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Preface

Both novice users and those familiar with the SunOS operating system can use online man pages to obtain information about the system and its features. A man page is intended to answer concisely the question “What does it do?” The man pages in general comprise a reference manual. They are not intended to be a tutorial.

Overview

The following contains a brief description of each man page section and the information it references:

- Section 1 describes, in alphabetical order, commands available with the operating system.
- Section 1M describes, in alphabetical order, commands that are used chiefly for system maintenance and administration purposes.
- Section 2 describes all of the system calls. Most of these calls have one or more error returns. An error condition is indicated by an otherwise impossible returned value.
- Section 3 describes functions found in various libraries, other than those functions that directly invoke UNIX system primitives, which are described in Section 2.
- Section 4 outlines the formats of various files. The C structure declarations for the file formats are given where applicable.
- Section 5 contains miscellaneous documentation such as character-set tables.
- Section 6 contains available games and demos.
- Section 7 describes various special files that refer to specific hardware peripherals and device drivers. STREAMS software drivers, modules and the STREAMS-generic set of system calls are also described.
Section 9 provides reference information needed to write device drivers in the kernel environment. It describes two device driver interface specifications: the Device Driver Interface (DDI) and the Driver/Kernel Interface (DKI).

Section 9E describes the DDI/DKI, DDI-only, and DKI-only entry-point routines a developer can include in a device driver.

Section 9F describes the kernel functions available for use by device drivers.

Section 9S describes the data structures used by drivers to share information between the driver and the kernel.

Below is a generic format for man pages. The man pages of each manual section generally follow this order, but include only needed headings. For example, if there are no bugs to report, there is no BUGS section. See the intro pages for more information and detail about each section, and man(1) for more information about man pages in general.

NAME

This section gives the names of the commands or functions documented, followed by a brief description of what they do.

SYNOPSIS

This section shows the syntax of commands or functions. When a command or file does not exist in the standard path, its full path name is shown. Options and arguments are alphabetized, with single letter arguments first, and options with arguments next, unless a different argument order is required.

The following special characters are used in this section:

[ ] Brackets. The option or argument enclosed in these brackets is optional. If the brackets are omitted, the argument must be specified.

... Ellipses. Several values can be provided for the previous argument, or the previous argument can be specified multiple times, for example, "filename ...".

| Separator. Only one of the arguments separated by this character can be specified at a time.

{ } Braces. The options and/or arguments enclosed within braces are interdependent, such that everything enclosed must be treated as a unit.
PROTOCOL This section occurs only in subsection 3R to indicate the protocol description file.

DESCRIPTION This section defines the functionality and behavior of the service. Thus it describes concisely what the command does. It does not discuss OPTIONS or cite EXAMPLES. Interactive commands, subcommands, requests, macros, and functions are described under USAGE.

IOCTL This section appears on pages in Section 7 only. Only the device class that supplies appropriate parameters to the ioctl(2) system call is called ioctl and generates its own heading. ioctl calls for a specific device are listed alphabetically (on the man page for that specific device). ioctl calls are used for a particular class of devices all of which have an io ending, such as mtcio(7I).

OPTIONS This section lists the command options with a concise summary of what each option does. The options are listed literally and in the order they appear in the SYNOPSIS section. Possible arguments to options are discussed under the option, and where appropriate, default values are supplied.

OPERANDS This section lists the command operands and describes how they affect the actions of the command.

OUTPUT This section describes the output – standard output, standard error, or output files – generated by the command.

RETURN VALUES If the man page documents functions that return values, this section lists these values and describes the conditions under which they are returned. If a function can return only constant values, such as 0 or –1, these values are listed in tagged paragraphs. Otherwise, a single paragraph describes the return values of each function. Functions declared void do not return values, so they are not discussed in RETURN VALUES.

ERRORS On failure, most functions place an error code in the global variable errno indicating why they failed. This section lists alphabetically all error codes a function can generate and describes the conditions that cause each error. When more than
one condition can cause the same error, each condition is described in a separate paragraph under the error code.

**USAGE**

This section lists special rules, features, and commands that require in-depth explanations. The subsections listed here are used to explain built-in functionality:

- Commands
- Modifiers
- Variables
- Expressions
- Input Grammar

**EXAMPLES**

This section provides examples of usage or of how to use a command or function. Wherever possible a complete example including command-line entry and machine response is shown. Whenever an example is given, the prompt is shown as example%, or if the user must be superuser, example#. Examples are followed by explanations, variable substitution rules, or returned values. Most examples illustrate concepts from the SYNOPSIS, DESCRIPTION, OPTIONS, and USAGE sections.

**ENVIRONMENT VARIABLES**

This section lists any environment variables that the command or function affects, followed by a brief description of the effect.

**EXIT STATUS**

This section lists the values the command returns to the calling program or shell and the conditions that cause these values to be returned. Usually, zero is returned for successful completion, and values other than zero for various error conditions.

**FILES**

This section lists all file names referred to by the man page, files of interest, and files created or required by commands. Each is followed by a descriptive summary or explanation.

**ATTRIBUTES**

This section lists characteristics of commands, utilities, and device drivers by defining the attribute type and its corresponding value. See attributes(5) for more information.

**SEE ALSO**

This section lists references to other man pages, in-house documentation, and outside publications.
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<thead>
<tr>
<th>Section</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>DIAGNOSTICS</td>
<td>This section lists diagnostic messages with a brief explanation of the condition causing the error.</td>
</tr>
<tr>
<td>WARNINGS</td>
<td>This section lists warnings about special conditions which could seriously affect your working conditions. This is not a list of diagnostics.</td>
</tr>
<tr>
<td>NOTES</td>
<td>This section lists additional information that does not belong anywhere else on the page. It takes the form of an aside to the user, covering points of special interest. Critical information is never covered here.</td>
</tr>
<tr>
<td>BUGS</td>
<td>This section describes known bugs and, wherever possible, suggests workarounds.</td>
</tr>
</tbody>
</table>
Introduction
This section describes all of the system calls. Most of these calls return one or more error conditions. An error condition is indicated by an otherwise impossible return value. This is almost always \(-1\) or the null pointer; the individual descriptions specify the details. An error number is also made available in the external variable \texttt{errno}, which is not cleared on successful calls, so it should be tested only after an error has been indicated.

In the case of multithreaded applications, the \texttt{−mt} option must be specified on the command line at compilation time (see \texttt{threads(3THR)}). When the \texttt{−mt} option is specified, \texttt{errno} becomes a macro that enables each thread to have its own \texttt{errno}. This \texttt{errno} macro can be used on either side of the assignment as though it were a variable.

Applications should use bound threads rather than the \texttt{_lwp_*()} functions (see \texttt{thr_create(3THR)}). Using LWPs (lightweight processes) directly is not advised because libraries are only safe to use with threads, not LWPs.

Each system call description attempts to list all possible error numbers. The following is a complete list of the error numbers and their names as defined in \texttt{<errno.h>}.

1 EPERM Not superuser
   Typically this error indicates an attempt to modify a file in some way forbidden except to its owner or the super-user. It is also returned for attempts by ordinary users to do things allowed only to the super-user.

2 ENOENT No such file or directory
   A file name is specified and the file should exist but doesn’t, or one of the directories in a path name does not exist.

3 ESRCH No such process, LWP, or thread
   No process can be found in the system that corresponds to the specified \texttt{PID}, \texttt{LWPID_t}, or \texttt{thread_t}.

4 EINTR Interrupted system call
   An asynchronous signal (such as interrupt or quit), which the user has elected to catch, occurred during a system service function. If execution is resumed after processing the signal, it will appear as if the interrupted function call returned this error condition.
In a multithreaded application, **EINTR** may be returned whenever another thread or LWP calls `fork(2)`.

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>5 EIO</td>
<td>I/O error</td>
</tr>
<tr>
<td></td>
<td>Some physical I/O error has occurred. This error may in some cases occur on a call following the one to which it actually applies.</td>
</tr>
<tr>
<td>6 ENXIO</td>
<td>No such device or address</td>
</tr>
<tr>
<td></td>
<td>I/O on a special file refers to a subdevice which does not exist, or exists beyond the limit of the device. It may also occur when, for example, a tape drive is not on-line or no disk pack is loaded on a drive.</td>
</tr>
<tr>
<td>7 E2BIG</td>
<td>Arg list too long</td>
</tr>
<tr>
<td></td>
<td>An argument list longer than <code>ARG_MAX</code> bytes is presented to a member of the <code>exec</code> family of functions (see <code>exec(2)</code>). The argument list limit is the sum of the size of the argument list plus the size of the environment’s exported shell variables.</td>
</tr>
<tr>
<td>8 ENOEXEC</td>
<td>Exec format error</td>
</tr>
<tr>
<td></td>
<td>A request is made to execute a file which, although it has the appropriate permissions, does not start with a valid format (see <code>a.out(4)</code>).</td>
</tr>
<tr>
<td>9 EBADF</td>
<td>Bad file number</td>
</tr>
<tr>
<td></td>
<td>Either a file descriptor refers to no open file, or a <code>read(2)</code> (respectively, <code>write(2)</code>) request is made to a file that is open only for writing (respectively, reading).</td>
</tr>
<tr>
<td>10 ECHILD</td>
<td>No child processes</td>
</tr>
<tr>
<td></td>
<td>A <code>wait(2)</code> function was executed by a process that had no existing or unwaited-for child processes.</td>
</tr>
<tr>
<td>11 EAGAIN</td>
<td>No more processes, or no more LWPs</td>
</tr>
<tr>
<td></td>
<td>For example, the <code>fork(2)</code> function failed because the system’s process table is full or the user is not allowed to create any more processes, or a call failed because of insufficient memory or swap space.</td>
</tr>
<tr>
<td>12 ENOMEM</td>
<td>Not enough space</td>
</tr>
</tbody>
</table>

**Intro(2)**
During execution of `brk()` or `sbrk()` (see `brk(2)`), or one of the `exec` family of functions, a program asks for more space than the system is able to supply. This is not a temporary condition; the maximum size is a system parameter. On some architectures, the error may also occur if the arrangement of text, data, and stack segments requires too many segmentation registers, or if there is not enough swap space during the `fork(2)` function. If this error occurs on a resource associated with Remote File Sharing (RFS), it indicates a memory depletion which may be temporary, dependent on system activity at the time the call was invoked.

13 EACCES Permission denied

An attempt was made to access a file in a way forbidden by the protection system.

14 EFAULT Bad address

The system encountered a hardware fault in attempting to use an argument of a routine. For example, `errno` potentially may be set to `EFAULT` any time a routine that takes a pointer argument is passed an invalid address, if the system can detect the condition. Because systems will differ in their ability to reliably detect a bad address, on some implementations passing a bad address to a routine will result in undefined behavior.

15 ENOTBLK Block device required

A non-block device or file was mentioned where a block device was required (for example, in a call to the `mount(2)` function).

16 EBUSY Device busy

An attempt was made to mount a device that was already mounted or an attempt was made to unmount a device on which there is an active file (open file, current directory, mounted-on file, active text segment). It will also occur if an attempt is made to enable accounting when it is already enabled. The device or resource is currently unavailable. `EBUSY` is also used by mutexes, semaphores, condition variables, and r/w locks, to indicate that a lock is held, and by the processor control function `P_ONLINE`.

17 EEXIST File exists
<table>
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<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>18 EXDEV</td>
<td>Cross-device link</td>
</tr>
<tr>
<td></td>
<td>A hard link to a file on another device was attempted.</td>
</tr>
<tr>
<td>19 ENODEV</td>
<td>No such device</td>
</tr>
<tr>
<td></td>
<td>An attempt was made to apply an inappropriate operation to a device (for example, read a write-only device).</td>
</tr>
<tr>
<td>20 ENOTDIR</td>
<td>Not a directory</td>
</tr>
<tr>
<td></td>
<td>A non-directory was specified where a directory is required (for example, in a path prefix or as an argument to the chdir(2) function).</td>
</tr>
<tr>
<td>21 EISDIR</td>
<td>Is a directory</td>
</tr>
<tr>
<td></td>
<td>An attempt was made to write on a directory.</td>
</tr>
<tr>
<td>22 EINVAL</td>
<td>Invalid argument</td>
</tr>
<tr>
<td></td>
<td>An invalid argument was specified (for example, unmounting a non-mounted device), mentioning an undefined signal in a call to the signal(3C) or kill(2) function, or an unsupported operation related to extended attributes was attempted.</td>
</tr>
<tr>
<td>23 ENFILE</td>
<td>File table overflow</td>
</tr>
<tr>
<td></td>
<td>The system file table is full (that is, SYS_OPEN files are open, and temporarily no more files can be opened).</td>
</tr>
<tr>
<td>24 EMFILE</td>
<td>Too many open files</td>
</tr>
<tr>
<td></td>
<td>No process may have more than OPEN_MAX file descriptors open at a time.</td>
</tr>
<tr>
<td>25 ENOTTY</td>
<td>Inappropriate ioctl for device</td>
</tr>
<tr>
<td></td>
<td>A call was made to the ioctl(2) function specifying a file that is not a special character device.</td>
</tr>
<tr>
<td>26 ETXTBSY</td>
<td>Text file busy (obsolete)</td>
</tr>
<tr>
<td></td>
<td>An attempt was made to execute a pure-procedure program that is currently open for writing. Also an attempt to open for writing or to remove a pure-procedure program that is being executed. <em>(This message is obsolete.)</em></td>
</tr>
<tr>
<td>Error Code</td>
<td>Description</td>
</tr>
<tr>
<td>------------</td>
<td>-------------</td>
</tr>
<tr>
<td>EFBIG</td>
<td>File too large</td>
</tr>
<tr>
<td>ENOSPC</td>
<td>No space left on device</td>
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<tr>
<td>ESPipe</td>
<td>Illegal seek</td>
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<tr>
<td>EROFS</td>
<td>Read-only file system</td>
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<td>EMLINK</td>
<td>Too many links</td>
</tr>
<tr>
<td>EPIPE</td>
<td>Broken pipe</td>
</tr>
<tr>
<td>EDOM</td>
<td>Math argument out of domain of function</td>
</tr>
<tr>
<td>ERANGE</td>
<td>Math result not representable</td>
</tr>
<tr>
<td>ENOMSG</td>
<td>No message of desired type</td>
</tr>
</tbody>
</table>

The size of the file exceeded the limit specified by resource RLIMITFSIZE; the file size exceeds the maximum supported by the file system; or the file size exceeds the offset maximum of the file descriptor. See the File Descriptor subsection of the DEFINITIONS section below.

While writing an ordinary file or creating a directory entry, there is no free space left on the device. In the `fcntl(2)` function, the setting or removing of record locks on a file cannot be accomplished because there are no more record entries left on the system.

A call to the `lseek(2)` function was issued to a pipe.

An attempt to modify a file or directory was made on a device mounted read-only.

An attempt to make more than the maximum number of links, LINK_MAX, to a file.

A write on a pipe for which there is no process to read the data. This condition normally generates a signal; the error is returned if the signal is ignored.

The argument of a function in the math package (3M) is out of the domain of the function.

The value of a function in the math package (3M) is not representable within machine precision.

