**. The routine then creates and initializes a new window using that data. It returns a pointer to the new window.

The **delay_output( )** routine inserts an *ms* millisecond pause in output. This routine should not be used extensively because padding characters are used rather than a CPU pause.

The **flushinp( )** routine throws away any typeahead that has been typed by the user and has not yet been read by the program.

RETURN VALUES

Except for **flushinp( )**, routines that return an integer return **ERR** upon failure and an integer value other than **ERR** upon successful completion.

flushinp( ) always returns **OK**.

Routines that return pointers return **NULL** on error.

SEE ALSO

curs_initscr(3X), **curs_scr_dump**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

Note that **unctrl( )** is a macro, which is defined in <**unctrl.h**>.

NAME

curs_window, newwin, delwin, mvwin, subwin, derwin, mvderwin, dupwin, wsyncup, syncok, wcursyncup, wsyncdown – create curses windows

SYNOPSIS

cc [*flag* ...] *file* ... **–lcurses** [*library* ..]

#include <curses.h>

WINDOW ∗**newwin(int** *nlines***, int** *ncols***, int** *begin_y***, int** *begin_x***);**

int delwin(WINDOW ∗*win***);**

int mvwin(WINDOW ∗*win***, int** *y***, int** *x***);**

WINDOW ∗**subwin(WINDOW** ∗*orig***, int** *nlines***, int** *ncols***, int** *begin_y***, int** *begin_x***);**

WINDOW ∗**derwin(WINDOW** ∗*orig***, int** *nlines***, int** *ncols***, int** *begin_y***, int** *begin_x***);**

int mvderwin(WINDOW ∗*win***, int** *par_y***, int** *par_x***);**

WINDOW ∗**dupwin(WINDOW** ∗*win***);**

void wsyncup(WINDOW ∗*win***);**

int syncok(WINDOW ∗*win***, bool** *bf***);**

void wcursyncup(WINDOW ∗*win***);**

void wsyncdown(WINDOW ∗*win***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **newwin( )** routine creates and returns a pointer to a new window with the given number of lines, *nlines*, and columns, *ncols*. The upper left-hand corner of the window is at line *begin_y*, column *begin_x*. If either *nlines* or *ncols* is zero, they default to **LINES** — *begin_y* and **COLS** — *begin_x*. A new full-screen window is created by calling **newwin(0,0,0,0)**.

The **delwin( )** routine deletes the named window, freeing all memory associated with it. Subwindows must be deleted before the main window can be deleted.

The **mvwin( )** routine moves the window so that the upper left-hand corner is at position (*x*, *y*). If the move would cause the window to be off the screen, it is an error and the window is not moved. Moving subwindows is allowed, but should be avoided.

The **subwin( )** routine creates and returns a pointer to a new window with the given number of lines, *nlines*, and columns, *ncols*. The window is at position (*begin_y*, *begin_x*) on the screen. (This position is relative to the screen, and not to the window *orig*.) The window is made in the middle of the window *orig*, so that changes made to one window will affect both windows. The subwindow shares memory with the window *orig*. When using this routine, it is necessary to call **touchwin( )** or **touchline( )** on *orig* before calling **wrefresh( )** on the subwindow.

The **derwin( )** routine is the same as **subwin( )**, except that *begin_y* and *begin_x* are relative to the origin of the window *orig* rather than the screen. There is no difference between the subwindows and the derived windows.

The **mvderwin( )** routine moves a derived window (or subwindow) inside its parent window. The screen-relative parameters of the window are not changed. This routine is used to display different parts of the parent window at the same physical position on the screen.

The **dupwin( )** routine creates an exact duplicate of the window *win*.

Each **curses** window maintains two data structures: the character image structure and the status structure. The character image structure is shared among all windows in the window hierarchy (that is, the window with all subwindows). The status structure, which contains information about individual line changes in the window, is private to each window. The routine **wrefresh( )** uses the status data structure when performing screen updating. Since status structures are not shared, changes made to one window in the hierarchy may not be properly reflected on the screen.

The routine **wsyncup( )** causes the changes in the status structure of a window to be reflected in the status structures of its ancestors. If **syncok( )** is called with second argument **TRUE** then **wsyncup( )** is called automatically whenever there is a change in the window.

The routine **wcursyncup( )** updates the current cursor position of all the ancestors of the window to reflect the current cursor position of the window.

The routine **wsyncdown( )** updates the status structure of the window to reflect the changes in the status structures of its ancestors. Applications seldom call this routine because it is called automatically by **wrefresh( )**.

RETURN VALUES

Routines that return an integer return the integer **ERR** upon failure and an integer value other than **ERR** upon successful completion.

delwin( ) returns the integer **ERR** upon failure and **OK** upon successful completion.

Routines that return pointers return **NULL** on error.

SEE ALSO

curs_refresh(3X), **curs_touch**(3X), **curses**(3X)

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

If many small changes are made to the window, the **wsyncup( )** option could degrade performance.

Note that **syncok( )** may be a macro.

NAME

curses – CRT screen handling and optimization package

SYNOPSIS

cc [*flag* ...] *file* ... –**lcurses** [*library* ...]

#include <curses.h>

MT-LEVEL

Unsafe

DESCRIPTION

The **curses** library routines give the user a terminal-independent method of updating character screens with reasonable optimization.

The **curses** package allows: overall screen, window and pad manipulation; output to windows and pads; reading terminal input; control over terminal and **curses** input and output options; environment query routines; color manipulation; use of soft label keys; terminfo access; and access to low-level **curses** routines.

To initialize the routines, the routine **initscr( )** or **newterm( )** must be called before any of the other routines that deal with windows and screens are used. The routine **endwin( )** must be called before exiting. To get character-at-a-time input without echoing (most interactive, screen oriented programs want this), the following sequence should be used:

initscr,cbreak,noecho;

Most programs would additionally use the sequence:

nonl,intrflush(stdscr,FALSE),keypad(stdscr,TRUE);

Before a **curses** program is run, the tab stops of the terminal should be set and its initialization strings, if defined, must be output. This can be done by executing the **tput init** command after the shell environment variable **TERM** has been exported. (See **terminfo**(4) for further details.)

The **curses** library permits manipulation of data structures, called *windows*, which can be thought of as two-dimensional arrays of characters representing all or part of a CRT screen. A default window called **stdscr**, which is the size of the terminal screen, is supplied. Others may be created with **newwin**(3X).

Windows are referred to by variables declared as **WINDOW ***. These data structures are manipulated with routines described on 3X pages (whose names begin "curs_"). Among which the most basic routines are **move**(3X) and **addch**(3X). More general versions of these routines are included with names beginning with **w**, allowing the user to specify a window. The routines not beginning with **w** affect **stdscr**.

After using routines to manipulate a window, **refresh**(3X) is called, telling **curses** to make the user's CRT screen look like **stdscr**. The characters in a window are actually of type **chtype**, (character and attribute data) so that other information about the character may also be stored with each character.

Special windows called *pads* may also be manipulated. These are windows which are not constrained to the size of the screen and whose contents need not be completely displayed. See **curs_pad**(3X) for more information.

In addition to drawing characters on the screen, video attributes and colors may be included, causing the characters to show up in such modes as underlined, in reverse video, or in color on terminals that support such display enhancements. Line drawing characters may be specified to be output. On input, **curses** is also able to translate arrow and function keys that transmit escape sequences into single values. The video attributes, line drawing characters, and input values use names, defined in **<curses.h>**, such as **A_REVERSE**, **ACS_HLINE**, and **KEY_LEFT**.

If the environment variables **LINES** and **COLUMNS** are set, or if the program is executing in a window environment, line and column information in the environment will override information read by *terminfo*. This would effect a program running in an AT&T 630 layer, for example, where the size of a screen is changeable.

If the environment variable **TERMINFO** is defined, any program using **curses** checks for a local terminal definition before checking in the standard place. For example, if **TERM** is set to **att4424**, then the compiled terminal definition is found in

/usr/share/lib/terminfo/a/att4424.

(The '**a**' is copied from the first letter of **att4424** to avoid creation of huge directories.) However, if **TERMINFO** is set to **$HOME/myterms**, **curses** first checks

$HOME/myterms/a/att4424,

and if that fails, it then checks

/usr/share/lib/terminfo/a/att4424.

This is useful for developing experimental definitions or when write permission in **/usr/share/lib/terminfo** is not available.

The integer variables **LINES** and **COLS** are defined in **<curses.h>** and will be filled in by **initscr** with the size of the screen. The constants **TRUE** and **FALSE** have the values **1** and **0**, respectively.

The **curses** routines also define the **WINDOW** * variable **curscr** which is used for certain low-level operations like clearing and redrawing a screen containing garbage. The **curscr** can be used in only a few routines.

International Functions

The number of bytes and the number of columns to hold a character from the supplementary character set is locale-specific (locale category **LC_CTYPE**) and can be specified in the character class table.

For editing, operating at the character level is entirely appropriate. For screen formatting, arbitrary movement of characters on screen is not desirable.

Overwriting characters (**addch**, for example) operates on a screen level. Overwriting a character by a character that requires a different number of columns may produce *orphaned columns*. These orphaned columns are filled with background characters.

Inserting characters (**insch**, for example) operates on a character level (that is, at the character boundaries). The specified character is inserted right before the character, regardless of which column of a character the cursor points to. Before insertion, the cursor

position is adjusted to the first column of the character.

As with inserting characters, deleting characters (**delch**, for example) operates on a character level (that is, at the character boundaries). The character at the cursor is deleted whichever column of the character the cursor points to. Before deletion, the cursor position is adjusted to the first column of the character.

A *multi-column* character cannot be put on the last column of a line. When such attempts are made, the last column is set to the background character. In addition, when such an operation creates orphaned columns, the orphaned columns are filled with background characters.

Overlapping and overwriting a window follows the operation of overwriting characters around its edge. The orphaned columns, if any, are handled as in the character operations.

The cursor is allowed to be placed anywhere in a window. If the insertion or deletion is made when the cursor points to the second or later column position of a character that holds multiple columns, the cursor is adjusted to the first column of the character before the insertion or deletion.

Routine and Argument Names

Many **curses** routines have two or more versions. The routines prefixed with **w** require a window argument. The routines prefixed with **p** require a pad argument. Those without a prefix generally use **stdscr**.

The routines prefixed with **mv** require an *x* and *y* coordinate to move to before performing the appropriate action. The **mv** routines imply a call to **move**(3X) before the call to the other routine. The coordinate *y* always refers to the row (of the window), and *x* always refers to the column. The upper left-hand corner is always (0,0), not (1,1).

The routines prefixed with **mvw** take both a window argument and *x* and *y* coordinates. The window argument is always specified before the coordinates.

In each case, *win* is the window affected, and *pad* is the pad affected; *win* and *pad* are always pointers to type **WINDOW**

Option setting routines require a Boolean flag *bf* with the value **TRUE** or **FALSE**; *bf* is always of type **bool**. The variables *ch* and *attrs* below are always of type **chtype**. The types **WINDOW**, **SCREEN**, **bool**, and **chtype** are defined in **<curses.h>**. The type **TERMINAL** is defined in **<term.h>**. All other arguments are integers.

Routine Name Index

The following table lists each **curses** routine and the name of the manual page on which it is described.

curses Routine Name	Manual Page Name
addch	curs_addch(3X)
addchnstr	curs_addchstr(3X)
addchstr	curs_addchstr(3X)
addnstr	curs_addstr(3X)
addnwstr	curs_addwstr(3X)
addstr	curs_addstr(3X)
addwch	curs_addwch(3X)
addwchnstr	curs_addwchstr(3X)
addwchstr	curs_addwchstr(3X)
addwstr	curs_addwstr(3X)
adjcurspos	curs_alecompat(3X)
attroff	curs_attr(3X)
attron	curs_attr(3X)
attrset	curs_attr(3X)
baudrate	curs_termattrs(3X)
beep	curs_beep(3X)
bkgd	curs_bkgd(3X)
bkgdset	curs_bkgd(3X)
border	curs_border(3X)
box	curs_border(3X)
can_change_color	curs_color(3X)
cbreak	curs_inopts(3X)
clear	curs_clear(3X)
clearok	curs_outopts(3X)
clrtobot	curs_clear(3X)
clrtoeol	curs_clear(3X)
color_content	curs_color(3X)
copywin	curs_overlay(3X)
curs_set	curs_kernel(3X)
def_prog_mode	curs_kernel(3X)
def_shell_mode	curs_kernel(3X)
del_curterm	curs_terminfo(3X)
delay_output	curs_util(3X)
delch	curs_delch(3X)
deleteln	curs_deleteln(3X)
delscreen	curs_initscr(3X)
delwin	curs_window(3X)
derwin	curs_window(3X)
doupdate	curs_refresh(3X)
dupwin	curs_window(3X)

echo	**curs_inopts(3X)**
echochar	**curs_addch(3X)**
echowchar	**curs_addwch(3X)**
endwin	**curs_initscr(3X)**
erase	**curs_clear(3X)**
erasechar	**curs_termattrs(3X)**
filter	**curs_util(3X)**
flash	**curs_beep(3X)**
flushinp	**curs_util(3X)**
getbegyx	**curs_getyx(3X)**
getch	**curs_getch(3X)**
getmaxyx	**curs_getyx(3X)**
getnwstr	**curs_getwstr(3X)**
getparyx	**curs_getyx(3X)**
getstr	**curs_getstr(3X)**
getsyx	**curs_kernel(3X)**
getwch	**curs_getwch(3X)**
getwin	**curs_util(3X)**
getwstr	**curs_getwstr(3X)**
getyx	**curs_getyx(3X)**
halfdelay	**curs_inopts(3X)**
has_colors	**curs_color(3X)**
has_ic	**curs_termattrs(3X)**
has_il	**curs_termattrs(3X)**
idcok	**curs_outopts(3X)**
idlok	**curs_outopts(3X)**
immedok	**curs_outopts(3X)**
inch	**curs_inch(3X)**
inchnstr	**curs_inchstr(3X)**
inchstr	**curs_inchstr(3X)**
init_color	**curs_color(3X)**
init_pair	**curs_color(3X)**
initscr	**curs_initscr(3X)**
innstr	**curs_instr(3X)**
innwstr	**curs_inwstr(3X)**
insch	**curs_insch(3X)**
insdelln	**curs_deleteln(3X)**
insertln	**curs_deleteln(3X)**
insnstr	**curs_insstr(3X)**
insnwstr	**curs_inswstr(3X)**
insstr	**curs_insstr(3X)**
instr	**curs_instr(3X)**
inswch	**curs_inswch(3X)**
inswstr	**curs_inswstr(3X)**
intrflush	**curs_inopts(3X)**

inwch	**curs_inwch(3X)**
inwchnstr	**curs_inwchstr(3X)**
inwchstr	**curs_inwchstr(3X)**
inwstr	**curs_inwstr(3X)**
is_linetouched	**curs_touch(3X)**
is_wintouched	**curs_touch(3X)**
isendwin	**curs_initscr(3X)**
keyname	**curs_util(3X)**
keypad	**curs_inopts(3X)**
killchar	**curs_termattrs(3X)**
leaveok	**curs_outopts(3X)**
longname	**curs_termattrs(3X)**
meta	**curs_inopts(3X)**
move	**curs_move(3X)**
movenextch	**curs_alecompat(3X)**
moveprevch	**curs_alecompat(3X)**
mvaddch	**curs_addch(3X)**
mvaddchnstr	**curs_addchstr(3X)**
mvaddchstr	**curs_addchstr(3X)**
mvaddnstr	**curs_addstr(3X)**
mvaddnwstr	**curs_addwstr(3X)**
mvaddstr	**curs_addstr(3X)**
mvaddwch	**curs_addwch(3X)**
mvaddwchnstr	**curs_addwchstr(3X)**
mvaddwchstr	**curs_addwchstr(3X)**
mvaddwstr	**curs_addwstr(3X)**
mvcur	**curs_terminfo(3X)**
mvdelch	**curs_delch(3X)**
mvderwin	**curs_window(3X)**
mvgetch	**curs_getch(3X)**
mvgetnwstr	**curs_getwstr(3X)**
mvgetstr	**curs_getstr(3X)**
mvgetwch	**curs_getwch(3X)**
mvgetwstr	**curs_getwstr(3X)**
mvinch	**curs_inch(3X)**
mvinchnstr	**curs_inchstr(3X)**
mvinchstr	**curs_inchstr(3X)**
mvinnstr	**curs_instr(3X)**
mvinnwstr	**curs_inwstr(3X)**
mvinsch	**curs_insch(3X)**
mvinsnstr	**curs_insstr(3X)**
mvinsnwstr	**curs_inswstr(3X)**
mvinsstr	**curs_insstr(3X)**
mvinstr	**curs_instr(3X)**
mvinswch	**curs_inswch(3X)**

mvinswstr	**curs_inswstr(3X)**
mvinwch	**curs_inwch(3X)**
mvinwchnstr	**curs_inwchstr(3X)**
mvinwchstr	**curs_inwchstr(3X)**
mvinwstr	**curs_inwstr(3X)**
mvprintw	**curs_printw(3X)**
mvscanw	**curs_scanw(3X)**
mvwaddch	**curs_addch(3X)**
mvwaddchnstr	**curs_addchstr(3X)**
mvwaddchstr	**curs_addchstr(3X)**
mvwaddnstr	**curs_addstr(3X)**
mvwaddnwstr	**curs_addwstr(3X)**
mvwaddstr	**curs_addstr(3X)**
mvwaddwch	**curs_addwch(3X)**
mvwaddwchnstr	**curs_addwchstr(3X)**
mvwaddwchstr	**curs_addwchstr(3X)**
mvwaddwstr	**curs_addwstr(3X)**
mvwdelch	**curs_delch(3X)**
mvwgetch	**curs_getch(3X)**
mvwgetnwstr	**curs_getwstr(3X)**
mvwgetstr	**curs_getstr(3X)**
mvwgetwch	**curs_getwch(3X)**
mvwgetwstr	**curs_getwstr(3X)**
mvwin	**curs_window(3X)**
mvwinch	**curs_inch(3X)**
mvwinchnstr	**curs_inchstr(3X)**
mvwinchstr	**curs_inchstr(3X)**
mvwinnstr	**curs_instr(3X)**
mvwinnwstr	**curs_inwstr(3X)**
mvwinsch	**curs_insch(3X)**
mvwinsnstr	**curs_insstr(3X)**
mvwinsstr	**curs_insstr(3X)**
mvwinstr	**curs_instr(3X)**
mvwinswch	**curs_inswch(3X)**
mvwinswstr	**curs_inswstr(3X)**
mvwinwch	**curs_inwch(3X)**
mvwinwchnstr	**curs_inwchstr(3X)**
mvwinwchstr	**curs_inwchstr(3X)**
mvwinwstr	**curs_inwstr(3X)**
mvwprintw	**curs_printw(3X)**
mvwscanw	**curs_scanw(3X)**
napms	**curs_kernel(3X)**
newpad	**curs_pad(3X)**
newterm	**curs_initscr(3X)**
newwin	**curs_window(3X)**

nl	**curs_outopts(3X)**
nocbreak	**curs_inopts(3X)**
nodelay	**curs_inopts(3X)**
noecho	**curs_inopts(3X)**
nonl	**curs_outopts(3X)**
noqiflush	**curs_inopts(3X)**
noraw	**curs_inopts(3X)**
notimeout	**curs_inopts(3X)**
overlay	**curs_overlay(3X)**
overwrite	**curs_overlay(3X)**
pair_content	**curs_color(3X)**
pechochar	**curs_pad(3X)**
pechowchar	**curs_pad(3X)**
pnoutrefresh	**curs_pad(3X)**
prefresh	**curs_pad(3X)**
printw	**curs_printw(3X)**
putp	**curs_terminfo(3X)**
putwin	**curs_util(3X)**
qiflush	**curs_inopts(3X)**
raw	**curs_inopts(3X)**
redrawwin	**curs_refresh(3X)**
refresh	**curs_refresh(3X)**
reset_prog_mode	**curs_kernel(3X)**
reset_shell_mode	**curs_kernel(3X)**
resetty	**curs_kernel(3X)**
restartterm	**curs_terminfo(3X)**
ripoffline	**curs_kernel(3X)**
savetty	**curs_kernel(3X)**
scanw	**curs_scanw(3X)**
scr_dump	**curs_scr_dump(3X)**
scr_init	**curs_scr_dump(3X)**
scr_restore	**curs_scr_dump(3X)**
scr_set	**curs_scr_dump(3X)**
scroll	**curs_scroll(3X)**
scrollok	**curs_outopts(3X)**
set_curterm	**curs_terminfo(3X)**
set_term	**curs_initscr(3X)**
setscrreg	**curs_outopts(3X)**
setsyx	**curs_kernel(3X)**
setterm	**curs_terminfo(3X)**
setupterm	**curs_terminfo(3X)**
slk_attroff	**curs_slk(3X)**
slk_attron	**curs_slk(3X)**
slk_attrset	**curs_slk(3X)**
slk_clear	**curs_slk(3X)**

slk_init	**curs_slk(3X)**
slk_label	**curs_slk(3X)**
slk_noutrefresh	**curs_slk(3X)**
slk_refresh	**curs_slk(3X)**
slk_restore	**curs_slk(3X)**
slk_set	**curs_slk(3X)**
slk_touch	**curs_slk(3X)**
srcl	**curs_scroll(3X)**
standend	**curs_attr(3X)**
standout	**curs_attr(3X)**
start_color	**curs_color(3X)**
subpad	**curs_pad(3X)**
subwin	**curs_window(3X)**
syncok	**curs_window(3X)**
termattrs	**curs_termattrs(3X)**
termname	**curs_termattrs(3X)**
tgetent	**curs_termcap(3X)**
tgetflag	**curs_termcap(3X)**
tgetnum	**curs_termcap(3X)**
tgetstr	**curs_termcap(3X)**
tgoto	**curs_termcap(3X)**
tigetflag	**curs_terminfo(3X)**
tigetnum	**curs_terminfo(3X)**
tigetstr	**curs_terminfo(3X)**
timeout	**curs_inopts(3X)**
touchline	**curs_touch(3X)**
touchwin	**curs_touch(3X)**
tparm	**curs_terminfo(3X)**
tputs	**curs_terminfo(3X)**
typeahead	**curs_inopts(3X)**
unctrl	**curs_util(3X)**
ungetch	**curs_getch(3X)**
ungetwch	**curs_getwch(3X)**
untouchwin	**curs_touch(3X)**
use_env	**curs_util(3X)**
vidattr	**curs_terminfo(3X)**
vidputs	**curs_terminfo(3X)**
vwprintw	**curs_printw(3X)**
vwscanw	**curs_scanw(3X)**
waddch	**curs_addch(3X)**
waddchnstr	**curs_addchstr(3X)**
waddchstr	**curs_addchstr(3X)**
waddnstr	**curs_addstr(3X)**
waddnwstr	**curs_addwstr(3X)**
waddstr	**curs_addstr(3X)**

waddwch	**curs_addwch(3X)**
waddwchnstr	**curs_addwchstr(3X)**
waddwchstr	**curs_addwchstr(3X)**
waddwstr	**curs_addwstr(3X)**
wadjcurspos	**curs_alecompat(3X)**
wattroff	**curs_attr(3X)**
wattron	**curs_attr(3X)**
wattrset	**curs_attr(3X)**
wbkgd	**curs_bkgd(3X)**
wbkgdset	**curs_bkgd(3X)**
wborder	**curs_border(3X)**
wclear	**curs_clear(3X)**
wclrtobot	**curs_clear(3X)**
wclrtoeol	**curs_clear(3X)**
wcursyncup	**curs_window(3X)**
wdelch	**curs_delch(3X)**
wdeleteln	**curs_deleteln(3X)**
wechochar	**curs_addch(3X)**
wechowchar	**curs_addwch(3X)**
werase	**curs_clear(3X)**
wgetch	**curs_getch(3X)**
wgetnstr	**curs_getstr(3X)**
wgetnwstr	**curs_getwstr(3X)**
wgetstr	**curs_getstr(3X)**
wgetwch	**curs_getwch(3X)**
wgetwstr	**curs_getwstr(3X)**
whline	**curs_border(3X)**
winch	**curs_inch(3X)**
winchnstr	**curs_inchstr(3X)**
winchstr	**curs_inchstr(3X)**
winnstr	**curs_instr(3X)**
winnwstr	**curs_inwstr(3X)**
winsch	**curs_insch(3X)**
winsdelln	**curs_deleteln(3X)**
winsertln	**curs_deleteln(3X)**
winsnstr	**curs_insstr(3X)**
winsnwstr	**curs_inswstr(3X)**
winsstr	**curs_insstr(3X)**
winstr	**curs_instr(3X)**
winswch	**curs_inswch(3X)**
winswstr	**curs_inswstr(3X)**
winwch	**curs_inwch(3X)**
winwchnstr	**curs_inwchstr(3X)**
winwchstr	**curs_inwchstr(3X)**
winwstr	**curs_inwstr(3X)**

wmove	**curs_move(3X)**
wmovenextch	**curs_alecompat(3X)**
wmoveprevch	**curs_alecompat(3X)**
wnoutrefresh	**curs_refresh(3X)**
wprintw	**curs_printw(3X)**
wredrawln	**curs_refresh(3X)**
wrefresh	**curs_refresh(3X)**
wscanw	**curs_scanw(3X)**
wscrl	**curs_scroll(3X)**
wsetscrreg	**curs_outopts(3X)**
wstandend	**curs_attr(3X)**
wstandout	**curs_attr(3X)**
wsyncdown	**curs_window(3X)**
wsyncup	**curs_window(3X)**
wtimeout	**curs_inopts(3X)**
wtouchln	**curs_touch(3X)**
wvline	**curs_border(3X)**

RETURN VALUES

Routines that return an integer return **ERR** upon failure and an integer value other than **ERR** upon successful completion, unless otherwise noted in the routine descriptions.

All macros return the value of the **w** version, except **setscrreg( )**, **wsetscrreg( )**, **getyx( )**, **getbegyx( )**, and **getmaxyx( )**. The return values of **setscrreg( )**, **wsetscrreg( )**, **getyx( ) , getbegyx( )**, and **getmaxyx( )** are undefined (that is, these should not be used as the right-hand side of assignment statements).

Routines that return pointers return **NULL** on error.

SEE ALSO

terminfo(4) and 3X pages whose names begin with ‘‘curs_’’ for detailed routine descriptions.

NOTES

The header **<curses.h>** automatically includes the headers **<stdio.h>** and **<unctrl.h>**.

NAME

cuserid – get character login name of the user

SYNOPSIS

```
#include <stdio.h>
char *cuserid(char *s);
```

MT-LEVEL

MT-Safe

DESCRIPTION

cuserid( ) generates a character-string representation of the login name that the owner of the current process is logged in under. If *s* is a **NULL** pointer, this representation is generated in an internal static area, the address of which is returned. Otherwise, *s* is assumed to point to an array of at least **L_cuserid** characters; the representation is left in this array. The constant **L_cuserid** is defined in the **<stdio.h>** header.

RETURN VALUES

If the login name cannot be found, **cuserid( )** returns a **NULL** pointer; if *s* is not a **NULL** pointer, a null character **`\0´** will be placed at *s***[0]**.

SEE ALSO

getlogin(3C), **getpwnam**(3C)

NOTES

In multi-thread applications, the caller must always supply an array *s* for the return value.

NAME

dbm, dbminit, dbmclose, fetch, store, delete, firstkey, nextkey – data base subroutines

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ... **–ldbm**

```
#include <dbm.h>

typedef struct {
   char *dptr;
   int dsize;
} datum;

int dbminit(file)
char *file;

int dbmclose()

datum fetch( key)
datum key;

int store(key, dat)
datum key, dat;

int delete(key)
datum key;

datum firstkey()

datum nextkey(key)
datum key;
```

DESCRIPTION

The **dbm( )** library has been superseded by **ndbm**(3).

These functions maintain key / content pairs in a data base. The functions will handle very large (a billion blocks) databases and will access a keyed item in one or two file system accesses.

key/dat and their content are described by the **datum** typedef. A **datum** specifies a string of *dsize* bytes pointed to by *dptr*. Arbitrary binary data, as well as normal ASCII strings, are allowed. The data base is stored in two files. One file is a directory containing a bit map and has **.dir** as its suffix. The second file contains all data and has **.pag** as its suffix.

Before a database can be accessed, it must be opened by **dbminit( )**. At the time of this call, the files *file***.dir** and *file***.pag** must exist. An empty database is created by creating zero-length **.dir** and **.pag** files.

A database may be closed by calling **dbmclose( )**. You must close a database before opening a new one.

Once open, the data stored under a key is accessed by **fetch( )** and data is placed under a key by **store**. A key (and its associated contents) is deleted by **delete( )**. A linear pass through all keys in a database may be made, in an (apparently) random order, by use of **firstkey( )** and **nextkey( )**. **firstkey( )** will return the first key in the database.
With any key **nextkey( )** will return the next key in the database. This code will traverse the data base:

for (key = firstkey; key.dptr != NULL; key = nextkey(key))

RETURN VALUES

All functions that return an **int** indicate errors with negative values. A zero return indicates no error. Routines that return a **datum** indicate errors with a NULL (0) *dptr*.

SEE ALSO

ar(1), **cat**(1), **cp**(1), **tar**(1), **ndbm**(3)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

The **.pag** file will contain holes so that its apparent size may be larger than its actual content. Older versions of the UNIX operating system may create real file blocks for these holes when touched. These files cannot be copied by normal means (**cp**(1), **cat**(1), **tar**(1), **ar**(1)) without filling in the holes.

dptr pointers returned by these subroutines point into static storage that is changed by subsequent calls.

The sum of the sizes of a key/content pair must not exceed the internal block size (currently 1024 bytes). Moreover all key/content pairs that hash together must fit on a single block. **store** will return an error in the event that a disk block fills with inseparable data.

delete() does not physically reclaim file space, although it does make it available for reuse.

The order of keys presented by **firstkey()** and **nextkey()** depends on a hashing function, not on anything interesting.

There are no interlocks and no reliable cache flushing; thus concurrent updating and reading is risky.

The database files (*file***.dir** and *file***.pag**) are binary and are architecture-specific (for example, they depend on the architecture's byte order.) These files are not guaranteed to be portable across architectures.

NAME

decimal_to_floating, decimal_to_single, decimal_to_double, decimal_to_extended, decimal_to_quadruple – convert decimal record to floating-point value

SYNOPSIS

```
#include <floatingpoint.h>

void decimal_to_single(single *px, decimal_mode *pm, decimal_record *pd,
        fp_exception_field_type *ps);

void decimal_to_double(double *px, decimal_mode *pm, decimal_record *pd,
        fp_exception_field_type *ps);

void decimal_to_extended(extended *px, decimal_mode *pm, decimal_record *pd,
        fp_exception_field_type *ps);

void decimal_to_quadruple(quadruple *px, decimal_mode *pm, decimal_record *pd,
        fp_exception_field_type *ps);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **decimal_to_floating( )** functions convert the decimal record at **pd* into a floating-point value at **px*, observing the modes specified in **pm* and setting exceptions in **ps*. If there are no IEEE exceptions, **ps* will be zero.

pd->sign and *pd->fpclass* are always taken into account. *pd->exponent*, *pd->ds* and *pd->ndigits* are used when *pd->fpclass* is *fp_normal* or *fp_subnormal*. In these cases *pd->ds* must contain one or more ascii digits followed by a NULL and *pd->ndigits* is assumed to be the length of the string *pd->ds*. Notice that for efficiency reasons, the assumption that *pd->ndigits* == strlen(*pd->ds*) is NEVER verified.

On output, **px* is set to a correctly rounded approximation to

(pd->sign)*(pd->ds)*10(pd->exponent)**

Thus if *pd->exponent* == –2 and *pd->ds* == "1234", **px* will get 12.34 rounded to storage precision. *pd->ds* cannot have more than **DECIMAL_STRING_LENGTH-1** significant digits because one character is used to terminate the string with a NULL. If *pd->more != 0* on input then additional nonzero digits follow those in *pd->ds; fp_inexact* is set accordingly on output in **ps*.

**px* is correctly rounded according to the IEEE rounding modes in *pm->rd*. **ps* is set to contain *fp_inexact*, *fp_underflow*, or *fp_overflow* if any of these arise.

pm->df and *pm->ndigits* are not used.

strtod(3C), **scanf**(3S), **fscanf**(3S), and **sscanf**(3S) all use **decimal_to_double( )**.

SEE ALSO

fscanf(3S), **scanf**(3S), **sscanf**(3S), **strtod**(3C)

NAME

demangle – decode a C++ encoded symbol name

SYNOPSIS

CC[**flag**...**]**file **... [**library**...]**
#include<demangle.h>
cpp_demangle(const char **in*, **char** **out*, **size_t size**);

DESCRIPTION

cplus_demangle(|) decodes the string *in*, and copies the result to *out*. *in* points to a string representing a name mangled by the C++ compiler. *out* is a buffer of *size* that you specify, which contains the byte size. If the output buffer is too small to contain the demangled name, **cplus_demangle( )** returns **DEMANGLE_ESPACE**, and the contents of *out* are undefined. Otherwise, if the name is a valid C++ name, **cplus_demangle( )** returns **0**. If *in* is not a valid C++ mangled name, it is copied unchanged to *out* and the function returns **DEMANGLE_ENAME**.

cplus_demangle( ) operates on mangled names generated by C++ 3.0.1 and all versions of C++ 4.0 and above.

SEE ALSO

cc(1B) in the C++ Library Reference Manual.

NAME

dial – establish an outgoing terminal line connection

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <dial.h>

int dial(CALL *call***);**

void undial(int *fd***);**

MT-LEVEL

Unsafe

DESCRIPTION

dial( ) returns a file-descriptor for a terminal line open for read / write. The argument to **dial( )** is a **CALL** structure (defined in the header **<dial.h>**).

When finished with the terminal line, the calling program must invoke **undial( )** to release the semaphore that has been set during the allocation of the terminal device.

CALL is defined in the header **<dial.h>** and has the following members:

```
struct termio   *attr;    /* pointer to termio attribute struct */
int             baud;     /* transmission data rate */
int             speed;    /* 212A modem: low=300, high=1200 */
char            *line;    /* device name for out-going line */
char            *telno;   /* pointer to tel-no digits string */
int             modem;    /* specify modem control for direct lines */
char            *device;  /* unused */
int             dev_len;  /* unused */
```

The **CALL** element **speed** is intended only for use with an outgoing dialed call, in which case its value should be the desired transmission baud rate. The **CALL** element **baud** is no longer used.

If the desired terminal line is a direct line, a string pointer to its device-name should be placed in the **line** element in the **CALL** structure. Legal values for such terminal device names are kept in the **Devices** file. In this case, the value of the **baud** element should be set to -1. This value will cause **dial** to determine the correct value from the **<Devices>** file.

The **telno** element is for a pointer to a character string representing the telephone number to be dialed. Such numbers may consist only of these characters:

0-9	dial 0-9
*	dial *
#	dial #
=	wait for secondary dial tone
–	delay for approximately 4 seconds

The **CALL** element **modem** is used to specify modem control for direct lines. This element should be non-zero if modem control is required. The **CALL** element **attr** is a pointer to a **termio** structure, as defined in the header **<termio.h>**. A **NULL** value for this pointer element may be passed to the **dial** function, but if such a structure is included, the elements specified in it will be set for the outgoing terminal line before the connection is established. This setting is often important for certain attributes such as parity and baud-rate.

The **CALL** elements **device** and **dev_len** are no longer used. They are retained in the **CALL** structure for compatibility reasons.

RETURN VALUES

On failure, a negative value indicating the reason for the failure will be returned. Mnemonics for these negative indices as listed here are defined in the header **<dial.h>**.

INTRPT	**–1**	/* interrupt occurred */
D_HUNG	**–2**	/* dialer hung (no return from write) */
NO_ANS	**–3**	/* no answer within 10 seconds */
ILL_BD	**–4**	/* illegal baud-rate */
A_PROB	**–5**	/* acu problem (open( ) failure) */
L_PROB	**–6**	/* line problem (open( ) failure) */
NO_Ldv	**–7**	/* can't open Devices file */
DV_NT_A	**–8**	/* requested device not available */
DV_NT_K	**–9**	/* requested device not known */
NO_BD_A	**–10**	/* no device available at requested baud */
NO_BD_K	**–11**	/* no device known at requested baud */
DV_NT_E	**–12**	/* requested speed does not match */
BAD_SYS	**–13**	/* system not in Systems file*/

FILES

/etc/uucp/Devices
/etc/uucp/Systems
/var/spool/uucp/LCK..*tty-device*

SEE ALSO

uucp(1C), **alarm**(2), **read**(2), **write**(2), **termio**(7I)

NOTES

Including the header **<dial.h>** automatically includes the header **<termio.h>**.

An **alarm**(2) system call for 3600 seconds is made (and caught) within the **dial** module for the purpose of ''touching'' the **LCK..** file and constitutes the device allocation semaphore for the terminal device. Otherwise, **uucp**(1C) may simply delete the **LCK..** entry on its 90-minute clean-up rounds. The alarm may go off while the user program is in a **read**(2) or **write**(2) function, causing an apparent error return. If the user program expects to be around for an hour or more, error returns from **read( )**s should be checked for **(errno==EINTR)**, and the **read( )** possibly reissued.

This interface is unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

difftime – computes the difference between two calendar times

SYNOPSIS

```
#include <time.h>
double difftime(time_t time1, time_t time0);
```

MT-LEVEL

MT-Safe

DESCRIPTION

difftime( ) computes the difference between two calendar times. **difftime( )** returns the difference *(time1-time0)* expressed in seconds as a **double**. This function is provided because there are no general arithmetic properties defined for type **time_t**.

SEE ALSO

ctime(3C)

NAME

directory, opendir, readdir, readdir_r, telldir, seekdir, rewinddir, closedir – directory operations

SYNOPSIS

#include <dirent.h>

DIR ∗**opendir(const char** ∗*filename*);

struct dirent ∗**readdir(DIR** ∗*dirp*);

long telldir(DIR ∗*dirp*);

void seekdir(DIR ∗*dirp*, **long** *loc*);

void rewinddir(DIR ∗*dirp*);

int closedir(DIR ∗*dirp*);

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struct dirent ∗**readdir_r(DIR** ∗*dirp*,**struct dirent** ∗*res*);

POSIX

cc [*flag*... **]** *file* ... –**D_POSIX_PTHREAD_SEMANTICS [** *library*... **]**

int ∗**readdir_r(DIR** ∗*dirp*,**struct dirent** ∗*entry* **struct dirent** ∗∗*result*);

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

opendir( ) opens the directory named by *filename* and associates a directory stream with it. **opendir( )** returns a pointer to be used to identify the directory stream in subsequent operations. The directory stream is positioned at the first entry. A null pointer is returned if *filename* cannot be accessed or is not a directory, or if it cannot **malloc**(3C) enough memory to hold a **DIR** structure or a buffer for the directory entries.

readdir( ) returns a pointer to a structure representing the directory entry at the current position in the directory stream to which *dirp* refers, and positions the directory stream at the next entry, except on read-only filesystems. It returns a **NULL** pointer upon reaching the end of the directory stream, or upon detecting an invalid location in the directory. **readdir( )** shall not return directory entries containing empty names. It is unspecified whether entries are returned for dot or dot-dot. The pointer returned by **readdir( )** points to data that may be overwritten by another call to **readdir( )** on the same directory stream. This data shall not be overwritten by another call to **readdir( )** on a different directory stream. **readdir( )** may buffer several directory entries per actual read operation; **readdir( )** marks for updata the *st_atime* field of the directory each time the directory is actually read.

readdir_r( ) has the equivalent functionality as **readdir( )** except that a buffer *res* must be supplied by the caller to store the result. To allocate *res* correctly a *struct dirent res* is not sufficient, thus the size should be *sizeof(struct dirent)* + **_POSIX_PATH_MAX** (defined in **<limits.h>**).

The POSIX **readdir_r( )** function initializes the structure referenced by *entry* and stores a pointer to this structure in *result*.

telldir( ) returns the current location associated with the named directory stream.

seekdir( ) sets the position of the next **readdir( )** operation on the directory stream. The new position reverts to the position associated with the directory stream at the time the **telldir( )** operation that provides *loc* was performed. Values returned by **telldir( )** are good only for the lifetime of the DIR pointer from which they are derived. If the directory is closed and then reopened, the **telldir( )** value may be invalidated due to undetected directory compaction. It is safe to use a previous **telldir( )** value immediately after a call to **opendir( )** and before any calls to **readdir**.

rewinddir( ) resets the position of the named directory stream to the beginning of the directory. It also causes the directory stream to refer to the current state of the corresponding directory, as a call to **opendir( )** would.

closedir( ) closes the named directory stream and frees the **DIR** structure.

RETURN VALUES

opendir( ), **readdir( )**, and **readdir_r( )** return **NULL** on failure and set **errno** to indicate the error. The POSIX **readdir_r( )** returns zero if successful, or an error number to indicate failure. **telldir( )**, **seekdir( )**, and **closedir( )** return –1 on failure and set **errno** to indicate the error.

ERRORS

opendir( ) will fail if one or more of the following are true:

EACCES	Read permission is denied on the specified directory.
EFAULT	*filename* points outside the allocated address space.
ELOOP	Too many symbolic links were encountered in translating *filename*.
ENOTDIR	A component of *filename* is not a directory.
EMFILE	The maximum number of file descriptors are currently open.
ENFILE	The system file table is full.
ENAMETOOLONG	The length of the *filename* argument exceeds **{PATH_MAX}**, or the length of a *filename* component exceeds **{NAME_MAX}** while **{_POSIX_NO_TRUNC}** is in effect.
ENOENT	A component of *filename* does not exist or is a null pathname.
EACCES	A component of *filename* denies search permission.

readdir( ) and **readdir_r( )** will fail if one or more of the following are true:

EAGAIN	Mandatory file / record locking was set, **O_NDELAY** or **O_NONBLOCK** was set, and there was a blocking record lock.
EAGAIN	Total amount of system memory available when reading using raw I / O is temporarily insufficient.
EAGAIN	No data is waiting to be read on a file associated with a tty device and **O_NONBLOCK** was set.
EAGAIN	No message is waiting to be read on a stream and **O_NDELAY** or **O_NONBLOCK** was set.
EBADF	The file descriptor determined by the **DIR** stream is no longer

valid. This results if the **DIR** stream has been closed.

EBADMSG Message waiting to be read on a stream is not a data message.

EDEADLK The **read( )** was going to go to sleep and cause a deadlock to occur.

EFAULT *buf* points to an illegal address.

EINTR A signal was caught during the **read( )** or **readv( )** function.

EINVAL Attempted to read from a stream linked to a multiplexor.

EIO A physical I/O error has occurred, or the process is in a background process group and is attempting to read from its controlling terminal, and either the process is ignoring or blocking the **SIGTTIN** signal or the process group of the process is orphaned.

ENOENT The current file pointer for the directory is not located at a valid entry.

ENOLCK The system record lock table was full, so the **read( )** or **readv( )** could not go to sleep until the blocking record lock was removed.

ENOLINK *fildes* is on a remote machine and the link to that machine is no longer active.

ENXIO The device associated with *fildes* is a block special or character special file and the value of the file pointer is out of range.

telldir( ), **seekdir( )**, and **closedir( )** return 0 on success and will fail if one or more of the following are true:

EBADF The file descriptor determined by the **DIR** stream is no longer valid. This results if the **DIR** stream has been closed.

EXAMPLES

Here is a sample program that prints the names of all the files in the current directory:

```
#include <stdio.h>
#include <dirent.h>

main( )
{
        DIR *dirp;
        struct dirent *direntp;

        dirp = opendir( "." );
        while ( (direntp = readdir( dirp )) != NULL )
                (void)printf( "%s\n", direntp->d_name );
        (void)closedir( dirp );
        return (0);
}
```

SEE ALSO

getdents(2), **dirent**(4)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

readdir( ) is unsafe in multi-thread applications. **readdir_r( )** is safe, and should be used instead. **closedir( )**, **directory( )**, **opendir( )**, **rewinddir( )**, **seekdir( )**, and **telldir( )** are safe in multi-thread applications.

The POSIX **readdir_r( )** interface is as specified in POSIX 1003.1c Draft #10.

NAME

dirname – report the parent directory name of a file path name

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

char ∗**dirname(char** ∗*path***);**

MT-LEVEL

MT-Safe

DESCRIPTION

Given a pointer to a null-terminated character string that contains a file system path name, **dirname( )** returns a string that is the parent directory of that file. In doing this, it may place a null byte in the path name after the next to last element, so the content of *path* must be disposable. The returned string should not be deallocated by the caller. Trailing ‘‘/’’ characters in the path are not counted as part of the path.

If *path* or ∗*path* is zero, a pointer to a static constant ‘‘**.**’’ is returned.

dirname( ) and **basename( )** together yield a complete path name. **dirname (***path***)** is the directory where **basename (***path***)** is found.

EXAMPLES

A simple file name and the strings ‘‘**.**’’ and ‘‘**..**’’ all have ‘‘**.**’’ as their return value.

Input string	Output pointer
/usr/lib	**/usr**
/usr/	**/**
usr	**.**
/	**/**
.	**.**
..	**.**

The following code reads a path name, changes directory to the parent directory of the named file (see **chdir**(2)), and opens the file.

```
char path[100], *pathcopy;
int fd;
gets (path);
pathcopy = strdup (path);
chdir (dirname (pathcopy) );
free (pathcopy);
fd = open (basename (path), O_RDONLY);
```

SEE ALSO

basename(1), **chdir**(2), **basename**(3G)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

div, ldiv, lldiv – compute the quotient and remainder

SYNOPSIS

```
#include <stdlib.h>
div_t div(int numer, int denom);
ldiv_t ldiv(long int numer, long int denom);
lldiv_t lldiv(long long numer, long long denom);
```

MT-LEVEL

MT-Safe

DESCRIPTION

div( ) computes the quotient and remainder of the division of the numerator *numer* by the denominator *denom*. This function provides a well-defined semantics for the signed integral division and remainder operations, unlike the implementation-defined semantics of the built-in operations. The sign of the resulting quotient is that of the algebraic quotient, and, if the division is inexact, the magnitude of the resulting quotient is the largest integer less than the magnitude of the algebraic quotient. If the result cannot be represented, the behavior is undefined; otherwise, *quotient* * *denom* + *remainder* will equal *numer*.

ldiv( ) and **lldiv( )** are similar to **div( )**, except that the arguments and the members of the returned structure are different. **ldiv( )** returns a structure of type **ldiv_t** and has type **long int**. **lldiv( )** returns a structure of type **lldiv_t** and has type **long long**.

RETURN VALUES

div( ) returns a structure of type **div_t**, comprising both the quotient and remainder:

```
int quot; /*quotient*/
int rem; /*remainder*/
```

ldiv( ) returns a structure of type **ldiv_t** and **lldiv( )** returns a structure of type **lldiv_t**, comprising both the quotient and remainder:

```
long int quot; /*quotient*/
long int rem; /*remainder*/
```

NAME

dladdr – translate address to symbolic information

SYNOPSIS

cc [*flag* ...] *file* ... **–ldl** [*library* ...]

#include <dlfcn.h>

int dladdr(void * *address***, Dl_info** * *dlip***);**

MT-LEVEL

MT-Safe

DESCRIPTION

dladdr( ) is one of a family of routines that give the user direct access to the dynamic linking facilities. (SEE *Linker and Libraries Guide*). These routines are made available via the library loaded when the option **-ldl** is passed to the link-editor.

These routines are available to dynamically linked processes ONLY.

dladdr( ) determines if the specified *address* is located within one of the mapped objects that make up the current applications address space. An address is deemed to fall within a mapped object when it is between the base address, and the *_end* address of that object. If a mapped object fits this criteria, the symbol table made available to the run-time linker is searched to locate the nearest symbol to the specified address. The nearest symbol is one that has a value less than or equal to the required address.

The **Dl_info** structure must be preallocated by the user. The structure members are filled in by **dladdr( )** based on the specified *address*. The **Dl_info** structure includes the following members:

```
const char *    dli_fname;
void *          dli_fbase;
const char *    dli_sname;
void *          dli_saddr;
```

Descriptions of these members appear below.

dli_fname Contains a pointer to the filename of the containing object.

dli_fbase Contains the base address of the containing object.

dli_sname Contains a pointer to the nearest symbol name to the specified address. This symbol either has the same address, or is the nearest symbol with a lower address.

dli_saddr Contains the actual address of the above symbol.

RETURN VALUES

If the specified *address* cannot be matched to a mapped object, a 0 is returned. Otherwise a non-zero return is made and the associated **Dl_info** elements are filled.

SEE ALSO

ld(1), **dlclose**(3X), **dlerror**(3X), **dlopen**(3X), **dlsym**(3X)
Linker and Libraries Guide

NOTES

The **Dl_info** pointer elements point to addresses within the mapped objects, these may become invalid if objects are removed prior to these elements being used (see **dlclose( )**).

If no symbol is found to describe the specified address, both the **dli_sname** and **dli_saddr** members are set to 0.

NAME

dlclose – close a shared object

SYNOPSIS

cc [*flag* ...] *file* ... **–ldl** [*library* ...]

#include <dlfcn.h>

int dlclose(void ∗*handle***);**

MT-LEVEL

MT-Safe

DESCRIPTION

dlclose( ) is one of a family of routines that give the user direct access to the dynamic linking facilities. (SEE *Linker and Libraries Guide*). These routines are made available via the library loaded when the option **-ldl** is passed to the link-editor.

These routines are available to dynamically linked processes ONLY.

dlclose( ) disassociates a shared object previously opened by **dlopen( )** from the current process. Once an object has been closed using **dlclose( )**, its symbols are no longer available to **dlsym( )**. All objects loaded automatically as a result of invoking **dlopen( )** on the referenced object are also closed. *handle* is the value returned by a previous invocation of **dlopen( )**.

RETURN VALUES

If the referenced object was successfully closed, **dlclose( )** returns 0. If the object could not be closed, or if *handle* does not refer to an open object, **dlclose( )** returns a non-0 value. More detailed diagnostic information will be available through **dlerror( )**.

SEE ALSO

ld(1), **dladdr**(3X), **dlerror**(3X), **dlopen**(3X), **dlsym**(3X)
Linker and Libraries Guide

NOTES

A successful invocation of **dlclose( )** does not guarantee that the objects associated with *handle* will actually be removed from the address space of the process. Objects loaded by one invocation of **dlopen( )** may also be loaded by another invocation of **dlopen( )**. The same object may also be opened multiple times. An object will not be removed from the address space until all references to that object through an explicit **dlopen( )** invocation have been closed and all other objects implicitly referencing that object have also been closed.

Once an object has been closed by **dlclose( )**, referencing symbols contained in that object can cause undefined behavior.

NAME

dlerror – get diagnostic information

SYNOPSIS

cc [*flag* ...] *file* ... **–ldl** [*library* ...]

#include <dlfcn.h>

char ∗**dlerror(void);**

MT-LEVEL

MT-Safe

DESCRIPTION

dlerror( ) is one of a family of routines that give the user direct access to the dynamic linking facilities. (SEE *Linker and Libraries Guide*). These routines are made available via the library loaded when the option **-ldl** is passed to the link-editor.

These routines are available to dynamically linked processes ONLY.

dlerror( ) returns a null-terminated character string (with no trailing newline) that describes the last error that occurred during dynamic linking processing. If no dynamic linking errors have occurred since the last invocation of **dlerror( )**, **dlerror( )** returns **NULL**. Thus, invoking **dlerror( )** a second time, immediately following a prior invocation, will result in **NULL** being returned.

SEE ALSO

ld(1), **dladdr**(3X), **dlclose**(3X), **dlopen**(3X), **dlsym**(3X)
Linker and Libraries Guide

NOTES

The messages returned by **dlerror( )** may reside in a static buffer that is overwritten on each call to **dlerror( )**. Application code should not write to this buffer. Programs wishing to preserve an error message should make their own copies of that message.

NAME

dlopen – open a shared object

SYNOPSIS

cc [*flag* ...] *file* ... **–ldl** [*library* ...]

#include <dlfcn.h>

void * **dlopen(const char** * *pathname***, int** *mode***);**

MT-LEVEL

MT-Safe

DESCRIPTION

dlopen( ) is one of a family of routines that give the user direct access to the dynamic linking facilities. (see *Linker and Libraries Guide*). These routines are made available via the library loaded when the option **-ldl** is passed to the link-editor.

These routines are available to dynamically linked processes ONLY.

dlopen( ) makes a shared object available to a running process. **dlopen( )** returns to the process a "handle" which the process may use on subsequent calls to **dlsym( )** and **dlclose( )**. The value of this *handle* should not be interpreted in any way by the process. *pathname* is the path name of the object to be opened. A path name containing an embedded **'/'** is interpreted as an absolute path or relative to the current directory, otherwise the set of search paths currently in effect by the run-time linker will be used to locate the specified file (see *NOTES* section below).

If the value of *pathname* is 0, **dlopen( )** makes the symbols contained in the original **a.out**, any objects loaded at program startup with the **a.out**, and any objects that were loaded using **dlopen( )** together with the **RTLD_GLOBAL** flag, available through **dlsym( )**.

When a shared object is brought into the address space of a process, it may contain references to symbols whose addresses are not known until the object is loaded. These references must be relocated before the symbols can be accessed. The *mode* parameter governs when these relocations take place and may have the following values:

RTLD_LAZY Only references to data symbols are relocated when the object is first loaded. References to functions are not relocated until a given function is invoked for the first time. This *mode* should improve performance, since a process may not reference all of the functions in any given shared object. This behavior mimics the normal loading of shared object dependencies by a dynamic executable during process initialization.

RTLD_NOW All necessary relocations are performed when the object is first loaded. This may waste some processing, if relocations are performed for functions that are never referenced. This behavior may be useful for applications that need to know as soon as an object is loaded that all symbols referenced during execution will be available.

Any object loaded by **dlopen( )** that requires relocations against global symbols can reference the symbols in the **a.out**, any objects loaded at program startup, from the object itself, and from any dependencies the object references. By default, the relocations of an

object loaded by one **dlopen( )** invocation may not reference symbols from objects loaded by a different **dlopen( )** invocation. However, the *mode* parameter may also be **ored** with the following value to effect the scope of symbol availability:

RTLD_GLOBAL The objects symbols are made available for the relocation processing of any other object. In addition, symbol lookup using **dlopen(0,** *mode***)** and an associated **dlsym( )**, allows objects loaded with this *mode* to be searched.

RETURN VALUES

If *pathname* cannot be found, cannot be opened for reading, is not a shared or relocatable object, or if an error occurs during the process of loading *pathname* or relocating its symbolic references, **dlopen( )** will return **NULL**. More detailed diagnostic information will be available through **dlerror( )**.

SEE ALSO

ld(1), **dladdr**(3X), **dlclose**(3X), **dlerror**(3X), **dlsym**(3X)
Linker and Libraries Guide

NOTES

If other shared objects were link edited with *pathname* when *pathname* was built (i.e., the *pathname* has dependencies on other shared objects), those objects will automatically be loaded by **dlopen( )**. The directory search path used to find both *pathname* and the other *needed* objects may be affected by setting the environment variable **LD_LIBRARY_PATH** (which is analyzed once at process startup), or from a run-path setting within the application or the shared object from which the **dlopen( )** originated. These search rules will only be applied to pathnames that do not contain an embedded **'/'**. Objects whose names resolve to the same absolute or relative path name may be opened any number of times using **dlopen( )**, however, the object referenced will only be loaded once into the address space of the current process.

Some symbols defined in dynamic executables or shared objects may not be available to the runtime linker. The symbol table created by **ld** for use by the runtime linker might contain only a subset of the symbols defined in the object.

NAME

dlsym – get the address of a symbol in a shared object

SYNOPSIS

cc [*flag* ...] *file* ... **–ldl** [*library* ...]

#include <dlfcn.h>

void **handle*, **const char** **name*);

MT-LEVEL

MT-Safe

DESCRIPTION

dlsym( ) is one of a family of routines that give the user direct access to the dynamic linking facilities. (SEE *Linker and Libraries Guide*). These routines are made available via the library loaded when the option **-ldl** is passed to the link-editor.

These routines are available to dynamically linked processes ONLY.

dlsym( ) allows a process to obtain the address of a symbol defined within a shared object. *handle* is either the value returned from a call to **dlopen( )**, or the special flag **RTLD_NEXT**. *name* is the symbol's name as a character string.

In the former case the corresponding shared object must not have been closed using **dlclose( )**. **dlsym( )** will search for the named symbol in all shared objects loaded automatically as a result of loading the object referenced by *handle* (see **dlopen( )).**

In the latter case **dlsym( )** will search for the named symbol in the objects that were loaded following the object from which the **dlsym( )** call is being made. If these subsequent objects were loaded from **dlopen( )** calls, **dlsym( )** will search the object only if the caller is part of the same **dlopen( )** dependency hierarchy, or the object was given global search access (see **dlopen( )** with reference to the mode **RTLD_GLOBAL).**

RETURN VALUES

If *handle* does not refer to a valid object opened by **dlopen( )**, is not the special flag **RTLD_NEXT**, or if the named symbol cannot be found within any of the objects associated with *handle*, **dlsym( )** will return **NULL**. More detailed diagnostic information will be available through **dlerror( )**.

EXAMPLES

The following example shows how one can use **dlopen( )** and **dlsym( )** to access either function or data objects. For simplicity, error checking has been omitted.

```
void    *handle;
int     *iptr, (*fptr)(int);

/* open the needed object */
handle = dlopen("/usr/home/me/libfoo.so.1", RTLD_LAZY);

/* find the address of function and data objects */
fptr = (int (*)(int))dlsym(handle, "my_function");
iptr = (int *)dlsym(handle, "my_object");

/* invoke function, passing value of integer as a parameter */
(*fptr)(*iptr);
```

SEE ALSO

ld(1), **dladdr**(3X), **dlclose**(3X), **dlerror**(3X), **dlopen**(3X)
Linker and Libraries Guide

NAME

doconfig – execute a configuration script

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

 # include <sac.h>

int doconfig(int *fildes***, char** **script***, long** *rflag***);**

MT-LEVEL

Unsafe

DESCRIPTION

doconfig( ) is a Service Access Facility library function that interprets the configuration scripts contained in the files **</etc/saf/***pmtag***/_config>**, **</etc/saf/_sysconfig>**, and **</etc/saf/***pmtag* / *svctag*>, where *pmtag* specifies the tag associated with the port monitor, and *svctag* specifies the service tag associated with a given service. See **pmadm**(1M) and **sacadm**(1M).

script is the name of the configuration script; *fildes* is a file descriptor that designates the stream to which stream manipulation operations are to be applied; *rflag* is a bitmask that indicates the mode in which *script* is to be interpreted. If *rflag* is zero, all commands in the configuration script are eligible to be interpreted. If *rflag* has the **NOASSIGN** bit set, the **assign** command is considered illegal and will generate an error return. If *rflag* has the **NORUN** bit set, the **run** and **runwait** commands are considered illegal and will generate error returns.

The configuration language in which *script* is written consists of a sequence of commands, each of which is interpreted separately. The following reserved keywords are defined: **assign**, **push**, **pop**, **runwait**, and **run**. The comment character is **#**; when a **#** occurs on a line, everything from that point to the end of the line is ignored. Blank lines are not significant. No line in a command script may exceed 1024 characters.

assign *variable=value*

Used to define environment variables. *variable* is the name of the environment variable and *value* is the value to be assigned to it. The value assigned must be a string constant; no form of parameter substitution is available. *value* may be quoted. The quoting rules are those used by the shell for defining environment variables. **assign** will fail if space cannot be allocated for the new variable or if any part of the specification is invalid.

push *module1*[, *module2, module3*, . . .]

Used to push STREAMS modules onto the stream designated by *fildes*. *module1* is the name of the first module to be pushed, *module2* is the name of the second module to be pushed, etc. The command will fail if any of the named modules cannot be pushed. If a module cannot be pushed, the subsequent modules on the same command line will be ignored and modules that have already been pushed will be popped.

pop [*module*]

Used to pop STREAMS modules off the designated stream. If **pop** is invoked with no arguments, the top module on the stream is popped. If an argument is given,

modules will be popped one at a time until the named module is at the top of the stream. If the named module is not on the designated stream, the stream is left as it was and the command fails. If *module* is the special keyword ALL, then all modules on the stream will be popped. Note that only modules above the top-most driver are affected.

runwait *command*

The **runwait** command runs a command and waits for it to complete. *command* is the pathname of the command to be run. The command is run with **/usr/bin/sh −c** prepended to it; shell scripts may thus be executed from configuration scripts. The **runwait** command will fail if *command* cannot be found or cannot be executed, or if *command* exits with a non-zero status.

run *command*

The **run** command is identical to **runwait** except that it does not wait for *command* to complete. *command* is the pathname of the command to be run. **run** will not fail unless it is unable to create a child process to execute the command.

Although they are syntactically indistinguishable, some of the commands available to **run** and **runwait** are interpreter built-in commands. Interpreter built-ins are used when it is necessary to alter the state of a process within the context of that process. The **doconfig( )** interpreter built-in commands are similar to the shell special commands and, like these, they do not spawn another process for execution. See **sh**(1). The built-in commands are:

cd
ulimit
umask

RETURN VALUES

doconfig( ) returns **0** if the script was interpreted successfully. If a command in the script fails, the interpretation of the script ceases at that point and a positive number is returned; this number indicates which line in the script failed. If a system error occurs, a value of **−1** is returned. When a script fails, the process whose environment was being established should *not* be started.

SEE ALSO

sh(1), **pmadm**(1M), **sacadm**(1M)

NOTES

This interface is unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

drand48, erand48, lrand48, nrand48, mrand48, jrand48, srand48, seed48, lcong48 – generate uniformly distributed pseudo-random numbers

SYNOPSIS

```
#include <stdlib.h>
double drand48(void);
double erand48(unsigned short xsubi[3]);
long lrand48(void);
long nrand48(unsigned short xsubi[3]);
long mrand48(void);
long jrand48(unsigned short xsubi[3]);
void srand48(long seedval);
unsigned short *seed48(unsigned short seed16v[3]);
void lcong48(unsigned short param[7]);
```

MT-LEVEL

Safe

DESCRIPTION

This family of functions generates pseudo-random numbers using the well-known linear congruential algorithm and 48-bit integer arithmetic.

Functions **drand48( )** and **erand48( )** return non-negative double-precision floating-point values uniformly distributed over the interval [0.0, 1.0).

Functions **lrand48( )** and **nrand48( )** return non-negative long integers uniformly distributed over the interval $[0, 2^{31})$.

Functions **mrand48( )** and **jrand48( )** return signed long integers uniformly distributed over the interval $[-2^{31}, 2^{31})$.

Functions **srand48( )**, **seed48( )**, and **lcong48( )** are initialization entry points, one of which should be invoked before either **drand48( )**, **lrand48( )**, or **mrand48( )** is called. (Although it is not recommended practice, constant default initializer values will be supplied automatically if **drand48( )**, **lrand48( )**, or **mrand48( )** is called without a prior call to an initialization entry point.) Functions **erand48( ), nrand48( ),** and **jrand48( )** do not require an initialization entry point to be called first.

All the routines work by generating a sequence of 48-bit integer values, X_i, according to the linear congruential formula

$$X_{n+1} = (aX_n + c)_{\mathrm{mod}\ m} \qquad n \geq 0.$$

The parameter $m = 2^{48}$; hence 48-bit integer arithmetic is performed. Unless **lcong48( )** has been invoked, the multiplier value a and the addend value c are given by

$$a = 5\mathrm{DEECE66D}_{16} = 273673163155_8$$
$$c = \mathrm{B}_{16} = 13_8.$$

The value returned by any of the functions **drand48()**, **erand48()**, **lrand48()**, **nrand48()**, **mrand48()**, or **jrand48()** is computed by first generating the next 48-bit X_i in the sequence. Then the appropriate number of bits, according to the type of data item to be returned, are copied from the high-order (leftmost) bits of X_i and transformed into the returned value.

The functions **drand48()**, **lrand48()**, and **mrand48()** store the last 48-bit X_i generated in an internal buffer. X_i must be initialized prior to being invoked. The functions **erand48()**, **nrand48()**, and **jrand48()** require the calling program to provide storage for the successive X_i values in the array specified as an argument when the functions are invoked. These routines do not have to be initialized; the calling program must place the desired initial value of X_i into the array and pass it as an argument. By using different arguments, functions **erand48()**, **nrand48()**, and **jrand48()** allow separate modules of a large program to generate several *independent* streams of pseudo-random numbers, that is, the sequence of numbers in each stream will *not* depend upon how many times the routines have been called to generate numbers for the other streams.

The initializer function **srand48()** sets the high-order 32 bits of X_i to the 32 bits contained in its argument. The low-order 16 bits of X_i are set to the arbitrary value $330E_{16}$.

The initializer function **seed48()** sets the value of X_i to the 48-bit value specified in the argument array. In addition, the previous value of X_i is copied into a 48-bit internal buffer, used only by **seed48()**, and a pointer to this buffer is the value returned by **seed48()**. This returned pointer, which can just be ignored if not needed, is useful if a program is to be restarted from a given point at some future time — use the pointer to get at and store the last X_i value, and then use this value to reinitialize using **seed48()** when the program is restarted.

The initialization function **lcong48()** allows the user to specify the initial X_i, the multiplier value *a*, and the addend value *c*. Argument array elements *param[0-2]* specify X_i, *param[3-5]* specify the multiplier *a*, and *param[6]* specifies the 16-bit addend *c*. After **lcong48()** has been called, a subsequent call to either **srand48()** or **seed48()** will restore the ''standard'' multiplier and addend values, *a* and *c*, specified above.

SEE ALSO **rand**(3C)

NAME

dup2 – duplicate an open file descriptor

SYNOPSIS

```
#include <unistd.h>
int dup2(int fildes, int fildes2);
```

MT-LEVEL

Unsafe

Async-Signal-Safe

DESCRIPTION

dup2( ) causes the file descriptor *fildes2* to refer to the same file as *fildes*. *fildes* is a file descriptor referring to an open file, and *fildes2* is a non-negative integer less than the current value for the maximum number of open file descriptors allowed the calling process (see **getrlimit**(2)). If *fildes2* already referred to an open file, not *fildes,* it is closed first. If *fildes2* refers to *fildes*, or if *fildes* is not a valid open file descriptor, *fildes2* will not be closed first.

RETURN VALUES

Upon successful completion a non-negative integer, namely, the file descriptor, is returned. Otherwise, a value of **–1** is returned and **errno** is set to indicate the error.

ERRORS

dup2( ) will fail if one or more of the following are true:

EBADF	*fildes* is not a valid open file descriptor.
EINTR	a signal was caught during the **dup2( )** call.
EMFILE	the process has too many open files (see **fcntl**(2)).

SEE ALSO

close(2), **creat**(2), **exec**(2), **fcntl**(2), **getrlimit**(2), **open**(2), **pipe**(2), **lockf**(3C)

NAME

econvert, fconvert, gconvert, seconvert, sfconvert, sgconvert, qeconvert, qfconvert, qgconvert, ecvt, fcvt, gcvt – output conversion

SYNOPSIS

```
#include <floatingpoint.h>
char *econvert(double value, int ndigit, int *decpt, int *sign, char *buf);
char *fconvert(double value, int ndigit, int *decpt, int *sign, char *buf);
char *gconvert(double value, int ndigit, int trailing, char *buf);
char *seconvert(single *value, int ndigit, int *decpt, int *sign, char *buf);
char *sfconvert(single *value, int ndigit, int *decpt, int *sign, char *buf);
char *sgconvert(single *value, int ndigit, int trailing, char *buf );
char *qeconvert(quadruple *value, int ndigit, int *decpt, int *sign, char *buf);
char *qfconvert(quadruple *value, int ndigit, int *decpt, int *sign, char *buf);
char *qgconvert(quadruple *value, int ndigit, int trailing, char *buf );
char *ecvt(double value, int ndigit, int *decpt, int *sign);
char *fcvt(double value, int ndigit, int *decpt, int *sign);
char *gcvt(double value, int ndigit, char *buf);
```

MT-LEVEL

MT-Safe

DESCRIPTION

econvert( ) converts the *value* to a NULL-terminated string of *ndigit* ASCII digits in *buf* and returns a pointer to *buf*. *buf* should contain at least *ndigit+1* characters. The position of the decimal point relative to the beginning of the string is stored indirectly through *decpt*. Thus *buf* == "314" and **decpt* == 1 corresponds to the numerical value 3.14, while *buf* == "314" and **decpt* == –1 corresponds to the numerical value .0314. If the sign of the result is negative, the word pointed to by *sign* is nonzero; otherwise it is zero. The least significant digit is rounded.

fconvert works much like **econvert**, except that the correct digit has been rounded as if for **sprintf(%w.nf)** output with *n=ndigit* digits to the right of the decimal point. *ndigit* can be negative to indicate rounding to the left of the decimal point. The return value is a pointer to *buf*. *buf* should contain at least *310+max(0,ndigit)* characters to accomodate any double-precision *value*.

gconvert( ) converts the *value* to a NULL-terminated ASCII string in *buf* and returns a pointer to *buf*. It produces *ndigit* significant digits in fixed-decimal format, like **sprintf(%w.nf)**, if possible, and otherwise in floating-decimal format, like **sprintf(%w.ne)**; in either case *buf* is ready for printing, with sign and exponent. The result corresponds to that obtained by

(void) sprintf(buf,"%w.ng",value) ;

If *trailing*= 0, trailing zeros and a trailing point are suppressed, as in **sprintf(%g)**. If *trailing*!= 0, trailing zeros and a trailing point are retained, as in **sprintf(%#g)**.

seconvert, **sfconvert**, and **sgconvert()** are single-precision versions of these functions, and are more efficient than the corresponding double-precision versions. A pointer rather than the value itself is passed to avoid C's usual conversion of single-precision arguments to double.

qeconvert, **qfconvert**, and **qgconvert()** are quadruple-precision versions of these functions. **qfconvert()** can overflow the *decimal_record* field *ds* if *value* is too large. In that case, *buf[0]* is set to zero.

ecvt() and **fcvt()** are obsolete versions of **econvert()** and **fconvert()** that create a string in a static data area, overwritten by each call, and return values that point to that static data. These functions are therefore not reentrant.

gcvt() is an obsolete version of **gconvert()** that always suppresses trailing zeros and point.

IEEE Infinities and NaNs are treated similarly by these functions. ''NaN'' is returned for NaN, and ''Inf'' or ''Infinity'' for Infinity. The longer form is produced when *ndigit* >= 8.

SEE ALSO

sprintf(3S)

NAME

ecvt, fcvt, gcvt – convert floating-point number to string

SYNOPSIS

```
#include <stdlib.h>
char *ecvt(double value, int ndigit, int *decpt, int *sign);
char *fcvt(double value, int ndigit, int *decpt, int *sign);
char *gcvt(double value, int ndigit, char *buf);
```

MT-LEVEL

Unsafe

DESCRIPTION

ecvt( ) converts *value* to a null-terminated string of *ndigit* digits and returns a pointer thereto. The high-order digit is non-zero, unless the value is zero. The low-order digit is rounded. The position of the decimal point relative to the beginning of the string is stored indirectly through *decpt* (negative means to the left of the returned digits). The decimal point is not included in the returned string. If the sign of the result is negative, the word pointed to by *sign* is non-zero, otherwise it is zero.

fcvt( ) is identical to **ecvt( )**, except that the correct digit has been rounded for **printf %f** output of the number of digits specified by *ndigit*.

gcvt( ) converts the *value* to a null-terminated string in the array pointed to by *buf* and returns *buf*. It attempts to produce *ndigit* significant digits in **%f** format if possible, otherwise **%e** format (scientific notation), ready for printing. A minus sign, if there is one, or a decimal point will be included as part of the returned string. Trailing zeros are suppressed.

SEE ALSO

printf(3S)

NOTES

The values returned by **ecvt( )** and **fcvt( )** point to a single static data array whose content is overwritten by each call.

NAME

elf – object file access library

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

MT-LEVEL

Unsafe

DESCRIPTION

Functions in the ELF access library let a program manipulate ELF (Executable and Linking Format) object files, archive files, and archive members. The header provides type and function declarations for all library services.

Programs communicate with many of the higher-level routines using an *ELF descriptor*. That is, when the program starts working with a file, **elf_begin**(3E) creates an ELF descriptor through which the program manipulates the structures and information in the file. These ELF descriptors can be used both to read and to write files. After the program establishes an ELF descriptor for a file, it may then obtain *section descriptors* to manipulate the sections of the file (see **elf_getscn**(3E)). Sections hold the bulk of an object file's real information, such as text, data, the symbol table, and so on. A section descriptor ''belongs'' to a particular ELF descriptor, just as a section belongs to a file. Finally, *data descriptors* are available through section descriptors, allowing the program to manipulate the information associated with a section. A data descriptor ''belongs'' to a section descriptor.

Descriptors provide private handles to a file and its pieces. In other words, a data descriptor is associated with one section descriptor, which is associated with one ELF descriptor, which is associated with one file. Although descriptors are private, they give access to data that may be shared. Consider programs that combine input files, using incoming data to create or update another file. Such a program might get data descriptors for an input and an output section. It then could update the output descriptor to reuse the input descriptor's data. That is, the descriptors are distinct, but they could share the associated data bytes. This sharing avoids the space overhead for duplicate buffers and the performance overhead for copying data unnecessarily.

FILE CLASSES

ELF provides a framework in which to define a family of object files, supporting multiple processors and architectures. An important distinction among object files is the *class*, or capacity, of the file. The 32-bit class supports architectures in which a 32-bit object can represent addresses, file sizes, etc., as in the following.

Name	Purpose
Elf32_Addr	Unsigned address
Elf32_Half	Unsigned medium integer
Elf32_Off	Unsigned file offset
Elf32_Sword	Signed large integer
Elf32_Word	Unsigned large integer
unsigned char	Unsigned small integer

Other classes will be defined as necessary, to support larger (or smaller) machines. Some library services deal only with data objects for a specific class, while others are class-independent. To make this distinction clear, library function names reflect their status, as described below.

DATA REPRESENTATION

Conceptually, two parallel sets of objects support cross compilation environments. One set corresponds to file contents, while the other set corresponds to the native memory image of the program manipulating the file. Type definitions supplied by the headers work on the native machine, which may have different data encodings (size, byte order, etc.) than the target machine. Although native memory objects should be at least as big as the file objects (to avoid information loss), they may be bigger if that is more natural for the host machine.

Translation facilities exist to convert between file and memory representations. Some library routines convert data automatically, while others leave conversion as the program's responsibility. Either way, programs that create object files must write file-typed objects to those files; programs that read object files must take a similar view. See **elf32_xlatetof**(3E) and **elf32_fsize**(3E) for more information.

Programs may translate data explicitly, taking full control over the object file layout and semantics. If the program prefers not to have and exercise complete control, the library provides a higher-level interface that hides many object file details. **elf_begin()** and related functions let a program deal with the native memory types, converting between memory objects and their file equivalents automatically when reading or writing an object file.

ELF VERSIONS

Object file versions allow ELF to adapt to new requirements. *Three independent versions* can be important to a program. First, an application program knows about a particular version by virtue of being compiled with certain headers. Second, the access library similarly is compiled with header files that control what versions it understands. Third, an ELF object file holds a value identifying its version, determined by the ELF version known by the file's creator. Ideally, all three versions would be the same, but they may differ.

If a program's version is newer than the access library, the program might use information unknown to the library. Translation routines might not work properly, leading to undefined behavior. This condition merits installing a new library.

The library's version might be newer than the program's and the file's. The library understands old versions, thus avoiding compatibility problems in this case.

Finally, a file's version might be newer than either the program or the library understands. The program might or might not be able to process the file properly, depending on whether the file has extra information and whether that information can be safely ignored. Again, the safe alternative is to install a new library that understands the file's version.

To accommodate these differences, a program must use **elf_version**(3E) to pass its version to the library, thus establishing the *working version* for the process. Using this, the library accepts data from and presents data to the program in the proper representations. When the library reads object files, it uses each file's version to interpret the data. When writing files or converting memory types to the file equivalents, the library uses the program's working version for the file data.

SYSTEM SERVICES

As mentioned above, **elf_begin( )** and related routines provide a higher-level interface to ELF files, performing input and output on behalf of the application program. These routines assume a program can hold entire files in memory, without explicitly using temporary files. When reading a file, the library routines bring the data into memory and perform subsequent operations on the memory copy. Programs that wish to read or write large object files with this model must execute on a machine with a large process virtual address space. If the underlying operating system limits the number of open files, a program can use **elf_cntl**(3E) to retrieve all necessary data from the file, allowing the program to close the file descriptor and reuse it.

Although the **elf_begin( )** interfaces are convenient and efficient for many programs, they might be inappropriate for some. In those cases, an application may invoke the **elf32_xlatetom**(3E) or **elf32_xlatetof**(3E) data translation routines directly. These routines perform no input or output, leaving that as the application's responsibility. By assuming a larger share of the job, an application controls its input and output model.

LIBRARY NAMES

Names associated with the library take several forms.

elf_*name*	These class-independent names perform some service, *name*, for the program.
elf32_*name*	Service names with an embedded class, **32** here, indicate they work only for the designated class of files.
Elf_*Type*	Data types can be class-independent as well, distinguished by *Type*.
Elf32_*Type*	Class-dependent data types have an embedded class name, **32** here.
ELF_C_*CMD*	Several functions take commands that control their actions. These values are members of the **Elf_Cmd** enumeration; they range from zero through **ELF_C_NUM**–1**.**
ELF_F_*FLAG*	Several functions take flags that control library status and / or actions. Flags are bits that may be combined.
ELF32_FSZ_*TYPE*	These constants give the file sizes in bytes of the basic ELF types for the 32-bit class of files. See **elf32_fsize( )** for more information.
ELF_K_*KIND*	The function **elf_kind( )** identifies the *KIND* of file associated with an ELF descriptor. These values are members of the **Elf_Kind** enumeration; they range from zero through **ELF_K_NUM**–1**.**

ELF_T_*TYPE* When a service function, such as **elf32_xlatetom( )** or **elf32_xlatetof( )**, deals with multiple types, names of this form specify the desired *TYPE*. Thus, for example, **ELF_T_EHDR** is directly related to **Elf32_Ehdr**. These values are members of the **Elf_Type** enumeration; they range from zero through **ELF_T_NUM**–1**.**

EXAMPLES

The basic interpretation of an ELF file consists of:

- opening an ELF object file
- obtaining an ELF descriptor
- analyzing the file using the descriptor.

The following example opens the file, obtains the ELF descriptor, and prints out the names of each section in the file.

```
#include     <fcntl.h>
#include     <stdio.h>
#include     <libelf.h>
#include     <stdlib.h>
#include     <string.h>
static void failure(void);

void
main(int argc, char ** argv)
{
        Elf32_Shdr *    shdr;
        Elf32_Ehdr *    ehdr;
        Elf *           elf;
        Elf_Scn *       scn;
        Elf_Data *      data;
        int             fd;
        unsigned int    cnt;

                /* Open the input file */
        if ((fd = open(argv[1], O_RDONLY)) == -1)
                exit(1);

                /* Obtain the ELF descriptor */
        (void) elf_version(EV_CURRENT);
        if ((elf = elf_begin(fd, ELF_C_READ, NULL)) == NULL)
                failure();

                /* Obtain the .shstrtab data buffer */
        if (((ehdr = elf32_getehdr(elf)) == NULL) ||
           ((scn = elf_getscn(elf, ehdr->e_shstrndx)) == NULL) ||
           ((data = elf_getdata(scn, NULL)) == NULL))
                failure();
```

```
                /* Traverse input filename, printing each section */
        for (cnt = 1, scn = NULL; scn = elf_nextscn(elf, scn); cnt++) {
                if ((shdr = elf32_getshdr(scn)) == NULL)
                        failure();

        (void) printf("[%d]%s0, cnt,
                (char *)data->d_buf + shdr->sh_name);
        }
}               /* end main */

static void
failure()
{
        (void) fprintf(stderr, "%s0, elf_errmsg(elf_errno()));
        exit(1);
}
```

Below is sample output from compiling and executing the above code, which prints the names of the sections using itself as the input file

```
% cc -o elfprint example.c -lelf
% elfprint elfprint
 [1]    .interp
 [2]    .hash
 [3]    .dynsym
 [4]    .dynstr
 [5]    .rela.ex_shared
 [6]    .rela.bss
 [7]    .rela.plt
 [8]    .text
 [9]    .init
 [10]   .fini
 [11]   .exception_ranges
 ...
```

SEE ALSO

elf32_fsize(3E), **elf32_getshdr**(3E), **elf32_xlatetof**(3E), **elf_begin**(3E), **elf_cntl**(3E), **elf_errmsg**(3E), **elf_fill**(3E), **elf_getarhdr**(3E), **elf_getarsym**(3E), **elf_getbase**(3E), **elf_getdata**(3E), **elf_getident**(3E), **elf_getscn**(3E), **elf_hash**(3E), **elf_kind**(3E), **elf_memory**(3E), **elf_rawfile**(3E), **elf_strptr**(3E), **elf_update**(3E), **elf_version**(3E), **ar**(4)

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NOTES

Information in the ELF headers is separated into common parts and processor-specific parts. A program can make a processor's information available by including the appropriate header: **<sys/elf_***NAME***.h>** where *NAME* matches the processor name as used in the ELF file header.

Name	Processor
M32	AT&T WE 32100
SPARC	SPARC
386	Intel 80386
486	Intel 80486
860	Intel 80860
68K	Motorola 68000
88K	Motorola 88000

Other processors will be added to the table as necessary.

To illustrate, a program could use the following code to ''see'' the processor-specific information for the SPARC.

```
#include <libelf.h>
#include <sys/elf_SPARC.h>
```

Without the **<sys/elf_SPARC.h>** definition, only the common ELF information would be visible.

A program could use the following code to ''see'' the processor-specific information for the Intel 80386:

```
#include <libelf.h>
#include <sys/elf_386.h>
```

Without the **<sys/elf_386.h>** definition, only the common ELF information would be visible.

Although reading the objects is rather straightforward, writing / updating them can corrupt the shared offsets among sections. Upon creation, relationships are established among the sections that must be maintained even if the object's size is changed.

NAME

elf32_fsize – return the size of an object file type

SYNOPSIS

cc [*flag* ...] *file* ... –**lelf** [*library* ...]

#include <libelf.h>

size_t elf32_fsize(Elf_Type *type***, size_t** *count***, unsigned** *ver***);**

MT-LEVEL

Unsafe

DESCRIPTION

elf32_fsize( ) gives the size in bytes of the 32-bit file representation of *count* data objects with the given *type*. The library uses version *ver* to calculate the size (see **elf**(3E) and **elf_version**(3E)).

Constant values are available for the sizes of fundamental types:

Elf_Type	File Size	Memory Size
ELF_T_ADDR	**ELF32_FSZ_ADDR**	**sizeof(Elf32_Addr)**
ELF_T_BYTE	**1**	**sizeof(unsigned char)**
ELF_T_HALF	**ELF32_FSZ_HALF**	**sizeof(Elf32_Half)**
ELT_T_OFF	**ELF32_FSZ_OFF**	**sizeof(Elf32_Off)**
ELF_T_SWORD	**ELF32_FSZ_SWORD**	**sizeof(Elf32_Sword)**
ELF_T_WORD	**ELF32_FSZ_WORD**	**sizeof(Elf32_Word)**

elf32_fsize( ) returns zero if the value of *type* or *ver* is unknown. See **elf32_xlatetof**(3E) for a list of the *type* values.

SEE ALSO

elf(3E), **elf32_xlatetof**(3E), **elf_version**(3E)

NAME

elf32_getehdr, elf32_newehdr – retrieve class-dependent object file header

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

Elf32_Ehdr ∗**elf32_getehdr(Elf** ∗*elf***);**

Elf32_Ehdr ∗**elf32_newehdr(Elf** ∗*elf***);**

MT-LEVEL

Unsafe

DESCRIPTION

For a 32-bit class file, **elf32_getehdr( )** returns a pointer to an ELF header, if one is available for the ELF descriptor *elf*. If no header exists for the descriptor, **elf32_newehdr( )** allocates a ''clean'' one, but it otherwise behaves the same as **elf32_getehdr( )**. It does not allocate a new header if one exists already. If no header exists (for **elf32_getehdr( )**), one cannot be created (for **elf32_newehdr( )**), a system error occurs, the file is not a 32-bit class file, or *elf* is null, both functions return a null pointer.

The header includes the following members.

```
unsigned char   e_ident[EI_NIDENT];
Elf32_Half      e_type;
Elf32_Half      e_machine;
Elf32_Word      e_version;
Elf32_Addr      e_entry;
Elf32_Off       e_phoff;
Elf32_Off       e_shoff;
Elf32_Word      e_flags;
Elf32_Half      e_ehsize;
Elf32_Half      e_phentsize;
Elf32_Half      e_phnum;
Elf32_Half      e_shentsize;
Elf32_Half      e_shnum;
Elf32_Half      e_shstrndx;
```

elf32_newehdr( ) automatically sets the **ELF_F_DIRTY** bit (see **elf_flagdata**(3E)). A program may use **elf_getident( )** to inspect the identification bytes from a file.

SEE ALSO

elf(3E), **elf_begin**(3E), **elf_flagdata**(3E), **elf_getident**(3E)

NAME

elf32_getphdr, elf32_newphdr – retrieve class-dependent program header table

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

Elf32_Phdr ∗**elf32_getphdr(Elf** ∗*elf***);**

Elf32_Phdr ∗**elf32_newphdr(Elf** ∗*elf***, size_t** *count***);**

MT-LEVEL

Unsafe

DESCRIPTION

For a 32-bit class file, **elf32_getphdr( )** returns a pointer to the program execution header table, if one is available for the ELF descriptor *elf*.

elf32_newphdr( ) allocates a new table with *count* entries, regardless of whether one existed previously, and sets the **ELF_F_DIRTY** bit for the table (see **elf_flagdata**(3E)). Specifying a zero *count* deletes an existing table. Note this behavior differs from that of **elf32_newehdr( )** (see **elf32_getehdr**(3E)), allowing a program to replace or delete the program header table, changing its size if necessary.

If no program header table exists, the file is not a 32-bit class file, an error occurs, or *elf* is null, both functions return a null pointer. Additionally, **elf32_newphdr( )** returns a null pointer if *count* is zero.

The table is an array of **Elf32_Phdr** structures, each of which includes the following members.

```
Elf32_Word    p_type;
Elf32_Off     p_offset;
Elf32_Addr    p_vaddr;
Elf32_Addr    p_paddr;
Elf32_Word    p_filesz;
Elf32_Word    p_memsz;
Elf32_Word    p_flags;
Elf32_Word    p_align;
```

The ELF header's **e_phnum** member tells how many entries the program header table has (see **elf32_getehdr**(3E)). A program may inspect this value to determine the size of an existing table; **elf32_newphdr( )** automatically sets the member's value to *count*. If the program is building a new file, it is responsible for creating the file's ELF header before creating the program header table.

SEE ALSO

elf(3E), **elf32_getehdr**(3E), **elf_begin**(3E), **elf_flagdata**(3E)

NAME

elf32_getshdr – retrieve class-dependent section header

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

Elf32_Shdr ∗**elf32_getshdr(Elf_Scn** ∗*scn***);**

MT-LEVEL

Unsafe

DESCRIPTION

For a 32-bit class file, **elf32_getshdr( )** returns a pointer to a section header for the section descriptor *scn*. Otherwise, the file is not a 32-bit class file, *scn* was null, or an error occurred; **elf32_getshdr( )** then returns null.

The header includes the following members.

```
Elf32_Word      sh_name;
Elf32_Word      sh_type;
Elf32_Word      sh_flags;
Elf32_Addr      sh_addr;
Elf32_Off       sh_offset;
Elf32_Word      sh_size;
Elf32_Word      sh_link;
Elf32_Word      sh_info;
Elf32_Word      sh_addralign;
Elf32_Word      sh_entsize;
```

If the program is building a new file, it is responsible for creating the file's ELF header before creating sections.

SEE ALSO

elf(3E), **elf_flagdata**(3E), **elf_getscn**(3E), **elf_strptr**(3E)

NAME

elf32_xlatetof, elf32_xlatetom – class-dependent data translation

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

Elf_Data ∗**elf32_xlatetof(Elf_Data** ∗*dst*, **const Elf_Data** ∗*src*, **unsigned** *encode*);

Elf_Data ∗**elf32_xlatetom(Elf_Data** ∗*dst*, **const Elf_Data** ∗*src*, **unsigned** *encode*);

MT-LEVEL

Unsafe

DESCRIPTION

elf32_xlatetom( ) translates various data structures from their 32-bit class file representations to their memory representations; **elf32_xlatetof( )** provides the inverse. This conversion is particularly important for cross development environments. *src* is a pointer to the source buffer that holds the original data; *dst* is a pointer to a destination buffer that will hold the translated copy. *encode* gives the byte encoding in which the file objects are to be represented and must have one of the encoding values defined for the ELF header's **e_ident[EI_DATA]** entry (see **elf_getident**(3E)). If the data can be translated, the functions return *dst*. Otherwise, they return null because an error occurred, such as incompatible types, destination buffer overflow, etc.

elf_getdata(3E) describes the **Elf_Data** descriptor, which the translation routines use as follows.

d_buf Both the source and destination must have valid buffer pointers.

d_type This member's value specifies the type of the data to which **d_buf** points and the type of data to be created in the destination. The program supplies a **d_type** value in the source; the library sets the destination's **d_type** to the same value. These values are summarized below.

d_size This member holds the total size, in bytes, of the memory occupied by the source data and the size allocated for the destination data. If the destination buffer is not large enough, the routines do not change its original contents. The translation routines reset the destination's **d_size** member to the actual size required, after the translation occurs. The source and destination sizes may differ.

d_version This member holds version number of the objects (desired) in the buffer. The source and destination versions are independent.

Translation routines allow the source and destination buffers to coincide. That is, **dst→d_buf** may equal **src→d_buf**. Other cases where the source and destination buffers overlap give undefined behavior.

Elf_Type	32-Bit Memory Type
ELF_T_ADDR	Elf32_Addr
ELF_T_BYTE	unsigned char
ELF_T_DYN	Elf32_Dyn
ELF_T_EHDR	Elf32_Ehdr
ELF_T_HALF	Elf32_Half
ELT_T_OFF	Elf32_Off
ELF_T_PHDR	Elf32_Phdr
ELF_T_REL	Elf32_Rel
ELF_T_RELA	Elf32_Rela
ELF_T_SHDR	Elf32_Shdr
ELF_T_SWORD	Elf32_Sword
ELF_T_SYM	Elf32_Sym
ELF_T_WORD	Elf32_Word

Translating buffers of type **ELF_T_BYTE** does not change the byte order.

SEE ALSO

elf(3E), **elf32_fsize**(3E), **elf_getdata**(3E), **elf_getident**(3E)

NAME

elf_begin, elf_end, elf_memory, elf_next, elf_rand – process ELF object files

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

Elf ∗**elf_begin(int** *fildes***, Elf_Cmd** *cmd***, Elf** ∗*ref***);**

int elf_end(Elf ∗*elf***);**

Elf ∗**elf_memory(char** ∗*image***, size_t***sz***);**

Elf_Cmd elf_next(Elf ∗*elf***);**

size_t elf_rand(Elf ∗*elf***, size_t** *offset***);**

MT-LEVEL

Unsafe

DESCRIPTION

elf_begin( ), **elf_end( )**, **elf_memory( )**, **elf_next( )**, and **elf_rand( )** work together to process Executable and Linking Format (ELF) object files, either individually or as members of archives. After obtaining an ELF descriptor from **elf_begin( )** or **elf_memory( )**, the program may read an existing file, update an existing file, or create a new file. *fildes* is an open file descriptor that **elf_begin( )** uses for reading or writing. *elf* is an ELF descriptor previously returned from **elf_begin( )**. The initial file offset (see **lseek**(2)) is unconstrained, and the resulting file offset is undefined.

cmd may have the following values.

ELF_C_NULL When a program sets *cmd* to this value, **elf_begin( )** returns a null pointer, without opening a new descriptor. *ref* is ignored for this command. See the examples below for more information.

ELF_C_READ When a program wishes to examine the contents of an existing file, it should set *cmd* to this value. Depending on the value of *ref*, this command examines archive members or entire files. Three cases can occur.

First, if *ref* is a null pointer, **elf_begin( )** allocates a new ELF descriptor and prepares to process the entire file. If the file being read is an archive, **elf_begin( )** also prepares the resulting descriptor to examine the initial archive member on the next call to **elf_begin( )**, as if the program had used **elf_next( )** or **elf_rand( )** to ''move'' to the initial member.

Second, if *ref* is a non-null descriptor associated with an archive file, **elf_begin( )** lets a program obtain a separate ELF descriptor associated with an individual member. The program should have used **elf_next( )** or **elf_rand( )** to position *ref* appropriately (except for the initial member, which **elf_begin( )** prepares; see the example below). In this case, *fildes* should be the same file descriptor used for the parent archive.

Finally, if *ref* is a non-null ELF descriptor that is not an archive, **elf_begin( )** increments the number of activations for the descriptor and returns *ref*, without allocating a new descriptor and without changing the descriptor's read / write permissions. To terminate the descriptor for *ref*, the program must call **elf_end( )** once for each activation. See the examples below for more information.

ELF_C_RDWR This command duplicates the actions of **ELF_C_READ** and additionally allows the program to update the file image (see **elf_update**(3E)). That is, using **ELF_C_READ** gives a read-only view of the file, while **ELF_C_RDWR** lets the program read *and* write the file. **ELF_C_RDWR** is not valid for archive members. If *ref* is non-null, it must have been created with the **ELF_C_RDWR** command.

ELF_C_WRITE If the program wishes to ignore previous file contents, presumably to create a new file, it should set *cmd* to this value. *ref* is ignored for this command.

elf_begin( ) ''works'' on all files (including files with zero bytes), providing it can allocate memory for its internal structures and read any necessary information from the file. Programs reading object files thus may call **elf_kind**(3E) or **elf32_getehdr**(3E) to determine the file type (only object files have an ELF header). If the file is an archive with no more members to process, or an error occurs, **elf_begin( )** returns a null pointer. Otherwise, the return value is a non-null ELF descriptor.

Before the first call to **elf_begin( )**, a program must call **elf_version( )** to coordinate versions.

elf_end( ) is used to terminate an ELF descriptor, *elf*, and to deallocate data associated with the descriptor. Until the program terminates a descriptor, the data remain allocated. A null pointer is allowed as an argument, to simplify error handling. If the program wishes to write data associated with the ELF descriptor to the file, it must use **elf_update( )** before calling **elf_end( )**.

Calling **elf_end( )** removes one activation and returns the remaining activation count. The library does not terminate the descriptor until the activation count reaches zero. Consequently, a 0 return value indicates the ELF descriptor is no longer valid.

elf_memory( ) returns a pointer to an Elf descriptor, the Elf image has read operations enabled (**ELF_C_READ**). *image* is a pointer to an image of the Elf file mapped into memory, *sz* is the size of the Elf image. You may read and modify an Elf image that is mapped in with **elf_memory( )**, but you may not change the Elf image size.

elf_next( ) provides sequential access to the next archive member. That is, having an ELF descriptor, *elf*, associated with an archive member, **elf_next( )** prepares the containing archive to access the following member when the program calls **elf_begin( )**. After successfully positioning an archive for the next member, **elf_next( )** returns the value **ELF_C_READ**. Otherwise, the open file was not an archive, *elf* was null, or an error occurred, and the return value is **ELF_C_NULL**. In either case, the return value may be passed as an argument to **elf_begin( )**, specifying the appropriate action.

elf_rand( ) provides random archive processing, preparing *elf* to access an arbitrary archive member. *elf* must be a descriptor for the archive itself, not a member within the archive. *offset* gives the byte offset from the beginning of the archive to the archive header of the desired member. See **elf_getarsym**(3E) for more information about archive member offsets. When **elf_rand( )** works, it returns *offset*. Otherwise it returns 0, because an error occurred, *elf* was null, or the file was not an archive (no archive member can have a zero offset). A program may mix random and sequential archive processing.

SYSTEM SERVICES

When processing a file, the library decides when to read or write the file, depending on the program's requests. Normally, the library assumes the file descriptor remains usable for the life of the ELF descriptor. If, however, a program must process many files simultaneously and the underlying operating system limits the number of open files, the program can use **elf_cntl( )** to let it reuse file descriptors. After calling **elf_cntl( )** with appropriate arguments, the program may close the file descriptor without interfering with the library.

All data associated with an ELF descriptor remain allocated until **elf_end( )** terminates the descriptor's last activation. After the descriptors have been terminated, the storage is released; attempting to reference such data gives undefined behavior. Consequently, a program that deals with multiple input (or output) files must keep the ELF descriptors active until it finishes with them.

EXAMPLES

A prototype for reading a file appears
on the next page. If the file is a simple object file, the program executes the loop one time, receiving a null descriptor in the second iteration. In this case, both **elf** and **arf** will have the same value, the activation count will be two, and the program calls **elf_end( )** twice to terminate the descriptor.

If the file is an archive, the loop processes each archive member in turn, ignoring those that are not object files.

```
if (elf_version(EV_CURRENT) == EV_NONE)
{
        /* library out of date */
        /* recover from error */
}
cmd = ELF_C_READ;
arf = elf_begin(fildes, cmd, (Elf *)0);
while ((elf = elf_begin(fildes, cmd, arf)) != 0)
{
        if ((ehdr = elf32_getehdr(elf)) != 0)
        {
                /* process the file ... */
        }
        cmd = elf_next(elf);
        elf_end(elf);
}
elf_end(arf);
```

Alternatively, the next example illustrates random archive processing. After identifying the file as an archive, the program repeatedly processes archive members of interest. For clarity, this example omits error checking and ignores simple object files. Additionally, this fragment preserves the ELF descriptors for all archive members, because it does not call **elf_end( )** to terminate them.

```
elf_version(EV_CURRENT);
arf = elf_begin(fildes, ELF_C_READ, (Elf *)0);
if (elf_kind(arf) != ELF_K_AR)
{
        /* not an archive */
}
/* initial processing */
/* set offset = ... for desired member header */
while (elf_rand(arf, offset) == offset)
{
        if ((elf = elf_begin(fildes, ELF_C_READ, arf)) == 0)
                break;
        if ((ehdr = elf32_getehdr(elf)) != 0)
        {
                /* process archive member ... */
        }
        /* set offset = ... for desired member header */
}
```

An archive starts with a ''magic string'' that has **SARMAG** bytes; the initial archive member follows immediately. An application could thus provide the following function to rewind an archive (the function returns –1 for errors and 0 otherwise).

```
#include <ar.h>
#include <libelf.h>

int
rewindelf(Elf *elf)
{
    if (elf_rand(elf, (size_t)SARMAG) == SARMAG)
        return 0;
    return –1;
}
```

The following outline shows how one might create a new ELF file. This example is simplified to show the overall flow.

```
elf_version(EV_CURRENT);
fildes = open("path/name", O_RDWR | O_TRUNC | O_CREAT, 0666);
if ((elf = elf_begin(fildes, ELF_C_WRITE, (Elf *)0)) == 0)
        return;
ehdr = elf32_newehdr(elf);
phdr = elf32_newphdr(elf, count);
scn = elf_newscn(elf);
shdr = elf32_getshdr(scn);
data = elf_newdata(scn);
elf_update(elf, ELF_C_WRITE);
elf_end(elf);
```

Finally, the following outline shows how one might update an existing ELF file. Again, this example is simplified to show the overall flow.

```
elf_version(EV_CURRENT);
fildes = open("path/name", O_RDWR);
elf = elf_begin(fildes, ELF_C_RDWR, (Elf *)0);

/* add new or delete old information */
...

/* ensure that the memory image of the file is complete */
elf_update(elf, ELF_C_NULL);

elf_update(elf, ELF_C_WRITE);   /* update file */
elf_end(elf);
```

Notice that both file creation examples open the file with write *and* read permissions. On systems that support **mmap**, the library uses it to enhance performance, and **mmap** requires a readable file descriptor. Although the library can use a write-only file

descriptor, the application will not obtain the performance advantages of **mmap**.

SEE ALSO

creat(2), **lseek**(2), **mmap**(2), **open**(2), **elf**(3E), **elf32_getehdr**(3E), **elf_cntl**(3E), **elf_getarhdr**(3E), **elf_getarsym**(3E), **elf_getbase**(3E), **elf_getdata**(3E), **elf_getscn**(3E), **elf_kind**(3E), **elf_rawfile**(3E), **elf_update**(3E), **elf_version**(3E), **ar**(4)

NAME

elf_cntl – control an elf file descriptor

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

int elf_cntl(Elf **elf*, **Elf_Cmd** *cmd*);

MT-LEVEL

Unsafe

DESCRIPTION

elf_cntl( ) instructs the library to modify its behavior with respect to an ELF descriptor, *elf*. As **elf_begin**(3E) describes, an ELF descriptor can have multiple activations, and multiple ELF descriptors may share a single file descriptor. Generally, **elf_cntl( )** commands apply to all activations of *elf*. Moreover, if the ELF descriptor is associated with an archive file, descriptors for members within the archive will also be affected as described below. Unless stated otherwise, operations on archive members do not affect the descriptor for the containing archive.

The *cmd* argument tells what actions to take and may have the following values.

ELF_C_FDDONE

This value tells the library not to use the file descriptor associated with *elf*. A program should use this command when it has requested all the information it cares to use and wishes to avoid the overhead of reading the rest of the file. The memory for all completed operations remains valid, but later file operations, such as the initial **elf_getdata( )** for a section, will fail if the data are not in memory already.

ELF_C_FDREAD

This command is similar to **ELF_C_FDDONE**, except it forces the library to read the rest of the file. A program should use this command when it must close the file descriptor but has not yet read everything it needs from the file. After **elf_cntl( )** completes the **ELF_C_FDREAD** command, future operations, such as **elf_getdata( )**, will use the memory version of the file without needing to use the file descriptor.

If **elf_cntl( )** succeeds, it returns zero. Otherwise *elf* was null or an error occurred, and the function returns –1.

SEE ALSO

elf(3E), **elf_begin**(3E), **elf_getdata**(3E), **elf_rawfile**(3E)

NOTES

If the program wishes to use the ‘‘raw’’ operations (see **elf_rawdata( )**, which **elf_getdata**(3E) describes, and **elf_rawfile**(3E)) after disabling the file descriptor with **ELF_C_FDDONE** or **ELF_C_FDREAD**, it must execute the raw operations explicitly beforehand. Otherwise, the raw file operations will fail. Calling **elf_rawfile( )** makes the entire image available, thus supporting subsequent **elf_rawdata( )** calls.

NAME

elf_errmsg, elf_errno – error handling

SYNOPSIS

cc [*flag* ...] *file* ... –**lelf** [*library* ...]

```
#include <libelf.h>
const char *elf_errmsg (int err);
int elf_errno(void);
```

MT-LEVEL

Unsafe

DESCRIPTION

If an ELF library function fails, a program may call **elf_errno( )** to retrieve the library's internal error number. As a side effect, this function resets the internal error number to zero, which indicates no error.

elf_errmsg( ) takes an error number, *err*, and returns a null-terminated error message (with no trailing new-line) that describes the problem. A zero *err* retrieves a message for the most recent error. If no error has occurred, the return value is a null pointer (not a pointer to the null string). Using *err* of –1 also retrieves the most recent error, except it guarantees a non-null return value, even when no error has occurred. If no message is available for the given number, **elf_errmsg( )** returns a pointer to an appropriate message. This function does not have the side effect of clearing the internal error number.

EXAMPLES

The following fragment clears the internal error number and checks it later for errors. Unless an error occurs after the first call to **elf_errno( )**, the next call will return zero.

```
(void)elf_errno( );
/* processing ... */
while (more_to_do)
{
        if ((err = elf_errno( )) != 0)
        {
                /* print msg */
                msg = elf_errmsg(err);
        }
}
```

SEE ALSO

elf(3E)

NAME

elf_fill – set fill byte

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

void elf_fill(int *fill***);**

MT-LEVEL

Unsafe

DESCRIPTION

Alignment constraints for ELF files sometimes require the presence of ''holes.'' For example, if the data for one section are required to begin on an eight-byte boundary, but the preceding section is too ''short,'' the library must fill the intervening bytes. These bytes are set to the *fill* character. The library uses zero bytes unless the application supplies a value. See **elf_getdata**(3E) for more information about these holes.

SEE ALSO

elf(3E), **elf_flagdata**(3E), **elf_getdata**(3E), **elf_update**(3E)

NOTES

An application can assume control of the object file organization by setting the **ELF_F_LAYOUT** bit (see **elf_flagdata**(3E)). When this is done, the library does *not* fill holes.

NAME

elf_flagdata, elf_flagehdr, elf_flagelf, elf_flagphdr, elf_flagscn, elf_flagshdr – manipulate flags

SYNOPSIS

cc [*flag* ...] *file* ... –**lelf** [*library* ...]

#include <libelf.h>

unsigned elf_flagdata(Elf_Data **data*, **Elf_Cmd** *cmd*, **unsigned** *flags*);

unsigned elf_flagehdr(Elf **elf*, **Elf_Cmd** *cmd*, **unsigned** *flags*);

unsigned elf_flagelf(Elf **elf*, **Elf_Cmd** *cmd*, **unsigned** *flags*);

unsigned elf_flagphdr(Elf **elf*, **Elf_Cmd** *cmd*, **unsigned** *flags*);

unsigned elf_flagscn(Elf_Scn **scn*, **Elf_Cmd** *cmd*, **unsigned** *flags*);

unsigned elf_flagshdr(Elf_Scn **scn*, **Elf_Cmd** *cmd*, **unsigned** *flags*);

MT-LEVEL

Unsafe

DESCRIPTION

These functions manipulate the flags associated with various structures of an ELF file. Given an ELF descriptor (*elf*), a data descriptor (*data*), or a section descriptor (*scn*), the functions may set or clear the associated status bits, returning the updated bits. A null descriptor is allowed, to simplify error handling; all functions return zero for this degenerate case.

cmd may have the following values:

ELF_C_CLR The functions clear the bits that are asserted in *flags*. Only the non-zero bits in *flags* are cleared; zero bits do not change the status of the descriptor.

ELF_C_SET The functions set the bits that are asserted in *flags*. Only the non-zero bits in *flags* are set; zero bits do not change the status of the descriptor.

Descriptions of the defined *flags* bits appear below:

ELF_F_DIRTY When the program intends to write an ELF file, this flag asserts the associated information needs to be written to the file. Thus, for example, a program that wished to update the ELF header of an existing file would call **elf_flagehdr( )** with this bit set in *flags* and *cmd* equal to **ELF_C_SET**. A later call to **elf_update( )** would write the marked header to the file.

ELF_F_LAYOUT Normally, the library decides how to arrange an output file. That is, it automatically decides where to place sections, how to align them in the file, etc. If this bit is set for an ELF descriptor, the program assumes responsibility for determining all file positions. This bit is meaningful only for **elf_flagelf( )** and applies to the entire file associated with the descriptor.

When a flag bit is set for an item, it affects all the subitems as well. Thus, for example, if the program sets the **ELF_F_DIRTY** bit with **elf_flagelf( )**, the entire logical file is ''dirty.''

EXAMPLES

The following fragment shows how one might mark the ELF header to be written to the output file:

```
/* dirty ehdr ... */
ehdr = elf32_getehdr(elf);
elf_flagehdr(elf, ELF_C_SET, ELF_F_DIRTY);
```

SEE ALSO

elf(3E), **elf32_getehdr**(3E), **elf_getdata**(3E), **elf_update**(3E)

NAME

elf_getarhdr – retrieve archive member header

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

Elf_Arhdr ∗**elf_getarhdr(Elf** ∗*elf***);**

MT-LEVEL

Unsafe

DESCRIPTION

elf_getarhdr( ) returns a pointer to an archive member header, if one is available for the ELF descriptor *elf*. Otherwise, no archive member header exists, an error occurred, or *elf* was null; **elf_getarhdr( )** then returns a null value. The header includes the following members.

```
char          *ar_name;
time_t        ar_date;
long          ar_uid;
long          ar_gid;
unsigned long ar_mode;
off_t         ar_size;
char          *ar_rawname;
```

An archive member name, available through **ar_name**, is a null-terminated string, with the **ar** format control characters removed. The **ar_rawname** member holds a null-terminated string that represents the original name bytes in the file, including the terminating slash and trailing blanks as specified in the archive format.

In addition to ''regular'' archive members, the archive format defines some special members. All special member names begin with a slash (/), distinguishing them from regular members (whose names may not contain a slash). These special members have the names (**ar_name**) defined below.

/ This is the archive symbol table. If present, it will be the first archive member. A program may access the archive symbol table through **elf_getarsym( )**. The information in the symbol table is useful for random archive processing (see **elf_rand( )** on **elf_begin**(3E)).

// This member, if present, holds a string table for long archive member names. An archive member's header contains a 16-byte area for the name, which may be exceeded in some file systems. The library automatically retrieves long member names from the string table, setting **ar_name** to the appropriate value.

Under some error conditions, a member's name might not be available. Although this causes the library to set **ar_name** to a null pointer, the **ar_rawname** member will be set as usual.

SEE ALSO

elf(3E), **elf_begin**(3E), **elf_getarsym**(3E), **ar**(4)

NAME

elf_getarsym – retrieve archive symbol table

SYNOPSIS

cc [*flag* ...] *file* ... –**lelf** [*library* ...]

#include <libelf.h>

Elf_Arsym ∗**elf_getarsym(Elf** ∗*elf***, size_t** ∗*ptr***);**

MT-LEVEL

Unsafe

DESCRIPTION

elf_getarsym( ) returns a pointer to the archive symbol table, if one is available for the ELF descriptor *elf*. Otherwise, the archive doesn't have a symbol table, an error occurred, or *elf* was NULL; **elf_getarsym( )** then returns a NULL value. The symbol table is an array of structures that include the following members.

```
char          *as_name;
size_t        as_off;
unsigned long as_hash;
```

These members have the following semantics.

as_name A pointer to a NULL-terminated symbol name resides here.

as_off This value is a byte offset from the beginning of the archive to the member's header. The archive member residing at the given offset defines the associated symbol. Values in **as_off** may be passed as arguments to **elf_rand( )**. See **elf_begin**(3E) to access the desired archive member.

as_hash This is a hash value for the name, as computed by **elf_hash( )**.

If *ptr* is non-NULL, the library stores the number of table entries in the location to which *ptr* points. This value is set to zero when the return value is NULL. The table's last entry, which is included in the count, has a NULL **as_name**, a zero value for **as_off**, and **~0UL** for **as_hash**.

The hash value returned is guaranteed not to be the bit pattern of all ones (**~0UL**).

SEE ALSO

elf(3E), **elf_begin**(3E), **elf_getarhdr**(3E), **elf_hash**(3E), **ar**(4)

NAME

elf_getbase – get the base offset for an object file

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

off_t elf_getbase(Elf ∗*elf***);**

MT-LEVEL

Unsafe

DESCRIPTION

elf_getbase( ) returns the file offset of the first byte of the file or archive member associated with *elf*, if it is known or obtainable, and –1 otherwise. A null *elf* is allowed, to simplify error handling; the return value in this case is –1. The base offset of an archive member is the beginning of the member's information, *not* the beginning of the archive member header.

SEE ALSO

elf(3E), **elf_begin**(3E), **ar**(4)

NAME

elf_getdata, elf_newdata, elf_rawdata – get section data

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

Elf_Data ∗**elf_getdata(Elf_Scn** ∗*scn*, **Elf_Data** ∗*data*);

Elf_Data ∗**elf_newdata(Elf_Scn** ∗*scn*);

Elf_Data ∗**elf_rawdata(Elf_Scn** ∗*scn*, **Elf_Data** ∗*data*);

MT-LEVEL

Unsafe

DESCRIPTION

These functions access and manipulate the data associated with a section descriptor, *scn*. When reading an existing file, a section will have a single data buffer associated with it. A program may build a new section in pieces, however, composing the new data from multiple data buffers. For this reason, the data for a section should be viewed as a list of buffers, each of which is available through a data descriptor.

elf_getdata( ) lets a program step through a section's data list. If the incoming data descriptor, *data*, is null, the function returns the first buffer associated with the section. Otherwise, *data* should be a data descriptor associated with *scn*, and the function gives the program access to the next data element for the section. If *scn* is null or an error occurs, **elf_getdata( )** returns a null pointer.

elf_getdata( ) translates the data from file representations into memory representations (see **elf32_xlatetof**(3E)) and presents objects with memory data types to the program, based on the file's *class* (see **elf**(3E)). The working library version (see **elf_version**(3E)) specifies what version of the memory structures the program wishes **elf_getdata( )** to present.

elf_newdata( ) creates a new data descriptor for a section, appending it to any data elements already associated with the section. As described below, the new data descriptor appears empty, indicating the element holds no data. For convenience, the descriptor's type (**d_type** below) is set to **ELF_T_BYTE**, and the version (**d_version** below) is set to the working version. The program is responsible for setting (or changing) the descriptor members as needed. This function implicitly sets the **ELF_F_DIRTY** bit for the section's data (see **elf_flagdata**(3E)). If *scn* is null or an error occurs, **elf_newdata( )** returns a null pointer.

elf_rawdata( ) differs from **elf_getdata( )** by returning only uninterpreted bytes, regardless of the section type. This function typically should be used only to retrieve a section image from a file being read, and then only when a program must avoid the automatic data translation described below. Moreover, a program may not close or disable (see **elf_cntl**(3E)) the file descriptor associated with *elf* before the initial raw operation, because **elf_rawdata( )** might read the data from the file to ensure it doesn't interfere with **elf_getdata( )**. See **elf_rawfile**(3E) for a related facility that applies to the entire file. When **elf_getdata( )** provides the right translation, its use is recommended over **elf_rawdata( )**. If *scn* is null or an error occurs, **elf_rawdata( )** returns a null pointer.

The **Elf_Data** structure includes the following members:

```
void        *d_buf;
Elf_Type    d_type;
size_t      d_size;
off_t       d_off;
size_t      d_align;
unsigned    d_version;
```

These members are available for direct manipulation by the program. Descriptions appear below.

d_buf A pointer to the data buffer resides here. A data element with no data has a null pointer.

d_type This member's value specifies the type of the data to which **d_buf** points. A section's type determines how to interpret the section contents, as summarized below.

d_size This member holds the total size, in bytes, of the memory occupied by the data. This may differ from the size as represented in the file. The size will be zero if no data exist. (See the discussion of **SHT_NOBITS** below for more information.)

d_off This member gives the offset, within the section, at which the buffer resides. This offset is relative to the file's section, not the memory object's.

d_align This member holds the buffer's required alignment, from the beginning of the section. That is, **d_off** will be a multiple of this member's value. For example, if this member's value is four, the beginning of the buffer will be four-byte aligned within the section. Moreover, the entire section will be aligned to the maximum of its constituents, thus ensuring appropriate alignment for a buffer within the section and within the file.

d_version This member holds the version number of the objects in the buffer. When the library originally read the data from the object file, it used the working version to control the translation to memory objects.

DATA ALIGNMENT

As mentioned above, data buffers within a section have explicit alignment constraints. Consequently, adjacent buffers sometimes will not abut, causing ''holes'' within a section. Programs that create output files have two ways of dealing with these holes.

First, the program can use **elf_fill()** to tell the library how to set the intervening bytes. When the library must generate gaps in the file, it uses the fill byte to initialize the data there. The library's initial fill value is zero, and **elf_fill()** lets the application change that.

Second, the application can generate its own data buffers to occupy the gaps, filling the gaps with values appropriate for the section being created. A program might even use different fill values for different sections. For example, it could set text sections' bytes to no-operation instructions, while filling data section holes with zero. Using this technique, the library finds no holes to fill, because the application eliminated them.

SECTION AND MEMORY TYPES

elf_getdata() interprets sections' data according to the section type, as noted in the section header available through **elf32_getshdr()**. The following table shows the section types and how the library represents them with memory data types for the 32-bit file class. Other classes would have similar tables. By implication, the memory data types control translation by **elf32_xlatetof**(3E)**.**

Section Type	**Elf_Type**	32-Bit Type
SHT_DYNAMIC	**ELF_T_DYN**	**Elf32_Dyn**
SHT_DYNSYM	**ELF_T_SYM**	**Elf32_Sym**
SHT_HASH	**ELF_T_WORD**	**Elf32_Word**
SHT_NOBITS	**ELF_T_BYTE**	**unsigned char**
SHT_NOTE	**ELF_T_BYTE**	**unsigned char**
SHT_NULL	*none*	*none*
SHT_PROGBITS	**ELF_T_BYTE**	**unsigned char**
SHT_REL	**ELF_T_REL**	**Elf32_Rel**
SHT_RELA	**ELF_T_RELA**	**Elf32_Rela**
SHT_STRTAB	**ELF_T_BYTE**	**unsigned char**
SHT_SYMTAB	**ELF_T_SYM**	**Elf32_Sym**
SHT_SUNW_verdef	**ELF_T_VDEF**	**Elf32_Verdef**
SHT_SUNW_verneed		**ELF_T_VNEEDElf32_Verneed**
SHT_SUNW_versym	**ELF_T_HALF**	**Elf32_Versym**
other	**ELF_T_BYTE**	**unsigned char**

elf_rawdata() creates a buffer with type **ELF_T_BYTE**.

As mentioned above, the program's working version controls what structures the library creates for the application. The library similarly interprets section types according to the versions. If a section type belongs to a version newer than the application's working version, the library does not translate the section data. Because the application cannot know the data format in this case, the library presents an untranslated buffer of type **ELF_T_BYTE**, just as it would for an unrecognized section type.

A section with a special type, **SHT_NOBITS**, occupies no space in an object file, even when the section header indicates a non-zero size. **elf_getdata()** and **elf_rawdata()** work on such a section, setting the *data* structure to have a null buffer pointer and the type indicated above. Although no data are present, the **d_size** value is set to the size from the section header. When a program is creating a new section of type **SHT_NOBITS**, it should use **elf_newdata()** to add data buffers to the section. These empty data buffers should have the **d_size** members set to the desired size and the **d_buf** members set to null.

EXAMPLES

The following fragment obtains the string table that holds section names (ignoring error checking). See **elf_strptr**(3E) for a variation of string table handling.

```
ehdr = elf32_getehdr(elf);
scn = elf_getscn(elf, (size_t)ehdr→e_shstrndx);
shdr = elf32_getshdr(scn);
if (shdr→sh_type != SHT_STRTAB)
{
        /* not a string table */
}
data = 0;
if ((data = elf_getdata(scn, data)) == 0 || data→d_size == 0)
{
        /* error or no data */
}
```

The **e_shstrndx** member in an ELF header holds the section table index of the string table. The program gets a section descriptor for that section, verifies it is a string table, and then retrieves the data. When this fragment finishes, **data→d_buf** points at the first byte of the string table, and **data→d_size** holds the string table's size in bytes.

SEE ALSO

elf(3E), **elf32_getehdr**(3E), **elf32_getshdr**(3E), **elf32_xlatetof**(3E), **elf_cntl**(3E), **elf_fill**(3E), **elf_flagdata**(3E), **elf_getscn**(3E), **elf_rawfile**(3E), **elf_strptr**(3E) **elf_version**(3E)

NAME

elf_getident – retrieve file identification data

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

char ∗**elf_getident(Elf** ∗*elf***, size_t** ∗*ptr***);**

MT-LEVEL

Unsafe

DESCRIPTION

As **elf**(3E) explains, ELF provides a framework for various classes of files, where basic objects may have 32 bits, 64 bits, etc. To accommodate these differences, without forcing the larger sizes on smaller machines, the initial bytes in an ELF file hold identification information common to all file classes. Every ELF header's **e_ident** has **EI_NIDENT** bytes with the following interpretation.

e_ident Index	Value	Purpose
EI_MAG0	**ELFMAG0**	File identification
EI_MAG1	**ELFMAG1**	
EI_MAG2	**ELFMAG2**	
EI_MAG3	**ELFMAG3**	
EI_CLASS	**ELFCLASSNONE**	File class
	ELFCLASS32	
	ELFCLASS64	
EI_DATA	**ELFDATANONE**	Data encoding
	ELFDATA2LSB	
	ELFDATA2MSB	
EI_VERSION	**EV_CURRENT**	File version
7-15	**0**	Unused, set to zero

Other kinds of files (see **elf_kind**(3E)) also may have identification data, though they would not conform to **e_ident**.

elf_getident( ) returns a pointer to the file's ''initial bytes.'' If the library recognizes the file, a conversion from the file image to the memory image may occur. In any case, the identification bytes are guaranteed not to have been modified, though the size of the unmodified area depends on the file type. If *ptr* is non-null, the library stores the number of identification bytes in the location to which *ptr* points. If no data are present, *elf* is null, or an error occurs, the return value is a null pointer, with zero stored through *ptr*, if *ptr* is non-null.

SEE ALSO **elf**(3E), **elf32_getehdr**(3E), **elf_begin**(3E), **elf_kind**(3E), **elf_rawfile**(3E)

NAME

elf_getscn, elf_ndxscn, elf_newscn, elf_nextscn – get section information

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

Elf_Scn *elf_getscn(Elf ****elf*, size_t** *index***);**

size_t elf_ndxscn(Elf_Scn ****scn*);**

Elf_Scn *elf_newscn(Elf ****elf*);**

Elf_Scn *elf_nextscn(Elf ****elf*, Elf_Scn ****scn***);**

MT-LEVEL

Unsafe

DESCRIPTION

These functions provide indexed and sequential access to the sections associated with the ELF descriptor *elf*. If the program is building a new file, it is responsible for creating the file's ELF header before creating sections; see **elf32_getehdr**(3E).

elf_getscn( ) returns a section descriptor, given an *index* into the file's section header table. Note the first ''real'' section has index 1. Although a program can get a section descriptor for the section whose *index* is 0 (**SHN_UNDEF**, the undefined section), the section has no data and the section header is ''empty'' (though present). If the specified section does not exist, an error occurs, or *elf* is null, **elf_getscn( )** returns a null pointer.

elf_newscn( ) creates a new section and appends it to the list for *elf*. Because the **SHN_UNDEF** section is required and not ''interesting'' to applications, the library creates it automatically. Thus the first call to **elf_newscn( )** for an ELF descriptor with no existing sections returns a descriptor for section 1. If an error occurs or *elf* is null, **elf_newscn( )** returns a null pointer.

After creating a new section descriptor, the program can use **elf32_getshdr( )** to retrieve the newly created, ''clean'' section header. The new section descriptor will have no associated data (see **elf_getdata**(3E)). When creating a new section in this way, the library updates the **e_shnum** member of the ELF header and sets the **ELF_F_DIRTY** bit for the section (see **elf_flagdata**(3E)). If the program is building a new file, it is responsible for creating the file's ELF header (see **elf32_getehdr**(3E)) before creating new sections.

elf_nextscn( ) takes an existing section descriptor, *scn*, and returns a section descriptor for the next higher section. One may use a null *scn* to obtain a section descriptor for the section whose index is 1 (skipping the section whose index is **SHN_UNDEF**). If no further sections are present or an error occurs, **elf_nextscn( )** returns a null pointer.

elf_ndxscn( ) takes an existing section descriptor, *scn*, and returns its section table index. If *scn* is null or an error occurs, **elf_ndxscn( )** returns **SHN_UNDEF**.

EXAMPLES

An example of sequential access appears below. Each pass through the loop processes the next section in the file; the loop terminates when all sections have been processed.

```
scn = 0;
while ((scn = elf_nextscn(elf, scn)) != 0)
{
        /* process section */
}
```

SEE ALSO

elf(3E), **elf32_getehdr**(3E), **elf32_getshdr**(3E), **elf_begin**(3E), **elf_flagdata**(3E), **elf_getdata**(3E)

NAME

elf_hash – compute hash value

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

unsigned long elf_hash(const char ∗*name***);**

MT-LEVEL

Unsafe

DESCRIPTION

elf_hash( ) computes a hash value, given a null terminated string, *name*. The returned hash value, *h*, can be used as a bucket index, typically after computing *h* mod *x* to ensure appropriate bounds.

Hash tables may be built on one machine and used on another because **elf_hash( )** uses unsigned arithmetic to avoid possible differences in various machines' signed arithmetic. Although *name* is shown as **char∗** above, **elf_hash( )** treats it as **unsigned char∗** to avoid sign extension differences. Using **char∗** eliminates type conflicts with expressions such as **elf_hash(***name***)**.

ELF files' symbol hash tables are computed using this function (see **elf_getdata**(3E) and **elf32_xlatetof**(3E)). The hash value returned is guaranteed not to be the bit pattern of all ones **(**˜0UL**).**

SEE ALSO

elf(3E), **elf_getdata**(3E), **elf32_xlatetof**(3E)

NAME

elf_kind – determine file type

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

Elf_Kind elf_kind(Elf ∗*elf***);**

MT-LEVEL

Unsafe

DESCRIPTION

This function returns a value identifying the kind of file associated with an ELF descriptor (*elf*). Defined values are below:

ELF_K_AR The file is an archive [see **ar**(4)]. An ELF descriptor may also be associated with an archive *member*, not the archive itself, and then **elf_kind( )** identifies the member's type.

ELF_K_COFF The file is a COFF object file. **elf_begin**(3E) describes the library's handling for COFF files.

ELF_K_ELF The file is an ELF file. The program may use **elf_getident( )** to determine the class. Other functions, such as **elf32_getehdr( )**, are available to retrieve other file information.

ELF_K_NONE This indicates a kind of file unknown to the library.

Other values are reserved, to be assigned as needed to new kinds of files. *elf* should be a value previously returned by **elf_begin( )**. A null pointer is allowed, to simplify error handling, and causes **elf_kind( )** to return **ELF_K_NONE**.

SEE ALSO

elf(3E), **elf32_getehdr**(3E), **elf_begin**(3E), **elf_getident**(3E), **ar**(4)

NAME

elf_rawfile – retrieve uninterpreted file contents

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

char ∗**elf_rawfile(Elf** ∗*elf***, size_t** ∗*ptr***);**

MT-LEVEL

Unsafe

DESCRIPTION

elf_rawfile( ) returns a pointer to an uninterpreted byte image of the file. This function should be used only to retrieve a file being read. For example, a program might use **elf_rawfile( )** to retrieve the bytes for an archive member.

A program may not close or disable (see **elf_cntl**(3E)) the file descriptor associated with *elf* before the initial call to **elf_rawfile( ) ,** because **elf_rawfile( )** might have to read the data from the file if it does not already have the original bytes in memory. Generally, this function is more efficient for unknown file types than for object files. The library implicitly translates object files in memory, while it leaves unknown files unmodified. Thus, asking for the uninterpreted image of an object file may create a duplicate copy in memory.

elf_rawdata( ) (see **elf_getdata**(3E)) is a related function, providing access to sections within a file.

If *ptr* is non-null, the library also stores the file's size, in bytes, in the location to which *ptr* points. If no data are present, *elf* is null, or an error occurs, the return value is a null pointer, with zero stored through *ptr*, if *ptr* is non-null.

SEE ALSO

elf(3E), **elf32_getehdr**(3E), **elf_begin**(3E), **elf_cntl**(3E), **elf_getdata**(3E), **elf_getident**(3E), **elf_kind**(3E)

NOTES

A program that uses **elf_rawfile( )** and that also interprets the same file as an object file potentially has two copies of the bytes in memory. If such a program requests the raw image first, before it asks for translated information (through such functions as **elf32_getehdr( )**, **elf_getdata( )**, and so on), the library ''freezes'' its original memory copy for the raw image. It then uses this frozen copy as the source for creating translated objects, without reading the file again. Consequently, the application should view the raw file image returned by **elf_rawfile( )** as a read-only buffer, unless it wants to alter its own view of data subsequently translated. In any case, the application may alter the translated objects without changing bytes visible in the raw image.

Multiple calls to **elf_rawfile( )** with the same ELF descriptor return the same value; the library does not create duplicate copies of the file.

NAME

elf_strptr – make a string pointer

SYNOPSIS

cc [*flag* ...] *file* ... –**lelf** [*library* ...]

#include <libelf.h>

char ∗**elf_strptr(Elf** ∗*elf***, size_t** *section***, size_t** *offset***);**

MT-LEVEL

Unsafe

DESCRIPTION

This function converts a string section *offset* to a string pointer. *elf* identifies the file in which the string section resides, and *section* identifies the section table index for the strings. **elf_strptr( )** normally returns a pointer to a string, but it returns a null pointer when *elf* is null, *section* is invalid or is not a section of type **SHT_STRTAB**, the section data cannot be obtained, *offset* is invalid, or an error occurs.

EXAMPLES

A prototype for retrieving section names appears below. The file header specifies the section name string table in the **e_shstrndx** member. The following code loops through the sections, printing their names.

```
/* handle the error */
if ((ehdr = elf32_getehdr(elf)) == 0) {
        return;
}
ndx = ehdr→e_shstrndx;
scn = 0;
while ((scn = elf_nextscn(elf, scn)) != 0) {
        char  *name = 0;
        if ((shdr = elf32_getshdr(scn)) != 0)
                name = elf_strptr(elf, ndx, (size_t)shdr→sh_name);
        printf("'%s'\n", name? name: "(null)");
}
```

SEE ALSO

elf(3E), **elf32_getshdr**(3E), **elf32_xlatetof**(3E), **elf_getdata**(3E)

NOTES

A program may call **elf_getdata( )** to retrieve an entire string table section. For some applications, that would be both more efficient and more convenient than using **elf_strptr( )**.

NAME

elf_update – update an ELF descriptor

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

off_t elf_update(Elf **elf***, Elf_Cmd** *cmd***);**

MT-LEVEL

Unsafe

DESCRIPTION

elf_update( ) causes the library to examine the information associated with an ELF descriptor, *elf*, and to recalculate the structural data needed to generate the file's image.

cmd may have the following values:

ELF_C_NULL This value tells **elf_update( )** to recalculate various values, updating only the ELF descriptor's memory structures. Any modified structures are flagged with the **ELF_F_DIRTY** bit. A program thus can update the structural information and then reexamine them without changing the file associated with the ELF descriptor. Because this does not change the file, the ELF descriptor may allow reading, writing, or both reading and writing (see **elf_begin**(3E)).

ELF_C_WRITE If *cmd* has this value, **elf_update( )** duplicates its **ELF_C_NULL** actions and also writes any ''dirty'' information associated with the ELF descriptor to the file. That is, when a program has used **elf_getdata**(3E) or the **elf_flagdata**(3E) facilities to supply new (or update existing) information for an ELF descriptor, those data will be examined, coordinated, translated if necessary (see **elf32_xlatetof**(3E)), and written to the file. When portions of the file are written, any **ELF_F_DIRTY** bits are reset, indicating those items no longer need to be written to the file (see **elf_flagdata**(3E)). The sections' data are written in the order of their section header entries, and the section header table is written to the end of the file.

When the ELF descriptor was created with **elf_begin( )**, it must have allowed writing the file. That is, the **elf_begin( )** command must have been either **ELF_C_RDWR** or **ELF_C_WRITE**.

If **elf_update( )** succeeds, it returns the total size of the file image (not the memory image), in bytes. Otherwise an error occurred, and the function returns –1.

When updating the internal structures, **elf_update( )** sets some members itself. Members listed below are the application's responsibility and retain the values given by the program.

The following table shows ELF Header members:

Member	Notes
e_ident[EI_DATA]	Library controls other **e_ident** values
e_type	
e_machine	
e_version	
e_entry	
e_phoff	Only when **ELF_F_LAYOUT** asserted
e_shoff	Only when **ELF_F_LAYOUT** asserted
e_flags	
e_shstrndx	

The following table shows the Program Header members:

Member	Notes
p_type	The application controls all program header entries
p_offset	
p_vaddr	
p_paddr	
p_filesz	
p_memsz	
p_flags	
p_align	

The following table shows the Section Header members:

Member	Notes
sh_name	
sh_type	
sh_flags	
sh_addr	
sh_offset	Only when **ELF_F_LAYOUT** asserted
sh_size	Only when **ELF_F_LAYOUT** asserted
sh_link	
sh_info	
sh_addralign	Only when **ELF_F_LAYOUT** asserted
sh_entsize	

The following table shows the Data Descriptor members:

Member	Notes
d_buf	
d_type	
d_size	
d_off	Only when **ELF_F_LAYOUT** asserted
d_align	
d_version	

Note that the program is responsible for two particularly important members (among others) in the ELF header. The **e_version** member controls the version of data structures written to the file. If the version is **EV_NONE**, the library uses its own internal version. The **e_ident[EI_DATA]** entry controls the data encoding used in the file. As a special case, the value may be **ELFDATANONE** to request the native data encoding for the host machine. An error occurs in this case if the native encoding doesn't match a file encoding known by the library.

Further note that the program is responsible for the **sh_entsize** section header member. Although the library sets it for sections with known types, it cannot reliably know the correct value for all sections. Consequently, the library relies on the program to provide the values for unknown section types. If the entry size is unknown or not applicable, the value should be set to zero.

When deciding how to build the output file, **elf_update()** obeys the alignments of individual data buffers to create output sections. A section's most strictly aligned data buffer controls the section's alignment. The library also inserts padding between buffers, as necessary, to ensure the proper alignment of each buffer.

SEE ALSO

elf(3E), **elf32_fsize**(3E), **elf32_getehdr**(3E), **elf32_getshdr**(3E), **elf32_xlatetof**(3E), **elf_begin**(3E), **elf_flagdata**(3E), **elf_getdata**(3E)

NOTES

As mentioned above, the **ELF_C_WRITE** command translates data as necessary, before writing them to the file. This translation is *not* always transparent to the application program. If a program has obtained pointers to data associated with a file (for example, see **elf32_getehdr**(3E) and **elf_getdata**(3E)), the program should reestablish the pointers after calling **elf_update()**.

NAME

elf_version – coordinate ELF library and application versions

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <libelf.h>

unsigned elf_version(unsigned *ver***);**

MT-LEVEL

Unsafe

DESCRIPTION

As **elf**(3E) explains, the program, the library, and an object file have independent notions of the latest ELF version. **elf_version( )** lets a program query the ELF library's *internal version*. It further lets the program specify what memory types it uses by giving its own *working version*, *ver*, to the library. Every program that uses the ELF library must coordinate versions as described below.

The header **<libelf.h>** supplies the version to the program with the macro **EV_CURRENT**. If the library's internal version (the highest version known to the library) is lower than that known by the program itself, the library may lack semantic knowledge assumed by the program. Accordingly, **elf_version( )** will not accept a working version unknown to the library.

Passing *ver* equal to **EV_NONE** causes **elf_version( )** to return the library's internal version, without altering the working version. If *ver* is a version known to the library, **elf_version( )** returns the previous (or initial) working version number. Otherwise, the working version remains unchanged and **elf_version( )** returns **EV_NONE**.

EXAMPLES

The following excerpt from an application program protects itself from using an older library.

```
if (elf_version(EV_CURRENT) == EV_NONE) {
        /* library out of date */
        /* recover from error */
}
```

SEE ALSO

elf(3E), **elf32_xlatetof**(3E), **elf_begin**(3E)

NOTES

The working version should be the same for all operations on a particular elf descriptor. Changing the version between operations on a descriptor will probably not give the expected results.

NAME

end, _end, etext, _etext, edata, _edata – last locations in program

SYNOPSIS

```
extern _etext;
extern _edata;
extern _end;
```

DESCRIPTION

These names refer neither to routines nor to locations with interesting contents; only their addresses are meaningful.

_etext The address of **_etext** is the first location after the program text.

_edata The address of **_edata** is the first location after the initialized data region.

_end The address of **_end** is the first location after the uninitialized data region.

SEE ALSO

cc(1B), **brk**(2), **malloc**(3C), **stdio**(3S)

NOTE

When execution begins, the program break (the first location beyond the data) coincides with **_end**, but the program break may be reset by the routines **brk( )**, **malloc( )**, the standard input / output library (see **stdio**(3S)), by the profile (**–p**) option of **cc**(1B), and so on. Thus, the current value of the program break should be determined by **sbrk ((char *)0)** (see **brk**(2)).

References to **end**, **etext**, and **edata**, without a preceding underscore can be made by the user; if this case is detected the symbol will be aliased to the associated symbol which begins with the underscore.

NAME

erf, erfc – error functions

SYNOPSIS

cc [*flag* ...] *file* ... **–lm** [*library* ...]

#include <math.h>

double erf(double *x***);**

double erfc(double *x***);**

MT-LEVEL

MT-Safe

DESCRIPTION

erf(*x***)** returns the error function of *x*; where **erf** (*x*)**:= (2/**$\sqrt{\pi}$**)** $\int_0^x$ **exp (–t^2) dt.**

erfc(*x***)** returns 1.0–erf (*x*), computed however by other methods that avoid cancellation for large *x*.

NAME

ethers, ether_ntoa, ether_aton, ether_ntohost, ether_hostton, ether_line – Ethernet address mapping operations

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket -lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>
#include <net/if.h>
#include <netinet/in.h>
#include <netinet/if_ether.h>

char *ether_ntoa (struct ether_addr *e);

struct ether_addr *ether_aton (char *s);

int ether_ntohost (char *hostname, struct ether_addr *e);

int ether_hostton (char *hostname, struct ether_addr *e);

int ether_line (char *l, struct ether_addr *e, char *hostname);
```

MT-LEVEL

MT-Safe

DESCRIPTION

These routines are useful for mapping 48 bit Ethernet numbers to their ASCII representations or their corresponding host names, and vice versa.

The function **ether_ntoa( )** converts a 48 bit Ethernet number pointed to by *e* to its standard ASCII representation; it returns a pointer to the ASCII string. The representation is of the form *x* : *x* : *x* : *x* : *x* : *x* where *x* is a hexadecimal number between **0** and **ff**. The function **ether_aton( )** converts an ASCII string in the standard representation back to a 48 bit Ethernet number; the function returns **NULL** if the string cannot be scanned successfully.

The function **ether_ntohost( )** maps an Ethernet number (pointed to by *e*) to its associated hostname. The string pointed to by hostname must be long enough to hold the hostname and a **NULL** character. The function returns zero upon success and non-zero upon failure. Inversely, the function **ether_hostton( )** maps a hostname string to its corresponding Ethernet number; the function modifies the Ethernet number pointed to by *e*. The function also returns zero upon success and non-zero upon failure. In order to do the mapping, both these functions may lookup one or more of the following sources: the ethers file, the NIS maps ‘‘ethers.byname’’ and ‘‘ethers.byaddr’’ and the NIS+ table ‘‘ethers’’. The sources and their lookup order are specified in the **/etc/nsswitch.conf** file (see **nsswitch.conf**(4) for details).

The function **ether_line( )** scans a line (pointed to by *l*) and sets the hostname and the Ethernet number (pointed to by *e*). The string pointed to by hostname must be long enough to hold the hostname and a **NULL** character. The function returns zero upon success and non-zero upon failure. The format of the scanned line is described by **ethers**(4).

FILES

/etc/ethers
/etc/nsswitch.conf

SEE ALSO

ethers(4), **nsswitch.conf**(4)

BUGS

Programs that call **ether_hostton( )** or **ether_ntohost( )** routines cannot be linked statically since the implementation of these routines requires dynamic linker functionality to access shared objects at run time.

NAME

euclen, euccol, eucscol – get byte length and display width of EUC characters

SYNOPSIS

```
#include <euc.h>
int euclen(const unsigned char *s);
int euccol(const unsigned char *s);
int eucscol(const unsigned char *str);
```

MT-LEVEL

Safe

DESCRIPTION

euclen( ) returns the length in bytes of the Extended Unix Code (EUC) character pointed to by *s*, including single-shift characters, if present.

euccol( ) returns the screen column width of the EUC character pointed to by *s*.

eucscol( ) returns the screen column width of the EUC string pointed to by *str*.

For the **euclen( )** and **euccol( )**, routines, *s* points to the first byte of the character. This byte is examined to determine its codeset. The character type table for the current *locale* is used for codeset byte length and display width information.

FILES

/usr/lib/locale/*locale***/LC_CTYPE**

SEE ALSO

getwidth(3I)

NAME

exit – terminate process

SYNOPSIS

#include <stdlib.h>

void exit(int *status***);**

MT-LEVEL

Safe

DESCRIPTION

The C function **exit( )** first calls any functions registered through the **atexit**(3C) function in the reverse order of their registration. **exit( )** then calls **_exit**, which circumvents all such functions and cleanup.

The symbols **EXIT_SUCCESS** and **EXIT_FAILURE** are defined in the header **<stdlib.h>** and may be used as the value of *status* to indicate successful or unsuccessful termination, respectively.

SEE ALSO

exit(2), **atexit**(3C)

NAME

exp, expm1, log, log1p, log10, pow – exponential, logarithm, power

SYNOPSIS

cc [*flag* ...] *file* ... **–lm** [*library* ...]

#include <math.h>

double exp(double *x***);**

double expm1(double *x***);**

double log(double *x***);**

double log1p(double *x***);**

double log10(double *x***);**

double pow(double *x***, double** *y***);**

MT-LEVEL

MT-Safe

DESCRIPTION

exp(*x***)** computes the exponential function *e**x*.

expm1(*x***)** computes *(e**x)–1* accurately even for tiny *x*.

log(*x***)** computes the natural logarithm of *x*.

log1p(*x***)** computes *log(1+x)* accurately even for tiny *x*.

log10(*x***)** computes the base-10 logarithm of *x*.

pow(*x***,** *y***)** computes *x* raised to the power *y*.

RETURN VALUES

For exceptional cases, **matherr**(3M) tabulates the values to be returned as dictated by various Standards.

SEE ALSO

matherr(3M)

DIAGNOSTICS

In IEEE 754 mode (i.e. the **–xlibmieee** cc compilation option), **log(**±0**)** returns –∞ and raises the division by zero exception; if *x*<0, **log(***x***)** returns a NaN and raises the invalid operation exception; if *x* **==** +∞ or a quiet NaN, **log(***x***)** returns *x* and raises no exception; if *x* is a signaling NaN, **log(***x***)** returns a quiet NaN and raises the invalid operation exception; **log(**1**)** returns 0 and raises no exception; for all other positive *x*, **log(***x***)** returns a normalized number and raises the inexact exception.

NAME

fattach – attach a STREAMS-based file descriptor to an object in the file system name space

SYNOPSIS

int fattach(int *fildes***, const char** **path***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **fattach( )** routine attaches a STREAMS-based file descriptor to an object in the file system name space, effectively associating a name with *fildes*. *fildes* must be a valid open file descriptor representing a STREAMS file. *path* is a path name of an existing object and the user must have appropriate privileges or be the owner of the file and have write permissions. All subsequent operations on *path* will operate on the STREAMS file until the STREAMS file is detached from the node. *fildes* can be attached to more than one *path*, that is, a stream can have several names associated with it.

The attributes of the named stream (see **stat**(2)), are initialized as follows: the permissions, user ID, group ID, and times are set to those of *path*, the number of links is set to 1, and the size and device identifier are set to those of the streams device associated with *fildes*. If any attributes of the named stream are subsequently changed (for example, **chmod**(2)), the attributes of the underlying object are not affected.

RETURN VALUES

If successful, **fattach( )** returns 0; otherwise it returns -1 and sets **errno** to indicate an error.

ERRORS

Under the following conditions, the function **fattach( )** fails and sets **errno** to:

EACCES	The user is the owner of *path* but does not have write permissions on *path* or *fildes* is locked.
EBADF	*fildes* is not a valid open file descriptor.
EBUSY	*path* is currently a mount point or has a STREAMS file descriptor attached it.
EINVAL	*path* is a file in a remotely mounted directory.
EINVAL	*fildes* does not represent a STREAMS file.
ELOOP	Too many symbolic links were encountered in translating *path*.
ENAMETOOLONG	The size of *path* exceeds **{PATH_MAX}**, or the component of a path name is longer than **{NAME_MAX}** while **{_POSIX_NO_TRUNC}** is in effect.
ENOENT	*path* does not exist.

ENOTDIR	A component of a path prefix is not a directory.
EPERM	The effective user ID is not the owner of *path* or a user with the appropriate privileges.

SEE ALSO

fdetach(1M), **chmod**(2), **stat**(2), **fdetach**(3C), **isastream**(3C), **streamio**(7I)

STREAMS Programming Guide

NAME

fclose, fflush – close or flush a stream

SYNOPSIS

```
#include <stdio.h>
int fclose(FILE *stream);
int fflush(FILE *stream);
```

MT-LEVEL

MT-Safe

DESCRIPTION

fclose( ) causes any buffered data waiting to be written for the named *stream* (see **intro**(3)) to be written out, and the *stream* to be closed. If the underlying file pointer is not already at end of file, and the file is one capable of seeking, the file pointer is adjusted so that the next operation on the open file pointer deals with the byte after the last one read from or written to the file being closed.

fclose( ) is performed automatically for all open files upon calling **exit( )**.

If *stream* points to an output stream or an update stream on which the most recent operation was not input, **fflush( )** causes any buffered data waiting to be written for the named *stream* to be written to that file. Any unread data buffered in *stream* is discarded. The *stream* remains open. If *stream* is open for reading, the underlying file pointer is not already at end of file, and the file is one capable of seeking, the file pointer is adjusted so that the next operation on the open file pointer deals with the byte after the last one read from or written to the stream.

When calling **fflush( )**, if *stream* is a null pointer, all files open for writing are flushed.

RETURN VALUES

Upon successful completion these functions return a value of zero. Otherwise **EOF** is returned.

SEE ALSO

close(2), **exit**(2), **intro**(3), **fopen**(3S), **setbuf**(3S), **stdio**(3S)

NAME

fdatasync – synchronize a file's data

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

#include <unistd.h>

int fdatasync(int *fildes***);**

MT-LEVEL

Async-Signal-Safe

DESCRIPTION

fdatasync( ) forces all currently queued I / O operations associated with the file descriptor *fildes* to synchronized I / O data integrity completion. See **fcntl**(5) definition of **O_DSYNC**.

RETURN VALUES

fdatasync( ) returns **0** upon success; otherwise, it returns **-1** and sets **errno** to indicate the error condition.

ERRORS

EBADF *fildes* is not a valid file descriptor.

EINVAL This implementation does not support synchronized I / O for this file.

ENOSYS **fdatasync( )** is not supported by this implementation.

In the event that any of the queued I / O operations fail, **fdatasync( )** returns the error conditions defined for **read**(2) and **write**(2).

SEE ALSO

fcntl(2), **open**(2), **read**(2), **write**(2), **fsync**(3C), **aio_fsync**(3R), **fcntl**(5)

NOTES

If **fdatasync( )** fails, outstanding I / O operations are not guaranteed to have been completed.

NAME

fdetach – detach a name from a STREAMS-based file descriptor

SYNOPSIS

int fdetach(const char **path***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **fdetach( )** routine detaches a STREAMS-based file descriptor from a name in the file system. *path* is the path name of the object in the file system name space, which was previously attached (see **fattach**(3C)). The user must be the owner of the file or a user with the appropriate privileges. All subsequent operations on *path* will operate on the file system node and not on the STREAMS file. The permissions and status of the node are restored to the state the node was in before the STREAMS file was attached to it.

RETURN VALUES

If successful, **fdetach( )** returns 0; otherwise it returns -1 and sets **errno** to indicate an error.

ERRORS

Under the following conditions, the function **fdetach( )** fails and sets **errno** to:

EINVAL	*path* is not attached to a STREAMS file.
ELOOP	Too many symbolic links were encountered in translating *path*.
ENAMETOOLONG	The size of *path* exceeds **{PATH_MAX}**, or a path name component is longer than **{NAME_MAX}** while **{_POSIX_NO_TRUNC}** is in effect.
ENOENT	*path* does not exist.
ENOTDIR	A component of the path prefix is not a directory.
EPERM	The effective user ID is not the owner of *path* or is not a user with appropriate permissions.

SEE ALSO

fdetach(1M), **fattach**(3C), **streamio**(7I)

STREAMS Programming Guide

NAME

ferror, feof, clearerr, fileno – stream status inquiries

SYNOPSIS

```
#include <stdio.h>
int ferror(FILE *stream);
int feof(FILE *stream);
void clearerr(FILE *stream);
int fileno(FILE *stream);
```

MT-LEVEL

MT-Safe

DESCRIPTION

ferror() returns non-zero when an error has previously occurred reading from or writing to the named *stream* (see **intro**(3)), otherwise zero.

feof() returns non-zero when **EOF** has previously been detected reading the named input *stream*, otherwise zero.

clearerr() resets the error indicator and **EOF** indicator to zero on the named *stream*.

fileno() returns the integer file descriptor associated with the named *stream*; see **open**(2).

SEE ALSO

open(2), **intro**(3), **fopen**(3S), **stdio**(3S)

NAME

ffs – find first set bit

SYNOPSIS

```
#include <string.h>
int ffs(const int i);
```

MT-LEVEL

MT-Safe

DESCRIPTION

ffs( ) finds the first bit set in the argument passed it and returns the index of that bit. Bits are numbered starting at 1 from the low order bit. A return value of zero indicates that the value passed is zero.

NAME

floating_to_decimal, single_to_decimal, double_to_decimal, extended_to_decimal, quadruple_to_decimal – convert floating-point value to decimal record

SYNOPSIS

```
#include <floatingpoint.h>

void single_to_decimal(single *px, decimal_mode *pm, decimal_record *pd,
        fp_exception_field_type *ps);

void double_to_decimal(double *px, decimal_mode *pm, decimal_record *pd,
        fp_exception_field_type *ps);

void extended_to_decimal(extended *px, decimal_mode *pm, decimal_record *pd,
        fp_exception_field_type *ps);

void quadruple_to_decimal(quadruple *px, decimal_mode *pm, decimal_record *pd,
        fp_exception_field_type *ps);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **floating_to_decimal( )** functions convert the floating-point value at **px* into a decimal record at **pd*, observing the modes specified in **pm* and setting exceptions in **ps*. If there are no IEEE exceptions, **ps* will be zero.

If **px* is zero, infinity, or NaN, then only *pd->sign* and *pd->fpclass* are set. Otherwise *pd->exponent* and *pd->ds* are also set so that

(pd->sign)*(pd->ds)*10(pd->exponent)**

is a correctly rounded approximation to **px*. *pd->ds* has at least one and no more than **DECIMAL_STRING_LENGTH–1** significant digits because one character is used to terminate the string with a NULL.

pd->ds is correctly rounded according to the IEEE rounding modes in *pm->rd*. **ps* has *fp_inexact* set if the result was inexact, and has *fp_overflow* set if the string result does not fit in *pd->ds* because of the limitation **DECIMAL_STRING_LENGTH**.

If *pm->df* == *floating_form*, then *pd->ds* always contains *pm->ndigits* significant digits. Thus if **px* == 12.34 and *pm->ndigits* == 8, then *pd->ds* will contain 12340000 and *pd->exponent* will contain –6.

If *pm->df* == *fixed_form* and *pm->ndigits* >= 0, then *pd->ds* always contains *pm->ndigits* after the point and as many digits as necessary before the point. Since the latter is not known in advance, the total number of digits required is returned in *pd->ndigits*; if that number >= **DECIMAL_STRING_LENGTH**, then *ds* is undefined. *pd->exponent* always gets *–pm->ndigits*. Thus if **px* == 12.34 and *pm->ndigits* == 1, then *pd->ds* gets 123, *pd->exponent* gets –1, and *pd->ndigits* gets 3.

If *pm->df* == *fixed_form* and *pm->ndigits* < 0, then *pd->ds* always contains *–pm->ndigits* trailing zeros; in other words, rounding occurs *–pm->ndigits* to the left of the decimal point, but the digits rounded away are retained as zeros. The total number of digits required is in *pd->ndigits*. *pd->exponent* always gets 0. Thus if **px* == 12.34 and

pm->ndigits == –1, then *pd->ds* gets 10, *pd->exponent* gets 0, and *pd->ndigits* gets 2. *pd->more* is not used.

econvert(3), **fconvert**(3), **gconvert**(3), **printf**(3S), and **sprintf**(3S) all use **double_to_decimal( )**.

SEE ALSO

econvert(3), **fconvert**(3), **gconvert**(3), **printf**(3S), **sprintf**(3S)

NAME

flock – apply or remove an advisory lock on an open file

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <sys/file.h>

int flock(*fd, operation***)**
int *fd, operation***;**

DESCRIPTION

flock() applies or removes an *advisory* lock on the file associated with the file descriptor *fd*. The compatibility version of **flock()** has been implemented on top of **fcntl**(2) locking. It does not provide complete binary compatibility.

Advisory locks allow cooperating processes to perform consistent operations on files, but do not guarantee exclusive access (that is, processes may still access files without using advisory locks, possibly resulting in inconsistencies).

The locking mechanism allows two types of locks: shared locks and exclusive locks. More than one process may hold a shared lock for a file at any given time, but multiple exclusive, or both shared and exclusive, locks may not exist simultaneously on a file.

A lock is applied by specifying an *operation* parameter **LOCK_SH** for a shared lock or **LOCK_EX** for an exclusive lock. The *operation* paramerer may be ORed with **LOCK_NB** to make the operation non-blocking. To unlock an existing lock, the *operation* should be **LOCK_UN**.

Read permission is required on a file to obtain a shared lock, and write permission is required to obtain an exclusive lock. Locking a segment that is already locked by the calling process causes the old lock type to be removed and the new lock type to take effect.

Requesting a lock on an object that is already locked normally causes the caller to block until the lock may be acquired. If **LOCK_NB** is included in *operation*, then this will not happen; instead, the call will fail and the error **EWOULDBLOCK** will be returned.

RETURN VALUES

flock() returns:

0 on success.

–1 on failure and sets **errno** to indicate the error.

ERRORS

EBADF	The argument *fd* is an invalid descriptor.
EINVAL	*operation* is not a valid argument.
EOPNOTSUPP	The argument *fd* refers to an object other than a file.
EWOULDBLOCK	The file is locked and the **LOCK_NB** option was specified.

SEE ALSO

lockd(1M), **chmod**(2), **close**(2), **dup**(2), **exec**(2), **fcntl**(2), **fork**(2), **open**(2), **lockf**(3C)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

Locks are on files, not file descriptors. That is, file descriptors duplicated through **dup**(2) or **fork**(2) do not result in multiple instances of a lock, but rather multiple references to a single lock. If a process holding a lock on a file forks and the child explicitly unlocks the file, the parent will lose its lock. Locks are not inherited by a child process.

Processes blocked awaiting a lock may be awakened by signals.

Mandatory locking may occur, depending on the mode bits of the file. See **chmod**(2).

Locks obtained through the **flock( )** mechanism under SunOS 4.1 were known only within the system on which they were placed. This is no longer true.

NAME

flockfile, funlockfile, ftrylockfile – acquire and release stream lock

SYNOPSIS

```
#include <stdio.h>
void flockfile(FILE *stream);
void funlockfile(FILE *stream);
```

POSIX

```
cc [ flag... ] file ... –D_POSIX_PTHREAD_SEMANTICS [ library... ]
int ftrylockfile(FILE *stream);
```

MT-LEVEL

MT-Safe

DESCRIPTION

flockfile( ) acquires an internal lock of a stream *stream*. If the lock is already acquired by another thread, the thread calling **flockfile( )** is suspended until it can acquire the lock. In the case that the stream lock is available, **flockfile( )** not only acquires the lock, but keeps track of the number of times it is being called by the current thread. This implies that the stream lock can be acquired more than once by the same thread.

funlockfile( ) releases the lock being held by the current thread. In the case of recursive locking, this function must be called the same number of times **flockfile( )** was called. After the number of **funlockfile( )** calls is equal to the number of **flockfile( )** calls, the stream lock is available for other threads to acquire.

ftrylockfile( ) acquires an internal lock of a stream *stream*, only if that object is available. In essence **ftrylockfile( )** is a non-blocking version of **flockfile( )**.

RETURN VALUES

ftrylockfile( ) returns zero on success or non-zero to indicate a lock cannot be acquired.

EXAMPLES

For example:

```
FILE iop;
...
flockfile(iop);
fprintf(iop, "hello ");
fprintf(iop, "world0);
fputc(iop, 'a');
funlockfile(iop);
```

will print everything out together, blocking other threads that might want to write to the same file between **fprintf( )**s.

An unlocked interface is available in case performance is an issue. For example:

```
flockfile(iop);
while (!feof(iop)) {
        *c++ = getc_unlocked(iop);
}
funlockfile(iop);
```

SEE ALSO

intro(3), **ferror**(3S), **getc**(3S), **putc**(3S), **stdio**(3S), **ungetc**(3S)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

The interfaces on this page are as specified in POSIX 1003.1c Draft #10.

NAME

floor, ceil, rint – round to integral value in floating-point format

SYNOPSIS

cc [*flag* ...] *file* ... **–lm** [*library* ...]

#include <math.h>

double ceil(double *x***);**

double floor(double *x***);**

double rint(double *x***);**

MT-LEVEL

MT-Safe

DESCRIPTION

ceil( ), **floor( )** and **rint( )** convert a double value into an integral value in **double** format. They vary in how they choose the result when the argument is not already an integral value. Here an "integral value" means a value of a mathematical integer, which however might be too large to fit in a particular computer's **int** format. All sufficiently large values in a particular floating-point format are already integral; in IEEE 754 double-precision format, that means all values >= 2**52. Zeros, infinities, and quiet NaNs are treated as integral values by these functions, which always preserve their argument's sign.

ceil( ) returns the least integral value greater than or equal to *x*. This corresponds to IEEE 754 rounding toward positive infinity.

floor( ) returns the greatest integral value less than or equal to *x*. This corresponds to IEEE 754 rounding toward negative infinity.

rint( ) rounds *x* to an integral value according to the current IEEE 754 rounding direction.

NAME

fmtmsg – display a message on stderr or system console

SYNOPSIS

```
#include <fmtmsg.h>

int fmtmsg(long classification, const char *label, int severity, const char *text,
    const char *action, const char *tag);
```

MT-LEVEL

Safe

DESCRIPTION

Based on a message's classification component, **fmtmsg( )** writes a formatted message to **stderr**, to the console, or to both.

fmtmsg( ) can be used instead of the traditional **printf**(3S) interface to display messages to **stderr**. **fmtmsg( ),** in conjunction with **gettxt( ),** provides a simple interface for producing language-independent applications.

A formatted message consists of up to five standard components as defined below. The component, *classification*, is not part of the standard message displayed to the user, but rather defines the source of the message and directs the display of the formatted message.

classification

Contains identifiers from the following groups of major classifications and subclassifications. Any one identifier from a subclass may be used in combination by ORing the values together with a single identifier from a different subclass. Two or more identifiers from the same subclass should not be used together, with the exception of identifiers from the display subclass. (Both display subclass identifiers may be used so that messages can be displayed to both **stderr** and the system console).

- "Major classifications" identify the source of the condition. Identifiers are: **MM_HARD** (hardware), **MM_SOFT** (software), and **MM_FIRM** (firmware).
- "Message source subclassifications" identify the type of software in which the problem is spotted. Identifiers are: **MM_APPL** (application), **MM_UTIL** (utility), and **MM_OPSYS** (operating system).
- "Display subclassifications" indicate where the message is to be displayed. Identifiers are: **MM_PRINT** to display the message on the standard error stream, **MM_CONSOLE** to display the message on the system console. Neither, either, or both identifiers may be used.
- "Status subclassifications" indicate whether the application will recover from the condition. Identifiers are: **MM_RECOVER** (recoverable) and **MM_NRECOV** (non-recoverable).
- An additional identifier, **MM_NULLMC**, indicates that no classification component is supplied for the message.

label Identifies the source of the message. The format of this component is two fields separated by a colon. The first field is up to 10 characters long; the second is up to 14 characters. Suggested usage is that *label* identifies the package in which the application resides as well as the program or application name. For example, the

label **UX:cat** indicates the UNIX System V package and the **cat** application.

severity
Indicates the seriousness of the condition. Identifiers for the standard levels of *severity* are:

- **MM_HALT** indicates that the application has encountered a severe fault and is halting. Produces the print string **HALT**.
- **MM_ERROR** indicates that the application has detected a fault. Produces the print string **ERROR**.
- **MM_WARNING** indicates a condition out of the ordinary that might be a problem and should be watched. Produces the print string **WARNING**.
- **MM_INFO** provides information about a condition that is not in error. Produces the print string **INFO**.
- **MM_NOSEV** indicates that no severity level is supplied for the message.

Other severity levels may be added by using the **addseverity( )** routine.

text Describes the condition that produced the message. The *text* string is not limited to a specific size.

action Describes the first step to be taken in the error recovery process. **fmtmsg( )** precedes each action string with the prefix: **TO**FIX**:.** The *action* string is not limited to a specific size.

tag An identifier which references on-line documentation for the message. Suggested usage is that *tag* includes the *label* and a unique identifying number. A sample *tag* is **UX:cat:146**.

Environment Variables

There are two environment variables that control the behavior of **fmtmsg( )**: **MSGVERB** and **SEV_LEVEL**.

MSGVERB tells **fmtmsg( )** which message components it is to select when writing messages to **stderr**. The value of **MSGVERB** is a colon-separated list of optional keywords. **MSGVERB** can be set as follows:

MSGVERB=[*keyword*[**:***keyword*[**:**. . .]]]
export MSGVERB

Valid *keywords* are: **label**, **severity**, **text**, **action**, and **tag**. If **MSGVERB** contains a keyword for a component and the component's value is not the component's null value, **fmtmsg( )** includes that component in the message when writing the message to **stderr**. If **MSGVERB** does not include a keyword for a message component, that component is not included in the display of the message. The keywords may appear in any order. If **MSGVERB** is not defined, if its value is the null-string, if its value is not of the correct format, or if it contains keywords other than the valid ones listed above, **fmtmsg( )** selects all components.

The first time **fmtmsg( )** is called, it examines the **MSGVERB** environment variable to see which message components it is to select when generating a message to write to the standard error stream, **stderr**. The values accepted on the initial call are saved for future calls.

MSGVERB affects only which components are selected for display to the standard error stream. All message components are included in console messages.

SEV_LEVEL defines severity levels and associates print strings with them for use by **fmtmsg( ) .** The standard severity levels shown below cannot be modified. Additional severity levels can also be defined, redefined, and removed using **addseverity( )** (see **addseverity**(3C)). If the same severity level is defined by both **SEV_LEVEL** and **addseverity( ) ,** the definition by **addseverity( )** is controlling.

0 (no severity is used)
1 **HALT**
2 **ERROR**
3 **WARNING**
4 **INFO**

SEV_LEVEL can be set as follows:

SEV_LEVEL=[*description*[**:***description*[**:** . . .]]]
export SEV_LEVEL

description is a comma-separated list containing three fields:

*description***=***severity_keyword***,***level***,***printstring*

severity_keyword is a character string that is used as the keyword on the **–s** *severity* option to the **fmtmsg** command. (This field is not used by the **fmtmsg( )** function.)

level is a character string that evaluates to a positive integer (other than 0, 1, 2, 3, or 4, which are reserved for the standard severity levels). If the keyword *severity_keyword* is used, *level* is the severity value passed on to the **fmtmsg( )** function.

printstring is the character string used by **fmtmsg( )** in the standard message format whenever the severity value *level* is used.

If a *description* in the colon list is not a three-field comma list, or, if the second field of a comma list does not evaluate to a positive integer, that *description* in the colon list is ignored.

The first time **fmtmsg( )** is called, it examines the **SEV_LEVEL** environment variable, if defined, to see whether the environment expands the levels of severity beyond the five standard levels and those defined using **addseverity( ) .** The values accepted on the initial call are saved for future calls.

Use in Applications

One or more message components may be systematically omitted from messages generated by an application by using the null value of the argument for that component.

The table below indicates the null values and identifiers for **fmtmsg( )** arguments.

Argument	Type	Null-Value	Identifier
label	**char***	**(char*) NULL**	**MM_NULLLBL**
severity	**int**	**0**	**MM_NULLSEV**
class	**long**	**0L**	**MM_NULLMC**
text	**char***	**(char*) NULL**	**MM_NULLTXT**
action	**char***	**(char*) NULL**	**MM_NULLACT**
tag	**char***	**(char*) NULL**	**MM_NULLTAG**

Another means of systematically omitting a component is by omitting the component keyword(s) when defining the **MSGVERB** environment variable (see the ‘‘Environment Variables’’ section).

RETURN VALUES

The exit codes for **fmtmsg( )** are the following:

MM_OK The function succeeded.

MM_NOTOK The function failed completely.

MM_NOMSG The function was unable to generate a message on the standard error stream, but otherwise succeeded.

MM_NOCON The function was unable to generate a console message, but otherwise succeeded.

EXAMPLES

Example 1:

The following example of **fmtmsg( )**:

fmtmsg(MM_PRINT, "UX:cat", MM_ERROR, "invalid syntax", "refer to manual", "UX:cat:001")

produces a complete message in the standard message format:

```
UX:cat: ERROR: invalid syntax
        TO FIX: refer to manual   UX:cat:001
```

Example 2:

When the environment variable **MSGVERB** is set as follows:

MSGVERB=severity:text:action

and the Example 1 is used, **fmtmsg( )** produces:

```
ERROR: invalid syntax
TO FIX: refer to manual
```

Example 3:

When the environment variable **SEV_LEVEL** is set as follows:

SEV_LEVEL=note,5,NOTE

the following call to **fmtmsg( ) :**

fmtmsg(MM_UTIL | MM_PRINT, "UX:cat", 5, "invalid syntax", "refer to manual", "UX:cat:001")

produces:

UX:cat: NOTE: invalid syntax
TO FIX: refer to manual UX:cat:001

SEE ALSO

fmtmsg(1), **addseverity**(3C), **gettxt**(3C), **printf**(3S)

NAME

fn_attr_get – return specified attribute associated with name

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_attribute_t ∗**fn_attr_get(FN_ctx_t** ∗*ctx***, const FN_composite_name_t** ∗*name***,**
const FN_identifier_t ∗*attribute_id***, FN_status_t** ∗*status***);**

MT-LEVEL

Safe

DESCRIPTION

This operation returns the identifier, syntax and values of a specified attribute for the object named *name* relative to *ctx*. If *name* is empty, the attribute associated with *ctx* is returned.

RETURN VALUE

fn_attr_get returns a pointer to an **FN_attribute_t** object if the operation succeeds; it returns a **NULL** pointer (**0**) if the operation fails.

ERRORS

fn_attr_get( ) sets *status* as described in **FN_status_t**(3N) and **xfn_status_codes**(3N).

APPLICATION USAGE

fn_attr_get_values( ) and its related operations are used for getting individual values of an attribute. They should be used if the combined size of all the values are expected to be too large to be returned in a single invocation of **fn_attr_get( ).**

SEE ALSO

FN_attribute_t(3N), **FN_composite_name_t**(3N), **FN_ctx_t**(3N), **FN_identifier_t**(3N), **FN_status_t**(3N), **fn_attr_get_values**(3N), **xfn_attributes**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_attr_get_ids – get a list of the identifiers of all attributes associated with named object

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_attrset_t ∗**fn_attr_get_ids(FN_ctx_t** ∗*ctx***, const FN_composite_name_t** ∗*name***,**
FN_status_t ∗*status***);**

MT-LEVEL

Safe.

DESCRIPTION

This operation returns a list of the attribute identifiers of all attributes associated with the object named by *name* relative to the context *ctx*. If *name* is empty, the attribute identifiers associated with *ctx* are returned.

RETURN VALUE

This operation returns a pointer to an object of type **FN_attrset_t**; if the operation fails, a **NULL** pointer (**0**) is returned.

ERRORS

This operation sets *status* as described in **FN_status_t**(3N) and **xfn_status_codes**(3N).

APPLICATION USAGE

The attributes in the returned set do not contain the syntax or values of the attributes, only their identifiers.

SEE ALSO

FN_attrset_t(3N), **FN_attribute_t**(3N), **FN_composite_name_t**(3N), **FN_ctx_t**(3N), **FN_status_t**(3N), **fn_attr_get**(3N), **fn_attr_multi_get**(3N)**, xfn_attributes**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_attr_get_values, FN_valuelist_t, fn_valuelist_next, fn_valuelist_destroy – return values of an attribute

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_valuelist_t ***fn_attr_get_values(FN_ctx_t** ***ctx*, **const FN_composite_name_t** ***name*, **const FN_identifier_t** ***attribute_id*, **FN_status_t** ***status*);

FN_attrvalue_t ***fn_valuelist_next(FN_valuelist_t** ***vl*, **FN_identifier_t** ****attr_syntax*, **FN_status_t** ***status*);

void fn_valuelist_destroy(FN_valuelist_t*vl*, **FN_status_t** ***status*);

MT-LEVEL

Safe.

DESCRIPTION

This set of operations is used to obtain the values of a single attribute, identified by *attribute_id*, associated with the object named *name*, resolved in the context *ctx*. If *name* is empty, the attribute values associated with *ctx* are obtained.

The operation **fn_attr_get_values( )** initiates the enumeration process. It returns a handle to an **FN_valuelist_t** object that can be used to enumerate the values of the specified attribute.

The operation **fn_valuelist_next( )** returns a new **FN_attrvalue_t** object containing the next value in the attribute and may be called multiple times until all values are retrieved. The syntax of the attribute is returned in *attr_syntax*.

The operation **fn_valuelist_destroy( )** is used to release the resources used during the enumeration. This may be invoked before the enumeration has completed to terminate the enumeration.

These operations work in a similar fashion as the **fn_ctx_list_names( )** operations.

RETURN VALUE

fn_attr_get_values( ) returns a pointer to an **FN_valuelist_t** object if the enumeration process is successfully initiated; it returns a **NULL** pointer if the process failed.

fn_valuelist_next ( ) returns a **NULL** pointer if no more attribute value can be returned.

In the case of a failure, these operations set *status* to indicate the nature of the failure.

ERRORS

Each successful call to **fn_valuelist_next( )** returns an attribute value. *status* is set to **FN_SUCCESS.**

When **fn_valuelist_next( )** returns a **NULL** pointer, it indicates that no more values can be returned. *status* is set in the following way:

FN_SUCCESS
The enumeration has completed successfully.

FN_E_INVALID_ENUM_HANDLE
The given enumeration handle is not valid. Possible reasons could be that the handle was from another enumeration, or the context being enumerated no

longer accepts the handle (due to such events as handle expiration or updates to the context).

FN_E_PARTIAL_RESULT
The enumeration is not yet complete but cannot be continued.

In addition to these status codes, other status codes are also possible in calls to these operations. In such cases, *status* is set as described in **FN_status_t**(3N) and **xfn_status_codes**(3N).

APPLICATION USAGE

This interface should be used instead of **fn_attr_get( )** if the combined size of all the values is expected to be too large to be returned by **fn_attr_get( ).**

There may be a relationship between the *ctx* argument supplied to **fn_attr_get_values( )** and the **FN_valuelist_t** object it returns. For example, some implementations may store the context handle *ctx* within the **FN_valuelist_t** object for subsequent **fn_valuelist_next( )** calls. In general, a **fn_ctx_handle_destroy( )** should not be invoked on *ctx* until the enumeration has terminated.

SEE ALSO

FN_attribute_t(3N), **FN_attrvalue_t**(3N), **FN_composite_name_t**(3N), **FN_ctx_t**(3N), **FN_identifier_t**(3N), **FN_status_t**(3N), **fn_attr_get**(3N), **fn_ctx_list_names**(3N), **xfn_attributes**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_attr_modify – modify specified attribute associated with name

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

```
#include <xfn/xfn.h>
int fn_attr_modify(FN_ctx_t *ctx, const FN_composite_name_t *name,
        unsigned int mod_op, const FN_attribute_t *attr, FN_status_t *status);
```

MT-LEVEL

Safe.

DESCRIPTION

This operation modifies according to *mod_op* the attribute *attr* associated with the object named *name* relative to *ctx*. If *name* is empty, the attribute associated with *ctx* is modified.

The modification is made on the attribute identified by the attribute identifier of *attr*. The syntax and values of *attr* are use according to the modification operation.

The modification operations are as follows:

FN_ATTR_OP_ADD Add an attribute with given attribute identifier and set of values. If an attribute with this identifier already exists, replace the set of values with those in the given set. The set of values may be empty if the target naming system permits.

FN_ATTR_OP_ADD_EXCLUSIVE Add an attribute with the given attribute identifier and set of values. The operation fails if an attribute with this identifier already exists. The set of values may be empty if the target naming system permits.

FN_ATTR_OP_REMOVE Remove the attribute with the given attribute identifier and all of its values. The operation succeeds even if the attribute does not exist. The values of the attribute supplied with this operation are ignored.

FN_ATTR_OP_ADD_VALUES Add the given values to those of the given attribute (resulting in the attribute having the union of its prior value set with the set given). Create the attribute if it does not exist already. The set of values may be empty if the target naming system permits.

FN_ATTR_OP_REMOVE_VALUES Remove the given values from those of the given attribute (resulting in the attribute having the set difference of its prior value set and the set given). This succeeds even if some of the given values are not in the set of values that the attribute has. In naming systems that require an attribute to have at least one value, removing the last value will remove the attribute as well.

RETURN VALUE

This operation returns **1** if the operation succeeds, **0** if the operation fails.

ERRORS

fn_attr_modify( ) sets *status* as described in **FN_status_t**(3N) and **xfn_status_codes**(3N).

SEE ALSO

FN_composite_name_t(3N), **FN_ctx_t**(3N), **FN_attribute_t**(3N), **FN_status_t**(3N), **fn_attr_multi_modify**(3N), **xfn_attributes**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_attr_multi_get, FN_multigetlist_t, fn_multigetlist_next, fn_multigetlist_destroy – return multiple attributes associated with named object

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_multigetlist_t ***fn_attr_multi_get(FN_ctx_t** ***ctx*,
const FN_composite_name_t ***name*, **const FN_attrset_t** ***attr_ids*,
FN_status_t ***status*);

FN_attribute_t ***fn_multigetlist_next(FN_multigetlist_t** ***ml*, **FN_status_t** ***status*);

void fn_multigetlist_destroy(FN_multigetlist_t ***ml*, **FN_status_t** ***status*);

MT-LEVEL

Safe.

DESCRIPTION

This set of operations returns one or more attributes associated with the object named by *name* relative to the context *ctx*. If *name* is empty, the attributes associated with *ctx* are returned.

The attributes returned are those specified in *attr_ids*. If the value of *attr_ids* is **0**, all attributes associated with the named object are returned. Any attribute values in *attr_ids* provided by the caller are ignored; only the attribute identifiers are relevant for this operation. Each attribute (identifier, syntax, values) is returned one at a time using an enumeration scheme similar to that for listing a context.

fn_attr_multi_get( ) initiates the enumeration process. It returns a handle to an **FN_multigetlist_t** object that can be used for the enumeration.

The operation **fn_multigetlist_next( )** returns a new **FN_attribute_t** object containing the next attribute (identifiers, syntaxes, and values) requested and updates *ml* to indicate the state of the enumeration.

The operation **fn_multigetlist_destroy( )** releases the resources used during the enumeration. It may be invoked before the enumeration has completed to terminate the enumeration.

RETURN VALUE

fn_attr_multi_get( ) returns a pointer to an **FN_multigetlist_t** object if the enumeration has been initiated successfully; a **NULL** pointer is returned if it failed.

fn_multigetlist_next( ) returns a pointer to an **FN_attribute_t** object if an attribute was returned, a **NULL** pointer (**0**) if no attribute was returned.

In the case of a failure, these operations set *status* to indicate the nature of the failure.

ERRORS

Each call to **fn_multigetlist_next( )** sets *status* as follows:

FN_SUCCESS

If an attribute was returned, there are more attributes to be enumerated. If no attribute was returned, the enumeration has completed successfully.

FN_E_ATTR_NO_PERMISSION

The caller did not have permission to read this attribute.

FN_E_INSUFFICIENT_RESOURCES
Insufficient resources are available to return the attribute's values.

FN_E_INVALID_ATTR_IDENTIFIER
This attribute identifier was not in a format acceptable to the naming system, or its contents was not valid for the format specified for the identifier.

FN_E_INVALID_ENUM_HANDLE
(No attribute should be returned with this status code). The given enumeration handle is not valid. Possible reasons could be that the handle was from another enumeration, or the object being processed no longer accepts the handle (due to such events as handle expiration or updates to the object's attribute set).

FN_E_NO_SUCH_ATTRIBUTE
The object did not have an attribute with the given identifier.

FN_E_PARTIAL_RESULT
(No attribute should be returned with this status code). The enumeration is not yet complete but cannot be continued.

For **FN_E_ATTR_NO_PERMISSION, FN_E_INVALID_ATTR_IDENTIFIER, FN_E_INSUFFICIENT_RESOURCES,** or **FN_E_NO_SUCH_ATTRIBUTE,** the returned attribute contains only the attribute identifier (no value or syntax). For these four status codes and **FN_SUCCESS** (when an attribute was returned), **fn_multigetlist_next( )** can be called again to return another attribute. All other status codes indicate that no more attributes can be returned by **fn_multigetlist_next( ).**

Other status codes, such as **FN_E_COMMUNICATION_FAILURE,** are also possible, in which case, no attribute is returned. In such cases, *status* is set as described in **FN_status_t**(3N) and **xfn_status_codes**(3N).

EXAMPLES

The following code fragment illustrates to obtain all attributes associated with a given name using the **fn_attr_multi_get( )** operations.