An attempt was made to receive a message of a type that does not exist on the specified message queue (see `msgget(2)`).
<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>36 EIDRM</td>
<td>Identifier removed</td>
</tr>
<tr>
<td></td>
<td>This error is returned to processes that resume execution due to the removal of an identifier from the file system's name space (see <code>msgctl(2)</code>, <code>semctl(2)</code>, and <code>shmctl(2)</code>).</td>
</tr>
<tr>
<td>37 ECHRNG</td>
<td>Channel number out of range</td>
</tr>
<tr>
<td>38 EL2NSYNC</td>
<td>Level 2 not synchronized</td>
</tr>
<tr>
<td>39 EL3HLT</td>
<td>Level 3 halted</td>
</tr>
<tr>
<td>40 EL3RST</td>
<td>Level 3 reset</td>
</tr>
<tr>
<td>41 ELNRNG</td>
<td>Link number out of range</td>
</tr>
<tr>
<td>42 EUNATCH</td>
<td>Protocol driver not attached</td>
</tr>
<tr>
<td>43 ENOCSI</td>
<td>No CSI structure available</td>
</tr>
<tr>
<td>44 EL2HLT</td>
<td>Level 2 halted</td>
</tr>
<tr>
<td>45 EDEADLK</td>
<td>Deadlock condition</td>
</tr>
<tr>
<td></td>
<td>A deadlock situation was detected and avoided. This error pertains to file and record locking, and also applies to mutexes, semaphores, condition variables, and r/w locks.</td>
</tr>
<tr>
<td>46 ENOLCK</td>
<td>No record locks available</td>
</tr>
<tr>
<td></td>
<td>There are no more locks available. The system lock table is full (see <code>fcntl(2)</code>).</td>
</tr>
<tr>
<td>47 ECANCELED</td>
<td>Operation canceled</td>
</tr>
<tr>
<td></td>
<td>The associated asynchronous operation was canceled before completion.</td>
</tr>
<tr>
<td>48 ENOTSUP</td>
<td>Not supported</td>
</tr>
<tr>
<td></td>
<td>This version of the system does not support this feature. Future versions of the system may provide support.</td>
</tr>
<tr>
<td>49 EDQUOT</td>
<td>Disc quota exceeded</td>
</tr>
<tr>
<td></td>
<td>A <code>write(2)</code> to an ordinary file, the creation of a directory or symbolic link, or the creation of a directory entry failed because the user's quota of disk blocks was exhausted, or the allocation of an inode for a newly created file failed because the user's quota of inodes was exhausted.</td>
</tr>
<tr>
<td>Error Code</td>
<td>Description</td>
</tr>
<tr>
<td>------------</td>
<td>-------------</td>
</tr>
<tr>
<td>58-59</td>
<td>Reserved</td>
</tr>
<tr>
<td>60 ENOSTR</td>
<td>Device not a stream</td>
</tr>
<tr>
<td></td>
<td>A <code>putmsg(2)</code> or <code>getmsg(2)</code> call was attempted on a file descriptor that is not a STREAMS device.</td>
</tr>
<tr>
<td>61 ENODATA</td>
<td>No data available</td>
</tr>
<tr>
<td>62 ETIME</td>
<td>Timer expired</td>
</tr>
<tr>
<td></td>
<td>The timer set for a STREAMS <code>ioctl(2)</code> call has expired. The cause of this error is device-specific and could indicate either a hardware or software failure, or perhaps a timeout value that is too short for the specific operation. The status of the <code>ioctl()</code> operation is indeterminate. This is also returned in the case of <code>_lwp_cond_timedwait(2)</code> or <code>cond_timedwait(3THR)</code>.</td>
</tr>
<tr>
<td>63 ENOSR</td>
<td>Out of stream resources</td>
</tr>
<tr>
<td></td>
<td>During a STREAMS <code>open(2)</code> call, either no STREAMS queues or no STREAMS head data structures were available. This is a temporary condition; one may recover from it if other processes release resources.</td>
</tr>
<tr>
<td>64 ENONET</td>
<td>Machine is not on the network</td>
</tr>
<tr>
<td></td>
<td>This error is Remote File Sharing (RFS) specific. It occurs when users try to advertise, unadvertise, mount, or unmount remote resources while the machine has not done the proper startup to connect to the network.</td>
</tr>
<tr>
<td>65 ENOPKG</td>
<td>Package not installed</td>
</tr>
<tr>
<td></td>
<td>This error occurs when users attempt to use a call from a package which has not been installed.</td>
</tr>
<tr>
<td>66 EREMOTE</td>
<td>Object is remote</td>
</tr>
<tr>
<td></td>
<td>This error is RFS-specific. It occurs when users try to advertise a resource which is not on the local machine, or try to mount/unmount a device (or pathname) that is on a remote machine.</td>
</tr>
<tr>
<td>67 ENOLINK</td>
<td>Link has been severed</td>
</tr>
<tr>
<td></td>
<td>This error is RFS-specific. It occurs when the link (virtual circuit) connecting to a remote machine is gone.</td>
</tr>
<tr>
<td>68 EADV</td>
<td>Advertise error</td>
</tr>
<tr>
<td>Error Code</td>
<td>Description</td>
</tr>
<tr>
<td>------------</td>
<td>-------------</td>
</tr>
<tr>
<td>69 ESRMNT</td>
<td>Srmount error</td>
</tr>
<tr>
<td>70 ECOMM</td>
<td>Communication error on send</td>
</tr>
<tr>
<td>71 EPROTO</td>
<td>Protocol error</td>
</tr>
<tr>
<td>76 EDOTDOT</td>
<td>Error 76</td>
</tr>
<tr>
<td>77 EBADMSG</td>
<td>Not a data message</td>
</tr>
<tr>
<td>78 ENAMETOOLONG</td>
<td>File name too long</td>
</tr>
<tr>
<td>79 EOVERFLOW</td>
<td>Value too large for defined data type</td>
</tr>
</tbody>
</table>

This error is RFS-specific. It occurs when users try to advertise a resource which has been advertised already, or try to stop RFS while there are resources still advertised, or try to force unmount a resource when it is still advertised.

This error is RFS-specific. It occurs when an attempt is made to stop RFS while resources are still mounted by remote machines, or when a resource is readvertised with a client list that does not include a remote machine that currently has the resource mounted.

This error is RFS-specific. It occurs when the current process is waiting for a message from a remote machine, and the virtual circuit fails.

Some protocol error occurred. This error is device-specific, but is generally not related to a hardware failure.

This error is RFS-specific. A way for the server to tell the client that a process has transferred back from mount point.

During a read(2), getmsg(2), or ioctl(2) I_RECVFD call to a STREAMS device, something has come to the head of the queue that cannot be processed. That something depends on the call:

- read(): control information or passed file descriptor.
- getmsg(): passed file descriptor.
- ioctl(): control or data information.

The length of the path argument exceeds PATH_MAX, or the length of a path component exceeds NAME_MAX while _POSIX_NO_TRUNC is in effect; see limits(4).
<table>
<thead>
<tr>
<th>Code</th>
<th>Description</th>
</tr>
</thead>
</table>
| 80 ENOTUNIQ | Name not unique on network  
Given log name not unique. |
| 81 EBADFD  | File descriptor in bad state  
Either a file descriptor refers to no open file or a read request was made to a file that is open only for writing. |
| 82 EREMCHG | Remote address changed |
| 83 ELIBACC | Cannot access a needed share library  
Trying to exec an a.out that requires a static shared library and the static shared library does not exist or the user does not have permission to use it. |
| 84 ELIBBAD | Accessing a corrupted shared library  
Trying to exec an a.out that requires a static shared library (to be linked in) and exec could not load the static shared library. The static shared library is probably corrupted. |
| 85 ELIBSCN | .lib section in a.out corrupted  
Trying to exec an a.out that requires a static shared library (to be linked in) and there was erroneous data in the .lib section of the a.out. The .lib section tells exec what static shared libraries are needed. The a.out is probably corrupted. |
| 86 ELIBMAX | Attempting to link in more shared libraries than system limit  
Trying to exec an a.out that requires more static shared libraries than is allowed on the current configuration of the system. See System Administration Guide: IP Services |
| 87 ELIBEXEC | Cannot exec a shared library directly  
Attempting to exec a shared library directly. |
| 88 EILSEQ  | Error 88  
Illegal byte sequence. Handle multiple characters as a single character. |
<p>| 89 ENOSYS  | Operation not applicable |
| 90 ELOOP   | Number of symbolic links encountered during path name traversal exceeds MAXSYMLINKS |</p>
<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ESTART</td>
<td>Restartable system call. Interrupted system call should be restarted.</td>
</tr>
<tr>
<td>ESTRPIPE</td>
<td>If pipe/FIFO, don’t sleep in stream head. Streams pipe error (not externally visible).</td>
</tr>
<tr>
<td>ENOTEMPTY</td>
<td>Directory not empty.</td>
</tr>
<tr>
<td>EUSERS</td>
<td>Too many users.</td>
</tr>
<tr>
<td>ENOTSOCK</td>
<td>Socket operation on non-socket.</td>
</tr>
<tr>
<td>EDESTADDRREQ</td>
<td>Destination address required. A required address was omitted from an operation on a transport endpoint. Destination address required.</td>
</tr>
<tr>
<td>EMGSIZE</td>
<td>Message too long. A message sent on a transport provider was larger than the internal message buffer or some other network limit.</td>
</tr>
<tr>
<td>EPROTOTYPE</td>
<td>Protocol wrong type for socket. A protocol was specified that does not support the semantics of the socket type requested.</td>
</tr>
<tr>
<td>ENOPROTOOPT</td>
<td>Protocol not available. A bad option or level was specified when getting or setting options for a protocol.</td>
</tr>
<tr>
<td>EPROTONOSUPPORT</td>
<td>Protocol not supported. The protocol has not been configured into the system or no implementation for it exists.</td>
</tr>
<tr>
<td>ESOCKTNOSUPPORT</td>
<td>Socket type not supported. The support for the socket type has not been configured into the system or no implementation for it exists.</td>
</tr>
<tr>
<td>EOPNOTSUPP</td>
<td>Operation not supported on transport endpoint. For example, trying to accept a connection on a datagram transport endpoint.</td>
</tr>
<tr>
<td>EPFNOSUPPORT</td>
<td>Protocol family not supported.</td>
</tr>
</tbody>
</table>
The protocol family has not been configured into the system or no implementation for it exists. Used for the Internet protocols.

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>124 EAFNOSUPPORT</td>
<td>Address family not supported by protocol family</td>
</tr>
<tr>
<td>125 EADDRINUSE</td>
<td>Address already in use</td>
</tr>
<tr>
<td>126 EADDRNOTAVAIL</td>
<td>Cannot assign requested address</td>
</tr>
<tr>
<td>127 ENETDOWN</td>
<td>Network is down</td>
</tr>
<tr>
<td>128 ENETUNREACH</td>
<td>Network is unreachable</td>
</tr>
<tr>
<td>129 ENETRESET</td>
<td>Network dropped connection because of reset</td>
</tr>
<tr>
<td>130 ECONNABORTED</td>
<td>Software caused connection abort</td>
</tr>
<tr>
<td>131 ECONNRESET</td>
<td>Connection reset by peer</td>
</tr>
<tr>
<td>132 ENOBUFS</td>
<td>No buffer space available</td>
</tr>
<tr>
<td>133 EISCONN</td>
<td>Transport endpoint is already connected</td>
</tr>
</tbody>
</table>
A connect request was made on an already connected transport endpoint; or, a sendto(3SOCKET) or sendmsg(3SOCKET) request on a connected transport endpoint specified a destination when already connected.

134 ENOTCONN Transport endpoint is not connected

A request to send or receive data was disallowed because the transport endpoint is not connected and (when sending a datagram) no address was supplied.

143 ESHUTDOWN Cannot send after transport endpoint shutdown

A request to send data was disallowed because the transport endpoint has already been shut down.

144 ETOOMANYREFS Too many references: cannot splice

145 ETIMEDOUT Connection timed out

A connect(3SOCKET) or send(3SOCKET) request failed because the connected party did not properly respond after a period of time; or a write(2) or fsync(3C) request failed because a file is on an NFS file system mounted with the soft option.

146 ECONNREFUSED Connection refused

No connection could be made because the target machine actively refused it. This usually results from trying to connect to a service that is inactive on the remote host.

147 EHOSTDOWN Host is down

A transport provider operation failed because the destination host was down.

148 EHOSTUNREACH No route to host

A transport provider operation was attempted to an unreachable host.

149 EALREADY Operation already in progress

An operation was attempted on a non-blocking object that already had an operation in progress.

150 EINPROGRESS Operation now in progress
An operation that takes a long time to complete (such as a `connect()` was attempted on a non-blocking object.

| 151 ESTALE | Stale NFS file handle |

| **Background Process Group** | Any process group that is not the foreground process group of a session that has established a connection with a controlling terminal. |
| **Controlling Process Controlling Terminal** | A session leader that established a connection to a controlling terminal. |
| **Directory** | A terminal that is associated with a session. Each session may have, at most, one controlling terminal associated with it and a controlling terminal may be associated with only one session. Certain input sequences from the controlling terminal cause signals to be sent to process groups in the session associated with the controlling terminal; see `termio(7)`. |
| **Downstream Driver** | Directories organize files into a hierarchical system where directories are the nodes in the hierarchy. A directory is a file that catalogs the list of files, including directories (sub-directories), that are directly beneath it in the hierarchy. Entries in a directory file are called links. A link associates a file identifier with a filename. By convention, a directory contains at least two links, . (dot) and .. (dot-dot). The link called dot refers to the directory itself while dot-dot refers to its parent directory. The root directory, which is the top-most node of the hierarchy, has itself as its parent directory. The pathname of the root directory is `/` and the parent directory of the root directory is `/`. |
| **Effective User ID and Effective Group ID** | In a stream, the direction from stream head to driver. |
| **Effective User ID and Effective Group ID** | In a stream, the driver provides the interface between peripheral hardware and the stream. A driver can also be a pseudo-driver, such as a multiplexor or log driver (see `log(7D)`), which is not associated with a hardware device. |
| **Effective User ID and Effective Group ID** | An active process has an effective user ID and an effective group ID that are used to determine file access permissions (see below). The effective user ID and effective group ID are equal to the process’s real user ID and real group ID, respectively, unless the process or one of its ancestors evolved from a file that had the set-user-ID bit or set-group-ID bit set (see `exec(2)`). |
| **File Access Permissions** | Read, write, and execute/search permissions for a file are granted to a process if one or more of the following are true: |
| **File Access Permissions** | The effective user ID of the process is super-user. |
| **File Access Permissions** | The effective user ID of the process matches the user ID of the owner of the file and the appropriate access bit of the “owner” portion (0700) of the file mode is set. |
| **File Access Permissions** | The effective user ID of the process does not match the user ID of the owner of the file, but either the effective group ID or one of the supplementary group IDs of the process match the group ID of the file and the appropriate access bit of the “group” portion (0070) of the file mode is set. |
The effective user ID of the process does not match the user ID of the owner of the file, and neither the effective group ID nor any of the supplementary group IDs of the process match the group ID of the file, but the appropriate access bit of the “other” portion (0007) of the file mode is set.

Otherwise, the corresponding permissions are denied.