```
/* list all attributes associated with given name */

extern FN_string_t *input_string;
FN_ctx_t  *ctx;
FN_composite_name_t *target_name = fn_composite_name_from_string(input_string);
FN_multigetlist_t *ml;
FN_status_t  *status = fn_status_create();
FN_attribute_t *attr;
int done = 0;

ctx = fn_ctx_handle_from_initial(status);
/* error checking on 'status' */

/* attr_ids == 0 indicates all attributes are to be returned */
if ((ml=fn_attr_multi_get(ctx, target_name, 0, status)) == 0) {
```

```
        /* report 'status' and exit */
}

while ((attr=fn_multigetlist_next(ml, status)) && !done) {
        switch (fn_status_code(status)) {
        case FN_SUCCESS:
                /* do something with 'attr' */
                break;
        case FN_E_ATTR_NO_PERMISSION:
        case FN_E_ATTR_INVALID_ATTR_IDENTIFIER:
        case FN_E_NO_SUCH_ATTRIBUTE:
                /* report error using identifier in 'attr' */
                break;
        default:
                /* other error handling */
                done = 1;
        }
        if (attr)
                fn_attribute_destroy(attr);
}

/* check 'status' for reason for end of enumeration and report if necessary */

/* clean up */
fn_multigetlist_destroy(ml, status);

/* report 'status' */
}
```

APPLICATION USAGE

Implementations are not required to return all attributes requested by *attr_ids*. Some may choose to return only the attributes found successfully, followed by a status of **FN_E_PARTIAL_RESULT**; such implementations may not necessarily return attributes identifying those that could not be read. Implementations are not required to return the attributes in any order.

There may be a relationship between the *ctx* argument supplied to **fn_attr_multi_get( )** and the **FN_multigetlist_t** object it returns. For example, some implementations may store the context handle *ctx* within the **FN_multigetlist_t** object for subsequent **fn_multigetlist_next( )** calls. In general, a **fn_ctx_handle_destroy( )** should not be invoked on *ctx* until the enumeration has terminated.

SEE ALSO

FN_attrset_t(3N), **FN_attribute_t**(3N), **FN_composite_name_t**(3N), **FN_ctx_t**(3N), **FN_identifier_t**(3N), **FN_status_t**(3N), **fn_attr_get**(3N), **fn_ctx_list_names**(3N), **xfn_attributes**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_attr_multi_modify – modify multiple attributes associated with named object

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

int fn_attr_multi_modify(FN_ctx_t **ctx***, const FN_composite_name_t** **name***,**
const FN_attrmodlist_t **mods***, FN_attrmodlist_t** ***unexecuted_mods***,**
FN_status_t **status***);**

MT-LEVEL

Safe.

DESCRIPTION

This operation modifies the attributes associated with the object named *name* relative to *ctx*. If *name* is empty, the attributes associated with *ctx* are modified.

In the *mods* parameter, the caller specifies a sequence of modifications that are to be done in order on the attributes. Each modification in the sequence specifies a modification operation code (see **fn_attr_modify**(3N)) and an attribute on which to operate.

The **FN_attrmodlist_t** type is described in **FN_attrmodlist_t**(3N).

RETURN VALUE

fn_attr_multi_modify( ) returns **1** if all the modification operations were performed successfully. The function returns **0** if it any error occurs. If the operation fails, *status* and *unexecuted_mods* are set as described below.

ERRORS

If an error is encountered while performing the list of modifications, *status* indicates the type of error and *unexecuted_mods* is set to a list of unexecuted modifications. The contents of *unexecuted_mods* do not share any state with *mods;* items in *unexecuted_mods* are copies of items in *mods* and appear in the same order in which they were originally supplied in *mods*. The first operation in *unexecuted_mods* is the first one that failed and the code in *status* applies to this modification operation in particular. If *status* indicates failure and a NULL pointer is returned in *unexecuted_mods*, that indicates no modifications were executed.

SEE ALSO

FN_attrmodlist_t(3N), **FN_composite_name_t**(3N), **FN_ctx_t**(3N), **FN_status_t**(3N), **fn_attr_modify**(3N), **xfn_attributes**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_ctx_bind – bind a reference to a name

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

int fn_ctx_bind(FN_ctx_t **ctx***, const FN_composite_name_t** **name***,**
const FN_ref_t **ref***, unsigned int** *exclusive***, FN_status_t** **status***);**

MT-LEVEL

Safe.

DESCRIPTION

This operation binds the supplied reference *ref* to the supplied composite name *name* relative to *ctx*. The binding is made in the target context — that context named by all but the terminal atomic part of *name*. The operation binds the terminal atomic name to the supplied reference in the target context. The target context must already exist.

The value of *exclusive* determines what happens if the terminal atomic part of the name is already bound in the target context. If *exclusive* is nonzero and *name* is already bound, the operation fails. If *exclusive* is zero, the new binding replaces any existing binding.

RETURN VALUE

When the bind operation is successful it returns **1**; on error it returns **0**.

ERRORS

fn_ctx_bind sets *status* as described in **FN_status_t**(3N) and **xfn_status_codes.** Of special relevance for this operation is the following status code:

FN_E_NAME_IN_USE
The supplied name is already in use.

APPLICATION USAGE

The value of *ref* cannot be **NULL**. If the intent is to reserve a name using **fn_ctx_bind( ),** a reference containing no address should be supplied. This reference may be name service-specific or it may be the conventional NULL reference defined in the X/Open registry (see **fns_references**(5)).

If multiple sources are updating a reference, they must synchronize amongst each other when adding, modifying, or removing from the address list of a bound reference.

SEE ALSO

FN_composite_name_t(3N), **FN_ctx_t**(3N), **FN_ref_t**(3N), **FN_status_t**(3N), **fn_ctx_lookup**(3N), **fn_ctx_unbind**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_ctx_create_subcontext – create a subcontext in a context

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_ref_t ***fn_ctx_create_subcontext(FN_ctx_t** ****ctx*****,**
const FN_composite_name_t ****name*****, FN_status_t** ****status*****);**

MT-LEVEL

Safe.

DESCRIPTION

This operation creates a new XFN context of the same type as the target context — that named by all but the terminal atomic component of *name* — and binds it to the supplied composite name.

As with **fn_ctx_bind(),** the target context must already exist. The new context is created and bound in the target context using the terminal atomic name in *name*. The operation returns a reference to the newly created context.

RETURN VALUE

fn_ctx_create_subcontext() returns a reference to the newly created context; if the operation fails, it returns a **NULL** pointer (**0**).

ERRORS

fn_ctx_create_subcontext() sets *status* as described in **FN_status_t**(3N) and **xfn_status_codes**(3N). Of special relevance for this operation is the following status code:

FN_E_NAME_IN_USE
The terminal atomic name already exists in the target context.

APPLICATION USAGE

The new subcontext is an XFN context and is created in the same naming system as the target context. The new subcontext also inherits the same syntax attributes as the target context. XFN does not specify any further properties of the new subcontext. The target context and its naming system determine these.

SEE ALSO

FN_composite_name_t(3N), **FN_ctx_t**(3N), **FN_ref_t**(3N), **FN_status_t**(3N), **fn_ctx_bind**(3N), **fn_ctx_lookup**(3N), **fn_ctx_destroy_subcontext**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_ctx_destroy_subcontext – destroy the named context and remove its binding from the parent context

SYNOPSIS

cc
[*flag* ...] *file* ...
–lxfn
[*library* ...]

#include <xfn/xfn.h>

int fn_ctx_destroy_subcontext(FN_ctx_t ∗*ctx*, **const FN_composite_name_t** ∗*name*,
FN_status_t ∗*status*);

MT-LEVEL

Safe.

DESCRIPTION

This operation destroys the subcontext named by *name* relative to *ctx*, and unbinds the name.

As with **fn_ctx_unbind( ),** this operation succeeds even if the terminal atomic name is not bound in the target context — the context named by all but the terminal atomic name in *name*.

RETURN VALUE

fn_ctx_destroy_subcontext( ) returns **1** on success and **0** on failure.

ERRORS

fn_ctx_destroy_subcontext( ) sets *status* as described in **FN_status_t**(3N) and **xfn_status_codes**(3N). Of special relevance for **fn_ctx_destroy_subcontext( )** are the following status codes:

FN_E_CTX_NOT_A_CONTEXT
name does not name a context.

FN_E_CTX_NOT_EMPTY
The naming system being asked to do the destroy does not support removal of a context that still contains bindings.

APPLICATION USAGE

Some aspects of this operation are not specified by XFN, but are determined by the target context and its naming system. For example, XFN does not specify what happens if the named subcontext is non-empty when the operation is invoked.

In naming systems that support attributes, and store the attributes along with names or contexts, this operation removes the name, the context, and its associated attributes.

Normal resolution always follows links. In a **fn_ctx_destroy_subcontext( )** operation, resolution of *name* continues to the target context; the terminal atomic name is not resolved. If the terminal atomic name is bound to a link, the link is not followed and the operation fails with **FN_E_CTX_NOT_A_CONTEXT** because the name is not bound to a context.

SEE ALSO

FN_ctx_t(3N), **FN_composite_name_t**(3N), **FN_status_t**(3N), **fn_ctx_create_subcontext**(3N), **fn_ctx_unbind**(3N)**, xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_ctx_get_ref – return a context's reference

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_ref_t ∗**fn_ctx_get_ref(const FN_ctx_t** ∗*ctx*, **FN_status_t** ∗*status*);

MT-LEVEL

Safe.

DESCRIPTION

This operation returns a reference to the supplied context object.

RETURN VALUE

fn_ctx_get_ref( ) returns a pointer to an **FN_ref_t** object if the operation succeeds, it returns **0** if the operation fails.

ERRORS

fn_ctx_get_ref( ) sets *status* as described in **FN_status_t**(3N) and **xfn_status_codes**(3N). The following status code is of particular relevance to this operation:

FN_E_OPERATION_NOT_SUPPORTED
Using the **fn_ctx_get_ref( )** operation on the Initial Context returns this status code.

APPLICATION USAGE

fn_ctx_get_ref( ) cannot be used on the Initial Context. **fn_ctx_get_ref( )** can be used on contexts bound in the Initial Context (in other words, the bindings in the Initial Context have references).

If the context handle was created earlier using the **fn_ctx_handle_from_ref( )** operation, the reference returned by the **fn_ctx_get_ref( )** operation may not necessarily be exactly the same in content as that originally supplied. For example, **fn_ctx_handle_from_ref( )** may construct the context handle from one address from the list of addresses. The context implementation may return with a call to **fn_ctx_get_ref( )** only that address, or a more complete list of addresses than what was supplied in **fn_ctx_handle_from_ref( ).**

SEE ALSO

FN_ctx_t(3N), **FN_ref_t**(3N), **FN_status_t**(3N), **fn_ctx_handle_from_initial**(3N), **fn_ctx_handle_from_ref**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_ctx_get_syntax_attrs – return syntax attributes associated with named context

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_attrset_t ***fn_ctx_get_syntax_attrs(FN_ctx_t** ***ctx*,
const FN_composite_name_t ***name*, **FN_status_t** ***status*);

MT-LEVEL

Safe.

DESCRIPTION

Each context has an associated set of syntax-related attributes. This operation returns the syntax attributes associated with the context named by *name* relative to the context *ctx*.

The attributes must contain the attribute **fn_syntax_type** (**FN_ID_STRING** format). If the context supports a syntax that conforms to the XFN standard syntax model, **fn_syntax_type** is set to "standard" (ASCII attribute syntax) and the attribute set contains the rest of the relevant syntax attributes described in **xfn_compound_names**(3N).

This operation is different from other XFN attribute operations in that these syntax attributes could be obtained directly from the context. Attributes obtained through other XFN attribute operations may not necessarily be associated with the context; they may be associated with the reference of context, rather than the context itself (see **xfn_attributes**(3N)).

RETURN VALUE

fn_ctx_get_syntax_attrs( ) returns an attribute set if successful; it returns a **NULL** pointer (**0**) if the operation fails.

ERRORS

fn_ctx_get_syntax_attrs( ) sets *status* as described in **FN_status_t**(3N) and **xfn_status_codes**(3N).

APPLICATION USAGE

Implementations may choose to support other syntax types in addition to, or in place of, the XFN standard syntax model, in which case, the value of the **fn_syntax_type** attribute would be set to an implementation-specific string, and different or additional syntax attributes will be in the set.

Syntax attributes of a context may be generated automatically by a context, in response to **fn_ctx_get_syntax_attrs( ),** or they may be created and updated using the base attribute operations. This is implementation-dependent.

SEE ALSO

FN_attrset_t(3N), **FN_composite_name_t**(3N), **FN_compound_name_t**(3N), **FN_ctx_t**(3N), **FN_status_t**(3N), **fn_attr_get**(3N), **fn_attr_multi_get**(3N), **xfn_compound_names**(3N), **xfn_attributes**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_ctx_handle_destroy – release storage associated with context handle

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

void fn_ctx_handle_destroy(FN_ctx_t **ctx****);**

MT-LEVEL

Safe.

DESCRIPTION

This operation destroys the context handle *ctx* and allows the implementation to free resources associated with the context handle. This operation does not affect the state of the context itself.

SEE ALSO

FN_ctx_t(3N), **fn_ctx_handle_from_initial**(3N), **fn_ctx_handle_from_ref**(3N), **xfn**(3N)

NAME

fn_ctx_handle_from_initial – return a handle to the Initial Context

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_ctx_t ∗**fn_ctx_handle_from_initial(FN_status_t** ∗*status***);**

MT-LEVEL

MT-Safe.

DESCRIPTION

This operation returns a handle to the caller's Initial Context. On successful return, the handle points to a context which meets the specification of the XFN Initial Context (see **fns_initial_context**(5)).

RETURN VALUE

fn_ctx_handle_from_initial( ) returns a pointer to an **FN_ctx_t** object if the operation succeeds; it returns a **NULL** pointer (**0**) otherwise.

ERRORS

fn_ctx_handle_from_initial( ) sets only the status code portion of the status object *status.*

SEE ALSO

FN_ctx_t(3N), **FN_status_t**(3N), **fns_initial_context**(5), **fn_ctx_get_ref**(3N), **fn_ctx_handle_from_ref**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_ctx_handle_from_ref – construct a handle to a context object using the given reference

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_ctx_t ∗**fn_ctx_handle_from_ref(const FN_ref_t** ∗*ref***, FN_status_t** ∗*status***);**

MT-LEVEL

Safe.

DESCRIPTION

This operation creates a handle to an **FN_ctx_t** object using an **FN_ref_t** object for that context.

RETURN VALUE

This operations returns a pointer to an **FN_ctx_t** object if the operation succeeds, otherwise, it returns a **NULL** pointer (**0**).

ERRORS

fn_ctx_handle_from_ref( ) sets *status* as described in **FN_status_t**(3N) and **xfn_status_codes**(3N). The following status code is of particular relevance to this operation.

FN_E_NO_SUPPORTED_ADDRESS
A context object could not be constructed from a particular reference. The reference contained no address type over which the context interface was supported.

SEE ALSO

FN_ctx_t(3N), **FN_ref_t**(3N), **FN_status_t**(3N), **fn_ctx_handle_destroy**(3N), **fn_ctx_get_ref**(3N), **fns_references**(5), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_ctx_list_bindings, FN_bindinglist_t, fn_bindinglist_next, fn_bindinglist_destroy – list the atomic names and references bound in a context

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_bindinglist_t ***fn_ctx_list_bindings(FN_ctx_t** ****ctx*****,**
const FN_composite_name_t ****name*****, FN_status_t** ****status*****);**

FN_string_t ***fn_bindinglist_next(FN_bindinglist_t** ****bl*****, FN_ref_t** ******ref*****,**
FN_status_t ****status*****);**

void fn_bindinglist_destroy(FN_bindinglist_t ****bl*****, FN_status_t** ****status*****);**

MT-LEVEL

Safe.

DESCRIPTION

This set of operations is used to list the names and bindings in the context named by *name* relative to the context *ctx*. Note that *name* must name a context. If the intent is to list the contents of *ctx*, *name* should be an empty composite name.

The semantics of these operations are similar to those for listing names See **fn_ctx_list_names**(3N). In addition to a name string being returned, **fn_bindinglist_next( )** also returns the reference of the binding for each member of the enumeration.

SEE ALSO

FN_composite_name_t(3N), **FN_ctx_t**(3N), **FN_ref_t**(3N), **FN_status_t**(3N), **FN_string_t**(3N), **fn_ctx_list_names**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_ctx_list_names, FN_namelist_t, fn_namelist_next, fn_namelist_destroy – list the atomic names bound in a context

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_namelist_t ∗**fn_ctx_list_names(FN_ctx_t** ∗*ctx***, const FN_composite_name_t** ∗*name***, FN_status_t** ∗*status***);**

FN_string_t ∗**fn_namelist_next(FN_namelist_t** ∗*nl***, FN_status_t** ∗*status***);**

void fn_namelist_destroy(FN_namelist_t ∗*nl***, FN_status_t** ∗*status***);**

MT-LEVEL

Safe.

DESCRIPTION

This set of operations is used to list the names bound in the target context named *name* relative to the context *ctx* . Note that *name* must name a context. If the intent is to list the contents of *ctx*, *name* should be an empty composite name.

The call to **fn_ctx_list_names( )** initiates the enumeration process. It returns a handle to an **FN_namelist_t** object that can be used to enumerate the names in the target context.

The operation **fn_namelist_next( )** returns the next name in the enumeration identified by *nl* and updates *nl* to indicate the state of the enumeration. Successive calls to **fn_namelist_next( )** using *nl* return successive names in the enumeration and further update the state of the enumeration. **fn_namelist_next( )** returns a **NULL** pointer when the enumeration has been completed.

fn_namelist_destroy( ) is used to release resources used during the enumeration. This may be invoked at any time to terminate the enumeration.

RETURN VALUE

fn_ctx_list_names( ) returns a pointer to an **FN_namelist_t** object if the enumeration is successfully initiated; otherwise it returns a **NULL** pointer.

fn_namelist_next( ) returns a **NULL** pointer if no more names can be returned in the enumeration.

In the case of a failure, these operations return in *status* a code indicating the nature of the failure.

ERRORS

Each successful call to **fn_namelist_next( )** returns a name and sets *status* to **FN_SUCCESS.**

When **fn_namelist_next( )** returns a **NULL** pointer, it indicates that no more names can be returned. *status* is set in the following way:

FN_SUCCESS
The enumeration has completed successfully.

FN_E_INVALID_ENUM_HANDLE
The supplied enumeration handle is not valid. Possible reasons could be that the handle was from another enumeration, or the context being enumerated no

longer accepts the handle (due to such events as handle expiration or updates to the context).

FN_E_PARTIAL_RESULT
The enumeration is not yet complete but cannot be continued.

Other status codes, such as **FN_E_COMMUNICATION_FAILURE,** are also possible in calls to **fn_ctx_list_names( ), fn_namelist_next( )** and **fn_namelist_destroy( ).** These functions set *status* for these other status codes as described in **FN_status_t**(3N) and **xfn_status_codes**(3N).

EXAMPLE

The following code fragment illustrates a how the list names operations may be used.

```
extern FN_string_t *user_input;
FN_ctx_t *ctx;
FN_composite_name_t *target_name = fn_composite_name_from_string(user_input);
FN_status_t *status = fn_status_create();
FN_string_t *name;
FN_namelist_t *nl;

ctx = fn_ctx_handle_from_initial(status);
/* error checking on 'status' */

if ((nl=fn_ctx_list_names(ctx, target_name, status)) == 0) {
        /* report 'status' and exit */
}

while (name=fn_namelist_next(nl, status)) {
        /* do something with 'name' */
        fn_string_destroy(name);
}

/* check 'status' for reason for end of enumeration and report if necessary */

/* clean up */
fn_namelist_destroy(nl, status);

/* report 'status' */
```

APPLICATION USAGE

The names enumerated using **fn_namelist_next( )** are not ordered in any way. There is no guaranteed relation between the order in which names are added to a context and the order of names obtained by enumeration. The specification does *not* guarantee that any two series of enumerations will return the names in the same order.

When a name is added to or removed from a context, this may or may not invalidate the enumeration handle that the client holds for that context. If the enumeration handle becomes invalid, the status code **FN_E_INVALID_ENUM_HANDLE** is returned in *status.* If the enumeration handle remains valid, the update may or may not be visible to the client.

In addition, there may be a relationship between the *ctx* argument supplied to **fn_ctx_list_names( )** and the **FN_namelist_t** object it returns. For example, some implementations may store the context handle *ctx* within the **FN_namelist_t** object for subsequent **fn_namelist_next( )** calls. In general, a **fn_ctx_handle_destroy( )** should not be invoked on *ctx* until the enumeration has terminated.

SEE ALSO

FN_composite_name_t(3N), **FN_ctx_t**(3N), **FN_status_t**(3N), **FN_string_t**(3N), **fn_ctx_handle_destroy**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_ctx_lookup – look up name in context

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_ref_t ∗**fn_ctx_lookup(FN_ctx_t** ∗*ctx*, **const FN_composite_name_t** ∗*name*, **FN_status_t** ∗*status*);

MT-LEVEL

Safe.

DESCRIPTION

This operation returns the reference bound to *name* relative to the context *ctx*.

RETURN VALUE

If the operation succeeds, the **fn_ctx_lookup( )** function returns a handle to the reference bound to *name*. Otherwise, **0** is returned and *status* is set appropriately.

ERRORS

fn_ctx_lookup( ) sets *status* as described **FN_status_t**(3N) and **xfn_status_codes**(3N).

APPLICATION USAGE

Some naming services may not always have reference information for all names in their contexts; for such names, such naming services may return a special reference whose type indicates that the name is not bound to any address. This reference may be name service specific or it may be the conventional NULL reference defined in the X/Open registry. See **fns_references**(5).

SEE ALSO

FN_composite_name_t(3N), **FN_ctx_t**(3N), **FN_ref_t**(3N), **FN_status_t**(3N), **fns_references**(5), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_ctx_lookup_link – look up the link reference bound to a name

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

FN_ref_t ∗**fn_ctx_lookup_link(FN_ctx_t** ∗*ctx*, **const FN_composite_name_t** ∗*name*, **FN_status_t** ∗*status*);

MT-LEVEL

Safe.

DESCRIPTION

This operation returns the XFN link bound to *name*. The terminal atomic part of *name* must be bound to an XFN link.

The normal **fn_ctx_lookup**(3N) operation follows all links encountered, including any bound to the terminal atomic part of *name*. This operation differs from the normal lookup in that when the terminal atomic part of *name* is an XFN link, this link is not followed, and the operation returns the link.

RETURN VALUE

If **fn_ctx_lookup_link( )** fails, a **NULL** pointer (**0**) is returned.

ERRORS

fn_ctx_lookup_link( ) sets *status* as described in **FN_status_t**(3N) and **xfn_status_codes**(3N). Of special relevance for **fn_ctx_lookup_link( )** is the following status code:

FN_E_MALFORMED_LINK
name resolved to a reference that was not a link.

SEE ALSO

FN_composite_name_t(3N), **FN_ctx_t**(3N), **FN_ref_t**(3N), **FN_status_t**(3N), **fn_ctx_lookup**(3N), **xfn_status_codes**(3N), **xfn_links**(3N), **xfn**(3N)

NAME

fn_ctx_rename – rename the name of a binding

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

int fn_ctx_rename(FN_ctx_t **ctx***, const FN_composite_name_t** **oldname***,**
const FN_composite_name_t **newname***, unsigned int** *exclusive***,**
FN_status_t **status***);**

MT-LEVEL

Safe.

DESCRIPTION

The **fn_ctx_rename( )** operation binds the reference currently bound to *oldname* relative to *ctx*, to the name *newname*, and unbinds *oldname*. *newname* is resolved relative to the target context (that named by all but the terminal atomic part of *oldname)*.

If *exclusive* is zero, the operation overwrites any old binding of *newname*. If *exclusive* is nonzero, the operation fails if *newname* is already bound.

RETURN VALUE

fn_ctx_rename( ) returns **1** if the operation is successful, **0** otherwise.

ERRORS

fn_ctx_rename( ) sets *status* as described **FN_status_t**(3N) and **xfn_status_codes**(3N).

APPLICATION USAGE

The only restriction that XFN places on *newname* is that it be resolved relative to the target context. XFN does not specify further restrictions on *newname*. For example, in some implementations, *newname* might be restricted to be a name in the same naming system as the terminal component of *oldname*. In another implementation, *newname* might be restricted to be an atomic name.

Normal resolution always follows links. In a **fn_ctx_rename()** operation, resolution of *oldname* continues to the target context; the terminal atomic name is not resolved. If the terminal atomic name is bound to a link, the link is not followed and the operation binds *newname* to the link and unbinds the terminal atomic name of *oldname*.

In naming systems that support attributes and store the attributes along with the names, the unbind of the terminal atomic name of *oldname* also removes its associated attributes. It is implementation dependent whether these attributes become associated with *newname*.

SEE ALSO

FN_composite_name_t(3N), **FN_ctx_t**(3N), **FN_ref_t**(3N), **FN_status_t**(3N), **fn_ctx_bind**(3N)**, fn_ctx_unbind**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fn_ctx_unbind – unbind a name from a context

SYNOPSIS

cc [*flag* ...] *file* ... **–lxfn** [*library* ...]

#include <xfn/xfn.h>

int fn_ctx_unbind(FN_ctx_t ∗*ctx*, **const FN_composite_name_t** ∗*name*, **FN_status_t** ∗*status*);

MT-LEVEL

Safe.

DESCRIPTION

This operation removes the terminal atomic name in *name* from the the target context — that named by all but the terminal atomic part of *name*.

This operation is successful even if the terminal atomic name was not bound in target context, but fails if any of the intermediate names are not bound. **fn_ctx_unbind( )** is idempotent.

RETURN VALUE

The operation returns **1** if successful, and **0** otherwise.

ERRORS

fn_ctx_unbind( ) sets *status* as described in **FN_status_t** and **xfn_status_codes**(3N).

Certain naming systems may disallow unbinding a name if the name is bound to an existing context in order to avoid orphan contexts that cannot be reached via any name. In such situations, the status code **FN_E_OPERATION_NOT_SUPPORTED** is returned.

APPLICATION USAGE

In naming systems that support attributes, and store the attributes along with the names, the unbind operation removes the name and its associated attributes.

Normal resolution always follows links. In an **fn_ctx_unbind( )** operation, resolution of *name* continues to the target context; the terminal atomic name is not resolved. If the terminal atomic name is bound to a link, the link is not followed and the link itself is unbound from the terminal atomic name.

SEE ALSO

FN_composite_name_t(3N), **FN_ctx_t**(3N), **FN_ref_t**(3N), **FN_status_t**(3N), **fn_ctx_bind**(3N), **fn_ctx_lookup**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

fnmatch – match filename or path name

SYNOPSIS

```
#include <fnmatch.h>
int fnmatch( const char *pattern, const char *string, int flags);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **fnmatch( )** function matches patterns as described on the **fnmatch**(5) manual page. It checks the *string* argument to see if it matches the *pattern* argument.

The *flags* argument modifies the interpretation of *pattern* and *string*. It is the bitwise inclusive OR of zero or more of the following flags defined in the header **<fnmatch.h>**.

FNM_PATHNAME If set, a slash (/) character in *string* will be explicitly matched by a slash in *pattern*; it will not be matched by either the asterisk (*) or question-mark (**?**) special characters, nor by a bracket (**[ ]**) expression.

If not set, the slash character is treated as an ordinary character.

FNM_NOESCAPE If not set, a backslash character (\) in *pattern* followed by any other character will match that second character in *string*. In particular, "\\" will match a backslash in *string*.

If set, a backslash character will be treated as an ordinary character.

FNM_PERIOD If set, a leading period in *string* will match a period in *pattern*; where the location of ''leading'' is indicated by the value of **FNM_PATHNAME**:

- If **FNM_PATHNAME** is set, a period is ''leading'' if it is the first character in *string* or if it immediately follows a slash.
- If **FNM_PATHNAME** is not set, a period is ''leading'' only if it is the first character of *string*.

If not set, no special restrictions are placed on matching a period.

RETURN VALUES

The following values are returned:

0	*string* matches the pattern specified by *pattern*.
FNM_NOMATCH	there is no match. **FNM_NOMATCH** is defined in the header **<fnmatch.h>**.
non-zero	an error has occurred.

USAGE

The **fnmatch( )** function has two major uses. It could be used by an application or utility that needs to read a directory and apply a pattern against each entry. The **find**(1) utility is an example of this. It can also be used by the **pax** utility to process its *pattern* operands, or by applications that need to match strings in a similar manner.

The name **fnmatch( )** is intended to imply *filename* match, rather than *pathname* match. The default action of this function is to match filenames, rather than path names, since it gives no special significance to the slash character. With the **FNM_PATHNAME** flag, **fnmatch( )** does match path names, but without tilde expansion, parameter expansion, or special treatment for period at the beginning of a filename.

SEE ALSO

find(1), **glob**(3C), **wordexp**(3C), **fnmatch**(5)

NAME

fopen, freopen – open a stream

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <stdio.h>

FILE ∗**fopen(***file***,** *mode***)**
const char ∗*file***,** ∗*mode***;**

FILE ∗**freopen(***file, mode, iop***)**
const char ∗*file, ∗mode***;**
register FILE ∗*iop***;**

DESCRIPTION

fopen( ) opens the file named by *file* and associates a stream with it. If the open succeeds, **fopen( )** returns a pointer to be used to identify the stream in subsequent operations.

file points to a character string that contains the name of the file to be opened.

mode is a character string having one of the following values:

- **r** open for reading
- **w** truncate or create for writing
- **a** append: open for writing at end of file, or create for writing
- **r+** open for update (reading and writing)
- **w+** truncate or create for update
- **a+** append; open or create for update at EOF

freopen( ) opens the file named by *file* and associates the stream pointed to by *iop* with it. The *mode* argument is used just as in **fopen( )**. The original stream is closed, regardless of whether the open ultimately succeeds. If the open succeeds, **freopen( )** returns the original value of *iop*.

freopen( ) is typically used to attach the preopened streams associated with **stdin**, **stdout**, and **stderr** to other files.

When a file is opened for update, both input and output may be done on the resulting stream. However, output may not be directly followed by input without an intervening **fseek**(3S) or **rewind**(3S), and input may not be directly followed by output without an intervening **fseek**(3S) or **rewind**(3S). An input operation which encounters EOF will fail.

RETURN VALUES

fopen( ) and **freopen( )** return a **NULL** pointer on failure.

SEE ALSO

open(2), **fclose**(3S), **fopen**(3S), **freopen**(3S), **fseek**(3S), **malloc**(3C), **rewind**(3S)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

In order to support the same number of open files that the system does, **fopen()** must allocate additional memory for data structures using **malloc**(3C) after 64 files have been opened. This confuses some programs which use their own memory allocators.

The interfaces of **fopen()** and **freopen()** differ from the Standard I/O Functions **fopen**(3S) and **freopen**(3S). The Standard I/O Functions distinguish binary from text files with an additional use of '**b**' as part of the *mode*. This enables portability of **fopen**(3S) and **freopen**(3S) beyond SunOS 4.*X* systems.

NAME

fopen, freopen, fdopen – open a stream

SYNOPSIS

#include <stdio.h>

FILE ∗**fopen(const char** ∗*filename*, **const char** ∗*type*);

FILE ∗**freopen(const char** ∗*filename*, **const char** ∗*type*, **FILE** ∗*stream*);

FILE ∗**fdopen(int** *fildes*, **const char** ∗*type*);

MT-LEVEL

MT-Safe

DESCRIPTION

fopen( ) opens the file named by *filename* and associates a *stream* with it. **fopen( )** returns a pointer to the **FILE** structure associated with the *stream*.

filename points to a character string that contains the name of the file to be opened.

type is a character string beginning with one of the following sequences:

"r" or **"rb"**	open for reading
"w" or **"wb"**	truncate to zero length or create for writing
"a" or **"ab"**	append; open for writing at end of file, or create for writing
"r+", **"r+b"** or **"rb+"**	open for update (reading and writing)
"w+", **"w+b"** or **"wb+"**	truncate or create for update
"a+", **"a+b"** or **"ab+"**	append; open or create for update at end-of-file

The ''**b**'' is ignored in the above *type*s. The ''**b**'' exists to distinguish binary files from text files. However, there is no distinction between these types of files on a UNIX system.

freopen( ) substitutes the named file in place of the open *stream*. A flush is first attempted, and then the original *stream* is closed, regardless of whether the open ultimately succeeds. Failure to flush or close *stream* successfully is ignored. **freopen( )** returns a pointer to the **FILE** structure associated with *stream*.

freopen( ) is typically used to attach the preopened *streams* associated with **stdin**, **stdout**, and **stderr** to other files. **stderr** is by default unbuffered, but the use of **freopen( )** will cause it to become buffered or line-buffered.

fdopen( ) associates a *stream* with a file descriptor. File descriptors are obtained from **open**(2), **dup**(2), **creat**(2), or **pipe**(2), which open files but do not return pointers to a **FILE** structure *stream*. Streams are necessary input for almost all of the Section 3S library routines. The *type* of *stream* must agree with the mode of the open file. The file position indicator associated with *stream* is set to the position indicated by the file offset associated with *fildes*.

When a file is opened for update, both input and output may be done on the resulting *stream*. However, output may not be directly followed by input without an intervening **fflush( )**, **fseek( )**, **fsetpos( )**, or **rewind( )**, and input may not be directly followed by output without an intervening **fseek( )**, **fsetpos( )**, or **rewind( )**, or an input operation that encounters end-of-file.

When a file is opened for append (that is, when *type* is **"a"**, **"ab"**, **"a+"**, or **"ab+"**), it is impossible to overwrite information already in the file. **fseek()** may be used to reposition the file pointer to any position in the file, but when output is written to the file, the current file pointer is disregarded. All output is written at the end of the file and causes the file pointer to be repositioned at the end of the output. If two separate processes open the same file for append, each process may write freely to the file without fear of destroying output being written by the other. The output from the two processes will be intermixed in the file in the order in which it is written.

When opened, a *stream* is fully buffered if and only if it can be determined not to refer to an interactive device. The error and end-of-file indicators are cleared for the *stream*.

RETURN VALUES

The functions **fopen()** and **freopen()** return a null pointer if *path* cannot be accessed, or if *type* is invalid, or if the file cannot be opened.

The function **fdopen()** returns a null pointer if *fildes* is not an open file descriptor, or if *type* is invalid, or if the file cannot be opened.

The functions **fopen()** or **fdopen()** may fail and not set **errno** if there are no free **stdio** streams.

File descriptors used by **fdopen()** must be less than 255.

SEE ALSO

close(2), **creat**(2), **dup**(2), **open**(2), **pipe**(2), **write**(2), **fclose**(3S), **fseek**(3S), **setbuf**(3S), **stdio**(3S)

NAME

form_cursor, pos_form_cursor – position forms window cursor

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lform -lcurses** [*library* . .]

#include <form.h>

int pos_form_cursor(FORM ∗*form***);**

MT-LEVEL

Unsafe

DESCRIPTION

pos_form_cursor( ) moves the form window cursor to the location required by the form driver to resume form processing. This may be needed after the application calls a **curses** library I/O routine.

RETURN VALUES

pos_form_cursor( ) returns one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.
E_BAD_ARGUMENT	–	An argument is incorrect.
E_NOT_POSTED	–	The form is not posted.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_data, data_ahead, data_behind – tell if forms field has off-screen data ahead or behind

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lform -lcurses** [*library* . .]

#include <form.h>

int data_ahead(FORM ∗*form***);**

int data_behind(FORM ∗*form***);**

MT-LEVEL

Unsafe

DESCRIPTION

data_ahead( ) returns **TRUE** (1) if the current field has more off-screen data ahead; otherwise it returns **FALSE** (0).

data_behind( ) returns **TRUE** (1) if the current field has more off-screen data behind; otherwise it returns **FALSE** (0).

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_driver – command processor for the forms subsystem

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lform -lcurses** [*library* . .]

#include <form.h>

int form_driver(FORM **form***, int** *c***);**

MT-LEVEL

Unsafe

DESCRIPTION

form_driver( ) is the workhorse of the **forms** subsystem; it checks to determine whether the character *c* is a **forms** request or data. If it is a request, the form driver executes the request and reports the result. If it is data (a printable ASCII character), it enters the data into the current position in the current field. If it is not recognized, the form driver assumes it is an application-defined command and returns **E_UNKNOWN_COMMAND**. Application defined commands should be defined relative to **MAX_COMMAND**, the maximum value of a request listed below.

Form driver requests:

REQ_NEXT_PAGE	Move to the next page.
REQ_PREV_PAGE	Move to the previous page.
REQ_FIRST_PAGE	Move to the first page.
REQ_LAST_PAGE	Move to the last page.
REQ_NEXT_FIELD	Move to the next field.
REQ_PREV_FIELD	Move to the previous field.
REQ_FIRST_FIELD	Move to the first field.
REQ_LAST_FIELD	Move to the last field.
REQ_SNEXT_FIELD	Move to the sorted next field.
REQ_SPREV_FIELD	Move to the sorted prev field.
REQ_SFIRST_FIELD	Move to the sorted first field.
REQ_SLAST_FIELD	Move to the sorted last field.
REQ_LEFT_FIELD	Move left to field.
REQ_RIGHT_FIELD	Move right to field.
REQ_UP_FIELD	Move up to field.
REQ_DOWN_FIELD	Move down to field.
REQ_NEXT_CHAR	Move to the next character in the field.
REQ_PREV_CHAR	Move to the previous character in the field.
REQ_NEXT_LINE	Move to the next line in the field.
REQ_PREV_LINE	Move to the previous line in the field.
REQ_NEXT_WORD	Move to the next word in the field.
REQ_PREV_WORD	Move to the previous word in the field.
REQ_BEG_FIELD	Move to the first char in the field.
REQ_END_FIELD	Move after the last char in the field.
REQ_BEG_LINE	Move to the beginning of the line.
REQ_END_LINE	Move after the last char in the line.
REQ_LEFT_CHAR	Move left in the field.

REQ_RIGHT_CHAR	Move right in the field.
REQ_UP_CHAR	Move up in the field.
REQ_DOWN_CHAR	Move down in the field.
REQ_NEW_LINE	Insert / overlay a new line.
REQ_INS_CHAR	Insert the blank character at the cursor.
REQ_INS_LINE	Insert a blank line at the cursor.
REQ_DEL_CHAR	Delete the character at the cursor.
REQ_DEL_PREV	Delete the character before the cursor.
REQ_DEL_LINE	Delete the line at the cursor.
REQ_DEL_WORD	Delete the word at the cursor.
REQ_CLR_EOL	Clear to the end of the line.
REQ_CLR_EOF	Clear to the end of the field.
REQ_CLR_FIELD	Clear the entire field.
REQ_OVL_MODE	Enter overlay mode.
REQ_INS_MODE	Enter insert mode.
REQ_SCR_FLINE	Scroll the field forward a line.
REQ_SCR_BLINE	Scroll the field backward a line.
REQ_SCR_FPAGE	Scroll the field forward a page.
REQ_SCR_BPAGE	Scroll the field backward a page.
REQ_SCR_FHPAGE	Scroll the field forward half a page.
REQ_SCR_BHPAGE	Scroll the field backward half a page.
REQ_SCR_FCHAR	Horizontal scroll forward a character.
REQ_SCR_BCHAR	Horizontal scroll backward a character.
REQ_SCR_HFLINE	Horizontal scroll forward a line.
REQ_SCR_HBLINE	Horizontal scroll backward a line.
REQ_SCR_HFHALF	Horizontal scroll forward half a line.
REQ_SCR_HBHALF	Horizontal scroll backward half a line.
REQ_VALIDATION	Validate field.
REQ_PREV_CHOICE	Display the previous field choice.
REQ_NEXT_CHOICE	Display the next field choice.

RETURN VALUES

form_driver() returns one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.
E_BAD_ARGUMENT	–	An argument is incorrect.
E_NOT_POSTED	–	The form is not posted.
E_INVALID_FIELD	–	The field contents are invalid.
E_BAD_STATE	–	The routine was called from an initialization or termination function.
E_REQUEST_DENIED	–	The form driver request failed.
E_UNKNOWN_COMMAND	–	An unknown request was passed to

the the form driver.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_field, set_form_fields, form_fields, field_count, move_field – connect fields to forms

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lform -lcurses** [*library* . .]

#include <form.h>

int set_form_fields(FORM ∗*form***, FIELD** ∗∗*field***);**

FIELD ∗∗**form_fields(FORM** ∗*form***);**

int field_count(FORM ∗*form***);**

int move_field(FIELD ∗*field***, int** *frow***, int** *fcol***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_form_fields( ) changes the fields connected to *form* to *fields*. The original fields are disconnected.

form_fields( ) returns a pointer to the field pointer array connected to *form*.

field_count( ) returns the number of fields connected to *form*.

move_field( ) moves the disconnected *field* to the location *frow*, *fcol* in the **forms** subwindow.

RETURN VALUES

form_fields( ) returns **NULL** on error.

field_count( ) returns **-1** on error.

set_form_fields( ) and **move_field( )** return one of the following:

E_OK	–	The function returned successfully.
E_CONNECTED	–	The field is already connected to a form.
E_SYSTEM_ERROR	–	System error.
E_BAD_ARGUMENT	–	An argument is incorrect.
E_POSTED	–	The form is posted.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_field_attributes, set_field_fore, field_fore, set_field_back, field_back, set_field_pad, field_pad – format the general display attributes of forms

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lform -lcurses** [*library* . .]

#include <form.h>

int set_field_fore(FIELD **field***, chtype** *attr***);**

chtype field_fore(FIELD **field***);**

int set_field_back(FIELD **field***, chtype** *attr***);**

chtype field_back(FIELD **field***);**

int set_field_pad(FIELD **field***, int** *pad***);**

int field_pad(FIELD **field***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_field_fore( ) sets the foreground attribute of *field*. The foreground attribute is the low-level **curses** display attribute used to display the field contents. **field_fore( )** returns the foreground attribute of *field*.

set_field_back( ) sets the background attribute of *field*. The background attribute is the low-level **curses** display attribute used to display the extent of the field. **field_back( )** returns the background attribute of *field*.

set_field_pad( ) sets the pad character of *field* to *pad*. The pad character is the character used to fill within the field. **field_pad( )** returns the pad character of *field*.

RETURN VALUES

field_fore( ), **field_back( )**, and **field_pad( )** return default values if *field* is **NULL**. If *field* is not **NULL** and is not a valid **FIELD** pointer, the return value from these routines is undefined.

set_field_fore( ), **set_field_back( )**, and **set_field_pad( )** return one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.
E_BAD_ARGUMENT	–	An argument is incorrect.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_field_buffer, set_field_buffer, field_buffer, set_field_status, field_status, set_max_field – set and get forms field attributes

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lform -lcurses** [*library* . .]

#include <form.h>

int set_field_buffer(FIELD ∗*field***, int** *buf***, char** ∗*value***);**

char ∗**field_buffer(FIELD** ∗*field***, int** *buf***);**

int set_field_status(FIELD ∗*field***, int** *status***);**

int field_status(FIELD ∗*field***);**

int set_max_field(FIELD ∗*field***, int** *max***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_field_buffer( ) sets buffer *buf* of *field* to *value*. Buffer 0 stores the displayed contents of the field. Buffers other than 0 are application specific and not used by the **forms** library routines. **field_buffer( )** returns the value of *field* buffer *buf*.

Every field has an associated status flag that is set whenever the contents of field buffer 0 changes. **set_field_status( )** sets the status flag of *field* to *status*. **field_status( )** returns the status of *field*.

set_max_field( ) sets a maximum growth on a dynamic field, or if *max*=**0** turns off any maximum growth.

RETURN VALUES

field_buffer( ) returns **NULL** on error.

field_status( ) returns **TRUE** or **FALSE**.

set_field_buffer( ), **set_field_status( )**, and **set_max_field( )** return one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.
E_BAD_ARGUMENT	–	An argument is incorrect.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_field_info, field_info, dynamic_field_info – get forms field characteristics

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lform -lcurses** [*library* . .]

#include <form.h>

int field_info(FIELD ∗*field***, int** ∗*rows***, int** ∗*cols***, int** ∗*frow***, int** ∗*fcol***, int** ∗*nrow***, int** ∗*nbuf***);**

int dynamic_field_info(FIELD ∗*field***, int** ∗*drows***, int** ∗*dcols***, int** ∗*max***);**

MT-LEVEL

Unsafe

DESCRIPTION

field_info( ) returns the size, position, and other named field characteristics, as defined in the original call to **new_field( )**, to the locations pointed to by the arguments *rows*, *cols*, *frow*, *fcol*, *nrow*, and *nbuf*.

dynamic_field_info( ) returns the actual size of the *field* in the pointer arguments *drows*, *dcols* and returns the maximum growth allowed for *field* in *max*. If no maximum growth limit is specified for *field*, *max* will contain 0. A field can be made dynamic by turning off the field option **O_STATIC**.

RETURN VALUES

These routines return one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.
E_BAD_ARGUMENT	–	An argument is incorrect.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_field_just, set_field_just, field_just – format the general appearance of forms

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lform -lcurses** [*library* . .]

#include <form.h>

int set_field_just(FIELD ∗*field***, int** *justification***);**

int field_just(FIELD ∗*field***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_field_just( ) sets the justification for *field*. Justification may be one of:
NO_JUSTIFICATION, **JUSTIFY_RIGHT**, **JUSTIFY_LEFT**, or **JUSTIFY_CENTER**.

The field justification will be ignored if *field* is a dynamic field.

field_just( ) returns the type of justification assigned to *field*.

RETURN VALUES

field_just( ) returns one of the following:

NO_JUSTIFICATION, **JUSTIFY_RIGHT**,
JUSTIFY_LEFT, or **JUSTIFY_CENTER**.

set_field_just( ) returns one of the following:

E_OK	–The function returned successfully.
E_SYSTEM_ERROR	–System error.
E_BAD_ARGUMENT	–An argument is incorrect.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_field_new, new_field, dup_field, link_field, free_field, – create and destroy forms fields

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lform -lcurses** [*library* . .]

#include <form.h>

FIELD ∗**new_field(int** *r***, int** *c***, int** *frow***, int** *fcol***, int** *nrow***, int** *ncol***);**

FIELD ∗**dup_field(FIELD** ∗*field***, int** *frow***, int** *fcol***);**

FIELD ∗**link_field(FIELD** ∗*field***, int** *frow***, int** *fcol***);**

int free_field(FIELD ∗*field***);**

MT-LEVEL

Unsafe

DESCRIPTION

new_field( ) creates a new field with *r* rows and *c* columns, starting at *frow*, *fcol*, in the subwindow of a form. *nrow* is the number of off-screen rows and *nbuf* is the number of additional working buffers. This routine returns a pointer to the new field.

dup_field( ) duplicates *field* at the specified location. All field attributes are duplicated, including the current contents of the field buffers.

link_field( ) also duplicates *field* at the specified location. However, unlike **dup_field( )**, the new field shares the field buffers with the original field. After creation, the attributes of the new field can be changed without affecting the original field.

free_field( ) frees the storage allocated for *field.*

RETURN VALUES

Routines that return pointers return **NULL** on error. **free_field( )** returns one of the following:

E_OK	–	The function returned successfully.
E_CONNECTED	–	The field is already connected to a form.
E_SYSTEM_ERROR	–	System error.
E_BAD_ARGUMENT	–	An argument is incorrect.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_field_opts, set_field_opts, field_opts_on, field_opts_off, field_opts – forms field option routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lform -lcurses** [*library* . .]

#include <form.h>

int set_field_opts(FIELD ∗*field***, OPTIONS** *opts***);**

int set_field_opts(FIELD ∗*field***, OPTIONS** *opts***);**

int field_opts_on(FIELD ∗*field***, OPTIONS** *opts***);**

int field_opts_off(FIELD ∗*field***, OPTIONS** *opts***);**

OPTIONS field_opts(FIELD ∗*field***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_field_opts( ) turns on the named options of *field* and turns off all remaining options. Options are boolean values that can be OR-ed together.

field_opts_on( ) turns on the named options; no other options are changed.

field_opts_off( ) turns off the named options; no other options are changed.

field_opts( ) returns the options set for *field*.

Field Options:

O_VISIBLE	The field is displayed.
O_ACTIVE	The field is visited during processing.
O_PUBLIC	The field contents are displayed as data is entered.
O_EDIT	The field can be edited.
O_WRAP	Words not fitting on a line are wrapped to the next line.
O_BLANK	The whole field is cleared if a character is entered in the first position.
O_AUTOSKIP	Skip to the next field when the current field becomes full.
O_NULLOK	A blank field is considered valid.
O_STATIC	The field buffers are fixed in size.
O_PASSOK	Validate field only if modified by user.

RETURN VALUES

set_field_opts, **field_opts_on** and **field_opts_off** return one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.
E_CURRENT	–	The field is the current field.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_field_userptr, set_field_userptr, field_userptr – associate application data with forms

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lform -lcurses** [*library* . .]

#include <form.h>

int set_field_userptr(FIELD ∗*field***, char** ∗*ptr***);**

char ∗**field_userptr(FIELD** ∗*field***);**

MT-LEVEL

Unsafe

DESCRIPTION

Every field has an associated user pointer that can be used to store pertinent data. **set_field_userptr( )** sets the user pointer of *field*. **field_userptr( )** returns the user pointer of *field*.

RETURN VALUES

field_userptr( ) returns **NULL** on error. **set_field_userptr( )** returns one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_field_validation, set_field_type, field_type, field_arg – forms field data type validation

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lform -lcurses** [*library* . .]

#include <form.h>

int set_field_type(FIELD ∗*field***, FIELDTYPE** ∗*type***, . . .);**

FIELDTYPE ∗**field_type(FIELD** ∗*field***);**

char ∗**field_arg(FIELD** ∗*field***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_field_type( ) associates the specified field type with *field*. Certain field types take additional arguments. **TYPE_ALNUM**, for instance, requires one, the minimum width specification for the field. The other predefined field types are: **TYPE_ALPHA**, **TYPE_ENUM**, **TYPE_INTEGER**, **TYPE_NUMERIC**, and **TYPE_REGEXP**.

field_type( ) returns a pointer to the field type of *field*. **NULL** is returned if no field type is assigned.

field_arg( ) returns a pointer to the field arguments associated with the field type of *field*. **NULL** is returned if no field type is assigned.

RETURN VALUES

field_type( ) and **field_arg( )** return **NULL** on error.

set_field_type( ) returns one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_fieldtype, new_fieldtype, free_fieldtype, set_fieldtype_arg, set_fieldtype_choice, link_fieldtype – forms fieldtype routines

SYNOPSIS

cc [*flag* ...] *file* ... –**lform -lcurses** [*library* ..]

#include <form.h>

FIELDTYPE ∗**new_fieldtype(int (**∗ **field_check) (FIELD** ∗**, char** ∗**),**
int (∗ **char_check)(int, char** ∗**));**

int free_fieldtype(FIELDTYPE ∗*fieldtype***);**

int set_fieldtype_arg(FIELDTYPE ∗*fieldtype***, char** ∗**(**∗ **mak_arg)(va_list** ∗**),**
char ∗**(**∗ **copy_arg) (char** ∗**), void (**∗ **free_arg)(char** ∗**));"**

int set_fieldtype_choice(FIELDTYPE ∗*fieldtype***, int (**∗ **next_choice)(FIELD** ∗**, char** ∗**),**
int (∗ **prev_choice)(FIELD** ∗**, char** ∗**));**

FIELDTYPE ∗**link_fieldtype(FIELDTYPE** ∗*type1***, FIELDTYPE** ∗*type2***);**

MT-LEVEL

Unsafe

DESCRIPTION

new_fieldtype( ) creates a new field type. The application programmer must write the function *field_check*, which validates the field value, and the function *char_check*, which validates each character. **free_fieldtype( )** frees the space allocated for the field type.

By associating function pointers with a field type, **set_fieldtype_arg( )** connects to the field type additional arguments necessary for a **set_field_type( )** call. Function *mak_arg* allocates a structure for the field specific parameters to **set_field_type( )** and returns a pointer to the saved data. Function *copy_arg* duplicates the structure created by *make_arg*. Function *free_arg* frees any storage allocated by *make_arg* or *copy_arg*.

The **form_driver( )** requests **REQ_NEXT_CHOICE** and **REQ_PREV_CHOICE** let the user request the next or previous value of a field type comprising an ordered set of values. **set_fieldtype_choice( )** allows the application programmer to implement these requests for the given field type. It associates with the given field type those application-defined functions that return pointers to the next or previous choice for the field.

link_fieldtype( ) returns a pointer to the field type built from the two given types. The constituent types may be any application-defined or pre-defined types.

RETURN VALUES

Routines that return pointers always return **NULL** on error. Routines that return an integer return one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.
E_BAD_ARGUMENT	–	An argument is incorrect.
E_CONNECTED	–	Type is connected to one or more fields.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_hook, set_form_init, form_init, set_form_term, form_term, set_field_init, field_init, set_field_term, field_term – assign application-specific routines for invocation by forms

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lform -lcurses** [*library* . .]

#include <form.h>

int set_form_init(FORM ∗*form***, void (**∗**func)(FORM** ∗**));**

void (∗**form_init)(FORM** ∗*form***);**

int set_form_term(FORM ∗*form***, void (**∗**func)(FORM** ∗**));**

void (∗**form_term)(FORM** ∗*form***);**

int set_field_init(FORM ∗*form***, void (**∗**func)(FORM** ∗**));**

void (∗**field_init)(FORM** ∗*form***);**

int set_field_term(FORM ∗*form***, void (**∗**func)(FORM** ∗**));**

void (∗**field_term)(FORM** ∗*form***);**

MT-LEVEL

Unsafe

DESCRIPTION

These routines allow the programmer to assign application specific routines to be executed automatically at initialization and termination points in the **forms** application. The user need not specify any application-defined initialization or termination routines at all, but they may be helpful for displaying messages or page numbers and other chores.

set_form_init( ) assigns an application-defined initialization function to be called when the *form* is posted and just after a page change. **form_init( )** returns a pointer to the initialization function, if any.

set_form_term( ) assigns an application-defined function to be called when the *form* is unposted and just before a page change. **form_term( )** returns a pointer to the function, if any.

set_field_init( ) assigns an application-defined function to be called when the *form* is posted and just after the current field changes. **field_init( )** returns a pointer to the function, if any.

set_field_term( ) assigns an application-defined function to be called when the *form* is unposted and just before the current field changes. **field_term( )** returns a pointer to the function, if any.

RETURN VALUES

Routines that return pointers always return **NULL** on error. Routines that return an integer return one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.

SEE ALSO **curses**(3X), **forms**(3X)

NOTES The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_new, new_form, free_form – create and destroy forms

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lform -lcurses** [*library* . .]

#include <form.h>

FORM ***new_form(FIELD** ****fields*****);**

int free_form(FORM ***form*****);**

MT-LEVEL

Unsafe

DESCRIPTION

new_form( ) creates a new form connected to the designated fields and returns a pointer to the form.

free_form( ) disconnects the *form* from its associated field pointer array and deallocates the space for the form.

RETURN VALUES

new_form( ) always returns **NULL** on error. **free_form( )** returns one of the following:

E_OK	–	The function returned successfully.
E_BAD_ARGUMENT	–	An argument is incorrect.
E_POSTED	–	The form is posted.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_new_page, set_new_page, new_page – forms pagination

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lform -lcurses** [*library* . .]

#include <form.h>

int set_new_page(FIELD ∗*field***, int** *bool***);**

int new_page(FIELD ∗*field***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_new_page( ) marks *field* as the beginning of a new page on the form.

new_page( ) returns a boolean value indicating whether or not *field* begins a new page of the form.

RETURN VALUES

new_page returns **TRUE** or **FALSE**.

set_new_page( ) returns one of the following:

E_OK	–	The function returned successfully.
E_CONNECTED	–	The field is already connected to a form.
E_SYSTEM_ERROR	–	System error.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_opts, set_form_opts, form_opts_on, form_opts_off – forms option routines

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lform -lcurses** [*library* . .]

#include <form.h>

int set_form_opts(FORM ∗*form***, OPTIONS** *opts***);**

int form_opts_on(FORM ∗*form***, OPTIONS** *opts***);**

int form_opts_off(FORM ∗*form***, OPTIONS** *opts***);**

OPTIONS form_opts(FORM ∗*form***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_form_opts( ) turns on the named options for *form* and turns off all remaining options. Options are boolean values which can be OR-ed together.

form_opts_on( ) turns on the named options; no other options are changed.

form_opts_off( ) turns off the named options; no other options are changed.

form_opts( ) returns the options set for *form*.

Form Options:

O_NL_OVERLOAD	Overload the **REQ_NEW_LINE** form driver request.
O_BS_OVERLOAD	Overload the **REQ_DEL_PREV** form driver request.

RETURN VALUES

set_form_opts( ), **form_opts_on( )**, and **form_opts_off( )** return one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_page, set_form_page, set_current_field, current_field, field_index – set forms current page and field

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lform -lcurses** [*library* . .]

#include <form.h>

int set_form_page(FORM ∗*form***, int** *page***);**

int form_page(FORM ∗*form***);**

int set_current_field(FORM ∗*form***, FIELD** ∗*field***);**

FIELD ∗**current_field(FORM** ∗*form***);**

int field_index(FIELD ∗*field***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_form_page( ) sets the page number of *form* to *page*. **form_page( )** returns the current page number of *form*.

set_current_field( ) sets the current field of *form* to *field*. **current_field( )** returns a pointer to the current field of *form*.

field_index( ) returns the index in the field pointer array of *field*.

RETURN VALUES

form_page( ) returns **-1** on error.

current_field( ) returns **NULL** on error.

field_index( ) returns -1 on error.

set_form_page( ) and **set_current_field( )** return one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.
E_BAD_ARGUMENT	–	An argument is incorrect.
E_BAD_STATE	–	The routine was called from an initialization or termination function.
E_INVALID_FIELD	–	The field contents are invalid.
E_REQUEST_DENIED	–	The form driver request failed.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_post, post_form, unpost_form – write or erase forms from associated subwindows

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lform -lcurses** [*library* . .]

#include <form.h>

int post_form(FORM ∗*form***);**

int unpost_form(FORM ∗*form***);**

MT-LEVEL

Unsafe

DESCRIPTION

post_form( ) writes *form* into its associated subwindow. The application programmer must use **curses** library routines to display the form on the physical screen or call **update_panels( )** if the **panels** library is being used.

unpost_form( ) erases *form* from its associated subwindow.

RETURN VALUES

These routines return one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.
E_BAD_ARGUMENT	–	An argument is incorrect.
E_POSTED	–	The form is posted.
E_NOT_POSTED	–	The form is not posted.
E_NO_ROOM	–	The form does not fit in the subwindow.
E_BAD_STATE	–	The routine was called from an initialization or termination function.
E_NOT_CONNECTED	–	The field is not connected to a form.

SEE ALSO

curses(3X), **forms**(3X), **panel_update**(3X), **panels**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_userptr, set_form_userptr – associate application data with forms

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lform -lcurses** [*library* . .]

```
#include <form.h>
int set_form_userptr(FORM *form, char *ptr);
char *form_userptr(FORM *form);
```

MT-LEVEL

Unsafe

DESCRIPTION

Every form has an associated user pointer that can be used to store pertinent data. **set_form_userptr( )** sets the user pointer of *form*. **form_userptr( )** returns the user pointer of *form*.

RETURN VALUES

form_userptr( ) returns **NULL** on error. **set_form_userptr( )** returns one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

form_win, set_form_win, set_form_sub, form_sub, scale_form – forms window and subwindow association routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lform -lcurses** [*library* . .]

#include <form.h>

int set_form_win(FORM ∗*form***, WINDOW** ∗*win***);**

WINDOW ∗**form_win(FORM** ∗*form***);**

int set_form_sub(FORM ∗*form***, WINDOW** ∗*sub***);**

WINDOW ∗**form_sub(FORM** ∗*form***);**

int scale_form(FORM ∗*form***, int** ∗*rows***, int** ∗*cols***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_form_win( ) sets the window of *form* to *win*. **form_win( )** returns a pointer to the window associated with *form*.

set_form_sub( ) sets the subwindow of *form* to *sub*. **form_sub( )** returns a pointer to the subwindow associated with *form*.

scale_form( ) returns the smallest window size necessary for the subwindow of *form*. *rows* and *cols* are pointers to the locations used to return the number of rows and columns for the form.

RETURN VALUES

Routines that return pointers always return **NULL** on error. Routines that return an integer return one of the following:

E_OK	–	The function returned successfully.
E_SYSTEM_ERROR	–	System error.
E_BAD_ARGUMENT	–	An argument is incorrect.
E_NOT_CONNECTED	–	The field is not connected to a form.
E_POSTED	–	The form is posted.

SEE ALSO

curses(3X), **forms**(3X)

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

forms – character based forms package

SYNOPSIS

#include <form.h>

MT-LEVEL

Unsafe

DESCRIPTION

The **form** library is built using the **curses** library, and any program using **forms** routines must call one of the **curses** initialization routines such as **initscr**. A program using these routines must be compiled with **–lform** and **–lcurses** on the **cc** command line.

The **forms** package gives the applications programmer a terminal-independent method of creating and customizing forms for user-interaction. The **forms** package includes: field routines, which are used to create and customize fields, link fields and assign field types; fieldtype routines, which are used to create new field types for validating fields; and form routines, which are used to create and customize forms, assign pre / post processing functions, and display and interact with forms.

Current Default Values for Field Attributes

The **forms** package establishes initial current default values for field attributes. During field initialization, each field attribute is assigned the current default value for that attribute. An application can change or retrieve a current default attribute value by calling the appropriate set or retrieve routine with a **NULL** field pointer. If an application changes a current default field attribute value, subsequent fields created using **new_field( )** will have the new default attribute value. (The attributes of previously created fields are not changed if a current default attribute value is changed.)

Routine Name Index

The following table lists each **forms** routine and the name of the manual page on which it is described.

forms Routine Name	Manual Page Name
current_field	**form_page(3X)**
data_ahead	**form_data(3X)**
data_behind	**form_data(3X)**
dup_field	**form_field_new(3X)**
dynamic_field_info	**form_field_info(3X)**
field_arg	**form_field_validation(3X)**
field_back	**form_field_attributes(3X)**
field_buffer	**form_field_buffer(3X)**
field_count	**form_field(3X)**
field_fore	**form_field_attributes(3X)**
field_index	**form_page(3X)**
field_info	**form_field_info(3X)**
field_init	**form_hook(3X)**
field_just	**form_field_just(3X)**
field_opts	**form_field_opts(3X)**
field_opts_off	**form_field_opts(3X)**

field_opts_on	**form_field_opts(3X)**
field_pad	**form_field_attributes(3X)**
field_status	**form_field_buffer(3X)**
field_term	**form_hook(3X)**
field_type	**form_field_validation(3X)**
field_userptr	**form_field_userptr(3X)**
form_driver	**form_driver(3X)**
form_fields	**form_field(3X)**
form_init	**form_hook(3X)**
form_opts	**form_opts(3X)**
form_opts_off	**form_opts(3X)**
form_opts_on	**form_opts(3X)**
form_page	**form_page(3X)**
form_sub	**form_win(3X)**
form_term	**form_hook(3X)**
form_userptr	**form_userptr(3X)**
form_win	**form_win(3X)**
free_field	**form_field_new(3X)**
free_fieldtype	**form_fieldtype(3X)**
free_form	**form_new(3X)**
link_field	**form_field_new(3X)**
link_fieldtype	**form_fieldtype(3X)**
move_field	**form_field(3X)**
new_field	**form_field_new(3X)**
new_fieldtype	**form_fieldtype(3X)**
new_form	**form_new(3X)**
new_page	**form_new_page(3X)**
pos_form_cursor	**form_cursor(3X)**
post_form	**form_post(3X)**
scale_form	**form_win(3X)**
set_current_field	**form_page(3X)**
set_field_back	**form_field_attributes(3X)**
set_field_buffer	**form_field_buffer(3X)**
set_field_fore	**form_field_attributes(3X)**
set_field_init	**form_hook(3X)**
set_field_just	**form_field_just(3X)**
set_field_opts	**form_field_opts(3X)**
set_field_pad	**form_field_attributes(3X)**
set_field_status	**form_field_buffer(3X)**
set_field_term	**form_hook(3X)**
set_field_type	**form_field_validation(3X)**
set_field_userptr	**form_field_userptr(3X)**
set_fieldtype_arg	**form_fieldtype(3X)**
set_fieldtype_choice	**form_fieldtype(3X)**
set_form_fields	**form_field(3X)**

set_form_init	**form_hook(3X)**
set_form_opts	**form_opts(3X)**
set_form_page	**form_page(3X)**
set_form_sub	**form_win(3X)**
set_form_term	**form_hook(3X)**
set_form_userptr	**form_userptr(3X)**
set_form_win	**form_win(3X)**
set_max_field	**form_field_buffer(3X)**
set_new_page	**form_new_page(3X)**
unpost_form	**form_post(3X)**

RETURN VALUES

Routines that return a pointer always return **NULL** on error. Routines that return an integer return one of the following:

E_OK	–	The function returned successfully.
E_CONNECTED	–	The field is already connected to a form.
E_SYSTEM_ERROR	–	System error.
E_BAD_ARGUMENT	–	An argument is incorrect.
E_CURRENT	–	The field is the current field.
E_POSTED	–	The form is posted.
E_NOT_POSTED	–	The form is not posted.
E_INVALID_FIELD	–	The field contents are invalid.
E_NOT_CONNECTED	–	The field is not connected to a form.
E_NO_ROOM	–	The form does not fit in the subwindow.
E_BAD_STATE	–	The routine was called from an initialization or termination function.
E_REQUEST_DENIED	–	The form driver request failed.
E_UNKNOWN_COMMAND	–	An unknown request was passed to the form driver.

SEE ALSO

curses(3X), and 3X pages whose names begin "form_" for detailed routine descriptions.

NOTES

The header **<form.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

fpgetround, fpsetround, fpgetmask, fpsetmask, fpgetsticky, fpsetsticky – IEEE floating-point environment control

SYNOPSIS

```
#include <ieeefp.h>
fp_rnd fpgetround(void);
fp_rnd fpsetround(fp_rnd rnd_dir);
fp_except fpgetmask(void);
fp_except fpsetmask(fp_except mask);
fp_except fpgetsticky(void);
fp_except fpsetsticky(fp_except sticky);
```

MT-LEVEL

MT-Safe

DESCRIPTION

There are five floating-point exceptions: divide-by-zero, overflow, underflow, imprecise (inexact) result, and invalid operation. When a floating-point exception occurs, the corresponding sticky bit is set (1), and if the mask bit is enabled (1), the trap takes place. These routines let the user change the behavior on occurrence of any of these exceptions, as well as change the rounding mode for floating-point operations.

The following floating-point exception masks are OR-ed together to form *mask*.

```
FP_X_INV    /* invalid operation exception */
FP_X_OFL    /* overflow exception */
FP_X_UFL    /* underflow exception */
FP_X_DZ     /* divide-by-zero exception */
FP_X_IMP    /* imprecise (loss of precision) */
```

The following floating-point rounding modes are passed to **fpsetround ( )** and returned by **fpgetround( )**.

```
FP_RN       /* round to nearest representative number */
FP_RP       /* round to plus infinity */
FP_RM       /* round to minus infinity */
FP_RZ       /* round to zero (truncate) */
```

The default environment is rounding mode set to nearest (**FP_RN**) and all traps disabled.

Individual bits may be examined using the constants defined in **<ieeefp.h>**.

RETURN VALUES

fpgetround( ) returns the current rounding mode.

fpsetround( ) sets the rounding mode and returns the previous rounding mode.

fpgetmask( ) returns the current exception masks.

fpsetmask( ) sets the exception masks and returns the previous setting.

fpgetsticky( ) returns the current exception sticky flags.

fpsetsticky( ) sets (clears) the exception sticky flags and returns the previous setting.

SEE ALSO

isnan(3C)

NOTES

fpsetsticky( ) modifies all sticky flags. **fpsetmask( )** changes all mask bits. **fpsetmask( )** clears the sticky bit corresponding to any exception being enabled.

C requires truncation (round to zero) for floating point to integral conversions. The current rounding mode has no effect on these conversions.

One must clear the sticky bit to recover from the trap and to proceed. If the sticky bit is not cleared before the next trap occurs, a wrong exception type may be signaled.

NAME

fread, fwrite – buffered binary input/output

SYNOPSIS

```
#include <stdio.h>
size_t fread(void *ptr, size_t size, size_t nitems, FILE *stream);
size_t fwrite(const void *ptr, size_t size, size_t nitems, FILE *stream);
```

MT-LEVEL

MT-Safe

DESCRIPTION

fread() reads into an array pointed to by *ptr* up to *nitems* items of data from *stream*, where an item of data is a sequence of bytes (not necessarily terminated by a null byte) of length *size*. **fread()** stops reading bytes if an end-of-file or error condition is encountered while reading *stream*, or if *nitems* items have been read. **fread()** increments the data pointer in *stream* to point to the byte following the last byte read if there is one. **fread()** does not change the contents of *stream*. **fread()** returns the number of items read.

fwrite() writes to the named output *stream* at most *nitems* items of data from the array pointed to by *ptr*, where an item of data is a sequence of bytes (not necessarily terminated by a null byte) of length *size*. **fwrite()** stops writing when it has written *nitems* items of data or if an error condition is encountered on *stream*. **fwrite()** does not change the contents of the array pointed to by *ptr*. **fwrite()** increments the data-pointer in *stream* by the number of bytes written. **fwrite()** returns the number of items written.

If *size* or *nitems* is 0, then **fread()** and **fwrite()** return 0 and do not effect the state of *stream*.

The **ferror()** or **feof()** routines must be used to distinguish between an error condition and end-of-file condition.

RETURN VALUES

If an error occurs, **fread()** and **fwrite()** return 0 and set the error indicator for *stream*.

SEE ALSO

read(2), **write**(2), **fclose**(3S), **fopen**(3S), **getc**(3S), **gets**(3S), **printf**(3S), **putc**(3S), **puts**(3S), **scanf**(3S), **stdio**(3S)

NAME

frexp, ldexp, logb, modf, modff, nextafter, scalb – manipulate parts of floating-point numbers

SYNOPSIS

cc [*flag* ...] *file* ... **–lm** [*library* ...]

#include <math.h>

double frexp(double *value***, int** ∗*eptr***);**

double ldexp(double *value***, int** *exp***);**

double logb(double *value***);**

double modf(double *value***, double** ∗*iptr***);**

float modff(float *value***, float** ∗*iptr***);**

double nextafter(double *value1***, double** *value2***);**

double scalb(double *value***, double** *exp***);**

MT-LEVEL

MT-Safe

DESCRIPTION

Every non-zero number can be written uniquely as $x * 2^n$, where the ''mantissa'' (fraction) x is in the range $0.5 \leq |x| < 1.0$, and the ''exponent'' n is an integer. **frexp( )** returns the mantissa of a double *value*, and stores the exponent indirectly in the location pointed to by *eptr*. If *value* is zero, both results returned by **frexp( )** are zero.

ldexp( ) and **scalb( )** return the quantity $value * 2^{exp}$. The only difference between the two is that **scalb( )** of a signaling NaN will result in the invalid operation exception being raised.

logb( ) returns the unbiased exponent of its floating-point argument as a double-precision floating-point value.

modf( ) and **modff( )** (single-precision version) return the signed fractional part of *value* and store the integral part indirectly in the location pointed to by *iptr*.

nextafter( ) returns the next representable double-precision floating-point value following *value1* in the direction of *value2*. Thus, if *value2* is less than *value1*, **nextafter( )** returns the largest representable floating-point number less than *value1*.

RETURN VALUES

If **ldexp( )** would cause overflow, **±HUGE** (defined in **<math.h>**) is returned (according to the sign of *value*), and **errno** is set to **ERANGE**. If **ldexp( )** would cause underflow, zero is returned and **errno** is set to **ERANGE**. If the input *value* to **ldexp( )** is NaN or infinity, that input is returned and **errno** is set to **EDOM**. The same error conditions apply to **scalb( )** except that a signaling NaN as input will result in the raising of the invalid operation exception.

logb( ) of NaN returns that NaN, **logb( )** of infinity returns positive infinity, and **logb( )** of zero returns negative infinity and results in the raising of the divide by zero exception. In each of these conditions **errno** is set to **EDOM**.

If input *value1* to **nextafter( )** is positive or negative infinity, that input is returned and **errno** is set to **EDOM**. The overflow and inexact exceptions are signalled when input *value1* is finite, but **nextafter(***value1***,** *value2***)** is not. The underflow and inexact exceptions are signalled when **nextafter(***value1***,** *value2***)** lies strictly between $+2^{-1022}$ and -2^{-1022}. In both cases **errno** is set to **ERANGE**.

NAME

fseek, rewind, ftell – reposition a file pointer in a stream

SYNOPSIS

```
#include <stdio.h>
int fseek(FILE *stream, long offset, int ptrname);
void rewind(FILE *stream);
long ftell(FILE *stream);
```

MT-LEVEL

MT-Safe

DESCRIPTION

fseek( ) sets the position of the next input or output operation on the *stream* (see **intro**(3)). The new position is at the signed distance *offset* bytes from the beginning, from the current position, or from the end of the file, according to a *ptrname* value of **SEEK_SET**, **SEEK_CUR**, or **SEEK_END** (defined in **<stdio.h>**) as follows:

SEEK_SET set position equal to *offset* bytes.

SEEK_CUR set position to current location plus *offset*.

SEEK_END set position to **EOF** plus *offset*.

fseek( ) allows the file position indicator to be set beyond the end of the existing data in the file. If data is later written at this point, subsequent reads of data in the gap will return zero until data is actually written into the gap. **fseek( )**, by itself, does not extend the size of the file.

rewind(stream) is equivalent to:

```
(void) fseek (stream, 0L, SEEK_SET);
```

except that **rewind( )** also clears the error indicator on *stream*.

fseek( ) and **rewind( )** clear the **EOF** indicator and undo any effects of **ungetc( )** on *stream*. After **fseek( )** or **rewind( )**, the next operation on a file opened for update may be either input or output.

If *stream* is writable and buffered data has not been written to the underlying file, **fseek( )** and **rewind( )** cause the unwritten data to be written to the file.

ftell( ) returns the offset of the current byte relative to the beginning of the file associated with the named *stream*.

RETURN VALUES

fseek( ) returns –1 for improper seeks, otherwise zero. An improper seek can be, for example, an **fseek( )** done on a file that has not been opened via **fopen( )**; in particular, **fseek( )** may not be used on a terminal or on a file opened via **popen( )**. After a stream is closed, no further operations are defined on that stream.

SEE ALSO

lseek(2), **write**(2), **intro**(3), **fopen**(3S), **popen**(3S), **stdio**(3S), **ungetc**(3S)

NOTES

Although on the UNIX system an offset returned by **ftell()** is measured in bytes, and it is permissible to seek to positions relative to that offset, portability to non-UNIX systems requires that an offset be used by **fseek()** directly. Arithmetic may not meaningfully be performed on such an offset, which is not necessarily measured in bytes.

NAME

fsetpos, fgetpos – reposition a file pointer in a stream

SYNOPSIS

```
#include <stdio.h>
int fsetpos(FILE *stream, const fpos_t *pos);
int fgetpos(FILE *stream, fpos_t *pos);
```

MT-LEVEL

MT-Safe

DESCRIPTION

fsetpos( ) sets the position of the next input or output operation on the *stream* according to the value of the object pointed to by *pos*. The object pointed to by *pos* must be a value returned by an earlier call to **fgetpos( )** on the same stream.

fsetpos( ) clears the end-of-file indicator for the stream and undoes any effects of the **ungetc( )** function on the same stream. After **fsetpos( )**, the next operation on a file opened for update may be either input or output.

fgetpos( ) stores the current value of the file position indicator for *stream* in the object pointed to by *pos*. The value stored contains information usable by **fsetpos( )** for repositioning the stream to its position at the time of the call to **fgetpos( )**.

RETURN VALUES

If successful, both **fsetpos( )** and **fgetpos( )** return zero. Otherwise, they both return nonzero.

SEE ALSO

lseek(2), **fseek**(3S), **ungetc**(3S)

NAME

fsync – synchronize a file's in-memory state with that on the physical medium

SYNOPSIS

```
#include <unistd.h>
int fsync(int fildes);
```

MT-LEVEL

Async-Signal-Safe

DESCRIPTION

fsync( ) moves all modified data and attributes of the file descriptor *fildes* to a storage device. When **fsync( )** returns, all in-memory modified copies of buffers associated with *fildes* have been written to the physical medium. **fsync( )** is different from **sync( )**, which schedules disk I/O for all files but returns before the I/O completes. **fsync( )** forces all outstanding data operations to synchronized file integrity completion (see **fcntl**(5) definition of **O_SYNC**.)

fsync( ) should be used by programs that require that a file be in a known state. For example, a program that contains a simple transaction facility might use **fsync( )** to ensure that all changes to a file or files caused by a given transaction were recorded on a storage medium.

RETURN VALUES

Upon successful completion, a value of **0** is returned. Otherwise, a value of **-1** is returned and **errno** is set to indicate the error.

ERRORS

fsync( ) fails if one or more of the following are true:

EBADF	*fildes* is not a valid file descriptor open for writing.
EINTR	A signal was caught during execution of the **fsync( )** function.
EIO	An I/O error occurred while reading from or writing to the file system.
ENOLINK	*fildes* is on a remote machine and the link to that machine is no longer active.

SEE ALSO

sync(2), **fdatasync**(3R), **fcntl**(5)

NOTES

The way the data reach the physical medium depends on both implementation and hardware. **fsync( )** returns when the device driver tells it that the write has taken place.

NAME

ftime – get date and time

SYNOPSIS

```
#include <sys/types.h>
#include <sys/timeb.h>

int ftime(struct timeb *tp);
```

DESCRIPTION

The **ftime( )** entry fills in a structure pointed to by its argument. The structure is defined in **<sys/timeb.h>** and contains the following members:

```
time_t          time;
unsigned short  millitm;
short           timezone;
short           dstflag;
```

The structure contains the time since the epoch in seconds, up to 1000 milliseconds of more-precise interval, the local time zone, and a flag that, if nonzero, indicates that Daylight Saving time applies locally during the appropriate part of the year.

The contents of the **timezone** and **dstflag** members of **tp** after a call to **ftime( )** are unspecified.

SEE ALSO

date(1), **time**(2), **gettimeofday**(3C), **ctime**(3C), **timezone**(4)

NAME

ftw, nftw – walk a file tree

SYNOPSIS

```
#include <ftw.h>

int ftw(const char *path, int (*fn) (const char *, const struct stat *, int), int depth);

int nftw(const char *path, int (*fn) (const char *, const struct stat *, int, struct FTW*),
        int depth, int flags);
```

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

ftw( ) recursively descends the directory hierarchy rooted in *path*. For each object in the hierarchy, **ftw( )** calls the user-defined function *fn*, passing it a pointer to a null-terminated character string containing the name of the object, a pointer to a **stat** structure (see **stat**(2)) containing information about the object, and an integer. Possible values of the integer, defined in the **<ftw.h>** header, are:

FTW_F The object is a file.

FTW_D The object is a directory.

FTW_DNR The object is a directory that cannot be read. Descendants of the directory will not be processed.

FTW_NS **stat** failed on the object because of lack of appropriate permission or the object is a symbolic link that points to a non-existent file. The stat buffer passed to *fn* is undefined.

ftw( ) visits a directory before visiting any of its descendants.

The tree traversal continues until the tree is exhausted, an invocation of *fn* returns a nonzero value, or some error is detected within **ftw( )** (such as an I/O error). If the tree is exhausted, **ftw( )** returns zero. If *fn* returns a nonzero value, **ftw( )** stops its tree traversal and returns whatever value was returned by *fn*.

The function **nftw( )** is similar to **ftw( )** except that it takes an additional argument, *flags*. The *flags* field is used to specify:

FTW_PHYS Physical walk, does not follow symbolic links. Otherwise, **nftw( )** will follow links but will not walk down any path that crosses itself.

FTW_MOUNT The walk will not cross a mount point.

FTW_DEPTH All subdirectories will be visited before the directory itself.

FTW_CHDIR The walk will change to each directory before reading it.

The function **nftw( )** calls *fn* with four arguments at each file and directory. The first argument is the pathname of the object, the second is a pointer to the **stat** buffer, the third is an integer giving additional information, and the fourth is a **struct FTW** that contains the following members:

```
int base;
int level;
```

base is the offset into the pathname of the base name of the object. **level** indicates the depth relative to the rest of the walk, where the root level is zero.

The values of the third argument are as follows:

FTW_F The object is a file.

FTW_D The object is a directory.

FTW_DP The object is a directory and subdirectories have been visited.

FTW_SL The object is a symbolic link.

FTW_SLN The object is a symbolic link that points to a non-existent file.

FTW_DNR The object is a directory that cannot be read. *fn* will not be called for any of its descendants.

FTW_NS **stat** failed on the object because of lack of appropriate permission. The stat buffer passed to *fn* is undefined. **stat** failure other than lack of appropriate permission. **EACCES** is considered an error and **nftw( )** will return –1.

Both **ftw( )** and **nftw( )** use one file descriptor for each level in the tree. The *depth* argument limits the number of file descriptors so used. If *depth* is zero or negative, the effect is the same as if it were 1. *depth* must not be greater than the number of file descriptors currently available for use. **ftw( )** will run faster if *depth* is at least as large as the number of levels in the tree. When **ftw( )** and **nftw( )** return, they close any file descriptors they have opened; they do not close any file descriptors that may have been opened by *fn*.

RETURN VALUES

If successful, **ftw( )** and **nftw( )** return **0**. If either function detects an error other than **EACCES**, it returns –1, and sets the error type in **errno**.

SEE ALSO

stat(2), **longjmp**(3C), **malloc**(3C)

NOTES

Because **ftw( )** is recursive, it is possible for it to terminate with a memory fault when applied to very deep file structures.

ftw( ) uses **malloc**(3C) to allocate dynamic storage during its operation. If **ftw( )** is forcibly terminated, such as by **longjmp**(3C) being executed by *fn* or an interrupt routine, **ftw( )** will not have a chance to free that storage, so it will remain permanently allocated. A safe way to handle interrupts is to store the fact that an interrupt has occurred, and arrange to have *fn* return a nonzero value at its next invocation.

ftw( ) is safe in multi-thread applications. **nftw( )** is safe in multi-thread applications when the **FTW_CHDIR** flag is not set.

NAME

getacinfo, getacdir, getacflg, getacmin, getacna, setac, endac – get audit control file information

SYNOPSIS

cc [*flag* ...] *file* ... **–lbsm –lsocket –lnsl –lintl** [*library* ...]

#include <bsm/libbsm.h>

int getacdir(char **dir,* **int** *len***);**

int getacmin(int **min_val***);**

int getacflg(char **auditstring,* **int** *len***);**

int getacna(char **auditstring,* **int** *len***);**

void setac(void);

void endac(void);

MT-LEVEL

Safe.

AVAILABILITY

The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See **bsmconv**(1M) for more information.

DESCRIPTION

When first called, **getacdir( )** provides information about the first audit directory in the **audit_control** file; thereafter, it returns the next directory in the file. Successive calls list all the directories listed in **audit_control**(4) The parameter *len* specifies the length of the buffer *dir*. On return, *dir* points to the directory entry.

getacmin( ) reads the minimum value from the **audit_control** file and returns the value in **min_val**. The minimum value specifies how full the file system to which the audit files are being written can get before the script **audit_warn**(1M) is invoked.

getacflg( ) reads the system audit value from the **audit_control** file and returns the value in *auditstring*. The parameter *len* specifies the length of the buffer *auditstring*.

getacna( ) reads the system audit value for non-attributable audit events from the **audit_control** file and returns the value in *auditstring*. The parameter *len* specifies the length of the buffer *auditstring*. Non-attributable events are events that cannot be attributed to an individual user. **inetd**(1M) and several other daemons record non-attributable events.

Calling *setac* rewinds the **audit_control** file to allow repeated searches.

Calling *endac* closes the **audit_control** file when processing is complete.

FILES

/etc/security/audit_control contains default parameters read by the audit daemon, **auditd**(1M)

RETURN VALUES

getacdir( ), **getacflg( )**, **getacna( )** and **getacmin( )** return:

0 on success.

–2 on failure and set **errno** to indicate the error.

getacmin( ) and **getacflg( )** return:

1 on EOF.

getacdir( ) returns:

–1 on EOF.

2 if the directory search had to start from the beginning because one of the other functions was called between calls to **getacdir( )**.

These functions return:

–3 if the directory entry format in the **audit_control** file is incorrect.

getacdir( ), **getacflg( )** and **getacna( )** return:

–3 if the input buffer is too short to accommodate the record.

SEE ALSO

audit_warn(1M), **bsmconv**(1M), **inetd**(1M), **audit_control**(4)

NAME

getauclassnam, getauclassent, setauclass, endauclass, getauclassnam_r, getauclassent_r – get audit_class entry

SYNOPSIS

cc [*flag* ...] *file* ... **–lbsm –lsocket –lnsl –lintl** [*library* ...]

#include <sys/param.h>
#include <bsm/libbsm.h>

struct au_class_ent ∗**getauclassnam(const char** ∗*name*);

struct au_class_ent ∗**getauclassnam_r(au_class_ent_t** ∗ *class_int*, **const char** ∗*name*);

struct au_class_ent ∗**getauclassent(void);**

struct au_class_ent ∗**getauclassent_r(au_class_ent_t** ∗ *class_int*);

void setauclass(void);

void endauclass(void);

AVAILABILITY

The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See **bsmconv**(1M) for more information.

MT-LEVEL

MT-Safe with exceptions.

All of the functions described in this man-page are MT-Safe except **getauclassent( )** and **getauclassnam( ).** The two functions, **getauclassent_r( )** and **getauclassnam_r( )** have the same functionality as the unsafe functions, but have a slightly different function call interface in order to make them MT-Safe.

DESCRIPTION

getauclassent( ) and **getauclassnam( )** each return an audit_class entry.

getauclassnam( ) searches for an audit_class entry with a given class name *name.*

getauclassent( ) enumerates audit_class entries: successive calls to **getauclassent( )** will return either successive audit_class entries or NULL.

setauclass( ) ''rewinds'' to the beginning of the enumeration of audit_class entries. Calls to **getauclassnam( )** may leave the enumeration in an indeterminate state, so **setauclass( )** should be called before the first **getauclassent( )**.

endauclass( ) may be called to indicate that audit_class processing is complete; the system may then close any open audit_class file, deallocate storage, and so forth.

getauclassent_r( ) and **getauclassnam_r( )** both return a pointer to an audit_class entry as do their similarly named counterparts. They each take an additional argument, a pointer to pre-allocated space for an **au_class_ent_t,** which is returned if the call is successful. To assure there is enough space for the information returned, the applications programmer should be sure to allocate **AU_CLASS_NAME_MAX** and **AU_CLASS_DESC_MAX** bytes for the ac_name and ac_desc elements of the **au_class_ent_t** data structure.

The internal representation of an audit_user entry is an **au_class_ent** structure defined in **<bsm/libbsm.h>** with the following members:

```
char        *ac_name;
au_class_t  ac_class;
char        *ac_desc;
```

RETURN VALUES

getauclassnam( ) and **getauclassnam_r( )** return a pointer to a **struct au_class_ent** if they successfully locate the requested entry; otherwise they return NULL.

getauclassent( ) and **getauclassent_r( )** return a pointer to a **struct au_class_ent** if they successfully enumerate an entry; otherwise they return NULL, indicating the end of the enumeration.

FILES

/etc/security/audit_class	Maps audit class numbers to audit class names

SEE ALSO

bsmconv(1M), **audit_class**(4), **audit_event**(4)

NOTES

All information is contained in a static area, so it must be copied if it is to be saved.

NAME

getauditflags, getauditflagsbin, getauditflagschar – convert audit flag specifications

SYNOPSIS

cc [*flag* ...] *file* ... **–lbsm –lsocket –lnsl –lintl** [*library* ...]

```
#include <sys/param.h>
#include <bsm/libbsm.h>

int getauditflagsbin(char *auditstring, au_mask_t *masks);

int getauditflagschar(char *auditstring, au_mask_t *masks, int verbose);
```

MT-LEVEL

MT-Safe.

AVAILABILITY

The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See **bsmconv**(1M) for more information.

DESCRIPTION

getauditflagsbin( ) converts the character representation of audit values pointed to by *auditstring* into **au_mask_t** fields pointed to by *masks*. These fields indicate which events are to be audited when they succeed and which are to be audited when they fail. The character string syntax is described in **audit_control**(4).

getauditflagschar( ) converts the **au_mask_t** fields pointed to by *masks* into a string pointed to by *auditstring*. If *verbose* is zero, the short (2-character) flag names are used. If *verbose* is non-zero, the long flag names are used. *auditstring* should be large enough to contain the ASCII representation of the events.

auditstring contains a series of event names, each one identifying a single audit class, separated by commas. The **au_mask_t** fields pointed to by *masks* correspond to binary values defined in **<bsm/audit.h>**, which is read by **<bsm/libbsm.h>**.

RETURN VALUES

getauditflagsbin( ) and **getauditflagschar( )**: **–1** is returned on error and **0** on success.

SEE ALSO

bsmconv(1M), **audit.log**(4), **audit_control**(4)

BUGS

This is not a very extensible interface.

NAME

getauevent, getauevnam, getauevnum, getauevnonam, setauevent, endauevent, getauevent_r, getauevnam_r, getauevnum_r – get audit_user entry

SYNOPSIS

cc [*flag* ...] *file* ... **–lbsm –lsocket –lnsl –lintl** [*library* ...]

```
#include <sys/param.h>
#include <bsm/libbsm.h>

struct au_event_ent *getauevent( void);

struct au_event_ent *getauevnam( char *name);

struct au_event_ent *getauevnum( au_event_t event_number);

au_event_t *getauevnonam( char *event_name);

void setauevent( void);

void endauevent( void);

struct au_event_ent *getauevent_r( au_event_ent_t *e, void);

struct au_event_ent *getauevnam_r( au_event_ent_t *e, char *name);

struct au_event_ent *getauevnum_r( au_event_ent_t *e, au_event_t event_number);
```

MT-LEVEL

MT-Safe with exceptions.

The functions **getauevent**(), **getauevnam**(), and **getauevnum**() are not MT-Safe; but, there are equivalent functions: **getauevent_r**(), **getauevnam_r**(), and **getauevnum_r**() - all of which provide the same functionality and a MT-Safe function call interface.

AVAILABILITY

The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See **bsmconv**(1M) for more information.

DESCRIPTION

getauevent( ), **getauevnam( ) , getauevnum( ) , getauevent( ) , getauevnam( ) ,** and **getauevnum( ) ,** each return a pointer to an **audit_event** structure.

getauevent( ) and **getauevent_r( )** enumerate audit_event entries: successive calls to these functions will return either successive audit_event entries or NULL.

getauevnam( ) and **getauevnam_r( )** search for an audit_event entry with a given event name *name.*

getauevnum( ) and **getauevnum_r( )** search for an audit_event entry with a given event number *number*.

getauevnonum( ) searches for an audit_event entry with a given event name *name* and returns the corresponding event number.

setauevent( ) ''rewinds'' to the beginning of the enumeration of audit_event entries. Calls to **getauevnam( )**, **getauevnum( ) , getauevnonum( ) , getauevnam_r( ) ,** or **getauevnum_r( )** may leave the enumeration in an indeterminate state; so, **setauevent( )** should be called before the first **getauevent( )** or **getauevent_r( ) .**

endauevent() may be called to indicate that **audit_event** processing is complete; the system may then close any open **audit_event file**, deallocate storage, and so forth.

The three functions **getauevent_r()**, **getauevnam_r()**, and **getauevnum_r()**, each take an argument *e* which is a pointer to an **au_event_ent_t.** This pointer is returned on a successful function call. To assure there is enough space for the information returned, the applications programmer should be sure to allocate **AU_EVENT_NAME_MAX** and **AU_EVENT_DESC_MAX** bytes for the **ae_name** and **ac_desc** elements of the **au_event_ent_t** data structure.

The internal representation of an audit_event entry is an **struct au_event_ent** structure defined in **<bsm/libbsm.h>** with the following members:

```
au_event_t   ae_number;
char         *ae_name;
char         *ae_desc;
au_class_t   ae_class;
```

RETURN VALUES

getauevent(), **getauevnam()** and **getauevnum()** return a pointer to a **struct au_event_ent** if it successfully locates the requested entry; otherwise it returns **NULL**.

getauevnonam() returns an event number of type **au_event_t** if it successfully enumerates an entry; otherwise it returns **NULL**, indicating it could not find the requested event name.

FILES

/etc/security/audit_event	Maps audit event numbers to audit event names
/etc/passwd	Stores user-id to username mappings

SEE ALSO

bsmconv(1M), **getpwnam**(3C), **getauclassent**(3), **audit_class**(4), **audit_event**(4), **passwd**(4)

NOTES

All information for the functions **getauevent ( ) , getauevnam ( ) ,** and **getauevnum ( )** is contained in a static area, so it must be copied if it is to be saved.

NAME

getauusernam, getauuserent, setauuser, endauuser – get audit_user entry

SYNOPSIS

cc [*flag* ...] *file* ... **–lbsm –lsocket –lnsl –lintl** [*library* ...]

#include <sys/param.h>
#include <bsm/libbsm.h>

struct au_user_ent ∗**getauusernam(const char** ∗*name*);

struct au_user_ent ∗**getauuserent(void);**

void setauuser(void);

void endauuser(void);

struct au_user_ent ∗**getauusernam_r(au_user_ent** ∗ **u, const char** ∗*name*);

struct au_user_ent ∗**getauuserent_r(au_user_ent** ∗**u);**

AVAILABILITY

The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See **bsmconv**(1M) for more information.

MT-LEVEL

MT-Safe with exceptions.

The functions **getauusernam( )** and **getauuserent( )** are not MT-safe. However, the functions **getauusernam_r( )** and **getauuserent_r( ) provide the same functionality with** MT-Safe interfaces.

DESCRIPTION

getauuserent( ) , getauusernam( ) , getauuserent_r( ) , and **getauusernam_r( )** each return an audit_user entry.

getauusernam( ) and **getauusernam_r( )** searche for an audit_user entry with a given login name *name.*

getauuserent( ) and **getauuserent_r( )** enumerate audit_user entries: successive calls to these functions will return either successive audit_user entries or **NULL**.

setauuser( ) ‘‘rewinds’’ to the beginning of the enumeration of audit_user entries. Calls to **getauusernam( )** and **getauusernam_r( )** may leave the enumeration in an indeterminate state, so **setauuser( )** should be called before the first **getauuserent( )** or **getauuserent_r( ) .**

endauuser( ) may be called to indicate that audit_user processing is complete; the system may then close any open audit_user file, deallocate storage, and so forth.

The two functions **getauuserent_r( )** and **getauusernam_r( )** both take an argument *u* , which is a pointer to an au_user_ent. This is the pointer that is returned on successful function calls.

The internal representation of an audit_user entry is an **au_user_ent** structure defined in **<bsm/libbsm.h>** with the following members:

```
char        *au_name;
au_mask_t   au_always;
au_mask_t   au_never;
```

RETURN VALUES

getauusernam( ) returns a pointer to a **struct au_user_ent** if it successfully locates the requested entry; otherwise it returns **NULL**.

getauuserent( ) returns a pointer to a **struct au_user_ent** if it successfully enumerates an entry; otherwise it returns **NULL**. indicating the end of the enumeration.

FILES

/etc/security/audit_user	Stores per-user audit event mask
/etc/passwd	Stores user-id to username mappings

SEE ALSO

bsmconv(1M), **getpwnam**(3C), **audit_user**(4), **passwd**(4)

NOTES

All information for the functions **getauuserent( )** and **getauusernam( )** is contained in a static area, so it must be copied if it is to be saved.

NAME

getc, getc_unlocked, getchar, getchar_unlocked, fgetc, getw – get character or word from a stream

SYNOPSIS

```
#include <stdio.h>
int getc(FILE *stream);
int getc_unlocked(FILE *stream);
int getchar(void);
int getchar_unlocked(void);
int fgetc(FILE *stream);
int getw(FILE *stream);
```

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

getc( ) returns the next character (that is, byte) from the named input *stream* (see **intro**(3)) as an **unsigned char** converted to an **int**. It also moves the file pointer, if defined, ahead one character in *stream*. **getchar( )** is defined as **getc(stdin)**. **getc( )** and **getchar( )** are macros.

getc_unlocked( ) and **getchar_unlocked( )** are respectively variants of **getc( )** and **getchar( )** that do not lock the stream. It is the caller's responsibility to acquire the stream lock before calling these functions and releasing the lock afterwards; see **flockfile**(3S) and **stdio**(3S).

fgetc( ) behaves like **getc( )**, but is a function rather than a macro. **fgetc( )** runs more slowly than **getc( ) ,** but it takes less space per invocation and its name can be passed as an argument to a function.

getw( ) returns the next word (that is, integer) from the named input *stream*. **getw( )** increments the associated file pointer, if defined, to point to the next word. The size of a word is the size of an integer and varies from machine to machine. **getw( )** assumes no special alignment in the file.

RETURN VALUES

These functions return the constant **EOF** at end-of-file or upon an error and set the **EOF** or error indicator of *stream*, respectively. Because is a valid integer, **ferror( )** should be used to detect **getw( )** errors.

SEE ALSO

intro(3), **fclose**(3S), **ferror**(3S), **flockfile**(3S), **fopen**(3S), **fread**(3S), **gets**(3S), **putc**(3S), **scanf**(3S), **stdio**(3S), **ungetc**(3S)

NOTES

If the integer value returned by **getc( )**, **getchar( )**, or **fgetc( )** is stored into a character variable and then compared against the integer constant **EOF**, the comparison may never succeed, because sign-extension of a character on widening to integer is implementation dependent.

The macro version of **getc()** evaluates a *stream* argument more than once and may treat side effects incorrectly. In particular, **getc(*f++)** does not work sensibly. Use **fgetc()** instead.

Because of possible differences in word length and byte ordering, files written using **putw()** are implementation dependent, and may not be read using **getw()** on a different processor.

Functions exist for all the above-defined macros. To get the function form, the macro name must be undefined (for example, **#undef getc**).

fgetc(), getc(), getchar(), getw(), and **ungetc()** are MT-Safe in multi-thread applications. **getc_unlocked()** and **getchar_unlocked()** are unsafe in multi-thread applications.

NAME

getcwd – get pathname of current working directory

SYNOPSIS

```
#include <unistd.h>
extern char *getcwd(char *buf, size_t size);
```

MT-LEVEL

Safe

DESCRIPTION

getcwd() returns a pointer to the current directory pathname. The value of *size* must be at least one greater than the length of the pathname to be returned.

If *buf* is not **NULL**, the pathname will be stored in the space pointed to by *buf*.

If *buf* is a **NULL** pointer, **getcwd()** will obtain *size* bytes of space using **malloc**(3C). In this case, the pointer returned by **getcwd()** may be used as the argument in a subsequent call to **free()**.

RETURN VALUES

getcwd() returns **NULL** with **errno** set if *size* is not large enough, or if an error occurs in a lower-level function.

ERRORS

getcwd() will fail if one or more of the following are true:

EACCES A parent directory cannot be read to get its name.

EINVAL *size* is equal to 0.

ERANGE *size* is greater than 0 and less than the length of the pathname plus 1.

EXAMPLE

Here is a program that prints the current working directory.

```
#include <unistd.h>
#include <stdio.h>

main()
{
        char *cwd;
        if ((cwd = getcwd(NULL, 64)) == NULL) {
                perror("pwd");
                exit(2);
        }
        (void)printf("%s\n", cwd);
        return(0);
}
```

SEE ALSO

chdir(2), **malloc**(3C)

NOTES

Using **chdir**(2) in conjunction with **getcwd** can give unpredictable results.

NAME

getdate – convert user format date and time

SYNOPSIS

```
#include <time.h>
struct tm *getdate(const char *string);
extern int getdate_err;
```

MT-LEVEL

MT-Safe

DESCRIPTION

getdate( ) converts user-definable date and / or time specifications pointed to by *string* into a **tm** structure. The **tm** structure declaration is in the **<time.h>** header file.

User-supplied templates are used to parse and interpret the input string. The templates are text files created by the user and identified via the environment variable **DATEMSK**. Each line in the template represents an acceptable date and / or time specification using conversion specifications similar to those used by **strftime**(3C) and **strptime**(3C). The first line in the template that matches the input specification is used for interpretation and conversion into the internal time format. If successful, the function **getdate( )** returns a pointer to a **tm** structure; otherwise, it returns **NULL** and sets the global variable **getdate_err** to indicate the error.

The following conversion specifications are supported:

%%	same as **%**
%a	locale's abbreviated weekday name
%A	locale's full weekday name
%b	locale's abbreviated month name
%B	locale's full month name
%c	locale's appropriate date and time representation
%C	century number [0,99]; leading zero is permitted but not required
%d	day of month [01,31]; leading zero is permitted but not required
%D	date as **%m** / **%d** / **%y**
%e	same as **%d**
%h	locale's abbreviated month name
%H	hour (24-hour clock) [0,23]; leading zero is permitted but not required
%I	hour (12-hour clock) [1,12]; leading zero is permitted but not required
%j	day number of the year [1,366]; leading zeros are permitted but not required
%m	month number [1,12]; leading zero is permitted but not required
%M	minute [0,59]; leading zero is permitted but not required
%n	any white space
%p	locale's equivalent of either a.m. or p.m.
%r	appropriate time representation in the 12-hour clock format with **%p**
%R	time as **%H**:**%M**
%S	seconds [0,61]; leading zero is permitted but not required
%t	any white space
%T	time as **%H**:**%M**:**%S**

%U week number of the year as a decimal number [0,53], with Sunday as the first day of the week; leading zero is permitted but not required
%w weekday as a decimal number [0,6], with 0 representing Sunday
%W week number of the year as a decimal number [0,53], with Monday as the first day of the week; leading zero is permitted but not required
%x locale's appropriate date representation
%X locale's appropriate time representation
%y year within the century [0,99]; leading zero is permitted but not required
%Y year, including the century (for example, 1993)
%Z time zone name or no characters if no time zone exists

Modified Conversion Specifications

Some conversion specifications can be modified by the **E** and **O** modifier characters to indicate that an alternative format or specification should be used rather than the one normally used by the unmodified specification. If the alternative format or specification does not exist in the current locale, the behaviour will be as if the unmodified conversion specification were used.

%Ec locale's alternative appropriate date and time representation
%EC name of the base year (period) in the locale's alternative representation
%Ex locale's alternative date representation
%EX locale's alternative time representation
%Ey offset from **%EC** (year only) in the locale's alternative representation
%EY full alternative year representation
%Od day of the month using the locale's alternative numeric symbols; leading zeros are permitted but not required
%Oe same as **%Od**
%OH hour (24-hour clock) using the locale's alternative numeric symbols
%OI hour (12-hour clock) using the locale's alternative numeric symbols
%Om month using the locale's alternative numeric symbols
%OM minutes using the locale's alternative numeric symbols
%OS seconds using the locale's alternative numeric symbols
%OU week number of the year (Sunday as the first day of the week) using the locale's alternative numeric symbols
%Ow number of the weekday (Sunday=0) using the locale's alternative numeric symbols
%OW week number of the year (Monday as the first day of the week) using the locale's alternative numeric symbols
%Oy year (offset from **%C**) in the locale's alternative representation and using the locale's alternative numeric symbols

Internal Format Conversion

The following rules are applied for converting the input specification into the internal format:

If only the weekday is given, today is assumed if the given day is equal to the current day and next week if it is less.

If only the month is given, the current month is assumed if the given month is equal to the current month and next year if it is less and no year is given. (The first day of month is assumed if no day is given.)

If no hour, minute, and second are given, the current hour, minute, and second are assumed.

If no date is given, today is assumed if the given hour is greater than the current hour and tomorrow is assumed if it is less.

General Specifications

A conversion specification that is an ordinary character is executed by scanning the next character from the buffer. If the character scanned from the buffer differs from the one comprising the conversion specification, the specification fails, and the differing and subsequent characters remain unscanned.

A series of conversion specifications composed of **%n**, **%t**, white space characters, or any combination is executed by scanning up to the first character that is not white space (which remains unscanned), or until no more characters can be scanned.

Any other conversion specification is executed by scanning characters until a character matching the next conversion specification is scanned, or until no more characters can be scanned. These characters, except the one matching the next conversion specification, are then compared to the locale values associated with the conversion specifier. If a match is found, values for the appropriate *tm* structure members are set to values corresponding to the locale information. If no match is found, **getdate( )** fails and no more characters are scanned.

The month names, weekday names, era names, and alternative numeric symbols can consist of any combination of upper and lower case letters. The user can request that the input date or time specification be in a specific language by setting the **LC_TIME** category using **setlocale**(3C).

RETURN VALUES

On failure **getdate( )** returns **NULL** and sets the variable **getdate_err** to indicate the error.

The following is a complete list of the **getdate_err** settings and their meanings.

1	The **DATEMSK** environment variable is null or undefined.
2	The template file cannot be opened for reading.
3	Failed to get file status information.
4	The template file is not a regular file.
5	An error is encountered while reading the template file.
6	**malloc( )** failed (not enough memory is available).
7	There is no line in the template that matches the input.
8	The input specification is invalid (for example, **February 31**).

EXAMPLES

The following example shows the possible contents of a template:

```
%m
%A %B %d %Y, %H:%M:%S
%A
%B
%m/%d/%y %I %p
%d,%m,%Y %H:%M
at %A the %dst of %B in %Y
run job at %I %p,%B %dnd
%A den %d. %B %Y %H.%M Uhr
```

The following are examples of valid input specifications for the above template:

```
getdate("10/1/87 4 PM")
getdate("Friday")
getdate("Friday September 19 1987, 10:30:30")
getdate("24,9,1986 10:30")
getdate("at monday the 1st of december in 1986")
getdate("run job at 3 PM, december 2nd")
```

If the **LANG** environment variable is set to **de** (German), the following is valid:

```
getdate("freitag den 10. oktober 1986 10.30 Uhr")
```

Local time and date specification are also supported. The following examples show how local date and time specification can be defined in the template.

Invocation	Line in Template
getdate("11/27/86")	**%m/%d/%y**
getdate("27.11.86")	**%d.%m.%y**
getdate("86-11-27")	**%y-%m-%d**
getdate("Friday 12:00:00")	**%A %H:%M:%S**

The following examples illustrate the Internal Format Conversion rules. Assume that the current date is Mon Sep 22 12:19:47 EDT 1986 and the **LANG** environment variable is not set.

Input	Line in Template	Date
Mon	**%a**	Mon Sep 22 12:19:48 EDT 1986
Sun	**%a**	Sun Sep 28 12:19:49 EDT 1986
Fri	**%a**	Fri Sep 26 12:19:49 EDT 1986
September	**%B**	Mon Sep 1 12:19:49 EDT 1986
January	**%B**	Thu Jan 1 12:19:49 EST 1987
December	**%B**	Mon Dec 1 12:19:49 EST 1986
Sep Mon	**%b %a**	Mon Sep 1 12:19:50 EDT 1986
Jan Fri	**%b %a**	Fri Jan 2 12:19:50 EST 1987
Dec Mon	**%b %a**	Mon Dec 1 12:19:50 EST 1986
Jan Wed 1989	**%b %a %Y**	Wed Jan 4 12:19:51 EST 1989
Fri 9	**%a %H**	Fri Sep 26 09:00:00 EDT 1986
Feb 10:30	**%b %H:%S**	Sun Feb 1 10:00:30 EST 1987
10:30	**%H:%M**	Tue Sep 23 10:30:00 EDT 1986
13:30	**%H:%M**	Mon Sep 22 13:30:00 EDT 1986

FILES

/usr/lib/locale/*locale***/LC_TIME/time**
locale specific date and time information

/usr/lib/locale/*locale***/LC_CTYPE/ctype**
character characterization information

SEE ALSO

setlocale(3C), **strftime**(3C), **strptime**(3C), **environ**(5)

NOTES

Subsequent calls to **getdate()** alter the contents of **getdate_err**.

Dates before 1970 and after 2037 are illegal.

The range of values for **%S** is [00,61] rather than [00,59] to allow for the occasional leap second and even more occasional double leap second.

getdate() makes explicit use of macros described in **ctype**(3C).

NAME

getdtablesize – get file descriptor table size

SYNOPSIS

```
#include <unistd.h>
int getdtablesize(void);
```

DESCRIPTION

Each process has a file descriptor table which is guaranteed to have at least 20 slots. The entries in the descriptor table are numbered with small integers starting at 0. The **getdtablesize( )** function returns the current maximum size of this table by calling the **getrlimit( )** function.

SEE ALSO

close(2), **getrlimit**(2), **open**(2), **pipe**(2), **select**(3C) **sysconf**(3C)

NAME

getenv – return value for environment name

SYNOPSIS

```
#include <stdlib.h>
char *getenv(const char *name);
```

MT-LEVEL

Safe

DESCRIPTION

getenv( ) searches the environment list (see **environ**(5)) for a string of the form *name*=*value* and, if the string is present, returns a pointer to the *value* in the current environment.

RETURN VALUES

If successful, **getenv( )** returns a pointer to the *value* in the current environment; otherwise, it returns a null pointer.

SEE ALSO

exec(2), **putenv**(3C), **environ**(5)

NOTES

getenv( ) can be safely called from a multi-thread program. However, care must still be taken when using **getenv( )** and **putenv**(3C) in a multi-thread program. These routines examine and modify the environment list. This list is shared by all threads in a program. The system prevents the list from being accessed simultaneously by two different threads. However, it does not prevent two threads from successively accessing the environment list using **getenv( )** or **putenv**(3C).

NAME

getfauditflags – generates the process audit state

SYNOPSIS

cc [*flag* ...] *file* ... **–lbsm –lsocket –lnsl –lintl** [*library* ...]

```
#include <sys/param.h>
#include <bsm/libbsm.h>

int getfauditflags(au_mask_t *usremasks, au_mask_t *usrdmasks, au_mask_t *lastmasks);
```

MT-LEVEL

MT-Safe.

AVAILABILITY

The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See **bsmconv**(1M) for more information.

DESCRIPTION

getfauditflags( ) generates a process audit state by combining the audit masks passed as parameters with the system audit masks specified in the **audit_control**(4) file. **getfauditflags( )** obtains the system audit value by calling **getacflg( )** (see **getacinfo**(3)).

usremasks points to **au_mask_t** fields which contains two values. The first value defines which events are *always* to be audited when they succeed. The second value defines which events are always to be audited when they fail.

usrdmasks also points to **au_mask_t** fields which contains two values. The first value defines which events are *never* to be audited when they succeed. The second value defines which events are never to be audited when they fail.

The structures pointed to by *usremasks* and *usrdmasks* may be obtained from the **audit_user**(4) file by calling **getauusernam( )** which returns a pointer to a strucure containing all **audit_user**(4) fields for a user.

The output of this function is stored in *lastmasks* which is a pointer of type **au_mask_t** as well. The first value defines which events are to be audited when they succeed and the second defines which events are to be audited when they fail.

Both *usremasks* and *usrdmasks* override the values in the system audit values.

RETURN VALUES

-1 is returned on error and **0** on success.

SEE ALSO

bsmconv(1M), **getacinfo**(3), **getauditflags**(3), **getauusernam**(3), **audit.log**(4), **audit_control**(4), **audit_user**(4)

NAME

getgrnam, getgrnam_r, getgrent, getgrent_r, getgrgid, getgrgid_r, setgrent, endgrent, fgetgrent, fgetgrent_r – get group entry

SYNOPSIS

```
#include <grp.h>

struct group *getgrnam(const char *name);

struct group *getgrnam_r(const char *name, struct group *result,
        char *buffer, int buflen);

struct group *getgrent(void);

struct group *getgrent_r(struct group *result, char *buffer, int buflen);

struct group *getgrgid(gid_t gid);

struct group *getgrgid_r(gid_t gid, struct group *result, char *buffer, int buflen);

void setgrent(void);

void endgrent(void);

struct group *fgetgrent(FILE *f);

struct group *fgetgrent_r(FILE *f, struct group *result, char *buffer, int buflen);
```

POSIX

```
cc [ flag... ] file ... -D_POSIX_PTHREAD_SEMANTICS [ library... ]

int getgrnam_r(const char *name, struct group *grp, char *buffer, size_t bufsize,
        struct group **result);

int getgrgid_r(gid_t gid, struct group *grp, char *buffer, size_t bufsize,
        struct group **result);
```

MT-LEVEL

See the subsection ‘‘Reentrant Interfaces’’ in the **DESCRIPTION** section of this page.

DESCRIPTION

These functions are used to obtain entries describing user groups. Entries can come from any of the sources for **group** specified in the **/etc/nsswitch.conf** file (see **nsswitch.conf**(4)).

getgrnam( ) searches for an entry with the group name specified by the character string parameter *name*.

getgrgid( ) searches for an entry with the (numeric) group id specified by *gid*.

The functions **setgrent( )**, **getgrent( )**, and **endgrent( )** are used to enumerate group entries from the database. **setgrent( )** sets (or resets) the enumeration to the beginning of the set of group entries. This function should be called before the first call to **getgrent( )**. Calls to **getgrnam( )** and **getgrgid( )** leave the enumeration position in an indeterminate state. Successive calls to **getgrent( )** return either successive entries or NULL, indicating the end of the enumeration.

endgrent( ) may be called to indicate that the caller expects to do no further group entry retrieval operations; the system may then close the group file, deallocate resources it was using, and so forth. It is still allowed, but possibly less efficient, for the process to call more group functions after calling **endgrent( )**.

fgetgrent( ), unlike the other functions above, does not use **nsswitch.conf**; it reads and parses the next line from the stream *f*, which is assumed to have the format of the **group** file (see **group**(4)).

Reentrant Interfaces

The functions **getgrnam( )**, **getgrgid( )**, **getgrent( )**, and **fgetgrent( )** use static storage that is re-used in each call, making them unsafe for multithreaded applications.

The parallel functions **getgrnam_r( )**, **getgrgid_r( )**, **getgrent_r( )**, and **fgetgrent_r( )** provide reentrant interfaces for these operations.

Each reentrant interface performs the same operation as its non-reentrant counterpart, named by removing the ''**_r**'' suffix. The reentrant interfaces, however, use buffers supplied by the caller to store returned results, and are safe for use in both single-threaded and multithreaded applications.

Each reentrant interface takes the same parameters as its non-reentrant counterpart, as well as the following additional parameters. The parameter *result* (or *grp* for the POSIX versions) must be a pointer to a **struct group** structure allocated by the caller. On successful completion, the function returns the group entry in this structure. The parameter *buffer* is a pointer to a buffer supplied by the caller, used as storage space for the group data. All of the pointers within the returned **struct group** *result* (or *grp*) point to data stored within this buffer; see **RETURN VALUES**. The buffer must be large enough to hold all the data associated with the group entry. The parameter *buflen* (or *bufsize* for the POSIX versions) should give the size in bytes of *buffer*. The POSIX versions place a pointer to the modified *grp* structure in the *result* parameter, instead of returning a pointer to this structure.

For enumeration in multithreaded applications, the position within the enumeration is a process-wide property shared by all threads. **setgrent( )** may be used in a multithreaded application but resets the enumeration position for all threads. If multiple threads interleave calls to **getgrent_r( )**, the threads will enumerate disjoint subsets of the group database.

Like their non-reentrant counterparts, **getgrnam_r( )** and **getgrgid_r( )** leave the enumeration position in an indeterminate state.

RETURN VALUES

Group entries are represented by the **struct group** structure defined in **<grp.h>**:

```
struct group {
        char *gr_name;      /* the name of the group */
        char *gr_passwd;    /* the encrypted group password */
        gid_t gr_gid;       /* the numerical group ID */
        char **gr_mem;      /* vector of pointers to member names */
};
```

The functions **getgrnam( )**, **getgrnam_r( )**, **getgrgid( )**, and **getgrgid_r( )** each return a pointer to a **struct group** if they successfully locate the requested entry; otherwise they return NULL. The POSIX functions **getgrnam_r( )** and **getgrgid_r( )** return zero upon success, or the error number in case of failure.

The functions **getgrent( )**, **getgrent_r( )**, **fgetgrent( )**, and **fgetgrent_r( )** each return a pointer to a **struct group** if they successfully enumerate an entry; otherwise they return NULL, indicating the end of the enumeration.

The functions **getgrnam( )**, **getgrgid( )**, **getgrent( )**, and **fgetgrent( )** use static storage, so returned data must be copied before a subsequent call to any of these functions if the data is to be saved.

When the pointer returned by the reentrant functions **getgrnam_r( )**, **getgrgid_r( )**, **getgrent_r( )**, and **fgetgrent_r( )** is non-NULL, it is always equal to the *result* pointer that was supplied by the caller.

ERRORS

The reentrant functions **getgrnam_r( )**, **getgrgid_r( )**, **getgrent_r( )**, and **fgetgrent_r( )** will return NULL and set *errno* to **ERANGE** (or in the case of POSIX functions **getgrnam_r( )** and **getgrgid_r( )** return the **ERANGE** error) if the length of the buffer supplied by caller is not large enough to store the result. See **Intro**(2) for the proper usage and interpretation of *errno* in multithreaded applications.

FILES

/etc/group
/etc/nsswitch.conf

SEE ALSO

getpwnam(3C), **group**(4), **nsswitch.conf**(4), **passwd**(4)

NOTES

Programs that use the interfaces described in this manual page cannot be linked statically since the implementations of these functions employ dynamic loading and linking of shared objects at run time.

When compiling multithreaded applications, see **Intro**(3), *Notes On Multithread Applications*, for information about the use of the **_REENTRANT** flag.

Use of the enumeration interfaces **getgrent( )** and **getgrent_r( )** is discouraged; enumeration is supported for the group file, NIS, and NIS+, but in general is not efficient and may not be supported for all database sources. The semantics of enumeration are discussed further in **nsswitch.conf**(4).

Previous releases allowed the use of ‘‘+’’ and ‘‘-’’ entries in **/etc/group** to selectively include and exclude entries from NIS. The primary usage of these ‘‘+/-’’ entries is superseded by the name service switch, so *the ‘‘+/-’’ form may not be supported in future releases*.

If required, the ‘‘+/-’’ functionality can still be obtained for NIS by specifying **compat** as the source for **group**.

If the ‘‘+/-’’ functionality is required in conjunction with NIS+, specify both **compat** as the source for **group** and **nisplus** as the source for the pseudo-database **group_compat**. See **group**(4), and **nsswitch.conf**(4) for details.

The reentrant interfaces **getgrnam_r( )**, **getgrgid_r( )**, **getgrent_r( )**, and **fgetgrent_r( )** are as specified in POSIX 1003.1c Draft #10.

NAME

gethostbyname, gethostbyname_r, gethostbyaddr, gethostbyaddr_r, gethostent, gethostent_r, sethostent, endhostent – get network host entry

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>
#include <netinet/in.h>
#include <arpa/inet.h>
#include <netdb.h>

struct hostent *gethostbyname(const char *name);

struct hostent *gethostbyname_r(const char *name, struct hostent *result,
        char *buffer, int buflen, int *h_errnop);

struct hostent *gethostbyaddr(const char *addr, int len, int type);

struct hostent *gethostbyaddr_r(const char *addr, int length, int type,
        struct hostent *result, char *buffer, int buflen, int *h_errnop);

struct hostent *gethostent(void);

struct hostent *gethostent_r(struct hostent *result, char *buffer, int buflen,
        int *h_errnop);

int sethostent(int stayopen);

int endhostent(void);
```

MT-LEVEL

See the subsection “Reentrant Interfaces” in the **DESCRIPTION** section of this page.

DESCRIPTION

These functions are used to obtain entries describing hosts. An entry may come from any of the sources for **hosts** specified in the **/etc/nsswitch.conf** file (see **nsswitch.conf**(4)).

gethostbyname( ) searches for information for a host with the hostname specified by the character-string parameter *name*.

gethostbyaddr( ) searches for information for a host with a given host address. The parameter *type* specifies the family of the address. This should be one of the address families defined in **<sys/socket.h>**. The parameter *addr* must be a pointer to a buffer containing the address. The address is given in a form specific to the address family. See the **NOTES** section below for more information. Also see the **EXAMPLES** section below on how to convert a ‘‘.’’ separated Internet IP address notation into the *addr* parameter. The parameter *len* specifies the length of the buffer indicated by *addr*.

The functions **sethostent( )**, **gethostent( )**, and **endhostent( )** are used to enumerate host entries from the database.

sethostent( ) sets (or resets) the enumeration to the beginning of the set of host entries. This function should be called before the first call to **gethostent( )**. Calls to **gethostbyname( )** and **gethostbyaddr( )** leave the enumeration position in an indeterminate state. If the *stayopen* flag is non-zero, the system may keep allocated resources such as open file descriptors until a subsequent call to **endhostent( )**.

Successive calls to **gethostent()** return either successive entries or **NULL**, indicating the end of the enumeration.

endhostent() may be called to indicate that the caller expects to do no further host entry retrieval operations; the system may then deallocate resources it was using. It is still allowed, but possibly less efficient, for the process to call more host retrieval functions after calling **endhostent()**.

Reentrant Interfaces

The functions **gethostbyname()**, **gethostbyaddr()**, and **gethostent()** use static storage that is re-used in each call, making these functions unsafe for use in multithreaded applications.

The functions:

gethostbyname_r(),
gethostbyaddr_r(),

and

gethostent_r()

provide reentrant interfaces for these operations.

Each reentrant interface performs the same operation as its non-reentrant counterpart, named by removing the ''**_r**'' suffix. The reentrant interfaces, however, use buffers supplied by the caller to store returned results, and are safe for use in both single-threaded and multithreaded applications.

Each reentrant interface takes the same parameters as its non-reentrant counterpart, as well as the following additional parameters. The parameter *result* must be a pointer to a **struct hostent** structure allocated by the caller. On successful completion, the function returns the host entry in this structure. The parameter *buffer* must be a pointer to a buffer supplied by the caller. This buffer is used as storage space for the host data. All of the pointers within the returned **struct hostent** *result* point to data stored within this buffer (see **RETURN VALUES**). The buffer must be large enough to hold all of the data associated with the host entry. The parameter *buflen* should give the size in bytes of the buffer indicated by *buffer*. The parameter *h_errnop* should be a pointer to an integer. An integer error status value is stored there on certain error conditions. (see **ERRORS**).

For enumeration in multithreaded applications, the position within the enumeration is a process-wide property shared by all threads. **sethostent()** may be used in a multithreaded application but resets the enumeration position for all threads. If multiple threads interleave calls to **gethostent_r()**, the threads will enumerate disjoint subsets of the host database.

Like their non-reentrant counterparts, **gethostbyname_r()** and **gethostbyaddr_r()** leave the enumeration position in an indeterminate state.

RETURN VALUES

Host entries are represented by the **struct hostent** structure defined in **<netdb.h>**:

```
struct hostent {
        char    *h_name;        /* canonical name of host */
        char    **h_aliases;    /* alias list */
        int     h_addrtype;     /* host address type */
        int     h_length;       /* length of address */
        char    **h_addr_list;  /* list of addresses */
};
```

See the **EXAMPLES** section below on how to retrieve a ''.'' separated Internet IP address string from the *h_addr_list* field of **struct hostent .**

The functions **gethostbyname( )**, **gethostbyname_r( )**, **gethostbyaddr( )**, and **gethostbyaddr_r( )** each return a pointer to a **struct hostent** if they successfully locate the requested entry; otherwise they return **NULL**.

The functions **gethostent( )** and **gethostent_r( )** each return a pointer to a **struct hostent** if they successfully enumerate an entry; otherwise they return **NULL**, indicating the end of the enumeration.

The functions **gethostbyname( )**, **gethostbyaddr( )**, and **gethostent( )** use static storage, so returned data must be copied before a subsequent call to any of these functions if the data is to be saved.

When the pointer returned by the reentrant functions **gethostbyname_r( )**, **gethostbyaddr_r( )**, and **gethostent_r( )** is not **NULL**, it is always equal to the *result* pointer that was supplied by the caller.

The functions **sethostent( )** and **endhostent( )** return **0** on success.

ERRORS

The reentrant functions **gethostbyname_r( )**, **gethostbyaddr_r( )** and **gethostent_r( )** will return **NULL** and set *errno* to **ERANGE** if the length of the buffer supplied by caller is not large enough to store the result. See **intro**(2) for the proper usage and interpretation of **errno** in multithreaded applications.

On failures, the non-reentrant functions **gethostbyname( )** and **gethostbyaddr( )** set a global integer *h_errno* to indicate one of these error codes (defined in **<netdb.h>**): **HOST_NOT_FOUND**, **TRY_AGAIN**, **NO_RECOVERY**, **NO_DATA**, and **NO_ADDRESS**. The reentrant functions **gethostbyname_r( )** and **gethostbyaddr_r( )** set the integer pointed to by *h_errnop* to one of these values in case of error.

EXAMPLES

Here is a sample program that gets the canonical name, aliases, and ''.'' separated Internet IP addresses for a given ''.'' separated IP address:

```
#include <stdio.h>
#include <string.h>
#include <sys/types.h>
#include <sys/socket.h>
#include <netinet/in.h>
#include <arpa/inet.h>
```

```
#include <netdb.h>

main(int argc, const char **argv)
{
      u_long addr;
      struct hostent *hp;
      char **p;

      if (argc != 2) {
         (void) printf("usage: %s IP-address\n", argv[0]);
         exit (1);
      }
      if ((int)(addr = inet_addr(argv[1])) == -1) {
         (void) printf("IP-address must be of the form a.b.c.d\n");
         exit (2);
      }

      hp = gethostbyaddr((char *)&addr, sizeof (addr), AF_INET);
      if (hp == NULL) {
         (void) printf("host information for %s not found\n", argv[1]);
         exit (3);
      }

      for (p = hp->h_addr_list; *p != 0; p++) {
         struct in_addr in;
         char **q;

         (void) memcpy(&in.s_addr, *p, sizeof (in.s_addr));
         (void) printf("%s\t%s", inet_ntoa(in), hp–>h_name);
         for (q = hp->h_aliases; *q != 0; q++)
            (void) printf(" %s", *q);
         (void) putchar('\n');
      }
      exit (0);
}
```

Note that the above sample program is unsafe for use in multithreaded applications.

FILES

/etc/hosts
/etc/netconfig
/etc/nsswitch.conf

SEE ALSO

inet(3N), **netdir**(3N), **hosts**(4), **netconfig**(4), **nsswitch.conf**(4)

WARNINGS

The reentrant interfaces **gethostbyname_r( )**, **gethostbyaddr_r( )**, and **gethostent_r( )** are included in this release on an uncommitted basis only, and are subject to change or removal in future minor releases.

NOTES

Programs that use the interfaces described in this manual page cannot be linked statically since the implementations of these functions employ dynamic loading and linking of shared objects at run time.

In order to ensure that they all return consistent results, **gethostbyname( ), gethostbyname_r( ),** and **netdir_getbyname( )** are implemented in terms of the same internal library function. This function obtains the system-wide source lookup policy based on the **inet** family entries in **netconfig**(4) and the **hosts:** entry in **nsswitch.conf**(4). Similarly, **gethostbyaddr( ), gethostbyaddr_r( ),** and **netdir_getbyaddr( )** are implemented in terms of the same internal library function. If the **inet** family entries in **netconfig**(4) have a ''-'' in the last column for nametoaddr libraries, then the entry for **hosts** in **nsswitch.conf** will be used; otherwise the nametoaddr libraries in that column will be used, and **nsswitch.conf** will not be consulted.

There is no analogue of **gethostent( )** and **gethostent_r( )** in the netdir functions, so these enumeration functions go straight to the **hosts** entry in **nsswitch.conf**. Thus enumeration may return results from a different source than that used by **gethostbyname( ), gethostbyname_r( ), gethostbyaddr( ),** and **gethostbyaddr_r( ).**

All the functions that return a **struct hostent** must always return the *canonical name* in the *h_name* field. This name, by definition, is the well-known and official hostname shared between all aliases and all addresses. The underlying source that satisfies the request determines the mapping of the input name or address into the set of names and addresses in **hostent.** Different sources might do that in different ways. If there is more than one alias and more than one address in **hostent,** no pairing is implied between them.

The system will strive to put the addresses on the same subnet as that of the caller first.

When compiling multithreaded applications, see **Intro**(3), *Notes On Multithread Applications*, for information about the use of the **_REENTRANT** flag.

Use of the enumeration interfaces **gethostent( )** and **gethostent_r( )** is discouraged; enumeration may not be supported for all database sources. The semantics of enumeration are discussed further in **nsswitch.conf**(4).

The current implementations of these functions only return or accept addresses for the Internet address family (type **AF_INET).**

The form for an address of type **AF_INET** is a **struct in_addr** defined in **<netinet/in.h>.** The functions described in **inet**(3N), and illustrated in the **EXAMPLES** section above, are helpful in constructing and manipulating addresses in this form.

NAME

gethostid – get unique identifier of current host

SYNOPSIS

```
#include <unistd.h>
long gethostid(void);
```

DESCRIPTION

gethostid( ) returns the 32-bit identifier for the current host, which should be unique across all hosts. This number is usually taken from the CPU board's ID PROM.

SEE ALSO

hostid(1), **sysinfo**(2)

NAME

gethostname, sethostname – get / set name of current host

SYNOPSIS

int gethostname(char **name*, **int** *namelen*);

int sethostname(char **name*, **int** *namelen*);

DESCRIPTION

gethostname( ) returns the standard host name for the current processor, as previously set by **sethostname**. The parameter *namelen* specifies the size of the array pointed to by *name*. The returned name is null-terminated unless insufficient space is provided.

sethostname( ) sets the name of the host machine to be *name*, which has length *namelen*. This call is restricted to the privileged user and is normally used only when the system is bootstrapped.

RETURN VALUES

If the call succeeds a value of 0 is returned. If the call fails, then a value of –1 is returned and an error code is placed in the global location **errno**.

ERRORS

The following error may be returned by these calls:

EFAULT	The *name* or *namelen* parameter gave an invalid address.
EPERM	The caller was not the privileged user. Note: this error only applies to **sethostname( )**.

SEE ALSO

uname(2), **sysinfo**(2), **gethostid**(3C)

NOTES

Host names are limited to **MAXHOSTNAMELEN** characters, currently 256. (See the **<sys/param.h>** header.)

NAME

gethrtime, gethrvtime – get high resolution time

SYNOPSIS

```
#include <sys/time.h>
hrtime_t gethrtime(void);
hrtime_t gethrvtime(void);
```

MT-LEVEL

MT-Safe

DESCRIPTION

gethrtime( ) returns the current high-resolution real time. Time is expressed as nanoseconds since some arbitrary time in the past; it is not correlated in any way to the time of day, and thus is *not* subject to resetting, drifting, etc. via **adjtime**(2) or **settimeofday**(3C). The hi-res timer is ideally suited to performance measurement tasks, where cheap, accurate interval timing is required.

gethrvtime( ) returns the current high-resolution LWP virtual time, expressed as total nanoseconds of execution time.

gethrtime( ) and **gethrvtime( )** both return an **hrtime_t,** which is a 64-bit (**long long**) signed integer.

EXAMPLE

The following code fragment measures the average cost of **getpid**(2):

```
hrtime_t start, end;
int i, iters = 100;

start = gethrtime();
for (i = 0; i < iters; i++)
        getpid();
end = gethrtime();

printf("Avg getpid() time = %lld nsec\n", (end - start) / iters);
```

SEE ALSO

adjtime(2), **gettimeofday**(3C), **settimeofday**(3C)

NOTES

Although the units of hi-res time are always the same (nanoseconds), the actual resolution is hardware dependent. Hi-res time is guaranteed to be monotonic (it won't go backward, it won't periodically wrap) and linear (it won't occasionally speed up or slow down for adjustment, like the time of day can), but not necessarily unique: two sufficiently proximate calls may return the same value.

NAME

getlogin, getlogin_r – get login name

SYNOPSIS

```
#include <stdlib.h>
char *getlogin(void);
char *getlogin_r(char *name, int namelen);
```

POSIX

```
cc [ flag... ] file ... -D_POSIX_PTHREAD_SEMANTICS [ library... ]
int getlogin_r(char *name, size_t namesize);
```

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

getlogin( ) returns a pointer to the login name as found in **/var/adm/utmp**. It may be used in conjunction with **getpwnam( )** to locate the correct password file entry when the same user id is shared by several login names.

If **getlogin( )** is called within a process that is not attached to a terminal, it returns a null pointer. The correct procedure for determining the login name is to call **cuserid( )**, or to call **getlogin( )** and if it fails to call **getpwuid( )**.

getlogin_r( ) has the same functionality as **getlogin( )** except that the caller must supply a buffer *name* with length *namelen* to store the result. The *name* buffer must be at least **LOGNAME_MAX** bytes in size (defined in **<limits.h>**). The POSIX version of **getlogin_r( )** takes a *namesize* parameter of type **size_t**.

RETURN VALUES

Returns a null pointer if the login name is not found.

ERRORS

getlogin_r( ) will fail if the following is true:

ERANGE The size of the buffer is smaller than the result to be returned.

The POSIX **getlogin_r( )** returns zero if successful, or the error number upon failure.

FILES

/var/adm/utmp

SEE ALSO

cuserid(3S), **getgrnam**(3C), **getpwnam**(3C), **utmp**(4)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

The return values point to static data whose content is overwritten by each call.

getlogin( ) is unsafe in multi-thread applications. **getlogin_r( )** should be used instead.

The new **getlogin_r( )** interface is as specified in POSIX 1003.1c Draft #10.

NAME

getmntent, getmntany, hasmntopt, putmntent – get mnttab file information

SYNOPSIS

```
#include <stdio.h>
#include <sys/mnttab.h>

int getmntent(FILE *fp, struct mnttab *mp);

int getmntany(FILE *fp, struct mnttab *mp, struct mnttab *mpref);

char *hasmntopt(struct mnttab *mnt, char *opt);

int putmntent(FILE *iop, struct mnttab *mp);
```

MT-LEVEL

Safe

DESCRIPTION

getmntent( ) and **getmntany( )** each fill in the structure pointed to by *mp* with the broken-out fields of a line in the **/etc/mnttab** file. Each line in the file contains a **mnttab** structure, which is declared in the **<sys/mnttab.h>** header. The structure contains the following members:

```
char    *mnt_special;
char    *mnt_mountp;
char    *mnt_fstype;
char    *mnt_mntopts;
char    *mnt_time;
```

The fields have meanings described in **mnttab**(4).

getmntent( ) returns a pointer to the next **mnttab** structure in the file; so successive calls can be used to search the entire file. **getmntany( )** searches the file referenced by *fp* until a match is found between a line in the file and *mpref*. *mpref* matches the line if all non-null entries in *mpref* match the corresponding fields in the file. Note that these routines do not open, close, or rewind the file.

hasmntopt( ) scans the **mnt_mntopts** field of the **mnttab** structure *mnt* for a substring that matches *opt*. It returns the address of the substring if a match is found, otherwise it returns 0.

The **putmntent( )** macro formats the contents of the **mnttab** structure according to the layout required for the **/etc/mnttab** file and writes the entry to the file. Note: the file should be opened in append mode (**fopen**(3S) with an "a" mode) so that the entry is appended to the file.

RETURN VALUES

If the next entry is successfully read by **getmntent( )** or a match is found with **getmntany( )**, 0 is returned. If an EOF is encountered on reading, these functions return –1. If an error is encountered, a value greater than 0 is returned. The possible error values are:

MNT_TOOLONG	A line in the file exceeded the internal buffer size of **MNT_LINE_MAX**.
MNT_TOOMANY	A line in the file contains too many fields.
MNT_TOOFEW	A line in the file contains too few fields.

On success, **putmntent()** returns the number of bytes printed to the specified file and on failure returns EOF.

FILES

/etc/mnttab

SEE ALSO

mnttab(4)

NOTES

The members of the **mnttab** structure point to information contained in a static area, so it must be copied if it is to be saved.

NAME

getnetbyname, getnetbyname_r, getnetbyaddr, getnetbyaddr_r, getnetent, getnetent_r, setnetent, endnetent – get network entry

SYNOPSIS

cc [*flag* ...] *file* ... –**lsocket -lnsl** [*library* ...]

#include <netdb.h>

struct netent ∗**getnetbyname(const char** ∗*name*);

struct netent ∗**getnetbyname_r(const char** ∗*name*, **struct netent** ∗*result*, **char** ∗*buffer*, **int** *buflen*);

struct netent ∗**getnetbyaddr(long** *net*, **int** *type*);

struct netent ∗**getnetbyaddr_r(long** *net*, **int** *type*, **struct netent** ∗*result*, **char** ∗*buffer*, **int** *buflen*);

struct netent ∗**getnetent(void);**

struct netent ∗**getnetent_r(struct netent** ∗*result*, **char** ∗*buffer*, **int** *buflen*);

int setnetent(int *stayopen*);

int endnetent(void);

MT-LEVEL

See the subsection “Reentrant Interfaces” in the **DESCRIPTION** section of this page.

DESCRIPTION

These functions are used to obtain entries for networks. An entry may come from any of the sources for **networks** specified in the **/etc/nsswitch.conf** file (see **nsswitch.conf**(4)).

getnetbyname( ) searches for a network entry with the network name specified by the character string parameter *name*.

getnetbyaddr( ) searches for a network entry with the network address specified by *net*. The parameter *type* specifies the family of the address. This should be one of the address families defined in **<sys/socket.h>**. See the **NOTES** section below for more information.

The functions **setnetent( )**, **getnetent( )**, and **endnetent( )** are used to enumerate network entries from the database.

setnetent( ) sets (or resets) the enumeration to the beginning of the set of network entries. This function should be called before the first call to **getnetent( )**. Calls to **getnetbyname( )** and **getnetbyaddr( )** leave the enumeration position in an indeterminate state. If the *stayopen* flag is non-zero, the system may keep allocated resources such as open file descriptors until a subsequent call to **endnetent( )**.

Successive calls to **getnetent( )** return either successive entries or NULL, indicating the end of the enumeration.

endnetent( ) may be called to indicate that the caller expects to do no further network entry retrieval operations; the system may then deallocate resources it was using. It is still allowed, but possibly less efficient, for the process to call more network entry retrieval functions after calling **endnetent( )**.

Reentrant Interfaces

The functions **getnetbyname( )**, **getnetbyaddr( )**, and **getnetent( )** use static storage that is re-used in each call, making these routines unsafe for use in multithreaded applications.

The functions:

getnetbyname_r( ),
getnetbyaddr_r( ),

and

getnetent_r( )

provide reentrant interfaces for these operations.

Each reentrant interface performs the same operation as its non-reentrant counterpart, named by removing the ''**_r**'' suffix. The reentrant interfaces, however, use buffers supplied by the caller to store returned results, and are safe for use in both single-threaded and multithreaded applications.

Each reentrant interface takes the same parameters as its non-reentrant counterpart, as well as the following additional parameters. The parameter *result* must be a pointer to a **struct netent** structure allocated by the caller. On successful completion, the function returns the network entry in this structure. The parameter *buffer* must be a pointer to a buffer supplied by the caller. This buffer is used as storage space for the network entry data. All of the pointers within the returned **struct netent** *result* point to data stored within this buffer (see **RETURN VALUES**). The buffer must be large enough to hold all of the data associated with the network entry. The parameter *buflen* should give the size in bytes of the buffer indicated by *buffer* .

For enumeration in multithreaded applications, the position within the enumeration is a process-wide property shared by all threads. **setnetent( )** may be used in a multithreaded application but resets the enumeration position for all threads. If multiple threads interleave calls to **getnetent_r( )**, the threads will enumerate disjoint subsets of the network database.

Like their non-reentrant counterparts, **getnetbyname_r( )** and **getnetbyaddr_r( )** leave the enumeration position in an indeterminate state.

RETURN VALUES

Network entries are represented by the **struct netent** structure defined in **<netdb.h>**:

```
struct netent {
        char    *n_name;
        char    **n_aliases;
        int     n_addrtype;
        long    n_net;
};
```

The functions **getnetbyname( )**, **getnetbyname_r( )**, **getnetbyaddr( )**, and **getnetbyaddr_r( )** each return a pointer to a **struct netent** if they successfully locate the requested entry; otherwise they return NULL.

The functions **getnetent( )** and **getnetent_r( )** each return a pointer to a **struct netent** if they successfully enumerate an entry; otherwise they return NULL, indicating the end of the enumeration.

The functions **getnetbyname( )**, **getnetbyaddr( )**, and **getnetent( )** use static storage, so returned data must be copied before a subsequent call to any of these functions if the data is to be saved.

When the pointer returned by the reentrant functions **getnetbyname_r( )**, **getnetbyaddr_r( )**, and **getnetent_r( )** is non-NULL, it is always equal to the *result* pointer that was supplied by the caller.

The functions **setnetent( )** and **endnetent( )** return **0** on success.

ERRORS

The reentrant functions **getnetbyname_r( )**, **getnetbyaddr_r( )** and **getnetent_r( )** will return NULL and set *errno* to **ERANGE** if the length of the buffer supplied by caller is not large enough to store the result. See **intro**(2) for the proper usage and interpretation of *errno* in multithreaded applications.

FILES

/etc/networks
/etc/nsswitch.conf

SEE ALSO

inet(3N), **networks**(4), **nsswitch.conf**(4)

WARNINGS

The reentrant interfaces **getnetbyname_r( )**, **getnetbyaddr_r( )**, and **getnetent_r( )** are included in this release on an uncommitted basis only, and are subject to change or removal in future minor releases.

NOTES

The current implementation of these functions only return or accept network numbers for the Internet address family (type **AF_INET).** The functions described in **inet**(3N) may be helpful in constructing and manipulating addresses and network numbers in this form.

Programs that use the interfaces described in this manual page cannot be linked statically since the implementations of these functions employ dynamic loading and linking of shared objects at run time.

When compiling multithreaded applications, see **Intro**(3), *Notes On Multithread Applications*, for information about the use of the **_REENTRANT** flag.

Use of the enumeration interfaces **getnetent( )** and **getnetent_r( )** is discouraged; enumeration may not be supported for all database sources. The semantics of enumeration are discussed further in **nsswitch.conf**(4).

NAME

getnetconfig, setnetconfig, endnetconfig, getnetconfigent, freenetconfigent, nc_perror, nc_sperror – get network configuration database entry

SYNOPSIS

```
#include <netconfig.h>

struct netconfig *getnetconfig(void *handlep);

void *setnetconfig(void);

int endnetconfig(void *handlep);

struct netconfig *getnetconfigent(const char *netid);

void freenetconfigent(struct netconfig *netconfigp);

void nc_perror(const char *msg);

char *nc_sperror(void);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The library routines described on this page are part of the Network Selection component. They provide the application access to the system network configuration database, **/etc/netconfig**. In addition to the routines for accessing the **netconfig** database, Network Selection includes the environment variable **NETPATH** (see **environ**(5)) and the **NETPATH** access routines described in **getnetpath**(3N).

getnetconfig( ) returns a pointer to the current entry in the **netconfig** database, formatted as a **struct netconfig**. Successive calls will return successive **netconfig** entries in the **netconfig** database. **getnetconfig( )** can be used to search the entire **netconfig** file. **getnetconfig( )** returns NULL at the end of the file. *handlep* is the handle obtained through **setnetconfig( )**.

A call to **setnetconfig( )** has the effect of ''binding'' to or ''rewinding'' the **netconfig** database. **setnetconfig( )** must be called before the first call to **getnetconfig( )** and may be called at any other time. **setnetconfig( )** need *not* be called before a call to **getnetconfigent( )**. **setnetconfig( )** returns a unique handle to be used by **getnetconfig( ) .**

endnetconfig( ) should be called when processing is complete to release resources for reuse. *handlep* is the handle obtained through **setnetconfig( )**. Programmers should be aware, however, that the last call to **endnetconfig( )** frees all memory allocated by **getnetconfig( )** for the **struct netconfig** data structure. **endnetconfig( )** may not be called before **setnetconfig( )**.

getnetconfigent( ) returns a pointer to the **struct netconfig** structure corresponding to *netid*. It returns NULL if *netid* is invalid (that is, does not name an entry in the **netconfig** database).

freenetconfigent( ) frees the netconfig structure pointed to by *netconfigp* (previously returned by **getnetconfigent( )**).

nc_perror( ) prints a message to the standard error indicating why any of the above routines failed. The message is prepended with the string *msg* and a colon. A NEWLINE is appended at the end of the message.

nc_sperror( ) is similar to **nc_perror( )** but instead of sending the message to the standard error, will return a pointer to a string that contains the error message.

nc_perror( ) and **nc_sperror( )** can also be used with the **NETPATH** access routines defined in **getnetpath**(3N).

RETURN VALUES

setnetconfig( ) returns a unique handle to be used by **getnetconfig( )**. In the case of an error, **setnetconfig( )** returns NULL and **nc_perror( )** or **nc_sperror( )** can be used to print the reason for failure.

getnetconfig( ) returns a pointer to the current entry in the **netconfig( )** database, formatted as a **struct netconfig. getnetconfig( )** returns NULL at the end of the file, or upon failure.

endnetconfig( ) returns **0** on success and **–1** on failure (for example, if **setnetconfig( )** was not called previously).

On success, **getnetconfigent( )** returns a pointer to the **struct netconfig** structure corresponding to *netid*; otherwise it returns NULL.

nc_sperror( ) returns a pointer to a buffer which contains the error message string. This buffer is overwritten on each call. In multithreaded applications, this buffer is implemented as thread-specific data.

SEE ALSO

getnetpath(3N), **netconfig**(4), **environ**(5)

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NAME

getnetgrent, getnetgrent_r, setnetgrent, endnetgrent, innetgr – get network group entry

SYNOPSIS

int getnetgrent(char ∗∗*machinep*, **char** ∗∗*userp*, **char** ∗∗*domainp*);

int getnetgrent_r(char ∗∗*machinep*, **char** ∗∗*userp*, **char** ∗∗*domainp*, **char** ∗*buffer*, **int** *buflen*);

void setnetgrent(const char ∗*netgroup*);

void endnetgrent(void);

int innetgr(const char ∗*netgroup*, **const char** ∗*machine*, **const char** ∗*user*, **const char** ∗*domain*);

MT-LEVEL

See the **DESCRIPTION** section of this page.

DESCRIPTION

These functions are used to test membership in and enumerate members of ‘‘netgroup’’ network groups defined in a system database. Netgroups are sets of (machine,user,domain) triples (see **netgroup**(4)).

These functions consult the source specified for **netgroup** in the **/etc/nsswitch.conf** file (see **nsswitch.conf**(4)).

The function **innetgr( )** returns 1 if there is a netgroup *netgroup* that contains the specified *machine, user, domain* triple as a member; otherwise it returns 0. Any of the supplied pointers *machine, user,* and *domain* may be **NULL,** signifying a "wild card" that matches all values in that position of the triple.

The **innetgr( )** function is safe for use in single-threaded and multi-threaded applications.

The functions **setnetgrent( )**, **getnetgrent( )**, and **endnetgrent( )** are used to enumerate the members of a given network group.

The function **setnetgrent( )** establishes the network group specified in the parameter *net-group* as the current group whose members are to be enumerated.

Successive calls to the function **getnetgrent( )** will enumerate the members of the group established by calling **setnetgrent( )**; each call returns 1 if it succeeds in obtaining another member of the network group, or 0 if there are no further members of the group.

When calling either **getnetgrent( )** or **getnetgrent_r( ) ,** addresses of the three character pointers are used as arguments; i.e.:

char ∗*mp*, ∗*up*, ∗*dp*;

getnetgrent(&*mp*, &*up*, &*dp*);

Upon successful return from **getnetgrent( ),** the pointer *mp* points to a string containing the name of the machine part of the member triple, *up* points to a string containing the user name and *dp* points to a string containing the domain name. If the pointer returned for *mp, up,* or *dp* is **NULL,** it signifies that the element of the netgroup contains wild card specifier in that position of the triple.

The pointers returned by **getnetgrent( )** point into a buffer allocated by **setnetgrent( )** that is re-used by in each call. This space is released when an **endnetgrent( )** call is made, and should not be released by the caller. This implementation is not safe for use in multi-threaded applications.

The function **getnetgrent_r( )** is similar to **getnetgrent( )** but uses a buffer supplied by the caller for the space needed to store the results. The parameter *buffer* should be a pointer to a buffer allocated by the caller and the length of this buffer should be specified by the parameter *buflen*. The buffer must be large enough to hold the data associated with the triple. The **getnetgrent_r( )** function is safe for use both in single-threaded and multi-threaded applications.

The function **endnetgrent( )** frees the space allocated by the previous **setnetgrent( )** call. The equivalent of an **endnetgrent( )** implicitly performed whenever a **setnetgrent( )** call is made to a new network group.

Note that while **setnetgrent( )** and **endnetgrent( )** are safe for use in multi-threaded applications, the effect of each is process-wide. Calling **setnetgrent( )** resets the enumeration position for all threads. If multiple threads interleave calls to **getnetgrent_r( )** each will enumerate a disjoint subset of the netgroup. Thus the effective use of these functions in multi-threaded applications may require coordination by the caller.

ERRORS

The function **getnetgrent_r( )** will return 0 and set **errno** to **ERANGE** if the length of the buffer supplied by caller is not large enough to store the result. See **intro**(2) for the proper usage and interpretation of **errno** in multi-threaded applications.

FILES

/etc/nsswitch.conf

SEE ALSO

netgroup(4) **nsswitch.conf**(4),

WARNINGS

The function **getnetgrent_r( )** is included in this release on an uncommitted basis only, and is subject to change or removal in future minor releases.

NOTES

Only the Network Information Services, NIS and NIS+, are supported as sources for the **netgroup** database.

Programs that use the interfaces described in this manual page cannot be linked statically since the implementations of these functions employ dynamic loading and linking of shared objects at run time.

When compiling multi-threaded applications, see **Intro**(3), *Notes On Multithread Applications*, for information about the use of the **_REENTRANT** flag.

NAME

getnetpath, setnetpath, endnetpath – get /etc/netconfig entry corresponding to NETPATH component

SYNOPSIS

```
#include <netconfig.h>
struct netconfig *getnetpath(void *handlep);
void *setnetpath(void);
int endnetpath(void *handlep);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The routines described on this page are part of the Network Selection component. They provide the application access to the system network configuration database, **/etc/netconfig**, as it is ''filtered'' by the **NETPATH** environment variable (see **environ**(5)). See **getnetconfig**(3N) for other routines that also access the network configuration database directly. The **NETPATH** variable is a list of colon-separated network identifiers.

getnetpath( ) returns a pointer to the **netconfig** database entry corresponding to the first valid **NETPATH** component. The **netconfig** entry is formatted as a **struct netconfig**. On each subsequent call, **getnetpath( )** returns a pointer to the **netconfig** entry that corresponds to the next valid **NETPATH** component. **getnetpath( )** can thus be used to search the **netconfig** database for all networks included in the **NETPATH** variable. When **NETPATH** has been exhausted, **getnetpath( )** returns NULL.

A call to **setnetpath( )** ''binds'' to or ''rewinds'' **NETPATH**. **setnetpath( )** must be called before the first call to **getnetpath( )** and may be called at any other time. It returns a handle that is used by **getnetpath( )**.

getnetpath( ) silently ignores invalid **NETPATH** components. A **NETPATH** component is invalid if there is no corresponding entry in the **netconfig** database.

If the **NETPATH** variable is *unset*, **getnetpath( )** behaves as if **NETPATH** were set to the sequence of ''default'' or ''visible'' networks in the **netconfig** database, in the order in which they are listed.

endnetpath( ) may be called to ''unbind'' from **NETPATH** when processing is complete, releasing resources for reuse. Programmers should be aware, however, that **endnetpath( )** frees all memory allocated by **getnetpath( )** for the **struct netconfig** data structure. **endnetpath( )** returns **0** on success and **-1** on failure (for example, if **setnetpath( )** was not called previously).

RETURN VALUES

setnetpath( ) returns a handle that is used by **getnetpath( )**. In case of an error, **setnetpath( )** returns NULL. **nc_perror( )** or **nc_sperror( )** can be used to print out the reason for failure. See **getnetconfig**(3N).

When first called, **getnetpath( )** returns a pointer to the **netconfig** database entry corresponding to the first valid **NETPATH** component. When **NETPATH** has been exhausted, **getnetpath( )** returns NULL.

endnetpath( ) returns **0** on success and **-1** on failure (for example, if **setnetpath( )** was not called previously).

SEE ALSO

getnetconfig(3N), **netconfig**(4), **environ**(5)

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NAME

getopt – get option letter from argument vector

SYNOPSIS

```
#include <stdlib.h>
int getopt(int argc, char * const *argv, const char *optstring);
extern char *optarg;
extern int optind, opterr, optopt;
```

MT-LEVEL

Unsafe

DESCRIPTION

getopt( ) returns the next option letter in *argv* that matches a letter in *optstring*. It supports all the rules of the command syntax standard (see **intro**(1)). Since all new commands are intended to adhere to the command syntax standard, they should use **getopts**(1), **getopt**(3C) or **getsubopt**(3C) to parse positional parameters and check for options that are legal for that command.

optstring must contain the option letters the command using **getopt( )** will recognize; if a letter is followed by a colon, the option is expected to have an argument, or group of arguments, which may be separated from it by white space. *optarg* is set to point to the start of the option argument on return from **getopt( )**.

getopt( ) places in *optind* the *argv* index of the next argument to be processed. *optind* is external and is initialized to 1 before the first call to **getopt( )**. When all options have been processed (that is, up to the first non-option argument), **getopt( )** returns **EOF**. The special option "––" (two hyphens) may be used to delimit the end of the options; when it is encountered, **EOF** is returned and "––"' is skipped. This is useful in delimiting non-option arguments that begin with ''–'' (hyphen).

RETURN VALUES

getopt( ) prints an error message on the standard error and returns a "**?**" (question mark) when it encounters an option letter not included in *optstring* or no argument after an option that expects one. This error message may be disabled by setting **opterr** to 0. The value of the character that caused the error is in **optopt**.

EXAMPLES

The following code fragment shows how one might process the arguments for a command that can take the mutually exclusive options **a** and **b**, and the option **o**, which requires an argument:

```
#include <stdlib.h>
#include <stdio.h>
main (int argc, char **argv)
{
        int c;
        extern char *optarg;
        extern int optind;
        int aflg = 0;
        int bflg = 0;
        int errflg = 0;
```

```
        char *ofile = NULL;

        while ((c = getopt(argc, argv, "abo:")) != EOF)
                switch (c) {
                case 'a':
                        if (bflg)
                                errflg++;
                        else
                                aflg++;
                        break;
                case 'b':
                        if (aflg)
                                errflg++;
                        else
                                bflg++;
                        break;
                case 'o':
                        ofile = optarg;
                        (void)printf("ofile = %s\n", ofile);
                        break;
                case '?':
                        errflg++;
                }
        if (errflg) {
                (void)fprintf(stderr,
                        "usage: cmd [–a|–b] [–o <filename>] files...\n");
                exit (2);
        }
        for ( ; optind < argc; optind++)
                (void)printf("%s\n", argv[optind]);
        return 0;
}
```

SEE ALSO **intro**(1), **getopts**(1), **getopt**(3C), **getsubopt**(3C), **setlocale**(3C), **gettext**(3I)

NOTES If the application is linked with **–lintl**, then messages printed from this function are in the native language specified by the **LC_MESSAGES** locale category; see **setlocale**(3C).

The library routine **getopt( )** does not fully check for mandatory arguments. That is, given an option string **a:b** and the input **–a –b**, **getopt( )** assumes that **–b** is the mandatory argument to the **–a** option and not that **–a** is missing a mandatory argument.

It is a violation of the command syntax standard (see **intro**(1)) for options with arguments to be grouped with other options, as in **cmd** –**abo** *filename* , where **a** and **b** are options, **o** is an option that requires an argument, and *filename* is the argument to **o**. Although this syntax is permitted in the current implementation, it should not be used because it may not be supported in future releases. The correct syntax to use is:

cmd –**ab** –**o** *filename.*

NAME

getpagesize – get system page size

SYNOPSIS

```
#include <unistd.h>
int getpagesize(void);
```

DESCRIPTION

getpagesize( ) returns the number of bytes in a page. Page granularity is the granularity of many of the memory management calls.

The page size is a system page size and need not be the same as the underlying hardware page size.

SEE ALSO

pagesize(1), **brk**(2), **mmap**(2), **sysconf**(3C)

NAME

getpass – read a password

SYNOPSIS

```
#include <stdlib.h>
char *getpass(const char *prompt);
```

MT-LEVEL

Unsafe

DESCRIPTION

getpass( ) reads up to a newline or **EOF** from the file **/dev/tty**, after prompting on the standard error output with the null-terminated string *prompt* and disabling echoing. A pointer is returned to a null-terminated string of at most 8 characters. If **/dev/tty** cannot be opened, a null pointer is returned. An interrupt will terminate input and send an interrupt signal to the calling program before returning.

FILES

/dev/tty

NOTES

The return value points to static data whose content is overwritten by each call.

NAME

getpeername – get name of connected peer

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

int getpeername(int *s***, struct sockaddr** ∗*name***, int** ∗*namelen***);**

MT-LEVEL

Safe

DESCRIPTION

getpeername( ) returns the name of the peer connected to socket *s*. The **int** pointed to by the *namelen* parameter should be initialized to indicate the amount of space pointed to by *name*. On return it contains the actual size of the name returned (in bytes). The name is truncated if the buffer provided is too small.

RETURN VALUES

If successful, **getpeername( )** returns 0; otherwise it returns –1 and sets **errno** to indicate the error.

ERRORS

The call succeeds unless:

EBADF	The argument *s* is not a valid descriptor.
ENOMEM	There was insufficient user memory for the operation to complete.
ENOSR	There were insufficient STREAMS resources available for the operation to complete.
ENOTCONN	The socket is not connected.
ENOTSOCK	The argument *s* is not a socket.

SEE ALSO

accept(3N), **bind**(3N), **getsockname**(3N), **socket**(3N)

NAME

getpriority, setpriority – get / set scheduling priority for process, process group or user

SYNOPSIS

#include <sys/resource.h>

int getpriority(int *which***, id_t** *who***);**

int setpriority(int *which***, id_t** *who***, int** *prio***);**

DESCRIPTION

The scheduling priority of the process, process group, or user, as indicated by *which* and *who* is obtained with **getpriority( )** and set with **setpriority( )** The default priority is 0; lower priorities cause more favorable scheduling.

which is one of **PRIO_PROCESS**, **PRIO_PGRP**, or **PRIO_USER**, and *who* is interpreted relative to *which* (a process identifier for **PRIO_PROCESS**, process group identifier for **PRIO_PGRP**, and a user ID for **PRIO_USER**). A zero value of *who* denotes the current process, process group, or user.

getpriority( ) returns the highest priority (lowest numerical value) enjoyed by any of the specified processes. **setpriority( )** sets the priorities of all of the specified processes to the value specified by *prio*. If *prio* is less than –20, a value of –20 is used; if it is greater than 20, a value of 20 is used. Only the privileged user may lower priorities.

RETURN VALUES

Since **getpriority( )** can legitimately return the value –1, it is necessary to clear the external variable **errno** prior to the call, then check it afterward to determine if a –1 is an error or a legitimate value. The **setpriority( )** call returns 0 if there is no error, or –1 if there is.

ERRORS

getpriority( ) and **setpriority( )** may return one of the following errors:

EINVAL *which* was not one of **PRIO_PROCESS**, **PRIO_PGRP**, or **PRIO_USER**.

ESRCH No process was located using the *which* and *who* values specified.

In addition to the errors indicated above, **setpriority( )** may fail with one of the following errors returned:

EPERM A process was located, but one of the following is true:

- Neither its effective nor real user ID matched the effective user ID of the caller, and neither the effective nor the real user ID of the process executing the **setpriority( )** was the privileged user.
- The call to **getpriority( )** would have changed a process' priority to a value lower than its current value, and the effective user ID of the process executing the call was not that of the privileged user.

SEE ALSO

nice(1), **renice**(1), **fork**(2)

NOTES

It is not possible for the process executing **setpriority( )** to lower any other process down to its current priority, without requiring privileged user permissions.

NAME

getprotobyname, getprotobyname_r, getprotobynumber, getprotobynumber_r, getprotoent, getprotoent_r, setprotoent, endprotoent – get protocol entry

SYNOPSIS

cc [*flag* ...] *file* ... –**lsocket -lnsl** [*library* ...]

#include <netdb.h>

struct protoent ∗**getprotobyname(const char** ∗*name*);

struct protoent ∗**getprotobyname_r(const char** ∗*name,* **struct protoent** ∗*result,* **char** ∗*buffer,* **int** *buflen*);

struct protoent ∗**getprotobynumber(int** *proto*);

struct protoent ∗**getprotobynumber_r(int** *proto,* **struct protoent** ∗*result,* **char** ∗*buffer,* **int** *buflen*);

struct protoent ∗**getprotoent(void);**

struct protoent ∗**getprotoent_r(struct protoent** ∗*result,* **char** ∗*buffer,* **int** *buflen*);

int setprotoent(int *stayopen*);

int endprotoent(void);

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

These routines return a protocol entry. Two types of interfaces are supported: reentrant (**getprotobyname_r( )**, **getprotobynumber_r( )**, and **getprotoent_r( )**) and non-reentrant (**getprotobyname( )**, **getprotobynumber( )**, and **getprotoent( )**). The reentrant routines may be used in single-threaded applications and are safe for multi-threaded applications, making them the preferred interfaces.

The reentrant routines require additional parameters which are used to return results data. *result* is a pointer to a **struct protoent** structure and will be where the returned results will be stored. *buffer* is used as storage space for elements of the returned results. *buflen* is the size of *buffer* and should be large enough to contain all returned data. *buflen* must be at least 1024 bytes.

getprotobyname_r( ), **getprotobynumber_r( )**, and **getprotoent_r( )** each return a protocol entry.

The entry may come from one of the following sources: the protocols file (see **protocols**(4)), the NIS maps ‘‘protocols.byname’’ and ‘‘protocols.bynumber’’, and the NIS+ table ‘‘protocols’’. The sources and their lookup order are specified in the **/etc/nsswitch.conf** file (see **nsswitch.conf**(4) for details). Some name services such as NIS will return only one name for a host, whereas others such as NIS+ or DNS will return all aliases.

getprotobyname_r( ) and **getprotobynumber_r( )** sequentially search from the beginning of the file until a matching protocol name or protocol number is found, or until an EOF is encountered.

getprotobyname( ) and **getprotobynumber( )** have the same functionality as **getprotobyname_r( )** and **getprotobynumber_r( )** except that a static buffer is used to store returned results. These routines are unsafe in a multi-threaded application.

getprotoent_r( ) enumerates protocol entries: successive calls to **getprotoent_r( )** will return either successive protocol entries or NULL. Enumeration may not be supported by some sources. Note that if multiple threads call **getprotoent_r( )**, each will retrieve a subset of the protocol database.

getprotent( ) has the same functionality as **getprotent_r( )** except that a static buffer is used to store returned results. This routine is unsafe in a multi-threaded application.

setprotoent( ) "rewinds" to the beginning of the enumeration of protocol entries. If the *stayopen* flag is non-zero, resources such as open file descriptors are not deallocated after each call to **getprotobynumber_r( )** and **getprotobyname_r( )**. Calls to **getprotobyname_r( ) , getprotobyname( ) , getprotobynumber_r( )** and **getprotobynumber( )** may leave the enumeration in an indeterminate state, so **setprotoent( )** should be called before the first **getprotoent_r( )** or **getprotoent( )**. Note that **setprotoent( )** has process-wide scope, and ''rewinds'' the protocol entries for all threads calling **getprotoent_r( )** as well as main-thread calls to **getprotoent( )**.

endprotoent( ) may be called to indicate that protocol processing is complete; the system may then close any open protocols file, deallocate storage, and so forth. It is legitimate, but possibly less efficient, to call more protocol routines after **endprotoent( )**.

The internal representation of a protocol entry is a **protoent** structure defined in **<netdb.h>** with the following members:

```
char    *p_name;
char    **p_aliases;
int     p_proto;
```

RETURN VALUES

getprotobyname_r( ), getprotobyname( ), getprotobynumber_r( ), and **getprotobynumber( )** return a pointer to a **struct protoent** if they successfully locate the requested entry; otherwise they return NULL.

getprotoent_r( ) and **getprotoent( )** return a pointer to a **struct protoent** if they successfully enumerate an entry; otherwise they return NULL, indicating the end of the enumeration.

ERRORS

getprotobyname_r( ), getprotobynumber_r( ), and **getprotoent_r( )** will fail if the following is true:

ERANGE	length of the buffer supplied by caller is not large enough to store the result.

FILES

/etc/protocols
/etc/nsswitch.conf

SEE ALSO

intro(3), **nsswitch.conf**(4), **protocols**(4)

NOTES

Although **getprotobyname_r( )**, **getprotobynumber_r( )**, and **getprotoent_r( )** are not mentioned by POSIX.4a Draft 6, they were added to complete the functionality provided by similar thread-safe functions. These interfaces are subject to change to be compatible with the "spirit" of POSIX.4a when it is approved as a standard.

When compiling multithreaded applications, see **intro**(3), *Notes On Multithread Applications*, for information about the use of the **_REENTRANT** flag.

The routines **getprotobyname_r( )**, **getprotobynumber_r( )**, and **getprotoent_r( )** are reentrant and multi-thread safe. The reentrant interfaces can be used in single-threaded as well as multi-threaded applications and are therefore the preferred interfaces.

The routines **getprotobyname( )**, **getprotobyaddr( )**, and **getprotoent( )** use static storage, so returned data must be copied if it is to be saved. Because of their use of static storage for returned data, these routines are not safe for multi-threaded applications.

setprotoent( ) and **endprotoent( )** have process-wide scope, and are therefore not safe in multi-threaded applications.

Use of **getprotoent_r()** and **getprotoent()** is discouraged; enumeration is well-defined for the protocols file and is supported (albeit inefficiently) for NIS and NIS+, but in general may not be well-defined. The semantics of enumeration are discussed in **nsswitch.conf**(4).

BUGS

Only the Internet protocols are currently understood.

Programs that call **getprotobyname_r( )** or **getprotobynumber_r( )** routines cannot be linked statically since the implementation of these routines requires dynamic linker functionality to access shared objects at run time.

NAME

getpublickey, getsecretkey, publickey – retrieve public or secret key

SYNOPSIS

```
#include <rpc/rpc.h>
#include <rpc/key_prot.h>

int getpublickey(const char netname[MAXNETNAMELEN],
        char publickey[HEXKEYBYTES+1]);

int getsecretkey(const char netname[MAXNETNAMELEN],
        char secretkey[HEXKEYBYTES+1], const char *passwd);
```

MT-LEVEL

Safe

DESCRIPTION

getpublickey( ) and **getsecretkey( )** get public and secret keys for *netname*. The key may come from one of the following sources: the **/etc/publickey** file (see **publickey**(4)) or the NIS map ''publickey.byname'' or the NIS+ table ''cred.org_dir''. The sources and their lookup order are specified in the **/etc/nsswitch.conf** file (see **nsswitch.conf**(4)).

getsecretkey( ) has an extra argument, *passwd*, used to decrypt the encrypted secret key stored in the database.

RETURN VALUES

Both routines return **1** if they are successful in finding the key, **0** otherwise. The keys are returned as NULL-terminated, hexadecimal strings. If the password supplied to **getsecretkey( )** fails to decrypt the secret key, the routine will return 1 but the *secretkey* [0] will be set to NULL.

SEE ALSO

secure_rpc(3N), **nsswitch.conf**(4), **publickey**(4)

WARNINGS

If **getpublickey( )** gets the public key from any source other than NIS+, all authenticated NIS+ operations may fail. To ensure that this does not happen, edit the **nsswitch.conf**(4) file to make sure that the public key is obtained from NIS+.

NAME

getpw – get passwd entry from UID

SYNOPSIS

```
#include <stdlib.h>
int getpw(uid_t uid, char *buf);
```

MT-LEVEL

Safe

DESCRIPTION

getpw( ) searches the user data base for a user id number that equals *uid*, copies the line of the password file in which *uid* was found into the array pointed to by *buf*, and returns 0. **getpw( )** returns non-zero if *uid* cannot be found.

This routine is included only for compatibility with prior systems and should not be used; see **getpwnam**(3C) for routines to use instead.

RETURN VALUES

getpw( ) returns non-zero on error.

FILES

/etc/passwd

SEE ALSO

getpwnam(3C), **passwd**(4)

NOTES

If the **/etc/passwd** and the **/etc/group** files have the ‘‘+’’ for the NIS entry, then **getpwent( )** and **getgwent( )** will not return **NULL** when the end of file is reached.

NAME

getpwnam, getpwnam_r, getpwent, getpwent_r, getpwuid, getpwuid_r, setpwent, endpwent, fgetpwent, fgetpwent_r – get password entry

SYNOPSIS

```
#include <pwd.h>

struct passwd *getpwnam(const char *name);

struct passwd *getpwnam_r(const char *name, struct passwd *result,
        char *buffer, int buflen);

struct passwd *getpwent(void);

struct passwd *getpwent_r(struct passwd *result, char *buffer, int buflen);

struct passwd *getpwuid(uid_t uid);

struct passwd *getpwuid_r(uid_t uid, struct passwd * result, char *buffer, int buflen);

void setpwent(void);

void endpwent(void);

struct passwd *fgetpwent(FILE *f);

struct passwd *fgetpwent_r(FILE *f, struct passwd *result, char *buffer, int buflen);
```

POSIX

```
cc [ flag... ] file ... –D_POSIX_PTHREAD_SEMANTICS [ library... ]

int getpwnam_r(const char *name,struct passwd *pwd, char *buffer, size_t bufsize
        struct passwd **result);

int getpwuid_r(uid_t uid, struct passwd *pwd, char *buffer, size_t bufsize
        struct passwd **result);
```

MT-LEVEL

See the subsection “Reentrant Interfaces” in the **DESCRIPTION** section of this page.

DESCRIPTION

These functions are used to obtain password entries. Entries can come from any of the sources for **passwd** specified in the **/etc/nsswitch.conf** file (see **nsswitch.conf**(4)).

getpwnam( ) searches for a password entry with the login name specified by the character string parameter *name*.

getpwuid( ) searches for a password entry with the (numeric) user id specified by the parameter *uid*.

The functions **setpwent( )**, **getpwent( )**, and **endpwent( )** are used to enumerate password entries from the database. **setpwent( )** sets (or resets) the enumeration to the beginning of the set of password entries. This function should be called before the first call to **getpwent( )**. Calls to **getpwnam( )** and **getpwuid( )** leave the enumeration position in an indeterminate state. Successive calls to **getpwent( )** return either successive entries or NULL, indicating the end of the enumeration.

endpwent( ) may be called to indicate that the caller expects to do no further password retrieval operations; the system may then close the password file, deallocate resources it was using, and so forth. It is still allowed, but possibly less efficient, for the process to call more password functions after calling **endpwent( )**.

fgetpwent(), unlike the other functions above, does not use **nsswitch.conf**; it reads and parses the next line from the stream *f*, which is assumed to have the format of the **passwd** file. See **passwd**(4).

Reentrant Interfaces

The functions **getpwnam()**, **getpwuid()**, **getpwent()**, and **fgetpwent()** use static storage that is re-used in each call, making these routines unsafe for use in multithreaded applications.

The parallel functions **getpwnam_r()**, **getpwuid_r()**, **getpwent_r()**, and **fgetpwent_r()** provide reentrant interfaces for these operations.

Each reentrant interface performs the same operation as its non-reentrant counterpart, named by removing the "_r" suffix. The reentrant interfaces, however, use buffers supplied by the caller to store returned results, and are safe for use in both single-threaded and multithreaded applications.

Each reentrant interface takes the same parameters as its non-reentrant counterpart, as well as the following additional parameters. The parameter *result* (or *pwd* for the POSIX versions) must be a pointer to a **struct passwd** structure allocated by the caller. On successful completion, the function returns the password entry in this structure. The parameter *buffer* is a pointer to a buffer supplied by the caller, used as storage space for the password data. All of the pointers within the returned **struct passwd** *result* (or *grp*) point to data stored within this buffer; see **RETURN VALUES**. The buffer must be large enough to hold all the data associated with the password entry. The parameter *buflen* (or *bufsize* for the POSIX versions) should give the size in bytes of *buffer*. The POSIX versions place a pointer to the modified *grp* structure in the *result* parameter, instead of returning a pointer to this structure.

For enumeration in multithreaded applications, the position within the enumeration is a process-wide property shared by all threads. **setpwent()** may be used in a multithreaded application but resets the enumeration position for all threads. If multiple threads interleave calls to **getpwent_r()**, the threads will enumerate disjoint subsets of the password database.

Like their non-reentrant counterparts, **getpwnam_r()** and **getpwuid_r()** leave the enumeration position in an indeterminate state.

RETURN VALUES

Password entries are represented by the **struct passwd** structure defined in **<pwd.h>**:

```
struct passwd {
  char *pw_name;    /* user's login name */
  char *pw_passwd;  /* no longer used */
  uid_t pw_uid;     /* user's uid */
  gid_t pw_gid;     /* user's gid */
  char *pw_age;     /* not used */
  char *pw_comment; /* not used */
  char *pw_gecos;   /* typically user's full name */
  char *pw_dir;     /* user's home dir */
  char *pw_shell;   /* user's login shell */
};
```

The functions **getpwnam()**, **getpwnam_r()**, **getpwuid()**, and **getpwuid_r()** each return a pointer to a **struct passwd** if they successfully locate the requested entry; otherwise they return NULL. The POSIX functions **getpwnam_r()** and **getpwuid_r()** return zero upon success, or the error number in case of failure.

The functions **getpwent()**, **getpwent_r()**, **fgetpwent()**, and **fgetpwent_r()** each return a pointer to a **struct passwd** if they successfully enumerate an entry; otherwise they return NULL, indicating the end of the enumeration.

The functions **getpwnam()**, **getpwuid()**, **getpwent()**, and **fgetpwent()** use static storage, so returned data must be copied before a subsequent call to any of these functions if the data is to be saved.

When the pointer returned by the reentrant functions **getpwnam_r()**, **getpwuid_r()**, **getpwent_r()**, and **fgetpwent_r()** is non-NULL, it is always equal to the *result* pointer that was supplied by the caller.

ERRORS

The reentrant functions **getpwnam_r()**, **getpwuid_r()**, **getpwent_r()**, and **fgetpwent_r()** will return NULL and set **errno** to **ERANGE** (or in the case of POSIX functions **getpwnam_r()** and **getpwuid_r()** return the **ERANGE** error) if the length of the buffer supplied by caller is not large enough to store the result. See **Intro**(2) for the proper usage and interpretation of **errno** in multithreaded applications.

FILES

/etc/passwd
/etc/shadow
/etc/nsswitch.conf

SEE ALSO

nispasswd(1), **passwd**(1), **yppasswd**(1), **Intro**(2), **Intro**(3), **cuserid**(3S), **getgrnam**(3C), **getlogin**(3C), **getspnam**(3C), **nsswitch.conf**(4), **passwd**(4), **shadow**(4)

NOTES

The **pw_passwd** field in the **passwd** structure should not be used as the encrypted password for the user; use **getspnam()** or **getspnam_r()** instead. See **getspnam**(3C).

Programs that use the interfaces described in this manual page cannot be linked statically since, the implementations of these functions employ dynamic loading and linking of shared objects at run time.

When compiling multithreaded applications, see **Intro**(3), *Notes On Multithread Applications*, for information about the use of the **_REENTRANT** flag.

Use of the enumeration interfaces **getpwent()** and **getpwent_r()** is discouraged; enumeration is supported for the passwd file, NIS, and NIS+, but in general is not efficient and may not be supported for all database sources. The semantics of enumeration are discussed further in **nsswitch.conf**(4).

Previous releases allowed the use of '+' and '-' entries in **/etc/passwd** to selectively include and exclude NIS entries. The primary usage of these '+/-' entries is superseded by the name service switch, so the '+/-' form may not be supported in future releases.

If required, the '+/-' functionality can still be obtained for NIS by specifying **compat** as the source for **passwd**.

If the '+/-' functionality is required in conjunction with NIS+, specify both **compat** as the source for **passwd** and **nisplus** as the source for the pseudo-database **passwd_compat**. See **passwd**(4), **shadow**(4), and **nsswitch.conf**(4) for details.

If the '+/-' is used, both **/etc/shadow** and **/etc/passwd** should have the same '+' and '-' entries to ensure consistency between the password and shadow databases.

If a password entry from any of the sources contains an empty *uid* or *gid* field, that entry will be ignored by the files, NIS, and NIS+ name service switch backends. This will cause the user to appear unknown to the system.

If a password entry contains an empty *gecos*, *home directory*, or *shell* field, **getpwnam()** and **getpwnam_r()** return a pointer to a null string in the respective field of the **passwd** structure.

If the shell field is empty, **login**(1) automatically assigns the default shell. See **login**(1).

Reentrant interfaces **getpwnam_r()**, **getpwent_r()**, **getpwuid_r()**, and **fgetpwent_r()** are as specified in POSIX 1003.1c Draft #10.

NAME

getrpcbyname, getrpcbyname_r, getrpcbynumber, getrpcbynumber_r, getrpcent, getrpcent_r, setrpcent, endrpcent – get RPC entry

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <rpc/rpcent.h>

struct rpcent ∗**getrpcbyname(const char** ∗ *name*);

struct rpcent ∗**getrpcbyname_r(const char** ∗ *name,* **struct rpcent** ∗*result,* **char** ∗*buffer,* **int** *buflen*);

struct rpcent ∗**getrpcbynumber(const int** *number*);

struct rpcent ∗**getrpcbynumber_r(const int** *number,* **struct rpcent** ∗*result,* **char** ∗*buffer,* **int** *buflen*);

struct rpcent ∗**getrpcent(void);**

struct rpcent ∗**getrpcent_r(struct rpcent** ∗*result,* **char** ∗*buffer,* **int** *buflen*);

void setrpcent(const int *stayopen*);

void endrpcent(void);

MT-LEVEL

See the subsection “Reentrant Interfaces” in the **DESCRIPTION** section of this page.

DESCRIPTION

These functions are used to obtain entries for RPC (Remote Procedure Call) services. An entry may come from any of the sources for **rpc** specified in the **/etc/nsswitch.conf** file (see **nsswitch.conf**(4)).

getrpcbyname( ) searches for an entry with the RPC service name specified by the parameter *name*.

getrpcbynumber( ) searches for an entry with the RPC program number *number*.

The functions **setrpcent( )**, **getrpcent( )**, and **endrpcent( )** are used to enumerate RPC entries from the database.

setrpcent( ) sets (or resets) the enumeration to the beginning of the set of RPC entries. This function should be called before the first call to **getrpcent( )**. Calls to **getrpcbyname( )** and **getrpcbynumber( )** leave the enumeration position in an indeterminate state. If the *stayopen* flag is non-zero, the system may keep allocated resources such as open file descriptors until a subsequent call to **endrpcent( )**.

Successive calls to **getrpcent( )** return either successive entries or NULL, indicating the end of the enumeration.

endrpcent( ) may be called to indicate that the caller expects to do no further RPC entry retrieval operations; the system may then deallocate resources it was using. It is still allowed, but possibly less efficient, for the process to call more RPC entry retrieval functions after calling **endrpcent( )**.

Reentrant Interfaces

The functions **getrpcbyname()**, **getrpcbynumber()**, and **getrpcent()** use static storage that is re-used in each call, making these routines unsafe for use in multithreaded applications.

The functions:

getrpcbyname_r(),
getrpcbynumber_r(),
and
getrpcent_r()
provide reentrant interfaces for these operations.

Each reentrant interface performs the same operation as its non-reentrant counterpart, named by removing the ''**_r**'' suffix. The reentrant interfaces, however, use buffers supplied by the caller to store returned results, and are safe for use in both single-threaded and multithreaded applications.

Each reentrant interface takes the same parameters as its non-reentrant counterpart, as well as the following additional parameters. The parameter *result* must be a pointer to a **struct rpcent** structure allocated by the caller. On successful completion, the function returns the RPC entry in this structure. The parameter *buffer* must be a pointer to a buffer supplied by the caller. This buffer is used as storage space for the RPC entry data. All of the pointers within the returned **struct rpcent** *result* point to data stored within this buffer (see **RETURN VALUES**). The buffer must be large enough to hold all of the data associated with the RPC entry. The parameter *buflen* should give the size in bytes of the buffer indicated by *buffer*.

For enumeration in multithreaded applications, the position within the enumeration is a process-wide property shared by all threads. **setrpcent()** may be used in a multithreaded application but resets the enumeration position for all threads. If multiple threads interleave calls to **getrpcent_r()**, the threads will enumerate disjoint subsets of the RPC entry database.

Like their non-reentrant counterparts, **getrpcbyname_r()** and **getrpcbynumber_r()** leave the enumeration position in an indeterminate state.

RETURN VALUES

RPC entries are represented by the **struct rpcent** structure defined in **<rpc/rpcent.h>**:

```
struct rpcent {
        char *r_name;       /* name of this rpc service */
        char **r_aliases;   /* zero-terminated list of
                               alternate names */
        long r_number;      /* rpc program number */
};
```

The functions **getrpcbyname()**, **getrpcbyname_r()**, **getrpcbynumber()**, and **getrpcbynumber_r()** each return a pointer to a **struct rpcent** if they successfully locate the requested entry; otherwise they return NULL.

The functions **getrpcent()** and **getrpcent_r()** each return a pointer to a **struct rpcent** if they successfully enumerate an entry; otherwise they return NULL, indicating the end of the enumeration.

The functions **getrpcbyname( )**, **getrpcbynumber( )**, and **getrpcent( )** use static storage, so returned data must be copied before a subsequent call to any of these functions if the data is to be saved.

When the pointer returned by the reentrant functions **getrpcbyname_r( )**, **getrpcbynumber_r( )**, and **getrpcent_r( )** is non-NULL, it is always equal to the *result* pointer that was supplied by the caller.

ERRORS

The reentrant functions **getrpcyname_r( )**, **getrpcbynumber_r( )** and **getrpcent_r( )** will return NULL and set **errno** to **ERANGE** if the length of the buffer supplied by caller is not large enough to store the result. See **intro**(2) for the proper usage and interpretation of **errno** in multithreaded applications.

FILES

/etc/rpc
/etc/nsswitch.conf

SEE ALSO

rpcinfo(1M), **rpc**(3N), **nsswitch.conf**(4), **rpc**(4)

WARNINGS

The reentrant interfaces **getrpcbyname_r( )**, **getrpcbynumber_r( )**, and **getrpcent_r( )** are included in this release on an uncommitted basis only, and are subject to change or removal in future minor releases.

NOTES

Programs that use the interfaces described in this manual page cannot be linked statically since the implementations of these functions employ dynamic loading and linking of shared objects at run time.

When compiling multithreaded applications, see **Intro**(3), *Notes On Multithread Applications*, for information about the use of the **_REENTRANT** flag.

Use of the enumeration interfaces **getrpcent( )** and **getrpcent_r( )** is discouraged; enumeration may not be supported for all database sources. The semantics of enumeration are discussed further in **nsswitch.conf**(4).

NAME

getrusage – get information about resource utilization

SYNOPSIS

```
#include <sys/resource.h>

int getrusage(int who, struct rusage *rusage);
```

DESCRIPTION

getrusage( ) returns information about the resources utilized by the current process, or all its terminated child processes. The interpretation for some values reported, such as **ru_idrss**, are dependent on the clock tick interval. This interval is an implementation dependent value.

The **who** parameter is one of **RUSAGE_SELF** or **RUSAGE_CHILDREN**. The buffer to which *rusage* points will be filled in with a structure with the following members:

```
struct timeval ru_utime;    /* user time used */
struct timeval ru_stime;    /* system time used */
int    ru_maxrss;           /* maximum resident set size */
int    ru_idrss;            /* integral resident set size */
int    ru_minflt;           /* page faults not requiring physical I/O */
int    ru_majflt;           /* page faults requiring physical I/O */
int    ru_nswap;            /* swaps */
int    ru_inblock;          /* block input operations */
int    ru_oublock;          /* block output operations */
int    ru_msgsnd;           /* messages sent */
int    ru_msgrcv;           /* messages received */
int    ru_nsignals;         /* signals received */
int    ru_nvcsw;            /* voluntary context switches */
int    ru_nivcsw;           /* involuntary context switches */
```

The fields are interpreted as follows:

ru_utime The total amount of time spent executing in user mode. Time is given in seconds and microseconds.

ru_stime The total amount of time spent executing in system mode. Time is given in seconds and microseconds.

ru_maxrss The maximum resident set size. Size is given in pages (the size of a page, in bytes, is given by the **getpagesize**(3C) function). See the **NOTES** section of this page.

ru_idrss An “integral” value indicating the amount of memory in use by a process while the process is running. This value is the sum of the resident set sizes of the process running when a clock tick occurs. The value is given in pages times clock ticks. It does not take sharing into account. See the **NOTES** section of this page.

ru_minflt The number of page faults serviced which did not require any physical I/O activity. See the **NOTES** section of this page.

ru_majflt The number of page faults serviced which required physical I/O activity. This could include page ahead operations by the kernel. See the **NOTES** section of this page.

ru_nswap The number of times a process was swapped out of main memory.

ru_inblock The number of times the file system had to perform input in servicing a **read**(2) request.

ru_oublock The number of times the file system had to perform output in servicing a **write**(2) request.

ru_msgsnd The number of messages sent over sockets.

ru_msgrcv The number of messages received from sockets.

ru_nsignals The number of signals delivered.

ru_nvcsw The number of times a context switch resulted due to a process voluntarily giving up the processor before its time slice was completed (usually to await availability of a resource).

ru_nivcsw The number of times a context switch resulted due to a higher priority process becoming runnable or because the current process exceeded its time slice.

RETURN VALUES

If successful, the value of the appropriate structure is filled in, and **0** is returned. If the call fails, –**1** is returned.

ERRORS

getrusage() will fail if:

EFAULT The address specified by the *rusage* argument is not in a valid portion of the process's address space.

EINVAL The **who** parameter is not a valid value.

SEE ALSO

sar(1M), **read**(2), **times**(2), **wait**(2), **write**(2), **getpagesize**(3C), **gettimeofday**(3C)

NOTES

Only the *timeval* fields of **struct rusage** are supported in this implementation.

The numbers **ru_inblock** and **ru_oublock** account only for real I/O, and are approximate measures at best. Data supplied by the cache mechanism is charged only to the first process to read and the last process to write the data.

The way resident set size is calculated is an approximation, and could misrepresent the true resident set size.

Page faults can be generated from a variety of sources and for a variety of reasons. The customary cause for a page fault is a direct reference by the program to a page which is not in memory. Now, however, the kernel can generate page faults on behalf of the user, for example, servicing **read**(2) and **write**(2) functions. Also, a page fault can be caused by an absent hardware translation to a page, even though the page is in physical memory.

In addition to hardware detected page faults, the kernel may cause pseudo page faults in order to perform some housekeeping. For example, the kernel may generate page faults, even if the pages exist in physical memory, in order to lock down pages involved in a raw I/O request.

By definition, major page faults require physical I/O, while minor page faults do not require physical I/O. For example, reclaiming the page from the free list would avoid I/O and generate a minor page fault. More commonly, minor page faults occur during process startup as references to pages which are already in memory. For example, if an address space faults on some “hot” executable or shared library, this results in a minor page fault for the address space. Also, any one doing a **read**(2) or **write**(2) to something that is in the page cache will get a minor page fault(s) as well.

There is no way to obtain information about a child process which has not yet terminated.

NAME

gets, fgets – get a string from a stream

SYNOPSIS

```
#include <stdio.h>
char *gets(char *s);
char *fgets(char *s, int n, FILE *stream);
```

MT-LEVEL

MT-Safe

DESCRIPTION

gets( ) reads characters from the standard input stream (see **intro**(3)), **stdin**, into the array pointed to by *s*, until a newline character is read or an end-of-file condition is encountered. The newline character is discarded and the string is terminated with a null character.

fgets( ) reads characters from the *stream* into the array pointed to by *s*, until *n*−1 characters are read, or a newline character is read and transferred to *s*, or an end-of-file condition is encountered. The string is then terminated with a null character.

When using **gets( )**, if the length of an input line exceeds the size of *s*, indeterminate behavior may result. For this reason, it is strongly recommended that **gets( )** be avoided in favor of **fgets( )**.

RETURN VALUES

If end-of-file is encountered and no characters have been read, no characters are transferred to *s* and a null pointer is returned. If a read error occurs, such as trying to use these functions on a file that has not been opened for reading, a null pointer is returned and the error indicator for the stream is set. If end-of-file is encountered, the **EOF** indicator for the stream is set. Otherwise *s* is returned.

SEE ALSO

lseek(2), **read**(2), **ferror**(3S), **fopen**(3S), **fread**(3S), **getc**(3S), **scanf**(3S), **stdio**(3S), **ungetc**(3S)

NAME

getservbyname, getservbyname_r, getservbyport, getservbyport_r, getservent, getservent_r, setservent, endservent – get service entry

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket -lnsl** [*library* ...]

#include <netdb.h>

struct servent ∗**getservbyname(const char** ∗*name*, **const char** ∗*proto*);

struct servent ∗**getservbyname_r(const char** ∗*name*, **const char** ∗*proto*, **struct servent** ∗*result*, **char** ∗*buffer*, **int** *buflen*);

struct servent ∗**getservbyport(int** *port*, **const char** ∗*proto*);

struct servent ∗**getservbyport_r(int** *port*, **const char** ∗*proto*, **struct servent** ∗*result*, **char** ∗*buffer*, **int** *buflen*);

struct servent ∗**getservent(void);**

struct servent ∗**getservent_r(struct servent** ∗*result*, **char** ∗*buffer*, **int** *buflen*);

int setservent(int *stayopen*);

int endservent(void);

MT-LEVEL

See the subsection “Reentrant Interfaces” in the **DESCRIPTION** section of this man page.

DESCRIPTION

These functions are used to obtain entries for Internet services. An entry may come from any of the sources for **services** specified in the **/etc/nsswitch.conf** file. See **nsswitch.conf**(4).

getservbyname( ) and **getservbyport( )** sequentially search from the beginning of the file until a matching protocol name or port number is found, or until end-of-file is encountered. If a protocol name is also supplied (non-NULL), searches must also match the protocol.

getservbyname( ) searches for an entry with the Internet service name specified by the parameter *name*.

getservbyport( ) searches for an entry with the Internet port number *port*.

The string *proto* is used by both **getservbyname( )** and **getservbyport( )** to restrict the search to entries with the specified protocol. If *proto* is **NULL,** entries with any protocol may be returned.

The functions **setservent( )**, **getservent( )**, and **endservent( )** are used to enumerate entries from the services database.

setservent( ) sets (or resets) the enumeration to the beginning of the set of service entries. This function should be called before the first call to **getservent( )**. Calls to the functions **getservbyname( )** and **getservbyport( )** leave the enumeration position in an indeterminate state. If the *stayopen* flag is non-zero, the system may keep allocated resources such as open file descriptors until a subsequent call to **endservent( )**.

getservent() reads the next line of the file, opening the file if necessary. **getservent()** opens and rewinds the file. If the *stayopen* flag is non-zero, the net data base will not be closed after each call to **getservent()** (either directly, or indirectly through one of the other "getserv" calls).

Successive calls to **getservent()** return either successive entries or **NULL,** indicating the end of the enumeration.

endservent() closes the file. **endservent()** may be called to indicate that the caller expects to do no further service entry retrieval operations; the system may then deallocate resources it was using. It is still allowed, but possibly less efficient, for the process to call more service entry retrieval functions after calling **endservent()**.

Reentrant Interfaces

The functions **getservbyname()**, **getservbyport()**, and **getservent()** use static storage that is re-used in each call, making these functions unsafe for use in multithreaded applications.

The functions:

getservbyname_r(),
getservbyport_r(),

and

getservent_r()

provide reentrant interfaces for these operations.

Each reentrant interface performs the same operation as its non-reentrant counterpart, named by removing the "**_r**" suffix. The reentrant interfaces, however, use buffers supplied by the caller to store returned results, and are safe for use in both single-threaded and multithreaded applications.

Each reentrant interface takes the same parameters as its non-reentrant counterpart, as well as the following additional parameters. The parameter *result* must be a pointer to a **struct servent** structure allocated by the caller. On successful completion, the function returns the service entry in this structure. The parameter *buffer* must be a pointer to a buffer supplied by the caller. This buffer is used as storage space for the service entry data. All of the pointers within the returned **struct servent** *result* point to data stored within this buffer. See the **RETURN VALUES** section of this man page. The buffer must be large enough to hold all of the data associated with the service entry. The parameter *buflen* should give the size in bytes of the buffer indicated by *buffer*.

For enumeration in multithreaded applications, the position within the enumeration is a process-wide property shared by all threads. **setservent()** may be used in a multithreaded application but resets the enumeration position for all threads. If multiple threads interleave calls to **getservent_r()**, the threads will enumerate disjoint subsets of the service database.

Like their non-reentrant counterparts, **getservbyname_r()** and **getservbyport_r()** leave the enumeration position in an indeterminate state.

RETURN VALUES

Service entries are represented by the **struct servent** structure defined in **<netdb.h>**:

```
struct servent {
        char    *s_name;        /* official name of service */
        char    **s_aliases;    /* alias list */
        int     s_port;         /* port service resides at */
        char    *s_proto;       /* protocol to use */
};
```

The members of this structure are:

s_name	The official name of the service.
s_aliases	A zero terminated list of alternate names for the service.
s_port	The port number at which the service resides. Port numbers are returned in network byte order.
s_proto	The name of the protocol to use when contacting the service.

The functions **getservbyname()**, **getservbyname_r()**, **getservbyport()**, and **getservbyport_r()** each return a pointer to a **struct servent** if they successfully locate the requested entry; otherwise they return **NULL.**

The functions **getservent()** and **getservent_r()** each return a pointer to a **struct servent** if they successfully enumerate an entry; otherwise they return **NULL,** indicating the end of the enumeration.

The functions **getservbyname()**, **getservbyport()**, and **getservent()** use static storage, so returned data must be copied before a subsequent call to any of these functions if the data is to be saved.

When the pointer returned by the reentrant functions **getservbyname_r()**, **getservbyport_r()**, and **getservent_r()** is non-null, it is always equal to the *result* pointer that was supplied by the caller.

ERRORS

The reentrant functions **getservbyname_r()**, **getservbyport_r()** and **getservent_r()** will return **NULL** and set **errno** to **ERANGE** if the length of the buffer supplied by caller is not large enough to store the result. See **intro**(2) for the proper usage and interpretation of **errno** in multithreaded applications.

FILES

/etc/services	Internet network services
/etc/netconfig	network configuration file
/etc/nsswitch.conf	configuration file for the name-service switch

SEE ALSO

intro(2), **intro**(3), **netdir**(3N), **netconfig**(4), **nsswitch.conf**(4), **services**(4)

WARNINGS

The reentrant interfaces **getservbyname_r()**, **getservbyport_r()**, and **getservent_r()** are included in this release on an uncommitted basis only, and are subject to change or removal in future minor releases.

NOTES

The functions that return **struct servent** return the least significant 16-bits of the *s_port* field in *network byte order*. **getservbyport( )** and **getservbyport_r( )** also expect the input parameter *port* in the *network byte order* . See **htons**(3N) for more details on converting between host and network byte orders.

Programs that use the interfaces described in this manual page cannot be linked statically since the implementations of these functions employ dynamic loading and linking of shared objects at run time.

In order to ensure that they all return consistent results, **getservbyname( )**, **getservbyname_r( )**, and **netdir_getbyname( )** are implemented in terms of the same internal library function. This function obtains the system-wide source lookup policy based on the **inet** family entries in **netconfig**(4) and the **services:** entry in **nsswitch.conf**(4). Similarly, **getservbyport( )**, **getservbyport_r( )**, and **netdir_getbyaddr( )** are implemented in terms of the same internal library function. If the **inet** family entries in **netconfig**(4) have a ''-'' in the last column for nametoaddr libraries, then the entry for **services** in **nsswitch.conf** will be used; otherwise the nametoaddr libraries in that column will be used, and **nsswitch.conf** will not be consulted.

There is no analogue of **getservent( )** and **getservent_r( )** in the netdir functions, so these enumeration functions go straight to the **services** entry in **nsswitch.conf**. Thus enumeration may return results from a different source than that used by **getservbyname( )**, **getservbyname_r( )**, **getservbyport( )**, and **getservbyport_r( )**.

When compiling multithreaded applications, see **intro**(3), *Notes On Multithread Applications*, for information about the use of the **_REENTRANT** flag.

Use of the enumeration interfaces **getservent( )** and **getservent_r( )** is discouraged; enumeration may not be supported for all database sources. The semantics of enumeration are discussed further in **nsswitch.conf**(4).

NAME

getsockname – get socket name

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>

int getsockname(int s, struct sockaddr *name, int *namelen);
```

MT-LEVEL

Safe

DESCRIPTION

getsockname( ) returns the current *name* for socket *s*. The *namelen* parameter should be initialized to indicate the amount of space pointed to by *name*. On return it contains the actual size in bytes of the name returned.

RETURN VALUES

If successful, **getsockname( )** returns 0; otherwise it returns –1 and sets *errno* to indicate the error.

ERRORS

The call succeeds unless:

EBADF	The argument *s* is not a valid file descriptor.
ENOMEM	There was insufficient memory available for the operation to complete.
ENOSR	There were insufficient STREAMS resources available for the operation to complete.
ENOTSOCK	The argument *s* is not a socket.

SEE ALSO

bind(3N), **getpeername**(3N), **socket**(3N)

NAME

getsockopt, setsockopt – get and set options on sockets

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>

int getsockopt(int s, int level, int optname, char *optval, int *optlen);
int setsockopt(int s, int level, int optname, const char *optval, int optlen);
```

MT-LEVEL

Safe

DESCRIPTION

getsockopt( ) and **setsockopt( )** manipulate options associated with a socket. Options may exist at multiple protocol levels; they are always present at the uppermost "socket" level.

When manipulating socket options, the level at which the option resides and the name of the option must be specified. To manipulate options at the "socket" level, *level* is specified as **SOL_SOCKET**. To manipulate options at any other level, *level* is the protocol number of the protocol that controls the option. For example, to indicate that an option is to be interpreted by the TCP protocol, *level* is set to the TCP protocol number (see **getprotobyname**(3N)).

The parameters *optval* and *optlen* are used to access option values for **setsockopt( )**. For **getsockopt( ),** they identify a buffer in which the value(s) for the requested option(s) are to be returned. For **getsockopt( )**, *optlen* is a value-result parameter, initially containing the size of the buffer pointed to by *optval*, and modified on return to indicate the actual size of the value returned. Use a 0 *optval* if no option value is to be supplied or returned.

optname and any specified options are passed uninterpreted to the appropriate protocol module for interpretation. The include file **<sys/socket.h>** contains definitions for the socket-level options described below. Options at other protocol levels vary in format and name.

Most socket-level options take an int for *optval*. For **setsockopt( )**, the *optval* parameter should be non-zero to enable a boolean option, or zero if the option is to be disabled. **SO_LINGER** uses a **struct linger** parameter that specifies the desired state of the option and the linger interval (see below). **struct linger** is defined in **<sys/socket.h>**. **struct linger** contains the following members:

l_onoff	on = 1 / off = 0
l_linger	linger time, in seconds

The following options are recognized at the socket level. Except as noted, each may be examined with **getsockopt( )** and set with **setsockopt( )**.

SO_DEBUG	enable / disable recording of debugging information
SO_REUSEADDR	enable / disable local address reuse
SO_KEEPALIVE	enable / disable keep connections alive
SO_DONTROUTE	enable / disable routing bypass for outgoing messages
SO_LINGER	linger on close if data is present

SO_BROADCAST	enable/disable permission to transmit broadcast messages
SO_OOBINLINE	enable/disable reception of out-of-band data in band
SO_SNDBUF	set buffer size for output
SO_RCVBUF	set buffer size for input
SO_TYPE	get the type of the socket (get only)
SO_ERROR	get and clear error on the socket (get only)

SO_DEBUG enables debugging in the underlying protocol modules. **SO_REUSEADDR** indicates that the rules used in validating addresses supplied in a **bind**(3N) call should allow reuse of local addresses. **SO_KEEPALIVE** enables the periodic transmission of messages on a connected socket. If the connected party fails to respond to these messages, the connection is considered broken and processes using the socket are notified using a **SIGPIPE** signal. **SO_DONTROUTE** indicates that outgoing messages should bypass the standard routing facilities. Instead, messages are directed to the appropriate network interface according to the network portion of the destination address.

SO_LINGER controls the action taken when unsent messages are queued on a socket and a **close**(2) is performed. If the socket promises reliable delivery of data and **SO_LINGER** is set, the system will block the process on the **close( )** attempt until it is able to transmit the data or until it decides it is unable to deliver the information (a timeout period, termed the linger interval, is specified in the **setsockopt( )** call when **SO_LINGER** is requested). If **SO_LINGER** is disabled and a **close( )** is issued, the system will process the **close( )** in a manner that allows the process to continue as quickly as possible.

The option **SO_BROADCAST** requests permission to send broadcast datagrams on the socket. With protocols that support out-of-band data, the **SO_OOBINLINE** option requests that out-of-band data be placed in the normal data input queue as received; it will then be accessible with **recv( )** or **read( )** calls without the **MSG_OOB** flag.

SO_SNDBUF and **SO_RCVBUF** are options that adjust the normal buffer sizes allocated for output and input buffers, respectively. The buffer size may be increased for high-volume connections or may be decreased to limit the possible backlog of incoming data. The Internet protocols place an absolute limit of 64 Kbytes on these values for **UDP** and **TCP** sockets.

Finally, **SO_TYPE** and **SO_ERROR** are options used only with **getsockopt( )**. **SO_TYPE** returns the type of the socket (for example, **SOCK_STREAM**). It is useful for servers that inherit sockets on startup. **SO_ERROR** returns any pending error on the socket and clears the error status. It may be used to check for asynchronous errors on connected datagram sockets or for other asynchronous errors.

RETURN VALUES

If successful, **getsockopt( )** returns **0**; otherwise, it returns **–1** and sets **errno** to indicate the error.

ERRORS

The call succeeds unless:

EBADF The argument *s* is not a valid file descriptor.

ENOMEM There was insufficient memory available for the operation to complete.

ENOPROTOOPT The option is unknown at the level indicated.

ENOSR There were insufficient STREAMS resources available for the operation to complete.

ENOTSOCK The argument *s* is not a socket.

SEE ALSO

close(2), **ioctl**(2), **bind**(3N), **getprotobyname**(3N), **socket**(3N)

NAME

getspnam, getspnam_r, getspent, getspent_r, setspent, endspent, fgetspent, fgetspent_r – get password entry

SYNOPSIS

```
#include <shadow.h>

struct spwd *getspnam(const char *name);

struct spwd *getspnam_r(const char *name, struct spwd *result, char *buffer,
        int buflen);

struct spwd *getspent(void);

struct spwd *getspent_r(struct spwd *result, char *buffer, int buflen);

void setspent(void);

void endspent(void);

struct spwd *fgetspent(FILE *fp);

struct spwd *fgetspent_r(FILE *fp, struct spwd *result, char *buffer, int buflen);
```

MT-LEVEL

See the subsection “Reentrant Interfaces” in the **DESCRIPTION** section of this page.

DESCRIPTION

These functions are used to obtain shadow password entries. An entry may come from any of the sources for **shadow** specified in the **/etc/nsswitch.conf** file (see **nsswitch.conf**(4)).

getspnam( ) searches for a shadow password entry with the login name specified by the character string parameter *name*.

The functions **setspent( )**, **getspent( )**, and **endspent( )** are used to enumerate shadow password entries from the database.

setspent( ) sets (or resets) the enumeration to the beginning of the set of shadow password entries. This function should be called before the first call to **getspent( )**. Calls to **getspnam( )** leave the enumeration position in an indeterminate state.

Successive calls to **getspent( )** return either successive entries or NULL, indicating the end of the enumeration.

endspent( ) may be called to indicate that the caller expects to do no further shadow password retrieval operations; the system may then close the shadow password file, deallocate resources it was using, and so forth. It is still allowed, but possibly less efficient, for the process to call more shadow password functions after calling **endspent( )**.

fgetspent( ), unlike the other functions above, does not use **nsswitch.conf**; it reads and parses the next line from the stream *f*, which is assumed to have the format of the **shadow** file (see **shadow**(4)).

Reentrant Interfaces

The functions **getspnam( )**, **getspent( )**, and **fgetspent( )** use static storage that is re-used in each call, making these routines unsafe for use in multithreaded applications.

The functions:

getspnam_r( ),
getspent_r( ),
and
fgetspent_r( )
provide reentrant interfaces for these operations.

Each reentrant interface performs the same operation as its non-reentrant counterpart, named by removing the ''**_r**'' suffix. The reentrant interfaces, however, use buffers supplied by the caller to store returned results, and are safe for use in both single-threaded and multithreaded applications.

Each reentrant interface takes the same parameters as its non-reentrant counterpart, as well as the following additional parameters. The parameter *result* must be a pointer to a **struct spwd** structure allocated by the caller. On successful completion, the function returns the shadow password entry in this structure. The parameter *buffer* must be a pointer to a buffer supplied by the caller. This buffer is used as storage space for the shadow password data. All of the pointers within the returned **struct spwd** *result* point to data stored within this buffer (see **RETURN VALUES**). The buffer must be large enough to hold all of the data associated with the shadow password entry. The parameter *buflen* should give the size in bytes of the buffer indicated by *buffer*.

For enumeration in multithreaded applications, the position within the enumeration is a process-wide property shared by all threads. **setspent( )** may be used in a multithreaded application but resets the enumeration position for all threads. If multiple threads interleave calls to **getspent_r( )**, the threads will enumerate disjoint subsets of the shadow password database.

Like its non-reentrant counterpart, **getspnam_r( )** leaves the enumeration position in an indeterminate state.

RETURN VALUES

Password entries are represented by the **struct spwd** structure defined in **<shadow.h>**:

```
struct spwd{
    char *sp_namp;      /* login name */
    char *sp_pwdp;      /* encrypted passwd */
    long sp_lstchg;     /* date of last change */
    long sp_min;        /* min days to passwd change */
    long sp_max;                /* max days to passwd change*/
    long sp_warn;       /* warning period */
    long sp_inact;      /* max days inactive */
    long sp_expire;     /* account expiry date */
    unsigned long sp_flag;  /* not used */
};
```

See **shadow**(4) for more information on the interpretation of this information.

The functions **getspnam( )**and **getspnam_r( )** each return a pointer to a **struct spwd** if they successfully locate the requested entry; otherwise they return NULL.

The functions **getspent( )**,**getspent_r( )**, **fgetspent( )**, and **fgetspent( )** each return a pointer to a **struct spwd** if they successfully enumerate an entry; otherwise they return NULL, indicating the end of the enumeration.

The functions **getspnam( )**, **getspent( )**, and **fgetspent( )** use static storage, so returned data must be copied before a subsequent call to any of these functions if the data is to be saved.

When the pointer returned by the reentrant functions **getspnam_r( )**, **getspent_r( )**, and **fgetspent_r( )** is non-NULL, it is always equal to the *result* pointer that was supplied by the caller.

ERRORS

The reentrant functions **getspnam_r( )**, **getspent_r( )**, and **fgetspent_r( )** will return NULL and set **errno** to **ERANGE** if the length of the buffer supplied by caller is not large enough to store the result. See **intro**(2) for the proper usage and interpretation of **errno** in multithreaded applications.

FILES

/etc/shadow
/etc/nsswitch.conf
/etc/passwd

SEE ALSO

nispasswd(1), **passwd**(1), **yppasswd**(1), **intro**(3) **getlogin**(3C), **getpwnam**(3C), **nsswitch.conf**(4), **passwd**(4), **shadow**(4)

WARNINGS

The reentrant interfaces **getspnam_r( )**, **getspent_r( )**, and **fgetspent_r( )** are included in this release on an uncommitted basis only, and are subject to change or removal in future minor releases.

NOTES

Programs that use the interfaces described in this manual page cannot be linked statically since the implementations of these functions employ dynamic loading and linking of shared objects at run time.

When compiling multithreaded applications, see **intro**(3), *Notes On Multithread Applications*, for information about the use of the **_REENTRANT** flag.

Use of the enumeration interfaces **getspent( )** and **getspent_r( )** is not recommended; enumeration is supported for the shadow file, NIS, and NIS+, but in general is not efficient and may not be supported for all database sources. The semantics of enumeration are discussed further in **nsswitch.conf**(4).

Access to shadow password information may be restricted in a manner depending on the database source being used. Access to the **/etc/shadow** file is generally restricted to processes running as the super-user (root). Other database sources may impose stronger or less stringent restrictions.

When NIS is used as the database source, the information for the shadow password entries is obtained from the ‘‘passwd.byname’’ map. This map stores only the information for the **sp_namp** and **sp_pwdp** fields of the **struct spwd** structure. Shadow password entries obtained from NIS will contain the value -1 in the remainder of the fields.

When NIS+ is used as the database source, and the caller lacks the permission needed to retrieve the encrypted password from the NIS+ ‘‘passwd.org_dir’’ table, the NIS+ service returns the string ‘‘*NP*’’ instead of the actual encrypted password string. The functions described on this page will then return the string ‘‘*NP*’’ to the caller as the value of the member **sp_pwdp** in the returned shadow password structure.

NAME

getsubopt – parse suboptions from a string

SYNOPSIS

```
#include <stdlib.h>
int getsubopt(char **optionp, const char * const *tokens, char **valuep);
```

MT-LEVEL

MT-Safe

DESCRIPTION

getsubopt( ) parses suboptions in a flag argument that was initially parsed by **getopt**(3C). These suboptions are separated by commas and may consist of either a single token or a token-value pair separated by an equal sign. Since commas delimit suboptions in the option string, they are not allowed to be part of the suboption or the value of a suboption. A command that uses this syntax is **mount**(1M), which allows the user to specify mount parameters with the **-o** option as follows:

```
mount –o rw,hard,bg,wsize=1024 speed:/usr /usr
```

In this example there are four suboptions: **rw**, **hard**, **bg**, and **wsize**, the last of which has an associated value of 1024.

getsubopt( ) takes the address of a pointer to the option string, a vector of possible tokens, and the address of a value string pointer. It returns the index of the token that matched the suboption in the input string or –1 if there was no match. If the option string at *optionp* contains only one subobtion, **getsubopt( )** updates *optionp* to point to the null character at the end of the string; otherwise it isolates the suboption by replacing the comma separator with a null character, and updates *optionp* to point to the start of the next suboption. If the suboption has an associated value, **getsubopt( )** updates *valuep* to point to the value's first character. Otherwise it sets *valuep* to **NULL**.

The token vector is organized as a series of pointers to null strings. The end of the token vector is identified by a null pointer.

When **getsubopt( )** returns, if *valuep* is not **NULL**, then the suboption processed included a value. The calling program may use this information to determine if the presence or lack of a value for this subobtion is an error.

Additionally, when **getsubopt( )** fails to match the suboption with the tokens in the *tokens* array, the calling program should decide if this is an error, or if the unrecognized option should be passed to another program.

RETURN VALUES

getsubopt( ) returns –1 when the token it is scanning is not in the token vector. The variable addressed by *valuep* contains a pointer to the first character of the token that was not recognized rather than a pointer to a value for that token.

The variable addressed by *optionp* points to the next option to be parsed, or a null character if there are no more options.

EXAMPLE

The following code fragment shows how to process options to the **mount**(1M) command using **getsubopt( )**.

```
#include <stdlib.h>

char *myopts[] = {
#define READONLY    0
                    "ro",
#define READWRITE   1
                    "rw",
#define WRITESIZE   2
                    "wsize",
#define READSIZE    3
                    "rsize",
                    NULL};

main(argc, argv)
        int  argc;
        char **argv;
{
        int sc, c, errflag;
        char *options, *value;
        extern char *optarg;
        extern int optind;
        .
        .
        .
        while((c = getopt(argc, argv, "abf:o:")) != -1) {
                switch (c) {
                case 'a': /* process a option */
                        break;
                case 'b': /* process b option */
                        break;
                case 'f':
                        ofile = optarg;
                        break;
                case '?':
                        errflag++;
                        break;
                case 'o':
                        options = optarg;
                        while (*options != '\0') {
                                switch(getsubopt(&options,myopts,&value) {
                                case READONLY : /* process ro option */
                                        break;
                                case READWRITE : /* process rw option */
                                        break;
```

```
                        case WRITESIZE : /* process wsize option */
                                if (value == NULL) {
                                        error_no_arg();
                                        errflag++;
                                } else
                                        write_size = atoi(value);
                                break;
                        case READSIZE : /* process rsize option */
                                if (value == NULL) {
                                        error_no_arg();
                                        errflag++;
                                } else
                                        read_size = atoi(value);
                                break;
                        default :
                                /* process unknown token */
                                error_bad_token(value);
                                errflag++;
                                break;
                        }
                    }
                    break;
                }
            }
            if (errflag) {
                /* print usage instructions etc. */
            }
            for (; optind<argc; optind++) {
                /* process remaining arguments */
            }
            .
            .
            .
}
```

SEE ALSO

mount(1M), **getopt**(3C)

NOTES

During parsing, commas in the option input string are changed to null characters. White space in tokens or token-value pairs must be protected from the shell by quotes.

NAME

gettext, dgettext, dcgettext, textdomain, bindtextdomain – message handling functions

SYNOPSIS

cc [*flag* ...] *file* ... **–lintl** [*library* ...]

```
#include <libintl.h>
#include <locale.h>      /* needed for dcgettext() only */

char *gettext(const char *msgid);

char *dgettext(const char *domainname, const char *msgid);

char *dcgettext(const char *domainname, const char *msgid, int category);

char *textdomain(const char *domainname);

char *bindtextdomain(const char *domainname, const char *dirname);
```

MT-LEVEL

Safe with exceptions

DESCRIPTION

gettext( ), **dgettext( )**, and **dcgettext( )** attempt to retrieve a target string based on the specified *msgid* argument within the context of a specific domain and the current locale. The length of strings returned by **gettext( )**, **dgettext( )**, and **dcgettext( )** is undetermined until the function is called. The *msgid* argument is a null-terminated string.

NLSPATH is searched first for the location of the **LC_MESSAGES** catalogue. The setting of the **LC_MESSAGES** category of the current locale determines the locale used by **gettext( )** and **dgettext( )** for string retrieval. *category* determines the locale used by **dcgettext( ).** If **NLSPATH** is not defined and the current locale is "C", **gettext( )**, **dgettext( )**, and **dcgettext( )** simply return the message string that was passed. In a locale other than "C", if **NLSPATH** is not defined or if a message catalogue is not found in any of the components specified by **NLSPATH**, the routines search for the message catalogue *dirname*/*locale*/*category*/*domainname***.mo**, after querying **bindtextdomain( )** for *dirname*.

For **gettext( )**, the domain used is set by the last valid call to **textdomain( )**. If a valid call to **textdomain( )** has not been made, the default domain (called **messages**) is used.

For **dgettext( )** and **dcgettext( )**, the domain used is specified by the *domainname* argument. The *domainname* argument is equivalent in syntax and meaning to the *domainname* argument to **textdomain( )**, except that the selection of the domain is valid only for the duration of the **dgettext( )** or **dcgettext( )** call.

textdomain( ) sets or queries the name of the current domain of the active **LC_MESSAGES** locale category. The *domainname* argument is a null-terminated string that can contain only the characters allowed in legal filenames.

The *domainname* argument is the unique name of a domain on the system. If there are multiple versions of the same domain on one system, namespace collisions can be avoided by using **bindtextdomain( )**. If **textdomain( )** is not called, a default domain is selected. The setting of domain made by the last valid call to **textdomain( )** remains valid across subsequent calls to **setlocale**(3C), and **gettext( )**.

The *domainname* argument is applied to the currently active **LC_MESSAGES** locale.

The current setting of the domain can be queried without affecting the current state of the domain by calling **textdomain()** with *domainname* set to the null pointer. Calling **textdomain()** with a *domainname* argument of a null string sets the domain to the default domain (**messages**).

bindtextdomain() binds the path predicate for a message domain *domainname* to the value contained in *dirname*. If *domainname* is a non-empty string and has not been bound previously, **bindtextdomain()** binds *domainname* with *dirname*.

If *domainname* is a non-empty string and has been bound previously, **bindtextdomain()** replaces the old binding with *dirname*. *dirname* can be an absolute or relative pathname being resolved when **gettext()**, **dgettext()**, or **dcgettext()** are called. If *domainname* is a null pointer or an empty string, **bindtextdomain()** returns **NULL**. User defined domain names cannot begin with the string **SYS_**. Domain names beginning with this string are reserved for system use.

RETURN VALUES

The individual bytes of the string returned by **gettext()**, **dgettext()**, or **dcgettext()** can contain any value other than null. If *msgid* is a null pointer, the return value is undefined. The string returned must not be modified by the program, and can be invalidated by a subsequent call to **gettext()**, **dgettext()**, **dcgettext()**, or **setlocale**(3C). If the *domainname* argument to **dgettext()** or **dcgettext()** is a null pointer, the results are undefined.

If the target string cannot be found in the current locale and selected domain, **gettext()**, **dgettext()**, and **dcgettext()** return *msgid*.

The normal return value from **textdomain()** is a pointer to a string containing the current setting of the domain. If *domainname* is a null pointer, **textdomain()** returns a pointer to the string containing the current domain. If **textdomain()** was not previously called and *domainname* is a null string, the name of the default domain is returned. The name of the default domain is **messages.**

The return value from **bindtextdomain()** is a null-terminated string containing *dirname* or the directory binding associated with *domainname* if *dirname* is **NULL**. If no binding is found, the default return value is **/usr/lib/locale**. If *domainname* is a null pointer or an empty string, **bindtextdomain()** takes no action and returns a null pointer. The string returned must not be modified by the caller.

FILES

/usr/lib/locale
: The default path predicate for message domain files.

/usr/lib/locale/*locale***/LC_MESSAGES/***domainname***.mo**
: system default location for file containing messages for language *locale* and *domainname*

/usr/lib/locale/*locale***/LC_XXX/***domainname***.mo**
: system default location for file containing messages for language *locale* and *domainname* for **dcgettext()** calls where **LC_XXX** is **LC_CTYPE**, **LC_NUMERIC**, **LC_TIME**, **LC_COLLATE**, **LC_MONETARY**, or **LC_MESSAGES.**

dirname/*locale*/**LC_MESSAGES**/*domainname*.**mo**
location for file containing messages for domain *domainname* and path predicate *dirname* after a successful call to **bindtextdomain( )**

dirname/*locale*/**LC_XXX**/*domainname*.**mo**
location for files containing messages for domain *domainname*, language *locale*, and path predicate *dirname* after a successful call to **bindtextdomain( )** for **dcgettext( )** calls where **LC_XXX** is one of **LC_CTYPE**, **LC_NUMERIC**, **LC_TIME**, **LC_COLLATE**, **LC_MONETARY**, or **LC_MESSAGES.**

SEE ALSO

msgfmt(1), **xgettext**(1), **setlocale**(3C), **environ**(5)

NOTES

These routines impose no limit on message length. However, a text *domainname* is limited to **TEXTDOMAINMAX** (256) bytes.

gettext, **dgettext**, **dcgettext**, **textdomain** and **bindtextdomain** can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

gettimeofday, settimeofday – get or set the date and time

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <sys/time.h>

int gettimeofday(*tp, tzp***)**
struct timeval ∗*tzp*;
struct timezone ∗*tzp*;

int settimeofday(*tp, tzp***)**
struct timeval ∗*tzp*;
struct timezone ∗*tzp*;

DESCRIPTION

The system's notion of the current Greenwich time is obtained with the **gettimeofday()** call, and set with the **settimeofday()** call. The current time is expressed in elapsed seconds and microseconds since 00:00 GMT, January 1, 1970 (zero hour). The resolution of the system clock is hardware dependent; the time may be updated continuously, or in clock ticks.

tp points to a **timeval** structure, which includes the following members:

```
long  tv_sec;   /* seconds since Jan. 1, 1970 */
long  tv_usec;  /* and microseconds */
```

If *tp* is a **NULL** pointer, the current time information is not returned or set.

tzp is an obsolete pointer formerly used to get and set timezone information. *tzp* is now ignored. Timezone information is now handled using the **TZ** environment variable; see **TIMEZONE**(4).

Only the privileged user may set the time of day.

RETURN VALUES

A –1 return value indicates an error occurred; in this case an error code is stored in the global variable **errno**.

ERRORS

The following error codes may be set in **errno**:

EINVAL *tp* specifies an invalid time.

EPERM A user other than the privileged user attempted to set the time.

SEE ALSO

adjtime(2), **ctime**(3C), **gettimeofday**(3C), **TIMEZONE**(4)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

tzp is ignored in SunOS 5.*X* releases.

tv_usec is always 0.

NAME

gettimeofday, settimeofday – get or set the date and time

SYNOPSIS

```
#include <sys/time.h>
int gettimeofday(struct timeval *tp, void * );
int settimeofday(struct timeval *tp, void * );
```

MT-LEVEL

MT-Safe

DESCRIPTION

gettimeofday( ) gets and **settimeofday( )** sets the system's notion of the current time. The current time is expressed in elapsed seconds and microseconds since 00:00 Universal Coordinated Time, January 1, 1970. The resolution of the system clock is hardware dependent; the time may be updated continuously or in clock ticks.

tp points to a **timeval** structure, which includes the following members:

```
long    tv_sec;   /* seconds since Jan. 1, 1970 */
long    tv_usec;  /* and microseconds */
```

If *tp* is a null pointer, the current time information is not returned or set.

The **TZ** environment variable holds time zone information. See **TIMEZONE**(4).

The second argument to **gettimeofday( )** and **settimeofday( )** should be a pointer to **NULL.**

Only the privileged user may set the time of day.

SEE ALSO

adjtime(2), **ctime**(3C), **TIMEZONE**(4)

RETURN VALUES

A **–1** return value indicates that an error occurred and **errno** has been set.

ERRORS

The following error codes may be set in **errno**:

EINVAL	*tp* specifies an invalid time.
EPERM	A user other than the privileged user attempted to set the time or time zone.

NOTES

The implementation of **settimeofday( )** ignores the **tv_usec** field of **tp**. If the time needs to be set with better than one second accuracy, call **settimeofday( )** for the seconds and then **adjtime( )** for finer accuracy.

NAME

gettxt – retrieve a text string

SYNOPSIS

```
#include <nl_types.h>
char *gettxt(const char *msgid, const char *dflt_str);
```

MT-LEVEL

Safe with exceptions

DESCRIPTION

gettxt( ) retrieves a text string from a message file. The arguments to the function are a message identification *msgid* and a default string *dflt_str* to be used if the retrieval fails.

The text strings are in files created by the **mkmsgs** utility (see **mkmsgs**(1)) and installed in directories in **/usr/lib/locale/***locale***/LC_MESSAGES**.

The directory *locale* can be viewed as the language in which the text strings are written. The user can request that messages be displayed in a specific language by setting the environment variable **LC_MESSAGES**. If **LC_MESSAGES** is not set, the environment variable **LANG** will be used. If **LANG** is not set, the files containing the strings are in **/usr/lib/locale/C/LC_MESSAGES/***.

The user can also change the language in which the messages are displayed by invoking the **setlocale( )** function with the appropriate arguments.

If **gettxt( )** fails to retrieve a message in a specific language it will try to retrieve the same message in U.S. English. On failure, the processing depends on what the second argument *dflt_str* points to. A pointer to the second argument is returned if the second argument is not the null string. If *dflt_str* points to the null string, a pointer to the U.S. English text string "**Message not found!!\n**" is returned.

The following depicts the acceptable syntax of *msgid* for a call to **gettxt( )**.

<*msgid*> **=** <*msgfilename*>**:**<*msgnumber*>

The first field is used to indicate the file that contains the text strings and must be limited to 14 characters. These characters must be selected from the set of all character values excluding \0 (null) and the ASCII code for **/** (slash) and **:** (colon). The names of message files must be the same as the names of files created by **mkmsgs** and installed in **/usr/lib/locale/***locale***/LC_MESSAGES/***. The numeric field indicates the sequence number of the string in the file. The strings are numbered from *1* to *n* where *n* is the number of strings in the file.

On failure to pass the correct **msgid** or a valid message number to **gettxt( )** a pointer to the text string "**Message not found!!\n**" is returned.

EXAMPLES

```
gettxt("UX:10", "hello world\n")
gettxt("UX:10", "")
```

UX is the name of the file that contains the messages. **10** is the message number.

FILES

/usr/lib/locale/C/LC_MESSAGES/*	contains default message files created by **mkmsgs**
/usr/lib/locale/*locale***/LC_MESSAGES/***	contains message files for different languages created by **mkmsgs**

SEE ALSO

exstr(1), **mkmsgs**(1), **srchtxt**(1), **gettext**(3I), **fmtmsg**(3C), **setlocale**(3C), **environ**(5)

NOTES

It is recommended that **gettext**(3I) be used in place of this routine.

NAME

getusershell, setusershell, endusershell – get legal user shells

SYNOPSIS

```
char *getusershell()
void setusershell()
void endusershell()
```

DESCRIPTION

getusershell() returns a pointer to a legal user shell as defined by the system manager in the file **/etc/shells**. If **/etc/shells** does not exist, a list of the ten locations of the standard system shells: **/usr/bin/sh**, **/usr/bin/csh**, **/usr/bin/ksh**, **/usr/bin/jsh**, **/bin/sh**, **/bin/csh**, **/bin/ksh**, **/bin/jsh**, **/sbin/sh**, **/sbin/jsh**, are used instead of the file.

getusershell() (opens the file **/etc/shells** if it exists) returns the next entry in the list of shells.

setusershell() rewinds the file, or the list.

endusershell() closes the file, and frees any memory used by **getusershell()** and **setusershell()**. As a side effect, **endusershell()** rewinds the file **/etc/shells**.

FILES

/etc/shells
/usr/bin/sh
/usr/bin/csh
/usr/bin/ksh
/usr/bin/jsh
/bin/sh
/bin/csh
/bin/ksh
/bin/jsh
/sbin/sh
/sbin/jsh

RETURN VALUES

getusershell() returns a **NULL** pointer on EOF.

BUGS

All information is contained in memory that may be freed with a call to **endusershell()**, so it must be copied if it is to be saved.

NAME

getutent, getutid, getutline, pututline, setutent, endutent, utmpname – access utmp file entry

SYNOPSIS

```
#include <utmp.h>

struct utmp *getutent(void);
struct utmp *getutid(const struct utmp *id);
struct utmp *getutline(const struct utmp *line);
struct utmp *pututline(const struct utmp *utmp);
void setutent(void);
void endutent(void); int utmpname(const char *file);
```

MT-LEVEL

Unsafe

DESCRIPTION

getutent( ), **getutid( )**, **getutline( )**, and **pututline( )** each return a pointer to a **utmp** structure with the following members:

```
char                ut_user[8];     /* user login name */
char                ut_id[4];       /* /sbin/inittab id */
                                    /* (usually line #) */
char                ut_line[12];    /* device name (console, lnxx) */
short               ut_pid;         /* process id */
short               ut_type;        /* type of entry */
struct exit_status  ut_exit;        /* exit status of a process */
                                    /* marked as DEAD_PROCESS */
time_t              ut_time;        /* time entry was made */
```

The structure exit status includes the following members:

```
short   e_termination;  /* termination status */
short   e_exit;         /* exit status */
```

getutent( ) reads in the next entry from a **utmp**-like file. If the file is not already open, it opens it. If it reaches the end of the file, it fails.

getutid( ) searches forward from the current point in the **utmp** file until it finds an entry with a **ut_type** matching *id*→**ut_type** if the type specified is **RUN_LVL**, **BOOT_TIME**, **OLD_TIME**, or **NEW_TIME**. If the type specified in *id* is **INIT_PROCESS**, **LOGIN_PROCESS**, **USER_PROCESS**, or **DEAD_PROCESS**, then **getutid( )** will return a pointer to the first entry whose type is one of these four and whose **ut_id** field matches *id*→**ut_id**. If the end of file is reached without a match, it fails.

getutline( ) searches forward from the current point in the **utmp** file until it finds an entry of the type **LOGIN_PROCESS** or **ut_line** string matching the *line*→**ut_line** string. If the end of file is reached without a match, it fails.

pututline( ) writes out the supplied **utmp** structure into the **utmp** file. It uses **getutid( )** to search forward for the proper place if it finds that it is not already at the proper place. It is expected that normally the user of **pututline( )** will have searched for the proper entry using one of the these routines. If so, **pututline( )** will not search. If **pututline( )** does not

find a matching slot for the new entry, it will add a new entry to the end of the file. It returns a pointer to the **utmp** structure. When called by a non-root user, **pututline()** invokes a **setuid()** root program to verify and write the entry, since **/etc/utmp** is normally writable only by root. In this event, the *ut_name* field must correspond to the actual user name associated with the process; the *ut_type* field must be either **USER_PROCESS** or **DEAD_PROCESS**; and the *ut_line* field must be a device special file and be writable by the user.

setutent() resets the input stream to the beginning of the file. This reset should be done before each search for a new entry if it is desired that the entire file be examined.

endutent() closes the currently open file.

utmpname() allows the user to change the name of the file examined, from **/var/adm/utmp** to any other file. It is most often expected that this other file will be **/var/adm/wtmp**. If the file does not exist, this will not be apparent until the first attempt to reference the file is made. **utmpname()** does not open the file. It just closes the old file if it is currently open and saves the new file name.

RETURN VALUES

A null pointer is returned upon failure to read, whether for permissions or having reached the end of file, or upon failure to write. If the file name given is longer than 79 characters, **utmpname()** returns 0. Otherwise, it returns 1.

FILES

/var/adm/utmp
/var/adm/wtmp

SEE ALSO

getutxent(3C), **ttyslot**(3C), **utmp**(4)

NOTES

The most current entry is saved in a static structure. Multiple accesses require that it be copied before further accesses are made. On each call to either **getutid()** or **getutline()**, the routine examines the static structure before performing more I/O. If the contents of the static structure match what it is searching for, it looks no further. For this reason, to use **getutline()** to search for multiple occurrences, it would be necessary to zero out the static area after each success, or **getutline()** would just return the same structure over and over again. There is one exception to the rule about emptying the structure before further reads are done. The implicit read done by **pututline()** (if it finds that it is not already at the correct place in the file) will not hurt the contents of the static structure returned by the **getutent()**, **getutid()** or **getutline()** routines, if the user has just modified those contents and passed the pointer back to **pututline()**.

These routines use buffered standard I/O for input, but **pututline()** uses an unbuffered non-standard write to avoid race conditions between processes trying to modify the **utmp** and **wtmp** files.

NAME

getutxent, getutxid, getutxline, pututxline, setutxent, endutxent, utmpxname, getutmp, getutmpx, updwtmp, updwtmpx – access utmpx file entry

SYNOPSIS

```
#include <utmpx.h>

struct utmpx *getutxent(void);
struct utmpx *getutxid(const struct utmpx *id);
struct utmpx *getutxline(const struct utmpx *line);
struct utmpx *pututxline(const struct utmpx *utmpx);
void setutxent(void);
void endutxent(void);
int utmpxname(const char *file);
void getutmp(struct utmpx *utmpx, struct utmp *utmp);
void getutmpx(struct utmp *utmp, struct utmpx *utmpx);
void updwtmp(char *wfile, struct utmp *utmp);
void updwtmpx(char *wfilex, struct utmpx *utmpx);
```

MT-LEVEL

Unsafe

DESCRIPTION

getutxent(), **getutxid()**, and **getutxline()** each return a pointer to a **utmpx** structure with the following members:

```
char    ut_user[32];          /* user login name */
char    ut_id[4];             /* /etc/inittab id */
                              /* (usually line #) */
char    ut_line[32];          /* device name (console, lnxx) */
pid_t   ut_pid;               /* process id */
short   ut_type;              /* type of entry */
struct  exit_status ut_exit;  /* exit status of a process */
                              /* marked as DEAD_PROCESS */
struct  timeval ut_tv;        /* time entry was made */
long    ut_session;           /* session ID, used for windowing */
long    pad[5];               /* reserved for future use */
short   ut_syslen;            /* significant length of ut_host */
                              /* including terminating null */
char    ut_host[257];         /* host name, if remote */
```

The structure exit status includes the following members:

```
short   e_termination;  /* termination status */
short   e_exit;         /* exit status */
```

getutxent()

Reads in the next entry from a **utmpx**-like file. If the file is not already open, it opens it. If it reaches the end of the file, it fails.

getutxid()

Searches forward from the current point in the **utmpx** file until it finds an entry with a **ut_type** matching *id*→**ut_type** if the type specified is **RUN_LVL**, **BOOT_TIME**, **OLD_TIME**, or **NEW_TIME**. If the type specified in *id* is **INIT_PROCESS**, **LOGIN_PROCESS**,

USER_PROCESS, or **DEAD_PROCESS**, then **getutxid()** will return a pointer to the first entry whose type is one of these four and whose *ut_id* field matches *id*→**ut_id**. If the end of file is reached without a match, it fails.

getutxline() Searches forward from the current point in the **utmpx** file until it finds an entry of the type **LOGIN_PROCESS** or **USER_PROCESS** which also has a *ut_line* string matching the *line*→**ut_line** string. If the end of file is reached without a match, it fails.

pututxline() Writes out the supplied **utmpx** structure into the **utmpx** file. It uses **getutxid()** to search forward for the proper place if it finds that it is not already at the proper place. It is expected that normally the user of **pututxline()** will have searched for the proper entry using one of the **getutx()** routines. If so, **pututxline()** will not search. If **pututxline()** does not find a matching slot for the new entry, it will add a new entry to the end of the file. It returns a pointer to the **utmpx** structure. When called by a non-root user, **pututxline()** invokes a **setuid()** root program to verify and write the entry, since **/etc/utmpx** is normally writable only by root. In this event, the *ut_name* field must correspond to the actual user name associated with the process; the *ut_type* field must be either **USER_PROCESS** or **DEAD_PROCESS**; and the *ut_line* field must be a device special file and be writable by the user.

setutxent() Resets the input stream to the beginning of the file. This should be done before each search for a new entry if it is desired that the entire file be examined.

endutxent() Closes the currently open file.

utmpxname() Allows the user to change the name of the file examined, from **/var/adm/utmpx** to any other file. It is most often expected that this other file will be **/var/adm/wtmpx**. If the file does not exist, this will not be apparent until the first attempt to reference the file is made. **utmpxname()** does not open the file. It just closes the old file if it is currently open and saves the new file name. The new file name must end with the ''**x**'' character to allow the name of the corresponding **utmp** file to be easily obtainable; otherwise, an error code of **1** is returned.

getutmp() Copies the information stored in the fields of the **utmpx** structure to the corresponding fields of the **utmp** structure. If the information in any field of **utmpx** does not fit in the corresponding **utmp** field, the data is truncated. (See **getutent**(3C) for **utmp** structure)

getutmpx() Copies the information stored in the fields of the **utmp** structure to the corresponding fields of the **utmpx** structure. (See **getutent**(3C) for **utmp** structure)

updwtmp() Checks the existence of *wfile* and its parallel file, whose name is obtained by appending an ''**x**'' to *wfile*. If only one of them exists, the second one is created and initialized to reflect the state of the existing file. *utmp* is written to *wfile* and the corresponding **utmpx** structure is written to the parallel file.

updwtmpx() Checks the existence of *wfilex* and its parallel file, whose name is obtained by truncating the final "**x**" from *wfilex*. If only one of them exists, the second one is created and initialized to reflect the state of the existing file. *utmpx* is written to *wfilex*, and the corresponding **utmp** structure is written to the parallel file.

RETURN VALUES A null pointer is returned upon failure to read, whether for permissions or having reached the end of file, or upon failure to write.

FILES

/var/adm/utmp	contains current user access and adminstrative information (old format)
/var/adm/utmpx	contains current user access and adminstration information (new format)
/var/adm/wtmp	contains a history of user access and adminstrative information.
/var/adm/wtmpx	contains a history of user access and adminstrative information.

SEE ALSO **getutent**(3C), **ttyslot**(3C), **utmp**(4), **utmpx**(4)

NOTES The most current entry is saved in a static structure. Multiple accesses require that it be copied before further accesses are made. On each call to either **getutxid()** or **getutxline()**, the routine examines the static structure before performing more I/O. If the contents of the static structure match what it is searching for, it looks no further. For this reason, to use **getutxline()** to search for multiple occurrences it would be necessary to zero out the static after each success, or **getutxline()** would just return the same structure over and over again. There is one exception to the rule about emptying the structure before further reads are done. The implicit read done by **pututxline()** (if it finds that it is not already at the correct place in the file) will not hurt the contents of the static structure returned by the **getutxent()**, **getutxid()**, or **getutxline()** routines, if the user has just modified those contents and passed the pointer back to **pututxline()**.

These routines use buffered standard I/O for input, but **pututxline()** uses an unbuffered write to avoid race conditions between processes trying to modify the **utmpx** and **wtmpx** files.

NAME

getvfsent, getvfsfile, getvfsspec, getvfsany – get vfstab file entry

SYNOPSIS

```
#include <stdio.h>
#include <sys/vfstab.h>

int getvfsent(FILE *fp, struct vfstab *vp);

int getvfsfile(FILE *fp, struct vfstab *vp, char *file);

int getvfsspec(FILE *, struct vfstab *vp, char *spec);

int getvfsany(FILE *, struct vfstab *vp, vfstab *vref);
```

MT-LEVEL

Safe

DESCRIPTION

getvfsent( ), **getvfsfile( )**, **getvfsspec( )**, and **getvfsany( )** each fill in the structure pointed to by *vp* with the broken-out fields of a line in the **/etc/vfstab** file. Each line in the file contains a **vfstab** structure, declared in the **<sys/vfstab.h>** header:

```
char    *vfs_special;
char    *vfs_fsckdev;
char    *vfs_mountp;
char    *vfs_fstype;
char    *vfs_fsckpass;
char    *vfs_automnt;
char    *vfs_mntopts;
```

The fields have meanings described in **vfstab**(4).

getvfsent( ) returns a pointer to the next **vfstab** structure in the file; so successive calls can be used to search the entire file. **getvfsfile( )** searches the file referenced by *fp* until a mount point matching *file* is found and fills *vp* with the fields from the line in the file. **getvfsspec( )** searches the file referenced by *fp* until a special device matching *spec* is found and fills *vp* with the fields from the line in the file. *spec* will try to match on device type (block or character special) and major and minor device numbers. If it cannot match in this manner, then it compares the strings. **getvfsany( )** searches the file referenced by *fp* until a match is found between a line in the file and *vref*. *vref* matches the line if all non-null entries in *vref* match the corresponding fields in the file.

Note that these routines do not open, close, or rewind the file.

RETURN VALUES

If the next entry is successfully read by **getvfsent( )** or a match is found with **getvfsfile( )**, **getvfsspec( )**, or **getvfsany( )**, 0 is returned. If an end-of-file is encountered on reading, these functions return –1. If an error is encountered, a value greater than 0 is returned. The possible error values are:

VFS_TOOLONG	A line in the file exceeded the internal buffer size of **VFS_LINE_MAX**.
VFS_TOOMANY	A line in the file contains too many fields.
VFS_TOOFEW	A line in the file contains too few fields.

NOTES

The members of the **vfstab** structure point to information contained in a static area, so it must be copied if it is to be saved.

FILES

/etc/vfstab

SEE ALSO

vfstab(4)

NAME

getwc, getwchar, fgetwc – convert EUC character from the stream to Process Code

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

```
#include <stdio.h>
#include <widec.h>

wint_t getwc(FILE *stream);

wint_t getwchar(void);

wint_t fgetwc (FILE *stream);
```

MT-LEVEL

MT-Safe

DESCRIPTION

getwc( ) and **fgetwc( )** convert the next Extended Unix Code (EUC) character from the named input *stream* into a *wchar_t* Process Code character, and return it as an integer. They also move the file pointer, if defined, ahead one EUC character in the *stream*. **getwchar( )** is defined as **fgetwc(stdin)**. **getwc( )** and **getwchar( )** are macros.

RETURN VALUES

These functions return the integer constant **EOF** at end-of-file or upon an error. The end-of-file condition is remembered, even on a terminal, and all subsequent attempts to read will return **EOF** until the condition is cleared with **clearerr( )** (see **ferror**(3S)).

SEE ALSO

ferror(3S), **fopen**(3S), **fread**(3S), **getws**(3I), **putwc**(3I), **scanf**(3S), **ungetwc**(3I)

WARNING

If the integer value returned by **getwc( )**, **getwchar( )**, or **fgetwc( )** is stored into a *wchar_t* variable and then compared against the integer constant **EOF**, the comparison will not succeed, because *wchar_t* is defined as unsigned.

NAME

getwd – get current working directory pathname

SYNOPSIS

```
#include <sys/param.h>
#include <unistd.h>

char *getwd(char *pathname);
```

DESCRIPTION

getwd( ) copies the absolute pathname of the current working directory to *pathname* and returns a pointer to the result.

RETURN VALUES

getwd( ) returns zero and places a message in *pathname* if an error occurs. *pathname* points to a character array of at least **MAXPATHLEN** length.

SEE ALSO

getcwd(3C)

NAME

getwidth – get codeset information

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <euc.h>
#include <getwidth.h>

void getwidth(eucwidth_t ∗*ptr***);**

MT-LEVEL

MT-Safe with exceptions

DESCRIPTION

The **getwidth( )** function reads the character class table for the current locale, generated by **wchrtbl( )** (see **chrtbl**(1M)) to get information on the supplementary codesets. **getwidth( )** sets this information into the struct **eucwidth_t**. This struct is defined in **<euc.h>** and has the following members:

short int _eucw1,_eucw2,_eucw3;
short int _scrw1,_scrw2,_scrw3;
short int _pcw;
char _multibyte;

Codeset width values for supplementary codesets 1, 2, and 3 are set in **_eucw1**, **_eucw2**, and **_eucw3**, respectively. Screen width values for supplementary codesets 1, 2, and 3 are set in **_scrw1**, **_scrw2**, and **_scrw3**, respectively.

The width of Extended Unix Code (EUC) Process Code is set in **_pcw**. The **_multibyte** entry is set to **1** if multibyte characters are used, and set to **0** if only single-byte characters are used.

SEE ALSO

chrtbl(1M), **euclen**(3I), **setlocale**(3C)

NOTES

getwidth can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

getws, fgetws – convert a string of EUC characters from the stream to Process Code

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

```
#include <stdio.h>
#include <widec.h>

wchar_t *getws(wchar_t *s);

wchar_t *fgetws(wchar_t *s, int n, FILE *stream);
```

MT-LEVEL

MT-Safe

DESCRIPTION

getws( ) reads a string of Extended Unix Code (EUC) characters from the standard input stream, **stdin**, converts it to process code, and writes it to the array pointed to by *s*, until a new-line character is read or an end-of-file condition is encountered. The new-line character is discarded and the string is terminated with a *wchar_t* **NULL** character. **getws( )** returns its argument.

fgetws( ) reads EUC characters from the *stream*, converts them to Process Code, and writes them to the array pointed to by *s*. It stops when either *n*–1 characters are read, a new-line character is read and transferred to *s*, or an end-of-file condition is encountered. The string is then terminated with a *wchar_t* **NULL** character. **fgetws( )** returns its first argument.

RETURN VALUES

If end-of-file is encountered and no characters have been read, no characters are transferred to *s* and a **NULL** pointer is returned. If a read error occurs, such as trying to use these functions on a file that has not been opened for reading, a **NULL** pointer is returned. Otherwise *s* is returned.

SEE ALSO

ferror(3S), **fread**(3S), **getwc**(3I), **putws**(3I), **scanf**(3S)

NAME

glob, globfree – generate path names matching a pattern

SYNOPSIS

```
#include <glob.h>
int glob( const char *pattern, int flags, int(*errfunc)(const char *epath, int eerrno),
    glob_t *pglob);
void globfree(glob_t *pglob);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **glob( )** function is a path name generator.

The **globfree( )** function frees any memory allocated by **glob( )** associated with *pglob*.

pattern **Argument**

The argument *pattern* is a pointer to a path name pattern to be expanded. The **glob( )** function matches all accessible path names against this pattern and develops a list of all path names that match. In order to have access to a path name, **glob( )** requires search permission on every component of a path except the last, and read permission on each directory of any filename component of *pattern* that contains any of the following special characters:

* ? [

pglob **Argument**

The structure type **glob_t** is defined in the header **<glob.h>** and includes at least the following members:

size_t gl_pathc	Count of paths matched by *pattern*.
char **gl_pathv	Pointer to a list of matched path names.
size_t gl_offs	Slots to reserve at the beginning of **gl_pathv**.

The **glob( )** function stores the number of matched path names into *pglob*–>**gl_pathc** and a pointer to a list of pointers to path names into *pglob*–>**gl_pathv**. The path names are in sort order as defined by the current setting of the **LC_COLLATE** category. The first pointer after the last path name is a **NULL** pointer. If the pattern does not match any path names, the returned number of matched paths is set to zero, and the contents of *pglob*–>**gl_pathv** are implementation-dependent.

It is the caller's responsibility to create the structure pointed to by *pglob*. The **glob( )** function allocates other space as needed, including the memory pointed to by **gl_pathv**. The **globfree( )** function frees any space associated with *pglob* from a previous call to **glob( )**.

flags **Argument**

The *flags* argument is used to control the behavior of **glob( )**. The value of *flags* is a bitwise inclusive OR of zero or more of the following constants, which are defined in the header **<glob.h>**:

GLOB_APPEND	Append path names generated to the ones from a previous call to **glob( )**.
GLOB_DOOFFS	Make use of *pglob*–>**gl_offs**. If this flag is set, *pglob*–>**gl_offs** is used to specify how many **NULL** pointers to add to the beginning of *pglob*–>**gl_pathv**. In other words, *pglob*–>**gl_pathv** will point to

pglob–>**gl_offs** **NULL** pointers, followed by *pglob*–>**gl_pathc** path name pointers, followed by a **NULL** pointer.

GLOB_ERR Causes **glob( )** to return when it encounters a directory that it cannot open or read. Ordinarily, **glob( )** continues to find matches.

GLOB_MARK Each path name that is a directory that matches *pattern* has a slash appended.

GLOB_NOCHECK If *pattern* does not match any path name, then **glob( )** returns a list consisting of only *pattern*, and the number of matched path names is 1.

GLOB_NOESCAPE Disable backslash escaping.

GLOB_NOSORT Ordinarily, **glob( )** sorts the matching path names according to the current setting of the **LC_COLLATE** category. When this flag is used the order of path names returned is unspecified.

The **GLOB_APPEND** flag can be used to append a new set of path names to those found in a previous call to **glob( )**. The following rules apply when two or more calls to **glob( )** are made with the same value of *pglob* and without intervening calls to **globfree( )**:

1. The first such call must not set **GLOB_APPEND**. All subsequent calls must set it.
2. All the calls must set **GLOB_DOOFFS**, or all must not set it.
3. After the second call, *pglob*–>**gl_pathv** points to a list containing the following:
   a. Zero or more **NULL** pointers, as specified by **GLOB_DOOFFS** and *pglob*–>**gl_offs**.
   b. Pointers to the path names that were in the *pglob*–>**gl_pathv** list before the call, in the same order as before.
   c. Pointers to the new path names generated by the second call, in the specified order.
4. The count returned in *pglob*–>**gl_pathc** will be the total number of path names from the two calls.
5. The application can change any of the fields after a call to **glob( )**. If it does, it must reset them to the original value before a subsequent call, using the same *pglob* value, to **globfree( )** or **glob( )** with the **GLOB_APPEND** flag.

errfunc and _epath_ Arguments

If, during the search, a directory is encountered that cannot be opened or read and *errfunc* is not a **NULL** pointer, **glob( )** calls **(****errfunc***)** with two arguments:

1. The *epath* argument is a pointer to the path that failed.
2. The *eerrno* argument is the value of *errno* from the failure, as set by the **opendir**(3C), **readdir**(3C) or **stat**(2) functions. (Other values may be used to report other errors not explicitly documented for those functions.)

The following constants are defined as error return values for **glob( )**:

GLOB_ABORTED The scan was stopped because **GLOB_ERR** was set or **(*errfunc)** returned non-zero.

GLOB_NOMATCH The pattern does not match any existing path name, and **GLOB_NOCHECK** was not set in flags.

GLOB_NOSPACE An attempt to allocate memory failed.

If **(****errfunc***)** is called and returns non-zero, or if the **GLOB_ERR** flag is set in *flags*, **glob( )** stops the scan and returns **GLOB_ABORTED** after setting *gl_pathc* and *gl_pathv* in *pglob* to reflect the paths already scanned. If **GLOB_ERR** is not set and either *errfunc* is a **NULL** pointer or **(****errfunc***)** returns zero, the error is ignored.

RETURN VALUES

The following values are returned by **glob( )**:

0 successful completion. The argument *pglob*–>**gl_pathc** returns the number of matched path names and the argument *pglob*–>**gl_pathv** contains a pointer to a null-terminated list of matched and sorted path names. However, if *pglob*–>**gl_pathc** is zero, the content of *pglob*–>**gl_pathv** is undefined.

non-zero an error has occurred. Non-zero constants are defined in **<glob.h>**. The arguments *pglob*–>**gl_pathc** and *pglob*–>**gl_pathv** are still set as defined above.

The **globfree( )** function returns no value.

USAGE

This function is not provided for the purpose of enabling utilities to perform path name expansion on their arguments, as this operation is performed by the shell, and utilities are explicitly not expected to redo this. Instead, it is provided for applications that need to do path name expansion on strings obtained from other sources, such as a pattern typed by a user or read from a file.

If a utility needs to see if a path name matches a given pattern, it can use **fnmatch**(3C).

Note that **gl_pathc** and **gl_pathv** have meaning even if **glob( )** fails. This allows **glob( )** to report partial results in the event of an error. However, if **gl_pathc** is zero, **gl_pathv** is unspecified even if **glob( )** did not return an error.

The **GLOB_NOCHECK** option could be used when an application wants to expand a path name if wildcards are specified, but wants to treat the pattern as just a string otherwise.

The new path names generated by a subsequent call with **GLOB_APPEND** are not sorted together with the previous path names. This mirrors the way that the shell handles path name expansion when multiple expansions are done on a command line.

Applications that need tilde and parameter expansion should use the **wordexp( )** function.

EXAMPLES

One use of the **GLOB_DOOFFS** flag is by applications that build an argument list for use with the **execv**(2), **execve( )** or **execvp( )** functions. Suppose, for example, that an application wants to do the equivalent of:

ls -l *.c

but for some reason:

system("ls -l *.c")

is not acceptable. The application could obtain approximately the same result using the sequence:

```
globbuf.gl_offs = 2;
glob ("*.c", GLOB_DOOFFS, NULL, &globbuf);
globbuf.gl_pathv[0] = "ls";
globbuf.gl_pathv[1] = "-l";
execvp ("ls", &globbuf.gl_pathv[0]);
```

Using the same example:

```
ls -l *.c *.h
```

could be approximately simulated using **GLOB_APPEND** as follows:

```
globbuf.gl_offs = 2;
glob ("*.c", GLOB_DOOFFS, NULL, &globbuf);
glob ("*.h", GLOB_DOOFFS | GLOB_APPEND, NULL, &globbuf);
...
```

SEE ALSO **execv**(2), **stat**(2), **fnmatch**(3C), **opendir**(3C), **readdir**(3C), **wordexp**(3C)

NAME

gmatch – shell global pattern matching

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

int gmatch(const char ∗*str***, const char** ∗*pattern***);**

MT-LEVEL

MT-Safe

DESCRIPTION

gmatch( ) checks whether the null-terminated string *str* matches the null-terminated pattern string *pattern*. See the **sh**(1) section **File Name Generation** for a discussion of pattern matching. A backslash (\) is used as an escape character in pattern strings.

RETURN VALUES

gmatch( ) returns non-zero if the pattern matches the string, zero if the pattern does not.

EXAMPLE

In the following example, **gmatch( )** returns non-zero (true) for all strings with "**a**" or "**-**" as their last character.

```
char *s;

gmatch (s, "*[a\-]" )
```

SEE ALSO

sh(1)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

grantpt – grant access to the slave pseudo-terminal device

SYNOPSIS

int grantpt(int *fildes***);**

MT-LEVEL

Safe

DESCRIPTION

The function **grantpt( )** changes the mode and ownership of the slave pseudo-terminal device associated with its master pseudo-terminal counter part. *fildes* is the file descriptor returned from a successful open of the master pseudo-terminal device. A *setuid* root program (see **setuid**(2)) is invoked to change the permissions. The user ID of the slave is set to the real UID of the calling process and the group ID is set to a reserved group. The permission mode of the slave pseudo-terminal is set to readable, writable by the owner and writeable by the group.

RETURN VALUES

Upon successful completion, the function **grantpt( )** returns **0**; otherwise it returns **-1**. Failure could occur if *fildes* is not an open file descriptor, if *fildes* is not associated with a master pseudo-terminal device, or if the corresponding slave device could not be accessed. **grantpt( )** will also fail if it is unable to successfully invoke the *setuid* root program.

SEE ALSO

open(2), **setuid**(2), **ptsname**(3C), **unlockpt**(3C)

STREAMS Programming Guide

NAME

hsearch, hcreate, hdestroy – manage hash search tables

SYNOPSIS

```
#include <search.h>
ENTRY *hsearch(ENTRY item, ACTION action);
int hcreate (size_t mekments);
void hdestroy(void);
```

MT-LEVEL

Safe

DESCRIPTION

hsearch( ) is a hash-table search routine generalized from Knuth (6.4) Algorithm D. It returns a pointer into a hash table indicating the location at which an entry can be found. The comparison function used by **hsearch( )** is **strcmp( )** (see **string**(3C)). *item* is a structure of type **ENTRY** (defined in the **<search.h>** header) containing two pointers: *item.key* points to the comparison key, and *item.data* points to any other data to be associated with that key. (Pointers to types other than void should be cast to pointer-to-void.) *action* is a member of an enumeration type **ACTION** (defined in **<search.h>**) indicating the disposition of the entry if it cannot be found in the table. **ENTER** indicates that the item should be inserted in the table at an appropriate point. Given a duplicate of an existing item, the new item is not entered and **hsearch( )** returns a pointer to the existing item. **FIND** indicates that no entry should be made. Unsuccessful resolution is indicated by the return of a null pointer.

hcreate( ) allocates sufficient space for the table, and must be called before **hsearch( )** is used. *nel* is an estimate of the maximum number of entries that the table will contain. This number may be adjusted upward by the algorithm in order to obtain certain mathematically favorable circumstances.

hdestroy( ) destroys the search table, and may be followed by another call to **hcreate( )**.

RETURN VALUES

hsearch( ) returns a null pointer if either the action is **FIND** and the item could not be found or the action is **ENTER** and the table is full.

hcreate( ) returns zero if it cannot allocate sufficient space for the table.

EXAMPLE

The following example will read in strings followed by two numbers and store them in a hash table, discarding duplicates. It will then read in strings and find the matching entry in the hash table and print it out.

```
#include <stdio.h>
#include <search.h>
#include <string.h>
#include <stdlib.h>

struct info {                   /* this is the info stored in table */
        int age, room;          /* other than the key */
};
```

```
#define NUM_EMPL    5000    /* # of elements in search table */

main( )
{
                    /* space to store strings */
        char string_space[NUM_EMPL*20];
                    /* space to store employee info */
        struct info info_space[NUM_EMPL];
                    /* next avail space in string_space */
        char *str_ptr = string_space;
                    /* next avail space in info_space */
        struct info *info_ptr = info_space;
        ENTRY item, *found_item;
                    /* name to look for in table */
        char name_to_find[30];
        int i = 0;

                    /* create table */
        (void) hcreate(NUM_EMPL);
        while (scanf("%s%d%d", str_ptr, &info_ptr->age,
              &info_ptr->room) != EOF && i++ < NUM_EMPL) {
                    /* put info in structure, and structure in item */
            item.key = str_ptr;
            item.data = (void *)info_ptr;
            str_ptr += strlen(str_ptr) + 1;
            info_ptr++;
                    /* put item into table */
            (void) hsearch(item, ENTER);
        }

                    /* access table */
        item.key = name_to_find;
        while (scanf("%s", item.key) != EOF) {
            if ((found_item = hsearch(item, FIND)) != NULL) {
                    /* if item is in the table */
                (void)printf("found %s, age = %d, room = %d\n",
                    found_item->key,
                    ((struct info *)found_item->data)->age,
                    ((struct info *)found_item->data)->room);
            } else {
                (void)printf("no such employee %s\n",
                    name_to_find)
            }
        }
        return 0;
```

```
}
```

SEE ALSO

bsearch(3C), **lsearch**(3C), **malloc**(3C), **string**(3C), **tsearch**(3C), **malloc**(3X)

The Art of Computer Programming, Volume 3, Sorting and Searching by Donald E. Knuth, published by Addison-Wesley Publishing Company, 1973.

NOTES

hsearch( ) and **hcreate( )** use **malloc**(3C) to allocate space.

Only one hash search table may be active at any given time.

NAME

hyperbolic, sinh, cosh, tanh, asinh, acosh, atanh – hyperbolic functions

SYNOPSIS

```
#include <math.h>
double sinh(double x);
double cosh(double x);
double tanh(double x);
double asinh(double x);
double acosh(double x);
double atanh(double x);
```

MT-LEVEL

MT-Safe

DESCRIPTION

These functions compute the designated direct and inverse hyperbolic functions for real arguments.

RETURN VALUES

For exceptional cases, **matherr**(3M) tabulates the values to be returned as dictated by various Standards.

SEE ALSO

matherr(3M)

DIAGNOSTICS

In IEEE 754 mode (i.e. the **–xlibmieee** cc compilation option), **sinh( )** and **cosh( )** return ±∞ as appropriate on overflow and raise the overflow exception; **acosh( )** returns a NaN and raises the invalid operation exception if its argument is less than 1; **atanh( )** returns a NaN and raises the invalid operation exception if its argument has absolute value greater than 1; **atanh(**±1**)** returns ±∞ and raises the division by zero exception.

NAME

hypot – Euclidean distance

SYNOPSIS

cc [*flag* ...] *file* ... **–lm** [*library* ...]

#include <math.h>

double hypot(double *x***, double** *y***);**

MT-LEVEL

MT-Safe

DESCRIPTION

hypot(*x***,***y***)** returns

sqrt(*x*∗*x* **+***y*∗*y***),**

taking precautions against unwarranted IEEE exceptions. On IEEE overflow, **hypot(***x***,***y***)** may also set **errno** and call **matherr**(3M). **hypot(**±∞**,***y***)** is +∞ for any *y*, even a *NaN*, and is exceptional only for a signaling *NaN*.

hypot(*x***,***y***)** and **atan2(***y***,***x***)** (see **trig**(3M)) convert rectangular coordinates (*x*,*y*) to polar (*r*,θ); **hypot(***x***,***y***)** computes *r*, the modulus or radius.

SEE ALSO

matherr(3M), **trig**(3M)

NAME

iconv – code conversion function

SYNOPSIS

```
#include <iconv.h>

size_t iconv(iconv_t cd, const char **inbuf, size_t *inbytesleft, char **outbuf,
        size_t *outbytesleft);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **iconv( )** function converts the sequence of characters from one codeset, in the array specified by *inbuf*, into a sequence of corresponding characters in another codeset, in the array specified by *outbuf*. The codesets are those specified in the *iconv_open*( ) call that returned the conversion descriptor, *cd*. The *inbuf* argument points to a variable that points to the first character in the input buffer and *inbytesleft* indicates the number of bytes to the end of the buffer to be converted. The *outbuf* argument points to a variable that points to the first available byte in the output buffer and *outbytesleft* indicates the number of the available bytes to the end of the buffer.

For state-dependent encodings, the conversion descriptor *cd* is placed into its initial shift state by a call for which *inbuf* is a null pointer, or for which *inbuf* points to a null pointer. When **iconv( )** is called in this way, and if *outbuf* is not a null pointer or a pointer to a null pointer, and *outbytesleft* points to a positive value, **iconv( )** will place, into the output buffer, the byte sequence to change the output buffer to its initial shift state. If the output buffer is not large enough to hold the entire reset sequence, **iconv( )** will fail and set **errno** to **E2BIG**. Subsequent calls with *inbuf* as other than a null pointer or a pointer to a null pointer cause the conversion to take place from the current state of the conversion descriptor.

If a sequence of input bytes does not form a valid character in the specified codeset, conversion stops after the previous successfully converted character. If the input buffer ends with an incomplete character or shift sequence, conversion stops after the previous successfully converted bytes. If the output buffer is not large enough to hold the entire converted input, conversion stops just prior to the input bytes that would cause the output buffer to overflow. The variable pointed to by *inbuf* is updated to point to the byte following the last byte successfully used in the conversion. The value pointed to by *inbytesleft* is decremented to reflect the number of bytes still not converted in the input buffer. The variable pointed to by *outbuf* is updated to point to the byte following the last byte of converted output data. The value pointed to by *outbytesleft* is decremented to reflect the number of bytes still available in the output buffer. For state-dependent encodings, the conversion descriptor is updated to reflect the shift state in effect at the end of the last successfully converted byte sequence.

If **iconv( )** encounters a character in the input buffer that is legal, but for which an identical character does not exist in the target codeset, **iconv( )** performs an implementation-defined conversion on this character.

RETURN VALUES

The **iconv()** function updates the variables pointed to by the arguments to reflect the extent of the conversion and returns the number of non-identical conversions performed. If the entire string in the input buffer is converted, the value pointed to by *inbytesleft* will be **0**. If the input conversion is stopped due to any conditions mentioned above, the value pointed to by *inbytesleft* will be non-zero and **errno** is set to indicate the condition. If an error occurs **iconv()** returns **(size_t) –1** and sets **errno** to indicate the error.

ERRORS

The **iconv()** function will fail if:

EILSEQ Input conversion stopped due to an input byte that does not belong to the input codeset.

E2BIG Input conversion stopped due to lack of space in the output buffer.

EINVAL Input conversion stopped due to an incomplete character or shift sequence at the end of the input buffer.

The **iconv()** function may fail if:

EBADF The *cd* argument is not a valid open conversion descriptor.

FILES

/usr/lib/iconv/*.so conversion modules

SEE ALSO

iconv(1), **iconv_open**(3), **iconv_close**(3), **iconv**(5)

NAME

iconv_close – code conversion deallocation function

SYNOPSIS

```
#include <iconv.h>
int iconv_close(iconv_t cd);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **iconv_close( )** function deallocates the conversion descriptor *cd* and all other associated resources allocated by the **iconv_open**(3) function.

If a file descriptor is used to implement the type **iconv_t**, that file descriptor will be closed.

RETURN VALUES

Upon successful completion, **iconv_close( )** returns **0**; otherwise, it returns **-1** and sets **errno** to indicate the error.

ERRORS

The **iconv_close( )** function may fail if:

EBADF The conversion descriptor is invalid.

SEE ALSO

iconv(3), **iconv_open**(3)

NAME

iconv_open – code conversion allocation function

SYNOPSIS

```
#include <iconv.h>
iconv_t iconv_open(const char *tocode, const char *fromcode);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **iconv_open()** function returns a conversion descriptor that describes a conversion from the codeset specified by the string pointed to by the *fromcode* argument to the codeset specified by the string pointed to by the *tocode* argument. For state-dependent encodings, the conversion descriptor will be in a codeset-dependent initial shift state, ready for immediate use with the **iconv**(3) function.

Settings of *fromcode* and *tocode* and their permitted combinations are implementation-dependent.

A conversion descriptor remains valid in a process until that process closes it.

RETURN VALUES

Upon successful completion **iconv_open()** returns a conversion descriptor for use on subsequent calls to **iconv().** Otherwise, **iconv_open()** returns **(iconv_t) –1** and sets **errno** to indicate the error.

ERRORS

The **iconv_open** function may fail if:

EMFILE	**{OPEN_MAX}** files descriptors are currently open in the calling process.
ENFILE	Too many files are currently open in the system.
ENOMEM	Insufficient storage space is available.
EINVAL	The conversion specified by *fromcode* and *tocode* is not supported by the implementation.

SEE ALSO

iconv(3), **iconv_close**(3), **malloc**(3C)

NOTES

iconv_open() uses **malloc**(3C) to allocate space for internal buffer areas. **iconv_open()** may fail if there is insufficient storage space to accommodate these buffers.

Portable applications must assume that conversion descriptors are not valid after a call to one of the **exec** functions.

NAME

ieee_functions, ilogb, isnan, copysign, fabs, fmod, nextafter, remainder, scalbn – appendix and related miscellaneous functions for IEEE arithmetic

SYNOPSIS

cc [*flag* ...] *file* ... **–lm** [*library* ...]

#include <math.h>

int ilogb(double *x***);**

int isnan(double *x***);**

double copysign(double *x***, double** *y***);**

double fabs(double *x***);**

double fmod(double *x***, double** *y***);**

double nextafter(double *x***, double** *y***);**

double remainder(double *x***, double** *y***);**

double scalbn(double *x***, int** *n***);**

MT-LEVEL

MT-Safe

DESCRIPTION

Most of these functions provide capabilities required by ANSI/IEEE Std 754-1985 or suggested in its appendix.

ilogb(*x***)** returns the unbiased exponent of *x* in integer format. **ilogb(±∞) = +MAXINT** and **ilogb(0) = –MAXINT**; **<values.h>** defines MAXINT as the largest int. **ilogb(***x***)** never generates an exception. When *x* is subnormal, **ilogb(***x***)** returns an exponent computed as if *x* were first normalized.

isnan(*x***)** returns 1 if *x* is *NaN*; otherwise it returns 0.

copysign(*x***,***y***)** returns a value with the magnitude of *x* and with the sign bit of *y*.

fabs(*x***)** returns the absolute value of *x*.

nextafter(*x***,***y***)** returns the next machine representable number from *x* in the direction *y*.

remainder(*x***,***y***)** and **fmod(***x***,***y***)** return a remainder of *x* with respect to *y*; that is, the result *r* is one of the numbers that differ from *x* by an integral multiple of *y*. Thus $(x-r)/y$ is an integral value, even though it might exceed MAXINT, the largest **int** defined in **<values.h>** if it were explicitly computed as an **int**. Both functions return one of the two such *r* smallest in magnitude. **remainder(***x***,***y***)** is the operation specified in ANSI/IEEE Std 754-1985; the result of **fmod(***x***,***y***)** may differ from **remainder**'*s* result by $\pm y$. The magnitude of **remainder**'*s* result cannot exceed half that of *y*; its sign might not agree with either *x* or *y*. The magnitude of **fmod**'*s* result is less than that of *y*; its sign agrees with that of *x*. Neither function can generate an exception as long as both arguments are normal or subnormal.

remainder(*x***,***0***), fmod(***x***,***0***), remainder(∞,***y***),** and **fmod(∞,***y***)** are invalid operations; in IEEE 754 mode (i.e. the **–xlibmieee cc** compilation option), a *NaN* is returned.

scalbn(*x***,***n***)** returns *x*2**n* computed by exponent manipulation rather than by actually performing an exponentiation or a multiplication. Thus

1 ≤ scalbn(fabs(*x***),–ilogb(***x***)) < 2**

for every *x* except 0, ∞, and *NaN*.

RETURN VALUES

For exceptional cases, **matherr**(3M) tabulates the values to be returned as dictated by various Standards.

SEE ALSO

matherr(3M)

NAME

ieee_test, logb, scalb, significand – IEEE test functions for verifying standard compliance

SYNOPSIS

cc [*flag* ...] *file* ... **–lm** [*library* ...]

#include <math.h>

double logb(double *x***);**

double scalb(double *x***, double** *y***);**

double significand(double *x***);**

MT-LEVEL

MT-Safe

DESCRIPTION

These functions allow users to verify compliance to ANSI/IEEE Std 754-1985 by running certain test vectors distributed by the University of California.

logb(*x***)** returns the unbiased exponent of *x* in floating-point format, for exercising the logb(L) test vector. **logb(±∞) = +∞**; **logb(0) =** –∞ with a division by 0 exception.

scalb(*x***, (double)***n***)** returns *x* ∗ *2**n* computed by exponent manipulation rather than by actually performing an exponentiation or a multiplication, for exercising the scalb(S) test vector. Thus

0 ≤ scalb(fabs(*x***),–logb(***x***)) < 2**

for every *x* except 0, ∞ and *NaN*. **scalb(***x***,***y***)** is not defined when *y* is not an integral value.

If *x* equals *sig* ∗ *2**n* with *1 ≤ sig < 2*, then **significand(***x***)** returns *sig* for exercising the fraction-part(F) test vector. **significand(***x***)** is not defined when *x* is either 0, ±∞ or *NaN*.

RETURN VALUES

For exceptional cases, **matherr**(3M) tabulates the values to be returned as dictated by various Standards.

SEE ALSO

matherr(3M)

NAME

index, rindex – string operations

SYNOPSIS

```
#include <strings.h>
char *index(const char *s, int c);
char *rindex(const char *s, int c);
```

DESCRIPTION

These functions operate on null-terminated strings.

index() returns a pointer to the first occurrence of character *c* in string *s*, and **rindex()** returns a pointer to the last occurrence of character *c* in string *s*. Both **index()** and **rindex()** return a null pointer if *c* does not occur in the string. The null character terminating a string is considered to be part of the string.

SEE ALSO

string(3C), **bstring**(3C), **malloc**(3C)

NOTES

On most modern computer systems, you can *not* use a null pointer to indicate a null string. A null pointer is an error and results in an abort of the program. If you wish to indicate a null string, you must have a pointer that points to an explicit null string. On some implementations of the C language on some machines, a null pointer, if dereferenced, would yield a null string; this highly non-portable trick was used in some programs. Programmers using a null pointer to represent an empty string should be aware of this portability issue; even on machines where dereferencing a null pointer does not cause an abort of the program, it does not necessarily yield a null string.

NAME

inet, inet_addr, inet_network, inet_makeaddr, inet_lnaof, inet_netof, inet_ntoa – Internet address manipulation

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>
#include <netinet/in.h>
#include <arpa/inet.h>

unsigned long inet_addr(const char *cp);

unsigned long inet_network(const char *cp);

struct in_addr inet_makeaddr(const int net, const int lna);

int inet_lnaof(const struct in_addr in);

int inet_netof(const struct in_addr in);

char *inet_ntoa(const struct in_addr in);
```

MT-LEVEL

Safe

DESCRIPTION

The **inet_addr( )** and **inet_network( )** routines interpret character strings representing numbers expressed in the Internet standard '**.**' notation, returning numbers suitable for use as Internet addresses and Internet network numbers, respectively. The routine **inet_makeaddr( )** takes an Internet network number and a local network address and constructs an Internet address from it. The routines **inet_netof( )** and **inet_lnaof( )** break apart Internet host addresses, returning the network number and local network address part, respectively.

The routine **inet_ntoa( )** returns a pointer to a string in the base 256 notation "d.d.d.d" described below.

All Internet addresses are returned in network order (bytes ordered from left to right). All network numbers and local address parts are returned as machine format integer values.

INTERNET ADDRESSES

Values specified using the '**.**' notation take one of the following forms:

a.b.c.d
a.b.c
a.b
a

When four parts are specified, each is interpreted as a byte of data and assigned, from left to right, to the four bytes of an Internet address.

When a three part address is specified, the last part is interpreted as a 16-bit quantity and placed in the right most two bytes of the network address. This makes the three part address format convenient for specifying Class B network addresses as "128.net.host".

When a two part address is supplied, the last part is interpreted as a 24-bit quantity and placed in the right most three bytes of the network address. This makes the two part address format convenient for specifying Class A network addresses as "net.host".

When only one part is given, the value is stored directly in the network address without any byte rearrangement.

All numbers supplied as "parts" in a '**.**' notation may be decimal, octal, or hexadecimal, as specified in the C language (that is, a leading 0x or 0X implies hexadecimal; otherwise, a leading 0 implies octal; otherwise, the number is interpreted as decimal).

RETURN VALUES

The value –**1** is returned by **inet_addr()** and **inet_network()** for malformed requests.

The routines **inet_netof()** and **inet_lnaof()** break apart Internet host addresses, returning the network number and local network address part, respectively.

The routine **inet_ntoa()** returns a pointer to a string in the base 256 notation "d.d.d.d" described below.

SEE ALSO

gethostbyname(3N), **getnetbyname**(3N), **hosts**(4), **networks**(4)

NOTES

The return value from **inet_ntoa()** points to a buffer which is overwritten on each call. This buffer is implemented as thread-specific data in multithreaded applications.

BUGS

The problem of host byte ordering versus network byte ordering is confusing. A simple way to specify Class C network addresses in a manner similar to that for Class B and Class A is needed.

NAME

initgroups – initialize the supplementary group access list

SYNOPSIS

```
#include <grp.h>
#include <sys/types.h>
int initgroups(const char *name, gid_t basegid);
```

MT-LEVEL

Unsafe

DESCRIPTION

initgroups( ) reads the group database to get the group membership for the user specified by *name* and then initializes the supplementary group access list of the calling process (see **getgrnam**(3C) and **getgroups**(2)). The *basegid* group id is also included in the supplementary group access list. This is typically the real group id from the user database.

While scanning the group database, if the number of groups, including the *basegid* entry, exceeds {**NGROUPS_MAX**}, subsequent group entries are ignored.

RETURN VALUES

Upon successful completion, a value of 0 is returned. Otherwise, a value of –1 is returned and **errno** is set to indicate the error.

ERRORS

initgroups( ) will fail and not change the supplementary group access list if:

EPERM The effective user id is not superuser.

SEE ALSO

getgroups(2), **getgrnam**(3C)

NAME

insque, remque – insert / remove element from a queue

SYNOPSIS

```
include <search.h>
void insque(struct qelem *elem, struct qelem *pred);
void remque(struct qelem *elem);
```

MT-LEVEL

Unsafe

DESCRIPTION

insque( ) and **remque( )** manipulate queues built from doubly linked lists. Each element in the queue must be in the following form:

```
struct qelem {
        struct  qelem *q_forw;
        struct  qelem *q_back;
        char    q_data[ ];
};
```

insque( ) inserts *elem* in a queue immediately after *pred.* **remque( )** removes an entry *elem* from a queue.

NAME

isastream – test a file descriptor

SYNOPSIS

```
#include <stropts.h>
int isastream(int fildes);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The function **isastream( )** determines if a file descriptor represents a STREAMS file. *fildes* refers to an open file descriptor.

RETURN VALUES

If successful, **isastream( )** returns **1** if *fildes* represents a STREAMS file, and **0** if not. On failure, **isastream( )** returns **–1** with **errno** set to indicate an error.

ERRORS

Under the following conditions, **isastream( )** fails and sets **errno** to:

EBADF *fildes* is not a valid file descriptor.

SEE ALSO

streamio(7I)

STREAMS Programming Guide

NAME

isencrypt – determine whether a buffer of characters is encrypted

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

int isencrypt(const char ∗*fbuf***, size_t** *ninbuf***);**

MT-LEVEL

MT-Safe

DESCRIPTION

isencrypt( ) uses heuristics to determine whether a buffer of characters is encrypted. It requires two arguments: a pointer to an array of characters and the number of characters in the buffer.

isencrypt( ) assumes that the file is not encrypted if all the characters in the first block are ASCII characters. If there are non-ASCII characters in the first *ninbuf* characters, **isencrypt( )** assumes that the buffer is encrypted if the **setlocale( ) LC_CTYPE** category is set to **C** or **ascii**.

If the **LC_CTYPE** category is set to a value other than **C** or **ascii**, then **isencrypt( )** uses a combination of heuristics to determine if the buffer is encrypted. If *ninbuf* has at least 64 characters, a chi-square test is used to determine if the bytes in the buffer have a uniform distribution; and **isencrypt( )** assumes the buffer is encrypted if it does. If the buffer has less than 64 characters, a check is made for null characters and a terminating new-line to determine whether the buffer is encrypted.

RETURN VALUES

If the buffer is encrypted, 1 is returned; otherwise zero is returned.

SEE ALSO

setlocale(3C)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

isnan, isnand, isnanf, finite, fpclass, unordered – determine type of floating-point number

SYNOPSIS

```
#include <ieeefp.h>
int isnand(double dsrc);
int isnanf(float fsrc);
int finite(double dsrc);
fpclass_t fpclass(double dsrc);
int unordered(double dsrc1, double dsrc2);
#include <math.h>
int isnan(double dsrc);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The functionalty of **isnan( )** is identical to that of **isnand( )**.

isnanf( ) is implemented as a macro included in the **<ieeefp.h>** header.

fpclass( ) returns the class the *dsrc* belongs to. The 10 possible classes are as follows:

FP_SNAN	signaling NaN
FP_QNAN	quiet NaN
FP_NINF	negative infinity
FP_PINF	positive infinity
FP_NDENORM	negative denormalized non-zero
FP_PDENORM	positive denormalized non-zero
FP_NZERO	negative zero
FP_PZERO	positive zero
FP_NNORM	negative normalized non-zero
FP_PNORM	positive normalized non-zero

None of these routines generate any exception, even for signaling NaNs.

RETURN VALUES

isnan( ), **isnand( )**, and **isnanf( )** return true (1) if the argument *dsrc* or *fsrc* is a NaN; otherwise they return false (0).

finite( ) returns true (1) if the argument *dsrc* is neither infinity nor NaN; otherwise it returns false (0).

unordered( ) returns true (1) if one of its two arguments is unordered with respect to the other argument. This is equivalent to reporting whether either argument is NaN. If neither of the arguments is NaN, false (0) is returned.

SEE ALSO

fpgetround(3C)

NAME

iswalpha, iswupper, iswlower, iswdigit, iswxdigit, iswalnum, iswspace, iswpunct, iswprint, iswcntrl, iswascii, iswgraph, isphonogram, isideogram, isenglish, isnumber, isspecial – Process Code character classification macros and functions

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <widec.h>
#include <wctype.h>

int iswalpha(wint_t *c***);**

MT-LEVEL

MT-Safe with exceptions

WIDE CHARACTER CLASSIFICATION

These functions classify Process Code characters (*wchar_t*) from the primary and supplementary codesets by table lookup. Each is a predicate returning nonzero for true, zero for false. The lookup table, generated by **wchrtbl( )**, contains values for both ASCII and supplementary codesets.

iswalpha(*c***)**	*c* is a Latin alphabet Process Code character, from either the primary or supplementary codesets.
iswupper(*c***)**	*c* is an upper case Latin alphabet Process Code character, from either the primary or supplementary codesets.
iswlower(*c***)**	*c* is a lower case Latin alphabet Process Code character, from either the primary or supplementary codesets.
iswdigit(*c***)**	*c* is a Process Code digit [0-9], from either the primary or supplementary codesets.
iswxdigit(*c***)**	*c* is an ASCII hexadecimal Process Code digit [0-9], [A-F], or [a-f].
iswalnum(*c***)**	*c* is a Process Code Latin letter or a digit, from either the primary or supplementary codesets.
iswspace(*c***)**	*c* is a Process Code space, tab, carriage return, newline, vertical tab, or formfeed, from either the primary or supplementary codesets.
iswpunct(*c***)**	*c* is an ASCII Process Code punctuation character (neither control nor alphanumeric).
iswprint(*c***)**	*c* is a Process Code printing character, from either the primary or supplementary codesets. It includes the *space* character.
iswgraph(*c***)**	*c* is a Process Code visible graphic character, from the primary or supplementary codesets. It does not include the *space* character.
iswcntrl(*c***)**	*c* is a Process Code ASCII delete character or ordinary control character, or a control character from a supplementary codeset.
iswascii(*c***)**	*c* is a Process Code ASCII character.
isphonogram(*c***)**	*c* is a Process Code phonetic language character from a supplementary codeset.

isideogram(*c*)	*c* is a Process Code ideographic language character from a supplementary codeset.
isenglish(*c*)	*c* is a Process Code English language character from a supplementary codeset.
isnumber(*c*)	*c* is a Process Code digit [0-9] from a supplementary codeset.
isspecial(*c*)	*c* is a Process Code special language character from a supplementary codeset.

SEE ALSO **setlocale**(3C), **stdio**(3S), **wconv**(3I), **ascii**(5)

NOTES **iswalpha( )**, **iswupper( )**, **iswlower( )**, **iswdigit( )**, **iswxdigit( )**, **iswalnum( )**, **iswspace( )**, **iswpunct( )**, **iswprint( )**, **iswcntrl( )**, **iswascii( )**, **iswgraph( )**, **isphonogram( )**, **isideogram( )**, **isenglish( )**, **isnumber( )** and **isspecial( )** can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

iswctype – test character for specified class

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <wchar.h>

int iswctype(wint_t *wc,* **wctype_t** *charclass***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **iswctype( )** function determines whether the wide-character code *wc* has the character class *charclass*, returning **TRUE** or **FALSE**. **iswctype( )** is defined on WEOF and wide-character codes corresponding to the valid character encodings in the current locale. If the *wc* argument is not in the domain of the function, the result is undefined. If the value of *charclass* is invalid (that is, not obtained by a call to **wctype**(3I) or *charclass* is invalidated by a subsequent call to **setlocale**(3C) that has affected category **LC_CTYPE**), the result is implementation-dependent.

RETURN VALUE

iswctype( ) returns **0** for **FALSE** and non-zero for **TRUE**.

USAGE

The twelve strings — "alnum", "alpha", "blank" "cntrl", "digit", "graph", "lower", "print", "punct", "space", "upper" and "xdigit" — are reserved for the standard character classes. In the table below, the functions in the left column are equivalent to the functions in the right column.

iswalnum(*wc***)**	**iswctype(***wc*, **wctype(**"alnum"**))**
iswalpha(*wc***)**	**iswctype(***wc*, **wctype(**"alpha"**))**
iswcntrl(*wc***)**	**iswctype(***wc*, **wctype(**"cntrl"**))**
iswdigit(*wc***)**	**iswctype(***wc*, **wctype(**"digit"**))**
iswgraph(*wc***)**	**iswctype(***wc*, **wctype(**"graph"**))**
iswlower(*wc***)**	**iswctype(***wc*, **wctype(**"lower"**))**
iswprint(*wc***)**	**iswctype(***wc*, **wctype(**"print"**))**
iswpunct(*wc***)**	**iswctype(***wc*, **wctype(**"punct"**))**
iswspace(*wc***)**	**iswctype(***wc*, **wctype(**"space"**))**
iswupper(*wc***)**	**iswctype(***wc*, **wctype(**"upper"**))**
iswxdigit(*wc***)**	**iswctype(***wc*, **wctype(**"xdigit"**))**

The call

iswctype(*wc*, **wctype(**"blank"**))**

does not have an equivalent **isw( )** function.

SEE ALSO

iswalnum(3I), **iswalpha**(3I), **iswcntrl**(3I), **iswdigit**(3I), **iswgraph**(3I), **iswlower**(3I), **iswprint**(3I), **iswpunct**(3I), **iswspace**(3I), **iswupper**(3I), **iswxdigit**(3I), **setlocale**(3C), **wctype**(3I)

NAME

kerberos, krb_mk_req, krb_rd_req, krb_kntoln, krb_set_key, krb_get_cred, krb_mk_safe, krb_rd_safe, krb_mk_err, krb_rd_err – Kerberos authentication library

SYNOPSIS

cc [*flag* ...] *file* ... **–lkrb** [*library* ...]

#include <kerberos/krb.h>

extern char ∗**krb_err_txt[ ];**

int krb_mk_req(KTEXT *authent***, const char** ∗*service***, const char** ∗*instance***,**
const char ∗*realm***, const u_long** *checksum***);**

int krb_rd_req(const KTEXT *authent***, const char** ∗ *service***, char** ∗ *instance***,**
const u_long *from_addr***, AUTH_DAT** ∗*ad***, const char** ∗*fn***);**

int krb_kntoln(const AUTH_DAT ∗*ad***, char** ∗*lname***);**

int krb_set_key(const char ∗*key***, const int** *cvt***);**

int krb_get_cred(const char ∗*service***, const char** ∗*instance***, const char** ∗*realm***,**
CREDENTIALS ∗*c***);**

long krb_mk_safe(const u_char ∗*in***, u_char** ∗*out***, const u_long** *in_length***,**
const des_cblock *key***, const struct sockaddr_in** *sender***,**
const struct sockaddr_in *receiver***);**

long krb_rd_safe(const u_char ∗*in***, const u_long** *length***, const des_cblock** *key***,**
const struct sockaddr_in *sender***, const struct sockaddr_in** *receiver***,**
MSG_DAT ∗*msg_data***);**

long krb_mk_err(u_char ∗*out***, const long** *code***, const char** ∗*string***);**

long krb_rd_err(const u_char ∗*in***, const u_long** *length***, long** ∗*code***,**
MSG_DAT *msg_data***);**

MT-LEVEL

Unsafe

DESCRIPTION

This library supports network authentication and various related operations. The library contains many routines beyond those described in this man page, but they are not intended to be used directly. Instead, they are called by the routines that are described, the authentication server and the login program.

krb_err_txt[ ] contains text string descriptions of various Kerberos error codes returned by some of the routines below.

krb_mk_req( ) takes a pointer to a text structure in which an authenticator is to be built. It also takes the name, instance, and realm of the service to be used and an optional checksum. It is up to the application to decide how to generate the checksum. **krb_mk_req( )** then retrieves a ticket for the desired service and creates an authenticator. The authenticator is built in *authent* and is accessible to the calling procedure.

It is up to the application to get the authenticator to the service where it will be read by **krb_rd_req( )**. Unless an attacker possesses the session key contained in the ticket, it will be unable to modify the authenticator. Thus, the checksum can be used to verify the authenticity of the other data that will pass through a connection.

krb_mk_req( ) returns KSUCCESS if successful, otherwise a Kerberos error code as defined in **<kerberos/krb.h>**.

krb_rd_req( ) takes an authenticator of type KTEXT, a service name, an instance, the address of the host originating the request, and a pointer to a structure of type AUTH_DAT which is filled in with information obtained from the authenticator. It also optionally takes the name of the file in which it will find the secret key(s) for the service. If the supplied *instance* is "*", then the first service key with the same service name found in the service key file will be used, and the *instance* argument will be filled in with the chosen instance. This means that the caller must provide space for such an instance name.

If the last argument is the null string (""), **krb_rd_req( )** will use the file **/etc/srvtab** to find its keys. If the last argument is NULL, it will assume that the key has been set by **krb_set_key( )** and will not bother looking further.

krb_rd_req( ) is used to find out information about the principal when a request has been made to a service. It is up to the application protocol to get the authenticator from the client to the service. The authenticator is then passed to **krb_rd_req( )** to extract the desired information.

krb_rd_req( ) returns zero (RD_AP_OK) upon successful authentication. If a packet was forged, modified, or replayed, authentication will fail. If the authentication fails, a non-zero value is returned indicating the particular problem encountered. See **<kerberos/krb.h>** for the list of error codes.

krb_kntoln( ) converts a Kerberos name to a local name. It takes a structure of type AUTH_DAT and uses the name, instance, and realm to determine the corresponding local name. A valid local name is returned if the instance is NULL and the realm is the same as the local realm. The local name returned is the Kerberos name and can be used by an application to change uids, directories, or other parameters. This routine is not an integral part of Kerberos, but is provided to support the use of Kerberos in existing utilities. This routine returns KSUCCESS or KFAILURE.

krb_set_key( ) takes as an argument a DES key. It then creates a key schedule from it and saves the original key to be used as an initialization vector. It is used to set the server's key which must be used to decrypt tickets.

If called with a non-zero second argument, **krb_set_key( )** will first convert the input from a string of arbitrary length to a DES key by encrypting it with a one-way function.

In most cases it should not be necessary to call **krb_set_key( )**. The necessary keys will usually be obtained and set inside **krb_rd_req( )**. **krb_set_key( )** is provided for those applications that do not wish to place the application keys on disk. It returns 0 for success, otherwise a non-zero value.

krb_get_cred( ) searches the caller's ticket file for a ticket for the given *service*, *instance*, and *realm*. If a ticket is found, the given CREDENTIALS structure is filled in with the ticket information.

If the ticket was found, **krb_get_cred( )** returns GC_OK. If the ticket file cannot be found, cannot be read, does not belong to the user (other than root), is not a regular file, or is in the wrong mode, the error GC_TKFIL is returned.

krb_mk_safe( ) creates an authenticated, but unencrypted message from any arbitrary application data, pointed to by *in* and *in_length* bytes long. The private session key, pointed to by *key*, is used to seed the **quad_cksum( )** checksum algorithm used as part of the authentication. *sender* and *receiver* point to the Internet address of the two parties. This message does not provide privacy, but does protect (via detection) against modifications, insertions or replays. The encapsulated message and header are placed in the area pointed to by *out* and the routine returns the length of the output, or –1 indicating an error.

krb_rd_safe( ) authenticates a received **krb_mk_safe( )** message. *in* points to the beginning of the received message, whose length is specified in *in_length*. The private session key, pointed to by *key*, is used to seed the **quad_cksum( )** routine as part of the authentication. *msg_data* is a pointer to a MSG_DAT struct, defined in **<kerberos/krb.h>**. The routine fills in these MSG_DAT fields: the *app_data* field with a pointer to the application data, *app_length* with the length of the *app_data* field, *time_sec* and *time_5ms* with the timestamps in the message, and *swap* with a 1 if the byte order of the receiver is different than that of the sender. (The application must still determine if it is appropriate to byte-swap application data; the Kerberos protocol fields are already taken care of.)

The routine returns zero if successful, or a Kerberos error code. Modified messages and old messages cause errors, but it is up to the caller to check the time sequence of messages, and to check against recently replayed messages.

krb_mk_err( ) constructs an application level error message that may be used along with **krb_mk_safe( )**. *out* is a pointer to the output buffer, *code* is an application specific error code, and *string* is an application specific error string. This routine returns the length of the error reply.

krb_rd_err( ) unpacks a received **krb_mk_err( )** message. *in* points to the beginning of the received message, whose length is specified in *in_length*. *code* is a pointer to a value to be filled in with the error value provided by the application. *msg_data* is a pointer to a MSG_DAT struct, defined in **<kerberos/krb.h>**. The routine fills in these MSG_DAT fields: the *app_data* field with a pointer to the application error text, *app_length* with the length of the *app_data* field, and *swap* with a 1 if the byte order of the receiver is different than that of the sender. (The application must still determine if it is appropriate to byte-swap application data; the Kerberos protocol fields are already taken care of).

The routine returns zero if the error message has been successfully received, or a Kerberos error code.

The KTEXT structure is used to pass around text of varying lengths. It consists of a buffer for the data, and a length. **krb_rd_req( )** takes an argument of this type containing the authenticator, and **krb_mk_req( )** returns the authenticator in a structure of this type. KTEXT itself is really a pointer to the structure. The actual structure is of type KTEXT_ST.

The AUTH_DAT structure is filled in by **krb_rd_req( )**. It must be allocated before calling **krb_rd_req( )**, and a pointer to it is passed. The structure is filled in with data obtained from Kerberos. The MSG_DAT structure is filled in by either **krb_rd_safe( )** or **krb_rd_err( )**. It must be allocated before the call and a pointer to it is passed. The structure is filled in with data obtained from Kerberos.

FILES

/usr/lib/libkrb.*
/etc/aname
/etc/srvtab
/tmp/tkt*uid*

SEE ALSO

kerberos(1), **kerberos_rpc**(3N), **krb_realmofhost**(3N), **krb_sendauth**(3N), **krb_set_tkt_string**(3N), **krb.conf**(4), **krb.realms**(4)

NOTES

These interfaces are unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

BUGS

The caller of **krb_rd_req( )** and **krb_rd_safe( )** must check time order and for replay attempts.

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RESTRICTIONS

NAME

kerberos_rpc, authkerb_getucred, authkerb_seccreate, svc_kerb_reg – library routines for remote procedure calls using Kerberos authentication

MT-LEVEL

Unsafe

DESCRIPTION

RPC library routines allow C programs to make procedure calls on other machines across the network.

RPC supports various authentication flavors. Among them are:

AUTH_NONE	(none) no authentication.
AUTH_SYS	Traditional UNIX-style authentication.
AUTH_DES	DES encryption-based authentication.
AUTH_KERB	Kerberos encryption-based authentication.

The **authkerb_getucred( )**, **authkerb_seccreate( )**, and **svc_kerb_reg( )** routines implement the **AUTH_KERB** authentication flavor. The **kerbd** daemon (see **kerbd**(1M)) must be running for the **AUTH_KERB** authentication system to work for kernel based services such as NFS, and **kinit**(1) must have been run by the user in all cases. Only the **AUTH_KERB** style of authentication is discussed here. For information about the **AUTH_NONE** and **AUTH_SYS** styles of authentication, refer to **rpc_clnt_auth**(3N). For information about the **AUTH_DES** style of authentication, refer to **secure_rpc**(3N).

Routines

See **rpc**(3N) for the definition of the **AUTH** data structure.

cc [*flag* ...] *file* ... **–lkrb** [*library* ...]

#include <rpc/rpc.h>

#include <sys/types.h>

int authkerb_getucred(const struct svc_rqst ∗*rqst***, uid_t** ∗*uidp***, gid_t** ∗*gidp***, short** ∗*gidlenp***, int** *gidlist***[NGROUPS]);**

authkerb_getucred( ) is used on the server side for converting an **AUTH_KERB** credential received in an RPC request, which is operating system independent, into an **AUTH_SYS** credential. This routine returns **1** if it succeeds, **0** if it fails.

∗*uidp* is set to the numerical ID of the user associated with the RPC request referenced by *rqst*. ∗*gidp* is set to the numerical ID of the user's group. The numerical IDs of the other groups to which the user belongs are stored in *gidlist*[ ]. ∗*gidlenp* is set to the number of valid group ID entries returned in *gidlist*[ ]. All information returned by this routine is based on the Kerberos principal name contained in *rqst*. This principal name is taken to be the login name of the user, and the IDs returned are the same as if that user had physically logged in to the system.

AUTH ∗**authkerb_seccreate(const char** ∗*service*, **const char** ∗*srv_inst*, **const char** ∗*realm*, **const unsigned int** *window*, **const char** ∗*timehost*, **int** ∗*status*);

authkerb_seccreate( ) is used on the client side to return an authentication handle that will enable the use of the Kerberos authentication system. The first parameter *service* is the Kerberos principal name of the service to be used. This name is generally a constant with respect to the service being used. *srv_instance* is the instance of the service to be called, and may be NULL to indicate any instance. *realm* is the Kerberos realm name of the desired service. If it is NULL, then the local default realm will be used.

The fourth parameter is the *window* on the validity of the client credential, given in seconds. If the difference in time between the client's clock and the server's clock exceeds *window*, the server will reject the client's credentials, and the clock will have to be resynchronized. A small window is more secure than a large one, but choosing too small of a window will increase the frequency of resynchronizations because of clock drift.

The fifth parameter, *timehost*, is optional. If it is NULL, then the authentication system will assume that the local clock is always in sync with the *timehost* clock, and will not attempt resynchronizations. If a timehost is supplied, however, then the system will consult with the remote time service whenever resynchronization is required. This parameter is usually the name of the host on which the server is running.

The final parameter *status* is also optional. If *status* is supplied, then it will be used to return a Kerberos error status codes if an error occurs. If *status* is NULL, then no detailed error codes will be returned.

If **authkerb_seccreate( )** fails, it returns NULL.

int svc_kerb_reg(const SVCXPRT ∗*xprt*, **const char** ∗*name*, **const char** ∗*inst*, **const char** ∗*realm*);

svc_kerb_reg( ) performs registration tasks in the server which are required before **AUTH_KERB** requests can be processed. *xprt* is the RPC tranport to which this information is to be associated. If *xprt* is NULL then this registration will be effective for any requests arriving on transports that have not been specifically registered.

The other parameters describe the Kerberos principal identity that this server will take on. This must be the same identity that the clients will use when requesting Kerberos tickets for authentication. *name* is the principal name of the service and must be provided. *inst* is the instance. This parameter may be NULL to specify the NULL instance of the service. Most common would be for *inst* to be "∗" which allows the Kerberos library to determine the correct instance to use, such as the hostname that the service is running on. *realm* is the Kerberos realm name to use in validating tickets. If it is NULL, then the local default realm will be used.

svc_kerb_reg( ) should generally be called immediately before **svc_run( )**. It returns 0 if it succeeds, and –1 if it fails.

SEE ALSO

kerberos(1), **kinit**(1), **kerbd**(1M), **rpc**(3N), **rpc_clnt_auth**(3N), **secure_rpc**(3N)

NOTES

These interfaces are unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

killpg – send signal to a process group

SYNOPSIS

```
#include <signal.h>
int killpg(pid_t pgrp, int sig);
```

DESCRIPTION

killpg( ) sends the signal *sig* to the process group *pgrp*. See **signal**(5) for a list of signals.

The real or effective user ID of the sending process must match the real or saved set-user ID of the receiving process, unless the effective user ID of the sending process is the privileged user. A single exception is the signal **SIGCONT**, which may always be sent to any descendant of the current process.

RETURN VALUES

Upon successful completion, a value of 0 is returned. Otherwise, a value of –1 is returned and the global variable **errno** is set to indicate the error.

ERRORS

killpg( ) will fail and no signal will be sent if any of the following occur:

EINVAL *sig* is not a valid signal number.

EPERM The effective user ID of the sending process is not privileged user, and neither its real nor effective user ID matches the real or saved set-user ID of one or more of the target processes.

ESRCH No processes were found in the specified process group.

SEE ALSO

kill(2), **setpgrp**(2), **sigaction**(2), **signal**(5)

NAME

krb_realmofhost, krb_get_phost, krb_get_krbhst, krb_get_admhst, krb_get_lrealm – additional Kerberos utility routines

SYNOPSIS

cc [*flag* ...] *file* ... **–lkrb** [*library* ...]

#include <kerberos/krb.h>
#include <netinet/in.h>

char ∗**krb_realmofhost(const char** ∗*host***);**

char ∗**krb_get_phost(const char** ∗*alias***);**

int krb_get_krbhst(char ∗*host***, const char** ∗*realm***, const int** *n***);**

int krb_get_admhst(char ∗*host***, const char** ∗*realm***, const int** *n***);**

int krb_get_lrealm(char ∗*realm***, const int** *n***);**

MT-LEVEL

Unsafe

DESCRIPTION

krb_realmofhost( ) returns the Kerberos realm of the host *host*, as determined by the translation table **/etc/krb.realms**. *host* should be the fully-qualified domain-style primary host name of the host in question. In order to prevent certain security attacks, this routine must either have a prior knowledge of a host's realm, or obtain such information securely.

The format of the translation file is described by **krb.realms**(4). If *host* exactly matches a host_name line, the corresponding realm is returned. Otherwise, if the domain portion of *host* matches a domain_name line, the corresponding realm is returned. If *host* contains a domain, but no translation is found, *host*'s domain is converted to upper-case and returned. If *host* contains no discernible domain, or an error occurs, the local realm name, as supplied by **krb_get_lrealm( )**, is returned.

krb_get_phost( ) converts the hostname *alias* (which can be either an official name or an alias) into the instance name to be used in obtaining Kerberos tickets for most services, including the Berkeley rcmd suite (rlogin, rcp, rsh). The current convention is to return the first segment of the official domain-style name after conversion to lower case.

krb_get_krbhst( ) fills in *host* with the hostname of the *n*th host running a Kerberos key distribution center (KDC) for realm *realm*, as specified in the configuration file (**/etc/krb.conf** or **krb.conf** NIS map). The configuration format is described by **krb.conf**(4). If the host is successfully filled in, the routine returns KSUCCESS. If the file (or NIS map) cannot be accessed, and *n* equals 1, then the hostname **kerberos** is filled in, and KSUCCESS is returned. If there are fewer than *n* hosts running a Kerberos KDC for the requested realm, or the configuration file is malformed, the routine returns KFAILURE.

When there is both a local **/etc/krb.conf** and a **krb.conf** NIS map, then the entries are counted starting first with the local file, then continuing with the NIS map. For example, if the local **/etc/krb.conf** file contains two entries which match *realm*, and the NIS map contains one matching entry, then there are three possible matches that **krb_get_krbhst( )** can return. The first two (for *n* values 1 and 2) come from the file, and the third (for *n*

equal to 3) comes from the map.

krb_get_admhst( ) fills in *host* with the hostname of the *n*th host running a Kerberos KDC database administration server for realm *realm*, as specified in **/etc/krb.conf**. If the file cannot be opened or is malformed, or there are fewer than *n* hosts running a Kerberos KDC database administration server, the routine returns KFAILURE.

The character arrays used as return values for **krb_get_krbhst( )** and **krb_get_admhst( )** should be large enough to hold any hostname.

krb_get_lrealm( ) fills in *realm* with the *n*th realm of the local host, as specified in the configuration file. *realm* should be at least REALM_SZ (from **<kerberos/krb.h>**) characters long. The return value is either KSUCCESS or KFAILURE.

SEE ALSO

kerberos(3N), **krb.conf**(4), **krb.realms**(4)

FILES

/etc/krb.realms	translation file for host-to-realm mapping.
/etc/krb.conf	local realm-name and realm / server configuration file.

NOTES

These interfaces are unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

BUGS

The current convention for instance names is too limited; the full domain name should be used.

krb_get_lrealm( ) currently only supports *n* equal to 1. It should really consult the user's ticket cache to determine the user's current realm, rather than consulting a file on the host.

NAME

krb_sendauth, krb_recvauth, krb_net_write, krb_net_read – Kerberos routines for sending authentication via network stream sockets

SYNOPSIS

cc [*flag* ...] *file* ... **–lkrb** [*library* ...]

#include <kerberos/krb.h>
#include <netinet/in.h>

int krb_sendauth(const long *options***, const int** *fd***, KTEXT** *ktext***, const char** ∗*service***,**
const char ∗*inst***, const char** ∗*realm***, const u_long** *checksum,*
MSG_DATA ∗*msg_data***, CREDENTIALS** ∗*cred***, Key_schedule** *schedule***,**
const struct sockaddr_in ∗*laddr***, const struct sockaddr_in** ∗*faddr***,**
const char ∗*version***);**

int krb_recvauth(const long *options***, const int** *fd***, KTEXT** *ktext***, const char** ∗*service***,**
char ∗*inst***, const struct sockaddr_in** ∗*faddr***, const struct sockaddr_in** ∗*laddr,*
AUTH_DAT ∗*auth_data***, const char** ∗*filename***, Key_version** *schedule***,**
char ∗*version***);**

int krb_net_write(const int *fd***, const char** ∗*buf***, const int** *len***);**

int krb_net_read(const int *fd***, char** ∗*buf***, const int** *len***);**

MT-LEVEL

Unsafe

DESCRIPTION

These functions, which are built on top of the core Kerberos library, provide a convenient means for client and server programs to send authentication messages to one another through network connections.

The **krb_sendauth( )** function sends an authenticated ticket from the client program to the server program by writing the ticket to a network socket.

The **krb_recvauth( )** function receives the ticket from the client by reading from a network socket.

krb_sendauth( )

This function writes the ticket to the network socket specified by the file descriptor *fd*, returning KSUCCESS if the write proceeds successfully, and an error code if it does not.

The *ktext* argument should point to an allocated KTEXT_ST structure. The *service*, *inst*, and *realm* arguments specify the server program's Kerberos principal name, instance, and realm. If you are writing a client that uses the local realm exclusively, you can set the *realm* argument to NULL.

The *version* argument allows the client program to pass an application-specific version string that the server program can then match against its own version string. The *version* string can be up to KSEND_VNO_LEN (see **<kerberos/krb.h>**) characters in length.

The *checksum* argument can be used to pass checksum information to the server program. The client program is responsible for specifying this information. This checksum information is difficult to corrupt because **krb_sendauth( )** passes it over the network in encrypted form. The *checksum* argument is passed as the checksum argument to **krb_mk_req( )** (see **kerberos**(3N)).

You can set **krb_sendauth( )**'s other arguments to NULL unless you want the client and server programs to mutually authenticate themselves. In the case of mutual authentication, the client authenticates itself to the server program, and demands that the server in turn authenticate itself to the client.

krb_sendauth( ) and Mutual Authentication

If you want mutual authentication, make sure that you read all pending data from the local socket before calling **krb_sendauth( )**. Set **krb_sendauth( )**'s *options* argument to KOPT_DO_MUTUAL (this macro is defined in **<kerberos/krb.h>**); make sure that the *laddr* argument points to the address of the local socket, and that *faddr* points to the foreign socket's network address.

krb_sendauth( ) fills in the other arguments — *msg_data*, *cred*, and *schedule* — before sending the ticket to the server program. You must, however, allocate space for these arguments before calling the function.

krb_sendauth( ) supports two other options: KOPT_DONT_MK_REQ and KOPT_DONT_CANON. If called with *options* set as KOPT_DONT_MK_REQ, **krb_sendauth( )** will not use the **krb_mk_req( )** (see **kerberos**(3N)) function to retrieve the ticket from the Kerberos server. The *ktext* argument must point to an existing ticket and authenticator (such as would be created by **krb_mk_req( )**), and the *service*, *inst*, and *realm* arguments can be set to NULL.

If called with *options* set as KOPT_DONT_CANON, **krb_sendauth( )** will not convert the service's instance to canonical form using **krb_get_phost( )** (see **krb_realmofhost**(3N)).

If you want to call **krb_sendauth( )** with a multiple *options* specification, construct *options* as a bitwise-OR of the options you want to specify.

krb_recvauth( )

The **krb_recvauth( )** function reads a ticket / authenticator pair from the socket pointed to by the *fd* argument. Set the *options* argument as a bitwise-OR of the options desired. Currently only KOPT_DO_MUTUAL is useful to the receiver.

The *ktext* argument should point to an allocated KTEXT_ST structure. **krb_recvauth( )** fills *ktext* with the ticket / authenticator pair read from *fd*, then passes it to **krb_rd_req( )** (see **kerberos**(3N)).

The *service* and *inst* arguments specify the expected service and instance for which the ticket was generated. They are also passed to **krb_rd_req( )** (see **kerberos**(3N)). The *inst* argument may be set to "*" if the caller wishes **krb_mk_req( )** (see **kerberos**(3N)) to fill in the instance used (note that there must be space in the *inst* argument to hold a full instance name, see **krb_mk_req( )** on **kerberos**(3N)).

The *faddr* argument should point to the address of the peer which is presenting the ticket. It is also passed to **krb_rd_req( )** (see **kerberos**(3N)).

If the client and server plan to mutually authenticate one another, the *laddr* argument should point to the local address of the file descriptor. Otherwise you can set this argument to NULL.

The *auth_data* argument should point to an allocated AUTH_DAT area. It is passed to and filled in by **krb_rd_req( )** (see **kerberos**(3N)). The checksum passed to the corresponding **krb_sendauth( )** is available as part of the filled-in AUTH_DAT area.

The *filename* argument specifies the filename which the service program should use to obtain its service key. **krb_recvauth( )** passes *filename* to the **krb_rd_req( )** function, see **kerberos**(3N), If you set this argument to "", **krb_rd_req( )** looks for the service key in the file **/etc/srvtab**.

If the client and server are performing mutual authentication, the *schedule* argument should point to an allocated Key_schedule. Otherwise it is ignored and may be NULL.

The *version* argument should point to a character array of at least KSEND_VNO_LEN characters. It is filled in with the version string passed by the client to **krb_sendauth( )**.

krb_net_write( ) and krb_net_read( )

The **krb_net_write( )** function emulates the **write**(2) system call, but guarantees that all data specified is written to *fd* before returning, unless an error condition occurs.

The **krb_net_read( )** function emulates the **read**(2) system call, but guarantees that the requested amount of data is read from *fd* before returning, unless an error condition occurs.

SEE ALSO

read(2), **write**(2), **kerberos**(3N), **kerberos_rpc**(3N), **krb_realmofhost**(3N)

NOTES

These interfaces are unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

BUGS

krb_sendauth( ), **krb_recvauth( )**, **krb_net_write( )**, and **krb_net_read( )** will not work properly on sockets set to non-blocking I/O mode.

AUTHOR

John T. Kohl, MIT Project Athena

RESTRICTIONS

NAME

krb_set_tkt_string – set Kerberos ticket cache file name

SYNOPSIS

cc [*flag* ...] *file* ... **–lkrb** [*library* ...]

#include <kerberos/krb.h>

void krb_set_tkt_string(const char ∗*filename*);

MT-LEVEL

Unsafe

DESCRIPTION

krb_set_tkt_string( ) sets the name of the file that holds the user's cache of Kerberos server tickets and associated session keys.

The string *filename* passed in is copied into local storage. Only MAXPATHLEN-1 (see **<sys/param.h>**) characters of the filename are copied in for use as the cache file name.

This routine should be called during initialization, before other Kerberos routines are called; otherwise the routines which fetch the ticket cache file name may be called and return an undesired ticket file name until this routine is called.

FILES

/tmp/tkt*uid* default ticket file name, unless the environment variable KRBTKFILE is set. *uid* denotes the user's uid, in decimal.

SEE ALSO

kerberos(3N)

NOTES

This interface is unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

kstat – kernel statistics facility

DESCRIPTION

The **kstat** facility is a general-purpose mechanism for providing kernel statistics to users.

The kstat model

The kernel maintains a linked list of statistics structures, or kstats. Each kstat has a common header section and a type-specific data section. The header section is defined by the **kstat_t** structure:

kstat header

```
typedef long        kid_t;                      /* unique kstat id */
typedef struct kstat {
    /*
     * Fields relevant to both kernel and user
     */
    hrtime_t        ks_crtime;                  /* creation time */
    struct kstat    *ks_next;                   /* kstat chain linkage */
    kid_t           ks_kid;                     /* unique kstat ID */
    char            ks_module[KSTAT_STRLEN];    /* module name */
    uchar_t         ks_resv;                    /* reserved */
    int             ks_instance;                /* module's instance */
    char            ks_name[KSTAT_STRLEN];      /* kstat name */
    uchar_t         ks_type;                    /* kstat data type */
    char            ks_class[KSTAT_STRLEN];     /* kstat class */
    uchar_t         ks_flags;                   /* kstat flags */
    void            *ks_data;                   /* kstat type-specific data */
    ulong_t         ks_ndata;                   /* # of data records */
    ulong_t         ks_data_size;               /* size of kstat data section */
    hrtime_t        ks_snaptime;                /* time of last data snapshot */

    /*
     * Fields relevant to kernel only
     */
    int             (*ks_update)(struct kstat *, int);
    void            *ks_private;
    int             (*ks_snapshot)(struct kstat *, void *, int);
    void            *ks_lock;
} kstat_t;
```

The fields that are of significance to the user are:

ks_crtime The time the kstat was created. This allows you to compute the rates of various counters since the kstat was created; "rate since boot" is replaced by the more general concept of "rate since kstat creation".

All times associated with kstats (e.g. creation time, last snapshot time, **kstat_timer_t** and **kstat_io_t** timestamps, etc.) are 64-bit nanosecond values. The accuracy of kstat timestamps is machine dependent, but the precision (units) is the same across all platforms. See **gethrtime**(3C) for

general information about high-resolution timestamps.

ks_next kstats are stored as a linked list, or chain. **ks_next** points to the next kstat in the chain.

ks_kid A unique identifier for the kstat.

ks_module, **ks_instance**
contain the name and instance of the the module that created the kstat. In cases where there can only be one instance, **ks_instance** is 0.

ks_name gives a meaningful name to a kstat. The full kstat namespace is <**ks_module**, **ks_instance**, **ks_name** so the name only need be unique within a module.

ks_type The type of data in this kstat. kstat data types are discussed below.

ks_class Each kstat can be characterized as belonging to some broad class of statistics, e.g. disk, tape, net, vm, streams, etc. This field can be used as a filter to extract related kstats. The following values are currently in use: **disk**, **tape**, **controller**, **net**, **rpc**, **vm**, **kvm**, **hat**, **kmem**, **kmem_cache**, **kstat**, and **misc**. (The kstat class encompasses things like *kstat_types*.)

ks_data, **ks_ndata**, **ks_data_size**
ks_data is a pointer to the kstat's data section. The type of data stored there depends on **ks_type**.

ks_ndata indicates the number of data records. Only some kstat types support multiple data records. Currently, **KSTAT_TYPE_RAW, KSTAT_TYPE_NAMED** and **KSTAT_TYPE_TIMER** kstats support multiple data records. **KSTAT_TYPE_INTR** and **KSTAT_TYPE_IO** kstats support only one data record.

ks_data_size is the total size of the data section, in bytes.

ks_snaptime The timestamp for the last data snapshot. This allows you to compute activity rates:

rate = (new_count - old_count) / (new_snaptime - old_snaptime);

kstat data types

The following types of kstats are currently available:

```
#define KSTAT_TYPE_RAW      0   /* can be anything */
#define KSTAT_TYPE_NAMED    1   /* name/value pairs */
#define KSTAT_TYPE_INTR     2   /* interrupt statistics */
#define KSTAT_TYPE_IO       3   /* I/O statistics */
#define KSTAT_TYPE_TIMER    4   /* event timers */
```

To get a list of all kstat types currently supported in the system, tools can read out the standard system kstat *kstat_types* (full name spec is <''*unix*'', *0*, ''*kstat_types*''>). This is a **KSTAT_TYPE_NAMED** kstat in which the *name* field describes the type of kstat, and the *value* field is the kstat type number (e.g., **KSTAT_TYPE_IO** is type 3 -- see above).

Raw kstat

KSTAT_TYPE_RAW: raw data

The "raw" kstat type is just treated as an array of bytes. This is generally used to export well-known structures, like *sysinfo*.

Name=value kstat

KSTAT_TYPE_NAMED: A list of arbitrary *name=value* statistics.

```
typedef struct kstat_named {
    char        name[KSTAT_STRLEN];   /* name of counter */
    uchar_t     data_type;            /* data type */
    union {
        char            c[16];        /* enough for 128-bit ints */
        long            l;
        ulong_t         ul;
        longlong_t      ll;
        u_longlong_t    ull;
        float           f;
        double          d;
        } value;                      /* value of counter */
} kstat_named_t;

#define KSTAT_DATA_CHAR         0
#define KSTAT_DATA_LONG         1
#define KSTAT_DATA_ULONG        2
#define KSTAT_DATA_LONGLONG     3
#define KSTAT_DATA_ULONGLONG    4
#define KSTAT_DATA_FLOAT        5
#define KSTAT_DATA_DOUBLE       6
```

Interrupt kstat

KSTAT_TYPE_INTR: Interrupt statistics.

An interrupt is a hard interrupt (sourced from the hardware device itself), a soft interrupt (induced by the system via the use of some system interrupt source), a watchdog interrupt (induced by a periodic timer call), spurious (an interrupt entry point was entered but there was no interrupt to service), or multiple service (an interrupt was detected and serviced just prior to returning from any of the other types).

```
#define KSTAT_INTR_HARD         0
#define KSTAT_INTR_SOFT         1
#define KSTAT_INTR_WATCHDOG     2
#define KSTAT_INTR_SPURIOUS     3
#define KSTAT_INTR_MULTSVC      4
#define KSTAT_NUM_INTRS         5

typedef struct kstat_intr {
    ulong_t intrs[KSTAT_NUM_INTRS]; /* interrupt counters */
} kstat_intr_t;
```

Event timer kstat

KSTAT_TYPE_TIMER: Event timer statistics.

These provide basic counting and timing information for any type of event.

```
typedef struct kstat_timer {
    char            name[KSTAT_STRLEN];   /* event name */
    uchar_t         resv;                 /* reserved */
    u_longlong_t    num_events;           /* number of events */
    hrtime_t        elapsed_time;         /* cumulative elapsed time */
    hrtime_t        min_time;             /* shortest event duration */
    hrtime_t        max_time;             /* longest event duration */
    hrtime_t        start_time;           /* previous event start time */
    hrtime_t        stop_time;            /* previous event stop time */
} kstat_timer_t;
```

I/O kstat

KSTAT_TYPE_IO: I/O statistics.

```
typedef struct kstat_io {
    /*
     * Basic counters.
     */
    u_longlong_t    nread;                /* number of bytes read */
    u_longlong_t    nwritten;             /* number of bytes written */
    ulong_t         reads;                /* number of read operations */
    ulong_t         writes;               /* number of write operations */
    /*
    * Accumulated time and queue length statistics.
    *
    * Time statistics are kept as a running sum of "active" time.
    * Queue length statistics are kept as a running sum of the
    * product of queue length and elapsed time at that length --
    * i.e., a Riemann sum for queue length integrated against time.
    *
    *                    ^
    *                    |                                 _________
    *                    8                                | i4      |
    *                    |                                |         |
    *          Queue     6                                |         |
    *          Length    |          _________             |         |
    *                    4         | i2      |_______|    |         |
    *                    |         |           i3          |         |
    *                    2_______|                         |
    *                    |   i1                            |
    *                    |_____________________________________|
    *                    Time->  t1        t2        t3        t4
    *
    * At each change of state (entry or exit from the queue),
```

```
 * we add the elapsed time (since the previous state change)
 * to the active time if the queue length was non-zero during
 * that interval; and we add the product of the elapsed time
 * times the queue length to the running length*time sum.
 *
 * This method is generalizable to measuring residency
 * in any defined system: instead of queue lengths, think
 * of "outstanding RPC calls to server X".
 *
 * A large number of I/O subsystems have at least two basic
 * "lists" of transactions they manage: one for transactions
 * that have been accepted for processing but for which processing
 * has yet to begin, and one for transactions which are actively
 * being processed (but not done). For this reason, two cumulative
 * time statistics are defined here: pre-service (wait) time,
 * and service (run) time.
 *
 * The units of cumulative busy time are accumulated nanoseconds.
 * The units of cumulative length*time products are elapsed time
 * times queue length.
 */
    hrtime_t    wtime;          /* cumulative wait (pre-service) time */
    hrtime_t    wlentime;       /* cumulative wait length*time product */
    hrtime_t    wlastupdate;    /* last time wait queue changed */
    hrtime_t    rtime;          /* cumulative run (service) time */
    hrtime_t    rlentime;       /* cumulative run length*time product */
    hrtime_t    rlastupdate;    /* last time run queue changed */
    ulong_t     wcnt;           /* count of elements in wait state */
    ulong_t     rcnt;           /* count of elements in run state */
} kstat_io_t;
```

Using libkstat

The kstat library, **libkstat**, defines the user interface (API) to the system's kstat facility.

You begin by opening libkstat with **kstat_open**(3K), which returns a pointer to a fully initialized kstat control structure. This is your ticket to subsequent libkstat operations:

```
typedef struct kstat_ctl {
    kid_t       kc_chain_id;    /* current kstat chain ID */
    kstat_t     *kc_chain;      /* pointer to kstat chain */
    int         kc_kd;          /* /dev/kstat descriptor */
} kstat_ctl_t;
```

Only the first two fields, **kc_chain_id** and **kc_chain**, are of interest to **libkstat** clients. (*kc_kd* is the descriptor for **/dev/kstat**, the kernel statistics driver. libkstat functions are built on top of **/dev/kstat ioctl**(2) primitives. Direct interaction with **/dev/kstat** is strongly discouraged, since it is NOT a public interface.)

kc_chain points to your copy of the kstat chain. You typically walk the chain to find and process a certain kind of kstat. For example, to display all I/O kstats:

```
kstat_ctl_t     *kc;
kstat_t         *ksp;
kstat_io_t      kio;

kc = kstat_open();
for (ksp = kc->kc_chain; ksp != NULL; ksp = ksp->ks_next) {
        if (ksp->ks_type == KSTAT_TYPE_IO) {
                kstat_read(kc, ksp, &kio);
                my_io_display(kio);
        }
}
```

kc_chain_id is the kstat chain ID, or KCID, of your copy of the kstat chain. See **kstat_chain_update**(3K) for an explanation of KCIDs.

FILES

/dev/kstat kernel statistics driver

/usr/include/kstat.h

/usr/include/sys/kstat.h

SEE ALSO

ioctl(2), **gethrtime**(3C), **kstat_chain_update**(3K), **kstat_close**(3K), **kstat_data_lookup**(3K), **kstat_lookup**(3K), **kstat_open**(3K), **kstat_read**(3K), **kstat_write**(3K)

NAME

kstat_chain_update – update the kstat header chain

SYNOPSIS

cc [*flag* . . .] *file*. . . **-lkstat** [*library*. . .]

#include <kstat.h>

kid_t ∗**kstat_chain_update(kstat_ctl_t** ∗*kc***);**

DESCRIPTION

kstat_chain_update( ) brings the user's kstat header chain in sync with the kernel's. The kstat chain is a linked list of kstat headers (**kstat_t**'s), pointed to by *kc->kc_chain*, which is initialized by **kstat_open**(3K). This chain constitutes a list of all kstats currently in the system. During normal operation, the kernel will occasionally create new kstats and delete old ones, as various device instances come and go. When this happens, the user's copy of the kstat chain becomes out of date.

kstat_chain_update( ) detects this by comparing the kernel's current kstat chain ID (KCID), which is incremented every time the kstat chain changes, to the user's KCID, *kc->kc_chain_id*. If the KCID's match, **kstat_chain_update( )** does nothing. Otherwise, it deletes any invalid kstat headers from the user's kstat chain, and adds any new ones, and sets *kc->kc_chain_id* to the new KCID. All other kstat headers in the user's kstat chain are unmodified.

RETURN VALUES

kstat_chain_update( ) returns the new KCID if the kstat chain has changed, **0** if it hasn't, or –**1** on failure.

FILES

/dev/kstat kernel statistics driver

SEE ALSO

kstat(3K), **kstat_close**(3K), **kstat_data_lookup**(3K), **kstat_lookup**(3K), **kstat_open**(3K), **kstat_read**(3K), **kstat_write**(3K)

NAME

kstat_lookup, kstat_data_lookup – find a kstat by name

SYNOPSIS

cc [*flag* . . .] *file*. . . **-lkstat** [*library*. . .]

#include <kstat.h>

kstat_t ∗**kstat_lookup(kstat_ctl_t** ∗*kc***, char** ∗*ks_module***, int** *ks_instance***, char** ∗*ks_name***);**

void ∗**kstat_data_lookup(kstat_t** ∗*ksp***, char** ∗*name***);**

DESCRIPTION

kstat_lookup( ) walks down the kstat chain, *kc->kc_chain*, looking for a kstat with the same *ks_module*, *ks_instance*, and *ks_name* fields; this triplet uniquely identifies a kstat. If *ks_module* is **NULL**, *ks_instance* is **-1 ,** or *ks_name* is **NULL**, then those fields will be ignored in the search. For example, **kstat_lookup(NULL, -1, "foo")** will simply find the first kstat with name "foo".

kstat_data_lookup( ) searches the kstat's data section for the record with the specified *name*. This operation is only valid for kstat types which have named data records. Currently, only the **KSTAT_TYPE_NAMED** and **KSTAT_TYPE_TIMER** kstats have named data records.

RETURN VALUES

kstat_lookup( ) returns a pointer to the requested kstat if it is found, or **NULL** if it isn't.

kstat_data_lookup( ) returns a pointer to the requested data record if it is found. If the requested record is not found, or if the kstat type is invalid, **kstat_data_lookup( )** returns **NULL**.

FILES

/dev/kstat kernel statistics driver

SEE ALSO

kstat(3K), **kstat_chain_update**(3K), **kstat_close**(3K), **kstat_open**(3K), **kstat_read**(3K), **kstat_write**(3K)

NAME

kstat_open, kstat_close – initialize kernel statistics facility

SYNOPSIS

cc [*flag* . . .] *file*. . . **-lkstat** [*library*. . .]

#include <kstat.h>

kstat_ctl_t ∗**kstat_open(void);**

int kstat_close(kstat_ctl_t ∗*kc***);**

DESCRIPTION

kstat_open( ) initializes a kstat control structure, which provides access to the kernel statistics library. It returns a pointer to this structure, which must be supplied as the *kc* argument in subsequent **libkstat** function calls.

kstat_close( ) frees all resources that were associated with *kc*. This is done automatically on **exit**(2) and **execve( )** (see **exec**(2)).

RETURN VALUES

kstat_open( ) returns a pointer to a kstat control structure. On failure, it returns **NULL** and no resources are allocated.

kstat_close( ) returns **0** on success, –**1** on failure.

FILES

/dev/kstat kernel statistics driver

SEE ALSO

kstat(3K), **kstat_chain_update**(3K), **kstat_data_lookup**(3K), **kstat_lookup**(3K), **kstat_read**(3K), **kstat_write**(3K)

NAME

kstat_read, kstat_write – read or write kstat data

SYNOPSIS

cc [*flag* . . .] *file*. . . **-lkstat** [*library*. . .]

#include <kstat.h>

kid_t kstat_read(kstat_ctl_t ∗*kc***, kstat_t** ∗*ksp***, void** ∗*buf***);**

kid_t kstat_write(kstat_ctl_t ∗*kc***, kstat_t** ∗*ksp***, void** ∗*buf***);**

DESCRIPTION

kstat_read( ) gets data from the kernel for the kstat pointed to by *ksp*. *ksp->ks_data* is automatically allocated (or reallocated) to be large enough to hold all of the data. *ksp->ks_ndata* is set to the number of data fields, *ksp->ks_data_size* is set to the total size of the data, and *ksp->ks_snaptime* is set to the high-resolution time at which the data snapshot was taken. If *buf* is non-**NULL**, the data is copied from *ksp->ks_data* into *buf*.

kstat_write( ) writes data from *buf*, or from *ksp->ks_data* if *buf* is **NULL**, to the corresponding kstat in the kernel. Only the superuser can use **kstat_write( )**.

RETURN VALUES

On success, **kstat_read( )** and **kstat_write( )** return the current kstat chain ID (KCID). On failure, they return **–1**.

FILES

/dev/kstat kernel statistics driver

SEE ALSO

kstat(3K), **kstat_chain_update**(3K), **kstat_close**(3K), **kstat_data_lookup**(3K), **kstat_lookup**(3K), **kstat_open**(3K)

NAME

kvm_getu, kvm_getcmd – get the u-area or invocation arguments for a process

SYNOPSIS

```
#include <kvm.h>
#include <sys/param.h>
#include <sys/user.h>
#include <sys/proc.h>

struct user *kvm_getu(kvm_t *kd, struct proc *proc);

int kvm_getcmd(kvm_t *kd, struct proc *proc, struct user *u,
        char ***arg, char ***env);
```

MT-LEVEL

Unsafe

DESCRIPTION

kvm_getu() reads the u-area of the process specified by *proc* to an area of static storage associated with *kd* and returns a pointer to it. Subsequent calls to **kvm_getu()** will overwrite this static area.

kd is a pointer to a kernel descriptor returned by **kvm_open**(3K). *proc* is a pointer to a copy (in the current process' address space) of a *proc* structure (obtained, for instance, by a prior **kvm_nextproc**(3K) call).

kvm_getcmd() constructs a list of string pointers that represent the command arguments and environment that were used to initiate the process specified by *proc*.

kd is a pointer to a kernel descriptor returned by **kvm_open**(3K). *u* is a pointer to a copy (in the current process' address space) of a *user* structure (obtained, for instance, by a prior **kvm_getu()** call). If *arg* is not NULL, then the command line arguments are formed into a null-terminated array of string pointers. The address of the first such pointer is returned in *arg*. If *env* is not NULL, then the environment is formed into a null-terminated array of string pointers. The address of the first of these is returned in *env*.

The pointers returned in *arg* and *env* refer to data allocated by **malloc**(3C) and should be freed (by a call to **free()** (see **malloc**(3C)) when no longer needed. Both the string pointers and the strings themselves are deallocated when freed.

Since the environment and command line arguments may have been modified by the user process, there is no guarantee that it will be possible to reconstruct the original command at all. Thus, **kvm_getcmd()** will make the best attempt possible, returning –1 if the user process data is unrecognizable.

RETURN VALUES

On success, **kvm_getu()** returns a pointer to a copy of the u-area of the process specified by *proc*. On failure, it returns NULL.

kvm_getcmd() returns:

0 on success.

–1 on failure.

SEE ALSO **kvm_nextproc**(3K), **kvm_open**(3K), **kvm_read**(3K), **malloc**(3C)

NOTES If **kvm_getcmd( )** returns –1, the caller still has the option of using the command line fragment that is stored in the u-area.

NAME

kvm_nextproc, kvm_getproc, kvm_setproc – read system process structures

SYNOPSIS

```
#include <kvm.h>
#include <sys/param.h>
#include <sys/time.h>
#include <sys/proc.h>

struct proc *kvm_getproc(kvm_t *kd, int pid);

struct proc *kvm_nextproc(kvm_t *kd);

int kvm_setproc (kvm_t *kd);
```

MT-LEVEL

Unsafe

DESCRIPTION

kvm_nextproc() may be used to sequentially read all of the system process structures from the kernel identified by *kd* (see **kvm_open**(3K)). Each call to **kvm_nextproc()** returns a pointer to the static memory area that contains a copy of the next valid process table entry. There is no guarantee that the data will remain valid across calls to **kvm_nextproc()**, **kvm_setproc()**, or **kvm_getproc()**. Therefore, if the process structure must be saved, it should be copied to non-volatile storage.

For performance reasons, many implementations will cache a set of system process structures. Since the system state is liable to change between calls to **kvm_nextproc()**, and since the cache may contain obsolete information, there is no guarantee that *every* process structure returned refers to an active process, nor is it certain that *all* processes will be reported.

kvm_setproc() rewinds the process list, enabling **kvm_nextproc()** to rescan from the beginning of the system process table. **kvm_setproc()** will always flush the process structure cache, allowing an application to re-scan the process table of a running system.

kvm_getproc() locates the **proc** structure of the process specified by *pid* and returns a pointer to it. **kvm_getproc()** does not interact with the process table pointer manipulated by **kvm_nextproc()**, however, the restrictions regarding the validity of the data still apply.

RETURN VALUES

On success, **kvm_nextproc()** returns a pointer to a copy of the next valid process table entry. On failure, it returns NULL.

On success, **kvm_getproc()** returns a pointer to the **proc** structure of the process specified by *pid*. On failure, it returns NULL.

kvm_setproc() returns:

0 on success.

–1 on failure.

SEE ALSO

kvm_getu(3K), **kvm_open**(3K), **kvm_read**(3K)

NAME

kvm_nlist – get entries from kernel symbol table

SYNOPSIS

```
#include <kvm.h>
#include <nlist.h>

int kvm_nlist(kvm_t *kd, struct nlist *nl);
```

MT-LEVEL

Unsafe

DESCRIPTION

kvm_nlist( ) examines the symbol table from the kernel image identified by *kd* (see **kvm_open**(3K)) and selectively extracts a list of values and puts them in the array of **nlist** structures pointed to by *nl*. The name list pointed to by *nl* consists of an array of structures containing names, types and values. The *n_name* field of each such structure is taken to be a pointer to a character string representing a symbol name. The list is terminated by an entry with a NULL pointer (or a pointer to a null string) in the *n_name* field. For each entry in *nl*, if the named symbol is present in the kernel symbol table, its value and type are placed in the *n_value* and *n_type* fields. If a symbol cannot be located, the corresponding *n_type* field of *nl* is set to zero.

RETURN VALUES

kvm_nlist( ) returns the value of **nlist**(3B) or **nlist**(3E), depending on the library used.

SEE ALSO

nlist(3B), **nlist**(3E), **kvm_open**(3K), **kvm_read**(3K)

NAME

kvm_open, kvm_close – specify a kernel to examine

SYNOPSIS

```
#include <kvm.h>
#include <fcntl.h>

kvm_t *kvm_open(char *namelist, char *corefile, char *swapfile, int flag, char *errstr);

int kvm_close(kvm_t *kd);
```

MT-LEVEL

Unsafe

DESCRIPTION

kvm_open( ) initializes a set of file descriptors to be used in subsequent calls to kernel VM routines. It returns a pointer to a kernel identifier that must be used as the *kd* argument in subsequent kernel VM function calls.

The *namelist* argument specifies an unstripped executable file whose symbol table will be used to locate various offsets in *corefile*. If *namelist* is NULL, the symbol table of the currently running kernel is used to determine offsets in the core image. In this case, it is up to the implementation to select an appropriate way to resolve symbolic references (for instance, using **/dev/ksyms** as a default *namelist* file).

corefile specifies a file that contains an image of physical memory, for instance, a kernel crash dump file (see **savecore**(1M)) or the special device **/dev/mem**. If *corefile* is NULL, the currently running kernel is accessed (using **/dev/mem** and **/dev/kmem**).

swapfile specifies a file that represents the swap device. If both *corefile* and *swapfile* are NULL, the swap device of the ''currently running kernel'' is accessed. Otherwise, if *swapfile* is NULL, **kvm_open( )** may succeed but subsequent **kvm_getu**(3K) function calls may fail if the desired information is swapped out.

flag is used to specify read or write access for *corefile* and may have one of the following values:

O_RDONLY	open for reading
O_RDWR	open for reading and writing

errstr is used to control error reporting. If it is a NULL pointer, no error messages will be printed. If it is non-NULL, it is assumed to be the address of a string that will be used to prefix error messages generated by **kvm_open**. Errors are printed to **stderr**. A useful value to supply for *errstr* would be **argv**[0]. This has the effect of printing the process name in front of any error messages.

kvm_close( ) closes all file descriptors that were associated with *kd*. These files are also closed on **exit**(2) and **execve**( ) (see **exec**(2)). **kvm_close( )** also resets the **proc** pointer associated with **kvm_nextproc**(3K) and flushes any cached kernel data.

RETURN VALUES

kvm_open( ) returns a non-NULL value suitable for use with subsequent kernel VM function calls. On failure, it returns NULL and no files are opened.

kvm_close( ) returns:

0 on success.

−1 on failure.

FILES

/dev/kmem
/dev/ksyms
/dev/mem

SEE ALSO

savecore(1M), **exec**(2), **exit**(2), **kvm_getu**(3K), **kvm_nextproc**(3K), **kvm_nlist**(3K), **kvm_read**(3K)

NOTES

Programs using **libkvm** are likely to be platform and release dependent.

Kernel core dumps should be examined on the same platform they were created on.

NAME

kvm_read, kvm_write, kvm_uread, kvm_uwrite, kvm_kread, kvm_kwrite – copy data to or from a kernel image or running system

SYNOPSIS

```
#include <kvm.h>
int kvm_read(kvm_t *kd, unsigned long addr, char *buf, unsigned nbytes);
int kvm_write(kvm_t *kd, unsigned long addr, char *buf, unsigned nbytes);
int kvm_uread(kvm_t *kd, unsigned long addr, char *buf, unsigned nbytes);
int kvm_uwrite(kvm_t *kd, unsigned long addr, char *buf, unsigned nbytes);
int kvm_kread(kvm_t *kd, unsigned long addr, char *buf, unsigned nbytes);
int kvm_kwrite(kvm_t *kd, unsigned long addr, char *buf, unsigned nbytes);
```

MT-LEVEL

Unsafe

DESCRIPTION

kvm_read( ) transfers data from the kernel image specified by *kd* (see **kvm_open**(3K)) to the address space of the process. *nbytes* bytes of data are copied from the kernel virtual address given by *addr* to the buffer pointed to by *buf*.

kvm_write( ) is like **kvm_read( )**, except that the direction of data transfer is reversed. In order to use this function, the **kvm_open**(3K) call that returned *kd* must have specified write access. If a user virtual address is given, it is resolved in the address space of the process specified in the most recent **kvm_getu**(3K) call.

kvm_uread( ) transfers data from the address space of the processes specified in the most recent **kvm_getu**(3K) call. *nbytes* bytes of data are copied from the user virtual address given by addr to the buffer pointed to by *buf*.

kvm_uwrite( ) is like **kvm_uread( )**, except that the direction of the transfer is reversed. In order to use this function, the **kvm_open**(3K) call that returned *kd* must have specified write access. The address is resolved in the address space of the process specified in the most recent **kvm_getu**(3K) call.

kvm_kread( ) transfers data from the kernel address space to the address space of the process. *nbytes* bytes of data are copied from the kernel virtual address given by *addr* to the buffer pointed to by *buf*.

kvm_kwrite( ) is like **kvm_kread( )**, except that the direction of the transfer is reversed. In order to use this function, the **kvm_open**(3K) call that returned *kd* must have specified write access.

Note: The use of **kvm_uread( )**, **kvm_uwrite( )**, **kvm_kread( )** and **kvm_kwrite( )** is encouraged over the use of **kvm_read( )** and **kvm_write( )** since these are more clearly defined interfaces.

RETURN VALUES

All the above functions return the following values:

<number of bytes actually transferred>	Success.
–1	Failure.

SEE ALSO **kvm_getu**(3K), **kvm_nlist**(3K), **kvm_open**(3K)

NAME

lckpwdf, ulckpwdf – manipulate shadow password database lock file

SYNOPSIS

```
#include <shadow.h>
int lckpwdf(void);
int ulckpwdf(void);
```

MT-LEVEL

MT-Safe

DESCRIPTION

lckpwdf( ) and **ulckpwdf( )** are routines that are used to gain modification access to the password databases, through the lock file. A process first uses **lckpwdf( )** to lock the lock file, thereby gaining exclusive rights to modify the **/etc/passwd** or **/etc/shadow** password database. Upon completing modifications, a process should release the lock on the lock file using **ulckpwdf( )**. This mechanism prevents simultaneous modification of the password databases. **/etc/.pwd.lock** is the lock file. It is used to coordinate modification access to the password databases **/etc/passwd** and **/etc/shadow**.

lckpwdf( ) attempts to lock the file **/etc/.pwd.lock** within 15 seconds. If unsuccessful, for example, **/etc/.pwd.lock** is already locked, it returns –1. If successful, a return code other than –1 is returned.

ulckpwdf( ) attempts to unlock the file **/etc/.pwd.lock**. If unsuccessful, for example, **/etc/.pwd.lock** is already unlocked, it returns –1. If successful, it returns 0.

RETURN VALUES

lckpwdf( ) and **ulckpwdf( )** return –1 on failure, and 0 otherwise.

FILES

/etc/shadow
/etc/passwd
/etc/.pwd.lock

SEE ALSO

getpwnam(3C), **getspnam**(3C)

NOTES

These routines are for internal use only; compatibility is not guaranteed.

NAME

lfmt – display error message in standard format and pass to logging and monitoring services

SYNOPSIS

#include <pfmt.h>
int lfmt(FILE ∗*stream*, **long** *flags*, **char** ∗*format*, **... /**∗ *arg* ∗**/);**

MT-LEVEL

MT-safe

DESCRIPTION

lfmt() retrieves a format string from a locale-specific message database (unless **MM_NOGET** is specified) and uses it for **printf()** style formatting of *args*. The output is displayed on *stream*. If *stream* is **NULL**, no output is displayed.

lfmt() encapsulates the output in the standard error message format (unless **MM_NOSTD** is specified, in which case the output is simply **printf()** like).

lfmt() forwards its output to the logging and monitoring facility, even if *stream* is **NULL**. Optionnally, **lfmt()** will display the output on the console, with a date and time stamp.

If the **printf()** format string is to be retrieved from a message database, the *format* argument must have the following structure:

<*catalog*>**:**<*msgnum*>**:**<*defmsg*>.

If **MM_NOGET** is specified, only the <*defmsg*> part must be specified.

<*catalog*> is used to indicate the message database that contains the localized version of the format string. <*catalog*> must be limited to 14 characters. These characters must be selected from a set of all characters values, excluding **\0** (null) and the ASCII codes for **/** (slash) and **:** (colon).

<*msgnum*> is a positive number that indicates the index of the string into the message database.

If the catalog does not exist in the locale (specified by the last call to **setlocale()** using the **LC_ALL** or **LC_MESSAGES** categories), or if the message number is out of bound, **lfmt()** will attempt to retrieve the message from the C locale. If this second retrieval fails, **lfmt()** uses the <*defmsg*> part of the *format* argument.

If <*catalog*> is omitted, **lfmt()** will attempt to retrieve the string from the default catalog specified by the last call to **setcat()**. In this case, the *format* argument has the following structure:

:<*msgnum*>**:**<*defmsg*>.

lfmt() will output **Message not found!!\n** as format string if <*catalog*> is not a valid catalog name, if no catalog is specified (either explicitely or via **setcat()**), if <*msgnum*> is not a valid number, or if no message could be retrieved from the message databases, and <*defmsg*> was omitted.

The *flags* determine the type of output (i.e. whether the *format* should be interpreted as is or encapsulated in the standard message format), and the access to message catalogs to retrieve a localized version of *format*.

The *flags* are composed of several groups, and can take the following values (one from each group): *Output format control*

MM_NOSTD Do not use the standard message format, interpret *format* as a **printf()** *format*. Only *catalog access control flags*, *console display control* and *logging information* should be specified if **MM_NOSTD** is used; all other flags will be ignored

MM_STD Output using the standard message format (default, value 0).

Catalog access control

MM_NOGET Do not retrieve a localized version of *format*. In this case, only the <*defmsg*> part of the *format* is specified.

MM_GET Retrieve a localized version of *format*, from the <*catalog*>, using <*msgid*> as the index and <*defmsg*> as the default message (default, value 0).

Severity (standard message format only)

MM_HALT generates a localized version of 4HALT, but does not halt the machine.

MM_ERROR generates a localized version of **ERROR** (default, value 0).

MM_WARNING generates a localized version of **WARNING**.

MM_INFO generates a localized version of **INFO**.

Additional severities can be defined. Add-on severities can be defined with number-string pairs with numeric values from the range [5-255], using **addsev()**. The numeric value ORed with other *flags* will generate the specified severity.

If the severity is not defined, **lfmt()** used the string **SEV=***N* where *N* is replaced by the integer severity value passed in *flags*.

Multiple severities passed inf *flags* will not be detected as an error. Any combination of severities will be summed and the numeric value will cause the display of either a severity string (if defined) or the string **SEV=***N* (if undefined).

Action

MM_ACTION specifies an action message. Any severity value is superseded and replaced by a localized version of **TO FIX**.

Console display control

MM_CONSOLE display the message to the console in addition to the specified *stream*.

MM_NOCONSOLE do not display the message to the console in addition to the specified *stream* (default, value 0).

Logging information

Major classification

Identifies the source of the condition. Identifiers are: **MM_HARD** (hardware), **MM_SOFT** (software), and **MM_FIRM** (firmware).

Message source subclassification

Identifies the type of software in which the problem is spotted. Identifiers are: **MM_APPL** (application), **MM_UTIL** (utility), and **MM_OPSYS** (operating system).

STANDARD ERROR MESSAGE FORMAT

lfmt() displays error messages in the following format:

label: *severity*: *text*

If no *label* was defined by a call to **setlabel()**, the message is displayed in the format:

severity: *text*

If **lfmt()** is called twice to display an error message and a helpful *action* or recovery message, the output can look like:

label: *severity*: *text*
label: **TO FIX:** *text*

RETURN VALUE

Upon success, **lfmt()** returns the number of bytes transmitted. Upon failure, it returns a negative value:

-1 write error to *stream*.

-2 cannot log and / or display at console.

EXAMPLES

Example 1:

```
setlabel("UX:test");
lfmt(stderr, MM_ERROR|MM_CONSOLE|MM_SOFT|MM_UTIL,
        "test:2:Cannot open file: %s\n", strerror(errno));
```

displays the message to *stderr* and to the console and makes it available for logging:

```
UX:test: ERROR: Cannot open file: No such file or directory
```

Example 2:

```
setlabel("UX:test");
lfmt(stderr, MM_INFO|MM_SOFT|MM_UTIL,
        "test:23:test facility is enabled\n");
```

displays the message to *stderr* and makes it available for logging:

```
UX:test: INFO: test facility enabled
```

NOTES

Since **lfmt()** uses **gettxt**(3C), it is recommended that **lfmt()** not be used.

SEE ALSO

addsev(3C), **gettxt**(3C), **pfmt**(3C), **printf**(3S), **setcat**(3C), **setlabel**(3C), **setlocale**(3C), **environ**(5)

NAME

lgamma, lgamma_r, gamma, gamma_r – log gamma function

SYNOPSIS

cc [*flag* ...] *file* ... **–lm** [*library* ...]

#include <math.h>

extern int signgam;

double lgamma(double *x***);**

double lgamma_r(double *x***, int** ∗*signgamp***);**

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

Both **lgamma( )** and **lgamma_r( )** return

ln | Γ(x) |

where

$$\Gamma(x) = \int_0^\infty \mathbf{t^{x-1}} e^{-t}\mathbf{dt}$$

for x > 0 and

$$\Gamma(x) = \pi / (\Gamma(1-x)\sin(\pi x))$$

for x < 1.

lgamma( ) uses the external integer **signgam** to return the sign of Γ(x) while **lgamma_r( )** uses the user-allocated space addressed by *signgamp*.

IDIOSYNCRASIES

In the case of **lgamma( )**, do *not* use the expression **signgam**∗**exp(lgamma(x))** to compute '**g := Γ(x)**'. Instead compute **lgamma( )** first:

lg = lgamma(x); g = signgam∗**exp(lg);**

only after **lgamma( )** has returned can **signgam** be correct. Note that Γ(x) must overflow when *x* is large enough, underflow when –*x* is large enough, and generate a division by 0 exception at the singularities *x* a nonpositive integer.

RETURN VALUES

For exceptional cases, **matherr**(3M) tabulates the values to be returned as dictated by various Standards.

SEE ALSO

matherr(3M)

NOTES

Although **lgamma_r( )** is not mentioned by POSIX.4a Draft 6, it was added to complete the functionality provided by similar thread-safe functions. This interface is subject to change to be compatible with the "spirit" of POSIX.4a when it is approved as a standard.

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

lgamma( ) is unsafe in multithreaded applications. **lgamma_r( )** should be used instead.

NAME

libthread_db, td_log, td_ta_new, td_ta_delete, td_init, td_ta_get_ph, td_ta_get_nthreads, td_ta_tsd_iter, td_ta_thr_iter, td_thr_validate, td_thr_tsd, td_thr_get_info, td_thr_getfpregs, td_thr_getxregsize, td_thr_getxregs, td_thr_sigsetmask, td_thr_setprio, td_thr_setsigpending, td_thr_setfpregs, td_thr_setxregs, td_ta_map_id2thr, td_ta_map_lwp2thr, td_thr_getgregs, td_thr_setgregs – interface to libthread threads information

SYNOPSIS

cc [*flag* ...] *file* ... **/lib/libthread_db.so.1** [*library* ...]

#include <proc_service.h>
#include <thread_db.h>

void td_log(const int *on_off***);**

td_err_e td_ta_new(const struct ps_prochandle **ph_p***,* **td_thragent_t** ***ta_pp***);**

td_err_e td_ta_delete(const td_thragent_t **ta_p***);*

td_err_e td_init();

td_err_e td_ta_get_ph(const td_thragent_t **ta_p***,* **struct ps_prochandle** ***ph_pp***);**

td_err_e td_ta_get_nthreads(const td_thragent_t **ta_p***,* **int** **nthread_p***);*

td_err_e td_ta_tsd_iter(const td_thragent_t **ta_p***,* **td_key_iter_f** **cb***,* **void** **cbdata_p***);*

td_err_e td_ta_thr_iter(const td_thragent_t **ta_p***,* **td_thr_iter_f** **cb***,* **void** **cbdata_p***,*
td_thr_state_e *state*, **int** *ti_pri,* **sigset_t** **ti_sigmask_p***,* **unsigned** *ti_user_flags*);

td_err_e td_thr_validate(const td_thrhandle_t **th_p***);*

td_err_e td_thr_tsd(const td_thrhandle_t **th_p***,* **const thread_key_t** *key*,
void ***data_pp***);**

td_err_e td_thr_get_info(const td_thrhandle_t **th_p***,* **td_thrinfo_t** **ti_p***);*

td_err_e td_thr_getfpregs(const td_thrhandle_t **th_p***,* **prfpregset_t** **fpregset***);*

td_err_e td_thr_getxregsize(const td_thrhandle_t **th_p***,* **int** **xregsize***);*

td_err_e td_thr_getxregs(const td_thrhandle_t **th_p***,* **const caddr_t** **xregset***);*

td_err_e td_thr_sigsetmask(const td_thrhandle_t **th_p***,* **const sigset_t** *ti_sigmask*);

td_err_e td_thr_setprio(const td_thrhandle_t **th_p***,* **const int** *ti_pri*);

td_err_e td_thr_setsigpending(const td_thrhandle_t **th_p***,*
const uchar_t *ti_pending_flag*, **const sigset_t** *ti_pending*);

td_err_e td_thr_setfpregs(const td_thrhandle_t **th_p***,* **const prfpregset_t** **fpregset***);*

td_err_e td_thr_setxregs(const td_thrhandle_t **th_p***,* **const caddr_t** **xregset***);*

td_err_e td_ta_map_id2thr(const td_thragent_t **ta_p***,* **thread_t** *tid*,
td_thrhandle_t **th_p***);*

td_err_e td_ta_map_lwp2thr(const td_thragent_t **ta_p***,* **lwpid_t** *lwpid*,
td_thrhandle_t **th_p***);*

td_err_e td_thr_getgregs(const td_thrhandle_t **th_p*, **prgregset_t** *regset*);

td_err_e td_thr_setgregs(const td_thrhandle_t **th_p*, **const prgregset_t** *regset*);

DESCRIPTION

libthread_db is a library of functions for accessing information about threads in a process under inspection (PUI). The typical use for this library is by a debugger that requires information about the threads in a process. To some extent it can also be used for self inspection of threads within the same process. libthread_db utilizes functions provided by the user for interrogating a process (i.e., reading, writing, stopping, continuing, etc. a process). The prototypes for these user provided functions, referred to as the proc_service functions are described in <proc_service.h>. libthread_db should be accessed by an explicit version of the library (e.g., libthread_db.so.1) to ensure the correct version is used.

Support Level

All interfaces in libthread_db are UNCOMMITTED interfaces and are subject to change in future releases.

FUNCTIONS

td_log( ) turns logging on and off. A function in the proc_service interface is provided by the user for logging execution points in libthread_db. If logging is turned on, this proc_service function is called. The nominal name of the function is **ps_plog**. *on_off* = 0 turns logging off, non-zero turns logging on. proc_service functions called: none.

td_ta_new( ) allocates a thread agent for the given process handle and return a pointer to it. proc_service functions called: ps_pglobal_lookup.

td_ta_delete( ) deallocates the thread agent. proc_service functions called: none.

td_init( ) performs initialization for libthread_db interface. proc_service functions called: none

td_ta_get_ph( ) gets the process handle out of a thread agent and returns it. proc_service functions called: none

td_ta_get_nthreads( ) gets the total number of threads in a process. This number includes both user and system threads. proc_service functions called: **ps_pglobal_lookup**, **ps_pdread**, **ps_pstop**, **ps_pcontinue**.

td_ta_tsd_iter( ) iterates over the set of global TSD keys. The call back function is called with three arguments, a key, a pointer to a destructor function, and an extra pointer which can be NULL depending on the call back. The call back function is called once for each key. If return value of *cb* is non-zero, terminate iterations. proc_service functions called: *ps_pglobal_lookup*, *ps_pdread*, *ps_pstop*, *ps_pcontinue*.

td_ta_thr_iter( ) iterates over all threads. For each thread call the function pointed to by *cb* with a pointer to a thread handle, and a pointer to data which can be NULL. Only call *cb on threads* which match the properties of *state*, *ti_pri*, *ti_sigmask_p*, and *ti_user_flags*. If *cb* returns a non-zero value, iterations terminate. The call back function is defined by user. **td_thr_iter_f** takes a thread handle and *cbdata_p* as a parameter. *state* is the state of threads of interest. A value of **TD_THR_ANY_STATE** from enum **td_thr_state_e** does not restrict iterations by state. *ti_pri* is the lower bound of priorities of threads of interest. A

value of **TD_THR_LOWEST_PRIORITY** defined in thread_db.h does not restrict iterations by priority. A thread with priority less than *ti_pri* will NOT be passed to the callback function. *ti_sigmask_p* is the signal mask of threads of interest. A value of **TD_SIGNO_MASK** defined in thread_db.h does not restrict iterations by signal mask. *ti_user_flags* is the user flags of threads of interest. A value of **TD_THR_ANY_USER_FLAGS** defined in thread_db.h does not restrict iterations by user flags. proc_service functions called: *ps_pdread*, *ps_pstop*, *ps_pcontinue*.

td_thr_validate() validates the thread handle. **td_thr_validate()** checks that a thread exists in the thread agent/process that corresponds to thread with handle **th_p*. Return value: **TD_OK** implies thread handle is valid. **TD_NOTHR** implies thread handle not valid. proc_service functions called: *ps_pdread*, *ps_pstop*, *ps_pcontinue*.

td_thr_tsd() gets a thread's private binding to a given thread specific data (TSD) key (see **thr_getspecific**(3T)). If the thread doesn't have a binding for a particular key, then NULL is returned. proc_service functions called: *ps_pglobal_lookup*, *ps_pdread*, *ps_pstop*, *ps_pcontinue*.

td_thr_get_info() updates the thread information struct. All fields in a thread information structure (**td_thrinfo_t**) will be updated to be consistent with properties of its respective thread. proc_service functions called: **ps_pdread**, **ps_pstop**, **ps_pcontinue**.

td_thr_getfpregs() gets the floating point registers for the given thread. proc_service functions called: **ps_pdread**, **ps_pstop**, **ps_pcontinue**, **ps_lgetfpregs**.

td_thr_getxregsize() gets the size of the extra register set for the given thread. proc_service functions called: **ps_pdread**, **ps_pstop**, **ps_pcontinue**, **ps_lgetxregsize**.

td_thr_getxregs() gets the extra registers for the given thread. proc_service functions called: **ps_pdread**, **ps_pstop**, **ps_pcontinue**, **ps_lgetxregs**.

td_thr_sigsetmask() changes a thread's signal mask to the value specified by *ti_sigmask*. proc_service functions called: **ps_pstop**, **ps_pcontinue**, **ps_pdwrite**.

td_thr_setprio() changes a thread's priority to the value specified by **ti_pri.** proc_service functions called: *ps_pdwrite*.

td_thr_setsigpending() changes a thread's pending signal state to that specified by *ti_pending_flag* and *ti_pending*. A null value for *ti_pending_flag* indicates that there are no pending signals for the thread. proc_service functions called: **ps_pstop**, **ps_pcontinue**, **ps_pdwrite**.

td_thr_setfpregs() sets the floating pointing registers for a given thread. proc_service functions called: **ps_pdread**, **ps_pstop**, **ps_lsetfpregs**, **ps_pcontinue**.

td_thr_setxregs() sets the extra registers for the given thread. proc_service functions called: **ps_pcontinue**, **ps_pstop**, **ps_pdread**, **ps_pdwrite**, **ps_lsetxregs**.

td_ta_map_id2thr() returns the thread handle corresponding to the given thread identifier tid. proc_service functions called: **ps_pdread**, **ps_pstop**, **ps_pcontinue**.

td_ta_map_lwp2thr() returns a thread handle for the given thread agent and lwp id. proc_service functions called: **ps_pdread**, **ps_lgetgregs**.

td_thr_getgregs() gets the general registers for a given thread. For a thread that is currently executing on an LWP, (**td_thr_state_e**) **TD_THR_ACTIVE**, all registers in *regset* will be read for the thread. For a thread not executing on an LWP, only the following registers will be read.

SPARC %i0-%i7,
%l0-%l7,
%g7, %pc, %sp (%o6).

x86 %ebp, %edi, %esi, %ebx, %sp, %pc.

%pc and %sp will be the program counter and stack pointer at the point where the thread will resume execution when it becomes active,
(**td_thr_state_e**) **TD_THR_ACTIVE**. proc_service functions called: **ps_pstop**, **ps_pcontinue**, **ps_pdread**, **ps_lgetregs**.

td_thr_setgregs() sets the general registers for a given thread. For a thread that is currently executing on an LWP, (**td_thr_state_e**) **TD_THR_ACTIVE**, all registers in regset will be written for the thread. For a thread not executing on an LWP, only the following registers will be written.

SPARC %i0-%i7,
%l0-%l7,
%pc, %sp (%o6).

x86 %ebp, %edi, %esi, %ebx, %sp, %pc.

%pc and %sp will be the program counter and stack pointer at the point where the thread will resume execution when it becomes active, (**td_thr_state_e**) **TD_THR_ACTIVE**. proc_service functions called: **ps_pstop**, **ps_pcontinue**, **ps_pdread**, **ps_lsetregs**.

RETURN VALUES

TD_OK	The call completed successfully.
TD_ERR	The call failed and no other error code applies.
TD_NOTHR	The call failed because the thread handle passed as a parameter does not correspond to any thread known to libthread_db. The thread may no longer exist or the user may be using a corrupted thread handle.
TD_NOSV	If the synchronization variable passed as a parameter does not correspond to any synchronization variable known to libthread_db the call will return this value. The synchronization variable may no longer exist or the user may be using a corrupted synchronization handle.
TD_NOLWP	If there is no LWP that is part of the process that corresponds to the given lwpid the call will fail and return this value. This can occur if the lwpid is corrupted or if the lwp has been returned to the OS.
TD_BADPH	The call failed because the process handle is invalid. This generally indicates that a NULL values was passed as a process handle.

TD_BADTH The call failed because the thread handle is invalid.

TD_BADSH The call failed because the synchronization handle is invalid.

TD_BADKEY The call failed because the key for the thread specific data area is invalid.

TD_NOMSG If there is not event message information, this value is returned.

TD_NOFPREGS If the floating point registers are not available for this thread the call returns this value.

TD_NOLIBTHREAD
If the applications is not threaded (i.e., not linked with libthread) this value is returned.

TD_NOEVENT This value is returned if an address was requested for an event which is not supported.

TD_NOCAPAB This value indicates that the requested service is not supported. This generally indicates a limitation with the current library that will be rectified in a future release.

TD_DBERR The call failed because an error occurred during a request for service from the user's proc_service interface.

TD_NOAPPLIC This value is returned if the requested function does not apply to the given variable.

TD_NOTSD The call failed because there is no thread specific data for this thread. A key may exist without an entry for every thread.

TD_MALLOC The call failed because a malloc failed.

TD_PARTIALREG
An operation on a register set was performed on only part of the register set. For example, at some points in the execution of a process not all the general purpose registers may be available.

TD_NOXREGS If the extra registers are not available for this thread the call returns this value.

EXAMPLES

```
/*
*  Specify libthread_db.so.1 on link line to access correct version.
*  cc thisfile.c /lib/libthread_db.so.1
*/
#include <stdio.h>
#include <sys/types.h>
#include <thread_db.h>
#include <dlfcn.h>

static int thread_cb( const td_thrhandle_t *th_p, void *s );
struct ps_prochandle {
  pid_t pid;
};
```

```
struct ps_prochandle ph = {1};
td_thragent_t *ta_p;

/*
*  libthread_db example.
*    Initialize libthread_db
*    Create a thread agent
*    Call thread iterator
*/
int main()
{
  td_err_e td_return;
  /*
  *  td_init()
  */
  td_return = td_init();
  if ( td_return != TD_OK ) {
    printf("Initialization error on td_init() call\n");
    return 0;
  }
  /*
  *  td_ta_new()
  */
  td_return = td_ta_new(&ph, &ta_p);
  if ( td_return == TD_OK ) {
    /*
    * td_ta_thr_iter()
    */
    (void) td_ta_thr_iter(ta_p, thread_cb, "Import calls test",
      TD_THR_ANY_STATE, TD_THR_LOWEST_PRIORITY,
      TD_SIGNO_MASK, TD_THR_ANY_USER_FLAGS);
  }

  return 0;
}

static int thread_cb( const td_thrhandle_t *th_p, void *s )
/*
*  Description:
*    Call back function for iterator
*
*  Input:
*    *th_p - thread handle
*    *s - data
*
```

```
* Output:
*    none
*
*/
{
  int return_val = 0;
  td_err_e td_return;
  td_thrinfo_t to;
  td_return = td_thr_get_info(th_p, &(to));
  if ( td_return == TD_ERR ) {
    printf("Thread update failed\n");
    return_val = 1;
  }
  return return_val;
}

/*
*  Use dlopen and dlsym to access libthread_db functions.
*/
#include <stdio.h>
#include <sys/types.h>
#include <thread_db.h>
#include <dlfcn.h>

static int thread_cb( const td_thrhandle_t *th_p, void *s );
struct ps_prochandle {
  pid_t pid;
};
struct ps_prochandle ph = {1};
td_thragent_t *ta_p;

td_err_e (*init)(void);
td_err_e (*ta_new)(const struct ps_prochandle *ph_p, td_thragent_t **ta_pp);
td_err_e (*thr_iter)(const td_thragent_t *ta_p, td_thr_iter_f *cb,
            void *cbdata_p, td_thr_state_e state, int ti_pri,
            sigset_t *ti_sigmask_p, unsigned ti_user_flags);
td_err_e (*get_info)(const td_thrhandle_t *th_p, td_thrinfo_t *ti_p);

/*
*  libthread_db example.
*    Initialize libthread_db
*    Create a thread agent
*    Call thread iterator
*/
int main()
```

```
{
  td_err_e td_return;
  void *handle;
  /*
  *   Access a specific version of libthread_db.
  */
  handle = dlopen("libthread_db.so.1", RTLD_LAZY);
  if ( !handle ) {
    printf("dlopen error: %s\n", dlerror());
    return 1;
  }
  else {
    init = (td_err_e (*)(void))dlsym(handle, "td_init");
    ta_new = (td_err_e (*)(const struct ps_prochandle *, td_thragent_t **))
          dlsym(handle, "td_ta_new");
    thr_iter = (td_err_e (*)(const td_thragent_t *, td_thr_iter_f *,
          void *, td_thr_state_e , int , sigset_t *, unsigned ))
          dlsym(handle, "td_ta_thr_iter");
    get_info = (td_err_e (*)(const td_thrhandle_t *, td_thrinfo_t *))
          dlsym(handle, "td_thr_get_info");
  }

  /*
  *   td_init()
  */
  td_return = (*init)();
  if ( td_return != TD_OK ) {
    printf("Initialization error on td_init() call\n");
    return 0;
  }
  /*
  *   td_ta_new()
  */
  td_return = (*ta_new)(&ph, &ta_p);
  if ( td_return == TD_OK ) {
    /*
    * td_ta_thr_iter()
    */
    (void) (*thr_iter)(ta_p, thread_cb, "Import calls test",
      TD_THR_ANY_STATE, TD_THR_LOWEST_PRIORITY,
      TD_SIGNO_MASK, TD_THR_ANY_USER_FLAGS);
  }

  return 0;
}
```

```
static int thread_cb( const td_thrhandle_t *th_p, void *s )
/*
*  Description:
*     Call back function for iterator
*
*  Input:
*     *th_p - thread handle
*     *s - data
*
*  Output:
*     none
*
*/
{
   int return_val = 0;
   td_err_e td_return;
   td_thrinfo_t to;
   td_return = (*get_info)(th_p, &(to));
   if ( td_return == TD_ERR ) {
      printf("Thread update failed\n");
      return_val = 1;
   }
   return return_val;
}
```

FILES **/usr/lib/libthread_db.so.1**

SEE ALSO **proc_service**(3T), **thr_getspecific**(3T)

NAME

lio_listio – list directed I / O

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <aio.h>

int lio_listio(int mode, struct aiocb * const list[], int nent, struct sigevent *sig);

struct aiocb {
  int              aio_fildes;     /* file descriptor */
  volatile void    *aio_buf;       /* buffer location */
  size_t           aio_nbytes;     /* length of transfer */
  off_t            aio_offset;     /* file offset */
  int              aio_reqprio;    /* request priority offset */
  struct sigevent  aio_sigevent;   /* signal number and offset */
  int              aio_lio_opcode; /* listio operation */
};

struct sigevent {
  int              sigev_notify;   /* notification mode */
  int              sigev_signo;    /* signal number */
  union sigval     sigev_value;    /* signal value */
};

union sigval {
  int              sival_int;      /* integer value */
  void             *sival_ptr;     /* pointer value */
};
```

MT-LEVEL

MT-Safe

DESCRIPTION

lio_listio( ) allows the calling process, LWP, or thread, to initiate a list of I / O requests within a single function call.

If *mode* is set to **LIO_WAIT, lio_listio( )** behaves synchronously, waiting until all I / O is completed, and the *sig* argument is ignored. If *mode* is set to **LIO_NOWAIT, lio_listio( )** behaves asynchronously; returning immediately, and signal delivery will occur, according to the *sig* argument, when all the I / O operations from this function complete. If *sig* is **NULL,** or the **sigev_signo** member of the **sigevent** structure referenced by *sig* is zero, then no signal delivery will occur. Otherwise, the signal number indicated by **sigev_signo** will be delivered when all the requests in *list* have completed.

list is an array of pointers to **aiocb** structures. This array consists of *nent* elements. The array may contain **NULL** pointers, which will be ignored.

The **aio_lio_opcode** field of each **aiocb** structure in *list* specifies the operation to be performed (see **/usr/include/aio.h**).

LIO_READ
requests **aio_read**(3R).

LIO_WRITE
requests **aio_write**(3R).

LIO_NOP
causes the *list* entry to be ignored.

nent specifies the length of the array (number of members of the list).

When *mode* has the value **LIO_WAIT**, a pointer to a signal control structure, *sig*, is used to define both the signal to be generated and how the calling process will be notified upon I/O completion. If *sig*->**sigev_notify** is **SIGEV_NONE**, then no signal will be posted upon I/O completion, but the error status and the return status for the operation will be set appropriately. If *sig*->**sigev_notify** is **SIGEV_SIGNAL**, then the signal specified in *sig*->**sigev_signo** will be sent to the process. If the **SA_SIGINFO** flag is set for that signal number, then the signal will be queued to the process and the value specified in *sig*->**sigev_value** will be the **si_value** component of the generated signal (see **siginfo**(5)).

The behavior of this function is altered according to the definitions of synchronized I/O data integrity completion and synchronized I/O file integrity completion if synchronized I/O is enabled on the file associated with **aio_fildes.** (see **fcntl**(5) definitions of **O_DSYNC** and **O_SYNC**.)

RETURN VALUES

If the *mode* argument has the value **LIO_NOWAIT**, and the I/O operations are successfully queued, **lio_listio( )** returns **0;** otherwise, it returns **-1,** and sets **errno** to indicate the error condition.

If the *mode* argument has the value **LIO_WAIT**, and when all the indicated I/O has completed successfully, **lio_listio( )** returns **0;** otherwise, it returns **-1,** and sets **errno** to indicate the error condition.

In either case, the return value only indicates the success or failure of the **lio_listio( )** call itself, not the status of the individual I/O requests. In some cases, one or more of the I/O requests contained in the list may fail. Failure of an individual request does not prevent completion of any other individual request. To determine the outcome of each I/O request, the application must examine the error status associated with each *aiocb* control block. Each error status so returned is identical to that returned as a result of an **aio_read**(3R) or **aio_write**(3R) function.

ERRORS

EAGAIN The resources necessary to queue all the I/O requests were not available. The error status for each request is recorded in the **aio_error** member of the corresponding **aiocb** structure, and can be retrieved using **aio_error**(3R).

nent entries exceed the system-wide limit, **AIO_MAX**.

EINVAL The *mode* argument is an improper value.

The value of *nent* is greater than **AIO_LISTIO_MAX**.

EINTR A signal was delivered while waiting for all I / O requests to complete during an **LIO_WAIT** operation. However, the outstanding I / O requests are not canceled. Use **aio_fsync**(3R) to determine if any request was initiated; **aio_return**(3R) to determine if any request has completed; or **aio_error**(3R) to determine if any request was canceled.

EIO One or more of the individual I / O operations failed. Using **aio_error**(3R) with each **aiocb** structure will determine the individual request(s) that failed.

ENOSYS **lio_listio( )** is not supported by this implementation.

If either **lio_listio( )** succeeds in queuing all of its requests, or **errno** is set to **EAGAIN**, **EINTR**, or **EIO**, then some of the I / O specified from the list may have been initiated. In this event, each **aiocb** structure contains errors specific to the **read**(2) or **write**(2) function being performed; i.e.:

EAGAIN The requested I / O operation was not queued due to resource limitations.

ECANCELED The requested I / O was canceled before the I / O completed due to an explicit **aio_cancel**(3R) request.

EINPROGRESS Thee requested I / O is in progress.

SEE ALSO

close(2), **exec**(2), **exit**(2), **fork**(2), **lseek**(2), **read**(2), **write**(2), **aio_cancel**(3R), **aio_read**(3R), **aio_return**(3R), **fcntl**(5), **siginfo**(5)

NOTES

Applications compiled under Solaris 2.3 and 2.4 and using POSIX aio must be recompiled to work correctly when Solaris supports the Asynchronous Input and Output option.

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Asynchronous Input and Output option. It is our intention to provide support for these interfaces in future releases.

NAME

listen – listen for connections on a socket

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>

int listen(int s, int backlog);
```

MT-LEVEL

Safe

DESCRIPTION

To accept connections, a socket is first created with **socket**(3N), a backlog for incoming connections is specified with **listen( )** and then the connections are accepted with **accept**(3N). The **listen( )** call applies only to sockets of type **SOCK_STREAM** or **SOCK_SEQPACKET**.

The *backlog* parameter defines the maximum length the queue of pending connections may grow to.

If a connection request arrives with the queue full, the client will receive an error with an indication of **ECONNREFUSED** for AF_UNIX sockets. If the underlying protocol supports retransmission, the connection request may be ignored so that retries may succeed. For AF_INET sockets, the tcp will retry the connection. If the *backlog* is not cleared by the time the tcp times out, the connect will fail with **ETIMEDOUT**.

RETURN VALUES

A **0** return value indicates success; **–1** indicates an error.

ERRORS

The call fails if:

EBADF	The argument *s* is not a valid file descriptor.
ENOTSOCK	The argument *s* is not a socket.
EOPNOTSUPP	The socket is not of a type that supports the operation **listen( )**.

SEE ALSO

accept(3N), **connect**(3N), **socket**(3N)

NOTES

There is currently no *backlog* limit.

NAME

localeconv – get numeric formatting information

SYNOPSIS

```
#include <locale.h>
struct lconv *localeconv(void);
```

MT-LEVEL

Safe with exceptions

DESCRIPTION

localeconv( ) sets the components of an object with type **struct lconv** (defined in **<locale.h>**) with the values appropriate for the formatting of numeric quantities (monetary and otherwise) according to the rules of the current locale (see **setlocale**(3C)). The definition of **struct lconv** is given below (the values for the fields in the "C" locale are given in comments).

```
char *decimal_point;        /* "." */
char *thousands_sep;        /* "" (zero length string) */
char *grouping;             /* "" */
char *int_curr_symbol;      /* "" */
char *currency_symbol;      /* "" */
char *mon_decimal_point;    /* "" */
char *mon_thousands_sep;    /* "" */
char *mon_grouping;         /* "" */
char *positive_sign;        /* "" */
char *negative_sign;        /* "" */
char int_frac_digits;       /* CHAR_MAX */
char frac_digits;           /* CHAR_MAX */
char p_cs_precedes;         /* CHAR_MAX */
char p_sep_by_space;        /* CHAR_MAX */
char n_cs_precedes;         /* CHAR_MAX */
char n_sep_by_space;        /* CHAR_MAX */
char p_sign_posn;           /* CHAR_MAX */
char n_sign_posn;           /* CHAR_MAX */
```

The members of the structure with type **char *** are strings, any of which (except **decimal_point**) can point to a null string (""), to indicate that the value is not available in the current locale or is of zero length. The members with type **char** are nonnegative numbers, any of which can be **CHAR_MAX** (defined in the **<limits.h>** header) to indicate that the value is not available in the current locale. The members are the following:

char *decimal_point
: The decimal-point character used to format non-monetary quantities.

char *thousands_sep
: The character used to separate groups of digits to the left of the decimal-point character in formatted non-monetary quantities.

char *grouping
A string in which each element is taken as an integer that indicates the number of digits that comprise the current group in a formatted non-monetary quantity. The elements of **grouping** are interpreted according to the following:

CHAR_MAX	No further grouping is to be performed.
0	The previous element is to be repeatedly used for the remainder of the digits.
other	The value is the number of digits that comprise the current group. The next element is examined to determine the size of the next group of digits to the left of the current group.

char *int_curr_symbol
The international currency symbol applicable to the current locale, left-justified within a four-character space-padded field. The character sequences should match with those specified in *ISO 4217 Codes for the Representation of Currency and Funds*.

char *currency_symbol
The local currency symbol applicable to the current locale.

char *mon_decimal_point
The decimal point used to format monetary quantities.

char *mon_thousands_sep
The separator for groups of digits to the left of the decimal point in formatted monetary quantities.

char *mon_grouping
A string in which each element is taken as an integer that indicates the number of digits that comprise the current group in a formatted monetary quantity. The elements of **mon_grouping** are interpreted according to the rules described under **grouping**.

char *positive_sign
The string used to indicate a nonnegative-valued formatted monetary quantity.

char *negative_sign
The string used to indicate a negative-valued formatted monetary quantity.

char int_frac_digits
The number of fractional digits (those to the right of the decimal point) to be displayed in an internationally formatted monetary quantity.

char frac_digits
The number of fractional digits (those to the right of the decimal point) to be displayed in a formatted monetary quantity.

char p_cs_precedes
Set to 1 or 0 if the **currency_symbol** respectively precedes or succeeds the value for a nonnegative formatted monetary quantity.

char p_sep_by_space
Set to 1 or 0 if the **currency_symbol** respectively is or is not separated by a space from the value for a nonnegative formatted monetary quantity.

char n_cs_precedes
Set to 1 or 0 if the **currency_symbol** respectively precedes or succeeds the value for a negative formatted monetary quantity.

char n_sep_by_space
Set to 1 or 0 if the **currency_symbol** respectively is or is not separated by a space from the value for a negative formatted monetary quantity.

char p_sign_posn
Set to a value indicating the positioning of the **positive_sign** for a nonnegative formatted monetary quantity. The value of **p_sign_posn** is interpreted according to the following:

0 Parentheses surround the quantity and **currency_symbol**.

1 The sign string precedes the quantity and **currency_symbol**.

2 The sign string succeeds the quantity and **currency_symbol**.

3 The sign string immediately precedes the **currency_symbol**.

4 The sign string immediately succeeds the **currency_symbol**.

char n_sign_posn
Set to a value indicating the positioning of the **negative_sign** for a negative formatted monetary quantity. The value of **n_sign_posn** is interpreted according to the rules described under **p_sign_posn**.