**File Descriptor**

A file descriptor is a small integer used to perform I/O on a file. The value of a file descriptor is from 0 to (NOFILES-1). A process may have no more than NOFILES file descriptors open simultaneously. A file descriptor is returned by calls such as open(2) or pipe(2). The file descriptor is used as an argument by calls such as read(2), write(2), ioctl(2), and close(2).

Each file descriptor has a corresponding offset maximum. For regular files that were opened without setting the O_LARGEFILE flag, the offset maximum is 2 Gbyte – 1 byte (2^{31} –1 bytes). For regular files that were opened with the O_LARGEFILE flag set, the offset maximum is 2^{63} –1 bytes.

**File Name**

Names consisting of 1 to NAME_MAX characters may be used to name an ordinary file, special file or directory.

These characters may be selected from the set of all character values excluding \0 (null) and the ASCII code for / (slash).

Note that it is generally unwise to use *, ?, [, or ] as part of file names because of the special meaning attached to these characters by the shell (see sh(1), csh(1), and ksh(1)). Although permitted, the use of unprintable characters in file names should be avoided.

A file name is sometimes referred to as a pathname component. The interpretation of a pathname component is dependent on the values of NAME_MAX and _POSIX_NO_TRUNC associated with the path prefix of that component. If any pathname component is longer than NAME_MAX and _POSIX_NO_TRUNC is in effect for the path prefix of that component (see fpathconf(2) and limits(4)), it shall be considered an error condition in that implementation. Otherwise, the implementation shall use the first NAME_MAX bytes of the pathname component.

**Foreground Process Group**

Each session that has established a connection with a controlling terminal will distinguish one process group of the session as the foreground process group of the controlling terminal. This group has certain privileges when accessing its controlling terminal that are denied to background process groups.

**IOV_MAX**

Maximum number of entries in a struct iovec array.

**LIMIT**

The braces notation, {LIMIT}, is used to denote a magnitude limitation imposed by the implementation. This indicates a value which may be defined by a header file (without the braces), or the actual value may be obtained at runtime by a call to the configuration inquiry pathconf(2) with the name argument _PC_LIMIT.
### Masks
The file mode creation mask of the process used during any create function calls to turn off permission bits in the `mode` argument supplied. Bit positions that are set in `umask (cmask)` are cleared in the mode of the created file.

### Message
In a stream, one or more blocks of data or information, with associated STREAMS control structures. Messages can be of several defined types, which identify the message contents. Messages are the only means of transferring data and communicating within a stream.

### Message Queue
In a stream, a linked list of messages awaiting processing by a module or driver.

### Message Queue Identifier
A message queue identifier (`msqid`) is a unique positive integer created by a `msgget(2)` call. Each `msqid` has a message queue and a data structure associated with it. The data structure is referred to as `msqid_ds` and contains the following members:

```c
struct ipc_perm msg_perm;
struct msg *msg_first;
struct msg *msg_last;
ulong_t msg_cbytes;
ulong_t msg_qnum;
ulong_t msg_qbytes;
pid_t msg_lspid;
pid_t msg_lrpid;
time_t msg_stime;
time_t msg_rtime;
time_t msg_ctime;
```

The following are descriptions of the `msqid_ds` structure members:

- **msg_perm** member is an `ipc_perm` structure that specifies the message operation permission (see below). This structure includes the following members:
  ```c
  uid_t cuid; /* creator user id */
gid_t cgid; /* creator group id */
uid_t uid; /* user id */
gid_t gid; /* group id */
mode_t mode; /* r/w permission */
ulong_t seq; /* slot usage sequence # */
key_t key; /* key */
  ```

- The `*msg_first` member is a pointer to the first message on the queue.
- The `*msg_last` member is a pointer to the last message on the queue.
- The `msg_cbytes` member is the current number of bytes on the queue.
- The `msg_qnum` member is the number of messages currently on the queue.
- The `msg_qbytes` member is the maximum number of bytes allowed on the queue.
- The `msg_lspid` member is the process ID of the last process that performed a `msgsnd()` operation.
The msg_lrpid member is the process id of the last process that performed a msgrcv() operation.

The msg_stime member is the time of the last msgsnd() operation.

The msg_rtime member is the time of the last msgrcv() operation.

The msg_ctime member is the time of the last msgctl() operation that changed a member of the above structure.

In the msgctl(2), msgget(2), msgrcv(2), and msgsnd(2) function descriptions, the permission required for an operation is given as {token}, where token is the type of permission needed, interpreted as follows:

<table>
<thead>
<tr>
<th>Token</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>00400</td>
<td>READ by user</td>
</tr>
<tr>
<td>00200</td>
<td>WRITE by user</td>
</tr>
<tr>
<td>00040</td>
<td>READ by group</td>
</tr>
<tr>
<td>00020</td>
<td>WRITE by group</td>
</tr>
<tr>
<td>00004</td>
<td>READ by others</td>
</tr>
<tr>
<td>00002</td>
<td>WRITE by others</td>
</tr>
</tbody>
</table>

Read and write permissions for a msqid are granted to a process if one or more of the following are true:

- The effective user ID of the process is super-user.
- The effective user ID of the process matches msg_perm.cuid or msg_perm.uid in the data structure associated with msqid and the appropriate bit of the “user” portion (0600) of msg_perm.mode is set.
- Any group ID in the process credentials from the set (cr_gid, cr_groups) matches msg_perm.cgid or msg_perm.gid and the appropriate bit of the “group” portion (060) of msg_perm.mode is set.
- The appropriate bit of the “other” portion (006) of msg_perm.mode is set.

Otherwise, the corresponding permissions are denied.

Module

A module is an entity containing processing routines for input and output data. It always exists in the middle of a stream, between the stream’s head and a driver. A module is the STREAMS counterpart to the commands in a shell pipeline except that a module contains a pair of functions which allow independent bidirectional (downstream and upstream) data flow and processing.

Multiplexor

A multiplexor is a driver that allows streams associated with several user processes to be connected to a single driver, or several drivers to be connected to a single user process. STREAMS does not provide a general multiplexing driver, but does provide the facilities for constructing them and for connecting multiplexed configurations of streams.

Offset Maximum

An offset maximum is an attribute of an open file description representing the largest value that can be used as a file offset.
A process group in which the parent of every member in the group is either itself a member of the group, or is not a member of the process group’s session.

A path name is a null-terminated character string starting with an optional slash (/), followed by zero or more directory names separated by slashes, optionally followed by a file name.

If a path name begins with a slash, the path search begins at the root directory. Otherwise, the search begins from the current working directory.

A slash by itself names the root directory.

Unless specifically stated otherwise, the null path name is treated as if it named a non-existent file.

Each process in the system is uniquely identified during its lifetime by a positive integer called a process ID. A process ID may not be reused by the system until the process lifetime, process group lifetime, and session lifetime ends for any process ID, process group ID, and session ID equal to that process ID. Within a process, there are threads with thread id’s, called thread_t and LWID_t. These threads are not visible to the outside process.

A new process is created by a currently active process (see fork(2)). The parent process ID of a process is the process ID of its creator.

Having appropriate privilege means having the capability to override system restrictions.

Each process in the system is a member of a process group that is identified by a process group ID. Any process that is not a process group leader may create a new process group and become its leader. Any process that is not a process group leader may join an existing process group that shares the same session as the process. A newly created process joins the process group of its parent.

A process group leader is a process whose process ID is the same as its process group ID.

Each active process is a member of a process group and is identified by a positive integer called the process group ID. This ID is the process ID of the group leader. This grouping permits the signaling of related processes (see kill(2)).

A process lifetime begins when the process is forked and ends after it exits, when its termination has been acknowledged by its parent process. See wait(2).

A process group lifetime begins when the process group is created by its process group leader, and ends when the lifetime of the last process in the group ends or when the last process in the group leaves the group.
The processors in a system may be divided into subsets, known as processor sets. A process bound to one of these sets will run only on processors in that set, and the processors in the set will normally run only processes that have been bound to the set. Each active processor set is identified by a positive integer. See pset_create(2).

In a stream, the message queue in a module or driver containing messages moving upstream.

Each user allowed on the system is identified by a positive integer (0 to MAXUID) called a real user ID.

Each user is also a member of a group. The group is identified by a positive integer called the real group ID.

An active process has a real user ID and real group ID that are set to the real user ID and real group ID, respectively, of the user responsible for the creation of the process.

Each process has associated with it a concept of a root directory and a current working directory for the purpose of resolving path name searches. The root directory of a process need not be the root directory of the root file system.

Saved resource limits is an attribute of a process that provides some flexibility in the handling of unrepresentable resource limits, as described in the exec family of functions and setrlimit(2).

The saved user ID and saved group ID are the values of the effective user ID and effective group ID just after an exec of a file whose set user or set group file mode bit has been set (see exec(2)).

A semaphore identifier (semid) is a unique positive integer created by a semget(2) call. Each semid has a set of semaphores and a data structure associated with it. The data structure is referred to as semid_ds and contains the following members:

```c
struct ipc_perm sem_perm; /* operation permission struct */
struct sem *sem_base; /* ptr to first semaphore in set */
ushort_t sem_nsems; /* number of sems in set */
time_t sem_otime; /* last operation time */
time_t sem_ctime; /* last change time */
/* Times measured in secs since */
/* 00:00:00 GMT, Jan. 1, 1970 */
```

The following are descriptions of the semid_ds structure members:

The sem_perm member is an ipc_perm structure that specifies the semaphore operation permission (see below). This structure includes the following members:

```c
uid_t uid; /* user id */
gid_t gid; /* group id */
uid_t cuid; /* creator user id */
gid_t cgid; /* creator group id */
mode_t mode; /* r/a permission */
```
ulong_t seq; /* slot usage sequence number */
key_t key; /* key */

The \texttt{sem_nsems} member is equal to the number of semaphores in the set. Each semaphore in the set is referenced by a nonnegative integer referred to as a \texttt{sem_num}. \texttt{sem_num} values run sequentially from 0 to the value of \texttt{sem_nsems} minus 1.

The \texttt{sem_otime} member is the time of the last \texttt{semop}(2) operation.

The \texttt{sem_ctime} member is the time of the last \texttt{semctl}(2) operation that changed a member of the above structure.

A semaphore is a data structure called \texttt{sem} that contains the following members:

\begin{verbatim}
ushort_t semval; /* semaphore value */
pid_t sempid; /* pid of last operation */
ushort_t semncnt; /* # awaiting semval > cval */
ushort_t semzcnt; /* # awaiting semval = 0 */
\end{verbatim}

The following are descriptions of the \texttt{sem} structure members:

The \texttt{semval} member is a non-negative integer that is the actual value of the semaphore.

The \texttt{sempid} member is equal to the process ID of the last process that performed a semaphore operation on this semaphore.

The \texttt{semncnt} member is a count of the number of processes that are currently suspended awaiting this semaphore’s \texttt{semval} to become greater than its current value.

The \texttt{semzcnt} member is a count of the number of processes that are currently suspended awaiting this semaphore’s \texttt{semval} to become 0.

In the \texttt{semop}(2) and \texttt{semctl}(2) function descriptions, the permission required for an operation is given as \{\texttt{token}\}, where \texttt{token} is the type of permission needed interpreted as follows:

\begin{verbatim}
00400 READ by user
00200 ALTER by user
00040 READ by group
00020 ALTER by group
00004 READ by others
00002 ALTER by others
\end{verbatim}

Read and alter permissions for a \texttt{semid} are granted to a process if one or more of the following are true:

- The effective user ID of the process is super-user.

---

\textbf{Semaphore Operation Permissions}

\begin{verbatim}
00400 READ by user
00200 ALTER by user
00040 READ by group
00020 ALTER by group
00004 READ by others
00002 ALTER by others
\end{verbatim}
The effective user ID of the process matches \texttt{sem_perm.cuid} or \texttt{sem_perm.uid} in the data structure associated with \texttt{semid} and the appropriate bit of the “user” portion (0600) of \texttt{sem_perm.mode} is set.

- The effective group ID of the process matches \texttt{sem_perm.cgid} or \texttt{sem_perm.gid} and the appropriate bit of the “group” portion (060) of \texttt{sem_perm.mode} is set.

- The appropriate bit of the “other” portion (06) of \texttt{sem_perm.mode} is set.

Otherwise, the corresponding permissions are denied.

### Session
A session is a group of processes identified by a common ID called a session ID, capable of establishing a connection with a controlling terminal. Any process that is not a process group leader may create a new session and process group, becoming the session leader of the session and process group leader of the process group. A newly created process joins the session of its creator.

### Session ID
Each session in the system is uniquely identified during its lifetime by a positive integer called a session ID, the process ID of its session leader.

### Session Leader
A session leader is a process whose session ID is the same as its process and process group ID.

### Session Lifetime
A session lifetime begins when the session is created by its session leader, and ends when the lifetime of the last process that is a member of the session ends, or when the last process that is a member in the session leaves the session.

### Shared Memory Identifier
A shared memory identifier (\texttt{shmid}) is a unique positive integer created by a \texttt{shmget(2)} call. Each \texttt{shmid} has a segment of memory (referred to as a shared memory segment) and a data structure associated with it. (Note that these shared memory segments must be explicitly removed by the user after the last reference to them is removed.) The data structure is referred to as \texttt{shmid_ds} and contains the following members:

```c
struct ipc_perm  shm_perm; /* operation permission struct */
size_t            shm_segsz; /* size of segment */
struct anon_map   *shm_amp; /* ptr to region structure */
char              pad[4];    /* for swap compatibility */
pid_t             shm_lpid; /* pid of last operation */
pid_t             shm_cpid; /* creator pid */
shmatt_t          shm_nattch; /* number of current attaches */
ulong_t           shm_cnattch; /* used only for shminfo */
time_t            shm_atime; /* last attach time */
time_t            shm_dtime; /* last detach time */
time_t            shm_ctime; /* last change time */
```

*Times measured in secs since* 00:00:00 GMT, Jan. 1, 1970

The following are descriptions of the \texttt{shmid_ds} structure members:

- The \texttt{shm_perm} member is an \texttt{ipc_perm} structure that specifies the shared memory operation permission (see below). This structure includes the following members:
shared memory segment in bytes.

The `shm_cpid` member is the process ID of the process that created the shared memory identifier.

The `shm_lpid` member is the process ID of the last process that performed a `shmat()` or `shmdt()` operation (see `shmop(2)`).

The `shm_nattch` member is the number of processes that currently have this segment attached.

The `shm_atime` member is the time of the last `shmat()` operation (see `shmop(2)`).

The `shm_dtime` member is the time of the last `shmdt()` operation (see `shmop(2)`).

The `shm_ctime` member is the time of the last `shmctl(2)` operation that changed one of the members of the above structure.

In the `shmctl(2), shmat(), and shmdt()` (see `shmop(2)` function descriptions, the permission required for an operation is given as `token`, where `token` is the type of permission needed interpreted as follows:

```
00400    READ by user
00200    WRITE by user
00040    READ by group
00020    WRITE by group
00004    READ by others
00002    WRITE by others
```

Read and write permissions for a `shmid` are granted to a process if one or more of the following are true:

- The effective user ID of the process is super-user.
- The effective user ID of the process matches `shm_perm.cuid` or `shm_perm.uid` in the data structure associated with `shmid` and the appropriate bit of the "user" portion (0600) of `shm_perm.mode` is set.
- The effective group ID of the process matches `shm_perm.cgid` or `shm_perm.gid` and the appropriate bit of the "group" portion (060) of `shm_perm.mode` is set.
- The appropriate bit of the "other" portion (06) of `shm_perm.mode` is set.