RETURN VALUES

localeconv( ) returns a pointer to the filled-in object. The structure pointed to by the return value may be overwritten by a subsequent call to **localeconv( )**.

EXAMPLES

The following table illustrates the rules used by four countries to format monetary quantities.

Country	Positive format	Negative format	International format
Italy	L.1.234	-L.1.234	ITL.1.234
Netherlands	F 1.234,56	F -1.234,56	NLG 1.234,56
Norway	kr1.234,56	kr1.234,56-	NOK 1.234,56
Switzerland	SFrs.1,234.56	SFrs.1,234.56C	CHF 1,234.56

For these four countries, the respective values for the monetary members of the structure returned by **localeconv** are as follows:

	Italy	Netherlands	Norway	Switzerland
int_curr_symbol	"ITL."	"NLG "	"NOK "	"CHF "
currency_symbol	"L."	"F"	"kr"	"SFrs."
mon_decimal_point	""	","	","	"."
mon_thousands_sep	"."	"."	"."	","
mon_grouping	"\3"	"\3"	"\3"	"\3"
positive_sign	""	""	""	""

negative_sign	"-"	"-"	"-"	"C"
int_frac_digits	0	2	2	2
frac_digits	0	2	2	2
p_cs_precedes	1	1	1	1
p_sep_by_space	0	1	0	0
n_cs_precedes	1	1	1	1
n_sep_by_space	0	1	0	0
p_sign_posn	1	1	1	1
n_sign_posn	1	4	2	2

FILES

/usr/lib/locale/*locale***/LC_MONETARY/monetary**
LC_MONETARY database for *locale*

/usr/lib/locale/*locale***/LC_NUMERIC/numeric**
LC_NUMERIC database for *locale*

SEE ALSO

chrtbl(1M), **montbl**(1M), **setlocale**(3C)

NOTES

localeconv( ) can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

LC_MONETARY
Determines how monetary formats are handled. In the "C" locale, monetary handling follows the U.S. rules.

LC_NUMERIC
Determines how numeric formats are handled. In the "C" locale, numeric handling follows the U.S. rules.

NAME

lockf – record locking on files

SYNOPSIS

```
#include <unistd.h>
int lockf(int fildes, int function, long size);
```

MT-LEVEL

MT-Safe

DESCRIPTION

lockf( ) allows sections of a file to be locked; advisory or mandatory write locks depending on the mode bits of the file (see **chmod**(2)). Locking calls from other processes that attempt to lock the locked file section will either return an error value or be put to sleep until the resource becomes unlocked. All the locks for a process are removed when the process terminates. (See **fcntl**(2) for more information about record locking.)

fildes is an open file descriptor. The file descriptor must have **O_WRONLY** or **O_RDWR** permission in order to establish locks with this function call.

function is a control value that specifies the action to be taken. The permissible values for *function* are defined in **<unistd.h>** as follows:

```
#define    F_ULOCK    0    /* unlock previously locked section */
#define    F_LOCK     1    /* lock section for exclusive use */
#define    F_TLOCK    2    /* test & lock section for exclusive use */
#define    F_TEST     3    /* test section for other locks */
```

All other values of *function* are reserved for future extensions and will result in an error return if not implemented.

F_TEST is used to detect if a lock by another process is present on the specified section. **F_LOCK** and **F_TLOCK** both lock a section of a file if the section is available. **F_ULOCK** removes locks from a section of the file.

size is the number of contiguous bytes to be locked or unlocked. The resource to be locked or unlocked starts at the current offset in the file and extends forward for a positive size and backward for a negative size (the preceding bytes up to but not including the current offset). If *size* is zero, the section from the current offset through the largest file offset is locked (that is, from the current offset through the present or any future end-of-file). An area need not be allocated to the file in order to be locked as such locks may exist past the end-of-file.

The sections locked with **F_LOCK** or **F_TLOCK** may, in whole or in part, contain or be contained by a previously locked section for the same process. Locked sections will be unlocked starting at the the point of the offset through *size* bytes or to the end of file if *size* is (**off_t**) 0. When this situation occurs, or if this situation occurs in adjacent sections, the sections are combined into a single section. If the request requires that a new element be added to the table of active locks and this table is already full, an error is returned, and the new section is not locked.

F_LOCK and **F_TLOCK** requests differ only by the action taken if the resource is not available. **F_LOCK** will cause the calling process to sleep until the resource is available. **F_TLOCK** will cause the function to return a –1 and set **errno** to **EAGAIN** if the section is already locked by another process.

F_ULOCK requests may, in whole or in part, release one or more locked sections controlled by the process. When sections are not fully released, the remaining sections are still locked by the process. Releasing the center section of a locked section requires an additional element in the table of active locks. If this table is full, an **errno** is set to **EDEADLK** and the requested section is not released.

A potential for deadlock occurs if a process controlling a locked resource is put to sleep by requesting another process's locked resource. Thus calls to **lockf( )** or **fcntl**(2) scan for a deadlock prior to sleeping on a locked resource. An error return is made if sleeping on the locked resource would cause a deadlock.

Sleeping on a resource is interrupted with any signal. The **alarm**(2) function may be used to provide a timeout facility in applications that require this facility.

RETURN VALUES

Upon successful completion, a value of 0 is returned. Otherwise, a value of –1 is returned and **errno** is set to indicate the error.

ERRORS

lockf( ) will fail if one or more of the following are true:

EAGAIN	*cmd* is **F_TLOCK** or **F_TEST** and the section is already locked by another process.
EBADF	*fildes* is not a valid open descriptor.
ECOMM	*fildes* is on a remote machine and the link to that machine is no longer active.
EDEADLK	*cmd* is **F_LOCK** and a deadlock occurred.
EDEADLK	*cmd* is **F_LOCK**, **F_TLOCK**, or **F_ULOCK** and the number of entries in the lock table exceeded the number allocated on the system.

SEE ALSO

intro(2), **alarm**(2), **chmod**(2), **close**(2), **creat**(2), **fcntl**(2), **open**(2), **read**(2), **write**(2)

NOTES

Unexpected results may occur in processes that do buffering in the user address space. The process may later read / write data that is / was locked. The standard I/O package is the most common source of unexpected buffering.

In the past, the variable **errno** was set to **EACCES** rather than **EAGAIN.** When a section of a file is already locked by another process, portable application programs should expect and test for either value.

NAME

lsearch, lfind – linear search and update

SYNOPSIS

```
#include <search.h>

void *lsearch(const void *key, void *base, size_t *nelp, size_t width,
        int (*compar) (const void *, const void *));

void *lfind(const void *key, const void *base, size_t *nelp, size_t width,
        int (*compar)(const void *, const void *));
```

MT-LEVEL

Safe

DESCRIPTION

lsearch( ) is a linear search routine generalized from Knuth (6.1) Algorithm S. (See *The Art of Computer Programming, Volume 3, Section 6.1, by Donald E. Knuth.*) It returns a pointer into a table indicating where a datum may be found. If the datum does not occur, it is added at the end of the table. *key* points to the datum to be sought in the table. *base* points to the first element in the table. *nelp* points to an integer containing the current number of elements in the table. The integer is incremented if the datum is added to the table. *width* is the size of an element in bytes. *compar* is a pointer to the comparison function that the user must supply (**strcmp**, for example). It is called with two arguments that point to the elements being compared. The function must return zero if the elements are equal and non-zero otherwise.

lfind( ) is the same as **lsearch( )** except that if the datum is not found, it is not added to the table. Instead, a null pointer is returned.

Note that:

- the pointers to the key and the element at the base of the table may be pointers to any type.
- The comparison function need not compare every byte, so arbitrary data may be contained in the elements in addition to the values being compared.
- The value returned should be cast into type pointer-to-element.

EXAMPLES

This program will read in less than **TABSIZE** strings of length less than **ELSIZE** and store them in a table, eliminating duplicates, and then will print each entry.

```
#include <search.h>
#include <string.h>
#include <stdlib.h>
#include <stdio.h>

#define TABSIZE 50
#define ELSIZE 120

main()
{
        char line[ELSIZE];      /* buffer to hold input string */
        char tab[TABSIZE][ELSIZE];   /* table of strings */
```

```
        size_t nel = 0;             /* number of entries in tab */
        int i;

        while (fgets(line, ELSIZE, stdin) != NULL &&
                nel < TABSIZE)
                (void) lsearch(line, tab, &nel, ELSIZE, mycmp);
        for( i = 0; i < nel; i++ )
                (void)fputs(tab[i], stdout);
        return 0;
}
```

SEE ALSO

bsearch(3C), **hsearch**(3C), **string**(3C), **tsearch**(3C)

The Art of Computer Programming, Volume 3, Sorting and Searching by Donald E. Knuth, published by Addison-Wesley Publishing Company, 1973.

NOTES

If the searched-for datum is found, both **lsearch( )** and **lfind( )** return a pointer to it. Otherwise, **lfind( )** returns **NULL** and **lsearch( )** returns a pointer to the newly added element.

Undefined results can occur if there is not enough room in the table to add a new item.

NAME

madvise – provide advice to VM system

SYNOPSIS

```
#include <sys/types.h>
#include <sys/mman.h>

int madvise(caddr_t addr, size_t len, int advice);
```

MT-LEVEL

MT-Safe

DESCRIPTION

madvise( ) advises the kernel that a region of user mapped memory in the range [*addr*, *addr* + *len*) will be accessed following a type of pattern. The kernel uses this information to optimize the procedure for manipulating and maintaining the resources associated with the specified mapping range.

Values for *advice* are defined in **<sys/mman.h>** as:

```
#define MADV_NORMAL      0x0    /* No further special treatment */
#define MADV_RANDOM      0x1    /* Expect random page references */
#define MADV_SEQUENTIAL  0x2    /* Expect sequential page references */
#define MADV_WILLNEED    0x3    /* Will need these pages */
#define MADV_DONTNEED    0x4    /* Don't need these pages */
```

MADV_NORMAL

The default system characteristic where accessing memory within the address range causes the system to read data from the mapped file. The kernel reads all data from files into pages which are retained for a period of time as a "cache." System pages can be a scarce resource, so the kernel steals pages from other mappings when needed. This is a likely occurrence, but adversely affects system performance only if a large amount of memory is accessed.

MADV_RANDOM

Tells the kernel to read in a minimum amount of data from a mapped file on any single particular access. If **MADV_NORMAL** is in effect when an address of a mapped file is accessed, the system tries to read in as much data from the file as reasonable, in anticipation of other accesses within a certain locality.

MADV_SEQUENTIAL

Tells the system that addresses in this range are likely to be accessed only once, so the system will free the resources mapping the address range as quickly as possible. This is used in the **cat**(1) and **cp**(1) utilities.

MADV_WILLNEED

Tells the system that a certain address range is definitely needed so the kernel will start reading the specified range into memory. This can benefit programs wanting to minimize the time needed to access memory the first time, as the kernel would need to read in from the file.

MADV_DONTNEED
Tells the kernel that the specified address range is no longer needed, so the system starts to free the resources associated with the address range.

madvise( ) should be used by programs with specific knowledge of their access patterns over a memory object, such as a mapped file, to increase system performance.

RETURN VALUES

madvise( ) returns:

0 on success.

–1 on failure and sets **errno** to indicate the error.

ERRORS

EINVAL *addr* is not a multiple of the page size as returned by **sysconf**(3C).
The length of the specified address range is less than or equal to 0, or the advice was invalid.

EIO An I/O error occurred while reading from or writing to the file system.

ENOMEM Addresses in the range [*addr, addr + len*) are outside the valid range for the address space of a process, or specify one or more pages that are not mapped.

ESTALE Stale nfs file handle.

SEE ALSO

cat(1), **cp**(1), **mmap**(2), **sysconf**(3C)

NAME

maillock – manage lockfile for user's mailbox

SYNOPSIS

cc [*flag* ...] *file* ... **–lmail** [*library* ...]

#include <maillock.h>

int maillock(const char ∗*user***, int** *retrycnt***);**

int mailunlock(void);

MT-LEVEL

Unsafe

DESCRIPTION

The **maillock( )** function attempts to create a lockfile for the user's mailfile. If a lockfile already exists, and it has not been modified in the last 5 minutes, **maillock( )** will remove the lockfile and set it's own lockfile.

It is crucial that programs locking mail files refresh their locks at least every minute. Lockfiles are refreshed by calling the routine **touchlock( )** with no arguments. It should be called before doing anything that might keep the system busy for a long time. It should be called at least every 3 minutes or the process may lose its lock on the mail file. In practice, arranging to call it once a minute works well.

The algorithm used to determine the age of the lockfile takes into account clock drift between machines using a network file system. A zero is written into the lockfile so that the lock will be respected by systems running the standard version of System V.

If the lockfile has been modified in the last 5 minutes the process will sleep until the lock is available. The sleep algorithm is to sleep for 5 seconds times the attempt number. That is, the first sleep will be for 5 seconds, the next sleep will be for 10 seconds, etc. until the number of attempts reaches *retrycnt*.

When the lockfile is no longer needed, it should be removed by calling **mailunlock( )**.

user is the login name of the user for whose mailbox the lockfile will be created. **maillock( )** assumes that users' mailfiles are in the ''standard'' place as defined in **<maillock.h>**.

RETURN VALUES

The following return code definitions are contained in **<maillock.h>**.

```
#define L_SUCCESS   0   /* Lockfile created or removed */
#define L_NAMELEN   1   /* Recipient name > 13 chars */
#define L_TMPLOCK   2   /* Can't create tmp file */
#define L_TMPWRITE  3   /* Can't write pid into lockfile */
#define L_MAXTRYS   4   /* Failed after retrycnt attempts */
#define L_ERROR     5   /* Check errno for reason */
```

FILES

LIBDIR/llib-mail.ln
LIBDIR/mail.a
/var/mail/∗
/var/mail/∗.lock

NOTES

mailunlock( ) will only remove the lockfile created from the most previous call to **maillock( )**. Calling **maillock( )** for different users without intervening calls to **mailunlock( )** will cause the initially created lockfile(s) to remain, potentially blocking subsequent message delivery until the current process finally terminates.

NAME

makecontext, swapcontext – manipulate user contexts

SYNOPSIS

```
#include <ucontext.h>
void makecontext(ucontext_t *ucp, void(*func)(), int argc, ...);
int swapcontext(ucontext_t *oucp, ucontext_t *ucp);
```

MT-LEVEL

MT-Safe

DESCRIPTION

These functions are useful for implementing user-level context switching between multiple threads of control within a process.

makecontext( ) modifies the context specified by *ucp*, which has been initialized using **getcontext( )**; when this context is resumed using **swapcontext( )** or **setcontext( )** (see **getcontext**(2)), program execution continues by calling the function *func*, passing it the arguments that follow *argc* in the **makecontext( )** call. The integer value of *argc* must be one-greater-than the number of arguments that follow *argc*; otherwise, the behavior is undefined. For 5 arguments, the value of *argc* must be 6.

swapcontext( ) saves the current context in the context structure pointed to by *oucp* and sets the context to the context structure pointed to by *ucp*.

RETURN VALUES

On successful completion, **swapcontext** return a value of zero. Otherwise, a value of –1 is returned and **errno** is set to indicate the error.

ERRORS

These functions will fail if either of the following is true:

EFAULT *ucp* or *oucp* points to an invalid address.

ENOMEM *ucp* does not have enough stack left to complete the operation.

SEE ALSO

exit(2), **getcontext**(2), **sigaction**(2), **sigprocmask**(2), **ucontext**(5)

NOTES

The size of the **ucontext_t** structure may change in future releases. To remain binary compatible, users of these features must always use **makecontext( )** or **getcontext( )** to create new instances of them.

NAME

makedev, major, minor – manage a device number

SYNOPSIS

```
#include <sys/types.h>
#include <sys/mkdev.h>

dev_t makedev(major_t maj, minor_t min);

major_t major(dev_t device);

minor_t minor(dev_t device);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **makedev( )** routine returns a formatted device number on success and **NODEV** on failure. *maj* is the major number. *min* is the minor number. **makedev( )** can be used to create a device number for input to **mknod**(2).

The **major( )** routine returns the major number component from *device.*

The **minor( )** routine returns the minor number component from *device.*

RETURN VALUES

On failure, **NODEV** is returned and **errno** is set to indicate the error.

ERRORS

makedev( ) will fail if one or more of the following are true:

EINVAL One or both of the arguments *maj* and *min* is too large.

EINVAL The *device* number created from *maj* and *min* is **NODEV**.

major( ) will fail if one or more of the following are true:

EINVAL The *device* argument is **NODEV**.

EINVAL The major number component of *device* is too large.

minor( ) will fail if the following is true:

EINVAL The *device* argument is **NODEV**.

SEE ALSO

mknod(2), **stat**(2)

NAME

malloc, calloc, free, memalign, realloc, valloc, alloca – memory allocator

SYNOPSIS

```
#include <stdlib.h>
void *malloc(size_t size);
void *calloc(size_t nelem, size_t elsize);
void free(void *ptr);
void *memalign(size_t alignment, size_t size);
void *realloc(void *ptr, size_t size);
void *valloc(size_t size);
#include <alloca.h>
void *alloca(size_t size);
```

MT-LEVEL

Safe

DESCRIPTION

malloc( ) and **free( )** provide a simple general-purpose memory allocation package. **malloc( )** returns a pointer to a block of at least *size* bytes suitably aligned for any use.

The argument to **free( )** is a pointer to a block previously allocated by **malloc( ) , calloc( )** or **realloc( )**. After **free( )** is performed this space is made available for further allocation. If *ptr* is a **NULL** pointer, no action occurs.

Undefined results will occur if the space assigned by **malloc( )** is overrun or if some random number is handed to **free( )**.

calloc( ) allocates space for an array of *nelem* elements of size *elsize*. The space is initialized to zeros.

memalign( ) allocates *size* bytes on a specified alignment boundary, and returns a pointer to the allocated block. The value of the returned address is guaranteed to be an even multiple of *alignment*. Note: the value of *alignment* must be a power of two, and must be greater than or equal to the size of a word.

realloc( ) changes the size of the block pointed to by *ptr* to *size* bytes and returns a pointer to the (possibly moved) block. The contents will be unchanged up to the lesser of the new and old sizes. If *ptr* is **NULL**, **realloc( )** behaves like **malloc( )** for the specified size. If *size* is zero and *ptr* is not a null pointer, the object pointed to is freed.

valloc( ) is equivalent to **memalign(sysconf(_SC_PAGESIZE),size)**.

Each of the allocation routines returns a pointer to space suitably aligned (after possible pointer coercion) for storage of any type of object.

malloc( ), **realloc( )**, **memalign( )**, and **valloc( )** will fail if there is not enough available memory.

alloca( ) allocates *size* bytes of space in the stack frame of the caller, and returns a pointer to the allocated block. This temporary space is automatically freed when the caller returns. Note: if the allocated block is beyond the current stack limit, the resulting behavior is undefined.

RETURN VALUES

If there is no available memory, **malloc( )**, **realloc( )**, **memalign( )**, **valloc( )**, and **calloc( )** return a null pointer. When **realloc( )** returns **NULL**, the block pointed to by *ptr* is left intact. If *size*, *nelem*, or *elsize* is 0, a unique pointer to the arena is returned.

ERRORS

If **malloc( )**, **calloc( )**, or **realloc( )** returns unsuccessfully, **errno** will be set to indicate the following:

ENOMEM *size* bytes of memory exceeds the physical limits of your system, and cannot be allocated.

EAGAIN There is not enough memory available AT THIS POINT IN TIME to allocate *size* bytes of memory; but the application could try again later.

SEE ALSO

brk(2), **getrlimit**(2), **bsdmalloc**(3X), **malloc**(3X), **mapmalloc**(3X)

WARNINGS

Undefined results will occur if the size requested for a block of memory exceeds the maximum size of a process's heap, which may be obtained with **getrlimit( )**.

alloca( ) is machine-, compiler-, and most of all, system-dependent. Its use is strongly discouraged.

NOTES

Comparative Features of **malloc**(3C), **bsdmalloc**(3X), and **malloc**(3X)**:**

- The **bsdmalloc**(3X) routines afford better performance, but are space-inefficient.
- The **malloc**(3X) routines are space-efficient, but have slower performance.
- The standard, fully SCD-compliant **malloc** routines are a trade-off between performance and space-efficiency.

free( ) does not set **errno.**

NAME

malloc, free, realloc, calloc, mallopt, mallinfo – memory allocator

SYNOPSIS

cc [*flag* ...] *file* ... **–lmalloc** [*library* ...]

#include <stdlib.h>

void **malloc(size_t** *size*);

void free(void **ptr*);

void **realloc(void** **ptr*, **size_t** *size*);

void **calloc(size_t** *nelem*, **size_t** *elsize*);

#include <malloc.h>

int mallopt(int *cmd*, **int** *value*);

struct mallinfo mallinfo(void);

MT-LEVEL

Safe

DESCRIPTION

malloc( ) and **free( )** provide a simple general-purpose memory allocation package.

malloc( ) returns a pointer to a block of at least *size* bytes suitably aligned for any use.

The argument to **free( )** is a pointer to a block previously allocated by **malloc( )**; after **free( )** is performed this space is made available for further allocation, and its contents have been destroyed (but see **mallopt( )** below for a way to change this behavior). If *ptr* is a null pointer, no action occurs.

Undefined results occur if the space assigned by **malloc( )** is overrun or if some random number is handed to **free( ).**

realloc( ) changes the size of the block pointed to by *ptr* to *size* bytes and returns a pointer to the (possibly moved) block. The contents are unchanged up to the lesser of the new and old sizes. If *ptr* is a null pointer, **realloc( )** behaves like **malloc( )** for the specified size. If *size* is zero and *ptr* is not a null pointer, the object it points to is freed.

calloc( ) allocates space for an array of *nelem* elements of size *elsize*. The space is initialized to zeros.

mallopt( ) provides for control over the allocation algorithm. The available values for *cmd* are:

M_MXFAST Set *maxfast* to *value.* The algorithm allocates all blocks below the size of *maxfast* in large groups and then doles them out very quickly. The default value for *maxfast* is 24.

M_NLBLKS Set *numlblks* to *value*. The above mentioned ''large groups'' each contain *numlblks* blocks. *numlblks* must be greater than 0. The default value for *numlblks* is 100.

M_GRAIN Set *grain* to *value*. The sizes of all blocks smaller than *maxfast* are considered to be rounded up to the nearest multiple of *grain*. *grain* must be greater than 0. The default value of *grain* is the smallest number of bytes that will allow

alignment of any data type. Value will be rounded up to a multiple of the default when *grain* is set.

M_KEEP Preserve data in a freed block until the next **malloc( )**, **realloc( )**, or **calloc( )**. This option is provided only for compatibility with the old version of **malloc( )** and is not recommended.

These values are defined in the **<malloc.h>** header.

mallopt( ) may be called repeatedly, but may not be called after the first small block is allocated.

mallinfo( ) provides instrumentation describing space usage. It returns the **mallinfo** structure with the following members:

```
int arena;      /* total space in arena */
int ordblks;    /* number of ordinary blocks */
int smblks;     /* number of small blocks */
int hblkhd;     /* space in holding block headers */
int hblks;      /* number of holding blocks */
int usmblks;    /* space in small blocks in use */
int fsmblks;    /* space in free small blocks */
int uordblks;   /* space in ordinary blocks in use */
int fordblks;   /* space in free ordinary blocks */
int keepcost;   /* space penalty if keep option
                   is used */
```

The **mallinfo** structure is defined in the **<malloc.h>** header.

Each of the allocation routines returns a pointer to space suitably aligned (after possible pointer coercion) for storage of any type of object.

RETURN VALUES

malloc( ), **realloc( )**, and **calloc( )** return a **NULL** pointer if there is not enough available memory. When **realloc( )** returns **NULL**, the block pointed to by *ptr* is left intact. If **mallopt( )** is called after any allocation or if *cmd* or *value* are invalid, non-zero is returned. Otherwise, it returns zero.

ERRORS

If **malloc( )**, **calloc( )**, or **realloc( )** returns unsuccessfully, **errno** will be set to indicate the following:

ENOMEM *size* bytes of memory exceeds the physical limits of your system, and cannot be allocated.

EAGAIN There is not enough memory available AT THIS POINT IN TIME to allocate *size* bytes of memory; but the application could try again later.

SEE ALSO

brk(2), **malloc**(3C), **bsdmalloc**(3X)

NOTES

Note that unlike **malloc**(3C), this package does not preserve the contents of a block when it is freed, unless the **M_KEEP** option of **mallopt( )** is used.

Undocumented features of **malloc**(3C) have not been duplicated.

Function prototypes for **malloc( )**, **realloc( )**, **calloc( )**, and **free( )** are also defined in the **<malloc.h>** header for compatibility with old applications. New applications should include **<stdlib.h>** to access the prototypes for these functions.

Comparative Features of **malloc**(3X), **bsdmalloc**(3X), and **malloc**(3C)**:**

- These **malloc**(3X) routines are space-efficient, but have slower performance.
- The **bsdmalloc**(3X) routines afford better performance, but are space-inefficient.
- The standard, fully SCD-compliant **malloc**(3C) routines are a trade-off between performance and space-efficiency.

free( ) does not set **errno.**

NAME

mapmalloc, calloc, cfree, free, realloc, – memory allocator

SYNOPSIS

cc [*flag* ...] *file* ... **–lmapmalloc** [*library* ...]

#include <stdlib.h>

void *__malloc(size_t__ *size*__);__

void *__calloc(size_t__ *nelem*__, size_t__ *elsize*__);__

void cfree(void *__ptr__, **unsigned** *num*__, unsigned__ *size*__);__

void free(void * *ptr*__);__

void *__realloc(void__ *__ptr__, **size_t** *size*__);__

MT-LEVEL

Safe

DESCRIPTION

The collection of **malloc** routines in this library use **mmap**(2) instead of **sbrk**(2) for acquiring new heap space. The routines in this library are intended to be used only if necessary, when applications must call **sbrk( ),** but need to call other library routines that might call **malloc.** The algorithms used by these routines are not sophisticated. There is no reclaiming of memory.

malloc( ) and **free( )** provide a simple general-purpose memory allocation package. **malloc( )** returns a pointer to a block of at least *size* bytes suitably aligned for any use.

The argument to **free( )** is a pointer to a block previously allocated by **malloc( )**, **calloc( )** or **realloc( )**. If *ptr* is a **NULL** pointer, no action occurs.

Undefined results will occur if the space assigned by **malloc( )** is overrun or if some random number is handed to **free( )**.

calloc( ) allocates space for an array of *nelem* elements of size *elsize*. The space is initialized to zeros.

realloc( ) changes the size of the block pointed to by *ptr* to *size* bytes and returns a pointer to the (possibly moved) block. The contents will be unchanged up to the lesser of the new and old sizes. If *ptr* is **NULL**, **realloc( )** behaves like **malloc( )** for the specified size. If *size* is zero and *ptr* is not a null pointer, the object pointed to is freed.

Each of the allocation routines returns a pointer to space suitably aligned (after possible pointer coercion) for storage of any type of object.

malloc( ) and **realloc**( ) will fail if there is not enough available memory.

Entry points for **malloc_debug**( )**, mallocmap**( )**, mallopt**( )**, mallinfo**( )**, memalign**( )**,** and **valloc**( )**,** are empty routines, and are provided only to protect the user from mixing **malloc**( ) functions from different implementations.

RETURN VALUES

If there is no available memory, **malloc( )**, **realloc( )**, and **calloc( )** return a null pointer. When **realloc( )** returns **NULL**, the block pointed to by *ptr* is left intact. If *size*, *nelem*, or *elsize* is 0, a unique pointer to the arena is returned.

FILES **/usr/lib/libmapmalloc**

SEE ALSO **brk**(2), **getrlimit**(2), **mmap**(2), **realloc**(3C)

NAME

matherr – math library exception-handling function

SYNOPSIS

```
#include <math.h>
int matherr(struct exception *exc);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The SVID3 (*System V Interface Definition Third Edition*) specifies that certain **libm** functions call **matherr( )** when exceptions are detected. Users may define their own mechanisms for handling exceptions, by including a function named **matherr( )** in their programs. **matherr( )** is of the form described above. When an exception occurs, a pointer to the exception structure *exc* will be passed to the user-supplied **matherr( )** function. This structure, which is defined in the **<math.h>** header file, is as follows:

```
struct exception {
        int type;
        char *name;
        double arg1, arg2, retval;
};
```

The element **type** is an integer describing the type of exception that has occurred, from the following list of constants (defined in the header file):

DOMAIN	argument domain exception
SING	argument singularity
OVERFLOW	overflow range exception
UNDERFLOW	underflow range exception
TLOSS	total loss of significance
PLOSS	partial loss of significance

Note that both **TLOSS** and **PLOSS** reflect limitations of particular algorithms for trigonometric functions that suffer abrupt declines in accuracy at definite boundaries. Since the Sun implementation does not suffer such abrupt declines, **PLOSS** is never signaled. **TLOSS** is signaled for Bessel functions *only* to satisfy SVID3 requirements.

The element **name** points to a string containing the name of the function that incurred the exception. The elements **arg1** and **arg2** are the arguments with which the function was invoked. **retval** is set to the default value that will be returned by the function unless the user's **matherr( )** sets it to a different value.

If the user's **matherr( )** function returns non-zero, no exception message will be printed, and **errno** will not be set.

SVID3 STANDARD CONFORMANCE

If **matherr( )** is not supplied by the user, the default matherr exception-handling mechanisms, summarized in the table below, will be invoked upon exception:

DOMAIN

0.0 is usually returned, **errno** is set to **EDOM**, and a message is usually printed on standard error.

SING The largest finite single-precision number, **HUGE** of appropriate sign is returned, **errno** is set to **EDOM**, and a message is printed on standard error.

OVERFLOW
The largest finite single-precision number, **HUGE** of appropriate sign is usually returned, **errno** is set to **ERANGE**.

UNDERFLOW
0.0 is returned, and **errno** is set to **ERANGE**.

TLOSS 0.0 is returned, **errno** is set to **ERANGE**, and a message is printed on standard error.

In general, **errno** is not a reliable error indicator in that it may be unexpectedly set by a function in a handler for an asynchronous signal.

DEFAULT ERROR HANDLING PROCEDURES (SVID3)

	Types of Errors				
<math.h> type	DOMAIN	SING	OVERFLOW	UNDERFLOW	TLOSS
errno	EDOM	EDOM	ERANGE	ERANGE	ERANGE
IEEE Exception	Invalid Operation	Division by Zero	Overflow	Underflow	–
fp_exception_type	fp_invalid	fp_division	fp_overflow	fp_underflow	–
ACOS, ASIN (|x| > 1):	Md, 0.0	–	–	–	–
ACOSH (x < 1), ATANH (|x| > 1):	NaN	–	–	–	–
ATAN2 (0,0):	Md, 0.0	–	–	–	–
COSH, SINH:	–	–	±HUGE	–	–
EXP:	–	–	+HUGE	0.0	–
FMOD (x,0):	x	–	–	–	–
HYPOT:	–	–	+HUGE	–	–
J0, J1, JN (|x| > X_TLOSS):	–	–	–	–	Mt, 0.0
LGAMMA:					
usual cases	–	–	+HUGE	–	–
(x = 0 or –integer)	–	Ms, +HUGE	–	–	–
LOG, LOG10:					
(x < 0)	Md, –HUGE	–	–	–	–
(x = 0)	–	Ms, –HUGE	–	–	–
POW:					
usual cases	–	–	±HUGE	±0.0	–
(x < 0) ** (y not an integer)	Md, 0.0	–	–	–	–
0 ** 0	Md, 0.0	–	–	–	–
0 ** (y < 0)	Md, 0.0	–	–	–	–
REMAINDER (x,0):	NaN	–	–	–	–
SCALB:	–	–	±HUGE_VAL	±0.0	–
SQRT (x < 0):	Md, 0.0	–	–	–	–
Y0, Y1, YN:					
(x < 0)	Md, –HUGE	–	–	–	–
(x = 0)	–	Md, –HUGE	–	–	–
(x > X_TLOSS)	–	–	–	–	Mt, 0.0

ABBREVIATIONS	
Md	Message is printed (DOMAIN error).
Ms	Message is printed (SING error).
Mt	Message is printed (TLOSS error).
NaN	IEEE NaN result and invalid operation exception.
HUGE	Maximum finite single-precision floating-point number.
HUGE_VAL	IEEE ∞ result and division-by-zero exception.
X_TLOSS	The value X_TLOSS is defined in <values.h>.

The interaction of IEEE arithmetic and **matherr( )** is not defined when executing under IEEE rounding modes other than the default round to nearest: **matherr( )** may not be called on overflow or underflow, and the SUN-provided **matherr( )** may return results that differ from those in this table.

X/OPEN (XPG3) STANDARD CONFORMANCE

XPG3 (*X/Open Portability Guide Issue 3*) no longer sanctions the use of the **matherr( )** interface. The following table summarizes the values returned in the exceptional cases. In general, XPG3 dictates that as long as one of the input argument(s) is a NaN, NaN shall be returned. In particular, **pow(**NaN,0**)** = NaN.

X / Open (XPG3) ERROR HANDLING PROCEDURES (compile with "cc –Xa")					
	Types of Errors				
<math.h> type	DOMAIN	SING	OVERFLOW	UNDERFLOW	TLOSS
errno	EDOM	EDOM	ERANGE	ERANGE	ERANGE
ACOS, ASIN (| x | > 1):	0.0	–	–	–	–
ATAN2 (0,0):	0.0	–	–	–	–
COSH, SINH:	–	–	{±HUGE_VAL}	–	–
EXP:	–	–	{+HUGE_VAL}	{0.0}	–
FMOD (x,0):	{NaN}	–	–	–	–
HYPOT:	–	–	{+HUGE_VAL}	–	–
J0, J1, JN (| x | > X_TLOSS):	–	–	–	–	{0.0}
LGAMMA:					
usual cases	–	–	{+HUGE_VAL}	–	–
(x = 0 or –integer)	–	+HUGE_VAL	–	–	–
LOG, LOG10:					
(x < 0)	–HUGE_VAL	–	–	–	–
(x = 0)	–	–HUGE_VAL	–	–	–
POW:					
usual cases	–	–	±HUGE_VAL	±0.0	–
(x < 0) ** (y not an integer)	0.0	–	–	–	–
0 ** 0	{1.0}	–	–	–	–
0 ** (y < 0)	{–HUGE_VAL}	–	–	–	–
SQRT (x < 0):	0.0	–	–	–	–
Y0, Y1, YN:					
(x < 0)	{–HUGE_VAL}	–	–	–	–
(x = 0)	–	{–HUGE_VAL}	–	–	–
(x > X_TLOSS)	–	–	–	–	0.0

ABBREVIATIONS	
{...}	**errno** is not to be relied upon in all braced cases.
NaN	IEEE NaN result and invalid operation exception.
HUGE_VAL	IEEE ∞ result and division-by-zero exception.
X_TLOSS	The value X_TLOSS is defined in <values.h>.

ANSI/ISO-C STANDARD CONFORMANCE

The ANSI/ISO-C standard covers a small subset of XPG3.

The following table summarizes the values returned in the exceptional cases.

ANSI/ISO-C ERROR HANDLING PROCEDURES (compile with "cc –Xc")				
	Types of Errors			
<math.h> type	DOMAIN	SING	OVERFLOW	UNDERFLOW
errno	EDOM	EDOM	ERANGE	ERANGE
ACOS, ASIN (| x | > 1):	0.0	–	–	–
ATAN2 (0,0):	0.0	–	–	–
EXP:	–	–	+HUGE_VAL	0.0
FMOD (x,0):	NaN	–	–	–
LOG, LOG10:				
(x < 0)	–HUGE_VAL	–	–	–
(x = 0)	–	–HUGE_VAL	–	–
POW:				
usual cases	–	–	±HUGE_VAL	±0.0
(x < 0) ** (y not an integer)	0.0	–	–	–
0 ** (y < 0)	–HUGE_VAL	–	–	–
SQRT (x < 0):	0.0	–	–	–

ABBREVIATIONS	
NaN	IEEE NaN result and invalid operation exception.
HUGE_VAL	IEEE ∞ result and division-by-zero exception.

EXAMPLES

```
#include <stdio.h>
#include <stdlib.h>
#include <math.h>

int
matherr(struct exception *x) {
        switch (x->type) {
        case DOMAIN:
                /* change sqrt to return sqrt(-arg1), not NaN */
                if (!strcmp(x->name, "sqrt")) {
                        x->retval = sqrt(-x->arg1);
                        return (0); /* print message and set errno */
                } /* FALLTHRU */
        case SING:
                /* all other domain or sing exceptions, print message and */
                /* abort */
                fprintf(stderr, "domain exception in %s\n", x->name);
```

```
                abort( );
                break;
        }
        return (0); /* all other exceptions, execute default procedure */
}
```

NAME

mbchar, mbtowc, mblen, wctomb – multibyte character handling

SYNOPSIS

```
#include <stdlib.h>
#include <limits.h>
int mbtowc(wchar_t *pwc, const char *s, size_t n);
int mblen(const char *s, size_t n);
int wctomb(char *s, wchar_t wchar);
```

MT-LEVEL

MT-Safe with exceptions

DESCRIPTION

Multibyte characters are used to represent characters in an extended character set. This is needed for locales where 8 bits are not enough to represent all the characters in the character set.

The multibyte character handling functions provide the means of translating multibyte characters into wide characters and back again. Wide characters have type **wchar_t** (defined in **<stdlib.h>**), which is an integral type whose range of values can represent distinct codes for all members of the largest extended character set specified among the supported locales.

A maximum of 3 extended character sets are supported for each locale. The number of bytes in an extended character set is defined by the **LC_CTYPE** category of the locale (see **setlocale**(3C)). However, the maximum number of bytes in any multibyte character will never be greater than **MB_LEN_MAX**. which is defined in **<limits.h>**. The maximum number of bytes in a character in an extended character set in the **current** locale is given by the macro, **MB_CUR_MAX**, defined in **<stdlib.h>**.

mbtowc( ) determines the number of bytes that comprise the multibyte character pointed to by *s*. Also, if *pwc* is not a null pointer, **mbtowc( )** converts the multibyte character to a wide character and places the result in the object pointed to by *pwc*. (The value of the wide character corresponding to the null character is zero.) At most *n* bytes will be examined, starting at the byte pointed to by *s*.

If *s* is a null pointer, **mbtowc( )** simply returns 0. If *s* is not a null pointer, then, if *s* points to the null character, **mbtowc( )** returns 0; if the next *n* or fewer bytes form a valid multibyte character, **mbtowc( )** returns the number of bytes that comprise the converted multibyte character; otherwise, *s* does not point to a valid multibyte character and **mbtowc( )** returns –1.

mblen( ) determines the number of bytes comprising the multibyte character pointed to by *s*. It is equivalent to

```
mbtowc((wchar_t *)0, s, n);
```

wctomb( ) determines the number of bytes needed to represent the multibyte character corresponding to the code whose value is *wchar*, and, if *s* is not a null pointer, stores the multibyte character representation in the array pointed to by *s*. At most **MB_CUR_MAX** bytes are stored.

If *s* is a null pointer, **wctomb( )** simply returns 0. If *s* is not a null pointer, **wctomb( )** returns –1 if the value of *wchar* does not correspond to a valid multibyte character; otherwise it returns the number of bytes that comprise the multibyte character corresponding to the value of *wchar*.

SEE ALSO

chrtbl(1M), **mbstring**(3C), **setlocale**(3C), **environ**(5)

NOTES

mbchar, **mbtowc**, **mblen** and **wctomb** can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

mbstring, mbstowcs, wcstombs – multibyte string functions

SYNOPSIS

```
#include <stdlib.h>

size_t mbstowcs(wchar_t *pwcs, const char *s, size_t n);
size_t wcstombs(char *s, const wchar_t *pwcs, size_t n);
```

MT-LEVEL

MT-Safe with exceptions

DESCRIPTION

mbstowcs( ) converts a sequence of multibyte characters from the array pointed to by *s* into a sequence of corresponding wide character codes and stores these codes into the array pointed to by *pwcs*, stopping after *n* codes are stored or a code with value zero (a converted null character) is stored. If an invalid multibyte character is encountered, **mbstowcs( )** returns **(size_t) –1**; otherwise, **mbstowcs( )** returns the number of array elements modified, not including the terminating zero code, if any.

wcstombs( ) converts a sequence of wide character codes from the array pointed to by *pwcs* into a sequence of multibyte characters and stores these multibyte characters into the array pointed to by *s*, stopping if a multibyte character would exceed the limit of *n* total bytes or if a null character is stored. If a wide character code is encountered that does not correspond to a valid multibyte character, **wcstombs( )** returns **(size_t) –1**; otherwise, **wcstombs( )** returns the number of bytes modified, not including a terminating null character, if any. If *s* is a null pointer, **wcstombs( )** returns the number of bytes required for the character array.

SEE ALSO

chrtbl(1M), **mbchar**(3C), **setlocale**(3C), **environ**(5)

NOTES

mbstowcs and **wcstombs** can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

mctl – memory management control

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <sys/types.h>
#include <sys/mman.h>

int mctl(*addr, len, function, arg***)**
caddr_t *addr***;**
size_t *len***;**
int *function***;**
int *arg***;**

DESCRIPTION

mctl( ) applies a variety of control functions over pages identified by the mappings established for the address range [*addr, addr + len*). The function to be performed is identified by the argument *function*. Valid functions are defined in **mman.h** as follows:

MC_LOCK

Lock the pages in the range in memory. This function is used to support **mlock( )**. See **mlock**(3C) for semantics and usage. *arg* is ignored.

MC_LOCKAS

Lock the pages in the address space in memory. This function is used to support **mlockall( )**. See **mlockall**(3C) for semantics and usage. *addr* and *len* are ignored. *arg* is an integer built from the flags:

MCL_CURRENT	Lock current mappings
MCL_FUTURE	Lock future mappings

MC_SYNC

Synchronize the pages in the range with their backing storage. Optionally invalidate cache copies. This function is used to support **msync( )**. See **msync**(3C) for semantics and usage. *arg* is used to represent the *flags* argument to **msync( )**. It is constructed from an OR of the following values:

MS_SYNC	Synchronized write
MS_ASYNC	Return immediately
MS_INVALIDATE	Invalidate mappings

MS_ASYNC returns after all I/O operations are scheduled. **MS_SYNC** does not return until all I/O operations are complete. Specify exactly one of **MS_ASYNC** or **MS_SYNC**. **MS_INVALIDATE** invalidates all cached copies of data from memory, requiring them to be re-obtained from the object's permanent storage location upon the next reference.

MC_UNLOCK

Unlock the pages in the range. This function is used to support **munlock( )**. *arg* is ignored.

MC_UNLOCKAS
Remove address space memory lock, and locks on all current mappings. This function is used to support **munlockall()**. *addr* and *len* must have the value 0. *arg* is ignored.

RETURN VALUES

mctl() returns 0 on success, –1 on failure.

ERRORS

mctl() fails if:

EAGAIN	Some or all of the memory identified by the operation could not be locked due to insufficient system resources.
EBUSY	**MS_INVALIDATE** was specified and one or more of the pages is locked in memory.
EINVAL	*addr* is not a multiple of the page size as returned by **getpagesize()**.
EINVAL	*addr* and/or *len* do not have the value 0 when **MC_LOCKAS** or **MC_UNLOCKAS** are specified.
EINVAL	*arg* is not valid for the function specified.
EIO	An I/O error occurred while reading from or writing to the file system.
ENOMEM	Addresses in the range [*addr*, *addr* + *len*) are invalid for the address space of a process, or specify one or more pages which are not mapped.
EPERM	The process's effective user ID is not super-user and one of **MC_LOCK MC_LOCKAS**, **MC_UNLOCK**, or **MC_UNLOCKAS** was specified.

SEE ALSO

mmap(2), **memcntl**(2), **getpagesize**(3C), **mlock**(3C), **mlockall**(3C), **msync**(3C)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

NAME

media_findname – convert a supplied name into an absolute pathname that can be used to access removable media

SYNOPSIS

cc [*flag* . . .] *file*. . . –**lvolmgt** [*library*. . . .]

#include <volmgt.h>

char ∗**media_findname(char** ∗*start***);**

MT-LEVEL

MT-Unsafe

DESCRIPTION

media_findname( ) converts the supplied *start* string into an absolute pathname that can then be used to access a particular piece of media.

The *start* parameter can be one of the following types of specifications:

/dev/. . .	An absolute pathname in **/dev**, such as **/dev/rdiskette0**, in which case a copy of that string is returned (see **NOTES** on this page).
/vol/. . .	An absolute Volume Management pathname, such as **/vol/dev/aliases/floppy0** or **/vol/dsk/fred**. If this supplied pathname is not a symbolic link, then a copy of that pathname is returned. If the supplied pathname is a symbolic link then it is dereferenced and a copy of that dereferenced pathname is returned.
volume_name	The Volume Management volume name for a particular volume, such as **fred** (see **fdformat**(1) for a description of how to label floppies). In this case a pathname in the Volume Management namespace is returned.
volmgt_symname	The Volume Management symbolic name for a device, such as **floppy0** or **cdrom2** (see **volfs**(7FS) for more information on Volume Management symbolic names), in which case a pathname in the Volume Management namespace is returned.
media_type	The Volume Management generic media type name. For example, **floppy** or **cdrom**. In this case **media_findname( )** looks for the first piece of media that matches that media type, starting at 0 (zero) and continuing on until a match is found (or some fairly large maximum number is reached). In this case, if a match is found, a copy of the pathname to the volume found is returned.

RETURN VALUES

Upon successful completion **media_findname( )** returns a pointer to the pathname found. In the case of an error a null pointer is returned.

ERRORS

For cases where the supplied *start* parameter is an absolute pathname, **media_findname( )** can fail, returning a null string pointer, if an **lstat**(2) of that supplied pathname fails. Also, if the supplied absolute pathname is a symbolic link, **media_findname( )** can fail if a **readlink**(2) of that symbolic link fails, or if a **stat**(2) of the pathname pointed to by that symbolic link fails, or if any of the following is true:

ENXIO The specified absolute pathname was not a character special device, and it was not a directory with a character special device in it.

EXAMPLES

The following example attempts to find what the Volume Management pathname is to a piece of media called fred. Notice that a **volmgt_check( )** is done first (see the **NOTES** section on this page).

```
(void) volmgt_check(NULL);
if ((nm = media_findname("fred")) != NULL) {
    (void) printf("\"fred\" is at \"%s\"\n", nm);
} else {
    (void) printf("\"fred\" is at \"%s\"\n", nm);
}
```

This example looks for whatever volume is in the first floppy drive, letting **media_findname( )** call **volmgt_check( )** if and only if no floppy is currently known to be the first floppy drive.

```
if ((nm = media_findname("floppy0")) != NULL) {
    (void) printf("path to floppy0 is \"%s\"\n", nm);
} else {
    (void) printf("nothing in floppy0\n");
}
```

SEE ALSO

cc(1B), **fdformat**(1), **vold**(1M), **lstat**(2), **readlink**(2), **stat**(2), **free**(3C), **malloc**(3C), **volmgt_check**(3X), **volmgt_inuse**(3X), **volmgt_root**(3X), **volmgt_running**(3X), **volmgt_symname**(3X), **volfs**(7FS)

NOTES

If **media_findname( )** cannot find a match for the supplied name, it performs a **volmgt_check**(3X) and tries again, so it can be more efficient to perform **volmgt_check( )** before calling **media_findname( )**.

Upon success **media_findname( )** returns a pointer to string which has been allocated; this should be freed when no longer in use (see **free**(3C)).

NAME

media_getattr, media_setattr – get and set media attributes

SYNOPSIS

cc [*flag* . . .] *file*. . . **–lvolmgt** [*library*. . . .]

#include <volmgt.h>

char ∗**media_getattr(char** ∗*vol_path***, char** ∗*attr***);**

int media_setattr(char ∗*vol_path***, char** ∗*attr***, char** ∗*value***);**

MT-LEVEL

MT-Safe

DESCRIPTION

media_setattr( ) and media_getattr( ) respectively set and get attribute-value pairs (called properties) on a per-volume basis.

Volume Management supports system properties and user properties. System properties are ones that Volume Management predefines. Some of these system properties are writable, but only by the user that owns the volume being specified, and some system properties are read only:

Attribute	Writable	Value	Description
s-access	RO	"seq", "rand"	sequential or random access
s-density	RO	"low", "medium", "high"	media density
s-parts	RO	comma separated list of slice numbers	list of partitions on this volume
s-location	RO	*pathname*	Volume Management pathname to media
s-mejectable	RO	"true", "false"	whether or not media is manually ejectable
s-rmoneject	R / W	"true", "false"	should media access points be removed from database upon ejection
s-enxio	R / W	"true", "false"	if set return **ENXIO** when media access attempted

Properties can also be defined by the user. In this case the value can be any string the user wishes.

RETURN VALUES

Upon successful completion **media_getattr( )** returns a pointer to the value corresponding to the specified attribute. A null pointer is returned if the specified volume doesn't exist, if the specified attribute for that volume doesn't exist, if the specified attribute is boolean and its value is false, or if **malloc**(3C) fails to allocate space for the return value.

media_setattr( ) returns **1** upon success, and **0** upon failure.

ERRORS

Both **media_getattr( )** and **media_setattr( )** can fail returning a null pointer if an **open**(2) of the specified *vol_path* fails, if an **fstat**(2) of that pathname fails, or if that pathname is not a block or character special device.

media_getattr( ) can also fail if the specified attribute was not found, and **media_setattr( )** can also fail if the caller doesn't have permission to set the attribute, either because it's is a system attribute, or because the caller doesn't own the specified volume.

Additionally, either routine can fail returning the following error values:

ENXIO The Volume Management daemon, **vold**, is not running

EINTR The routine was interrupted by the user before finishing

EXAMPLES

The following example checks to see if the volume called *fred* that Volume Management is managing can be ejected via software, or if it can only be manually ejected:

```
if (media_getattr("/vol/rdsk/fred", "s-mejectable") != NULL) {
    (void) printf("\"fred\" must be manually ejected\n");
} else {
    (void) printf("software can eject \"fred\"\n");
}
```

This example shows setting the *s-enxio* property for the floppy volume currently in the first floppy drive:

```
int    res;

if ((res = media_setattr("/vol/dev/aliases/floppy0", "s-enxio",
   "true")) == 0) {
    (void) printf("can't set s-enxio flag for floppy0\n");
}
```

SEE ALSO

cc(1B), **vold**(1M), **lstat**(2), **open**(2), **readlink**(2), **stat**(2), **free**(3C), **malloc**(3C), **media_findname**(3X), **volmgt_check**(3X), **volmgt_inuse**(3X), **volmgt_root**(3X), **volmgt_running**(3X), **volmgt_symname**(3X)

NOTES

Upon success **media_getattr( )** returns a pointer to a string which has been allocated, and should be freed when no longer in use (see **free**(3C)).

NAME

memory, memccpy, memchr, memcmp, memcpy, memmove, memset – memory operations

SYNOPSIS

```
#include <string.h>
void *memccpy(void *s1, const void *s2, int c, size_t n);
void *memchr(const void *s, int c, size_t n);
int memcmp(const void *s1, const void *s2, size_t n);
void *memcpy(void *s1, const void *s2, size_t n);
void *memmove(void *s1, const void *s2, size_t n);
void *memset(void *s, int c, size_t n);
```

MT-LEVEL

MT-Safe

DESCRIPTION

These functions operate as efficiently as possible on memory areas (arrays of bytes bounded by a count, not terminated by a null character). They do not check for the overflow of any receiving memory area.

memccpy( ) copies bytes from memory area *s2* into *s1*, stopping after the first occurrence of *c* (converted to an **unsigned char**) has been copied, or after *n* bytes have been copied, whichever comes first. It returns a pointer to the byte after the copy of *c* in *s1*, or a null pointer if *c* was not found in the first *n* bytes of *s2*.

memchr( ) returns a pointer to the first occurrence of *c* (converted to an **unsigned char**) in the first *n* bytes (each interpreted as an **unsigned char**) of memory area *s*, or a null pointer if *c* does not occur.

memcmp( ) compares its arguments, looking at the first *n* bytes (each interpreted as an **unsigned char**), and returns an integer less than, equal to, or greater than 0, according as *s1* is lexicographically less than, equal to, or greater than *s2* when taken to be unsigned characters.

memcpy( ) copies *n* bytes from memory area *s2* to *s1*. It returns *s1*.

memmove( ) copies *n* bytes from memory areas *s2* to *s1*. Copying between objects that overlap will take place correctly. It returns *s1*.

memset( ) sets the first *n* bytes in memory area *s* to the value of *c* (converted to an **unsigned char**). It returns *s*.

SEE ALSO

string(3C)

NAME

menu_attributes, set_menu_fore, menu_fore, set_menu_back, menu_back, set_menu_grey, menu_grey, set_menu_pad, menu_pad – control menus display attributes

SYNOPSIS

cc [*flag* ...] *file* ... –**lmenu -lcurses** [*library* ..]

#include <menu.h>

int set_menu_fore(MENU ∗*menu***, chtype** *attr***);**

chtype menu_fore(MENU ∗*menu***);**

int set_menu_back(MENU ∗*menu***, chtype** *attr***);**

chtype menu_back(MENU ∗*menu***);**

int set_menu_grey(MENU ∗*menu***, chtype** *attr***);**

chtype menu_grey(MENU ∗*menu***);**

int set_menu_pad(MENU ∗*menu***, int** *pad***);**

int menu_pad(MENU ∗*menu***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_menu_fore( ) sets the foreground attribute of *menu* — the display attribute for the current item (if selectable) on single-valued menus and for selected items on multi-valued menus. This display attribute is a **curses** library visual attribute. **menu_fore( )** returns the foreground attribute of *menu*.

set_menu_back( ) sets the background attribute of **menu** — the display attribute for unselected, yet selectable, items. This display attribute is a **curses** library visual attribute.

set_menu_grey( ) sets the grey attribute of *menu* — the display attribute for nonselectable items in multi-valued menus. This display attribute is a **curses** library visual attribute. **menu_grey( )** returns the grey attribute of *menu*.

The pad character is the character that fills the space between the name and description of an item. **set_menu_pad( )** sets the pad character for *menu* to *pad*. **menu_pad( )** returns the pad character of *menu*.

RETURN VALUES

These routines return one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_cursor, pos_menu_cursor – correctly position a menus cursor

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

int pos_menu_cursor(MENU ∗*menu***);**

MT-LEVEL

Unsafe

DESCRIPTION

pos_menu_cursor( ) moves the cursor in the window of *menu* to the correct position to resume menu processing. This is needed after the application calls a **curses** library I/O routine.

RETURN VALUES

This routine returns one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.
E_NOT_POSTED	The menu has not been posted.

SEE ALSO

curses(3X), **menus**(3X), **panel_update**(3X), **panels**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_driver – command processor for the menus subsystem

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

int menu_driver(MENU ∗*menu***, int** *c***);**

MT-LEVEL

Unsafe

DESCRIPTION

menu_driver( ) is the workhorse of the **menus** subsystem. It checks to determine whether the character *c* is a menu request or data. If *c* is a request, the menu driver executes the request and reports the result. If *c* is data (a printable ASCII character), it enters the data into the pattern buffer and tries to find a matching item. If no match is found, the menu driver deletes the character from the pattern buffer and returns **E_NO_MATCH**. If the character is not recognized, the menu driver assumes it is an application-defined command and returns **E_UNKNOWN_COMMAND**.

Menu driver requests:

REQ_LEFT_ITEM	Move left to an item.
REQ_RIGHT_ITEM	Move right to an item.
REQ_UP_ITEM	Move up to an item.
REQ_DOWN_ITEM	Move down to an item.
REQ_SCR_ULINE	Scroll up a line.
REQ_SCR_DLINE	Scroll down a line.
REQ_SCR_DPAGE	Scroll up a page.
REQ_SCR_UPAGE	Scroll down a page.
REQ_FIRST_ITEM	Move to the first item.
REQ_LAST_ITEM	Move to the last item.
REQ_NEXT_ITEM	Move to the next item.
REQ_PREV_ITEM	Move to the previous item.
REQ_TOGGLE_ITEM	Select / de-select an item.
REQ_CLEAR_PATTERN	Clear the menu pattern buffer.
REQ_BACK_PATTERN	Delete the previous character from pattern buffer.
REQ_NEXT_MATCH	Move the next matching item.
REQ_PREV_MATCH	Move to the previous matching item.

RETURN VALUES

menu_driver( ) returns one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.
E_BAD_STATE	The routine was called from an initialization or termination

	function.
E_NOT_POSTED	The menu has not been posted.
E_UNKNOWN_COMMAND	An unknown request was passed to the menu driver.
E_NO_MATCH	The character failed to match.
E_NOT_SELECTABLE	The item cannot be selected.
E_REQUEST_DENIED	The menu driver could not process the request.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

Application defined commands should be defined relative to (greater than) **MAX_COMMAND**, the maximum value of a request listed above.

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_format, set_menu_format – set and get maximum numbers of rows and columns in menus

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

int set_menu_format(MENU ∗*menu***, int** *rows***, int** *cols***);**

void menu_format(MENU ∗*menu***, int** ∗*rows***, int** ∗*cols***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_menu_format( ) sets the maximum number of rows and columns of items that may be displayed at one time on a menu. If the menu contains more items than can be displayed at once, the menu will be scrollable.

menu_format( ) returns the maximum number of rows and columns that may be displayed at one time on *menu*. *rows* and *cols* are pointers to the variables used to return these values.

RETURN VALUES

set_menu_format( ) returns one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.
E_POSTED	The menu is already posted.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_hook, set_item_init, item_init, set_item_term, item_term, set_menu_init, menu_init, set_menu_term, menu_term – assign application-specific routines for automatic invocation by menus

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

int set_item_init(MENU ∗*menu***, void (**∗**func)(MENU** ∗**));**

void (∗**item_init)(MENU** ∗*menu***);**

int set_item_term(MENU ∗*menu***, void (**∗**func)(MENU** ∗**));**

void (∗**item_term)(MENU** ∗*menu***);**

int set_menu_init(MENU ∗*menu***, void (**∗**func)(MENU** ∗**));**

void (∗**menu_init)(MENU** ∗*menu***);**

int set_menu_term(MENU ∗*menu***, void (**∗**func)(MENU** ∗**));**

void (∗**menu_term)(MENU** ∗*menu***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_item_init( ) assigns the application-defined function to be called when the *menu* is posted and just after the current item changes. **item_init( )** returns a pointer to the item initialization routine, if any, called when the *menu* is posted and just after the current item changes.

set_item_term( ) assigns an application-defined function to be called when the *menu* is unposted and just before the current item changes. **item_term( )** returns a pointer to the termination function, if any, called when the *menu* is unposted and just before the current item changes.

set_menu_init( ) assigns an application-defined function to be called when the *menu* is posted and just after the top row changes on a posted menu. **menu_init( )** returns a pointer to the menu initialization routine, if any, called when the *menu* is posted and just after the top row changes on a posted menu.

set_menu_term( ) assigns an application-defined function to be called when the *menu* is unposted and just before the top row changes on a posted menu. **menu_term( )** returns a pointer to the menu termination routine, if any, called when the *menu* is unposted and just before the top row changes on a posted menu.

RETURN VALUES

Routines that return pointers always return NULL on error. Routines that return an integer return one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.

SEE ALSO **curses**(3X), **menus**(3X)

NOTES The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_item_current, set_current_item, current_item, set_top_row, top_row, item_index – set and get current menus items

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

int set_current_item(MENU ∗*menu***, ITEM** ∗*item***);**

ITEM ∗**current_item(MENU** ∗*menu***);**

int set_top_row(MENU ∗*menu***, int** *row***);**

int top_row(MENU ∗*menu***);**

int item_index(ITEM ∗*item***);**

MT-LEVEL

Unsafe

DESCRIPTION

The current item of a menu is the item where the cursor is currently positioned. **set_current_item( )** sets the current item of *menu* to *item*. **current_item( )** returns a pointer to the the current item in *menu*.

set_top_row( ) sets the top row of *menu* to *row*. The left-most item on the new top row becomes the current item. **top_row( )** returns the number of the menu row currently displayed at the top of *menu*.

item_index( ) returns the index to the *item* in the item pointer array. The value of this index ranges from **0** through *N*-**1**, where *N* is the total number of items connected to the menu.

RETURN VALUES

current_item( ) returns NULL on error.

top_row( ) and **index_item( )** return **-1** on error.

set_current_item( ) and **set_top_row( )** return one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.
E_BAD_STATE	The routine was called from an initialization or termination function.
E_NOT_CONNECTED	No items are connected to the menu.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_item_name, item_name, item_description – get menus item name and description

SYNOPSIS

cc [*flag* ...] *file* ... **–lmenu -lcurses** [*library* ..]
#include <menu.h>
char ∗**item_name(ITEM** ∗*item***);**
char ∗**item_description(ITEM** ∗*item***);**

MT-LEVEL

Unsafe

DESCRIPTION

item_name( ) returns a pointer to the name of *item*.
item_description( ) returns a pointer to the description of *item*.

RETURN VALUES

These routines return NULL on error.

SEE ALSO

curses(3X), **menus**(3X), **menu_new**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_item_new, new_item, free_item – create and destroy menus items

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

ITEM ∗**new_item(char** ∗*name***, char** ∗*desc***);**

int free_item(ITEM ∗*item***);**

MT-LEVEL

Unsafe

DESCRIPTION

new_item( ) creates a new item from *name* and *description*, and returns a pointer to the new item.

free_item( ) frees the storage allocated for *item*. Once an item is freed, the user can no longer connect it to a menu.

RETURN VALUES

new_item( ) returns NULL on error.

free_item( ) returns one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.
E_CONNECTED	One or more items are already connected to another menu.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_item_opts, set_item_opts, item_opts_on, item_opts_off, item_opts – menus item option routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

int set_item_opts(ITEM ∗*item***, OPTIONS** *opts***);**

int item_opts_on(ITEM ∗*item***, OPTIONS** *opts***);**

int item_opts_off(ITEM ∗*item***, OPTIONS** *opts***);**

OPTIONS item_opts(ITEM ∗*item***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_item_opts( ) turns on the named options for *item* and turns off all other options. Options are boolean values that can be OR-ed together.

item_opts_on( ) turns on the named options for *item*; no other option is changed.

item_opts_off( ) turns off the named options for *item*; no other option is changed.

item_opts( ) returns the current options of *item*.

Item Options:

O_SELECTABLE The item can be selected during menu processing.

RETURN VALUES

Except for **item_opts( )**, these routines return one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_item_userptr, set_item_userptr, item_userptr – associate application data with menus items

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

int set_item_userptr(ITEM ∗*item***, char** ∗*userptr***);**

char ∗**item_userptr(ITEM** ∗*item***);**

MT-LEVEL

Unsafe

DESCRIPTION

Every item has an associated user pointer that can be used to store relevant information. **set_item_userptr( )** sets the user pointer of *item*. **item_userptr( )** returns the user pointer of *item*.

RETURN VALUES

item_userptr( ) returns NULL on error. **set_item_userptr( )** returns one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_item_value, set_item_value, item_value – set and get menus item values

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

int set_item_value(ITEM ∗*item***, int** *bool***);**

int item_value(ITEM ∗*item***);**

MT-LEVEL

Unsafe

DESCRIPTION

Unlike single-valued menus, multi-valued menus enable the end-user to select one or more items from a menu. **set_item_value( )** sets the selected value of the *item* — **TRUE** (selected) or **FALSE** (not selected). **set_item_value( )** may be used only with multi-valued menus. To make a menu multi-valued, use **set_menu_opts** or **menu_opts_off( )** to turn off the option **O_ONEVALUE**. (See **menu_opts**(3X)).

item_value( ) returns the select value of *item*, either **TRUE** (selected) or **FALSE** (unselected).

RETURN VALUES

set_item_value( ) returns one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_REQUEST_DENIED	The menu driver could not process the request.

SEE ALSO

curses(3X), **menus**(3X), **menu_opts**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_item_visible, item_visible – tell if menus item is visible

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

int item_visible(ITEM ∗*item***);**

MT-LEVEL

Unsafe

DESCRIPTION

A menu item is visible if it currently appears in the subwindow of a posted menu. **item_visible( )** returns **TRUE** if *item* is visible, otherwise it returns **FALSE**.

SEE ALSO

curses(3X), **menus**(3X), **menu_new**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_items, set_menu_items, item_count – connect and disconnect items to and from menus

SYNOPSIS

cc [*flag* ...] *file* ... –**lmenu -lcurses** [*library* ..]

#include <menu.h>

int set_menu_items(MENU ∗*menu***, ITEM** ∗∗*items***);**

ITEM ∗∗**menu_items(MENU** ∗*menu***);**

int item_count(MENU ∗*menu***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_menu_items( ) changes the item pointer array connected to *menu* to the item pointer array *items*.

menu_items( ) returns a pointer to the item pointer array connected to *menu*.

item_count( ) returns the number of items in *menu*.

RETURN VALUES

menu_items( ) returns NULL on error.

item_count( ) returns -1 on error.

set_menu_items( ) returns one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.
E_POSTED	The menu is already posted.
E_CONNECTED	One or more items are already connected to another menu.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_mark, set_menu_mark – menus mark string routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

int set_menu_mark(MENU ∗*menu***, char** ∗*mark***);**

char ∗**menu_mark(MENU** ∗*menu***);**

MT-LEVEL

Unsafe

DESCRIPTION

menus displays mark strings to distinguish selected items in a menu (or the current item in a single-valued menu). **set_menu_mark( )** sets the mark string of *menu* to *mark*. **menu_mark( )** returns a pointer to the mark string of *menu*.

RETURN VALUES

menu_mark( ) returns NULL on error. **set_menu_mark( )** returns one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_new, new_menu, free_menu – create and destroy menus

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

MENU ∗**new_menu(ITEM** ∗∗*items***);**

int free_menu(MENU ∗*menu***);**

MT-LEVEL

Unsafe

DESCRIPTION

new_menu( ) creates a new menu connected to the item pointer array *items* and returns a pointer to the new menu.

free_menu( ) disconnects *menu* from its associated item pointer array and frees the storage allocated for the menu.

RETURN VALUES

new_menu( ) returns NULL on error.

free_menu( ) returns one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.
E_POSTED	The menu is already posted.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_opts, set_menu_opts, menu_opts_on, menu_opts_off – menus option routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

OPTIONS menu_opts(MENU ∗*menu***);**

int set_menu_opts(MENU ∗*menu***, OPTIONS** *opts***);**

int menu_opts_on(MENU ∗*menu***, OPTIONS** *opts***);**

int menu_opts_off(MENU ∗*menu***, OPTIONS** *opts***);**

MT-LEVEL

Unsafe

DESCRIPTION

Menu Options

set_menu_opts( ) turns on the named options for *menu* and turns off all other options. Options are boolean values that can be OR-ed together.

menu_opts_on( ) turns on the named options for *menu*; no other option is changed.

menu_opts_off( ) turns off the named options for *menu*; no other option is changed.

menu_opts( ) returns the current options of *menu.*

The following values can be OR'd together to create *opts*.

O_ONEVALUE	Only one item can be selected from the menu.
O_SHOWDESC	Display the description of the items.
O_ROWMAJOR	Display the menu in row major order.
O_IGNORECASE	Ignore the case when pattern matching.
O_SHOWMATCH	Place the cursor within the item name when pattern matching.
O_NONCYCLIC	Make certain menu driver requests non-cyclic.

RETURN VALUES

Except for **menu_opts( )**, these routines return one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_POSTED	The menu is already posted.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_pattern, set_menu_pattern – set and get menus pattern match buffer

SYNOPSIS

cc [*flag* ...] *file* ... –**lmenu -lcurses** [*library* ..]

#include <menu.h>

char *__**menu_pattern(MENU** *__*menu***);**

int set_menu_pattern(MENU *__*menu***, char** *__*pat***);**

MT-LEVEL

Unsafe

DESCRIPTION

Every menu has a pattern buffer to match entered data with menu items. **set_menu_pattern( )** sets the pattern buffer to *pat* and tries to find the first item that matches the pattern. If it does, the matching item becomes the current item. If not, the current item does not change. **menu_pattern( )** returns the string in the pattern buffer of *menu*.

RETURN VALUES

menu_pattern( ) returns **NULL** on error. **set_menu_pattern( )** returns one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.
E_NO_MATCH	The character failed to match.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_post, post_menu, unpost_menu – write or erase menus from associated subwindows

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lmenu -lcurses** [*library* . .]

#include <menu.h>

int post_menu(MENU ∗*menu***);**

int unpost_menu(MENU ∗*menu***);**

MT-LEVEL

Unsafe

DESCRIPTION

post_menu( ) writes *menu* to the subwindow. The application programmer must use **curses** library routines to display the menu on the physical screen or call **update_panels( )** if the **panels** library is being used.

unpost_menu( ) erases *menu* from its associated subwindow.

RETURN VALUES

These routines return one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.
E_POSTED	The menu is already posted.
E_BAD_STATE	The routine was called from an initialization or termination function.
E_NO_ROOM	The menu does not fit within its subwindow.
E_NOT_POSTED	The menu has not been posted.
E_NOT_CONNECTED	No items are connected to the menu.

SEE ALSO

curses(3X), **menus**(3X), **panels**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_userptr, set_menu_userptr – associate application data with menus

SYNOPSIS

cc [*flag* ...] *file* ... –**lmenu -lcurses** [*library* ..]

#include <menu.h>

char *__**menu_userptr(MENU** *__*menu**);**

int set_menu_userptr(MENU *__*menu**, char** *__*userptr**);**

MT-LEVEL

Unsafe

DESCRIPTION

Every menu has an associated user pointer that can be used to store relevant information. **set_menu_userptr( )** sets the user pointer of *menu*. **menu_userptr( )** returns the user pointer of *menu*.

RETURN VALUES

menu_userptr( ) returns **NULL** on error.

set_menu_userptr( ) returns one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menu_win, set_menu_win, set_menu_sub, menu_sub, scale_menu – menus window and subwindow association routines

SYNOPSIS

cc [*flag* ...] *file* ... –**lmenu -lcurses** [*library* ..]

#include <menu.h>

int set_menu_win(MENU ∗*menu***, WINDOW** ∗*win***);**

WINDOW ∗**menu_win(MENU** ∗*menu***);**

int set_menu_sub(MENU ∗*menu***, WINDOW** ∗*sub***);**

WINDOW ∗**menu_sub(MENU** ∗*menu***);**

int scale_window(MENU ∗*menu***, int** ∗*rows***, int** ∗*cols***);**

MT-LEVEL

Unsafe

DESCRIPTION

set_menu_win( ) sets the window of *menu* to *win*. **menu_win( )** returns a pointer to the window of *menu*.

set_menu_sub( ) sets the subwindow of *menu* to *sub*. **menu_sub( )** returns a pointer to the subwindow of *menu*.

scale_window( ) returns the minimum window size necessary for the subwindow of *menu*. *rows* and *cols* are pointers to the locations used to return the values.

RETURN VALUES

Routines that return pointers always return **NULL** on error. Routines that return an integer return one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.
E_POSTED	The menu is already posted.
E_NOT_CONNECTED	No items are connected to the menu.

SEE ALSO

curses(3X), **menus**(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

menus – character based menus package

SYNOPSIS

#include <menu.h>

MT-LEVEL

Unsafe

DESCRIPTION

The **menu** library is built using the **curses** library, and any program using menus routines must call one of the **curses** initialization routines, such as **initscr**. A program using these routines must be compiled with **–lmenu** and **–lcurses** on the **cc** command line.

The menus package gives the applications programmer a terminal-independent method of creating and customizing menus for user interaction. The menus package includes: item routines, which are used to create and customize menu items; and menu routines, which are used to create and customize menus, assign pre- and post-processing routines, and display and interact with menus.

Current Default Values for Item Attributes

The menus package establishes initial current default values for item attributes. During item initialization, each item attribute is assigned the current default value for that attribute. An application can change or retrieve a current default attribute value by calling the appropriate set or retrieve routine with a **NULL** item pointer. If an application changes a current default item attribute value, subsequent items created using **new_item( )** will have the new default attribute value. The attributes of previously created items are not changed if a current default attribute value is changed.

Routine Name Index

The following table lists each menus routine and the name of the manual page on which it is described.

menus Routine Name	Manual Page Name
current_item	**menu_item_current(3X)**
free_item	**menu_item_new(3X)**
free_menu	**menu_new(3X)**
item_count	**menu_items(3X)**
item_description	**menu_item_name(3X)**
item_index	**menu_item_current(3X)**
item_init	**menu_hook(3X)**
item_name	**menu_item_name(3X)**
item_opts	**menu_item_opts(3X)**
item_opts_off	**menu_item_opts(3X)**
item_opts_on	**menu_item_opts(3X)**
item_term	**menu_hook(3X)**
item_userptr	**menu_item_userptr(3X)**
item_value	**menu_item_value(3X)**
item_visible	**menu_item_visible(3X)**
menu_back	**menu_attributes(3X)**
menu_driver	**menu_driver(3X)**
menu_fore	**menu_attributes(3X)**

menu_format	**menu_format(3X)**
menu_grey	**menu_attributes(3X)**
menu_init	**menu_hook(3X)**
menu_items	**menu_items(3X)**
menu_mark	**menu_mark(3X)**
menu_opts	**menu_opts(3X)**
menu_opts_off	**menu_opts(3X)**
menu_opts_on	**menu_opts(3X)**
menu_pad	**menu_attributes(3X)**
menu_pattern	**menu_pattern(3X)**
menu_sub	**menu_win(3X)**
menu_term	**menu_hook(3X)**
menu_userptr	**menu_userptr(3X)**
menu_win	**menu_win(3X)**
new_item	**menu_item_new(3X)**
new_menu	**menu_new(3X)**
pos_menu_cursor	**menu_cursor(3X)**
post_menu	**menu_post(3X)**
scale_menu	**menu_win(3X)**
set_current_item	**menu_item_current(3X)**
set_item_init	**menu_hook(3X)**
set_item_opts	**menu_item_opts(3X)**
set_item_term	**menu_hook(3X)**
set_item_userptr	**menu_item_userptr(3X)**
set_item_value	**menu_item_value(3X)**
set_menu_back	**menu_attributes(3X)**
set_menu_fore	**menu_attributes(3X)**
set_menu_format	**menu_format(3X)**
set_menu_grey	**menu_attributes(3X)**
set_menu_init	**menu_hook(3X)**
set_menu_items	**menu_items(3X)**
set_menu_mark	**menu_mark(3X)**
set_menu_opts	**menu_opts(3X)**
set_menu_pad	**menu_attributes(3X)**
set_menu_pattern	**menu_pattern(3X)**
set_menu_sub	**menu_win(3X)**
set_menu_term	**menu_hook(3X)**
set_menu_userptr	**menu_userptr(3X)**
set_menu_win	**menu_win(3X)**
set_top_row	**menu_item_current(3X)**
top_row	**menu_item_current(3X)**
unpost_menu	**menu_post(3X)**

RETURN VALUES

Routines that return pointers always return **NULL** on error. Routines that return an integer return one of the following:

E_OK	The routine returned successfully.
E_SYSTEM_ERROR	System error.
E_BAD_ARGUMENT	An incorrect argument was passed to the routine.
E_POSTED	The menu is already posted.
E_CONNECTED	One or more items are already connected to another menu.
E_BAD_STATE	The routine was called from an initialization or termination function.
E_NO_ROOM	The menu does not fit within its subwindow.
E_NOT_POSTED	The menu has not been posted.
E_UNKNOWN_COMMAND	An unknown request was passed to the menu driver.
E_NO_MATCH	The character failed to match.
E_NOT_SELECTABLE	The item cannot be selected.
E_NOT_CONNECTED	No items are connected to the menu.
E_REQUEST_DENIED	The menu driver could not process the request.

SEE ALSO

curses(3X)

NOTES

The header **<menu.h>** automatically includes the headers **<eti.h>** and **<curses.h>**.

NAME

mkdirp, rmdirp – create, remove directories in a path

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

int mkdirp(const char ∗*path*, **mode_t** *mode*);

int rmdirp(char ∗*dir*, **char** ∗*dir1*);

MT-LEVEL

MT-Safe

DESCRIPTION

mkdirp( ) creates all the missing directories in the given *path* with the given *mode*. See **chmod**(2) for the values of *mode*.

rmdirp( ) removes directories in path *dir*. This removal starts at the end of the path and moves back toward the root as far as possible. If an error occurs, the remaining path is stored in *dir1*. **rmdirp( )** returns a 0 only if it is able to remove every directory in the path.

EXAMPLES

```
/* create scratch directories */
if(mkdirp("/tmp/sub1/sub2/sub3", 0755) == –1) {
        fprintf(stderr, "cannot create directory");
        exit(1);
}
chdir("/tmp/sub1/sub2/sub3");
.
.
.
/* cleanup */
chdir("/tmp");
rmdirp("sub1/sub2/sub3");
```

RETURN VALUES

If a needed directory cannot be created, **mkdirp( )** returns –1 and sets **errno** to one of the **mkdir( )** error numbers. If all the directories are created, or existed to begin with, it returns zero.

SEE ALSO

mkdir(2), **rmdir**(2)

NOTES

mkdirp( ) uses **malloc**(3C) to allocate temporary space for the string.

rmdirp( ) returns –2 if a ‘‘**.**’’ or ‘‘**..**’’ is in the path and –3 if an attempt is made to remove the current directory. If an error occurs other than one of the above, –1 is returned.

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

mkfifo – create a new FIFO

SYNOPSIS

```
#include <sys/types.h>
#include <sys/stat.h>

int mkfifo(const char *path, mode_t mode);
```

MT-LEVEL

MT-Safe

Async-Signal-Safe

DESCRIPTION

The **mkfifo( )** routine creates a new FIFO special file named by the pathname pointed to by *path*. The mode of the new FIFO is initialized from *mode*. The file permission bits of the *mode* argument are modified by the process's file creation mask (see **umask**(2)).

The FIFO's owner id is set to the process's effective user id. The FIFO's group id is set to the process's effective group id, or if the **S_ISGID** bit is set in the parent directory then the group id of the FIFO is inherited from the parent directory.

mkfifo( ) calls the **mknod**(2) function to make the file.

RETURN VALUES

Upon successful completion a value of 0 is returned. Otherwise, a value of –1 is returned and **errno** is set to indicate the error.

SEE ALSO

mkdir(1), **chmod**(2), **exec**(2), **mknod**(2), **umask**(2), **fs_ufs**(4), **stat**(5)

NOTES

Bits other than the file permission bits in *mode* are ignored.

NAME

mktemp – make a unique file name

SYNOPSIS

```
#include <stdlib.h>
char *mktemp(char *template);
```

MT-LEVEL

Safe

DESCRIPTION

mktemp( ) replaces the contents of the string pointed to by *template* with a unique file name, and returns *template*. The string in *template* should look like a file name with six trailing **X**s; **mktemp( )** will replace the **X**s with a character string that can be used to create a unique file name.

RETURN VALUES

mktemp( ) will assign to *template* the empty string if it cannot create a unique name.

SEE ALSO

tmpfile(3S), **tmpnam**(3S)

NOTES

mktemp( ) can create only 26 unique file names per thread for each unique *template*.

NAME

mktime – converts a tm structure to a calendar time

SYNOPSIS

```
#include <time.h>
time_t mktime(struct tm *timeptr);
```

MT-LEVEL

Unsafe

DESCRIPTION

mktime( ) converts the time represented by the **tm** structure pointed to by *timeptr* into a calendar time (the number of seconds since 00:00:00 UTC, January 1, 1970).

The **tm** structure contains the following members:

```
int     tm_sec;     /* seconds after the minute [0, 61]  */
int     tm_min;     /* minutes after the hour [0, 59] */
int     tm_hour;    /* hour since midnight [0, 23] */
int     tm_mday;    /* day of the month [1, 31] */
int     tm_mon;     /* months since January [0, 11] */
int     tm_year;    /* years since 1900 */
int     tm_wday;    /* days since Sunday [0, 6] */
int     tm_yday;    /* days since January 1 [0, 365] */
int     tm_isdst;   /* flag for daylight savings time */
```

In addition to computing the calendar time, **mktime( )** normalizes the supplied **tm** structure. The original values of the **tm_wday** and **tm_yday** components of the structure are ignored, and the original values of the other components are not restricted to the ranges indicated in the definition of the structure. On successful completion, the values of the **tm_wday** and **tm_yday** components are set appropriately, and the other components are set to represent the specified calendar time, but with their values forced to be within the appropriate ranges. The final value of **tm_mday** is not set until **tm_mon** and **tm_year** are determined.

The original values of the components may be either greater than or less than the specified range. For example, a **tm_hour** of –1 means 1 hour before midnight, **tm_mday** of 0 means the day preceding the current month, and **tm_mon** of –2 means 2 months before January of **tm_year**.

If **tm_isdst** is positive, the original values are assumed to be in the alternate timezone. If it turns out that the alternate timezone is not valid for the computed calendar time, then the components are adjusted to the main timezone. Likewise, if **tm_isdst** is zero, the original values are assumed to be in the main timezone and are converted to the alternate timezone if the main timezone is not valid. If **tm_isdst** is negative, the correct timezone is determined and the components are not adjusted.

Local timezone information is used as if **mktime( )** had called **tzset( )** (see **ctime**(3C)).

mktime( ) returns the specified calendar time. If the calendar time cannot be represented, the function returns the value (**time_t**)–1.

EXAMPLES

What day of the week is July 4, 2001?

```
#include <stdio.h>
#include <time.h>

static char *const wday[ ] = {
        "Sunday", "Monday", "Tuesday", "Wednesday",
        "Thursday", "Friday", "Saturday", "-unknown-"
};
struct tm time_str;
/* ...*/
time_str.tm_year        = 2001 - 1900;
time_str.tm_mon         = 7 - 1;
time_str.tm_mday        = 4;
time_str.tm_hour        = 0;
time_str.tm_min         = 0;
time_str.tm_sec         = 1;
time_str.tm_isdst       = –1;
if (mktime(&time_str)== –1)
   time_str.tm_wday=7;
printf("%s\n", wday[time_str.tm_wday]);
```

SEE ALSO

ctime(3C), **getenv**(3C), **TIMEZONE**(4)

NOTES

tm_year of the **tm** structure must be for year 1970 or later. Calendar times before 00:00:00 UTC, January 1, 1970 or after 03:14:07 UTC, January 19, 2038 cannot be represented.

NAME

mlock, munlock – lock (or unlock) pages in memory

SYNOPSIS

```
#include <sys/types.h>
int mlock(caddr_t addr, size_t len);
int munlock(caddr_t addr, size_t len);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The function **mlock( )** uses the mappings established for the address range [*addr, addr + len*) to identify pages to be locked in memory. If the page identified by a mapping changes, such as occurs when a copy of a writable **MAP_PRIVATE** page is made upon the first store, the lock will be transferred to the newly copied private page.

munlock( ) removes locks established with **mlock( )**.

A given page may be locked multiple times by executing an **mlock( )** through different mappings. That is, if two different processes lock the same page, then the page will remain locked until both processes remove their locks. However, within a given mapping, page locks do not nest – multiple **mlock( )** operations on the same address in the same process will all be removed with a single **munlock( )**. Of course, a page locked in one process and mapped in another (or visible through a different mapping in the locking process) is still locked in memory. This fact can be used to create applications that do nothing other than lock important data in memory, thereby avoiding page I/O faults on references from other processes in the system.

If the mapping through which an **mlock( )** has been performed is removed, an **munlock( )** is implicitly performed. An **munlock( )** is also performed implicitly when a page is deleted through file removal or truncation.

Locks established with **mlock( )** are not inherited by a child process after a **fork( )** and are not nested.

Because of the impact on system resources, the use of **mlock( )** and **munlock( )** is restricted to the super-user.

Attempts to **mlock( )** more memory than a system-specific limit will fail.

RETURN VALUES

Upon successful completion, the functions **mlock( )** and **munlock( )** return 0; otherwise, they return –1 and set **errno** to indicate the error.

ERRORS

EAGAIN	**mlock( )** only. Some or all of the memory identified by the range [*addr, addr + len*) could not be locked because of insufficient system resources.
EINVAL	*addr* is not a multiple of the page size as returned by **sysconf**(3C).
ENOMEM	Addresses in the range [*addr, addr + len*) are invalid for the address space of a process, or specify one or more pages which are not mapped.
EPERM	The process's effective user ID is not super-user.

SEE ALSO **fork**(2), **memcntl**(2), **mmap**(2), **plock**(3C), **mlockall**(3C), **sysconf**(3C)

NOTES **mlock** and **munlock** require super-user privileges.

NAME

mlockall, munlockall – lock or unlock address space

SYNOPSIS

```
#include <sys/mman.h>
int mlockall(int flags);
int munlockall(void);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The function **mlockall( )** causes all pages mapped by an address space to be locked in memory.

The value of *flags* determines whether the pages to be locked are those currently mapped by the address space, those that will be mapped in the future, or both:

MCL_CURRENT	Lock current mappings
MCL_FUTURE	Lock future mappings

If **MCL_FUTURE** is specified to **mlockall( )**, then as mappings are added to the address space (or existing mappings are replaced) they will also be locked, provided sufficient memory is available.

Mappings locked via **mlockall( )** with any option may be explicitly unlocked with a **munlock( )** call.

The function **munlockall( )** removes address space locks and locks on mappings in the address space.

All conditions and constraints on the use of locked memory as exist for **mlock( )** apply to **mlockall( )**.

Locks established with **mlockall( )** are not inherited by a child process after a **fork( )** and are not nested.

RETURN VALUES

Upon successful completion, the functions **mlockall( )** and **munlockall( )** return 0; otherwise, they return –1 and set **errno** to indicate the error.

ERRORS

EAGAIN	**mlockall( )** only. Some or all of the memory in the address space could not be locked due to sufficient resources.
EINVAL	*flags* contains values other than **MCL_CURRENT** and **MCL_FUTURE**.
EPERM	The process's effective user ID is not super-user.

SEE ALSO

fork(2), **memcntl**(2), **mmap**(2), **plock**(3C), **mlock**(3C), **sysconf**(3C)

NOTES

mlockall( ) and **munlockall( )** require super-user privileges.

NAME

monitor – prepare process execution profile

SYNOPSIS

```
#include <mon.h>
void monitor(int (*lowpc)(), int (*highpc)(), WORD *buffer, size_t bufsize,
        size_t nfunc);
```

MT-LEVEL

Safe

DESCRIPTION

monitor( ) is an interface to **profil( )**, and is called automatically with default parameters by any program created by **cc**(1B) **–p.** Except to establish further control over profiling activity, it is not necessary to explicitly call **monitor( )**.

When used, **monitor( )** is called at least at the beginning and the end of a program. The first call to **monitor( )** initiates the recording of two different kinds of execution-profile information: execution-time distribution and function call count. Execution-time distribution data is generated by **profil( )** and the function call counts are generated by code supplied to the object file (or files) by **cc**(1B) **–p.** Both types of information are collected as a program executes. The last call to **monitor( )** writes this collected data to the output file **mon.out**.

The name of the file written by **monitor( )** is controlled by the environment variable **PROFDIR**. If **PROFDIR** does not exist, the file **mon.out** is created in the current directory. If **PROFDIR** exists but has no value, **monitor( )** does no profiling and creates no output file. If **PROFDIR** is *dirname*, and **monitor( )** is called automatically by compilation with **cc –p**, the file created is *dirname/pid.progname* where *progname* is the name of the program.

lowpc and *highpc* are the beginning and ending addresses of the region to be profiled.

buffer is the address of a user-supplied array of **WORD** (**WORD** is defined in the header **<mon.h>**). *buffer* is used by **monitor( )** to store the histogram generated by **profil( )** and the call counts.

bufsize identifies the number of array elements in *buffer*.

nfunc is the number of call count cells that have been reserved in *buffer*. Additional call count cells will be allocated automatically as they are needed.

bufsize should be computed using the following formula:

```
size_of_buffer =
        sizeof(struct hdr) +
        nfunc * sizeof(struct cnt) +
        ((highpc-lowpc) / BARSIZE) * sizeof(WORD) +
        sizeof(WORD) – 1 ;
bufsize = (size_of_buffer / sizeof(WORD)) ;
```

where:

lowpc, *highpc*, *nfunc* are the same as the arguments to **monitor( )**;

BARSIZE is the number of program bytes that correspond to each histogram bar, or cell, of the **profil( )** buffer;

the **hdr** and **cnt** structures and the type **WORD** are defined in the header **<mon.h>**.

The default call to **monitor( )** is shown below:

monitor (&eprol, &etext, wbuf, wbufsz, 600);

where:

eprol is the beginning of the user's program when linked with **cc –p** (see **end**(3C));

etext is the end of the user's program (see **end**(3C));

wbuf is an array of **WORD** with *wbufsz* elements;

wbufsz is computed using the *bufsize* formula shown above with *BARSIZE* of 8;

600 is the number of call count cells that have been reserved in *buffer*.

These parameter settings establish the computation of an execution-time distribution histogram that uses **profil( )** for the entire program, initially reserves room for 600 call count cells in *buffer*, and provides for enough histogram cells to generate significant distribution-measurement results. For more information on the effects of *bufsize* on execution-distribution measurements, see **profil**(2).

EXAMPLES

To stop execution monitoring and write the results to a file, use the following:

monitor((int (*)( ))0, (int (*)( ))0, (WORD *)0, 0, 0);

Use **prof** to examine the results.

FILES

mon.out

SEE ALSO

cc(1B), **profil**(2), **end**(3C), **prof**(5)

NOTE

Additional calls to **monitor( )** after **main( )** has been called and before **exit( )** has been called will add to the function-call count capacity, but such calls will also replace and restart the **profil( )** histogram computation.

NAME

mq_close – close a message queue

SYNOPSIS

cc [*flag* ...] *file* ... –**lposix4** [*library* ...]

#include <mqueue.h>

int mq_close(mqd_t *mqdes***);**

MT-LEVEL

MT-Safe

DESCRIPTION

mq_close( ) removes the association between the message queue descriptor, *mqdes*, and its message queue.

If the process (or thread) has registered a notification request to the message queue via this *mqdes*, this registration is removed and the message queue is available for another process to attach for notification.

RETURN VALUES

Upon successful completion, **mq_close( )** returns **0**; otherwise, the function returns **-1** and sets **errno** to indicate the error condition.

ERRORS

EBADF *mqdes* is an invalid message queue descriptor.

ENOSYS **sem_open( )** is not supported by this implementation.

SEE ALSO

mq_notify(3R), **mq_open**(3R), **mq_unlink**(3R)

BUGS

In Solaris 2.5, these functions always return –**1** and set **errno** to **ENOSYS,** because this release does not support the Message Passing option. It is our intention to provide support for these interfaces in future releases.

NAME

mq_notify – notify process (or thread) that a message is available on a queue

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <mqueue.h>

int mq_notify(mqd_t mqdes, const struct sigevent *notification);

struct sigevent {
        int             sigev_notify    /* notification type */
        int             sigev_signo;    /* signal number */
        union sigval    sigev_value;    /* signal value */
};

union sigval {
        int             sival_int;      /* integer value */
        void            *sival_ptr;     /* pointer value */
};
```

MT-LEVEL

MT-Safe

DESCRIPTION

mq_notify( ) provides an asynchronous mechanism for processes to receive notice that messages are available in a message queue, rather than synchronously blocking (waiting) in **mq_receive**(3R).

If *notification* is not **NULL**, this function registers the calling process to be notified of message arrival at an empty message queue associated with the message queue descriptor, *mqdes*. The notification specified by *notification* will be sent to the process when the message queue becomes non-empty. At any time, only one process may be registered for notification by a specific message queue. Also, if the calling process or any other process has already registered for notification of message arrival at the specified message queue, subsequent attempts to register for that message queue will fail.

notification points to a structure that defines both the signal to be generated and how the calling process will be notified upon I/O completion. If *notification*->**sigev_notify** is **SIGEV_NONE**, then no signal will be posted upon I/O completion, but the error status and the return status for the operation will be set appropriately. If *notification*->**sigev_notify** is **SIGEV_SIGNAL**, then the signal specified in *notification*->**sigev_signo** will be sent to the process. If the **SA_SIGINFO** flag is set for that signal number, then the signal will be queued to the process and the value specified in *notification*->**sigev_value** will be the **si_value** component of the generated signal (see **siginfo**(5)).

If *notification* is **NULL** and the process is currently registered for notification by the specified message queue, the existing registration is removed. The message queue is then available for future registration.

When the notification is sent to the registered process, its registration is removed. The message queue is then be available for registration.

If a process has registered for notification of message arrival at a message queue and some processes is blocked in **mq_receive**(3R) waiting to receive a message when a message arrives at the queue, the arriving message will be received by the appropriate **mq_receive**(3R), and no notification will be sent to the registered process. The resulting behavior is as if the message queue remains empty, and this notification will not be sent until the next arrival of a message at this queue.

Any notification registration is removed if the calling process either closes the message queue or exits.

RETURN VALUES

Upon successful completion, **mq_notify( )** returns **0**; otherwise, it returns a value of **-1** and sets **errno** to indicate the error condition.

ERRORS

EBADF	*mqdes* is not a valid message queue descriptor.
EBUSY	A process is already registered for notification by the message queue.
ENOSYS	**mq_notify( )** is not supported by this implementation.

SEE ALSO

mq_close(3R), **mq_open**(3R), **mq_receive**(3R), **mq_send**(3R), **siginfo**(5)

BUGS

In Solaris 2.5, these functions always return –**1** and set **errno** to **ENOSYS,** because this release does not support the Message Passing option. It is our intention to provide support for these interfaces in future releases.

NAME

mq_open – open a message queue

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <mqueue.h>

mqd_t mq_open(const char *name, int oflag,
        /* unsigned long mode, mq_attr attr */ ... );

struct      mq_attr {
    long    mq_flags;       /* message queue flags */
    long    mq_maxmsg;      /* maximum number of messages */
    long    mq_msgsize;     /* maximum message size */
    long    mq_curmsgs;     /* number of messages currently queued */
    ...
};
```

MT-LEVEL

MT-Safe

DESCRIPTION

mq_open( ) establishes a connection to a named message queue, *name,* returning the address of the message queue descriptor to the caller for subsequent calls to **mq_send**(3R) or **mq_receive**(3R). The message queue once opened remains usable by this process until the message queue is closed by a successful call to **mq_close**(3R), **exit**(2), or **exec**(2).

name points to a string naming a message queue. The *name* argument must conform to the construction rules for a path-name. If *name* is not the name of an existing message queue and its creation is not requested, **mq_open( )** fails and returns an error.

oflag requests the desired receive and / or send access to the message queue. The requested access permission to receive messages or send messages is granted if the calling process would be granted read or write access, respectively, to a file with the equivalent permissions.

The value of *oflag* is the bitwise inclusive OR of values from the following list. Applications must specify exactly one of the first three values (access modes) below in the value of *oflag :*

O_RDONLY Open the message queue for receiving messages. The process can use the returned message queue descriptor with **mq_receive**(3R), but not **mq_send**(3R). A message queue may be open multiple times in the same or different processes for receiving messages.

O_WRONLY Open the queue for sending messages. The process can use the returned message queue descriptor with **mq_send**(3R) but not **mq_receive**(3R). A message queue may be open multiple times in the same or different processes for sending messages.

O_RDWR Open the queue for both receiving and sending messages. The process can use any of the functions allowed for **O_RDONLY** and **O_WRONLY.** A message queue may be open multiple times in the same or different processes for sending messages.

Any combination of the remaining flags may additionally be specified in the value of *oflag*:

O_CREAT This option is used to create a message queue, and it requires two additional arguments: *mode*, which is of type **mode_t**, and *attr*, which is pointer to a **mq_attr** structure. If the pathname, *name*, has already been used to create a message queue that still exists, then this flag has no effect, unless combined with **O_EXCL** (see below). Otherwise, a message queue is created without any messages in it.

The message queue's user ID is set to the process's effective user ID, and the message queue's group ID is set to the process's effective group ID. The message queue's permission bits will be set to the value of *mode*, and modified by clearing all bits set in the file mode creation mask of the process (see **umask**(2)). "AND-NOT") those already set in the file mode creation mask of the process.

If *attr* is **NULL,** the message queue is created with the default message queue attributes, (**mq_maxmsg** = 128 and **mq_maxsize** = 1024). If *attr* is non-**NULL**, the message queue **mq_maxmsg** and **mq_msgsize** attributes are set to the values of the corresponding members in the **mq_attr** structure referred to by *attr*.

O_EXCL If both **O_EXCL** and **O_CREAT** are set, **mq_open( )** will fail if the message queue *name* exists. The check for the existence of the message queue and the creation of the message queue if it does not exist are atomic with respect to other processes executing **mq_open( )** naming the same *name* with both **O_EXCL** and **O_CREAT** set.

O_NONBLOCK The setting of this flag is associated with the open message queue descriptor and determines whether a calling **mq_send**(3R) waits for message buffer space or a calling **mq_receive**(3R) waits for messages that are not currently available; or whether the calling function fails, thereby setting **errno** to **EAGAIN.**

RETURN VALUES

Upon successful completion, **mq_open( )** returns a message queue descriptor; otherwise the function returns **(mqd_t)(-1)** and sets **errno** to indicate the error condition.

ERRORS

EACCESS The message queue exists and the permissions specified by *oflag* are denied, or the message queue does not exist and permission to create the message queue is denied.

EEXIST **O_CREAT** and **O_EXCL** are set and the named message queue already exists.

EINTR The **mq_open( )** operation was interrupted by a signal.

EINVAL *name* is not a valid name.

O_CREAT was specified in *oflag*, the value of *attr* is not **NULL,** and either **mq_maxmsg** or **mq_msgsize** was less than or equal to zero.

EMFILE The number of open message queue descriptors in this process exceeds **{MQ_OPEN_MAX}.**

The number of open file descriptors in this process exceeds **{OPEN_MAX}.**

ENAMETOOLONG
The length of the *name* string exceeds **PATH_MAX,** or a pathname component is longer than **NAME_MAX** while **_POSIX_NO_TRUNC** is in effect.

ENFILE The system file table is full

ENOENT **O_CREAT** is not set and the named message queue, *name*, does not exist.

ENOSPC There is insufficient space for the creation of the new message queue.

ENOSYS **mq_open( )** is not supported by this implementation.

SEE ALSO

exec(2), **exit**(2), **umask**(2), **sysconf**(3C), **mq_close**(3R), **mq_receive**(3R), **mq_send**(3R), **mq_setattr**(3R), **mq_unlink**(3R),

NOTES

In Solaris, message queues are based on shared memory. Although permissions to send and receive messages are checked by the **mq_receive**() and **mq_send**() interfaces, any application which can open the message queue can directly access the shared memory to examine and manipulate messages in the queue. Thus message queues should not be considered secure.

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Message Passing option. It is our intention to provide support for these interfaces in future releases.

NAME

mq_receive – receive a message from a message queue

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <mqueue.h>

ssize_t mq_receive(mqd_t mqdes, char *msg_ptr, size_t msg_len,
        unsigned int msg_prio);

struct mq_attr {
        long    mq_flags;       /* message queue flags */
        long    mq_maxmsg;      /* maximum number of messages */
        long    mq_msgsize;     /* maximum message size */
        long    mq_curmsgs;     /* number of messages currently queued */
        ...
};
```

MT-LEVEL

MT-Safe

DESCRIPTION

mq_receive( ) is used to receive the oldest of the highest priority message(s) from the message queue specified by *mqdes*. If the size of the buffer in bytes, specified by *msg_len*, is less than the **mq_msgsize** attribute of the message queue, the function fails and returns an error. Otherwise, the selected message is removed from the queue and copied to the buffer pointed to by *msg_ptr*.

If *msg_prio* is not **NULL**, the priority of the selected message is stored in the location referenced by *msg_prio*.

If the specified message queue is empty and **O_NONBLOCK** is not set in the message queue description associated with *mqdes*, (see **mq_open**(3R) and **mq_setattr**(3R)), **mq_receive( )** blocks, waiting until a message is enqueued on the message queue, or until **mq_receive( )** is interrupted by a signal. If more than one process (or thread) is waiting to receive a message when a message arrives at an empty queue, then the process of highest priority that has been waiting the longest is selected to receive the message. If the specified message queue is empty and **O_NONBLOCK** is set in the message queue description associated with *mqdes*, no message is removed from the queue, and **mq_receive( )** returns an error.

RETURN VALUES

Upon successful completion, **mq_receive( )** returns the length of the selected message in bytes and the message will have been removed from the queue. Otherwise, no message is removed from the queue, the function returns a value of **-1**, and sets **errno** to indicate the error condition.

ERRORS

EAGAIN	**O_NONBLOCK** was set in the message description associated with *mqdes*, and the specified message queue is empty.
EBADF	*mqdes* is not a valid message queue descriptor open for reading.
EMSGSIZE	*msg_len* is less than the message size attribute of the message queue.
EINTR	**mq_receive( )** operation was interrupted by a signal.

ENOSYS **mq_receive( )** is not supported by this implementation.

SEE ALSO **mq_open**(3R), **mq_send**(3R), **mq_setattr**(3R)

BUGS In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Message Passing option. It is our intention to provide support for these interfaces in future releases.

NAME

mq_send – send a message to a message queue

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <mqueue.h>

int mq_send(mqd_t mqdes, const char *msg_ptr,
        size_t msg_len, unsigned int msg_prio);

struct mq_attr {
    long    mq_flags;       /* message queue flags */
    long    mq_maxmsg;      /* maximum number of messages */
    long    mq_msgsize;     /* maximum message size */
    long    mq_curmsgs;     /* number of messages currently queued */
    ...
};
```

MT-LEVEL

MT-Safe

DESCRIPTION

mq_send( ) adds the message pointed to by *msg_ptr* to the message queue specified by *mqdes*. *msg_len* specifies the length of the message in bytes pointed to by *msg_ptr*. The value of *msg_len* must be less than or equal to the **mq_msgsize** attribute of the message queue, or **mq_send( )** will fail.

If the specified message queue is not full, **mq_send( )** behaves as if the message is inserted into the message queue at the position indicated by *msg_prio*. A message with a larger numeric value of *msg_prio* is inserted before messages with lower values of *msg_prio*. A message is inserted after other messages in the queue, if any, with equal *msg_prio* priority. The value of *msg_prio* must be greater than **0**, and less than or equal to **{MQ_PRIO_MAX}**.

If the specified message queue is full and if **O_NONBLOCK** is not set in the message queue description associated with *mqdes* (see **mq_open**(3R) and **mq_setattr**(3R)), **mq_send( )** blocks, waiting until space becomes available to enqueue the message, or until **mq_send( )** is interrupted by a signal. If more than one process (or thread) is waiting to send when space becomes available in the message queue, then the process of the highest priority which has been waiting the longest is unblocked to send its message. If the specified message queue is full and **O_NONBLOCK** is set in the message queue description associated with *mqdes*, the message is not queued, and **mq_send( )** returns an error.

RETURN VALUES

Upon successful completion, **mq_send( )** returns a value of **0**; otherwise, no message is enqueued, the function returns **-1**, and sets **errno** to indicate the error condition.

ERRORS

EAGAIN	The **O_NONBLOCK** flag is set in the message queue description associated with *mqdes*, and the specified message queue is full.
EBADF	*mqdes* is not a valid message queue descriptor open for writing.
EINTR	A signal interrupted the call to **mq_send( )**

EMSGSIZE The specified message length, *msg_len*, exceeds the message size attribute of the message queue.

ENOSYS **mq_send( )** is not supported by this implementation.

SEE ALSO **mq_open**(3R), **mq_receive**(3R), **mq_setattr**(3R), **sysconf**(3C)

BUGS In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Message Passing option. It is our intention to provide support for these interfaces in future releases.

NAME

mq_setattr, mq_getattr – set / get message queue attributes

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <mqueue.h>

int mq_setattr(mqd_t mqdes, const struct mq_attr *mqstat,
    struct mq_attr* omqstat);

int mq_getattr(mqd_t mqdes, struct mq_attr *mqstat);

struct mq_attr {
    long    mq_flags;       /* message queue flags */
    long    mq_maxmsg;      /* maximum number of messages */
    long    mq_msgsize;     /* maximum message size */
    long    mq_curmsgs;     /* number of messages currently queued */
    ...
};
```

MT-LEVEL

MT-Safe

DESCRIPTION

mq_setattr( ) is used to set attributes associated with the message queue specified by *mqdes*.

The message queue attributes corresponding to the following members defined in the *mq_attr* structure are set to the specified values upon successful completion of **mq_setattr( )**:

mq_flags The value of this member is either **0** or **O_NONBLOCK**.

The values of **mq_maxmsg**, **mq_msgsize**, and **mq_curmsgs** are ignored by **mq_setattr( ).**

If *omqstat* is non-NULL, **mq_setattr( )** stores, in the location referenced by *omqstat*, the previous message queue attributes and the current queue status. These values are the same as would be returned by a call to **mq_getattr( )** at that point. **mq_getattr( )** is used to get status information and attributes associated with the message queue specified in *mqdes*. Upon return, the **mq_flags** member of the **mq_attr** structure referenced by *mqstat* has the value that was set when the message queue was created but also with modifications made by subsequent **mq_setattr( )** calls.

The following attributes were set at message queue creation:

mq_maxmsg

mq_msgsize

Upon return, the **mq_curmsgs** (the number of messages currently on the queue) member of the **mq_attr** structure referenced by *mqstat* is set according to the current state of the message queue.

RETURN VALUES

Upon successful completion, these function(s) return **0**; otherwise, they return **-1**, and set **errno** to indicate the error condition.

mq_setattr( ), if successful, also changes the attributes of the message queue as specified.

ERRORS

EBADF *mqdes* is not a valid message queue descriptor.

ENOSYS **mq_setattr( )** and **mq_getattr( )** are not supported by this implementation.

SEE ALSO

mq_open(3R), **mq_receive**(3R), **mq_send**(3R)

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Message Passing option. It is our intention to provide support for these interfaces in future releases.

NAME

mq_unlink – remove a message queue

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

#include <mqueue.h>

int mq_unlink(const char ∗*name***);**

MT-LEVEL

MT-Safe

DESCRIPTION

mq_unlink( ) removes the message queue named by *name*. After a successful call to **mq_unlink( )** with *name*, a call to **mq_open**(3R) with the same *name* will fail if the flag **O_CREAT** is not set in *flags*. If one or more processes have the message queue open when **mq_unlink( )** is called, destruction of the message queue is postponed until all references to the message queue have been closed. Calls to **mq_open**(3R) to re-create the message queue may fail until the message queue is actually removed. However, **mq_unlink( )** does not block (wait) until all references have been closed; it returns immediately.

RETURN VALUES

Upon successful completion, **mq_unlink( )** returns a value of **0**; otherwise, the named message queue is not changed by this function call, the function returns a value of **-1** and sets **errno** to indicate the error condition.

ERRORS

EACCESS Permission is denied to unlink the named message queue.

ENAMETOOLONG
The length of the *name* string exceeds **{PATH_MAX}**, or a pathname component is longer than **{NAME_MAX}** while **_POSIX_NO_TRUNC** is in effect.

ENOENT The named message queue, *name*, does not exist.

ENOSYS **mq_unlink( )** is not supported by this implementation.

SEE ALSO

mq_close(3R), **mq_open**(3R)

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Message Passing option. It is our intention to provide support for these interfaces in future releases.

NAME

msync – synchronize memory with physical storage

SYNOPSIS

```
#include <sys/types.h>
#include <sys/mman.h>

int msync(caddr_t addr, size_t len, int flags);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The function **msync( )** writes all modified copies of pages over the range [*addr, addr + len*) to their backing storage locations. **msync( )** optionally invalidates any copies so that further references to the pages will be obtained by the system from their backing storage locations. The backing storage for a modified **MAP_SHARED** mapping is the file the page is mapped to; the backing storage for a modified **MAP_PRIVATE** mapping is its swap area.

flags is a bit pattern built from the following values:

MS_ASYNC	perform asynchronous writes
MS_SYNC	perform synchronous writes
MS_INVALIDATE	invalidate mappings

If **MS_ASYNC** is set, **msync( )** returns immediately once all write operations are scheduled; if **MS_SYNC** is set, **msync( )** does not return until all write operations are completed.

MS_INVALIDATE invalidates all cached copies of data in memory, so that further references to the pages will be obtained by the system from their backing storage locations.

RETURN VALUES

Upon successful completion, the function **msync( )** returns 0; otherwise, it returns –1 and sets **errno** to indicate the error.

ERRORS

EBUSY	Some or all of the addresses in the range [*addr, addr + len*) are locked and **MC_SYNC** with the **MS_INVALIDATE** option is specified.
EINVAL	*addr* is not a multiple of the page size as returned by **sysconf**(3C).
	flags is not some combination of **MS_ASYNC** and **MS_INVALIDATE.**
EIO	An I/O error occurred while reading from or writing to the file system.
ENOMEM	Addresses in the range [*addr, addr + len*) are outside the valid range for the address space of a process, or specify one or more pages that are not mapped.
EPERM	**MS_INVALIDATE** was specified and one or more of the pages is locked in memory.

SEE ALSO

mmap(2), **memcntl**(2), **sysconf**(3C)

NOTES

msync( ) should be used by programs that require a memory object to be in a known state, such as in building transaction facilities.

NAME

mutex, pthread_mutex_init, pthread_mutex_lock, pthread_mutex_trylock, pthread_mutex_unlock, pthread_mutex_destroy, mutex_init, mutex_lock, mutex_trylock, mutex_unlock, mutex_destroy – mutual exclusion locks

SYNOPSIS

POSIX

cc [*flag* ...] *file* ... –**lpthread** [*library* ...]

```
#include <pthread.h>
int pthread_mutex_init(pthread_mutex_t *mp, const pthread_mutexattr_t *attr);
pthread_mutex_t mutex = PTHREAD_MUTEX_INITIALIZER;
int pthread_mutex_lock(pthread_mutex_t *mp);
int pthread_mutex_trylock(pthread_mutex_t *mp);
int pthread_mutex_unlock(pthread_mutex_t *mp);
int pthread_mutex_destroy(pthread_mutex_t *mp);
```

Solaris

cc [*flag* ...] *file* ... –**lthread** [*library* ...]

```
#include <thread.h>
#include <synch.h>
int mutex_init(mutex_t *mp, int type, void * arg);
int mutex_lock(mutex_t *mp);
int mutex_trylock(mutex_t *mp);
int mutex_unlock(mutex_t *mp);
int mutex_destroy(mutex_t *mp);
```

MT-LEVEL

MT-Safe

DESCRIPTION

Mutual exclusion locks (mutexes) prevent multiple threads from simultaneously executing critical sections of code which access shared data (i.e., mutexes are used to serialize the execution of threads). All mutexes must be global. A successful call for a mutex lock via **pthread_mutex_lock()** or **mutex_lock()** will cause another thread that is also trying to lock the same mutex to block until the owner thread unlocks it via **pthread_mutex_unlock()** or **mutex_unlock().** Threads within the same process or within other processes can share mutexes.

Mutexes can synchronize threads within the same process or in other processes. Mutexes can be used to synchronize threads between processes if the mutexes are allocated in writable memory and shared among the cooperating processes (see **mmap**(2)), and have been initialized for this task.

Initialize

Mutexes are either intra-process or inter-process, depending upon the argument passed implicitly or explicitly to the initialization of that mutex. A statically allocated mutex does not need to be explicitly initialized; by default, a statically allocated mutex is initialized with all zeros and its scope is set to be within the calling process. For POSIX portability of statically allocated mutexes, use the **pthread_mutex_initializer** macro (see below).

For inter-process synchronization, a mutex needs to be allocated in memory shared between these processes. Since the memory for such a mutex must be allocated dynamically, the mutex needs to be explicitly initialized using **mutex_init( )** or **pthread_mutex_init( )** with the appropriate attribute that indicates inter-process use.

POSIX Initialize

POSIX mutexes, threads, and condition variables use attributes objects in the same manner; they are initialized with the configuration of an attributes object (see **pthread_mutexattr_init**(3T)). The **pthread_mutex_init( )** function initializes the mutex referenced by *mp* with attributes specified by *attr*. If *attr* is **NULL**, the default mutex attributes are used, which is the same as passing the address of a default mutex attributes object. Upon initialization, the state of the mutex is initialized and unlocked. If default mutex attributes are used, then only threads created within the same process can operate on the initialized mutex variable.

In POSIX, the attributes of a mutex may be specified via the attribute object created via **pthread_mutexattr_init( )** and modified using the **pthread_mutexattr_*( )** functions. To explicitly specify whether a mutex is or is not shared between processes, it can be initialized with an attribute object modified via **pthread_mutexattr_setpshared**(3T). The second argument to this function can be either of the following:

PTHREAD_PROCESS_PRIVATE The mutex can synchronize threads within this process. The **PTHREAD_PROCESS_PRIVATE** POSIX mutex type for process scope is equivalent to the **USYNC_THREAD** flag to **mutex_init( )** in the Solaris API (see below).

PTHREAD_PROCESS_SHARED The mutex can synchronize threads in this process and other processes. Only one process should initialize the mutex. The **PTHREAD_PROCESS_SHARED** POSIX mutex type for system-wide scope is equivalent to the **USYNC_PROCESS** flag to **mutex_init( )** in the Solaris API (see below).

Initializing mutexes can also be accomplished by allocating in zeroed memory (default), in which case, **PTHREAD_PROCESS_PRIVATE** is assumed. The same mutex must not be simultaneously initialized by multiple threads, nor should a mutex lock be re-initialized while in use by other threads.

If default mutex attributes are used, statically allocated mutexes can be initialized by the macro **PTHREAD_MUTEX_INITIALIZER.** The effect is the same as a dynamic initialization by a call to **pthread_mutex_init( )** with parameter *attr* specified as **NULL ,** except error checks are not performed.

Default mutex initialization (intra-process):

```
pthread_mutex_t mp;
pthread_mutexattr_t mattr;

pthread_mutex_init(&mp, NULL);
```

OR

```
pthread_mutexattr_init(&mattr);
pthread_mutex_init(&mp, &mattr);
```

OR

```
pthread_mutexattr_setpshared(&mattr, PTHREAD_PROCESS_PRIVATE);
pthread_mutex_init(&mp, &mattr);
```

OR

```
pthread_mutex_t mp = PTHREAD_MUTEX_INITIALIZER;
```

OR

```
pthread_mutex_t mp;
mp = calloc (1, sizeof (pthread_mutex_t));
```

Customized mutex initialization (inter-process):

```
pthread_mutexattr_init(&mattr);
pthread_mutexattr_setpshared(&mattr, PTHREAD_PROCESS_SHARED);
pthread_mutex_init(&mp, &mattr);
```

Solaris Initialize

The equivalent Solaris API used to initialize a mutex so that it has several different types of behavior is the *type* argument passed to **mutex_init( )**. No current type uses *arg* although a future type may specify additional behavior parameters via *arg*. *type* may be one of the following:

USYNC_THREAD The mutex can synchronize threads only in this process. *arg* is ignored. The **USYNC_THREAD** Solaris mutex type for process scope is equivalent to the POSIX mutex attribute setting **PTHREAD_PROCESS_PRIVATE.**

USYNC_PROCESS The mutex can synchronize threads in this process and other processes. Only one process should initialize the mutex. *arg* is ignored. The **USYNC_PROCESS** Solaris mutex type for process scope is equivalent to the POSIX mutex attribute setting **PTHREAD_PROCESS_SHARED.**

Initializing mutexes can also be accomplished by allocating in zeroed memory (default), in which case, a *type* of **USYNC_THREAD** is assumed. The same mutex must not be simultaneously initialized by multiple threads. A mutex lock must not be re-initialized while in use by other threads.

If default mutex attributes are used, the macro **DEFAULTMUTEX** can be used to initialize mutexes that are statically allocated.

Default mutex initialization (intra-process):

```
mutex_t mp;

mutex_init(&mp, NULL, NULL);
```

OR

```
mutex_init(&mp, USYNC_THREAD, NULL);
```

OR

mutex_t mp = DEFAULTMUTEX;

OR

mutex_t mp;

mp = calloc(1, sizeof (mutex_t));

OR

mutex_t mp;

mp = malloc(sizeof (mutex_t));

memset(mp, 0, sizeof (mutex_t));

Customized mutex initialization (inter-process):

mutex_init(&mp, USYNC_PROCESS, NULL);

Lock and Unlock

A critical section of code is enclosed by a the call to lock the mutex and the call to unlock the mutex to protect it from simultaneous access by multiple threads. Only one thread at a time may possess mutually exclusive access to the critical section of code that is enclosed by the mutex-locking call and the mutex-unlocking call, whether the mutex's scope is intra-process or inter-process. A thread calling to lock the mutex either gets exclusive access to the code starting from the successful locking until its call to unlock the mutex, or it waits until the mutex is unlocked by the thread that locked it.

Mutexes have ownership, unlike semaphores. Although any thread, within the scope of a mutex, can get an unlocked mutex and lock access to the same critical section of code, only the thread that locked a mutex can unlock it.

If a thread waiting for a mutex receives a signal, upon return from the signal handler, the thread resumes waiting for the mutex as if there was no interrupt. A mutex protects code, not data; therefore, strongly bind a mutex with the data by putting both within the same structure, or at least within the same procedure.

POSIX/Solaris Locking

A call to **pthread_mutex_lock( )** or **mutex_lock( )** locks the mutex object referenced by *mp*. If the mutex is already locked, the calling thread blocks until the mutex is freed; this will return with the mutex object referenced by *mp* in the locked state with the calling thread as its owner. If the current owner of a mutex tries to relock the mutex, it will result in deadlock.

pthread_mutex_trylock( ) and **mutex_trylock( )** is the same as **pthread_mutex_lock( )** and **mutex_lock( ) ,** respectively, except that if the mutex object referenced by *mp* is locked (by any thread, including the current thread), the call returns immediately with an error.

pthread_mutex_unlock( ) or **mutex_unlock( )** are called by the owner of the mutex object referenced by *mp* to release it. The mutex must be locked and the calling thread must be the one that last locked the mutex (the owner). If there are threads blocked on the mutex

object referenced by *mp* when **pthread_mutex_unlock( )** is called, the *mp* is freed, and the scheduling policy will determine which thread gets the mutex. If the calling thread is not the owner of the lock, no error status is returned, and the behavior of the program is undefined.

Destroy

Either **pthread_mutex_destroy( )** or **mutex_destroy( )** destroys the mutex object referenced by *mp*; the mutex object becomes uninitialized. The space used by the destroyed mutex variable is not freed. It needs to be explicitly reclaimed.

RETURN VALUES

If successful, all of these functions return **0**; otherwise, an error number is returned.

pthread_mutex_trylock( ) or **mutex_trylock( )** returns **0** if a lock on the mutex object referenced by *mp* is obtained; otherwise, an error number is returned.

ERRORS

These functions fail and return the corresponding value if any of the following conditions are detected:

EFAULT *mp* or *attr* points to an illegal address.

pthread_mutex_init( ) or **mutex_init( )** fails and returns the corresponding value if any of the following conditions are detected:

EINVAL The value specified by *mp* or *attr* is invalid.

pthread_mutex_trylock( ) or **mutex_trylock( )** fails and returns the corresponding value if any of the following conditions occur:

EBUSY The mutex pointed to by *mp* was already locked.

SEE ALSO

mmap(2), **pthread_create**(3T), **pthread_mutexattr_init**(3T)

EXAMPLES

Single Gate

The following example uses one global mutex as a gate-keeper to permit each thread exclusive sequential access to the code within the user-defined function "change_global_data." This type of synchronization will protect the state of shared data, but it also prohibits parallelism.

```
/* cc thisfile.c -lthread */

#define _REENTRANT
#include <stdio.h>
#include <thread.h>
#define NUM_THREADS 12
void *change_global_data(void *);    /*  for thr_create()  */
main(int argc,char * argv[])      {
        int i=0;
        for (i=0; i< NUM_THREADS; i++)        {
                thr_create(NULL, 0, change_global_data, NULL, 0, NULL);
        }
        while ((thr_join(NULL, NULL, NULL) == 0));
}
```

```
void * change_global_data(void *null)  {
        static mutex_t   Global_mutex;
        static int       Global_data = 0;
        mutex_lock(&Global_mutex);
        Global_data++;
        sleep(1);
        printf("%d is global data\n",Global_data);
        mutex_unlock(&Global_mutex);
         return NULL;
}
```

Multiple Instruction Single Data

The previous example, the mutex, the code it owns, and the data it protects was enclosed in one function. The next example uses C++ features to accommodate many functions that use just one mutex to protect one data:

```
/* CC thisfile.c -lthread   use C++ to compile*/
#define _REENTRANT
#include <stdlib.h>
#include <stdio.h>
#include <thread.h>
#include <errno.h>
#include <iostream.h>
#define NUM_THREADS 16
void *change_global_data(void *);    /*  for thr_create()   */

class Mutected {
        private:
                static mutex_t   Global_mutex;
                static int       Global_data;
        public:
                static int       add_to_global_data(void);
                static int       subtract_from_global_data(void);
};

int Mutected::Global_data = 0;
mutex_t Mutected::Global_mutex;

int Mutected::add_to_global_data()      {
        mutex_lock(&Global_mutex);
        Global_data++;
        mutex_unlock(&Global_mutex);
        return Global_data;
}


int Mutected::subtract_from_global_data()       {
```

```
        mutex_lock(&Global_mutex);
        Global_data--;
        mutex_unlock(&Global_mutex);
        return Global_data;
}

void
main(int argc,char * argv[])    {
        int i=0;
        for (i=0;i< NUM_THREADS;i++){
                thr_create(NULL,0,change_global_data,NULL,0,NULL);
        }
        while ((thr_join(NULL,NULL,NULL) == 0));
}

void * change_global_data(void *)       {
        static int switcher = 0;
        if ((switcher++ % 3) == 0)  /* one-in-three threads subtracts */
                cout << Mutected::subtract_from_global_data() << endl;
        else
                cout << Mutected::add_to_global_data() << endl;
        return NULL;
}
```

Interprocess Locking

A mutex can protect data that is shared among processes. The mutex would need to be initialized as either **PTHREAD_PROCESS_SHARED,** for POSIX, or **USYNC_PROCESS,** for Solaris threads. (see **mutex_init**(3T) and **pthread_mutexattr_init**(3T)). One process initializes the process-shared mutex and writes it to a file to be mapped into memory by all cooperating processes (see **mmap**(2)). Afterwards, other independent processes can run the same program (whether concurrently or not) and share mutex-protected data.

```
/* cc thisfile.c -lthread */
/* To execute, run the command line "a.out 0 & a.out 1" */

#define _REENTRANT
#include <sys/types.h>
#include <sys/mman.h>
#include <sys/stat.h>
#include <fcntl.h>
#include <stdio.h>
#include <thread.h>
#define INTERPROCESS_FILE "ipc-sharedfile"
#define NUM_ADDTHREADS 12
#define NUM_SUBTRACTTHREADS 10
```

```
#define INCREMENT '0'
#define DECREMENT '1'
typedef struct {
        mutex_t         Interprocess_mutex;
        int             Interprocess_data;
} buffer_t;
buffer_t *buffer;


void *add_interprocess_data(), *subtract_interprocess_data();
void create_shared_memory(), test_argv();
int zeroed[sizeof(buffer_t)];
int ipc_fd, i=0;


void
main(int argc,char * argv[])    {
        test_argv(argv[1]);

        switch (*argv[1]) {
         case INCREMENT:
                create_shared_memory();
                ipc_fd = open(INTERPROCESS_FILE, O_RDWR);
                buffer = (buffer_t *)mmap(NULL, sizeof(buffer_t),
                        PROT_READ | PROT_WRITE, MAP_SHARED, ipc_fd, 0);
                buffer->Interprocess_data = 0;
                mutex_init(&buffer->Interprocess_mutex, USYNC_PROCESS,0);
                for (i=0; i< NUM_ADDTHREADS; i++)
                thr_create(NULL, 0, add_interprocess_data, argv[1],
                                0, NULL);
                break;

         case DECREMENT:
                while((ipc_fd = open(INTERPROCESS_FILE, O_RDWR)) == -1)
                        sleep(1);
                buffer = (buffer_t *)mmap(NULL, sizeof(buffer_t),
                        PROT_READ | PROT_WRITE, MAP_SHARED, ipc_fd, 0);
                for (i=0; i< NUM_SUBTRACTTHREADS; i++)
                 thr_create(NULL, 0, subtract_interprocess_data, argv[1],
                        0, NULL);
                break;
        } /* end switch */

        while ((thr_join(NULL,NULL,NULL) == 0));
} /* end main */
```

```
void *add_interprocess_data(char argv_1[])      {
        mutex_lock(&buffer->Interprocess_mutex);
        buffer->Interprocess_data++;
        sleep(2);
        printf("%d is add-interprocess data, and %c is argv1\n",
              buffer->Interprocess_data, argv_1[0]);
        mutex_unlock(&buffer->Interprocess_mutex);
        return NULL;
}

void *subtract_interprocess_data(char argv_1[]){
        mutex_lock(&buffer->Interprocess_mutex);
        buffer->Interprocess_data--;
        sleep(2);
        printf("%d is subtract-interprocess data, and %c is argv1\n",
              buffer->Interprocess_data, argv_1[0]);
        mutex_unlock(&buffer->Interprocess_mutex);
        return NULL;
}

void create_shared_memory()   {
        int i;
        ipc_fd = creat(INTERPROCESS_FILE, O_CREAT|O_RDWR );
        for (i=0; i<sizeof(buffer_t); i++){
                zeroed[i] = 0;
                write(ipc_fd, &zeroed[i],2);
        }
        close(ipc_fd);
        chmod(INTERPROCESS_FILE, S_IRWXU|S_IRWXG|S_IRWXO);
}

void test_argv(char argv1[])    {
  if (argv1 == NULL)  {
    printf("use 0 as arg1 for initial process\n \
    or use 1 as arg1 for the second process\n");
    exit(NULL);
  }
}
```

In this example, run the commandline

```
a.out 0 & a.out 1
```

NOTES

Currently, the only supported policy is **SCHED_OTHER.** In Solaris, under the **SCHED_OTHER** policy, there is no established order in which threads are unblocked.

In the current implementation of threads, **pthread_mutex_lock( )**, **pthread_mutex_unlock( )**, **mutex_lock( )**, **mutex_unlock( )**, **pthread_mutex_trylock( )**, and **mutex_trylock( )** do not validate the mutex type. Therefore, an uninitialized mutex or a mutex with an invalid type does not return **EINVAL.** Interfaces for mutexes with an invalid type have unspecified behavior.

Uninitialized mutexes which are allocated locally may contain junk data. Such mutexes need to be initialized using **pthread_mutex_init( )** or **mutex_init( )**.

By default, if multiple threads are waiting for a mutex, the order of acquisition is undefined.

NAME

nanosleep – high resolution sleep

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <time.h>

int nanosleep(const struct timespec *rqtp, struct timespec *rmtp);

struct timespec {
        time_t      tv_sec;   /* seconds */
        long        tv_nsec;  /* and nanoseconds */
};
```

MT-LEVEL

MT-Safe

DESCRIPTION

nanosleep( ) suspends the current thread from execution until either the time interval specified by *rqtp* has elapsed, or a signal is delivered to the calling thread and its action is to invoke a signal-catching function or to terminate the thread. The suspension time may be longer than requested because the argument value is rounded up to an integer multiple of the sleep resolution or because of the scheduling of other activity by the system. Except for the case of being interrupted by a signal, the suspension time will not be less than the time specified by *rqtp*, as measured by the system clock, **CLOCK_REALTIME**.

nanosleep( ) will not block nor effect the action of any signal.

RETURN VALUES

If **nanosleep( )** returns because the requested time has elapsed, it returns **0.** Otherwise, if it returns because it has been interrupted by a signal:

it returns **-1** and sets **errno** to indicate the interruption.

If *rmtp* is non-**NULL**, the **timespec** structure referenced by *rmtp* will be updated to contain the remaining amount of time between *rqtp* and the time actually slept.

If any of the following error conditions occur, **nanosleep( )** returns **-1** and sets **errno** to indicate the error condition.

ERRORS

EINTR	**nanosleep( )** was interrupted by a signal.
EINVAL	*rqtp* specified a nanosecond value less than zero or greater than or equal to 1,000,000,000.
ENOSYS	**nanosleep( )** is not supported by this implementation.

SEE ALSO

sleep(3C)

NAME

ndbm, dbm_clearerr, dbm_close, dbm_delete, dbm_error, dbm_fetch, dbm_firstkey, dbm_nextkey, dbm_open, dbm_store – data base subroutines

SYNOPSIS

#include <ndbm.h>

int dbm_clearerr(DBM ∗*db***);**

void dbm_close(DBM ∗*db***);**

int dbm_delete (DBM ∗*db***, datum** *key***);**

int dbm_error(DBM ∗*db***);**

datum dbm_fetch(DBM ∗*db***, datum** *key***);**

datum dbm_firstkey(DBM ∗*db***);**

datum dbm_nextkey(DBM ∗*db***);**

DBM ∗**dbm_open(char** ∗*file***, int** *flags***, int** *mode***);**

int dbm_store(DBM ∗*db***, datum** *key***, datum** *content***, int** *flags***);**

MT-LEVEL

Unsafe

DESCRIPTION

These functions maintain *key* / *content* pairs in a data base. The functions will handle very large (a billion blocks) data base and will access a keyed item in one or two file system accesses. This package replaces the earlier **dbm**(3B) library, which managed only a single data base.

*key*s and *content*s are described by the **datum** typedef. A **datum** specifies a string of *dsize* bytes pointed to by *dptr*. Arbitrary binary data, as well as normal ASCII strings, are allowed. The data base is stored in two files. One file is a directory containing a bit map and has **.dir** as its suffix. The second file contains all data and has **.pag** as its suffix.

Before a data base can be accessed, it must be opened by **dbm_open( )**. This will open and / or create the files *file***.dir** and *file***.pag** depending on the flags parameter (see **open**(2)).

A data base is closed by calling **dbm_close( )**.

Once open, the data stored under a key is accessed by **dbm_fetch( )** and data is placed under a key by **dbm_store( )**. The *flags* field can be either **DBM_INSERT** or **DBM_REPLACE**. **DBM_INSERT** will only insert new entries into the data base and will not change an existing entry with the same key. **DBM_REPLACE** will replace an existing entry if it has the same key. A key (and its associated contents) is deleted by **dbm_delete( )**. A linear pass through all keys in a data base may be made, in an (apparently) random order, by use of **dbm_firstkey( )** and **dbm_nextkey( ). dbm_firstkey( )** will return the first key in the data base. **dbm_nextkey( )** will return the next key in the data base. This code will traverse the data base:

for (key = dbm_firstkey(db); key.dptr != NULL; key = dbm_nextkey(db))

dbm_error( ) returns non-zero when an error has occurred reading or writing the data base. **dbm_clearerr( )** resets the error condition on the named data base.

RETURN VALUES

All functions that return an **int** indicate errors with negative values. A return value of **0** indicates no error. Routines that return a **datum** indicate errors with a **NULL** (0) *dptr*. If **dbm_store( )** is called with a *flags* value of **DBM_INSERT** and finds an existing entry with the same key, it returns **1**.

EXAMPLES

The following example stores and retrieves a phone number, using the name as the key. Note that this example does not include error checking.

```
#include <ndbm.h>
#include <stdio.h>
#include <fcntl.h>

#define NAME        "Bill"
#define PHONE_NO    "123-4567"
#define DB_NAME     "phones"

main()
{
        DBM *db;
        datum name = {NAME, sizeof (NAME)};
        datum put_phone_no = {PHONE_NO, sizeof (PHONE_NO)};
        datum get_phone_no;

        /* Open the database and store the record */
        db = dbm_open(DB_NAME, O_RDWR | O_CREAT, 0660);
        (void) dbm_store(db, name, put_phone_no, DBM_INSERT);

        /* Retrieve the record */
        get_phone_no = dbm_fetch(db, name);
        (void) printf("Name: %s, Phone Number: %s0, name.dptr,
           get_phone_no.dptr);

        /* Close the database */
        dbm_close(db);
        return (0);
}
```

SEE ALSO

ar(1), **cat**(1), **cp**(1), **tar**(1), **open**(2), **dbm**(3B), **netconfig**(4)

NOTES

The **.pag** file will contain holes so that its apparent size may be larger than its actual content. Older versions of the UNIX operating system may create real file blocks for these holes when touched. These files cannot be copied by normal means (**cp**(1), **cat**(1), **tar**(1), **ar**(1)) without filling in the holes.

dptr pointers returned by these subroutines point into static storage that is changed by subsequent calls.

The sum of the sizes of a *key*/*content* pair must not exceed the internal block size (currently 1024 bytes). Moreover all *key*/*content* pairs that hash together must fit on a single block. **dbm_store()** will return an error in the event that a disk block fills with inseparable data.

dbm_delete() does not physically reclaim file space, although it does make it available for reuse.

The order of keys presented by **dbm_firstkey()** and **dbm_nextkey()** depends on a hashing function.

There are no interlocks and no reliable cache flushing; thus concurrent updating and reading is risky.

The database files (*file***.dir** and *file***.pag**) are binary and are architecture-specific (for example, they depend on the architecture's byte order.) These files are not guaranteed to be portable across architectures.

NAME

netdir, netdir_getbyname, netdir_getbyaddr, netdir_free, netdir_options, taddr2uaddr, uaddr2taddr, netdir_perror, netdir_sperror, netdir_mergeaddr – generic transport name-to-address translation

SYNOPSIS

```
#include <netdir.h>

int netdir_getbyname(const struct netconfig *config,
        const struct nd_hostserv *service, struct nd_addrlist **addrs);

int netdir_getbyaddr(const struct netconfig *config,
        struct nd_hostservlist **service, const struct netbuf *netaddr);

void netdir_free(void *ptr, const int struct_type);

int netdir_options(const struct netconfig *config, const int option,
        const int fildes, char *point_to_args);

char *taddr2uaddr(const struct netconfig *config, const struct netbuf *addr);

struct netbuf *uaddr2taddr(const struct netconfig *config, const char *uaddr);

void netdir_perror(char *s);

char *netdir_sperror(void);
```

MT-LEVEL

MT-Safe

DESCRIPTION

These routines provide a generic interface for name-to-address mapping that will work with all transport protocols. This interface provides a generic way for programs to convert transport specific addresses into common structures and back again. The **netconfig** structure, described on the **netconfig**(4) manual page, identifies the transport.

The **netdir_getbyname( )** routine maps the machine name and service name in the **nd_hostserv** structure to a collection of addresses of the type understood by the transport identified in the **netconfig** structure. This routine returns all addresses that are valid for that transport in the **nd_addrlist** structure. The **nd_hostserv** structure contains the following members:

```
char *h_host;     /* host name */
char *h_serv;     /* service name */
```

The **nd_addrlist** structure contains the following members:

```
int             n_cnt;      /* number of addresses */
struct netbuf   *n_addrs;
```

netdir_getbyname( ) accepts some special-case host names. The host names are defined in **<netdir.h>**. The currently defined host names are:

HOST_SELF Represents the address to which local programs will bind their endpoints. **HOST_SELF** differs from the host name provided by **gethostname**(3C), which represents the address to which *remote* programs will bind their endpoints.

HOST_ANY Represents any host accessible by this transport provider. **HOST_ANY** allows applications to specify a required service without specifying a particular host name.

HOST_SELF_CONNECT Represents the host address that can be used to connect to the local host.

HOST_BROADCAST Represents the address for all hosts accessible by this transport provider. Network requests to this address will be received by all machines.

All fields of the **nd_hostserv** structure must be initialized.

To find the address of a given host and service on all available transports, call the **netdir_getbyname( )** routine with each **struct netconfig** structure returned by **getnetconfig**(3N).

The **netdir_getbyaddr( )** routine maps addresses to service names. This routine returns *service*, a list of host and service pairs that would yield this address. If more than one tuple of host and service name is returned, then the first tuple contains the preferred host and service names:

```
struct nd_hostservlist {
        int             *h_cnt;         /* number of hostservs found */
        struct hostserv *h_hostservs;
}
```

The **netdir_free( )** structure is used to free the structures allocated by the name to address translation routines. *ptr* points to the structure that has to be freed. The **struct_type** identifies the structure:

struct netbuf	**ND_ADDR**
struct nd_addrlist	**ND_ADDRLIST**
struct hostserv	**ND_HOSTSERV**
struct nd_hostservlist	**ND_HOSTSERVLIST**

The universal address returned by **taddr2uaddr( )** should be freed by **free**( ).

The **netdir_options( )** routine is used to do all transport-specific setups and option management. *fildes* is the associated file descriptor. *option*, *fildes*, and *pointer_to_args* are passed to the **netdir_options( )** routine for the transport specified in *config*. Currently four values are defined for *option*:

ND_SET_BROADCAST
ND_SET_RESERVEDPORT
ND_CHECK_RESERVEDPORT
ND_MERGEADDR

The **taddr2uaddr( )** and **uaddr2taddr( )** routines support translation between universal addresses and TLI type **netbufs**. The **taddr2uaddr( )** routine takes a **struct netbuf** data structure and returns a pointer to a string that contains the universal address. It returns **NULL** if the conversion is not possible. This is not a fatal condition as some transports may not suppose a universal address form.

uaddr2taddr( ) is the reverse of **taddr2uaddr( )**. It returns the **struct netbuf** data structure for the given universal address.

If a transport provider does not support an option, **netdir_options** returns **-1** and the error message can be printed through **netdir_perror( )** or **netdir_sperror( )**.

The specific actions of each option follow.

ND_SET_BROADCAST Sets the transport provider up to allow broadcast, if the transport supports broadcast. *fildes* is a file descriptor into the transport (i.e., the result of a **t_open** of **/dev/udp**). *pointer_to_args* is not used. If this completes, broadcast operations may be performed on file descriptor *fildes*.

ND_SET_RESERVEDPORT Allows the application to bind to a reserved port, if that concept exists for the transport provider. *fildes* is an unbound file descriptor into the transport. If *pointer_to_args* is **NULL,** *fildes* will be bound to a reserved port. If *pointer_to_args* is a pointer to a **netbuf** structure, an attempt will be made to bind to any reserved port on the specified address.

ND_CHECK_RESERVEDPORT Used to verify that the address corresponds to a reserved port, if that concept exists for the transport provider. *fildes* is not used. *pointer_to_args* is a pointer to a **netbuf** structure that contains the address. This option returns **0** only if the address specified in *pointer_to_args* is reserved.

ND_MERGEADDR Used to take a ''local address'' (like the **0.0.0.0** address that TCP uses) and return a ''real address'' that client machines can connect to. *fildes* is not used. *pointer_to_args* is a pointer to a **struct nd_mergearg**, which has the following members:

```
char s_uaddr;   /* server's universal address */
char c_uaddr;   /* client's universal address */
char m_uaddr;   /* the result */
```

If **s_uaddr** is something like **0.0.0.0.1.12**, and, if the call is successful, **m_uaddr** will be set to something like **192.11.109.89.1.12**. For most transports, **m_uaddr** is exactly what **s_uaddr** is.

RETURN VALUES

The **netdir_perror( )** routine prints an error message on the standard output stating why one of the name-to-address mapping routines failed. The error message is preceded by the string given as an argument.

The **netdir_sperror( )** routine returns a string containing an error message stating why one of the name-to-address mapping routines failed.

netdir_sperror( ) returns a pointer to a buffer which contains the error message string. This buffer is overwritten on each call. In multithreaded applications, this buffer is implemented as thread-specific data.

SEE ALSO

gethostname(3C), **getnetconfig**(3N), **getnetpath**(3N), **netconfig**(4)

NAME

nice – change priority of a process

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

int nice(*incr***)**
int *incr***;**

DESCRIPTION

The scheduling priority of the process is augmented by *incr*. Positive priorities get less service than normal. Priority 10 is recommended to users who wish to execute long-running programs without undue impact on system performance.

Negative increments are illegal, except when specified by the privileged user. The priority is limited to the range –20 (most urgent) to 20 (least). Requests for values above or below these limits result in the scheduling priority being set to the corresponding limit.

The priority of a process is passed to a child process by **fork**(2). For a privileged process to return to normal priority from an unknown state, **nice( )** should be called successively with arguments –40 (goes to priority –20 because of truncation), 20 (to get to 0), then 0 (to maintain compatibility with previous versions of this call).

RETURN VALUES

Upon successful completion, **nice( )** returns 0. Otherwise, a value of –1 is returned and **errno** is set to indicate the error.

ERRORS

The priority is not changed if:

EPERM The value of *incr* specified was negative, and the effective user ID is not the privileged user.

SEE ALSO

nice(1), **renice**(1), **fork**(2), **priocntl**(2), **getpriority**(3C)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

NAME

nis_db, db_initialize, db_create_table, db_destroy_table, db_first_entry, db_next_entry, db_reset_next_entry, db_list_entries, db_remove_entry, db_add_entry, db_table_exists, db_unload_table, db_checkpoint, db_standby, db_free_result – NIS+ Database access functions

SYNOPSIS

cc [*flag* . . .] *file*. . . **–lnisdb –lnsl** [*library*. . .]

#include <rpcsvc/nis.h>
#include <rpcsvc/nis_db.h>

bool db_initialize(const char ∗*dictionary_pathname*);

db_status db_create_table(const char ∗*table_name*, **const table_obj** ∗*table*);

db_status db_destroy_table(const char ∗*table_name*);

db_result ∗**db_first_entry(const char** ∗*table_name*, **const int** *numattrs*,
const nis_attr ∗*attrs*);

db_result ∗**db_next_entry(const char** ∗*table_name*, **const db_next_desc** ∗*next_handle*);

db_result ∗**db_reset_next_entry(const char** ∗*table_name*,
const db_next_desc ∗*next_handle*);

db_result ∗**db_list_entries(const char** ∗*table_name*, **const int** *numattrs*,
const nis_attr ∗*attrs*);

db_result ∗**db_remove_entry(const char** ∗*table_name*, **const int** *numattrs*,
const nis_attr ∗*attrs*);

db_result ∗**db_add_entry(const char** ∗*table_name*, **const int** *numattrs*,
const nis_attr ∗*attrs*, **const entry_obj** ∗*entry*);

db_status db_table_exists(const char ∗*table_name*);

db_status db_unload_table(const char ∗*table_name*);

db_status db_checkpoint(const char ∗*table_name*);

db_status db_standby(const char ∗*table_name*);

void db_free_result(db_result ∗);

MT-LEVEL

Unsafe

DESCRIPTION

These functions describe the interface between the NIS+ server and the underlying database. They are defined in the shared library **/usr/lib/libnisdb.so**.

The interface is a simple subset of a complete relational database and provides just those items that are needed by the NIS+ server daemon. When you replace the database, your interface routines should match these exactly. Also note that the database is responsible for verifying that the objects passed do not exceed the internal limits of the database being used.

The database's performance will directly affect the performance of the server. The default information base that is provided with NIS+ is the Structured Storage Manager (SSM). This is a memory based database that has been tuned for NIS+.

These routines should not be invoked by any NIS+ client. NIS+ clients should use the NIS+ tables API described in **nis_tables**(3N).

These routines only use the **table_obj**, **entry_obj** and the **nis_attr** structures defined in **<rpcsvc/nis.h>**. The NIS+ directory is itself stored in a table by the service daemon. This table has two columns, one searchable with the name of the object in it, the other non-searchable with binary XDRed data in it. The NIS+ server converts directory lookup requests in the namespace into table searches. The table it searches in response to these requests will have the same name as the directory of the name it is searching for.

The structure returned by the DB access routines is defined as:

```
enum db_status {DB_SUCCESS, DB_NOTFOUND, DB_NOTUNIQUE, DB_BADTABLE,
    DB_BADQUERY, DB_BADOBJECT, DB_MEMORY_LIMIT, DB_STORAGE_LIMIT,
    DB_INTERNAL_ERROR };

struct db_result {
    db_status       status;                 /* Result status */
    db_next_desc    nextinfo;               /* descriptor */
    struct {
                    u_int       objects_len;
                    entry_obj   *objects_val;
    } objects;                              /* A variable list
                                               of objects */
    long            ticks;                  /* execution time in
                                               microseconds */
};
```

For a complete description of NIS+ objects, see **nis_objects**(3N).

The structure **db_next_desc** should be used as an opaque handle for **db_next_entry( )** and **db_reset_next_entry( )**.

The **nis_attr** structure used in **db_first_entry** and other related functions is defined as follows:

```
struct nis_attr {
    char    *zattr_ndx;
    struct {
            u_int   zattr_val_len;
            char    *zattr_val_val;
    } zattr_val;
};
```

zattr_ndx is the name of the attribute. **zattr_val_len** is the value of the attribute **zattr_val_val**.

In **db_result**, the *objects* array contains objects if and only if the result returned in the *status* variable is **DB_SUCCESS**. A null pointer, instead of a pointer to a **db_result** structure, is returned if there is insufficient memory to create the structure.

db_initialize( ) is called prior to any interaction with the database. It takes as argument the pathname of the file that contains, or will contain, catalog information associated with the database.

db_create_table( ) creates a new table using the given table name and the table object. It returns TRUE if the table was successfully created; FALSE otherwise.

db_destroy_table( ) destroys the table of the given name. It returns TRUE if the destruction was successful; FALSE otherwise.

db_first_entry( ) returns a copy of the first entry in the specified table that satisfies the given attributes. If no attributes are supplied, a copy of the first entry in the table is returned. **attrs** is an array of **nis_attr** structure with *numattrs* number of elements. The returned structure, **db_result**, contains a structure, **db_next_desc**, to be used as an argument to **db_next_entry( )** or **db_reset_next_entry( )**. **db_next_desc** should only be used only as an opaque handle. **db_free_result( )** can be used to free up the returned **db_result** structure.

db_next_entry( ) returns a copy of the next entry as indicated by the *next_handle*. An initial call to **db_first_entry( )**, followed by a sequence of calls to **db_next_entry( )**, can be used to successfully obtain entries of an entire table or entries that satisfy the attributes supplied to **db_first_entry( )**. **db_free_result( )** can be used to free up the returned **db_result** structure.

db_reset_next_entry( ) terminates the **db_first_entry( )** / **db_next_entry( )** sequence as indicated by *next_handle*, freeing any resources that have been used to maintain the sequence. After a call to **db_reset_next_entry( )**, a call to **db_next_entry( )** using the same *next_handle* would fail, returning a **DB_BADQUERY** reply. **db_free_result( )** can be used to free up the returned **db_result** structure.

db_list_entries( ) returns copies of entries that satisfy the given attributes. **db_free_result( )** can be used to free up the returned **db_result** structure. **attrs** is an array of **nis_attr** structure with *numattrs* number of elements.

db_remove_entry( ) removes all entries that satisfy the given attributes. **db_free_result( )** can be used to free up the returned **db_result** structure. **attrs** is an array of **nis_attr** structure with *numattrs* number of elements.

db_add_entry( ) adds a copy of the given object to the specified table, replacing the one identified by the given attributes. If the given attributes identify more than one object, **DB_NOTUNIQUE** is returned. If no object is identified by the given attributes, the object is added. **attrs** is an array of **nis_attr** structure with *numattrs* number of elements. **db_free_result( )** can be used to free up the returned **db_result** structure.

db_table_exists( ) provides an efficient way for the NIS+ service to detect that a table exists. This increases response time to the client and lowers the load on the server.

db_unload_table( ) is used by the service to unload or deactivate tables that are not currently being used. The service internally keeps track of access patterns to tables and will unload those tables that have not been accessed for a while. By unloading infrequently accessed tables, the service can minimize the amount of system resources for efficient operation.

db_checkpoint( ) organizes the contents of the table in a more efficient manner. Checkpointing may mean different things to different types of databases. It does not affect the logical contents of the table — operations and queries should return the same result before and after a checkpoint. For example, in a log-based system, checkpointing may mean incorporating log entries of updates accumulated since the previous checkpoint into the table.

db_free_result( ) frees up the space allocated by various functions listed on this manual page that return a **db_result** structure.

db_standby( ) is an advisory call to the database manager. This call informs the database that activity has slowed down and it can free up unnecessary resources such as file descriptors.

PROGRAMMING

Most of the routines in this library use an NIS+ *name* to identify the object that the user desires. The name must be in canonical form before being passed to the database because one server may be serving several namespaces and discrimination of the requested objects is accomplished by comparing the domain names.

DIAGNOSTICS

DB_SUCCESS The query or operation completed successfully and returned status.

DB_NOTFOUND
The name or entry that was named in the argument did not exist.

DB_NOTUNIQUE
An attempt was made to remove an entry from a table that is not uniquely specified.

DB_BADQUERY
The query that was submitted to the database was invalid (for example, it might name some nonexistent fields).

DB_BADTABLE
The table was corrupted.

DB_BADOBJECT
The fields of the object does not conform to the fields of the table to which it is being added.

DB_MEMORY_LIMIT
There is insufficient memory to complete the operation requested.

DB_STORAGE_LIMIT
There is insufficient file storage available to complete the operation requested.

DB_INTERNAL_ERROR
An internal error was encountered during the execution of the operation requested (either a programming error or an unrecoverable exception).

SEE ALSO

rpc.nisd(1M), **nis_objects**(3N), **nisfiles**(4)

NAME

nis_error, nis_sperrno, nis_perror, nis_lerror, nis_sperror, nis_sperror_r – display NIS+ error messages

SYNOPSIS

cc [*flag* . . .] *file*. . . **–lnsl** [*library*. . .]

#include <rpcsvc/nis.h>

char ∗**nis_sperrno(const nis_error** *status***);**

void nis_perror(const nis_error *status***, const char** ∗*label***);**

void nis_lerror(const nis_error *status***, const char** ∗*label***);**

char ∗**nis_sperror_r(nis_error** *status***, char** ∗*label***, char** ∗ *buf***);**

char ∗**nis_sperror(const nis_error** *status***, const char** ∗*label***);**

MT-LEVEL

Safe

DESCRIPTION

These functions convert NIS+ status values into text strings.

nis_sperrno( ) simply returns a pointer to a string constant which is the error string.

nis_perror( ) prints the error message corresponding to *status* as ‘‘*label*: **error message**’’ on standard error.

nis_lerror( ) sends the error text to **syslog**(3) at level LOG_ERR.

The function **nis_sperror_r( )**, returns a pointer to a string that can be used or copied using the **strdup( )** function (see **string**(3C).) The caller must supply a string buffer, **buf**, large enough to hold the error string (a buffer size of 128 bytes is guaranteed to be sufficiently large).

The last function, **nis_sperror( )**, is similar to **nis_sperror_r( )** except that the string is returned as a pointer to a buffer that is reused on each call. **nis_sperror_r( )** is the preferred interface, since it is suitable for single-threaded and multi-threaded programs.

SEE ALSO

niserror(1), **string**(3C), **syslog**(3)

NOTES

When compiling multithreaded applications, see **Intro**(3), *Notes On Multithread Applications*, for information about the use of the **_REENTRANT** flag.

NAME

nis_groups, nis_ismember, nis_addmember, nis_removemember, nis_creategroup, nis_destroygroup, nis_verifygroup, nis_print_group_entry, nis_map_group, __nis_map_group – NIS+ group manipulation functions

SYNOPSIS

cc [*flag* . . .] *file*. . . **-lnsl** [*library*. . .]

#include <rpcsvc/nis.h>

bool_t nis_ismember(const nis_name *principal***, const nis_name** *group***);**

nis_error nis_addmember(const nis_name *member***, const nis_name** *group***);**

nis_error nis_removemember(const nis_name *member***, const nis_name** *group***);**

nis_error nis_creategroup(const nis_name *group***, const u_long** *flags***);**

nis_error nis_destroygroup(const nis_name *group***);**

void nis_print_group_entry(const nis_name *group***);**

nis_error nis_verifygroup(const nis_name *group***);**

MT-LEVEL

MT-Safe

DESCRIPTION

These functions manipulate NIS+ groups. They are used by NIS+ clients and servers, and are the interfaces to the group authorization object.

The names of NIS+ groups are syntactically similar to names of NIS+ objects but they occupy a separate namespace. A group named "a.b.c.d." is represented by a NIS+ group object named "a.groups_dir.b.c.d."; the functions described here all expect the name of the group, not the name of the corresponding group object.

There are three types of group members:

- An *explicit* member is just a NIS+ principal-name, for example "wickedwitch.west.oz."
- An *implicit* ("domain") member, written "*.west.oz.", means that all principals in the given domain belong to this member. No other forms of wildcarding are allowed: "wickedwitch.*.oz." is invalid, as is "wickedwitch.west.*.". Note that principals in subdomains of the given domain are *not* included.
- A *recursive* ("group") member, written "@cowards.oz.", refers to another group; all principals that belong to that group are considered to belong here.

Any member may be made *negative* by prefixing it with a minus sign ('–'). A group may thus contain explicit, implicit, recursive, negative explicit, negative implicit, and negative recursive members.

A principal is considered to belong to a group if it belongs to at least one non-negative group member of the group and belongs to no negative group members.

The **nis_ismember( )** function returns TRUE if it can establish that *principal* belongs to *group*; otherwise it returns FALSE.

The **nis_addmember( )** and **nis_removemember( )** functions add or remove a member. They do not check whether the member is valid. The user must have read and modify rights for the group in question.

The **nis_creategroup()** and **nis_destroygroup()** functions create and destroy group objects. The user must have create or destroy rights, respectively, for the *groups_dir* directory in the appropriate domain. The parameter *flags* to **nis_creategroup()** is currently unused and should be set to zero.

The **nis_print_group_entry()** function lists a group's members on the standard output.

The **nis_verifygroup()** function returns **NIS_SUCCESS** if the given group exists, otherwise it returns an error code.

EXAMPLES

Simple Memberships

Given a group **sadsouls.oz.** with members **tinman.oz.**, **lion.oz.**, and **scarecrow.oz.**, the function call

bool_var = nis_ismember("lion.oz.", "sadsouls.oz.");

will return 1 (TRUE) and the function call

bool_var = nis_ismember("toto.oz.", "sadsouls.oz.");

will return 0 (FALSE).

Implicit Memberships

Given a group **baddies.oz.**, with members **wickedwitch.west.oz.** and ∗**.monkeys.west.oz.**, the function call

bool_var = nis_ismember("hogan.monkeys.west.oz.", "baddies.oz.");

will return 1 (TRUE) because any principal from the **monkeys.west.oz.** domain belongs to the implicit group ∗**.monkeys.west.oz.**, but the function call

bool_var = nis_ismember("hogan.big.monkeys.west.oz.", "baddies.oz.");

will return 0 (FALSE).

Recursive Memberships

Given a group **goodandbad.oz.**, with members **toto.kansas**, **@sadsouls.oz.**, and **@baddies.oz.**, and the groups **sadsouls.oz.** and **baddies.oz.** defined above, the function call

bool_var = nis_ismember("wickedwitch.west.oz.", "goodandbad.oz.");

will return 1 (TRUE), because **wickedwitch.west.oz.** is a member of the **baddies.oz.** group which is recursively included in the **goodandbad.oz.** group.

SEE ALSO

nisgrpadm(1), **nis_objects**(3N)

NOTES

These functions only accept fully-qualified NIS+ names.

A group is represented by a NIS+ object (see **nis_objects**(3N)) with a variant part that is defined in the **group_obj** structure. It contains the following fields:

```
u_long          gr_flags;           /* Interpretation Flags
                                       (currently unused) */
struct {
        u_int       gr_members_len;
        nis_name    *gr_members_val;
} gr_members;                       /* Array of members */
```

NIS+ servers and clients maintain a local cache of expanded groups to enhance their performance when checking for group membership. Should the membership of a group change, servers and clients with that group cached will not see the change until either the group cache has expired or it is explicitly flushed. A server's cache may be flushed programmatically by calling the **nis_servstate( )** function with tag **TAG_GCACHE** and a value of 1.

There are currently no known methods for **nis_ismember( )**, **nis_print_group_entry( )**, and **nis_verifygroup( )** to get their answers from only the master server.

NAME

nis_local_names, nis_local_directory, nis_local_host, nis_local_group, nis_local_principal – NIS+ local names

SYNOPSIS

cc [*flag* . . .] *file*. . . **–lnsl** [*library*. . .]

#include <rpcsvc/nis.h>

nis_name nis_local_directory(void);

nis_name nis_local_host(void);

nis_name nis_local_group(void);

nis_name nis_local_principal(void);

MT-LEVEL

MT-Safe

DESCRIPTION

These functions return several default NIS+ names associated with the current process.

nis_local_directory( ) returns the name of the NIS+ domain for this machine. This is currently the same as the Secure RPC domain returned by the **sysinfo**(2) system call.

nis_local_host( ) returns the NIS+ name of the current machine. This is the fully qualified name for the host and is either the value returned by the **gethostname**(3C) function or, if the host name is only partially qualified, the concatenation of that value and the name of the NIS+ directory. Note that if a machine's name and address cannot be found in the local NIS+ directory, its hostname must be fully qualified.

nis_local_group( ) returns the name of the current NIS+ group name. This is currently set by setting the environment variable **NIS_GROUP** to the groupname.

nis_local_principal( ) returns the NIS+ principal name for the user associated with the effective UID of the calling process. This function maps the effective uid into a principal name by looking for a LOCAL type credential (see **nisaddcred**(1M)) in the table named *cred.org_dir* in the default domain.

Note: The result returned by these routines is a pointer to a data structure with the NIS+ library, and should be considered a "read-only" result and should not be modified.

ENVIRONMENT

NIS_GROUP This variable contains the name of the local NIS+ group. If the name is not fully qualified, the value returned by **nis_local_directory( )** will be concatenated to it.

SEE ALSO

nisdefaults(1), **nisaddcred**(1M), **sysinfo**(2), **gethostname**(3C), **nis_names**(3N), **nis_objects**(3N)

NAME

nis_names, nis_lookup, nis_add, nis_remove, nis_modify, nis_freeresult – NIS+ namespace functions

SYNOPSIS

cc [*flag* . . .] *file*. . . **-lnsl** [*library*. . .]

#include <rpcsvc/nis.h>

nis_result ∗**nis_lookup(const nis_name** *name***, const u_long** *flags***);**

nis_result ∗**nis_add(const nis_name** *name***, const nis_object** ∗*obj***);**

nis_result ∗**nis_remove(const nis_name** *name***, const nis_object** ∗*obj***);**

nis_result ∗**nis_modify(const nis_name** *name***, const nis_object** ∗*obj***);**

void nis_freeresult(nis_result ∗*result***);**

MT-LEVEL

MT-Safe

DESCRIPTION

These functions are used to locate and manipulate all NIS+ objects (see **nis_objects**(3N)) except the NIS+ entry objects. To look up the NIS+ entry objects within a NIS+ table, refer to **nis_subr**(3N).

nis_lookup( ) resolves a NIS+ name and returns a copy of that object from a NIS+ server. **nis_add( )** and **nis_remove( )** add and remove objects to the NIS+ namespace, respectively. **nis_modify( )** can change specific attributes of an object that already exists in the namespace.

These functions should be used only with names that refer to an NIS+ Directory, NIS+ Table, NIS+ Group, or NIS+ Private object. If a name refers to an NIS+ entry object, the functions listed in **nis_subr**(3N) should be used.

nis_freeresult( ) frees all memory associated with a **nis_result** structure. This function must be called to free the memory associated with a NIS+ result. **nis_lookup( ), nis_add( ), nis_remove( ),** and **nis_modify( )** all return a pointer to a **nis_result structure** which *must* be freed by calling **nis_freeresult( )** when you have finished using it. If one or more of the objects returned in the structure need to be retained, they can be copied with **nis_clone_object**(3N) (see **nis_subr**(3N)).

nis_lookup( ) takes two parameters, the name of the object to be resolved in *name*, and a flags parameter, *flags*, which is defined below. The object name is expected to correspond to the syntax of a non-indexed NIS+ name (see **nis_tables**(3N)). The **nis_lookup( )** function is the *only* function from this group that can use a non-fully qualified name. If the parameter *name* is not a fully qualified name, then the flag **EXPAND_NAME** *must* be specified in the call. If this flag is not specified, the function will fail with the error NIS_BADNAME.

The *flags* parameter is constructed by logically ORing zero or more flags from the following list.

FOLLOW_LINKS When specified, the client library will ‘‘follow’’ links by issuing another NIS+ lookup call for the object named by the link. If the linked object is itself a link, then this process will iterate until the

either a object is found that is not a *LINK* type object, or the library has followed 16 links.

HARD_LOOKUP When specified, the client library will retry the lookup until it is answered by a server. Using this flag will cause the library to block until at least one NIS+ server is available. If the network connectivity is impaired, this can be a relatively long time.

NO_CACHE When specified, the client library will bypass any object caches and will get the object from either the master NIS+ server or one of its replicas.

MASTER_ONLY When specified, the client library will bypass any object caches and any domain replicas and fetch the object from the NIS+ master server for the object's domain. This insures that the object returned is up to date at the cost of a possible performance degradation and failure if the master server is unavailable or physically distant.

EXPAND_NAME When specified, the client library will attempt to expand a partially qualified name by calling the function **nis_getnames( )** (see **nis_subr**(3N)) which uses the environment variable **NIS_PATH**.

The status value may be translated to ascii text using the function **nis_sperrno( )** (see **nis_error**(3N)).

On return, the *objects* array in the result will contain one and possibly several objects that were resolved by the request. If the FOLLOW_LINKS flag was present, on success the function could return several entry objects if the link in question pointed within a table. If an error occurred when following a link, the objects array will contain a copy of the link object itself.

The function **nis_add( )** will take the object *obj* and add it to the NIS+ namespace with the name *name*. This operation will fail if the client making the request does not have the *create* access right for the domain in which this object will be added. The parameter *name* must contain a fully qualified NIS+ name. The object members *zo_name* and *zo_domain* will be constructed from this name. This operation will fail if the object already exists. This feature prevents the accidental addition of objects over another object that has been added by another process.

The function **nis_remove( )** will remove the object with name *name* from the NIS+ namespace. The client making this request must have the *destroy* access right for the domain in which this object resides. If the named object is a link, the link is removed and *not* the object that it points to. If the parameter *obj* is not NULL, it is assumed to point to a copy of the object being removed. In this case, if the object on the server does not have the same object identifier as the object being passed, the operation will fail with the NIS_NOTSAMEOBJ error. This feature allows the client to insure that it is removing the desired object. The parameter *name* must contain a fully qualified NIS+ name.

The function **nis_modify( )** will modify the object named by *name* to the field values in the object pointed to by *obj*. This object should contain a copy of the object from the name space that is being modified. This operation will fail with the error NIS_NOTSAMEOBJ if the object identifier of the passed object does not match that of the object being modified

in the namespace.

Note: Normally the contents of the member *zo_name* in the *nis_object* structure would be constructed from the name passed in the *name* parameter. However, if it is non-NULL the client library will use the name in the *zo_name* member to perform a rename operation on the object. This name *must not* contain any unquoted '.'(dot) characters. If these conditions are not met the operation will fail and return the NIS_BADNAME error code.

Results

These functions return a pointer to a structure of type **nis_result**:

```
struct nis_result {
        nis_error status;
        struct {
                        u_int objects_len;
                        nis_object *objects_val;
        } objects;
        netobj          cookie;
        u_long          zticks;
        u_long          dticks;
        u_long          aticks;
        u_long          cticks;
};
```

The *status* member contains the error status of the the operation. A text message that describes the error can be obtained by calling the function **nis_sperrno()** (see **nis_error**(3N)).

The *objects* structure contains two members. *objects_val* is an array of *nis_object* structures; *objects_len* is the number of cells in the array. These objects will be freed by the call to **nis_freeresult()**. If you need to keep a copy of one or more objects, they can be copied with the function **nis_clone_object()** and freed with the function **nis_destroy_object()** (see **nis_server**(3N)). Refer to **nis_objects**(3N) for a description of the **nis_object** structure.

The various ticks contain details of where the time was taken during a request. They can be used to tune one's data organization for faster access and to compare different database implementations (see **nis_db**(3N)).

zticks The time spent in the NIS+ service itself. This count starts when the server receives the request and stops when it sends the reply.

dticks The time spent in the database backend. This time is measured from the time a database call starts, until the result is returned. If the request results in multiple calls to the database, this is the sum of all the time spent in those calls.

aticks The time spent in any ''accelerators'' or caches. This includes the time required to locate the server needed to resolve the request.

cticks The total time spent in the request. This clock starts when you enter the client library and stops when a result is returned. By subtracting the sum of the other ticks values from this value, you can obtain the local

overhead of generating a NIS+ request.

Subtracting the value in *dticks* from the value in *zticks* will yield the time spent in the service code itself. Subtracting the sum of the values in *zticks* and *aticks* from the value in *cticks* will yield the time spent in the client library itself. Note: all of the tick times are measured in microseconds.

RETURN VALUES

The client library can return a variety of error returns and diagnostics. The more salient ones are documented below.

NIS_SUCCESS The request was successful.

NIS_S_SUCCESS The request was successful, however the object returned came from an object cache and not directly from the server. If you do not wish to see objects from object caches you must specify the flag **NO_CACHE** when you call the lookup function.

NIS_NOTFOUND The named object does not exist in the namespace.

NIS_CACHEEXPIRED The object returned came from an object cache taht has *expired*. The time to live value has gone to zero and the object may have changed. If the flag **NO_CACHE** was passed to the lookup function then the lookup function will retry the operation to get an unexpired copy of the object.

NIS_NAMEUNREACHABLE A server for the directory of the named object could not be reached. This can occur when there is a network partition or all servers have crashed. See the **HARD_LOOKUP** flag.

NIS_UNKNOWNOBJ The object returned is of an unknown type.

NIS_TRYAGAIN The server connected to was too busy to handle your request. For the *add*, *remove*, and *modify* operations this is returned when either the master server for a directory is unavailable or it is in the process of checkpointing its database. It can also be returned when the server is updating it's internal state. And in the case of **nis_list( )** if the client specifies a callback and the server does not have enough resources to handle the callback.

NIS_SYSTEMERROR A generic system error occurred while attempting the request. Most commonly the server has crashed or the database has become corrupted. Check the syslog record for error messages from the server.

NIS_NOT_ME A request was made to a server that does not serve the name in question. Normally this will not occur, however if you are not using the built in location mechanism for servers you may see this if your mechanism is broken.

NIS_NOMEMORY Generally a fatal result. It means that the service ran out of

heap space.

NIS_NAMEEXISTS An attempt was made to add a name that already exists. To add the name, first remove the existing name and then add the new object or modify the existing named object.

NIS_NOTMASTER An attempt was made to update the database on a replica server.

NIS_INVALIDOBJ The object pointed to by *obj* is not a valid **NIS+** object.

NIS_BADNAME The name passed to the function is not a legal **NIS+** name.

NIS_LINKNAMEERROR The name passed resolved to a *LINK* type object and the contents of the link pointed to an invalid name.

NIS_NOTSAMEOBJ An attempt to remove an object from the namespace was aborted because the object that would have been removed was not the same object that was passed in the request.

NIS_NOSUCHNAME This hard error indicates that the named directory of the table object does not exist. This occurs when the server that should be the parent of the server that serves the table, does not know about the directory in which the table resides.

NIS_NOSUCHTABLE The named table does not exist.

NIS_MODFAIL The attempted modification failed.

NIS_FOREIGNNS The name could not be completely resolved. When the name passed to the function would resolve in a namespace that is outside the NIS+ name tree, this error is returned with a NIS+ object of type **DIRECTORY**, which contains the type of namespace and contact information for a server within that namespace.

NIS_RPCERROR This fatal error indicates the RPC subsystem failed in some way. Generally there will be a **syslog**(3) message indicating why the RPC request failed.

ENVIRONMENT

NIS_PATH If the flag **EXPAND_NAME** is set, this variable is the search path used by **nis_lookup()**.

SEE ALSO

nis_error(3N), **nis_objects**(3N), **nis_server**(3N), **nis_subr**(3N), **nis_tables**(3N)

NOTES

You cannot modify the name of an object if that modification would cause the object to reside in a different domain.

You cannot modify the schema of a table object.

NAME

nis_objects – NIS+ object formats

SYNOPSIS

cc [*flag* . . .] *file*. . . **-lnsl** [*library*. . .]

/usr/include/rpcsvc/nis_objects.x

DESCRIPTION

Common Attributes

The NIS+ service uses a variant record structure to hold the contents of the objects that are used by the NIS+ service. These objects all share a common structure which defines a set of attributes that all objects possess. The **nis_object** structure contains the following members:

```
typedef  char          *nis_name;
struct   nis_object {
         nis_oid       zo_oid;
         nis_name      zo_name;
         nis_name      zo_owner;
         nis_name      zo_group;
         nis_name      zo_domain;
         u_long        zo_access;
         u_long        zo_ttl;
         objdata       zo_data;
};
```

In this structure, the first member **zo_oid**, is a 64 bit number that uniquely identifies this instance of the object on this server. This member is filled in by the server when the object is created and changed by the server when the object is modified. When used in conjunction with the object's name and domain it uniquely identifies the object in the entire NIS+ namespace.

The second member, **zo_name**, contains the leaf name of the object. This name is *never* terminated with a '.' (dot). When an object is created or added to the namespace, the client library will automatically fill in this field and the domain name from the name that was passed to the function.

zo_domain contains the name of the NIS+ domain to which this object belongs. This information is useful when tracking the parentage of an object from a cache. When used in conjunction with the members **zo_name** and **zo_oid**, it uniquely identifies an object. This makes it possible to always reconstruct the name of an object by using the code fragment

```
sprintf(buf,"%s.%s", obj→zo_name, obj→zo_domain);
```

The **zo_owner** and **zo_group** members contain the NIS+ names of the object's principal owner and group owner, respectively. Both names *must be* NIS+ fully qualified names. However, neither name can be used directly to identify the object they represent. This stems from the condition that NIS+ uses itself to store information that it exports.

The **zo_owner** member contains a fully qualified NIS+ name of the form *principal.domain*. This name is called a NIS+ principal name and is used to identify authentication information in a credential table. When the server constructs a search query of the form

[cname=*principal***],cred.org_dir.***domain***.**

The query will return to the server credential information about *principal* for all flavors of RPC authentication that are in use by that principal. When an RPC request is made to the server, the authentication flavor is extracted from the request and is used to find out the NIS+ principal name of the client. For example, if the client is using the AUTH_DES authentication flavor, it will include in the authentication credentials the network name or *netname* of the user making the request. This netname will be of the form

unix.*UID@domain*

The NIS+ server will then construct a query on the credential database of the form

[auth_name=*netname***,auth_type=AUTH_DES],cred.org_dir.***domain***.**

This query will return an entry which contains a principal name in the first column. This NIS+ principal name is used to control access to NIS+ objects.

The group owner for the object is treated differently. The group owner member is optional (it should be the null string if not present) but must be fully qualified if present. A group name takes the form

group.domain.

which the server then maps into a name of the form

*group***.groups_dir.***domain*.

The purpose of this mapping is to prevent NIS+ group names from conflicting with user specified domain or table names. For example, if a domain was called *engineering.foo.com*., then without the mapping a NIS+ group of the same name to represent members of engineering would not be possible. The contents of groups are lists of NIS+ principal names which are used exactly like the **zo_owner** name in the object. See **nis_groups**(3N) for more details.

The **zo_access** member contains the bitmask of access rights assigned to this object. There are four access rights defined, and four are reserved for future use and must be zero. This group of 8 access rights can be granted to four categories of client. These categories are the object's owner, the object's group owner, all authenticated clients (world), and all unauthenticated clients (nobody). Note that access granted to ''nobody'' is really access granted to everyone, authenticated and unauthenticated clients.

The **zo_ttl** member contains the number of seconds that the object can ''live'' in a cache before it is expired. This value is called the time to live for this object. This number is particularly important on group and directory (domain) objects. When an object is cached, the current time is added to the value in **zo_ttl**. Then each time the cached object is used, the time in **zo_ttl** is compared with the current time. If the current time is later than the time in **zo_ttl** the object is said to have expired and the cached copy should not be used.

Setting the TTL is somewhat of an art. You can think of it as the ''half life'' of the object, or half the amount of time you believe will pass before the object changes. The benefit of setting the ttl to a large number is that the object will stay in a cache for long periods of time. The problem with setting it to a large value is that when the object changes it will take a long time for the caches to flush out old copies of that object. The problems and benefits are reversed for setting the time to a small value. Generally setting the value to 43200 (12 hrs) is reasonable for things that change day to day, and 3024000 is good for things that change week to week. Setting the value to 0 will prevent the object from ever being cached since it would expire immediately.

The **zo_data** member is a discriminated union with the following members:

```
zotypes zo_type;
union {
        struct directory_obj    di_data;
        struct group_obj        gr_data;
        struct table_obj        ta_data;
        struct entry_obj        en_data;
        struct link_obj         li_data;
        struct {
                u_int           po_data_len;
                char            *po_data_val;
        } po_data;
} objdata_u;
```

The union is discriminated based on the type value contained in **zo_type**. There six types of objects currently defined in the NIS+ service. These types are the directory, link, group, table, entry, and private types.

```
enum zotypes {
        BOGUS_OBJ       = 0,
        NO_OBJ          = 1,
        DIRECTORY_OBJ   = 2,
        GROUP_OBJ       = 3,
        TABLE_OBJ       = 4,
        ENTRY_OBJ       = 5,
        LINK_OBJ        = 6,
        PRIVATE_OBJ     = 7
};
typedef enum zotypes zotypes;
```

All object types define a structure that contains data specific to that type of object. The simplest are private objects which are defined to contain a variable length array of octets. Only the owner of the object is expected to understand the contents of a private object. The following section describe the other five object types in more significant detail.

Directory Objects

The first type of object is the *directory* object. This object's variant part is defined as follows:

```
enum nstype {
        UNKNOWN     = 0,
        NIS         = 1,
        SUNYP       = 2,
        DNS         = 4,
        X500        = 5,
        DNANS       = 6,
        XCHS        = 7,
}
typedef enum nstype nstype;

struct oar_mask {
            u_long      oa_rights;
            zotypes     oa_otype;
}
typedef struct oar_mask oar_mask;

struct endpoint {
            char        *uaddr;
            char        *family;
            char        *proto;
}
typedef struct endpoint endpoint;

struct nis_server {
            nis_name  name;
            struct {
                        u_int       ep_len;
                        endpoint    *ep_val;
            } ep;
            u_long      key_type;
            netobj      pkey;
}
typedef struct nis_server nis_server;

struct directory_obj {
            nis_name   do_name;
            nstype      do_type;
            struct {
                        u_int       do_servers_len;
                        nis_server  *do_servers_val;
            } do_servers;
            u_long      do_ttl;
            struct {
```

```
                u_int       do_armask_len;
                oar_mask    *do_armask_val;
            } do_armask;
        }
        typedef struct directory_obj directory_obj;
```

The main structure contains five primary members: **do_name**, **do_type**, **do_servers**, **do_ttl**, and **do_armask**. The information in the **do_servers** structure is sufficient for the client library to create a network connection with the named server for the directory.

The **do_name** member contains the name of the directory or domain represented in a format that is understandable by the type of nameservice serving that domain. In the case of NIS+ domains, this is the same as the name that can be composed using the **zo_name** and **zo_domain** members. For other name services, this name will be a name that they understand. For example, if this were a directory object describing an X.500 namespace that is ''under'' the NIS+ directory *eng.sun.com.*, this name might contain ''/C=US, /O=Sun Microsystems, /OU=Engineering/''. The type of nameservice that is being described is determined by the value of the member **do_type**.

The **do_servers** structure contains two members. **do_servers_val** is an array of *nis_server* structures; **do_servers_len** is the number of cells in the array. The *nis_server* structure is designed to contain enough information such that machines on the network providing name services can be contacted without having to use a name service. In the case of NIS+ servers, this information is the name of the machine in *name*, its public key for authentication in *pkey*, and a variable length array of endpoints, each of which describes the network endpoint for the **rpcbind** daemon on the named machine. The client library uses the addresses to contact the server using a transport that both the client and server can communicate on and then queries the **rpcbind** daemon to get the actual transport address that the server is using.

Note that the first server in the *do_servers* list is always the master server for the directory.

The *key_type* field describes the type of key stored in the *pkey* netobj (see **/usr/include/rpc/xdr.h** for a definition of the network object structure). Currently supported types are **NIS_PK_NONE** for no public key and **NIS_PK_DH** for a Diffie-Hellman type public key.

The **do_ttl** member contains a copy of the **zo_ttl** member from the common attributes. This is the duplicated because the cache manager only caches the variant part of the directory object.

The **do_armask** structure contains two members. **do_armask_val** is an array of **oar_mask** structures; **do_armask_len** is the number of cells in the array. The **oar_mask** structure contains two members: **oa_rights** specifies the access rights allowed for objects of type **oa_otype**. These access rights are used for objects of the given type in the directory when they are present in this array.

The granting of access rights for objects contained within a directory is actually two-tiered. If the directory object itself grants a given access right (using the **zo_access** member in the **nis_object** structure representing the directory), then all objects within the directory are allowed that access. Otherwise, the **do_armask** structure is examined to see

if the access is allowed specifically for that type of structure. This allows the administrator of a namespace to set separate policies for different object types, for example, one policy for the creation of tables and another policy for the creation of other directories. See **nis+**(1) for more details.

Link Objects

Link objects provide a means of providing *aliases* or symbolic links within the namespace. Their variant part is defined as follows.

```
struct link_obj {
        zotypes     li_rtype;
        struct {
                    u_int       li_attrs_len;
                    nis_attr    *li_attrs_val;
        } li_attrs;
        nis_name li_name;
}
```

The **li_rtype** member contains the object type of the object pointed to by the link. This is only a hint, since the object which the link points to may have changed or been removed. The fully qualified name of the object (table or otherwise) is specified in the member **li_name**.

NIS+ links can point to either other objects within the NIS+ namespace, or to entries within a NIS+ table. If the object pointed to by the link is a table and the member **li_attrs** has a nonzero number of attributes (index name / value pairs) specified, the table is searched when this link is followed. All entries which match the specified search pattern are returned. Note, that unless the flag **FOLLOW_LINKS** is specified, the **nis_lookup**(3N) function will always return non-entry objects.

Group Objects

Group objects contain a membership list of NIS+ principals. The group objects' variant part is defined as follows.

```
struct group_obj {
        u_long      gr_flags;
        struct {
                    u_int       gr_members_len;
                    nis_name    *gr_members_val;
        } gr_members;
}
```

The **gr_flags** member contains flags that are currently unused. The **gr_members** structure contains the list of principals. For a complete description of how group objects are manipulated see **nis_groups**(3N).

Table Objects

The NIS+ table object is analogous to a YP map. The differences stem from the access controls, and the variable schemas that NIS+ allows. The table objects data structure is defined as follows:

```
#define TA_BINARY          1
#define TA_CRYPT           2
#define TA_XDR             4
#define TA_SEARCHABLE      8
#define TA_CASE            16

struct table_col {
        char        *tc_name;
        u_long      tc_flags;
        u_long      tc_rights;
}
typedef struct table_col table_col;

struct table_obj {
        char        *ta_type;
        u_int       ta_maxcol;
        u_char      ta_sep;
        struct {
                    u_int       ta_cols_len;
                    table_col   *ta_cols_val;
        } ta_cols;
        char        *ta_path;
}
```

The **ta_type** member contains a string that identifies the type of entries in this table. NIS+ does not enforce any policies as to the contents of this string. However, when entries are added to the table, the NIS+ service will check to see that they have the same ''type'' as the table as specified by this member.

The structure **ta_cols** contains two members. **ta_cols_val** is an array of **table_col** structures. The length of the array depends on the number of columns in the table; it is defined when the table is created and is stored in **ta_cols_len**. **ta_maxcol** also contains the number of columns in the table and always has the same value as **ta_cols_len**. Once the table is created, this length field cannot be changed.

The **ta_sep** character is used by client applications that wish to print out an entry from the table. Typically this is either space (‘‘ ’’) or colon (‘‘:’’).

The **ta_path** string defines a concatenation path for tables. This string contains an ordered list of fully qualified table names, separated by colons, that are to be searched if a search on this table fails to match any entries. This path is only used with the flag **FOLLOW_PATH** with a **nis_list( )** call. See **nis_tables**(3N) for information on these flags.

In addition to checking the type, the service will check that the number of columns in an entry is the same as those in the table before allowing that entry to be added.

Each column has associated with it a name in **tc_name**, a set of flags in **tc_flags**, and a set of access rights in **tc_rights**. The name should be indicative of the contents of that column.

The **TA_BINARY** flag indicates that data in the column is binary (rather than text). Columns that are searchable cannot contain binary data. The **TA_CRYPT** flag specifies that the information in this column should be encrypted prior to sending it over the network. This flag has no effect in the export version of NIS+. The **TA_XDR** flag is used to tell the client application that the data in this column is encoded using the XDR protocol. The **TA_BINARY** flag must be specified with the XDR flag. Further, by convention, the name of a column that has the **TA_XDR** flag set is the name of the XDR function that will decode the data in that column.

The **TA_SEARCHABLE** flag specifies that values in this column can be searched. Searchable columns must contain textual data and must have a name associated with them. The flag **TA_CASE** specifies that searches involving this column ignore the case of the value in the column. At least one of the columns in the table should be searchable. Also, the combination of all searchable column values should uniquely select an entry within the table.

Entry Objects

Entry objects are stored in tables. The structure used to define the entry data is as follows.