Otherwise, the corresponding permissions are denied.
## Special Processes

The process with ID 0 and the process with ID 1 are special processes referred to as proc0 and proc1; see `kill(2)`. proc0 is the process scheduler. proc1 is the initialization process (`init`); proc1 is the ancestor of every other process in the system and is used to control the process structure.

## STREAMS

A set of kernel mechanisms that support the development of network services and data communication drivers. It defines interface standards for character input/output within the kernel and between the kernel and user level processes. The STREAMS mechanism is composed of utility routines, kernel facilities and a set of data structures.

## Stream

A stream is a full-duplex data path within the kernel between a user process and driver routines. The primary components are a stream head, a driver, and zero or more modules between the stream head and driver. A stream is analogous to a shell pipeline, except that data flow and processing are bidirectional.

## Stream Head

In a stream, the stream head is the end of the stream that provides the interface between the stream and a user process. The principal functions of the stream head are processing STREAMS-related system calls and passing data and information between a user process and the stream.

## Super-user

A process is recognized as a super-user process and is granted special privileges, such as immunity from file permissions, if its effective user ID is 0.

## Upstream

In a stream, the direction from driver to stream head.

## Write Queue

In a stream, the message queue in a module or driver containing messages moving downstream.
System Calls
access(2)

NAME
  access – determine accessibility of a file

SYNOPSIS
  #include <unistd.h>
  int access(const char *path, int amode);

DESCRIPTION
  The access() function checks the file named by the pathname pointed to by the path
  argument for accessibility according to the bit pattern contained in amode, using the
  real user ID in place of the effective user ID and the real group ID in place of the
  effective group ID. This allows a setuid process to verify that the user running it
  would have had permission to access this file.

  The value of amode is either the bitwise inclusive OR of the access permissions to be
  checked (R_OK, W_OK, X_OK) or the existence test, F_OK.

  These constants are defined in <unistd.h> as follows:

    R_OK   Test for read permission.
    W_OK   Test for write permission.
    X_OK   Test for execute or search permission.
    F_OK   Check existence of file

  See intro(2) for additional information about "File Access Permission".

  If any access permissions are to be checked, each will be checked individually, as
  described in intro(2). If the process has appropriate privileges, an implementation
  may indicate success for X_OK even if none of the execute file permission bits are set.

RETURN VALUES
  If the requested access is permitted, access() succeeds and returns 0. Otherwise, −1
  is returned and errno is set to indicate the error.

ERRORS
  The access() function will fail if:

    EACCESS  Permission bits of the file mode do not permit the
             requested access, or search permission is denied on a
             component of the path prefix.

    EFAULT   path points to an illegal address.

    EINTR    A signal was caught during the access() function.

    ELOOP    Too many symbolic links were encountered in
             resolving path.

    ENAMETOOLONG The length of the path argument exceeds PATH_MAX, or
             a pathname component is longer than NAME_MAX while
             _POSIX_NO_TRUNC is in effect.

    ENOENT   A component of path does not name an existing file or
             path is an empty string.
ENOLINK: *path* points to a remote machine and the link to that machine is no longer active.

ENOTDIR: A component of the path prefix is not a directory.

EROFS: Write access is requested for a file on a read-only file system.

The *access()* function may fail if:

EINVAL: The value of the *amode* argument is invalid.

ENAMETOOLONG: Pathname resolution of a symbolic link produced an intermediate result whose length exceeds *PATH_MAX*.

ETXTBSY: Write access is requested for a pure procedure (shared text) file that is being executed.

**USAGE**

Additional values of *amode* other than the set defined in the description may be valid, for example, if a system has extended access controls.

**ATTRIBUTES**

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO**

intro(2), chmod(2), stat(2), attributes(5)
acct(2)

NAME acct – enable or disable process accounting

SYNOPSIS #include <unistd.h>

    int acct(const char *path);

DESCRIPTION The acct() function enables or disables the system process accounting routine. If the routine is enabled, an accounting record will be written in an accounting file for each process that terminates. The termination of a process can be caused by either an exit(2) call or a signal(3C). The effective user ID of the process calling acct() must be super-user.

The path argument points to the pathname of the accounting file, whose file format is described on the acct(3HEAD) manual page.

The accounting routine is enabled if path is non-zero and no errors occur during the function. It is disabled if path is (char*)NULL and no errors occur during the function.

RETURN VALUES Upon successful completion, 0 is returned. Otherwise, -1 is returned and errno is set to indicate the error.

ERRORS The acct() function will fail if:

EACCES The file named by path is not an ordinary file.

EBUSY An attempt is being made to enable accounting using the same file that is currently being used.

EFAULT The path argument points to an illegal address.

ELOOP Too many symbolic links were encountered in translating path.

ENAMETOOLONG The length of the path argument exceeds PATH_MAX, or the length of a path argument exceeds NAME_MAX while _POSIX_NO_TRUNC is in effect.

ENOENT One or more components of the accounting file pathname do not exist.

ENOTDIR A component of the path prefix is not a directory.

EPERM The effective user of the calling process is not super-user.

EROFS The named file resides on a read-only file system.

SEE ALSO exit(2), signal(3C), acct(3HEAD)
NAME acl, facl – get or set a file’s Access Control List (ACL)

SYNOPSIS

```c
#include <sys/acl.h>

int acl(char *pathp, int cmd, int nentries, aclent_t *aclbufp);
int facl(int fildes, int cmd, int nentries, aclent_t *aclbufp);
```

DESCRIPTION

The acl() and facl() functions get or set the ACL of a file whose name is given by `pathp` or referenced by the open file descriptor `fildes`. The `nentries` argument specifies how many ACL entries fit into buffer `aclbufp`. The acl() function is used to manipulate ACL on file system objects.

The following values for `cmd` are supported:

- **SETACL**: `nentries` ACL entries, specified in buffer `aclbufp`, are stored in the file’s ACL. All directories in the path name must be searchable.
- **GETACL**: Buffer `aclbufp` is filled with the file’s ACL entries. Read access to the file is not required, but all directories in the path name must be searchable.
- **GETACLCNT**: The number of entries in the file’s ACL is returned. Read access to the file is not required, but all directories in the path name must be searchable.

RETURN VALUES

Upon successful completion, `acl()` and `facl()` return 0 if `cmd` is `SETACL`. If `cmd` is `GETACL` or `GETACLCNT`, the number of ACL entries is returned. Otherwise, -1 is returned and `errno` is set to indicate the error.

ERRORS

The acl() function will fail if:

- **EACCESS**: The caller does not have access to a component of the pathname.
- **EFAULT**: The `pathp` or `aclbufp` argument points to an illegal address.
- **EINVAL**: The `cmd` argument is not `GETACL`, `SETACL`, or `GETACLCNT`; the `cmd` argument is `SETACL` and `nentries` is less than 3; or the `cmd` argument is `SETACL` and the ACL specified in `aclbufp` is not valid.
- **EIO**: A disk I/O error has occurred while storing or retrieving the ACL.
- **ENOENT**: A component of the path does not exist.
- **ENOSPC**: The `cmd` argument is `GETACL` and `nentries` is less than the number of entries in the file’s ACL, or the `cmd` argument is `SETACL` and there is insufficient space in the file system to store the ACL.
- **ENOTDIR**: A component of the path specified by `pathp` is not a directory, or the `cmd` argument is `SETACL` and an attempt is made to set a default ACL on a file type other than a directory.
The `cmd` argument is `SETACL` and the file specified by `pathp` resides on a file system that does not support ACLs, or the `acl()` function is not supported by this implementation.

The effective user ID does not match the owner of the file and the process does not have appropriate privilege.

The `cmd` argument is `SETACL` and the file specified by `pathp` resides on a file system that is mounted read-only.

See `attributes(5)` for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Evolving</td>
</tr>
</tbody>
</table>

`getfacl(1), setfacl(1), aclcheck(3SEC), acls(3SEC)`
include <sys/time.h>

int adjtime(struct timeval *delta, struct timeval *olddelta);

The adjtime() function adjusts the system’s notion of the current time as returned by gettimeofday(3C), advancing or retarding it by the amount of time specified in the struct timeval pointed to by delta. The adjustment is effected by speeding up (if that amount of time is positive) or slowing down (if that amount of time is negative) the system’s clock by some small percentage, generally a fraction of one percent. The time is always a monotonically increasing function. A time correction from an earlier call to adjtime() may not be finished when adjtime() is called again.

If delta is 0, then olddelta returns the status of the effects of the previous adjtime() call with no effect on the time correction as a result of this call. If olddelta is not a null pointer, then the structure it points to will contain, upon successful return, the number of seconds and/or microseconds still to be corrected from the earlier call. If olddelta is a null pointer, the corresponding information will not be returned.

This call may be used in time servers that synchronize the clocks of computers in a local area network. Such time servers would slow down the clocks of some machines and speed up the clocks of others to bring them to the average network time.

Only the super-user may adjust the time of day.

The adjustment value will be silently rounded to the resolution of the system clock.

Upon successful completion, adjtime() returns 0. Otherwise, it returns −1 and sets errno to indicate the error.

The adjtime() function will fail if:

EFAULT The delta or olddelta argument points outside the process’s allocated address space, or olddelta points to a region of the process’s allocated address space that is not writable.

EINVAL The tv_usec member of delta is not within valid range (−1000000 to 1000000).

EPERM The effective user of the calling process is not super-user.

Additionally, the adjtime() function will fail for 32-bit interfaces if:

EOVERFLOW The size of the tv_sec member of the timeval structure pointed to by olddelta is too small to contain the correct number of seconds.

SEE ALSO date(1), gettimeofday(3C)
alarm(2)

NAME  alarm – schedule an alarm signal

SYNOPSIS  #include <unistd.h>

          unsigned int alarm(unsigned int sec);

DESCRIPTION  The alarm() function causes the system to generate a SIGALRM signal for the
          process after the number of real-time seconds specified by seconds have elapsed (see
          signal(3HEAD)). Processor scheduling delays may prevent the process from
          handling the signal as soon as it is generated.

          If seconds is 0, a pending alarm request, if any, is cancelled.

          Alarm requests are not stacked; only one SIGALRM generation can be scheduled in
          this manner; if the SIGALRM signal has not yet been generated, the call will result in
          rescheduling the time at which the SIGALRM signal will be generated.

          The fork(2) function clears pending alarms in the child process. A new process image
          created by one of the exec functions inherits the time left to an alarm signal in the old
          process’s image.

RETURN VALUES  If there is a previous alarm request with time remaining, alarm() returns a non-zero
          value that is the number of seconds until the previous request would have generated a
          SIGALRM signal. Otherwise, alarm() returns 0.

ERRORS  The alarm() function is always successful; no return value is reserved to indicate an
          error.

ATTRIBUTES  See attributes(5) for descriptions of the following attributes:

          ATTRIBUTE TYPE | ATTRIBUTE VALUE
          ---------------|-----------------|
              MT-Level   |     Async-Signal-Safe

SEE ALSO  exec(2), fork(2), signal(3HEAD), attributes(5), standards(5)
The `audit()` function is used to write a record to the system audit log. The data pointed to by `record` is written to the log after a minimal consistency check, with the `length` parameter specifying the size of the record in bytes. The data should be a well-formed audit record as described by `audit.log(4)`.

The kernel validates the record header token type and length, and sets the time stamp value before writing the record to the audit log. The kernel does not do any preselection for user-level generated events. If the audit policy is set to include sequence or trailer tokens, the kernel will append them to the record.

Upon successful completion, 0 is returned. Otherwise, -1 is returned and `errno` is set to indicate the error.

The `audit()` function will fail if:

- `EFAULT` The `record` argument points outside the process’s allocated address space.
- `EINVAL` The record header token ID is invalid or the length is either less than the header token size or greater than `MAXAUDITDATA`.
- `EPERM` The process’s effective user ID is not superuser. Only the superuser can successfully execute this call.

See `attributes(5)` for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Stable</td>
</tr>
<tr>
<td>MT-Level</td>
<td>MT-Safe</td>
</tr>
</tbody>
</table>

The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See `bsmconv(1M)` for more information.
auditon(2)

NAME
auditon – manipulate auditing

SYNOPSIS
cc [ flag ... ] file ... -lbsm -lsoket -lnsl -lint1 [ library ... ]
#include <sys/param.h>
#include <bsm/audit.h>

int auditon(int cmd, caddr_t data, int length);

DESCRIPTION
The auditon() function performs various audit subsystem control operations. The
cmd argument designates the particular audit control command. The data argument is
a pointer to command-specific data. The length argument is the length in bytes of the
command-specific data.

The following commands are supported:

A_GETCOND
Return the system audit on/off/disabled condition in
the integer long pointed to by data. The following
values may be returned:
AUC_AUDITING Auditing has been turned on.
AUC_DISABLED Auditing system has not been
enabled.
AUC_NOAUDIT Auditing has been turned off.
AUC_NOSPACE Auditing has blocked due to lack of
space in audit partition.

A_SETCOND
Set the system’s audit on/off condition to the value in
the integer long pointed to by data. The BSM audit
module must be enabled by bsmconv(1M) before
auditing can be turned on. The following audit states
may be set:
AUC_AUDITING Turns on audit record generation.
AUC_NOAUDIT Turns off audit record generation.

A_GETCLASS
Return the event to class mapping for the designated
audit event. The data argument points to the
au_evclass_map structure containing the event
number. The preselection class mask is returned in the
same structure.

A_SETCLASS
Set the event class preselection mask for the designated
audit event. The data argument points to the
au_evclass_map structure containing the event
number and class mask.

A_GETKMASK
Return the kernel preselection mask in the au_mask
structure pointed to by data. This is the mask used to
preselect non-attributable audit events.
Set the kernel preselection mask. The `data` argument points to the `au_mask` structure containing the class mask. This is the mask used to preselect non-attributable audit events.

Return the audit ID, preselection mask, terminal ID and audit session ID of the specified process in the `auditpinfo` structure pointed to by `data`.

Note that `A_GETPINFO` may fail if the terminal ID contains a network address longer than 32 bits. In this case, the `A_GETPINFO_ADDR` command should be used.

Returns the audit ID, preselection mask, terminal ID and audit session ID of the specified process in the `auditpinfo_addr` structure pointed to by `data`.

Set the preselection mask of the specified process. The `data` argument points to the `auditpinfo` structure containing the process ID and the preselection mask. The other fields of the structure are ignored and should be set to `NULL`.

Set the preselection mask for all processes with the specified audit ID. The `data` argument points to the `auditinfo` structure containing the audit ID and the preselection mask. The other fields of the structure are ignored and should be set to `NULL`.

Set the preselection mask for all processes with the specified audit session ID. The `data` argument points to the `auditinfo` structure containing the audit session ID and the preselection mask. The other fields of the structure are ignored and should be set to `NULL`.

Return the kernel audit queue control parameters. These control the high and low water marks of the number of audit records allowed in the audit queue. The high water mark is the maximum allowed number of undelivered audit records. The low water mark determines when threads blocked on the queue are wakened. Another parameter controls the size of the data buffer used by `auditsvc(2)` to write data to the audit trail. There is also a parameter that specifies a maximum delay before data is attempted to be written to the audit trail. The audit queue parameters are returned in the `au_qctrl` structure pointed to by `data`. 

---

**auditon(2)**

### A_SETKMASK

Set the kernel preselection mask. The `data` argument points to the `au_mask` structure containing the class mask. This is the mask used to preselect non-attributable audit events.

### A_GETPINFO

Return the audit ID, preselection mask, terminal ID and audit session ID of the specified process in the `auditpinfo` structure pointed to by `data`.

Note that `A_GETPINFO` may fail if the terminal ID contains a network address longer than 32 bits. In this case, the `A_GETPINFO_ADDR` command should be used.

### A_GETPINFO_ADDR

Returns the audit ID, preselection mask, terminal ID and audit session ID of the specified process in the `auditpinfo_addr` structure pointed to by `data`.

### A_SETPMASK

Set the preselection mask of the specified process. The `data` argument points to the `auditpinfo` structure containing the process ID and the preselection mask. The other fields of the structure are ignored and should be set to `NULL`.

### A_SETUMASK

Set the preselection mask for all processes with the specified audit ID. The `data` argument points to the `auditinfo` structure containing the audit ID and the preselection mask. The other fields of the structure are ignored and should be set to `NULL`.

### A_SETSMASK

Set the preselection mask for all processes with the specified audit session ID. The `data` argument points to the `auditinfo` structure containing the audit session ID and the preselection mask. The other fields of the structure are ignored and should be set to `NULL`.

### A_GETQCTRL

Return the kernel audit queue control parameters. These control the high and low water marks of the number of audit records allowed in the audit queue. The high water mark is the maximum allowed number of undelivered audit records. The low water mark determines when threads blocked on the queue are wakened. Another parameter controls the size of the data buffer used by `auditsvc(2)` to write data to the audit trail. There is also a parameter that specifies a maximum delay before data is attempted to be written to the audit trail. The audit queue parameters are returned in the `au_qctrl` structure pointed to by `data`. 

---

**System Calls**

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### A_SETQCTRL
Set the kernel audit queue control parameters as described above in the A_GETQCTRL command. The `data` argument points to the `au_qctrl` structure containing the audit queue control parameters. The default and maximum values 'A/B' for the audit queue control parameters are:

- **High water**: 100/10000 (audit records)
- **Low water**: 10/1024 (audit records)
- **Output buffer size**: 1024/1048576 (bytes)
- **Delay**: 20/20000 (hundredths second)

### A_GETCWD
Return the current working directory as kept by the audit subsystem. This path is anchored on the real root, rather than on the active root. The `data` argument points to a buffer into which the path is copied. The `length` argument is the length of the buffer.

### A_GETCAR
Return the current active root as kept by the audit subsystem. This path may be used to anchor an absolute path for a path token generated by an application. The `data` argument points to a buffer into which the path is copied. The `length` argument is the length of the buffer.

### A_GETSTAT
Return the system audit statistics in the `audit_stat` structure pointed to by `data`.

### A_SETSTAT
Reset system audit statistics values. The kernel statistics value is reset if the corresponding field in the statistics structure pointed to by the `data` argument is `CLEAR_VAL`. Otherwise, the value is not changed.

### A_SETFSIZE
Set the maximum size of an audit trail file. When the audit file reaches the designated size, it is closed and a new file started. If the maximum size is unset, the audit trail file generated by `auditsvc()` will grow to the size of the file system. The `data` argument points to the `au_fstat_t` structure containing the maximum audit file size in bytes. The size cannot be set less than 0x80000 bytes.

### A_GETFSIZE
Return the maximum audit file size and current file size in the `au_fstat_t` structure pointed to by the `data` argument.
**A_GETPOLICY**
Return the audit policy flags in the integer long pointed to by `data`.

**A_SETPOLICY**
Set the audit policy flags to the values in the integer long pointed to by `data`. The following policy flags are recognized:

- **AUDIT_CNT**
  Do not suspend processes when audit storage is full or inaccessible. The default action is to suspend processes until storage becomes available.

- **AUDIT_AHLT**
  Halt the machine when a non-attributable audit record can not be delivered. The default action is to count the number of events that could not be recorded.

- **AUDIT_ARGV**
  Include in the audit record the argument list for a member of the `exec(2)` family of functions. The default action is not to include this information.

- **AUDIT_ARGE**
  Include the environment variables for the `execv(2)` function in the audit record. The default action is not to include this information.

- **AUDIT_SEQ**
  Add a `sequence` token to each audit record. The default action is not to include it.

- **AUDIT_TRAIL**
  Append a `trailer` token to each audit record. The default action is not to include it.

- **AUDIT_GROUP**
  Include the supplementary groups list in audit records. The default action is not to include it.

- **AUDIT_PATH**
  Include secondary paths in audit records. Examples of secondary paths are dynamically loaded shared library modules and the command shell path for executable scripts. The default action is to include only the primary path from the system call.
auditon(2)

RETURN VALUES  Upon successful completion, auditon() returns 0. Otherwise, –1 is returned and errno is set to indicate the error.

ERRORS  The auditon() function will fail if:

E2BIG  The length field for the command was too small to hold the returned value.

EFAULT  The copy of data to/from the kernel failed.

EINVAL  One of the arguments was illegal, or BSM has not been installed.

EPERM  The process’s effective user ID is not superuser.

USAGE  The auditon() function can be invoked only by processes with superuser privileges.

ATTRIBUTES  See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Stable</td>
</tr>
<tr>
<td>MT-Level</td>
<td>MT-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO  auditconfig(1M), auditd(1M), bsmconv(1M), audit(2), auditsvc(2), exec(2), audit.log(4), attributes(5)

NOTES  The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See bsmconv(1M) for more information.
auditsvc(2)

NAME
auditsvc – write audit log to specified file descriptor

SYNOPSIS
c{ [ flag ... ] file ... -l bsm -l socket -l nsl -l int1 [ library ... ]
#include <sys/param.h>
#include <bsm/audit.h>

int auditsvc(int fd, int limit);

DESCRIPTION
The auditsvc() function specifies the audit log file to the kernel. The kernel writes audit records to this file until an exceptional condition occurs and then the call returns. The fd argument is a file descriptor that identifies the audit file. Applications should open this file for writing before calling auditsvc().

The limit argument specifies the number of free blocks that must be available in the audit file system, and causes auditsvc() to return when the free disk space on the audit filesystem drops below this limit. Thus, the invoking program can take action to avoid running out of disk space.

The auditsvc() function does not return until one of the following conditions occurs:

- The process receives a signal that is not blocked or ignored.
- An error is encountered writing to the audit log file.
- The minimum free space (as specified by limit), has been reached.

RETURN VALUES
The auditsvc() function returns only on an error.

ERRORS
The auditsvc() function will fail if:

EAGAIN The descriptor referred to a stream, was marked for System V-style non-blocking I/O, and no data could be written immediately.

EBADF The fd argument is not a valid descriptor open for writing.

EBUSY A second process attempted to perform this call.

EFBIG An attempt was made to write a file that exceeds the process’s file size limit or the maximum file size.

EINTR The call is forced to terminate prematurely due to the arrival of a signal whose SV_INTERRUPT bit in sv_flags is set (see sigvec(3UCB)). The signal(3C) function sets this bit for any signal it catches.

EINVAL Auditing is disabled (see auditon(2)), or the fd argument does not refer to a file of an appropriate type (regular files are always appropriate.)

EIO An I/O error occurred while reading from or writing to the file system.
The user’s quota of disk blocks on the file system containing the file has been exhausted; audit filesystem space is below the specified limit; or there is no free space remaining on the file system containing the file.

A hangup occurred on the stream being written to.

The process’s effective user ID is not superuser.

The file was marked for 4.2 BSD-style non-blocking I/O, and no data could be written immediately.

Only processes with an effective user ID of superuser can execute this call successfully.

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Stable</td>
</tr>
<tr>
<td>MT-Level</td>
<td>MT-Safe</td>
</tr>
</tbody>
</table>

The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See bsmconv(1M) for more information.
NAME
brk, sbrk – change the amount of space allocated for the calling process’s data segment

SYNOPSIS
#include <unistd.h>

int brk(void *endds);

void *sbrk(intptr_t incr);

DESCRIPTION
The brk() and sbrk() functions are used to change dynamically the amount of space allocated for the calling process’s data segment (see exec(2)). The change is made by resetting the process’s break value and allocating the appropriate amount of space. The break value is the address of the first location beyond the end of the data segment. The amount of allocated space increases as the break value increases. Newly allocated space is set to zero. If, however, the same memory space is reallocated to the same process its contents are undefined.

When a program begins execution using execve() the break is set at the highest location defined by the program and data storage areas.

The getrlimit(2) function may be used to determine the maximum permissible size of the data segment; it is not possible to set the break beyond the rlim_max value returned from a call to getrlimit(), that is to say, “end + rlim.rlim_max.” See end(3C).

The brk() function sets the break value to endds and changes the allocated space accordingly.

The sbrk() function adds incr function bytes to the break value and changes the allocated space accordingly. The incr function can be negative, in which case the amount of allocated space is decreased.

RETURN VALUES
Upon successful completion, brk() returns 0. Otherwise, it returns -1 and sets errno to indicate the error.

Upon successful completion, sbrk() returns the prior break value. Otherwise, it returns (void *)-1 and sets errno to indicate the error.

ERRORS
The brk() and sbrk() functions will fail and no additional memory will be allocated if:

ENOMEM The data segment size limit as set by setrlimit() (see getrlimit(2)) would be exceeded; the maximum possible size of a data segment (compiled into the system) would be exceeded; insufficient space exists in the swap area to support the expansion; or the new break value would extend into an area of the address space defined by some previously established mapping (see mmap(2)).

EAGAIN Total amount of system memory available for private pages is temporarily insufficient. This may occur even though the space requested was less than the maximum data segment size (see ulimit(2)).
The behavior of `brk()` and `sbrk()` is unspecified if an application also uses any other memory functions (such as `malloc(3C)`, `mmap(2)`, `free(3C)`). The `brk()` and `sbrk()` functions have been used in specialized cases where no other memory allocation function provided the same capability. The use of `mmap(2)` is now preferred because it can be used portably with all other memory allocation functions and with any function that uses other allocation functions.

It is unspecified whether the pointer returned by `sbrk()` is aligned suitably for any purpose.

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>MT-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO `exec(2)`, `getrlimit(2)`, `mmap(2)`, `shmop(2)`, `ulimit(2)`, `end(3C)`, `free(3C)`, `malloc(3C)`

The value of `incr` may be adjusted by the system before setting the new break value. Upon successful completion, the implementation guarantees a minimum of `incr` bytes will be added to the data segment if `incr` is a positive value. If `incr` is a negative value, a maximum of `incr` bytes will be removed from the data segment. This adjustment may not be necessary for all machine architectures.

The value of the arguments to both `brk()` and `sbrk()` are rounded up for alignment with eight-byte boundaries.

Setting the break may fail due to a temporary lack of swap space. It is not possible to distinguish this from a failure caused by exceeding the maximum size of the data segment without consulting `getrlimit()`.
#chdir, fchdir – change working directory

##SYNOPSIS
```
#include <unistd.h>

int chdir(const char *path);
int fchdir(int fildes);
```

##DESCRIPTION
The `chdir()` and `fchdir()` functions cause a directory pointed to by `path` or `fildes` to become the current working directory. The starting point for path searches for path names not beginning with `/` (slash). The `path` argument points to the path name of a directory. The `fildes` argument is an open file descriptor of a directory.

For a directory to become the current directory, a process must have execute (search) access to the directory.

##RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, -1 is returned, the current working directory is unchanged, and `errno` is set to indicate the error.

##ERRORS
The `chdir()` function will fail if:

- **EACCES** Search permission is denied for any component of the path name.
- **EFAULT** The `path` argument points to an illegal address.
- **EINTR** A signal was caught during the execution of the `chdir()` function.
- **EIO** An I/O error occurred while reading from or writing to the file system.
- **ELOOP** Too many symbolic links were encountered in translating `path`.
- **ENAMETOOLONG** The length of the `path` argument exceeds `PATH_MAX`, or the length of a `path` component exceeds `NAME_MAX` while `_POSIX_NO_TRUNC` is in effect.
- **ENOENT** Either a component of the path prefix or the directory named by `path` does not exist or is a null pathname.
- **ENOLINK** The `path` argument points to a remote machine and the link to that machine is no longer active.
- **ENOTDIR** A component of the path name is not a directory.

The `fchdir()` function will fail if:

- **EACCES** Search permission is denied for `fildes`.
- **EBADF** The `fildes` argument is not an open file descriptor.
- **EINTR** A signal was caught during the execution of the `fchdir()` function.
chdir(2)

**EIO**
An I/O error occurred while reading from or writing to the file system.

**ENOLINK**
The *fdes* argument points to a remote machine and the link to that machine is no longer active.

**ENOTDIR**
The open file descriptor *fdes* does not refer to a directory.

**ATTRIBUTES**
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>chdir() is Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO**
chroot(2), attributes(5)
chmod, fchmod – change access permission mode of file

#include <sys/types.h>
#include <sys/stat.h>

int chmod(const char *path, mode_t mode);
int fchmod(int fildes, mode_t mode);

DESCRIPTION
The chmod() and fchmod() functions set the access permission portion of the mode of the file whose name is given by path or referenced by the open file descriptor fildes to the bit pattern contained in mode. Access permission bits are interpreted as follows:

<table>
<thead>
<tr>
<th>Access Permission Bit</th>
<th>Bit Pattern</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>S_ISUID</td>
<td>04000</td>
<td>Set user ID on execution.</td>
</tr>
<tr>
<td>S_ISGID</td>
<td>020#0</td>
<td>Set group ID on execution if # is 7, 5, 3, or 1. Enable mandatory file/record locking if # is 6, 4, 2, or 0.</td>
</tr>
<tr>
<td>S_ISVTX</td>
<td>01000</td>
<td>Save text image after execution.</td>
</tr>
<tr>
<td>S_IRWXU</td>
<td>00700</td>
<td>Read, write, execute by owner.</td>
</tr>
<tr>
<td>S_IRUSR</td>
<td>00400</td>
<td>Read by owner.</td>
</tr>
<tr>
<td>S_IWUSR</td>
<td>00200</td>
<td>Write by owner.</td>
</tr>
<tr>
<td>S_IXUSR</td>
<td>00100</td>
<td>Execute (search if a directory) by owner.</td>
</tr>
<tr>
<td>S_IRWXG</td>
<td>00070</td>
<td>Read, write, execute by group.</td>
</tr>
<tr>
<td>S_IROGRP</td>
<td>00010</td>
<td>Read by group.</td>
</tr>
<tr>
<td>S_IWGRP</td>
<td>00020</td>
<td>Write by group.</td>
</tr>
<tr>
<td>S_IXGRP</td>
<td>00010</td>
<td>Execute by group.</td>
</tr>
<tr>
<td>S_IRWXO</td>
<td>00007</td>
<td>Read, write, execute (search) by others.</td>
</tr>
<tr>
<td>S_IROTH</td>
<td>00004</td>
<td>Read by others.</td>
</tr>
<tr>
<td>S_IWOTH</td>
<td>00002</td>
<td>Write by others.</td>
</tr>
<tr>
<td>S_IXOTH</td>
<td>00001</td>
<td>Execute by others.</td>
</tr>
</tbody>
</table>

Modes are constructed by the bitwise OR operation of the access permission bits.