```
#define EN_BINARY       1
#define EN_CRYPT        2
#define EN_XDR          4
#define EN_MODIFIED     8

struct entry_col {
        u_long      ec_flags;
        struct {
                    u_int       ec_value_len;
                    char        *ec_value_val;
        } ec_value;
}
typedef struct entry_col entry_col;

struct entry_obj {
        char        *en_type;
        struct {
                    u_int       en_cols_len;
                    entry_col   *en_cols_val;
        } en_cols;
}
```

The **en_type** member contains a string that specifies the type of data this entry represents. The NIS+ server will compare this string to the type string specified in the table object and disallow any updates or modifications if they differ.

The **en_cols** structure contains two members: **en_cols_len** and **en_cols_val**. **en_cols_val** is an array of **entry_col** structures. **en_cols_len** contains a count of the number of cells in the **en_cols_val** array and reflects the number of columns in the table -- it always contains the same value as the **table_obj.ta_cols.ta_cols_len** member from the table which

contains the entry.

The **entry_col** structure contains information about the entry's per-column values. **ec_value** contains information about a particular value. It has two members: **ec_value_val**, which is the value itself, and **ec_value_len**, which is the length (in bytes) of the value. **entry_col** also contains the member **ec_flags**, which contains a set of flags for the entry.

The flags in **ec_flags** are primarily used when adding or modifying entries in a table. All columns that have the flag **EN_CRYPT** set will be encrypted prior to sending them over the network. Columns with **EN_BINARY** set are presumed to contain binary data. The server will ensure that the column in the table object specifies binary data prior to allowing the entry to be added. When modifying entries in a table, only those columns that have changed need be sent to the server. Those columns should each have the **EN_MODIFIED** flag set to indicate this to the server.

SEE ALSO

nis+(1), **nis_groups**(3N), **nis_names**(3N), **nis_server**(3N), **nis_subr**(3N), **nis_tables**(3N)

NAME

nis_ping, nis_checkpoint – misc NIS+ log administration functions

SYNOPSIS

cc [*flag* . . .] *file*. . . **–lnsl** [*library*. . .]

#include <rpcsvc/nis.h>

void nis_ping(const nis_name *dirname***, const u_long** *utime***, const nis_object** ∗*dirobj***);**

nis_result ∗**nis_checkpoint(const nis_name** *dirname***);**

MT-LEVEL

MT-Safe

DESCRIPTION

nis_ping( ) is called by the master server for a directory when a change has occurred within that directory. The parameter *dirname* identifies the directory with the change. If the parameter *dirobj* is **NULL**, this function looks up the directory object for *dirname* and uses the list of replicas it contains. The parameter *utime* contains the timestamp of the last change made to the directory. This timestamp is used by the replicas when retrieving updates made to the directory.

The effect of calling **nis_ping( )** is to schedule an update on the replica. A short time after a ping is received, typically about two minutes, the replica compares the last update time for its databases to the timestamp sent by the ping. If the ping timestamp is later, the replica establishes a connection with the master server and request all changes from the log that occurred after the last update that it had recorded in its local log.

nis_checkpoint( ) is used to force the service to checkpoint information that has been entered in the log but has not been checkpointed to disk. When called, this function checkpoints the database for each table in the directory, the database containing the directory and the transaction log. Care should be used in calling this function since directories that have seen a lot of changes may take several minutes to checkpoint. During the checkpointing process, the service will be unavailable for updates for all directories that are served by this machine as master.

nis_checkpoint( ) returns a pointer to a *nis_result* structure (described in **nis_tables**(3N)). This structure should be freed with **nis_freeresult( )** (see **nis_names**(3N)). The only items of interest in the returned result are the status value and the statistics.

SEE ALSO

nislog(1M), **nis_names**(3N), **nis_tables**(3N), **nisfiles**(4)

NAME

nis_server, nis_mkdir, nis_rmdir, nis_servstate, nis_stats, nis_getservlist, nis_freeservlist, nis_freetags – miscellaneous NIS+ functions

SYNOPSIS

cc [*flag* . . .] *file*. . . **–lnsl** [*library*. . .]

#include <rpcsvc/nis.h>

nis_error nis_mkdir(const nis_name *dirname,* **const nis_server** ∗*machine);*

nis_error nis_rmdir(const nis_name *dirname,* **const nis_server** ∗*machine);*

nis_error nis_servstate(const nis_server ∗*machine,* **const nis_tag** ∗*tags,*
const int *numtags***, nis_tag** ∗∗*result***);**

nis_error nis_stats(const nis_server ∗*machine***, const nis_tag** ∗*tags,* **const int** *numtags,*
nis_tag ∗∗*result***);**

void nis_freetags(nis_tag ∗*tags***, const int** *numtags***);**

nis_server ∗∗**nis_getservlist(const nis_name** *dirname***);**

void nis_freeservlist(nis_server ∗∗*machines***);**

MT-LEVEL

MT-Safe

DESCRIPTION

These functions provide a variety of services for NIS+ applications.

nis_mkdir( ) is used to create the necessary databases to support NIS+ service for a directory, *dirname*, on a server, *machine*. If this operation is successful, it means that the directory object describing *dirname* has been updated to reflect that server *machine* is serving the named directory. For a description of the **nis_server** structure, refer to **nis_objects**(3N).

nis_rmdir( ) is used to delete the directory, *dirname*, from the specified machine. The *machine* parameter cannot be **NULL**. For a description of the **nis_server** structure, refer to **nis_objects**(3N).

nis_servstate( ) is used to set and read the various state variables of the NIS+ servers. In particular the internal debugging state of the servers may be set and queried.

The **nis_stats( )** function is used to retrieve statistics about how the server is operating. Tracking these statistics can help administrators determine when they need to add additional replicas or to break up a domain into two or more subdomains. For more information on reading statistics, see **nisstat**(1M)**.**

nis_servstate( ) and **nis_stats( )** use the tag list. This tag list is a variable length array of *nis_tag* structures whose length is passed to the function in the *numtags* parameter. The set of legal tags are defined in the file **<rpcsvc/nis_tags.h>** which is included in **<rpcsvc/nis.h>**. Because these tags can and do vary between implementations of the NIS+ service, it is best to consult this file for the supported list. Passing unrecognized tags to a server will result in their *tag_value* member being set to the string ‘‘unknown.’’ Both of these functions return their results in malloced tag structure, ∗*result*. If there is an error, ∗*result* is set to **NULL**. The *tag_value* pointers points to allocated string memory which contains the results. Use **nis_freetags( )** to free the tag structure.

nis_getservlist( ) returns a null terminated list of *nis_server* structures that represent the list of servers that serve the domain named *dirname*. Servers from this list can be used when calling functions that require the name of a NIS+ server. For a description of the **nis_server** structure, refer to **nis_objects**(3N). **nis_freeservlist( )** frees the list of servers returned by **nis_getservlist( )**. Note that this is the only legal way to free that list.

SEE ALSO

nisstat(1M), **nis_names**(3N), **nis_objects**(3N), **nis_subr**(3N)

NAME

nis_subr, nis_leaf_of, nis_name_of, nis_domain_of, nis_getnames, nis_freenames, nis_dir_cmp, nis_clone_object, nis_destroy_object, nis_print_object – NIS+ subroutines

SYNOPSIS

cc [*flag* . . .] *file*. . . **-lnsl** [*library*. . .]

#include <rpcsvc/nis.h>

nis_name nis_leaf_of(const nis_name *name***);**

nis_name nis_name_of(const nis_name *name***);**

nis_name nis_domain_of(const nis_name *name***);**

nis_name ∗**nis_getnames(const nis_name** *name***);**

void nis_freenames(nis_name ∗*namelist***);**

name_pos nis_dir_cmp(const nis_name *n1***, const nis_name** *n2***);**

nis_object ∗**nis_clone_object(const nis_object** ∗*src***, nis_object** ∗*dest***);**

void nis_destroy_object(nis_object ∗*obj***);**

void nis_print_object(const nis_object ∗*obj***);**

MT-LEVEL

Safe

DESCRIPTION

These subroutines are provided to assist in the development of NIS+ applications. They provide several useful operations on both NIS+ names and objects.

The first group, **nis_leaf_of( )**, **nis_domain_of( )**, and **nis_name_of( )** provide the functions for parsing NIS+ names. **nis_leaf_of( )** will return the first label in an NIS+ name. It takes into account the double quote character '"' which can be used to protect embedded '.' (dot) characters in object names. Note that the name returned will never have a trailing dot character. If passed the global root directory name ".", it will return the null string.

nis_domain_of( ) returns the name of the NIS+ domain in which an object resides. This name will always be a fully qualified NIS+ name and ends with a dot. By iteratively calling **nis_leaf_of( )** and **nis_domain_of( )** it is possible to break a NIS+ name into its individual components.

nis_name_of( ) is used to extract the unique part of a NIS+ name. This function removes from the tail portion of the name all labels that are in common with the local domain. Thus if a machine were in domain **foo.bar.baz.** and **nis_name_of( )** were passed a name **bob.friends.foo.bar.baz**, then **nis_name_of( )** would return the unique part, **bob.friends**. If the name passed to this function is not in either the local domain or one of its children, this function will return null.

nis_getnames( ) will return a list of candidate names for the name passed in as *name.* If this name is not fully qualified, **nis_getnames( )** will generate a list of names using the default NIS+ directory search path, or the environment variable **NIS_PATH** if it is set. The returned array of pointers is terminated by a NULL pointer, and the memory associated with this array should be freed by calling **nis_freenames( ).**

Though **nis_dir_cmp( )** can be used to compare any two NIS+ names, it is used primarily to compare domain names. This comparison is done in a case independent fashion, and the results are an enum of type **name_pos**. When the names passed to this function are identical, the function returns a value of SAME_NAME. If the name *n1* is a direct ancestor of name *n2*, then this function returns the result HIGHER_NAME. Similarly, if the name *n1* is a direct descendant of name *n2*, then this function returns the result LOWER_NAME. When the name *n1* is neither a direct ancestor nor a direct descendant of *n2*, as it would be if the two names were siblings in separate portions of the namespace, then this function returns the result NOT_SEQUENTIAL. Finally, if either name cannot be parsed as a legitimate name then this function returns the value BAD_NAME.

The second set of functions, consisting of **nis_clone_object( )** and **nis_destroy_object( )**, are used for manipulating objects. **nis_clone_object( )** creates an exact duplicate of the NIS+ object *src*. If the value of *dest* is non-null, it creates the clone of the object into this object structure and allocate the necessary memory for the variable length arrays. If this parameter is null, a pointer to the cloned object is returned. Refer to **nis_objects**(3N) for a description of the **nis_object** structure.

nis_destroy_object( ) can be used to destroy an object created by **nis_clone_object( )**. This will free up all memory associated with the object and free the pointer passed. If the object was cloned into an array (using the *dest* parameter to **nis_clone_object( )**) then the object *cannot* be freed with this function. Instead, the function **xdr_free(xdr_nis_object, *dest*)** must be used.

nis_print_object( ) prints out the contents of a NIS+ object structure on the standard output. Its primary use is for debugging NIS+ programs.

ENVIRONMENT

NIS_PATH This variable overrides the default NIS+ directory search path used by **nis_getnames( ).** It contains an ordered list of directories separated by ':' (colon) characters. The '$' (dollar sign) character is treated specially. Directory names that end in '$' have the default domain appended to them, and a '$' by itself is replaced by the list of directories between the default domain and the global root that are at least two levels deep. The default NIS+ directory search path is '$'.

SEE ALSO

nis_names(3N), **nis_objects**(3N), **nis_tables**(3N)

NOTES

nis_leaf_of( ), **nis_name_of( )**and **nis_clone_object( )** return their results as thread-specific data in multithreaded applications.

NAME

nis_tables, nis_list, nis_add_entry, nis_remove_entry, nis_modify_entry, nis_first_entry, nis_next_entry – NIS+ table functions

SYNOPSIS

cc [*flag* . . .] *file*. . . **–lnsl** [*library*. . .]

#include <rpcsvc/nis.h>

nis_result ∗**nis_list(const nis_name** *name***, const u_long** *flags***,**
int (∗**callback)(const nis_name** *table_name***, const nis_object** ∗*object***,**
const void ∗*userdata***), const void** ∗*userdata***);**

nis_result ∗**nis_add_entry(const nis_name** *table_name***, const nis_object** ∗*object***,**
const u_long *flags***);**

nis_result ∗**nis_remove_entry(const nis_name** *name***, const nis_object** ∗*object***,**
const u_long *flags***);**

nis_result ∗**nis_modify_entry(const nis_name** *name***, const nis_object** ∗*object***,**
const u_long *flags***);**

nis_result ∗**nis_first_entry(const nis_name** *table_name***);**

nis_result ∗**nis_next_entry(const nis_name** *table_name***, const netobj** ∗*cookie***);**

void nis_freeresult(nis_result ∗*result***);**

MT-LEVEL

MT-Safe

DESCRIPTION

These functions are used to search and modify NIS+ tables. **nis_list( )** is used to search a table in the NIS+ namespace. **nis_first_entry( )** and **nis_next_entry( )** are used to enumerate a table one entry at a time. **nis_add_entry( )**, **nis_remove_entry( )**, and **nis_modify_entry( )** are used to change the information stored in a table. **nis_freeresult( )** is used to free the memory associated with the **nis_result** structure.

Entries within a table are named by NIS+ indexed names. An indexed name is a compound name that is composed of a search criteria and a simple NIS+ name that identifies a table object. A search criteria is a series of column names and their associated values enclosed in bracket '[ ]' characters. Indexed names have the following form:

[*colname***=***value***, . . .],***tablename*

The list function, **nis_list( )**, takes an indexed name as the value for the *name* parameter. Here, the tablename should be a fully qualified NIS+ name unless the **EXPAND_NAME** flag (described below) is set. The second parameter, *flags*, defines how the function will respond to various conditions. The value for this parameter is created by logically **OR**ing together one or more flags from the following list.

FOLLOW_LINKS
: If the table specified in *name* resolves to be a **LINK** type object (see **nis_objects**(3N)), this flag specifies that the client library follow that link and do the search at that object. If this flag is not set and the name resolves to a link, the error **NIS_NOTSEARCHABLE** will be returned.

FOLLOW_PATH This flag specifies that if the entry is not found within this table, the list

operation should follow the path specified in the table object. When used in conjunction with the **ALL_RESULTS** flag below, it specifies that the path should be followed regardless of the result of the search. When used in conjunction with the **FOLLOW_LINKS** flag above, named tables in the path that resolve to links will be followed until the table they point to is located. If a table in the path is not reachable because no server that serves it is available, the result of the operation will be either a "soft" success or a "soft" failure to indicate that not all tables in the path could be searched. If a name in the path names is either an invalid or non-existent object then it is silently ignored.

HARD_LOOKUP
This flag specifies that the operation should continue trying to contact a server of the named table until a definitive result is returned (such as NIS_NOTFOUND).

ALL_RESULTS This flag can only be used in conjunction with **FOLLOW_PATH** and a callback function. When specified, it forces all of the tables in the path to be searched. If *name* does not specify a search criteria (imply that all entries are to be returned), then this flag will cause all of the entries in all of the tables in the path to be returned.

NO_CACHE This flag specifies that the client library should bypass any client object caches and get its information directly from either the master server or a replica server for the named table.

MASTER_ONLY This flag is even stronger than **NO_CACHE** in that it specifies that the client library should *only* get its information from the master server for a particular table. This guarantees that the information will be up to date. However, there may be severe performance penalties associated with contacting the master server directly on large networks. When used in conjunction with the **HARD_LOOKUP** flag, this will block the list operation until the master server is up and available.

EXPAND_NAME
When specified, the client library will attempt to expand a partially qualified name by calling **nis_getnames( )** (see **nis_local_names**(3N)) which uses the environment variable **NIS_PATH.**

RETURN_RESULT
This flag is used to specify that a copy of the returning object be returned in the **nis_result** structure if the operation was successful.

The third parameter to **nis_list( )**, *callback*, is an optional pointer to a function that will process the **ENTRY** type objects that are returned from the search. If this pointer is **NULL**, then all entries that match the search criteria are returned in the *nis_result* structure, otherwise this function will be called once for each entry returned. When called, this function should return **0** when additional objects are desired and **1** when it no longer wishes to see any more objects. The fourth parameter, *userdata*, is simply passed to callback function along with the returned entry object. The client can use this pointer to pass

state information or other relevant data that the callback function might need to process the entries.

nis_add_entry( ) will add the NIS+ object to the NIS+ *table_name*. The *flags* parameter is used to specify the failure semantics for the add operation. The default (*flags* equal 0) is to fail if the entry being added already exists in the table. The **ADD_OVERWRITE** flag may be used to specify that existing object is to be overwritten if it exists, (a modify operation) or added if it does not exist. With the **ADD_OVERWRITE** flag, this function will fail with the error **NIS_PERMISSION** if the existing object does not allow modify privileges to the client.

If the flag **RETURN_RESULT** has been specified, the server will return a copy of the resulting object if the operation was successful.

nis_remove_entry( ) removes the identified entry from the table or a set of entries identified by *table_name*. If the parameter *object* is non-null, it is presumed to point to a cached copy of the entry. When the removal is attempted, and the object that would be removed is not the same as the cached object pointed to by *object* then the operation will fail with an **NIS_NOTSAMEOBJ** error. If an object is passed with this function, the search criteria in name is optional as it can be constructed from the values within the entry. However, if no object is present, the search criteria must be included in the *name* parameter. If the flags variable is null, and the search criteria does not uniquely identify an entry, the **NIS_NOTUNIQUE** error is returned and the operation is aborted. If the flag parameter **REM_MULTIPLE** is passed, and if remove permission is allowed for each of these objects, then all objects that match the search criteria will be removed. Note that a null search criteria and the **REM_MULTIPLE** flag will remove all entries in a table.

nis_modify_entry( ) modifies an object identified by *name*. The parameter *object* should point to an entry with the **EN_MODIFIED** flag set in each column that contains new information.

The owner, group, and access rights of an entry are modified by placing the modified information into the respective fields of the parameter, *object*: **zo_owner**, **zo_group**, and **zo_access**.

These columns will replace their counterparts in the entry that is stored in the table. The entry passed must have the same number of columns, same type, and valid data in the modified columns for this operation to succeed.

If the flags parameter contains the flag **MOD_SAMEOBJ** then the object pointed to by *object* is assumed to be a cached copy of the original object. If the OID of the object passed is different than the OID of the object the server fetches, then the operation fails with the **NIS_NOTSAMEOBJ** error. This can be used to implement a simple read-modify-write protocol which will fail if the object is modified before the client can write the object back.

If the flag **RETURN_RESULT** has been specified, the server will return a copy of the resulting object if the operation was successful.

nis_first_entry( ) fetches entries from a table one at a time. This mode of operation is extremely inefficient and callbacks should be used instead wherever possible. The table containing the entries of interest is identified by *name*. If a search criteria is present in *name* it is ignored. The value of *cookie* within the 34nis_result structure must be copied by

the caller into local storage and passed as an argument to **nis_next_entry( )**.

nis_next_entry( ) retrieves the "next" entry from a table specified by *table_name*. The order in which entries are returned is not guaranteed. Further, should an update occur in the table between client calls to **nis_next_entry( )** there is no guarantee that an entry that is added or modified will be seen by the client. Should an entry be removed from the table that would have been the "next" entry returned, the error **NIS_CHAINBROKEN** is returned instead.

RETURN VALUES

These functions return a pointer to a structure of type **nis_result**:

```
struct nis_result {
        nis_error   status;
        struct {
                    u_int       objects_len;
                    nis_object  *objects_val;
        } objects;
        netobj      cookie;
        u_long      zticks;
        u_long      dticks;
        u_long      aticks;
        u_long      cticks;
};
```

The *status* member contains the error status of the the operation. A text message that describes the error can be obtained by calling the function **nis_sperrno( )** (see **nis_error**(3N)).

The **objects** structure contains two members. *objects_val* is an array of *nis_object* structures; *objects_len* is the number of cells in the array. These objects will be freed by a call to **nis_freeresult( )** (see **nis_names**(3N)). If you need to keep a copy of one or more objects, they can be copied with the function **nis_clone_object( )** and freed with the function **nis_destroy_object( )** (see **nis_server**(3N)).

The various ticks contain details of where the time (in microseconds) was taken during a request. They can be used to tune one's data organization for faster access and to compare different database implementations (see **nis_db**(3N)).

zticks The time spent in the NIS+ service itself, this count starts when the server receives the request and stops when it sends the reply.

dticks The time spent in the database backend, this time is measured from the time a database call starts, until a result is returned. If the request results in multiple calls to the database, this is the sum of all the time spent in those calls.

aticks The time spent in any ''accelerators'' or caches. This includes the time required to locate the server needed to resolve the request.

cticks The total time spent in the request, this clock starts when you enter the client library and stops when a result is returned. By subtracting the sum of the other ticks values from this value you can obtain the local overhead of generating a NIS+ request.

Subtracting the value in *dticks* from the value in *zticks* will yield the time spent in the service code itself. Subtracting the sum of the values in *zticks* and *aticks* from the value in *cticks* will yield the time spent in the client library itself. Note: all of the tick times are measured in microseconds.

ERRORS

The client library can return a variety of error returns and diagnostics. The more salient ones are documented below.

NIS_BADATTRIBUTE The name of an attribute did not match up with a named column in the table, or the attribute did not have an associated value.

NIS_BADNAME The name passed to the function is not a legal NIS+ name.

NIS_BADREQUEST A problem was detected in the request structure passed to the client library.

NIS_CACHEEXPIRED The entry returned came from an object cache that has *expired*. This means that the time to live value has gone to zero and the entry may have changed. If the flag NO_CACHE was passed to the lookup function then the lookup function will retry the operation to get an unexpired copy of the object.

NIS_CBERROR An RPC error occurred on the server while it was calling back to the client. The transaction was aborted at that time and any unsent data was discarded.

NIS_CBRESULTS Even though the request was successful, all of the entries have been sent to your callback function and are thus not included in this result.

NIS_FOREIGNNS The name could not be completely resolved. When the name passed to the function would resolve in a namespace that is outside the NIS+ name tree, this error is returned with a NIS+ object of type **DIRECTORY**. The returned object contains the type of namespace and contact information for a server within that namespace.

NIS_INVALIDOBJ The object pointed to by *object* is not a valid NIS+ entry object for the given table. This could occur if it had a mismatched number of columns, or a different data type (for example, binary or text) than the associated column in the table.

NIS_LINKNAMEERROR The name passed resolved to a **LINK** type object and the contents of the object pointed to an invalid name.

NIS_MODFAIL The attempted modification failed for some reason.

NIS_NAMEEXISTS An attempt was made to add a name that already exists. To add the name, first remove the existing name and then add the new name or modify the existing named object.

NIS_NAMEUNREACHABLE
This soft error indicates that a server for the desired directory of the named table object could not be reached. This can occur when there is a network partition or the server has crashed. Attempting the operation again may succeed. See the **HARD_LOOKUP** flag.

NIS_NOCALLBACK The server was unable to contact the callback service on your machine. This results in no data being returned.

NIS_NOMEMORY Generally a fatal result. It means that the service ran out of heap space.

NIS_NOSUCHNAME This hard error indicates that the named directory of the table object does not exist. This occurs when the server that should be the parent of the server that serves the table, does not know about the directory in which the table resides.

NIS_NOSUCHTABLE The named table does not exist.

NIS_NOT_ME A request was made to a server that does not serve the given name. Normally this will not occur, however if you are not using the built in location mechanism for servers, you may see this if your mechanism is broken.

NIS_NOTFOUND No entries in the table matched the search criteria. If the search criteria was null (return all entries) then this result means that the table is empty and may safely be removed by calling the **nis_remove( )**.

If the FOLLOW_PATH flag was set, this error indicates that none of the tables in the path contain entries that match the search criteria.

NIS_NOTMASTER A change request was made to a server that serves the name, but it is not the master server. This can occur when a directory object changes and it specifies a new master server. Clients that have cached copies of the directory object in the **/var/nis/NIS_SHARED_DIRCACHE** file will need to have their cache managers restarted (use **nis_cachemgr -i**) to flush this cache.

NIS_NOTSAMEOBJ An attempt to remove an object from the namespace was aborted because the object that would have been removed was not the same object that was passed in the request.

NIS_NOTSEARCHABLE
The table name resolved to a NIS+ object that was not searchable.

NIS_PARTIAL This result is similar to **NIS_NOTFOUND** except that it means the request succeeded but resolved to zero entries. When this occurs, the server returns a copy of the table object instead of an entry so that the client may then process the path or implement some other

local policy.

NIS_RPCERROR This fatal error indicates the RPC subsystem failed in some way. Generally there will be a **syslog**(3) message indicating why the RPC request failed.

NIS_S_NOTFOUND The named entry does not exist in the table, however not all tables in the path could be searched, so the entry may exist in one of those tables.

NIS_S_SUCCESS Even though the request was successful, a table in the search path was not able to be searched, so the result may not be the same as the one you would have received if that table had been accessible.

NIS_SUCCESS The request was successful.

NIS_SYSTEMERROR Some form of generic system error occurred while attempting the request. Check the **syslog**(3) record for error messages from the server.

NIS_TOOMANYATTRS The search criteria passed to the server had more attributes than the table had searchable columns.

NIS_TRYAGAIN The server connected to was too busy to handle your request. **add_entry( )**, **remove_entry( )**, and **modify_entry( )** return this error when the master server is currently updating its internal state. It can be returned to **nis_list( )** when the function specifies a callback and the server does not have the resources to handle callbacks.

NIS_TYPEMISMATCH An attempt was made to add or modify an entry in a table, and the entry passed was of a different type than the table.

ENVIRONMENT

NIS_PATH When set, this variable is the search path used by **nis_list( )** if the flag **EXPAND_NAME** is set.

SEE ALSO

niscat(1), **niserror**(1), **nismatch**(1), **nis_cachemgr**(1M), **nis_error**(3N), **nis_local_names**(3N), **nis_names**(3N), **nis_objects**(3N), **syslog**(3)

WARNINGS

Use the flag **HARD_LOOKUP** carefully since it can cause the application to block indefinitely during a network partition.

NOTES

The path used when the flag **FOLLOW_PATH** is specified, is the one present in the *first* table searched. The path values in tables that are subsequently searched are ignored.

It is legal to call functions that would access the nameservice from within a list callback. However, calling a function that would itself use a callback, or calling **nis_list( )** with a callback from within a list callback function is not currently supported.

There are currently no known methods for **nis_first_entry( )** and **nis_next_entry( )** to get their answers from only the master server.

NAME

nl_langinfo – language information

SYNOPSIS

```
#include <nl_types.h>
#include <langinfo.h>

char *nl_langinfo(nl_item item);
```

MT-LEVEL

Safe with exceptions

DESCRIPTION

nl_langinfo( ) returns a pointer to a null-terminated string containing information relevant to a particular language or cultural area defined in the programs locale. The manifest constant names and values of *item* are defined by **<langinfo.h>**.

Since yes and no strings are implemented using **gettext**(3I), link with the **–lintl** library.

RETURN VALUES

If **setlocale**(3C) has not been called successfully, or if data for a supported language is either not available, or if *item* is not defined therein, then **nl_langinfo( )** returns a pointer to the corresponding string in the C locale. In all locales, **nl_langinfo( )** returns a pointer to an empty string if *item* contains an invalid setting.

EXAMPLES

For example:

```
nl_langinfo (ABDAY_1);
```

would return a pointer to the string "**Dim**" if the identified language was French and a French locale was correctly installed; or "**Sun**" if the identified language was English.

SEE ALSO

gettext(3I), **localeconv**(3C), **setlocale**(3C), **strftime**(3C), **langinfo**(5), **nl_types**(5)

WARNINGS

The array pointed to by the return value should not be modified by the program. Subsequent calls to **nl_langinfo( )** may overwrite the array.

This function is built upon the functions **localeconv( )**, **strftime( )**, and **gettext( )**. Where possible users are advised to use these interfaces to the required data instead of using calls to **nl_langinfo( )**.

NOTES

nl_langinfo can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

nlist – get entries from symbol table

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <nlist.h>

int nlist(*filename, nl***)**
char ∗*filename***;**
struct nlist ∗*nl***;**

DESCRIPTION

nlist( ) examines the symbol table from the executable image whose name is pointed to by *filename*, and selectively extracts a list of values and puts them in the array of **nlist** structures pointed to by *nl*. The name list pointed to by *nl* consists of an array of structures containing names, types and values. The **n_name** field of each such structure is taken to be a pointer to a character string representing a symbol name. The list is terminated by an entry with a **NULL** pointer (or a pointer to a **NULL** string) in the **n_name** field. For each entry in *nl*, if the named symbol is present in the executable image's symbol table, its value and type are placed in the **n_value** and **n_type** fields. If a symbol cannot be located, the corresponding **n_type** field of *nl* is set to zero.

RETURN VALUES

Upon normal completion, **nlist( )** returns the number of symbols that were not located in the symbol table. If an error occurs, **nlist( )** returns –1 and sets all of the **n_type** fields in members of the array pointed to by *nl* to zero.

SEE ALSO

nlist(3E), **a.out**(4)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

Only the **n_value** field is compatibly set. Other fields in the **nlist** structure are filled with the ELF (Executable and Linking Format) values (see **nlist**(3E) and **a.out**(4)).

NAME

nlist – get entries from name list

SYNOPSIS

cc [*flag* ...] *file* ... **–lelf** [*library* ...]

#include <nlist.h>

int nlist(const char ∗*filename***, struct nlist** ∗*nl***);**

MT-LEVEL

Safe

DESCRIPTION

nlist( ) examines the name list in the executable file whose name is pointed to by *filename*, and selectively extracts a list of values and puts them in the array of **nlist( )** structures pointed to by *nl*. The name list *nl* consists of an array of structures containing names of variables, types, and values. The list is terminated with a null name, that is, a null string is in the name position of the structure. Each variable name is looked up in the name list of the file. If the name is found, the type, value, storage class, and section number of the name are inserted in the other fields. The type field may be set to 0 if the file was not compiled with the **–g** option to **cc**(1B).

nlist( ) will always return the information for an external symbol of a given name if the name exists in the file. If an external symbol does not exist, and there is more than one symbol with the specified name in the file (such as static symbols defined in separate files), the values returned will be for the last occurrence of that name in the file. If the name is not found, all fields in the structure except **n_name** are set to 0.

This function is useful for examining the system name list kept in the file **/dev/ksyms**. In this way programs can obtain system addresses that are up to date.

RETURN VALUES

All value entries are set to 0 if the file cannot be read or if it does not contain a valid name list.

nlist( ) returns 0 on success, –1 on error.

SEE ALSO

cc(1B), **elf**(3E), **kvm_nlist**(3K), **kvm_open**(3K), **a.out**(4), **ksyms**(7D), **mem**(7D)

NAME

nlsgetcall – get client's data passed via the listener

SYNOPSIS

#include <sys/tiuser.h>

struct t_call ∗**nlsgetcall(int** *fildes***);**

MT-LEVEL

Unsafe

DESCRIPTION

nlsgetcall( ) allows server processes started by the listener process to access the client's **t_call** structure, that is, the *sndcall* argument of **t_connect**(3N).

The **t_call** structure returned by **nlsgetcall( )** can be released using **t_free**(3N).

nlsgetcall( ) returns the address of an allocated **t_call** structure or NULL if a **t_call** structure cannot be allocated. If the **t_alloc( )** succeeds, undefined environment variables are indicated by a negative *len* field in the appropriate **netbuf** structure. A *len* field of zero in the **netbuf** structure is valid and means that the original buffer in the listener's **t_call** structure was NULL.

WARNING

The *len* field in the **netbuf** structure is defined as being unsigned. In order to check for error returns, it should first be cast to an int.

The listener process limits the amount of user data (*udata*) and options data (*opt*) to 128 bytes each. Address data *addr* is limited to 64 bytes. If the original data was longer, no indication of overflow is given.

RETURN VALUES

A NULL pointer is returned if a **t_call** structure cannot be allocated by **t_alloc( )**. **t_errno** can be inspected for further error information. Undefined environment variables are indicated by a negative length field (*len*) in the appropriate **netbuf** structure.

FILES

/usr/lib/libnsl_s.a
/usr/lib/libslan.a
/usr/lib/libnls.a

SEE ALSO

nlsadmin(1M), **getenv**(3C), **t_alloc**(3N), **t_connect**(3N), **t_error**(3N), **t_free**(3N), **t_sync**(3N)

NOTES

Server processes must call **t_sync**(3N) before calling this routine.

This interface is unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

nlsprovider – get name of transport provider

SYNOPSIS

char *nlsprovider(void);

MT-LEVEL

Unsafe

DESCRIPTION

nlsprovider( ) returns a pointer to a null terminated character string which contains the name of the transport provider as placed in the environment by the listener process. If the variable is not defined in the environment, a NULL pointer is returned.

The environment variable is only available to server processes started by the listener process.

RETURN VALUES

If the variable is not defined in the environment, a NULL pointer is returned.

FILES

/usr/lib/libslan.a (7300)
/usr/lib/libnls.a (3B2 Computer)
/usr/lib/libnsl_s.a

SEE ALSO

nlsadmin(1M)

NOTES

This interface is unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

nlsrequest – format and send listener service request message

SYNOPSIS

#include <listen.h>

int nlsrequest(int *fildes***, char** ∗*service_code***);**

extern int _nlslog, t_errno;
extern char ∗**_nlsrmsg;**

MT-LEVEL

Unsafe

DESCRIPTION

Given a virtual circuit to a listener process (*fildes*) and a service code of a server process, **nlsrequest( )** formats and sends a *service request message* to the remote listener process requesting that it start the given service. **nlsrequest( )** waits for the remote listener process to return a *service request response message*, which is made available to the caller in the static, null terminated data buffer pointed to by **_nlsrmsg**. The *service request response message* includes a success or failure code and a text message. The entire message is printable.

RETURN VALUES

The success or failure code is the integer return code from **nlsrequest( )**. Zero indicates success, other negative values indicate **nlsrequest( )** failures as follows:

–1 Error encountered by **nlsrequest( )**, see **t_errno**.

Positive values are error return codes from the *listener* process. Mnemonics for these codes are defined in **<listen.h>**.

2 Request message not interpretable.

3 Request service code unknown.

4 Service code known, but currently disabled.

If non-null, **_nlsrmsg** contains a pointer to a static, null terminated character buffer containing the *service request response message.* Note that both **_nlsrmsg** and the data buffer are overwritten by each call to **nlsrequest( )**.

If **_nlslog** is non-zero, **nlsrequest( )** prints error messages on stderr. Initially, **_nlslog** is zero.

FILES

/usr/lib/libnls.a
/usr/lib/libslan.a
/usr/lib/libnsl_s.a

SEE ALSO

nlsadmin(1M), **t_error**(3N)

WARNINGS

nlsrequest( ) cannot always be certain that the remote server process has been successfully started. In this case, **nlsrequest( )** returns with no indication of an error and the caller will receive notification of a disconnect event via a **T_LOOK** error before or during the first **t_snd( )** or **t_rcv( )** call.

NOTES

These interfaces are unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

offsetof – offset of structure member

SYNOPSIS

```
#include <stddef.h>
size_t offsetof(type, member-designator);
```

MT-LEVEL

MT-Safe

DESCRIPTION

offsetof( ) is a macro defined in **<stddef.h>** which expands to an integral constant expression that has type **size_t**, the value of which is the offset in bytes, to the structure member (designated by *member-designator*), from the beginning of its structure (designated by *type*).

NAME

p2open, p2close – open, close pipes to and from a command

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

int p2open(const char ∗*cmd***, FILE** ∗*fp***[2]);**

int p2close(FILE ∗*fp***[2]);**

MT-LEVEL

Unsafe

DESCRIPTION

p2open( ) forks and execs a shell running the command line pointed to by *cmd*. On return, **fp[0]** points to a **FILE** pointer to write the command's standard input and **fp[1]** points to a **FILE** pointer to read from the command's standard output. In this way the program has control over the input and output of the command.

The function returns **0** if successful; otherwise, it returns **–1**.

p2close( ) is used to close the file pointers that **p2open( )** opened. It waits for the process to terminate and returns the process status. It returns **0** if successful; otherwise, it returns **–1**.

RETURN VALUES

A common problem is having too few file descriptors. **p2close( )** returns **–1** if the two file pointers are not from the same **p2open( )**.

EXAMPLES

```
#include <stdio.h>
#include <libgen.h>

main(argc,argv)
int argc;
char **argv;
{
        FILE *fp[2];
        pid_t pid;
        char buf[16];

        pid=p2open("/usr/bin/cat", fp);
        if ( pid == –1 ) {
                fprintf(stderr, "p2open failed\n");
                exit(1);
        }
        write(fileno(fp[0]),"This is a test\n", 16);
        if(read(fileno(fp[1]), buf, 16) <=0)
                fprintf(stderr, "p2open failed\n");
        else
                write(1, buf, 16);
        (void)p2close(fp);
}
```

SEE ALSO **fclose**(3S), **popen**(3S), **setbuf**(3S)

NOTES Buffered writes on **fp[0]** can make it appear that the command is not listening. Judiciously placed **fflush( )** calls or unbuffering **fp[0]** can be a big help; see **fclose**(3S).

Many commands use buffered output when connected to a pipe. That, too, can make it appear as if things are not working.

Usage is not the same as for **popen( )**, although it is closely related.

NAME

panel_above, panel_below – panels deck traversal primitives

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lpanel -lcurses** [*library* . .]

#include <panel.h>

PANEL **panel_above(**PANEL** *panel**);**

PANEL **panel_below(**PANEL** *panel**);**

MT-LEVEL

Unsafe

DESCRIPTION

panel_above( ) returns a pointer to the panel just above *panel*, or NULL if *panel* is the top panel. **panel_below( )** returns a pointer to the panel just below *panel*, or NULL if *panel* is the bottom panel.

If NULL is passed for *panel*, **panel_above( )** returns a pointer to the bottom panel in the deck, and **panel_below( )** returns a pointer to the top panel in the deck.

RETURN VALUES

NULL is returned if an error occurs.

SEE ALSO

curses(3X), **panels**(3X)

NOTES

These routines allow traversal of the deck of currently visible panels.

The header **<panel.h>** automatically includes the header **<curses.h>**.

NAME

panel_move, move_panel – move a panels window on the virtual screen

SYNOPSIS

cc [*flag* ...] *file* ... –**lpanel -lcurses** [*library* ..]

#include <panel.h>

int move_panel(PANEL **panel***, int** *starty***, int** *startx***);**

MT-LEVEL

Unsafe

DESCRIPTION

move_panel() moves the **curses** window associated with *panel* so that its upper left-hand corner is at *starty*, *startx*. See usage note, below.

RETURN VALUES

OK is returned if the routine completes successfully, otherwise **ERR** is returned.

SEE ALSO

curses(3X), **panel_update**(3X), **panels**(3X)

NOTES

For **panels** windows, use **move_panel()** instead of the **mvwin() curses** routine. Otherwise, **update_panels()** will not properly update the virtual screen.

The header **<panel.h>** automatically includes the header **<curses.h>**.

NAME

panel_new, new_panel, del_panel – create and destroy panels

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lpanel -lcurses** [*library* . .]

#include <panel.h>

PANEL ∗**new_panel(WINDOW** ∗*win***);**

int del_panel(PANEL ∗ *panel***);**

MT-LEVEL

Unsafe

DESCRIPTION

new_panel( ) creates a new panel associated with *win* and returns the panel pointer. The new panel is placed on top of the panel deck.

del_panel( ) destroys *panel*, but not its associated window.

RETURN VALUES

new_panel( ) returns **NULL** if an error occurs.

del_win( ) returns **OK** if successful, **ERR** otherwise.

SEE ALSO

curses(3X), **panel_update**(3X), **panels**(3X)

NOTES

The header **<panel.h>** automatically includes the header **<curses.h>**.

NAME

panel_show, show_panel, hide_panel, panel_hidden – panels deck manipulation routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lpanel -lcurses** [*library* . .]

#include <panel.h>

int show_panel(PANEL **panel***);**

int hide_panel(PANEL **panel***);**

int panel_hidden(PANEL **panel***);**

MT-LEVEL

Unsafe

DESCRIPTION

show_panel( ) makes *panel*, previously hidden, visible and places it on top of the deck of panels.

hide_panel( ) removes *panel* from the panel deck and, thus, hides it from view. The internal data structure of the panel is retained.

panel_hidden( ) returns **TRUE (1) or FALSE (0)** indicating whether or not *panel* is in the deck of panels.

RETURN VALUES

show_panel( ) and **hide_panel( )** return the integer **OK** upon successful completion or **ERR** upon error.

SEE ALSO

curses(3X), **panel_update**(3X), **panels**(3X)

NOTES

The header **<panel.h>** automatically includes the header **<curses.h>**.

NAME

panel_top, top_panel, bottom_panel – panels deck manipulation routines

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lpanel -lcurses** [*library* . .]

#include <panel.h>

int top_panel(PANEL **panel***);**

int bottom_panel(PANEL **panel***);**

MT-LEVEL

Unsafe

DESCRIPTION

top_panel( ) pulls *panel* to the top of the desk of panels. It leaves the size, location, and contents of its associated window unchanged.

bottom_panel( ) puts *panel* at the bottom of the deck of panels. It leaves the size, location, and contents of its associated window unchanged.

RETURN VALUES

All of these routines return the integer **OK** upon successful completion or **ERR** upon error.

SEE ALSO

curses(3X), **panel_update**(3X), **panels**(3X)

NOTES

The header **<panel.h>** automatically includes the header **<curses.h>**.

NAME

panel_update, update_panels – panels virtual screen refresh routine

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lpanel -lcurses** [*library* . .]

#include <panel.h>

void update_panels(void);

MT-LEVEL

Unsafe

DESCRIPTION

update_panels( ) refreshes the virtual screen to reflect the depth relationships between the panels in the deck. The user must use the curses library call **doupdate( )** (see **curs_refresh**(3X)) to refresh the physical screen.

SEE ALSO

curs_refresh(3X), **curses**(3X), **panels**(3X)

NOTES

The header **<panel.h>** automatically includes the header **<curses.h>**.

NAME

panel_userptr, set_panel_userptr – associate application data with a panels panel

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lpanel -lcurses** [*library* . .]
#include <panel.h>
int set_panel_userptr(PANEL **panel***, char** **ptr***);**
char * **panel_userptr(PANEL** **panel***);**

MT-LEVEL

Unsafe

DESCRIPTION

Each panel has a user pointer available for maintaining relevant information.
set_panel_userptr( ) sets the user pointer of *panel* to *ptr*.
panel_userptr( ) returns the user pointer of *panel*.

RETURN VALUES

set_panel_userptr returns **OK** if successful, **ERR** otherwise.
panel_userptr returns **NULL** if there is no user pointer assigned to *panel*.

SEE ALSO

curses(3X), **panels**(3X)

NOTES

The header **<panel.h>** automatically includes the header **<curses.h>**.

NAME

panel_window, replace_panel – get or set the current window of a panels panel

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lpanel -lcurses** [*library* . .]
#include <panel.h>
WINDOW ∗**panel_window(PANEL** ∗*panel***);**
int replace_panel(PANEL ∗*panel***, WINDOW** ∗*win***);**

MT-LEVEL

Unsafe

DESCRIPTION

panel_window( ) returns a pointer to the window of *panel*.
replace_panel( ) replaces the current window of *panel* with *win*.

RETURN VALUES

panel_window( ) returns **NULL** on failure.
replace_panel( ) returns **OK** on successful completion, **ERR** otherwise.

SEE ALSO

curses(3X), **panels**(3X)

NOTES

The header **<panel.h>** automatically includes the header **<curses.h>**.

NAME

panels – character based panels package

SYNOPSIS

#include <panel.h>

MT-LEVEL

Unsafe

DESCRIPTION

The **panel** library is built using the **curses** library, and any program using **panels** routines must call one of the **curses** initialization routines such as **initscr**. A program using these routines must be compiled with **–lpanel** and **–lcurses** on the **cc** command line.

The **panels** package gives the applications programmer a way to have depth relationships between **curses** windows; a **curses** window is associated with every panel. The **panels** routines allow **curses** windows to overlap without making visible the overlapped portions of underlying windows. The initial **curses** window, **stdscr**, lies beneath all panels. The set of currently visible panels is the *deck* of panels.

The **panels** package allows the applications programmer to create panels, fetch and set their associated windows, shuffle panels in the deck, and manipulate panels in other ways.

Routine Name Index

The following table lists each **panels** routine and the name of the manual page on which it is described.

panels Routine Name	Manual Page Name
bottom_panel	panel_top(3X)
del_panel	panel_new(3X)
hide_panel	panel_show(3X)
move_panel	panel_move(3X)
new_panel	panel_new(3X)
panel_above	panel_above(3X)
panel_below	panel_above(3X)
panel_hidden	panel_show(3X)
panel_userptr	panel_userptr(3X)
panel_window	panel_window(3X)
replace_panel	panel_window(3X)
set_panel_userptr	panel_userptr(3X)
show_panel	panel_show(3X)
top_panel	panel_top(3X)
update_panels	panel_update(3X)

RETURN VALUES

Each **panels** routine that returns a pointer to an object returns NULL if an error occurs. Each panel routine that returns an integer, returns **OK** if it executes successfully and **ERR** if it does not.

SEE ALSO

curses(3X), and 3X pages whose names begin “panel_” for detailed routine descriptions.

NOTES

The header **<panel.h>** automatically includes the header **<curses.h>**.

NAME

pathfind – search for named file in named directories

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

char ∗**pathfind(const char** ∗*path***, const char** ∗*name***, const char** ∗*mode***);**

MT-LEVEL

MT-Safe

DESCRIPTION

pathfind( ) searches the directories named in *path* for the file *name*. The directories named in *path* are separated by semicolons. *mode* is a string of option letters chosen from the set [**rwxfbcdpugks**]:

Letter	Meaning
r	readable
w	writable
x	executable
f	normal file
b	block special
c	character special
d	directory
p	FIFO (pipe)
u	set user ID bit
g	set group ID bit
k	sticky bit
s	size nonzero

Options read, write, and execute are checked relative to the real (not the effective) user ID and group ID of the current process.

If the file *name*, with all the characteristics specified by *mode*, is found in any of the directories specified by *path*, then **pathfind( )** returns a pointer to a string containing the member of *path*, followed by a slash character (*/*), followed by *name*.

If *name* begins with a slash, it is treated as an absolute path name, and *path* is ignored.

An empty *path* member is treated as the current directory. **/** is not prepended at the occurrence of the first match; rather, the unadorned *name* is returned.

EXAMPLES

To find the **ls** command using the **PATH** environment variable:

```
pathfind (getenv ("PATH"), "ls", "rx")
```

RETURN VALUES

If no match is found, **pathname** returns a null pointer, **((char *) 0)**.

SEE ALSO

sh(1), **test**(1), **access**(2), **mknod**(2), **stat**(2), **getenv**(3C)

NOTES

The string pointed to by the returned pointer is stored in an area that is reused on subsequent calls to **pathfind( )**. The string should not be deallocated by the caller.

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

perror, errno – print system error messages

SYNOPSIS

```
#include <stdio.h>
void perror(const char *s);

#include <errno.h>
int errno;
```

MT-LEVEL

MT-Safe

DESCRIPTION

perror( ) produces a message on the standard error output (file descriptor 2), describing the last error encountered during a call to a system or library function. The argument string *s* is printed first, then a colon and a blank, then the message and a newline. (However, if *s* is a null pointer or points to a null string, the colon is not printed.) To be of most use, the argument string should include the name of the program that incurred the error. The error number is taken from the external variable **errno**,
(see **intro**(2)), which is set when errors occur but not cleared when non-erroneous calls are made.

SEE ALSO

intro(2), **fmtmsg**(3C), **gettext**(3I), **setlocale**(3C), **strerror**(3C)

NOTES

If the application is linked with –**lintl**, then messages printed from this function are in the native language specified by the **LC_MESSAGES** locale category; see **setlocale**(3C).

NAME

pfmt – display error message in standard format

SYNOPSIS

```
#include <pfmt.h>
int pfmt(FILE *stream, long flags, char *format, ... /* arg */);
```

MT-LEVEL

MT-safe

DESCRIPTION

pfmt() retrieves a format string from a locale-specific message database (unless **MM_NOGET** is specified) and uses it for **printf()** style formatting of *args*. The output is displayed on *stream*.

pfmt() encapsulates the output in the standard error message format (unless **MM_NOSTD** is specified, in which case the output is simply **printf()** like).

If the **printf()** format string is to be retrieved from a message database, the *format* argument must have the following structure:

<*catalog*>:<*msgnum*>:<*defmsg*>.

If **MM_NOGET** is specified, only the <*defmsg*> part must be specified.

<*catalog*> is used to indicate the message database that contains the localized version of the format string. <*catalog*> must be limited to 14 characters. These characters must be selected from a set of all characters values, excluding **\0** (null) and the ASCII codes for **/** (slash) and **:** (colon).

<*msgnum*> is a positive number that indicates the index of the string into the message database.

If the catalog does not exist in the locale (specified by the last call to **setlocale()** using the **LC_ALL** or **LC_MESSAGES** categories), or if the message number is out of bound, **pfmt()** will attempt to retrieve the message from the C locale. If this second retrieval fails, **pfmt()** uses the <*defmsg*> part of the *format* argument.

If <*catalog*> is omitted, **pfmt()** will attempt to retrieve the string from the default catalog specified by the last call to **setcat()**. In this case, the *format* argument has the following structure:

:<*msgnum*>:<*defmsg*>.

pfmt() will output **Message not found!!\n** as format string if <*catalog*> is not a valid catalog name, if no catalog is specified (either explicitely or via **setcat()**), if <*msgnum*> is not a valid number, or if no message could be retrieved from the message databases, and <*defmsg*> was omitted.

The *flags* determine the type of output (i.e. whether the *format* should be interpreted as is or encapsulated in the standard message format), and the access to message catalogs to retrieve a localized version of *format*.

The *flags* are composed of several groups, and can take the following values (one from each group): *Output format control*

MM_NOSTD Do not use the standard message format, interpret *format* as a **printf()** *format*. Only *catalog access control flags* should be specified if **MM_NOSTD** is used; all other flags will be ignored

MM_STD Output using the standard message format (default, value 0).

Catalog access control

MM_NOGET Do not retrieve a localized version of *format*. In this case, only the <*defmsg*> part of the *format* is specified.

MM_GET Retrieve a localized version of *format*, from the <*catalog*>, using <*msgid*> as the index and <*defmsg*> as the default message (default, value 0).

Severity (standard message format only)

MM_HALT generates a localized version of **HALT**, but does not halt the machine.

MM_ERROR generates a localized version of **ERROR** (default, value 0).

MM_WARNING generates a localized version of **WARNING**.

MM_INFO generates a localized version of **INFO**.

Additional severities can be defined. Add-on severities can be defined with number-string pairs with numeric values from the range [5-255], using **addsev()**. The numeric value ORed with other *flags* will generate the specified severity.

If the severity is not defined, **pfmt()** used the string **SEV=***N* where *N* is replaced by the integer severity value passed in *flags*.

Multiple severities passed inf *flags* will not be detected as an error. Any combination of severities will be summed and the numeric value will cause the display of either a severity string (if defined) or the string **SEV=***N* (if undefined).

Action

MM_ACTION specifies an action message. Any severity value is superseded and replaced by a localized version of **TO FIX**.

STANDARD ERROR MESSAGE FORMAT

pfmt() displays error messages in the following format:

*label***:** *severity***:** *text*

If no *label* was defined by a call to **setlabel()**, the message is displayed in the format:

*severity***:** *text*

If **pfmt()** is called twice to display an error message and a helpful *action* or recovery

message, the output can look like:

label: *severity*: *text*
label: **TO FIX:** *text*

RETURN VALUE

Upon success, **pfmt()** returns the number of bytes transmitted. Upon failure, it returns a negative value:

-1 write error to *stream*.

EXAMPLES

Example 1:

```
setlabel("UX:test");
pfmt(stderr, MM_ERROR, "test:2:Cannot open file: %s\n", strerror(errno));
```

displays the message:

```
UX:test: ERROR: Cannot open file: No such file or directory
```

Example 2:

```
setlabel("UX:test");
setcat("test");
pfmt(stderr, MM_ERROR, ":10:Syntax error\n");
pfmt(stderr, MM_ACTION, "55:Usage ...\n");
```

displays the message

```
UX:test: ERROR: Syntax error
UX:test: TO FIX: Usage ...
```

NOTES

pfmt() uses **gettxt**(3C), it is recommended that **pfmt()** not be used.

SEE ALSO

addsev(3C), **gettxt**(3C), **lfmt**(3C), **printf**(3S), **setcat**(3C), **setlabel**(3C), **setlocale**(3C), **environ**(5)

NAME

plock – lock or unlock into memory process, text, or data

SYNOPSIS

```
#include <sys/lock.h>
int plock(int op);
```

DESCRIPTION

plock( ) allows the calling process to lock or unlock into memory its text segment (text lock), its data segment (data lock), or both its text and data segments (process lock). Locked segments are immune to all routine swapping. The effective user ID of the calling process must be super-user to use this call. **plock( )** performs the function specified by *op*:

PROCLOCK	Lock text and data segments into memory (process lock).
TXTLOCK	Lock text segment into memory (text lock).
DATLOCK	Lock data segment into memory (data lock).
UNLOCK	Remove locks.

RETURN VALUES

Upon successful completion, a value of 0 is returned to the calling process. Otherwise, a value of –1 is returned and **errno** is set to indicate the error.

ERRORS

plock( ) fails and does not perform the requested operation if one or more of the following are true:

EAGAIN	Not enough memory.
EINVAL	*op* is equal to **PROCLOCK** and a process lock, a text lock, or a data lock already exists on the calling process.
EINVAL	*op* is equal to **TXTLOCK** and a text lock, or a process lock already exists on the calling process.
EINVAL	*op* is equal to **DATLOCK** and a data lock, or a process lock already exists on the calling process.
EINVAL	*op* is equal to **UNLOCK** and no lock exists on the calling process.
EPERM	The effective user of the calling process is not super-user.

SEE ALSO

exec(2), **exit**(2), **fork**(2), **memcntl**(2), **mlock**(3C), **mlockall**(3C)

NOTES

mlock(3C) and **mlockall**(3C) are the preferred interfaces to process locking.

NAME

plot, arc, box, circle, closepl, closevt, cont, erase, label, line, linmod, move, openpl, openvt, point, space – graphics interface

SYNOPSIS

void arc(short *x0***, short** *y0***, short** *x1* **, short** *y1***, short** *x2***, short** *y2* **);**

void box(short *x0***, short** *y0***, short** *x1* **, short** *y1***);**

void circle(short *x***, short** *y* **, short** *r***);**

void closepl( );

void closevt( );

void cont(short *x***, short** *y***);**

void erase( );

void label(char ∗*s***);**

void line(short *x0***, short** *y0***, short** *x1* **, short** *y1***);**

void linmod(char ∗*s***);**

void move(short *x***, short** *y***);**

void openpl( );

void openvt( );

void point(short *x***, short** *y***);**

void space(short *x0***, short** *y0***, short** *x1* **, short** *y1***);**

MT-LEVEL

Safe

DESCRIPTION

These routines generate graphics output for a set of output devices. The format of the output is dependent upon which link editor option is used when the program is compiled and linked (see **Link Editor**).

The term "current point" refers to the current setting for the *x* and *y* coordinates.

arc() specifies a circular arc. The coordinates **(***x0***,** *y0)* specify the center of the arc. The coordinates **(***x1***,** *y1)* specify the starting point of the arc. The coordinates **(***x2***,** *y2)* specify the end point of the circular arc.

box() specifies a rectangle with coordinates **(***x0***,** *y0***)**, **(***x0***,** *y1***)**, **(***x1***,** *y0***)**, and **(***x1***,** *y1***)**. The current point is set to **(***x1***,** *y1***)**.

circle() specifies a circle with a center at the coordinates **(***x***,** *y***)** and a radius of *r*.

closevt() and **closepl()** flush the output.

cont() specifies a line beginning at the current point and ending at the coordinates **(***x***,** *y***)**. The current point is set to **(***x***,** *y***)**.

erase() starts another frame of output.

label() places the null terminated string *s* so that the first character falls on the current point. The string is then terminated by a NEWLINE character.

line() draws a line starting at the coordinates **(***x0***,** *y0***)** and ending at the coordinates **(***x1***,** *y1***)**. The current point is set to **(***x1***,** *y1***)**.

linmod() specifies the style for drawing future lines. *s* may contain one of the following: **dotted**, **solid**, **longdashed**, **shortdashed**, or **dotdashed**.

move() sets the current point to the coordinates **(***x***,** *y***)**.

openpl() or **openvt()** must be called to open the device before any other **plot** routines are called.

point() plots the point given by the coordinates **(***x***,** *y***)**. The current point is set to **(***x***,** *y***)**.

space() specifies the size of the plotting area. The plot will be reduced or enlarged as necessary to fit the area specified. The coordinates **(***x0***,** *y0***)** specify the lower left hand corner of the plotting area. The coordinates **(***x1***,** *y1***)** specify the upper right hand corner of the plotting area.

Link Editor

Various flavors of these routines exist for different output devices. They are obtained by using the following **ld**(1) options:

-lplot	device-independent graphics stream on standard output in the format described in **plot**(4B)
-l300	GSI 300 terminal
-l300s	GSI 300S terminal
-l4014	Tektronix 4014 terminal
-l450	GSI 450 terminal
-lvt0	

FILES

/usr/lib/libplot.a
/usr/lib/lib300.a
/usr/lib/lib300s.a
/usr/lib/lib4014.a
/usr/lib/lib450.a
/usr/lib/libvt0.a

SEE ALSO

graph(1), **ld**(1), **plot**(4B)

NAME

popen, pclose – initiate pipe to/from a process

SYNOPSIS

```
#include <stdio.h>
FILE *popen(const char *command, const char *type);
int pclose (FILE *stream);
```

MT-LEVEL

Unsafe

DESCRIPTION

popen() creates a pipe between the calling program and the command to be executed. The arguments to **popen()** are pointers to null-terminated strings. *command* consists of a shell command line. *type* is an I/O mode, either **r** for reading or **w** for writing. The value returned is a stream pointer such that one can write to the standard input of the command, if the I/O mode is **w**, by writing to the file *stream* (see **intro**(3)); and one can read from the standard output of the command, if the I/O mode is **r**, by reading from the file *stream*. Because open files are shared, a type **r** command may be used as an input filter and a type **w** as an output filter.

The environment of the executed command will be as if a child process were created within the **popen()** call using **fork**(2), and the child invoked the shell using the call:

Solaris **execl("/usr/bin/sh", "sh", "-c", *command*, (char *)0);**
XPG4 **execl("/usr/bin/ksh", "ksh", "-c", *command*, (char *)0);**

A stream opened by **popen()** should be closed by **pclose(),** which closes the pipe, and waits for the associated process to terminate and returns the termination status of the process running the command language interpreter. This is the value returned by **wait-pid**(2). See **wstat**(5) for more details on termination status.

RETURN VALUES

popen() returns a null pointer if files or processes cannot be created.

pclose() returns the termination status of the command. **pclose()** returns **-1** if *stream* is not associated with a **popen()** command and sets **errno** to indicate the error.

EXAMPLES

The following is an example of a typical call:

```
#include <stdio.h>
#include <stdlib.h>
main()
{
        char *cmd = "/usr/bin/ls *.c";
        char buf[BUFSIZ];
        FILE *ptr;

        if ((ptr = popen(cmd, "r")) != NULL)
                while (fgets(buf, BUFSIZ, ptr) != NULL)
                        (void) printf("%s", buf);
        return 0;
}
```

This program will print on the standard output (see **stdio**(3S)) all the file names in the current directory that have a **.c** suffix.

SEE ALSO

ksh(1), **pipe**(2), **wait**(2), **waitpid**(2), **fclose**(3S), **fopen**(3S), **stdio**(3S), **system**(3S), **wstat**(5), **xpg4**(5)

NOTES

If the original and **popen()** processes concurrently read or write a common file, neither should use buffered I/O. Problems with an output filter may be forestalled by careful buffer flushing, for example, with **fflush()** (see **fclose**(3S)). A security hole exists through the **IFS** and **PATH** environment variables. Full pathnames should be used (or **PATH** reset) and **IFS** should be set to space and tab (" **t**").

NAME

printf, fprintf, sprintf, vprintf, vfprintf, vsprintf – formatted output conversion

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <stdio.h>

int printf(*format***, ...)**
const char **format*;

int fprintf(*stream, format, va_list***)**
FILE **stream*;
char **format*;
va_dcl;

char ***sprintf(** *s, format, va_list***)**
char **s*, **format*;
va_dcl;

int vprintf(*format, ap***)**
char **format*;
va_list *ap*;

int vfprintf(*stream, format, ap***)**
FILE **stream*;
char **format*;
va_list *ap*;

char ***vsprintf(** *s, format, ap***)**
char **s*,**format*;
va_list *ap*;

DESCRIPTION

printf() places output on the standard output stream **stdout**. **fprintf()** places output on the named output *stream*. **sprintf()** places "output," followed by the **NULL** character (**\0**), in consecutive bytes starting at **s*; it is the user's responsibility to ensure that enough storage is available.

vprintf(), **vfprintf()**, and **vsprintf()** are the same as **printf()**, **fprintf()**, and **sprintf()** respectively, except that instead of being called with a variable number of arguments, they are called with an argument list as defined by **varargs**(5).

Each of these functions converts, formats, and prints its *arg*s under control of the *format*. The *format* is a character string which contains two types of objects: plain characters, which are simply copied to the output stream, and conversion specifications, each of which causes conversion and printing of zero or more *arg*s. The results are undefined if there are insufficient *arg*s for the format. If the format is exhausted while *arg*s remain, the excess *arg*s are simply ignored.

Each conversion specification is introduced by the character **%**. After the **%**, the following appear in sequence:

Zero or more *flags*, which modify the meaning of the conversion specification.

An optional decimal digit string specifying a minimum *field width*. If the converted value has fewer characters than the field width, it will be padded on the left (or right, if the left-adjustment flag '–', described below, has been given) to the field width. The padding is with blanks unless the field width digit string starts with a zero, in which case the padding is with zeros.

A *precision* that gives the minimum number of digits to appear for the **d**, **i**, **o**, **u**, **x**, or **X** conversions, the number of digits to appear after the decimal point for the **e**, **E**, and **f** conversions, the maximum number of significant digits for the **g** and **G** conversion, or the maximum number of characters to be printed from a string in **s** conversion. The precision takes the form of a period (**.**) followed by a decimal digit string; a **NULL** digit string is treated as zero. Padding specified by the precision overrides the padding specified by the field width.

An optional **l** (ell) specifying that a following **d**, **i**, **o**, **u**, **x**, or **X** conversion character applies to a long integer *arg*. An **l** before any other conversion character is ignored.

A character that indicates the type of conversion to be applied.

A field width or precision or both may be indicated by an asterisk (*****) instead of a digit string. In this case, an integer *arg* supplies the field width or precision. The *arg* that is actually converted is not fetched until the conversion letter is seen, so the *arg*s specifying field width or precision must appear *before* the *arg* (if any) to be converted. A negative field width argument is taken as a '–' flag followed by a positive field width. If the precision argument is negative, it will be changed to zero.

The flag characters and their meanings are:

– The result of the conversion will be left-justified within the field.

+ The result of a signed conversion will always begin with a sign (**+** or **–**).

blank If the first character of a signed conversion is not a sign, a blank will be prefixed to the result. This implies that if the blank and **+** flags both appear, the blank flag will be ignored.

This flag specifies that the value is to be converted to an "alternate form."For **c**, **d**, **i**, **s**, and **u** conversions, the flag has no effect. For **o** conversion, it increases the precision to force the first digit of the result to be a zero. For **x** or **X** conversion, a non-zero result will have **0x** or **0X** prefixed to it. For **e**, **E**, **f**, **g**, and **G** conversions, the result will always contain a decimal point, even if no digits follow the point (normally, a decimal point appears in the result of these conversions only if a digit follows it). For **g** and **G** conversions, trailing zeroes will *not* be removed from the result (which they normally are).

The conversion characters and their meanings are:

d,**i**,**o**,**u**,**x**,**X** The integer *arg* is converted to signed decimal (**d** or **i**), unsigned octal (**o**), unsigned decimal (**u**), or unsigned hexadecimal notation (**x** and **X**), respectively; the letters **abcdef** are used for **x** conversion and the letters **ABCDEF** for **X** conversion. The precision specifies the minimum number of digits to appear; if the value being converted can be represented in fewer digits, it will

be expanded with leading zeroes. (For compatibility with older versions, padding with leading zeroes may alternatively be specified by prepending a zero to the field width. This does not imply an octal value for the field width.) The default precision is 1. The result of converting a zero value with a precision of zero is a **NULL** string.

f The float or double *arg* is converted to decimal notation in the style [–]*ddd.ddd* where the number of digits after the decimal point is equal to the precision specification. If the precision is missing, 6 digits are given; if the precision is explicitly 0, no digits and no decimal point are printed.

e,**E** The float or double *arg* is converted in the style [–]*d.ddd***e**±*ddd*, where there is one digit before the decimal point and the number of digits after it is equal to the precision; when the precision is missing, 6 digits are produced; if the precision is zero, no decimal point appears. The **E** format code will produce a number with **E** instead of **e** introducing the exponent. The exponent always contains at least two digits.

g,**G** The float or double *arg* is printed in style **f** or **e** (or in style **E** in the case of a **G** format code), with the precision specifying the number of significant digits. The style used depends on the value converted: style **e** or **E** will be used only if the exponent resulting from the conversion is less than –4 or greater than the precision. Trailing zeroes are removed from the result; a decimal point appears only if it is followed by a digit.

The **e**, **E f**, **g,** and **G** formats print IEEE indeterminate values (infinity or not-a-number) as "Infinity" or "NaN" respectively.

c The character *arg* is printed.

s The *arg* is taken to be a string (character pointer) and characters from the string are printed until a **NULL** character (**\0**) is encountered or until the number of characters indicated by the precision specification is reached. If the precision is missing, it is taken to be infinite, so all characters up to the first **NULL** character are printed. A **NULL** value for *arg* will yield undefined results.

% Print a **%**; no argument is converted.

In no case does a non-existent or small field width cause truncation of a field; if the result of a conversion is wider than the field width, the field is simply expanded to contain the conversion result. Padding takes place only if the specified field width exceeds the actual width. Characters generated by **printf()** and **fprintf()** are printed as if **putc**(3S) had been called.

RETURN VALUES

Upon success, **printf()** and **fprintf()** return the number of characters transmitted, excluding the null character. **vprintf()** and **vfprintf()** return the number of characters transmitted. **sprintf()** and **vsprintf()** always return *s*. If an output error is encountered, **printf()**, **fprint()**, **vprintf()**, and **vfprintf()** return EOF.

EXAMPLES

To print a date and time in the form "Sunday, July 3, 10:02," where *weekday* and *month* are pointers to **NULL**-terminated strings:

printf("%s, %s %i, %d:%.2d", weekday, month, day, hour, min);

To print π to 5 decimal places:

printf("pi = %.5f", 4 * atan(1.0));

SEE ALSO

econvert(3), **putc**(3S), **scanf**(3S), **vprintf**(3S), **varargs**(5)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

Very wide fields (>128 characters) fail.

NAME

printf, fprintf, sprintf – print formatted output

SYNOPSIS

```
#include <stdio.h>
int printf(const char *format, /* args */ ...);
int fprintf(FILE *strm, const char *format, /* args */ ...);
int sprintf(char *s, const char *format, /* args */ ...);
```

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

printf( ) places output on the standard output stream **stdout**.

fprintf( ) places output on *strm*.

sprintf( ) places output, followed by the null character (**\0**), in consecutive bytes starting at *s*. It is the user's responsibility to ensure that enough storage is available. Each function returns the number of characters transmitted (not including the **\0** in the case of **sprintf( )**) or a negative value if an output error was encountered.

Each of these functions converts, formats, and prints its *args* under control of the *format*. The *format* is a character string that contains three types of objects defined below:

1. plain characters that are simply copied to the output stream;
2. escape sequences that represent non-graphic characters;
3. conversion specifications.

The following escape sequences produce the associated action on display devices capable of the action:

\a Alert. Ring the bell.

\b Backspace. Move the printing position to one character before the current position, unless the current position is the start of a line.

\f Form feed. Move the printing position to the initial printing position of the next logical page.

\n Newline. Move the printing position to the start of the next line.

\r Carriage return. Move the printing position to the start of the current line.

\t Horizontal tab. Move the printing position to the next implementation-defined horizontal tab position on the current line.

\v Vertical tab. Move the printing position to the start of the next implementation-defined vertical tab position.

All forms of the **printf( )** functions allow for the insertion of a language-dependent decimal-point character. The decimal-point character is defined by the program's locale (category **LC_NUMERIC**). In the C locale, or in a locale where the decimal-point character is not defined, the decimal-point character defaults to a period (**.**).

Each conversion specification is introduced by the character **%**. After the character **%**, the following appear in sequence:

An optional field, consisting of a decimal digit string followed by a **$**, specifying the next *args* to be converted. If this field is not provided, the *args* following the last *args* converted will be used.

Zero or more *flags*, which modify the meaning of the conversion specification.

An optional string of decimal digits to specify a minimum *field width*. If the converted value has fewer characters than the field width, it will be padded on the left (or right, if the left-adjustment flag (–), described below, has been given) to the field width. If the format is **%s** or **%ws** (wide-character string), then the field width should be interpreted as the minimum columns of screen display. E.g. **%10s** means if the converted value has a screen width of 7 columns, then 3 spaces would be padded on the right.

An optional precision that gives the minimum number of digits to appear for the **d**, **i**, **o**, **u**, **x**, or **X** conversions (the field is padded with leading zeros), the number of digits to appear after the decimal-point character for the **e**, **E**, and **f** conversions, the maximum number of significant digits for the **g** and **G** conversions, or the maximum number of characters to be printed from a string in **s** or **ws** conversions. The precision takes the form of a period (**.**) followed by a decimal digit string; a null digit string is treated as zero. Padding specified by the precision overrides the padding specified by the field width.

An optional **h** specifies that a following **d**, **i**, **o**, **u**, **x**, or **X** conversion specifier applies to a **short int** or **unsigned short int** argument (the argument will be promoted according to the integral promotions and its value converted to **short int** or **unsigned short int** before printing); an optional **h** specifies that a following **n** conversion specifier applies to a pointer to a **short int** argument. An optional **l** (ell) specifies that a following **d**, **i**, **o**, **u**, **x**, or **X** conversion specifier applies to a **long int** or **unsigned long int** argument; an optional **l** (ell) specifies that a following **n** conversion specifier applies to a pointer to a **long int** argument. An optional **ll** (ell ell) specifies that a following **d**, **i**, **o**, **u**, **x**, or **X** conversion specifier applies to a **long long int** or **unsigned long long int** argument; an optional **ll** (ell ell) specifies that a following **n** conversion specifier applies to a pointer to a **long long int** argument. An optional **L** specifies that a following **e**, **E**, **f**, **g**, or **G** conversion specifier applies to a **long double** argument. If an **h**, **l**, or **L** appears before any other conversion specifier, the behavior is undefined.

A conversion character (see below) that indicates the type of conversion to be applied.

A field width or precision may be indicated by an asterisk (*) instead of a digit string. In this case, an integer *args* supplies the field width or precision. The *args* that is actually converted is not fetched until the conversion letter is seen, so the *args* specifying field width or precision must appear before the *args* (if any) to be converted. If the *precision* argument is negative, it will be changed to zero. A negative field width argument is taken as a – flag, followed by a positive field width.

In format strings containing the *%digits*$ form of a conversion specification, a field width or precision may also be indicated by the sequence *%digits*$, giving the position in the argument list of an integer *args* containing the field width or precision.

When numbered argument specifications are used, specifying the *N*th argument requires that all the leading arguments, from the first to the (*N*–1)th, be specified in the format string.

The *flag* characters and their meanings are:

– The result of the conversion will be left-justified within the field. (It will be right-justified if this flag is not specified.)

+ The result of a signed conversion will always begin with a sign (+ or –). (It will begin with a sign only when a negative value is converted if this flag is not specified.)

space If the first character of a signed conversion is not a sign, a space will be placed before the result. This means that if the space and + flags both appear, the space flag will be ignored.

The value is to be converted to an alternate form. For **c**, **d**, **i**, **s**, and **u** conversions, the flag has no effect. For an **o** conversion, it increases the precision to force the first digit of the result to be a zero. For **x** (or **X**) conversion, a non-zero result will have **0x** (or **0X**) prepended to it. For **e**, **E**, **f**, **g**, and **G** conversions, the result will always contain a decimal-point character, even if no digits follow the point (normally, a decimal point appears in the result of these conversions only if a digit follows it). For **g** and **G** conversions, trailing zeros will not be removed from the result as they normally are.

0 For **d**, **i**, **o**, **u**, **x**, **X**, **e**, **E**, **f**, **g**, and **G** conversions, leading zeros (following any indication of sign or base) are used to pad to the field width; no space padding is performed. If the **0** and **–** flags both appear, the **0** flag will be ignored. For **d**, **i**, **o**, **u**, **x**, and **X** conversions, if a precision is specified, the **0** flag will be ignored. For other conversions, the behavior is undefined.

Each conversion character results in fetching zero or more *args*. The results are undefined if there are insufficient *args* for the format. If the format is exhausted while *args* remain, the excess *args* are ignored.

The conversion characters and their meanings are:

d,**i**,**o**,**u**,**x**,**X** The integer *arg* is converted to signed decimal (**d** or **i**), unsigned octal (**o**), unsigned decimal (**u**), or unsigned hexadecimal notation (**x** and **X**). The **x** conversion uses the letters **abcdef** and the **X** conversion uses the letters **ABCDEF**. The precision specifies the minimum number of digits to appear. If the value being converted can be represented in fewer digits than the specified minimum, it will be expanded with leading zeros. The default precision is 1. The result of converting a zero value with a precision of zero is no characters.

f The double *args* is converted to decimal notation in the style **[**–**]***ddd***.***ddd*, where the number of digits after the decimal-point character (see

setlocale(3C)) is equal to the precision specification. If the precision is omitted from *arg*, six digits are output; if the precision is explicitly zero and the **#** flag is not specified, no decimal-point character appears. If a decimal-point character appears, at least 1 digit appears before it. The value is rounded to the appropriate number of digits.

e,**E** The double *args* is converted to the style **[–]***d***.***ddd***e**±*dd*, where there is one digit before the decimal-point character (which is non-zero if the argument is non-zero) and the number of digits after it is equal to the precision. When the precision is missing, six digits are produced; if the precision is zero and the **#** flag is not specified, no decimal-point character appears. The **E** conversion character will produce a number with **E** instead of **e** introducing the exponent. The exponent always contains at least two digits. The value is rounded to the appropriate number of digits.

g,**G** The double *args* is printed in style **f** or **e** (or in style **E** in the case of a **G** conversion character), with the precision specifying the number of significant digits. If the precision is zero, it is taken as one. The style used depends on the value converted: style **e** (or **E**) will be used only if the exponent resulting from the conversion is less than –4 or greater than or equal to the precision. Trailing zeros are removed from the fractional part of the result. A decimal-point character appears only if it is followed by a digit.

c The **int** *args* is converted to an **unsigned char**, and the resulting character is printed.

wc The **int** *args* is converted to a wide character (**wchar_t**), and the resulting wide character is printed.

s The *args* is taken to be a string (character pointer) and characters from the string are written up to (but not including) a terminating null character; if the precision is specified, no more than that many characters are written. If the precision is not specified, it is taken to be infinite, so all characters up to the first null character are printed. A null value for *args* will yield undefined results.

ws The *args* is taken to be a wide character string (wide character pointer) and wide characters from the string are written up to (but not including) a terminating null character; if the precision is specified, no more than that many wide characters are written. If the precision is not specified, it is taken to be infinite, so all wide characters up to the first null character are printed. A null value for *args* will yield undefined results.

p The *args* should be a pointer to **void**. The value of the pointer is converted to an implementation-defined set of sequences of printable characters, which should be the same as the set of sequences that are matched by the **%p** conversion of the **scanf** function.

n The argument should be a pointer to an integer into which is written the

number of characters written to the output standard I/O stream so far by this call to **printf( )**, **fprintf( )**, or **sprintf( )**. No argument is converted.

% Print a **%**; no argument is converted.

If the character after the **%** or **%***digits***$** sequence is not a valid conversion character, the results of the conversion are undefined.

If a floating-point value is the internal representation for infinity, the output is [±]*Infinity*, where *Infinity* is either **Infinity** or **Inf**, depending on the desired output string length. Printing of the sign follows the rules described above.

If a floating-point value is the internal representation for ''not-a-number,'' the output is [±]*NaN*. Printing of the sign follows the rules described above.

In no case does a non-existent or small field width cause truncation of a field; if the result of a conversion is wider than the field width, the field is simply expanded to contain the conversion result. Characters generated by **printf** and **fprintf** are printed as if the **putc( )** routine had been called.

LC_NUMERIC
Determines how numeric formats are handled. In the "C" locale, numeric handling follows the U.S. rules.

RETURN VALUES

printf( ), **fprintf( )**, and **sprintf( )** return the number of characters transmitted, or return a negative value if an error was encountered.

EXAMPLES

To print a date and time in the form **Sunday, July 3, 10:02**, where **weekday** and **month** are pointers to null-terminated strings:

```
printf("%s, %s %i, %d:%.2d", weekday, month, day, hour, min);
```

To print π to 5 decimal places:

```
printf("pi = %.5f", 4 * atan(1.0));
```

To print a list of names in columns which are 20 characters wide:

```
printf("%20s%20s%20s", lastname, firstname, middlename );
```

FILES

/usr/lib/locale/*locale***/LC_NUMERIC/numeric**
LC_NUMERIC database for *locale*

SEE ALSO

exit(2), **lseek**(2), **write**(2), **abort**(3C), **ecvt**(3C), **putc**(3S), **scanf**(3S), **setlocale**(3C), **stdio**(3S)

NOTES

sprintf( ) is MT-Safe in multi-thread applications. **printf** and **fprintf** can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

proc_service, ps_pstop, ps_pcontinue, ps_lstop, ps_lcontinue, ps_pglobal_lookup, ps_pdread, ps_pdwrite, ps_ptread, ps_ptwrite, ps_lgetregs, ps_lsetregs, ps_plog, ps_lgetxregsize, ps_lgetxregs, ps_lsetxregs, ps_lgetfpregs, ps_lsetfpregs – process service interface

SYNOPSIS

```
#include <proc_service.h>

ps_err_e ps_pstop(const struct ps_prochandle *ph);
ps_err_e ps_pcontinue(const struct ps_prochandle *ph);
ps_err_e ps_lstop(const struct ps_prochandle *ph, lwpid_t lwpid);
ps_err_e ps_lcontinue(const struct ps_prochandle *ph, lwpid_t lwpid);

ps_err_e ps_pglobal_lookup(const struct ps_prochandle *ph,
        const char *ld_object_name, const char *ld_symbol_name,
        paddr_t *ld_symbol_addr);

ps_err_e ps_pdread(const struct ps_prochandle *ph, paddr_t addr,
        char *buf, int size);

ps_err_e ps_pdwrite(const struct ps_prochandle *ph, paddr_t addr,
        char *buf, int size);

ps_err_e ps_ptread(const struct ps_prochandle *ph, paddr_t addr,
        char *buf, int size);

ps_err_e ps_ptwrite(const struct ps_prochandle *ph, paddr_t addr,
        char *buf, int size);

ps_err_e ps_lgetregs(const struct ps_prochandle *ph, lwpid_t lwpid,
        prgregset_t gregset);

ps_err_e ps_lsetregs(const struct ps_prochandle *ph, lwpid_t lwpid,
        const prgregset_t gregset);

void ps_plog(const char *fmt, ... );

ps_err_e ps_lgetxregsize( const struct ps_prochandle *ph, lwpid_t lwpid,
        int *xregsize);

ps_err_e ps_lgetxregs( const struct ps_prochandle *ph, lwpid_t lwpid,
        prxregset_t *xregset);

ps_err_e ps_lsetxregs( const struct ps_prochandle *ph, lwpid_t lwpid,
        prxregset_t *xregset);

ps_err_e ps_lgetfpregs(const struct ps_prochandle *ph, lwpid_t lwpid,
        prfpregset_t *fpregset);

ps_err_e ps_lsetfpregs(const struct ps_prochandle *ph, lwpid_t lwpid,
        const prfpregset_t *fpregset);
```

x86 Only

ps_err_e ps_lgetLDT(const struct ps_prochandle **ph*, **lwpid_t** *lwpid*, **struct ssd** **ldt*);

DESCRIPTION

proc_service is a set of interfaces that provide operations on a process. These services are utilized by libraries such as libthread_db to examine a process under inspection (PUI). These functions are provided by the user of such libraries.

FUNCTIONS

ps_pstop( ) stops the process with the given process handle and returns after the process has stopped.

ps_continue( ) continues a process with the given process handle that was stopped by **ps_pstop( )**.

ps_lstop( ) stops an LWP on the process with the given lwp identifier and process handle and returns after the LWP has stopped.

ps_lcontinue( ) continues an LWP on the process with the given lwp identifier and process handle that was stopped by **ps_lstop( )**.

ps_pglobal_lookup( ) finds the absolute address of a symbol given a process handle, symbol name and load object name (e.g., "libthread.so"). Only the global name space is searched.

ps_pdread( ) reads from a process's data segment at the given address, addr, for size bytes. Contents are returned in buf.

ps_pdwrite( ) writes contents of buf for size bytes to a process's data segment at the given address, addr.

ps_ptread( ) reads from a process's text segment at the given address, addr, for the size bytes. Contents are returned in buf.

ps_ptwrite( ) writes contents of buf for size bytes to a process's text segment at the address, addr.

ps_lgetregs( ) reads an LWP's (process with process handle *ph and lwp with identifier *lwpid*) general register set into gregset.

ps_lsetregs( ) writes an LWP's (process with process handle *ph and lwp with identifier *lwpid*) general register set from gregset.

ps_plog( ) logs a message. ps_plog( ) may be used to write diagnostic messages. Input is in printf style format.

ps_lgetxregsize( ) reads an LWP's (process with process handle *ph and lwp with identifier *lwpid*) extra register set size.

ps_lgetxregs( ) reads an LWP's (process with process handle *ph and lwp with identifier *lwpid*) extra register set into xregset.

ps_lsetxregs( ) writes an LWP's (process with process handle *ph and lwp with identifier *lwpid*) extra register set from xregset.

ps_lgetfpregs( ) reads an LWP's (process with process handle *ph and lwp with identifier *lwpid*) floating point register set into fpregset.

ps_lsetfpregs( ) writes an LWP's (process with process handle *ph and lwp with identifier *lwpid*) floating point register set from fpregset.

x86 Only

ps_lgetLDT( ) reads an LWP's (process with process handle *ph and lwp with identifier *lwpid*) Local Descriptor Table.

RETURN VALUES

PS_OK	The call succeeded.
PS_ERR	The call failed without a specific reason.
PS_BADPID	This return value indicates that a bad process handled was passed to the function.
PS_BADLID	This return value indicates that a bad lwp identifier was passed to the function.
PS_BADADDR	This return value indicates that a bad lwp identifier was passed to the function.
PS_NOSYM	The symbol could not be found.
PS_NOFREGS	This return value indicates that an FPU register set was not available for the lwp.

SEE ALSO

libthread_db(3T)

NAME

psignal, sys_siglist – system signal messages

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

void psignal (*sig, s***)**
unsigned *sig*;
char **s*;

char ***sys_siglist[ ];**

DESCRIPTION

psignal( ) produces a short message on the standard error file describing the indicated signal. First the argument string *s* is printed, then a colon, then the name of the signal and a NEWLINE. Most usefully, the argument string is the name of the program which incurred the signal. The signal number should be from among those found in **<signal.h>**.

To simplify variant formatting of signal names, the vector of message strings **sys_siglist** is provided; the signal number can be used as an index in this table to get the signal name without the newline. The define **NSIG** defined in **<signal.h>** is the number of messages provided for in the table; it should be checked because new signals may be added to the system before they are added to the table.

SEE ALSO

perror(3C), **signal**(3C)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

NAME

psignal, psiginfo – system signal messages

SYNOPSIS

```
#include <siginfo.h>
void psignal(int sig, const char *s);
void psiginfo(siginfo_t *pinfo, char *s);
```

MT-LEVEL

Safe

DESCRIPTION

psignal() and **psiginfo()** produce messages on the standard error output describing a signal. *sig* is a signal that may have been passed as the first argument to a signal handler. *pinfo* is a pointer to a **siginfo** structure that may have been passed as the second argument to an enhanced signal handler (see **sigaction**(2)). The argument string *s* is printed first, then a colon and a blank, then the message and a newline.

SEE ALSO

sigaction(2), **gettext**(3I), **perror**(3C), **setlocale**(3C), **siginfo**(5), **signal**(5)

NOTES

If the application is linked with **–lintl**, then messages printed from these functions are in the native language specified by the **LC_MESSAGES** locale category; see **setlocale**(3C).

NAME

pthread_atfork – register fork handlers

SYNOPSIS

```
#include <sys/types.h>
int pthread_atfork (void (*prepare)(void), void (*parent)(void), void (*child)(void));
```

MT-LEVEL

MT-Safe

DESCRIPTION

pthread_atfork( ) declares fork handlers to be called prior to and following **fork( )**, within the thread that called **fork( )**. The order of calls to **pthread_atfork( )** is important.

Before **fork( )** processing begins, the *prepare* fork handler is called. The *prepare* handler is not called if its address is **NULL.**

The *parent* fork handler is called after **fork( )** processing finishes in the parent process, and the *child* fork handler is called after **fork( )** processing finishes in the child process. If the address of *parent* or *child* is **NULL,** then its handler is not called.

The *prepare* fork handler is called in LIFO (last-in first-out) order, whereas the *parent* and *child* fork handlers are called in FIFO (first-in first-out) order. This calling order allows applications to preserve locking order.

RETURN VALUES

Upon successful completion, **pthread_atfork( )** returns **0**; otherwise, an error number is returned.

ERRORS

ENOMEM Insufficient table space exists to record the fork handler addresses.

SEE ALSO

fork(2), **atexit**(3C)

NOTES

Solaris threads do not offer this functionality, although a call to this interface may be used by a Solaris thread program since the two thread APIs are interoperable.

EXAMPLES

All multi-threaded applications that call **fork()** in a POSIX threads program, or call **fork1()** in a Solaris threads program, and which do more than simply call **exec()** in the child of the fork, should ensure that the child is protected from deadlock.

The deadlock scenario: since the "fork-one" model results in cloning only the thread that called fork, it is possible that, at the time of the call, another thread in the parent owns a lock. In the child, this thread is not cloned, and so no thread will unlock this lock in the child. Now, if the single thread in the child needs this lock, there is a deadlock.

The problem is more serious with locks in libraries. Since a library writer does not know if the application that is using the library calls **fork()** or not, the library has to protect itself, for complete correctness, from such a deadlock scenario. If the application that links with this library calls **fork()** and does not call **exec()** in the child, and needs a library lock that may be held by some other thread in the parent which is inside the library at the time of the fork, then the application deadlocks inside the library. The problem may be solved by using **pthread_atfork()**.

The following is a brief and simple description of how to make a library safe with respect to **fork1()** by using **pthread_atfork()**.

- Identify all the locks used by the library. Let's say this list is {L1,...Ln}. Also identify the locking order for these locks. Let's say that this order is also L1...Ln.
- Add a call to **pthread_atfork(f1, f2, f3)** in the library's *.init* section. f1, f2, f3 are defined as follows:

```
f1()
{
        pthread_mutex_lock(L1);          |
        pthread_mutex_lock(...);         |  --> ordered in lock order
        pthread_mutex_lock(Ln);          |
}                                        V

f2()
{
        pthread_mutex_unlock(L1);
        pthread_mutex_unlock(...);
        pthread_mutex_unlock(Ln);
}

f3()
{
        pthread_mutex_unlock(L1);
        pthread_mutex_unlock(...);
        pthread_mutex_unlock(Ln);
}
```

NAME

pthread_attr_init, pthread_attr_destroy, pthread_attr_setscope, pthread_attr_getscope, pthread_attr_setdetachstate, pthread_attr_getdetachstate, pthread_attr_setstacksize, pthread_attr_getstacksize, pthread_attr_setstackaddr, pthread_attr_getstackaddr, pthread_attr_setschedparam, pthread_attr_getschedparam, pthread_attr_setschedpolicy, pthread_attr_getschedpolicy, pthread_attr_setinheritsched, pthread_attr_getinheritsched – thread creation attributes

SYNOPSIS

```
#include <pthread.h>

int pthread_attr_init(pthread_attr_t *attr);
int pthread_attr_destroy(pthread_attr_t *attr);
int pthread_attr_setscope(pthread_attr_t *attr, int contentionscope);
int pthread_attr_getscope(const pthread_attr_t *attr, int *contentionscope);
int pthread_attr_setdetachstate(pthread_attr_t *attr, int detachstate);
int pthread_attr_getdetachstate(const pthread_attr_t *attr, int *detachstate);
int pthread_attr_setstacksize(pthread_attr_t *attr, size_t stacksize);
int pthread_attr_getstacksize(const pthread_attr_t *attr, size_t *stacksize);
int pthread_attr_setstackaddr(pthread_attr_t *attr, void *stackaddr);
int pthread_attr_getstackaddr(const pthread_attr_t *attr, void **stackaddr);
int pthread_attr_setschedparam(pthread_attr_t *attr,
        const struct sched_param *param);
int pthread_attr_getschedparam(const pthread_attr_t *attr,
        struct sched_param *param);
int pthread_attr_setschedpolicy(pthread_attr_t *attr, int policy);
int pthread_attr_getschedpolicy(const pthread_attr_t *attr, int *policy);
int pthread_attr_setinheritsched(pthread_attr_t *attr, int inheritsched);
int pthread_attr_getinheritsched(const pthread_attr_t *attr, int *inheritsched);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The pthread approach to setting attributes for threads is to request the initialization of an attribute object, *attr*, and pass the initialized attribute object to **pthread_create**(3T). The convention in Solaris is to pass these attributes as flags to **thr_create**(3T).

All attributes in *attr* are independent of one another and may be singularly modified or retrieved. *attr*, itself, is independent of any thread and can be modified or used to create new threads. However, any change to *attr* after a thread is created will not affect that thread.

init

The **pthread_attr_init( )** function initializes a thread attributes object (*attr*) with the default value for each attribute as follows:

Attribute	Default	Value
contentionscope	**PTHREAD_SCOPE_PROCESS**	resource competition within process
detachstate	**PTHREAD_CREATE_JOINABLE**	joinable by other threads
stackaddr	**NULL**	stack allocated by system
stacksize	**NULL**	1 megabyte

priority	----	priority of parent (calling) thread
policy	**SCHED_OTHER**	determined by system
inheritsched	**PTHREAD_EXPLICIT_SCHED**	scheduling policy and parameters not inherited but explicitly defined by the attribute object

NOTE: Attribute objects should be destroyed before an initialized attribute object is re-initialized.

destroy

pthread_attr_destroy() destroys a thread attributes object (*attr*), which cannot be reused until it is reinitialized.

resource contentionscope

The **pthread_attr_setscope()** and **pthread_attr_getscope()** functions set and get the *contentionscope* thread attribute in the *attr* object. The *contentionscope* value may be set to the following:

PTHREAD_SCOPE_SYSTEM Indicates system scheduling contention scope. This thread is permanently "bound" to an LWP, and is also called a bound thread. This value is equivalent to **THR_BOUND** in Solaris threads (see **thr_create**(3T)).

PTHREAD_SCOPE_PROCESS Indicates process scheduling contention scope. This thread is not "bound" to an LWP, and is also called an unbound thread. **PTHREAD_SCOPE_PROCESS**, or unbound, is the default.

detachstate

The **pthread_attr_setdetachstate()** and **pthread_attr_getdetachstate()** functions set and get the *detachstate* attribute in the *attr* object. The *detachstate* attribute determines whether the thread is created in a detached state or not. The *detachstate* may be set to the following values:

PTHREAD_CREATE_DETACHED Creates a new detached thread. A detached thread disappears without leaving a trace. The thread ID and any of its resources are freed and ready for reuse. **pthread_join**(3T) and **thr_join**(3T) cannot wait for a detached thread.

PTHREAD_CREATE_JOINABLE Creates a new non-detached thread. The thread ID and its user-defined stack, if specified at thread creation time, is not freed until **pthread_join**(3T) or **thr_join**(3T) are called. **pthread_join**(3T) or **thr_join**(3T) must be called to release any resources associated with the terminated thread.

stacksize and stackaddr

The **pthread_attr_setstacksize()** and **pthread_attr_getstacksize()** functions set and get the *stacksize* thread attribute in the *attr* object. The *stacksize* default argument is **NULL**, and a thread default stack size is 1 megabyte.

The **pthread_attr_setstackaddr()** and **pthread_attr_getstackaddr()** functions set and get the *stackaddr* thread attribute in the *attr* object. The *stackaddr* default is **NULL**. (See **pthread_create**(3T).)

schedparam (priority)

The **pthread_attr_setschedparam()** and **pthread_attr_getschedparam()** functions set and get the scheduling parameter thread attributes in the *attr* argument, determined by the scheduling policy set in the *attr* object. The only required member of the *param* structure for the **SCHED_OTHER**, **SCHED_FIFO**, and **SCHED_RR** policies is *sched_priority* (see **NOTES** section below). You can use these functions to get and set the priority of the thread to be created. The **sched_priority** of the *param* structure is **NULL**, by default, which means the newly created thread inherits the priority of its parent thread.

schedpolicy

The **pthread_attr_setschedpolicy()** and **pthread_attr_getschedpolicy()** functions set and get the *schedpolicy* thread attribute in the *attr* argument.

Values for the *policy* attribute are **SCHED_FIFO**, **SCHED_RR**, or the default value **SCHED_OTHER** (see **NOTES** section below.

RETURN VALUES

Upon successful completion, the following functions return **0**; otherwise, an error number is returned to indicate the error: **pthread_attr_init()**, **pthread_attr_destroy()**, **pthread_attr_setstacksize()**, **pthread_attr_getstacksize()**, **pthread_attr_setstackaddr()**, **pthread_attr_getstackaddr()**, **pthread_attr_setdetachstate()**, **pthread_attr_getdetachstate()**, **pthread_attr_setscope**()**, pthread_attr_getscope()**, **pthread_attr_setinheritsched()**, **pthread_attr_getinheritsched()**, **pthread_attr_setschedpolicy()**, and **pthread_attr_getschedpolicy()**.

ERRORS

If any of the following conditions occur, **pthread_attr_init()** returns the corresponding error number:

ENOMEM Insufficient memory exists to create the thread attributes object.

If any of the following conditions occur, **pthread_attr_setstacksize()** returns the corresponding error number:

EINVAL The value of *stacksize* is less than **PTHREAD_STACK_MIN** or exceeds a system-imposed limit.

If any of the following conditions occur, **pthread_attr_destroy()**, **pthread_attr_setstacksize()**, **pthread_attr_getstacksize()**, **pthread_attr_setstackaddr()**, **pthread_attr_getstackaddr()**, **pthread_attr_setdetachstate()**, **pthread_attr_getdetachstate()**, **pthread_attr_setscope()**, **pthread_attr_getscope()**, **pthread_attr_setschedparam()**, **pthread_attr_getschedparam()**, **pthread_attr_setinheritsched()**, **pthread_attr_getinheritsched()**, **pthread_attr_setschedpolicy()**, and **pthread_attr_getschedpolicy()** return the corresponding error number:

EINVAL The value of *attr* is not valid.

If any of the following conditions occur, **pthread_attr_setstacksize()** returns the corresponding error number:

EINVAL The value of *stacksize* is less than **PTHREAD_STACK_MIN**.

If any of the following conditions occur, **pthread_attr_setdetachstate()** returns the corresponding error number:

EINVAL The value of *detachstate* is not valid.

If any of the following conditions occur, **pthread_attr_setscope()** returns the corresponding error number:

EINVAL The value of *contentionscope* is not valid.

If any of the following conditions occur, **pthread_attr_setschedparam()** returns the corresponding error number:

EINVAL The value of the **sched_priority** member of the *param* structure is less than or equal to **0**.

If any of the following conditions occur, **pthread_attr_getstacksize()** returns the corresponding error number:

EINVAL The value of *stacksize* is **NULL**.

If any of the following conditions occur, **pthread_attr_getstackaddr()** returns the corresponding error number:

EINVAL The value of *stackaddr* is **NULL**.

If any of the following conditions occur, **pthread_attr_getdetachstate()** returns the corresponding error number:

EINVAL The value of *detachstate* is **NULL**.

If any of the following conditions occur, **pthread_attr_getscope()** returns the corresponding error number:

EINVAL The value of *contentionscope* is **NULL**.

If any of the following conditions occur, either **pthread_attr_setschedparam()** and **pthread_attr_getschedparam()** returns the corresponding error number:

EINVAL The value of *param* is **NULL**.

For each of the following conditions, if the condition is detected, **pthread_attr_setinheritsched()** and **pthread_attr_setschedpolicy()** return the corresponding error number:

ENOTSUP An attempt was made to set the attribute to an unsupported *policy* or *inheritsched*.

For each of the following conditions, if the condition is detected, **pthread_attr_getinheritsched()** and **pthread_attr_getschedpolicy()** return the corresponding error number:

EINVAL *policy* or *inheritsched* is **NULL**.

SEE ALSO **pthread_create**(3T), **pthread_join**(3T), **thr_create**(3T).

NOTES

Currently, the only policy supported is **SCHED_OTHER**. Attempting to set policy as **SCHED_FIFO** or **SCHED_RR** will result in the error **ENOSUP**.

The attribute object is part of the POSIX threads interface. There is no Solaris threads counterpart to the POSIX threads attribute object.

NAME

pthread_condattr_init, pthread_condattr_setpshared, pthread_condattr_getpshared, pthread_condattr_destroy – condition variable initialization attributes

SYNOPSIS

```
#include <pthread.h>

int pthread_condattr_init(pthread_condattr_t *attr);
int pthread_condattr_setpshared(pthread_condattr_t *attr, int process-shared);
int pthread_condattr_getpshared(const pthread_condattr_t *attr, int *process-shared);
int pthread_condattr_destroy(pthread_condattr_t *attr);
```

MT-LEVEL

MT-Safe

DESCRIPTION

Initialize

The function **pthread_condattr_init( )** initializes a condition variable attributes object *attr* with the default value for all the attributes.

At present, the only attribute available is the scope of condition variables, specified by *process-shared*.

The default value of the *process-shared* attribute is **PTHREAD_PROCESS_PRIVATE**, which only allows the condition variable to be operated upon by threads created within the same process as the thread that initialized the condition variable. If threads from other processes try to operate on this condition variable, the behavior is undefined.

The *process-shared* attribute may be set to **PTHREAD_PROCESS_SHARED ,** which allows a condition variable to be operated upon by any thread with access to the memory allocated to the condition variable, even if the condition variable is allocated in memory that is shared by multiple processes.

Attempts to initialize previously initialized condition variable attributes object will leave the storage allocated by the previous initialization unallocated.

Once a condition variable attributes object initializes one or more condition variables, any function affecting the attributes object (including destruction) will not effect any previously initialized condition variables.

Set/Get Scope

pthread_condattr_setpshared( ) sets the *process-shared* attribute in an initialized attributes object referenced by *attr*. **pthread_condattr_getpshared( )** obtains the value of the *process-shared* attribute from the attributes object referenced by *attr*.

Destroy

pthread_condattr_destroy( ) destroys a condition variable attributes object; the object becomes uninitialized. A destroyed condition variable attributes object can be reinitialized with **pthread_condattr_init( )**; however, the results of referencing the object after it has been destroyed are undefined.

RETURN VALUES

pthread_condattr_init( ), **pthread_condattr_destroy( )**, and **pthread_condattr_setpshared( )** return **0** upon a successful return; otherwise, an error number is returned.

pthread_condattr_getpshared() returns **0** upon a successful return, and stores the value of the *process-shared* attribute of *attr* in the object referenced by the *process-shared* parameter; otherwise, an error number is returned.

ERRORS

pthread_condattr_init() returns an error number if any of the following conditions are detected:

ENOMEM Insufficient memory exists to initialize the condition variable attributes object.

pthread_condattr_destroy(), **pthread_condattr_getpshared()**, and **pthread_condattr_setpshared()** return an error number if the following condition is detected:

EINVAL The value specified by *attr* is invalid.

pthread_condattr_setpshared() returns an error number if the following condition is detected:

EINVAL The new value specified for the attribute is outside the range of legal values for that attribute.

SEE ALSO

cond_init(3T), **pthread_create**(3T), **pthread_cond_init**(3T), **pthread_mutex_init**(3T).

NAME

pthread_create, thr_create – thread creation

SYNOPSIS

POSIX

cc [*flag* ...] *file* ... **–lpthread** [*library* ...]

```
#include <pthread.h>
int pthread_create(pthread_t *new_thread_ID, const pthread_attr_t *attr,
        void * (*start_func)(void *), void *arg);
```

Solaris

cc [*flag* ...] *file* ... **–lthread** [*library* ...]

```
#include <thread.h>
int thr_create(void *stack_base, size_t stack_size, void *(*start_func)(void *),
        void *arg, long flags, thread_t *new_thread_ID);
```

MT-LEVEL

MT-Safe

DESCRIPTION

Thread creation adds a new thread of control to the current process. The procedure **main( )**, itself, is a single thread of control. Each thread executes simultaneously with all the other threads within the calling process, and with other threads from other active processes.

A newly created thread shares all of the calling process' global data with the other threads in this process; however, it has its own set of attributes and private execution stack. The new thread inherits the calling thread's signal mask, possibly, and scheduling priority. Pending signals for a new thread are not inherited and will be empty.

The call to create a thread takes the address of a user-defined function, specified by *start_func*, as one of its arguments, which is the complete execution routine for the new thread.

The lifetime of a thread begins with the successful return from **pthread_create( )** or **thr_create( )**, which calls *start_func*() and ends with either:

- the normal completion of *start_func*(),
- the return from an explicit call to **pthread_exit**(3T) or **thr_exit**(3T),
- a thread cancellation (see **pthread_cancel**(3T)). or
- the conclusion of the calling process (see **exit**(2)).

The new thread performs by calling the function defined by *start_func* with one argument, *arg*. If more than one argument needs to be passed to *start_func*, the arguments can be packed into a structure, and the address of that structure can be passed to *arg*.

If *start_func* returns, the thread will terminate with the exit status set to the *start_func* return value (see **pthread_exit**(3T) or **thr_exit**(3T)).

Note that when the thread returns in which **main( )** originated from, the effect is the same as if there were an implicit call to **exit( )** using the return value of **main( )** as the exit status. This differs from a *start_func* return. However, if **main( )** itself calls either **pthread_exit**(3T) or **thr_exit**(3T), only the **main** thread exits, not the entire process.

If the thread creation itself fails, a new thread is not created and the contents of the location referenced by the pointer to the new thread are undefined.

Attributes

The configuration of a set of attributes defines the behavior of a thread. At creation, each attribute of a new thread may be user-defined or set to the default. All attributes are defined upon thread creation, however, some may be dynamically modified after creation. Establishing these attributes varies depending upon whether POSIX or Solaris threads are used. Both implementations offer a few attributes the other does not.

The available attributes are:

Attribute	Description	API
contentionscope	Scheduled by threads library (local scope) or scheduled by the OS (global scope)	both
detachstate	Allows other threads to wait for a particular thread to terminate	both
stackaddr	Sets a pointer to the thread's stack	both
stacksize	Sets the size of the thread's stack	both
concurrency	Elevates concurrency, if possible	Solaris
priority	Sets ranking within the policy (scheduling class)	both
policy	Sets scheduling class; **SCHED_OTHER**	POSIX
inheritsched	Determines whether scheduling parameters are inherited or explicitly defined	POSIX
suspended	Sets thread to runnable vs. suspended	Solaris
daemon	Defines a thread's behavior to be like a daemon	Solaris

POSIX

pthread_create( ) creates a new thread within a process with attributes defined by *attr*. Default attributes are used if *attr* is **NULL**. If any attributes specified by *attr* are changed in the attribute object prior to the call to **pthread_create( )**, the new thread will acquire those changes. However, if any attributes specified by *attr* are changed after the call to **pthread_create( )**, the attributes of existing threads will not be affected. Since **pthread_create( )** can use an attribute object in its call, a user-defined thread creation must be preceded by a user-defined attribute object (see **pthread_attr_init**(3T)). Upon successful completion, and if the return value is not **NULL, pthread_create( )** will store the ID of the created thread in the location referenced by *new_thread_ID*.

It is recommended that for POSIX thread creation, all attribute objects, *attr*s, which will be used later during creation calls, be initialized and modified in the early stages of program execution.

The default creation attributes for **pthread_create**(3T) are:

Attribute	Default Value	Meaning of Default Value
contentionscope	**PTHREAD_SCOPE_PROCESS**	Resource competition within process
detachstate	**PTHREAD_CREATE_JOINABLE**	Joinable by other threads
stackaddr	**NULL**	Allocated by system
stacksize	**NULL**	1 megabyte
priority	**NULL**	Parent (calling) thread's priority
policy	**SCHED_OTHER**	Determined by system

inheritsched	**PTHREAD_EXPLICIT_SCHED**	Scheduling attributes explicitly set, e.g., policy is **SCHED_OTHER**.

Default thread creation:

```
pthread_t tid;
void *start_func(void *), *arg;

pthread_create(&tid, NULL, start_func, arg);
```

This would have the same effect as:

```
pthread_attr_t attr;

pthread_attr_init(&attr); /* initialize attr with default attributes */
pthread_create(&tid, &attr, start_func, arg);
```

User-defined thread creation:

To create a thread that is scheduled on a system-wide basis (i.e., a bound thread, as per the Solaris API), use:

```
pthread_attr_init(&attr); /* initialize attr with default attributes */
pthread_attr_setscope(&attr, PTHREAD_SCOPE_SYSTEM);  /* system-wide contention */
pthread_create(&tid, &attr, start_func, arg);
```

To customize the attributes for POSIX threads, see **pthread_attr_init**(3T).

A new thread created with **pthread_create( )** uses the stack specified by the *stackaddr* attribute, and the stack continues for the number of bytes specified by the *stacksize* attribute. By default, the stack size is 1 megabyte (see **pthread_attr_setstacksize**(3T)). If the default is used for both the *stackaddr* and *stacksize* attributes, **pthread_create( )** creates a stack for the new thread with at least 1 megabyte. (For customizing stack sizes, see NOTES).

Solaris

In the Solaris API, **thr_create( )** either results in the creation of a default thread or a thread whose attributes are defined by the *flags* passed to **thr_create( )**. There is no attribute object to configure, as there is in POSIX. The attributes are either the separate arguments, *stackaddr* or *stacksize*, or the result of bitwise inclusive OR-ing the possible values for *flags*.

The creation attributes for **thr_create**(3T) are:

Attribute	Default Value	Meaning of Default Value	Specified Via
contentionscope	NULL	Resource competition within process	flags
detachstate	NULL	Joinable by other threads	flags
stackaddr	NULL	Allocated by system	separate argument
stacksize	NULL	1 megabyte	separate argument
priority	NULL	Parent (calling) thread's priority	
concurrency	NULL	Determined by system	flags

suspended	NULL	Runnable, not suspended	flags
daemon	NULL	Not a daemon	flags

flags specifies which attributes are modifiable for the created thread. The value in *flags* is determined by the bitwise inclusive OR of the following:

THR_BOUND This flag affects the contentionscope attribute of the thread. The new thread is created permanently bound to an LWP (i.e. it is a *bound thread*). This thread will now contend among system-wide resources. The bind flag is equivalent to setting the *contentionscope* to the **PTHREAD_SCOPE_SYSTEM** in POSIX.

THR_DETACHED This flag affects the detachstate attribute of the thread. The new thread is created detached. The exit status of a detached thread is not accessible to other threads. Its thread ID and other resources may be re-used as soon as the thread terminates. **thr_join**(3T) (nor **pthread_join**(3T)) will not wait for a detached thread. This is equivalent to **PTHREAD_CREATE_DETACHED** in POSIX, which is the default for POSIX.

THR_NEW_LWP This flag affects the concurrency attribute of the thread. The desired concurrency level for unbound threads is increased by one. This is similar to incrementing concurrency by one via **thr_setconcurrency**(3T)**.** Typically, this adds a new LWP to the pool of LWPs running unbound threads.

THR_SUSPENDED This flag affects the suspended attribute of the thread. The new thread is created suspended and will not execute *start_func* until it is started by **thr_continue( )**.

THR_DAEMON This flag affects the daemon attribute of the thread. The thread is marked as a daemon. The process will exit when all non-daemon threads exit. **thr_join**(3T) will not wait for a daemon thread. Daemon threads do not interfere with the exit conditions for a process. A process will terminate when all regular threads exit or the process calls **exit( )**. Daemon threads are most useful in libraries that want to use threads.

Default thread creation:

```
thread_t tid;
void *start_func(void *), *arg;
thr_create(NULL, NULL, start_func, arg, NULL, &tid);
```

User-defined thread creation:

To create a thread scheduled on a system-wide basis (i.e., a bound thread), use:

```
thr_create(NULL, NULL, start_func, arg, THR_BOUND, &tid);
```

Another example of customization is, if both **THR_BOUND** and **THR_NEW_LWP** are specified then, typically, two LWPs are created, one for the bound thread and another for the pool of LWPs running unbound threads.

```
thr_create(NULL, NULL, start_func, arg, THR_BOUND | THR_NEW_LWP, &tid);
```

With **thr_create( ),** the new thread will use the stack starting at the address specified by *stack_base* and continuing for *stack_size* bytes. *stack_size* must be greater than the value returned by **thr_min_stack**(3T)**.** If *stack_base* is NULL then **thr_create( )** allocates a stack for the new thread with at least *stack_size* bytes. If *stack_size* is zero then a default size is used. If *stack_size* is not zero then it must be greater than the value returned by **thr_min_stack**(3T) (see NOTES).

When *new_thread_ID* is not NULL then it points to a location where the ID of the new thread is stored if **thr_create( )** is successful. The ID is only valid within the calling process.

RETURN VALUES

Zero indicates a successful return and a non-zero value indicates an error.

ERRORS

If any of the following conditions occur, these functions fail and return the corresponding value:

EAGAIN The system-imposed limit on the total number of threads in a process has been exceeded or some system resource has been exceeded (e.g., too many LWPs were created).

EINVAL The value specified by *attr* is invalid.

If any of the following conditions are detected, **pthread_create( )** fails and returns the corresponding value:

ENOMEM Not enough memory was available to create the new thread.

If any of the following conditions are detected, **thr_create( )** fails and returns the corresponding value:

EINVAL

- *stack_base* is not NULL and *stack_size* is less than the value returned by **thr_min_stack**(3T)**.**
- *stack_base* is NULL and *stack_size* is not zero and is less than the value returned by **thr_min_stack**(3T)**.**

EXAMPLES

This is an example of concurrency with multi-threading. Since POSIX threads and Solaris threads are fully compatible even within the same process, this example uses **pthread_create**(3T) if you execute **a.out 0**, or **thr_create**(3T) if you execute **a.out 1**.

Five threads are created that simultaneously perform a time-consuming function, **sleep(**10**)**. If the execution of this process is timed, the results will show that all five individual calls to sleep for ten-seconds completed in about ten seconds, even on a uniprocessor. If a single-threaded process calls **sleep(**10**)** five times, the execution time will be about 50-seconds.

The command-line to time this process is:

/usr/bin/time a.out 0 (for POSIX threading)

or

/usr/bin/time a.out 1 (for Solaris threading)

```
/* cc thisfile.c -lthread -lpthread */
#define _REENTRANT    /* basic 3-lines for threads */
#include <pthread.h>
#include <thread.h>

#define NUM_THREADS 5
#define SLEEP_TIME 10

void *sleeping(void *);   /* thread routine */
void test_argv();         /* optional */
int i;
thread_t tid[NUM_THREADS];      /* array of thread IDs */

main( int argc, char *argv[] ) {
  test_argv(argv[1]);

  switch (*argv[1]) {
    case '0':       /* POSIX */
      for ( i = 0; i < NUM_THREADS; i++)
        pthread_create(&tid[i], NULL, sleeping, SLEEP_TIME);
      for ( i = 0; i < NUM_THREADS; i++)
        pthread_join(tid[i], NULL);
      break;

    case '1':       /* Solaris */
      for ( i = 0; i < NUM_THREADS; i++)
        thr_create(NULL,0,sleeping,NULL,0,&tid[i]);
      while (thr_join(NULL, NULL, NULL) == 0);
      break;
  } /* switch */

  printf("main( ) reporting that all %d threads have terminated\n", i);
} /* main */

void *sleeping(int sleep_time *)    {
      printf("thread %d sleeping %d seconds ...\n", thr_self(), SLEEP_TIME);
      sleep(sleep_time);
      printf("\nthread %d awakening\n", thr_self());
}
```

```
void test_argv(char argv1[])      {   /* optional */
  if (argv1 == NULL)  {
    printf("use 0 as arg1 to use thr_create();\n \
    or use 1 as arg1 for use pthread_create()\n");
    exit(NULL);
  }
}
```

If **main( )** had not waited for the completion of the other threads (using **pthread_join**(3T) or **thr_join**(3T)), it would have continued to process concurrently until it reached the end of its routine and the entire process would have exited prematurely (see **exit**(2)).

The following example shows how to create a default thread with a new signal mask. **new_mask** is assumed to have a different value than the creator's signal mask (orig_mask). **new_mask** is set to block all signals except for SIGINT. The creator's signal mask is changed so that the new thread inherits a different mask, and is restored to its original value after **thr_create()** returns.

This example assumes that SIGINT is also unmasked in the creator. If it is masked by the creator, then unmasking the signal opens the creator up to this signal. The other alternative is to have the new thread set its own signal mask in its start routine.

```
thread_t tid;
sigset_t new_mask, orig_mask;
int error;

(void)sigfillset(&new_mask);
(void)sigdelset(&new_mask, SIGINT);
(void)thr_sigsetmask(SIG_SETMASK, &new_mask, &orig_mask):
error = thr_create(NULL, 0, do_func, NULL, 0, &tid);
(void)thr_sigsetmask(SIG_SETMASK, &orig_mask, NULL);
```

SEE ALSO

_lwp_create(2), **exit**(2), **exit**(3C), **pthread_attr_init**(3T), **pthread_cancel**(3T), **pthread_exit**(3T), **pthread_join**(3T), **thr_suspend**(3T), **thr_min_stack**(3T), **thr_setconcurrency**(3T), **threads**(3T)

NOTES

MT application threads execute independently of each other, thus their relative behavior is unpredictable. Therefore, it is possible for the thread executing **main( )** to finish before all other user application threads.

Using **thr_join**(3T) in the following syntax,

```
while (thr_join(NULL, NULL, NULL) == 0);
```

will cause the invoking thread (which may be **main( )**) to wait for the termination of all other undetached and non-daemon threads; however, the second and third arguments to **thr_join**(3T) need not necessarily be **NULL.**

pthread_join(3T), on the other hand, must specify the terminating thread (IDs) for which it will wait.

A thread has not terminated until **thr_exit**() has finished. The only way to determine this is by **thr_join**(). When **thr_join**() returns a departed thread, it means that this thread has terminated and its resources are reclaimable. For instance, if a user specified a stack to **thr_create**(), this stack can only be reclaimed after **thr_join**() has reported this thread as a departed thread. It is not possible to determine when a *detached* thread has terminated. A detached thread disappears without leaving a trace.

Typically, thread stacks allocated by **thr_create( )** begin on page boundaries and any specified (a red-zone) size is rounded up to the next page boundary. A page with no access permission is appended to the top of the stack so that most stack overflows will result in a **SIGSEGV** signal being sent to the offending thread. Thread stacks allocated by the caller are used as is.

Using a default stack size for the new thread, instead of passing a user-specified stack size, results in much better **thr_create( )** performance. The default stack size for a user-thread is 1 megabyte, in this implementation.

A user-specified stack size must be greater than the value **THR_MIN_STACK** or **PTHREAD_STACK_MIN.** A minimum stack size may not accommodate the stack frame for the user thread function *start_func*. If a stack size is specified, it must accommodate *start_func* requirements and the functions that it may call in turn, in addition to the minimum requirement.

It is usually very difficult to determine the runtime stack requirements for a thread. **THR_MIN_STACK** or **PTHREAD_STACK_MIN** specifies how much stack storage is required to execute a NULL *start_func*. The total runtime requirements for stack storage are dependent on the storage required to do runtime linking, the amount of storage required by library runtimes (like **printf( )**) that your thread calls. Since these storage parameters are not known before the program runs, it is best to use default stacks. If you know your runtime requirements or decide to use stacks that are larger than the default, then it makes sense to specify your own stacks.

NAME

pthread_detach – dynamically detaching a thread

SYNOPSIS

POSIX

cc [*flag* ...] *file* ... –**lpthread** [*library* ...]

```
#include <pthread.h>
int pthread_detach(pthread_t threadID);
```

MT-LEVEL

MT-Safe

DESCRIPTION

pthread_detach( ) can dynamically reset the detachstate attribute of a thread to **PTHREAD_CREATE_DETACHED.** For example, a thread could detach itself as follows:

```
pthread_detach(pthread_self());
```

RETURN VALUES

Upon successful completion, **0** is returned; otherwise, a non-zero value indicates an error.

ERRORS

These functions fail and return the corresponding value, if any of the following conditions are detected:

EINVAL The value specified by *threadID* is not a joinable thread.

ESRCH The value specified by *threadID* is not an existing thread ID.

SEE ALSO

pthread_create(3T), **pthread_join**(3T)

NAME

pthread_equal – compare thread IDs

SYNOPSIS

```
#include <pthread.h>
int pthread_equal(pthread_t t1, pthread_t t2);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **pthread_equal( )** function compares the thread IDs *t1* and *t2*.

RETURN VALUES

If *t1* and *t2* are equal, **pthread_equal( )** returns a non-zero value; otherwise, **0** is returned.

If either *t1* or *t2* is an invalid thread ID, the result is unpredictable.

SEE ALSO

pthread_create(3T), **pthread_self**(3T)

NOTES

Solaris thread IDs do not require an equal function because the **thread_t** structure is really an unsigned int.

NAME

pthread_exit, thr_exit – thread termination

SYNOPSIS

POSIX

cc [*flag* ...] *file* ... **–lpthread** [*library* ...]

```
#include <pthread.h>
void pthread_exit(void *status);
```

Solaris

cc [*flag* ...] *file* ... **–lthread** [*library* ...]

```
#include <thread.h>
void thr_exit(void *status);
```

MT-LEVEL

MT-Safe

DESCRIPTION

pthread_exit( ) and **thr_exit( )** terminates the calling threads, similar to how **exit**(3C) terminates calling processes. If the calling thread is not detached, then the thread's ID and the exit status specified by *status* are retained. The value *status* is then made available to any successful join with the terminating thread (see **pthread_join**(3T)); otherwise, *status* is disregarded allowing the thread's ID to be reclaimed immediately.

Upon thread termination, all thread-specific data bindings are released (see **pthread_key_create**(3T)), and its cancellation routines are called, but application visible process resources, including, but not limited to, mutexes and file descriptors are not released.

The cleanup handlers are called before the thread-specific data bindings are released (see **pthread_cancel**(3T)). Any cancellation cleanup handlers that have been pushed and not yet popped will be popped in reverse order of when they were pushed and then executed. If the thread still has any thread-specific data after all cancellation cleanup handlers have been executed, appropriate destructor functions will be called in an unspecified order. If any thread, including the **main( )**thread, calls **pthread_exit( )**, only that thread will exit.

If **main( )** returns or exits (either implicitly or explicitly), or any thread explicitly calls **exit( )**, the entire process will exit.

If any thread (except the **main( )** thread) implicitly or explicitly returns, the result is the same as if the thread called **pthread_exit( )** and it will return the value of *status* as the exit code.

The process will terminate with an exit status of **0** after the last thread has terminated (including the **main( )** thread). This action is the same as if the application had called **exit( )** with a zero argument at any time.

RETURN VALUES

pthread_exit( ) or **thr_exit( )** does not return to its caller.

SEE ALSO

exit(3C), **pthread_cancel**(3T), **pthread_create**(3T), **pthread_join**(3T), **pthread_key_create**(3T), **pthread_cancel**(3T).

NOTES

Although only POSIX implements cancellation, cancellation can be used with Solaris threads, due to their interoperability.

Do not call **pthread_exit()** from a cancellation cleanup handler or destructor function that will be invoked as a result of either an implicit or explicit call to **pthread_exit()**.

status should not reference any variables local to the calling thread.

NAME

pthread_join, thr_join – wait for thread termination

SYNOPSIS

POSIX

cc [*flag* ...] *file* ... **–lpthread** [*library* ...]

```
#include <pthread.h>
int pthread_join(pthread_t target_thread, void **status);
```

Solaris

cc [*flag* ...] *file* ... **–lthread** [*library* ...]

```
#include <thread.h>
int thr_join(thread_t target_thread, thread_t *departed, void **status);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **pthread_join( )** and **thr_join( )** functions suspend processing of the calling thread until the target *target_thread* completes. *target_thread* must be a member of the current process and it cannot be a detached or daemon thread (see **pthread_create**(3T)).

Several threads cannot wait for the same thread to complete; one thread will complete successfully and the others will terminate with an error of **ESRCH**. **pthread_join( )** or **thr_join( )** will not block processing of the calling thread if the target *target_thread* has already terminated.

pthread_join( ) or **thr_join( )** will return successfully when the target *target_thread* terminates.

POSIX

If a **pthread_join( )** call returns successfully with a non-null *status* argument, the value passed to **pthread_exit**(3T) by the terminating thread will be placed in the location referenced by *status*.

If the **pthread_join( )** calling thread is cancelled, then the target *target_thread* will remain joinable by **pthread_join( )**. However, the calling thread may set up a cancellation cleanup handler on *target_thread* prior to the join call, which may detach the target thread by calling **pthread_detach**(3T). (See **pthread_detach**(3T) and **pthread_cancel**(3T).)

pthread_join( ) does not return the *target_thread*'s ID, as does the Solaris threads' function **thr_join( )**, and it does not cause the calling thread to wait for detached threads. **pthread_join( )** returns **ESRCH** if the target is detached.

Solaris

If a **thr_join( )** call returns successfully with a non-null *status* argument, the value passed to **thr_exit**(3T) by the terminating thread will be placed in the location referenced by *status*.

If the target *target_thread* ID is **0**, **thr_join( )** waits for any undetached thread in the process to terminate.

If *departed* is not **NULL**, it points to a location that is set to the ID of the terminated thread if **thr_join( )** returns successfully.

RETURN VALUES

If successful, both **pthread_join( )** and **thr_join( )** would return **0**; otherwise, an error number is returned to indicate the error.

ERRORS

ESRCH No undetached thread could be found corresponding to that specified by the given thread ID.

If the target *target_thread* ID is **0**, **pthread_join( )** will return with error **ESRCH**.

EDEADLK A deadlock was detected or the value of *target_thread* specifies the calling thread. (See NOTES section below.)

SEE ALSO

wait(2), **pthread_create**(3T), **pthread_exit**(3T), **pthread_join**(3T)

NOTES

Using **thr_join**(3T) in the following syntax,

while (thr_join(NULL, NULL, NULL) == 0);

will wait for the termination of all other undetached and non-daemon threads; after which, **EDEADLK** will be returned.

pthread_join(3T), on the other hand, must specify the *target_thread* ID for whose termination it will wait.

Calling **pthread_join( )** also "detaches" the thread, that is, **pthread_join( )** includes the effect of **pthread_detach( )**. Hence, if a thread were to be cancelled when blocked in **pthread_join( )**, an explicit detach would have to be done in the cancellation cleanup handler. In fact, the routine **pthread_detach( )** exists mainly for this reason.

NAME

pthread_key_create, pthread_setspecific, pthread_getspecific, pthread_key_delete, thr_keycreate, thr_setspecific, thr_getspecific – thread-specific-data functions

SYNOPSIS

POSIX

cc [*flag* ...] *file* ... **–lpthread** [*library* ...]

```
#include <pthread.h>
int pthread_key_create(pthread_key_t *keyp, void (*destructor)(void *value));
int pthread_setspecific(pthread_key_t key, const void *value);
void *pthread_getspecific(pthread_key_t key);
int pthread_key_delete(pthread_key_t key);
```

Solaris

cc [*flag* ...] *file* ... **–lthread** [*library* ...]

```
#include <thread.h>
int thr_keycreate(thread_key_t *keyp, void (*destructor)(void *value));
int thr_setspecific(thread_key_t key, void *value);
int thr_getspecific(thread_key_t key, void **valuep);
```

MT-LEVEL

MT-Safe

DESCRIPTION

Create Key

In general, thread key creation allocates a key that locates data specific to each thread in the process. The key is global to all threads in the process, which allows each thread to bind a value to the key once the key has been created. The key independently maintains specific values for each binding thread. **pthread_key_create( )** or **thr_keycreate( )** allocates a global *key* namespace, pointed to by *keyp*, that is visible to all threads in the process. Each thread is initially bound to a private element of this *key*, which allows access to its thread-specific data.

Upon key creation, a new key is assigned the value **NULL** for all active threads. Additionally, upon thread creation, all previously created keys in the new thread are assigned the value **NULL.**

Optionally, a destructor function, *destructor*, may be associated with each *key*. Upon thread exit, if a *key* has a non-NULL *destructor* function and the thread has a non-NULL *value* associated with that *key*, the *destructor* function is called with the current associated *value*. If more than one *destructor* exists for a thread when it exits, the order of destructor calls is unspecified.

Set Value

Once a key has been created, each thread may bind a new *value* to the key using **pthread_setspecific( )** or **thr_setspecific( )**. The values are unique to the binding thread and are individually maintained. These values continue for the life of the calling thread.

Proper synchronization of *key* storage and access must be ensured by the caller. The *value* argument to either **pthread_setspecific( )** or **thr_setspecific( )** is generally a pointer to a block of dynamically allocated memory reserved by the calling thread for its own use. (see "Examples" section below).

At thread exit, the *destructor* function, which is associated at time of creation, is called and it uses the specific key value as its sole argument.

POSIX Get Value

pthread_getspecific() returns the current value bound to the designated *key* specified by the calling thread. If the key has no value bound to it, the value NULL is returned. (see "Warnings" section below).

Solaris Get Value

thr_getspecific() stores the current value bound to *key* for the calling thread into the location pointed to by *valuep*.

POSIX Delete Key

pthread_key_delete() deletes a thread-specific data key formerly created by **pthread_key_create()** or **thr_keycreate()**. At the time **pthread_key_delete()** is called, the thread-specific data values associated with *key* do not have to be **NULL.** It is the application's responsibility to perform cleanup actions related to the deleted key or associated thread-specific data in any threads. Cleanup can be done either before or after calling **pthread_key_delete()**. **pthread_key_delete()** does not invoke a destructor function.

Although **pthread_key_create()**'s or **thr_keycreate()**'s *destructor* function should cleanup the *key*'s thread-specific-data storage, **pthread_key_delete()** needs to be used to free the storage associated with the *key*.

Solaris threads do not have a similar delete function.

RETURN VALUES

POSIX/Solaris

If successful, **pthread_key_create()**, **pthread_setspecific()**, **pthread_key_delete()**, **thr_keycreate()**, **thr_setspecific()**, or **thr_getspecific()** returns **0**; otherwise, an error number is returned to indicate the error. **pthread_getspecific()** does not return any errors.

ERRORS

If the following conditions occur, **pthread_key_create()** or **thr_keycreate()** return the corresponding error number:

EAGAIN The system lacked the necessary resources to create another thread-specific data key, or the number of keys exceeds the pre-process limit of **PTHREAD_KEYS_MAX**.

ENOMEM Insufficient memory exists to create the key.

If the following conditions occur, **pthread_key_create()**, **pthread_setspecific()**, **thr_keycreate()**, or **thr_setspecific()** return the corresponding error number:

ENOMEM Insufficient memory exists to associate the value with the key.

For each of the following conditions, if the condition is detected, **pthread_setspecific()**, **thr_setspecific()**, or **pthread_key_delete()** return the corresponding error number:

EINVAL The *key* value is invalid.

EXAMPLES

In this example, the thread-specific data in this function can be called from more than one thread without special initialization. POSIX threads are used exclusively in this example.

For each argument you pass to the executable of this example, a thread is created and privately bound to the string-value of that argument.

```
/* cc thisfile.c -lpthread */

#define _REENTRANT
#include <pthread.h>
void *thread_specific_data(), free();
#define MAX_ARGC 20
pthread_t tid[MAX_ARGC];
int num_threads;

main( int argc, char *argv[] ) {
    int i;
    num_threads = argc - 1;
    for( i = 0; i < num_threads; i++)
        pthread_create(&tid[i], NULL, thread_specific_data, argv[i+1]);
    for( i = 0; i < num_threads; i++)
        pthread_join(tid[i], NULL);
} /* end main */

void *thread_specific_data(char private_data[])
{
 static pthread_mutex_t   keylock;  /* static ensures only one copy of keylock */
 static pthread_key_t     key;
 static int               once_per_keyname = 0;
 void *                   tsd = NULL;

 if (!once_per_keyname) {    /* see pthread_once(3T) */
       pthread_mutex_lock(&keylock);
       if (!once_per_keyname++)    /* retest with lock */
            pthread_key_create(&key, free);
       pthread_mutex_unlock(&keylock);
 }
 tsd = pthread_getspecific(key);
 if (tsd == NULL) {
       tsd = (void *)malloc(strlen(private_data) + 1);
       strcpy(tsd, private_data);
       pthread_setspecific(key, tsd);
       printf("tsd for %d = %s\n",thr_self(),(char *)pthread_getspecific(key));
       sleep(2);
       printf("tsd for %d remains %s\n",thr_self(),(char *)pthread_getspecific(key));
 }
} /* end thread_specific_data */
```

```
void
free(void *v)  {
  /* application-specific clean-up function */
}
```

SEE ALSO

pthread_exit(3T)

WARNINGS

pthread_setspecific( ), **pthread_getspecific( )**, **thr_setspecific( )**, and **thr_getspecific( )**, may be called either explicitly, or implicitly from a thread-specific data destructor function. However, calling **pthread_setspecific( )** or **thr_setspecific( )** from a destructor may result in lost storage or infinite loops.

NAME

pthread_kill, thr_kill – send a signal to a thread

SYNOPSIS

POSIX

cc [*flag* ...] *file* ... –**lpthread** [*library* ...]

```
#include <signal.h>
#include <pthread.h>
int pthread_kill(pthread_t thread, int sig);
```

Solaris

cc [*flag* ...] *file* ... –**lthread** [*library* ...]

```
#include <signal.h>
#include <thread.h>
int thr_kill(thread_t thread, int sig);
```

MT-LEVEL

MT-Safe

Async-Signal-Safe

DESCRIPTION

pthread_kill( ) sends the *sig* signal to the thread designated by *thread*. *thread* must be a member of the same process as the calling thread. *sig* must be one of the signals listed in **signal**(5); with the exception of **SIGLWP, SIGCANCEL,** and **SIGWAITING** being reserved and off limits to **thr_kill( )** or **pthread_kill( )**. If *sig* is **0**, a validity check is done for the existence of the target thread; no signal is sent.

thr_kill( ) performs the same function as **pthread_kill( )**.

RETURN VALUES

Upon successful completion, **pthread_kill( )** and **thr_kill( )** return **0**; otherwise, they return an error number. In the event of failure, no signal is sent.

ERRORS

ESRCH No thread was found that corresponded to the thread designated by *thread* ID.

EINVAL The *sig* argument value is not zero and is an invalid or an unsupported signal number.

SEE ALSO

kill(2), **sigaction**(2), **pthread_self**(3T), **pthread_sigmask**(3T), **raise**(3C), **signal**(5)

NOTES

Although **pthread_kill( )** is Async-Signal-Safe with respect to the Solaris environment, this safeness is not guaranteed to be portable to other POSIX domains.

NAME

pthread_mutex_setprioceiling, pthread_mutex_getprioceiling – change the priority ceiling of a mutex

SYNOPSIS

```
#include <pthread.h>
int pthread_mutex_setprioceiling(pthread_mutex_t *mutex, int prioceiling,
        int *old_ceiling);
int pthread_mutex_getprioceiling(const pthread_mutex_t *mutex, int *prioceiling);
```

MT-LEVEL

MT-Safe

DESCRIPTION

In the current implementation, **{_POSIX_THREAD_PRIO_PROTECT}** is undefined and the functions **pthread_mutex_setprioceiling( )** and **pthread_mutex_getprioceiling( )** return **ENOSYS**.

SEE ALSO

pthread_mutex_init(3T)

NAME

pthread_mutexattr_init, pthread_mutexattr_destroy, pthread_mutexattr_setpshared, pthread_mutexattr_getpshared, pthread_mutexattr_setprotocol, pthread_mutexattr_getprotocol, pthread_mutexattr_setprioceiling, pthread_mutexattr_getprioceiling – mutex initialization attributes

SYNOPSIS

```
#include <pthread.h>

int pthread_mutexattr_init(pthread_mutexattr_t *attr);
int pthread_mutexattr_destroy(pthread_mutexattr_t *attr);
int pthread_mutexattr_setpshared(pthread_mutexattr_t *attr, int process-shared);
int pthread_mutexattr_getpshared(const pthread_mutexattr_t *attr, int *process-shared);
```

MT-LEVEL

MT-Safe

DESCRIPTION

Initialize

pthread_mutexattr_init() initializes a mutex attributes object, *attr*, with the default value for its attribute, which is **PTHREAD_PROCESS_PRIVATE**. If the *process-shared* attribute is **PTHREAD_PROCESS_PRIVATE**, only threads created within the same process as the thread that initialized the mutex can access the mutex. If threads of differing processes attempt to access the mutex, the behavior is unpredictable.

Attempts to initialize an already initialized mutex variable attributes object will leave the storage allocated by the previous initialization unallocated.

Once a mutex attributes object is used to initialize one or more mutexes, any function that affects the attributes object (including destruction) will not affect any previously initialized mutexes.

Destroy

pthread_mutexattr_destroy() destroys a mutex attributes object; the object will then become uninitialized. A destroyed mutex attributes object can be reinitialized using **pthread_mutexattr_init()**. The results of referencing the object after it has been destroyed are undefined.

Set/Get Scope

pthread_mutexattr_setpshared() and **pthread_mutexattr_getpshared()** sets the *process-shared* attribute in an initialized attributes object pointed to by *attr*, and gets the value of the *process-shared* attribute from the attributes object pointed to by *attr*, respectively.

At present, only the attribute *process-shared* is defined.

Unsupported Interfaces

Currently, the following interfaces, which are optional under POSIX, are not supported:

```
int pthread_mutexattr_setprotocol (pthread_mutexattr_t *attr, int protocol);
```

RETURN VALUES

Upon successful completion, **pthread_mutexattr_init()**, **pthread_mutexattr_destroy()**, **pthread_mutexattr_setprotocol()**, **pthread_mutexattr_getprotocol()**, **pthread_mutexattr_setprioceiling()**, **pthread_mutexattr_getprioceiling()**, and **pthread_mutexattr_setpshared()** return **0**; otherwise, an error number is returned.

Upon successful completion, **pthread_mutexattr_getpshared( )** returns **0** and stores the value of the *process-shared* attribute of *attr* in the object pointed to by the *process-shared* parameter; otherwise, an error number is returned.

ERRORS

The function **pthread_mutexattr_init( )** returns an error number if the following condition is detected:

ENOMEM Insufficient memory exists to initialize the mutex attributes object.

The functions **pthread_mutexattr_destroy( )**, **pthread_mutexattr_getpshared( )**, and **pthread_mutexattr_setpshared( )** return an error number if the following condition is detected:

EINVAL The value specified by *attr* is invalid.

The function **pthread_mutexattr_setpshared( )** returns an error number if the following condition is detected:

EINVAL The new value specified for the attribute is outside the range of legal values for that attribute.

Currently, the functions **pthread_mutexattr_setprotocol( )**, **pthread_mutexattr_getprotocol( )**, **pthread_mutexattr_setprioceiling( )**, and **pthread_mutexattr_getprioceiling( )** always return the following error code:

ENOSYS These optional interfaces are not supported.

SEE ALSO

pthread_cond_init(3T), **pthread_create**(3T), **pthread_mutex_init**(3T)

NOTES

The functions **pthread_mutexattr_setprotocol( )**, **pthread_mutexattr_getprotocol( )**, **pthread_mutexattr_setprioceiling( )**, and **pthread_mutexattr_getprioceiling( )** return **ENOSYS** in the current implementation, i.e., this function is not currently implemented.

NAME

pthread_once – dynamic package initialization

SYNOPSIS

```
#include <pthread.h>
pthread_once_t once_control = PTHREAD_ONCE_INIT;
int pthread_once(pthread_once_t *once_control, void (*init_routine)(void));
```

MT-LEVEL

MT-Safe

DESCRIPTION

If any thread in a process with a *once_control* parameter makes a call to **pthread_once( )**, the first call will summon the **init_routine( )**, but subsequent calls will not. The *once_control* parameter determines whether the associated initialization routine has been called. The **init_routine( )** is complete upon return of **pthread_once( )**.

pthread_once( ) is not a cancellation point; however, if the function **init_routine( )** is a cancellation point and is canceled, the effect on *once_control* is the same as if **pthread_once( )** had never been called.

The constant **PTHREAD_ONCE_INIT** is defined in the **<pthread.h>** header.

If *once_control* has automatic storage duration or is not initialized by **PTHREAD_ONCE_INIT**, the behavior of **pthread_once( )** is undefined.

RETURN VALUES

pthread_once( ) returns **0** upon successful completion; otherwise, an error number is returned.

ERRORS

EINVAL *once_control* or *init_routine* is **NULL**.

NOTES

Solaris threads do not offer this functionality.

NAME

pthread_self, thr_self – get calling thread's ID

SYNOPSIS

POSIX

cc [*flag* ...] *file* ... **–lpthread** [*library* ...]

```
#include <pthread.h>
pthread_t pthread_self(void);
typedef unsigned int pthread_t;
```

Solaris

cc [*flag* ...] *file* ... **–lthread** [*library* ...]

```
#include <thread.h>
thread_t thr_self(void)
typedef unsigned int thread_t;
```

MT-LEVEL

MT-Safe

DESCRIPTION

thr_self( ) returns the thread ID of the calling thread.

pthread_self( ) performs the same function as **thr_self( )**.

SEE ALSO

pthread_create(3T), **pthread_equal**(3T)

NAME

pthread_setschedparam, pthread_getschedparam, thr_setprio, thr_getprio – dynamic access to thread scheduling

SYNOPSIS

POSIX

cc [*flag* ...] *file* ... **–lpthread** [*library* ...]
#include <pthread.h>
int pthread_setschedparam(pthread_t *target_thread,* **int** *policy,*
const struct sched_param ∗*param***);**
int pthread_getschedparam(pthread_t *target_thread,* **int** ∗*policy,*
struct sched_param ∗*param***);**

Solaris

cc [*flag* ...] *file* ... **–lthread** [*library* ...]
#include <thread.h>
int thr_setprio(thread_t *target_thread,* **int** *priority***);**
int thr_getprio(thread_t *target_thread,* **int** ∗*priority***);**

MT-LEVEL

MT-Safe

DESCRIPTION

Thread scheduling is controlled by three attributes: its scope of contention, being either inter-process or intra-process (bound vs. unbound), (see **priocntl**(2)); a relative scheduling priority; and a scheduling policy.

Contentionscope

Bound threads, which are inter-process, compete system-wide for scheduling resources and must be set at creation, for example:

pthread_attr_setscope(&attr, PTHREAD_SCOPE_SYSTEM);
pthread_create(NULL, &attr, thread_routine, arg);

OR

thr_create(NULL, NULL, thread_routine, arg, THR_BOUND, NULL);

A bound thread is bound to an LWP and its scheduling is dependent upon the scheduling of the LWP to which it is bound. LWPs compete with other LWPs in other processes, however, their scheduling may be dynamically controlled by **priocntl**(2), or **sched_setscheduler**(3R).

By default, the scope for newly-created threads are unbound, or intra-process, and their setting is **PTHREAD_SCOPE_PROCESS** or **NULL.** An unbound thread is scheduled by **libthread** or **libpthread** on an underlying LWP, which competes with other LWPs in the same process.

The following dynamic scheduling functions should be used only with unbound threads: **pthread_setschedparam( )**, **pthread_getschedparam( )**, **thr_setprio( )**, and **thr_getprio( ).**

Priority

Priority scheduling is determined as follows:

- Higher priority threads are scheduled before lower priority threads.
- Both POSIX and Solaris assume that the priority is inherited across a thread

create.

• POSIX can modify priority at creation time (see **pthread_attr_setschedparam**(3T)). Equivalently, a Solaris thread can be created suspended and its priority can be modified.

pthread_setschedparam( ) and **thr_setprio( )** can dynamically modify an unbound thread's priority, and **pthread_getschedparam( )** and **thr_getprio( )** can read an unbound thread's priority.

Policy

The scheduling *policy* setting is:

SCHED_OTHER *(system default, often time-sharing)*
Competing threads in this class are multiplexed according to their relative *priority*.

NOTE: POSIX specifies, under an option, the additional policies, SCHED_FIFO and SCHED_RR . Solaris has chosen to not implement these options at this time. Equivalent functionality may be obtained by creating bound threads (i.e., threads with the PTHREAD_SCOPE_SYSTEM value for the *contentionscope* attribute), which use **priocntl**(2). See **pthread_create**(3T) and **priocntl**(2).

POSIX Scheduling

The **pthread_setschedparam( )** and **pthread_getschedparam( )** functions allow the scheduling policy and scheduling priority parameters to be retrieved and set for individual threads within a multi-threaded process.

The **pthread_setschedparam( )** function sets the scheduling policy and related scheduling priority for the thread ID given by *target_thread* to the policy and associated priority provided in *policy*, and the **sched_priority** member of *param*, respectively.

No scheduling parameters are changed for the target thread if **pthread_setschedparam( )** fails.

For **SCHED_OTHER**, the affected scheduling parameter is the *sched_priority* member of the **sched_param** structure.

Presently, **SCHED_OTHER** is the only policy supported. An **ENOSUP** error will occur following an attempt to set policy as **SCHED_FIFO** or **SCHED_RR.** (The latter two policies are optional under POSIX.)

The **pthread_getschedparam( )** function retrieves the scheduling policy and scheduling priority parameters for the thread ID given by *target_thread*, and then stores the values in *policy* and the **sched_priority** member of *param*, respectively.

Solaris Scheduling

Solaris scheduling may only dynamically affect *priority*. There is no functionality to alter the *policy* of any thread; by default, a Solaris thread's schedule is equivalent to **SCHED_OTHER,** which is the only available Solaris policy.

thr_setprio( ) changes the priority of the thread, specified by *target_thread*, within the current process to the priority specified by *priority*. Currently, by default, threads are scheduled based on fixed priorities that range from zero, the least significant, to 127. The *target_thread* will preempt lower priority threads, and will yield to higher priority threads in their contention for LWPs, not CPUs.

The function **thr_getprio( )** stores the current priority for the thread specified by *target_thread* in the location pointed to by *priority*. Note that thread priorities regulate access to LWPs, not CPUs, and hence are different from real-time priorities, which regulate and enforce access to CPU resources. A thread's priority set via these functions is more like a hint in terms of guaranteed access to execution resources. Programs that need access to "real" priorities should use bound threads in the real-time class (see **priocntl**(2)).

RETURN VALUES

Zero is returned upon successful completion; otherwise, an error number is returned.

ERRORS

For each of the following conditions, these functions return an error number if the condition is detected.

ESRCH The value specified by *target_thread* does not refer to an existing thread.

For each of the following conditions, **pthread_setschedparam( )** and **pthread_getschedparam( )** return an error number if the condition is detected.

ENOSUP The only policy supported is **SCHED_OTHER.** Attempts to set policy as **SCHED_FIFO** or **SCHED_RR** will result in the error **ENOSUP.**

EINVAL The *policy* or *param* specified value is invalid.

For each of the following conditions, if the condition is detected, **thr_setprio( )** returns an error number.

EINVAL The value of *priority* makes no sense for the scheduling class associated with the *target_thread.*

SEE ALSO

priocntl(2), **sched_setparam**(3R), **sched_setscheduler**(3R), **pthread_attr_init**(3T), **pthread_create**(3T), **thr_suspend**(3T), **thr_yield**(3T)

NOTES

Currently, the only supported policy is **SCHED_OTHER**. Attempts to set policy as **SCHED_FIFO** or **SCHED_RR** will result in the error **ENOSUP.**

NAME

pthread_sigmask, thr_sigsetmask – change and / or examine calling thread's signal mask

SYNOPSIS

POSIX

cc [*flag* ...] *file* ... –**lpthread** [*library* ...]

```
#include <pthread.h>
#include <signal.h>
int pthread_sigmask(int how, const sigset_t *set, sigset_t *oset);
```

Solaris

cc [*flag* ...] *file* ... –**lthread** [*library* ...]

```
#include <thread.h>
#include <signal.h>
int thr_sigsetmask(int how, const sigset_t *set, sigset_t *oset);
```

MT-LEVEL

MT-Safe

Async-Signal-Safe

DESCRIPTION

pthread_sigmask( ) and **thr_sigsetmask( )** changes and / or examines a calling thread's signal mask. Each thread has its own signal mask. A new thread inherits the calling thread's signal mask and priority, however, pending signals are not inherited. Signals pending for a new thread will be empty.

If the value of the argument *set* is not **NULL,** *set* points to a set of signals that can modify the currently blocked set. If the value of *set* is **NULL**, the value of *how* is insignificant and the thread's signal mask is unmodified; thus, **pthread_sigmask( )** or **thr_sigsetmask( )** can be used to inquire about the currently blocked signals.

The value of the argument *how* specifies the method in which the set is changed. *how* takes one of the following values:

SIG_BLOCK *set* corresponds to a set of signals to block. They are added to the current signal mask.

SIG_UNBLOCK *set* corresponds to a set of signals to unblock. These signals are deleted from the current signal mask.

SIG_SETMASK *set* corresponds to the new signal mask. The current signal mask is replaced by *set*.

If the value of *oset* is not **NULL**, it points to the location where the previous signal mask is stored.

RETURN VALUES

Zero is returned upon successful completion; otherwise, a non-zero value indicates an error.

ERRORS

If any of the following conditions occur, **pthread_sigmask( )** or **thr_sigsetmask( )** fails and returns the corresponding value:

EINVAL *set* is not **NULL** and the value of *how* is not defined.

If any of the following conditions are detected, **pthread_sigmask()** or **thr_sigsetmask()** fails and returns the corresponding value:

EFAULT *set* or *oset* are not valid addresses.

EXAMPLES

The following example shows how to create a default thread that can serve as a signal catcher/handler with its own signal mask. *new* will have a different value than the creator's signal mask.

```
/* cc thisfile.c -lthread -lpthread */
#define _REENTRANT    /* basic first 3-lines for threads */
#include <pthread.h>
#include <thread.h>

thread_t user_threadID;
sigset_t new;
void *handler(), interrupt();

main( int argc, char *argv[] ) {
 test_argv(argv[1]);

 sigemptyset(&new);
 sigaddset(&new, SIGINT);
 switch(*argv[1]) {

  case '0':   /* POSIX */
   pthread_sigmask(SIG_BLOCK, &new, NULL);
   pthread_create(&user_threadID, NULL, handler, argv[1]);
   pthread_join(user_threadID, NULL);
   break;

  case '1':   /* Solaris */
   thr_sigsetmask(SIG_BLOCK, &new, NULL);
   thr_create(NULL, 0, handler, argv[1], 0, &user_threadID);
   thr_join(user_threadID, NULL, NULL);
   break;
  } /* switch */

  printf("thread handler, # %d, has exited\n",user_threadID);
  sleep(2);
  printf("main thread, # %d is done\n", thr_self());
} /* end main */

struct sigaction act;

void *
handler(char argv1[])
```

```
{
 act.sa_handler = interrupt;
 sigaction(SIGINT, &act, NULL);
 switch(*argv1) {
  case '0':    /* POSIX */
   pthread_sigmask(SIG_UNBLOCK, &new, NULL);
   break;
  case '1':  /* Solaris */
   thr_sigsetmask(SIG_UNBLOCK, &new, NULL);
   break;
 }
 printf("\n Press cntrl-C to deliver SIGINT signal to the process\n");
 sleep(8);  /* give user time to hit cntrl-C */
}

void
interrupt(int sig)
{
 printf("thread %d caught signal %d\n", thr_self(), sig);
}

void test_argv(char argv1[])   {
 if(argv1 == NULL)  {
   printf("use 0 as arg1 to use thr_create();\n \
   or use 1 as arg1 to use pthread_create()\n");
   exit(NULL);
 }
}
```

Since POSIX threads and Solaris threads are fully compatible even within the same process, this example uses **pthread_create**(3T) if you execute **a.out 0**, or **thr_create**(3T) if you execute **a.out 1**.

Here's an explanation of the above example:

- **sigemptyset**(3C) initializes a null signal set, *new*. **sigaddset**(3C) packs the signal, **SIGINT,** into that new set.
- Either **pthread_sigmask( )** or **thr_sigsetmask( )** is used to mask the signal, **SIGINT,** (cntrl-C), from the calling thread, which is **main( )**. The signal is masked to guarantee that only the new thread will receive this signal.
- **pthread_create( )** or **thr_create( )** creates the signal-handling thread.
- Using **pthread_join**(3T) or **thr_join**(3T), **main( )** then waits for the termination of that signal-handling thread, whose ID number is *user_threadID*; after which, **main( )** will **sleep**(3C) for 2 seconds, and then the program terminates.

- The signal-handling thread, *handler*:
  - Assigns the handler *interrupt( )* to handle the signal SIGINT , via the call to **sigaction**(2).
  - Resets its own signal set to NOT BLOCK the signal, **SIGINT.**
  - Sleeps for 8 seconds to allow time for the user to deliver the signal, **SIGINT,** by pressing the cntrl-C keys.

In the example, the *handler* thread served as a signal-handler while also taking care of activity of its own (in this case, sleeping, although it could have been some other activity). A thread could be completely dedicated to signal-handling, simply by waiting for the delivery of a selected signal by blocking with **sigwait**(2). Thus, the two subroutines in the previous example, *handler( )* and *interrupt( )*, could have been replaced with the following routine:

```
void *
handler()
{ int signal;
  printf("thread %d is waiting for you to press the cntrl-C keys\n", thr_self());
  sigwait(&new, &signal);
  printf("thread %d has received the signal %d \n", thr_self(), signal);
}
  /* pthread_create( ) and thr_create( ) would use NULL instead of argv[1]
     for the arg passed to handler( ) */
```

In this routine, one thread is dedicated to catching and handling the signal specified by the set, *new*, which allows **main( )** and all of its other sub-threads, (created AFTER **pthread_sigmask( )** or **thr_sigsetmask( )** masked that signal), to continue uninterrupted. In fact, any use of **sigwait**(2) should be such that all threads block the signals passed to **sigwait**(2), at all times. Only the thread that calls **sigwait( )** will get the signals. Note that the call to **sigwait**(2) takes two arguments. See **sigwait**(2).

For this type of background dedicated signal-handling routine, you may wish to use a Solaris daemon thread by passing the argument, **THR_DAEMON,** to **thr_create**(3T).

SEE ALSO

sigaction(2), **sigprocmask**(2), **sigwait**(2), **sigsetops**(3C), **pthread_cancel**(3T), **pthread_create**(3T), **pthread_exit**(3T), **pthread_join**(3T), **pthread_kill**(3T), **pthread_self**(3T)

NOTES

It is not possible to block signals that cannot be ignored (see **sigaction**(2)). If using the threads library, it is not possible to block the signals **SIGLWP** or **SIGCANCEL,** which are reserved by the threads library. Additionally, it is impossible to unblock the signal **SIGWAITING,** which is always blocked on all threads. This restriction is quietly enforced by the threads library.

Using **sigwait**(2) in a dedicated thread allows asynchronous signals to be managed synchronously; however, **sigwait**(2) should never be used to manage synchronous signals. Synchronous signals (i.e., exceptions or traps) are sent by the process itself, such as **SIGFPE, pthread_kill**(3T), **pthread_exit**(3T), **pthread_cancel**(3T), **thr_kill**(3T), or

thr_exit(3T), rather than device interrupts or signals sent by other processes.

Synchronous signals are exceptions that are generated by a thread and are directed at the thread causing the exception. Since **sigwait()** blocks waiting for signals, the blocking thread will not generate any synchronous signals.

If **sigprocmask**(2) is used in a multi-threaded program, it will be the same as if **thr_sigsetmask()** or **pthread_sigmask()** has been called. Note that POSIX leaves the semantics of the call to **sigprocmask**(2) unspecified in a multi-threaded process, so programs that care about POSIX portability should not depend on this semantic.

If a signal is delivered while a thread is waiting on a condition variable, the **cond_wait()** will be interrupted and the handler will be executed. The handler should assume that the lock protecting the condition variable is held.

Although **pthread_sigmask()** is Async-Signal-Safe with respect to the Solaris environment, this safeness is not guaranteed to be portable to other POSIX domains.

NAME

ptsname – get name of the slave pseudo-terminal device

SYNOPSIS

```
#include <stdio.h>
char *ptsname(int fildes);
```

MT-LEVEL

Safe

DESCRIPTION

The **ptsname( )** function returns the name of the slave pseudo-terminal device associated with a master pseudo-terminal device. *fildes* is a file descriptor returned from a successful open of the master device. **ptsname( )** returns a pointer to a string containing the null-terminated path name of the slave device of the form **/dev/pts/N**, where **N** is a non-negative integer.

RETURN VALUES

Upon successful completion, the function **ptsname( )** returns a pointer to a string which is the name of the pseudo-terminal slave device. This value points to a static data area that is overwritten by each call to **ptsname( )**. Upon failure, **ptsname( )** returns **NULL**. This could occur if *fildes* is an invalid file descriptor or if the slave device name does not exist in the file system.

SEE ALSO

open(2), **grantpt**(3C), **ttyname**(3C), **unlockpt**(3C)

STREAMS Programming Guide

NAME

putc, putc_unlocked, putchar, putchar_unlocked, fputc, putw – put character or word on a stream

SYNOPSIS

```
#include <stdio.h>
int putc(int c, FILE *stream);
int putc_unlocked(int c, FILE *stream);
int putchar(int c);
int putchar_unlocked(int c);
int fputc(int c, FILE *stream);
int putw(int w, FILE *stream);
```

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

putc( ) writes *c* (converted to an **unsigned char**) onto the output *stream* (see **intro**(3)) at the position where the file pointer (if defined) is pointing, and advances the file pointer appropriately. If the file cannot support positioning requests, or *stream* was opened with append mode, the character is appended to the output *stream*. **putchar(***c***)** is defined as **putc(***c***,** *stdout***)**. **putc( )** and **putchar( )** are macros.

putc_unlocked( ) and **putchar_unlocked( )** are respectively variants of **putc( )** and **putchar( )** that do not lock the stream. It is the caller's responsibility to acquire the stream lock before calling these functions and releasing the lock afterwards; see **flockfile**(3S) and **stdio**(3S).

fputc( ) behaves like **putc( )**, but is a function rather than a macro. **fputc( )** runs more slowly than **putc( )**, but it takes less space per invocation and its name can be passed as an argument to a function.

putw( ) writes the C **int** (word) *w* to the standard I/O output *stream* (at the position of the file pointer, if defined). The size of a word is the size of an integer and varies from machine to machine. **putw( )** neither assumes nor causes special alignment in the file.

RETURN VALUES

On success, **putc( )**, **fputc( )**, and **putchar( )** return the value that was written. On error, those functions return the constant **EOF**. **putw( )** returns **ferror(stream)**, so that it returns 0 on success and 1 on failure.

Failure will occur, for example, if the file *stream* is not open for writing or if the output file cannot grow.

SEE ALSO

write(2), **intro**(3), **fclose**(3S), **ferror**(3S), **flockfile**(3S), **fopen**(3S), **printf**(3S), **puts**(3S), **setbuf**(3S), **stdio**(3S)

NOTES

Because it is implemented as a macro, **putc( )** evaluates a *stream* argument more than once. In particular, **putc(***c***,** **f++***);** does not work sensibly. **fputc( )** should be used instead.

Because of possible differences in word length and byte ordering, files written using **putw( )** are machine-dependent, and may not be read using **getw( )** on a different processor.

Functions exist for all the above defined macros. To get the function form, the macro name must be undefined (for example, **#undef putc**).

fputc( ), putc( ), putchar( ), and **putw( )** are MT-Safe in multi-thread applications. **putc_unlocked( )** and **putchar_unlocked( )** are unsafe in multi-thread applications.

NAME

putenv – change or add value to environment

SYNOPSIS

```
#include <stdlib.h>
int putenv(const char *string);
```

MT-LEVEL

Safe

DESCRIPTION

putenv( ) makes the value of the environment variable *name* equal to *value* by altering an existing variable or creating a new one. In either case, the string pointed to by *string* becomes part of the environment, so altering the string will change the environment. *string* points to a string of the form *''name=value.''* The space used by *string* is no longer used once a new string-defining *name* is passed to **putenv( )**.

RETURN VALUES

putenv( ) returns non-zero if it was unable to obtain enough space using **malloc( )** for an expanded environment, otherwise zero is returned.

SEE ALSO

exec(2), **getenv**(3C), **malloc**(3C), **environ**(5)

NOTES

This routine uses **malloc**(3C) to enlarge the environment.

After **putenv( )** is called, environment variables are not in alphabetical order.

string should not be an automatic variable.

string should be declared static if it is declared within a function because it cannot be automatically declared.

A potential error is to call the function **putenv( )** with a pointer to an automatic variable as the argument and to then exit the calling function while *string* is still part of the environment.

putenv( ) can be safely called from a multi-thread program. However, care must still be taken when using **putenv( )** and **getenv**(3C) in a multi-thread program. These routines examine and modify the environment list. This list is shared by all threads in a program. The system prevents the list from being accessed simultaneously by two different threads. However, it does not prevent two threads from successively accessing the environment list using **putenv( )** or **getenv**(3C).

NAME

putpwent – write password file entry

SYNOPSIS

```
#include <pwd.h>
int putpwent(const struct passwd *p, FILE *f);
```

MT-LEVEL

Unsafe

DESCRIPTION

putpwent( ) is the inverse of **getpwent( )**, (see **getpwnam**(3C)). Given a pointer to a **passwd** structure created by **getpwent( )** (or **getpwuid( )** or **getpwnam( )**), **putpwent( )** writes a line on the stream *f*, which matches the format of **/etc/passwd**.

RETURN VALUES

putpwent( ) returns non-zero if an error was detected during its operation, otherwise zero.

SEE ALSO

getpwnam(3C), **putspent**(3C)

NOTES

Do not use without also using **putspent( )** to update the shadow file.

The use of this function is discouraged.

BUGS

This routine is of limited utility, since most password files are maintained as Network Information Service (NIS) files, and cannot be updated with this routine.

NAME

puts, fputs – put a string on a stream

SYNOPSIS

```
#include <stdio.h>
int puts(const char *s);
int fputs(const char *s, FILE *stream);
```

MT-LEVEL

MT-Safe

DESCRIPTION

puts( ) writes the string pointed to by *s*, followed by a new-line character, to the standard output stream **stdout** (see **intro**(3)).

fputs( ) writes the null-terminated string pointed to by *s* to the named output *stream*.

Neither function writes the terminating null character.

RETURN VALUES

On success both routines return the number of characters written; otherwise they return **EOF**.

SEE ALSO

write(2), **intro**(3)), **fclose**(3S), **ferror**(3S), **fopen**(3S), **printf**(3S), **putc**(3S), **stdio**(3S)

NOTES

puts( ) appends a new-line character while **fputs( )** does not.

NAME

putspent – write shadow password file entry

SYNOPSIS

```
#include <shadow.h>
int putspent(const struct spwd *p, FILE *fp);
```

MT-LEVEL

Unsafe

DESCRIPTION

The **putspent( )** routine is the inverse of **getspent( )**. Given a pointer to a **spwd** structure created by the **getspent( )** routine (or the **getspnam( )** routine), the **putspent( )** routine writes a line on the stream *fp*, which matches the format of **/etc/shadow**. The **spwd** structure contains the following members:

```
char     *sp_namp;
char     *sp_pwdp;
long     sp_lstchg;
long     sp_min;
long     sp_max;
long     sp_warn;
long     sp_inact;
long     sp_expire;
unsigned long  sp_flag;
```

If the **sp_min**, **sp_max**, **sp_lstchg**, **sp_warn**, **sp_inact**, or **sp_expire** field of the **spwd** structure is –1, or if **sp_flag** is 0, the corresponding **/etc/shadow** field is cleared.

RETURN VALUES

The **putspent( )** routine returns non-zero if an error was detected during its operation, otherwise zero.

SEE ALSO

getpwnam(3C), **getspnam**(3C), **putpwent**(3C)

NOTES

This routine is for internal use only, compatibility is not guaranteed.

Do not use without also using **putpwent( )** to update the password file.

The use of this function is discouraged.

NAME

putwc, putwchar, fputwc – convert Process Code character to EUC and put on a stream

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <stdio.h>
#include <widec.h>

wchar_t putwc(int *c***, FILE** ∗*stream***);**

wchar_t putwchar(int *c***);**

wchar_t fputwc(int *c***, FILE** ∗*stream***);**

MT-LEVEL

MT-Safe

DESCRIPTION

putwc( ) and **fputwc( )** convert the Process Code (*wchar_t*) character *c* to Extended Unix Code (EUC) and write it onto the named output *stream* (at the position where the file pointer, if defined, is pointing). It returns the character written.

putwchar(**c**) is defined as **fputwc**(**c**,**stdout**). **putwc( )** and **putwchar( )** are macros.

See **stdio**(3S) for a discussion of output streams.

RETURN VALUES

On success, these functions each return the value passed. On error, these functions return the constant **EOF**.

SEE ALSO

fclose(3S), **ferror**(3S), **fopen**(3S), **fread**(3S), **getwc**(3I), **printf**(3S), **putc**(3S), **putws**(3I), **setbuf**(3S), **stdio**(3S)

NAME

putws, fputws – convert a string of Process Code characters to EUC characters and put it on a stream

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <stdio.h>
#include <widec.h>

int putws(wchar_t ∗*s***);**

int fputws(wchar_t ∗*s***, FILE** ∗*stream***);**

MT-LEVEL

MT-Safe

DESCRIPTION

putws( ) converts the Process Code string (terminated by a (*wchar_t*)NULL) pointed to by *s*, to an Extended Unix Code (EUC) string followed by a NEWLINE character, and writes it to the standard output stream **stdout**.

fputws( ) writes the (*wchar_t*)NULL-terminated string pointed to by *s* to the named output stream, and does *not* append a NEWLINE.

Neither function writes the terminal NULL character.

RETURN VALUES

Both routines return the number of Process Code characters transformed and written. Both routines return **EOF** on error. This will happen if the routines try to write on a file that has not been opened for writing.

SEE ALSO

ferror(3S), **fopen**(3S), **fread**(3S), **getws**(3I), **printf**(3S), **putwc**(3I)

NAME

qsort – quick sort

SYNOPSIS

```
#include <stdlib.h>

void qsort(void *base, size_t nel, size_t width,
        int (*compar) (const void *, const void *));
```

MT-LEVEL

Safe

DESCRIPTION

qsort( ) is an implementation of the quick-sort algorithm. It sorts a table of data in place. The contents of the table are sorted in ascending order according to the user-supplied comparison function.

base points to the element at the base of the table. *nel* is the number of elements in the table. *width* specifies the size of each element in bytes. *compar* is the name of the comparison function, which is called with two arguments that point to the elements being compared. The function must return an integer less than, equal to, or greater than zero to indicate if the first argument is to be considered less than, equal to, or greater than the second argument.

The contents of the table are sorted in ascending order according to the user supplied comparison function.

EXAMPLES

The following program sorts a simple array:

```
static  int intcompare(int *i, int *j)
{
        if (*i > *j)
                return (1);
        if (*i < *j)
                return (–1);
        return (0);
}

main( )
{
        int a[10];
        int i;

        a[0] = 9;
        a[1] = 8;
        a[2] = 7;
        a[3] = 6;
        a[4] = 5;
        a[5] = 4;
        a[6] = 3;
        a[7] = 2;
        a[8] = 1;
        a[9] = 0;
```

```
        qsort((char *) a, 10, sizeof(int), intcompare);
        for (i=0; i<10; i++) printf(" %d",a[i]);
        printf("\n");
}
```

SEE ALSO

sort(1), **bsearch**(3C), **lsearch**(3C), **string**(3C)

NOTES

The comparison function need not compare every byte, so arbitrary data may be contained in the elements in addition to the values being compared.

The relative order in the output of two items that compare as equal is unpredictable.

NAME

raise – send signal to program

SYNOPSIS

```
#include <signal.h>
int raise(int sig);
```

MT-LEVEL

MT-Safe

DESCRIPTION

raise( ) sends the signal *sig* to the executing program.

raise( ) uses **kill( )** to send the signal to the executing program:

```
kill(getpid( ), sig);
```

See **kill**(2) for a detailed list of failure conditions. See **signal**(3C) for a list of signals.

RETURN VALUES

raise( ) returns zero if the operation succeeds. Otherwise, **raise( )** returns –1 and **errno** is set to indicate the error.

SEE ALSO

getpid(2), **kill**(2), **signal**(3C)

NAME

rand, srand – simple random number generator

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

int rand( )

int srand(*seed***)**
unsigned *seed***;**

DESCRIPTION

rand( ) uses a multiplicative congruential random number generator with period 2^{32} to return successive pseudo-random numbers in the range from 0 to "$2^{31} - 1$."

srand( ) can be called at any time to reset the random-number generator to a random starting point. The generator is initially seeded with a value of 1.

SEE ALSO

drand48(3C), **rand**(3C), **random**(3C)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

The spectral properties of **rand( )** leave a great deal to be desired. **drand48**(3C) and **random**(3C) provide much better, though more elaborate, random-number generators.

The low bits of the numbers generated are not very random; use the middle bits. In particular the lowest bit alternates between 0 and 1.

NAME

rand, srand, rand_r – simple random-number generator

SYNOPSIS

```
#include <stdlib.h>
int rand(void);
void srand(unsigned int seed);
int rand_r(unsigned int *seed);
```

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

rand( ) uses a multiplicative congruential random-number generator with period 2^{32} that returns successive pseudo-random numbers in the range from 0 to **RAND_MAX** (defined in **<stdlib.h>**).

The function **srand( )** uses the argument *seed* as a seed for a new sequence of pseudo-random numbers to be returned by subsequent calls to the function **rand( ).** If the function **srand( )** is then called with the same *seed* value, the sequence of pseudo-random numbers will be repeated. If the function **rand( )** is called before any calls to **srand( )** have been made, the same sequence will be generated as when **srand( )** is first called with a *seed* value of 1.

rand_r( ) has the same functionality as **rand( )** except that a pointer to a seed *seed* must be supplied by the caller. The seed to be supplied is not the same seed as in **srand( ).**

SEE ALSO

drand48(3C)

NOTES

The **rand_r( )** interface is as proposed in the POSIX.4a Draft #6 document, and is subject to change to be compliant to the standard when it is accepted.

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

The spectral properties of **rand( )** are limited. **drand48**(3C) provides a much better, though more elaborate, random-number generator.

rand( ) is unsafe in multi-thread applications. **rand_r( )** is MT-Safe, and should be used instead. **srand( )** is unsafe in multi-thread applications.

NAME

random, srandom, initstate, setstate – better random number generator; routines for changing generators

SYNOPSIS

long random( );

int srandom(unsigned *seed***);**

char ∗**initstate(unsigned** *seed***, char** ∗*state***, int** *n***);**

char ∗**setstate(char** ∗*state***);**

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

random( ) uses a non-linear additive feedback random number generator employing a default table of size 31 long integers to return successive pseudo-random numbers in the range from 0 to $2^{31}-1$. The period of this random number generator is very large, approximately $16\times(2^{31}-1)$.

random( ) and **srandom( )** have (almost) the same calling sequence and initialization properties as **rand( )** and **srand( )** (see **rand**(3C)). The difference is that **rand**(3C) produces a much less random sequence—in fact, the low dozen bits generated by rand go through a cyclic pattern. All the bits generated by **random( )** are usable. For example,

random( )&01

will produce a random binary value.

Unlike **srand( )**, **srandom( )** does not return the old seed because the amount of state information used is much more than a single word. Two other routines are provided to deal with restarting / changing random number generators. Like **rand**(3C), however, **random( )** will, by default, produce a sequence of numbers that can be duplicated by calling **srandom( )** with **1** as the seed.

The **initstate( )** routine allows a state array, passed in as an argument, to be initialized for future use. *n* specifies the size of *state* in bytes. **initstate( )** uses *n* to decide how sophisticated a random number generator it should use—the more state, the better the random numbers will be. Current "optimal" values for the amount of state information are 32, 64, 128, and 256 bytes. If the amount of state information is less than 32 bytes, a simple linear congruential random number generator is used. Using less than 8 bytes causes an error. The seed for the initialization (which specifies a starting point for the random number sequence, and provides for restarting at the same point) is also an argument. **initstate( )** returns a pointer to the previous state information array.

Once a state has been initialized, the **setstate( )** routine provides for rapid switching between states. **setstate( )** returns a pointer to the previous state array; its argument state array is used for further random number generation until the next call to **initstate( )** or **setstate( )**.

Once a state array has been initialized, it may be restarted at a different point either by calling **initstate( )** (with the desired seed, the state array, and its size) or by calling both **setstate( )** (with the state array) and **srandom( )** (with the desired seed). The advantage of calling both **setstate( )** and **srandom( )** is that the size of the state array does not have to

be remembered after it is initialized.

With 256 bytes of state information, the period of the random number generator is greater than 2^{69}, which should be sufficient for most purposes.

RETURN VALUES

If **initstate( )** is called with less than 8 bytes of state information, or if **setstate( )** detects that the state information has been garbled, error messages are printed on the standard error output.

EXAMPLES

```
/* Initialize an array and pass it in to initstate. */
static long state1[32] = {
        3,
        0x9a319039, 0x32d9c024, 0x9b663182, 0x5da1f342,
        0x7449e56b, 0xbeb1dbb0, 0xab5c5918, 0x946554fd,
        0x8c2e680f, 0xeb3d799f, 0xb11ee0b7, 0x2d436b86,
        0xda672e2a, 0x1588ca88, 0xe369735d, 0x904f35f7,
        0xd7158fd6, 0x6fa6f051, 0x616e6b96, 0xac94efdc,
        0xde3b81e0, 0xdf0a6fb5, 0xf103bc02, 0x48f340fb,
        0x36413f93, 0xc622c298, 0xf5a42ab8, 0x8a88d77b,
        0xf5ad9d0e, 0x8999220b, 0x27fb47b9
        };
main( ) {
        unsigned seed;
        int n;
        seed = 1;
        n = 128;
        initstate(seed, state1, n);
        setstate(state1);
        printf("%d0,random( ));
}
```

SEE ALSO

drand48(3C), **rand**(3C)

NOTES

random( ) and **srandom( )** are unsafe in multi-thread applications.

Use of these interfaces in multi-thread applications is unsupported.

random( ) and **srandom( )** function at about two-thirds the speed of **rand**(3C).

NAME

rcmd, rresvport, ruserok – routines for returning a stream to a remote command

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket -lnsl** [*library* ...]

int rcmd(char ***ahost*, **unsigned short** *inport*, **const char** **luser*, **const char** **ruser*, **const char** **cmd*, **int** **fd2p);*

int rresvport(int **port);*

int ruserok(const char **rhost*, **int** *suser*, **const char** **ruser*, **const char** **luser);*

MT-LEVEL

Unsafe

DESCRIPTION

rcmd( ) is a routine used by the super-user to execute a command on a remote machine using an authentication scheme based on reserved port numbers. **rresvport( )** is a routine which returns a descriptor to a socket with an address in the privileged port space. **ruserok( )** is a routine used by servers to authenticate clients requesting service with **rcmd**. All three functions are present in the same file and are used by the **in.rshd**(1M) server (among others).

rcmd( ) looks up the host **ahost* using **gethostbyname**(3N), returning –1 if the host does not exist. Otherwise **ahost* is set to the standard name of the host and a connection is established to a server residing at the well-known Internet port *inport*.

If the connection succeeds, a socket in the Internet domain of type **SOCK_STREAM** is returned to the caller, and given to the remote command as its standard input (file descriptor 0) and standard output (file descriptor 1). If *fd2p* is non-zero, then an auxiliary channel to a control process will be set up, and a descriptor for it will be placed in **fd2p*. The control process will return diagnostic output from the command (file descriptor 2) on this channel, and will also accept bytes on this channel as signal numbers, to be forwarded to the process group of the command. If *fd2p* is 0, then the standard error (file descriptor 2) of the remote command will be made the same as its standard output and no provision is made for sending arbitrary signals to the remote process, although you may be able to get its attention by using out-of-band data.

The protocol is described in detail in **in.rshd**(1M).

The **rresvport( )** routine is used to obtain a socket bound to a privileged port number. This socket is suitable for use by **rcmd( )** and several other routines. Privileged Internet ports are those in the range 1 to 1023. Only the super-user is allowed to bind a socket to a privileged port number. The application must pass in *port*, which must be in the range 512 to 1023. The system first tries to bind to that port number. If it fails, it then tries to bind to port numbers less than *port* until either it succeeds or port number 512 is reached.

ruserok( ) takes a remote host's name, as returned by a **gethostbyaddr( )** (see **gethostbyname**(3N)) routine, two user names and a flag indicating whether the local user's name is that of the super-user. It then checks the files **/etc/hosts.equiv** and possibly **.rhosts** in the local user's home directory to see if the request for service is allowed. **0** is returned if the machine name is listed in the **/etc/hosts.equiv** file, or the host and remote user name are found in the **.rhosts** file; otherwise **ruserok( )** returns **–1**. If the

super-user flag is **1**, the checking of the **/etc/hosts.equiv** file is bypassed.

RETURN VALUES

rcmd( ) returns a valid socket descriptor on success. It returns **–1** on error and prints a diagnostic message on the standard error.

rresvport( ) returns a valid, bound socket descriptor on success. It returns **–1** on error with the global value **errno** set according to the reason for failure.

FILES

/etc/hosts.equiv	system trusted hosts and users
~/.rhosts	user's trusted hosts and users

SEE ALSO

rlogin(1), **rsh**(1), **in.rexecd**(1M), **in.rshd**(1M), **intro**(2), **gethostbyname**(3N), **rexec**(3N)

NOTES

The error code **EAGAIN** is overloaded to mean "All network ports in use."

These interfaces are unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

read_vtoc, write_vtoc – read and write a disk's VTOC

SYNOPSIS

```
#include <sys/vtoc.h>
cc [ flag ... ] file ... –ladm [ library ... ]
int read_vtoc(int fd, struct vtoc *vtoc);
int write_vtoc(int fd, struct vtoc *vtoc);
```

MT-LEVEL

Unsafe

DESCRIPTION

read_vtoc( ) returns the VTOC structure that is stored on the disk associated with the open file descriptor *fd*.

write_vtoc( ) stores the VTOC structure on at disk associated with the open file descriptor *fd*.

fd refers to any slice on a raw disk.

RETURN VALUES

read_vtoc returns:

positive number	Success. The positive number is the slice index associated with the open file descriptor.
negative number	There are two possible error returns. **VT_EIO** indicates an I/O error occurred and **VT_ERROR** indicates an unknown error.

write_vtoc returns:

0	Success
negative number	There are three possible error returns. **VT_EIO** indicates an I/O error occurred, **VT_ERROR** indicates an unknown error, and **VT_EINVAL** indicates an incorrect field within the VTOC.

SEE ALSO

format(1M), **fmthard**(1M), **prtvtoc**(1M), **ioctl**(2), **dkio**(7I)

BUGS

write_vtoc cannot write a VTOC on an unlabeled disk. Use **format**(1M) for this purpose.

NAME

readdir – read a directory entry

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

```
#include <sys/types.h>
#include <sys/dir.h>
struct direct *readdir(dirp);
DIR *dirp;
```

DESCRIPTION

readdir() returns a pointer to a structure representing the directory entry at the current position in the directory stream to which *dirp* refers, and positions the directory stream at the next entry, except on read-only filesystems. It returns a **NULL** pointer upon reaching the end of the directory stream, or upon detecting an invalid location in the directory. **readdir()** shall not return directory entries containing empty names. It is unspecified whether entries are returned for dot or dot-dot. The pointer returned by **readdir()** points to data that may be overwritten by another call to **readdir()** on the same directory stream. This data shall not be overwritten by another call to **readdir()** on a different directory stream. **readdir()** may buffer several directory entries per actual read operation. **readdir()** marks for update the *st_atime* field of the directory each time the directory is actually read.

RETURN VALUES

readdir() returns **NULL** on failure and sets **errno** to indicate the error.

ERRORS

readdir() will fail if one or more of the following are true:

EAGAIN	Mandatory file/record locking was set, **O_NDELAY** or **O_NONBLOCK** was set, and there was a blocking record lock.
EAGAIN	Total amount of system memory available when reading using raw I/O is temporarily insufficient.
EAGAIN	No data is waiting to be read on a file associated with a tty device and **O_NONBLOCK** was set.
EAGAIN	No message is waiting to be read on a stream and **O_NDELAY** or **O_NONBLOCK** was set.
EBADF	The file descriptor determined by the **DIR** stream is no longer valid. This results if the **DIR** stream has been closed.
EBADMSG	Message waiting to be read on a stream is not a data message.
EDEADLK	The **read()** was going to go to sleep and cause a deadlock to occur.
EFAULT	*buf* points to an illegal address.
EINTR	A signal was caught during the **read()** or **readv()** function.
EINVAL	Attempted to read from a stream linked to a multiplexor.
EIO	A physical I/O error has occurred, or the process is in a background process group and is attempting to read from its controlling terminal, and either the process is ignoring or blocking the **SIGTTIN** signal or the

process group of the process is orphaned.

ENOENT The current file pointer for the directory is not located at a valid entry.

ENOLCK The system record lock table was full, so the **read()** or **readv()** could not go to sleep until the blocking record lock was removed.

ENOLINK *fildes* is on a remote machine and the link to that machine is no longer active.

ENXIO The device associated with *fildes* is a block special or character special file and the value of the file pointer is out of range.

SEE ALSO **getdents**(2), **scandir**(3B), **directory**(3C)

NOTES Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

NAME

realpath – returns the real file name

SYNOPSIS

```
#include <stdlib.h>
#include <sys/param.h>

char *realpath(char *file_name, char *resolved_name);
```

MT-LEVEL

MT-Safe

DESCRIPTION

realpath( ) resolves all links and references to ‘.’’ and ‘..’’ in *file_name* and stores it in *resolved_name.*

It can handle both relative and absolute path names. For absolute path names and the relative names whose resolved name cannot be expressed relatively (for example, **../../reldir**), it returns the *resolved absolute* name. For the other relative path names, it returns the *resolved relative* name.

resolved_name must be big enough (**MAXPATHLEN**) to contain the fully resolved path name.

RETURN VALUES

If there is no error, **realpath( )** returns a pointer to the *resolved_name.* Otherwise it returns a null pointer and places the name of the offending file in *resolved_name.* The global variable **errno** is set to indicate the error.

SEE ALSO

getcwd(3C), **sysconf**(3C)

NOTES

realpath( ) operates on null-terminated strings.

One should have execute permission on all the directories in the given and the resolved path.

realpath( ) may fail to return to the current directory if an error occurs.

NAME

reboot – reboot system or halt processor

SYNOPSIS

#include <sys/reboot.h>

int reboot(int *howto***, char** ∗*bootargs***);**

DESCRIPTION

reboot( ) reboots the system. *howto* is an option passed to specify the behaviour of the system while rebooting. The function interface permits only one of **RB_HALT, RB_ASKNAME** or **RB_AUTOBOOT** to be passed. **RB_AUTOBOOT** is the default.

The *howto* options are:

RE_AUTOBOOT The machine is rebooted from the root filesystem on the default boot device. See **boot**(1M) and **kernel**(1M).

RB_HALT the processor is simply halted; no reboot takes place. **RB_HALT** should be used with caution.

RB_ASKNAME Interpreted by the bootstrap program and kernel, causing the user to be asked for pathnames during the bootstrap.

The interpretation of the *bootargs* argument is platform dependent.

RETURN VALUES

If successful, this call never returns. Otherwise, a –1 is returned and an error is returned in the global variable **errno**.

ERRORS

EPERM The caller is not the super-user.

SEE ALSO

intro(1M), **boot**(1M), **halt**(1M), **init**(1M), **kernel**(1M), **reboot**(1M), **uadmin**(2)

NOTES

Any other *howto* argument causes the kernel file to boot.

Only the super-user may **reboot( )** a machine.