The effective user ID of the process must match the owner of the file or the process must have the appropriate privilege to change the mode of a file.

If the process is not a privileged process and the file is not a directory, mode bit 01000 (save text image on execution) is cleared.

If neither the process is privileged, nor the file’s group is a member of the process’s supplementary group list, and the effective group ID of the process does not match the group ID of the file, mode bit 02000 (set group ID on execution) is cleared.
If a directory is writable and has S_ISVTX (the sticky bit) set, files within that
directory can be removed or renamed only if one or more of the following is true (see
unlink(2) and rename(2)):

- the user owns the file
- the user owns the directory
- the file is writable by the user
- the user is a privileged user

If a directory has the set group ID bit set, a given file created within that directory will
have the same group ID as the directory, if that group ID is part of the group ID set of
the process that created the file. Otherwise, the newly created file’s group ID will be
set to the effective group ID of the creating process.

If the mode bit 02000 (set group ID on execution) is set and the mode bit 00010
(execute or search by group) is not set, mandatory file/record locking will exist on a
regular file. This may affect future calls to open(2), creat(2), read(2), and write(2)
on this file.

Upon successful completion, chmod() and fchmod() mark for update the st_ctime
field of the file.

**RETURN VALUES**

Upon successful completion, 0 is returned. Otherwise, -1 is returned, the file mode is
unchanged, and errno is set to indicate the error.

**ERRORS**

The chmod() function will fail if:

- **EACCES**
  Search permission is denied on a component of the path prefix of path.

- **EFAULT**
  The path argument points to an illegal address.

- **EINTR**
  A signal was caught during execution of the function.

- **EIO**
  An I/O error occurred while reading from or writing to the file system.

- **ELOOP**
  Too many symbolic links were encountered in translating path.

- **ENAMETOOLONG**
  The length of the path argument exceeds PATH_MAX, or the length of a path component exceeds NAME_MAX
  while _POSIX_NO_TRUNC is in effect.

- **ENOENT**
  Either a component of the path prefix or the file referred to by path does not exist or is a null pathname.

- **ENOLINK**
  The filedes argument points to a remote machine and the link to that machine is no longer active.

- **ENOTDIR**
  A component of the prefix of path is not a directory.

- **EPERM**
  The effective user ID does not match the owner of the file and is not super-user.
The file referred to by *path* resides on a read-only file system.

The `fchmod()` function will fail if:

**EBADF**
- The *fildes* argument is not an open file descriptor

**EIO**
- An I/O error occurred while reading from or writing to the file system.

**EINTR**
- A signal was caught during execution of the `fchmod()` function.

**ENOLINK**
- The *path* argument points to a remote machine and the link to that machine is no longer active.

**EPERM**
- The effective user ID does not match the owner of the file and the process does not have appropriate privilege.

**EROFS**
- The file referred to by *fildes* resides on a read-only file system.

### ATTRIBUTES

See `attributes(5)` for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td><code>chmod()</code> is Async-Signal-Safe</td>
</tr>
</tbody>
</table>

### SEE ALSO

`chmod(1), chown(2), creat(2), fcntl(2), mknod(2), open(2), read(2), rename(2), stat(2), write(2), mkfifo(3C), attributes(5), stat(3HEAD)`

*Programming Interfaces Guide*

### NOTES

If you use `chadm()` to change the file group owner permissions on a file with ACL entries, both the file group owner permissions and the ACL mask are changed to the new permissions. Be aware that the new ACL mask permissions may change the effective permissions for additional users and groups who have ACL entries on the file.
The `chown()` function sets the owner ID and group ID of the file specified by `path` or referenced by the open file descriptor `fd` to `owner` and `group` respectively. If `owner` or `group` is specified as −1, `chown()` does not change the corresponding ID of the file.

The `lchown()` function sets the owner ID and group ID of the named file in the same manner as `chown()`, unless the named file is a symbolic link. In this case, `lchown()` changes the ownership of the symbolic link file itself, while `chown()` changes the ownership of the file or directory to which the symbolic link refers.

The `fchownat()` function sets the owner ID and group ID of the named file in the same manner as `chown()`. If, however, the `path` argument is relative, the path is resolved relative to the `fd` argument rather than the current working directory. If the `fd` argument has the special value `FDCWD`, the path path resolution reverts back to current working directory relative. If the `flag` argument is set to `SYMLNK`, the function behaves like `lchown()` with respect to symbolic links. If the `path` argument is absolute, the `fd` argument is ignored. If the `path` argument is a null pointer, the function behaves like `fchown()`.

If `chown()`, `lchown()`, `fchown()`, or `fchownat()` is invoked by a process other than super-user, the set-user-ID and set-group-ID bits of the file mode, `S_ISUID` and `S_ISGID` respectively, are cleared (see `chmod(2)`).

The operating system provides a configuration option, `_POSIX_CHOWN_RESTRICTED`, to restrict ownership changes for the `chown()`, `lchown()`, and `fchown()` functions. When `_POSIX_CHOWN_RESTRICTED` is not in effect, either the effective user ID of the process must match the owner of the file or the process must be the super-user to change the ownership of a file. When `_POSIX_CHOWN_RESTRICTED` is in effect (the default behavior), the `chown()`, `lchown()`, and `fchown()` functions, for users other than super-user, prevent the owner of the file from changing the owner ID of the file and restrict the change of the group of the file to the list of supplementary group IDs. To set this configuration option, include the following line in `/etc/system`:

```
s set rstchown = 1
```

To disable this option, include the following line in `/etc/system`:

```
s set rstchown = 0
```
set rstchown = 0

See system(4) and fpathconf(2).

Upon successful completion, chown(), fchown() and lchown() mark for update the st_ctime field of the file.

RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, −1 is returned, the owner and group of the named file remain unchanged, and errno is set to indicate the error.

ERRORS
The chown(), lchown(), and fchownat() functions will fail if:

- `EACCES`: Search permission is denied on a component of the path prefix of path.
- `EFAULT`: The path argument points to an illegal address and for fchownat(), the file descriptor has the value AT_FDCWD.
- `EINVAL`: The group or owner argument is out of range.
- `EIO`: An I/O error occurred while reading from or writing to the file system.
- `ENOENT`: Either a component of the path prefix or the file referred to by path does not exist or is a null pathname.
- `EPERM`: The effective user ID does not match the owner of the file or the process is not the super-user and _POSIX_CHOWN_RESTRICTED indicates that such privilege is required.
- `EROS`: The named file resides on a read-only file system.

The fchown() and fchownat() functions will fail if:
chown(2)

EBADF  For `fchown()` the `fildes` argument is not an open file descriptor and.

EBADF  For `fchownat()` the `path` argument is not absolute and the `fildes` argument is not AT_FDCWD or an open file descriptor.

EIO     An I/O error occurred while reading from or writing to the file system.

EINTR   A signal was caught during execution of the function.

ENOLINK The `fildes` argument points to a remote machine and the link to that machine is no longer active.

EINVAL  The `group` or `owner` argument is out of range.

EPERM   The effective user ID does not match the owner of the file, or the process is not the super-user and _POSIX_CHOWN_RESTRICTED indicates that such privilege is required.

EROFS   The named file referred to by `fildes` resides on a read-only file system.

ATTRIBUTES  See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td><code>chown()</code> is Standard; <code>fchownat()</code> is Evolving</td>
</tr>
<tr>
<td>MT-Level</td>
<td><code>chown()</code> and <code>fchownat()</code> are Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO  `chgrp(1)`, `chown(1)`, `chmod(2)`, `fpathconf(2)`, `system(4)`, `attributes(5)`
chroot, fchroot – change root directory

#include <unistd.h>

int chroot(const char *path);
int fchroot(int fildes);

DESCRIPTION

The chroot() and fchroot() functions cause a directory to become the root directory, the starting point for path searches for path names beginning with / (slash). The user’s working directory is unaffected by the chroot() and fchroot() functions.

The path argument points to a path name naming a directory. The fildes argument to fchroot() is the open file descriptor of the directory which is to become the root.

The effective user ID of the process must be super-user to change the root directory. While it is always possible to change to the system root using the fchroot() function, it is not guaranteed to succeed in any other case, even should fildes be valid in all respects.

The “..” entry in the root directory is interpreted to mean the root directory itself. Therefore, “..” cannot be used to access files outside the subtree rooted at the root directory. Instead, fchroot() can be used to reset the root to a directory that was opened before the root directory was changed.

RETURN VALUES

Upon successful completion, 0 is returned. Otherwise, –1 is returned, the root directory remains unchanged, and errno is set to indicate the error.

ERRORS

The chroot() function will fail if:

EACCES Search permission is denied for a component of the path prefix of dirname, or search permission is denied for the directory referred to by dirname.
EBADF The descriptor is not valid.
EFAULT The path argument points to an illegal address.
EINVAL The fchroot() function attempted to change to a directory that is not the system root and external circumstances do not allow this.
EINTR A signal was caught during the execution of the chroot() function.
EIO An I/O error occurred while reading from or writing to the file system.
ELOOP Too many symbolic links were encountered in translating path.
### ENAMETOOLONG
The length of the *path* argument exceeds PATH_MAX, or the length of a *path* component exceeds NAME_MAX while _POSIX_NO_TRUNC is in effect.

### ENOENT
The named directory does not exist or is a null pathname.

### ENOLINK
The *path* argument points to a remote machine and the link to that machine is no longer active.

### ENOTDIR
Any component of the path name is not a directory.

### EPERM
The effective user of the calling process is not super-user.

**SEE ALSO**
- chroot(1M), chdir(2)

**WARNINGS**
The only use of fchroot() that is appropriate is to change back to the system root.
The `close()` function will deallocate the file descriptor indicated by `fdes`. To deallocate means to make the file descriptor available for return by subsequent calls to `open(2)` or other functions that allocate file descriptors. All outstanding record locks owned by the process on the file associated with the file descriptor will be removed (that is, unlocked).

If `close()` is interrupted by a signal that is to be caught, it will return −1 with `errno` set to `EINTR` and the state of `fdes` is unspecified.

When all file descriptors associated with a pipe or FIFO special file are closed, any data remaining in the pipe or FIFO will be discarded.

When all file descriptors associated with an open file description have been closed the open file description will be freed.

If the link count of the file is 0, when all file descriptors associated with the file are closed, the space occupied by the file will be freed and the file will no longer be accessible.

If a STREAMS-based (see `intro(2)`)`fdes` is closed and the calling process was previously registered to receive a `SIGPOLL` signal (see `signal(3C)`) for events associated with that STREAM (see `I_SETSIG in streamio(7I)`), the calling process will be unregistered for events associated with the STREAM. The last `close()` for a STREAM causes the STREAM associated with `fdes` to be dismantled. If `O_NONBLOCK` and `O_NDELAY` are not set and there have been no signals posted for the STREAM, and if there is data on the module's write queue, `close()` waits up to 15 seconds (for each module and driver) for any output to drain before dismantling the STREAM. The time delay can be changed via an `I_SETCLTIME ioctl(2)` request (see `streamio(7I)`). If the `O_NONBLOCK` or `O_NDELAY` flag is set, or if there are any pending signals, `close()` does not wait for output to drain, and dismantles the STREAM immediately.

If `fdes` is associated with one end of a pipe, the last `close()` causes a hangup to occur on the other end of the pipe. In addition, if the other end of the pipe has been named by `fattach(3C)`, then the last `close()` forces the named end to be detached by `fdetach(3C)`. If the named end has no open file descriptors associated with it and gets detached, the STREAM associated with that end is also dismantled.

If `fdes` refers to the master side of a pseudo-terminal, a `SIGHUP` signal is sent to the process group, if any, for which the slave side of the pseudo-terminal is the controlling terminal. It is unspecified whether closing the master side of the pseudo-terminal flushes all queued input and output.
close(2)

If *fildes* refers to the slave side of a STREAMS-based pseudo-terminal, a zero-length message may be sent to the master.

If *fildes* refers to a socket, `close()` causes the socket to be destroyed. If the socket is connection-mode, and the `SOCK_LINGER` option is set for the socket, and the socket has untransmitted data, then `close()` will block for up to the current linger interval until all data is transmitted.

**RETURN VALUES**

Upon successful completion, 0 is returned. Otherwise, −1 is returned and *errno* is set to indicate the error.

**ERRORS**

The `close()` function will fail if:

- **EBADF**       The *fildes* argument is not a valid file descriptor.
- **EINTR**       The `close()` function was interrupted by a signal.
- **ENOLINK**     The *fildes* argument is on a remote machine and the link to that machine is no longer active.
- **ENOSPC**      There was no free space remaining on the device containing the file.

The `close()` function may fail if:

- **EIO**        An I/O error occurred while reading from or writing to the file system.

**USAGE**

An application that used the `stdio` function `fopen(3C)` to open a file should use the corresponding `fclose(3C)` function rather than `close()`.

**ATTRIBUTES**

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO**

`intro(2)`, `creat(2)`, `dup(2)`, `exec(2)`, `fcntl(2)`, `ioctl(2)`, `open(2)`, `pipe(2)`, `fattach(3C)`, `fclose(3C)`, `fdetach(3C)`, `fopen(3C)`, `signal(3C)`, `attributes(5)`, `signal(3HEAD)`, `streamio(7I)`
NAME
creat – create a new file or rewrite an existing one

SYNOPSIS
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>

int creat(const char *path, mode_t mode);

DESCRIPTION
The function call
creat(path, mode)

is equivalent to:
open(path, O_WRONLY | O_CREAT | O_TRUNC, mode)

RETURN VALUES
Refer to open(2).

ERRORS
Refer to open(2).

EXAMPLES
EXAMPLE 1 Creating a File
The following example creates the file /tmp/file with read and write permissions for the file owner and read permission for group and others. The resulting file descriptor is assigned to the fd variable.

#include <fcntl.h>
...
int fd;
mode_t mode = S_IRUSR | S_IWUSR | S_IRGRP | S_IROTH;
char *filename = "/tmp/file";
...
fd = creat(filename, mode);
...

USAGE
The creat() function has a transitional interface for 64-bit file offsets. See lf64(5).

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

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<tr>
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<tbody>
<tr>
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</tr>
</tbody>
</table>

SEE ALSO
open(2), attributes(5), largefile(5), lf64(5)
dup(2)

NAME
dup – duplicate an open file descriptor

SYNOPSIS
#include <unistd.h>

int dup(int fildes);

DESCRIPTION
The dup() function returns a new file descriptor having the following in common
with the original open file descriptor fildes:

- same open file (or pipe)
- same file pointer (that is, both file descriptors share one file pointer)
- same access mode (read, write or read/write).

The new file descriptor is set to remain open across exec functions (see fcntl(2)).

The file descriptor returned is the lowest one available.

The dup(fildes) function call is equivalent to:

fcntl(fildes, F_DUPFD, 0)

RETURN VALUES
Upon successful completion, a non-negative integer representing the file descriptor is
returned. Otherwise, −1 is returned and errno is set to indicate the error.

ERRORS
The dup() function will fail if:

EBADF The fildes argument is not a valid open file descriptor.
EINTR A signal was caught during the execution of the dup() function.
EMFILE The process has too many open files (see getrlimit(2)).
ENOLINK The fildes argument is on a remote machine and the link to that
machine is no longer active.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

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<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
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<tr>
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</table>

SEE ALSO
close(2), creat(2), exec(2), fcntl(2), getrlimit(2), open(2), pipe(2), dup2(3C),
lockf(3C), attributes(5)
exec, execcl, execle, execlp, execv, execve, execvp - execute a file

#include <unistd.h>

int exec(const char *path, const char *arg0, ..., const char *argvn, char */*NULL*/);
int execv(const char *path, char *const argv[]);
int execle(const char *path, const char *arg0, ..., const char *argvn, char */*NULL*/., char *const envp[]);
int execve(const char *path, char *const argv[], char *const envp[]);
int execlp(const char *file, const char *arg0, ..., const char *argvn, char */*NULL*/.);
int execvp(const char *file, char *const argv[]);

DESCRIPTION

Each of the functions in the exec family replaces the current process image with a new process image. The new image is constructed from a regular, executable file called the new process image file. This file is either an executable object file or a file of data for an interpreter. There is no return from a successful call to one of these functions because the calling process image is overlaid by the new process image.

An interpreter file begins with a line of the form

```bash
#! pathname [arg]
```

where `pathname` is the path of the interpreter, and `arg` is an optional argument. When an interpreter file is executed, the system invokes the specified interpreter. The pathname specified in the interpreter file is passed as `argv0` to the interpreter. If `arg` was specified in the interpreter file, it is passed as `argv1` to the interpreter. The remaining arguments to the interpreter are `argv0` through `argvn` of the originally exec’d file. The interpreter named by `pathname` must not be an interpreter file.

When a C-language program is executed as a result of this call, it is entered as a C-language function call as follows:

```c
int main (int argc, char *argv[], char *envp[]);
```

where `argc` is the argument count, `argv` is an array of character pointers to the arguments themselves, and `envp` is an array of character pointers to the environment strings. The `argv` and `environ` arrays are each terminated by a null pointer. The null pointer terminating the `argv` array is not counted in `argc`. The value of `argc` is non-negative, and if greater than 0, `argv[0]` points to a string containing the name of the file. If `argc` is 0, `argv[0]` is a null pointer, in which case there are no arguments. Applications should verify that `argc` is greater than 0 or that `argv[0]` is not a null pointer before dereferencing `argv[0]`. 
The arguments specified by a program with one of the `exec` functions are passed on to the new process image in the `main()` arguments.

The `path` argument points to a path name that identifies the new process image file.

The `file` argument is used to construct a pathname that identifies the new process image file. If the `file` argument contains a slash character, it is used as the pathname for this file. Otherwise, the path prefix for this file is obtained by a search of the directories passed in the `PATH` environment variable (see `environ(5)`). The environment is supplied typically by the shell. If the process image file is not a valid executable object file, `exec1p()` and `execvp()` use the contents of that file as standard input to the shell. In this case, the shell becomes the new process image. In a standard-conforming application (see `standards(5)`), the `exec` family of functions use `/usr/xpg4/bin/sh` (see `ksh(1)`) if `sh` is available; otherwise, they use `/usr/bin/sh` (see `sh(1)`).

The arguments represented by `arg0`... are pointers to null-terminated character strings. These strings constitute the argument list available to the new process image. The list is terminated by a null pointer. The `arg0` argument should point to a filename that is associated with the process being started by one of the `exec` functions.

The `argv` argument is an array of character pointers to null-terminated strings. The last member of this array must be a null pointer. These strings constitute the argument list available to the new process image. The value in `argv[0]` should point to a filename that is associated with the process being started by one of the `exec` functions.

The `envp` argument is an array of character pointers to null-terminated strings. These strings constitute the environment for the new process image. The `envp` array is terminated by a null pointer. For `exec1()`, `execv()`, `execvp()`, and `exec1p()`, the C-language run-time start-off routine places a pointer to the environment of the calling process in the global object `extern char **environ`, and it is used to pass the environment of the calling process to the new process image.

The number of bytes available for the new process’s combined argument and environment lists is `ARG_MAX`. It is implementation-dependent whether null terminators, pointers, and/or any alignment bytes are included in this total.

File descriptors open in the calling process image remain open in the new process image, except for those whose close-on-exec flag `FD_CLOEXEC` is set; (see `fcntl(2)`). For those file descriptors that remain open, all attributes of the open file description, including file locks, remain unchanged.

The preferred hardware address translation size (see `memcntl(2)`) for the stack and heap of the new process image are set to the default system page size.

Directory streams open in the calling process image are closed in the new process image.

The state of conversion descriptors and message catalogue descriptors in the new process image is undefined. For the new process, the equivalent of:
setlocale(LC_ALL, "C") is executed at startup.

Signals set to the default action (SIG_DFL) in the calling process image are set to the default action in the new process image (see signal(3C)). Signals set to be ignored (SIG_IGN) by the calling process image are set to be ignored by the new process image. Signals set to be caught by the calling process image are set to the default action in the new process image (see signal(3HEAD)). After a successful call to any of the exec functions, alternate signal stacks are not preserved and the SA_ONSTACK flag is cleared for all signals.

After a successful call to any of the exec functions, any functions previously registered by atexit(3C) are no longer registered.

The saved resource limits in the new process image are set to be a copy of the process’s corresponding hard and soft resource limits.

If the ST_NOSUID bit is set for the file system containing the new process image file, then the effective user ID and effective group ID are unchanged in the new process image. If the set-user-ID mode bit of the new process image file is set (see chmod(2)), the effective user ID of the new process image is set to the owner ID of the new process image file. Similarly, if the set-group-ID mode bit of the new process image file is set, the effective group ID of the new process image is set to the group ID of the new process image file. The real user ID and real group ID of the new process image remain the same as those of the calling process image. The effective user ID and effective group ID of the new process image are saved (as the saved set-user-ID and the saved set-group-ID for use by setuid(2).

If the effective user-ID is root or superuser, the set-user-ID and set-group-ID bits will be honored when the process is being controlled by ptrace().

Any shared memory segments attached to the calling process image will not be attached to the new process image (see shmat(2)). Any mappings established through mmap() are not preserved across an exec. Memory mappings created in the process are unmapped before the address space is rebuilt for the new process image. See mmap(2).

Memory locks established by the calling process via calls to mlockall(3C) or mlock(3C) are removed. If locked pages in the address space of the calling process are also mapped into the address spaces the locks established by the other processes will be unaffected by the call by this process to the exec function. If the exec function fails, the effect on memory locks is unspecified.

If _XOPEN_REALTIME is defined and has a value other than −1, any named semaphores open in the calling process are closed as if by appropriate calls to sem_close(3RT).

Profiling is disabled for the new process; see profil(2).
Timers created by the calling process with `timer_create(3RT)` are deleted before replacing the current process image with the new process image.

For the `SCHED_FIFO` and `SCHED_RR` scheduling policies, the policy and priority settings are not changed by a call to an `exec` function.

All open message queue descriptors in the calling process are closed, as described in `mq_close(3RT)`.

Any outstanding asynchronous I/O operations may be cancelled. Those asynchronous I/O operations that are not canceled will complete as if the `exec` function had not yet occurred, but any associated signal notifications are suppressed. It is unspecified whether the `exec` function itself blocks awaiting such I/O completion. In no event, however, will the new process image created by the `exec` function be affected by the presence of outstanding asynchronous I/O operations at the time the `exec` function is called.

The new process also inherits the following attributes from the calling process:

- nice value (see `nice(2)`)
- scheduler class and priority (see `priocntl(2)`)
- process ID
- parent process ID
- process group ID
- task ID
- supplementary group IDs
- `semadj` values (see `semop(2)`)
- session membership (see `exit(2)` and `signal(3C)`)
- real user ID
- real group ID
- project ID
- trace flag (see `ptrace(2)` request 0)
- time left until an alarm clock signal (see `alarm(2)`)
- current working directory
- root directory
- file mode creation mask (see `umask(2)`)
- file size limit (see `ulimit(2)`)
- resource limits (see `getrlimit(2)`)
- `tms_utime`, `tms_stime`, `tms_cutime`, and `tms_cstime` (see `times(2)`)
- file-locks (see `fcntl(2)` and `lockf(3C)`)
- controlling terminal
- process signal mask (see `sigprocmask(2)`) and pending signals (see `sigpending(2)`)
- processor bindings (see `processor_bind(2)`)
- processor set bindings (see `pset_bind(2)`)

A call to any `exec` function from a process with more than one thread results in all threads being terminated and the new executable image being loaded and executed. No destructor functions will be called.
Upon successful completion, each of the functions in the exec family marks for update the st_atime field of the file. If an exec function failed but was able to locate the process image file, whether the st_atime field is marked for update is unspecified. Should the function succeed, the process image file is considered to have been opened with open(2). The corresponding close(2) is considered to occur at a time after this open, but before process termination or successful completion of a subsequent call to one of the exec functions. The argv[ ] and envp[ ] arrays of pointers and the strings to which those arrays point will not be modified by a call to one of the exec functions, except as a consequence of replacing the process image.

The saved resource limits in the new process image are set to be a copy of the process’s corresponding hard and soft limits.

RETURN VALUES
If a function in the exec family returns to the calling process image, an error has occurred; the return value is −1 and errno is set to indicate the error.

ERRORS
The exec functions will fail if:

E2BIG
The number of bytes in the new process’s argument list is greater than the system-imposed limit of [ARG_MAX] bytes. The argument list limit is sum of the size of the argument list plus the size of the environment’s exported shell variables.

EACCES
Search permission is denied for a directory listed in the new process file’s path prefix; the new process file is not an ordinary file; or the new process file mode denies execute permission.

EAGAIN
Total amount of system memory available when reading using raw I/O is temporarily insufficient.

EFAULT
An argument points to an illegal address.

EINTR
A signal was caught during the execution of one of the functions in the exec family.

ELOOP
Too many symbolic links were encountered in translating path or file.

ENAMETOOLONG
The length of the file or path argument exceeds [PATH_MAX], or the length of a file or path component exceeds [NAME_MAX] while [POSIX_NO_TRUNC] is in effect.

ENOENT
One or more components of the new process path name of the file do not exist or is a null pathname.

ENOLINK
The path argument points to a remote machine and the link to that machine is no longer active.

ENOTDIR
A component of the new process path of the file prefix is not a directory.
The exec functions, except for execlp() and execvp(), will fail if:

- **ENOEXEC** The new process image file has the appropriate access permission but is not in the proper format.

The exec functions may fail if:

- **ENAMETOOLONG** Pathname resolution of a symbolic link produced an intermediate result whose length exceeds \([\text{PATH_MAX}]\).
- **ENOMEM** The new process image requires more memory than is allowed by the hardware or system-imposed by memory management constraints. See \texttt{brk(2)}.
- **ETXTBSY** The new process image file is a pure procedure (shared text) file that is currently open for writing by some process.

**USAGE**

As the state of conversion descriptors and message catalogue descriptors in the new process image is undefined, portable applications should not rely on their use and should close them prior to calling one of the exec functions.

Applications that require other than the default POSIX locale should call \texttt{setlocale(3C)} with the appropriate parameters to establish the locale of the new process.

The \texttt{environ} array should not be accessed directly by the application.

**ATTRIBUTES**

See \texttt{attributes(5)} for descriptions of the following attributes:

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<thead>
<tr>
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<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Standard</td>
</tr>
<tr>
<td>MT-Level</td>
<td>\texttt{execle()} and \texttt{execve()} are Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO**

\texttt{ksh(1), ps(1), sh(1), alarm(2), brk(2), chmod(2), exit(2), fcntl(2), fork(2), getrlimit(2), memcntl(2), mmap(2), nice(2), priocntl(2), profi1(2), semop(2), shmp(2), sigpending(2), sigprocmask(2), times(2), umask(2), lockf(3C), ptrace(2), setlocale(3C), signal(3C), system(3C), timer_create(3RT), a.out(4), attributes(5), environ(5), standards(5)}

**WARNINGS**

If a program is setuid to a user ID other than the superuser, and the program is executed when the real user ID is super-user, then the program has some of the powers of a super-user as well.
NAME
  exit, _exit – terminate process

SYNOPSIS

#include <stdlib.h>

void exit(int status);

#include <unistd.h>

void _exit(int status);

DESCRIPTION

The exit() function first calls all functions registered by atexit(3C), in the reverse
color of their registration. Each function is called as many times as it was registered.

If a function registered by a call to atexit(3C) fails to return, the remaining registered
functions are not called and the rest of the exit() processing is not completed. If
exit() is called more than once, the effects are undefined.

The exit() function then flushes all output streams, closes all open streams, and
removes all files created by tmpfile(3C).

The _exit() and exit() functions terminate the calling process with the following
consequences:

- All of the file descriptors, directory streams, conversion descriptors and message
catalogue descriptors open in the calling process are closed.
- If the parent process of the calling process is executing a wait(2), wait3(3C),
  waitid(2) or waitpid(2), and has neither set its SA_NOCLEWDATA flag nor set
  SIGCHLD to SIG_IGN, it is notified of the calling process’s termination and the
  low-order eight bits (that is, bits 0377) of status are made available to it. If the
  parent is not waiting, the child’s status will be made available to it when the parent
  subsequently executes wait(2), wait3(3C), waitid(2) or waitpid(2).
- If the parent process of the calling process is not executing a wait(2), wait3(3C),
  waitid(2) or waitpid(2), and has not set its SA_NOCLEWDATA flag, or set SIGCHLD
to SIG_IGN, the calling process is transformed into a zombie process. A
zombie process is an inactive process and it will be deleted at some later time when
its parent process executes wait(2), wait3(3C), waitid(2) or waitpid(2). A
zombie process only occupies a slot in the process table; it has no other space
allocated either in user or kernel space. The process table slot that it occupies is
partially overlaid with time accounting information (see <sys/proc.h>) to be
used by the times(2) function.
- Termination of a process does not directly terminate its children. The sending of a
  SIGHUP signal as described below indirectly terminates children in some
  circumstances.
- A SIGCHLD will be sent to the parent process.
- The parent process ID of all of the calling process’s existing child processes and
  zombie processes is set to 1. That is, these processes are inherited by the
  initialization process (see intro(2)).
- Each mapped memory object is unmapped.
Each attached shared-memory segment is detached and the value of \texttt{shm_nattch} (see \texttt{shmget(2)}) in the data structure associated with its shared memory ID is decremented by 1.

For each semaphore for which the calling process has set a \texttt{semadj} value (see \texttt{semop(2)}), that value is added to the \texttt{semval} of the specified semaphore.

If the process is a controlling process, the \texttt{SIGHUP} signal will be sent to each process in the foreground process group of the controlling terminal belonging to the calling process.

If the process is a controlling process, the controlling terminal associated with the session is disassociated from the session, allowing it to be acquired by a new controlling process.

If the exit of the process causes a process group to become orphaned, and if any member of the newly-orphaned process group is stopped, then a \texttt{SIGHUP} signal followed by a \texttt{SIGCONT} signal will be sent to each process in the newly-orphaned process group.

If the parent process has set its \texttt{SA_NOCLDWAIT} flag, or set \texttt{SIGCHLD} to \texttt{SIG_IGN}, the status will be discarded, and the lifetime of the calling process will end immediately.