NAME

recv, recvfrom, recvmsg – receive a message from a socket

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>
#include <sys/uio.h>

int recv(int s, char *buf, int len, int flags);

int recvfrom(int s, char *buf, int len, int flags, struct sockaddr *from, int *fromlen);

int recvmsg(int s, struct msghdr *msg, int flags);
```

MT-LEVEL

Safe

DESCRIPTION

recv( ), **recvfrom( )**, and **recvmsg( )** are used to receive messages from another socket. **recv( )** may be used only on a *connected* socket (see **connect**(3N)), while **recvfrom( )** and **recvmsg( )** may be used to receive data on a socket whether it is in a connected state or not. *s* is a socket created with **socket**(3N).

If *from* is not a NULL pointer, the source address of the message is filled in. *fromlen* is a value-result parameter, initialized to the size of the buffer associated with *from*, and modified on return to indicate the actual size of the address stored there. The length of the message is returned. If a message is too long to fit in the supplied buffer, excess bytes may be discarded depending on the type of socket the message is received from (see **socket**(3N)).

If no messages are available at the socket, the receive call waits for a message to arrive, unless the socket is nonblocking (see **fcntl**(2)) in which case **-1** is returned with the external variable **errno** set to **EWOULDBLOCK**.

The **select( )** call may be used to determine when more data arrives.

The *flags* parameter is formed by ORing one or more of the following:

MSG_OOB Read any “out-of-band” data present on the socket rather than the regular “in-band” data.

MSG_PEEK “Peek” at the data present on the socket; the data is returned, but not consumed, so that a subsequent receive operation will see the same data.

The **recvmsg( )** call uses a **msghdr** structure to minimize the number of directly supplied parameters. This structure is defined in **<sys/socket.h>** and includes the following members:

```
caddr_t      msg_name;        /* optional address */
int          msg_namelen;     /* size of address */
struct iovec *msg_iov;        /* scatter/gather array */
int          msg_iovlen;      /* # elements in msg_iov */
caddr_t      msg_accrights;   /* access rights sent/received */
int          msg_accrightslen;
```

Here **msg_name** and **msg_namelen** specify the destination address if the socket is unconnected; **msg_name** may be given as a NULL pointer if no names are desired or required. The **msg_iov** and **msg_iovlen** describe the scatter-gather locations, as described in **read**(2). A buffer to receive any access rights sent along with the message is specified in **msg_accrights**, which has length **msg_accrightslen**.

RETURN VALUES

These calls return the number of bytes received, or –1 if an error occurred.

ERRORS

The calls fail if:

EBADF	*s* is an invalid file descriptor.
EINTR	The operation was interrupted by delivery of a signal before any data was available to be received.
EIO	An I/O error occurred while reading from or writing to the file system.
ENOMEM	There was insufficient user memory available for the operation to complete.
ENOSR	There were insufficient STREAMS resources available for the operation to complete.
ENOTSOCK	*s* is not a socket.
ESTALE	A stale NFS file handle exists.
EWOULDBLOCK	The socket is marked non-blocking and the requested operation would block.

SEE ALSO

fcntl(2), **ioctl**(2), **read**(2), **connect**(3N), **getsockopt**(3N), **send**(3N), **socket**(3N)

NAME

regcmp, regex – compile and execute regular expression

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

char *__**regcmp(const char** *__*string1*__**, /*** **char** *__*string2* *__**/ ..., int /*(char** *__**)0*__**/);**

char *__**regex(const char** *__*re*__**, const char** *__*subject*__**, /*** **char** *__*ret0* *__**/ ...);**

extern char *__**__loc1;**

MT-LEVEL

MT-Safe

DESCRIPTION

regcmp( ) compiles a regular expression (consisting of the concatenated arguments) and returns a pointer to the compiled form. **malloc**(3C) is used to create space for the compiled form. It is the user's responsibility to free unneeded space so allocated. A **NULL** return from **regcmp( )** indicates an incorrect argument. **regcmp**(1) has been written to generally preclude the need for this routine at execution time.

regex( ) executes a compiled pattern against the subject string. Additional arguments are passed to receive values back. **regex( )** returns **NULL** on failure or a pointer to the next unmatched character on success. A global character pointer **__loc1** points to where the match began. **regcmp( )** and **regex( )** were mostly borrowed from the editor, **ed**(1); however, the syntax and semantics have been changed slightly. The following are the valid symbols and associated meanings.

[ ] * . ^ This group of symbols retains its meaning as described on the **regexp**(5) manual page.

$ Matches the end of the string; **\n** matches a newline.

– Within brackets the minus means *through*. For example, **[a–z]** is equivalent to **[abcd...xyz]**. The **–** can appear as itself only if used as the first or last character. For example, the character class expression **[]–]** matches the characters **]** and **–**.

+ A regular expression followed by **+** means *one or more times*. For example, **[0–9]+** is equivalent to **[0–9][0–9]*.**

{*m***}** **{***m***,}** **{***m***,***u***}** Integer values enclosed in **{ }** indicate the number of times the preceding regular expression is to be applied. The value *m* is the minimum number and *u* is a number, less than 256, which is the maximum. If only *m* is present (that is, **{***m***}**), it indicates the exact number of times the regular expression is to be applied. The value **{***m***,}** is analogous to **{***m***,***infinity***}.** The plus (**+**) and star (*****) operations are equivalent to **{1,}** and **{0,}** respectively.

(...)$*n* The value of the enclosed regular expression is to be returned. The value will be stored in the (*n*+1)th argument following the subject argument. At most, ten enclosed regular expressions are allowed. **regex( )** makes its assignments unconditionally.

(. . .) Parentheses are used for grouping. An operator, for example, *****, **+**, **{}**, can work on a single character or a regular expression enclosed in parentheses. For example, **(a*(cb+)*)$0**.

By necessity, all the above defined symbols are special. They must, therefore, be escaped with a \ (backslash) to be used as themselves.

EXAMPLES

The following example matches a leading newline in the subject string pointed at by cursor.

```
char *cursor, *newcursor, *ptr;
        . . .
newcursor = regex((ptr = regcmp("^\n", (char *)0)), cursor);
free(ptr);
```

The following example matches through the string **Testing3** and returns the address of the character after the last matched character (the "**4**"). The string **Testing3** is copied to the character array **ret0**.

```
char ret0[9];
char *newcursor, *name;
        . . .
name = regcmp("([A–Za–z][A–za–z0–9]{0,7})$0", (char *)0);
newcursor = regex(name, "012Testing345", ret0);
```

The following example applies a precompiled regular expression in **file.i** (see **regcmp**(1)) against *string*.

```
#include "file.i"
char *string, *newcursor;
        . . .
newcursor = regex(name, string);
```

FILES

/usr/ccs/lib/libgen.a

SEE ALSO

ed(1), **regcmp**(1), **malloc**(3C), **regexp**(5)

NOTES

The user program may run out of memory if **regcmp( )** is called iteratively without freeing the vectors no longer required.

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

regcomp, regexec, regerror, regfree – regular expression matching

SYNOPSIS

```
#include <sys/types.h>
#include <regex.h>

int regcomp(regex_t *preg, const char *pattern, int cflags);

int regexec(const regex_t *preg, const char *string, size_t nmatch, regmatch_t pmatch[],
        int eflags);

size_t regerror(int errcode, const regex_t *preg, char *errbuf, size_t errbuf_size);

void regfree(regex_t *preg);
```

MT-LEVEL

MT-Safe

DESCRIPTION

These functions interpret *basic* and *extended* regular expressions (described on the **regex**(5) manual page).

The structure type **regex_t** contains at least the following member:

size_t re_nsub	Number of parenthesised subexpressions.

The structure type **regmatch_t** contains at least the following members:

regoff_t rm_so	Byte offset from start of *string* to start of substring.
regoff_t rm_eo	Byte offset from start of *string* of the first character after the end of substring.

regcomp()

The **regcomp( )** function will compile the regular expression contained in the string pointed to by the *pattern* argument and place the results in the structure pointed to by *preg*. The *cflags* argument is the bitwise inclusive OR of zero or more of the following flags, which are defined in the header **<regex.h>**:

REG_EXTENDED	Use Extended Regular Expressions.
REG_ICASE	Ignore case in match.
REG_NOSUB	Report only success / fail in **regexec( )**.
REG_NEWLINE	Change the handling of NEWLINE characters, as described in the text.

The default regular expression type for *pattern* is a Basic Regular Expression. The application can specify Extended Regular Expressions using the **REG_EXTENDED** *cflags* flag.

If the **REG_NOSUB** flag was not set in *cflags*, then **regcomp( )** will set *re_nsub* to the number of parenthesised subexpressions (delimited by \(\) in basic regular expressions or ( ) in extended regular expressions) found in *pattern*.

regexec()

The **regexec( )** function compares the null-terminated string specified by *string* with the compiled regular expression *preg* initialized by a previous call to **regcomp( )**. The *eflags* argument is the bitwise inclusive OR of zero or more of the following flags, which are defined in the header **<regex.h>**:

REG_NOTBOL The first character of the string pointed to by *string* is not the beginning of the line. Therefore, the circumflex character (^), when taken as a special character, will not match the beginning of *string*.

REG_NOTEOL The last character of the string pointed to by *string* is not the end of the line. Therefore, the dollar sign ($), when taken as a special character, will not match the end of *string*.

If *nmatch* is zero or **REG_NOSUB** was set in the *cflags* argument to **regcomp( )**, then **regexec( )** will ignore the *pmatch* argument. Otherwise, the *pmatch* argument must point to an array with at least *nmatch* elements, and **regexec( )** will fill in the elements of that array with offsets of the substrings of *string* that correspond to the parenthesised subexpressions of *pattern*: *pmatch***[***i***].***rm_so* will be the byte offset of the beginning and *pmatch***[***i***].***rm_eo* will be one greater than the byte offset of the end of substring *i*. (Subexpression *i* begins at the *i*th matched open parenthesis, counting from 1.) Offsets in *pmatch***[0]** identify the substring that corresponds to the entire regular expression. Unused elements of *pmatch* up to *pmatch***[***nmatch*–**1]** will be filled with –**1**. If there are more than *nmatch* subexpressions in *pattern* (*pattern* itself counts as a subexpression), then **regexec( )** will still do the match, but will record only the first *nmatch* substrings.

When matching a basic or extended regular expression, any given parenthesised subexpression of *pattern* might participate in the match of several different substrings of *string*, or it might not match any substring even though the pattern as a whole did match. The following rules are used to determine which substrings to report in *pmatch* when matching regular expressions:

1. If subexpression *i* in a regular expression is not contained within another subexpression, and it participated in the match several times, then the byte offsets in *pmatch***[***i***]** will delimit the last such match.
2. If subexpression *i* is not contained within another subexpression, and it did not participate in an otherwise successful match, the byte offsets in *pmatch***[***i***]** will be –**1**. A subexpression does not participate in the match when:

   * or **\{ \}** appears immediately after the subexpression in a basic regular expression, or *, **?**, or **{}** appears immediately after the subexpression in an extended regular expression, and the subexpression did not match (matched zero times)

   or

   | is used in an extended regular expression to select this subexpression or another, and the other subexpression matched.
3. If subexpression *i* is contained within another subexpression *j*, and *i* is not contained within any other subexpression that is contained within *j*, and a match of subexpression *j* is reported in *pmatch***[***j***]**, then the match or non-match of subexpression *i* reported in *pmatch***[***i***]** will be as described in 1. and 2. above, but within the substring reported in *pmatch***[***j***]** rather than the whole string.
4. If subexpression *i* is contained in subexpression *j*, and the byte offsets in *pmatch***[***j***]** are –**1**, then the pointers in *pmatch***[***i***]** also will be –**1**.

5. If subexpression *i* matched a zero-length string, then both byte offsets in *pmatch***[***i***]** will be the byte offset of the character or **NULL** terminator immediately following the zero-length string.

If, when **regexec( )** is called, the locale is different from when the regular expression was compiled, the result is undefined.

If **REG_NEWLINE** is not set in *cflags*, then a NEWLINE character in *pattern* or *string* will be treated as an ordinary character. If **REG_NEWLINE** is set, then newline will be treated as an ordinary character except as follows:

1. A NEWLINE character in *string* will not be matched by a period outside a bracket expression or by any form of a non-matching list.
2. A circumflex (^) in *pattern*, when used to specify expression anchoring will match the zero-length string immediately after a newline in *string*, regardless of the setting of **REG_NOTBOL**.
3. A dollar-sign ($) in *pattern*, when used to specify expression anchoring, will match the zero-length string immediately before a newline in *string*, regardless of the setting of **REG_NOTEOL**.

regfree( )

The **regfree( )** function frees any memory allocated by **regcomp( )** associated with *preg*.

The following constants are defined as error return values:

REG_NOMATCH	**regexec( )** failed to match.
REG_BADPAT	Invalid regular expression.
REG_ECOLLATE	Invalid collating element referenced.
REG_ECTYPE	Invalid character class type referenced.
REG_EESCAPE	Trailing \ in pattern.
REG_ESUBREG	Number in *digit* invalid or in error.
REG_EBRACK	**[ ]** imbalance.
REG_ENOSYS	The function is not supported.
REG_EPAREN	**\(\)** or **( )** imbalance.
REG_EBRACE	\{ \} imbalance.
REG_BADBR	Content of \{ \} invalid: not a number, number too large, more than two numbers, first larger than second.
REG_ERANGE	Invalid endpoint in range expression.
REG_ESPACE	Out of memory.
REG_BADRPT	?, * or + not preceded by valid regular expression.

regerror( )

The **regerror( )** function provides a mapping from error codes returned by **regcomp( )** and **regexec( )** to unspecified printable strings. It generates a string corresponding to the value of the *errcode* argument, which must be the last non-zero value returned by **regcomp( )** or **regexec( )** with the given value of *preg*. If *errcode* is not such a value, an error message indicating that the error code is invalid is returned.

If *preg* is a **NULL** pointer, but *errcode* is a value returned by a previous call to **regexec( )** or **regcomp( )**, the **regerror( )** still generates an error string corresponding to the value of *errcode*.

If the *errbuf_size* argument is not zero, **regerror( )** will place the generated string into the buffer of size *errbuf_size* bytes pointed to by *errbuf*. If the string (including the terminating **NULL**) cannot fit in the buffer, **regerror( )** will truncate the string and null-terminate the result.

If *errbuf_size* is zero, **regerror( )** ignores the *errbuf* argument, and returns the size of the buffer needed to hold the generated string.

If the *preg* argument to **regexec( )** or **regfree( )** is not a compiled regular expression returned by **regcomp( )**, the result is undefined. A *preg* is no longer treated as a compiled regular expression after it is given to **regfree( )**.

RETURN VALUES

The following values are returned by **regcomp( )**:

0	successful completion
non-zero	an error has occurred. The value returned is described in **<regex.h>**, and the content of *preg* is undefined.

The following values are returned by **regexec( )**:

0	successful completion.
REG_NOMATCH	no match
REG_ENOSYS	the function is not supported.

The following values are returned by **regerror( )**:

0	the function is not implemented.

Upon successful completion, the function returns the number of bytes needed to hold the entire generated string.

The **regfree( )** function returns no value.

USAGE

An application could use:

regerror(code,preg,(char *)NULL,(size_t)0)

to find out how big a buffer is needed for the generated string, **malloc** a buffer to hold the string, and then call **regerror( )** again to get the string (see **malloc**(3C)). Alternately, it could allocate a fixed, static buffer that is big enough to hold most strings, and then use **malloc**( ) to allocate a larger buffer if it finds that this is too small.

EXAMPLES

```
#include <regex.h>

/*
 * Match string against the extended regular expression in
 * pattern, treating errors as no match.
 *
 * return 1 for match, 0 for no match
 */

int
match(const char *string, char *pattern)
{
        int     status;
        regex_t re;

        if (regcomp(&re, pattern, REG_EXTENDED | REG_NOSUB) != 0) {
                return(0);      /* report error */
        }
        status = regexec(&re, string, (size_t) 0, NULL, 0);
        regfree(&re);
        if (status != 0) {
                return(0);      /* report error */
        }
        return(1);
}
```

The following demonstrates how the **REG_NOTBOL** flag could be used with **regexec( )** to find all substrings in a line that match a pattern supplied by a user. (For simplicity of the example, very little error checking is done.)

```
(void) regcomp (&re, pattern, 0);
/* this call to regexec( ) finds the first match on the line */
error = regexec (&re, &buffer[0], 1, &pm, 0);
while (error == 0) {       /* while matches found */
        /* substring found between pm.rm_so and pm.rm_eo */
        /* This call to regexec( ) finds the next match */
        error = regexec (&re, buffer + pm.rm_eo, 1, &pm, REG_NOTBOL);
}
```

SEE ALSO

fnmatch(3C), **glob**(3C), **malloc**(3C), **regex**(5)

NAME

regex, re_comp, re_exec – regular expression handler

SYNOPSIS

```
#include <re_comp.h>
char *re_comp(const char *sp);
int re_exec(const char *p1);
```

DESCRIPTION

re_comp( ) compiles a string into an internal form suitable for pattern matching. **re_exec( )** checks the argument string against the last string passed to **re_comp( )**.

re_comp( ) returns a **NULL** pointer if the string *sp* was compiled successfully; otherwise a string containing an error message is returned. If **re_comp( )** is passed 0 or a **NULL** string, it returns without changing the currently compiled regular expression.

re_exec( ) returns 1 if the string *p1* matches the last compiled regular expression, 0 if the string *p1* failed to match the last compiled regular expression, and –1 if the compiled regular expression was invalid (indicating an internal error).

The strings passed to both **re_comp( )** and **re_exec( )** may have trailing or embedded NEWLINE characters; they are terminated by **NULL** characters. The regular expressions are described on the **regexp**(5) manual page.

RETURN VALUES

re_exec( ) returns –1 for an internal error.

re_comp( ) returns one of the following strings if an error occurs:

```
No previous regular expression
Regular expression too long
unmatched \(
missing ]
too many \(\) pairs
unmatched \)
```

SEE ALSO

grep(1), **regcmp**(1), **regcmp**(3G), **regexpr**(3G), **regexp**(5)

NOTES

The regular expressions of the form **\{**m**\}**, **\{**m,**\}**, or **\{**m,n**\}** are not supported.

NAME

regexpr, compile, step, advance – regular expression compile and match routines

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <regexpr.h>

char *__**compile(const char** *__*instring***, char** *__*expbuf***, char** *__*endbuf***);**

int step(const char *__*string***, char** *__*expbuf***);**

int advance(const char *__*string***, char** *__*expbuf***);**

extern char *loc1, *loc2, *locs;
extern int nbra, regerrno, reglength;
extern char *braslist[], *braelist[];

MT-LEVEL

MT-Safe

DESCRIPTION

These routines are used to compile regular expressions and match the compiled expressions against lines. The regular expressions compiled are in the form used by **ed**(1).

The parameter *instring* is a null-terminated string representing the regular expression.

The parameter *expbuf* points to the place where the compiled regular expression is to be placed. If *expbuf* is **NULL**, **compile( )** uses **malloc**(3C) to allocate the space for the compiled regular expression. If an error occurs, this space is freed. It is the user's responsibility to free unneeded space after the compiled regular expression is no longer needed.

The parameter *endbuf* is one more than the highest address where the compiled regular expression may be placed. This argument is ignored if *expbuf* is **NULL**. If the compiled expression cannot fit in (*endbuf–expbuf*) bytes, **compile( )** returns **NULL** and **regerrno** (see below) is set to 50.

The parameter *string* is a pointer to a string of characters to be checked for a match. This string should be null-terminated.

The parameter *expbuf* is the compiled regular expression obtained by a call of the function **compile( )**.

The function **step( )** returns non-zero if the given string matches the regular expression, and zero if the expressions do not match. If there is a match, two external character pointers are set as a side effect to the call to **step( ) .** The variables set in **step( )** are **loc1** and **loc2**. **loc1** is a pointer to the first character that matched the regular expression. The variable **loc2** points to the character after the last character that matches the regular expression. Thus if the regular expression matches the entire line, **loc1** points to the first character of *string* and **loc2** points to the null at the end of *string*.

The purpose of **step( )** is to step through the *string* argument until a match is found or until the end of *string* is reached. If the regular expression begins with ^, **step( )** tries to match the regular expression at the beginning of the string only.

The **advance( )** function is similar to **step( )**; but, it only sets the variable **loc2** and always restricts matches to the beginning of the string.

If one is looking for successive matches in the same string of characters, **locs** should be set equal to **loc2**, and **step( )** should be called with *string* equal to **loc2**. **locs** is used by commands like **ed** and **sed** so that global substitutions like **s/y*//g** do not loop forever, and is **NULL** by default.

The external variable **nbra** is used to determine the number of subexpressions in the compiled regular expression. **braslist** and **braelist** are arrays of character pointers that point to the start and end of the **nbra** subexpressions in the matched string. For example, after calling **step( )** or **advance( )** with string **sabcdefg** and regular expression **\(abcdef\)**, **braslist[0]** will point at **a** and **braelist[0]** will point at **g**. These arrays are used by commands like **ed** and **sed** for substitute replacement patterns that contain the *\n* notation for subexpressions.

Note that it is not necessary to use the external variables **regerrno**, **nbra**, **loc1**, **loc2 locs**, **braelist**, and **braslist** if one is only checking whether or not a string matches a regular expression.

EXAMPLES

The following is similar to the regular expression code from **grep**:

```
#include <regexpr.h>

. . .
if(compile(*argv, (char *)0, (char *)0) == (char *)0)
        regerr(regerrno);
. . .
if (step(linebuf, expbuf))
        succeed( );
```

RETURN VALUES

If **compile( )** succeeds, it returns a non-**NULL** pointer whose value depends on *expbuf*. If *expbuf* is non-**NULL**, **compile( )** returns a pointer to the byte after the last byte in the compiled regular expression. The length of the compiled regular expression is stored in **reglength**. Otherwise, **compile( )** returns a pointer to the space allocated by **malloc**.

The functions **step( )** and **advance( )** return non-zero if the given string matches the regular expression, and zero if the expressions do not match.

ERRORS

If an error is detected when compiling the regular expression, a **NULL** pointer is returned from **compile( )** and **regerrno** is set to one of the non-zero error numbers indicated below:

ERROR	MEANING
11	Range endpoint too large.
16	Bad number.
25	‘‘\digit’’ out of range.
36	Illegal or missing delimiter.
41	No remembered search string.
42	**\(˜\)** imbalance.
43	Too many **\(**.

44	More than 2 numbers given in **\{˜\}**.
45	**}** expected after \.
46	First number exceeds second in **\{˜\}**.
49	**[ ]** imbalance.
50	Regular expression overflow.

ENVIRONMENT

If any of the **LC_*** variables (**LC_CTYPE, LC_MESSAGES, LC_TIME, LC_COLLATE, LC_NUMERIC,** and **LC_MONETARY**) (see **environ**(5)) are not set in the environment, the operational behavior of **tar** for each corresponding locale category is determined by the value of the **LANG** environment variable. If **LC_ALL** is set, its contents are used to override both the **LANG** and the other **LC_*** variables. If none of the above variables is set in the environment, the "C" (U.S. style) locale determines how **tar** behaves.

LC_CTYPE

Determines how **tar** handles characters. When **LC_CTYPE** is set to a valid value, **tar** can display and handle text and filenames containing valid characters for that locale. **tar** can display and handle Extended Unix code (EUC) characters where any individual character can be 1, 2, or 3 bytes wide. **tar** can also handle EUC characters of 1, 2, or more column widths. In the "C" locale, only characters from ISO 8859-1 are valid.

SEE ALSO

ed(1), **grep**(1), **sed**(1), **malloc**(3C), **regexp**(5)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

remove – remove file

SYNOPSIS

#include <stdio.h>

int remove(const char **path****);**

MT-LEVEL

MT-Safe

DESCRIPTION

remove( ) causes the file or empty directory whose name is the string pointed to by *path* to be no longer accessible by that name. A subsequent attempt to open that file using that name will fail, unless the file is created anew.

For files, **remove( )** is identical to **unlink( )**. For directories, **remove( )** is identical to **rmdir( )**.

See **rmdir**(2) and **unlink**(2) for a detailed list of failure conditions.

RETURN VALUES

Upon successful completion, **remove( )** returns a value of 0; otherwise, it returns a value of –1 and sets **errno** to indicate an error.

SEE ALSO

rmdir(2), **unlink**(2)

NAME

resolver, res_init, res_mkquery, res_send, res_search, dn_comp, dn_expand – resolver routines

SYNOPSIS

cc [*flag* ...] *file* ... **–lresolv –lsocket –lnsl** [*library* ...]

```
#include <sys/types.h>
#include <netinet/in.h>
#include <arpa/nameser.h>
#include <resolv.h>

int res_init(void);

int res_mkquery(int op, char *dname, int class, int type, char *data, int datalen,
        struct rrec *newrr, char *buf, int buflen);

int res_send(char *msg, int msglen, char *answer, int anslen);

int res_search(char *dname, int class, int type, uchar *answer, int anslen);

int dn_comp(char *exp_dn, char *comp_dn, int length, char **dnptrs, char **lastdnptr);

int dn_expand(char *msg, char *comp_dn, char exp_dn, int  msglen, int length);
```

MT-LEVEL

Unsafe

DESCRIPTION

These routines are used for making, sending and interpreting packets to Internet domain name servers. Global information that is used by the resolver routines is kept in the variable **_res**. Most of the values have reasonable defaults and can be ignored. Options are a simple bit mask and are OR-ed in to enable. Options stored in **_res.options** are defined in **<resolv.h>** and are as follows.

RES_INIT	True if the initial name server address and default domain name are initialized (that is, **res_init( )** has been called).
RES_DEBUG	Print debugging messages.
RES_AAONLY	Accept authoritative answers only. **res_send( )** will continue until it finds an authoritative answer or finds an error. Currently this is not implemented.
RES_USEVC	Use TCP connections for queries instead of UDP.
RES_PRIMARY	Query primary server only.
RES_IGNTC	Unused currently (ignore truncation errors, that is, do not retry with TCP).
RES_RECURSE	Set the recursion desired bit in queries. This is the default. **res_send( )** does not do iterative queries and expects the name server to handle recursion.
RES_DEFNAMES	Append the default domain name to single label queries. This is the default.

RES_STAYOPEN Used with **RES_USEVC** to keep the TCP connection open between queries. This is useful only in programs that regularly do many queries. UDP should be the normal mode used.

RES_DNSRCH Search up local domain tree.

res_init( ) reads the initialization file to get the default domain name and the Internet address of the initial hosts running the name server. If this line does not exist, the host running the resolver is tried.

res_mkquery( ) makes a standard query message and places it in *buf*. **res_mkquery( )** will return the size of the query or –**1** if the query is larger than *buflen*. *op* is usually **QUERY** but can be any of the query types defined in **<arpa/nameser.h>**. *dname* is the domain name. If *dname* consists of a single label and the **RES_DEFNAMES** flag is enabled (the default), *dname* will be appended with the current domain name. The current domain name is defined in a system file and can be overridden by the environment variable **LOCALDOMAIN**. *newrr* is currently unused but is intended for making update messages. *class* and *type* define the class and type of query. ∗*data* is the resource record; and *datalen* is the length of the record.

res_send( ) sends a query to name servers and returns an answer. It will call **res_init( )** if **RES_INIT** is not set, send the query to the local name server, and handle timeouts and retries. *msg* is the query sent; *msglen* is its length. *answer* is the response returned. The length of the response is stored in *anslen*. **res_send( )** returns the length of the response or –**1** if there were errors.

res_search( ) formulates and sends a normal query (**QUERY**) message, and stores the response in a buffer supplied by the caller. **dname** is the domain name. **class** and **type** define the class and type of query (see **<arpa/nameser.h>**). The response is returned in the user-supplied buffer **answer**. **res_search** returns the length of **answer** in **anslen**. **res_search( )** will call **res_init( )** if the **RES_INIT** flag is not enabled. If **dname** consists of a single label and the **RES_DEFNAMES** flag is enabled (the default), **dname** will be appended with the current domain name. If the **RES_DNSRCH** flag is enabled, **res_search( )** will search up the local domain tree until an answer has been retrieved or an unrecoverable error has been encountered. **res_search( )** returns the length of **answer** on success and –**1** on error. Note that **res_search( )** is only useful for queries in the same name hierarchy as the local host.

dn_expand( ) expands the compressed domain name *comp_dn* to a full domain name. Expanded names are converted to upper case. *msg* is a pointer to the beginning of the message, *exp_dn* is a pointer to a buffer of size *length* for the result. The size of compressed name is returned or –**1** if there was an error.

dn_comp( ) compresses the domain name *exp_dn* and stores it in *comp_dn*. The size of the compressed name is returned or –**1** if there were errors. *length* is the size of the array pointed to by *comp_dn*. *dnptrs* is a list of pointers to previously compressed names in the current message. The first pointer points to to the beginning of the message and the list ends with **NULL**. *lastdnptr* is a pointer to the end of the array pointed to *dnptrs*. A side effect is to update the list of pointers for labels inserted into the message by **dn_comp( )** as the name is compressed. If *dnptrs* is **NULL**, do not try to compress names. If *lastdnptr*

is **NULL**, do not update the list.

FILES **/etc/resolv.conf**

SEE ALSO **in.named**(1M), **nstest**(1M), **resolv.conf**(4)

NOTES These interfaces are unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

rexec – return stream to a remote command

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

int rexec(char ∗∗*ahost*, **unsigned short** *inport*, **const char** ∗*user*, **const char** ∗*passwd*, **const char** ∗*cmd*, **int** ∗*fd2p*);

MT-LEVEL

Unsafe

DESCRIPTION

rexec( ) looks up the host ∗*ahost* using **gethostbyname**(3N), returning –1 if the host does not exist. Otherwise ∗*ahost* is set to the standard name of the host. If a username and password are both specified, then these are used to authenticate to the foreign host; otherwise the user's **.netrc** file in his home directory is searched for appropriate information. If all this fails, the user is prompted for the information.

The port *inport* specifies which well-known DARPA Internet port to use for the connection. The protocol for connection is described in detail in **in.rexecd**(1M).

If the call succeeds, a socket of type **SOCK_STREAM** is returned to the caller, and given to the remote command as its standard input and standard output. If *fd2p* is non-zero, then an auxiliary channel to a control process will be setup, and a file descriptor for it will be placed in ∗*fd2p*. The control process will return diagnostic output (file descriptor 2, the standard error) from the command on this channel, and will also accept bytes on this channel as signal numbers, to be forwarded to the process group of the command. If *fd2p* is 0, then the standard error (file descriptor 2 of the remote command) will be made the same as its standard output and no provision is made for sending arbitrary signals to the remote process, although you may be able to get its attention by using out-of-band data.

RETURN VALUES

If **rexec( )** succeeds, a file descriptor number, which is a socket of type **SOCK_STREAM**, is returned by the routine. ∗*ahost* is set to the standard name of the host, and if *fd2p* is not NULL, a file descriptor number is placed in ∗*fd2p* which represents the command's standard error stream.

If **rexec( )** fails, **–1** is returned.

SEE ALSO

in.rexecd(1M), **gethostbyname**(3N), **getservbyname**(3N)

NOTES

There is no way to specify options to the **socket( )** call that **rexec( )** makes.

This interface is unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

rpc – library routines for remote procedure calls

SYNOPSIS

cc [*flag* ...] *file* ... –**lnsl** [*library* ...]

#include <rpc/rpc.h>
#include <netconfig.h>

MT-LEVEL

MT-Safe with exceptions

DESCRIPTION

These routines allow C language programs to make procedure calls on other machines across a network. First, the client sends a request to the server. On receipt of the request, the server calls a dispatch routine to perform the requested service, and then sends back a reply.

All RPC routines require the header **<rpc/rpc.h>**. Routines that take a **netconfig** structure also require that **<netconfig.h>** be included. Applications using RPC and XDR routines should be linked with the **libnsl** library.

Multithread Considerations

In the case of multithreaded applications, the **_REENTRANT** flag must be defined on the command line at compilation time (**-D_REENTRANT**). Defining this flag enables a thread-specific version of **rpc_createerr** (see **rpc_clnt_create**(3N)).

Client-side routines are MT-Safe. CLIENT handles (see **rpc_clnt_create**(3N)) can be shared between threads, however in this implementation requests by different threads are serialized (that is, the first request will receive its results before the second request is sent).

Server-side routines are mostly MT-Unsafe. In this implementation the service transport handle, SVCXPRT (see **rpc_svc_create**(3N)), contains a single data area for decoding arguments and encoding results. Therefore, this structure cannot be freely shared between threads that call functions that do this. Routines that are affected by this restriction are marked as unsafe for MT applications (see **rpc_svc_calls**(3N)).

Nettype

Some of the high-level RPC interface routines take a *nettype* string as one of the parameters (for example, **clnt_create( )**, **svc_create( )**, **rpc_reg( )**, **rpc_call( )**). This string defines a class of transports which can be used for a particular application.

nettype can be one of the following:

netpath Choose from the transports which have been indicated by their token names in the **NETPATH** environment variable. If **NETPATH** is unset or **NULL**, it defaults to **visible**. **netpath** is the default *nettype*.

visible Choose the transports which have the visible flag (**v**) set in the **/etc/netconfig** file.

circuit_v This is same as **visible** except that it chooses only the connection oriented transports (semantics **tpi_cots** or **tpi_cots_ord**) from the entries in the **/etc/netconfig** file.

datagram_v This is same as **visible** except that it chooses only the connectionless

datagram transports (semantics **tpi_clts**) from the entries in the **/etc/netconfig** file.

circuit_n This is same as **netpath** except that it chooses only the connection oriented datagram transports (semantics **tpi_cots** or **tpi_cots_ord**).

datagram_n This is same as **netpath** except that it chooses only the connectionless datagram transports (semantics **tpi_clts**).

udp This refers to Internet UDP.

tcp This refers to Internet TCP.

If *nettype* is **NULL**, it defaults to **netpath**. The transports are tried in left to right order in the **NETPATH** variable or in top to down order in the **/etc/netconfig** file.

Data Structures

Some of the data structures used by the RPC package are shown below.

The AUTH Structure

```
union des_block {
    struct {
        u_int32 high;
        u_int32 low;
    } key;
    char c[8];
};
typedef union des_block des_block;
extern bool_t xdr_des_block( );

/*
* Authentication info. Opaque to client.
*/
struct opaque_auth {
    enum_t   oa_flavor;     /* flavor of auth */
    caddr_t  oa_base;       /* address of more auth stuff */
    u_int    oa_length;     /* not to exceed MAX_AUTH_BYTES */
};

/*
* Auth handle, interface to client authenticators.
*/
typedef struct {
    struct   opaque_auth   ah_cred;
    struct   opaque_auth   ah_verf;
    union    des_block     ah_key;
    struct auth_ops {
        void   (*ah_nextverf)( );
        int    (*ah_marshal)( );   /* nextverf & serialize */
        int    (*ah_validate)( );  /* validate verifier */
        int    (*ah_refresh)( );   /* refresh credentials */
```

```
        void    (*ah_destroy)( );    /* destroy this structure */
    } *ah_ops;
    caddr_t ah_private;
} AUTH;
```

The CLIENT Structure

```
/*
 * Client rpc handle.
 * Created by individual implementations.
 * Client is responsible for initializing
 */
typedef struct {
    AUTH            *cl_auth;           /* authenticator */
    struct clnt_ops {
      enum clnt_stat  (*cl_call)( );    /* call remote procedure */
      void            (*cl_abort)( );   /* abort a call */
      void            (*cl_geterr)( );  /* get specific error code */
      bool_t          (*cl_freeres)( ); /* frees results */
      void            (*cl_destroy)( ); /* destroy this structure */
      bool_t          (*cl_control)( ); /* the ioctl( ) of rpc */
    } *cl_ops;
    caddr_t         cl_private;         /* private stuff */
    char            *cl_netid;          /* network identifier */
    char            *cl_tp;             /* device name */
} CLIENT;
```

The SVCXPRT Structure

```
enum xprt_stat {
    XPRT_DIED,
    XPRT_MOREREQS,
    XPRT_IDLE
};

/*
 * Server side transport handle
 */
typedef struct {
    int             xp_fd;              /* file descriptor for the
                                           server handle */
    u_short         xp_port;            /* obsolete */
    struct xp_ops {
      bool_t          (*xp_recv)( );    /* receive incoming requests */
      enum xprt_stat  (*xp_stat)( );    /* get transport status */
      bool_t          (*xp_getargs)( ); /* get arguments */
      bool_t          (*xp_reply)( );   /* send reply */
      bool_t          (*xp_freeargs)( ); /* free mem allocated
                                            for args */
```

```
        void            (*xp_destroy)();  /* destroy this struct */
    } *xp_ops;
    int                 xp_addrlen;       /* length of remote addr.
                                             Obsolete */
    char                *xp_tp;           /* transport provider device
                                             name */
    char                *xp_netid;        /* network identifier */
    struct netbuf       xp_ltaddr;        /* local transport address */
    struct netbuf       xp_rtaddr;        /* remote transport address */
    char                xp_raddr[16];     /* remote address. Obsolete */
    struct opaque_auth  xp_verf;          /* raw response verifier */
    caddr_t             xp_p1;            /* private: for use
                                             by svc ops */
    caddr_t             xp_p2;            /* private: for use
                                             by svc ops */
    caddr_t             xp_p3;            /* private: for use
                                             by svc lib */
    int                 xp_type           /* transport type */
} SVCXPRT;
```

The svc_reg Structure

```
struct svc_req {
    u_long              rq_prog;      /* service program number */
    u_long              rq_vers;      /* service protocol version */
    u_long              rq_proc;      /* the desired procedure */
    struct opaque_auth  rq_cred;      /* raw creds from the wire */
    caddr_t             rq_clntcred;  /* read only cooked cred */
    SVCXPRT             *rq_xprt;     /* associated transport */
};
```

The XDR Structure

```
/*
 * XDR operations.
 * XDR_ENCODE causes the type to be encoded into the stream.
 * XDR_DECODE causes the type to be extracted from the stream.
 * XDR_FREE can be used to release the space allocated by an XDR_DECODE
 * request.
 */
enum xdr_op {
   XDR_ENCODE=0,
   XDR_DECODE=1,
   XDR_FREE=2
};
```

```
/*
 * This is the number of bytes per unit of external data.
 */
#define BYTES_PER_XDR_UNIT(4)
#define RNDUP(x)  ((((x) + BYTES_PER_XDR_UNIT - 1) /
              BYTES_PER_XDR_UNIT) \ * BYTES_PER_XDR_UNIT)

/*
 * A xdrproc_t exists for each data type which is to be encoded or
 * decoded.  The second argument to the xdrproc_t is a pointer to
 * an opaque pointer.  The opaque pointer generally points to a
 * structure of the data type to be decoded.  If this points to 0,
 * then the type routines should allocate dynamic storage of the
 * appropriate size and return it.
 * bool_t  (*xdrproc_t)(XDR *, caddr_t *);
 */
typedef  bool_t (*xdrproc_t)( );

/*
 * The XDR handle.
 * Contains operation which is being applied to the stream,
 * an operations vector for the particular implementation
 */
typedef struct {
    enum xdr_op    x_op;     /* operation; fast additional param */
    struct xdr_ops {
      bool_t    (*x_getlong)( );     /* get a long from underlying stream */
      bool_t    (*x_putlong)( );     /* put a long to underlying stream */
      bool_t    (*x_getbytes)( );    /* get bytes from underlying stream */
      bool_t    (*x_putbytes)( );    /* put bytes to underlying stream */
      u_int     (*x_getpostn)( );    /* returns bytes off from beginning */
      bool_t    (*x_setpostn)( );    /* lets you reposition the stream */
      long *    (*x_inline)( );      /* buf quick ptr to buffered data */
      void      (*x_destroy)( );     /* free privates of this xdr_stream */
    } *x_ops;
    caddr_t     x_public;            /* users' data */
    caddr_t     x_private;           /* pointer to private data */
    caddr_t     x_base;              /* private used for position info */
    int         x_handy;             /* extra private word */
} XDR;
```

Index to Routines

The following table lists RPC routines and the manual reference pages on which they are described:

RPC Routine	Manual Reference Page
auth_destroy	**rpc_clnt_auth**(3N)
authdes_create	**rpc_soc**(3N)
authdes_getucred	**secure_rpc**(3N)
authdes_seccreate	**secure_rpc**(3N)
authkerb_getucred	**kerberos_rpc**(3N)
authkerb_seccreate	**kerberos_rpc**(3N)
authnone_create	**rpc_clnt_auth**(3N)
authsys_create	**rpc_clnt_auth**(3N)
authsys_create_default	**rpc_clnt_auth**(3N)
authunix_create	**rpc_soc**(3N)
authunix_create_default	**rpc_soc**(3N)
callrpc	**rpc_soc**(3N)
clnt_broadcast	**rpc_soc**(3N)
clnt_call	**rpc_clnt_calls**(3N)
clnt_control	**rpc_clnt_create**(3N)
clnt_create	**rpc_clnt_create**(3N)
clnt_destroy	**rpc_clnt_create**(3N)
clnt_dg_create	**rpc_clnt_create**(3N)
clnt_freeres	**rpc_clnt_calls**(3N)
clnt_geterr	**rpc_clnt_calls**(3N)
clnt_pcreateerror	**rpc_clnt_create**(3N)
clnt_perrno	**rpc_clnt_calls**(3N)
clnt_perror	**rpc_clnt_calls**(3N)
clnt_raw_create	**rpc_clnt_create**(3N)
clnt_spcreateerror	**rpc_clnt_create**(3N)
clnt_sperrno	**rpc_clnt_calls**(3N)
clnt_sperror	**rpc_clnt_calls**(3N)
clnt_tli_create	**rpc_clnt_create**(3N)
clnt_tp_create	**rpc_clnt_create**(3N)
clnt_udpcreate	**rpc_soc**(3N)
clnt_vc_create	**rpc_clnt_create**(3N)
clntraw_create	**rpc_soc**(3N)
clnttcp_create	**rpc_soc**(3N)
clntudp_bufcreate	**rpc_soc**(3N)
get_myaddress	**rpc_soc**(3N)
getnetname	**secure_rpc**(3N)
host2netname	**secure_rpc**(3N)
key_decryptsession	**secure_rpc**(3N)
key_encryptsession	**secure_rpc**(3N)
key_gendes	**secure_rpc**(3N)
key_setsecret	**secure_rpc**(3N)

netname2host	**secure_rpc**(3N)
netname2user	**secure_rpc**(3N)
pmap_getmaps	**rpc_soc**(3N)
pmap_getport	**rpc_soc**(3N)
pmap_rmtcall	**rpc_soc**(3N)
pmap_set	**rpc_soc**(3N)
pmap_unset	**rpc_soc**(3N)
rac_drop	**rpc_rac**(3N)
rac_poll	**rpc_rac**(3N)
rac_recv	**rpc_rac**(3N)
rac_send	**rpc_rac**(3N)
registerrpc	**rpc_soc**(3N)
rpc_broadcast	**rpc_clnt_calls**(3N)
rpc_broadcast_exp	**rpc_clnt_calls**(3N)
rpc_call	**rpc_clnt_calls**(3N)
rpc_reg	**rpc_svc_calls**(3N)
svc_create	**rpc_svc_create**(3N)
svc_destroy	**rpc_svc_create**(3N)
svc_dg_create	**rpc_svc_create**(3N)
svc_dg_enablecache	**rpc_svc_calls**(3N)
svc_fd_create	**rpc_svc_create**(3N)
svc_fds	**rpc_soc**(3N)
svc_freeargs	**rpc_svc_reg**(3N)
svc_getargs	**rpc_svc_reg**(3N)
svc_getcaller	**rpc_soc**(3N)
svc_getreq	**rpc_soc**(3N)
svc_getreqset	**rpc_svc_calls**(3N)
svc_getrpccaller	**rpc_svc_calls**(3N)
svc_kerb_reg	**kerberos_rpc**(3N)
svc_raw_create	**rpc_svc_create**(3N)
svc_reg	**rpc_svc_calls**(3N)
svc_register	**rpc_soc**(3N)
svc_run	**rpc_svc_reg**(3N)
svc_sendreply	**rpc_svc_reg**(3N)
svc_tli_create	**rpc_svc_create**(3N)
svc_tp_create	**rpc_svc_create**(3N)
svc_unreg	**rpc_svc_calls**(3N)
svc_unregister	**rpc_soc**(3N)
svc_vc_create	**rpc_svc_create**(3N)
svcerr_auth	**rpc_svc_err**(3N)
svcerr_decode	**rpc_svc_err**(3N)
svcerr_noproc	**rpc_svc_err**(3N)
svcerr_noprog	**rpc_svc_err**(3N)
svcerr_progvers	**rpc_svc_err**(3N)
svcerr_systemerr	**rpc_svc_err**(3N)

svcerr_weakauth	**rpc_svc_err**(3N)
svcfd_create	**rpc_soc**(3N)
svcraw_create	**rpc_soc**(3N)
svctcp_create	**rpc_soc**(3N)
svcudp_bufcreate	**rpc_soc**(3N)
svcudp_create	**rpc_soc**(3N)
user2netname	**secure_rpc**(3N)
xdr_accepted_reply	**rpc_xdr**(3N)
xdr_authsys_parms	**rpc_xdr**(3N)
xdr_authunix_parms	**rpc_soc**(3N)
xdr_callhdr	**rpc_xdr**(3N)
xdr_callmsg	**rpc_xdr**(3N)
xdr_opaque_auth	**rpc_xdr**(3N)
xdr_rejected_reply	**rpc_xdr**(3N)
xdr_replymsg	**rpc_xdr**(3N)
xprt_register	**rpc_svc_calls**(3N)
xprt_unregister	**rpc_svc_calls**(3N)

FILES

/etc/netconfig

SEE ALSO

getnetconfig(3N), **getnetpath**(3N), **kerberos_rpc**(3N), **rpc_clnt_auth**(3N), **rpc_clnt_calls**(3N), **rpc_clnt_create**(3N), **rpc_svc_calls**(3N), **rpc_svc_create**(3N), **rpc_svc_err**(3N), **rpc_svc_reg**(3N), **rpc_xdr**(3N), **rpcbind**(3N), **secure_rpc**(3N), **xdr**(3N), **netconfig**(4), **rpc**(4), **environ**(5)

NAME

rpc_clnt_auth, auth_destroy, authnone_create, authsys_create, authsys_create_default – library routines for client side remote procedure call authentication

MT-LEVEL

MT-Safe

DESCRIPTION

These routines are part of the RPC library that allows C language programs to make procedure calls on other machines across the network, with desired authentication.

These routines are normally called after creating the **CLIENT** handle. The **cl_auth** field of the **CLIENT** structure should be initialized by the **AUTH** structure returned by some of the following routines. The client's authentication information is passed to the server when the RPC call is made.

Only the **NULL** and the **SYS** style of authentication is discussed here. For the **DES** style authentication, please refer to **secure_rpc**(3N). For the **Kerberos** style authentication, please refer to **kerberos_rpc**(3N).

The **NULL** and **SYS** style of authentication are safe in multithreaded applications. For the MT-level of the **DES** and **Kerberos** styles, see their respective pages.

Routines

The following routines require that the header **<rpc/rpc.h>** be included (see **rpc**(3N) for the definition of the **AUTH** data structure).

#include <rpc/rpc.h>

void auth_destroy(AUTH ∗*auth***);**

A function macro that destroys the authentication information associated with *auth*. Destruction usually involves deallocation of private data structures. The use of *auth* is undefined after calling **auth_destroy( )**.

AUTH ∗**authnone_create(void);**

Create and return an RPC authentication handle that passes nonusable authentication information with each remote procedure call. This is the default authentication used by RPC.

AUTH ∗**authsys_create(const char** ∗*host***, const uid_t** *uid***, const gid_t** *gid***, const int** *len***, const gid_t** ∗*aup_gids***);**

Create and return an RPC authentication handle that contains **AUTH_SYS** authentication information. The parameter *host* is the name of the machine on which the information was created; *uid* is the user's user ID; *gid* is the user's current group ID; *len* and *aup_gids* refer to a counted array of groups to which the user belongs.

AUTH ∗**authsys_create_default(void);**

Call **authsys_create( )** with the appropriate parameters.

SEE ALSO

kerberos_rpc(3N), **rpc**(3N), **rpc_clnt_calls**(3N), **rpc_clnt_create**(3N), **secure_rpc**(3N)

NAME

rpc_clnt_calls, clnt_call, clnt_freeres, clnt_geterr, clnt_perrno, clnt_perror, clnt_sperrno, clnt_sperror, rpc_broadcast, rpc_broadcast_exp, rpc_call – library routines for client side calls

MT-LEVEL

MT-Safe

DESCRIPTION

RPC library routines allow C language programs to make procedure calls on other machines across the network. First, the client calls a procedure to send a request to the server. Upon receipt of the request, the server calls a dispatch routine to perform the requested service, and then sends back a reply.

The **clnt_call( )**, **rpc_call( )**, and **rpc_broadcast( )** routines handle the client side of the procedure call. The remaining routines deal with error handling in the case of errors.

Some of the routines take a **CLIENT** handle as one of the parameters. A **CLIENT** handle can be created by an RPC creation routine such as **clnt_create( )** (see **rpc_clnt_create**(3N)).

These routines are safe for use in multithreaded applications. **CLIENT** handles can be shared between threads, however in this implementation requests by different threads are serialized (that is, the first request will receive its results before the second request is sent).

Routines

See **rpc**(3N) for the definition of the **CLIENT** data structure.

#include <rpc/rpc.h>

enum clnt_stat clnt_call(CLIENT ∗*clnt***, const u_long** *procnum***, const xdrproc_t** *inproc***, const caddr_t** *in***, const xdrproc_t** *outproc***, caddr_t** *out***, const struct timeval** *tout***);**

A function macro that calls the remote procedure *procnum* associated with the client handle, *clnt*, which is obtained with an RPC client creation routine such as **clnt_create( )** (see **rpc_clnt_create**(3N)). The parameter *inproc* is the XDR function used to encode the procedure's parameters, and *outproc* is the XDR function used to decode the procedure's results; *in* is the address of the procedure's argument(s), and *out* is the address of where to place the result(s). *tout* is the time allowed for results to be returned, which is overridden by a time-out set explicitly through **clnt_control( )**, see **rpc_clnt_create**(3N).

If the remote call succeeds, the status returned is **RPC_SUCCESS**, otherwise an appropriate status is returned.

bool_t clnt_freeres(CLIENT ∗*clnt***, const xdrproc_t** *outproc***, caddr_t** *out***);**

A function macro that frees any data allocated by the RPC/XDR system when it decoded the results of an RPC call. The parameter *out* is the address of the results, and *outproc* is the XDR routine describing the results. This routine returns **1** if the results were successfully freed, and **0** otherwise.

void clnt_geterr(const CLIENT **clnt*, **struct rpc_err** **errp*);

A function macro that copies the error structure out of the client handle to the structure at address *errp*.

void clnt_perrno(const enum clnt_stat *stat*);

Print a message to standard error corresponding to the condition indicated by *stat*. A newline is appended. Normally used after a procedure call fails for a routine for which a client handle is not needed, for instance **rpc_call( ).**

void clnt_perror(const CLIENT **clnt*, **const char** **s*);

Print a message to the standard error indicating why an RPC call failed; *clnt* is the handle used to do the call. The message is prepended with string *s* and a colon. A newline is appended. Normally used after a remote procedure call fails for a routine which requires a client handle, for instance **clnt_call( )**.

char *clnt_sperrno(const enum clnt_stat *stat*);

Take the same arguments as **clnt_perrno( )**, but instead of sending a message to the standard error indicating why an RPC call failed, return a pointer to a string which contains the message.

clnt_sperrno( ) is normally used instead of **clnt_perrno( )** when the program does not have a standard error (as a program running as a server quite likely does not), or if the programmer does not want the message to be output with **printf( )** (see **printf**(3S)), or if a message format different than that supported by **clnt_perrno( )** is to be used. Note: unlike **clnt_sperror( )** and **clnt_spcreaterror( )** (see **rpc_clnt_create**(3N)), **clnt_sperrno( )** does not return pointer to static data so the result will not get overwritten on each call.

char *clnt_sperror(const CLIENT **clnt*, **const char** **s*);

Like **clnt_perror( )**, except that (like **clnt_sperrno( )**) it returns a string instead of printing to standard error. However, **clnt_sperror( )** does not append a newline at the end of the message.

Warning: returns pointer to a buffer that is overwritten on each call. In multithread applications, this buffer is implemented as thread-specific data.

enum clnt_stat rpc_broadcast(const u_long *prognum*, **const u_long** *versnum*,
const u_long *procnum*, **const xdrproc_t** *inproc*, **const caddr_t** *in*,
const xdrproc_t *outproc*, **caddr_t** *out*, **const resultproc_t** *eachresult*,
const char **nettype*);

Like **rpc_call( )**, except the call message is broadcast to all the connectionless transports specified by *nettype*. If *nettype* is **NULL**, it defaults to "**netpath**. Each time it receives a response, this routine calls **eachresult( )**, whose form is:

bool_t eachresult(caddr_t *out***, const struct netbuf** ∗*addr***,**
const struct netconfig ∗*netconf***);**

where *out* is the same as *out* passed to **rpc_broadcast( )**, except that the remote procedure's output is decoded there; *addr* points to the address of the machine that sent the results, and *netconf* is the netconfig structure of the transport on which the remote server responded. If **eachresult( )** returns **0**, **rpc_broadcast( )** waits for more replies; otherwise it returns with appropriate status.

Warning: broadcast file descriptors are limited in size to the maximum transfer size of that transport. For Ethernet, this value is 1500 bytes. **rpc_broadcast( )** uses **AUTH_SYS** credentials by default (see **rpc_clnt_auth**(3N)).

enum clnt_stat rpc_broadcast_exp(const u_long *prognum***, const u_long** *versnum***,**
const u_long *procnum***, const xdrproc_t** *xargs***, caddr_t** *argsp***,**
const xdrproc_t *xresults***, caddr_t** *resultsp***, const resultproc_t** *eachresult***,**
const int *inittime***, const int** *waittime***, const char** ∗*nettype***);**

Like **rpc_broadcast( )**, except that the initial timeout, *inittime* and the maximum timeout, *waittime* are specified in milliseconds.

inittime is the initial time that **rpc_broadcast_exp( )** waits before resending the request. After the first resend, the re-transmission interval increases exponentially until it exceeds *waittime*.

enum clnt_stat rpc_call(const char ∗*host***, const u_long** *prognum***,**
const u_long *versnum***, const u_long** *procnum***, const xdrproc_t** *inproc***,**
const char ∗*in***, const xdrproc_t** *outproc***, char** ∗*out***, const char** ∗*nettype***);**

Call the remote procedure associated with *prognum*, *versnum*, and *procnum* on the machine, *host*. The parameter *inproc* is used to encode the procedure's parameters, and *outproc* is used to decode the procedure's results; *in* is the address of the procedure's argument(s), and *out* is the address of where to place the result(s). *nettype* can be any of the values listed on **rpc**(3N). This routine returns RPC_SUCCESS if it succeeds, or an appropriate status is returned. Use the **clnt_perrno( )** routine to translate failure status into error messages.

Warning: **rpc_call( )** uses the first available transport belonging to the class *nettype*, on which it can create a connection. You do not have control of timeouts or authentication using this routine.

SEE ALSO **printf**(3S), **rpc**(3N), **rpc_clnt_auth**(3N), **rpc_clnt_create**(3N)

NAME

rpc_clnt_create, clnt_control, clnt_create, clnt_create_timed, clnt_create_vers, clnt_destroy, clnt_dg_create, clnt_pcreateerror, clnt_raw_create, clnt_spcreateerror, clnt_tli_create, clnt_tp_create, clnt_tp_create_timed, clnt_vc_create, rpc_createerr – library routines for dealing with creation and manipulation of CLIENT handles

MT-LEVEL

MT-Safe

DESCRIPTION

RPC library routines allow C language programs to make procedure calls on other machines across the network. First a **CLIENT** handle is created and then the client calls a procedure to send a request to the server. On receipt of the request, the server calls a dispatch routine to perform the requested service, and then sends a reply.

These routines are MT-Safe. In the case of multithreaded applications, the **_REENTRANT** flag must be defined on the command line at compilation time **(-D_REENTRANT).** When the **_REENTRANT** flag is defined, **rpc_createerr** becomes a macro which enables each thread to have its own **rpc_createerr**.

Routines

See **rpc**(3N) for the definition of the **CLIENT** data structure.

#include <rpc/rpc.h>

bool_t clnt_control(CLIENT ∗*clnt*, **const u_int** *req*, **char** ∗*info*);

A function macro to change or retrieve various information about a client object. *req* indicates the type of operation, and *info* is a pointer to the information. For both connectionless and connection-oriented transports, the supported values of *req* and their argument types and what they do are:

CLSET_TIMEOUT	**struct timeval** ∗	set total timeout
CLGET_TIMEOUT	**struct timeval** ∗	get total timeout

Note: if you set the timeout using **clnt_control( )**, the timeout argument passed by **clnt_call( )** is ignored in all subsequent calls.

Note: If you set the timeout value to **0 clnt_control( )** immediately returns an error (**RPC_TIMEDOUT**). Set the timeout parameter to **0** for batching calls.

CLGET_FD	**int** ∗	get the associated file descriptor
CLGET_SVC_ADDR	**struct netbuf** ∗	get servers address
CLSET_FD_CLOSE	**void**	close the file descriptor when destroying the client handle (see **clnt_destroy( )**)
CLSET_FD_NCLOSE	**void**	do not close the file descriptor when destroying the client handle
CLGET_VERS	**unsigned long** ∗	get the RPC program's version number associated with the client handle

CLSET_VERS	**unsigned long** *	set the RPC program's version number associated with the client handle. This assumes that the RPC server for this new version is still listening at the address of the previous version.
CLGET_XID	**unsigned long** *	get the XID of the previous remote procedure call
CLSET_XID	**unsigned long** *	set the XID of the next remote procedure call

The following operations are valid for connectionless transports only:

CLSET_RETRY_TIMEOUT	**struct timeval** *	set the retry timeout
CLGET_RETRY_TIMEOUT	**struct timeval** *	get the retry timeout

The retry timeout is the time that RPC waits for the server to reply before retransmitting the request.

clnt_control( ) returns **TRUE** on success and **FALSE** on failure.

CLIENT *__clnt_create(const char__ *host__, const u_long__ *prognum__,__*
const u_long *versnum__, const char__ *nettype__);__*

Generic client creation routine for program *prognum* and version *versnum*. *host* identifies the name of the remote host where the server is located. *nettype* indicates the class of transport protocol to use. The transports are tried in left to right order in **NETPATH** variable or in top to bottom order in the netconfig database.

clnt_create( ) tries all the transports of the *nettype* class available from the **NETPATH** environment variable and the netconfig database, and chooses the first successful one. A default timeout is set and can be modified using **clnt_control( )**. This routine returns **NULL** if it fails. The **clnt_pcreateerror( )** routine can be used to print the reason for failure.

Note: **clnt_create( )** returns a valid client handle even if the particular version number supplied to **clnt_create( )** is not registered with the **rpcbind** service. This mismatch will be discovered by a **clnt_call** later (see **rpc_clnt_calls**(3N)).

CLIENT *__clnt_create_timed(const char__ *host__,__*
const u_long *prognum__, const u_long__ *versnum__,__*
const char *nettype__, const struct timeval__ *timeout__);__*

Generic client creation routine which is similar to **clnt_create( )** but which also has the additional parameter *timeout* that specifies the maximum amount of time allowed for each transport class tried. In all other respects, the **clnt_create_timed( )** call behaves exactly like the **clnt_create( )** call.

CLIENT ∗**clnt_create_vers(const char** ∗*host*, **const u_long** *prognum*, **u_long** ∗*vers_outp*, **const u_long** *vers_low*, **const u_long** *vers_high*, **char** ∗*nettype*);

Generic client creation routine which is similar to **clnt_create( )** but which also checks for the version availability. *host* identifies the name of the remote host where the server is located. *nettype* indicates the class transport protocols to be used. If the routine is successful it returns a client handle created for the highest version between *vers_low* and *vers_high* that is supported by the server. *vers_outp* is set to this value. That is, after a successful return *vers_low* <= ∗*vers_outp* <= *vers_high*. If no version between *vers_low* and *vers_high* is supported by the server then the routine fails and returns NULL. A default timeout is set and can be modified using **clnt_control( )**. This routine returns **NULL** if it fails. The **clnt_pcreateerror( )** routine can be used to print the reason for failure.

Note: **clnt_create( )** returns a valid client handle even if the particular version number supplied to **clnt_create( )** is not registered with the **rpcbind** service. This mismatch will be discovered by a **clnt_call** later (see **rpc_clnt_calls**(3N)). However, **clnt_create_vers( )** does this for you and returns a valid handle only if a version within the range supplied is supported by the server.

void clnt_destroy(CLIENT ∗*clnt*);

A function macro that destroys the client's RPC handle. Destruction usually involves deallocation of private data structures, including *clnt* itself. Use of *clnt* is undefined after calling **clnt_destroy( )**. If the RPC library opened the associated file descriptor, or **CLSET_FD_CLOSE** was set using **clnt_control( )**, the file descriptor will be closed.

The caller should call **auth_destroy(***clnt*→**cl_auth)** (before calling **clnt_destroy( )**) to destroy the associated AUTH structure (see **rpc_clnt_auth**(3N)).

CLIENT ∗**clnt_dg_create(const int** *fildes*, **const struct netbuf** ∗*svcaddr*, **const u_long** *prognum*, **const u_long** *versnum*, **const u_int** *sendsz*, **const u_int** *recvsz*);

This routine creates an RPC client for the remote program *prognum* and version *versnum*; the client uses a connectionless transport. The remote program is located at address *svcaddr*. The parameter *fildes* is an open and bound file descriptor. This routine will resend the call message in intervals of 15 seconds until a response is received or until the call times out. The total time for the call to time out is specified by **clnt_call( )** (see **clnt_call( )** in **rpc_clnt_calls**(3N)). The retry time out and the total time out periods can be changed using **clnt_control( )**. The user may set the size of the send and receive buffers with the parameters *sendsz* and *recvsz*; values of **0** choose suitable defaults. This routine returns **NULL** if it fails.

void clnt_pcreateerror(const char **s*);

Print a message to standard error indicating why a client RPC handle could not be created. The message is prepended with the string *s* and a colon, and appended with a newline.

CLIENT *clnt_raw_create(const u_long *prognum*, **const u_long** *versnum*);

This routine creates an RPC client handle for the remote program *prognum* and version *versnum*. The transport used to pass messages to the service is a buffer within the process's address space, so the corresponding RPC server should live in the same address space; (see **svc_raw_create( )** in **rpc_svc_create**(3N)). This allows simulation of RPC and measurement of RPC overheads, such as round trip times, without any kernel or networking interference. This routine returns **NULL** if it fails. **clnt_raw_create( )** should be called after **svc_raw_create( )**.

char *clnt_spcreateerror(const char **s*);

Like **clnt_pcreateerror( )**, except that it returns a string instead of printing to the standard error. A newline is not appended to the message in this case.

Warning: returns a pointer to a buffer that is overwritten on each call. In multithread applications, this buffer is implemented as thread-specific data.

CLIENT *clnt_tli_create(const int *fildes*, **const struct netconfig** **netconf*, **const struct netbuf** **svcaddr*, **const_long** *prognum*, **const u_long** *versnum*, **const u_int** *sendsz*, **const u_int** *recvsz*);

This routine creates an RPC client handle for the remote program *prognum* and version *versnum*. The remote program is located at address *svcaddr*. If *svcaddr* is **NULL** and it is connection-oriented, it is assumed that the file descriptor is connected. For connectionless transports, if *svcaddr* is **NULL**, **RPC_UNKNOWNADDR** error is set. *fildes* is a file descriptor which may be open, bound and connected. If it is **RPC_ANYFD**, it opens a file descriptor on the transport specified by *netconf*. If *fildes* is **RPC_ANYFD** and *netconf* is **NULL**, a **RPC_UNKNOWNPROTO** error is set. If *fildes* is unbound, then it will attempt to bind the descriptor. The user may specify the size of the buffers with the parameters *sendsz* and *recvsz*; values of **0** choose suitable defaults. Depending upon the type of the transport (connection-oriented or connectionless), **clnt_tli_create( )** calls appropriate client creation routines. This routine returns **NULL** if it fails. The **clnt_pcreateerror( )** routine can be used to print the reason for failure. The remote **rpcbind** service (see **rpcbind**(1M)) is not consulted for the address of the remote service.

CLIENT ∗**clnt_tp_create(const char** ∗*host*, **const u_long** *prognum*, **const u_long** *versnum*, **const struct netconfig** ∗*netconf*);

Like **clnt_create( )** except **clnt_tp_create( )** tries only one transport specified through *netconf*.

clnt_tp_create( ) creates a client handle for the program *prognum*, the version *versnum*, and for the transport specified by *netconf*. Default options are set, which can be changed using **clnt_control( )** calls. The remote **rpcbind** service on the host *host* is consulted for the address of the remote service. This routine returns **NULL** if it fails. The **clnt_pcreateerror( )** routine can be used to print the reason for failure.

CLIENT ∗**clnt_tp_create_timed(const char** ∗*host*, **const u_long** *prognum*, **const u_long** *versnum*, **const struct netconfig** ∗*netconf*, **const struct timeval** ∗*timeout*);

Like **clnt_tp_create( )** except **clnt_tp_create_timed( )** has the extra parameter *timeout* which specifies the maximum time allowed for the creation attempt to succeed. In all other respects, the **clnt_tp_create_timed( )** call behaves exactly like the **clnt_tp_create( )** call.

CLIENT ∗**clnt_vc_create(const int** *fildes*, **const struct netbuf** ∗*svcaddr*, **const u_long** *prognum*, **const u_long** *versnum*, **const u_int** *sendsz*, **const u_int** *recvsz*);

This routine creates an RPC client for the remote program *prognum* and version *versnum*; the client uses a connection-oriented transport. The remote program is located at address *svcaddr*. The parameter *fildes* is an open and bound file descriptor. The user may specify the size of the send and receive buffers with the parameters *sendsz* and *recvsz*; values of **0** choose suitable defaults. This routine returns **NULL** if it fails.

The address *svcaddr* should not be **NULL** and should point to the actual address of the remote program. **clnt_vc_create( )** does not consult the remote **rpcbind** service for this information.

struct rpc_createerr rpc_createerr;

A global variable whose value is set by any RPC client handle creation routine that fails. It is used by the routine **clnt_pcreateerror( )** to print the reason for the failure.

In multithreaded applications, **rpc_createerr** becomes a macro which enables each thread to have its own **rpc_createerr**.

SEE ALSO

rpc(3N), **rpc_clnt_auth**(3N), **rpc_clnt_calls**(3N), **rpcbind**(1M)

NAME

rpc_control – library routine for manipulating global RPC attributes for client and server applications

SYNOPSIS

bool_t rpc_control(int *op***,***void* ∗**info***);*

MT-LEVEL

MT-Safe

DESCRIPTION

This RPC library routine allows applications to set and modify global RPC attributes that apply to clients as well as servers. At present, it supports only server side operations. This function allows applications to set and modify global attributes that apply to client as well as server functions. *op* indicates the type of operation, and *info* is a pointer to the operation specific information. The supported values of *op* and their argument types, and what they do are:

RPC_SVC_MTMODE_SET	**int** ∗	set multithread mode
RPC_SVC_MTMODE_GET	**int** ∗	get multithread mode
RPC_SVC_THRMAX_SET	**int** ∗	set maximum number of threads
RPC_SVC_THRMAX_GET	**int** ∗	get maximum number of threads
RPC_SVC_THRTOTAL_GET	**int** ∗	get number of active threads
RPC_SVC_THRCREATES_GET	**int** ∗	get number of threads created
RPC_SVC_THRERRORS_GET	**int** ∗	get number of thread create errors

There are three multithread (MT) modes. These are:

RPC_SVC_MT_NONE	Single threaded mode	(default)
RPC_SVC_MT_AUTO	Automatic MT mode	
RPC_SVC_MT_USER	User MT mode	

Unless the application sets the Automatic or User MT modes, it will stay in the default (single threaded) mode. See the Network Interfaces Programming Guide for the meanings of these modes and programming examples. Once a mode is set, it cannot be changed.

By default, the maximum number of threads that the server will create at any time is 16. This allows the service developer to put a bound on thread resources consumed by a server. If a server needs to process more than 16 client requests concurrently, the maximum number of threads must be set to the desired number. This parameter may be set at any time by the server.

Set and get operations will succeed even in modes where the operations don't apply. For example, you can set the maximum number of threads in any mode, even though it makes sense only for the Automtatic MT mode. All of the get operations except **RPC_SVC_MTMODE_GET** apply only to the Automatic MT mode, so values returned in other modes may be undefined.

RETURN VALUES

This routine returns **TRUE** if the operation was successful, and **FALSE** otherwise.

SEE ALSO

rpc(3N), **rpc_svc_calls**(3N), **rpcbind**(1M)

NAME

rpc_rac, rac_drop, rac_poll, rac_recv, rac_send – remote asynchronous calls

SYNOPSIS

cc [*flag* ...] *file* ... **–lrac –lnsl** [*library* ...]

#include <rpc/rpc.h>
#include <rpc/rac.h>

MT-LEVEL

Unsafe

DESCRIPTION

The remote asynchronous calls (RAC) package is a special interface to the RPC library that allows messages to be sent using the RPC protocol without blocking during the time between when the message is sent and the reply is received. To RPC servers, RAC messages are indistinguishable from RPC messages.

A client establishes an RPC session in the usual way (see **rpc_clnt_create**(3N)). A RAC message is sent using **rac_send( )**. This routine returns immediately, allowing the client to conduct other processing. When the client wants to determine whether the returned value from the call has been received, **rac_poll( )** is used. **rac_recv( )** is used to collect the returned value; it can also be used to block while waiting for the returned value to arrive. **rac_drop( )** is used to inform the RPC library that the client is no longer interested in the results of a particular RAC message.

#include <rpc/rpc.h>

void rac_drop(CLIENT ∗*cl***, void** ∗*h***);**

rac_drop( ) should be called when the user is no longer interested in the result of a **rac_send( )** currently in progress. No message to the server is generated by this call, but any subsequent reply received for this handle will be silently dropped. It also frees any space occupied by the asynchronous call handle *h*.

After a call to **rac_drop( )** the handle referred to by *h* is invalid. It may no longer be used in any asynchronous operation.

enum clnt_stat rac_poll(CLIENT ∗*cl***, void** ∗*h***);**

rac_poll( ) returns the status of the call currently in progress on the <CLIENT, asynchronous handle> tuple referred to by *cl* and *h*.

rac_poll( ) return values are:

RPC_SUCCESS
A reply has been received and is available for reading by **rac_recv( )**.

RPC_INPROGRESS
No reply has been received. The call referred to by the given handle has not yet timed out.

RPC_TIMEDOUT
No reply has been received. The call referred to by the given handle has exceeded the maximum timeout value specified in **rac_send( )**.

RPC_STALERACHANDLE
Either the handle referred to by *h* is invalid or no call is currently in progress for the given <CLIENT, asynchronous handle> tuple.

RPC_CANTRECV
Either the file descriptor associated with the given CLIENT handle is bad, or an error occurred while attempting to receive a packet.

RPC_SYSTEMERROR
Space could not be allocated to receive a packet.

On unreliable transports, a call to **rac_poll( )** will trigger a retransmission when necessary (that is, if a **rac_send( )** is in progress, no reply has been received, the per-call timeout has expired, and the total timeout has not yet expired).

The return value for **rac_poll( )** is independent of the RPC return value in the reply packet. Although a combination of **clnt_control( )**'s CLGET_FD request and **poll**(2) may be used to extract the proper file descriptor and poll for packets, **rac_poll( )** is still useful since it will determine whether a reply is available for a specific <CLIENT, asynchronous handle> tuple.

enum clnt_stat rac_recv(CLIENT ∗*cl***, void** ∗ *h***);**

rac_recv( ) retrieves the results of a previous asynchronous RPC call, placing them in the buffer indicated in the **rac_send( )** call and using the XDR decode function supplied there. It depends on the application to have ensured that a reply is present (using **rac_poll()**). If **rac_recv( )** is called before a reply has been received, it will block awaiting a reply.

All errors normally returned by the RPC client call functions may be returned here. In addition:

RPC_STALERACHANDLE
Either the handle referred to by *h* is invalid or no call is currently in progress for the given <CLIENT, asynchronous handle> tuple.

Additionally, if a packet is present and its status is not RPC_SUCCESS, it is possible that the client credentials need refreshing. In this case, RPC_AUTHERROR is returned and the client should attempt to resend the call.

When a reply has been received, **rac_recv( )** will invoke the XDR decode procedure specified in the **rac_send( )** call. After a call to **rac_recv( )**, the handle referred to by *h* is invalid. It may no longer be used in any asynchronous operation.

void ∗**rac_send(**CLIENT ∗*cl*, **unsigned long** *proc*, **xdrproc_t** *xargs*, **void** ∗*argsp*, **xdrproc_t** *xresults*, **void** ∗*resultsp*, **struct timeval** *timeout*);

rac_send( ) initiates (sends to the server) an RPC call to the specified procedure. It does not await a reply from the server. *argsp* is the address of the procedure's arguments, *resultsp* is the address in which to place the results, *xargs* and *xresults* are XDR functions used to encode and decode respectively. Note: *resultsp* must be a valid pointer when **rac_recv( )** is called. *timeout* should contain the total amount of time the application is willing to wait for a reply.

Upon success, an opaque handle, known as the asynchronous handle, is returned. This handle is to be used in subsequent asynchronous calls to poll for the status of the call (**rac_poll( )**), receive the returned results of the call (**rac_recv( )**), or cancel the call (**rac_drop( )**).

On failure, *(void ∗) 0* is returned.

In case of failure, the application may retrieve the RPC failure code by calling **clnt_geterr( )** immediately after a **rac_send( )** failure (see **rpc**(3N)). Possible errors include both transient problems (such as transport failures) and permanent ones (such as XDR encoding failures).

Multiple **rac_send**s on the same client handle are permitted, but may introduce unpredictable perturbations to the current timeout and retry model used by the RPC library.

The interface imposes a limit on the amount of time a call may be in progress before it is considered to have failed. This method was chosen over limitations on the number of retries because of a desire for transport independence.

SEE ALSO

poll(2), **rpc**(3N), **rpc_clnt_create**(3N), **rpc_clnt_calls**(3N), **xdr**(3N)

WARNINGS

The RAC interface is not the recommended interface for having multiple RPC requests outstanding. The preferred method of accomplishing this in the Solaris environment is to use synchronous RPC calls with threads. The RAC interface is provided as a service to developers interested in porting RPC applications to Solaris 2.0. Use of this interface will degrade the performance of normal synchronous RPC calls (see **rpc_clnt_calls**(3N)). For these reasons, use of this interface is disparaged.

The library **librac** must be linked before **libnsl** to use RAC. If the libraries are not linked in the correct order, then the results are indeterminate.

NOTES

These interfaces are unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

rpc_soc, authdes_create, authunix_create, authunix_create_default, callrpc, clnt_broadcast, clntraw_create, clnttcp_create, clntudp_bufcreate, clntudp_create, get_myaddress, getrpcport, pmap_getmaps, pmap_getport, pmap_rmtcall, pmap_set, pmap_unset, registerrpc, svc_fds, svc_getcaller, svc_getreq, svc_register, svc_unregister, svcfd_create, svcraw_create, svctcp_create, svcudp_bufcreate, svcudp_create, xdr_authunix_parms – obsolete library routines for RPC

MT-LEVEL

Unsafe

DESCRIPTION

RPC routines allow C programs to make procedure calls on other machines across the network. First, the client calls a procedure to send a request to the server. Upon receipt of the request, the server calls a dispatch routine to perform the requested service, and then sends back a reply. Finally, the procedure call returns to the client.

The routines described in this manual page have been superseded by other routines. The preferred routine is given after the description of the routine. New programs should use the preferred routines, as support for the older interfaces may be dropped in future releases.

File Descriptors

Transport independent RPC uses TLI as its transport interface instead of sockets.

Some of the routines described in this section (such as **clnttcp_create( )**) take a pointer to a file descriptor as one of the parameters. If the user wants the file descriptor to be a socket, then the application will have to be linked with both **librpcsoc** and **libnsl**. If the user passed **RPC_ANYSOCK** as the file descriptor, and the application is linked with **libnsl** only, then the routine will return a TLI file descriptor and not a socket.

Routines

The following routines require that the header **<rpc/rpc.h>** be included. The symbol **PORTMAP** should be defined so that the appropriate function declarations for the old interfaces are included through the header files.

#define PORTMAP
#include <rpc/rpc.h>

AUTH * **authdes_create(char** **name***, **unsigned** *window***, struct sockaddr** **syncaddr***,*
*****des_block** **ckey***);*

authdes_create( ) is the first of two routines which interface to the RPC secure authentication system, known as DES authentication. The second is **authdes_getucred( )**, below. Note: the keyserver daemon **keyserv**(1M) must be running for the DES authentication system to work.

authdes_create( ), used on the client side, returns an authentication handle that will enable the use of the secure authentication system. The first parameter *name* is the network name, or *netname*, of the owner of the server process. This field usually represents a hostname derived from the utility routine **host2netname( )**, but could also represent a user name using **user2netname( )** (see **secure_rpc**(3N)). The second field is window on the validity of the client credential, given in seconds. A small window is more secure than a large one, but

choosing too small of a window will increase the frequency of resynchronizations because of clock drift. The third parameter *syncaddr* is optional. If it is NULL, then the authentication system will assume that the local clock is always in sync with the server's clock, and will not attempt resynchronizations. If an address is supplied, however, then the system will use the address for consulting the remote time service whenever resynchronization is required. This parameter is usually the address of the RPC server itself. The final parameter *ckey* is also optional. If it is NULL, then the authentication system will generate a random DES key to be used for the encryption of credentials. If it is supplied, however, then it will be used instead.

Warning: this routine exists for backward compatibility only, and is obsoleted by **authdes_seccreate( )** (see **secure_rpc**(3N)).

AUTH * **authunix_create(char** **host***, **int** *uid* , **int** *gid*, **int** *grouplen*, **int** *gidlistp***);**

Create and return an RPC authentication handle that contains authentication information. The parameter *host* is the name of the machine on which the information was created; *uid* is the user's user ID; *gid* is the user's current group ID; *grouplen* and *gidlistp* refer to a counted array of groups to which the user belongs.

Warning: it is not very difficult to impersonate a user.

Warning: this routine exists for backward compatibility only, and is obsoleted by **authsys_create( )** (see **rpc_clnt_auth**(3N)).

AUTH * **authunix_create_default(void)**

Call **authunix_create( )** with the appropriate parameters.

Warning: this routine exists for backward compatibility only, and is obsoleted by **authsys_create_default( )** (see **rpc_clnt_auth**(3N)).

callrpc(char **host***, **u_long** *prognum*, **u_long** *versnum*, **u_long** *procnum*,
xdrproc_t *inproc*, **char** **in***, **xdrproc_t** *outproc*, **char** **out***);*

Call the remote procedure associated with *prognum*, *versnum*, and *procnum* on the machine, *host*. The parameter *inproc* is used to encode the procedure's parameters, and *outproc* is used to decode the procedure's results; *in* is the address of the procedure's argument, and *out* is the address of where to place the result(s). This routine returns **0** if it succeeds, or the value of **enum clnt_stat** cast to an integer if it fails. The routine **clnt_perrno( )** (see **rpc_clnt_calls**(3N)) is handy for translating failure statuses into messages.

Warning: you do not have control of timeouts or authentication using this routine. This routine exists for backward compatibility only, and is obsoleted by **rpc_call( )** (see **rpc_clnt_calls**(3N)).

enum clnt_stat clnt_broadcast(u_long *prognum***, u_long** *versnum***, u_long** *procnum***, xdrproc_t** *inproc***, char** **in***, xdrproc_t** *outproc***, char** **out***, resultproc_t** *eachresult***);**

Like **callrpc( )**, except the call message is broadcast to all locally connected broadcast nets. Each time the caller receives a response, this routine calls **eachresult( )**, whose form is:

eachresult(char **out***, struct sockaddr_in** **addr***);**

where *out* is the same as *out* passed to **clnt_broadcast( )**, except that the remote procedure's output is decoded there; *addr* points to the address of the machine that sent the results. If **eachresult( )** returns **0 clnt_broadcast( )** waits for more replies; otherwise it returns with appropriate status. If **eachresult( )** is NULL, **clnt_broadcast( )** returns without waiting for any replies.

Warning: broadcast packets are limited in size to the maximum transfer unit of the transports involved. For Ethernet, the callers argument size is approximately 1500 bytes. Since the call message is sent to all connected networks, it may potentially lead to broadcast storms. **clnt_broadcast( )** uses SB AUTH_SYS credentials by default (see **rpc_clnt_auth**(3N)).

Warning: this routine exists for backward compatibility only, and is obsoleted by **rpc_broadcast( )** (see **rpc_clnt_calls**(3N)).

CLIENT *** **clntraw_create(u_long** *prognum***, u_long** *versnum***);**

This routine creates an internal, memory-based RPC client for the remote program *prognum*, version *versnum*. The transport used to pass messages to the service is actually a buffer within the process's address space, so the corresponding RPC server should live in the same address space; see **svcraw_create( )**. This allows simulation of RPC and acquisition of RPC overheads, such as round trip times, without any kernel interference. This routine returns NULL if it fails.

Warning: this routine exists for backward compatibility only, and has the same functionality as **clnt_raw_create( )** (see **rpc_clnt_create**(3N)), which obsoletes it.

CLIENT *** **clnttcp_create(struct sockaddr_in** **addr***, u_long** *prognum***, u_long** *versnum***, int** **fdp***, u_int** *sendsz***, u_int** *recvsz***);**

This routine creates an RPC client for the remote program *prognum*, version *versnum*; the client uses TCP/IP as a transport. The remote program is located at Internet address *addr*. If *addr*→*sin_port* is **0**,, then it is set to the actual port that the remote program is listening on (the remote **rpcbind** service is consulted for this information). The parameter **fdp* is a file descriptor, which may be open and bound; if it is **RPC_ANYSOCK**, then this routine opens a new one and sets **fdp*. Refer to the **File Descriptor** section for more information. Since TCP-based RPC uses buffered I/O, the user may specify the size of the send and receive buffers with the parameters *sendsz* and *recvsz*; values of **0** choose suitable defaults. This routine returns NULL if it fails.

Warning: this routine exists for backward compatibility only. **clnt_create( )**, **clnt_tli_create( )**, or **clnt_vc_create( )** (see **rpc_clnt_create**(3N)) should be used instead.

CLIENT * **clntudp_bufcreate(struct sockaddr_in** **addr***,* **u_long** *prognum***,** **u_long** *versnum***,** **struct timeval** *wait***,** **int** **fdp,* **u_int** *sendsz***,** **u_int** *recvsz***);**

Create a client handle for the remote program *prognum*, on *versnum*; the client uses UDP / IP as the transport. The remote program is located at the Internet address *addr*. If *addr*→*sin_port* is **0**, it is set to port on which the remote program is listening on (the remote **rpcbind** service is consulted for this information). The parameter **fdp* is a file descriptor, which may be open and bound; if it is **RPC_ANYSOCK**, then this routine opens a new one and sets **fdp*. Refer to the **File Descriptor** section for more information. The UDP transport resends the call message in intervals of *wait* time until a response is received or until the call times out. The total time for the call to time out is specified by **clnt_call( )** (see **rpc_clnt_calls**(3N)). If successful it returns a client handle, otherwise it returns NULL. The error can be printed using the **clnt_pcreateerror( )** (see **rpc_clnt_create**(3N)) routine.

The user can specify the maximum packet size for sending and receiving by using *sendsz* and *recvsz* arguments for UDP-based RPC messages.

Warning: if *addr*→*sin_port* is **0** and the requested version number *versnum* is not registered with the remote portmap service, it returns a handle if at least a version number for the given program number is registered. The version mismatch is discovered by a **clnt_call( )** later (see **rpc_clnt_calls**(3N)).

Warning: this routine exists for backward compatibility only. **clnt_tli_create( )** or **clnt_dg_create( )** (see **rpc_clnt_create**(3N)) should be used instead.

CLIENT * **clntudp_create(struct sockaddr_in** **addr***,* **u_long** *prognum***,** **u_long** *versnum***,** **struct timeval** *wait***,** **int** **fdp***);*

This routine creates an RPC client handle for the remote program *prognum*, version *versnum*; the client uses UDP / IP as a transport. The remote program is located at Internet address *addr*. If *addr*→*sin_port* is **0**, then it is set to actual port that the remote program is listening on (the remote **rpcbind** service is consulted for this information). The parameter **fdp* is a file descriptor, which may be open and bound; if it is **RPC_ANYSOCK**, then this routine opens a new one and sets **fdp*. Refer to the **File Descriptor** section for more information. The UDP transport resends the call message in intervals of **wait** time until a response is received or until the call times out. The total time for the call to time out is specified by **clnt_call( )** (see **rpc_clnt_calls**(3N)). **clntudp_create( )** returns a client handle on success, otherwise it returns NULL. The error can be printed using the **clnt_pcreateerror( )** (see **rpc_clnt_create**(3N)) routine.

Warning: since UDP-based RPC messages can only hold up to 8 Kbytes of encoded data, this transport cannot be used for procedures that take large arguments or return huge results.

Warning: this routine exists for backward compatibility only. **clnt_create( )**, **clnt_tli_create( )**, or **clnt_dg_create( )** (see **rpc_clnt_create**(3N)) should be used instead.

void get_myaddress(struct sockaddr_in ∗*addr***);**

Places the local system's IP address into ∗*addr*, without consulting the library routines that deal with **/etc/hosts**. The port number is always set to **htons(PMAPPORT)**.

Warning: this routine is only intended for use with the RPC library. It returns the local system's address in a form compatible with the RPC library, and should not be taken as the system's actual IP address. In fact, the ∗*addr* buffer's host address part is actually zeroed. This address may have only local significance and should **NOT** be assumed to be an address that can be used to connect to the local system by remote systems or processes.

Warning: this routine remains for backward compatibility only. The routine **netdir_getbyname( )** (see **netdir**(3N)) should be used with the name **HOST_SELF** to retrieve the local system's network address as a *netbuf* structure.

void getrpcport(char ∗*host***, int** *prognum***, int** *versnum***, int** *proto***)**

getrpcport( ) returns the port number for the version *versnum* of the RPC program *prognum* running on *host* and using protocol *proto*. **getrpcport( )** returns **0** if the RPC system failed to contact the remote portmap service, the program associated with *prognum* is not registered, or there is no mapping between the program and a port.

Warning: This routine exists for backward compatibility only. Enhanced functionality is provided by **rpcb_getaddr( )** (see **rpcbind**(3N)).

struct pmaplist ∗ **pmap_getmaps(struct sockaddr_in** ∗*addr***);**

A user interface to the **portmap** service, which returns a list of the current RPC program-to-port mappings on the host located at IP address *addr*. This routine can return NULL . The command '**rpcinfo** −**p**' uses this routine.

Warning: this routine exists for backward compatibility only, enhanced functionality is provided by **rpcb_getmaps( )** (see **rpcbind**(3N)).

u_short pmap_getport(struct sockaddr_in ∗*addr***, u_long** *prognum***, u_long** *versnum***, u_long** *protocol***);**

A user interface to the **portmap** service, which returns the port number on which waits a service that supports program *prognum*, version *versnum*, and speaks the

transport protocol associated with *protocol*. The value of *protocol* is most likely **IPPROTO_UDP** or **IPPROTO_TCP**. A return value of **0** means that the mapping does not exist or that the RPC system failured to contact the remote **portmap** service. In the latter case, the global variable **rpc_createerr** contains the RPC status.

Warning: this routine exists for backward compatibility only, enhanced functionality is provided by **rpcb_getaddr( )** (see **rpcbind**(3N)).

enum clnt_stat pmap_rmtcall(struct sockaddr_in ∗*addr***, u_long** *prognum***, u_long** *versnum***, u_long** *procnum***, char** ∗*in***, xdrproct_t** *inproc***, char** ∗*out***, xdrproct_t** *outproc***, struct timeval** *tout***, u_long** ∗*portp***);**

Request that the **portmap** on the host at IP address ∗*addr* make an RPC on the behalf of the caller to a procedure on that host. ∗*portp* is modified to the program's port number if the procedure succeeds. The definitions of other parameters are discussed in **callrpc( )** and **clnt_call( )** (see **rpc_clnt_calls**(3N)).

Note: this procedure is only available for the UDP transport.

Warning: if the requested remote procedure is not registered with the remote **portmap** then no error response is returned and the call times out. Also, no authentication is done.

Warning: this routine exists for backward compatibility only, enhanced functionality is provided by **rpcb_rmtcall( )** (see **rpcbind**(3N)).

bool_t pmap_set(u_long *prognum***, u_long** *versnum***, u_long** *protocol***, u_short** *port***);**

A user interface to the **portmap** service, that establishes a mapping between the triple [*prognum*, *versnum*, *protocol*] and *port* on the machine's **portmap** service. The value of *protocol* may be **IPPROTO_UDP** or **IPPROTO_TCP**. Formerly, the routine failed if the requested *port* was found to be in use. Now, the routine only fails if it finds that *port* is still bound. If *port is not bound, the routine* completes the requested registration. This routine returns **1** if it succeeds, **0** otherwise. Automatically done by **svc_register( )**.

Warning: this routine exists for backward compatibility only, enhanced functionality is provided by **rpcb_set( )** (see **rpcbind**(3N)).

bool_t pmap_unset(u_long *prognum***, u_long** *versnum***);**

A user interface to the **portmap** service, which destroys all mapping between the triple [*prognum*, *versnum*, *all-protocols*] and *port* on the machine's **portmap** service. This routine returns one if it succeeds, **0** otherwise.

Warning: this routine exists for backward compatibility only, enhanced functionality is provided by **rpcb_unset( )** (see **rpcbind**(3N)).

int svc_fds;

A global variable reflecting the RPC service side's read file descriptor bit mask; it is suitable as a parameter to the **select( )** call. This is only of interest if a service implementor does not call **svc_run( )**, but rather does his own asynchronous event processing. This variable is read-only (do not pass its address to **select( )**!), yet it may change after calls to **svc_getreq( )** or any creation routines. Similar to **svc_fdset**, but limited to 32 descriptors.

Warning: this interface is obsoleted by **svc_fdset** (see **rpc_svc_calls**(3N)).

struct sockaddr_in * **svc_getcaller(SVCXPRT** **xprt**);**

This routine returns the network address, represented as a **struct sockaddr_in**, of the caller of a procedure associated with the RPC service transport handle, *xprt*.

Warning: this routine exists for backward compatibility only, and is obsolete. The preferred interface is **svc_getrpccaller( )** (see **rpc_svc_reg**(3N)), which returns the address as a **struct netbuf**.

void svc_getreq(int *rdfds***);**

This routine is only of interest if a service implementor does not call **svc_run( )**, but instead implements custom asynchronous event processing. It is called when the **select( )** call has determined that an RPC request has arrived on some RPC file descriptors; *rdfds* is the resultant read file descriptor bit mask. The routine returns when all file descriptors associated with the value of *rdfds* have been serviced.
This routine is similar to **svc_getreqset( )** but is limited to 32 descriptors.

Warning: this interface is obsoleted by **svc_getreqset( ).**

SVCXPRT * **svcfd_create(int** *fd***, u_int** *sendsz***, u_int** *recvsz***);**

Create a service on top of any open and bound descriptor. Typically, this descriptor is a connected file descriptor for a stream protocol. Refer to the **File Descriptor** section for more information. *sendsz* and *recvsz* indicate sizes for the send and receive buffers. If they are **0**, a reasonable default is chosen.

Warning: this interface is obsoleted by **svc_fd_create( )** (see **rpc_svc_create**(3N)).

SVCXPRT * **svcraw_create(void);**

This routine creates an internal, memory-based RPC service transport, to which it returns a pointer. The transport is really a buffer within the process's address space, so the corresponding RPC client should live in the same address space; see **clntraw_create( )**. This routine allows simulation of RPC and acquisition of RPC overheads (such as round trip times), without any kernel interference. This routine returns NULL if it fails.

Warning: this routine exists for backward compatibility only, and has the same

functionality of **svc_raw_create( )** (see **rpc_svc_create**(3N)), which obsoletes it.

SVCXPRT * **svctcp_create(int** *fd***, u_int** *sendsz***, u_int** *recvsz***);**

This routine creates a TCP/IP-based RPC service transport, to which it returns a pointer. The transport is associated with the file descriptor *fd*, which may be **RPC_ANYSOCK**, in which case a new file descriptor is created. If the file descriptor is not bound to a local TCP port, then this routine binds it to an arbitrary port. Refer to the **File Descriptor** section for more information. Upon completion, *xprt→xp_fd* is the transport's file descriptor, and *xprt→xp_port* is the transport's port number. This routine returns NULL if it fails. Since TCP-based RPC uses buffered I/O, users may specify the size of buffers; values of **0** choose suitable defaults.

Warning: this routine exists for backward compatibility only. **svc_create( )**, **svc_tli_create( )**, or **svc_vc_create( )** (see **rpc_svc_create**(3N)) should be used instead.

SVCXPRT * **svcudp_bufcreate(int** *fd***, u_int** *sendsz***, u_int** *recvsz***);**

This routine creates a UDP/IP-based RPC service transport, to which it returns a pointer. The transport is associated with the file descriptor *fd*. If *fd* is **RPC_ANYSOCK ,** then a new file descriptor is created. If the file descriptor is not bound to a local UDP port, then this routine binds it to an arbitrary port. Upon completion, *xprt→xp_fd* is the transport's file descriptor, and *xprt→xp_port* is the transport's port number. Refer to the **File Descriptor** section for more information. This routine returns NULL if it fails.

The user specifies the maximum packet size for sending and receiving UDP-based RPC messages by using the *sendsz* and *recvsz* parameters.

Warning: this routine exists for backward compatibility only. **svc_tli_create( )**, or **svc_dg_create( )** (see **rpc_svc_create**(3N)) should be used instead.

SVCXPRT * **svcudp_create(int** *fd***);**

This routine creates a UDP/IP-based RPC service transport, to which it returns a pointer. The transport is associated with the file descriptor *fd*, which may be **RPC_ANYSOCK ,** in which case a new file descriptor is created. If the file descriptor is not bound to a local UDP port, then this routine binds it to an arbitrary port. Upon completion, *xprt→xp_fd* is the transport's file descriptor, and *xprt→xp_port* is the transport's port number. This routine returns NULL if it fails.

Warning: since UDP-based RPC messages can only hold up to 8 Kbytes of encoded data, this transport cannot be used for procedures that take large arguments or return huge results.

Warning: this routine exists for backward compatibility only. **svc_create( )**, **svc_tli_create( )**, or **svc_dg_create( )** (see **rpc_svc_create**(3N)) should be used instead.

registerrpc(u_long *prognum***, u_long** *versnum***, u_long** *procnum***, char** ∗**(**∗*procname*)**(),**
xdrproc_t *inproc***, xdrproc_t** *outproc***);**

Register program *prognum*, procedure *procname*, and version *versnum* with the RPC service package. If a request arrives for program *prognum*, version *versnum*, and procedure *procnum*, *procname* is called with a pointer to its parameter(s); *procname* should return a pointer to its static result(s); *inproc* is used to decode the parameters while *outproc* is used to encode the results. This routine returns **0** if the registration succeeded, –1 otherwise.

svc_run( ) must be called after all the services are registered.

Warning: this routine exists for backward compatibility only, and is obsoleted by **rpc_reg( )**.

svc_register(SVCXPRT ∗*xprt***, u_long** *prognum***, u_long** *versnum***, void (**∗*dispatch*)**(),**
u_long *protocol***);**

Associates *prognum* and *versnum* with the service dispatch procedure, *dispatch*. If *protocol* is **0**, the service is not registered with the **portmap** service. If *protocol* is non-zero, then a mapping of the triple [*prognum*, *versnum*, *protocol*] to *xprt→xp_port* is established with the local **portmap** service (generally *protocol* is **0**, **IPPROTO_UDP** or **IPPROTO_TCP**). The procedure *dispatch* has the following form:

dispatch(struct svc_req ∗*request***, SVCXPRT** ∗*xprt***);**

The **svc_register( )** routine returns one if it succeeds, and **0** otherwise.

Warning: this routine exists for backward compatibility only; enhanced functionality is provided by **svc_reg( )**.

void svc_unregister(u_long *prognum***, u_long** *versnum***);**

Remove all mapping of the double [*prognum*, *versnum*] to dispatch routines, and of the triple [*prognum*, *versnum*, *all-protocols*] to port number from **portmap**.

Warning: this routine exists for backward compatibility, enhanced functionality is provided by **svc_unreg( )**.

xdr_authunix_parms(XDR ∗*xdrs***, struct authunix_parms** ∗*aupp***);**

Used for describing UNIX credentials. This routine is useful for users who wish to generate these credentials without using the RPC authentication package.

Warning: this routine exists for backward compatibility only, and is obsoleted by **xdr_authsys_parms( )** (see **rpc_xdr**(3N)).

SEE ALSO

keyserv(1M), **rpcbind**(1M), **rpcinfo**(1M), **rpc**(3N), **rpc_clnt_auth**(3N), **rpc_clnt_calls**(3N), **rpc_clnt_create**(3N), **rpc_svc_calls**(3N), **rpc_svc_create**(3N), **rpc_svc_err**(3N), **rpc_svc_reg**(3N), **rpcbind**(3N), **secure_rpc**(3N), **select**(3C)

NOTES

These interfaces are unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

rpc_svc_calls, svc_dg_enablecache, svc_done, svc_exit, svc_fdset, svc_freeargs, svc_getargs, svc_getreq_common, svc_getreq_poll, svc_getreqset, svc_getrpccaller, svc_pollset, svc_run, svc_sendreply – library routines for RPC servers

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

These routines are part of the RPC library which allows C language programs to make procedure calls on other machines across the network.

These routines are associated with the server side of the RPC mechanism. Some of them are called by the server side dispatch function, while others (such as **svc_run( )**) are called when the server is initiated.

In the current implementation, the service transport handle **SVCXPRT** contains a single data area for decoding arguments and encoding results. Therefore, this structure cannot be freely shared between threads that call functions that do this. However, when a server is operating in the Automatic or User MT modes, a copy of this structure is passed to the service dispatch procedure in order to enable concurrent request processing. Under these circumstances, some routines which would otherwise be unsafe, become safe. These are marked as such. Also marked are routines that are unsafe for MT applications, and are not to be used by such applications.

Routines

#include <rpc/rpc.h>

int svc_dg_enablecache(SVCXPRT ∗*xprt*, **const unsigned long** *cache_size*);

This function allocates a duplicate request cache for the service endpoint *xprt*, large enough to hold *cache_size* entries. Once enabled, there is no way to disable caching. This routine returns **1** if space necessary for a cache of the given size was successfully allocated, and **0** otherwise.

This function is safe in MT applications.

int svc_done(SVCXPRT ∗*xprt*);

This function frees resources allocated to service a client request directed to the service endpoint *xprt*. This call pertains only to servers executing in the User MT mode. In the User MT mode, service procedures must invoke this call before returning, either after a client request has been serviced, or after an error or abnormal condition that prevents a reply from being sent. After **svc_done**() is invoked, the service endpoint *xprt* should not be referenced by the service procedure. Server multithreading modes and parameters can be set using the **rpc_control**() call.

This function is safe in MT applications. It will have no effect if invoked in modes other than the User MT mode.

void svc_exit(void);

This function when called by any of the RPC server procedure or otherwise, destroys all services registered by the server and causes **svc_run( )** to return.

If RPC server activity is to be resumed, services must be reregistered with the RPC library either through one of the **rpc_svc_create**(3N) functions, or using **xprt_register**(3N).

svc_exit( ) has global scope and ends all RPC server activity.

fd_set svc_fdset;

A global variable reflecting the RPC server's read file descriptor bit mask. This is only of interest if service implementors do not call **svc_run( )**, but rather do their own asynchronous event processing. This variable is read-only, and it may change after calls to **svc_getreqset( )** or any creation routines. Do not pass its address to **select**(3C)! Instead, pass the address of a copy.

MT applications executing in either the Automatic MT mode or the user MT mode should never read this variable. They should use auxiliary threads to do asynchronous event processing.

bool_t svc_freeargs(const SVCXPRT ∗*xprt***, const xdrproc_t** *inproc***, caddr_t** *in***);**

A function macro that frees any data allocated by the **RPC/XDR** system when it decoded the arguments to a service procedure using **svc_getargs( )**. This routine returns **TRUE** if the results were successfully freed, and **FALSE** otherwise.

This function macro is safe in MT applications utilizing the Automatic or User MT modes.

bool_t svc_getargs(const SVCXPRT ∗*xprt***, const xdrproc_t** *inproc***, caddr_t** *in***);**

A function macro that decodes the arguments of an RPC request associated with the RPC service transport handle *xprt*. The parameter *in* is the address where the arguments will be placed; *inproc* is the XDR routine used to decode the arguments. This routine returns **TRUE** if decoding succeeds, and **FALSE** otherwise.

This function macro is safe in MT applications utilizing the Automatic or User MT modes.

void svc_getreq_common(const int *fd***);**

This routine is called to handle a request on the given file descriptor.

void svc_getreq_poll(struct pollfd **pfdp***, const int** *pollretval***);**

This routine is only of interest if a service implementor does not call **svc_run( )**, but instead implements custom asynchronous event processing. It is called when **poll**(2) has determined that an RPC request has arrived on some RPC file descriptors; *pollretval* is the return value from **poll**(2) and *pfdp* is the array of *pollfd* structures on which the **poll**(2) was done. It is assumed to be an array large enough to contain the maximal number of descriptors allowed.

This function macro is unsafe in MT applications.

void svc_getreqset(fd_set **rdfds***);**

This routine is only of interest if a service implementor does not call **svc_run( )**, but instead implements custom asynchronous event processing. It is called when **select**(3C) has determined that an RPC request has arrived on some RPC file descriptors; *rdfds* is the resultant read file descriptor bit mask. The routine returns when all file descriptors associated with the value of *rdfds* have been serviced.

This function macro is unsafe in MT applications.

struct netbuf *svc_getrpccaller(const SVCXPRT **xprt***);**

The approved way of getting the network address of the caller of a procedure associated with the RPC service transport handle *xprt*.

This function macro is safe in MT applications.

void svc_run(void);

This routine never returns. In single threaded mode, it waits for RPC requests to arrive, and calls the appropriate service procedure using **svc_getreq_poll( )** when one arrives. This procedure is usually waiting for the **poll**(2) library call to return.

Applications executing in the Automatic or User MT modes should invoke this function exactly once. It the Automatic MT mode, it will create threads to service client requests. In the User MT mode, it will provide a framework for service developers to create and manage their own threads for servicing client requests.

bool_t svc_sendreply(const SVCXPRT **xprt*, **const xdrproc_t** *outproc*, **const caddr_t** *out*);

Called by an RPC service's dispatch routine to send the results of a remote procedure call. The parameter *xprt* is the request's associated transport handle; *outproc* is the XDR routine which is used to encode the results; and *out* is the address of the results. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

This function macro is safe in MT applications utilizing the Automatic or User MT modes.

SEE ALSO

rpcgen(1), **poll**(2), **rpc**(3N), **rpc_control**(3N), **rpc_svc_create**(3N), **rpc_svc_err**(3N), **rpc_svc_reg**(3N), **select**(3C), **xprt_register**(3N)

NOTES

svc_dg_enablecache( ) and **svc_getrpccaller( )** are safe in multithreaded applications. **svc_freeargs( )**, **svc_getargs( )**, and **svc_sendreply( )** are safe in MT applications utilizing the Automatic or User MT modes. **svc_getreq_common( )**, **svc_getreqset( )**, and **svc_getreq_poll( )** are unsafe in multithreaded applications and should be called only from the main thread.

NAME

rpc_svc_create, svc_control, svc_create, svc_destroy, svc_dg_create, svc_fd_create, svc_raw_create, svc_tli_create, svc_tp_create, svc_vc_create – library routines for the creation of server handles

MT-LEVEL

MT-Safe

DESCRIPTION

These routines are part of the RPC library which allows C language programs to make procedure calls on servers across the network. These routines deal with the creation of service handles. Once the handle is created, the server can be invoked by calling **svc_run( )**.

Routines

See **rpc**(3N) for the definition of the **SVCXPRT** data structure.

#include <rpc/rpc.h>

bool_t
svc_control (SVCXPRT ∗*svc*, **const u_int** *req*, **void** ∗*info*);

A function to change or retrieve various information about a service object. *req* indicates the type of operation and *info* is a pointer to the information. The supported values of *req*, their argument types, and what they do are:

SVCGET_VERSQUIET If a request is received for a program number served by this server but the version number is outside the range registered with the server, an RPC_PROGVERSMISMATCH error will normally be returned. *info* should be a pointer to an integer. Upon successful completion of the SVCGET_VERSQUIET request, ∗*info* contains an integer which describes the server's current behavior: 0 indicates normal server behavior (that is, an RPC_PROGVERSMISMATCH error will be returned); 1 indicates that the out of range request will be silently ignored.

SVCSET_VERSQUIET If a request is received for a program number served by this server but the version number is outside the range registered with the server, an RPC_PROGVERSMISMATCH error will normally be returned. It is sometimes desirable to change this behavior. *info* should be a pointer to an integer which is either 0 (indicating normal server behavior – an RPC_PROGVERSMISMATCH error will be returned), or 1 (indicating that the out of range request should be silently ignored).

int svc_create(const void (∗*dispatch*)(**const struct svc_req** ∗, **const SVCXPRT** ∗), **const u_long** *prognum*, **const u_long** *versnum*, **const char** ∗*nettype*);

svc_create( ) creates server handles for all the transports belonging to the class *nettype*.

nettype defines a class of transports which can be used for a particular

application. The transports are tried in left to right order in **NETPATH** variable or in top to bottom order in the netconfig database. If *nettype* is **NULL**, it defaults to **netpath**.

svc_create( ) registers itself with the **rpcbind** service (see **rpcbind**(1M)). *dispatch* is called when there is a remote procedure call for the given *prognum* and *versnum*; this requires calling **svc_run( )** (see **svc_run( )** in **rpc_svc_reg**(3N)). If **svc_create( )** succeeds, it returns the number of server handles it created, otherwise it returns **0** and an error message is logged.

void svc_destroy(SVCXPRT ∗*xprt***);**

A function macro that destroys the RPC service handle *xprt*. Destruction usually involves deallocation of private data structures, including *xprt* itself. Use of *xprt* is undefined after calling this routine.

SVCXPRT ∗**svc_dg_create(const int** *fildes***, const u_int** *sendsz***, const u_int** *recvsz***);**

This routine creates a connectionless RPC service handle, and returns a pointer to it. This routine returns **NULL** if it fails, and an error message is logged. *sendsz* and *recvsz* are parameters used to specify the size of the buffers. If they are **0**, suitable defaults are chosen. The file descriptor *fildes* should be open and bound. The server is not registered with **rpcbind**(1M).

Warning: since connectionless-based RPC messages can only hold limited amount of encoded data, this transport cannot be used for procedures that take large arguments or return huge results.

SVCXPRT ∗**svc_fd_create(const int** *fildes***, const u_int** *sendsz***, const u_int** *recvsz***);**

This routine creates a service on top of an open and bound file descriptor, and returns the handle to it. Typically, this descriptor is a connected file descriptor for a connection-oriented transport. *sendsz* and *recvsz* indicate sizes for the send and receive buffers. If they are **0**, reasonable defaults are chosen. This routine returns **NULL** if it fails, and an error message is logged.

SVCXPRT ∗**svc_raw_create(void);**

This routine creates an RPC service handle and returns a pointer to it. The transport is really a buffer within the process's address space, so the corresponding RPC client should live in the same address space; (see **clnt_raw_create( )** in **rpc_clnt_create**(3N)). This routine allows simulation of RPC and acquisition of RPC overheads (such as round trip times), without any kernel and networking interference. This routine returns **NULL** if it fails, and an error message is logged.

Note: **svc_run( )** should not be called when the raw interface is being used.

SVCXPRT ∗**svc_tli_create(const int** *fildes***, const struct netconfig** ∗*netconf***,**
const struct t_bind ∗*bindaddr***, const u_int** *sendsz***, const u_int** *recvsz***);**

This routine creates an RPC server handle, and returns a pointer to it. *fildes* is the file descriptor on which the service is listening. If *fildes* is **RPC_ANYFD**, it opens a file descriptor on the transport specified by *netconf*. If the file descriptor is unbound and *bindaddr* is non-null *fildes* is bound to the address specified by *bindaddr*, otherwise *fildes* is bound to a default address chosen by the transport. In the case where the default address is chosen, the number of outstanding connect requests is set to 8 for connection-oriented transports. The user may specify the size of the send and receive buffers with the parameters *sendsz* and *recvsz ;* values of **0** choose suitable defaults. This routine returns **NULL** if it fails, and an error message is logged. The server is not registered with the **rpcbind**(1M) service.

SVCXPRT ∗**svc_tp_create(const void (**∗*dispatch*)**(const struct svc_req** ∗**,**
const SVCXPRT ∗**), const u_long** *prognum***, const u_long** *versnum***,**
const struct netconfig ∗*netconf***);**

svc_tp_create( ) creates a server handle for the network specified by *netconf*, and registers itself with the **rpcbind** service. *dispatch* is called when there is a remote procedure call for the given *prognum* and *versnum*; this requires calling **svc_run( )**. **svc_tp_create( )** returns the service handle if it succeeds, otherwise a **NULL** is returned and an error message is logged.

SVCXPRT ∗**svc_vc_create(const int** *fildes***, const u_int** *sendsz***, const u_int** *recvsz***);**

This routine creates a connection-oriented RPC service and returns a pointer to it. This routine returns **NULL** if it fails, and an error message is logged. The users may specify the size of the send and receive buffers with the parameters *sendsz* and *recvsz*; values of **0** choose suitable defaults. The file descriptor *fildes* should be open and bound. The server is not registered with the **rpcbind**(1M) service.

SEE ALSO

rpcbind(1M), **rpc**(3N), **rpc_svc_calls**(3N), **rpc_svc_err**(3N), **rpc_svc_reg**(3N)

NAME

rpc_svc_err, svcerr_auth, svcerr_decode, svcerr_noproc, svcerr_noprog, svcerr_progvers, svcerr_systemerr, svcerr_weakauth – library routines for server side remote procedure call errors

MT-LEVEL

MT-Safe

DESCRIPTION

These routines are part of the RPC library which allows C language programs to make procedure calls on other machines across the network.

These routines can be called by the server side dispatch function if there is any error in the transaction with the client.

Routines

See **rpc**(3N) for the definition of the **SVCXPRT** data structure.

#include <rpc/rpc.h>

void svcerr_auth(const SVCXPRT **xprt*, **const enum auth_stat** *why*);

Called by a service dispatch routine that refuses to perform a remote procedure call due to an authentication error.

void svcerr_decode(const SVCXPRT **xprt*);

Called by a service dispatch routine that cannot successfully decode the remote parameters (see **svc_getargs( )** in **rpc_svc_reg**(3N)).

void svcerr_noproc(const SVCXPRT **xprt*);

Called by a service dispatch routine that does not implement the procedure number that the caller requests.

void svcerr_noprog(const SVCXPRT **xprt*);

Called when the desired program is not registered with the RPC package. Service implementors usually do not need this routine.

void svcerr_progvers(const SVCXPRT **xprt*, **u_long** *low_vers*, **u_long** *high_vers*);

Called when the desired version of a program is not registered with the RPC package. *low_vers* is the lowest version number, and *high_vers* is the highest version number. Service implementors usually do not need this routine.

void svcerr_systemerr(const SVCXPRT **xprt*);

Called by a service dispatch routine when it detects a system error not covered by any particular protocol. For example, if a service can no longer allocate storage, it may call this routine.

void svcerr_weakauth(const SVCXPRT **xprt****);**

Called by a service dispatch routine that refuses to perform a remote procedure call due to insufficient (but correct) authentication parameters. The routine calls **svcerr_auth(xprt, AUTH_TOOWEAK)**.

SEE ALSO

rpc(3N), **rpc_svc_calls**(3N), **rpc_svc_create**(3N), **rpc_svc_reg**(3N)

NAME

rpc_svc_reg, rpc_reg, svc_reg, svc_unreg, svc_auth_reg, xprt_register, xprt_unregister – library routines for registering servers

MT-LEVEL

MT-Safe

DESCRIPTION

These routines are a part of the RPC library which allows the RPC servers to register themselves with **rpcbind( )** (see **rpcbind**(1M)), and associate the given program and version number with the dispatch function. When the RPC server receives a RPC request, the library invokes the dispatch routine with the appropriate arguments.

Routines

See **rpc**(3N) for the definition of the **SVCXPRT** data structure.

#include <rpc/rpc.h>

bool_t rpc_reg(u_long *prognum***, u_long** *versnum***, u_long** *procnum***,**
char ∗ **const(**∗*procname***) (char** ∗*arg***), xdrproc_t** *inproc***, xdrproc_t** *outproc***,**
const char ∗*nettype***);**

Register program *prognum*, procedure *procname*, and version *versnum* with the RPC service package. If a request arrives for program *prognum*, version *versnum*, and procedure *procnum*, *procname* is called with a pointer to its parameter(s); *procname* should return a pointer to its **static** result(s). The *arg* parameter to *procname* is a pointer to the (decoded) procedure argument. *inproc* is the XDR function used to decode the parameters while *outproc* is the XDR function used to encode the results. Procedures are registered on all available transports of the class *nettype*. See **rpc**(3N). This routine returns **0** if the registration succeeded, **–1** otherwise.

int svc_reg(const SVCXPRT ∗*xprt***, const u_long** *prognum***, const u_long** *versnum***,**
const void (∗*dispatch***), const struct netconfig** ∗*netconf***);**

Associates *prognum* and *versnum* with the service dispatch procedure, *dispatch*. If *netconf* is **NULL**, the service is not registered with the *rpcbind* service. For example, if a service has already been registered using some other means, such as **inetd** (see **inetd**(1M)), it will not need to be registered again. If *netconf* is nonzero, then a mapping of the triple [*prognum*, *versnum*, *netconf→nc_netid*] to *xprt→xp_ltaddr* is established with the local **rpcbind** service.

The **svc_reg( )** routine returns **1** if it succeeds, and **0** otherwise.

void svc_unreg(const u_long *prognum***, const u_long** *versnum***);**

Remove from the **rpcbind** service, all mappings of the triple [*prognum*, *versnum*, *all-transports*] to network address and all mappings within the RPC service package of the double [*prognum*, *versnum*] to dispatch routines.

int svc_auth_reg(const int *cred_flavor***, const enum auth_stat (****handler***));**

Registers the service authentication routine *handler* with the dispatch mechanism so that it can be invoked to authenticate RPC requests received with authentication type *cred_flavor*. This interface allows developers to add new authentication types to their RPC applications without needing to modify the libraries. Service implementors usually do not need this routine.

Typical service application would call **svc_auth_reg( )** after registering the service and prior to calling **svc_run( )**. When needed to process an RPC credential of type *cred_flavor*, the *handler* procedure will be called with two parameters **(struct svc_req** **rqst***, struct rpc_msg** **msg***)** and is expected to return a valid **enum auth_stat** value. There is no provision to change or delete an authentication handler once registered.

The **svc_auth_reg( )** routine returns **0** if the registration is successful, **1** if *cred_flavor* already has an authentication handler registered for it, and **–1** otherwise.

void xprt_register(const SVCXPRT **xprt***);**

After RPC service transport handle *xprt* is created, it is registered with the RPC service package. This routine modifies the global variable **svc_fdset** (see **rpc_svc_calls**(3N)). Service implementors usually do not need this routine.

void xprt_unregister(const SVCXPRT **xprt***);**

Before an RPC service transport handle *xprt* is destroyed, it unregisters itself with the RPC service package. This routine modifies the global variable **svc_fdset** (see **rpc_svc_calls**(3N)). Service implementors usually do not need this routine.

SEE ALSO

inetd(1M), **rpcbind**(1M), **rpc**(3N), **rpc_svc_calls**(3N), **rpc_svc_create**(3N), **rpc_svc_err**(3N), **rpcbind**(3N), **select**(3C)

NAME

rpc_xdr, xdr_accepted_reply, xdr_authsys_parms, xdr_callhdr, xdr_callmsg, xdr_opaque_auth, xdr_rejected_reply, xdr_replymsg – XDR library routines for remote procedure calls

MT-LEVEL

Safe

DESCRIPTION

These routines are used for describing the RPC messages in XDR language. They should normally be used by those who do not want to use the RPC package directly. These routines return TRUE if they succeed, FALSE otherwise.

Routines

See **rpc**(3N) for the definition of the **XDR** data structure.

#include <rpc/rpc.h>

bool_t xdr_accepted_reply(XDR ∗*xdrs*, **const struct accepted_reply** ∗*ar*);

Used to translate between RPC reply messages and their external representation. It includes the status of the RPC call in the XDR language format. In the case of success, it also includes the call results.

bool_t xdr_authsys_parms(XDR ∗*xdrs*, **struct authsys_parms** ∗*aupp*);

Used for describing UNIX operating system credentials. It includes machine-name, uid, gid list, etc.

void xdr_callhdr(XDR ∗*xdrs*, **struct rpc_msg** ∗*chdr*);

Used for describing RPC call header messages. It encodes the static part of the call message header in the XDR language format. It includes information such as transaction ID, RPC version number, program and version number.

bool_t xdr_callmsg(XDR ∗*xdrs*, **struct rpc_msg** ∗*cmsg*);

Used for describing RPC call messages. This includes all the RPC call information such as transaction ID, RPC version number, program number, version number, authentication information, etc. This is normally used by servers to determine information about the client RPC call.

bool_t xdr_opaque_auth(XDR ∗*xdrs*, **struct opaque_auth** ∗*ap*);

Used for describing RPC opaque authentication information messages.

bool_t xdr_rejected_reply(XDR ∗*xdrs*, **const struct rejected_reply** ∗*rr*);

Used for describing RPC reply messages. It encodes the rejected RPC message in the XDR language format. The message could be rejected either because of version number mis-match or because of authentication errors.

bool_t xdr_replymsg(XDR **xdrs***, const struct rpc_msg** **rmsg***);**

Used for describing RPC reply messages. It translates between the RPC reply message and its external representation. This reply could be either an acceptance, rejection or **NULL**.

SEE ALSO

rpc(3N), **xdr**(3N)

NAME

rpcbind, rpcb_getmaps, rpcb_getaddr, rpcb_gettime, rpcb_rmtcall, rpcb_set, rpcb_unset – library routines for RPC bind service

MT-LEVEL

MT-Safe

DESCRIPTION

These routines allow client C programs to make procedure calls to the RPC binder service. **rpcbind** (see **rpcbind**(1M)) maintains a list of mappings between programs and their universal addresses.

Routines

#include <rpc/rpc.h>

struct rpcblist ∗**rpcb_getmaps(const struct netconfig** ∗*netconf*, **const char** ∗*host*);

An interface to the **rpcbind** service, which returns a list of the current RPC program-to-address mappings on *host*. It uses the transport specified through *netconf* to contact the remote **rpcbind** service on *host*. This routine will return **NULL**, if the remote **rpcbind** could not be contacted.

bool_t rpcb_getaddr(const u_long *prognum*, **const u_long** *versnum*,
const struct netconfig ∗*netconf*, **struct netbuf** ∗*svcaddr*, **const char** ∗*host*);

An interface to the **rpcbind** service, which finds the address of the service on *host* that is registered with program number *prognum*, version *versnum*, and speaks the transport protocol associated with *netconf*. The address found is returned in *svcaddr*. *svcaddr* should be preallocated. This routine returns **TRUE** if it succeeds. A return value of **FALSE** means that the mapping does not exist or that the RPC system failed to contact the remote **rpcbind** service. In the latter case, the global variable **rpc_createerr** (see **rpc_clnt_create**(3N)) contains the RPC status.

bool_t rpcb_gettime(const char ∗*host*, **time_t** ∗*timep*);

This routine returns the time on *host* in *timep*. If *host* is **NULL**, **rpcb_gettime( )** returns the time on its own machine. This routine returns **TRUE** if it succeeds, **FALSE** if it fails. **rpcb_gettime( )** can be used to synchronize the time between the client and the remote server. This routine is particularly useful for secure RPC.

enum clnt_stat rpcb_rmtcall(const struct netconfig ∗*netconf*, **const char** ∗*host*,
const u_long *prognum*, **const u_long** *versnum*, **const u_long** *procnum*,
const xdrproc_t *inproc*, **const caddr_t** *in*, **const xdrproc_t** *outproc*,
caddr_t *out*, **const struct timeval** *tout*, **struct netbuf** ∗*svcaddr*);

An interface to the **rpcbind** service, which instructs **rpcbind** on *host* to make an RPC call on your behalf to a procedure on that host. The **netconfig** structure should correspond to a connectionless transport. The parameter ∗*svcaddr* will be modified to the server's address if the procedure succeeds (see **rpc_call( )** and **clnt_call( )** in **rpc_clnt_calls**(3N) for the definitions of other parameters).

This procedure should normally be used for a "ping" and nothing else. This routine allows programs to do lookup and call, all in one step.

Note: Even if the server is not running **rpcbind** does not return any error messages to the caller. In such a case, the caller times out.

Note: **rpcb_rmtcall( )** is only available for connectionless transports.

bool_t rpcb_set(const u_long *prognum***, const u_long** *versnum***,**
const struct netconfig ∗*netconf***, const struct netbuf** ∗*svcaddr***);**

An interface to the **rpcbind** service, which establishes a mapping between the triple [*prognum*, *versnum*, *netconf→nc_netid*] and *svcaddr* on the machine's **rpcbind** service. The value of *nc_netid* must correspond to a network identifier that is defined by the netconfig database. This routine returns **TRUE** if it succeeds, **FALSE** otherwise. (See also **svc_reg( )** in **rpc_svc_calls**(3N)). If there already exists such an entry with **rpcbind**, **rpcb_set( )** will fail.

bool_t rpcb_unset(const u_long *prognum***, const u_long** *versnum***,**
const struct netconfig ∗*netconf***);**

An interface to the **rpcbind** service, which destroys the mapping between the triple [*prognum*, *versnum*, *netconf→nc_netid*] and the address on the machine's **rpcbind** service. If *netconf* is **NULL**, **rpcb_unset( )** destroys all mapping between the triple [*prognum*, *versnum*, *all-transports*] and the addresses on the machine's **rpcbind** service. This routine returns **TRUE** if it succeeds, **FALSE** otherwise. Only the owner of the service or the super-user can destroy the mapping. (See also **svc_unreg( )** in **rpc_svc_calls**(3N)).

SEE ALSO

rpcbind(1M), **rpcinfo**(1M), **rpc_clnt_calls**(3N), **rpc_svc_calls**(3N)

NAME

rstat, havedisk – get performance data from remote kernel

PROTOCOL

/usr/include/rpcsvc/rstat.x

SYNOPSIS

cc [*flag* ...] *file* ... **–lrpcsvc** [*library* ...]

```
#include <rpc/rpc.h>
#include <rpcsvc/rstat.h>

enum clnt_stat rstat(char *host, struct statstime *statp);

havedisk(char *host);
```

MT-LEVEL

MT-Safe

DESCRIPTION

These routines require that the **rpc.rstatd**(1M) daemon be configured and available on the remote system indicated by *host*. The **rstat( )** protocol is used to gather statistics from remote kernel. Statistics will be available on items such as paging, swapping, and cpu utilization.

rstat( ) fills in the **statstime** structure *statp* for *host*. *statp* must point to an allocated **statstime** structure. **rstat( )** returns **RPC_SUCCESS** if it was successful; otherwise a **enum clnt_stat** is returned which can be displayed using **clnt_perrno**(3N).

havedisk( ) returns **1** if *host* has disk, **0** if it does not, and **-1** if this cannot be determined.

The following XDR routines are available in **librpcsvc**:

xdr_statstime
xdr_statsvar

SEE ALSO

rpc.rstatd(1M), **rup**(1), **rpc_clnt_calls**(3N)

NAME

rusers, rnusers – return information about users on remote machines

PROTOCOL

/usr/include/rpcsvc/rusers.x

SYNOPSIS

cc [*flag* ...] *file* ... **–lrpcsvc** [*library* ...]

```
#include <rpc/rpc.h>
#include <rpcsvc/rusers.h>
enum clnt_stat rusers(char *host, struct utmpidlearr *up);
int rnusers(char *host);
```

MT-LEVEL

MT-Safe

DESCRIPTION

These routines require that the **rpc.rusersd**(1M) daemon be configured and available on the remote system indicated by *host*. The **rusers( )** protocol is used to retrieve information about users logged in on the remote system.

rusers( ) fills the **utmpidlearr** structure with data about *host*, and returns **0** if successful. *up* must point to an allocated **utmpidlearr** structure. If **rusers( )** returns successful it will have allocated data structures within the *up* structure, which should be freed with **xdr_free**(3N) when you no longer need them:

```
xdr_free(xdr_utimpidlearr, up);
```

On error, the returned value can be interpreted as an **enum clnt_stat** and can be displayed with **clnt_perror**(3N) or **clnt_sperrno**(3N).

See the header **<rpcsvc/rusers.h>** for a definition of struct **utmpidlearr**.

rnusers( ) returns the number of users logged on to *host* (**-1** if it cannot determine that number).

The following XDR routines are available in **librpcsvc**:

```
xdr_utmpidlearr.
```

SEE ALSO

rusers(1), **rpc.rusersd**(1M), **rpc_clnt_calls**(3N), **xdr_free**(3N)

NAME

rwall – write to specified remote machines

PROTOCOL

/usr/include/rpcsvc/rwall.x

SYNOPSIS

cc [*flag* ...] *file* ... **–lrpcsvc** [*library* ...]

```
#include <rpc/rpc.h>
#include <rpcsvc/rwall.h>
enum clnt_stat rwall(char *host, char *msg);
```

MT-LEVEL

MT-Safe

DESCRIPTION

These routines require that the **rpc.rwalld**(1M) daemon be configured and available on the remote system indicated by *host.*

rwall( ) executes **wall**(1M) on *host*. The **rpc.rwalld** process on *host* prints *msg* to all users logged on to that system. **rwall( )** returns **RPC_SUCCESS** if it was successful; otherwise a **enum clnt_stat** is returned which can be displayed using **clnt_perrno**(3N).

SEE ALSO

rpc.rwalld(1M), **wall**(1M), **rpc_clnt_calls**(3N)

NAME

rwlock, rwlock_init, rwlock_destroy, rw_rdlock, rw_wrlock, rw_tryrdlock, rw_trywrlock, rw_unlock – multiple readers, single writer locks

SYNOPSIS

cc [*flag* ...] *file* ... **–lthread –lc** [*library* ...]

#include <synch.h>

int rwlock_init(rwlock_t **rwlp***, int** *type***, void** ** arg***);**

int rwlock_destroy(rwlock_t **rwlp***);**

int rw_rdlock(rwlock_t **rwlp***);**

int rw_wrlock(rwlock_t **rwlp***);**

int rw_unlock(rwlock_t **rwlp***);**

int rw_tryrdlock(rwlock_t **rwlp***);**

int rw_trywrlock(rwlock_t **rwlp***);**

MT-LEVEL

MT-Safe

DESCRIPTION

Many threads can have simultaneous read-only access to data, while only one thread can have write access at any given time. Multiple read access with single write access is controlled by locks, which are generally used to protect data that is frequently searched.

Readers / writer locks can synchronize threads in this process and other processes if they are allocated in writable memory and shared among cooperating processes (see **mmap**(2)), and are initialized for this purpose.

Additionally, readers / writer locks must be initialized prior to use. **rwlock_init( )** The readers / writer lock pointed to by *rwlp* is initialized by **rwlock_init( )**. A readers / writer lock is capable of having several types of behavior, which is specified by *type*. *arg* is currently not used, although a future type may define new behavior parameters via *arg*.

type may be one of the following:

USYNC_PROCESS The readers / writer lock can synchronize threads in this process and other processes. The readers / writer lock should be initialized by only one lock. *arg* is ignored.

USYNC_THREAD The readers / writer lock can synchronize threads in this process, only. *arg* is ignored.

Additionally, readers / writer locks can be initialized by allocation in zeroed memory. A *type* of **USYNC_THREAD** is assumed in this case. Multiple threads must not simultaneously initialize the same readers / writer lock. And a readers / writer lock must not be reinitialized while in use by other threads.

The following are default readers / writer lock initialization (intra-process):

```
rwlock_t rwlp;

rwlock_init(&rwlp, NULL, NULL);
```

OR

rwlock_init(&rwlp, USYNC_THREAD, NULL);

OR

rwlock_t rwlp = DEFAULTRWLOCK;

The following is a customized readers / writer lock initialization (inter-process):

rwlock_init(&rwlp, USYNC_PROCESS, NULL);

Any state associated with the readers / writer lock pointed to by *rwlp* are destroyed by **rwlock_destroy( )** and the readers / writer lock storage space is not released.

rw_rdlock( ) gets a read lock on the readers / writer lock pointed to by *rwlp*. If the readers / writer lock is currently locked for writing, the calling thread blocks until the write lock is freed. Multiple threads may simultaneously hold a read lock on a readers / writer lock.

rw_tryrdlock( ) trys to get a read lock on the readers / writer lock pointed to by *rwlp*. If the readers / writer lock is locked for writing, it returns an error; otherwise, the read lock is acquired.

rw_wrlock( ) gets a write lock on the readers / writer lock pointed to by *rwlp*. If the readers / writer lock is currently locked for reading or writing, the calling thread blocks until all the read and write locks are freed. At any given time, only one thread may have a write lock on a readers / writer lock.

rw_trywrlock( ) trys to get a write lock on the readers / writer lock pointed to by *rwlp*. If the readers / writer lock is currently locked for reading or writing, it returns an error.

rw_unlock( ) unlocks a readers / writer lock pointed to by *rwlp*, if the readers / writer lock is locked and the calling thread holds the lock for either reading or writing. One of the other threads that is waiting for the readers / writer lock to be freed will be unblocked, provided there is other waiting threads. If the calling thread does not hold the lock for either reading or writing, no error status is returned, and the program's behavior is unknown.

RETURN VALUES

Upon successful completion, **0** is returned; otherwise, a non-zero value indicates an error.

ERRORS

These functions fail and return the corresponding value if any of the following conditions are detected.

EINVAL Invalid argument.

EFAULT *rwlp* or *arg* point to an illegal address.

rw_tryrdlock( ) or **rw_trywrlock( )** fails and returns the corresponding value if any of the following conditions are detected.

EBUSY The readers / writer lock pointed to by *rwlp* was already locked.

SEE ALSO **mmap**(2)

NOTES These interfaces also available via:

#include <thread.h>

If multiple threads are waiting for a readers / writer lock, the acquisition order is random by default. However, some implementations may bias acquisition order to avoid depriving writers. The current implementation favors writers over readers.

NAME

scandir, alphasort – scan a directory

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

```
#include <sys/types.h>
#include <sys/dir.h>

int scandir(dirname, namelist , select, dcomp)
char *dirname;
struct direct *(*namelist[ ]);
int (*select)(. ), (*dcomp)( );

int alphasort( d1, d2)
struct direct **d1, **d2;
```

DESCRIPTION

scandir( ) reads the directory *dirname* and builds an array of pointers to directory entries using **malloc**(3C). The second parameter is a pointer to an array of structure pointers. The third parameter is a pointer to a routine which is called with a pointer to a directory entry and should return a non zero value if the directory entry should be included in the array. If this pointer is **NULL**, then all the directory entries will be included. The last argument is a pointer to a routine which is passed to **qsort**(3C), which sorts the completed array. If this pointer is **NULL**, the array is not sorted. **alphasort( )** is a routine that sorts the array alphabetically.

scandir( ) returns the number of entries in the array and a pointer to the array through the parameter *namelist*.

RETURN VALUES

Returns –1 if the directory cannot be opened for reading or if **malloc**(3C) cannot allocate enough memory to hold all the data structures.

SEE ALSO

getdents(2), **readdir**(3B), **directory**(3C), **malloc**(3C), **qsort**(3C)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

NAME

scanf, fscanf, sscanf – convert formatted input

SYNOPSIS

```
#include <stdio.h>
int scanf(const char *format, ...);
int fscanf(FILE *strm, const char *format, ...);
int sscanf(const char *s, const char *format, ...);
```

MT-LEVEL

MT-Safe

DESCRIPTION

scanf( ) reads from the standard input stream, **stdin**.

fscanf( ) reads from the stream *strm*.

sscanf( ) reads from the character string *s*.

Each function reads characters, interprets them according to a format, and stores the results in its arguments. Each expects, as arguments, a control string, *format*, described below and a set of pointer arguments indicating where the converted input should be stored. If there are insufficient arguments for the format, the behavior is undefined. If the format is exhausted while arguments remain, the excess arguments are simply ignored.

The control string usually contains conversion specifications, which are used to direct interpretation of input sequences. The control string may contain:

1. White-space characters (blanks, tabs, new-lines, or form-feeds) that, except in two cases described below, cause input to be read up to the next non-white-space character.
2. An ordinary character (not **%**) that must match the next character of the input stream.
3. Conversion specifications consisting of the character **%** or the character sequence **%***digits***$**, an optional assignment suppression character *****, a decimal digit string that specifies an optional numerical maximum field width, an optional letter **l** (ell), **L**, or **h** indicating the size of the receiving object, and a conversion code:
   % or *digit***$**, *****, *decimal digit string*, **h** or **l** or **L**, *conversion code*

   The following defines which size indicators can used with which conversion codes, and the size they indicate.

Conversion Code	Size Indicator	Size
d, i, n	none	int
	h	short int
	l	long int
o, u, x	none	unsigned int
	h	unsigned short int
	l	unsigned long int
e, f, g	none	float
	l	double
	L	long double

The **h**, **l**, or **L** modifier is ignored with any other conversion codes.

A conversion specification directs the conversion of the next input field; the result is placed in the variable pointed to by the corresponding argument unless assignment suppression was indicated by the character *. The suppression of assignment provides a way of describing an input field that is to be skipped. An input field is defined as a string of non-space characters; it extends to the next inappropriate character or until the maximum field width, if one is specified, is exhausted. For all descriptors except the character **[** and the character **c**, white space leading an input field is ignored.

Conversions can be applied to the *nth* argument in the argument list, rather than to the next unused argument. In this case, the conversion character **%** (see above) is replaced by the sequence **%***digits***$** where *digits* is a decimal integer *n*, giving the position of the argument in the argument list. The first such argument, **%1$**, immediately follows *format*. The control string can contain either form of a conversion specification, that is, **%** or **%***digits***$**, although the two forms cannot be mixed within a single control string.

The conversion code indicates the interpretation of the input field; the corresponding pointer argument must usually be of a restricted type. For a suppressed field, no pointer argument is given. The following conversion codes are valid:

% A single **%** is expected in the input at this point; no assignment is done.

d Matches an optionally signed decimal integer, whose format is the same as expected for the subject sequence of the **strtol( )** function with the value 10 for the *base* argument. The corresponding argument should be a pointer to integer.

u Matches an optionally signed decimal integer, whose format is the same as expected for the subject sequence of the **strtoul( )** function (see **strtol**(3C)) with the value 10 for the *base* argument. The corresponding argument should be a pointer to unsigned integer.

o Matches an optionally signed octal integer, whose format is the same as expected for the subject sequence of the **strtoul( )** function with the value 8 for the *base* argument. The corresponding argument should be a pointer to unsigned integer.

x Matches an optionally signed hexadecimal integer, whose format is the

same as expected for the subject sequence of the **strtoul( )** function with the value 16 for the *base* argument. The corresponding argument should be a pointer to unsigned integer.

i Matches an optionally signed integer, whose format is the same as expected for the subject sequence of the **strtol( )** function with the value 0 for the *base* argument. The corresponding argument should be a pointer to integer.

n No input is consumed. The corresponding argument should be a pointer to integer into which is to be written the number of characters read from the input stream so far by the call to the function. Execution of a **%n** directive does not increment the assignment count returned at the completion of execution of the function.

e,**f**,**g** Matches an optionally signed floating point number, whose format is the same as expected for the subject string of the **strtod** function. The corresponding argument should be a pointer to floating.

s A character string is expected; the corresponding argument should be a character pointer pointing to an array of characters large enough to accept the string and a terminating **\0**, which will be added automatically. The input field is terminated by a white-space character.

ws A wide character string is expected; the corresponding argument should be a wide character pointer pointing to an array of wide characters large enough to accept the wide character string and a terminating **\0**, which will be added automatically. The input field is terminated by a white-space character.

c Matches a sequence of characters of the number specified by the field width (1 if no field width is present in the directive). The corresponding argument should be a pointer to the initial character of an array large enough to accept the sequence. No null character is added. The normal skip over white space is suppressed.

wc Matches a sequence of wide characters of the number specified by the field width (1 if no field width is present in the directive). The corresponding argument should be a pointer to the initial character of an array large enough to accept the sequence. No null character is added. The normal skip over white space is suppressed.

[Matches a nonempty sequence of characters from a set of expected characters (the *scanset*). The corresponding argument should be a pointer to the initial character of an array large enough to accept the sequence and a terminating null character, which will be added automatically. The conversion specifier includes all subsequent characters in the *format* string, up to and including the matching right bracket (**]**). The characters between the brackets (the *scanlist*) comprise the scanset, unless the character after the left bracket is a circumflex (^), in which case the scanset contains all characters that do not appear in the scanlist between the circumflex and the right bracket. If the conversion specifier begins with **[]** or **[^]**, the right bracket

character is in the scanlist and the next right bracket character is the matching right bracket that ends the specification; otherwise the first right bracket character is the one that ends the specification.

A range of characters in the scanset may be represented by the construct *first* – *last*; thus **[0123456789]** may be expressed **[0–9]**. Using this convention, *first* must be lexically less than or equal to *last*, or else the dash will stand for itself. The character – will also stand for itself whenever it is the first or the last character in the scanlist. To include the right bracket as an element of the scanset, it must appear as the first character (possibly preceded by a circumflex) of the scanlist and in this case it will not be syntactically interpreted as the closing bracket. At least one character must match for this conversion to be considered successful.

p Matches the set of implementation-defined sequences produced as output by the **%p** conversion of the **printf**(3S) function. The corresponding argument should be a pointer to **void**. If the input item is a value converted earlier during the same program execution, the pointer that results compares equal to that value; otherwise, the behavior of the **%p** conversion is undefined.

If an invalid conversion character follows the **%**, the results of the operation may not be predictable.

The conversion specifiers **E**, **G**, and **X** are also valid and, under the **-Xa** and **-Xc** compilation modes (see **cc**(1B)), behave the same as **e**, **g**, and **x**, respectively. Under the **-Xt** compilation mode, **E**, **G**, and **X** behave the same as **le**, **lg**, and **lx**, respectively.

Each function allows for detection of a language-dependent decimal point character in the input string. The decimal point character is defined by the program's locale (category **LC_NUMERIC**). In the "**C**" locale, or in a locale where the decimal point character is not defined, the decimal point character defaults to a period (**.**).

The **scanf( )** conversion terminates at end of file, at the end of the control string, or when an input character conflicts with the control string.

If end-of-file is encountered during input, conversion is terminated. If end-of-file occurs before any characters matching the current directive have been read (other than leading white space, where permitted), execution of the current directive terminates with an input failure; otherwise, unless execution of the current directive is terminated with a matching failure, execution of the following directive (if any) is terminated with an input failure.

If conversion terminates on a conflicting input character, the offending input character is left unread in the input stream. Trailing white space (including new-line characters) is left unread unless matched by a directive. The success of literal matches and suppressed assignments is not directly determinable other than via the **%n** directive.

LC_NUMERIC

Determines how numeric formats are handled. In the "C" locale, numeric

handling follows the U.S. rules.

RETURN VALUES

These routines return the number of successfully matched and assigned input items; this number can be 0 in the event of an early matching failure between an input character and the control string. If the input ends before the first matching failure or conversion, **EOF** is returned.

EXAMPLES

The call to the function **scanf( )**:

```
int i, n; float x; char name[50];
n = scanf ("%d%f%s", &i, &x, name);
```

with the input line:

```
25 54.32E–1 thompson
```

will assign to **n** the value **3**, to **i** the value **25**, to **x** the value **5.432**, and **name** will contain **thompson\0**.

The call to the function **scanf( )**:

```
int i; float x; char name[50];
(void) scanf ("%2d%f%*d %[0–9]", &i, &x, name);
```

with the input line:

```
56789 0123 56a72
```

will assign **56** to **i**, **789.0** to **x**, skip **0123**, and place the characters **56\0** in **name**. The next character read from **stdin** will be **a**.

FILES

/usr/lib/locale/*locale***/LC_NUMERIC/numeric**
LC_NUMERIC database for *locale*

SEE ALSO

cc(1B), **strtod**(3C), **strtol**(3C), **printf**(3S)

NAME

sched_get_priority_max, sched_get_priority_min, sched_rr_get_interval – get scheduling parameter limits

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <sched.h>

int sched_get_priority_max(int policy);

int sched_get_priority_min(int policy);

int sched_rr_get_interval(pid_t pid, struct timespec *interval);

struct timespec {
        time_t  tv_sec;     /* seconds */
        long    tv_nsec;    /* and nanoseconds */
};
```

MT-LEVEL

MT-Safe

DESCRIPTION

sched_get_priority_max( ) and **sched_get_priority_min( )** return the appropriate maximum or minimum values, respectively, for the scheduling policy specified by *policy*.

sched_rr_get_interval( ) updates the **timespec** structure referenced by *interval* to contain the current execution time limit (i.e., time quantum) for the process specified by *pid* under the **SCHED_RR** policy. After that time limit expires, when another process at the same priority is ready to execute, a scheduling decision will be made. If *pid* is zero, the current execution time limit for the calling process is stored in *interval*.

The value of *policy* must be one of the scheduling policy values defined in **<sched.h>**: **SCHED_FIFO, SCHED_RR,** or **SCHED_OTHER.**

RETURN VALUES

If successful, **sched_get_priority_max( )** or **sched_get_priority_min( )** returns the appropriate maximum or minimum values, respectively.

If successful, **sched_rr_get_interval( )** returns **0**.

If unsuccessful, these functions return **-1**, and set **errno** to indicate the error condition.

ERRORS

EINVAL	The value of *policy* does not represent a defined scheduling policy.
ENOSYS	**sched_get_priority_max( ), sched_get_priority_min( ),** and **sched_rr_get_interval( )** are not supported by this implementation.
ESRCH	No process can be found corresponding to that specified by *pid*.

SEE ALSO

sched_setparam(3R), **sched_setscheduler**(3R)

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Priority Scheduling option. It is our intention to provide support for these interfaces in future releases.

NAME

sched_setparam, sched_getparam – set / get scheduling parameters

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

#include <sched.h>

int sched_setparam(pid_t *pid***, const struct sched_param** ∗*param***);**

int sched_getparam(pid_t *pid***, struct sched_param** ∗*param***);**

```
struct sched_param {
        int             sched_priority;       /* process execution scheduling priority */
 ...
}
```

MT-LEVEL

MT-Safe

DESCRIPTION

sched_setparam( ) sets the scheduling parameters of the process specified by *pid* to the values specified by the **sched_param** structure referenced by *param.*

sched_getparam( ) stores the scheduling parameters of a process, specified by *pid,* in the **sched_param** structure pointed to by *param.*

If the target process has as its scheduling policy, **SCHED_FIFO** or **SCHED_RR**:

If *pid* is zero, the scheduling parameters are set / stored for the calling process. Otherwise, if a process specified by *pid* exists and if the calling process has permission, the scheduling parameters are set / stored for the process whose process ID is equal to *pid.* The real or effective user ID of the calling process must match the real or saved (from **exec**(2)) user ID of the target process unless the effective user ID of the calling process is **0**. See **intro**(2).

The target process, *pid,* whether it is running or not running, resumes execution after all other runnable processes of equal or greater priority have been scheduled to run.

If the priority of the process, *pid,* is set higher than that of the lowest priority running process, and if process *pid* is ready to run, then process *pid* preempts a lowest priority running process. Similarly, if the process calling **sched_setparam( )** sets its own priority lower than that of one or more other non-empty process lists, then the process that is the head of the highest priority list preempts the calling process. Thus, in either case, the originating process might not receive notification of the completion of the requested priority change until the higher priority process has executed.

The value of *param*->**sched_priority** must be an integer within the inclusive priority range for the current scheduling policy of the process specified by *pid.* Higher numerical values for the priority represent higher priorities.

RETURN VALUES

If successful, **sched_setparam( )** and **sched_getparam( )** returns **0;** otherwise, the priority remains unchanged, the function returns **-1,** and sets **errno** to indicate the error condition.

ERRORS

EINVAL	One or more of **sched_setparam( )**'s requested scheduling parameters is outside the range defined for the specified *pid*'s scheduling policy.
ENOSYS	**sched_setparam( )** and **sched_getparam( )** are not supported by this implementation.
EPERM	The requesting process does not have permission to set / get the scheduling parameters for the specified process, or does not have the appropriate privilege to invoke **sched_setparam( ).**
ESRCH	No process can be found corresponding to that specified by *pid.*

SEE ALSO

intro(2), **exec**(2), **sched_setscheduler**(3R)

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Priority Scheduling option. It is our intention to provide support for these interfaces in future releases.

NAME

sched_setscheduler, sched_getscheduler – set / get scheduling policy and scheduling parameters

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

#include <sched.h>

int sched_setscheduler(pid_t *pid***, int** *policy***, const struct sched_param** ∗*param***);**

int sched_getscheduler(pid_t *pid***);**

```
struct sched_param {
        int             sched_priority;      /* process execution scheduling priority */
        ...
}
```

MT-LEVEL

MT-Safe

DESCRIPTION

sched_setscheduler( ) sets the scheduling policy and scheduling parameters of the process specified by *pid* to *policy* and the parameters specified in the **sched_param** structure pointed to by *param*, respectively. The value of *param*->**sched_priority** must be any integer within the inclusive priority range for the scheduling policy specified by *policy*.

The possible values for the *policy* parameter are defined in the header file **<sched.h>**: **SCHED_FIFO, SCHED_RR,** or **SCHED_OTHER.**

If *pid* is zero, the scheduling policy and scheduling parameters are set for the calling process. Otherwise, if a process specified by *pid* exists and if the calling process has permission, the scheduling policy and scheduling parameters are set for the process whose process ID is equal to *pid*. The real or effective user ID of the calling process must match the real or saved (from **exec**(2)) user ID of the target process unless the effective user ID of the calling process is super-user. See **intro**(2).

To change the *policy* of any process to either of the real time policies **SCHED_FIFO** or **SCHED_RR,** the calling process must either have the **SCHED_FIFO,** or **SCHED_RR** policy or have an effective user ID of **0**.

sched_getscheduler( ) returns the scheduling policy of the process specified by *pid*. If *pid* is zero, the scheduling policy is returned for the calling process. Otherwise, if a process specified by *pid* exists and if the calling process has permission, the scheduling policy is returned for the process whose process ID is equal to *pid*.

RETURN VALUES

If successful, **sched_setscheduler( )** returns the former scheduling policy of the specified process *(pid)*, which will be one of the following values:

SCHED_FIFO (realtime),
First-In-First-Out; processes scheduled to this policy, if not pre-empted by a higher priority or interrupted by a signal, will proceed until completion.

SCHED_RR (realtime),
Round-Robin; processes scheduled to this policy, if not pre-empted by a higher priority or interrupted by a signal, will execute for a time period, returned by **sched_rr_get_interval**(3R) or by the system.

or

SCHED_OTHER (time-sharing).

Otherwise, the policy and sheduling parameters remain unchanged, **sched_setscheduler( )** returns **-1**, and sets **errno** to indicate the error condition.

If successful, **sched_getscheduler( )** returns the scheduling policy of the specified processr; otherwise, it returns **-1**, and sets **errno** to indicate the error condition.

ERRORS

EINVAL The value of *policy* is invalid, or one or more of the parameters contained in **param** is outside the valid range for the specified scheduling policy.

ENOSYS **sched_setscheduler( )** and **sched_getscheduler( )** are not supported by this implementation.

EPERM **sched_setscheduler( )** does not have permission to set either or both of the scheduling parameters or the scheduling policy of the specified process.

sched_getscheduler( ) does not have permission to determine the scheduling policy of the specified process.

ESRCH No process can be found corresponding to that specified by *pid.*

SEE ALSO

priocntl(1), **intro**(2), **exec**(2), **priocntl**(2), **sched_get_priority_max**(3R), **sched_setparam**(3R)

BUGS

In Solaris 2.5, these functions always return –**1** and set **errno** to **ENOSYS,** because this release does not support the Priority Scheduling option. It is our intention to provide support for these interfaces in future releases.

NAME

sched_yield – yield processor

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <sched.h>
int sched_yield(void);
```

MT-LEVEL

MT-Safe

DESCRIPTION

sched_yield( ) forces the running process to relinquish the processor until the process again becomes the head of its process list.

RETURN VALUES

If successful, **sched_yield( )** returns **0,** otherwise, it returns **-1,** and sets **errno** to indicate the error condition.

ERRORS

ENOSYS sched_yield( ) is not supported by this implementation.

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Priority Scheduling option. It is our intention to provide support for these interfaces in future releases.

NAME

secure_rpc, authdes_getucred, authdes_seccreate, getnetname, host2netname, key_decryptsession, key_encryptsession, key_gendes, key_setsecret, key_secretkey_is_set, netname2host, netname2user, user2netname – library routines for secure remote procedure calls

MT-LEVEL

MT-Safe

DESCRIPTION

RPC library routines allow C programs to make procedure calls on other machines across the network.

RPC supports various authentication flavors. Among them are:

AUTH_NONE	(none) no authentication.
AUTH_SYS	Traditional UNIX-style authentication.
AUTH_DES	DES encryption-based authentication.
AUTH_KERB	Kerberos encryption-based authentication.

The **authdes_getucred( )** and **authdes_seccreate( )** routines implement the **AUTH_DES** authentication flavor. The keyserver daemon **keyserv** (see **keyserv**(1M)) must be running for the **AUTH_DES** authentication system to work, and **keylogin**(1) must have been run. Only the **AUTH_DES** style of authentication is discussed here. For information about the **AUTH_NONE** and **AUTH_SYS** styles of authentication, refer to **rpc_clnt_auth**(3N). For information about the **AUTH_KERB** style of authentication, refer to **kerberos_rpc**(3N).

The routines documented on this page are MT-Safe. See the pages of the other authentication styles for their MT-level.

Routines

See **rpc**(3N) for the definition of the **AUTH** data structure.

#include <rpc/rpc.h>
#include <sys/types.h>

int authdes_getucred(const struct authdes_cred **adc***, uid_t** **uidp***, gid_t** **gidp***, short** **gidlenp***, gid_t** **gidlist***);**

authdes_getucred( ) is the first of the two routines which interface to the RPC secure authentication system known as **AUTH_DES**. The second is **authdes_seccreate( )**, below. **authdes_getucred( )** is used on the server side for converting an **AUTH_DES** credential, which is operating system independent, into an **AUTH_SYS** credential. This routine returns **1** if it succeeds, **0** if it fails.

uidp* is set to the user's numerical ID associated with *adc*. **gidp* is set to the numerical ID of the user's group. **gidlist* contains the numerical IDs of the other groups to which the user belongs. **gidlenp* is set to the number of valid group ID entries in **gidlist* (see **netname2user( ), below).

Warning: **authdes_getucred( )** will fail if the authdes_cred structure was created with the netname of a host. In such a case, **netname2host( )** should be used on the host netname in the authdes_cred structure to get the host name.

AUTH ∗**authdes_seccreate(const char** ∗*name***, const unsigned int** *window***, const char** ∗*timehost***, const des_block** ∗*ckey***);**

authdes_seccreate( ), the second of two **AUTH_DES** authentication routines, is used on the client side to return an authentication handle that will enable the use of the secure authentication system. The first parameter *name* is the network name, or *netname*, of the owner of the server process. This field usually represents a hostname derived from the utility routine **host2netname( )**, but could also represent a user name using **user2netname( )**, described below.

The second field is *window* on the validity of the client credential, given in seconds. If the difference in time between the client's clock and the server's clock exceeds *window*, the server will reject the client's credentials, and the clock will have to be resynchronized. A small window is more secure than a large one, but choosing too small of a window will increase the frequency of resynchronizations because of clock drift.

The third parameter, *timehost*, the host's name, is optional. If it is **NULL**, then the authentication system will assume that the local clock is always in sync with the *timehost* clock, and will not attempt resynchronizations. If a timehost is supplied, however, then the system will consult with the remote time service whenever resynchronization is required. This parameter is usually the name of the host on which the server is running.

The final parameter *ckey* is also optional. If it is **NULL,** then the authentication system will generate a random DES key to be used for the encryption of credentials. If *ckey* is supplied, then it will be used instead.

If **authdes_seccreate( )** fails, it returns **NULL.**

int getnetname(char *name***[MAXNETNAMELEN+1]);**

getnetname( ) returns the unique, operating system independent netname of the caller in the fixed-length array *name*. Returns **1** if it succeeds, and **0** if it fails.

int host2netname(char *name***[MAXNETNAMELEN+1], const char** ∗*host***, const char** ∗*domain***);**

Convert from a domain-specific hostname *host* to an operating system independent netname. Returns **1** if it succeeds, and **0** if it fails. Inverse of **netname2host( )**. If *domain* is **NULL**, **host2netname( )** uses the default domain name of the machine. If *host* is **NULL**, it defaults to that machine itself. If *domain* is **NULL** and *host* is a NIS name like "host1.ssi.sun.com," **host2netname( )** uses the domain "ssi.sun.com" rather than the default domain name of the machine.

int key_decryptsession(const char **remotename*, **des_block** **deskey*);

key_decryptsession( ) is an interface to the keyserver daemon, which is associated with RPC's secure authentication system (**AUTH_DES** authentication). User programs rarely need to call it, or its associated routines **key_encryptsession( )**, **key_gendes( )**, and **key_setsecret( )**.

key_decryptsession( ) takes a server netname *remotename* and a DES key *deskey*, and decrypts the key by using the the public key of the the server and the secret key associated with the effective UID of the calling process. It is the inverse of **key_encryptsession( )**.

int key_encryptsession(const char **remotename*, **des_block** **deskey*);

key_encryptsession( ) is a keyserver interface routine. It takes a server netname *remotename* and a DES key *deskey*, and encrypts it using the public key of the the server and the secret key associated with the effective UID of the calling process. It is the inverse of **key_decryptsession( )**. This routine returns **0** if it succeeds, **–1** if it fails.

int key_gendes(des_block **deskey*);

key_gendes( ) is a keyserver interface routine. It is used to ask the keyserver for a secure conversation key. Choosing one at random is usually not good enough, because the common ways of choosing random numbers, such as using the current time, are very easy to guess. This routine returns **0** if it succeeds, **–1** if it fails.

int key_setsecret(const char **key*);

key_setsecret( ) is a keyserver interface routine. It is used to set the key for the effective UID of the calling process. This routine returns **0** if it succeeds, **–1** if it fails.

int key_secretkey_is_set(void);

key_secretkey_is_set( ) is a keyserver interface routine that may be used to determine whether a key has been set for the effective UID of the calling process. If the keyserver has a key stored for the effective UID of the calling process, this routine returns **1**. Otherwise it returns **0**.

int netname2host(const char **name*, **char** **host*, **const int** *hostlen*);

Convert from an operating system independent netname *name* to a domain-specific hostname *host*. *hostlen* is the maximum size of *host*. Returns **1** if it succeeds, and **0** if it fails. Inverse of **host2netname( )**.

int netname2user(const char **name*, **uid_t** **uidp*, **gid_t** **gidp*,
int **gidlenp*, **gid_t** *gidlist***[NGROUPS]);**

Convert from an operating system independent netname to a domain-specific user ID. Returns **1** if it succeeds, and **0** if it fails. Inverse of **user2netname( )**.

**uidp* is set to the user's numerical ID associated with *name*. **gidp* is set to the numerical ID of the user's group. *gidlist* contains the numerical IDs of the other groups to which the user belongs. **gidlenp* is set to the number of valid group ID entries in *gidlist*.

int user2netname(char *name***[MAXNETNAMELEN+1], const uid_t** *uid***,**
const char **domain***);**

Convert from a domain-specific username to an operating system independent netname. Returns **1** if it succeeds, and **0** if it fails. Inverse of **netname2user( )**.

SEE ALSO

chkey(1), **keyserv**(1M), **newkey**(1M), **kerberos_rpc**(3N), **rpc**(3N), **rpc_clnt_auth**(3N)

NAME

select – synchronous I/O multiplexing

SYNOPSIS

```
#include <sys/time.h>
#include <sys/types.h>

int select(int nfds, fd_set *readfds, fd_set *writefds, fd_set *exceptfds,
        struct timeval *timeout);

void FD_SET(int fd, fd_set &fdset);

void FD_CLR(int fd, fd_set &fdset);

int FD_ISSET(int fd, fd_set &fdset);

void FD_ZERO(fd_set &fdset);
```

MT-LEVEL

MT-Safe

DESCRIPTION

select( ) examines the I/O file descriptor sets whose addresses are passed in *readfds*, *writefds*, and *exceptfds* to see if any of their file descriptors are ready for reading, are ready for writing, or have an exceptional condition pending, respectively. Out-of-band data is the only exceptional condition. *nfds* is the number of bits to be checked in each bit mask that represents a file descriptor; the file descriptors from **0** to *nfds* **–1** in the file descriptor sets are examined. On return, **select( )** replaces the given file descriptor sets with subsets consisting of those file descriptors that are ready for the requested operation. The return value from the call to **select( )** is the number of ready file descriptors.

The file descriptor sets are stored as bit fields in arrays of integers. The following macros are provided for manipulating such file descriptor sets: **FD_ZERO( )** initializes a file descriptor set *fdset* to the null set. **FD_SET( )** includes a particular file descriptor *fd* in *fdset*. **FD_CLR( )** removes *fd* from *fdset*. **FD_ISSET( )** is nonzero if *fd* is a member of *fdset*, zero otherwise. The behavior of these macros is undefined if a file descriptor value is less than zero or greater than or equal to **FD_SETSIZE**. **FD_SETSIZE** is a constant defined in **<sys/select.h>**.

If *timeout* is not a NULL pointer, it specifies a maximum interval to wait for the selection to complete. If *timeout* is a NULL pointer, the **select( )** blocks indefinitely. To effect a poll, the *timeout* argument should be a non-NULL pointer, pointing to a zero-valued **timeval** structure.

Any of *readfds*, *writefds*, and *exceptfds* may be given as NULL pointers if no file descriptors are of interest.

RETURN VALUES

select( ) returns the number of ready file descriptors contained in the file descriptor sets or **–1** if an error occurred. If the time limit expires, then **select( )** returns **0**.

ERRORS

The call fails if:

EBADF One of the I/O file descriptor sets specified an invalid I/O file descriptor.

EINTR A signal was delivered before any of the selected events occurred, or the time limit expired.

EINVAL A component of the pointed-to time limit is outside the acceptable range: **t_sec** must be between **0** and 10^8, inclusive. **t_usec** must be greater than or equal to **0**, and less than 10^6.

SEE ALSO

poll(2), **read**(2), **write**(2)

NOTES

The default value for **FD_SETSIZE** (currently 1024) is larger than the default limit on the number of open files. In order to accommodate programs that may use a larger number of open files with **select( )**, it is possible to increase this size within a program by providing a larger definition of **FD_SETSIZE** before the inclusion of **<sys/types.h>**.

The file descriptor sets are always modified on return, even if the call returns as the result of a timeout.

NAME

sem_close – close a named semaphore

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <semaphore.h>
int sem_close(sem_t *sem);
typedef struct {
        ...
} sem_t;                    /*opaque POSIX.4 semaphore*/
```

MT-LEVEL

MT-Safe

DESCRIPTION

sem_close( ) is used to indicate that the calling process is finished using the named semaphore *sem*. **sem_close( )** deallocates any system resources for use by this process for this semaphore. If the semaphore has not been removed with a successful call to **sem_unlink**(3R), then **sem_close( )** has no effect on the state of the semaphore. If **sem_unlink**(3R) has been successfully invoked for *name* after the most recent call to **sem_open**(3R) with **O_CREAT** for this semaphore, then when all processes that have opened the semaphore close it, the semaphore will no longer be accessible.

sem_close( ) should not be called for an unnamed semaphore initialized by **sem_init**(3R).

RETURN VALUES

If successful, **sem_close( )** returns **0,** otherwise it returns **-1** and sets **errno** to indicate the error condition.

ERRORS

EINVAL *sem* is not a valid semaphore descriptor.

ENOSYS **sem_close( )** is not supported by this implementation.

SEE ALSO

sem_init(3R), **sem_open**(3R), **sem_unlink**(3R)

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Semaphores option. It is our intention to provide support for these interfaces in future releases.

NAME

sem_destroy – destroy an unnamed semaphore

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <semaphore.h>
int sem_destroy(sem_t *sem);
typedef struct {
        ...
} sem_t; /*opaque POSIX.4 semaphore*/
```

MT-LEVEL

MT-Safe

DESCRIPTION

sem_destroy( ) is used to destroy the unnamed semaphore, *sem*, which was initialized by **sem_init**(3R).

RETURN VALUES

If successful, **sem_destroy( )** returns **0**, otherwise it returns **-1** and sets **errno** to indicate the error condition.

ERRORS

EINVAL	*sem* is not a valid semaphore.
ENOSYS	**sem_destroy( )** is not supported by this implementation.
EBUSY	Other processes (or LWPs or threads) are currently blocked on the semaphore.

SEE ALSO

sem_init(3R), **sem_open**(3R)

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Semaphores option. It is our intention to provide support for these interfaces in future releases.

NAME

sem_getvalue – get the value of a semaphore

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <semaphore.h>
int sem_getvalue(sem_t *sem, int *sval);
typedef struct {
 ...
} sem_t; /*opaque POSIX.4 semaphore*/
```

MT-LEVEL

MT-Safe

DESCRIPTION

sem_getvalue( ) updates the location referenced by *sval* to have the value of the semaphore referenced by *sem* without affecting the state of the semaphore. The updated value represents an actual semaphore value that occurred at some unspecified time during the call to **sem_getvalue( ),** but may not be the actual value of the semaphore when **sem_getvalue( )** is returned to the caller.

The value set in *sval* may be zero or positive. If *sval* is zero, there may be other processes (or LWPs or threads) waiting for the semaphore; if *sval* is positive, no one is waiting.

RETURN VALUES

If successful, **sem_getvalue( )** returns **0,** otherwise, it returns **-1,** and sets **errno** to indicate the error condition.

ERRORS

EINVAL *sem* does not refer to a valid semaphore.

ENOSYS **sem_getvalue( )** is not supported by this implementation.

SEE ALSO

sem_post(3R), **sem_wait**(3R)

BUGS

In Solaris 2.5, these functions always return –**1** and set **errno** to **ENOSYS,** because this release does not support the Semaphores option. It is our intention to provide support for these interfaces in future releases.

NAME

sem_init – initialize an unnamed semaphore

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <semaphore.h>
int sem_init(sem_t *sem, int pshared, unsigned int value );
typedef struct {
 ...
} sem_t; /*opaque POSIX.4 semaphore*/
```

MT-LEVEL

MT-Safe

DESCRIPTION

sem_init( ) is used to initialize the unnamed semaphore, referred to by *sem*, to *value*. This semaphore may be used in subsequent calls to **sem_wait**(3R), **sem_trywait**(3R), **sem_post**(3R), and **sem_destroy**(3R). This semaphore remains usable until the semaphore is destroyed.

If *pshared* is non-zero, then the semaphore is sharable between processes. If the semaphore is not being shared between processes, the application should set *pshared* to **0**.

RETURN VALUES

If successful, **sem_init( )** returns **0** and initializes the semaphore in *sem;* otherwise it returns **-1** and sets **errno** to indicate the error condition.

ERRORS

EINVAL	*value* exceeds **SEM_VALUE_MAX**.
ENOSPC	A resource required to initialize the semaphore has been exhausted.
	The resources have reached the limit on semaphores, **SEM_NSEMS_MAX**.
ENOSYS	**sem_init( )** is not supported by this implementation.
EPERM	The calling process lacks the appropriate privileges to initialize the semaphore.

SEE ALSO

sem_destroy(3R), **sem_post**(3R), **sem_wait**(3R)

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Semaphores option. It is our intention to provide support for these interfaces in future releases.

NAME

sem_open – initialize / open a named semaphore

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <semaphore.h>

sem_t *sem_open(const char *name, int oflag,
        /* unsigned long mode, unsigned int value */ ... );

typedef struct {
 ...
} sem_t;                    /*opaque POSIX.4 semaphore*/
```

MT-LEVEL

MT-Safe

DESCRIPTION

sem_open( ) establishes a connection to a semaphore, *name,* returning the address of the semaphore to the calling process (or LWP or thread) for subsequent calls to **sem_wait**(3R), **sem_trywait**(3R), **sem_post**(3R), and **sem_close**(3R). The semaphore remains usable by this process until the semaphore is closed.

name points to a string naming a semaphore object. The *name* argument should conform to the construction rules for a pathname. If a process makes multiple successful calls to **sem_open( )** with the same value for *name,* the same semaphore address will be returned for each such successful call, provided that there have been no calls to **sem_unlink**(3R) for this semaphore.

oflag determines whether the semaphore is created or merely accessed by the call to **sem_open( ).** The three valid values for *oflag* are **0, O_CREAT**, or the bitwise inclusive OR of **O_CREAT** and **O_EXCL**. Setting the *oflag* bits to **O_CREAT** will create the semaphore if it does not already exist. Setting both **O_CREAT** and **O_EXCL** will fail if the semaphore already exists. The check for the existence of the semaphore and the creation of the semaphore if it does not exist is atomic with respect to other processes executing **sem_open( ).** After the semaphore named *name* has been created by **sem_open( )** with the **O_CREAT** flag, other processes can connect to this semaphore by calling **sem_open( )** with the same value of *name,* and nobits set in *oflag.*

Using the **O_CREAT** flag requires a third and a fourth argument: *mode* and *value.* The semaphore is created with an initial count of *value. value* must be less than or equal to **{SEM_VALUE_MAX}.** The semaphore's user ID acquires the effective user ID of the process; the semaphore's group ID is set to a system default group ID or to the effective group ID of the process. The semaphore's permission bits is set to the value of *mode*, modified by clearing all bits set in the file creation mask of the process (see **umask**(2)).

RETURN VALUES

If successful, **sem_open( )** returns the address of the semaphore, otherwise it returns **-1** and sets **errno** to indicate the error condition.

ERRORS

EACCES The named semaphore exists and the **O_RDWR** permissions are denied, or the named semaphore does not exist and permission to create the named semaphore is denied.

EEXIST **O_CREAT** and **O_EXCL** are set and the named semaphore already exists.

EINTR **sem_open( )** was interrupted by a signal.

EINVAL *name* is not a valid name.

O_CREAT was set in *oflag* and *value* is greater than **{SEM_VALUE_MAX}**.

EMFILE The number of open semaphore descriptors in this process exceeds **{SEM_NSEMS_MAX}**.

The number of open file descriptors in this process exceeds **{OPEN_MAX}**.

ENAMETOOLONG
The string-length of *name* exceeds **{PATH_MAX}**, or a pathname component is longer than **{NAME_MAX}** while **_POSIX_NO_TRUNC** is in effect.

ENFILE The system file table is full.

ENOENT **O_CREAT** is not set and the named semaphore does not exist.

ENOSPC There is insufficient space for the creation of the new named semaphore.

ENOSYS **sem_open( )** is not supported by this implementation.

SEE ALSO

exec(2), **exit**(2), **umask**(2), **sysconf**(3C), **sem_close**(3R), **sem_post**(3R), **sem_unlink**(3R), **sem_wait**(3R)

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Semaphores option. It is our intention to provide support for these interfaces in future releases.

NAME

sem_post – increment the count of a semaphore

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

#include <semaphore.h>

int sem_post(sem_t ∗*sem***);**

typedef struct {
...
} sem_t /∗opaque POSIX.4 semaphore∗/

MT-LEVEL

Async-Signal-Safe

DESCRIPTION

If, prior to the call to **sem_post( ),** the value of *sem* was **0,** and other processes (or LWPs or threads) were blocked waiting for the semaphore, then one of them will be allowed to return successfully from its call to **sem_wait**(3R). The process to be unblocked will be chosen in a manner appropriate to the scheduling policies and parameters in effect for the blocked processes. In the case of the policies **SCHED_FIFO** and **SCHED_RR**, the highest priority waiting process is unblocked, and if there is more than one highest-priority process blocked waiting for the semaphore, then the highest priority process which has been waiting the longest is unblocked.

If, prior to the call to **sem_post( ),** no other processes (or LWPs or thread) were blocked for the semaphore, then its value is incremented by one.

sem_post( ) is reentrant with respect to signals (ASYNC-SAFE), and may be invoked from a signal-catching function. The semaphore functionality described on this man page is for the POSIX threads implementation. For the documentation of the Solaris threads interface, see **semaphore**(3T)).

RETURN VALUES

If successful, **sem_post( )** returns **0,** otherwise it returns **-1,** and sets **errno** to indicate the error condition.

ERRORS

EINVAL *sem* does not refer to a valid semaphore.

ENOSYS **sem_post( )** is not supported by this implementation.

EXAMPLES

(see **sem_wait**(3R))

SEE ALSO

sched_setscheduler(3R), **sem_wait**(3R), **semaphore**(3T)

NOTES

sem_wait(3R) and **sem_trywait**(3R) decrement the semaphore upon their successful return.

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Semaphores option. It is our intention to provide support for these interfaces in future releases.

NAME

sem_unlink – remove a named semaphore

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

#include <semaphore.h>

int sem_unlink(const char ∗*name***);**

MT-LEVEL

MT-Safe

DESCRIPTION

sem_unlink( ) removes the semaphore named by the string *name.* If the semaphore, *name,* is currently referenced by other processes, then **sem_unlink( )** has no effect on the state of the semaphore. If one or more processes have the semaphore open when **sem_unlink( )** is called, destruction of the semaphore is postponed until all references to the semaphore have been destroyed by calls to **sem_close**(3R), **exit**(2), or **exec**(2). Calls to **sem_open**(3R) to re-create or re-connect to the semaphore will refer to a new semaphore after **sem_unlink( )** is called. **sem_unlink( )** does not block until all references have been destroyed; rather, it returns immediately.

RETURN VALUES

If successful, **sem_unlink( )** returns **0;** otherwise, the function returns **-1,** sets **errno** to indicate the error condition, and the semaphore is left unchanged.

ERRORS

EACCES Permission is denied to unlink the named semaphore.

ENAMETOOLONG
The string-length of *name* exceeds **{PATH_MAX}**, or a pathname component is longer than **{NAME_MAX}** while **_POSIX_NO_TRUNC** is in effect.

ENOENT The named semaphore does not exist.

ENOSYS **sem_unlink( )** is not supported by this implementation.

SEE ALSO

exec(2), **exit**(2), **sem_close**(3R), **sem_open**(3R)

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Semaphores option. It is our intention to provide support for these interfaces in future releases.

NAME

sem_wait, sem_trywait – acquire or wait for a semaphore

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

#include <semaphore.h>

int sem_wait(sem_t ∗*sem***);**

int sem_trywait(sem_t ∗*sem***);**

typedef struct {

...

} sem_t; /∗opaque POSIX.4 semaphore∗/

MT-LEVEL

MT-Safe

DESCRIPTION

sem_wait( ) and **sem_trywait( )** are the functions by which a calling thread waits or proceeds depending upon the state of a semaphore. A synchronizing process can proceed only if the value of the semaphore it accesses is currently greater than **0.**

If at the time of a call to either **sem_wait( )** or **sem_trywait( ),** the value of *sem* is positive, these functions decrement the value of the semaphore, return immediately, and allow the calling process to continue.

If the semaphore's value is **0:**

sem_wait( ) blocks, awaiting the semaphore to be released by another process (or LWP or thread).

sem_trywait( ) fails, returning immediately.

RETURN VALUES

If at the time of a call to either **sem_wait( )** or **sem_trywait( ),** the value of *sem* is positive, these functions return **0** on success. If the call was unsuccessful, the state of the semaphore is unchanged, the calling function returns **-1,** and sets **errno** to indicate the error condition.

ERRORS

EAGAIN	The value of *sem* was **0** when **sem_trywait( )** was called.
EINVAL	*sem* does not refer to a valid semaphore.
EINTR	**sem_wait( )** was interrupted by a signal.
ENOSYS	**sem_wait( )** and **sem_trywait( )** are not supported by this implementation.
EDEADLK	A deadlock condition was detected; i.e., two separate processes are waiting for an available resource to be released via a semaphore "held" by the other process.

EXAMPLES

The customer waiting-line in a bank may be analogous to the synchronization scheme of a semaphore utilizing **sem_wait( )** and **sem_trywait( ):**

```
/* cc [ flag ... ] file ... –lposix4 –lthread [ library ... ] */
#include <errno.h>
#define TELLERS 10
sem_t bank_line;        /* semaphore */
int banking_hours(), deposit_withdrawal;
void *customer(), do_business(), skip_banking_today();
thread_t tid;
...

sem_init(&bank_line,TRUE,TELLERS);/* 10 tellers available */
while(banking_hours())
        thr_create(NULL, NULL, customer, (void *)deposit_withdrawal,
                THREAD_NEW_LWP, &tid);
...

void *
customer(deposit_withdrawal)
void *deposit_withdrawal;
{
        int this_customer, in_a_hurry = 50;
        this_customer = rand() % 100;
        if (this_customer == in_a_hurry) {
                if (sem_trywait(&bank_line) != 0)
                 if (errno == EAGAIN) { /* no teller available */
                                skip_banking_today(this_customer);
                                return;
                 } /*else go immediately to available teller & decrement bank_line*/
        }
        else
                sem_wait(&bank_line); /* wait for next teller, then proceed,
                                         and decrement bank_line */


        do_business((int *)deposit_withdrawal);
        sem_post(&bank_line); /* increment bank_line;
                                 this_customer's teller
                                 is now available */
}
```

SEE ALSO

sem_post(3R)

NOTES

sem_wait( ) can be interrupted by a signal, which may result in its premature return.

sem_post(3R) increments the semaphore upon its successful return.

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Semaphores option. It is our intention to provide support for these interfaces in future releases.

NAME

semaphore, sema_init, sema_destroy, sema_wait, sema_trywait, sema_post – semaphores

SYNOPSIS

cc [*flag* ...] *file* ... **–lthread –lc** [*library* ...]

#include <synch.h>

int sema_init(sema_t **sp***, unsigned int** *count***, int** *type***, void** ** arg***);**

int sema_destroy(sema_t **sp***);**

int sema_wait(sema_t **sp***);**

int sema_trywait(sema_t **sp***);**

int sema_post(sema_t **sp***);**

MT-LEVEL

MT-Safe

sema_post( ) is Async-Signal-Safe

DESCRIPTION

A semaphore is a non-negative integer count and is generally used to coordinate access to resources. The initial semaphore count is set to the number of free resources, then threads slowly increment and decrement the count as resources are added and removed. If the semaphore count drops to zero, which means no available resources, threads attempting to decrement the semaphore will block until the count is greater than zero.

Semaphores can synchronize threads in this process and other processes if they are allocated in writable memory and shared among the cooperating processes (see **mmap**(2)), and have been initialized for this purpose.

Semaphores must be initialized before use; semaphores pointed to by *sp* to *count* are initialized by **sema_init( )**. *type* can assign several different types of behavior to a semaphore. No current type uses *arg* although it may be used in the future.

type may be one of the following:

USYNC_PROCESS The semaphore can synchronize threads in this process and other processes. Initializing the semaphore should be done by only one process. *arg* is ignored.

USYNC_THREAD The semaphore can synchronize threads only in this process. *arg* is ignored.

A semaphore must not be simultaneously initialized by multiple threads, nor re-initialized while in use by other threads.

Default semaphore initialization (intra-process):

```
sema_t sp;

sema_init(&sp, NULL, NULL);
 OR
sema_init(&sp, USYNC_THREAD, NULL);
 OR
```

sema_t sp = DEFAULTSEMA;

Customized semaphore initialization (inter-process):

sema_init(&sp, USYNC_PROCESS, NULL);

sema_destroy( ) destroys any state related to the semaphore pointed to by *sp*. The semaphore storage space is not released.

sema_wait( ) blocks the calling thread until the semaphore count pointed to by *sp* is greater than zero, and then it atomically decrements the count.

sema_trywait( ) atomically decrements the semaphore count pointed to by *sp*, if the count is greater than zero; otherwise, it returns an error.

sema_post( ) atomically increments the semaphore count pointed to by *sp*. If there are any threads blocked on the semaphore, one will be unblocked.

The semaphore functionality described on this man page is for the Solaris threads implementation. For the POSIX-compliant semaphore interface documentation, see **sem_open**(3R), **sem_init**(3R), **sem_wait**(3R), **sem_post**(3R), **sem_getvalue**(3R), **sem_unlink**(3R), **sem_close**(3R), **sem_destroy**(3R)).

RETURN VALUES

Upon successful completion, **0** is returned; otherwise, a non-zero value indicates an error.

ERRORS

These functions fail and return the corresponding value if any of the following conditions are detected:

EINVAL Invalid argument.

EFAULT *sp* or *arg* points to an illegal address.

sema_wait( ) fails and returns the corresponding value if any of the following conditions are detected:

EINTR The wait was interrupted by a signal or **fork( )**.

sema_trywait( ) fails and returns the corresponding value if any of the following conditions are detected:

EBUSY The semaphore pointed to by *sp* has a zero count.

SEE ALSO

mmap(2), **sem_open**(3R), **sem_init**(3R), **sem_wait**(3R), **sem_post**(3R), **sem_getvalue**(3R), **sem_unlink**(3R), **sem_close**(3R), **sem_destroy**(3R)

EXAMPLES

The customer waiting-line in a bank is analogous to the synchronization scheme of a semaphore using **sema_wait( )** and **sema_trywait( ):**

```
/* cc [ flag ... ] file ... –lthread [ library ... ] */

#include <errno.h>
#define TELLERS 10
sema_t      tellers;    /* semaphore */
int   banking_hours(), deposit_withdrawal;
void  *customer(), do_business(), skip_banking_today();
```

```
...

sema_init(&tellers, TELLERS, USYNC_THREAD, NULL);
        /* 10 tellers available */
while(banking_hours())
        pthread_create(NULL, NULL, customer, deposit_withdrawal);
...

void *
customer(int deposit_withdrawal)
{
        int this_customer, in_a_hurry = 50;
        this_customer = rand() % 100;

        if (this_customer == in_a_hurry)  {
                if (sema_trywait(&tellers) != 0)
                    if (errno == EAGAIN) { /* no teller available */
                                skip_banking_today(this_customer);
                                return;
                    } /* else go immediately to available teller & decrement tellers */
        }
        else
                sema_wait(&tellers); /* wait for next teller, then proceed,
                                        and decrement tellers */

        do_business(deposit_withdrawal);
        sema_post(&tellers); /* increment tellers;
                        this_customer's teller
                        is now available */
}
```

NOTES

These interfaces are also available via:

#include <thread.h>

If multiple threads are waiting for a semaphore, by default, there is no defined order of unblocking.

NAME

send, sendto, sendmsg – send a message from a socket

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>

int send(int s, const char *msg, int len, int flags);

int sendto(int s, const char *msg, int len, int flags, const struct sockaddr *to,
        int tolen);

int sendmsg(int s, const struct msghdr *msg, int flags);
```

MT-LEVEL

Safe

DESCRIPTION

send( ), **sendto( )**, and **sendmsg( )** are used to transmit a message to another transport end-point. **send( )** may be used only when the socket is in a *connected* state, while **sendto( )** and **sendmsg( )** may be used at any time. *s* is a socket created with **socket**(3N).

The address of the target is given by *to* with *tolen* specifying its size. The length of the message is given by *len*. If the message is too long to pass atomically through the underlying protocol, then the error **EMSGSIZE** is returned, and the message is not transmitted.

A return value of –1 indicates locally detected errors only. It does not implicitly mean the message was not delivered.

If the socket does not have enough buffer space available to hold the message being sent, **send( )** blocks, unless the socket has been placed in non-blocking I/O mode (see **fcntl**(2)). The **select**(3C) or **poll**(2) call may be used to determine when it is possible to send more data.

The *flags* parameter is formed from the bitwise OR of zero or more of the following:

MSG_OOB	Send "out-of-band" data on sockets that support this notion. The underlying protocol must also support "out-of-band" data. Only **SOCK_STREAM** sockets created in the **AF_INET** address family support out-of-band data.
MSG_DONTROUTE	The **SO_DONTROUTE** option is turned on for the duration of the operation. It is used only by diagnostic or routing programs.

See **recv**(3N) for a description of the **msghdr** structure.

RETURN VALUES

These calls return the number of bytes sent, or –1 if an error occurred.

ERRORS

The calls fail if:

EBADF	*s* is an invalid file descriptor.
EINTR	The operation was interrupted by delivery of a signal before any data could be buffered to be sent.
EINVAL	*tolen* is not the size of a valid address for the specified address family.

EMSGSIZE	The socket requires that message be sent atomically, and the message was too long.
ENOMEM	There was insufficient memory available to complete the operation.
ENOSR	There were insufficient STREAMS resources available for the operation to complete.
ENOTSOCK	*s* is not a socket.
EWOULDBLOCK	The socket is marked non-blocking and the requested operation would block.

SEE ALSO **fcntl**(2), **poll**(2), **write**(2), **connect**(3N), **getsockopt**(3N), **recv**(3N), **select**(3C), **socket**(3N)

NAME

setbuf, setvbuf – assign buffering to a stream

SYNOPSIS

```
#include <stdio.h>
void setbuf(FILE *stream, char *buf);
int setvbuf(FILE *stream, char *buf, int type, size_t size);
```

MT-LEVEL

MT-Safe

DESCRIPTION

setbuf() may be used after a *stream* (see **intro**(3)) has been opened but before it is read or written. It causes the array pointed to by *buf* to be used instead of an automatically allocated buffer. If *buf* is the **NULL** pointer input/output will be completely unbuffered. The constant **BUFSIZ**, defined in the **<stdio.h>** header, indicates how large the array pointed to by *buf* should be.

```
char buf[BUFSIZ];
```

setvbuf() may be used after a stream has been opened but before it is read or written. *type* determines how *stream* will be buffered. Legal values for *type* (defined in **<stdio.h>**) are:

_IOFBF	causes input/output to be fully buffered.
_IOLBF	causes output to be line buffered; the buffer will be flushed when a newline is written, the buffer is full, or input is requested.
_IONBF	causes input/output to be completely unbuffered.

If *buf* is not the **NULL** pointer, the array it points to will be used for buffering, instead of an automatically allocated buffer. *size* specifies the size of the buffer to be used. If input/output is unbuffered, *buf* and *size* are ignored.

For a further discussion of buffering, see **stdio**(3S).

RETURN VALUES

If an illegal value for *type* is provided, **setvbuf()** returns a non-zero value. Otherwise, it returns zero.

SEE ALSO

fopen(3S), **getc**(3S), **malloc**(3C), **putc**(3S), **stdio**(3S)

NOTES

A common source of error is allocating buffer space as an ‘‘automatic’’ variable in a code block, and then failing to close the stream in the same block.

When using **setbuf()**, *buf* should always be sized using **BUFSIZ**. If the array pointed to by *buf* is larger than **BUFSIZ**, a portion of *buf* will not be used. If *buf* is smaller than **BUFSIZ**, other memory may be unexpectedly overwritten.

Parts of **buf** will be used for internal bookkeeping of the stream and, therefore, **buf** will contain less than *size* bytes when full. It is recommended that **stdio**(3S) be used to handle buffer allocation when using **setvbuf()**.

NAME

setbuffer, setlinebuf – assign buffering to a stream

SYNOPSIS

#include <stdio.h>

void setbuffer(FILE **iop*, **char** **abuf,* **size_t** *asize*);

void setlinebuf(FILE **iop*);

DESCRIPTION

setbuffer, setlinebuf – assign buffering to a stream The three types of buffering available are unbuffered, block buffered, and line buffered. When an output stream is unbuffered, information appears on the destination file or terminal as soon as written; when it is block buffered many characters are saved up and written as a block; when it is line buffered characters are saved up until a NEWLINE is encountered or input is read from **stdin**. **fflush**(3S) may be used to force the block out early. Normally all files are block buffered. A buffer is obtained from **malloc**(3C) upon the first **getc**(3S) or **putc**(3S) on the file. If the standard stream **stdout** refers to a terminal it is line buffered. The standard stream **stderr** is unbuffered by default.

setbuffer( ) can be used after a stream, *iop,* has been opened but before it is read or written. It uses the character array *abuf* whose size is determined by the *asize* argument instead of an automatically allocated buffer. If *abuf* is the **NULL** pointer, input/output will be completely unbuffered. A manifest constant **BUFSIZ ,** defined in the **<stdio.h>** header, tells how big an array is needed:

char buf[BUFSIZ];

setlinebuf( ) is used to change the buffering on a stream from block buffered or unbuffered to line buffered. Unlike **setbuffer( ),** it can be used at any time that the stream, *iop,* is active.

A stream can be changed from unbuffered or line buffered to block buffered by using **freopen**(3S). A stream can be changed from block buffered or line buffered to unbuffered by using **freopen**(3S) followed by **setbuf**(3S) with a buffer argument of **NULL**.

SEE ALSO

malloc(3C), **fclose**(3S), **fopen**(3S), **fread**(3S), **getc**(3S), **printf**(3S), **putc**(3S), **puts**(3S), **setbuf**(3S), **setvbuf**(3S)

NOTES

A common source of error is allocating buffer space as an “automatic” variable in a code block, and then failing to close the stream in the same block.

NAME

setcat – define default catalog

SYNOPSIS

```
#include <pfmt.h>
char *setcat(const char *catalog);
```

MT-LEVEL

MT-safe

DESCRIPTION

The routine **setcat()** defines the default message catalog to be used by subsequent calls to **pfmt()**, **pfmt()** or **gettxt()** which do not explicitely specify a message catalog.

catalog must be limited to 14 characters. These characters must be selected from a set of all characters values, excluding **\0** (null) and the ASCII codes for **/** (slash) and **:** (colon).

setcat() assumes that the catalog exists. No checking is done on the argument.

A NULL pointer passed as an argument will result in the return of a pointer to the current default message catalog name. A pointer to an empty string passed as an argument will cancel the default catalog.

If no default catalog is specified, or if *catalog* is an invalid catalog name, Subsequent calls to **gettxt()**, **pfmt()** or **lfmt()** that do not explicitely specify a catalog name will use **Message not found!!\n** as default string.

RETURN VALUE

Upon success, **setcat()** returns a pointer to the catalog name. Upon failure, **setcat()** returns a NULL pointer.

EXAMPLE

```
setcat("test");
gettxt(":10", "hello world\n")
```

SEE ALSO

gettxt(3C), **lfmt**(3C), **pfmt**(3C), **setlocale**(3C), **environ**(5)

NAME

setjmp, longjmp, _setjmp, _longjmp – non-local goto

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <setjmp.h>

int setjmp(*env***)**
jmp_buf *env*;

void longjmp(*env***,** *val***)**
jmp_buf *env*;
int *val*;

int _setjmp(*env***)**
jmp_buf *env*;

void _longjmp(*env***,** *val***)**
jmp_buf *env*;
int *val*;

DESCRIPTION

setjmp( ) and **longjmp( )** are useful for dealing with errors and interrupts encountered in a low-level subroutine of a program.

setjmp( ) saves its stack environment in *env* for later use by **longjmp( )**. A normal call to **setjmp( )** returns zero. **setjmp( )** also saves the register environment. If a **longjmp( )** call will be made, the routine which called **setjmp( )** should not return until after the **longjmp( )** has returned control (see below).

longjmp( ) restores the environment saved by the last call of **setjmp( )**, and then returns in such a way that execution continues as if the call of **setjmp( )** had just returned the value *val* to the function that invoked **setjmp( )**; however, if *val* were zero, execution would continue as if the call of **setjmp( )** had returned one. This ensures that a ''return'' from **setjmp( )** caused by a call to **longjmp( )** can be distinguished from a regular return from **setjmp( )**. The calling function must not itself have returned in the interim, otherwise **longjmp( )** will be returning control to a possibly non-existent environment. All memory-bound data have values as of the time **longjmp( )** was called. The CPU and floating-point data registers are restored to the values they had at the time that **setjmp( )** was called. But, because the **register** storage class is only a hint to the C compiler, variables declared as **register** variables may not necessarily be assigned to machine registers, so their values are unpredictable after a **longjmp( )**. This is especially a problem for programmers trying to write machine-independent C routines.

setjmp( ) and **longjmp( )** save and restore the signal mask while **_setjmp( )** and **_longjmp( )** manipulate only the C stack and registers.

None of these functions save or restore any floating-point status or control registers.

EXAMPLES

The following example uses both **setjmp()** and **longjmp()** to return the flow of control to the appropriate instruction block:

```
#include <stdio.h>
#include <setjmp.h>
#include <signal.h>
#include <unistd.h>
jmp_buf env; static void signal_handler();

main() {
      int returned_from_longjump, processing = 1;
      unsigned int time_interval = 4;
      if ((returned_from_longjump = setjmp(env)) != 0)
         switch (returned_from_longjump)    {
          case SIGINT:
            printf("longjumped from interrupt %d\n",SIGINT);
            break;
          case SIGALRM:
            printf("longjumped from alarm %d\n",SIGALRM);
            break;
         }
      (void) signal(SIGINT, signal_handler);
      (void) signal(SIGALRM, signal_handler);
      alarm(time_interval);
      while (processing)      {
       printf(" waiting for you to INTERRUPT (cntrl-C) ...\n");
       sleep(1);
      }                              /* end while forever loop */
}

static void signal_handler(sig)
int sig; {
      switch (sig)    {
       case SIGINT:          ...  /* process for interrupt */
                             longjmp(env,sig);
                               /* break never reached */
       case SIGALRM:         ...  /* process for alarm */
                             longjmp(env,sig);
                               /* break never reached */
       default:              exit(sig);
      }
}
```

When this example is compiled and executed, and the user sends an interrupt signal, the output will be:

longjumped from interrupt

Additionally, every 4 seconds the alarm will expire, signalling this process, and the output will be:

longjumped from alarm

SEE ALSO

cc(1B), **sigvec**(3B), **setjmp**(3C), **signal**(3C)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

BUGS

setjmp( ) does not save the current notion of whether the process is executing on the signal stack. The result is that a **longjmp( )** to some place on the signal stack leaves the signal stack state incorrect.

On some systems **setjmp( )** also saves the register environment. Therefore, all data that are bound to registers are restored to the values they had at the time that **setjmp( )** was called. All memory-bound data have values as of the time **longjmp( )** was called. However, because the **register** storage class is only a hint to the C compiler, variables declared as **register** variables may not necessarily be assigned to machine registers, so their values are unpredictable after a **longjmp( )**. When using compiler options that specify automatic register allocation (see **cc**(1B)), the compiler will not attempt to assign variables to registers in routines that call **setjmp( )**.

longjmp( ) never causes **setjmp( )** to return zero, so programmers should not depend on **longjmp( )** being able to cause **setjmp( )** to return zero.

NAME

setjmp, sigsetjmp, longjmp, siglongjmp – non-local goto

SYNOPSIS

```
#include <setjmp.h>
int setjmp(jmp_buf env);
int sigsetjmp(sigjmp_buf env, int savemask);
void longjmp(jmp_buf env, int val);
void siglongjmp(sigjmp_buf env, int val);
```

MT-LEVEL

Unsafe

DESCRIPTION

These functions are useful for dealing with errors and interrupts encountered in a low-level subroutine of a program.

setjmp( ) saves its stack environment in *env* for later use by **longjmp( )**.

sigsetjmp( ) saves the calling process's registers and stack environment (see **sigaltstack**(2)) in *env* for later use by **siglongjmp( )**. If *savemask* is non-zero, the calling process's signal mask (see **sigprocmask**(2)) and scheduling parameters (see **priocntl**(2)) are also saved.

longjmp( ) restores the environment saved by the last call of **setjmp( )** with the corresponding *env* argument. After **longjmp( )** is completed, program execution continues as if the corresponding call of **setjmp( )** had just returned the value *val*. The caller of **setjmp( )** must not have returned in the interim. **longjmp( )** cannot cause **setjmp( )** to return the value 0. If **longjmp( )** is invoked with a second argument of 0, **setjmp( )** will return 1. At the time of the second return from **setjmp( )**, all external and static variables have values as of the time **longjmp( )** is called (see example).

siglongjmp( ) restores the environment saved by the last call of **sigsetjmp( )** with the corresponding *env* argument. After **siglongjmp( )** is completed, program execution continues as if the corresponding call of **sigsetjmp( )** had just returned the value *val*. **siglongjmp( )** cannot cause **sigsetjmp( )** to return the value 0. If **siglongjmp( )** is invoked with a second argument of 0, **sigsetjmp( )** will return 1. At the time of the second return from **sigsetjmp( )**, all external and static variables have values as of the time **siglongjmp( )** is called.

If a signal-catching function interrupts **sleep( )** and calls **siglongjmp( )** to restore an environment saved prior to the **sleep( )** call, the action associated with **SIGALRM** and time it is scheduled to be generated are unspecified. It is also unspecified whether the **SIGALRM** signal is blocked, unless the process's signal mask is restored as part of the environment.

The function **siglongjmp( )** restores the saved signal mask if and only if the *env* argument was initialized by a call to the **sigsetjmp( )** function with a non-zero *savemask* argument.

The values of register and automatic variables are undefined. Register or automatic variables whose value must be relied upon must be declared as **volatile**.

EXAMPLES

The following example uses both **setjmp( )** and **longjmp( )** to return the flow of control to the appropriate instruction block:

```
#include <stdio.h>
#include <setjmp.h>
#include <signal.h>
#include <unistd.h>
jmp_buf env; static void signal_handler();

main( ) {
      int returned_from_longjump, processing = 1;
      unsigned int time_interval = 4;
      if ((returned_from_longjump = setjmp(env)) != 0)
         switch (returned_from_longjump)    {
          case SIGINT:
            printf("longjumped from interrupt %d\n",SIGINT);
            break;
          case SIGALRM:
            printf("longjumped from alarm %d\n",SIGALRM);
            break;
         }
      (void) signal(SIGINT, signal_handler);
      (void) signal(SIGALRM, signal_handler);
      alarm(time_interval);
      while (processing)      {
       printf(" waiting for you to INTERRUPT (cntrl-C) ...\n");
       sleep(1);
      }                              /* end while forever loop */
}

static void signal_handler(sig)
int sig; {
      switch (sig)    {
       case SIGINT:          ...  /* process for interrupt */
                             longjmp(env,sig);
                               /* break never reached */
       case SIGALRM:         ...  /* process for alarm */
                             longjmp(env,sig);
                               /* break never reached */
       default:              exit(sig);
      }
}
```

When this example is compiled and executed, and the user sends an interrupt signal, the output will be:

longjumped from interrupt

Additionally, every 4 seconds the alarm will expire, signalling this process, and the output will be:

longjumped from alarm

RETURN VALUES

If **longjmp( )** or **siglongjmp( )** are invoked with a second argument of 0, **setjmp( )** and **sigsetjmp( ) ,** respectively, return 1. Otherwise, **setjmp( )** and **sigsetjmp( )** return 0.

SEE ALSO

getcontext(2), **priocntl**(2), **sigaction**(2), **sigaltstack**(2), **sigprocmask**(2), **signal**(3C)

WARNINGS

If **longjmp( )** or **siglongjmp( )** are called even though *env* was never primed by a call to **setjmp( )** or **sigsetjmp( )**, or when the last such call was in a function that has since returned, absolute chaos is guaranteed.

NAME

setlabel – define the label for pfmt() and lfmt().

MT-LEVEL

MT-safe

SYNOPSIS

```
#include <pfmt.h>
int setlabel(const char *label);
```

DESCRIPTION

The routine **setlabel()** defines the label for messages produced in standard format by subsequent calls to **pfmt()** and **lfmt()**.

label is a character string no more than 25 characters in length.

No label is defined before **setlabel()** is called. A NULL pointer or an empty string passed as argument will reset the definition of the label.

RETURN VALUE

setlabel() returns **0** in case of success, non-zero otherwise.

EXAMPLE

The following code (without previous call to **setlabel()**):

```
pfmt(stderr, MM_ERROR, "test:2:Cannot open file\n");
setlabel("UX:test");
pfmt(stderr, MM_ERROR, "test:2:Cannot open file\n");
```

will produce the following output:

```
ERROR: Cannot open file
UX:test: ERROR: Cannot open file
```

USAGE

The label should be set once at the beginning of a utility and remain constant.

SEE ALSO

getopt(3C), **lfmt**(3C), **pfmt**(3C)

NAME

setlocale – modify and query a program's locale

SYNOPSIS

#include <locale.h>

char ∗**setlocale(int** *category***, const char** ∗*locale***);**

MT-LEVEL

Safe with exceptions

DESCRIPTION

setlocale( ) selects the appropriate piece of the program's locale as specified by the *category* and *locale* arguments. The *category* argument may have the following values: **LC_CTYPE**, **LC_NUMERIC**, **LC_TIME**, **LC_COLLATE**, **LC_MONETARY**, **LC_MESSAGES**, and **LC_ALL**. These names are defined in the **<locale.h>** header. **LC_ALL** names all of a program's locale categories.

LC_CTYPE affects the behavior of character handling functions such as **isdigit( )** and **tolower( )**, and multibyte character functions such as **mbtowc( )** and **wctomb( )**.

LC_NUMERIC affects the decimal point character and thousands separator character for the formatted input / output functions and string conversion functions.

LC_TIME affects the date and time format as delivered by **ascftime( )**, **cftime( )**, **getdate( )**, and **strftime( )**.

LC_COLLATE affects the sort order produced by **strcoll( )** and **strxfrm( )**.

LC_MONETARY affects the monetary formatted information returned by **localeconv( )**.

LC_MESSAGES affects the behavior of **dgettext( )**, **gettext( )**, and **gettxt( )**.

Each category corresponds to a set of databases which contain the relevant information for each defined locale. The location of a database is given by the following path, **/usr/lib/locale/***locale* / *category*, where *locale* and *category* are the names of locale and category, respectively. For example, the database for the **LC_CTYPE** category of the **de** (Deutsch or German) locale would be found in **/usr/lib/locale/de/LC_CTYPE**

A value of "C" for *locale* specifies the traditional UNIX system behavior. At program startup, the equivalent of

setlocale(LC_ALL, "C")

is executed. This has the effect of initializing each category to the locale described by the environment "C".

A value of "" for *locale* specifies that the locale should be taken from environment variables. The order in which the environment variables are checked for the various categories is given below:

Category	1st Env. Var.	2nd Env. Var.	3rd Env. Var
LC_CTYPE:	**LC_ALL**	**LC_CTYPE**	**LANG**
LC_COLLATE:	**LC_ALL**	**LC_COLLATE**	**LANG**
LC_TIME:	**LC_ALL**	**LC_TIME**	**LANG**
LC_NUMERIC:	**LC_ALL**	**LC_NUMERIC**	**LANG**
LC_MONETARY:	**LC_ALL**	**LC_MONETARY**	**LANG**
LC_MESSAGES:	**LC_ALL**	**LC_MESSAGES**	**LANG**

If a pointer to a string is given for *locale*, **setlocale( )** attempts to set the locale for the given category to *locale*. If **setlocale( )** succeeds, *locale* is returned. If **setlocale( )** fails, a null pointer is returned and the program's locale is not changed.

For category **LC_ALL**, the behavior is slightly different. If a pointer to a string is given for *locale* and **LC_ALL** is given for *category*, **setlocale( )** attempts to set the locale for all the categories to *locale*. The *locale* may be a simple locale, consisting of a single locale, or a composite locale. A composite locale is a string returned by a prior call to **setlocale(LC_ALL,0)**. This string will restore each category to the previous locale. If **setlocale( )** fails to set the locale for any category, a null pointer is returned and the program's locale for all categories is not changed. Otherwise, locale is returned.

A null pointer for *locale* causes **setlocale( )** to return the current locale associated with the *category*. The program's locale is not changed.

FILES

/usr/lib/locale/*locale***/LC_CTYPE**	character type database for *locale*
/usr/lib/locale/*locale***/LC_NUMERIC**	numeric format data for *locale*
/usr/lib/locale/*locale***/LC_TIME**	date and time formats for *locale*
/usr/lib/locale/*locale***/LC_COLLATE**	sort ordering information for *locale*
/usr/lib/locale/*locale***/LC_MESSAGES**	message catalogs for *locale*
/usr/lib/locale/*locale***/LC_MONETARY**	currency format data for *locale*

SEE ALSO

ctype(3C), **localeconv**(3C), **mbchar**(3C), **strcoll**(3C), **strftime**(3C), **gettext**(3I), **environ**(5)

NOTES

To change locale in a multi-thread application **setlocale** should be called prior to using any locale sensitive routine. Using **setlocale** to query the current locale is safe and can be used anywhere in a multi-thread application.

It is the user's responsibility to ensure that mixed locale categories are compatible. For example, setting **LC_CTYPE=C** and **LC_TIME=ja** (where **ja** indicates Japanese) will not work, because Japanese time cannot be represented in the "C" locale's ASCII codeset.

NAME

shm_open – open a shared memory object

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

#include <sys/mman.h>

int shm_open(const char **name*, **int** *oflag*, **mode_t** *mode* **);**

MT-LEVEL

MT-Safe

DESCRIPTION

shm_open( ) either opens a file descriptor for the shared memory object with the name referenced by *name.* If successful, **shm_open( )** returns a file descriptor for the shared memory object that is the lowest numbered file descriptor not currently open for that process. Since the open file description is new, the new file descriptor is not as yet shared with any other processes.

The file status flags and file access modes of the open file descriptor are set according to the value of *oflag:* the bitwise inclusive OR of the following flags, defined in the header <fcntl.h>. (Applications must specify exactly one of the first two values below in the value of *oflag):*

O_RDONLY Open for read access only.

O_RDWR Open for read or write access.

Any combination of the remaining flags may be bitwise inclusive OR- ed with the value of *oflag:*

O_CREAT If *name* does not exist, the shared memory object is created, it's user ID is set to the effective user ID of the process, and it's group ID is set to a system default group ID or to the effective group ID of the process. The shared memory object's permission bits will be set to the value of *mode*, modified by clearing all bits set in the file mode creation mask of the process (see **umask**(2)).

mode does not affect whether the shared memory object is opened for reading, for writing, or for both. The new shared memory object has a size of zero.

If the shared memory object does exist, this flag will have no effect, except as specified under **O_EXCL** below.

O_EXCL If both **OEXCL** and **O_CREAT** are set, **shm_open( )** fails if the shared memory object, *name,* exists. The check for the existence of the shared memory object and the creation of the object if it does not exist is atomic with respect to other processes executing **shm_open( )** naming the same shared memory object with
OEXCL and **O_CREAT** set.

O_TRUNC If the shared memory object exists, and it is successfully opened **O_RDWR,** the object is truncated to zero length and the mode and ownership are unchanged by this function call.

RETURN VALUES

If successful, **shm_open( )** returns a nonnegative integer representing the lowest numbered unused file descriptor, otherwise it returns **-1** and sets **errno** to indicate the error condition.

ERRORS

EACCES The shared memory object exists and the permissions specified by *oflag* are denied, or the shared memory object does not exist and permission to create the shared memory object is denied, or **O_TRUNC** is specified and write permission is denied.

EEXIST **O_CREAT** and **O_EXCL** are set and the named shared memory object already exists.

EINTR The **shm_open( )** operation was interrupted by a signal.

EINVAL *name* is an invalid file description.

EMFILE The number of open file descriptors in this process exceeds **{OPEN_MAX}**.

ENAMETOOLONG The length of the *name* string exceeds **{PATH_MAX}**, or a pathname component is longer than **{NAME_MAX}** while **_POSIX_NO_TRUNC** is in effect.

ENFILE The system file table is full

ENOENT **O_CREAT** is not set and the named shared memory object does not exist.

ENOSPC There is insufficient space for the creation of the new shared memory object.

ENOSYS **shm_open( )** is not supported by this implementation.

FILES

/usr/include/fcntl.h

SEE ALSO

close(2), **dup**(2), **exec**(2), **fcntl**(2), **mmap**(2), **umask**(2), **sysconf**(3C), **shm_unlink**(3R), **fcntl**(5)

NOTES

When a shared memory object is created, the state of the shared memory object, including all data associated with the shared memory object, persists until the shared memory object is unlinked and all other references are gone.

BUGS

In Solaris 2.5, these functions always return –**1** and set **errno** to **ENOSYS,** because this release does not support the Shared Memory Objects option. It is our intention to provide support for these interfaces in future releases.

NAME

shm_unlink – remove a shared memory object

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

int shm_unlink(const char ∗*name***);**

MT-LEVEL

MT-Safe

DESCRIPTION

shm_unlink( ) removes the name of the shared memory object named by the string pointed to by *name* . If one or more references to the shared memory object exists when the object is unlinked, the name is removed before **shm_unlink( )** returns, but the removal of the memory object contents will be postponed until all open and mapped references to the shared memory object have been removed.

RETURN VALUES

If successful, **shm_unlink( )** returns **0,** otherwise it returns **-1** and sets **errno** to indicate the error condition, and the named shared memory object is not affected by this function.

ERRORS

EACCES Permission is denied to unlink the named shared memory object.

ENAMETOOLONG
The length of the *name* string exceeds **{PATH_MAX}**,or a pathname component is longer than **{NAME_MAX}** while **_POSIX_NO_TRUNC** is in effect.

ENOENT The named shared memory object does not exist.

ENOSYS **shm_unlink( )** is not supported by this implementation.

SEE ALSO

close(2), **mmap**(2), **mlock**(3C), **shm_open**(3R)

BUGS

In Solaris 2.5, these functions always return **–1** and set **errno** to **ENOSYS,** because this release does not support the Shared Memory Objects option. It is our intention to provide support for these interfaces in future releases.

NAME

shutdown – shut down part of a full-duplex connection

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

int shutdown(int *s***, int** *how***);**

MT-LEVEL

Safe

DESCRIPTION

The **shutdown( )** call shuts down all or part of a full-duplex connection on the socket associated with *s*. If *how* is **0**, then further receives will be disallowed. If *how* is **1**, then further sends will be disallowed. If *how* is **2**, then further sends and receives will be disallowed.

RETURN VALUES

A 0 is returned if the call succeeds, –1 if it fails.

ERRORS

The call succeeds unless:

EBADF	*s* is not a valid file descriptor.
ENOMEM	There was insufficient user memory available for the operation to complete.
ENOSR	There were insufficient STREAMS resources available for the operation to complete.
ENOTCONN	The specified socket is not connected.
ENOTSOCK	*s* is not a socket.

SEE ALSO

connect(3N), **socket**(3N)

NOTES

The *how* values should be defined constants.

NAME

sigblock, sigmask, sigpause, sigsetmask – block signals

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <signal.h>

int sigblock(*mask***)**
int *mask*;

int sigmask(*signum***)**
int *signum*;

int sigpause(int *mask***)**
int *mask*;

int sigsetmask(*mask***)**
int *mask*;

DESCRIPTION

sigblock, sigmask, sigpause, sigsetmask – block signals

sigblock( ) adds the signals specified in *mask* to the set of signals currently being blocked from delivery. Signals are blocked if the appropriate bit in *mask* is a 1; the macro **sigmask** is provided to construct the mask for a given *signum*. **sigblock( )** returns the previous mask. The previous mask may be restored using **sigsetmask( )**.

sigpause( ) assigns *mask* to the set of masked signals and then waits for a signal to arrive; on return the set of masked signals is restored. *mask* is usually 0 to indicate that no signals are now to be blocked. **sigpause( )** always terminates by being interrupted, returning –1 and setting **errno** to EINTR.

sigsetmask( ) sets the current signal mask (those signals that are blocked from delivery). Signals are blocked if the corresponding bit in *mask* is a 1; the macro **sigmask** is provided to construct the mask for a given *signum*.

In normal usage, a signal is blocked using **sigblock( )**. To begin a critical section, variables modified on the occurrence of the signal are examined to determine that there is no work to be done, and the process pauses awaiting work by using **sigpause( )** with the mask returned by **sigblock( )**.

It is not possible to block **SIGKILL**, **SIGSTOP**, or **SIGCONT**, this restriction is silently imposed by the system.

RETURN VALUES

sigblock( ) and **sigsetmask( )** return the previous set of masked signals. **sigpause( )** returns –1 and sets **errno** to EINTR.

SEE ALSO

kill(2), **sigaction**(2), **signal**(3B), **sigvec**(3B)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

NAME

sigfpe – signal handling for specific SIGFPE codes

SYNOPSIS

#include <floatingpoint.h>

#include <siginfo.h>

sigfpe_handler_type sigfpe(sigfpe_code_type *code***, sigfpe_handler_type** *hdl***);**

MT-LEVEL

Safe

DESCRIPTION

This function allows signal handling to be specified for particular **SIGFPE** codes. A call to **sigfpe( )** defines a new handler *hdl* for a particular **SIGFPE** *code* and returns the old handler as the value of the function **sigfpe( ).** Normally handlers are specified as pointers to functions; the special cases **SIGFPE_IGNORE**, **SIGFPE_ABORT**, and **SIGFPE_DEFAULT** allow ignoring, dumping core using **abort**(3C), or default handling respectively. Default handling is to dump core using **abort**(3C).

code is usually one of the five IEEE 754-related **SIGFPE** codes:

FPE_FLTRES	fp_inexact – floating-point inexact result
FPE_FLTDIV	fp_division – floating-point division by zero
FPE_FLTUND	fp_underflow – floating-point underflow
FPE_FLTOVF	fp_overflow – floating-point overflow
FPE_FLTINV	fp_invalid – floating-point invalid operation

Three steps are required to intercept an IEEE 754-related **SIGFPE** code with **sigfpe( )**:

1) Set up a handler with **sigfpe( )**.

2) Enable the relevant IEEE 754 trapping capability in the hardware, perhaps by using assembly-language instructions.

3) Perform a floating-point operation that generates the intended IEEE 754 exception.

sigfpe( ) never changes floating-point hardware mode bits affecting IEEE 754 trapping. No IEEE 754-related **SIGFPE** signals will be generated unless those hardware mode bits are enabled.

SIGFPE signals can be handled using **sigfpe**( ), **sigaction**(2) or **signal**(3C). In a particular program, to avoid confusion, use only one of these interfaces to handle **SIGFPE** signals.

EXAMPLES

A user-specified signal handler might look like this:

```
#include <floatingpoint.h>
#include <siginfo.h>
#include <ucontext.h>
/*
 * The sample_handler prints out a message then commits suicide.
 */
void
sample_handler(int sig, siginfo_t *sip, ucontext_t *uap) {
        char *label;

        switch (sip->si_code) {
        case FPE_FLTINV: label = "invalid operand"; break;
        case FPE_FLTRES: label = "inexact"; break;
        case FPE_FLTDIV: label = "division-by-zero"; break;
        case FPE_FLTUND: label = "underflow"; break;
        case FPE_FLTOVF: label = "overflow"; break;
        default: label = "???"; break;
        }

        fprintf(stderr, "FP exception %s (0x%x) occurred at address %p.\n",
                        label, sip->si_code, (void *) sip->si_addr);
        abort();
}
```

and it might be set up like this:

```
#include <floatingpoint.h>
#include <siginfo.h>
#include <ucontext.h>
extern void sample_handler(int, siginfo_t *, ucontext_t *);
main(void) {
        sigfpe_handler_type hdl, old_handler1, old_handler2;
/*
 * save current fp_overflow and fp_invalid handlers; set the new
 * fp_overflow handler to sample_handler() and set the new
 * fp_invalid handler to SIGFPE_ABORT (abort on invalid)
 */
        hdl = (sigfpe_handler_type) sample_handler;
        old_handler1 = sigfpe(FPE_FLTOVF, hdl);
        old_handler2 = sigfpe(FPE_FLTINV, SIGFPE_ABORT);
        ...
/*
 * restore old fp_overflow and fp_invalid handlers
 */
```

```
        sigfpe(FPE_FLTOVF, old_handler1);
        sigfpe(FPE_FLTINV, old_handler2);
}
```

FILES

/usr/include/floatingpoint.h
/usr/include/siginfo.h

SEE ALSO

sigaction(2), **abort**(3C), **signal**(3C), **floatingpoint**(5)

DIAGNOSTICS

sigfpe( ) returns BADSIG if *code* is not zero or a defined **SIGFPE** code.

NAME

siginterrupt – allow signals to interrupt functions

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

int siginterrupt(*sig, flag***)**
int *sig, flag***;**

DESCRIPTION

siginterrupt( ) is used to change the function restart behavior when a function is interrupted by the specified signal. If the flag is false (0), then functions will be restarted if they are interrupted by the specified signal and no data has been transferred yet. System call restart is the default behavior when the **signal**(3C) routine is used.

If the flag is true (1), then restarting of functions is disabled. If a function is interrupted by the specified signal and no data has been transferred, the function will return –1 with **errno** set to **EINTR**. Interrupted functions that have started transferring data will return the amount of data actually transferred.

Issuing a **siginterrupt( )** call during the execution of a signal handler will cause the new action to take place on the next signal to be caught.

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

This library routine uses an extension of the **sigvec**(3B) function that is not available in 4.2BSD, hence it should not be used if backward compatibility is needed.

RETURN VALUES

A 0 value indicates that the call succeeded. A –1 value indicates that the call failed and **errno** is set to indicate the error.

ERRORS

siginterrupt( ) may return the following error:

EINVAL *sig* is not a valid signal.

SEE ALSO

sigblock(3B), **sigvec**(3B), **signal**(3C)

NAME

signal – simplified software signal facilities

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <signal.h>

void (*signal(*sig,func***))()**
int *sig***;**
void (**func***)()**;

DESCRIPTION

signal( ) is a simplified interface to the more general **sigvec**(3B) facility. Programs that use **signal( )** in preference to **sigvec( )** are more likely to be portable to all systems.

A signal is generated by some abnormal event, initiated by a user at a terminal (quit, interrupt, stop), by a program error (bus error, etc.), by request of another program (kill), or when a process is stopped because it wishes to access its control terminal while in the background (see **termio**(7I)). Signals are optionally generated when a process resumes after being stopped, when the status of child processes changes, or when input is ready at the control terminal. Most signals cause termination of the receiving process if no action is taken; some signals instead cause the process receiving them to be stopped, or are simply discarded if the process has not requested otherwise. Except for the **SIGKILL** and **SIGSTOP** signals, the **signal( )** call allows signals either to be ignored or to interrupt to a specified location. See **sigvec**(3B) for a complete list of the signals.

If *func* is **SIG_DFL**, the default action for signal *sig* is reinstated; this default is termination (with a core image for starred signals) except for signals marked with • or †. Signals marked with • are discarded if the action is **SIG_DFL**; signals marked with † cause the process to stop. If *func* is **SIG_IGN** the signal is subsequently ignored and pending instances of the signal are discarded. Otherwise, when the signal occurs further occurrences of the signal are automatically blocked and *func* is called.

A return from the function unblocks the handled signal and continues the process at the point it was interrupted.

If a caught signal occurs during certain functions, terminating the call prematurely, the call is automatically restarted. In particular this can occur during a **read**(2) or **write**(2) on a slow device (such as a terminal; but not a file) and during a **wait**(2).

The value of **signal( )** is the previous (or initial) value of *func* for the particular signal.

After a **fork**(2) or **vfork**(2) the child inherits all signals. An **exec**(2) resets all caught signals to the default action; ignored signals remain ignored.

RETURN VALUES

The previous action is returned on a successful call. Otherwise, –1 is returned and **errno** is set to indicate the error.

ERRORS

signal( ) will fail and no action will take place if the following occurs:

EINVAL *sig* is not a valid signal number, or is **SIGKILL** or **SIGSTOP**.

SEE ALSO

kill(1), **exec**(2), **fcntl**(2), **fork**(2), **getitimer**(2), **getrlimit**(2), **kill**(2), **ptrace**(2), **read**(2), **sigaction**(2), **wait**(2), **write**(2), **abort**(3C), **setjmp**(3B), **sigblock**(3B), **sigstack**(3B), **sigvec**(3B), **wait**(3B), **setjmp**(3C), **signal**(3C), **signal**(5), **termio**(7I)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

The handler routine, *func*, can be declared:

void handler(*signum***)**
int *signum***;**

Here *signum* is the signal number. See **sigvec**(3B) for more details.

NAME

signal, sigset, sighold, sigrelse, sigignore, sigpause – simplified signal management for application processes

SYNOPSIS

#include <signal.h>

void (*signal (int *sig***, void (***disp***)(int)))(int);**

void (*sigset (int *sig***, void (***disp***)(int)))(int);**

int sighold(int *sig***);**

int sigrelse(int *sig***);**

int sigignore(int *sig***);**

int sigpause(int *sig***);**

DESCRIPTION

These functions provide simplified signal management for application processes. See **signal**(5) for an explanation of general signal concepts.

signal( ) and **sigset( )** are used to modify signal dispositions. *sig* specifies the signal, which may be any signal except **SIGKILL** and **SIGSTOP**. *disp* specifies the signal's disposition, which may be **SIG_DFL**, **SIG_IGN**, or the address of a signal handler. If **signal( )** is used, *disp* is the address of a signal handler, and *sig* is not **SIGILL**, **SIGTRAP**, or **SIGPWR**, the system first sets the signal's disposition to **SIG_DFL** before executing the signal handler. If **sigset( )** is used and *disp* is the address of a signal handler, the system adds *sig* to the calling process's signal mask before executing the signal handler; when the signal handler returns, the system restores the calling process's signal mask to its state prior to the delivery of the signal. In addition, if **sigset( )** is used and *disp* is equal to **SIG_HOLD**, *sig* is added to the calling process's signal mask and the signal's disposition remains unchanged.

sighold( ) adds *sig* to the calling process's signal mask.

sigrelse( ) removes *sig* from the calling process's signal mask.

sigignore( ) sets the disposition of *sig* to **SIG_IGN**.

sigpause( ) removes *sig* from the calling process's signal mask and suspends the calling process until a signal is received.

RETURN VALUES

On success, **signal( )** returns the signal's previous disposition. On failure, it returns **SIG_ERR** and sets **errno** to indicate the error.

On success, **sigset( )** returns **SIG_HOLD** if the signal had been blocked or the signal's previous disposition if it had not been blocked. On failure, it returns **SIG_ERR** and sets **errno** to indicate the error.

All other functions return zero on success. On failure, they return –1 and set **errno** to indicate the error.

ERRORS

These functions fail if any of the following are true:

EINTR A signal was caught during the function **sigpause( )**.

EINVAL The value of the *sig* argument is not a valid signal or is equal to **SIGKILL** or **SIGSTOP**.

SEE ALSO

kill(2), **pause**(2), **sigaction**(2), **sigsend**(2), **wait**(2), **waitid**(2), **signal**(5)

NOTES

sighold( ) in conjunction with **sigrelse( )** or **sigpause( )** may be used to establish critical regions of code that require the delivery of a signal to be temporarily deferred.

If **signal( )** or **sigset( )** is used to set **SIGCHLD**'s disposition to a signal handler, **SIGCHLD** will not be sent when the calling process's children are stopped or continued.

If any of the above functions are used to set **SIGCHLD**'s disposition to **SIG_IGN**, the calling process's child processes will not create zombie processes when they terminate (see **exit**(2)). If the calling process subsequently waits for its children, it blocks until all of its children terminate; it then returns a value of –1 with **errno** set to **ECHILD** (see **wait**(2), **waitid**(2)).

NAME

sigqueue – queue a signal to a process

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <signal.h>

int sigqueue(pid_t pid, int signo, const union sigval value);

union sigval {
    int     sival_int;      /* integer value */
    void    *sival_ptr;     /* pointer value */
};
```

MT-LEVEL

Async-Signal-Safe

DESCRIPTION

sigqueue( ) causes the signal, *signo* to be sent with the value, *value* to the process, *pid.* If *signo* is zero (the null signal), error checking is performed, but no signal is actually sent. The null signal can be used to check the validity of *pid.*

The conditions required for a process to have permission to queue a signal to another process are the same as for **kill**(2).

If resources were not available to queue the signal, **sigqueue( )** exits and returns immediately. If **SA_SIGINFO** is set for *signo* in the receiving process, and if the resources were available, the signal is left queued and pending. If **SA_SIGINFO** is not set for *signo*, then *signo* is sent at least once to the receiving process.

If the value of *pid* causes *signo* to be generated for the sending process, and if *signo* is not blocked, either *signo* or at least the pending, unblocked signal with the lowest number will be delivered to the sending process before **sigqueue( )** returns.

RETURN VALUES

If successful, **sigqueue( )** returns **0**, and queues the specified signal. Otherwise, **sigqueue( )** returns **-1** and sets **errno** to indicate the error condition.

ERRORS

EAGAIN	No resources are available to queue the signal. The process has already queued **{SIGQUEUE_MAX}** signals that are still pending at the receiver(s), or a system wide resource limit has been exceeded.
EINVAL	The value of *signo* is an invalid or unsupported signal number.
ENOSYS	**sigqueue( )** is not supported by this implementation.
EPERM	The process does not have the appropriate privilege to send the signal to the receiving process.
ESRCH	The process *pid* does not exist.

SEE ALSO

kill(2), **sigwaitinfo**(3R), **siginfo**(5), **signal**(5)

NAME

sigsetops, sigemptyset, sigfillset, sigaddset, sigdelset, sigismember – manipulate sets of signals

SYNOPSIS

```
#include <signal.h>
int sigemptyset(sigset_t *set);
int sigfillset(sigset_t *set);
int sigaddset(sigset_t *set, int signo);
int sigdelset(sigset_t *set, int signo);
int sigismember(sigset_t *set, int signo);
```

MT-LEVEL

MT-Safe

Async-Signal-Safe

DESCRIPTION

These functions manipulate *sigset_t* data types, representing the set of signals supported by the implementation.

sigemptyset( ) initializes the set pointed to by *set* to exclude all signals defined by the system.

sigfillset( ) initializes the set pointed to by *set* to include all signals defined by the system.

sigaddset( ) adds the individual signal specified by the value of *signo* to the set pointed to by *set*.

sigdelset( ) deletes the individual signal specified by the value of *signo* from the set pointed to by *set*.

sigismember( ) checks whether the signal specified by the value of *signo* is a member of the set pointed to by *set*.

Any object of type *sigset_t* must be initialized by applying either **sigemptyset( )** or **sigfillset( )** before applying any other operation.

RETURN VALUES

Upon successful completion, the **sigismember( )** function returns a value of one if the specified signal is a member of the specified set, or a value of 0 if it is not. Upon successful completion, the other functions return a value of 0. Otherwise a value of –1 is returned and **errno** is set to indicate the error.

ERRORS

sigaddset( ), **sigdelset( )**, and **sigismember( )** will fail if the following is true:

EINVAL The value of the *signo* argument is not a valid signal number.

sigfillset( ) will fail if the following is true:

EFAULT The *set* argument specifies an invalid address.

SEE ALSO

sigaction(2), **sigpending**(2), **sigprocmask**(2), **sigsuspend**(2), **signal**(5)

NAME

sigstack – set and / or get signal stack context

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <signal.h>

int sigstack(*nss, oss***)**
struct sigstack **nss, *oss*;*

DESCRIPTION

sigstack( ) allows users to define an alternate stack, called the ''signal stack'', on which signals are to be processed. When a signal's action indicates its handler should execute on the signal stack (specified with a **sigvec**(3B) call), the system checks to see if the process is currently executing on that stack. If the process is not currently executing on the signal stack, the system arranges a switch to the signal stack for the duration of the signal handler's execution.

A signal stack is specified by a **sigstack( )** structure, which includes the following members:

```
char    *ss_sp;      /* signal stack pointer */
int     ss_onstack;  /* current status */
```

ss_sp is the initial value to be assigned to the stack pointer when the system switches the process to the signal stack. Note that, on machines where the stack grows downwards in memory, this is *not* the address of the beginning of the signal stack area. **ss_onstack** field is zero or non-zero depending on whether the process is currently executing on the signal stack or not.

If *nss* is not a **NULL** pointer, **sigstack( )** sets the signal stack state to the value in the **sigstack( )** structure pointed to by *nss*. Note: if **ss_onstack** is non-zero, the system will think that the process is executing on the signal stack. If *nss* is a **NULL** pointer, the signal stack state will be unchanged. If *oss* is not a **NULL** pointer, the current signal stack state is stored in the **sigstack( )** structure pointed to by *oss*.

RETURN VALUES

Upon successful completion, a value of 0 is returned. Otherwise, a value of –1 is returned and **errno** is set to indicate the error.

ERRORS

sigstack( ) will fail and the signal stack context will remain unchanged if one of the following occurs.

EFAULT Either *nss* or *oss* points to memory that is not a valid part of the process address space.

SEE ALSO

sigaltstack(2), **sigvec**(3B), **signal**(3C)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

WARNINGS

Signal stacks are not "grown" automatically, as is done for the normal stack. If the stack overflows unpredictable results may occur.

NAME

sigvec – software signal facilities

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <signal.h>

int sigvec(*sig, nvec, ovec***)**
int *sig***;**
struct sigvec **nvec, *ovec***;**

DESCRIPTION

The system defines a set of signals that may be delivered to a process. Signal delivery resembles the occurrence of a hardware interrupt: the signal is blocked from further occurrence, the current process context is saved, and a new one is built. A process may specify a *handler* to which a signal is delivered, or specify that a signal is to be *blocked* or *ignored*. A process may also specify that a default action is to be taken by the system when a signal occurs. Normally, signal handlers execute on the current stack of the process. This may be changed, on a per-handler basis, so that signals are taken on a special *signal stack*.

All signals have the same *priority*. Signal routines execute with the signal that caused their invocation to be *blocked*, but other signals may yet occur. A global *signal mask* defines the set of signals currently blocked from delivery to a process. The signal mask for a process is initialized from that of its parent (normally 0). It may be changed with a **sigblock( )** or **sigsetmask( )** call, or when a signal is delivered to the process.

A process may also specify a set of *flags* for a signal that affect the delivery of that signal.

When a signal condition arises for a process, the signal is added to a set of signals pending for the process. If the signal is not currently *blocked* by the process then it is delivered to the process. When a signal is delivered, the current state of the process is saved, a new signal mask is calculated (as described below), and the signal handler is invoked. The call to the handler is arranged so that if the signal handling routine returns normally the process will resume execution in the context from before the signal's delivery. If the process wishes to resume in a different context, then it must arrange to restore the previous context itself.

When a signal is delivered to a process a new signal mask is installed for the duration of the process' signal handler (or until a **sigblock( )** or **sigsetmask( )** call is made). This mask is formed by taking the current signal mask, adding the signal to be delivered, and ORing in the signal mask associated with the handler to be invoked.

The action to be taken when the signal is delivered is specified by a **sigvec( )** structure, which includes the following members:

```
void    (*sv_handler)();    /* signal handler */
int     sv_mask;            /* signal mask to apply */
int     sv_flags;           /* see signal options */

#define SV_ONSTACK          /* take signal on signal stack */
#define SV_INTERRUPT        /* do not restart system on signal return */
#define SV_RESETHAND        /* reset handler to SIG_DFL when signal taken */
```

If the **SV_ONSTACK** bit is set in the flags for that signal, the system will deliver the signal to the process on the signal stack specified with **sigstack**(3B) rather than delivering the signal on the current stack.

If *nvec* is not a **NULL** pointer, **sigvec( )** assigns the handler specified by **sv_handler( )**, the mask specified by **sv_mask( )**, and the flags specified by **sv_flags( )** to the specified signal. If *nvec* is a **NULL** pointer, **sigvec( )** does not change the handler, mask, or flags for the specified signal.

The mask specified in *nvec* is not allowed to block **SIGKILL**, **SIGSTOP**, or **SIGCONT**. The system enforces this restriction silently.

If *ovec* is not a **NULL** pointer, the handler, mask, and flags in effect for the signal before the call to **sigvec( )** are returned to the user. A call to **sigvec( )** with *nvec* a **NULL** pointer and *ovec* not a **NULL** pointer can be used to determine the handling information currently in effect for a signal without changing that information.

The following is a list of all signals with names as in the include file **<signal.h>**:

SIGHUP		hangup
SIGINT		interrupt
SIGQUIT	*	quit
SIGILL	*	illegal instruction
SIGTRAP	*	trace trap
SIGABRT	*	abort (generated by **abort**(3C) routine)
SIGEMT	*	emulator trap
SIGFPE	*	arithmetic exception
SIGKILL		kill (cannot be caught, blocked, or ignored)
SIGBUS	*	bus error
SIGSEGV	*	segmentation violation
SIGSYS	*	bad argument to function
SIGPIPE		write on a pipe or other socket with no one to read it
SIGALRM		alarm clock
SIGTERM		software termination signal
SIGURG	•	urgent condition present on socket
SIGSTOP	†	stop (cannot be caught, blocked, or ignored)
SIGTSTP	†	stop signal generated from keyboard
SIGCONT	•	continue after stop (cannot be blocked)
SIGCHLD	•	child status has changed
SIGTTIN	†	background read attempted from control terminal
SIGTTOU	†	background write attempted to control terminal
SIGIO	•	I/O is possible on a descriptor (see **fcntl**(2))
SIGXCPU		cpu time limit exceeded (see **getrlimit**(2))
SIGXFSZ		file size limit exceeded (see **getrlimit**(2))
SIGVTALRM		virtual time alarm (see **setitimer( )** on **getitimer**(2))
SIGPROF		profiling timer alarm (see **setitimer( )** on **getitimer**(2))

SIGWINCH	•	window changed (see **termio**(7I))
SIGLOST	*	resource lost (see **lockd**(1M))
SIGUSR1		user-defined signal 1
SIGUSR2		user-defined signal 2

The starred signals in the list above cause a core image if not caught or ignored.

Once a signal handler is installed, it remains installed until another **sigvec( )** call is made, or an **execve**(2) is performed, unless the **SV_RESETHAND** bit is set in the flags for that signal. In that case, the value of the handler for the caught signal will be set to **SIG_DFL** before entering the signal-catching function, unless the signal is **SIGILL**, **SIGPWR**, or **SIGTRAP**. Also, if this bit is set, the bit for that signal in the signal mask will not be set; unless the signal mask associated with that signal blocks that signal, further occurrences of that signal will not be blocked. The **SV_RESETHAND** flag is not available in 4.2BSD, hence it should not be used if backward compatibility is needed.

The default action for a signal may be reinstated by setting the signal's handler to **SIG_DFL**; this default is termination except for signals marked with • or †. Signals marked with • are discarded if the action is **SIG_DFL**; signals marked with † cause the process to stop. If the process is terminated, a "core image" will be made in the current working directory of the receiving process if the signal is one for which an asterisk appears in the above list (see **core**(4)).

If the handler for that signal is **SIG_IGN**, the signal is subsequently ignored, and pending instances of the signal are discarded.

If a caught signal occurs during certain functions, the call is normally restarted. The call can be forced to terminate prematurely with an **EINTR** error return by setting the **SV_INTERRUPT** bit in the flags for that signal. The **SV_INTERRUPT** flag is not available in 4.2BSD, hence it should not be used if backward compatibility is needed. The affected functions are **read**(2) or **write**(2) on a slow device (such as a terminal or pipe or other socket, but not a file) and during a **wait**(2).

After a **fork**(2) or **vfork**(2) the child inherits all signals, the signal mask, the signal stack, and the restart/interrupt and reset-signal-handler flags.

The **execve**(2) call resets all caught signals to default action and resets all signals to be caught on the user stack. Ignored signals remain ignored; the signal mask remains the same; signals that interrupt functions continue to do so.

The accuracy of *addr* is machine dependent. For example, certain machines may supply an address that is on the same page as the address that caused the fault. If an appropriate *addr* cannot be computed it will be set to **SIG_NOADDR**.

RETURN VALUES

A 0 value indicates that the call succeeded. A –1 return value indicates that an error occurred and **errno** is set to indicate the reason.

ERRORS

sigvec() will fail and no new signal handler will be installed if one of the following occurs:

EFAULT Either *nvec* or *ovec* is not a **NULL** pointer and points to memory that is not a valid part of the process address space.

EINVAL *sig* is not a valid signal number, or, **SIGKILL**, or **SIGSTOP**.

SEE ALSO

intro(2), **exec**(2), **fcntl**(2), **fork**(2), **getitimer**(2), **getrlimit**(2), **ioctl**(2), **kill**(2), **ptrace**(2), **read**(2), **umask**(2), **vfork**(2), **wait**(2), **write**(2), **setjmp**(3C) **sigblock**(3B), **sigstack**(3B), **signal**(3B), **wait**(3B), **signal**(3C), **core**(4), **streamio**(7I), **termio**(7I)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

SIGPOLL is a synonym for **SIGIO**. A **SIGIO** will be issued when a file descriptor corresponding to a STREAMS (see **intro**(2)) file has a "selectable" event pending. Unless that descriptor has been put into asynchronous mode (see **fcntl**(2)), a process may specifically request that this signal be sent using the **I_SETSIG ioctl**(2) call (see **streamio**(7I)). Otherwise, the process will never receive **SIGPOLLs0**.

The handler routine can be declared:

void handler(int *sig***, int** *code***, struct sigcontext** ∗*scp***, char** ∗*addr***);**

Here *sig* is the signal number; *code* is a parameter of certain signals that provides additional detail; *scp* is a pointer to the **sigcontext** structure (defined in **signal.h**), used to restore the context from before the signal; and *addr* is additional address information.

The signals **SIGKILL**, **SIGSTOP**, and **SIGCONT** cannot be ignored.

NAME

sigwaitinfo, sigtimedwait – wait for queued signals

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <signal.h>

int sigwaitinfo(const sigset_t *set, siginfo_t *info);

int sigtimedwait(const sigset_t *set, siginfo_t *info, const struct timespec *timeout);

typedef struct siginfo {
  int           si_signo;   /* signal from signal.h */
  int           si_code;    /* code from above     */
 ...
  int           si_value;
 ...
} siginfo_t;

struct timespec {
  time_t        tv_sec;     /* seconds */
  long          tv_nsec;    /* and nanoseconds */
};
```

MT-LEVEL

Async-Safe

DESCRIPTION

sigwaitinfo( ) and **sigtimedwait( )** select the pending signal from the set specified by *set*. When multiple signals are pending, the lowest numbered one will be selected. The selection order between realtime and non-realtime signals, or between multiple pending non-realtime signals, is unspecified.

If no signal in *set* is pending at the time of the call, **sigwaitinfo( )** suspends the calling process until one or more signals in *set* become pending or until it is interrupted by an unblocked, caught signal. **sigtimedwait( )**, on the other hand, suspends itself for the time interval specified in the **timespec** structure referenced by *timeout*. If the **timespec** structure pointed to by *timeout* is zero-valued, and if none of the signals specified by *set* are pending, then **sigtimedwait( )** returns immediately with the error **EAGAIN.**

If, while **sigwaitinfo( )** or **sigtimedwait( )** is waiting, a signal occurs which is eligible for delivery (i.e., not blocked by the process signal mask), that signal is handled asynchronously and the wait is interrupted.

If *info* is non-**NULL**, the selected signal number is stored in **si_signo**, and the cause of the signal is stored in the **si_code**. If any value is queued to the selected signal, the first such queued value is dequeued and, if *info* is non-**NULL**, the value is stored in the **si_value** member of *info*. The system resource used to queue the signal is released and made available to queue other signals.

If the value of the **si_code** member is **SI_NOINFO,** only the **si_signo** member of **siginfo_t** is meaningful, and the value of all other members is unspecified.

If no further signals are queued for the selected signal, the pending indication for that

signal is reset.

RETURN VALUES

If one of the signals specified by *set* is either pending or generated, **sigwaitinfo( )** or **sigtimedwait( )** returns the selected signal number. Otherwise, the function returns **-1** and sets **errno** to indicate the error condition.

ERRORS

EINTR The wait was interrupted by an unblocked, caught signal.

ENOSYS **sigwaitinfo( )** or **sigtimedwait( )** is not supported by this implementation.

The following errors relate to only **sigtimedwait( ):**

EAGAIN No signal specified by *set* was delivered within the specified timeout period.

EINVAL *timeout* specified a **tv_nsec** value less than 0 or greater than 1,000,000,000.

SEE ALSO

time(2), **sigqueue**(3R), **siginfo**(5), **signal**(5)

NAME

sleep – suspend execution for interval

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

int sleep(*seconds***)**
unsigned *seconds***;**

MT-LEVEL

Async-Signal-Safe

DESCRIPTION

sleep( ) suspends the current process from execution for the number of seconds specified by the argument. The actual suspension time may be up to 1 second less than that requested, because scheduled wakeups occur at fixed 1-second intervals, and may be an arbitrary amount longer because of other activity in the system.

sleep( ) is implemented by setting an interval timer and pausing until it expires. The previous state of this timer is saved and restored. If the sleep time exceeds the time to the expiration of the previous value of the timer, the process sleeps only until the timer would have expired, and the signal which occurs with the expiration of the timer is sent one second later.

SEE ALSO

alarm(2), **getitimer**(2), **longjmp**(3C), **siglongjmp**(3C), **sleep**(3C), **usleep**(3C)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

SIGALRM should *not* be blocked or ignored during a call to **sleep( )**. Only a prior call to **alarm**(2) should generate SIGALRM for the calling process during a call to **sleep( )**. A signal-catching function should *not* interrupt a call to **sleep( )** to call **siglongjmp**(3C) or **longjmp**(3C) to restore an environment saved prior to the **sleep( )** call.

WARNINGS

sleep( ) is slightly incompatible with **alarm**(2). Programs that do not execute for at least one second of clock time between successive calls to **sleep( )** indefinitely delay the alarm signal. Use **sleep**(3C). Each **sleep**(3C) call postpones the alarm signal that would have been sent during the requested sleep period to occur one second later.

NAME

sleep – suspend execution for interval

SYNOPSIS

#include <unistd.h>

unsigned sleep(unsigned *seconds***);**

MT-LEVEL

Safe

DESCRIPTION

The current process is suspended from execution for the number of *seconds* specified by the argument. The actual suspension time may be less than that requested because any caught signal will terminate the **sleep( )** following execution of that signal's catching routine. Also, the suspension time may be longer than requested by an arbitrary amount because of the scheduling of other activity in the system. The value returned by **sleep( )** will be the ''unslept'' amount (the requested time minus the time actually slept) in case the caller had an alarm set to go off earlier than the end of the requested **sleep( )** time, or premature arousal because of another caught signal.

The routine is implemented by setting an alarm signal and pausing until it (or some other signal) occurs. The previous state of the alarm signal is saved and restored. The calling program may have set up an alarm signal before calling **sleep( )**. If the **sleep( )** time exceeds the time until such alarm signal, the process sleeps only until the alarm signal would have occurred. The caller's alarm catch routine is executed just before the **sleep( )** routine returns. But if the **sleep( )** time is less than the time till such alarm, the prior alarm time is reset to go off at the same time it would have without the intervening **sleep( )**.

SEE ALSO

alarm(2), **pause**(2), **signal**(3C)

NOTES

SIGALRM should *not* be blocked or ignored during a call to **sleep( )**. Only a prior call to **alarm**(2) should generate **SIGALRM** for the calling process during a call to **sleep( )**.

In a multithreaded program, only the invoking thread is suspended from execution.

NAME

socket – create an endpoint for communication

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>

int socket(int domain, int type, int protocol);
```

MT-LEVEL

Safe

DESCRIPTION

socket( ) creates an endpoint for communication and returns a descriptor.

The *domain* parameter specifies a communications domain within which communication will take place; this selects the protocol family which should be used. The protocol family generally is the same as the address family for the addresses supplied in later operations on the socket. These families are defined in the include file **<sys/socket.h>**. There must be an entry in the **netconfig**(4) file for at least each protocol family and type required. If *protocol* has been specified, but no exact match for the tuplet family, type, protocol is found, then the first entry containing the specified family and type with zero for protocol will be used. The currently understood formats are:

PF_UNIX UNIX system internal protocols

PF_INET ARPA Internet protocols

The socket has the indicated *type*, which specifies the communication semantics. Currently defined types are:

SOCK_STREAM
SOCK_DGRAM
SOCK_RAW
SOCK_SEQPACKET
SOCK_RDM

A **SOCK_STREAM** type provides sequenced, reliable, two-way connection-based byte streams. An out-of-band data transmission mechanism may be supported. A **SOCK_DGRAM** socket supports datagrams (connectionless, unreliable messages of a fixed (typically small) maximum length). A **SOCK_SEQPACKET** socket may provide a sequenced, reliable, two-way connection-based data transmission path for datagrams of fixed maximum length; a consumer may be required to read an entire packet with each read system call. This facility is protocol specific, and presently not implemented for any protocol family. **SOCK_RAW** sockets provide access to internal network interfaces. The types **SOCK_RAW**, which is available only to the super-user, and **SOCK_RDM**, for which no implementation currently exists, are not described here.

protocol specifies a particular protocol to be used with the socket. Normally only a single protocol exists to support a particular socket type within a given protocol family. However, multiple protocols may exist, in which case a particular protocol must be specified in this manner. The protocol number to use is particular to the “communication domain” in which communication is to take place. If a protocol is specified by the caller, then it

will be packaged into a socket level option request and sent to the underlying protocol layers.

Sockets of type **SOCK_STREAM** are full-duplex byte streams, similar to pipes. A stream socket must be in a *connected* state before any data may be sent or received on it. A connection to another socket is created with a **connect**(3N) call. Once connected, data may be transferred using **read**(2) and **write**(2) calls or some variant of the **send**(3N) and **recv**(3N) calls. When a session has been completed, a **close**(2) may be performed. Out-of-band data may also be transmitted as described on the **send**(3N) manual page and received as described on the **recv**(3N) manual page.

The communications protocols used to implement a **SOCK_STREAM** insure that data is not lost or duplicated. If a piece of data for which the peer protocol has buffer space cannot be successfully transmitted within a reasonable length of time, then the connection is considered broken and calls will indicate an error with –1 returns and with **ETIMEDOUT** as the specific code in the global variable **errno**. The protocols optionally keep sockets "warm" by forcing transmissions roughly every minute in the absence of other activity. An error is then indicated if no response can be elicited on an otherwise idle connection for a extended period (for instance 5 minutes). A **SIGPIPE** signal is raised if a process sends on a broken stream; this causes naive processes, which do not handle the signal, to exit.

SOCK_SEQPACKET sockets employ the same system calls as **SOCK_STREAM** sockets. The only difference is that **read**(2) calls will return only the amount of data requested, and any remaining in the arriving packet will be discarded.

SOCK_DGRAM and **SOCK_RAW** sockets allow datagrams to be sent to correspondents named in **sendto**(3N) calls. Datagrams are generally received with **recvfrom**(3N), which returns the next datagram with its return address.

An **fcntl**(2) call can be used to specify a process group to receive a **SIGURG** signal when the out-of-band data arrives. It may also enable non-blocking I/O and asynchronous notification of I/O events with **SIGIO** signals.

The operation of sockets is controlled by socket level *options*. These options are defined in the file **<sys/socket.h>**. **setsockopt**(3N) and **getsockopt**(3N) are used to set and get options, respectively.

RETURN VALUES

A –1 is returned if an error occurs. Otherwise the return value is a descriptor referencing the socket.

ERRORS

The **socket( )** call fails if:

EACCES	Permission to create a socket of the specified type and/or protocol is denied.
EMFILE	The per-process descriptor table is full.
ENOMEM	Insufficient user memory is available.

ENOSR	There were insufficient STREAMS resources available to complete the operation.
EPROTONOSUPPORT	The protocol type or the specified protocol is not supported within this domain.

SEE ALSO

close(2), **fcntl**(2), **ioctl**(2), **read**(2), **write**(2), **accept**(3N), **bind**(3N), **connect**(3N), **getsockname**(3N), **getsockopt**(3N), **listen**(3N), **recv**(3N), **recvfrom**(3N), **setsockopt**(3N), **send**(3N), **shutdown**(3N), **socketpair**(3N)

NAME

socketpair – create a pair of connected sockets

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

```
#include <sys/types.h>
#include <sys/socket.h>

int socketpair(int domain, int type, int protocol, int sv[2]);
```

MT-LEVEL

Safe

DESCRIPTION

The **socketpair( )** library call creates an unnamed pair of connected sockets in the specified address family *d*, of the specified *type* , and using the optionally specified *protocol*. The descriptors used in referencing the new sockets are returned in *sv*[0] and *sv*[1]. The two sockets are indistinguishable.

RETURN VALUES

socketpair( ) returns **–1** on failure, and **0** on success.

ERRORS

The call succeeds unless:

EAFNOSUPPORT	The specified address family is not supported on this machine.
EMFILE	Too many descriptors are in use by this process.
ENOMEM	There was insufficient user memory for the operation to complete.
ENOSR	There were insufficient STREAMS resources for the operation to complete.
EOPNOSUPPORT	The specified protocol does not support creation of socket pairs.
EPROTONOSUPPORT	The specified protocol is not supported on this machine.

SEE ALSO

pipe(2), **read**(2), **write**(2)

NOTES

This call is currently implemented only for the **AF_UNIX** address family.

NAME

spray – scatter data in order to test the network

SYNOPSIS

cc [*flag* ...] *file* ... **–lsocket –lnsl** [*library* ...]

#include <rpcsvc/spray.h>

bool_t xdr_sprayarr(XDR ∗*xdrs***, sprayarr** ∗*objp***);**

bool_t xdr_spraycumul(XDR ∗*xdrs***, spraycumul** ∗*objp***);**

MT-LEVEL

Unsafe

DESCRIPTION

The spray program sends packets to a given machine to test communications with that machine.

The spray program is not a C function interface, per se, but can be accessed using the generic remote procedure calling interface **clnt_call( )** (see **rpc_clnt_calls**(3N)). The program sends a packet to the called host. The host acknowledges receipt of the packet. The program counts the number of acknowledgments and can return that count.

The spray program currently supports the following procedures, which should be called in the order given:

SPRAYPROC_CLEAR
This procedure clears the counter.

SPRAYPROC_SPRAY
This procedure sends the packet.

SPRAYPROC_GET
This procedure returns the count and the amount of time since the last **SPRAYPROC_CLEAR**.

EXAMPLES

The following code fragment demonstrates how the spray program is used:

```
#include <rpc/rpc.h>
#include <rpcsvc/spray.h>

. . .
        spraycumul      spray_result;
        sprayarr        spray_data;
        char            buf[100];               /* arbitrary data */
        int             loop = 1000;
        CLIENT          *clnt;
        struct timeval timeout0 = {0, 0};
        struct timeval timeout25 = {25, 0};

        spray_data.sprayarr_len = (u_int)100;
        spray_data.sprayarr_val = buf;

        clnt = clnt_create("somehost", SPRAYPROG, SPRAYVERS, "netpath");
        if (clnt == (CLIENT *)NULL) {
```

```
        /* handle this error */
    }
    if (clnt_call(clnt, SPRAYPROC_CLEAR,
        xdr_void, NULL, xdr_void, NULL, timeout25)) {
            /* handle this error */
    }
    while (loop-- > 0) {
        if (clnt_call(clnt, SPRAYPROC_SPRAY,
            xdr_sprayarr, &spray_data, xdr_void, NULL, timeout0)) {
                /* handle this error */
        }
    }
    if (clnt_call(clnt, SPRAYPROC_GET,
        xdr_void, NULL, xdr_spraycumul, &spray_result, timeout25)) {
            /* handle this error */
    }
    printf("Acknowledged %ld of 1000 packets in %d secs %d usecs\n",
        spray_result.counter,
        spray_result.clock.sec,
        spray_result.clock.usec);
```

SEE ALSO

spray(1M), **rpc_clnt_calls**(3N)

NOTES

This interface is unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

A spray program is not useful as a networking benchmark as it uses unreliable connectionless transports, (udp for example). It can report a large number of packets dropped when the drops were caused by the program sending packets faster than they can be buffered locally (before the packets get to the network medium).

NAME

sqrt, cbrt – square root, cube root

SYNOPSIS

cc [*flag* ...] *file* ... **–lm** [*library* ...]

#include <math.h>

double sqrt(double *x***);**

double cbrt(double *x***);**

MT-LEVEL

MT-Safe

DESCRIPTION

sqrt(*x***)** returns the square root of *x*, correctly rounded according to ANSI/IEEE 754-1985.

cbrt(*x***)** returns the cube root of *x*. **cbrt( )** is accurate to within 0.7 *ulp*s.

RETURN VALUES

For exceptional cases, **matherr**(3M) tabulates the values to be returned as dictated by various Standards.

SEE ALSO

matherr(3M)

NAME

ssignal, gsignal – software signals

SYNOPSIS

```
#include <signal.h>
void (*ssignal (int sig, int (*action) (int))) (int);
int gsignal(int sig);
```

MT-LEVEL

Unsafe

DESCRIPTION

ssignal( ) and **gsignal( )** implement a software facility similar to **signal**(3C). This facility is made available to users for their own purposes.

Software signals made available to users are associated with integers in the inclusive range 1 through 17. A call to **ssignal( )** associates a procedure, *action*, with the software signal *sig*; the software signal, *sig*, is raised by a call to **gsignal( )**. Raising a software signal causes the action established for that signal to be *taken*.

The first argument to **ssignal( )** is a number identifying the type of signal for which an action is to be established. The second argument defines the action; it is either the name of a (user-defined) *action function* or one of the manifest constants **SIG_DFL** (default) or **SIG_IGN** (ignore). **ssignal( )** returns the action previously established for that signal type; if no action has been established or the signal number is illegal, **ssignal( )** returns **SIG_DFL**.

gsignal( ) raises the signal identified by its argument, *sig*:

If an action function has been established for *sig*, then that action is reset to **SIG_DFL** and the action function is entered with argument *sig*. **gsignal( )** returns the value returned to it by the action function.

If the action for *sig* is **SIG_IGN**, **gsignal( )** returns the value 1 and takes no other action.

If the action for *sig* is **SIG_DFL**, **gsignal( )** returns the value 0 and takes no other action.

If *sig* has an illegal value or no action was ever specified for *sig*, **gsignal( )** returns the value 0 and takes no other action.

SEE ALSO

raise(3C), **signal**(3C)

NAME

stdio – standard buffered input / output package

SYNOPSIS

#include <stdio.h>

extern **FILE** ∗**stdin;**

extern **FILE** ∗**stdout;**

extern **FILE** ∗**stderr;**

DESCRIPTION

The functions described in the entries of section 3S of this manual constitute an efficient, user-level I / O buffering scheme. The in-line macros **getc( )** and **putc( )** handle characters quickly. The macros **getchar( )** and **putchar( )**, and the higher-level routines **fgetc( )**, **fgets( )**, **fprintf( )**, **fputc( )**, **fputs( )**, **fread( )**, **fscanf( )**, **fwrite( )**, **gets( )**, **getw( )**, **printf( )**, **puts( )**, **putw( )**, and **scanf( )** all use or act as if they use **getc( )** and **putc( )**; they can be freely intermixed.

A file with associated buffering is called a *stream* (see **intro**(3)) and is declared to be a pointer to a defined type **FILE**. **fopen( )** creates certain descriptive data for a stream and returns a pointer to designate the stream in all further transactions. Normally, there are three open streams with constant pointers declared in the **<stdio.h>** header and associated with the standard open files:

stdin	standard input file
stdout	standard output file
stderr	standard error file

The following symbolic values in **<unistd.h>** define the file descriptors that will be associated with the C-language *stdin*, *stdout* and *stderr* when the application is started:

STDIN_FILENO	Standard input value	0	**stdin**
STDOUT_FILENO	Standard output value	1	**stdout**
STDERR_FILENO	Standard error value	2	**stderr**

The constant **NULL** designates a null pointer.

The integer-constant **EOF** is returned upon end-of-file or error by most integer functions that deal with streams (see the individual descriptions for details).

The integer constant **BUFSIZ** specifies the size of the buffers used by the particular implementation.

The integer constant **FILENAME_MAX** specifies the number of bytes needed to hold the longest pathname of a file allowed by the implementation. If the system does not impose a maximum limit, this value is the recommended size for a buffer intended to hold a file's pathname.

The integer constant **FOPEN_MAX** specifies the minimum number of files that the implementation guarantees can be open simultaneously. Note that no more than 255 files may be opened using **fopen( )**, and only file descriptors 0 through 255 can be used in a stream.

The functions and constants mentioned in the entries of section 3S of this manual are declared in that header and need no further declaration. The constants and the following ''functions'' are implemented as macros (redeclaration of these names is perilous): **getc( )**,

getchar(), **putc()**, **putchar()**, **ferror()**, **feof()**, **clearerr()**, and **fileno()**. There are also function versions of **getc()**, **getchar()**, **putc()**, **putchar()**, **ferror**, **feof()**, **clearerr()**, and **fileno()**.

Output streams, with the exception of the standard error stream **stderr**, are by default buffered if the output refers to a file and line-buffered if the output refers to a terminal. The standard error output stream **stderr** is by default unbuffered, but use of **freopen()** (see **fopen**(3S)) will cause it to become buffered or line-buffered. When an output stream is unbuffered, information is queued for writing on the destination file or terminal as soon as written; when it is buffered, many characters are saved up and written as a block. When it is line-buffered, each line of output is queued for writing on the destination terminal as soon as the line is completed (that is, as soon as a new-line character is written or terminal input is requested). **setbuf()** or **setvbuf()** (both described in **setbuf**(3S)) may be used to change the stream's buffering strategy.

Interactions of Other FILE-Type C Functions

A single open file description can be accessed both through streams and through file descriptors. Either a file descriptor or a stream will be called a *handle* on the open file description to which it refers; an open file description may have several handles.

Handles can be created or destroyed by user action without affecting the underlying open file description. Some of the ways to create them include **fcntl(), dup(), fdopen(), fileno(),** and **fork()** (which duplicates existing ones into new processes). They can be destroyed by at least **fclose(), close(),** and the **exec** functions (which close some file descriptors and destroy streams).

A file descriptor that is never used in an operation, that could affect the file offset (for example **read**(), **write**(), or **lseek**()) is not considered a handle in this discussion, but could give rise to one (as a consequence of **fdopen**(), **dup**(), or **fork**(), for example). This exception does include the file descriptor underlying a stream, whether created with **fopen**() or **fdopen**(), as long as it is not used directly by the application to affect the file offset. (The **read**() and **write**() functions implicitly affect the file offset; **lseek**() explicitly affects it.)

If two or more handles are used, and any one of them is a stream, their actions shall be coordinated as described below. If this is not done, the result is undefined.

A handle that is a stream is considered to be closed when either an **fclose**() or **freopen**() is executed on it (the result of **freopen**() is a new stream for this discussion, which cannot be a handle on the same open file description as its previous value) or when the process owning that stream terminates the **exit**() or **abort().** A file descriptor is closed by **close**(), **_exit**(), or by one of the **exec** functions when FD_CLOEXEC is set on that file descriptor.

For a handle to become the active handle, the actions below must be performed between the last other user of the first handle (the current active handle) and the first other user of the second handle (the future active handle). The second handle then becomes the active handle. All activity by the application affecting the file offset on the first handle shall be suspended until it again becomes the active handle. (If a stream function has as an underlying function that affects the file offset, the stream function will be considered to affect the file offset. The underlying functions are described below.)

The handles need not be in the same process for these rules to apply. Note that after a **fork**( ), two handles exist where one existed before. The application shall assure that, if both handles will ever be accessed, that they will both be in a state where the other could become the active handle first. The application shall prepare for a **fork**( ) exactly as if it were a change of active handle. (If the only action performed by one of the processes is one of the **exec** functions or **_exit**( ), the handle is never accessed in that process.)

(1) For the first handle, the first applicable condition below shall apply. After the actions required below are taken, the handle may be closed if it is still open.

- (a) If it is a file descriptor, no action is required.
- (b) If the only further action to be performed on any handle to this open file description is to close it, no action need be taken.
- (c) If it is a stream that is unbuffered, no action need be taken.
- (d) If it is a stream that is line-buffered and the last character written to the stream was a newline (that is, as if a **putc('\n')** was the most recent operation on that stream), no action need be taken.
- (e) If it is a stream that is open for writing or append (but not also open for reading), either an **fflush**( ) shall occur or the stream shall be closed.
- (f) If the stream is open for reading and it is at the end of the file (**feof**( ) is true), no action need be taken.
- (g) If the stream is open with a mode that allows reading and the underlying open file description refers to a device that is capable of seeking, either an **fflush**( ) shall occur or the stream shall be closed.
- (h) Otherwise, the result is undefined.

(2) For the second handle: if any previous active handle has called a function that explicitly changed the file offset, except as required above for the first handle, the application shall perform an **lseek**( ) or an **fseek**( ) (as appropriate to the type of the handle) to an appropriate location.

(3) If the active handle ceases to be accessible before the requirements on the first handle above have been met, the state of the open file description becomes undefined. This might occur, for example, during a **fork**( ) or an **_exit**( ).

(4) The **exec** functions shall be considered to make inaccessible all streams that are open at the time they are called, independent of what streams or file descriptors may be available to the new process image.

(5) Implementation shall assure that an application, even one consisting of several processes, shall yield correct results (no data is lost or duplicated when writing, all data is written in order, except as requested by seeks) when the rules above are followed, regardless of the sequence of handles used. If the rules above are not followed, the result is unspecified. When these rules are followed, it is implementation defined whether, and under what conditions, all input is seen exactly once.

Use of stdio in Multithreaded Applications

All the stdio functions are safe unless they have the _unlocked suffix. Each **file** pointer has its own lock to guarantee that only one thread can access it. In the case that output needs to be synchronized, the lock for the **FILE** pointer can be acquired before performing a series of stdio operations. For example:

```
FILE iop;
.
.
flockfile(iop);
fprintf(iop, "hello ");
fprintf(iop, "world0);
fputc(iop, 'a');
funlockfile(iop);
```

will print everything out together, blocking other threads that might want to write to the same file between fprintf's.

An unlocked interface is available in case performace is an issue. For example:

```
flockfile(iop);
while (!feof(iop)) {
        *c++ = getc_unlocked(iop);
}
funlockfile(iop);
```

RETURN VALUES

Invalid stream pointers usually cause grave disorder, possibly including program termination. Individual function descriptions describe the possible error conditions.

SEE ALSO

close(2), **lseek**(2), **open**(2), **pipe**(2), **read**(2), **write**(2), **ctermid**(3S), **cuserid**(3S), **fclose**(3S), **ferror**(3S), **fopen**(3S), **fread**(3S), **fseek**(3S), **flockfile**(3S), **getc**(3S), **gets**(3S), **popen**(3S), **printf**(3S), **putc**(3S), **puts**(3S), **scanf**(3S), **setbuf**(3S), **system**(3S), **tmpfile**(3S), **tmpnam**(3S), **ungetc**(3S)

NAME

stdipc, ftok – standard interprocess communication package

SYNOPSIS

```
#include <sys/types.h>
#include <sys/ipc.h>

key_t ftok(const char *path, int id);
```

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

Certain interprocess communication facilities require the user to supply a key to be used by the **msgget**(2), **semget**(2), and **shmget**(2) functions to obtain interprocess communication identifiers. One suggested method for forming a key is to use the **ftok( )** subroutine described below. Another way to compose keys is to include the project ID in the most significant byte and to use the remaining portion as a sequence number. There are many other ways to form keys, but it is necessary for each system to define standards for forming them. If some standard is not adhered to, it will be possible for unrelated processes to unintentionally interfere with each other's operation. It is still possible to interfere intentionally. Therefore, it is strongly suggested that the most significant byte of a key in some sense refer to a project so that keys do not conflict across a given system.

ftok( ) returns a key based on *path* and *id* that is usable in subsequent **msgget( )**, **semget( )**, and **shmget( )** functions. *path* must be the path name of an existing file that is accessible to the process. *id* is a character that uniquely identifies a project. Note that **ftok( )** returns the same key for linked files when called with the same *id* and that it returns different keys when called with the same file name but different *ids*.

RETURN VALUES

ftok( ) returns **(key_t) –1** if *path* does not exist or if it is not accessible to the process.

SEE ALSO

intro(2), **msgget**(2), **semget**(2), **shmget**(2)

NOTES

If the file whose *path* is passed to **ftok( )** is removed when keys still refer to the file, future calls to **ftok( )** with the same *path* and *id* return an error. If the same file is recreated, then **ftok( )** is likely to return a different key from the original call.

ftok( ) is MT-Safe in mutli-thread applications.

NAME

str2sig, sig2str – translation between signal name and signal number

SYNOPSIS

```
#include <signal.h>

int str2sig(const char *str, int *signum);
int sig2str(int signum, char *str);
```

DESCRIPTION

The **str2sig()** function translates the signal name *str* to a signal number, and stores that result in the location referenced by *signum*. The name in *str* can be either the symbol for that signal, without the "SIG" prefix, or a decimal number. All the signal symbols defined in **<sys/signal.h>** are recognized. This means that both "CLD" and "CHLD" are recognized and return the same signal number, as do both "POLL" and "IO". For access to the signals in the range **SIGRTMIN** to **SIGRTMAX**, the first four signals match the strings "RTMIN", "RTMIN+1", "RTMIN+2", and "RTMIN+3" and the last four match the strings "RTMAX-3", "RTMAX-2", "RTMAX-1", and "RTMAX".

The **sig2str()** function translates the signal number *signum* to the symbol for that signal, without the "SIG" prefix, and stores that symbol at the location specified by *str*. The storage referenced by *str* should be large enough to hold the symbol and a terminating null byte. The symbol **SIG2STR_MAX** defined by **<signal.h>** gives the maximum size in bytes required.

RETURN VALUES

The **str2sig()** function returns **0** if it recognizes the signal name specified in *str*; otherwise, it returns **–1**.

The **sig2str()** function returns **0** if the value *signum* corresponds to a valid signal number; otherwise, it returns **–1**.

EXAMPLES

```
int i;
char buf[STR2SIG_MAX];      /* storage for symbol */

str2sig("KILL", &i);        /* stores 9 in i */
str2sig("9", &i);           /* stores 9 in i */
sig2str(SIGKILL, buf);      /* stores "KILL" in buf */
sig2str(9, buf);            /* stores "KILL" in buf */
```

SEE ALSO

kill(1), **strsignal**(3C)

NAME

strccpy, streadd, strcadd, strecpy – copy strings, compressing or expanding escape codes

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

char *__strccpy(__char** *__output__, __const char** *__input__);__

char *__strcadd(__char** *__output__, __const char** *__input__);__

char *__strecpy(__char** *__output__, __const char** *__input__, __const char** *__exceptions__);__

char *__streadd(__char** *__output__, __const char** *__input__, __const char** *__exceptions__);__

MT-LEVEL

MT-Safe

DESCRIPTION

strccpy( ) copies the *input* string, up to a null byte, to the *output* string, compressing the C-language escape sequences (for example, **\n**, **\001**) to the equivalent character. A null byte is appended to the output. The *output* argument must point to a space big enough to accommodate the result. If it is as big as the space pointed to by *input* it is guaranteed to be big enough. **strccpy( )** returns the *output* argument.

strcadd( ) is identical to **strccpy( )**, except that it returns the pointer to the null byte that terminates the output.

strecpy( ) copies the *input* string, up to a null byte, to the *output* string, expanding non-graphic characters to their equivalent C-language escape sequences (for example, **\n**, **\001**). The *output* argument must point to a space big enough to accommodate the result; four times the space pointed to by *input* is guaranteed to be big enough (each character could become \ and 3 digits). Characters in the *exceptions* string are not expanded. The *exceptions* argument may be zero, meaning all non-graphic characters are expanded. **strecpy( )** returns the *output* argument.

streadd( ) is identical to **strecpy( )**, except that it returns the pointer to the null byte that terminates the output.

EXAMPLES

```
/* expand all but newline and tab */
strecpy( output, input, "\n\t" );

/* concatenate and compress several strings */
cp = strcadd( output, input1 );
cp = strcadd( cp, input2 );
cp = strcadd( cp, input3 );
```

SEE ALSO

string(3C), **strfind**(3G)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

strcoll – string collation

SYNOPSIS

```
#include <string.h>
int strcoll(const char *s1, const char *s2);
```

MT-LEVEL

Safe with exceptions

DESCRIPTION

strcoll( ) returns an integer greater than, equal to, or less than zero in direct correlation to whether string *s1* is greater than, equal to, or less than the string *s2*. The comparison is based on strings interpreted as appropriate to the program's locale for category **LC_COLLATE** (see **setlocale**(3C)).

Both **strcoll( )** and **strxfrm( )** provide for locale-specific string sorting. **strcoll( )** is intended for applications in which the number of comparisons per string is small. When strings are to be compared a number of times, **strxfrm( )** is a more appropriate function because the transformation process occurs only once.

FILES

/usr/lib/*locale***/LC_COLLATE** **LC_COLLATE** database for *locale*

SEE ALSO

colltbl(1M), **setlocale**(3C), **string**(3C), **strxfrm**(3C), **wsxfrm**(3I), **environ**(5)

NOTES

strcoll can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

strerror – get error message string

SYNOPSIS

```
#include <string.h>
char *strerror(int errnum);
```

MT-LEVEL

Safe

DESCRIPTION

strerror( ) maps the error number in *errnum* to an error message string, and returns a pointer to that string. **strerror( )** uses the same set of error messages as **perror( )**. The returned string should not be overwritten.

ERRORS

strerror returns **NULL** if *errnum* is out-of-range.

SEE ALSO

gettext(3I), **perror**(3C), **setlocale**(3C)

NOTES

If the application is linked with **–lintl**, then messages returned from this function are in the native language specified by the **LC_MESSAGES** locale category; see **setlocale**(3C).

NAME

strfind, strrspn, strtrns, str – string manipulations

SYNOPSIS

cc [*flag* ...] *file* ... **–lgen** [*library* ...]

#include <libgen.h>

int strfind(const char **as1*, **const char** **as2*);

char ***__**strrspn(const char** **string*, **const char** **tc*);

char *** **strtrns(const char** **string*, **const char** **old*, **const char** **new*, **char** **result*);

MT-LEVEL

MT-Safe

DESCRIPTION

strfind( ) returns the offset of the first occurrence of the second string, *as2*, if it is a substring of string *as1*. If the second string is not a substring of the first string **strfind( )** returns **–1.**

strrspn( ) returns a pointer to the first character in the string that is not one of the characters in *tc*.

strtrns( ) transforms *string* and copies it into *result*. Any character that appears in *old* is replaced with the character in the same position in *new*. The *new* result is returned.

EXAMPLES

```
/* find offset to substring "hello" within as1 */
i = strfind(as1, "hello");

/* trim junk from end of string */
s2 = strrspn(s1, "*?#$%");
*s2 = '\0';

/* transform lower case to upper case */
a1[] = "abcdefghijklmnopqrstuvwxyz";
a2[] = "ABCDEFGHIJKLMNOPQRSTUVWXYZ";
s2 = strtrns(s1, a1, a2, s2);
```

SEE ALSO

string(3C)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

NAME

strfmon – convert monetary value to string

SYNOPSIS

```
#include <monetary.h>
ssize_t strfmon(char *s, size_t maxsize, const char *format, ...);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **strfmon()** function places characters into the array pointed to by *s* as controlled by the string pointed to by *format*. No more than *maxsize* bytes are placed into the array.

The format is a character string that contains two types of objects: plain characters, which are simply copied to the output stream, and conversion specifications, each of which results in the fetching of zero or more arguments which are converted and formatted. The results are undefined if there are insufficient arguments for the format. If the format is exhausted while arguments remain, the excess arguments are simply ignored.

A conversion specification consists of the following sequence:

- a % character
- optional flags
- optional field width
- optional left precision
- optional right precision
- a required conversion character that determines the conversion to be performed.

Flags

One or more of the following optional flags can be specified to control the conversion:

=f An = followed by a single character *f* which is used as the numeric fill character. The fill character must be representable in a single byte in order to work with precision and width counts. The default numeric fill character is the space character. This flag does not affect field width filling which always uses the space character. This flag is ignored unless a left precision (see below) is specified.

^ Do not format the currency amount with grouping characters. The default is to insert the grouping characters if defined for the current locale.

+ or (Specify the style of representing positive and negative currency amounts. Only one of '+' or '(' may be specified. If '+' is specified, the locale's equivalent of + and '–' are used (for example, in the U.S.A.: the empty string if positive and '–' if negative). If '(' is specified, negative amounts are enclosed within parentheses. If neither flag is specified, the '+' style is used.

! Suppress the currency symbol from the output conversion.

– Specify the alignment. If this flag is present all fields are left-justified (padded to the right) rather than right-justified.

Field Width

w A decimal digit string *w* specifying a minimum field width in bytes in which the result of the conversion is right-justified (or left-justified if the flag '–' is specified). The default is zero.

Left Precision

#n A '#' followed by a decimal digit string *n* specifying a maximum number of digits expected to be formatted to the left of the radix character. This option can be used to keep the formatted output from multiple calls to the **strfmon( )** aligned in the same columns. It can also be used to fill unused positions with a special character as in **$***123.45**. This option causes an amount to be formatted as if it has the number of digits specified by *n*. If more than *n* digit positions are required, this conversion specification is ignored. Digit positions in excess of those actually required are filled with the numeric fill character (see the *=f* flag above).

If grouping has not been suppressed with the '^' flag, and it is defined for the current locale, grouping separators are inserted before the fill characters (if any) are added. Grouping sparators are not applied to fill characters even if the fill character is a digit.

To ensure alignment, any characters appearing before or after the number in the formatted output such as currency or sign symbols are padded as necessary with space characters to make their positive and negative formats an equal length.

Right Precision

.p A period followed by a decimal digit string *p* specifying the number of digits after the radix character. If the value of the right precision *p* is zero, no radix character appears. If a right precision is not included, a default specified by the current locale is used. The amount being formatted is rounded to the specified number of digits prior to formatting.

Conversion Characters

The conversion characters and their meanings are:

i The **double** argument is formatted according to the locale's international currency format (for example, in the U.S.A.: USD 1,234.56**).**

n The **double** argument is formatted according to the locale's national currency format (for example, in the U.S.A.: $1,234.56).

% Convert to a %; no argument is converted. The entire conversion specification must be %%.

Locale Information

The **LC_MONETARY** category of the program's locale affects the behaviour of this function including the monetary radix character (which may be different from the numeric radix character affected by the **LC_NUMERIC** category), the grouping separator, the currency symbols and formats. The international currency symbol should be conformant with the ISO 4217: 1987 standard.

RETURN VALUES

If the total number of resulting bytes (including the terminating null byte) is not more than *maxsize*, **strfmon( )** returns the number of bytes placed into the array pointed to by *s*, not including the terminating null byte. Otherwise, **–1** is returned, the contents of the array are indeterminate, and **errno** is set to indicate the error.

ERRORS

strfmon( ) will fail if:

ENOSYS The function is not supported.

E2BIG Conversion stopped due to lack of space in the buffer.

EXAMPLES

Given a locale for the U.S.A. and the values 123.45, –123.45, and 3456.781:

Conversion Specification	Output	Comments
%n	$123.45 -$123.45 $3,456.78	default formatting
%11n	$123.45 -$123.45 $3,456.78	right align within an 11 character field
%#5n	$ 123.45 -$ 123.45 $ 3,456.78	aligned columns for values up to 99,999
%=*#5n	$***123.45 -$***123.45 $*3,456.78	specify a fill character
%=0#5n	$000123.45 -$000123.45 $03,456.78	fill characters do not use grouping even if the fill character is a digit
%^#5n	$ 123.45 -$ 123.45 $ 3456.78	disable the grouping separator
%^#5.0n	$ 123 -$ 123 $ 3457	round off to whole units
%^#5.4n	$ 123.4500 -$ 123.4500 $ 3456.7810	increase the precision
%(#5n	123.45 ($ 123.45) $ 3,456.78	use an alternative pos/neg style
%!(#5n	123.45 (123.45) 3,456.78	disable the currency symbol

SEE ALSO **localeconv**(3C)

NOTES This interface is expected to be mandatory in a future issue of this document.

Lower-case conversion characters are reserved for future use and upper-case for implementation-dependent use.

NAME

strftime, cftime, ascftime – convert date and time to string

SYNOPSIS

```
#include <time.h>

size_t strftime(const char *s, size_t maxsize, const char *format,
        const struct tm *timeptr);

int cftime(char *s, char *format, const time_t *clock);

int ascftime(char *s, const char *format, const struct tm *timeptr);
```

MT-LEVEL

MT-Safe

DESCRIPTION

strftime( ), **ascftime( )**, and **cftime( )** place bytes into the array pointed to by *s* as controlled by the string pointed to by *format*. The *format* string consists of zero or more conversion specifications and ordinary characters. A conversion specification consists of a '**%**' (percent) character and one or two terminating conversion characters that determine the conversion specification's behavior. All ordinary characters (including the terminating null byte) are copied unchanged into the array pointed to by *s*. If copying takes place between objects that overlap, the behavior is undefined. For **strftime( )**, no more than *maxsize* bytes are placed into the array.

If *format* is **(char *)0**, then the locale's default format is used. For **strftime( )** the default format is the same as **%c**; for **cftime( )** and **ascftime( )** the default format is the same as **%C**. **cftime( )** and **ascftime( )** first try to use the value of the environment variable **CFTIME**, and if that is undefined or empty, the default format is used.

Each conversion specification is replaced by appropriate characters as described in the following list. The appropriate characters are determined by the **LC_TIME** category of the program's locale and by the values contained in the structure pointed to by *timeptr* for **strftime( )** and **ascftime( )**, and by the time represented by *clock* for **cftime( )**.

	%%	same as %
	%a	locale's abbreviated weekday name
	%A	locale's full weekday name
	%b	locale's abbreviated month name
	%B	locale's full month name
	%c	locale's appropriate date and time representation
Solaris	**%C**	locale's date and time representation as produced by **date**(1)
XPG4	**%C**	century number (the year divided by 100 and truncated to an integer as a decimal number [1,99]); single digits are preceded by **0**
	%d	day of month [1,31]; single digits are preceded by **0**
	%D	date as **%m**/**%d**/**%y**
	%e	day of month [1,31]; single digits are preceded by a space
	%h	locale's abbreviated month name
	%H	hour (24-hour clock) [0,23]; single digits are preceded by **0**
	%I	hour (12-hour clock) [1,12]; single digits are preceded by **0**
	%j	day number of year [1,366]; single digits are preceded by **0**
	%k	hour (24-hour clock) [0,23]; single digits are preceded by a blank

%l hour (12-hour clock) [1,12]; single digits are preceded by a blank
%m month number [1,12]; single digits are preceded by **0**
%M minute [00,59]; leading zero is permitted but not required
%n insert a newline
%p locale's equivalent of either a.m. or p.m.
%r appropriate time representation in 12-hour clock format with **%p**
%R time as **%H**:**%M**
%S seconds [00,61]
%t insert a tab
%T time as **%H**:**%M**:**%S**
%u weekday as a decimal number [1,7], with 1 representing Sunday
%U week number of year as a decimal number [00,53], with Sunday as the first day of week 1
%V week number of the year as a decimal number [01,53], with Monday as the first day of the week. If the week containing 1 January has four or more days in the new year, then it is considered week 1; otherwise, it is week 53 of the previous year, and the next week is week 1.
%w weekday as a decimal number [0,6], with 0 representing Sunday
%W week number of year as a decimal number [00,53], with Monday as the first day of week 1
%x locale's appropriate date representation
%X locale's appropriate time representation
%y year within century [00,99]
%Y year, including the century (for example 1993)
%Z time zone name or abbreviation, or no bytes if no time zone information exists

If a conversion specification does not correspond to any of the above or to any of the modified conversion specifications listed below, the behavior is undefined and **0** is returned.

The difference between **%U** and **%W** (and also between modified conversion specifications **%OU** and **%OW**) lies in which day is counted as the first of the week. Week number 1 is the first week in January starting with a Sunday for **%U** or a Monday for **%W**. Week number 0 contains those days before the first Sunday or Monday in January for **%U** and **%W**, respectively.

Modified Conversion Specifications

Some conversion specifications can be modified by the **E** and **O** modifiers to indicate that an alternate format or specification should be used rather than the one normally used by the unmodified conversion specification. If the alternate format or specification does not exist in the current locale, the behavior will be as if the unmodified specification were used.

%Ec locale's alternate appropriate date and time representation
%EC name of the base year (period) in the locale's alternate representation
%Ex locale's alternate date representation
%EX locale's alternate time representation
%Ey offset from **%EC** (year only) in the locale's alternate representation

%EY	full alternate year representation
%Od	day of the month using the locale's alternate numeric symbols
%Oe	same as **%Od**
%OH	hour (24-hour clock) using the locale's alternate numeric symbols
%OI	hour (12-hour clock) using the locale's alternate numeric symbols
%Om	month using the locale's alternate numeric symbols
%OM	minutes using the locale's alternate numeric symbols
%OS	seconds using the locale's alternate numeric symbols
%OU	week number of the year (Sunday as the first day of the week) using the locale's alternate numeric symbols
%Ow	number of the weekday (Sunday=0) using the locale's alternate numeric symbols
%OW	week number of the year (Monday as the first day of the week) using the locale's alternate numeric symbols
%Oy	year (offset from **%C**) in the locale's alternate representation and using the locale's alternate numeric symbols

Selecting the Output Language

By default, the output of **strftime( )**, **cftime( )**, and **ascftime( )** appear in U.S. English. The user can request that the output of **strftime( )**, **cftime( )**, or **ascftime( )** be in a specific language by setting the **LC_TIME** category using **setlocale( )**.

Time Zone

Local time zone information is used as though **tzset**(3C) were called.

RETURN VALUES

strftime( ), **cftime( )**, and **ascftime( )** return the number of characters placed into the array pointed to by *s*, not including the terminating null character. If the total number of resulting characters including the terminating null character is more than *maxsize*, **strftime()** returns **0** and the contents of the array are indeterminate.

EXAMPLES

The following example illustrates the use of **strftime( )**. It shows what the string in *str* would look like if the structure pointed to by *tmptr* contains the values corresponding to Thursday, August 28, 1986 at 12:44:36.

```
strftime (str, strsize, "%A %b %d %j", tmptr)
```

This results in *str* containing "Thursday Aug 28 240".

FILES

/usr/lib/locale/*locale***/LC_TIME/time** locale specific date and time information

SEE ALSO

date(1), **ctime**(3C), **mktime**(3C), **setlocale**(3C), **strptime**(3C), **tzset**(3C), **TIMEZONE**(4), **strftime**(4), **environ**(5), **xpg4**(5)

NOTES

The range of values for **%S** is [00,61] rather than [00,59] to allow for the occasional leap second and even more occasional double leap second.

NAME

string, strcasecmp, strncasecmp, strcat, strncat, strchr, strrchr, strcmp, strncmp, strcpy, strncpy, strcspn, strspn, strdup, strlen, strpbrk, strstr, strtok, strtok_r – string operations

SYNOPSIS

```
#include <string.h>
int strcasecmp(const char *s1, const char *s2);
int strncasecmp(const char *s1, const char *s2, int n);
char *strcat(char *dst, const char *src);
char *strncat(char *dst, const char *src, size_t n);
char *strchr(const char *s, int c);
char *strrchr(const char *s, int c);
int strcmp(const char *s1, const char *s2);
int strncmp(const char *s1, const char *s2, size_t n);
char *strcpy(char *dst, const char *src);
char *strncpy(char *dst, const char *src, size_t n);
size_t strcspn(const char *s1, const char *s2);
size_t strspn(const char *s1, const char *s2);
char *strdup(const char *s1);
size_t strlen(const char *s);
char *strpbrk(const char *s1, const char *s2);
char *strstr(const char *s1, const char *s2);
char *strtok(char *s1, const char *s2);
char *strtok_r(char *s1, const char *s2, char **lasts);
```

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

string, strcasecmp, strncasecmp, strcat, strncat, strchr, strrchr, strcmp, strncmp, strcpy, strncpy, strcspn, strspn, strdup, strlen, strpbrk, strstr, strtok, strtok_r – string operations The arguments *s*, *s1*, *s2*, *src*, and *dst* point to strings (arrays of characters terminated by a null character). The functions **strcat( )**, **strncat( )**, **strcpy( )**, **strncpy( )**, **strtok( )**, and **strtok_r( )** all alter their first argument. These functions do not check for overflow of the array pointed to by the first argument.

strcasecmp( ) and **strncasecmp( )** are case-insensitive versions of **strcmp( )** and **strncmp( )** respectively, described below. **strcasecmp( )** and **strncasecmp( )** assume the ASCII character set and ignore differences in case when comparing lower and upper case characters.

strcat( ) appends a copy of string *src*, including the terminating null character, to the end of string *dst*. **strncat( )** appends at most *n* characters. Each returns a pointer to the null-terminated result. The initial character of *src* overrides the null character at the end of *dst*.

strchr( ) returns a pointer to the first occurrence of *c* (converted to a **char**) in string *s*, or a null pointer if *c* does not occur in the string. **strrchr( )** returns a pointer to the last occurrence of *c*. The null character terminating a string is considered to be part of the string.

strcmp( ) compares two strings byte-by-byte, according to the ordering of your machine's character set. The function returns an integer greater than, equal to, or less than 0, if the string pointed to by *s1* is greater than, equal to, or less than the string pointed to by *s2* respectively. The sign of a non-zero return value is determined by the sign of the difference between the values of the first pair of bytes that differ in the strings being compared. **strncmp( )** makes the same comparison but looks at a maximum of *n* bytes. Bytes following a null byte are not compared.

strcpy( ) copies string *src* to *dst* including the terminating null character, stopping after the null character has been copied. **strncpy( )** copies exactly *n* bytes, truncating *src* or adding null characters to *dst* if necessary. The result will not be null-terminated if the length of *src* is *n* or more. Each function returns *dst*.

strcspn( ) returns the length of the initial segment of string *s1* that consists entirely of characters not from string *s2*. **strspn( )** returns the length of the initial segment of string *s1* that consists entirely of characters from string *s2*.

strdup( ) returns a pointer to a new string that is a duplicate of the string pointed to by *s1*. The space for the new string is obtained using **malloc**(3C). If the new string cannot be created, a null pointer is returned.

strlen( ) returns the number of bytes in *s*, not including the terminating null character.

strpbrk( ) returns a pointer to the first occurrence in string *s1* of any character from string *s2*, or a null pointer if no character from *s2* exists in *s1*.

strstr( ) locates the first occurrence of the string *s2* (excluding the terminating null character) in string *s1*. **strstr( )** returns a pointer to the located string, or a null pointer if the string is not found. If *s2* points to a string with zero length (that is, the string ""), the function returns *s1*.

strtok( ) can be used to break the string pointed to by *s1* into a sequence of tokens, each of which is delimited by one or more characters from the string pointed to by *s2*. **strtok( )** considers the string *s1* to consist of a sequence of zero or more text tokens separated by spans of one or more characters from the separator string *s2*. The first call (with pointer *s1* specified) returns a pointer to the first character of the first token, and will have written a null character into *s1* immediately following the returned token. The function keeps track of its position in the string between separate calls, so that subsequent calls (which must be made with the first argument being a null pointer) will work through the string *s1* immediately following that token. In this way subsequent calls will work through the string *s1* until no tokens remain. The separator string *s2* may be different from call to call. When no token remains in *s1*, a null pointer is returned.

strtok_r( ) has the same functionality as **strtok( )** except that a pointer to a string placeholder *lasts* must be supplied by the caller. The **lasts** pointer is to keep track of the next substring in which to search for the next token.

SEE ALSO **malloc**(3C), **setlocale**(3C), **strxfrm**(3C)

NOTES The **strtok_r( )** interface is as proposed in the POSIX.4a Draft #6 document, and is subject to change to be compliant to the standard when it is accepted.

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

All of these functions assume the default locale ‘‘C.’’ For some locales, **strxfrm( )** should be applied to the strings before they are passed to the functions.

strtok( ) is unsafe in multi-thread applications. **strtok_r( )** should be used instead.

string( ), **strcasecmp( )**, **strcat( )**, **strchr( )**, **strcmp( )**, **strcpy( )**, **strcspn( )**, **strdup( )**, **strlen( ), strncasecmp( ), strncat( ), strncmp( ), strncpy( ), strpbrk( ), strrchr( )**, **strspn( )**, and **strstr( )**, are MT-Safe in multi-thread applications.

NAME

string_to_decimal, file_to_decimal, func_to_decimal – parse characters into decimal record

SYNOPSIS

```
#include <floatingpoint.h>

void string_to_decimal(char **pc, int nmax, int fortran_conventions,
        decimal_record *pd, enum decimal_string_form *pform, char **pechar);

void func_to_decimal(char **pc, int nmax, int fortran_conventions,
        decimal_record *pd, enum decimal_string_form *pform, char **pechar,
        int (*pget)(void), int *pnread, int (*punget)(int c));

#include <stdio.h>

void file_to_decimal(char **pc, int nmax, int fortran_conventions, decimal_record *pd,
        enum decimal_string_form *pform, char **pechar, FILE *pf, int *pnread);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **char_to_decimal** functions parse a numeric token from at most *nmax* characters in a string ***pc* or file ***pf* or function *(*pget)***()** into a decimal record ***pd*, classifying the form of the string in ***pform* and ***pechar*. The accepted syntax is intended to be sufficiently flexible to accomodate many languages:

whitespace value

or

whitespace sign value

where *whitespace* is any number of characters defined by *isspace* in **<ctype.h>**, *sign* is either of [+–], and *value* can be *number*, *nan*, or *inf*. *inf* can be INF (*inf_form*) or INFINITY (*infinity_form*) without regard to case. *nan* can be NAN (*nan_form*) or NAN(*nstring*) (*nanstring_form*) without regard to case; *nstring* is any string of characters not containing '**)**' or NULL; *nstring* is copied to *pd*–**>ds** and, currently, not used subsequently. *number* consists of

significand

or

significand efield

where *significand* must contain one or more digits and may contain one point; possible forms are

digits	*(int_form)*
digits.	*(intdot_form)*
.digits	*(dotfrac_form)*
digits.digits	*(intdotfrac_form)*

efield consists of

echar digits

or

echar sign digits

where *echar* is one of [Ee], and *digits* contains one or more digits.

When *fortran_conventions* is nonzero, additional input forms are accepted according to various Fortran conventions:

0 no Fortran conventions
1 Fortran list-directed input conventions
2 Fortran formatted input conventions, ignore blanks **(BN)**
3 Fortran formatted input conventions, blanks are zeros **(BZ)**

When *fortran_conventions* is nonzero, *echar* may also be one of [DdQq], and *efield* may also have the form

sign digits.

When *fortran_conventions*>= 2, blanks may appear in the *digits* strings for the integer, fraction, and exponent fields and may appear between *echar* and the exponent sign and after the infinity and NaN forms. If *fortran_conventions*== 2, the blanks are ignored. When *fortran_conventions*== 3, the blanks that appear in *digits* strings are interpreted as zeros, and other blanks are ignored.

When *fortran_conventions* is zero, the current locale's decimal point character is used as the decimal point; when *fortran_conventions* is nonzero, the period is used as the decimal point.

The form of the accepted decimal string is placed in ∗*peform*. If an *efield* is recognized, ∗*pechar* is set to point to the *echar*.

On input, ∗*pc* points to the beginning of a character string buffer of length >= *nmax*. On output, ∗*pc* points to a character in that buffer, one past the last accepted character. **string_to_decimal( )** gets its characters from the buffer; **file_to_decimal( )** gets its characters from ∗*pf* and records them in the buffer, and places a null after the last character read. **func_to_decimal( )** gets its characters from an int function *(∗pget)***( )**.

The scan continues until no more characters could possibly fit the acceptable syntax or until *nmax* characters have been scanned. If the *nmax* limit is not reached then at least one extra character will usually be scanned that is not part of the accepted syntax. **file_to_decimal( )** and **func_to_decimal( )** set ∗*pnread* to the number of characters read from the file; if greater than *nmax*, some characters were lost. If no characters were lost, **file_to_decimal( )** and **func_to_decimal( )** attempt to push back, with **ungetc**(3S) or *(∗punget)***( )**, as many as possible of the excess characters read, adjusting ∗*pnread* accordingly. If all unget calls are successful, then ∗∗*pc* will be NULL. No push back will be attempted if *(∗punget)***( )** is NULL.

Typical declarations for ∗*pget*() and ∗*punget*() are:

```
int xget(void)
{ . . . }
int (∗pget)(void) = xget;
int xunget(int c)
{ . . . }
int (∗punget)(int) = xunget;
```

If no valid number was detected, *pd*–>**fpclass** is set to **fp_signaling**, ∗*pc* is unchanged, and ∗*pform* is set to **invalid_form**.

atof(3C) and **strtod**(3C) use **string_to_decimal()**. **scanf**(3S) uses **file_to_decimal()**.

SEE ALSO **ctype**(3C), **localeconv**(3C), **scanf**(3S), **setlocale**(3C), **strtod**(3C), **ungetc**(3S)

NAME

strptime – date and time conversion

SYNOPSIS

#include <time.h>

char ******strptime(const char** **buf,* **const char** **format,* **struct tm** **tm***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **strptime( )** function converts the character string pointed to by *buf* to values which are stored in the **tm** structure pointed to by *tm*, using the format specified by *format*.

format is composed of zero or more conversion specifications. Each conversion specification is composed of a '%' (percent) character followed by one or two conversion characters which specify the replacement required. One or more white space characters (as specified by **isspace**(3C)) may precede or follow a conversion specification. There must be white-space or other non-alphanumeric characters between any two conversion specifications.

The following conversion specifications are supported:

%%	same as %
%a	day of week, using the locale's weekday names; either the abbreviated or full name may be specified
%A	same as **%a**
%b	month, using the locale's month names; either the abbreviated or full name may be specified
%B	same as **%b**
%c	locale's appropriate date and time representation
%C	century number [0,99]; leading zero is permitted but not required
%d	day of month [1,31]; leading zero is permitted but not required
%D	date as **%m**/**%d**/**%y**
%e	same as **%d**
%h	same as **%b**
%H	hour (24-hour clock) [0,23]; leading zero is permitted but not required
%I	hour (12-hour clock) [1,12]; leading zero is permitted but not required
%j	day number of the year [1,366]; leading zeros are permitted but not required
%m	month number [1,12]; leading zero is permitted but not required
%M	minute [0-59]; leading zero is permitted but not required
%n	any white space
%p	locale's equivalent of either a.m. or p.m.
%r	appropriate time representation in the 12-hour clock format with **%p**
%R	time as **%H**:**%M**
%S	seconds [0,61]; leading zero is permitted but not required
%t	any white space
%T	time as **%H**:**%M**:**%S**
%U	week number of the year as a decimal number [0,53], with Sunday as the first day of the week; leading zeros are permitted but not required
%w	weekday as a decimal number [0,6], with 0 representing Sunday;

%W week number of the year as a decimal number [0,53], with Monday as the first day of the week; leading zero is permitted but not required
%x locale's appropriate date representation
%X locale's appropriate time representation
%y year within the century [0,99]; leading zero is permitted but not required
%Y year, including the century (for example, 1993)
%Z timezone name or no characters if no time zone information exists

Modified Conversion Specifications

Some conversion specifications can be modified by the **E** and **O** modifier characters to indicate that an alternate format or specification should be used rather than the one normally used by the unmodified specification. If the alternate format or specification does not exist in the current locale, the behaviour will be as if the unmodified conversion specification were used.

%Ec locale's alternate appropriate date and time representation
%EC name of the base year (era) in the locale's alternate representation
%Ex locale's alternate date representation
%EX locale's alternate time representation
%Ey offset from **%EC** (year only) in the locale's alternate representation
%EY full alternate year representation
%Od day of the month using the locale's alternate numeric symbols
%Oe same as **%Od**
%OH hour (24-hour clock) using the locale's alternate numeric symbols
%OI hour (12-hour clock) using the locale's alternate numeric symbols
%Om month using the locale's alternate numeric symbols
%OM minutes using the locale's alternate numeric symbols
%OS seconds using the locale's alternate numeric symbols
%OU week number of the year (Sunday as the first day of the week) using the locale's alternate numeric symbols
%Ow number of the weekday (Sunday=0) using the locale's alternate numeric symbols
%OW week number of the year (Monday as the first day of the week) using the locale's alternate numeric symbols
%Oy year (offset from **%C**) in the locale's alternate representation and using the locale's alternate numeric symbols

General Specifications

A conversion specification that is an ordinary character is executed by scanning the next character from the buffer. If the character scanned from the buffer differs from the one comprising the specification, the specification fails, and the differing and subsequent characters remain unscanned.

A series of specifications composed of **%n**, **%t**, white-space characters or any combination is executed by scanning up to the first character that is not white space (which remains unscanned), or until no more characters can be scanned. White space is defined by **isspace**(3C).

Any other conversion specification is executed by scanning characters until a character matching the next specification is scanned, or until no more characters can be scanned. These characters, except the one matching the next specification, are then compared to the locale values associated with the conversion specifier. If a match is found, values for the appropriate *tm* structure members are set to values corresponding to the locale information. If no match is found, **strptime( )** fails and no more characters are scanned.

The month names, weekday names, era names, and alternate numeric symbols can consist of any combination of upper and lower case letters. The user can request that the input date or time specification be in a specific language by setting the **LC_TIME** category using **setlocale**(3C).

RETURN VALUES

Upon successful completion, **strptime( )** returns a pointer to the character following the last character parsed. Otherwise, a null pointer is returned.

FILES

/usr/lib/locale/*locale***/LC_TIME/time**
locale specific date and time information
/usr/lib/locale/*locale***/LC_CTYPE/ctype**
character characterization information

SEE ALSO

isspace(3C), **getdate**(3C), **setlocale**(3C), **strftime**(3C)

NOTES

Several “same as” formats, and the special processing of white-space characters are provided in order to ease the use of identical *format* strings for **strftime( )** and **strptime( )**.

The range of values for **%S** is [00,61] rather than [00,59] to allow for the occasional leap second and even more occasional double leap second.

NAME

strsignal – get error message string

SYNOPSIS

```
#include <string.h>
char *strsignal(int sig);
```

MT-LEVEL

Safe

DESCRIPTION

strsignal( ) maps the signal number in *sig* to a string describing the signal, and returns a pointer to that string. **strsignal( )** uses the same set of the messages as **psignal**(3C). The returned string should not be overwritten.

RETURN VALUES

strsignal( ) returns NULL if *sig* is not a valid signal number.

SEE ALSO

gettext(3I), **psignal**(3C), **setlocale**(3C), **str2sig**(3C)

NOTES

If the application is linked with **–lintl**, then messages returned from this function are in the native language specified by the **LC_MESSAGES** locale category; see **setlocale**(3C).

NAME

strtod, atof – convert string to double-precision number

SYNOPSIS

#include <stdlib.h>

double strtod(const char ∗*nptr*, **char** ∗∗*endptr*);

double atof(const char ∗*nptr*);

MT-LEVEL

MT-Safe

DESCRIPTION

strtod( ) returns as a double-precision floating-point number the value represented by the character string pointed to by *nptr*. The string is scanned up to the first unrecognized character.

strtod( ) recognizes an optional string of ''white-space'' characters (as defined by **isspace( )** in **ctype**(3C)), then an optional sign, then a string of digits optionally containing a decimal point character, then an optional exponent part including **e** or **E** followed by an optional sign, followed by an integer. The decimal point character is defined by the program's locale (category **LC_NUMERIC**). In the "**C**" locale, or in a locale where the decimal point character is not defined, the decimal point character defaults to a period (**.**).

If the value of *endptr* is not **(char** ∗∗**)NULL**, a pointer to the character terminating the scan is returned in the location pointed to by *endptr*. If no number can be formed, ∗*endptr* is set to *nptr*, and zero is returned.

atof(nptr) is equivalent to:

strtod(nptr, (char ∗∗**)NULL)**.

LC_NUMERIC

Determines how **strtod** and **atof** handle numeric formats. In the "C" locale, numeric handling follows the U.S. rules.

RETURN VALUES

If the correct value would cause overflow, ±**HUGE** is returned (according to the sign of the value), and **errno** is set to **ERANGE**.

If the correct value would cause underflow, **0** is returned and **errno** is set to **ERANGE**.

When the **–Xc** or **–Xa** compilation options are used, **HUGE_VAL** is returned instead of **HUGE**.

If *nptr* is NaN, then **atof( )** returns NaN.

FILES

/usr/lib/locale/*locale***/LC_NUMERIC/numeric**

LC_NUMERIC database for *locale*

SEE ALSO

ctype(3C), **matherr**(3M), **scanf**(3S), **strtol**(3C), **math**(5)

NAME

strtol, strtoll, strtoul, strtoull, atol, atoll, atoi, lltostr, ulltostr – conversion routines

SYNOPSIS

```
#include <stdlib.h>
long strtol(const char *str, char **ptr, int base);
long long strtoll(const char *str, char **ptr, int base);
unsigned long strtoul(const char *str, char **ptr, int base);
unsigned long long strtoull(const char *str, char **ptr, int base);
long atol(const char *str);
long long atoll(const char *str);
int atoi(const char *str);
char *lltostr(long long value, char *ptr);
char *ulltostr(unsigned long long value, char *ptr);
```

MT-LEVEL

MT-Safe

DESCRIPTION

strtol( ) returns as a **long** integer the value represented by the character string pointed to by **str**. The string is scanned up to the first character inconsistent with *base*. Leading ''white-space'' characters (as defined by **isspace( )** in **ctype**(3C)) are ignored.

strtoll( ) is similar to **strtol( )** except that the value is returned as a **long long**.

If the value of *ptr* is not **(char **)NULL**, a pointer to the character terminating the scan is returned in the location pointed to by *ptr*. If no integer can be formed, that location is set to **str**, and zero is returned.

If *base* is positive (and not greater than 36), it is used as the base for conversion. After an optional leading sign, leading zeros are ignored, and ''0x'' or ''0X'' is ignored if *base* is 16.

If *base* is zero, the string itself determines the base as follows: After an optional leading sign a leading zero indicates octal conversion, and a leading ''0x'' or ''0X'' hexadecimal conversion. Otherwise, decimal conversion is used.

Truncation from **long** to **int** can, of course, take place upon assignment or by an explicit cast.

If the value represented by *str* would cause overflow, **LONG_MAX** or **LONG_MIN** is returned (according to the sign of the value), and **errno** is set to the value, **ERANGE**.

strtoul( ) and **strtoull( )** are similar to **strtol( )** except that **strtoul( )** returns the value represented by *str* as an **unsigned long** integer and **strtoull( )** returns this value as an **unsigned long long**. If the value represented by *str* would cause overflow, **ULONG_MAX** is returned, and **errno** is set to the value, **ERANGE**.

Except for behavior on error, **atol( )** is equivalent to: **strtol(str, (char **)NULL, 10)**.

Except for behavior on error, **atoll( )** is equivalent to: **strtoll(str, (char **)NULL, 10)**.

Except for behavior on error, **atoi( )** is equivalent to: **(int) strtol(str, (char **)NULL, 10)**.

lltostr( ) returns a pointer to the string represented by the **long long** *value*.

ulltostr( ) is similar to **lltostr( )** except that *value* is an **unsigned long long**.

RETURN VALUES

If **strtol( )** is given a *base* greater than 36, it returns 0 and sets **errno** to **EINVAL**.

SEE ALSO

ctype(3C), **scanf**(3S), **strtod**(3C)

NOTES

strtol( ) no longer accepts values greater than **LONG_MAX** as valid input. Use **strtoul( )** instead.

NAME

strxfrm – string transformation

SYNOPSIS

```
#include <string.h>
size_t strxfrm(char *dst, const char *src, size_t n);
```

MT-LEVEL

Safe with exceptions

DESCRIPTION

strxfrm( ) transforms the string *src* and places the resulting string into the array *dst*. If **strcmp( )** is applied to two transformed strings, it will return the same result as **strcoll( )** applied to the same two original strings. The transformation is based on the program's locale for category **LC_COLLATE** (see **setlocale**(3C)).

No more than *n* bytes will be placed into the resulting array pointed to by *dst*, including the terminating null character. If *n* is 0 and *dst* is a NULL parameter, **strxfrm( )** returns the number of bytes required for the transformed string. If copying takes place between objects that overlap, the behavior is undefined.

strxfrm( ) returns the length of the transformed string (not including the terminating null character). If the value returned is *n* or more, the contents of the array *dst* are indeterminate.

RETURN VALUES

On failure, **strxfrm( )** returns **(size_t)–1**.

EXAMPLES

The value of the following expression is the size of the array needed to hold the transformation of the string pointed to by *s*.

```
1 + strxfrm(NULL, s, 0);
```

FILES

/usr/lib/locale/*locale***/LC_COLLATE** **LC_COLLATE** database for *locale*

SEE ALSO

colltbl(1M), **setlocale**(3C), **strcoll**(3C), **string**(3C), **wscoll**(3I), **environ**(5)

NOTES

strxfrm can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

swab – swap bytes

SYNOPSIS

```
#include <stdlib.h>
void swab(const char *from, char *to, int nbytes);
```

MT-LEVEL

MT-Safe

DESCRIPTION

swab( ) copies *nbytes* bytes pointed to by *from* to the array pointed to by *to*, exchanging adjacent even and odd bytes. *nbytes* should be even and non-negative. If *nbytes* is odd and positive, **swab( )** uses *nbytes*–1 instead. If *nbytes* is negative, **swab( )** does nothing.

NAME

syscall – indirect system call

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <sys/syscall.h>

int syscall(*number, arg, ...***)**

DESCRIPTION

syscall( ) performs the function whose assembly language interface has the specified *number*, and arguments *arg* .... Symbolic constants for functions can be found in the header **<sys/syscall.h>**.

RETURN VALUES

On error **syscall( )** returns –1 and sets the external variable **errno** (see **intro**(2)).

FILES

<sys/syscall.h>

SEE ALSO

intro(2), **pipe**(2)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

WARNINGS

There is no way to use **syscall( )** to call functions such as **pipe**(2)**,** which return values that do not fit into one hardware register.

Since many system calls are implemented as library wrappers around traps to the kernel, these calls may not behave as documented when called from **syscall( )**, which bypasses these wrappers. For these reasons, using **syscall( )** is not recommended.

NAME

sysconf – get configurable system variables

SYNOPSIS

#include <unistd.h>

long sysconf(int *name***);**

MT-LEVEL

MT-Safe, Async-Signal-Safe

DESCRIPTION

The **sysconf( )** function provides a method for an application to determine the current value of a configurable system limit or option (variable).

The *name* argument represents the system variable to be queried. The following table lists the minimal set of system variables from **<limits.h>** and **<unistd.h>** that can be returned by **sysconf( )**, and the symbolic constants, defined in **<unistd.h>** that are the corresponding values used for *name*.

Name	*Return Value*	*Meaning*
_SC_ARG_MAX	ARG_MAX	Max size of **argv[ ]** plus **envp[ ]**.
_SC_CHILD_MAX	CHILD_MAX	Max processes allowed to a UID.
_SC_CLK_TCK	CLK_TCK	Ticks per second (**clock_t**).
_SC_NGROUPS_MAX	NGROUPS_MAX	Max simultaneous groups to which one may belong.
_SC_OPEN_MAX	OPEN_MAX	Max open files per process.
_SC_PASS_MAX	PASS_MAX	
_SC_PAGESIZE	PAGESIZE	System memory page size.
_SC_JOB_CONTROL	_POSIX_JOB_CONTROL	Job control supported?
_SC_SAVED_IDS	_POSIX_SAVED_IDS	Saved IDs (**seteuid( )**) supported?
_SC_VERSION	_POSIX_VERSION	POSIX.1 version supported.
_SC_XOPEN_VERSION	_XOPEN_VERSION	
_SC_LOGNAME_MAX	LOGNAME_MAX	
_SC_NPROCESSORS_CONF		Number of processors configured.
_SC_NPROCESSORS_ONLN		Number of processors online.
_SC_PHYS_PAGES		Total number of pages of physical memory in system.
_SC_AVPHYS_PAGES		Number physical memory pages not currently in use by system.
_SC_AIO_LISTIO_MAX	AIO_LISTIO_MAX	Max number of I/O operations in a single list I/O call supported by implementation.
_SC_AIO_MAX	AIO_MAX	Max number of outstanding asynchronous I/O operations supported by implementation.
_SC_AIO_PRIO_DELTA_MAX	AIO_PRIO_DELTA_MAX	Max amount by which a process can decrease its asynchronous I/O priority level from its own scheduling priority.

_SC_DELAYTIMER_MAX	DELAYTIMER_MAX	Max number of timer expiration overruns.
_SC_MQ_OPEN_MAX	MQ_OPEN_MAX	Max number of open message queues a process may hold.
_SC_MQ_PRIO_MAX	MQ_PRIO_MAX	Max number of message priorities supported by implementation.
_SC_RTSIG_MAX	RTSIG_MAX	Max number of realtime signals reserved for application use in this implementation.
_SC_SEM_NSEMS_MAX	SEM_NSEMS_MAX	Max number of semaphores that a process may have.
_SC_SEM_VALUE_MAX	SEM_VALUE_MAX	Max value a semaphore may have.
_SC_SIGQUEUE_MAX	SIGQUEUE_MAX	Max number of queued signals that a process may send and have pending at receiver(s) at a time.
_SC_TIMER_MAX	TIMER_MAX	Max number of timers per process supported by implementation.
_SC_ASYNCHRONOUS_IO	_POSIX_ASYNCHRONOUS_IO	Supports Asynchronous I/O.
_SC_FSYNC	_POSIX_FSYNC	Supports File Synchronization.
_SC_MAPPED_FILES	_POSIX_MAPPED_FILES	Supports Memory Mapped Files.
_SC_MEMLOCK	_POSIX_MEMLOCK	Supports Process Memory Locking.
_SC_MEMLOCK_RANGE	_POSIX_MEMLOCK_RANGE	Supports Range Memory Locking.
_SC_MEMORY_PROTECTION	_POSIX_MEMORY_PROTECTION	Supports Memory Protection.
_SC_MESSAGE_PASSING	_POSIX_MESSAGE_PASSING	Supports Message Passing.
_SC_PRIORITIZED_IO	_POSIX_PRIORITIZED_IO	Supports Prioritized I/O.
_SC_PRIORITY_SCHEDULING	_POSIX_PRIORITY_SCHEDULING	Supports Process Scheduling
_SC_REALTIME_SIGNALS	_POSIX_REALTIME_SIGNALS	Supports Realtime Signals.
_SC_SEMAPHORES	_POSIX_SEMAPHORES	Supports Semaphores.
_SC_SHARED_MEMORY_ OBJECTS	_POSIX_SHARED_MEMORY_ OBJECTS	Supports Shared Memory Objects.
_SC_SYNCHRONIZED_IO	_POSIX_SYNCHRONIZED_IO	Supports Synchronized I/O.
_SC_TIMERS	_POSIX_TIMERS	Supports Timers.
_SC_GETGR_R_SIZE_MAX		
_SC_GETPW_R_SIZE_MAX		
_SC_LOGIN_NAME_MAX	LOGIN_NAME_MAX	Max length of login name.
_SC_THREAD_DESTRUCTOR_ ITERATIONS	PTHREAD_DESTRUCTOR_ ITERATIONS	Number attempts made to destroy thread-specific data upon exit.
_SC_THREAD_KEYS_MAX	PTHREAD_KEYS_MAX	Max number of data keys per process.
_SC_THREAD_STACK_MIN	PTHREAD_STACK_MIN	Min byte size of thread stack storage.
_SC_THREAD_THREADS_MAX	PTHREAD_THREADS_MAX	Max number of threads per process.
_SC_TTY_NAME_MAX	TTY_NAME_MAX	Max length of terminal device name.
_SC_THREADS	POSIX_THREADS	Supports Threads option.

_SC_THREAD_ATTR_STACKADDR	POSIX_THREAD_ATTR_STACKADDR Supports Thread Stack Address Attribute option.
_SC_THREAD_ATTR_STACKSIZE	POSIX_THREAD_ATTR_STACKSIZE Supports Thread Stack Size Attribute option.
_SC_THREAD_PRIORITY_SCHEDULING	POSIX_THREAD_PRIORITY_SCHEDULING Supports Thread Execution Scheduling option.
_SC_THREAD_PRIO_INHERIT	POSIX_THREAD_PRIO_INHERIT Supports Priority Inheritance option.
_SC_THREAD_PRIO_PROTECT	POSIX_THREAD_PRIO_PROTECT Supports Priority Protection option.
_SC_THREAD_PROCESS_SHARED	POSIX_THREAD_PROCESS_SHARED Supports Process-Shared Synchronization option.
_SC_THREAD_SAFE_FUNCTIONS	POSIX_THREAD_SAFE_FUNCTIONS Supports Thread-Safe Functions option.

RETURN VALUES

If *name* is an invalid value, **sysconf( )** will return –1 and set **errno** to indicate the error. If **sysconf( )** fails due to a value of *name* that is not defined on the system, the function will return a value of –1 without changing the value of **errno**.

SEE ALSO

fpathconf(2), **seteuid**(2), **setrlimit**(2)

NOTES

A call to **setrlimit( )** may cause the value of **OPEN_MAX** to change.

Multiplying **sysconf(_SC_PHYS_PAGES)** or **sysconf(_SC_AVPHYS_PAGES)** by **sysconf(_SC_PAGESIZE)** to determine memory amount in bytes can exceed the maximum values representable in a long or unsigned long.

_SC_PHYS_PAGES and **_SC_AVPHYS_PAGES** are specific to Solaris 2.3 and later releases.

The value of **CLK_TCK** may be variable and it should not be assumed that **CLK_TCK** is a compile-time constant.

NAME

syslog, openlog, closelog, setlogmask – control system log

SYNOPSIS

```
#include <syslog.h>
void openlog(char *ident, int logopt, int facility);
void syslog(int priority, char *logstring, /* parameters */ ...);
void closelog(void);
int setlogmask(int maskpri);
```

MT-LEVEL

Safe

DESCRIPTION

syslog() passes a message to **syslogd**(1M), which may append it to a log file, write it to the system console, or forward it (to either a list of users or **syslogd** on another host on the network), depending on the configuration of **/etc/syslog.conf**. *logstring* is tagged with a priority of *priority*, and looks like a **printf**(3B) string with one additional allowable format specification, **%m**, which is replaced with the error message string corresponding to the error number in **errno**. A trailing NEWLINE is added if needed. Options passed to **openlog()** may cause the size of the message to expand. The maximum size of the message passed to **syslogd** is 1024 bytes.

Priorities are encoded as a *facility* and a *level*. The facility describes the part of the system generating the message. The level is selected from the bitwise inclusive OR of zero or more of the following flags, defined in the header **<syslog.h>**.

LOG_EMERG	A panic condition. This is normally broadcast to all users.
LOG_ALERT	A condition that should be corrected immediately, such as a corrupted system database.
LOG_CRIT	Critical conditions, such as hard device errors.
LOG_ERR	Errors.
LOG_WARNING	Warning messages.
LOG_NOTICE	Conditions that are not error conditions, but that may require special handling.
LOG_INFO	Informational messages.
LOG_DEBUG	Messages that contain information normally of use only when debugging a program.

If special processing is needed, **openlog()** can be called to initialize the log file. The parameter *ident* is a string that is prepended to every message. *logopt* is a bit field indicating logging options. Values for *logopt* are:

LOG_PID	Log the process ID with each message. This is useful for identifying specific daemon processes (for daemons that fork).

LOG_CONS	Write messages to the system console if they cannot be sent to **syslogd**(1M). This option is safe to use in daemon processes that have no controlling terminal, since **syslog()** forks before opening the console.
LOG_NDELAY	Open the connection to **syslogd**(1M) immediately. Normally the open is delayed until the first message is logged. This is useful for programs that need to manage the order in which file descriptors are allocated.
LOG_NOWAIT	Do not wait for child processes that have been forked to log messages onto the console. This option should be used by processes that enable notification of child termination using **SIGCHLD**, since **syslog()** may otherwise block waiting for a child whose exit status has already been collected.

The *facility* parameter encodes a default facility to be assigned to all messages that do not have an explicit facility already encoded:

LOG_KERN	Messages generated by the kernel. These cannot be generated by any user processes.
LOG_USER	Messages generated by random user processes. This is the default facility identifier if none is specified.
LOG_MAIL	The mail system.
LOG_DAEMON	System daemons, such as **in.ftpd**(1M).
LOG_AUTH	The authorization system: **login**(1), **su**(1M), **getty**(1M), etc.
LOG_LPR	The line printer spooling system: **lpr**(1B), **lpc**(1B), etc.
LOG_NEWS	Reserved for the USENET network news system.
LOG_UUCP	Reserved for the UUCP system; it does not currently use **syslog**.
LOG_CRON	The **cron**/**at** facility; **crontab**(1), **at**(1), **cron**(1M), etc.
LOG_LOCAL0	Reserved for local use.
LOG_LOCAL1	Reserved for local use.
LOG_LOCAL2	Reserved for local use.
LOG_LOCAL3	Reserved for local use.
LOG_LOCAL4	Reserved for local use.
LOG_LOCAL5	Reserved for local use.
LOG_LOCAL6	Reserved for local use.
LOG_LOCAL7	Reserved for local use.

closelog() can be used to close the log file.

setlogmask() sets the log priority mask to *maskpri* and returns the previous mask. Calls to **syslog()** with a priority not set in *maskpri* are rejected. The mask for an individual priority *pri* is calculated by the macro **LOG_MASK(***pri***);** the mask for all priorities up to and including *toppri* is given by the macro **LOG_UPTO(***toppri***)**. The default allows all priorities to be logged.

EXAMPLES

This call logs a message at priority **LOG_ALERT**:

```
syslog(LOG_ALERT, "who: internal error 23");
```

The FTP daemon **ftpd** would make this call to **openlog()** to indicate that all messages it logs should have an identifying string of **ftpd**, should be treated by **syslogd**(1M) as other messages from system daemons are, should include the process ID of the process logging the message:

```
openlog("ftpd", LOG_PID, LOG_DAEMON);
```

Then it would make the following call to **setlogmask()** to indicate that messages at priorities from **LOG_EMERG** through **LOG_ERR** should be logged, but that no messages at any other priority should be logged:

```
setlogmask(LOG_UPTO(LOG_ERR));
```

Then, to log a message at priority **LOG_INFO**, it would make the following call to **syslog**:

```
syslog(LOG_INFO, "Connection from host %d", CallingHost);
```

A locally-written utility could use the following call to **syslog()** to log a message at priority **LOG_INFO** to be treated by **syslogd**(1M) as other messages to the facility **LOG_LOCAL2** are:

```
syslog(LOG_INFO|LOG_LOCAL2, "error: %m");
```

SEE ALSO

at(1), **crontab**(1), **logger**(1), **login**(1), **lpc**(1B), **lpr**(1B), **cron**(1M), **getty**(1M), **in.ftpd**(1M), **su**(1M), **syslogd**(1M), **printf**(3B), **syslog.conf**(4)

NAME

sysmem, asysmem – return physical memory information

SYNOPSIS

```
long sysmem(void);
long asysmem(void);
```

DESCRIPTION

These routines are obsolete and have been replaced by arguments to **sysconf**(3C). They were mistakenly published in the *System V Interface Definition*, Third Edition, (SVID) and corrected by the Errata: "The following routines were mistakenly include in SVID Edition 3 and were not designed as customer level interfaces: **sysmem(AS_LIB)**, **asysmem(AS_LIB)** , . . . They are therefore removed."

The routine **sysmem( )** determines the total amount of physical memory of the system. It returns a long integer representing the total amount of physical memory, in bytes. Because **sysmem( )** returns a long integer it cannot report the amount of memory for configurations with amounts of memory in bytes greater than the maximum positive value represented by a long integer. **sysconf(_SC_PHYS_PAGES)** should be used to avoid this limitation. (See **sysconf**(3C).)

The routine **asysmem( )** determines the total amount of memory not currently in use on the system. It returns a long integer representing the total amount of available memory, in bytes. Because **asysmem( )** returns a long integer it is limited similar to **sysmem( )**. **sysconf(_SC_AVPHYS_PAGES)** should be used to avoid this limitation. (See **sysconf**(3C).)

RETURN VALUES

Upon successful completion, these routines return the amount of memory in bytes; otherwise, they return -1.

SEE ALSO

sysconf(3C)

NOTES

sysmen( ) and **asysmem( )** are obsolete and should be replaced with **sysconf**(3C).

NAME

system – issue a shell command

SYNOPSIS

```
#include <stdlib.h>
int system(const char *string);
```

MT-LEVEL

MT-Safe

DESCRIPTION

system( ) causes the *string* to be given to the shell as input, as if the string had been typed as a command at a terminal. The invoker waits until the shell has completed, then returns the exit status of the shell in the format specified by **waitpid**(2).

If *string* is a null pointer, **system( )** checks if the shell exists and is executable. If the shell is available, **system( )** returns non-zero; otherwise it returns zero.

Solaris **system( )** uses **/usr/bin/sh** (see **sh**(1)).

XPG4 **system( )** uses the XPG4-compliant shell **/usr/bin/ksh** (see **ksh**(1)).

RETURN VALUES

system( ) forks to create a child process that in turn **exec**s the shell in order to execute *string*. If the **fork( )** or **exec( )** fails, **system( )** returns a value of **–1** and sets **errno**.

ERRORS

system( ) fails if one or more of the following are true:

EAGAIN The system-imposed limit on the total number of processes under execution by a single user would be exceeded.

EINTR **system( )** was interrupted by a signal.

ENOMEM The new process requires more memory than is available.

SEE ALSO

ksh(1), **sh**(1), **useradd**(1M), **exec**(2), **fork**(2), **setuid**(2), **waitpid**(2), **xpg4**(5)

NOTES

system( ) will fail to execute **setuid( )** or **setgid( )** if either the uid or gid of the application's owner / group is less than 100. (see **useradd**(1M) and **setuid**(2)).

NAME

t_accept – accept a connect request

SYNOPSIS

cc [*flag* ...] *file* ... –**lnsl** [*library* ...]

#include <tiuser.h>

int t_accept(int *fildes***, int** *resfd***, struct t_call** ∗*call***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function is issued by a transport user to accept a connect request. *fildes* identifies the local transport endpoint where the connect indication arrived, *resfd* specifies the local transport endpoint where the connection is to be established, and *call* contains information required by the transport provider to complete the connection. *call* points to a **t_call** structure that contains the following members:

```
struct netbuf   addr;
struct netbuf   opt;
struct netbuf   udata;
int             sequence;
```

struct netbuf contains the following members:

```
unsigned int    maxlen;
unsigned int    len;
char *      buf;
```

In *call*, **addr** is the address of the caller, **opt** indicates any protocol-specific parameters associated with the connection, **udata** points to any user data to be returned to the caller, and **sequence** is the value returned by **t_listen** that uniquely associates the response with a previously received connect indication.

A transport user may accept a connection on either the same, or on a different, local transport endpoint from the one on which the connect indication arrived. If the same endpoint is specified (that is, *resfd=fildes*), the connection can be accepted unless the following condition is true: The user has received other indications on that endpoint but has not responded to them (with **t_accept( )** or **t_snddis**(3N)**).** For this condition, **t_accept**(3N) will fail and set **t_errno** to **TBADF**.

If a different transport endpoint is specified (*resfd*!=*fildes*), the endpoint must be bound to a protocol address and must be in the **T_IDLE** state (see **t_getstate**(3N)) before the **t_accept**(3N) is issued.

For both types of endpoints, **t_accept( )** will fail and set **t_errno** to **TLOOK** if there are indications (for example, a connect or disconnect) waiting to be received on that endpoint.

The values of parameters specified by **opt** and the syntax of those values are protocol specific. The **udata** argument enables the called transport user to send user data to the caller and the amount of user data must not exceed the limits supported by the transport provider, as returned in the **connect** field of the *info* argument of **t_open**(3N) or **t_getinfo**(3N). If the **len** field of **udata** is zero, no data will be sent to the caller.

RETURN VALUES

Upon successful completion, a value of 0 is returned. Otherwise, a value of –1 is returned, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TACCES	The user does not have permission to accept a connection on the responding transport endpoint or to use the specified options.
TBADF	The specified file descriptor does not refer to a transport endpoint, or the user is illegally accepting a connection on the same transport endpoint on which the connect indication arrived.
TBADDATA	The amount of user data specified was not within the bounds allowed by the transport provider.
TBADOPT	The specified options were in an incorrect format or contained illegal information.
TBADSEQ	An invalid sequence number was specified.
TLOOK	An asynchronous event has occurred on the transport endpoint referenced by *fildes* and requires immediate attention.
TNOTSUPPORT	This function is not supported by the underlying transport provider.
TOUTSTATE	The function was issued in the wrong sequence on the transport endpoint referenced by *fildes*, or the transport endpoint referred to by *resfd* is not in the **T_IDLE** state.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_accept(3N), **t_connect**(3N), **t_getinfo**(3N), **t_getstate**(3N), **t_listen**(3N), **t_open**(3N), **t_snddis**(3N), **t_rcvconnect**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_alloc – allocate a library structure

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

char ∗**t_alloc(int** *fildes***, int** *struct_type***, int** *fields***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **t_alloc( )** function dynamically allocates memory for the various transport function argument structures as specified below. This function will allocate memory for the specified structure, and will also allocate memory for buffers referenced by the structure.

The structure to allocate is specified by *struct_type*, and can be one of the following:

T_BIND	**struct t_bind**
T_CALL	**struct t_call**
T_OPTMGMT	**struct t_optmgmt**
T_DIS	**struct t_discon**
T_UNITDATA	**struct t_unitdata**
T_UDERROR	**struct t_uderr**
T_INFO	**struct t_info**

where each of these structures may subsequently be used as an argument to one or more transport functions.

Each of the above structures, except **T_INFO**, contains at least one field of type **struct netbuf**. **netbuf** is described in **t_connect**(3N). For each field of this type, the user may specify that the buffer for that field should be allocated as well. The *fields* argument specifies this option, where the argument is the bitwise-OR of any of the following:

T_ADDR The **addr** field of the **t_bind**, **t_call**, **t_unitdata**, or **t_uderr** structures.

T_OPT The **opt** field of the **t_optmgmt**, **t_call**, **t_unitdata**, or **t_uderr** structures.

T_UDATA The **udata** field of the **t_call**, **t_discon**, or **t_unitdata** structures.

T_ALL All relevant fields of the given structure.

For each field specified in *fields*, **t_alloc( )** will allocate memory for the buffer associated with the field, and initialize the **buf** pointer and **maxlen** (see **netbuf** in **t_connect**(3N) for description of **buf** and **maxlen**) field accordingly. The length of the buffer allocated will be based on the same size information that is returned to the user on **t_open**(3N) and **t_getinfo**(3N). Thus, *fildes* must refer to the transport endpoint through which the newly allocated structure will be passed, so that the appropriate size information can be accessed. If the size value associated with any specified field is –1, the underlying service provider can support a buffer of unlimited size. If this is the case, **t_alloc( )** will allocate a buffer with the default size 1024 bytes. See the **NOTES** section for information regarding memory allocation for buffers other than 1024 bytes. If the size value is –2, **t_alloc( )** will set the buffer pointer to **NULL** and the buffer maximum size to 0, and then will return

success (see **t_open**(3N) or **t_getinfo**(3N)).
For any field not specified in *fields*, **buf** will be set to NULL and **maxlen** will be set to zero.

Use of **t_alloc( )** to allocate structures will help ensure the compatibility of user programs with future releases of the transport interface.

RETURN VALUES

On successful completion, **t_alloc( )** returns a pointer to the newly allocated structure. On failure, NULL is returned, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF The specified file descriptor does not refer to a transport endpoint.

TSYSERR A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_connect(3N), **t_free**(3N), **t_getinfo**(3N), **t_open**(3N)

Transport Interfaces Programming Guide

NOTES

If the underlying service provider supports a buffer of unlimited size in the **netbuf** structure (see **t_connect**(3N)), **t_alloc( )** will return a buffer of size 1024 bytes. If a larger size buffer is required, it will need to be allocated separately using a memory allocation routine such as **malloc**(3C). The **buf** and **maxlen** fields of the **netbuf** data structure can then be updated with the address of the new buffer and the 1024 byte buffer originally allocated by **t_alloc( )** can be freed using **free**(3C).

This interface is safe in multithreaded applications.

NAME

t_bind – bind an address to a transport endpoint

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_bind(int *fildes***, const struct t_bind** ∗*req***, struct t_bind** ∗*ret***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function associates a protocol address with the transport endpoint specified by *fildes* and activates that transport endpoint. In connection mode, the transport provider may begin accepting or requesting connections on the transport endpoint. In connectionless mode, the transport user may send or receive data units through the transport endpoint.

The *req* and *ret* arguments point to a **t_bind** structure containing the following members:

```
struct netbuf addr;
unsigned qlen;
```

netbuf is described in **t_connect**(3N). The **addr** field of the **t_bind** structure specifies a protocol address and the **qlen** field is used to indicate the maximum number of outstanding connect indications.

req is used to request that an address, represented by the **netbuf** structure, be bound to the given transport endpoint. **len** (see **netbuf** in **t_connect**(3N); also for **buf** and **maxlen**) specifies the number of bytes in the address and **buf** points to the address buffer. **maxlen** has no meaning for the *req* argument. On return, *ret* contains the address that the transport provider actually bound to the transport endpoint; this may be different from the address specified by the user in *req*. In *ret*, the user specifies **maxlen**, which is the maximum size of the address buffer, and **buf**, which points to the buffer where the address is to be placed. On return, **len** specifies the number of bytes in the bound address and **buf** points to the bound address. If **maxlen** is not large enough to hold the returned address, an error will result.

If the requested address is not available, or if no address is specified in *req* (the **len** field of **addr** in *req* is zero) the transport provider may assign an appropriate address to be bound, and will return that address in the **addr** field of *ret*. The user can compare the addresses in *req* and *ret* to determine whether the transport provider bound the transport endpoint to a different address than that requested.

req may be **NULL** if the user does not wish to specify an address to be bound. Here, the value of **qlen** is assumed to be zero, and the transport provider must assign an address to the transport endpoint. Similarly, *ret* may be **NULL** if the user does not care what address was bound by the provider and is not interested in the negotiated value of **qlen**. It is valid to set *req* and *ret* to **NULL** for the same call, in which case the provider chooses the address to bind to the transport endpoint and does not return that information to the user.

The **qlen** field has meaning only when initializing a connection-mode service. It specifies the number of outstanding connect indications the transport provider should support for the given transport endpoint. An outstanding connect indication is one that has been passed to the transport user by the transport provider. A value of **qlen** greater than zero is only meaningful when issued by a passive transport user that expects other users to call it. The value of **qlen** will be negotiated by the transport provider and may be changed if the transport provider cannot support the specified number of outstanding connect indications. On return, the **qlen** field in *ret* will contain the negotiated value.

This function allows more than one transport endpoint to be bound to the same protocol address (however, the transport provider must support this capability also), but it is not allowable to bind more than one protocol address to the same transport endpoint. If a user binds more than one transport endpoint to the same protocol address, only one endpoint can be used to listen for connect indications associated with that protocol address. In other words, only one **t_bind( )** for a given protocol address may specify a value of **qlen** greater than zero. In this way, the transport provider can identify which transport endpoint should be notified of an incoming connect indication. If a user attempts to bind a protocol address to a second transport endpoint with a value of **qlen** greater than zero, the transport provider will assign another address to be bound to that endpoint. If a user accepts a connection on the transport endpoint that is being used as the listening endpoint, the bound protocol address will be found to be busy for the duration of that connection. No other transport endpoints may be bound for listening while that initial listening endpoint is in the data transfer phase. This will prevent more than one transport endpoint bound to the same protocol address from accepting connect indications.

RETURN VALUES

t_bind( ) returns **0** on success. On failure, **t_bind( )** returns **–1**, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TACCES	The user does not have permission to use the specified address.
TBADADDR	The specified protocol address was in an incorrect format or contained illegal information.
TBADF	The specified file descriptor does not refer to a transport endpoint.
TBUFOVFLW	The number of bytes allowed for an incoming argument is not sufficient to store the value of that argument. The provider's state will change.
TNOADDR	The transport provider could not allocate an address. **T_IDLE** and the information to be returned in *ret* will be discarded.

TOUTSTATE	The function was issued in the wrong sequence.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_connect(3N), **t_open**(3N), **t_optmgmt**(3N), **t_unbind**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_close – close a transport endpoint

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_close(int *fildes***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **t_close( )** function informs the transport provider that the user is finished with the transport endpoint specified by *fildes*, and frees any local library resources associated with the endpoint. In addition, **t_close( )** closes the file associated with the transport endpoint.

t_close( ) should be called from the **T_UNBND** state (see **t_getstate**(3N)). However, this function does not check state information, so it may be called from any state to close a transport endpoint. If this occurs, the local library resources associated with the endpoint will be freed automatically. In addition, **close**(2) will be issued for that file descriptor; if no other process has the file descriptor open, the close will terminate any connection that may be associated with that endpoint. The connection termination will be abortive or orderly depending on the service type supported by the underlying transport provider.

RETURN VALUES

t_close returns 0 on success. On failure **t_close** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to the following:

TBADF The specified file descriptor does not refer to a transport endpoint.
TSYSERR A system error occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

close(2), **t_getstate**(3N), **t_open**(3N), **t_unbind**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_connect – establish a connection with another transport user

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_connect(int *fildes***, const struct t_call** ∗*sndcall***, struct t_call** ∗*rcvcall***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function enables a transport user to request a connection to the specified destination transport user. *fildes* identifies the local transport endpoint where communication will be established, while *sndcall* and *rcvcall* point to a **t_call** structure that contains the following members:

```
struct netbuf    addr;
struct netbuf    opt;
struct netbuf    udata;
int              sequence;
```

sndcall specifies information needed by the transport provider to establish a connection and *rcvcall* specifies information that is associated with the newly established connection.

The address is specified in the **netbuf** structure which has the following format:

```
struct netbuf {
        unsigned int maxlen;
        unsigned int len;
        char *buf;
}
```

where **maxlen** specifies the maximum length of the buffer in bytes, **len** specifies the bytes of data in the buffer, and **buf** points to the buffer that contains the data.

In *sndcall*, **addr** specifies the protocol address of the destination transport user, **opt** presents any protocol-specific information that might be needed by the transport provider, **udata** points to optional user data that may be passed to the destination transport user during connection establishment, and **sequence** has no meaning for this function.

On return in *rcvcall*, **addr** returns the protocol address associated with the responding transport endpoint, **opt** presents any protocol-specific information associated with the connection, **udata** points to optional user data that may be returned by the destination transport user during connection establishment, and **sequence** has no meaning for this function.

The **opt** argument implies no structure on the options that may be passed to the transport provider. The transport provider is free to specify the structure of any options passed to it. These options are specific to the underlying protocol of the transport provider. The user may choose not to negotiate protocol options by setting the **len** field of **opt** to zero. In this case, the provider may use default options.

The **udata** argument enables the caller to pass user data to the destination transport user and receive user data from the destination user during connection establishment. However, the amount of user data must not exceed the limits supported by the transport provider as returned in the **connect** field of the **info** argument of **t_open**(3N) or **t_getinfo**(3N). If the **len** (see **netbuf** in **t_connect**(3N)) field of **udata** is zero in *sndcall*, no data will be sent to the destination transport user.

On return, the **addr**, **opt**, and **udata** fields of *rcvcall* will be updated to reflect values associated with the connection. Thus, the **maxlen** (see **netbuf** in **t_connect**(3N)) field of each argument must be set before issuing this function to indicate the maximum size of the buffer for each. However, *rcvcall* may be NULL, in which case no information is given to the user on return from **t_connect( )**.

By default, **t_connect( )** executes in synchronous mode, and will wait for the destination user's response before returning control to the local user. A successful return (that is, return value of zero) indicates that the requested connection has been established. However, if **O_NDELAY** or **O_NONBLOCK** is set (using **t_open**(3N) or **fcntl**(2)**), t_connect( )** executes in asynchronous mode. In this case, the call will not wait for the remote user's response, but will return control immediately to the local user and return -1 with **t_errno** set to **TNODATA** to indicate that the connection has not yet been established. In this way, the function simply initiates the connection establishment procedure by sending a connect request to the destination transport user.

RETURN VALUES

t_connect( ) returns 0 on success. On failure **t_connect( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TACCES	The user does not have permission to use the specified address or options.
TBADADDR	The specified protocol address was in an incorrect format or contained illegal information.
TBADDATA	The amount of user data specified was not within the bounds allowed by the transport provider.
TBADF	The specified file descriptor does not refer to a transport endpoint.
TBADOPT	The specified protocol options were in an incorrect format or contained illegal information.
TBUFOVFLW	The number of bytes allocated for an incoming argument is not sufficient to store the value of that argument. If executed in synchronous mode, the provider's state, as seen by the user, changes to **T_DATAXFER**, and the connect indication information to be returned in *rcvcall* is discarded.
TLOOK	An asynchronous event has occurred on this transport endpoint and requires immediate attention.
TNODATA	**O_NDELAY** or **O_NONBLOCK** was set, so the function successfully initiated the connection establishment procedure, but did not wait

	for a response from the remote user.
TNOTSUPPORT	This function is not supported by the underlying transport provider.
TOUTSTATE	The function was issued in the wrong sequence.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

fcntl(2), **t_accept**(3N), **t_getinfo**(3N), **t_listen**(3N), **t_open**(3N), **t_optmgmt**(3N), **t_rcvconnect**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_error – produce error message

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

void t_error(const char ∗*errmsg***);**

extern int t_errno;
extern char ∗**t_errlist[ ];**
extern int t_nerr;

MT-LEVEL

MT-Safe

DESCRIPTION

t_error( ) produces a message on the standard error output which describes the last error encountered during a call to a transport function. The argument string *errmsg* is a user-supplied error message that gives context to the error.

t_error( ) prints the user-supplied error message followed by a colon and the standard transport function error message for the current value contained in **t_errno**. If **t_errno** is **TSYSERR**, **t_error** will also print the standard error message for the current value contained in **errno** (see **intro**(2)).

t_errlist is the array of message strings, to allow user message formatting. **t_errno** can be used as an index into this array to retrieve the error message string (without a terminating newline). **t_nerr** is the maximum index value for the **t_errlist** array.

t_errno is set when an error occurs and is not cleared on subsequent successful calls.

EXAMPLES

If a **t_connect( )** function fails on transport endpoint *fd2* because a bad address was given, the following call might follow the failure:

t_error("t_connect failed on fd2");

The diagnostic message would print as:

t_connect failed on fd2: Incorrect transport address format

where "**t_connect failed on fd2**" tells the user which function failed on which transport endpoint, and ‘‘Incorrect transport address format’’ identifies the specific error that occurred.

SEE ALSO

intro(2)

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WARNINGS

This interface is deprecated. Use **t_strerror**(3N) instead. For details on the ramifications of referencing **t_errlist[ ]** directly, please refer to the Copy Relocations section of Chapter 4 in the *Linker and Libraries Guide*.

NOTES

This interface is safe in multithreaded applications.

NAME

t_free – free a library structure

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_free(char **ptr***, int** *struct_type***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **t_free( )** function frees memory previously allocated by **t_alloc**(3N). This function will free memory for the specified structure, and will also free memory for buffers referenced by the structure.

ptr points to one of the six structure types described for **t_alloc**(3N), and *struct_type* identifies the type of that structure, which can be one of the following:

T_BIND	**struct t_bind**
T_CALL	**struct t_call**
T_OPTMGMT	**struct t_optmgmt**
T_DIS	**struct t_discon**
T_UNITDATA	**struct t_unitdata**
T_UDERROR	**struct t_uderr**
T_INFO	**struct t_info**

where each of these structures is used as an argument to one or more transport functions.

t_free( ) will check the **addr**, **opt**, and **udata** fields of the given structure (as appropriate), and free the buffers pointed to by the **buf** field of the **netbuf** (see **t_connect**(3N)) structure. If **buf** is NULL, **t_free( )** will not attempt to free memory. After all buffers are freed, **t_free( )** will free the memory associated with the structure pointed to by *ptr*.

Undefined results will occur if *ptr* or any of the **buf** pointers points to a block of memory that was not previously allocated by **t_alloc**(3N).

RETURN VALUES

t_free returns 0 on success. On failure **t_free** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to the following:

TSYSERR A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_connect(3N), **t_alloc**(3N)

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NOTES

This interface is safe in multithreaded applications.

NAME

t_getinfo – get protocol-specific service information

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_getinfo(int *fildes***, struct t_info** ∗*info***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function returns the current characteristics of the underlying transport protocol associated with file descriptor *fildes*. The *info* structure is used to return the same information returned by **t_open**(3N). This function enables a transport user to access this information during any phase of communication.

The *info* argument points to a **t_info** structure, which contains the following members:

```
long addr;      /* max size of transport protocol address */
long options;   /* max number of bytes of protocol-specific options */
long tsdu;      /* max size of transport service data unit (TSDU) */
long etsdu;     /* max size of expedited transport service data unit (ETSDU) */
long connect;   /* max amount of data allowed on connection establishment
                   functions */
long discon;    /* max amount of data allowed on t_snddis(3N) and
                   t_rcvdis(3N) functions */
long servtype;  /* service type supported by the transport provider */
```

The values of the fields have the following meanings:

addr A value greater than or equal to zero indicates the maximum size of a transport protocol address; a value of –1 specifies that there is no limit on the address size; and a value of –2 specifies that the transport provider does not provide user access to transport protocol addresses.

options A value greater than or equal to zero indicates the maximum number of bytes of protocol-specific options supported by the provider; a value of –1 specifies that there is no limit on the option size; and a value of –2 specifies that the transport provider does not support user-settable options.

tsdu A value greater than zero specifies the maximum size of a transport service data unit (**TSDU**); a value of zero specifies that the transport provider does not support the concept of **TSDU**, although it does support the sending across a connection of a data stream with no logical boundaries preserved; a value of –1 specifies that there is no limit on the size of a **TSDU**; and a value of –2 specifies that the transfer of normal data is not supported by the transport provider.

etsdu A value greater than zero specifies the maximum size of an expedited transport service data unit (**ETSDU**); a value of zero specifies that the transport provider does not support the concept of **ETSDU**, although it does support the sending across a connection of an expedited data stream with no logical boundaries preserved. A value of –1 specifies that there is no limit on the size of an **ETSDU**; and a value of –2 specifies that the transfer of expedited data is not supported by the transport provider.

connect A value greater than or equal to zero specifies the maximum amount of data that may be associated with connection establishment functions; a value of –1 specifies that there is no limit on the amount of data sent during connection establishment; and a value of –2 specifies that the transport provider does not allow data to be sent with connection establishment functions.

discon A value greater than or equal to zero specifies the maximum amount of data that may be associated with the **t_snddis**(3N) and **t_rcvdis**(3N) functions; a value of –1 specifies that there is no limit on the amount of data sent with these abortive release functions; and a value of –2 specifies that the transport provider does not allow data to be sent with the abortive release functions.

servtype This field specifies the service type supported by the transport provider, as described below.

If a transport user is concerned with protocol independence, the above sizes may be accessed to determine how large the buffers must be to hold each piece of information. Alternatively, the **t_alloc( )** function may be used to allocate these buffers. An error will result if a transport user exceeds the allowed data size on any function. The value of each field may change as a result of option negotiation, and **t_getinfo( )** enables a user to retrieve the current characteristics.

The **servtype** field of *info* may specify one of the following values on return:

T_COTS The transport provider supports a connection-mode service but does not support the optional orderly release facility.

T_COTS_ORD The transport provider supports a connection-mode service with the optional orderly release facility.

T_CLTS The transport provider supports a connectionless-mode service. For this service type, **t_open**(3N) will return –2 for **etsdu**, **connect**, and **discon**.

RETURN VALUES

t_getinfo( ) returns 0 on success. On failure **t_getinfo( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_open(3N), **t_rcvdis**(3N), **t_snddis**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_getstate – get the current state

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_getstate(int *fildes***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **t_getstate( )** function returns the current state of the provider associated with the transport endpoint specified by *fildes*.

RETURN VALUES

t_getstate( ) returns the current state on successful completion. On failure **t_getstate( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set. The current state may be one of the following:

T_UNBND	unbound
T_IDLE	idle
T_OUTCON	outgoing connection pending
T_INCON	incoming connection pending
T_DATAXFER	data transfer
T_OUTREL	outgoing orderly release (waiting for an orderly release indication)
T_INREL	incoming orderly release (waiting for an orderly release request)

If the provider is undergoing a state transition when **t_getstate( )** is called, the function will fail.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TSTATECHNG	The transport provider is undergoing a state change.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_open(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_listen – listen for a connect request

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_listen(int *fildes***, struct t_call** ∗*call***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function listens for a connect request from a calling transport user. *fildes* identifies the local transport endpoint where connect indications arrive, and on return, *call* contains information describing the connect indication. *call* points to a **t_call** structure, which contains the following members:

```
struct netbuf   addr;
struct netbuf   opt;
struct netbuf   udata;
int             sequence;
```

netbuf is described in **t_connect**(3N). In *call*, **addr** returns the protocol address of the calling transport user, **opt** returns protocol-specific parameters associated with the connect request, **udata** returns any user data sent by the caller on the connect request, and **sequence** is a number that uniquely identifies the returned connect indication. The value of **sequence** enables the user to listen for multiple connect indications before responding to any of them.

Since this function returns values for the **addr**, **opt**, and **udata** fields of *call*, the **maxlen** (see **netbuf** in **t_connect**(3N)) field of each must be set before issuing **t_listen( )** to indicate the maximum size of the buffer for each.

By default, **t_listen( )** executes in synchronous mode and waits for a connect indication to arrive before returning to the user. However, if **O_NDELAY** or **O_NONBLOCK** is set (using **t_open( )** or **fcntl( )**), **t_listen( )** executes asynchronously, reducing to a poll for existing connect indications. If none are available, it returns –1 and sets **t_errno** to TNODATA.

RETURN VALUES

t_listen( ) returns 0 on success. On failure **t_listen( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TBUFOVFLW	The number of bytes allocated for an incoming argument is not sufficient to store the value of that argument. The provider's state, as seen by the user, changes to **T_INCON**, and the connect indication information to be returned in *call* is discarded.

TLOOK	An asynchronous event has occurred on this transport endpoint and requires immediate attention.
TNODATA	**O_NDELAY** or **O_NONBLOCK** was set, but no connect indications had been queued.
TNOTSUPPORT	This function is not supported by the underlying transport provider.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_accept(3N), **t_bind**(3N), **t_connect**(3N), **t_open**(3N), **t_rcvconnect**(3N)

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WARNINGS

If a user issues **t_listen( )** in synchronous mode on a transport endpoint that was not bound for listening (that is, **qlen** was zero on **t_bind( )**), the call will wait forever because no connect indications will arrive on that endpoint.

NOTES

This interface is safe in multithreaded applications.

NAME

t_look – look at the current event on a transport endpoint

SYNOPSIS

cc [*flag* ...] *file* ... –**lnsl** [*library* ...]

#include <tiuser.h>

int t_look(int *fildes***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function returns the current event on the transport endpoint specified by *fildes*. This function enables a transport provider to notify a transport user of an asynchronous event when the user is issuing functions in synchronous mode. Certain events require immediate notification of the user and are indicated by a specific error, **TLOOK**, on the current or next function to be executed.

This function also enables a transport user to poll a transport endpoint periodically for asynchronous events.

RETURN VALUES

Upon success, **t_look( )** returns a value that indicates which of the allowable events has occurred. Otherwise, **t_look( )** returns zero if no event exists. One of the following events is returned:

T_LISTEN	connection indication received
T_CONNECT	connect confirmation received
T_DATA	normal data received
T_EXDATA	expedited data received
T_DISCONNECT	disconnect received
T_UDERR	datagram error indication
T_ORDREL	orderly release indication

On failure, –1 is returned, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_open(3N)

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NOTES

This interface is safe in multithreaded applications.

NAME

t_open – establish a transport endpoint

SYNOPSIS

cc [*flag* ...] *file* ... –**lnsl** [*library* ...]

```
#include <tiuser.h>
#include <fcntl.h>

int t_open(const char *path, int oflag, struct t_info *info);
```

MT-LEVEL

MT-Safe

DESCRIPTION

t_open( ) must be called as the first step in the initialization of a transport endpoint. This function establishes a transport endpoint by opening a file that identifies a particular transport provider (that is, transport protocol) and returning a file descriptor that identifies that endpoint. For example, opening the file **/dev/iso_cots** identifies an OSI connection-oriented transport layer protocol as the transport provider.

path points to the path name of the file to open, and *oflag* identifies any open flags (as in **open**(2)). *oflag* may be constructed from **O_NDELAY** or **O_NONBLOCK** **OR**-ed with **O_RDWR**. These flags are defined in the header **<fcntl.h>**. **t_open( )** returns a file descriptor that will be used by all subsequent functions to identify the particular local transport endpoint.

This function also returns various default characteristics of the underlying transport protocol by setting fields in the **t_info** structure. See **t_getinfo**(3N) for a description of the **t_info** structure.

If *info* is set to **NULL** by the transport user, no protocol information is returned by **t_open( )**.

RETURN VALUES

t_open( ) returns a valid file descriptor on success. On failure **t_open( )** returns **–1,** **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to the following:

TBADFLAG	An invalid flag is specified.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

open(2), **t_getinfo**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_optmgmt – manage options for a transport endpoint

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_optmgmt(int *fildes***, const struct t_optmgmt** ∗*req***, struct t_optmgmt** ∗*ret***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **t_optmgmt( )** function enables a transport user to retrieve, verify, or negotiate protocol options with the transport provider. *fildes* identifies a bound transport endpoint.

The *req* and *ret* arguments point to a **t_optmgmt** structure containing the following members:

struct netbuf opt;
long flags;

The **opt** field identifies protocol options and the **flags** field is used to specify the action to take with those options.

The options are represented by a **netbuf** (see **t_connect**(3N); also for **len**, **buf**, and **maxlen**) structure in a manner similar to the address in **t_bind**(3N). *req* is used to request a specific action of the provider and to send options to the provider. **len** specifies the number of bytes in the options, **buf** points to the options buffer, and **maxlen** has no meaning for the *req* argument. The transport provider may return options and flag values to the user through *ret*. For *ret*, **maxlen** specifies the maximum size of the options buffer and **buf** points to the buffer where the options are to be placed. On return, **len** specifies the number of bytes of options returned. **maxlen** has no meaning for the *req* argument, but must be set in the *ret* argument to specify the maximum number of bytes the options buffer can hold. The actual structure and content of the options is imposed by the transport provider.

The **flags** field of *req* can specify one of the following actions:

T_NEGOTIATE This action enables the user to negotiate the values of the options specified in *req* with the transport provider. The provider will evaluate the requested options and negotiate the values, returning the negotiated values through *ret*.

T_CHECK This action enables the user to verify whether the options specified in *req* are supported by the transport provider. On return, the **flags** field of *ret* will have either **T_SUCCESS** or **T_FAILURE** set to indicate to the user whether the options are supported. These flags are only meaningful for the **T_CHECK** request.

T_DEFAULT This action enables a user to retrieve the default options supported by the transport provider into the **opt** field of *ret*. In *req*, the **len** field of **opt** must be zero and the **buf** field may be NULL.

If issued as part of the connectionless-mode service, **t_optmgmt()** may block due to flow control constraints. The function will not complete until the transport provider has processed all previously sent data units.

RETURN VALUES

t_optmgmt() returns 0 on success. On failure **t_optmgmt()** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TACCES	The user does not have permission to negotiate the specified options.
TBADFLAG	An invalid flag was specified.
TBADOPT	The specified protocol options were in an incorrect format or contained illegal information.
TBUFOVFLW	The number of bytes allowed for an incoming argument is not sufficient to store the value of that argument. The information to be returned in *ret* will be discarded.
TOUTSTATE	The function was issued in the wrong sequence.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_bind(3N), **t_connect**(3N), **t_getinfo**(3N), **t_open**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_rcv – receive data or expedited data sent over a connection

SYNOPSIS

cc [*flag* ...] *file* ... –**lnsl** [*library* ...]

int t_rcv(int *fildes***, char** **buf***, unsigned** *nbytes***, int** **flags***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function receives either normal or expedited data. *fildes* identifies the local transport endpoint through which data will arrive, *buf* points to a receive buffer where user data will be placed, and *nbytes* specifies the size of the receive buffer. *flags* may be set on return from **t_rcv( )** and specifies optional flags as described below.

By default, **t_rcv( )** operates in synchronous mode and will wait for data to arrive if none is currently available. However, if **O_NDELAY** or **O_NONBLOCK** is set (using **t_open**(3N) or **fcntl**(2)), **t_rcv( )** will execute in asynchronous mode and will fail if no data is available. (See **TNODATA** below.)

On return from the call, if **T_MORE** is set in *flags*, this indicates that there is more data and the current transport service data unit (**TSDU**) or expedited transport service data unit (**ETSDU**) must be received in multiple **t_rcv( )** calls. Each **t_rcv( )** with the **T_MORE** flag set indicates that another **t_rcv( )** must follow to get more data for the current **TSDU**. The end of the **TSDU** is identified by the return of a **t_rcv( )** call with the **T_MORE** flag not set. If the transport provider does not support the concept of a **TSDU** as indicated in the *info* argument on return from **t_open**(3N) or **t_getinfo**(3N), the **T_MORE** flag is not meaningful and should be ignored.

On return, the data returned is expedited data if **T_EXPEDITED** is set in *flags*. If the number of bytes of expedited data exceeds *nbytes*, **t_rcv( )** will set **T_EXPEDITED** and **T_MORE** on return from the initial call. Subsequent calls to retrieve the remaining **ETSDU** will have **T_EXPEDITED** set on return. The end of the **ETSDU** is identified by the return of a **t_rcv( )** call with the **T_MORE** flag not set.

If expedited data arrives after part of a **ETSDU** has been retrieved, receipt of the remainder of the **TSDU** will be suspended until the **ETSDU** has been processed. Only after the full **ETSDU** has been retrieved (**T_MORE** not set) will the remainder of the **TSDU** be available to the user.

RETURN VALUES

On successful completion, **t_rcv( )** returns the number of bytes received. On failure **t_rcv( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TLOOK	An asynchronous event has occurred on this transport endpoint and requires immediate attention.
TNODATA	**O_NDELAY** or **O_NONBLOCK** was set, but no data is currently available from the transport provider.

TNOTSUPPORT	This function is not supported by the underlying transport provider.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

fcntl(2), **t_getinfo**(3N), **t_open**(3N), **t_snd**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_rcvconnect – receive the confirmation from a connect request

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_rcvconnect(int *fildes***, struct t_call** ∗*call***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function enables a calling transport user to determine the status of a previously sent connect request and is used in conjunction with **t_connect**(3N) to establish a connection in asynchronous mode. The connection will be established on successful completion of this function.

fildes identifies the local transport endpoint where communication will be established, and *call* contains information associated with the newly established connection. *call* points to a **t_call** structure which contains the following members:

```
struct netbuf     addr;
struct netbuf     opt;
struct netbuf     udata;
int               sequence;
```

netbuf is described in **t_connect**(3N). In *call*, **addr** returns the protocol address associated with the responding transport endpoint, **opt** presents any protocol-specific information associated with the connection, **udata** points to optional user data that may be returned by the destination transport user during connection establishment, and **sequence** has no meaning for this function.

The **maxlen** (see **netbuf** in **t_connect**(3N)) field of each argument must be set before issuing this function to indicate the maximum size of the buffer for each. However, *call* may be NULL, in which case no information is given to the user on return from **t_rcvconnect( )**. By default, **t_rcvconnect( )** executes in synchronous mode and waits for the connection to be established before returning. On return, the **addr**, **opt**, and **udata** fields reflect values associated with the connection.

If **O_NDELAY** or **O_NONBLOCK** is set (using **t_open**(3N) or **fcntl**(2)), **t_rcvconnect( )** executes in asynchronous mode, and reduces to a poll for existing connect confirmations. If none are available, **t_rcvconnect( )** fails and returns immediately without waiting for the connection to be established. (See **TNODATA** below.) **t_rcvconnect( )** must be re-issued at a later time to complete the connection establishment phase and retrieve the information returned in *call*.

RETURN VALUES

t_rcvconnect( ) returns 0 on success. On failure **t_rcvconnect( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TBUFOVFLW	The number of bytes allocated for an incoming argument is not sufficient to store the value of that argument and the connect information to be returned in *call* will be discarded. The provider's state, as seen by the user, will be changed to **DATAXFER**.
TLOOK	An asynchronous event has occurred on this transport connection and requires immediate attention.
TNODATA	**O_NDELAY** or **O_NONBLOCK** was set, but a connect confirmation has not yet arrived.
TNOTSUPPORT	This function is not supported by the underlying transport provider.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

fcntl(2), **t_accept**(3N), **t_bind**(3N), **t_connect**(3N), **t_listen**(3N), **t_open**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_rcvdis – retrieve information from disconnect

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_rcvdis(int *fildes***, struct t_discon** ∗*discon***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function is used to identify the cause of a disconnect, and to retrieve any user data sent with the disconnect. *fildes* identifies the local transport endpoint where the connection existed, and *discon* points to a **t_discon** structure containing the following members:

```
struct netbuf    udata;
int              reason;
int              sequence;

/*
    struct netbuf {
      unsigned int    maxlen;
      unsigned int    len;
      char            *buf;
    };
*/
```

reason specifies the reason for the disconnect through a protocol-dependent reason code, **udata** identifies any user data that was sent with the disconnect, and **sequence** may identify an outstanding connect indication with which the disconnect is associated. **sequence** is only meaningful when **t_rcvdis( )** is issued by a passive transport user who has executed one or more **t_listen**(3N) functions and is processing the resulting connect indications. If a disconnect indication occurs, **sequence** can be used to identify which of the outstanding connect indications is associated with the disconnect.

If a user does not care if there is incoming data and does not need to know the value of **reason** or **sequence**, *discon* may be NULL and any user data associated with the disconnect will be discarded. However, if a user has retrieved more than one outstanding connect indication (using **t_listen**(3N)) and *discon* is NULL, the user will be unable to identify which connect indication the disconnect is associated with.

RETURN VALUES

t_rcvdis returns 0 on success. On failure **t_rcvdis** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF The specified file descriptor does not refer to a transport endpoint.

TBUFOVFLW	The number of bytes allocated for incoming data is not sufficient to store the data. The provider's state, as seen by the user, will change to **T_IDLE**, and the disconnect indication information to be returned in *discon* will be discarded.
TNODIS	No disconnect indication currently exists on the specified transport endpoint.
TNOTSUPPORT	This function is not supported by the underlying transport provider.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_connect(3N), **t_listen**(3N), **t_open**(3N), **t_snddis**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_rcvrel – acknowledge receipt of an orderly release indication

SYNOPSIS

cc [*flag* ...] *file* ... –**lnsl** [*library* ...]

#include <tiuser.h>

int t_rcvrel(int *fildes***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function is used to acknowledge receipt of an orderly release indication. *fildes* identifies the local transport endpoint where the connection exists. After receipt of this indication, the user should not attempt to receive more data because such an attempt will block forever. However, the user may continue to send data over the connection if **t_sndrel( )** has not been issued by the user.

This function is an optional service of the transport provider, and is only supported if the transport provider returned service type **T_COTS_ORD** on **t_open( )** or **t_getinfo( )**.

RETURN VALUES

t_rcvrel( ) returns 0 on success. On failure **t_rcvrel( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TLOOK	An asynchronous event has occurred on this transport endpoint and requires immediate attention.
TNOREL	No orderly release indication currently exists on the specified transport endpoint.
TNOTSUPPORT	This function is not supported by the underlying transport provider.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_open(3N), **t_sndrel**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_rcvudata – receive a data unit

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_rcvudata(int *fildes***, struct t_unitdata** ∗*unitdata***, int** ∗*flags***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function is used in connectionless mode to receive a data unit from another transport user. *fildes* identifies the local transport endpoint through which data will be received, *unitdata* holds information associated with the received data unit, and *flags* is set on return to indicate that the complete data unit was not received. *unitdata* points to a **t_unitdata** structure containing the following members:

```
struct netbuf    addr;
struct netbuf    opt;
struct netbuf    udata;
```

The **maxlen** (see **netbuf** in **t_connect**(3N)) field of **addr**, **opt**, and **udata** must be set before issuing this function to indicate the maximum size of the buffer for each.

On return from this call, **addr** specifies the protocol address of the sending user, **opt** identifies protocol-specific options that were associated with this data unit, and **udata** specifies the user data that was received.

By default, **t_rcvudata( )** operates in synchronous mode and will wait for a data unit to arrive if none is currently available. However, if **O_NDELAY** or **O_NONBLOCK** is set (using **t_open**(3N) or **fcntl**(2)), **t_rcvudata( )** will execute in asynchronous mode and will fail if no data units are available.

If the buffer defined in the **udata** field of *unitdata* is not large enough to hold the current data unit, the buffer will be filled and **T_MORE** will be set in *flags* on return to indicate that another **t_rcvudata( )** should be issued to retrieve the rest of the data unit. Subsequent **t_rcvudata( )** call(s) will return zero for the length of the address and options until the full data unit has been received.

RETURN VALUES

t_rcvudata( ) returns 0 on successful completion. On failure **t_rcvudata( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TBUFOVFLW	The number of bytes allocated for the incoming protocol address or options is greater than zero but not sufficient to store the information. The unit data information to be returned in *unitdata* will be discarded.
TLOOK	An asynchronous event has occurred on this transport endpoint and requires immediate attention.

TNODATA	**O_NDELAY** or **O_NONBLOCK** was set, but no data units are currently available from the transport provider.
TNOTSUPPORT	This function is not supported by the underlying transport provider.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

fcntl(2), **t_connect**(3N), **t_rcvuderr**(3N), **t_sndudata**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_rcvuderr – receive a unit data error indication

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_rcvuderr(int *fildes***, struct t_uderr** ∗*uderr***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function is used in connectionless mode to receive information concerning an error on a previously sent data unit, and should be issued only after a unit data error indication. It informs the transport user that a data unit with a specific destination address and protocol options produced an error. *fildes* identifies the local transport endpoint through which the error report will be received, and *uderr* points to a **t_uderr** structure containing the following members:

```
struct netbuf    addr;
struct netbuf    opt;
long             error;
```

netbuf is described in **t_connect**(3N). The **maxlen** (see **netbuf** in **t_connect**(3N)) field of **addr** and **opt** must be set before issuing this function to indicate the maximum size of the buffer for each.

On return from this call, the **addr** structure specifies the destination protocol address of the erroneous data unit, the **opt** structure identifies protocol-specific options that were associated with the data unit, and **error** specifies a protocol-dependent error code.

If the user does not care to identify the data unit that produced an error, *uderr* may be set to NULL and **t_rcvuderr( )** will simply clear the error indication without reporting any information to the user.

RETURN VALUES

t_rcvuderr( ) returns 0 on successful completion. On failure **t_rcvuderr( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TBUFOVFLW	The number of bytes allocated for the incoming protocol address or options is not sufficient to store the information. The unit data error information to be returned in *uderr* will be discarded.
TNOTSUPPORT	This function is not supported by the underlying transport provider.
TNOUDERR	No unit data error indication currently exists on the specified transport endpoint.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_connect(3N), **t_rcvudata**(3N), **t_sndudata**(3N)
Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_snd – send data or expedited data over a connection

SYNOPSIS

cc [*flag* ...] *file* ... –**lnsl** [*library* ...]

#include <tiuser.h>

int t_snd(int *fildes***, char** ∗*buf***, unsigned** *nbytes***, int** *flags***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function is used to send either normal or expedited data. *fildes* identifies the local transport endpoint over which data should be sent, *buf* points to the user data, *nbytes* specifies the number of bytes of user data to be sent, and *flags* specifies any optional flags described below.

By default, **t_snd( )** operates in synchronous mode and may wait if flow control restrictions prevent the data from being accepted by the local transport provider at the time the call is made. However, if **O_NDELAY** or **O_NONBLOCK** is set (using **t_open**(3N) or **fcntl**(2)), **t_snd( )** will execute in asynchronous mode, and will fail immediately if there are flow control restrictions.

Even when there are no flow control restrictions, **t_snd( )** will wait if **STREAMS** internal resources are not available, regardless of the state of **O_NDELAY** or **O_NONBLOCK**.

On successful completion, **t_snd( )** returns the number of bytes accepted by the transport provider. Normally this will equal the number of bytes specified in *nbytes*. However, if **O_NDELAY** or **O_NONBLOCK** is set, it is possible that only part of the data will be accepted by the transport provider. In this case, **t_snd( )** will set **T_MORE** for the data that was sent (see below) and will return a value less than *nbytes*. If *nbytes* is zero and sending of zero bytes is not supported by the underlying transport provider, **t_snd( )** will return –1 with **t_errno** set to **TBADDATA**. A return value of zero indicates that the request to send a zero-length data message was sent to the provider.

If **T_EXPEDITED** is set in *flags*, the data will be sent as expedited data, and will be subject to the interpretations of the transport provider.

If **T_MORE** is set in *flags*, or is set as described above, an indication is sent to the transport provider that the transport service data unit (**TSDU**) or expedited transport service data unit (**ETSDU**) is being sent through multiple **t_snd( )** calls. Each **t_snd( )** with the **T_MORE** flag set indicates that another **t_snd( )** will follow with more data for the current **TSDU**. The end of the **TSDU** (or **ETSDU**) is identified by a **t_snd( )** call with the **T_MORE** flag not set. Use of **T_MORE** enables a user to break up large logical data units without losing the boundaries of those units at the other end of the connection. The flag implies nothing about how the data is packaged for transfer below the transport interface. If the transport provider does not support the concept of a **TSDU** as indicated in the *info* argument on return from **t_open**(3N) or **t_getinfo**(3N), the **T_MORE** flag is not meaningful and should be ignored.

The size of each **TSDU** or **ETSDU** must not exceed the limits of the transport provider as returned by **t_open**(3N) or **t_getinfo**(3N). If the size is exceeded, a **TSYSERR** with system error **EPROTO** will occur. However, the **t_snd( )** may not fail because **EPROTO** errors may not be reported immediately. In this case, a subsequent call that accesses the transport endpoint will fail with the associated **TSYSERR**.

If **t_snd( )** is issued from the **T_IDLE** state, the provider may silently discard the data. If **t_snd( )** is issued from any state other than **T_DATAXFER**, **T_INREL**, or **T_IDLE**, the provider will generate a **TSYSERR** with system error **EPROTO** (which may be reported in the manner described above).

RETURN VALUES

On successful completion, **t_snd( )** returns the number of bytes accepted by the transport provider. On failure **t_snd( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADDATA	*nbytes* is zero and sending zero bytes is not supported by the transport provider.
TBADF	The specified file descriptor does not refer to a transport endpoint.
TFLOW	**O_NDELAY** or **O_NONBLOCK** was set, but the flow control mechanism prevented the transport provider from accepting data at this time.
TNOTSUPPORT	This function is not supported by the underlying transport provider.
TSYSERR	A system error (see **intro**(2)) has been detected during execution of this function.

SEE ALSO

fcntl(2), **t_getinfo**(3N), **t_open**(3N), **t_rcv**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_snddis – send user-initiated disconnect request

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_snddis(int *fildes***, struct t_call** ∗*call***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function is used to initiate an abortive release on an already established connection or to reject a connect request. *fildes* identifies the local transport endpoint of the connection, and *call* specifies information associated with the abortive release. *call* points to a **t_call** structure that contains the following members:

```
struct netbuf   addr;
struct netbuf   opt;
struct netbuf   udata;
int             sequence;
```

netbuf is described in **t_connect**(3N). The values in *call* have different semantics, depending on the context of the call to **t_snddis( )**. When rejecting a connect request, *call* must be non-NULL and contain a valid value of **sequence** to identify uniquely the rejected connect indication to the transport provider. The **addr** and **opt** fields of *call* are ignored. In all other cases, *call* need only be used when data is being sent with the disconnect request. The **addr**, **opt**, and **sequence** fields of the **t_call** structure are ignored. If the user does not wish to send data to the remote user, the value of *call* may be NULL.

udata specifies the user data to be sent to the remote user. The amount of user data must not exceed the limits supported by the transport provider as returned in the **discon** field of the *info* argument of **t_open**(3N) or **t_getinfo**(3N). If the *len* field of **udata** is zero, no data will be sent to the remote user.

RETURN VALUES

t_snddis( ) returns 0 on success. On failure **t_snddis( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADDATA The amount of user data specified was not within the bounds allowed by the transport provider. The transport provider's outgoing queue will be flushed, so data may be lost.

TBADF The specified file descriptor does not refer to a transport endpoint.

TBADSEQ An invalid sequence number was specified, or a NULL call structure was specified when rejecting a connect request. The transport provider's outgoing queue will be flushed, so data may be lost.

TLOOK An asynchronous event has occurred on this transport endpoint and requires immediate attention.

TNOTSUPPORT	This function is not supported by the underlying transport provider.
TOUTSTATE	The function was issued in the wrong sequence. The transport provider's outgoing queue may be flushed, so data may be lost.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_connect(3N), **t_getinfo**(3N), **t_listen**(3N), **t_open**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_sndrel – initiate an orderly release

SYNOPSIS

cc [*flag* ...] *file* ... –**lnsl** [*library* ...]

#include <tiuser.h>

int t_sndrel(int *fildes***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function is used to initiate an orderly release of a transport connection and indicates to the transport provider that the transport user has no more data to send. *fildes* identifies the local transport endpoint where the connection exists. After issuing **t_sndrel( )**, the user may not send any more data over the connection. However, a user may continue to receive data if an orderly release indication has not been received.

This function is an optional service of the transport provider, and is only supported if the transport provider returned service type **T_COTS_ORD** on **t_open( )** or **t_getinfo( )**.

If **t_sndrel( )** is issued from an invalid state, the provider will generate an **EPROTO** protocol error; however, this error may not occur until a subsequent reference to the transport endpoint.

RETURN VALUES

t_sndrel( ) returns 0 on success. On failure **t_sndrel( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

TSYSERR – A system error has occurred during execution of this function, **errno** will be set to the specific error.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TFLOW	**O_NDELAY** or **O_NONBLOCK** was set, but the flow control mechanism prevented the transport provider from accepting the function at this time.
TNOTSUPPORT	This function is not supported by the underlying transport provider.

SEE ALSO

t_open(3N), **t_rcvrel**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_sndudata – send a data unit

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

int t_sndudata(int *fildes***, struct t_unitdata** ∗*unitdata***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This function is used in connectionless mode to send a data unit to another transport user. *fildes* identifies the local transport endpoint through which data will be sent, and *unitdata* points to a **t_unitdata** structure containing the following members:

```
struct netbuf    addr;
struct netbuf    opt;
struct netbuf    udata;
```

netbuf is described in **t_connect**(3N). In *unitdata*, **addr** specifies the protocol address of the destination user, **opt** identifies protocol-specific options that the user wants associated with this request, and **udata** specifies the user data to be sent. The user may choose not to specify what protocol options are associated with the transfer by setting the **len** field of **opt** to zero. In this case, the provider may use default options.

If the **len** field of **udata** is zero, and the sending of zero bytes is not supported by the underlying transport provider, **t_sndudata** will return –1 with **t_errno** set to **TBADDATA**.

By default, **t_sndudata( )** operates in synchronous mode and may wait if flow control restrictions prevent the data from being accepted by the local transport provider at the time the call is made. However, if **O_NDELAY** or **O_NONBLOCK** is set (using **t_open**(3N) or **fcntl**(2)), **t_sndudata( )** will execute in asynchronous mode and will fail under such conditions.

If **t_sndudata( )** is issued from an invalid state, or if the amount of data specified in **udata** exceeds the **TSDU** size as returned in the **tsdu** field of the *info* argument of **t_open**(3N) or **t_getinfo**(3N), the provider will generate an **EPROTO** protocol error. (See **TSYSERR** below.) If the state is invalid, this error may not occur until a subsequent reference is made to the transport endpoint.

RETURN VALUES

t_sndudata( ) returns 0 on successful completion. On failure **t_sndudata( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADDATA	**nbytes** is zero and sending zero bytes is not supported by the transport provider.
TBADF	The specified file descriptor does not refer to a transport endpoint.
TFLOW]	**O_NDELAY** or **O_NONBLOCK** was set, but the flow control mechanism prevented the transport provider from accepting data at this time.

TNOTSUPPORT	This function is not supported by the underlying transport provider.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

fcntl(2), **t_connect**(3N), **t_getinfo**(3N), **t_rcvudata**(3N), **t_rcvuderr**(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_strerror – get error message string

SYNOPSIS

cc [*flag* ...] *file* ... **–lnsl** [*library* ...]

#include <tiuser.h>

char ∗**t_strerror(int** *errnum***);**

MT-LEVEL

Unsafe

DESCRIPTION

The **t_strerror( )** function maps the error number in *errnum* that corresponds to a TLI error to a language-independent error message string and returns a pointer to that string. The string pointed to will not be modified by the program, but may be overwritten by a subsequent call to the **t_strerror( )** function. The string is not terminated by a newline character. The language for the error message strings written by **t_strerror( )** is implementation-defined. If an error code is unknown and the language is English, **t_strerror( )** returns the string:

"**<error>: error unknown**"

where **<error>** is the error number supplied as input. In other languages, an equivalent text is provided.

SEE ALSO

gettext(3I), **perror**(3C), **setlocale**(3C), **strerror**(3C), **t_error**(3N)

VALID STATES

All – apart from **T_UNINIT**.

NOTES

If the application is linked with **–lintl**, then messages returned from this function are in the native language specified by the **LC_MESSAGES** locale category; see **setlocale**(3C).

NAME

t_sync – synchronize transport library

SYNOPSIS

cc [*flag* ...] *file* ... –**lnsl** [*library* ...]

#include <tiuser.h>

int t_sync(int *fildes***);**

MT-LEVEL

MT-Safe

DESCRIPTION

For the transport endpoint specified by *fildes*, **t_sync( )** synchronizes the data structures managed by the transport library with information from the underlying transport provider. In doing so, it can convert a raw file descriptor (obtained using **open**(2), **dup**(2), or as a result of a **fork**(2) and **exec**(2)) to an initialized transport endpoint, assuming that file descriptor referenced a transport provider. This function also allows two cooperating processes to synchronize their interaction with a transport provider.

For example, if a process **fork**(2)S a new process and issues an **exec**(2), the new process must issue a **t_sync( )** to build the private library data structure associated with a transport endpoint and to synchronize the data structure with the relevant provider information.

It is important to remember that the transport provider treats all users of a transport endpoint as a single user. If multiple processes are using the same endpoint, they should coordinate their activities so as not to violate the state of the provider. **t_sync( )** returns the current state of the provider to the user, thereby enabling the user to verify the state before taking further action. This coordination is only valid among cooperating processes; it is possible that a process or an incoming event could change the provider's state *after* a **t_sync( )** is issued.

If the provider is undergoing a state transition when **t_sync( )** is called, the function will fail.

RETURN VALUES

t_sync( ) returns the state of the transport provider on successful completion. **t_sync( )** returns -1 on failure, **t_errno** is set to indicate the error, and possibly **errno** is set. The state returned may be one of the following:

T_UNBND	unbound
T_IDLE	idle
T_OUTCON	outgoing connection pending
T_INCON	incoming connection pending
T_DATAXFER	data transfer
T_OUTREL	outgoing orderly release (waiting for an orderly release indication)
T_INREL	incoming orderly release (waiting for an orderly release request)

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TSTATECHNG	The transport provider is undergoing a state change.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

dup(2), **exec**(2), **fork**(2), **open**(2)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

t_unbind – disable a transport endpoint

SYNOPSIS

cc [*flag* ...] *file* ... –**lnsl** [*library* ...]

#include <tiuser.h>

int t_unbind(int *fildes***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **t_unbind( )** function disables the transport endpoint specified by *fildes* which was previously bound by **t_bind**(3N). On completion of this call, no further data or events destined for this transport endpoint will be accepted by the transport provider.

RETURN VALUES

t_unbind( ) returns 0 on success. On failure **t_unbind( )** returns –1, **t_errno** is set to indicate the error, and possibly **errno** is set.

ERRORS

On failure, **t_errno** will be set to one of the following:

TBADF	The specified file descriptor does not refer to a transport endpoint.
TLOOK	An asynchronous event has occurred on this transport endpoint.
TOUTSTATE	The function was issued in the wrong sequence.
TSYSERR	A system error has occurred during execution of this function, **errno** will be set to the specific error.

SEE ALSO

t_bind(3N)

Transport Interfaces Programming Guide

NOTES

This interface is safe in multithreaded applications.

NAME

tcsetpgrp – set foreground process group id of terminal

SYNOPSIS

```
#include <unistd.h>
int tcsetpgrp(int fildes, pid_t pgid);
```

MT-LEVEL

MT-Safe

DESCRIPTION

tcsetpgrp( ) sets the foreground process group ID of the terminal specified by *fildes* to *pgid*. The file associated with *fildes* must be the controlling terminal of the calling process and the controlling terminal must be currently associated with the session of the calling process. The value of *pgid* must match a process group ID of a process in the same session as the calling process.

RETURN VALUES

Upon successful completion, **tcsetpgrp( )** returns a value of 0. Otherwise, a value of –1 is returned and **errno** is set to indicate the error.

ERRORS

tcsetpgrp( ) fails if one or more of the following is true:

EBADF The *fildes* argument is not a valid file descriptor.

EINVAL The *fildes* argument is a terminal that does not support **tcsetpgrp( )**, or *pgid* is not a valid process group ID.

ENOTTY The calling process does not have a controlling terminal, or the file is not the controlling terminal, or the controlling terminal is no longer associated with the session of the calling process.

EPERM *pgid* does not match the process group ID of an existing process in the same session as the calling process.

SEE ALSO

termio(7I)

NAME

termios, tcgetattr, tcsetattr, tcsendbreak, tcdrain, tcflush, tcflow, cfgetospeed, cfgetispeed, cfsetispeed, cfsetospeed, tcgetpgrp, tcsetpgrp, tcgetsid – general terminal interface

SYNOPSIS

#include <termios.h>

int tcgetattr(int *fildes***, struct termios** ∗*termios_p***);**

int tcsetattr(int *fildes***, int** *optional_actions***, const struct termios** ∗*termios_p***);**

int tcsendbreak(int *fildes***, int** *duration***);**

int tcdrain(int *fildes***);**

int tcflush(int *fildes***, int** *queue_selector***);**

int tcflow(int *fildes***, int** *action***);**

speed_t cfgetospeed(const struct termios ∗*termios_p***);**

int cfsetospeed(struct termios ∗*termios_p***, speed_t** *speed***);**

speed_t cfgetispeed(const struct termios ∗*termios_p***);**

int cfsetispeed(struct termios ∗*termios_p***, speed_t** *speed***);**

#include <sys/types.h>
#include <termios.h>

pid_t tcgetpgrp(int *fildes***);**

int tcsetpgrp(int *fildes***, pid_t** *pgid***);**

pid_t tcgetsid(int *fildes***);**

MT-LEVEL

MT-Safe

tcgetattr( ), **tcsetattr( )**, **tcsendbreak( )**, **tcdrain( )**, **tcflush( )**, **tcflow( )**, **cfgetospeed( )**, **cfgetispeed( )**, **cfsetispeed( )**, **cfsetospeed( )**, **tcgetpgrp( )**, and **tcsetpgrp( )** are Async-Signal-Safe

DESCRIPTION

These functions describe a general terminal interface for controlling asynchronous communications ports. A more detailed overview of the terminal interface can be found in **termio**(7I), which also describes an **ioctl**(2) interface that provides the same functionality. However, the function interface described here is the preferred user interface.

Many of the functions described here have a *termios_p* argument that is a pointer to a **ter-**

mios structure. This structure contains the following members:

```
tcflag_t    c_iflag;    /* input modes */
tcflag_t    c_oflag;    /* output modes */
tcflag_t    c_cflag;    /* control modes */
tcflag_t    c_lflag;    /* local modes */
cc_t        c_cc[NCCS]; /* control chars */
```

These structure members are described in detail in **termio**(7I).

Get and Set Terminal Attributes

The **tcgetattr( )** function gets the parameters associated with the object referred by *fildes* and stores them in the **termios** structure referenced by *termios_p*. This function may be invoked from a background process; however, the terminal attributes may be subsequently changed by a foreground process.

The **tcsetattr( )** function sets the parameters associated with the terminal (unless support is required from the underlying hardware that is not available) from the **termios** structure referenced by *termios_p* as follows:

If *optional_actions* is **TCSANOW**, the change occurs immediately.

If *optional_actions* is **TCSADRAIN**, the change occurs after all output written to *fildes* has been transmitted. This function should be used when changing parameters that affect output.

If *optional_actions* is **TCSAFLUSH**, the change occurs after all output written to the object referred by *fildes* has been transmitted, and all input that has been received but not read is discarded before the change is made.

The symbolic constants for the values of *optional_actions* are defined in **<termios.h>**.

Line Control

If the terminal is using asynchronous serial data transmission, the **tcsendbreak( )** function causes transmission of a continuous stream of zero-valued bits for a specific duration. If *duration* is zero, it causes transmission of zero-valued bits for at least 0.25 seconds, and not more than 0.5 seconds. If *duration* is not zero, it behaves in a way similar to **tcdrain( )**.

If the terminal is not using asynchronous serial data transmission, the **tcsendbreak( )** function sends data to generate a break condition or returns without taking any action.

The **tcdrain( )** function waits until all output written to the object referred to by *fildes* has been transmitted.

The **tcflush( )** function discards data written to the object referred to by *fildes* but not transmitted, or data received but not read, depending on the value of *queue_selector*:

If *queue_selector* is **TCIFLUSH**, it flushes data received but not read.

If *queue_selector* is **TCOFLUSH**, it flushes data written but not transmitted.

If *queue_selector* is **TCIOFLUSH**, it flushes both data received but not read, and data written but not transmitted.

The **tcflow( )** function suspends transmission or reception of data on the object referred to by *fildes*, depending on the value of *action*:

If *action* is **TCOOFF**, it suspends output.

If *action* is **TCOON**, it restarts suspended output.

If *action* if **TCIOFF**, the system transmits a STOP character, which causes the terminal device to stop transmitting data to the system.

If *action* is **TCION**, the system transmits a START character, which causes the terminal device to start transmitting data to the system.

Get and Set Baud Rate

The baud rate functions get and set the values of the input and output baud rates in the **termios** structure. The effects on the terminal device described below do not become effective until the **tcsetattr( )** function is successfully called.

The input and output baud rates are stored in the **termios** structure. The values shown in the table are supported. The names in this table are defined in <**termios.h**>.

Name	Description	Name	Description
B0	Hang up	**B4800**	4800 baud
B50	50 baud	**B9600**	9600 baud
B75	75 baud	**B19200**	19200 baud
B110	110 baud	**B38400**	38400 baud
B134	134.5 baud	**B57600**	57600 baud
B150	150 baud	**B76800**	76800 baud
B200	200 baud	**B115200**	115200 baud
B300	300 baud	**B153600**	153600 baud
B600	600 baud	**B230400**	230400 baud
B1200	1200 baud	**B307200**	307200 baud
B1800	1800 baud	**B460800**	460800 baud
B24000	24000 baud		

cfgetospeed( ) returns the output baud rate stored in the **termios** structure pointed to by *termios_p*.

cfsetospeed( ) sets the output baud rate stored in the **termios** structure pointed to by *termios_p* to *speed*. The zero baud rate, **B0**, is used to terminate the connection. If **B0** is specified, the modem control lines are no longer be asserted. Normally, this disconnects the line.

cfgetispeed( ) returns the input baud rate stored in the **termios** structure pointed to by *termios_p*.

cfsetispeed( ) sets the input baud rate stored in the **termios** structure pointed to by *termios_p* to *speed*. If the input baud rate is set to zero, the input baud rate is specified by the value of the output baud rate. Both **cfsetispeed( )** and **cfsetospeed( )** return a value of zero if successful and **–1** to indicate an error. This refers both to changes to baud rates not supported by the hardware, and to changes setting the input and output baud rates to

different values if the hardware does not support this.

Get and Set Terminal Foreground Process Group ID

tcsetpgrp( ) sets the foreground process group ID of the terminal specified by *fildes* to *pgid*. The file associated with *fildes* must be the controlling terminal of the calling process and the controlling terminal must be currently associated with the session of the calling process. *pgid* must match a process group ID of a process in the same session as the calling process.

tcgetpgrp( ) returns the foreground process group ID of the terminal specified by *fildes*. **tcgetpgrp( )** is allowed from a process that is a member of a background process group; however, the information may be subsequently changed by a process that is a member of a foreground process group.

Get Terminal Session ID

tcgetsid( ) returns the session ID of the terminal specified by *fildes*.

RETURN VALUES

On success, **tcgetpgrp( )** returns the process group ID of the foreground process group associated with the specified terminal. Otherwise, it returns **–1** and sets **errno** to indicate the error.

On success, **tcgetsid( )** returns the session ID associated with the specified terminal. Otherwise, it returns **–1** and sets **errno** to indicate the error.

On success, all other functions return a value of **0**. Otherwise, they return **–1** and set **errno** to indicate the error.

ERRORS

All of the functions fail if one of more of the following is true:

EBADF The *fildes* argument is not a valid file descriptor.

ENOTTY The file associated with *fildes* is not a terminal.

tcsetattr( ) also fails if the following is true:

EINVAL The *optional_actions* argument is not a proper value, or an attempt was made to change an attribute represented in the **termios** structure to an unsupported value.

tcsendbreak( ) also fails if the following is true:

EINVAL The device does not support the **tcsendbreak( )** function.

tcdrain( ) also fails if one or more of the following is true:

EINTR A signal interrupted the **tcdrain( )** function.

EINVAL The device does not support the **tcdrain( )** function.

tcflush( ) also fails if the following is true:

EINVAL The device does not support the **tcflush( )** function or the *queue_selector* argument is not a proper value.

tcflow( ) also fails if the following is true:

EINVAL The device does not support the **tcflow( )** function or the *action* argument is not a proper value.

tcgetpgrp() also fails if the following is true:

ENOTTY the calling process does not have a controlling terminal, or *fildes* does not refer to the controlling terminal.

tcsetpgrp() also fails if the following is true:

EINVAL *pgid* is not a valid process group ID .

ENOTTY the calling process does not have a controlling terminal, or *fildes* does not refer to the controlling terminal, or the controlling terminal is no longer associated with the session of the calling process.

EPERM *pgid* does not match the process group of an existing process in the same session as the calling process.

tcgetsid() also fails if the following is true:

EACCES *fildes* is a terminal that is not allocated to a session.

SEE ALSO **setpgid**(2), **getsid**(2), **termio**(7I)

NAME

thr_main – identify the main thread

SYNOPSIS

cc [*flag* ...] *file* ... **–lthread** [*library* ...]

```
#include <thread.h>
int thr_main(void)
```

MT-LEVEL

MT-Safe

DESCRIPTION

thr_main( ) returns:

- **1** if the calling thread is the main thread.
- **0** if the calling thread is not the main thread.
- **-1** if **libthread** is not linked in or thread initialization has not completed.

FILES

/lib/libthread

SEE ALSO

thr_self(3T)

NAME

thr_min_stack – returns the minimum-allowable size for a thread's stack

SYNOPSIS

cc [*flag* ...] *file* ... **–lthread** [*library* ...]

#include <thread.h>

size_t thr_min_stack(*void***);**

MT-LEVEL

MT-Safe

DESCRIPTION

When a thread is created with a user-supplied stack, the user must reserve enough space to run this thread. In a dynamically linked execution environment, it is very hard to know what the minimum stack requirments are for a thread. The function **thr_min_stack()** returns the amount of space needed to execute a null thread. This is a thread that was created to execute a null procedure. A thread that does something useful should have a stack size that is **thr_min_stack()** + <*some increment*>.

Most users should not be creating threads with user-supplied stacks. This functionality was provided to support applications that wanted complete control over their execution environment.

Typically, users should let the threads library manage stack allocation. The threads library provides default stacks which should meet the requirements of any created thread.

thr_min_stack( ) will return the unsigned int **THR_MIN_STACK,** which is the minimum-allowable size for a thread's stack.

In this implementation the default size for a user-thread's stack is one mega-byte. If the second argument to **thr_create**(3T) is NULL, then the default stack size for the newly-created thread will be used. Otherwise, you may specify a stack-size that is at least **THR_MIN_STACK,** yet less than the size of your machine's virtual memory.

It is recommended that the default stack size be used.

To determine the smallest-allowable size for a thread's stack, execute the following:

```
/* cc thisfile.c -lthread */
#define _REENTRANT
#include <thread.h>
#include <stdio.h>

main() {
     printf("thr_min_stack() returns %u \n",thr_min_stack());
}
```

SEE ALSO

pthread_attr_init(3T), **pthread_create**(3T)

NOTES

Although the POSIX threads implementation, pthreads, does not have a corresponding function to **thr_min_stack( ),** it does implement a minimum stack size, whose value is **PTHREAD_STACK_MIN,** which may be ascertained as follows:

```
/* cc thisfile.c -lpthread */
#define _REENTRANT
#include <pthread.h>
#include <stdio.h>

main() {
      printf("minimum POSIX stack size is  %u\n", PTHREAD_STACK_MIN);
}
```

NAME

thr_setconcurrency, thr_getconcurrency – get / set thread concurrency level

SYNOPSIS

```
#include <thread.h>
int thr_setconcurrency(int new_level);
int thr_getconcurrency(void );
```

MT-LEVEL

MT-Safe

DESCRIPTION

Unbound threads in a process (see **thr_create**(3T)) may or may not be required to be simultaneously active. By default, the threads system ensures that a sufficient number of threads are active so that the process can continue to make progress. While this conserves system resources, it may not produce the most effective level of concurrency. **thr_setconcurrency( )** permits the application to give the threads system a hint, specified by *new_level*, for the desired level of concurrency. The actual number of simultaneously active threads may be larger or smaller than this number. The value for the desired concurrency level may also be affected by creating threads with the THR_NEW_LWP flag set (see **thr_create**(3T)).

If *new_level* is zero, the threads system will only ensure that a sufficient number of threads are active so that the process can continue to make progress.

thr_getconcurrency( ) returns the current value for the desired concurrency level. The actual number of simultaneously active threads may be larger or smaller than this number.

RETURN VALUES

thr_setconcurrency( ) returns zero when successful. A non-zero value indicates an error code.

thr_getconcurrency( ) always returns the current value for the desired concurrency level.

ERRORS

If any of the following conditions are detected, **thr_setconcurrency( )** fails and returns the corresponding value:

EAGAIN the specified concurrency level would cause a system resource to be exceeded.

EINVAL *new_level* is negative.

SEE ALSO

thr_create(3T).

NOTES

The Solaris threads set / get concurrency functionality described on this man page are not implemented in the POSIX threads interface.

NAME

thr_stksegment – get thread stack bottom and stack size

SYNOPSIS

cc [*flag* ...] *file* ... **–lthread** [*library* ...]

#include <thread.h>

#include <sys/signal.h>

int thr_stksegment(stack_t*)

MT-LEVEL

MT-Safe

DESCRIPTION

The stack information provided by **thr_stksegment( )** is typically used by debuggers, garbage collectors, and similar applications. Most applications should not require such information. The bottom of the thread stack returned by **thr_stksegment( )** points to a part of the stack which may contain data maintained by **libthread**. The user's thread stack starts at a point below the bottom of the stack as returned by **thr_stksegment( ).**

RETURN VALUES

thr_stksegment( ) returns **0** if both the thread stack bottom and stack size were successfully retrieved; otherwise, it returns a non-zero error code.

ERRORS

If any of the following conditions are detected, **thr_stksegment( )** fails and returns the corresponding value:

EFAULT A system call used to get the stack information failed because a bad address was passed to it.

EAGAIN The stack information for the thread is not available because the thread's initialization is not yet complete.

SEE ALSO

thr_create(3T)

NAME

thr_suspend, thr_continue – suspend or continue thread execution

SYNOPSIS

```
#include <thread.h>
int thr_suspend(thread_t target_thread);
int thr_continue(thread_t target_thread);
```

MT-LEVEL

MT-Safe

DESCRIPTION

thr_suspend( ) immediately suspends the execution of the thread specified by *target_thread*. On successful return from **thr_suspend( )**, the suspended thread is no longer executing. Once a thread is suspended, subsequent calls to **thr_suspend( )** have no effect.

thr_continue( ) resumes the execution of a suspended thread. Once a suspended thread is continued, subsequent calls to **thr_continue( )** have no effect.

A suspended thread will not be awakened by a signal. The signal stays pending until the execution of the thread is resumed by **thr_continue( )**.

RETURN VALUES

Zero is returned when successful. A non-zero value indicates an error.

ERRORS

If any of the following conditions are detected, **thr_suspend( )** or **thr_continue( )** fails and returns the corresponding value:

ESRCH *target_thread* cannot be found in the current process.

SEE ALSO

thr_create(3T)

NOTES

The are no POSIX counterparts to the Solaris threads suspend and continue functionality described on this man page.

NAME

thr_yield – thread yield to another thread

SYNOPSIS

cc [*flag* ...] *file* ... **–lthread** [*library* ...]

```
#include <thread.h>
void thr_yield(void);
```

MT-LEVEL

MT-Safe

DESCRIPTION

thr_yield( ) causes the current thread to yield its execution in favor of another thread with the same or greater priority.

RETURN VALUES

thr_yield( ) returns nothing and does not set **errno**.

SEE ALSO

sched_yield(3R), **thr_setprio**(3T)

NOTES

Although there is no POSIX "pthreads" counterpart to **thr_yield( )**, the POSIX real-time function, **sched_yield**(3R) does correspond to **thr_yield( )**; thereby affording the POSIX portability to this functionality.

NAME

threads, pthreads, libpthread, libthread – thread libraries: libpthread and libthread

SYNOPSIS

POSIX

cc [*flag* ...] *file* ... –**lpthread** [**-lposix4** *library* ...]
#include <pthread.h>

Solaris

cc [*flag* ...] *file* ... –**lthread** [*library* ...]
#include <thread.h>
#include <sched.h>

MT-LEVEL

Fork1-Safe MT-Safe

DESCRIPTION

Two threads libraries are available, POSIX and Solaris. Both implementations are interoperable, their functionality similar, and can be used within the same application. However, only POSIX threads are guaranteed to be fully portable to other POSIX-compliant environments. As indicated by the "Synopsis" section above, their use requires different source include files and different linking libraries.

Similarities

Most of the functions in both libraries, **libpthread** and **libthread**, have a counterpart in the other's library. POSIX functions and Solaris functions, whose names have similar endings, usually have similar functionality, number of arguments, and use of arguments. i.e.:

POSIX	Solaris
pthread_kill( )	**thr_kill( )**
pthread_sigmask( )	**thr_sigsetmask( )**
pthread_mutex_lock( )	**mutex_lock( )**
sem_wait( )	**sema_wait( )**

All POSIX threads function names begin with the prefix "**pthread**", with semaphore names being the exception.

Differences

POSIX

- is more portable,
- establishes characteristics for each thread according to configurable attribute objects,
- implements thread cancellation,
- enforces scheduling algorithms, and
- allows for clean-up handlers for **fork**(2) calls.

Solaris

- threads can be suspended and continued,
- implements an optimized mutex, reader / writer locking.
- may increase the concurrency, and

- implements daemon threads, for whose demise the process does not wait.

IMPLEMENTATION	POSIX	Solaris
Creation	**pthread_create( )**	**thr_create( )**
	pthread_attr_init( )	—
	pthread_attr_setdetachstate( )	—
	pthread_attr_getdetachstate( )	—
	pthread_attr_setinheritsched( )	—
	pthread_attr_getinheritsched( )	—
	pthread_attr_setschedparam( )	—
	pthread_attr_getschedparam( )	—
	pthread_attr_setschedpolicy( )	—
	pthread_attr_getschedpolicy( )	—
	pthread_attr_setscope( )	—
	pthread_attr_getscope( )	—
	pthread_attr_setstackaddr( )	—
	pthread_attr_getstackaddr( )	—
	pthread_attr_setstacksize( )	—
	pthread_attr_getstacksize( )	—
	pthread_attr_destroy( )	—
	—	**thr_min_stack( )**
Exit	**pthread_exit( )**	**thr_exit( )**
	pthread_join( )	**thr_join( )**
	pthread_detach( )	—
Thread Specific Data	**pthread_key_create( )**	**thr_keycreate( )**
	pthread_setspecific( )	**thr_setspecific( )**
	pthread_getspecific( )	**thr_getspecific( )**
	pthread_key_delete( )	—
Signal	**pthread_sigmask( )**	**thr_sigsetmask( )**
	pthread_kill( )	**thr_kill( )**
ID	**pthread_self( )**	**thr_self( )**
	pthread_equal( )	—
	—	**thr_main( )**
Scheduling	—	**thr_yield( )**
	—	**thr_suspend( )**
	—	**thr_continue( )**
	—	**thr_setconcurrency( )**
	—	**thr_getconcurrency( )**
	pthread_setschedparam( )	**thr_setprio( )**
	pthread_getschedparam( )	**thr_getprio( )**

Cancellation	**pthread_cancel( )**	——
	pthread_setcancelstate( )	——
	pthread_setcanceltype( )	——
	pthread_testcancel( )	——
	pthread_cleanup_pop( )	——
	pthread_cleanup_push( )	——
Mutex	**pthread_mutex_init( )**	**mutex_init( )**
	pthread_mutexattr_init( )	——
	pthread_mutexattr_setpshared( )	——
	pthread_mutexattr_getpshared( )	——
	pthread_mutexattr_setprotocol( )	——
	pthread_mutexattr_getprotocol( )	——
	pthread_mutexattr_setprioceiling( )	——
	pthread_mutexattr_getprioceiling( )	——
	pthread_mutexattr_destroy( )	——
	pthread_mutex_setprioceiling( )	——
	pthread_mutex_getprioceiling( )	——
	pthread_mutex_lock( )	**mutex_lock( )**
	pthread_mutex_trylock( )	**mutex_trylock( )**
	pthread_mutex_unlock( )	**mutex_unlock( )**
	pthread_mutex_destroy( )	**mutex_destroy( )**
Condition Variable	**pthread_cond_init( )**	**cond_init( )**
	pthread_condattr_init( )	——
	pthread_condattr_setpshared( )	——
	pthread_condattr_getpshared( )	——
	pthread_condattr_destroy( )	——
	pthread_cond_wait( )	**cond_wait( )**
	pthread_cond_timedwait( )	**cond_timedwait( )**
	pthread_cond_signal( )	**cond_signal( )**
	pthread_cond_broadcast( )	**cond_broadcast( )**
	pthread_cond_destroy( )	**cond_destroy( )**
Reader/Writer Locking	——	**rwlock_init( )**
	——	**rw_rdlock( )**
	——	**rw_tryrdlock( )**
	——	**rw_wrlock( )**
	——	**rw_trywrlock( )**
	——	**rw_unlock( )**
	——	**rwlock_destroy( )**
Semaphore	**sem_init( )**	**sema_init( )**
	sem_open( )	——
	sem_close( )	——
	sem_wait( )	**sema_wait( )**

	sem_trywait( )	**sema_trywait( )**
	sem_post( )	**sema_post( )**
	sem_getvalue( )	—
	sem_unlink( )	—
	sem_destroy( )	**sema_destroy( )**
fork() Clean Up Handling	**pthread_atfork( )**	—
Limits	**pthread_once( )**	—
Debugging	—	**thr_stksegment( )**

LOCKING

Synchronization

Multi-threaded behavior is asynchronous, and therefore, optimized for concurrent and parallel processing. Since threads, always from within the same process and sometimes from multiple processes, share global data with each other, they are not guaranteed exclusive access to the shared data at any point in time. Securing mutually exclusive access to shared data requires synchronization among the threads. Solaris implements four synchronization mechanisms:

- mutex
- condition variable
- reader / writer locking *(optimized frequent-read occasional-write mutex)*
- semaphore

POSIX implements all but reader / writer locking.

Synchronizing multiple threads diminishes their concurrency. The coarser the grain of synchronization, that is, the larger the block of code that is locked, the lesser the concurrency.

MT fork()

If a multi-threaded program calls **fork**(2), it implicitly calls **fork1**(2), which replicates only the calling thread. Should there be any outstanding mutexes throughout the process, the application should call **pthread_atfork**(3T), to wait for and acquire those mutexes, prior to calling **fork( )**.

FILES

POSIX

/usr/include/pthread.h
/lib/libpthread.*
/lib/libposix4.*

Solaris

/usr/include/thread.h
/usr/include/sched.h
/lib/libthread.*

SEE ALSO **fork**(2), **intro**(3), **pthread_atfork**(3T), **pthread_create**(3T)

ERRORS In a multi-threaded application, linked with **libpthread** or **libthread**, **EINTR** may be returned whenever another thread calls **fork**(2), which calls **fork1**(2) instead.

NAME

timer_create – create a timer

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <signal.h>
#include <time.h>

int timer_create(clockid_t clock_id, struct sigevent *evp, timer_t *timerid);

struct sigevent {
        int             sigev_notify    /* notification type */
        int             sigev_signo;    /* signal number */
        union sigval    sigev_value;    /* signal value */
};

union sigval {
        int             sival_int;      /* integer value */
        void            *sival_ptr;     /* pointer value */
};
```

MT-LEVEL

MT-Safe with exceptions

DESCRIPTION

timer_create( ) creates a timer using the specified clock, *clock_id*, as the timing base. This timer ID is unique and meaningful only within the calling LWP until the timer is deleted. This timer is initially disarmed upon return from **timer_create( ).**

The timer may be created per-LWP or per-process. Expiration signals for a per-LWP timer will be sent to the creating LWP. Expiration signals for a per-process timer will be sent to the process. A per-LWP timer will be automatically deleted when the creating LWP exits. By default, timers are created per-LWP. If the symbol **_POSIX_PER_PROCESS_TIMERS** is defined or the symbol **_POSIX_C_SOURCE** is defined to have a value greater than **199500L** before the inclusion of **<time.h>**, timers will be created per-process.

If *evp* is non-**NULL**:

then *evp* points to a **sigevent** structure, allocated by the application, which defines the asynchronous notification that will occur when the timer expires.

If the **sigev_notify** member of *evp* is **SIGEV_SIGNAL,** then the structure also contains the signal number and the application specific data value to be sent to the process. If **SA_SIGINFO** is set for the expiration signal, then the signal and application-defined value specified in the structure will be queued to the process on timer expiration. If **SA_SIGINFO** is not set for the expiration signal, then the signal specified in the structure will be sent upon the timer expiration.

If the **sigev_notify** member is **SIGEV_NONE,** no notification will be sent.

If *evp* is **NULL,** and **SA_SIGINFO** is set for the expiration signal, then the default signal, **SIGALRM,** will be queued to the process and the signal data value will be set to the timer ID.

RETURN VALUES

timer_create( ) returns **0** upon success and creates a **timer_t,** *timerid*, which can be passed to the timer calls; otherwise it returns **-1** and sets **errno** to indicate the error condition.

ERRORS

EAGAIN	The system lacks sufficient signal queuing resources to honor the request.
	The calling process has already created all of the timers it is allowed by this implementation.
EINVAL	The specified clock ID, *clock_id*, is not defined.
ENOSYS	**timer_create( )** is not supported by this implementation.

SEE ALSO

exec(2), **fork**(2), **time**(2), **clock_settime**(3R), **timer_delete**(3R), **timer_settime**(3R)

NOTES

Timers are not inherited by a child process across a **fork**(2) and can be disarmed and deleted by an **exec**(2).

In a future release, the ability to create per-LWP timers will be removed, and all calls to **timer_create**() will result in per-process timers.

NAME

timer_delete – delete a per-LWP timer

SYNOPSIS

cc [*flag* ...] *file* ... –**lposix4** [*library* ...]

#include <time.h>

int timer_delete(timer_t *timerid***);**

MT-LEVEL

MT-Safe with exceptions

DESCRIPTION

timer_delete( ) deletes the specified timer, *timerid,* previously created by **timer_create**(3R). If the timer is armed when **timer_delete( )** is called, the behavior is as if the timer is automatically disarmed before removal.

RETURN VALUES

timer_delete( ) returns **0** upon success, otherwise it returns **-1** and sets **errno** to indicate the error condition.

ERRORS

EINVAL *timerid* does not refer to a valid timer.

ENOSYS **timer_delete( )** is not supported by this implementation.

SEE ALSO

timer_create(3R)

NAME

timer_settime, timer_gettime, timer_getoverrun – high-resolution timer operations

SYNOPSIS

cc [*flag* ...] *file* ... **–lposix4** [*library* ...]

```
#include <time.h>

int timer_settime(timer_t timerid, int flags, const struct itimerspec *value,
                  struct itimerspec *ovalue);

int timer_gettime(timer_t timerid, struct itimerspec *value);

int timer_getoverrun(timer_t timerid);

struct itimerspec       {
        struct timespec   it_interval;  /* timer period */
        struct timespec   it_value;     /* timer expiration */
};
struct timespec         {
        time_t            tv_sec;       /* seconds */
        long              tv_nsec;      /* and nanoseconds */
};
```

MT-LEVEL

Async-Signal-Safe

DESCRIPTION

If *value*->**it_value** is non-zero, **timer_settime( )** arms the timer, *timerid,* to next expire after the time designated by *value*->**it_value**. Upon expiration, an application-specified notification (see **timer_create**(3R)) or the default signal, **SIGALRM**, is queued for the calling LWP. If *timerid* was already armed when **timer_settime( )** is called, this call resets the time until the next expiration to the value of *value*->**it_value**. If *value*->**it_value** is zero, then the timer is disarmed.

value->**it_value** may be expressed as either an absolute or relative time. If *flags* is set to **TIMER_RELTIME**, then the timer will initially expire relative to when the call is made. If *flags* is set to **TIMER_ABSTIME**, then the initial expiration will be relative to 00:00 Universal Coordinated Time, January 1, 1970. If the specified (absolute) time has already passed, **timer_settime( )** succeeds and the expiration notification is made.

If *value*->**it_interval** is non-zero, then *timerid,* will be a “"periodic” timer, to be reloaded to expire every *value*->**it_interval** seconds (nanoseconds). Otherwise, if *value*->**it_interval** is zero and *value*->**it_value** is non-zero, then *timerid* is a “one-shot” timer, which will expire only at the time specified by *value*->**it_value**.

If *ovalue* is not **NULL**, and timer *timerid* had previously been used, then **timer_settime( )** will store the remaining time until the previous timer expires in *ovalue*->**it_value,** and the previous reload interval in *ovalue*->**it_interval.** (If the previous timer was disarmed, *ovalue*->**it_value** will be set to zero). The values stored in *ovalue* by **timer_settime( )** are the same values that would have been returned by a call to **timer_gettime(** *timerid,***...).**

timer_gettime( ) stores the amount of time until the specified timer, *timerid,* expires into *value*->**it_value,** and the timer’s reload value into *value*->**it_interval.**

Only a single signal can be queued to the LWP for a given timer at any point in time. When a timer, for which a signal is still pending expires, (from a previous interval), no signal will be queued, and a "timer overrun count" will be incremented. When a timer expiration signal is delivered to an LWP, **timer_overrun( )** may be used to determine the timer expiration overrun count for the specified timer. The overrun count returned contains the number of extra timer expirations which occurred between the time the signal was generated (queued) and when it was delivered, up to but not including a maximum of **{DELAYTIMER_MAX}**. If the number of such extra expirations is greater than or equal to **{DELAYTIMER_MAX}**, then the overrun count is set to **{DELAYTIMER_MAX}**. The value returned by **timer_getoverrun( )** applies to the most recent expiration signal delivery for the timer.

RETURN VALUES

timer_settime( ), and **timer_gettime( )** return **0** upon success. If **timer_getoverrun( )** succeeds, the number of extra timer expirations which occurred between the time the signal was queued and when it was delivered is returned. If these functions fail, they return **-1** and set **errno** to indicate the error condition.

ERRORS

EINVAL *timerid* does not correspond to a timer returned by **timer_create**(3R).

The timer, *timerid*, had already been deleted by **timer_delete**(3R).

A *value* structure specified a nanosecond value less than zero or greater than or equal to 1,000,000,000.

ENOSYS **timer_settime( ), timer_gettime( ),** or **timer_getoverrun( )** is not supported by this implementation.

SEE ALSO

clock_settime(3R), **timer_create**(3R), **timer_delete**(3R)

NAME

times – get process times

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

```
#include <sys/param.h>
#include <sys/types.h>
#include <sys/times.h>

int times(tmsp)
register struct tms *tmsp;
```

DESCRIPTION

times( ) returns time-accounting information for the current process and for the terminated child processes of the current process. All times are reported in clock ticks. The number of clock ticks per second is defined by the variable **CLK_TCK**, found in the header **<limits.h>**.

A structure with the following members is returned by **times( )**:

```
time_t   tms_utime;    /* user time */
time_t   tms_stime;    /* system time */
time_t   tms_cutime;   /* user time, children */
time_t   tms_cstime;   /* system time, children */
```

The children's times are the sum of the children's process times and their children's times.

RETURN VALUES

times( ) returns

0	on success.
–1	on failure.

SEE ALSO

time(1), **time**(2), **wait**(2), **getrusage**(3C)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

times( ) has been superseded by **getrusage**(3C).

NAME

tmpfile – create a temporary file

SYNOPSIS

```
#include <stdio.h>
FILE *tmpfile(void);
```

MT-LEVEL

Safe

DESCRIPTION

tmpfile( ) creates a temporary file using a name generated by the **tmpnam( )** routine and returns a corresponding **FILE** pointer. If the file cannot be opened, a **NULL** pointer is returned. The file is automatically deleted when the process using it terminates or when the file is closed. The file is opened for update (''w+'').

SEE ALSO

creat(2), **open**(2), **unlink**(2), **fopen**(3S), **mktemp**(3C), **perror**(3C), **stdio**(3S), **tmpnam**(3S)

NAME

tmpnam, tmpnam_r, tempnam – create a name for a temporary file

SYNOPSIS

```
#include <stdio.h>
char *tmpnam(char *s);
char *tmpnam_r(char *s);
char *tempnam(const char *dir, const char *pfx);
```

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

These functions generate file names that can safely be used for a temporary file.

tmpnam( ) always generates a file name using the path-prefix defined as **P_tmpdir** in the **<stdio.h>** header. If *s* is **NULL**, **tmpnam( )** leaves its result in an internal static area and returns a pointer to that area. The next call to **tmpnam( )** will destroy the contents of the area. If *s* is not **NULL**, it is assumed to be the address of an array of at least **L_tmpnam** bytes, where **L_tmpnam** is a constant defined in **<stdio.h>**; **tmpnam( )** places its result in that array and returns *s*.

tmpnam_r( ) has the same functionality as **tmpnam( )** except that if *s* is a **NULL** pointer, the function returns **NULL.**

tempnam( ) allows the user to control the choice of a directory. The argument *dir* points to the name of the directory in which the file is to be created. If *dir* is **NULL** or points to a string that is not a name for an appropriate directory, the path-prefix defined as **P_tmpdir** in the **<stdio.h>** header is used. If that directory is not accessible, **/tmp** will be used as a last resort. This entire sequence can be up-staged by providing an environment variable **TMPDIR** in the user's environment, whose value is the name of the desired temporary-file directory.

Many applications prefer their temporary files to have certain favorite initial letter sequences in their names. Use the *pfx* argument for this. This argument may be **NULL** or point to a string of up to five characters to be used as the first few characters of the temporary-file name.

tempnam( ) uses **malloc**(3C) to get space for the constructed file name, and returns a pointer to this area. Thus, any pointer value returned from **tempnam( )** may serve as an argument to **free**(3C) (see **malloc**(3C)). If **tempnam( )** cannot return the expected result for any reason—for example, **malloc**(3C) failed—or none of the above mentioned attempts to find an appropriate directory was successful, a **NULL** pointer will be returned.

tempnam( ) fails if there is not enough space.

SEE ALSO

creat(2), **unlink**(2), **fopen**(3S), **free**(3C), **malloc**(3C), **mktemp**(3C), **tmpfile**(3S)

NOTES

The **tmpnam_r()** interface is as proposed in the POSIX.4a Draft #6 document, and is subject to change to be compliant to the standard when it is accepted.

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

These functions generate a different file name each time they are called.

Files created using these functions and either **fopen**(3S) or **creat**(2) are temporary only in the sense that they reside in a directory intended for temporary use, and their names are unique. It is the user's responsibility to remove the file when its use is ended.

If called more than **TMP_MAX** (defined in **<stdio.h>**) times in a single process, these functions start recycling previously used names.

Between the time a file name is created and the file is opened, it is possible for some other process to create a file with the same name. This can never happen if that other process is using these functions or **mktemp**(3C) and the file names are chosen to render duplication by other means unlikely.

tempnam() is safe in multi-thread applications. **tmpnam()** is unsafe in multi-thread applications, **tmpnam_r()** should be used instead.

On Solaris systems, the default value for **P_tmpdir** is **/var/tmp.**

NAME

tnf_process_disable, tnf_process_enable, tnf_thread_disable, tnf_thread_enable – probe control internal interface

SYNOPSIS

cc [*flag* ...] *file* ... –**ltnfprobe** [*library* ...]

#include <tnf/probe.h>

void tnf_process_disable(void);

void tnf_process_enable(void);

void tnf_thread_disable(void);

void tnf_thread_enable(void);

AVAILABILITY

SUNWtnfd

MT-LEVEL

MT-Safe

DESCRIPTION

There are three levels of granularity for controlling tracing and probe functions (called probing from here on) — probing for the entire process, a particular thread, and the probe itself can be disabled / enabled. The first two (process and thread) are controlled by this interface. The probe is controlled via the application **prex**(1).

tnf_process_disable( ) turns off probing for the process. The default process state is to have probing enabled. **tnf_process_enable( )** turns on probing for the process.

tnf_thread_disable( ) turns off probing for the currently running thread. Threads are "born" or created with this state enabled. **tnf_thread_enable( )** turns on probing for the currently running thread. If the program is a non-threaded program, these two thread interfaces disable or enable probing for the process.

SEE ALSO

prex(1), **tnfdump**(1), **TNF_DECLARE_RECORD**(3X), **TNF_PROBE**(3X)

NOTES

A probe is considered enabled only if:

- **prex**(1) has enabled the probe AND
- the process has probing enabled — which is the default or could be set via **tnf_process_enable( )** AND
- the thread that hits the probe has probing enabled — which is every thread's default or could be set via **tnf_thread_enable( ).**

There is a run time cost associated with determining that the probe is disabled. To reduce the performance effect of probes, this cost should be minimized. The quickest way that a probe can be determined to be disabled is by the enable control that **prex**(1) uses. Therefore, to disable all the probes in a process use the **disable** command in **prex**(1) rather than **tnf_process_disable( ).**

tnf_proces_disable( ) and **tnf_process_enable( )** should only be used to toggle probing based on some internal program condition. **tnf_thread_disable( )** should be used to turn off probing for threads that are uninteresting.

NAME

trig, sin, cos, tan, asin, acos, atan, atan2 – trigonometric functions

SYNOPSIS

cc [*flag* ...] *file* ... **–lm** [*library* ...]

#include <math.h>

double sin(double *x***);**

double cos(double *x***);**

double tan(double *x***);**

double asin(double *x***);**

double acos(double *x***);**

double atan(double *x***);**

double atan2(double *y***, double** *x***);**

MT-LEVEL

MT-Safe

DESCRIPTION

sin(*x***), cos(***x***)** and **tan(***x***)** return trigonometric functions of radian arguments. Trigonometric argument reduction is carried out with respect to the infinitely precise π.

asin(*x***)** returns the arc sine of *x* in the range $-\pi/2$ to $\pi/2$.

acos(*x***)** returns the arc cosine of *x* in the range 0 to π.

atan(*x***)** returns the arc tangent of *x* in the range $-\pi/2$ to $\pi/2$.

atan2(*y***,***x***)** and **hypot(***x***,***y***)** (see **hypot**(3M)) convert rectangular coordinates **(***x***,***y***)** to polar **(***r***,**θ**)**; **atan2(***y***,***x***)** computes θ, the argument or phase, by computing an arc tangent of y/x in the range $-\pi$ to π.

RETURN VALUES

For exceptional cases, **matherr**(3M) tabulates the values to be returned as dictated by various Standards.

SEE ALSO

hypot(3M), **matherr**(3M)

DIAGNOSTICS

In IEEE 754 mode (i.e. the **–xlibmieee** cc compilation option), these functions handle exceptional arguments in the spirit of ANSI / IEEE Std 754-1985. **sin(**±∞**)**, **cos(**±∞**)** and **tan(**±∞**)** return NaN; **asin(***x***)** and **acos(***x***)** return NaN if $|x| > 1$. **atan(**±∞**)** returns $\pm\pi/2$. For **atan2( )**,

- **atan2(**±0, *x***)** returns ±0 for $x > 0$ or $x = +0$;
- **atan2(**±0, *x***)** returns ±π for $x < 0$ or $x = -0$;
- **atan2(***y*, ±0**)** returns $\pi/2$ for $y > 0$;
- **atan2(***y*, ±0**)** returns $-\pi/2$ for $y < 0$;
- **atan2(**±*y*, ∞**)** returns ±0 for finite $y > 0$;
- **atan2(**±∞, *x***)** returns $\pm\pi/2$ for finite *x*;
- **atan2(**±*y*, –∞**)** returns ±π for finite $y > 0$;
- **atan2(**±∞, ∞**)** returns $\pm\pi/4$;
- **atan2(**±∞, –∞**)** returns $\pm 3\pi/4$.

NAME

truncate, ftruncate – set a file to a specified length

SYNOPSIS

```
#include <unistd.h>
int truncate(const char *path, off_t length);
int ftruncate(int fildes, off_t length);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The file whose name is given by *path* or referenced by the descriptor *fildes* has its size set to *length* bytes.

If the file was previously longer than *length*, bytes past *length* will no longer be accessible. If it was shorter, bytes from the EOF before the call to the EOF after the call will be read in as zeros. The effective user ID of the process must have write permission for the file, and for **ftruncate( )** the file must be open for writing.

RETURN VALUES

Upon successful completion, a value of 0 is returned. Otherwise, a value of –1 is returned and **errno** is set to indicate the error.

ERRORS

truncate( ) fails if one or more of the following are true:

EACCES	Search permission is denied on a component of the path prefix.
EACCES	Write permission is denied for the file referred to by *path*.
EFAULT	*path* points outside the process's allocated address space.
EINTR	A signal was caught during execution of the **truncate** routine.
EINVAL	*path* is not an ordinary file.
EIO	An I/O error occurred while reading from or writing to the file system.
EISDIR	The file referred to by *path* is a directory.
ELOOP	Too many symbolic links were encountered in translating *path*.
EMFILE	The maximum number of file descriptors available to the process has been reached.
EMULTIHOP	Components of *path* require hopping to multiple remote machines and file system type does not allow it.
ENAMETOOLONG	The length of a *path* component exceeds **{NAME_MAX}** characters, or the length of *path* exceeds **{PATH_MAX}** characters.
ENFILE	Could not allocate any more space for the system file table.
ENOENT	Either a component of the path prefix or the file referred to by *path* does not exist.
ENOLINK	*path* points to a remote machine and the link to that machine is no longer active.
ENOTDIR	A component of the path prefix of *path* is not a directory.

EROFS	The file referred to by *path* resides on a read-only file system.

ftruncate( ) fails if one or more of the following are true:

EAGAIN	The file exists, mandatory file / record locking is set, and there are outstanding record locks on the file (see **chmod**(2)).
EBADF	*fildes* is not a file descriptor open for writing.
EINTR	A signal was caught during execution of the **ftruncate** routine.
EIO	An I / O error occurred while reading from or writing to the file system.
ENOLINK	*fildes* points to a remote machine and the link to that machine is no longer active.
EINVAL	*fildes* does not correspond to an ordinary file.

SEE ALSO **chmod**(2), **fcntl**(2), **open**(2)

NAME

tsearch, tfind, tdelete, twalk – manage binary search trees

SYNOPSIS

```
#include <search.h>

void *tsearch(const void *key, void **rootp,
        int (*compar)(const void *, const void *));

void *tfind(const void *key, void * const *rootp,
        int (*compar)(const void *, const void *));

void *tdelete(const void *key, void **rootp,
        int (*compar)(const void *, const void *));

void twalk(void *root, void(*action) (void *, VISIT, int));
```

MT-LEVEL

Safe

DESCRIPTION

tsearch( ), **tfind( )**, **tdelete( )**, and **twalk( )** are routines for manipulating binary search trees. They are generalized from *Knuth (6.2.2) Algorithms T and D.* All comparisons are done with a user-supplied routine. This routine is called with two arguments, the pointers to the elements being compared. It returns an integer less than, equal to, or greater than 0, according to whether the first argument is to be considered less than, equal to or greater than the second argument. The comparison function need not compare every byte, so arbitrary data may be contained in the elements in addition to the values being compared.

tsearch( ) is used to build and access the tree. *key* is a pointer to a datum to be accessed or stored. If there is a datum in the tree equal to **key* (the value pointed to by *key*), a pointer to this found datum is returned. Otherwise, **key* is inserted, and a pointer to it returned. Only pointers are copied, so the calling routine must store the data. *rootp* points to a variable that points to the root of the tree. A **NULL** value for the variable pointed to by *rootp* denotes an empty tree; in this case, the variable will be set to point to the datum which will be at the root of the new tree.

Like **tsearch( )**, **tfind( )** will search for a datum in the tree, returning a pointer to it if found. However, if it is not found, **tfind( )** will return a **NULL** pointer. The arguments for **tfind( )** are the same as for **tsearch( )**.

tdelete( ) deletes a node from a binary search tree. The arguments are the same as for **tsearch( )**. The variable pointed to by *rootp* will be changed if the deleted node was the root of the tree. **tdelete( )** returns a pointer to the parent of the deleted node, or a **NULL** pointer if the node is not found.

twalk( ) traverses a binary search tree. *root* is the root of the tree to be traversed. (Any node in a tree may be used as the root for a walk below that node.) *action* is the name of a routine to be invoked at each node. This routine is, in turn, called with three arguments. The first argument is the address of the node being visited. The second argument is a value from an enumeration data type *typedef enum { preorder, postorder, endorder, leaf } VISIT;* (defined in the <**search.h**> header), depending on whether this is the first, second or third time that the node has been visited (during a depth-first, left-to-right traversal of

the tree), or whether the node is a leaf. The third argument is the level of the node in the tree, with the root being level zero.

The pointers to the key and the root of the tree should be of type pointer-to-element, and cast to type pointer-to-character. Similarly, although declared as type pointer-to-character, the value returned should be cast into type pointer-to-element.

RETURN VALUES

A **NULL** pointer is returned by **tsearch( )** if there is not enough space available to create a new node. A **NULL** pointer is returned by **tfind( )** and **tdelete( )** if *rootp* is **NULL** on entry.

If the datum is found, both **tsearch( )** and **tfind( )** return a pointer to it. If not, **tfind( )** returns **NULL**, and **tsearch( )** returns a pointer to the inserted item.

EXAMPLES

The following code reads in strings and stores structures containing a pointer to each string and a count of its length. It then walks the tree, printing out the stored strings and their lengths in alphabetical order.

```
#include <string.h>
#include <stdio.h>
#include <search.h>

struct node {
        char *string;
        int length;
};
char string_space[10000];
struct node nodes[500];
void *root = NULL;

int node_compare(const void *node1, const void *node2) {
        return strcmp(((const struct node *) node1)->string,
                ((const struct node *) node2)->string);
}

void print_node(void **node, VISIT order, int level) {
        if (order == preorder || order == leaf) {
                printf("length=%d, string=%20s\n",
                (*(struct node **)node)->length,
                (*(struct node **)node)->string);
        }
}

main()
{
        char *strptr = string_space;
        struct node *nodeptr = nodes;
        int i = 0;
```

```
        while (gets(strptr) != NULL && i++ < 500) {
                nodeptr→string = strptr;
                nodeptr→length = strlen(strptr);
                (void) tsearch((void *)nodeptr,
                                &root, node_compare);
                strptr += nodeptr→length + 1;
                nodeptr++;
        }
        twalk(root, print_node);
}
```

SEE ALSO **bsearch**(3C), **hsearch**(3C), **lsearch**(3C)

NOTES The **root** argument to **twalk( )** is one level of indirection less than the *rootp* arguments to **tsearch( )** and **tdelete( )**.

There are two nomenclatures used to refer to the order in which tree nodes are visited. **tsearch** uses preorder, postorder and endorder to refer respectively to visiting a node before any of its children, after its left child and before its right, and after both its children. The alternate nomenclature uses preorder, inorder and postorder to refer to the same visits, which could result in some confusion over the meaning of postorder.

If the calling function alters the pointer to the root, results are unpredictable.

NAME

ttyname, ttyname_r, isatty – find name of a terminal

SYNOPSIS

#include <stdlib.h>

char *__**ttyname(int** *fildes***);**

char *__**ttyname_r(int** *fildes*, **char** *__buf*, **int** *len***);**

int isatty(int *fildes***);**

POSIX

cc [*flag***...]** *file* **... −D_POSIX_PTHREAD_SEMANTICS [** *library***...]**

int ttyname_r(int *fildes*, **char** *__name*, **size_t** *namesize***);**

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

ttyname( ) returns a pointer to a string containing the null-terminated path name of the terminal device associated with file descriptor *fildes*.

ttyname_r( ) has the same functionality as **ttyname( )** except that the caller must supply a buffer *buf* with length *len* to store the result; *buf* must be at least **_POSIX_PATH_MAX** in size (defined in **<limits.h>**). The POSIX version of **ttyname_r( )** takes a *namesize* parameter of type **size_t**.

isatty( ) returns 1 if *fildes* is associated with a terminal device, 0 otherwise.

RETURN VALUES

ttyname( ) and **ttyname_r( )** return a NULL pointer if *fildes* does not describe a terminal device in directory **/dev**. The POSIX **ttyname_r( )** returns zero if successful, or the error number upon failure.

ERRORS

ttyname_r( ) will fail if the following is true:

ERANGE The size of the buffer is smaller than the result to be returned.

FILES

/dev/*

SEE ALSO

gettext(3I), **setlocale**(3C)

NOTES

When compiling multi-thread applications, the **_REENTRANT** flag must be defined on the compile line. This flag should only be used in multi-thread applications.

If the application is linked with **−lintl**, then messages printed from this function are in the native language specified by the **LC_MESSAGES** locale category; see **setlocale**(3C).

The return value points to static data whose content is overwritten by each call.

ttyname( ) is unsafe in multi-thread applications. **ttyname_r( )** is MT-Safe, and should be used instead. **isatty( )** is MT-Safe in multi-thread applications.

The new **ttyname_r( )** interface is as specified in POSIX 1003.1c Draft #10.

NAME

ttyslot – find the slot in the utmp file of the current user

SYNOPSIS

```
#include <stdlib.h>
int ttyslot(void);
```

MT-LEVEL

Safe

DESCRIPTION

ttyslot( ) returns the index of the current user's entry in the **/var/adm/utmp** file. The returned index is accomplished by scanning files in **/dev** for the name of the terminal associated with the standard input, the standard output, or the standard error output (0, 1, or 2).

RETURN VALUES

A value of –1 is returned if an error was encountered while searching for the terminal name or if none of the above file descriptors are associated with a terminal device.

FILES

/var/adm/utmp

SEE ALSO

getutent(3C), **ttyname**(3C)

NAME

ualarm – schedule signal after interval in microseconds

SYNOPSIS

```
#include <unistd.h>
unsigned int ualarm(unsigned int usecs,
        unsigned int interval);
```

DESCRIPTION

ualarm( ) sends signal **SIGALRM** (see **signal**(3C)), to the caller in a number of microseconds given by the *usecs* argument. Unless caught or ignored, the signal terminates the caller.

If the *interval* argument is non-zero, the **SIGALRM** signal will be sent to the caller every *interval* microseconds after the timer expires (for instance, after *usecs* microseconds have passed).

Because of scheduling delays, resumption of execution when the signal is caught may be delayed an arbitrary amount.

Interactions between **ualarm**( ) and either **alarm**(2) or **sleep**(3C) are unspecified.

RETURN VALUES

The return value is the amount of time previously remaining in the alarm clock.

SEE ALSO

alarm(2), **getitimer**(2), **setitimer**(2), **sighold**(3C), **signal**(3C), **sleep**(3C), **usleep**(3C)

NOTES

ualarm( ) is a simplified interface to **setitimer**(2); (see **getitimer**(2)).

NAME

ungetc – push character back onto input stream

SYNOPSIS

```
#include <stdio.h>
int ungetc(int c, FILE *stream);
```

MT-LEVEL

MT-Safe

DESCRIPTION

ungetc( ) inserts the character specified by *c* (converted to an **unsigned char**) into the buffer associated with an input stream (see **intro**(3)). That character, *c*, will be returned by the next **getc**(3S) call on that stream. **ungetc( )** returns *c*, and leaves the file corresponding to *stream* unchanged. A successful call to **ungetc( )** clears the **EOF** indicator for *stream*.

Four bytes of pushback are guaranteed.

The value of the file position indicator for *stream* after reading or discarding all pushed-back characters will be the same as it was before the characters were pushed back.

If *c* equals **EOF**, **ungetc( )** does nothing to the buffer and returns **EOF**.

fseek( ), **rewind( )** (both described on **fseek**(3S)), and **fsetpos**(3C) erase the memory of inserted characters for the stream on which they are applied.

RETURN VALUES

ungetc( ) returns **EOF** if it cannot insert the character.

SEE ALSO

intro(3), **fseek**(3S), **fsetpos**(3C), **getc**(3S), **setbuf**(3S), **stdio**(3S)

NAME

ungetwc – push a Process Code character back into input stream

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <stdio.h>
#include <widec.h>

int ungetwc(int *c***, FILE** ∗*stream***);**

MT-LEVEL

MT-Safe

DESCRIPTION

ungetwc( ) pushes back the Process Code character *c* onto an input stream. That character will be returned by the next **getwc**(3I) call on that stream. **ungetwc( )** returns *c*, and leaves the file *stream* unchanged.

One character of pushback is guaranteed provided something has been read from the stream and the stream is actually buffered. In the case that *stream* is **stdin**, one character may be pushed back onto the buffer without a previous read statement.

If *c* equals **EOF**, **ungetwc( )** does nothing to the buffer and returns **EOF**.

An **fseek**(3S) erases all memory of pushed back characters.

RETURN VALUES

ungetwc( ) returns **EOF** if it can't push a character back.

SEE ALSO

fseek(3S), **getwc**(3I), **setbuf**(3S), **ungetc**(3S)

NAME

unlockpt – unlock a pseudo-terminal master/slave pair

SYNOPSIS

int unlockpt(int *fildes***);**

MT-LEVEL

Safe

DESCRIPTION

The function **unlockpt()** clears a lock flag associated with the slave pseudo-terminal device associated with its master pseudo-terminal counterpart so that the slave pseudo-terminal device can be opened. *fildes* is a file descriptor returned from a successful open of a master pseudo-terminal device.

RETURN VALUES

Upon successful completion, the function **unlockpt()** returns **0**; otherwise it returns **-1**. A failure may occur if *fildes* is not an open file descriptor or is not associated with a master pseudo-terminal device.

SEE ALSO

open(2), **grantpt**(3C), **ptsname**(3C)

STREAMS Programming Guide

NAME

usleep – suspend execution for interval in microseconds

SYNOPSIS

```
#include <unistd.h>
int usleep(unsigned int useconds);
```

DESCRIPTION

Suspend the caller for the number of microseconds specified by the argument. The actual suspension time may be an arbitrary amount longer because of other activity in the system, or because of the time spent in processing the call.

The routine is implemented by setting an interval timer and pausing until it occurs. The previous state of this timer is saved and restored. If the sleep time exceeds the time to the expiration of the previous timer, the process sleeps only until the signal would have occurred, and the signal is sent a short time later.

Interactions between **usleep**( ) and either **alarm**(2) or **sleep**(3C) are unspecified.

SEE ALSO

alarm(2), **getitimer**(2), **poll**(2), **sigprocmask**(2), **select**(3C), **sleep**(3C), **ualarm**(3C)

NOTES

A microsecond is .000001 seconds.

NAME

vlfmt – display error message in standard format and pass to logging and monitoring services

MT-LEVEL

MT-safe

SYNOPSIS

```
#include <stdarg.h>
#include <pfmt.h>

int vlfmt(FILE *stream, long flags, char *format, va_list ap);
```

DESCRIPTION

vlfmt() is the same as **lfmt()** except that instead of being called with a variable number of arguments, it is called with an argument list as defined by the **<stdarg.h>** header file.

The **<stdarg.h>** header file defines the type **va_list** and a set of macros for advancing through a list of arguments whose number and types may vary. The argument *ap* to **vlfmt()** is of type **va_list**. This argument is used with the **<stdarg.h>** header file macros **va_start()**, **va_arg()** and **va_end()**.
[see **va_start()**, **va_arg()**, and **va_end()** in **stdarg**(5)]. The EXAMPLE section below shows their use with **vlfmt()**.

The macro **va_alist** is used as the parameter list in a function definition as in the function called **errlog()** in the example below. The macro **va_start(ap,)**, where *ap* is of type **va_list**, must be called before any attempt to traverse and access unnamed arguments. Calls to **va_arg(***ap***,** *atype***)** traverse the argument list. Each execution of **va_arg()** expands to an expression with the value and type of the next argument in the list *ap*, which is the same object initialized by **va_start**. The argument *atype* is the type that the returned argument is expected to be. The **va_end(***ap***)** macro must be invoked when all desired arguments have been accessed. (The argument list in *ap* can be traversed again if **va_start()** is called again after **va_end()**.) In the example below, **va_arg()** is executed first to retrieve the format string passed to **errlog()**. The remainting **errlog()** arguments, *arg1*, *arg2*, ..., are given to *vlfmt()* in the argument *ap*.

RETURN VALUE

Upon success, **vlfmt()** returns the number of bytes transmitted. Upon failure, it returns a negative value:

–1 write error to *stream*.

–2 cannot log and / or display at console.

EXAMPLE

The following demonstrates how **vlfmt()** could be used to write an **errlog()** routine:

```
#include <pfmt.h>
#include <stdarg.h>
/*
 *  errlog should be called like
 *        errlog(log_info, format, arg1, ...);
 */
void errlog(long log_info, ...)
```

```
{
   va_list ap;
   char *format;

   va_start(ap, );
   format = va_arg(ap, char *);
   (void) vlfmt(stderr, log_info | MM_ERROR, format, ap);
   va_end(ap);
   (void) abort();
}
```

NOTES

Since **vlfmt()** uses **gettxt**(3C), it is recommended that **vlfmt()** not be used.

SEE ALSO

gettxt(3C), **lfmt**(3C), **stdarg**(5)

NAME

volmgt_check – have Volume Management check for media

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lvolmgt** [*library*. . . .]

#include <volmgt.h>

int volmgt_check(char **pathname***);**

MT-LEVEL

MT-Safe

DESCRIPTION

This routine asks Volume Management to check the specified *pathname* and determine if new media has been inserted in that drive.

If a null pointer is passed in, then Volume Management will check each device it is managing that can be checked.

If new media is found, **volmgt_check( )** tells Volume Management to initiate any "actions" specified in **/etc/vold.conf** (see **vold.conf**(4)).

RETURN VALUES

This routine returns **0** if no media was found, and a non-zero value if any media was found.

ERRORS

This routine can fail, returning **0**, if a **stat**(2) or **open**(2) of the supplied *pathname* fails, or if any of the following is true:

ENXIO Volume Management is not running.

EINTR An interrupt signal was detected while checking for media.

EXAMPLES

To check if any drive managed by Volume Management has any new media inserted in it:

```
if (volmgt_check(NULL)) {
    (void) printf("Volume Management found media\n");
}
```

This would also request Volume Management to take whatever action was specified in **/etc/vold.conf** for any media found.

SEE ALSO

cc(1B), **volcheck**(1), **vold**(1M), **open**(2), **stat**(2), **volmgt_inuse**(3X), **volmgt_running**(3X), **vold.conf**(4), **volfs**(7FS)

NOTES

Volume Management must be running for this routine to work.

Since **volmgt_check( )** returns **0** for two different cases (both when no media is found, and when an error occurs), it is up to the user to to check *errno* to differentiate the two, and to ensure that Volume Management is running.

NAME

volmgt_inuse – check whether or not Volume Management is managing a pathname

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lvolmgt** [*library*. . . .]

#include <volmgt.h>

int volmgt_inuse(char ∗*pathname***);**

MT-LEVEL

MT-Safe

DESCRIPTION

volmgt_inuse( ) checks whether Volume Management is managing the specified *pathname*.

RETURN VALUES

A non-zero value is returned if Volume Management is managing the specified *pathname*, otherwise **0** is returned.

ERRORS

This routine can fail, returning **0**, if a **stat**(2) of the supplied *pathname* or an **open**(2) of **/dev/volctl** fails, or if any of the following is true:

ENXIO Volume Management is not running.

EINTR An interrupt signal was detected while checking for the supplied *pathname* for use.

EXAMPLES

To see if Volume Management is managing the first floppy disk:

```
if (volmgt_inuse("/dev/rdiskette0") != 0) {
    (void) printf("volmgt is managing diskette 0\n");
} else {
    (void) printf("volmgt is NOT managing diskette 0\n");
}
```

SEE ALSO

cc(1B), **vold**(1M), **open**(2), **stat**(2), **errno**(3C), **volmgt_check**(3X), **volmgt_running**(3X), **volfs**(7FS)

NOTES

This routine requires Volume Management to be running.

Since **volmgt_inuse( )** returns **0** for two different cases (both when a volume is not in use, and when an error occurs), it is up to the user to to check **errno** to differentiate the two, and to ensure that Volume Management is running.

NAME

volmgt_root – return the Volume Management root directory

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lvolmgt** [*library*. . . .]

#include <volmgt.h>

char ∗**volmgt_root(void);**

MT-LEVEL

MT-Safe

DESCRIPTION

volmgt_root() returns the current Volume Management root directory, which by default is **/vol** but can be configured to be in a different location.

RETURN VALUES

A pointer to a static string containing the root directory for Volume Management is returned.

ERRORS

This routine may fail if an **open()** of **/dev/volctl** fails. If this occurs a pointer to the default Volume Management root directory is returned.

EXAMPLES

To find out where the Volume Management root directory is:

```
if ((path = volmgt_root()) != NULL) {
    (void) printf("Volume Management root dir=%s\n", path);
} else {
    (void) printf("can't find Volume Management root dir\n");
}
```

FILES

/vol Default location for the Volume Management root directory

SEE ALSO

cc(1B), **vold**(1M), **open**(2), **volmgt_check**(3X), **volmgt_inuse**(3X), **volmgt_running**(3X), **volfs**(7FS)

NOTES

This routine will return the default root directory location even when Volume Management is not running.

NAME

volmgt_running – return whether or not Volume Management is running

SYNOPSIS

cc [*flag* . . .] *file* . . . –**lvolmgt** [*library*. . . .]

```
#include <volmgt.h>
int volmgt_running(void);
```

MT-LEVEL

MT-Safe

DESCRIPTION

volmgt_running( ) tells whether or not Volume Management is running.

RETURN VALUES

A non-zero value is returned if Volume Management is running, else **0** is returned.

ERRORS

volmgt_running( ) will fail, returning **0**, if a **stat**(2) or **open**(2) of **/dev/volctl** fails, or if any of the following is true:

ENXIO Volume Management is not running.

EINTR An interrupt signal was detected while checking to see if Volume Management was running.

EXAMPLES

To see if Volume Management is running:

```
if (volmgt_running() != 0) {
    (void) printf("Volume Management is running\n");
} else {
    (void) printf("Volume Management is NOT running\n");
}
```

SEE ALSO

cc(1B), **vold**(1M), **open**(2), **stat**(2), **volmgt_check**(3X), **volmgt_inuse**(3X), **volfs**(7FS)

NOTES

Volume Management must be running for many of the Volume Management library routines to work.

NAME

volmgt_symname, volmgt_symdev – convert between Volume Management symbolic names, and the devices that correspond to them

SYNOPSIS

cc [*flag* . . .] *file* . . . **–lvolmgt** [*library*. . . .]

#include <volmgt.h>

char ∗**volmgt_symname(char** ∗*pathname***);**

char ∗**volmgt_symdev(char** ∗*symname***);**

MT-LEVEL

MT-Safe

DESCRIPTION

These two routines compliment each other, translating between Volume Management's symbolic name for a device, called a *symname*, and the **/dev** *pathname* for that same device.

volmgt_symname( ) converts a supplied **/dev** *pathname* to a **symname**, Volume Management's idea of that device's symbolic name (see **volfs**(7FS) for a description of Volume Management symbolic names).

volmgt_symdev( ) does the opposite conversion, converting between a *symname*, Volume Management's idea of a device's symbolic name for a volume, to the **/dev** *pathname* for that device.

RETURN VALUES

volmgt_symname( ) returns the symbolic name for the device pathname supplied, and **volmgt_symdev**( ) returns the device pathname for the supplied symbolic name.

These strings are allocated upon success, and therefore must be freed by the caller when they are no longer needed (see **free**(3C)).

ERRORS

volmgt_symname( ) can fail, returning a null string pointer, if a **stat**(2) of the supplied **pathname** fails, or if an **open**(2) of **/dev/volctl** fails, or if any of the following is true:

ENXIO Volume Management is not running.

EINTR An interrupt signal was detected while trying to convert the supplied *pathname* to a *symname*.

volmgt_symdev( ) can fail if an **open**(2) of **/dev/volctl** fails, or if any of the following is true:

ENXIO Volume Management is not running.

EINTR An interrupt signal was detected while trying to convert the supplied *symname* to a **/dev** *pathname*.

EXAMPLES

The following tests how many floppies Volume Management currently sees in floppy drives (up to 10):

```
for (i=0; i < 10; i++) {
    (void) sprintf(path, "floppy%d", i);
    if (volmgt_symdev(path) != NULL) {
        (void) printf("volume %s is in drive %d\n",
```

```
                path, i);
        }
    }
```

This code finds out what symbolic name (if any) Volume Management has for **/dev/rdsk/c0t6d0s2**:

```
    if ((nm = volmgt_symname("/dev/rdsk/c0t6d0s2")) == NULL) {
        (void) printf("path not managed\n");
    } else {
        (void) printf("path managed as %s\n", nm);
    }
```

SEE ALSO

cc(1B), **vold**(1M), **open**(2), **stat**(2), **free**(3C), **malloc**(3C), **volmgt_check**(3X), **volmgt_inuse**(3X), **volmgt_running**(3X), **volfs**(7FS)

NOTES

These routines only work when Volume Management is running.

BUGS

There should be a straightforward way to query Volume Management for a list of all media types it's managing, and how many of each type are being managed.

NAME

vpfmt – display error message in standard format and pass to logging and monitoring services

MT-LEVEL

MT-safe

SYNOPSIS

```
#include <stdarg.h>
#include <pfmt.h>

int vpfmt(FILE *stream, long flags, char *format, va_list ap);
```

DESCRIPTION

vpfmt() is the same as **lfmt()** except that instead of being called with a variable number of arguments, it is called with an argument list as defined by the **<stdarg.h>** header file.

The **<stdarg.h>** header file defines the type **va_list** and a set of macros for advancing through a list of arguments whose number and types may vary. The argument *ap* to **vpfmt()** is of type **va_list**. This argument is used with the **<stdarg.h>** header file macros **va_start()**, **va_arg()** and **va_end()**.
[see **va_start()**, **va_arg()**, and **va_end()** in **stdarg**(5)]. The EXAMPLE section below shows their use with **vpfmt()**.

The macro **va_alist** is used as the parameter list in a function definition as in the function called **error()** in the example below. The macro **va_start(ap,)**, where *ap* is of type **va_list**, must be called before any attempt to traverse and access unnamed arguments. Calls to **va_arg(***ap***,** *atype***)** traverse the argument list. Each execution of **va_arg()** expands to an expression with the value and type of the next argument in the list *ap*, which is the same object initialized by **va_start**. The argument *atype* is the type that the returned argument is expected to be. The **va_end(***ap***)** macro must be invoked when all desired arguments have been accessed. (The argument list in *ap* can be traversed again if **va_start()** is called again after **va_end()**.) In the example below, **va_arg()** is executed first to retrieve the format string passed to **error()**. The remaining **error()** arguments, *arg1*, *arg2*, ..., are given to *vpfmt()* in the argument *ap*.

RETURN VALUE

Upon success, **lfmt()** returns the number of bytes transmitted. Upon failure, it returns a negative value:

–1 write error to *stream*.

EXAMPLE

The following demonstrates how **vpfmt()** could be used to write an **error()** routine:

```
#include <pfmt.h>
#include <stdarg.h>
/*
 *   error should be called like
 *        error(format, arg1, ...);
 */
void error(...)

{
```

```
        va_list ap;
        char *format;

        va_start(ap, );
        format = va_arg(ap, char *);
        (void) vpfmt(stderr, MM_ERROR, format, ap);
        va_end(ap);
        (void) abort();
    }
```

NOTES

Since **vpfmt()** uses **gettxt**(3C), it is recommended that **vpfmt()** not be used.

SEE ALSO

pfmt(3C), **stdarg**(5)

NAME

vprintf, vfprintf, vsprintf – print formatted output of a variable argument list

SYNOPSIS

```
#include <stdio.h>
#include <stdarg.h>

int vprintf(const char *format, va_list ap);

int vfprintf(FILE *stream, const char *format, va_list ap);

int vsprintf(char *s, const char *format, va_list ap);
```

MT-LEVEL

See the **NOTES** section of this page.

DESCRIPTION

vprintf(), **vfprintf()**, and **vsprintf()** are the same as **printf()**, **fprintf()**, and **sprintf()** respectively, except that instead of being called with a variable number of arguments, they are called with an argument list as defined by the **<stdarg.h>** header.

The **<stdarg.h>** header defines the type **va_list** and a set of macros for advancing through a list of arguments whose number and types may vary. The argument *ap* to the vprint family of routines is of type **va_list**. This argument is used with the **<stdarg.h>** header file macros **va_start()**, **va_arg()**, and **va_end()** (see **stdarg**(5)). The **EXAMPLES** section below shows the use of **va_start()** and **va_end()** with **vprintf()**.

The macro **va_alist** is used as the parameter list in a function definition, as in the function called **error()** in the example below. The macro **va_start(***ap, parmN***)**, where *ap* is of type **va_list**, and *parmN* is the rightmost parameter (just before ...), must be called before any attempt to traverse and access unnamed arguments is made. The **va_end(***ap***)** macro must be invoked when all desired arguments have been accessed. (The argument list in *ap* can be traversed again if **va_start()** is called again after **va_end()**.) In the example below, the **error()** arguments, *arg1, arg2,* ..., are given to **vfprintf()** in the argument *ap*.

RETURN VALUES

vprintf() and **vfprintf()** return the number of characters transmitted, or return **–1** if an error was encountered.

EXAMPLES

The following demonstrates how **vfprintf()** could be used to write an **error** routine:

```
#include <stdio.h>
#include <stdarg.h>
...
/*
 *   error should be called like
 *        error(function_name, format, arg1, ...);
 */
void error(char *function_name, char *format, ...)
{
   va_list ap;
   va_start(ap, );
   /* print out name of function causing error */
   (void) fprintf(stderr, "ERR in %s: ", function_name);
```

```
        /* print out remainder of message */
        (void) vfprintf(stderr, format, ap);
        va_end(ap);
        (void) abort;
    }
```

SEE ALSO

printf(3S), **stdarg**(5)

NOTES

vprintf( ), **vfprintf( )**, and **vsprintf( )** are MT-Safe in multi-thread applications.

NAME

vsyslog – log message with a varargs argument list

SYNOPSIS

```
#include <syslog.h>
#include <varargs.h>

int vsyslog(int priority, const char *message, va_list ap);
```

MT-LEVEL

Safe

DESCRIPTION

vsyslog() is the same as **syslog**(3) except that instead of being called with a variable number of arguments, it is called with an argument list as defined by **varargs**(5).

EXAMPLES

The following demonstrates how **vsyslog()** could be used to write an error routine.

```
#include <syslog.h>
#include <varargs.h>
...
        /*
         * error should be called like:
         *      error(pri, function_name, format, arg1, arg2...);
         * Note that pri, function_name, and format cannot be declared
         * separately because of the definition of varargs.
         */

/*VARARGS0*/
void
error(va_alist)
        va_dcl; {
        va_list args;
        int pri;
        char *message;

        va_start(args);
        pri = va_arg(args, int);
        /* log name of function causing error */
        (void) syslog(pri, "ERROR in %s", va_arg(args, char *));
        message = va_arg(args, char *);
        /* log remainder of message */
        (void) vsyslog(pri, msg, args);
        va_end(args);
        (void) abort();
}
```

SEE ALSO

syslog(3), **varargs**(5)

NAME

wait, wait3, wait4, waitpid, WIFSTOPPED, WIFSIGNALED, WIFEXITED – wait for process to terminate or stop

SYNOPSIS

/usr/ucb/cc [*flag* ...] *file* ...

#include <sys/wait.h>

int wait(*statusp***)**
int **statusp*;

int waitpid(*pid, statusp, options***)**
int *pid*;
int **statusp*;
int *options*;

#include <sys/time.h>
#include <sys/resource.h>

int wait3(*statusp, options, rusage***)**
int **statusp*;
int *options*;
struct rusage **rusage*;

int wait4(*pid, statusp, options, rusage***)**
int *pid*;
int **statusp*;
int *options*;
struct rusage **rusage*;

WIFSTOPPED(*status***)**
int *status*;

WIFSIGNALED(*status***)**
int *status*;

WIFEXITED(*status***)**
int *status*;

DESCRIPTION

wait( ) delays its caller until a signal is received or one of its child processes terminates or stops due to tracing. If any child process has died or stopped due to tracing and this has not been reported using **wait( )**, return is immediate, returning the process ID and exit status of one of those children. If that child process has died, it is discarded. If there are no children, return is immediate with the value –1 returned. If there are only running or stopped but reported children, the calling process is blocked.

If *status* is not a **NULL** pointer, then on return from a successful **wait( )** call the status of the child process whose process ID is the return value of **wait( )** is stored in the **wait( )** union pointed to by *status*. The **w_status** member of that union is an **int**; it indicates the cause of termination and other information about the terminated process in the following manner:

- If the low-order 8 bits of **w_status** are equal to 0177, the child process has

stopped; the 8 bits higher up from the low-order 8 bits of **w_status** contain the number of the signal that caused the process to stop. See **ptrace**(2) and **sigvec**(3B).

- If the low-order 8 bits of **w_status** are non-zero and are not equal to 0177, the child process terminated due to a signal; the low-order 7 bits of **w_status** contain the number of the signal that terminated the process. In addition, if the low-order seventh bit of **w_status** (that is, bit 0200) is set, a ''core image'' of the process was produced; see **sigvec**(3B).
- Otherwise, the child process terminated due to an **exit()** call; the 8 bits higher up from the low-order 8 bits of **w_status** contain the low-order 8 bits of the argument that the child process passed to **exit()**; see **exit**(2).

waitpid() behaves identically to **wait()** if *pid* has a value of –1 and *options* has a value of zero. Otherwise, the behavior of **waitpid()** is modified by the values of *pid* and *options* as follows:

pid specifies a set of child processes for which status is requested. **waitpid()** only returns the status of a child process from this set.

- If *pid* is equal to –1, status is requested for any child process. In this respect, **waitpid()** is then equivalent to **wait()**.
- If *pid* is greater than zero, it specifies the process ID of a single child process for which status is requested.
- If *pid* is equal to zero, status is requested for any child process whose process group ID is equal to that of the calling process.
- If *pid* is less than –1, status is requested for any child process whose process group ID is equal to the absolute value of *pid*.

options is constructed from the bitwise inclusive OR of zero or more of the following flags, defined in the header **<sys/wait.h>**:

WNOHANG
: **waitpid()** does not suspend execution of the calling process if status is not immediately available for one of the child processes specified by *pid*.

WUNTRACED
: The status of any child processes specified by *pid* that are stopped, and whose status has not yet been reported since they stopped, are also reported to the requesting process.

wait3() is an alternate interface that allows both non-blocking status collection and the collection of the status of children stopped by any means. The *status* parameter is defined as above. The *options* parameter is used to indicate the call should not block if there are no processes that have status to report (**WNOHANG**), and/or that children of the current process that are stopped due to a **SIGTTIN**, **SIGTTOU**, **SIGTSTP**, or **SIGSTOP** signal are eligible to have their status reported as well (**WUNTRACED**). A terminated child is discarded after it reports status, and a stopped process will not report its status more than once. If *rusage* is not a **NULL** pointer, a summary of the resources used by the terminated process and all its children is returned. Only the user time used and the system time

used are currently available. They are returned in **rusage.ru_utime** and **rusage.ru_stime**, respectively.

When the **WNOHANG** option is specified and no processes have status to report, **wait3( )** returns 0. The **WNOHANG** and **WUNTRACED** options may be combined by ORing the two values.

wait4( ) is another alternate interface. With a *pid* argument of 0, it is equivalent to **wait3( )**. If *pid* has a nonzero value, then **wait4( )** returns status only for the indicated process ID, but not for any other child processes.

WIFSTOPPED, **WIFSIGNALED**, **WIFEXITED**, are macros that take an argument *status*, of type **int**, as returned by **wait( )**, or **wait3( )**, or **wait4( )**. **WIFSTOPPED** evaluates to true (1) when the process for which the **wait( )** call was made is stopped, or to false (0) otherwise. **WIFSIGNALED** evaluates to true when the process was terminated with a signal. **WIFEXITED** evaluates to true when the process exited by using an **exit**(2) call.

RETURN VALUES

If **wait( )** or **waitpid( )** returns due to a stopped or terminated child process, the process ID of the child is returned to the calling process. Otherwise, a value of –1 is returned and **errno** is set to indicate the error.

If **wait( )** or **waitpid( )** return due to the delivery of a signal to the calling process, a value of –1 is returned and **errno** is set to EINTR. If **waitpid( )** function was invoked with **WNOHANG** set in *options*, it has at least one child process specified by *pid* for which status is not available, and status is not available for any process specified by *pid*, a value of zero is returned. Otherwise, a value of –1 is returned, and **errno** is set to indicate the error.

wait3( ) and **wait4( )** returns 0 if **WNOHANG** is specified and there are no stopped or exited children, and returns the process ID of the child process if it returns due to a stopped or terminated child process. Otherwise, they returns a value of –1 and sets **errno** to indicate the error.

ERRORS

wait( ), **wait3( )** or **wait4( )** will fail and return immediately if one or more of the following are true:

ECHILD	The calling process has no existing unwaited-for child processes.
EFAULT	The *status* or *rusage* arguments point to an illegal address.

waitpid( ) may set **errno** to:

ECHILD	The process or process group specified by *pid* does not exist or is not a child of the calling process.
EINTR	The function was interrupted by a signal. The value of the location pointed to by *statusp* is undefined.
EINVAL	The value of *options* is not valid.

wait( ), and **wait3( )**, and **wait4( )** will terminate prematurely, return –1, and set **errno** to EINTR upon the arrival of a signal whose **SV_INTERRUPT** bit in its flags field is set (see **sigvec**(3B) and **siginterrupt**(3B)). **signal**(3B), sets this bit for any signal it catches.

SEE ALSO

exit(2), **ptrace**(2), **wait**(2), **waitpid**(2), **getrusage**(3C), **siginterrupt**(3B), **signal**(3B), **sigvec**(3B), **signal**(3C)

NOTES

Use of these interfaces should be restricted to only applications written on BSD platforms. Use of these interfaces with any of the system libraries or in multi-thread applications is unsupported.

If a parent process terminates without waiting on its children, the initialization process (process ID = 1) inherits the children.

wait( ), and **wait3( )**, and **wait4( )** are automatically restarted when a process receives a signal while awaiting termination of a child process, unless the **SV_INTERRUPT** bit is set in the flags for that signal.

Calls to **wait( )** with an argument of **0** should be cast to type '**int** ∗', as in:

wait((int ∗**)0)**

Previous SunOS releases used **union wait** ∗**statusp** and **union wait status** in place of **int** ∗**statusp** and **int** status. The union contained a member **w_status** that could be treated in the same way as *status*.

Other members of the **wait** union could be used to extract this information more conveniently:

- If the **w_stopval** member had the value **WSTOPPED**, the child process had stopped; the value of the **w_stopsig** member was the signal that stopped the process.
- If the **w_termsig** member was non-zero, the child process terminated due to a signal; the value of the **w_termsig** member was the number of the signal that terminated the process. If the **w_coredump** member was non-zero, a core dump was produced.
- Otherwise, the child process terminated due to a call to **exit( )**. The value of the **w_retcode** member was the low-order 8 bits of the argument that the child process passed to **exit( )**.

union wait is obsolete in light of the new specifications provided by *IEEE Std 1003.1-1988* and endorsed by *SVID89* and *XPG3*. SunOS Release 4.1 supports **union wait** for backward compatibility, but it will disappear in a future release.

NAME

wait3, wait4 – wait for process to terminate or stop

SYNOPSIS

```
#include <sys/wait.h>
#include <sys/time.h>
#include <sys/resource.h>

pid_t wait3(int *statusp, int options,
        struct rusage *rusage);

pid_t wait4(pid_t pid, int *statusp, int options,
        struct rusage *rusage);
```

DESCRIPTION

wait3( ) delays its caller until a signal is received or one of its child processes terminates or stops due to tracing. If any child process has died or stopped due to tracing and this has not already been reported, return is immediate, returning the process ID and status of one of those children. If that child process has died, it is discarded. If there are no children, –1 is returned immediately. If there are only running or stopped but reported children, the calling process is blocked.

If *statusp* is not a **NULL** pointer, then on return from a successful **wait3( )** call, the status of the child process is stored in the integer pointed to by *statusp*. **statusp* indicates the cause of termination and other information about the terminated process in the following manner:

- If the low-order 8 bits of ***statusp** are equal to 0177, the child process has stopped; the 8 bits higher up from the low-order 8 bits of ***statusp** contain the number of the signal that caused the process to stop. See **signal**(5).
- If the low-order 8 bits of ***statusp** are non-zero and are not equal to 0177, the child process terminated due to a signal; the low-order 7 bits of ***statusp** contain the number of the signal that terminated the process. In addition, if the low-order seventh bit of ***statusp** (that is, bit 0200) is set, a ''core image'' of the process was produced; see **signal**(5).
- Otherwise, the child process terminated due to an **exit( )** call; the 8 bits higher up from the low-order 8 bits of ***statusp** contain the low-order 8 bits of the argument that the child process passed to **exit( )**; see **exit**(2).

options is constructed from the bitwise inclusive OR of zero or more of the following flags, defined in the header **<sys/wait.h>**:

WNOHANG
Execution of the calling process is not suspended if status is not immediately available for any child process.

WUNTRACED
The status of any child processes that are stopped, and whose status has not yet been reported since they stopped, are also reported to the requesting process.

If *rusage* is not a **NULL** pointer, a summary of the resources used by the terminated process and all its children is returned. Only the user time used and the system time used are currently available. They are returned in the **ru_utime** and **ru_stime**, members of the rusage structure respectively.

When the **WNOHANG** option is specified and no processes have status to report, **wait3()** returns 0. The **WNOHANG** and **WUNTRACED** options may be combined by ORing the two values.

wait4() is an extended interface. With a *pid* argument of 0, it is equivalent to **wait3()**. If *pid* has a nonzero value, then **wait4()** returns status only for the indicated process ID, but not for any other child processes. The status can be evaluated using the macros defined by **wstat**(5).

RETURN VALUES

If **wait3()** or **wait4()** returns due to a stopped or terminated child process, the process ID of the child is returned to the calling process. Otherwise, a value of –1 is returned and **errno** is set to indicate the error.

If **wait3()** or **wait4()** return due to the delivery of a signal to the calling process, a value of –1 is returned and **errno** is set to EINTR. If **WNOHANG** was set in *options*, it has at least one child process specified by *pid* for which status is not available, and status is not available for any process specified by *pid*, a value of zero is returned. Otherwise, a value of –1 is returned, and **errno** is set to indicate the error.

wait3() and **wait4()** return 0 if **WNOHANG** is specified and there are no stopped or exited children, and return the process ID of the child process if they return due to a stopped or terminated child process. Otherwise, they return a value of –1 and sets **errno** to indicate the error.

ERRORS

wait3() or **wait4()** will fail and return immediately if one or more of the following are true:

ECHILD	The calling process has no existing unwaited-for child processes.
EFAULT	The *statusp* or *rusage* arguments point to an illegal address.
EINTR	The function was interrupted by a signal. The value of the location pointed to by *statusp* is undefined.
EINVAL	The value of *options* is not valid.

wait4() may set **errno** to:

ECHILD	The process specified by *pid* does not exist or is not a child of the calling process.

wait3(), and **wait4()** will terminate prematurely, return –1, and set **errno** to EINTR upon the arrival of a signal whose **SA_RESTART** bit in its flags field is not set (see **sigaction**(2)).

SEE ALSO

kill(1), **exit**(2), **wait**(2), **waitid**(2), **waitpid**(2), **getrusage**(3C), **signal**(3C), **proc**(4), **signal**(5), **wstat**(5)

NOTES

If a parent process terminates without waiting on its children, the initialization process (process ID = 1) inherits the children.

wait3( ), and **wait4( )** are automatically restarted when a process receives a signal while awaiting termination of a child process, unless the **SA_RESTART** bit is not set in the flags for that signal.

NAME

wconv, towupper, towlower – Process Code character conversion macros

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <widec.h>
#include <wctype.h>

int towlower(int *c***);**

int towupper(int *c***);**

MT-LEVEL

MT-Safe with exceptions

CHARACTER CONVERSION MACROS

These macros perform simple case conversions on Latin characters in Process Code, *wchar_t*, from the primary and supplementary codesets, by table lookup.

towupper(*c***)** converts the lower-case Latin character *c* to its upper-case equivalent. If *c* is not a lower-case Latin character *c* is returned.

towlower(*c***)** converts the upper-case character *c* to its lower-case equivalent. If *c* is not an upper-case Latin character *c* is returned.

SEE ALSO

ctype(3C), **setlocale**(3C), **iswalpha**(3I), **stdio**(3S)

NOTES

towupper and **towlower** can be used safely in a multi-thread application, as long as **setlocale**(3C) is not being called to change the locale.

NAME

wcscoll, wscoll – wide character string comparison using collating information

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <wchar.h>

int wcscoll(const wchar_t **ws1,* **const wchar_t** **ws2***);**
int wscoll(const wchar_t **ws1,* **const wchar_t** **ws2***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **wcscoll( )** and **wscoll( )** functions compare the wide character string pointed to by *ws1* to the wide character string pointed to by *ws2*, both interpreted as appropriate to the **LC_COLLATE** category of the current locale.

RETURN VALUES

Upon successful completion, **wcscoll( )** and **wscoll( )** return an integer greater than, equal to, or less than 0, depending upon whether the wide character string pointed to by *ws1* is greater than, equal to, or less than the wide character string pointed to by *ws2*, when both are interpreted as appropriate to the current locale. On error, **wcscoll( )** and **wscoll( )** may set **errno**, but no return value is reserved to indicate an error.

ERRORS

wcscoll( ) and **wscoll( )** may fail if:

EINVAL The *ws1* or *ws2* arguments contain wide character codes outside the domain of the collating sequence.

ENOSYS The function is not supported.

SEE ALSO

setlocale(3C), **wcscmp**(3I), **wcsxfrm**(3I)

NOTES

Because no return value is reserved to indicate an error, an application wishing to check for error situations should set **errno** to **0**, call either **wcscoll( )** or **wscoll( )**, then check **errno** and if it is non-zero, assume an error has occurred.

wcsxfrm(3I) and **wcscmp**(3I) should be used for sorting large lists.

NAME

wcsftime – convert date and time to wide character string

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <wchar.h>

size_t wcsftime(wchar_t ∗*wcs*, **size_t** *maxsize*, **const char** ∗*format*,
const struct tm ∗*timptr***);**

DESCRIPTION

The **wcsftime( )** function places wide-character codes into the array pointed to by *wcs* as controlled by the string pointed to by *format*.

This function behaves as if the character string generated by the **strftime**(3C) function is passed to the **mbstowcs**(3C) function as the character string argument, and **mbstowcs( )** places the result in the wide character string argument of the **wcsftime( )** function, up to a limit of *maxsize* wide-character codes.

If copying takes place between objects that overlap, the behaviour is undefined.

RETURN VALUES

If the total number of resulting wide character codes (including the terminating null wide-character code) is no more than *maxsize*, **wcsftime( )** returns the number of wide-character codes placed into the array pointed to by *wcs*, not including the terminating null wide-character code. Otherwise, **0** is returned and the contents of the array are indeterminate.

wcfstime( ) uses **malloc**(3C) and should **malloc( )** fail, **errno** will be set by **malloc( )**.

SEE ALSO

malloc(3C), **mbstowcs**(3C), **strftime**(3C)

NAME

wcstod, wstod, watof – convert wide character string to double-precision number

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <wchar.h>

double wcstod(const wchar_t ∗*nptr,* **wchar_t** ∗∗*endptr***);**
double wstod(const wchar_t ∗*nptr,* **wchar_t** ∗∗*endptr***);**
double watof(wchar_t ∗*nptr***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **wcstod( )** and **wstod( )** functions convert the initial portion of the wide character string pointed to by *nptr* to **double** representation. They first decompose the input wide character string into three parts: an initial, possibly empty, sequence of white-space wide character codes (as specified by **iswspace**(3I)); a subject sequence interpreted as a floating-point constant; and a final wide-character string of one or more unrecognised wide-character codes, including the terminating null wide character code of the input wide character string. They then attempt to convert the subject sequence to a floating-point number, and return the result.

The expected form of the subject sequence is an optional '+' or '–' sign, then a non-empty sequence of digits optionally containing a radix, then an optional exponent part. An exponent part consists of 'e' or 'E', followed by an optional sign, followed by one or more decimal digits. The subject sequence is defined as the longest initial subsequence of the input wide character string, starting with the first non-white-space wide-character code, that is of the expected form. The subject sequence contains no wide-character codes if the input wide character string is empty or consists entirely of white-space wide-character codes, or if the first wide-character code that is not white space other than a sign, a digit or a radix.

If the subject sequence has the expected form, the sequence of wide-character codes starting with the first digit or the radix (whichever occurs first) is interpreted as a floating constant as defined in the C language, except that the radix is used in place of a period, and that if neither an exponent part nor a radix appears, a radix is assumed to follow the last digit in the wide character string. If the subject sequence begins with a minus sign (-), the value resulting from the conversion is negated. A pointer to the final wide character string is stored in the object pointed to by *endptr*, provided that *endptr* is not a null pointer.

The radix is defined in the program's locale (category **LC_NUMERIC**). In the POSIX locale, or in a locale where the radix is not defined, the radix defaults to a period (.).

In other than the POSIX locale, other implementation-dependent subject sequence forms may be accepted.

If the subject sequence is empty or does not have the expected form, no conversion is performed; the value of *nptr* is stored in the object pointed to by *endptr*, provided that *endptr* is not a null pointer.

watof(str) is equivalent to **wstod(str, (wchar_t **)NULL)**.

RETURN VALUES

wcstod( ) and **wstod( )** return the converted value, if any. If no conversion could be performed, **0** is returned, and **errno** may be set to **EINVAL**.

If the correct value is outside the range of representable values, ±**HUGE_VAL** is returned (according to the sign of the value), and **errno** is set to **ERANGE**.

If the correct value would cause underflow, **0** is returned, and **errno** is set to **ERANGE**.

ERRORS

wcstod( ) and **wstod( )** will fail if:

ERANGE The value to be returned would cause overflow or underflow.

wcstod( ) and **wcstod( )** may fail if:

EINVAL No conversion could be performed.

SEE ALSO

iswspace(3I), **localeconv**(3C), **scanf**(3S), **setlocale**(3C), **wcstol**(3I)

NOTES

Because **0** is returned on error and is also a valid return on success, an application wishing to check for error situations should set **errno** to **0**, call **wcstod( )** or **wstod( )**, then check **errno** and if it is non-zero, assume an error has occurred.

NAME

wcstol, wstol, watol, watoll, watoi – convert wide character string to long integer

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <wchar.h>

long int wcstol(const wchar_t ∗*nptr,* **wchar_t** ∗∗*endptr,* **int** *base***);**
long int wstol(const wchar_t ∗*nptr,* **wchar_t** ∗∗*endptr,* **int** *base***);**
long watol(wchar_t ∗*nptr***);**
long long watoll(wchar_t ∗*nptr***);**
int watoi(wchar_t ∗*nptr***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **wcstol( )** and **wstol( )** functions convert the initial portion of the wide character string pointed to by *nptr* to **long int** representation. They first decompose the input wide character string into three parts: an initial, possibly empty, sequence of white-space wide-character codes (as specified by **iswspace**(3I)), a subject sequence interpreted as an integer represented in some radix determined by the value of *base*; and a final wide character string of one or more unrecognised wide character codes, including the terminating null wide-character code of the input wide character string. They then attempt to convert the subject sequence to an integer, and return the result.

If the value of *base* is **0**, the expected form of the subject sequence is that of a decimal constant, octal constant or hexadecimal constant, any of which may be preceded by a '+' or '–' sign. A decimal constant begins with a non-zero digit, and consists of a sequence of decimal digits. An octal constant consists of the prefix '0' optionally followed by a sequence of the digits '0' to '7' only. A hexadecimal constant consists of the prefix '0x' or '0X' followed by a sequence of the decimal digits and letters 'a' (or 'A') to 'f' (or 'F') with values 10 to 15 respectively.

If the value of *base* is between **2** and **36**, the expected form of the subject sequence is a sequence of letters and digits representing an integer with the radix specified by *base*, optionally preceded by a '+' or '–' sign, but not including an integer suffix. The letters from 'a' (or 'A') to 'z' (or 'Z') inclusive are ascribed the values 10 to 35; only letters whose ascribed values are less than that of *base* are permitted. If the value of *base* is **16**, the wide-character code representations of '0x' or '0X' may optionally precede the sequence of letters and digits, following the sign if present.

The subject sequence is defined as the longest initial subsequence of the input wide character string, starting with the first non-white-space wide-character code, that is of the expected form. The subject sequence contains no wide-character codes if the input wide character string is empty or consists entirely of white-space wide-character code, or if the first non-white-space wide-character code is other than a sign or a permissible letter or digit.

If the subject sequence has the expected form and the value of *base* is **0**, the sequence of wide-character codes starting with the first digit is interpreted as an integer constant. If the subject sequence has the expected form and the value of *base* is between **2** and **36**, it is

used as the base for conversion, ascribing to each letter its value as given above. If the subject sequence begins with a minus sign (-), the value resulting from the conversion is negated. A pointer to the final wide character string is stored in the object pointed to by *endptr*, provided that *endptr* is not a null pointer.

In other than the POSIX locale, additional implementation-dependent subject sequence forms may be accepted.

If the subject sequence is empty or does not have the expected form, no conversion is performed; the value of *nptr* is stored in the object pointed to by *endptr*, provided that *endptr* is not a null pointer.

watol( ) is equivalent to **wstol(***str***, (wchar_t** ****)NULL, 10)**.

watoll( ) is the long-long (double long) version of **watol( )**.

watoi( ) is equivalent to **(int)watol( )**.

RETURN VALUES

Upon successful completion, **wcstol( )** and **wstol( )** return the converted value, if any. If no conversion could be performed, **0** is returned, and **errno** may be set to indicate the error. If the correct value is outside the range of representable values, **{LONG_MAX}** or **{LONG_MIN}** is returned (according to the sign of the value), and **errno** is set to **ERANGE**.

ERRORS

wcstol( ) and **wstol( )** will fail if:

EINVAL The value of *base* is not supported.

ERANGE The value to be returned is not representable.

wcstol( ) and **wstol( )** may fail if:

EINVAL No conversion could be performed.

SEE ALSO

iswalpha(3I), **iswspace**(3I), **scanf**(3S), **wcstod**(3I),

NOTES

Because **0**, **{LONG_MIN}**, and **{LONG_MAX}** are returned on error and are also valid returns on success, an application wishing to check for error situations should set **errno** to **0**, call **wcstol( )** or **wstol( ),** then check **errno** and if it is non-zero assume an error has occurred.

Truncation from **long long** to **long** can take place upon assignment or by an explicit cast.

NAME

wcstoul – convert wide character string to unsigned long

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <wchar.h>

unsigned long int wcstoul(const wchar_t **nptr,* **wchar_t** ***endptr,* **int** *base***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **wcstoul( )** function converts the initial portion of the wide character string pointed to by *nptr* to **unsigned long int** representation. It first decomposes the input wide-character string into three parts: an initial, possibly empty, sequence of white-space wide-character codes (as specified by the function **iswspace**(3I)); a subject sequence interpreted as an integer represented in some radix determined by the value of *base*; and a final wide-character string of one or more unrecognized wide character codes, including the terminating null wide-character code of the input wide character string. It then attempts to convert the subject sequence to an unsigned integer, and returns the result.

If the value of *base* is 0, the expected form of the subject sequence is that of a decimal constant, an octal constant, or a hexadecimal constant, any of which may be preceded by a '+' or a '–' sign. A decimal constant begins with a non-zero digit, and consists of a sequence of decimal digits. An octal constant consists of the prefix '0', optionally followed by a sequence of the digits '0' to '7' only. A hexadecimal constant consists of the prefix '0x' or '0X', followed by a sequence of the decimal digits and letters 'a' (or 'A') to 'f' (or 'F'), with values 10 to 15, respectively.

If the value of *base* is between 2 and 36, the expected form of the subject sequence is a sequence of letters and digits representing an integer with the radix specified by *base*, optionally preceded by a '+' or a '–' sign, but not including an integer suffix. The letters from 'a' (or 'A') to 'z' (or 'Z') inclusive are ascribed the values 10 to 35; only letters whose ascribed values are less than that of *base* are permitted. If the value of *base* is 16, the wide-character codes '0x' or '0X' may optionally precede the sequence of letters and digits, following the sign, if present.

The subject sequence is defined as the longest initial subsequence of the input wide-character string, starting with the first wide-character code that is not a white space and is of the expected form. The subject sequence contains no wide-character codes if the input wide-character string is empty or consists entirely of white-space wide-character codes, or if the first wide-character code that is not a white space is other than a sign or a permissible letter or digit.

If the subject sequence has the expected form and the value of *base* is 0, the sequence of wide-character codes starting with the first digit is interpreted as an integer constant. If the subject sequence has the expected form and the value of *base* is between 2 and 36, it is used as the base for conversion, ascribing to each letter its value as given above. If the subject sequence begins with a minus sign, the value resulting from the conversion is negated. A pointer to the final wide character string is stored in the object pointed to by *endptr*, provided that *endptr* is not a null pointer.

In other than the POSIX locale, additional implementation-dependent subject sequence forms may be accepted.

If the subject sequence is empty or does not have the expected form, no conversion is performed; the value of *nptr* is stored in the object pointed to by *endptr*, provided that *endptr* is not a null pointer.

RETURN VALUE

Upon successful completion, **wcstoul( )** returns the converted value, if any. If no conversion could be performed, **0** is returned and **errno** may be set to indicate the error. If the correct value is outside the range of representable values, **{ULONG_MAX}** is returned and **errno** is set to **ERANGE**.

ERRORS

wcstoul( ) will fail if:

EINVAL The value of *base* is not supported.

ERANGE The value to be returned is not representable.

wcstoul( ) function may fail if:

EINVAL No conversion could be performed.

SEE ALSO

isspace(3C), **iswalpha**(3I), **scanf**(3S), **wcstod**(3I), **wcstol**(3I)

WARNINGS

Because **0** and **{ULONG_MAX}** are returned on error and **0** is also a valid return on success, an application wishing to check for error situations should set **errno** to **0**, call **wcstoul( )**, then check **errno** and if it is non-zero, assume an error has occurred.

Unlike **wcstod**(3I) and **wcstol**(3I), **wcstoul( )** must always return a non-negative number; so, using the return value of **wcstoul( )** for out-of-range numbers with **wcstoul( )**.

NAME

wcstring, wcscat, wscat, wcsncat, wsncat, wcscmp, wscmp, wcsncmp, wsncmp, wcscpy, wscpy, wcsncpy, wsncpy, wcslen, wslen, wcwidth, wcswidth, wcschr, wschr, wcsrchr, wsrchr, windex, wrindex, wcspbrk, wspbrk, wcswcs, wcsspn, wsspn, wcscspn, wscspn, wcstok, wstok – wide character string operations

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <wchar.h>

wchar_t ∗**wcscat(wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2***);**
wchar_t ∗**wscat(wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2***);**

wchar_t ∗**wcsncat(wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2,* **size_t** *n***);**
wchar_t ∗**wsncat(wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2,* **size_t** *n***);**

int wcscmp(const wchar_t ∗ws1, **const wchar_t** ∗ws2**);**
int wscmp(const wchar_t ∗ws1, **const wchar_t** ∗ws2**);**

int wcsncmp(const wchar_t ∗*ws1,* **const wchar_t** ∗*ws2,* **size_t** *n***);**
int wsncmp(const wchar_t ∗*ws1,* **const wchar_t** ∗*ws2,* **size_t** *n***);**

wchar_t ∗**wcscpy(wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2***);**
wchar_t ∗**wscpy(wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2***);**

wchar_t ∗**wcsncpy(wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2,* **size_t** *n***);**
wchar_t ∗**wsncpy(wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2,* **size_t** *n***);**

size_t wcslen(const wchar_t ∗*ws***);**
size_t wslen(const wchar_t ∗*ws***);**

int wcwidth(wint_t *wc***);**

int wcswidth(const wchar_t ∗*pwcs,* **size_t** *n***);**

wchar_t ∗**wcschr(const wchar_t** ∗*ws,* **wint_t** *wc***);**
wchar_t ∗**wschr(const wchar_t** ∗*ws,* **wint_t** *wc***);**

wchar_t ∗**wcsrchr(const wchar_t** ∗*ws,* **wint_t** *wc***);**
wchar_t ∗**wsrchr(const wchar_t** ∗*ws,* **wint_t** *wc***);**

wchar_t ∗**windex(const wchar_t** ∗*ws***, wchar_t** *wc***);**
wchar_t ∗**wrindex(const wchar_t** ∗*ws***, wchar_t** *wc***);**

wchar_t ∗**wcspbrk(const wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2***);**
wchar_t ∗**wspbrk(const wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2***);**

wchar_t ∗**wcswcs(const wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2***);**

size_t wcsspn(const wchar_t ∗*ws1,* **const wchar_t** ∗*ws2***);**
size_t wsspn(const wchar_t ∗*ws1,* **const wchar_t** ∗*ws2***);**

size_t wcscspn(const wchar_t ∗*ws1,* **const wchar_t** ∗*ws2***);**
size_t wscspn(const wchar_t ∗*ws1,* **const wchar_t** ∗*ws2***);**

wchar_t ∗**wcstok(wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2***);**
wchar_t ∗**wstok(wchar_t** ∗*ws1,* **const wchar_t** ∗*ws2***);**

MT-LEVEL MT-Safe

DESCRIPTION These functions operate on wide character strings terminated by **wchar_t NULL** characters. During appending or copying, these routines do not check for an overflow condition of the receiving string. In the following, *ws*, *ws1*, and *ws2* point to wide character strings terminated by a **wchar_t NULL**.

wcscat(), wscat() The **wcscat( )** and **wscat( )** functions append a copy of the wide character string pointed to by *ws2* (including the terminating null wide-character code) to the end of the wide character string pointed to by *ws1*. The initial wide-character code of *ws2* overwrites the null wide-character code at the end of *ws1*. If copying takes place between objects that overlap, the behaviour is undefined. Both functions return *s1*; no return value is reserved to indicate an error.

wcsncat(), wsncat() The **wcsncat( )** and **wsncat( )** functions append not more than *n* wide-character codes (a null wide-character code and wide character codes that follow it are not appended) from the array pointed to by *ws2* to the end of the wide character string pointed to by *ws1*. The initial wide-character code of *ws2* overwrites the null wide-character code at the end of *ws1*. A terminating null wide-character code is always appended to the result. Both functions return *ws1*; no return value is reserved to indicate an error.

wcscmp(), wscmp() The **wcscmp( )** and **wscmp( )** functions compare the wide character string pointed to by *ws1* to the wide character string pointed to by *ws2*. The sign of a non-zero return value is determined by the sign of the difference between the values of the first pair of wide-character codes that differ in the objects being compared. Upon completion, both functions return an integer greater than, equal to, or less than zero, if the wide character string pointed to by *ws1* is greater than, equal to, or less than the wide character string pointed to by *ws2*.

wcsncmp(), wsncmp() The **wcsncmp( )** and **wsncmp( )** functions compare not more than *n* wide-character codes (wide-character codes that follow a null wide character code are not compared) from the array pointed to by *ws1* to the array pointed to by *ws2*. The sign of a non-zero return value is determined by the sign of the difference between the values of the first pair of wide-character codes that differ in the objects being compared. Upon successful completion, both functions return an integer greater than, equal to, or less than zero, if the possibly null-terminated array pointed to by *ws1* is greater than, equal to, or less than the possibly null-terminated array pointed to by *ws2*.

wcscpy(), wscpy() The **wcscpy( )** and **wscpy( )** functions copy the wide character string pointed to by *ws2* (including the terminating null wide-character code) into the array pointed to by *ws1*. If copying takes place between objects that overlap, the behaviour is undefined. Both functions return *ws1*; no return value is reserved to indicate an error.

wcsncpy(), wsncpy() The **wcsncpy( )** and **wsncpy( )** functions copy not more than *n* wide-character codes (wide-character codes that follow a null wide character code are not copied) from the array pointed to by *ws2* to the array pointed to by *ws1*. If copying takes place between

objects that overlap, the behaviour is undefined. If the array pointed to by *ws2* is a wide character string that is shorter than *n* wide-character codes, null wide-character codes are appended to the copy in the array pointed to by *ws1*, until a total *n* wide-character codes are written. Both functions return *ws1*; no return value is reserved to indicate an error.

wcslen(), wslen() The **wcslen( )** and **wslen( )** functions compute the number of wide-character codes in the wide character string to which *ws* points, not including the terminating null wide-character code. Both functions return *ws*; no return value is reserved to indicate an error.

wcwidth() The **wcwidth( )** function determines the number of column positions required for the wide character *wc*. The value of *wc* must be a character representable as a **wchar_t**, and must be a wide-character code corresponding to a valid character in the current locale. The function returns **0** if *wc* is a null wide-character code; the number of column positions to be occupied by the wide-character code *wc*; or **–1** if *wc* does not correspond to a printing wide-character code.

wcswidth() The **wcswidth( )** function determines the number of column positions required for *n* wide-character codes (or fewer than *n* wide-character codes if a null wide-character code is encountered before *n* wide-character codes are exhausted) in the string pointed to by *pwcs*. The function returns **0** if *pwcs* points to a null wide-character code; the number of column positions to be occupied by the wide character string pointed to by *pwcs*; or **–1** if any wide-character code in the wide character string pointed to by *pwcs* is not a printing wide-character code.

wcschr(), wschr() The **wcschr( )** and **wschr( )** functions locate the first occurrence of *wc* in the wide character string pointed to by *ws*. The value of *wc* must be a character representable as a type **wchar_t** and must be a wide-character code corresponding to a valid character in the current locale. The terminating null wide-character code is considered to be part of the wide character string. Upon completion, both functions return a pointer to the wide-character code, or a null pointer if the wide-character code is not found.

wcsrchr(), wsrchr() The **wcsrchr( )** and **wsrchr( )** functions locate the last occurrence of *wc* in the wide character string pointed to by *ws*. The value of *wc* must be a character representable as a type **wchar_t** and must be a wide-character code corresponding to a valid character in the current locale. The terminating null wide-character code is considered to be part of the wide character string. Upon successful completion, both functions return a pointer to the wide-character code, or a null pointer if *wc* does not occur in the wide character string.

windex(), wrindex() The **windex( )** and **wrindex( )** functions behave the same as **wschr( )** and **wsrchr( )**, respectively.

wcspbrk(), wspbrk() The **wcspbrk( )** and **wspbrk( )** functions locate the first occurrence in the wide character string pointed to by *ws1* of any wide-character code from the wide character string pointed to by *ws2*. Upon successful completion, the function returns a pointer to the wide-character code, or a null pointer if no wide-character code from *ws2* occurs in *ws1*.

wcswcs()

The **wcswcs()** function locates the first occurrence in the wide character string pointed to by *ws1* of the sequence of wide-character codes (excluding the terminating null wide-character code) in the wide character string pointed to by *ws2*. Upon successful completion, the function returns a pointer to the located wide character string, or a null pointer if the wide character string is not found. If *ws2* points to a wide character string with zero length, the function returns *ws1*.

wcsspn(), wsspn()

The **wcsspn()** and **wsspn()** functions compute the length of the maximum initial segment of the wide character string pointed to by *ws1* which consists entirely of wide-character codes from the wide string pointed to by *ws2*. Both functions return *ws1*; no return value is reserved to indicate an error.

wcscspn(), wscspn()

The **wcscspn()** and **wscspn()** functions compute the length of the maximum initial segment of the wide character string pointed to by *ws1* which consists entirely of wide-character codes *not* from the wide character string pointed to by *ws2*. Both functions return *ws1*; no return value is reserved to indicate an error.

wcstok(), wstok()

A sequence of calls to the **wcstok()** and **wstok()** functions break the wide character string pointed to by *ws1* into a sequence of tokens, each of which is delimited by a wide-character code from the wide character string pointed to by *ws2*. The first call in the sequence has *ws1* as its first argument, and is followed by calls with a null pointer as their first argument. The separator string pointed to by *ws2* may be different from call to call.

The first call in the sequence searches the wide character string pointed to by *ws1* for the first wide-character code that is *not* contained in the current separator string pointed to by *ws2*. If no such wide-character code is found, then there are no tokens in the wide character string pointed to by *ws1*, and **wcstok()** and **wstok()** return a null pointer. If such a wide-character code is found, it is the start of the first token.

wcstok() and **wstok()** then search from that point for a wide-character code that *is* contained in the current separator string. If no such wide-character code is found, the current token extends to the end of the wide character string pointed to by *ws1*, and subsequent searches for a token will return a null pointer. If such a wide-character code is found, it is overwritten by a null wide character, which terminates the current token. **wcstok()** and **wstok()** save a pointer to the following wide-character code, from which the next search for a token will start.

Each subsequent call, with a null pointer as the value of the first argument, starts searching from the saved pointer and behaves as described above.

Upon successful completion, both functions return a pointer to the first wide-character code of a token. Otherwise, if there is no token, a null pointer is returned.

SEE ALSO

malloc(3C), **string**(3C), **wstring**(3I)

NAME

wcsxfrm, wsxfrm – wide character string transformation

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <wchar.h>

size_t wcsxfrm(wchar_t **ws1,* **const wchar_t** **ws2,* **size_t** *n***);**
size_t wsxfrm(wchar_t **ws1,* **const wchar_t** **ws2,* **size_t** *n***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **wcsxfrm( )** and **wcsxfrm( )** functions transform the wide character string pointed to by *ws2* and place the resulting wide character string into the array pointed to by *ws1*. The transformation is such that if either the **wcscmp**(3I) or **wscmp**(3I) functions are applied to two transformed wide strings, they return a value greater than, equal to, or less than 0, corresponding to the result of the **wcscoll**(3I) or **wscoll**(3I) function applied to the same two original wide character strings. No more than *n* wide-character codes are placed into the resulting array pointed to by *ws1*, including the terminating null wide-character code. If *n* is **0**, *ws1* is permitted to be a null pointer. If copying takes place between objects that overlap, the behaviour is undefined.

RETURN VALUES

wcsxfrm( ) and **wsxfrm( )** return the length of the transformed wide character string (not including the terminating null wide-character code). If the value returned is *n* or more, the contents of the array pointed to by *ws1* are indeterminate.

On error, **wcsxfrm()** and **wsxfrm()** return **(size_t)**–**1**, and set **errno** to indicate the error.

ERRORS

wcsxfrm( ) and **wsxfrm( )** may fail if:

EINVAL The wide character string pointed to by *ws2* contains wide-character codes outside the domain of the collating sequence.

ENOSYS The function is not supported.

NOTES

The transformation function is such that two transformed wide character strings can be ordered by the **wcscmp( )** or **wscmp( )** functions as appropriate to collating sequence information in the program's locale (category **LC_COLLATE**).

The fact that when *n* is **0**, *ws1* is permitted to be a null pointer, is useful to determine the size of the *ws1* array prior to making the transformation.

Because no return value is reserved to indicate an error, an application wishing to check for error situations should set **errno** to **0**, call **wcsxfrm( )** or **wsxfrm( )**, then check **errno** and if it is non-zero, assume an error has occurred.

SEE ALSO

wcscmp(3I), **wcscoll**(3I), **wscmp**(3I), **wscoll**(3I)

NAME

wctype – define character class

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <wchar.h>

wctype_t wctype(const char ∗*charclass***);**

MT-LEVEL

MT-Safe

DESCRIPTION

The **wctype( )** function is defined for valid character class names as defined in the current locale. The *charclass* is a string identifying a generic character class for which codeset-specific type information is required. The following character class names are defined in all locales — "alnum", "alpha", "blank" "cntrl", "digit", "graph", "lower", "print", "punct", "space", "upper", and "xdigit".

Additional character class names defined in the locale definition file (category **LC_CTYPE**) can also be specified.

The function returns a value of type **wctype_t**, which can be used as the second argument to subsequent calls of **iswctype**(3I). **wctype( )** determines values of **wctype_t** according to the rules of the coded character set defined by character type information in the program's locale (category **LC_CTYPE**). The values returned by **wctype( )** are valid until a call to **setlocale**(3C) that modifies the category **LC_CTYPE**.

RETURN VALUES

wctype( ) returns **0** if the given character class name is not valid for the current locale (category **LC_CTYPE**); otherwise it returns an object of type **wctype_t** that can be used in calls to **iswctype( )**.

SEE ALSO

iswctype(3I), **setlocale**(3C)

NAME

wordexp, wordfree – perform word expansions

SYNOPSIS

```
#include <wordexp.h>
int wordexp(const char *words, wordexp_t *pwordexp, int flags);
void wordfree(wordexp_t *pwordexp);
```

MT-LEVEL

MT-Safe

DESCRIPTION

The **wordexp()** function performs word expansions, subject to quoting, and places the list of expanded words into the structure pointed to by *pwordexp*.

The **wordfree()** function frees any memory allocated by **wordexp()** associated with *pwordexp*.

***words* Argument**

The *words* argument is a pointer to a string containing one or more words to be expanded. The expansions will be the same as would be performed by the shell if *words* were the part of a command line representing the arguments to a utility. Therefore, *words* must not contain an unquoted NEWLINE or any of the unquoted shell special characters:

| & ; < >

except in the context of command substitution. It also must not contain unquoted parentheses or braces, except in the context of command or variable substitution. If the argument *words* contains an unquoted comment character (number sign) that is the beginning of a token, **wordexp()** may treat the comment character as a regular character, or may interpret it as a comment indicator and ignore the remainder of *words*.

***pwordexp* Argument**

The structure type **wordexp_t** is defined in the header **<wordexp.h>** and includes at least the following members:

size_t we_wordc	Count of words matched by *words*.
char **we_wordv	Pointer to list of expanded words.
size_t we_offs	Slots to reserve at the beginning of *pwordexp*–>**we_wordv**.

The **wordexp()** function stores the number of generated words into *pwordexp*–>**we_wordc** and a pointer to a list of pointers to words in *pwordexp*–>**we_wordv**. Each individual field created during field splitting is a separate word in the *pwordexp*–>**we_wordv** list. The words are in order. The first pointer after the last word pointer will be a **NULL** pointer.

It is the caller's responsibility to allocate the storage pointed to by *pwordexp*. The **wordexp()** function allocates other space as needed, including memory pointed to by *pwordexp*–>**we_wordv**. The **wordfree()** function frees any memory associated with *pwordexp* from a previous call to **wordexp()**.

***flags* Argument**

The *flags* argument is used to control the behavior of **wordexp()**. The value of *flags* is the bitwise inclusive OR of zero or more of the following constants, which are defined in **<wordexp.h>**:

WRDE_APPEND Append words generated to the ones from a previous call to

wordexp( ).

WRDE_DOOFFS Make use of *pwordexp*–> **we_offs**. If this flag is set, *pwordexp*–> **we_offs** is used to specify how many **NULL** pointers to add to the beginning of *pwordexp*–> **we_wordv**. In other words, *pwordexp*–> **we_wordv** will point to *pwordexp*–> **we_offs NULL** pointers, followed by *pwordexp*–> **we_wordc** word pointers, followed by a **NULL** pointer.

WRDE_NOCMD Fail if command substitution is requested.

WRDE_REUSE The *pwordexp* argument was passed to a previous successful call to **wordexp( )**, and has not been passed to **wordfree( )**. The result will be the same as if the application had called **wordfree( )** and then called **wordexp( )** without **WRDE_REUSE**.

WRDE_SHOWERR Do not redirect **stderr** to **/dev/null**.

WRDE_UNDEF Report error on an attempt to expand an undefined shell variable.

The **WRDE_APPEND** flag can be used to append a new set of words to those generated by a previous call to **wordexp( )**. The following rules apply when two or more calls to **wordexp( )** are made with the same value of *pwordexp* and without intervening calls to **wordfree( )**:

1. The first such call must not set **WRDE_APPEND**. All subsequent calls must set it.
2. All of the calls must set **WRDE_DOOFFS**, or all must not set it.
3. After the second and each subsequent call, *pwordexp*–> **we_wordv** will point to a list containing the following:
   a. zero or more **NULL** pointers, as specified by **WRDE_DOOFFS** and *pwordexp*–> **we_offs**.
   b. pointers to the words that were in the *pwordexp*–> **we_wordv** list before the call, in the same order as before.
   c. pointers to the new words generated by the latest call, in the specified order.
4. The count returned in *pwordexp*–> **we_wordc** will be the total number of words from all of the calls.
5. The application can change any of the fields after a call to **wordexp( )**, but if it does it must reset them to the original value before a subsequent call, using the same *pwordexp* value, to **wordfree( )** or **wordexp( )** with the **WRDE_APPEND** or **WRDE_REUSE** flag.

If *words* contains an unquoted:

NEWLINE **| & ; < > ( ) { }**

in an inappropriate context, **wordexp( )** will fail, and the number of expanded words will be zero.

Unless **WRDE_SHOWERR** is set in *flags*, **wordexp( )** will redirect **stderr** to **/dev/null** for any utilities executed as a result of command substitution while expanding *words*. If **WRDE_SHOWERR** is set, **wordexp( )** may write messages to *stderr* if syntax errors are

detected while expanding *words*.

If **WRDE_DOOFFS** is set, then *pwordexp*–> **we_offs** must have the same value for each **wordexp( )** call and **wordfree( )** call using a given *pwordexp*.

The following constants are defined as error return values:

WRDE_BADCHAR	One of the unquoted characters: NEWLINE **| & ; < > ( ) { }** appears in *words* in an inappropriate context.
WRDE_BADVAL	Reference to undefined shell variable when **WRDE_UNDEF** is set in *flags*.
WRDE_CMDSUB	Command substitution requested when **WRDE_NOCMD** was set in flags.
WRDE_NOSPACE	Attempt to allocate memory failed.
WRDE_SYNTAX	Shell syntax error, such as unbalanced parentheses or unterminated string.

RETURN VALUES

The following values are returned by **wordexp( )**:

0	successful completion.
non-zero	an error has occurred.
WRDE_NOSPACE	*pwordexp*–> **we_wordc** and *pwordexp*–> **we_wordv** will be updated to reflect any words that were successfully expanded. In other cases, they will not be modified.

The **wordfree( )** function returns no value.

USAGE

This function is intended to be used by an application that wants to do all of the shell's expansions on a word or words obtained from a user. For example, if the application prompts for a filename (or list of filenames) and then uses **wordexp( )** to process the input, the user could respond with anything that would be valid as input to the shell.

The **WRDE_NOCMD** flag is provided for applications that, for security or other reasons, want to prevent a user from executing shell commands. Disallowing unquoted shell special characters also prevents unwanted side effects such as executing a command or writing a file.

SEE ALSO

fnmatch(3C), **glob**(3C)

NAME

wsprintf – formatted output conversion

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

```
#include <stdio.h>
#include <widec.h>

int wsprintf(wchar_t *s, const char *format, /* arg */ ... );
```

MT-LEVEL

MT-Safe

DESCRIPTION

wsprintf( ) outputs a Process Code string ending with a Process Code (*wchar_t*) NULL character. It is the user's responsibility to allocate enough space for this *wchar_t* string.

This returns the number of Process Code characters (excluding the NULL terminator) that have been written. The conversion specifications and behavior of **wsprintf( )** are the same as the regular **sprintf**(3S) function except that the result is a Process Code string for **wsprintf( ),** and on Extended Unix Code (EUC) character string for **sprintf**(3S).

RETURN VALUES

Upon success, **wsprintf( )** returns the number of characters printed. When an error condition is encountered, a negative value is returned.

SEE ALSO

wsscanf(3I), **printf**(3S), **scanf**(3S), **sprintf**(3S)

NAME

wsscanf – formatted input conversion

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

```
#include <stdio.h>
#include <widec.h>

int wsscanf(wchar_t *s, const char *format, /* pointer */ ... );
```

MT-LEVEL

MT-Safe

DESCRIPTION

wsscanf( ) reads Process Code characters from the Process Code string *s*, interprets them according to the *format*, and stores the results in its arguments. **wsscanf( )** expects, as arguments, a control string *format*, and a set of *pointer* arguments indicating where the converted input should be stored. The results are undefined if there are insufficient *arg*s for the format. If the format is exhausted while *arg*s remain, the excess *arg*s are simply ignored.

The conversion specifications and behavior of **wsscanf( )** are the same as the regular **sscanf**(3S) function except that the source is a Process Code string for **wsscanf( ),** and on Extended Unix Code (EUC) character string for **sscanf**(3S).

RETURN VALUES

wsscanf( ) returns the number of characters matched. On error **wsscanf( )** returns a negative value.

SEE ALSO

wsprintf(3I), **printf**(3S), **scanf**(3S)

NAME

wstring, wscasecmp, wsncasecmp, wsdup, wscol – Process Code string operations

SYNOPSIS

cc [*flag* ...] *file* ... **–lw** [*library* ...]

#include <widec.h>

int wscasecmp(const wchar_t ∗*s1***, const wchar_t** ∗*s2***);**
int wsncasecmp(const wchar_t ∗*s1***, const wchar_t** ∗*s2***, int** *n***);**
wchar_t ∗**wsdup(const wchar_t** ∗*s***);**
int wscol(const wchar_t ∗*s***);**

MT-LEVEL

MT-Safe

DESCRIPTION

These functions operate on Process Code strings terminated by **wchar_t NULL** characters. During appending or copying, these routines do not check for an overflow condition of the receiving string. In the following, *s*, *s1*, and *s2* point to Process Code strings terminated by a **wchar_t NULL**.

wscasecmp(), wsncasecmp()

The **wscasecmp( )** function compares its arguments, ignoring case, and returns an integer greater than, equal to, or less than 0, depending upon whether *s1* is lexicographically greater than, equal to, or less than *s2*. **wsncasecmp( )** makes the same comparison but compares at most *n* Process Code characters. The four Extended Unix Code (EUC) codesets are ordered from lowest to highest as 0, 2, 3, 1 when characters from different codesets are compared.

wsdup( )

The **wsdup( )** function returns a pointer to a new Process Code string, which is a duplicate of the string pointed to by *s*. The space for the new string is obtained using **malloc**(3C). If the new string cannot be created, a null pointer is returned.

wscol( )

The **wscol( )** function returns the screen display width (in columns) of the Process Code string *s*.

SEE ALSO

malloc(3C), **string**(3C), **wcstring**(3I)

NAME

xdr – library routines for external data representation

MT-LEVEL

Safe

DESCRIPTION

XDR routines allow C programmers to describe arbitrary data structures in a machine-independent fashion. Data for remote procedure calls (RPC) are transmitted using these routines.

Index to Routines

The following table lists XDR routines and the manual reference pages on which they are described:

XDR Routine	Manual Reference Page
xdr_array	**xdr_complex**(3N)
xdr_bool	**xdr_simple**(3N)
xdr_bytes	**xdr_complex**(3N)
xdr_char	**xdr_simple**(3N)
xdr_control	**xdr_admin**(3N)
xdr_destroy	**xdr_create**(3N)
xdr_double	**xdr_simple**(3N)
xdr_enum	**xdr_simple**(3N)
xdr_float	**xdr_simple**(3N)
xdr_free	**xdr_simple**(3N)
xdr_getpos	**xdr_admin**(3N)
xdr_hyper	**xdr_simple**(3N)
xdr_inline	**xdr_admin**(3N)
xdr_int	**xdr_simple**(3N)
xdr_long	**xdr_simple**(3N)
xdr_longlong_t	**xdr_simple**(3N)
xdr_opaque	**xdr_complex**(3N)
xdr_pointer	**xdr_complex**(3N)
xdr_quadruple	**xdr_simple**(3N)
xdr_reference	**xdr_complex**(3N)
xdr_setpos	**xdr_admin**(3N)
xdr_short	**xdr_simple**(3N)
xdr_sizeof	**xdr_admin**(3N)
xdr_string	**xdr_complex**(3N)
xdr_u_char	**xdr_simple**(3N)
xdr_u_hyper	**xdr_simple**(3N)
xdr_u_int	**xdr_simple**(3N)
xdr_u_long	**xdr_simple**(3N)
xdr_u_longlong_t	**xdr_simple**(3N)
xdr_u_short	**xdr_simple**(3N)
xdr_union	**xdr_complex**(3N)
xdr_vector	**xdr_complex**(3N)
xdr_void	**xdr_simple**(3N)

xdr_wrapstring	**xdr_complex**(3N)
xdrmem_create	**xdr_create**(3N)
xdrrec_create	**xdr_create**(3N)
xdrrec_endofrecord	**xdr_admin**(3N)
xdrrec_eof	**xdr_admin**(3N)
xdrrec_readbytes	**xdr_admin**(3N)
xdrrec_skiprecord	**xdr_admin**(3N)
xdrstdio_create	**xdr_create**(3N)

SEE ALSO

rpc(3N), **xdr_admin**(3N), **xdr_complex**(3N), **xdr_create**(3N), **xdr_simple**(3N)

NAME

xdr_admin, xdr_control, xdr_getpos, xdr_inline, xdrrec_endofrecord, xdrrec_eof, xdrrec_readbytes, xdrrec_skiprecord, xdr_setpos, xdr_sizeof – library routines for external data representation

MT-LEVEL

Safe

DESCRIPTION

XDR library routines allow C programmers to describe arbitrary data structures in a machine-independent fashion. Protocols such as remote procedure calls (RPC) use these routines to describe the format of the data.

These routines deal specifically with the management of the XDR stream.

Routines

See **rpc**(3N) for the definition of the **XDR** data structure. Note that any buffers passed to the XDR routines must be properly aligned. It is suggested that **malloc**(3C) be used to allocate these buffers or that the programmer insure that the buffer address is divisible evenly by four.

#include <rpc/xdr.h>

bool_t xdr_control(XDR ∗*xdrs***, int** *req***, void** ∗*info***);**

A function macro to change or retrieve various information about an XDR stream. *req* indicates the type of operation and *info* is a pointer to the information. The supported values of *req*, their argument types and what they do are:

XDR_GET_BYTES_AVAIL	**xdr_bytesrec ∗**	return number of bytes left unconsumed in the stream and a flag indicating whether or not this is the last fragment.

u_int xdr_getpos(const XDR ∗*xdrs***);**

A macro that invokes the get-position routine associated with the XDR stream, *xdrs*. The routine returns an unsigned integer, which indicates the position of the XDR byte stream. A desirable feature of XDR streams is that simple arithmetic works with this number, although the XDR stream instances need not guarantee this. Therefore, applications written for portability should not depend on this feature.

long ∗**xdr_inline(XDR** ∗*xdrs***, const int** *len***);**

A macro that invokes the in-line routine associated with the XDR stream, *xdrs*. The routine returns a pointer to a contiguous piece of the stream's buffer; *len* is the byte length of the desired buffer. Note: pointer is cast to **long ∗**.

Warning: **xdr_inline( )** may return **NULL** (**0**) if it cannot allocate a contiguous piece of a buffer. Therefore the behavior may vary among stream instances;

it exists for the sake of efficiency, and applications written for portability should not depend on this feature.

bool_t xdrrec_endofrecord(XDR * **xdrs , int** *sendnow***);**

This routine can be invoked only on streams created by **xdrrec_create( )** (see **xdr_create**(3N)). The data in the output buffer is marked as a completed record, and the output buffer is optionally written out if *sendnow* is non-zero. This routine returns TRUE if it succeeds, FALSE otherwise.

bool_t xdrrec_eof(XDR **xdrs***)***;**

This routine can be invoked only on streams created by **xdrrec_create( )**. After consuming the rest of the current record in the stream, this routine returns **TRUE** if there is no more data in the stream's input buffer. It returns **FALSE** if there is additional data in the stream's input buffer.

int xdrrec_readbytes(XDR **xdrs***)***, caddr_t** *addr***, u_int** *nbytes***);**

This routine can be invoked only on streams created by **xdrrec_create**( ). It attempts to read *nbytes* bytes from the XDR stream into the buffer pointed to by *addr*. On success this routine returns the number of bytes read, –1 on failure. A return value of 0 indicates an end of record.

bool_t xdrrec_skiprecord(XDR **xdrs***)***;**

This routine can be invoked only on streams created by **xdrrec_create( )** (see **xdr_create**(3N)). It tells the XDR implementation that the rest of the current record in the stream's input buffer should be discarded. This routine returns TRUE if it succeeds, FALSE otherwise.

bool_t xdr_setpos(XDR **xdrs***)***, const u_int** *pos***);**

A macro that invokes the set position routine associated with the XDR stream *xdrs*. The parameter *pos* is a position value obtained from **xdr_getpos( )**. This routine returns **TRUE** if the XDR stream was repositioned, and **FALSE** otherwise.

Warning: it is difficult to reposition some types of XDR streams, so this routine may fail with one type of stream and succeed with another. Therefore, applications written for portability should not depend on this feature.

unsigned long xdr_sizeof(xdrproc_t *func***, void** **data***)***;**

This routine returns the number of bytes required to encode *data* using the XDR filter function *func*, excluding potential overhead such as RPC headers or record markers. **O** is returned on error. This information might be used to select between transport protocols, or to determine the buffer size for various lower levels of RPC client and server creation routines, or to allocate storage when XDR is used outside of the RPC subsystem.

SEE ALSO

malloc(3C), **rpc**(3N), **xdr_complex**(3N), **xdr_create**(3N), **xdr_simple**(3N)

NAME

xdr_complex, xdr_array, xdr_bytes, xdr_opaque, xdr_pointer, xdr_reference, xdr_string, xdr_union, xdr_vector, xdr_wrapstring – library routines for external data representation

MT-LEVEL

Safe

DESCRIPTION

XDR library routines allow C programmers to describe complex data structures in a machine-independent fashion. Protocols such as remote procedure calls (RPC) use these routines to describe the format of the data. These routines are the XDR library routines for complex data structures. They require the creation of XDR stream (see **xdr_create**(3N)).

Routines

See **rpc**(3N) for the definition of the **XDR** data structure. Note that any buffers passed to the XDR routines must be properly aligned. It is suggested that **malloc( )** be used to allocate these buffers or that the programmer insure that the buffer address is divisible evenly by four.

#include <rpc/xdr.h>

bool_t xdr_array(XDR **xdrs***, caddr_t** **arrp***, u_int** **sizep***, const u_int** *maxsize***, const u_int** *elsize***, const xdrproc_t** *elproc***);**

xdr_array( ) translates between variable-length arrays and their corresponding external representations. The parameter *arrp* is the address of the pointer to the array, while *sizep* is the address of the element count of the array; this element count cannot exceed *maxsize*. The parameter *elsize* is the size of each of the array's elements, and *elproc* is an XDR routine that translates between the array elements' C form and their external representation. If **aarp* is null when decoding, **xdr_array( )** allocates memory and **aarp* points to it. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_bytes(XDR **xdrs***, char** ***sp***, u_int** **sizep***, const u_int** *maxsize***);**

xdr_bytes( ) translates between counted byte strings and their external representations. The parameter *sp* is the address of the string pointer. The length of the string is located at address *sizep*; strings cannot be longer than *maxsize*. If **sp* is null when decoding, **xdr_bytes( )** allocates memory and **sp* points to it. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_opaque(XDR **xdrs***, caddr_t** *cp***, const u_int** *cnt***);**

xdr_opaque( ) translates between fixed size opaque data and its external representation. The parameter *cp* is the address of the opaque object, and *cnt* is its size in bytes. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_pointer(XDR **xdrs***, char** ***objpp***, u_int** *objsize***, const xdrproc_t** *xdrobj***);**

Like **xdr_reference( )** except that it serializes **NULL** pointers, whereas **xdr_reference( )** does not. Thus, **xdr_pointer( )** can represent recursive data structures, such as binary trees or linked lists. If **objpp* is null when decoding, **xdr_pointer( )** allocates memory and **objpp* points to it.

bool_t xdr_reference(XDR **xdrs***, caddr_t** **pp***, u_int** *size***, const xdrproc_t** *proc***);**

xdr_reference( ) provides pointer chasing within structures. The parameter *pp* is the address of the pointer; *size* is the **sizeof** the structure that **pp* points to; and *proc* is an XDR procedure that translates the structure between its C form and its external representation. If **pp* is null when decoding, **xdr_reference( )** allocates memory and **pp* points to it. This routine returns **1** if it succeeds, **0** otherwise.

Warning: this routine does not understand **NULL** pointers. Use **xdr_pointer( )** instead.

bool_t xdr_string(XDR **xdrs***, char** ***sp***, const u_int** *maxsize***);**

xdr_string( ) translates between C strings and their corresponding external representations. Strings cannot be longer than *maxsize*. Note: *sp* is the address of the string's pointer. If **sp* is null when decoding, **xdr_string( )** allocates memory and **sp* points to it. This routine returns **TRUE** if it succeeds, **FALSE** otherwise. Note: **xdr_string( )** can be used to send an empty string (""), but not a NULL string.

bool_t xdr_union(XDR **xdrs***, enum_t** **dscmp***, char** **unp***,**
const struct xdr_discrim **choices***, const xdrproc_t (*** *defaultarm***);**

xdr_union( ) translates between a discriminated C **union** and its corresponding external representation. It first translates the discriminant of the union located at *dscmp*. This discriminant is always an **enum_t**. Next the union located at *unp* is translated. The parameter *choices* is a pointer to an array of **xdr_discrim** structures. Each structure contains an ordered pair of [*value, proc*]. If the union's discriminant is equal to the associated *value*, then the *proc* is called to translate the union. The end of the **xdr_discrim** structure array is denoted by a routine of value **NULL**. If the discriminant is not found in the *choices* array, then the *defaultarm* procedure is called (if it is not **NULL**). Returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_vector(XDR **xdrs***, char** **arrp***, const u_int** *size***, const u_int** *elsize***,**
const xdrproc_t *elproc***);**

xdr_vector( ) translates between fixed-length arrays and their corresponding external representations. The parameter *arrp* is the address of the pointer to the array, while *size* is the element count of the array. The parameter *elsize* is the **sizeof** each of the array's elements, and *elproc* is an XDR routine that translates

between the array elements' C form and their external representation. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_wrapstring(XDR ∗*xdrs***, char** ∗∗*sp***);**

A routine that calls **xdr_string(***xdrs***,** *sp***,** *maxuint***);** where *maxuint* is the maximum value of an unsigned integer.

Many routines, such as **xdr_array( )**, **xdr_pointer( )**, and **xdr_vector( )** take a function pointer of type **xdrproc_t( )**, which takes two arguments. **xdr_string( )**, one of the most frequently used routines, requires three arguments, while **xdr_wrapstring( )** only requires two. For these routines, **xdr_wrapstring( )** is desirable. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

SEE ALSO

rpc(3N), **xdr_admin**(3N), **xdr_create**(3N), **xdr_simple**(3N)

NAME

xdr_create, xdr_destroy, xdrmem_create, xdrrec_create, xdrstdio_create – library routines for external data representation stream creation

MT-LEVEL

MT-Safe

DESCRIPTION

XDR library routines allow C programmers to describe arbitrary data structures in a machine-independent fashion. Protocols such as remote procedure calls (RPC) use these routines to describe the format of the data.

These routines deal with the creation of XDR streams. XDR streams have to be created before any data can be translated into XDR format.

Routines

See **rpc**(3N) for the definition of the **XDR**, **CLIENT**, and **SVCXPRT** data structures. Note that any buffers passed to the XDR routines must be properly aligned. It is suggested that **malloc**(3C) be used to allocate these buffers or that the programmer insure that the buffer address is divisible evenly by four.

#include <rpc/xdr.h>

void xdr_destroy(XDR ∗*xdrs***);**

A macro that invokes the destroy routine associated with the XDR stream, *xdrs*. Destruction usually involves freeing private data structures associated with the stream. Using *xdrs* after invoking **xdr_destroy( )** is undefined.

void xdrmem_create(XDR ∗*xdrs***, const caddr_t** *addr***, const u_int** *size***,**
const enum xdr_op *op***);**

This routine initializes the XDR stream object pointed to by *xdrs*. The stream's data is written to, or read from, a chunk of memory at location *addr* whose length is no less than *size* bytes long. The *op* determines the direction of the XDR stream (either **XDR_ENCODE**, **XDR_DECODE**, or **XDR_FREE**).

void xdrrec_create(XDR ∗*xdrs***, const u_int** *sendsz***, const u_int** *recvsz***,**
const caddr_t *handle***, const int (**∗*readit***)(const void** ∗*read_handle***, char** ∗*buf***,**
const int *len***), const int (**∗*writeit***)(const void** ∗*write_handle***, const char** ∗*buf***,**
const int *len***));**

This routine initializes the read-oriented XDR stream object pointed to by *xdrs*. The stream's data is written to a buffer of size *sendsz*; a value of **0** indicates the system should use a suitable default. The stream's data is read from a buffer of size *recvsz*; it too can be set to a suitable default by passing a **0** value. When a stream's output buffer is full, *writeit* is called. Similarly, when a stream's input buffer is empty, *readit* is called. The behavior of these two routines is similar to the system calls **read( )** and **write( )** (see **read**(2) and **write**(2), respectively), except that an appropriate handle (*read_handle* or *write_handle*) is passed to the former routines as the first parameter instead of a file descriptor. Note: the XDR

stream's *op* field must be set by the caller.

Warning: this XDR stream implements an intermediate record stream. Therefore there are additional bytes in the stream to provide record boundary information.

void xdrstdio_create(XDR ∗*xdrs*, **FILE** ∗*file*, **const enum xdr_op** *op*);

This routine initializes the XDR stream object pointed to by *xdrs*. The XDR stream data is written to, or read from, the standard I/O stream *file*. The parameter *op* determines the direction of the XDR stream (either **XDR_ENCODE**, **XDR_DECODE**, or **XDR_FREE**).

Warning: the destroy routine associated with such XDR streams calls **fflush( )** on the *file* stream, but never **fclose( )** (see **fclose**(3S)).

Failure of any of these functions can be detected by first initializing the *x_ops* field in the **XDR** structure (*xdrs→x_ops*) to **NULL** before calling the xdr∗_create() function. After the return from the xdr∗_create() function, if the *x_ops* field is still **NULL ,** the call has failed. If the *x_ops* field contains some other value, the call can be assumed to have succeeded.

SEE ALSO

read(2), **write**(2), **malloc**(3C), **rpc**(3N), **xdr_admin**(3N), **xdr_complex**(3N), **xdr_simple**(3N), **fclose**(3S)

NAME

xdr_simple, xdr_bool, xdr_char, xdr_double, xdr_enum, xdr_float, xdr_free, xdr_hyper, xdr_int, xdr_long, xdr_longlong_t, xdr_quadruple, xdr_short, xdr_u_char, xdr_u_hyper, xdr_u_int, xdr_u_long, xdr_u_longlong_t, xdr_u_short, xdr_void – library routines for external data representation

MT-LEVEL

Safe

DESCRIPTION

XDR library routines allow C programmers to describe simple data structures in a machine-independent fashion. Protocols such as remote procedure calls (RPC) use these routines to describe the format of the data.

These routines require the creation of XDR streams (see **xdr_create**(3N)).

Routines

See **rpc**(3N) for the definition of the **XDR** data structure. Note that any buffers passed to the XDR routines must be properly aligned. It is suggested that **malloc**(3C) be used to allocate these buffers or that the programmer insure that the buffer address is divisible evenly by four.

#include <rpc/xdr.h>

bool_t xdr_bool(XDR **xdrs***, bool_t** **bp***);**

xdr_bool( ) translates between booleans (C integers) and their external representations. When encoding data, this filter produces values of either **1** or **0**. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_char(XDR **xdrs***, char** **cp***);**

xdr_char( ) translates between C characters and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise. Note: encoded characters are not packed, and occupy 4 bytes each. For arrays of characters, it is worthwhile to consider **xdr_bytes( )**, **xdr_opaque( )**, or **xdr_string( )** (see **xdr_complex**(3N)).

bool_t xdr_double(XDR **xdrs***, double** **dp***);**

xdr_double( ) translates between C **double** precision numbers and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_enum(XDR **xdrs***, enum_t** **ep***);**

xdr_enum( ) translates between C **enum**s (actually integers) and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_float(XDR **xdrs***, float** **fp***);**

xdr_float( ) translates between C **float**s and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

void xdr_free(xdrproc_t *proc***, char** ∗*objp***);**

Generic freeing routine. The first argument is the XDR routine for the object being freed. The second argument is a pointer to the object itself. Note: the pointer passed to this routine is not freed, but what it points to is freed (recursively, depending on the XDR routine).

bool_t xdr_hyper(XDR ∗*xdrs***, longlong_t** ∗*llp***);**

xdr_hyper( ) translates between ANSI C **long long** integers and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_int(XDR ∗*xdrs***, int** ∗*ip***);**

xdr_int( ) translates between C integers and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_long(XDR ∗*xdrs***, long** ∗*lp***);**

xdr_long( ) translates between C **long** integers and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_longlong_t(XDR ∗*xdrs***, longlong_t** ∗*llp***);**

xdr_longlong_t( ) translates between ANSI C **long long** integers and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise. This routine is identical to **xdr_hyper( )**.

bool_t xdr_quadruple(XDR ∗*xdrs***, long double** ∗*pq***);**

xdr_quadruple( ) translates between IEEE quadruple precision floating point numbers and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_short(XDR ∗*xdrs***, short** ∗*sp***);**

xdr_short( ) translates between C **short** integers and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_u_char(XDR ∗*xdrs***, unsigned char** ∗*ucp***);**

xdr_u_char( ) translates between **unsigned** C characters and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_u_hyper(XDR ∗*xdrs***, u_longlong_t** ∗*ullp***);**

xdr_u_hyper( ) translates between unsigned ANSI C **long long** integers and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_u_int(XDR **xdrs***, unsigned** **up***);**

A filter primitive that translates between a C **unsigned** integer and its external representation. This routine returns TRUE if it succeeds, FALSE otherwise.

bool_t xdr_u_long(XDR **xdrs***, unsigned long** **ulp***);**

xdr_u_long( ) translates between C **unsigned long** integers and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_u_longlong_t(XDR **xdrs***, u_longlong_t** **ullp***);**

xdr_u_longlong_t( ) translates between unsigned ANSI C **long long** integers and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise. This routine is identical to **xdr_u_hyper( )**.

bool_t xdr_u_short(XDR **xdrs***, unsigned short** **usp***);**

xdr_u_short( ) translates between C **unsigned short** integers and their external representations. This routine returns **TRUE** if it succeeds, **FALSE** otherwise.

bool_t xdr_void(void);

This routine always returns **TRUE**. It may be passed to RPC routines that require a function parameter, where nothing is to be done.

SEE ALSO

malloc(3C), **rpc**(3N), **xdr_admin**(3N), **xdr_complex**(3N), **xdr_create**(3N)

NAME

xfn – overview of the XFN interface

DESCRIPTION

The primary service provided by a federated naming system is to map a *composite name* to a *reference*. A composite name is composed of name components from one or more naming systems. A reference consists of one or more communication end points. An additional service provided by a federated naming system is to provide access to attributes associated with named objects. This extension is to satisfy most applications' additional naming service needs without cluttering the basic naming service model. XFN is a programming interface for a federated naming service.

To use the XFN interface, include the **xfn/xfn.h** header file and link the application with **-lxfn**.

The **xfn/xfn.h** header file contains the interface declarations for:

- the XFN base context interface,
- the XFN base attribute interface,
- status object and status codes used by operations in these two interfaces,
- abstract data types passed as parameters to and returned as values from operations in these two interfaces, and
- the interface for the XFN standard syntax model for parsing compound names.

FILES

/usr/include/xfn/xfn.h

SEE ALSO

FN_ctx_t(3N), **FN_status_t**(3N), **xfn_attributes**(3N), **xfn_composite_names**(3N), **xfn_compound_names**(3N), **xfn_status_codes**(3N), **fns**(5), **fns_policies**(5)

NAME

xfn_attributes – an overview of XFN attribute operations

DESCRIPTION

XFN assumes the following model for attributes. A set of zero or more attributes is associated with a named object. Each attribute in the set has a unique attribute identifier, an attribute syntax and a (possibly empty) set of distinct data values. Each attribute value has an opaque data type. The attribute identifier serves as a name for the attribute. The attribute syntax indicates the how the value is encoded in the buffer.

The operations of the base attribute interface may be used to examine and modify the settings of attributes associated with existing named objects. These objects may be contexts or other types of objects. The attribute operations do not create names or remove names from contexts.

The range of support for attribute operations may vary widely. Some naming systems may not support any attribute operations. Other naming systems may only support read operations, or operations on attributes whose identifiers are in some fixed set. A naming system may limit attributes to have a single value, or may require at least one value. Some naming systems may only associate attributes with context objects, while others may allow associating attributes with non-context objects.

These are the interfaces:

#include <xfn/xfn.h>

FN_attribute_t ∗**fn_attr_get(FN_ctx_t** ∗*ctx*, **const FN_composite_name_t** ∗*name*,
const FN_identifier_t ∗*attribute_id*, **FN_status_t** ∗*status*);

int fn_attr_modify(FN_ctx_t ∗*ctx*, **const FN_composite_name_t** ∗*name*,
unsigned int *mod_op*, **const FN_attribute_t** ∗*attr*, **FN_status_t** ∗*status*);

FN_attrset_t ∗**fn_attr_get_ids(FN_ctx_t** ∗*ctx*, **const FN_composite_name_t** ∗*name*,
FN_status_t ∗*status*);

FN_valuelist_t ∗**fn_attr_get_values(FN_ctx_t** ∗*ctx*, **const FN_composite_name_t** ∗*name*,
const FN_identifier_t ∗*attribute_id*, **FN_status_t** ∗*status*);

FN_attrvalue_t ∗**fn_valuelist_next(FN_valuelist_t** ∗*vl*, **FN_identifier_t** ∗∗*attr_syntax*,
FN_status_t ∗*status*);

void fn_valuelist_destroy(FN_valuelist_t ∗*vl*, **FN_status_t** ∗*status*);

FN_multigetlist_t ∗**fn_attr_multi_get(FN_ctx_t** ∗*ctx*,
const FN_composite_name_t ∗*name*, **const FN_attrset_t** ∗*attr_ids*,
FN_status_t ∗*status*);

FN_attribute_t ∗**fn_multigetlist_next(FN_multigetlist_t** ∗*ml*, **FN_status_t** ∗*status*);

void fn_multigetlist_destroy(FN_multigetlist_t ∗*ml*, **FN_status_t** ∗*status*);

int fn_attr_multi_modify(FN_ctx_t ∗*ctx*, **const FN_composite_name_t** ∗*name*,
const FN_attrmodlist_t ∗*mods*, **FN_status_t** ∗*status*,
FN_attrmodlist_t ∗∗*unexecuted_mods*);

FN_attrset_t *__fn_ctx_get_syntax_attrs(**FN_ctx_t** *__ctx__,
const FN_composite_name_t *__name__, **FN_status_t** *__status__);

The following describes briefly the operations in the base attribute interface. Detailed descriptions are given in the respective reference manual pages for these operations.

fn_attr_get() returns the attribute identified. **fn_attr_modify()** modifies the attribute identified as described by *mod_op*.

fn_attr_get_ids() returns the identifiers of the attributes of the named object.

fn_attr_get_values() and its set of related operations are used for returning the individual values of an attribute.

fn_attr_multi_get() and its set of related operations are used for returning the requested attributes associated with the named object. **fn_attr_multi_modify()** modifies multiple attributes associated with the named object in a single invocation.

fn_ctx_get_syntax_attrs() returns the syntax attributes associated with the named context.

ERRORS

status is set as described in **FN_status_t**(3N) and **xfn_status_codes**(3N). The following status codes are of special relevance to attribute operations:

FN_E_ATTR_VALUE_REQUIRED
The operation attempted to create an attribute without a value, and the specific naming system does not allow this.

FN_E_ATTR_NO_PERMISSION
The caller did not have permission to perform the attempted attribute operation.

FN_E_INSUFFICIENT_RESOURCES
There is insufficient resources to retrieve the requested attribute(s).

FN_E_INVALID_ATTR_IDENTIFIER
The attribute identifier was not in a format acceptable to the naming system, or its contents was not valid for the format specified for the identifier.

FN_E_INVALID_ATTR_VALUE
One of the values supplied was not in the appropriate form for the given attribute.

FN_E_NO_SUCH_ATTRIBUTE
The object did not have an attribute with the given identifier.

FN_E_TOO_MANY_ATTR_VALUES
The operation attempted to associate more values with an attribute than the naming system supported.

APPLICATION USAGE

Except for **fn_ctx_get_syntax_attrs(),** an attribute operation using a composite name is not necessarily equivalent to an independent **fn_ctx_lookup()** operation followed by an attribute operation in which the caller supplies the resulting reference and an empty name. This is because there are a range of attribute models in which an attribute is associated with a name in a context, or an attribute is associated with the object named, or both. XFN accommodates all of these alternatives. Invoking an attribute operation using the

target context and the terminal atomic name accesses either the attributes that are associated with the target name or target named object — this is dependent on the underlying attribute model. This document uses the term *attributes associated with a named object* to refer to all of these cases.

XFN specifies no guarantees about the relationship between the attributes and the reference associated with a given name. Some naming systems may store the reference bound to a name in one or more attributes associated with a name. Attribute operations might affect the information used to construct a reference.

To avoid undefined results, programmers must use the operations in the context interface and not attribute operations when the intention is to manipulate a reference. Programmers should avoid the use of specific knowledge about how an XFN context implementation over a particular naming system constructs references.

SEE ALSO

FN_attribute_t(3N), **FN_attrvalue_t**(3N), **FN_attrset_t**(3N), **FN_composite_name_t**(3N), **FN_ctx_t**(3N), **FN_identifier_t**(3N), **FN_status_t**(3N), **fn_attr_get**(3N), **fn_attr_get_ids**(3N), **fn_attr_get_values**(3N), **fn_attr_modify**(3N), **fn_attr_multi_get**(3N), **fn_attr_multi_modify**(3N), **fn_ctx_get_syntax_attrs**(3N), **fn_ctx_lookup**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

xfn_composite_names – XFN composite syntax: an overview of the syntax for XFN composite name

DESCRIPTION

An *XFN composite name* consists of an ordered list of zero or more components. Each component is a string name from the namespace of a single naming system. It may be an atomic or a compound name in that namespace.

XFN defines an abstract data type, **FN_composite_name_t**, for representing the structural form of a composite name. XFN also defines a standard string form for composite names. This form is the concatenation of the components of a composite name from left to right with the *XFN component separator* (' / ') character to separate each component.

These are the interfaces:

#include <xfn/xfn.h>

FN_composite_name_t ∗**fn_composite_name_from_string(const FN_string_t** ∗*str*);

FN_string_t ∗**fn_string_from_composite_name(const FN_composite_name_t** ∗*name*);

The function **fn_composite_name_from_string** parses the string representation of a composite name into its corresponding composite name object **FN_composite_name_t**. The function **fn_string_from_composite_name** composes the string representation of a composite name given its composite name object form **FN_composite_name_t**.

The details of the syntax and the semantics of these functions are described in

APPLICATION USAGE

Special characters used in the XFN composite name syntax, such as the separator or escape characters, have the same encoding as they would in **ISO 646**.

All XFN implementations are required to support the portable representation, **ISO 646**. All other representations are optional.

All characters of the string form of a XFN composite name use a single encoding. This does not preclude component names of a composite name in its structural form from having different encodings. Code set mismatches that occur during the process of coverting a composite name structure to its string form are resolved in an implementation-dependent way. When an implementation discovers that a composite name has components with incompatible code sets, it returns the error code **FN_E_INCOMPATIBLE_CODE_SETS.**

SEE ALSO

FN_string_t(3N), **FN_compound_name_t**(3N), **xfn**(3N)

NAME

xfn_compound_names – XFN compound syntax: an overview of XFN model for compound name parsing

DESCRIPTION

Each naming system in an XFN federation has a naming convention. XFN defines a standard model of expressing compound name syntax that covers a large number of specific name syntaxes and is expressed in terms of syntax properties of the naming convention.

The model uses the attributes in the following table to describe properties of the syntax. Unless otherwise qualified, these syntax attributes have attribute identifiers that use the **FN_ID_STRING** format. A context that supports the XFN standard syntax model has an attribute set containing the **fn_syntax_type** (**FN_ID_STRING** format) attribute with the value "standard" (ASCII attribute syntax).

These are the interfaces:

#include <xfn/xfn.h>

FN_attrset_t ∗**fn_ctx_get_syntax_attrs(FN_ctx_t** ∗*ctx*, **const FN_composite_name_t** ∗*name*, **FN_status_t** ∗*status*);

FN_compound_name_t ∗**fn_compound_name_from_syntax_attrs(** **const FN_attrset_t** ∗*aset*, **const FN_string_t** ∗*name*, **FN_status_t** ∗*status*);

fn_syntax_type
: Its value is the ASCII string "standard" if the context supports the XFN standard syntax model. Its value is an implementation-specific value if another syntax model is supported.

fn_std_syntax_direction
: Its value is an ASCII string, one of "left_to_right", "right_to_left", or "flat". This determines whether the order of components in a compound name string goes from left to right, right to left, or whether the namespace is flat (in other words, not hierarchical — all names are atomic).

fn_std_syntax_separator
: Its value is the separator string for this name syntax. This attribute is required unless the **fn_std_syntax_direction** is "flat".

fn_std_syntax_escape
: If present, its value is the escape string for this name syntax.

fn_std_syntax_case_insensitive
: If this attribute is present, it indicates that names that differ only in case are considered identical. If this attribute is absent, it indicates that case is significant. If a value is present, it is ignored.

fn_std_syntax_begin_quote
: If present, its value is the begin-quote string for this syntax. There can be multiple values for this attribute.

fn_std_syntax_end_quote
: If present, its value is the end-quote string for this syntax. There can be multiple values for this attribute.

fn_std_syntax_ava_separator
If present, its value is the attribute value assertion separator string for this syntax.

fn_std_syntax_typeval_separator
If present, its value is the attribute type-value separator string for this syntax.

fn_std_syntax_code_sets
If present, its value identifies the code sets of the string representation for this syntax. Its value consists of a structure containing an array of code sets supported by the context; the first member of the array is the preferred code set of the context. The values for the code sets are defined in the X/Open code set registry. If this attribute is not present, or if the value is empty, the default code set is **ISO 646** (same encoding as ASCII).

fn_std_syntax_locale_info
If present, identifies locale information, such as character set information, of the string representation for this syntax. The interpretation of its value is implementation-dependent.

The XFN standard syntax attributes are interpreted according to the following rules:

1. In a string without quotes or escapes, any instance of the separator string delimits two atomic names.
2. A separator, quotation or escape string is escaped if preceded immediately (on the left) by the escape string.
3. A non-escaped begin-quote which precedes a component must be matched by a non-escaped end-quote at the end of the component. Quotes embedded in non-quoted names are treated as simple characters and do not need to be matched. An unmatched quotation fails with the status code **FN_E_ILLEGAL_NAME.**
4. If there are multiple values for begin-quote and end-quote, a specific begin-quote value must be matched with its corresponding end-quote value.
5. When the separator appears between a (non-escaped) begin quote and the end quote, it is ignored.
6. When the separator is escaped, it is ignored. An escaped begin-quote or end-quote string is not treated as a quotation mark. An escaped escape string is not treated as an escape string.
7. A non-escaped escape string appearing within quotes is interpreted as an escape string. This can be used to embed an end-quote within a quoted string.

After constructing a compound name from a string, the resulting component atoms have one level of escape strings and quotations interpreted and consumed.

fn_ctx_get_syntax_attrs( ) is used to obtain the syntax attributes associated with a context.

fn_compound_name_from_syntax( ) is used to construct a compound name object using the string form of the name and the syntax attributes of the name.

ERRORS

FN_E_ILLEGAL_NAME
The name supplied to the operation was not a well-formed component according to the name syntax of the context.

FN_E_INCOMPATIBLE_CODE_SETS
Code set mismatches that occur during the construction of the compound name's string form are resolved in an implementation-dependent way. When an implementation discovers that a compound name has components with incompatible code sets, it returns the error code **FN_E_INCOMPATIBLE_CODE_SETS.**

FN_E_INVALID_SYNTAX_ATTRS
The syntax attributes supplied are invalid or insufficient to fully specify the syntax.

FN_E_SYNTAX_NOT_SUPPORTED
The syntax specified is not supported.

APPLICATION USAGE

Most applications treat names as opaque data and hence, the majority of clients of the XFN interface will not need to parse compound names from specific naming systems. Some applications, however, such as browsers, need such capabilities. These applications would use **fn_ctx_get_syntax_attrs( )** to obtain the syntax related attributes of a context and, if the context uses the XFN standard syntax model, it would examine these attributes to determine the name syntax of the context.

SEE ALSO

FN_attribute_t(3N), **FN_attrset_t**(3N), **FN_compound_name_t**(3N), **FN_identifier_t**(3N), **FN_string_t**(3N)**, fn_ctx_get_syntax_attrs**(3N), **xfn**(3N)

NAME

xfn_links – XFN links: an overview of XFN links

DESCRIPTION

An *XFN link* is a special form of reference that contains a composite name, the *link name*, and that may be bound to an atomic name in an XFN context. Because the link name is a composite name, it may span multiple namespaces.

Normal resolution of names in context operations always follows XFN links. If the first composite name component of the link name is the atomic name ".", the link name is resolved relative to the same context in which the link is bound, otherwise, the link name is resolved relative to the XFN Initial Context of the client. The link name may itself cause resolution to pass through other XFN links. This gives rise to the possibility of a cycle of links whose resolution could not terminate normally. As a simple means to avoid such non-terminating resolutions, implementations may define limits on the number of XFN links that may be resolved in any single operation invoked by the caller.

These are the interfaces:

#include <xfn/xfn.h>

FN_ref_t ∗**fn_ref_create_link(const FN_composite_name_t** ∗*link_name*);

int fn_ref_is_link(const FN_ref_t ∗*ref*);

FN_composite_name_t ∗**fn_ref_link_name(const FN_ref_t** ∗*link_ref*);

FN_ref_t ∗**fn_ctx_lookup_link(FN_ctx_t** ∗*ctx*, **const FN_composite_name_t** ∗*name*, **FN_status_t** ∗*status*);

unsigned int fn_status_link_code(const FN_status_t ∗*stat*);

const FN_composite_name_t ∗**fn_status_link_remaining_name(const FN_status_t** ∗*stat*);

const FN_composite_name_t ∗**fn_status_link_resolved_name(const FN_status_t** ∗*stat*);

const FN_ref_t ∗**fn_status_link_resolved_ref(const FN_status_t** ∗*stat*);

int fn_status_set_link_code(FN_status_t ∗*stat*, **unsigned int** *code*);

int fn_status_set_link_remaining_name(FN_status_t ∗*stat*, **const FN_composite_name_t** ∗*name*);

int fn_status_set_link_resolved_name(FN_status_t ∗*stat*, **const FN_composite_name_t** ∗*name*);

int fn_status_set_link_resolved_ref(FN_status_t ∗*stat*, **const FN_ref_t** ∗*ref*);

Links are bound to names using the normal **fn_ctx_bind( )** and unbound using the normal **fn_ctx_unbind( )** operation. The operation **fn_ref_create_link( )** is provided for constructing a link reference from a composite name. Since normal resolution always follows links, a separate operation, **fn_ctx_lookup_link( )** is provided to lookup the link itself.

In the case that an error occurred while resolving an XFN link, the status object set by the operation contains additional information about that error and sets the corresponding link status fields using **fn_status_set_link_code( )**, **fn_status_set_link_remaining_name( )**, **fn_status_set_link_resolved_name( )** and

fn_status_set_link_resolved_ref(). The link status fields can be retrieved using **fn_status_link_code()**, **fn_status_link_remaining_name()**, **fn_status_link_resolved_name()** and **fn_status_link_resolved_ref()**.

ERRORS

The following status codes are of special relevance when performing operations involving XFN links:

FN_E_LINK_ERROR
: There was an error encountered resolving an XFN link encountered during resolution of the supplied name. Check the link part of the status object to determine cause of the link error.

FN_E_LINK_LOOP_LIMIT
: A non-terminating loop (cycle) in the resolution can arise due to XFN links encountered during the resolution of a composite name. This code indicates either the definite detection of such a cycle, or that resolution exceeded an implementation-defined limit on the number of XFN links allowed for a single operation invoked by the caller.

FN_E_MALFORMED_LINK
: A malformed link reference was encountered. For the **fn_ctx_lookup_link()** operation, the name supplied resolved to a reference that was not a link.

APPLICATION USAGE

For the **fn_ctx_bind()**, **fn_ctx_unbind()**, **fn_ctx_rename()**, **fn_ctx_lookup_link()**, **fn_ctx_create_subcontext()** and **fn_ctx_destroy_subcontext()** operations, resolution of the given name continues to the target context — that named by all but the terminal atomic part of the given name; the terminal atomic name is not resolved. Consequently, for operations that involve unbinding the terminal atomic part such as **fn_ctx_unbind()**, if the terminal atomic name is bound to a link, the link is not followed and the link itself is unbound from the terminal atomic name.

Many naming systems support a native notion of link that may be used within the naming system itself. XFN does not determine whether there is any relationship between such native links and XFN links.

SEE ALSO

FN_composite_name_t(3N), **FN_ref_t**(3N), **FN_status_t**(3N), **fn_ctx_bind**(3N), **fn_ctx_destroy_subcontext**(3N), **fn_ctx_lookup**(3N), **fn_ctx_lookup_link**(3N), **fn_ctx_rename**(3N), **fn_ctx_unbind**(3N), **xfn_status_codes**(3N), **xfn**(3N)

NAME

xfn_status_codes – descriptions of XFN status codes

DESCRIPTION

The result status of operations in the context interface and the attribute interface is encapsulated in an **FN_status_t** object. This object contains information about how the operation completed: whether an error occurred in performing the operation, if so what kind of error, and information localizing where the error occurred. In the case that the error occurred while resolving an XFN link, the status object contains additional information about that error.

The context status object consists of several items of information. One of them is the primary status code, describing the disposition of the operation. In the case that an error occurred while resolving an XFN link, the primary status code has the value **FN_E_LINK_ERROR,** and link status code describes the error that occurred while resolving the XFN link.

Both the primary status code and the link status code are values of type **unsigned int** that are drawn from the same set of meaningful values. XFN reserves the values **0** through **127** for standard meanings. Currently values and interpretations for the following codes are determined by XFN.

FN_SUCCESS The operation succeeded.

FN_E_ATTR_NO_PERMISSIONS The caller did not have permission to perform the attempted attribute operation.

FN_E_ATTR_VALUE_REQUIRED The operation attempted to create an attribute without a value, and the specific naming system does not allow this.

FN_E_AUTHENTICATION_FAILURE
The identity of the client principal could not be verified.

FN_E_COMMUNICATION_FAILURE
An error occurred in communicating with one of the contexts involved in the operation.

FN_E_CONFIGURATION_ERROR A problem was detected that indicated an error in the installation of the XFN implementation.

FN_E_CONTINUE The operation should be continued using the remaining name and the resolved reference returned in the status.

FN_E_CTX_NO_PERMISSION The client did not have permission to perform the operation.

FN_E_CTX_NOT_EMPTY (Applies only to **fn_ctx_destroy_subcontext( )**.) The naming system required that the context be empty before its destruction, and it was not empty.

FN_E_CTX_UNAVAILABLE Service could not be obtained from one of the contexts involved in the operation. This may be because the naming system is busy, or is not providing service. In

some implementations this may not be distinguished from a communication failure.

FN_E_ILLEGAL_NAME The name supplied to the operation was not a well-formed XFN composite name, or one of the component names was not well-formed according to the syntax of the naming system(s) involved in its resolution.

FN_E_E_INCOMPATIBLE_CODE_SETS
The operation involved character strings of incompatible code sets; or the supplied code set is not supported by the implementation.

FN_E_INSUFFICIENT_RESOURCES
Either the client or one of the involved contexts could not obtain sufficient resources (for example, memory, file descriptors, communication ports, stable media space, and so on.) to complete the operation successfully.

FN_E_INVALID_ATTR_IDENTIFIER
The attribute identifier was not in a format acceptable to the naming system, or its contents was not valid for the format specified for the identifier.

FN_E_INVALID_ATTR_VALUE One of the values supplied was not in the appropriate form for the given attribute.

FN_E_INVALID_ENUM_HANDLE
The enumeration handle supplied was invalid, either because it was from another enumeration, or because an update operation occurred during the enumeration, or because of some other reason.

FN_E_INVALID_SYNTAX_ATTRS The syntax attributes supplied are invalid or insufficient to fully specify the syntax.

FN_E_LINK_ERROR There was an error encountered resolving an XFN link encountered during resolution of the supplied name.

FN_E_LINK_LOOP_LIMIT A non-terminating loop (cycle) in the resolution can arise due to XFN links encountered during the resolution of a composite name. This code indicates either the definite detection of such a cycle, or that resolution exceeded an implementation-defined limit on the number of XFN links allowed for a single operation invoked by the caller.

FN_E_MALFORMED_LINK A malformed link reference was encountered. For **fn_ctx_lookup_link( )**, the name supplied resolved to a reference that was not a link.

FN_E_MALFORMED_REFERENCE

A context object could not be constructed from the supplied reference, because the reference was not properly formed.

FN_E_NAME_IN_USE (Only for operations that bind names.) The supplied name was already in use.

FN_E_NAME_NOT_FOUND Resolution of the supplied composite name proceeded to a context in which the next atomic component of the name was not bound.

FN_E_NO_SUCH_ATTRIBUTE The object did not have an attribute with the given identifier.

FN_E_NO_SUPPORTED_ADDRESS
A context object could not be constructed from a particular reference. The reference contained no address type over which the context interface was supported.

FN_E_NOT_A_CONTEXT Either one of the intermediate atomic names did not name a context, and resolution could not proceed beyond this point, or the operation required that the caller supply the name of a context, and the name did not resolve to a reference for a context.

FN_E_OPERATION_NOT_SUPPORTED
The operation attempted is not supported.

FN_E_PARTIAL_RESULT The operation attempted is returning a partial result.

FN_E_SYNTAX_NOT_SUPPORTED
The syntax type specified is not supported.

FN_E_TOO_MANY_ATTR_VALUES
The operation attempted to associate more values with an attribute than the naming system supported.

FN_E_UNSPECIFIED_ERROR An error occurred that could not be classified by any of the other error codes.

FILES

#include <xfn/xfn.h>

SEE ALSO

FN_status_t(3N), **xfn**(3N)

NAME

yp_update – change NIS information

SYNOPSIS

```
#include <rpcsvc/ypclnt.h>

int yp_update(char *domain, char *map, unsigned ypop, char *key, int keylen,
        char *data, int datalen);
```

MT-LEVEL

Unsafe

DESCRIPTION

yp_update( ) is used to make changes to the NIS database. The syntax is the same as that of **yp_match( )** except for the extra parameter *ypop* which may take on one of four values. If it is **POP_CHANGE** then the data associated with the key will be changed to the new value. If the key is not found in the database, then **yp_update( )** will return **YPERR_KEY**. If *ypop* has the value **YPOP_INSERT** then the key-value pair will be inserted into the database. The error **YPERR_KEY** is returned if the key already exists in the database. To store an item into the database without concern for whether it exists already or not, pass *ypop* as **YPOP_STORE** and no error will be returned if the key already or does not exist. To delete an entry, the value of *ypop* should be **YPOP_DELETE**.

This routine depends upon secure RPC, and will not work unless the network is running secure RPC.

RETURN VALUES

If the value of *ypop* is **POP_CHANGE, yp_update( )** returns the error **YPERR_KEY** if the key is not found in the database.

If the value of *ypop* is **POP_INSERT, yp_update( )** returns the error **YPERR_KEY** if the key already exists in the database.

SEE ALSO

secure_rpc(3N), **ypclnt**(3N)

NOTES

This interface is unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

NAME

ypclnt, yp_get_default_domain, yp_bind, yp_unbind, yp_match, yp_first, yp_next, yp_all, yp_order, yp_master, yperr_string, ypprot_err – NIS Version 2 client interface

SYNOPSIS

cc [*flag* ...] *file* ... –**lnsl** [*library* ...]

#include <rpcsvc/ypclnt.h>
#include <rpcsvc/yp_prot.h>

MT-LEVEL

Unsafe

DESCRIPTION

This package of functions provides an interface to NIS, Network Information Service Version 2, formerly referred to as YP. In this version of SunOS, NIS version 2 is supported only for compatibility with previous versions. The recommended enterprise level information service is NIS+ or NIS version 3, see **nis+**(1). Moreover, this version of SunOS supports only the client interface to NIS version 2. It is expected that this client interface will be served either by an existing **ypserv** process running on another machine on the network that has an earlier version of SunOS or by an NIS+ server, see **rpc.nisd**(1M), running in "YP-compatibility mode". Refer to the **NOTES** section in **ypfiles**(4) for implications of being an NIS client of an NIS+ server in "YP-compatibility mode", and to **ypbind**(1M), **ypwhich**(1), **ypmatch**(1), and **ypcat**(1) for commands to access NIS from a client machine. The package can be loaded from the standard library, **/usr/lib/libnsl.so.1**.

All input parameter names begin with *in*. Output parameters begin with *out*. Output parameters of type **char **** should be addresses of uninitialized character pointers. Memory is allocated by the NIS client package using **malloc**(3C), and may be freed by the user code if it has no continuing need for it. For each *outkey* and *outval*, two extra bytes of memory are allocated at the end that contain NEWLINE and null, respectively, but these two bytes are not reflected in *outkeylen* or *outvallen*. *indomain* and *inmap* strings must be non-null and null-terminated. String parameters which are accompanied by a count parameter may not be null, but may point to null strings, with the count parameter indicating this. Counted strings need not be null-terminated.

All functions in this package of type *int* return **0** if they succeed, and a failure code (**YPERR_***xxxx*) otherwise. Failure codes are described under **ERRORS** below.

Routines

yp_bind (char **indomain***);**

To use the NIS name services, the client process must be “bound” to an NIS server that serves the appropriate domain using **yp_bind( )**. Binding need not be done explicitly by user code; this is done automatically whenever an NIS lookup function is called. **yp_bind( )** can be called directly for processes that make use of a backup strategy (for example, a local file) in cases when NIS services are not available. If a process calls **yp_bind( ),** it should call **yp_unbind( )** when it is done using NIS in order to free up resources.

void yp_unbind(char **indomain***);**

Each binding allocates (uses up) one client process socket descriptor; each bound domain costs one socket descriptor. However, multiple requests to the same domain use that same descriptor. **yp_unbind( )** is available at the client interface for processes that explicitly manage their socket descriptors while accessing multiple domains. The call to **yp_unbind( )** makes the domain *unbound*, and frees all per-process and per-node resources used to bind it.

If an RPC failure results upon use of a binding, that domain will be unbound automatically. At that point, the **ypclnt( )** layer will retry a few more times or until the operation succeeds, provided that **rpcbind**(1M) and **ypbind**(1M) are running, and either

- the client process cannot bind a server for the proper domain, or
- RPC requests to the server fail.

If an error is not RPC-related, or if **rpcbind** is not running, or if **ypbind** is not running, or if a bound **ypserv** process returns any answer (success or failure), the **ypclnt** layer will return control to the user code, either with an error code, or a success code and any results.

yp_get_default_domain (char ***outdomain***);**

The NIS lookup calls require a map name and a domain name, at minimum. It is assumed that the client process knows the name of the map of interest. Client processes should fetch the node's default domain by calling **yp_get_default_domain( )**, and use the returned *outdomain* as the *indomain* parameter to successive NIS name service calls. The domain thus returned is the same as that returned using the **SI_SRPC_DOMAIN** command to the **sysinfo**(2) system call.

yp_match(char **indomain***, char** **inmap***, char** **inkey***, int** *inkeylen***, char** ***outval***,**
int **outvallen***);**

yp_match( ) returns the value associated with a passed key. This key must be exact; no pattern matching is available. **yp_match( )** requires a full YP map name; for example, **hosts.byname** instead of the nickname **hosts**.

yp_first(char **indomain***, char** **inmap***, char** ***outkey***, int** **outkeylen***, char** ***outval***,**
int **outvallen***);**

yp_first( ) returns the first key-value pair from the named map in the named domain.

yp_next(char **indomain***, char** **inmap***, char** **inkey***, int** *inkeylen***, char** ***outkey***,**
int **outkeylen***, char** ***outval***, int** **outvallen***);**

yp_next() returns the next key-value pair in a named map. The *inkey* parameter must be the *outkey* returned from an initial call to **yp_first()** (to get the second key-value pair) or the one returned from the *n*th call to **yp_next()** (to get the *n*th + second key-value pair). Similarly, the *inkeylen* parameter must be the *outkeylen* returned from the earlier **yp_first()** or **yp_next()** call.

The concept of first (and, for that matter, of next) is particular to the structure of the NIS map being processing; there is no relation in retrieval order to either the lexical order within any original (non-NIS name service) data base, or to any obvious numerical sorting order on the keys, values, or key-value pairs. The only ordering guarantee made is that if the **yp_first()** function is called on a particular map, and then the **yp_next()** function is repeatedly called on the same map at the same server until the call fails with a reason of **YPERR_NOMORE**, every entry in the data base will be seen exactly once. Further, if the same sequence of operations is performed on the same map at the same server, the entries will be seen in the same order.

Under conditions of heavy server load or server failure, it is possible for the domain to become unbound, then bound once again (perhaps to a different server) while a client is running. This can cause a break in one of the enumeration rules; specific entries may be seen twice by the client, or not at all. This approach protects the client from error messages that would otherwise be returned in the midst of the enumeration. The next paragraph describes a better solution to enumerating all entries in a map.

yp_all(char ∗*indomain***, char** ∗*inmap***, struct ypall_callback** ∗*incallback***);**

yp_all() provides a way to transfer an entire map from server to client in a single request using TCP (rather than UDP as with other functions in this package). The entire transaction take place as a single RPC request and response. **yp_all()** can be used just like any other NIS name service procedure, identify the map in the normal manner, and supply the name of a function which will be called to process each key-value pair within the map. The call to **yp_all()** returns only when the transaction is completed (successfully or unsuccessfully), or the **foreach()** function decides that it does not want to see any more key-value pairs.

The third parameter to **yp_all()** is

```
struct ypall_callback *incallback {
        int (*foreach)();
        char *data;
};
```

The function **foreach()** is called

```
foreach(int instatus, char *inkey, int inkeylen, char *inval, int invallen,
    char *indata);
```

The *instatus* parameter will hold one of the return status values defined in **<rpcsvc/yp_prot.h** — either **YP_TRUE** or an error code. (See **ypprot_err( )**, below, for a function which converts an NIS name service protocol error code to a **ypclnt** layer error code.)

The key and value parameters are somewhat different than defined in the synopsis section above. First, the memory pointed to by the *inkey* and *inval* parameters is private to the **yp_all( )** function, and is overwritten with the arrival of each new key-value pair. It is the responsibility of the **foreach( )** function to do something useful with the contents of that memory, but it does not own the memory itself. Key and value objects presented to the **foreach( )** function look exactly as they do in the server's map — if they were not NEWLINE-terminated or null-terminated in the map, they will not be here either.

The *indata* parameter is the contents of the *incallback→data* element passed to **yp_all( )**. The **data** element of the callback structure may be used to share state information between the **foreach( )** function and the mainline code. Its use is optional, and no part of the NIS client package inspects its contents — cast it to something useful, or ignore it.

The **foreach( )** function is a Boolean. It should return **0** to indicate that it wants to be called again for further received key-value pairs, or non-zero to stop the flow of key-value pairs. If **foreach( )** returns a non-zero value, it is not called again; the functional value of **yp_all( )** is then **0**.

yp_order(char ∗*indomain***, char** ∗*inmap***, unsigned long** ∗*outorder***);**

yp_order( ) returns the order number for a map. This function is not supported if the **ypbind** process on the client's system is bound to an NIS+ server running in "YP-compatibility mode".

yp_master(char ∗*indomain***, char** ∗*inmap***, char** ∗∗*outname***);**

yp_master( ) returns the machine name of the master NIS server for a map.

char ∗**yperr_string(int** *incode***);**

yperr_string( ) returns a pointer to an error message string that is null-terminated but contains no period or NEWLINE.

ypprot_err (unsigned int *incode***);**

ypprot_err( ) takes an NIS name service protocol error code as input, and returns a ypclnt layer error code, which may be used in turn as an input to **yperr_string( )**.

RETURN VALUES

All integer functions return **0** if the requested operation is successful, or one of the following errors if the operation fails.

YPERR_ACCESS	**15**	**/* access violation */**
YPERR_BADARGS	**1**	**/* args to function are bad */**
YPERR_BADDB	**13**	**/* yp database is bad */**
YPERR_BUSY	**16**	**/* database busy */**
YPERR_DOMAIN	**3**	**/* can't bind to server on this domain */**
YPERR_KEY	**5**	**/* no such key in map */**
YPERR_MAP	**4**	**/* no such map in server's domain */**
YPERR_NODOM	**12**	**/* local domain name not set */**
YPERR_NOMORE	**8**	**/* no more records in map database */**
YPERR_PMAP	**9**	**/* can't communicate with rpcbinder */**
YPERR_RESRC	**7**	**/* resource allocation failure */**
YPERR_RPC	**2**	**/* RPC failure – domain has been unbound */**
YPERR_YPBIND	**10**	**/* can't communicate with ypbind */**
YPERR_YPERR	**6**	**/* internal yp server or client error */**
YPERR_YPSERV	**11**	**/* can't communicate with ypserv */**
YPERR_VERS	**14**	**/* yp version mismatch */**

FILES

/usr/lib/libnsl.so.1

SEE ALSO

nis+(1), **ypcat**(1), **ypmatch**(1), **ypwhich**(1), **rpcbind**(1M), **rpc.nisd**(1M), **ypbind**(1M), **sysinfo**(2), **malloc**(3C), **ypfiles**(4)

NOTES

This interface is unsafe in multithreaded applications. Unsafe interfaces should be called only from the main thread.

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