If the process has process, text or data locks, an \texttt{UNLOCK} is performed (see \texttt{plock(3C)} and \texttt{memcntl(2)}).

All open named semaphores in the process are closed as if by appropriate calls to \texttt{sem_close(3RT)}. All open message queues in the process are closed as if by appropriate calls to \texttt{mq_close(3RT)}. Any outstanding asynchronous I/O operations may be cancelled.

An accounting record is written on the accounting file if the system’s accounting routine is enabled (see \texttt{acct(2)}).

An extended accounting record is written to the extended process accounting file if the system’s extended process accounting facility is enabled (see \texttt{acctadm(1M)}).

If the current process is the last process within its task and if the system’s extended task accounting facility is enabled (see \texttt{acctadm(1M)}), an extended accounting record is written to the extended task accounting file.

These functions do not return.

No errors are defined.

Normally applications should use \texttt{exit()} rather than \texttt{_exit()}.

See \texttt{attributes(5)} for descriptions of the following attributes:

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<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>\texttt{_exit()} is Async-Signal Safe</td>
</tr>
</tbody>
</table>
SEE ALSO
acctadm(1M), intro(2), acct(2), close(2), memcntl(2), semop(2), shmget(2),
sigaction(2), times(2), wait(2), waitid(2), waitpid(2), atexit(3C),
fclose(3C), mq_close(3RT), plock(3C), signal(3HEAD), tmpfile(3C),
wait3(3C), attributes(5)
The `fcntl()` function provides for control over open files. The `fd` argument is an open file descriptor.

The `fcntl()` function may take a third argument, `arg`, whose data type, value and use depend upon the value of `cmd`. The `cmd` argument specifies the operation to be performed by `fcntl()`.

The available values for `cmd` are defined in the header `<fcntl.h>`, which include:

- **F_DUPFD**
  Return a new file descriptor which is the lowest numbered available (that is, not already open) file descriptor greater than or equal to the third argument, `arg`, taken as an integer of type `int`. The new file descriptor refers to the same open file description as the original file descriptor, and shares any locks. The `FD_CLOEXEC` flag associated with the new file descriptor is cleared to keep the file open across calls to one of the `exec(2)` functions.

- **F_DUP2FD**
  Similar to `F_DUPFD`, but always returns `arg`. `F_DUP2FD` closes `arg` if it is open and not equal to `fd`. `F_DUP2FD` is equivalent to `dup2(fd, arg)`.

- **F_FREESP**
  Free storage space associated with a section of the ordinary file `fd`. The section is specified by a variable of data type `struct flock` pointed to by `arg`. The data type `struct flock` is defined in the `<fcntl.h>` header (see `fcntl(3HEAD)`) and is described below. Note that all file systems might not support all possible variations of `F_FREESP` arguments. In particular, many file systems allow space to be freed only at the end of a file.

- **F_GETFD**
  Get the file descriptor flags defined in `<fcntl.h>` that are associated with the file descriptor `fd`. File descriptor flags are associated with a single file descriptor and do not affect other file descriptors that refer to the same file.

- **F_GETFL**
  Get the file status flags and file access modes, defined in `<fcntl.h>`, for the file descriptor specified by `fd`. The file access modes can be extracted from the return value using the mask `O_ACCMODE`, which is defined in `<fcntl.h>`. File status flags and file access modes do not affect other file descriptors that refer to the same file with different open file descriptions.

- **F_GETOWN**
  If `fd` refers to a socket, get the process or process group ID specified to receive SIGURG signals when out-of-band data is available. Positive values indicate a process ID; negative values,
other than −1, indicate a process group ID. If *fildes* does not refer to a socket, the results are unspecified.

**F_GETXFL**  
Get the file status flags, file access modes, and file creation and assignment flags, defined in `<fcntl.h>`, for the file descriptor specified by *fildes*. The file access modes can be extracted from the return value using the mask `O_ACCMODE`, which is defined in `<fcntl.h>`. File status flags, file access modes, and file creation and assignment flags do not affect other file descriptors that refer to the same file with different open file descriptions.

**F_SETFD**  
Set the file descriptor flags defined in `<fcntl.h>`, that are associated with *fildes*, to the third argument, *arg*, taken as type `int`. If the `FD_CLOEXEC` flag in the third argument is 0, the file will remain open across the `exec()` functions; otherwise the file will be closed upon successful execution of one of the `exec()` functions.

**F_SETFL**  
Set the file status flags, defined in `<fcntl.h>`, for the file descriptor specified by *fildes* from the corresponding bits in the *arg* argument, taken as type `int`. Bits corresponding to the file access mode and file creation and assignment flags that are set in *arg* are ignored. If any bits in *arg* other than those mentioned here are changed by the application, the result is unspecified.

**F_SETOWN**  
If *fildes* refers to a socket, set the process or process group ID specified to receive SIGURG signals when out-of-band data is available, using the value of the third argument, *arg*, taken as type `int`. Positive values indicate a process ID; negative values, other than −1, indicate a process group ID. If *fildes* does not refer to a socket, the results are unspecified.

The following commands are available for advisory record locking. Record locking is supported for regular files, and may be supported for other files.

**F_GETLK**  
Get the first lock which blocks the lock description pointed to by the third argument, *arg*, taken as a pointer to type `struct flock`, defined in `<fcntl.h>`. The information retrieved overwrites the information passed to `fcntl()` in the structure `flock`. If no lock is found that would prevent this lock from being created, then the structure will be left unchanged except for the lock type which will be set to `F_UNLCK`.

**F_GETLK64**  
Equivalent to `F_GETLK`, but takes a `struct flock64` argument rather than a `struct flock` argument.

**F_SETLK**  
Set or clear a file segment lock according to the lock description pointed to by the third argument, *arg*, taken as a pointer to type `struct flock`, defined in `<fcntl.h>`. `F_SETLK` is used to establish shared (or read) locks (`F_RDLCK`) or exclusive (or write)
locks (F_WRLCK), as well as to remove either type of lock
(F_RDLCK, F_WRLCK and F_UNLCK) are defined in
<fcntl.h>. If a shared or exclusive lock cannot be set, fcntl() will return immediately with a return value of -1.

F_SETLK64 Equivalent to F_SETLK, but takes a struct flock64 argument rather than a struct flock argument.

F_SETLKW This command is the same as F_SETLK except that if a shared or exclusive lock is blocked by other locks, the process will wait until the request can be satisfied. If a signal that is to be caught is received while fcntl() is waiting for a region, fcntl() will be interrupted. Upon return from the process' signal handler, fcntl() will return -1 with errno set to EINTR, and the lock operation will not be done.

F_SETLKW64 Equivalent to F_SETLKW, but takes a struct flock64 argument rather than a struct flock argument.

When a shared lock is set on a segment of a file, other processes will be able to set shared locks on that segment or a portion of it. A shared lock prevents any other process from setting an exclusive lock on any portion of the protected area. A request for a shared lock will fail if the file descriptor was not opened with read access.

An exclusive lock will prevent any other process from setting a shared lock or an exclusive lock on any portion of the protected area. A request for an exclusive lock will fail if the file descriptor was not opened with write access.

The flock structure contains at least the following elements:

short l_type; /* lock operation type */
short l_whence; /* lock base indicator */
off_t l_start; /* starting offset from base */
off_t l_len; /* lock length; l_len == 0 means until end of file */
int l_sysid; /* system ID running process holding lock */
pid_t l_pid; /* process ID of process holding lock */

The value of l_whence is SEEK_SET, SEEK_CUR, or SEEK_END, to indicate that the relative offset l_start bytes will be measured from the start of the file, current position or end of the file, respectively. The value of l_len is the number of consecutive bytes to be locked. The value of l_len may be negative (where the definition of off_t permits negative values of l_len). After a successful F_GETLK or F_GETLK64 request, that is, one in which a lock was found, the value of l_whence will be SEEK_SET.

The l_pid and l_sysid fields are used only with F_GETLK or F_GETLK64 to return the process ID of the process holding a blocking lock and to indicate which system is running that process.
If `l_len` is positive, the area affected starts at `l_start` and ends at `l_start + l_len - 1`. If `l_len` is negative, the area affected starts at `l_start + l_len` and ends at `l_start - 1`. Locks may start and extend beyond the current end of a file, but must not be negative relative to the beginning of the file. A lock will be set to extend to the largest possible value of the file offset for that file by setting `l_len` to 0. If such a lock also has `l_start` set to 0 and `l_whence` is set to `SEEK_SET`, the whole file will be locked.

If a process has an existing lock in which `l_len` is 0 and which includes the last byte of the requested segment, and an unlock (F_UNLCK) request is made in which `l_len` is non-zero and the offset of the last byte of the requested segment is the maximum value for an object of type `off_t`, then the F_UNLCK request will be treated as a request to unlock from the start of the requested segment with an `l_len` equal to 0. Otherwise, the request will attempt to unlock only the requested segment.

There will be at most one type of lock set for each byte in the file. Before a successful return from an F_SETLK, F_SETLKW, or F_SETLK64 request when the calling process has previously existing locks on bytes in the region specified by the request, the previous lock type for each byte in the specified region will be replaced by the new lock type. As specified above under the descriptions of shared locks and exclusive locks, an F_SETLK, F_SETLKW, or F_SETLK64 request will (respectively) fail or block when another process has existing locks on bytes in the specified region and the type of any of those locks conflicts with the type specified in the request.

All locks associated with a file for a given process are removed when a file descriptor for that file is closed by that process or the process holding that file descriptor terminates. Locks are not inherited by a child process created using fork(2).

A potential for deadlock occurs if a process controlling a locked region is put to sleep by attempting to lock another process' locked region. If the system detects that sleeping until a locked region is unlocked would cause a deadlock, fcntl() will fail with an EDEADLK error.

The following values for `cmd` are used for file share reservations. A share reservation is placed on an entire file to allow cooperating processes to control access to the file.

- **F_SHARE**: Sets a share reservation on a file with the specified access mode and designates which types of access to deny.
- **F_UNSHARE**: Remove an existing share reservation.

File share reservations are an advisory form of access control among cooperating processes, on both local and remote machines. They are most often used by DOS or Windows emulators and DOS based NFS clients. However, native UNIX versions of DOS or Windows applications may also choose to use this form of access control.

A share reservation is described by an fshare structure defined in `<sys/fcntl.h>`, which is included in `<fcntl.h>` as follows:

```c
struct fshare
```

System Calls  83
typedef struct fshare {
    short f_access;
    short f_deny;
    int f_id;
} fshare_t;

A share reservation specifies the type of access, f_access, to be requested on the open file descriptor. If access is granted, it further specifies what type of access to deny other processes, f_deny. A single process on the same file may hold multiple non-conflicting reservations by specifying an identifier, f_id, unique to the process, with each request.

An F_UNSHARE request releases the reservation with the specified f_id. The f_access and f_deny fields are ignored.

Valid f_access values are:
F_RDACC Set a file share reservation for read-only access.
F_WRACC Set a file share reservation for write-only access.
F_RWACC Set a file share reservation for read and write access.

Valid f_deny values are:
F_COMPAT Set a file share reservation to compatibility mode.
F_RDDNY Set a file share reservation to deny read access to other processes.
F_WRDNY Set a file share reservation to deny write access to other processes.
F_RWDNY Set a file share reservation to deny read and write access to other processes.
F_NODNY Do not deny read or write access to any other process.

RETURN VALUES
Upon successful completion, the value returned depends on cmd as follows:
F_DUPFD A new file descriptor.
F_FREESP Value of 0.
F_GETFD Value of flags defined in <fcntl.h>. The return value will not be negative.
F_GETFL Value of file status flags and access modes. The return value will not be negative.
F_GETLK Value other than −1.
F_GETLK64 Value other than −1.
F_GETOWN Value of the socket owner process or process group; this will not be −1.
F_GETXFL Value of file status flags, access modes, and creation and assignment flags. The return value will not be negative.
 otherwise, -1 is returned and errno is set to indicate the error.

**ERRORS**

The `fcntl()` function will fail if:

- **EAGAIN**
  - The `cmd` argument is `F_SETLK` or `F_SETLK64`, the type of lock (l_type) is a shared (`F_RDLCK`) or exclusive (`F_WRLCK`) lock, and the segment of a file to be locked is already exclusive-locked by another process; or the type is an exclusive lock and some portion of the segment of a file to be locked is already shared-locked or exclusive-locked by another process.

  The `cmd` argument is `F_FREESP`, the file exists, mandatory file/record locking is set, and there are outstanding record locks on the file; or the `cmd` argument is `F_SETLK`, `F_SETLK64`, `F_SETLKW`, or `F_SETLKW64`, mandatory file/record locking is set, and the file is currently being mapped to virtual memory using `mmap(2)`.

  The `cmd` argument is `F_SHARE` and `f_access` conflicts with an existing `f_deny` share reservation.

- **EBADF**
  - The `fd` argument is not a valid open file descriptor; or the `cmd` argument is `F_SETLK`, `F_SETLK64`, `F_SETLKW`, or `F_SETLKW64`, the type of lock, l_type, is a shared lock (`F_RDLCK`), and `fd` is not a valid file descriptor open for reading; or the type of lock l_type is an exclusive lock (`F_WRLCK`) and `fd` is not a valid file descriptor open for writing.

  The `cmd` argument is `F_FREESP` and `fd` is not a valid file descriptor open for writing.

  The `cmd` argument is `F_DUP2FD`, and `arg` is negative or is not less than the current resource limit for `RLIMIT_NOFILE`. 

---

`fcntl(2)`

<table>
<thead>
<tr>
<th><code>F_SETFD</code></th>
<th>Value other than -1.</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>F_SETFL</code></td>
<td>Value other than -1.</td>
</tr>
<tr>
<td><code>F_SETLK</code></td>
<td>Value other than -1.</td>
</tr>
<tr>
<td><code>F_SETLK64</code></td>
<td>Value other than -1.</td>
</tr>
<tr>
<td><code>F_SETLKW</code></td>
<td>Value other than -1.</td>
</tr>
<tr>
<td><code>F_SETLKW64</code></td>
<td>Value other than -1.</td>
</tr>
<tr>
<td><code>F_SETOWN</code></td>
<td>Value other than -1.</td>
</tr>
<tr>
<td><code>F_SHARE</code></td>
<td>Value other than -1.</td>
</tr>
<tr>
<td><code>F_UNSHARE</code></td>
<td>Value other than -1.</td>
</tr>
</tbody>
</table>
The `cmd` argument is `F_SHARE`, the `f_access` share reservation is for write access, and `fd` is not a valid file descriptor open for writing.

The `cmd` argument is `F_SHARE`, the `f_access` share reservation is for read access, and `fd` is not a valid file descriptor open for reading.

**EFAULT**
The `cmd` argument is `F_GETLK`, `F_GETLK64`, `F_SETLK`, `F_SETLK64`, `F_SETLK64`, or `F_FREESP` and the `arg` argument points to an illegal address.

The `cmd` argument is `FSHARE` or `F_UNSHARE` and `arg` points to an illegal address.

**EINTR**
The `cmd` argument is `F_SETLKW` or `F_SETLKW64` and the function was interrupted by a signal.

**EINVAL**
The `cmd` argument is invalid; or the `cmd` argument is `F_DUPFD` and `arg` is negative or greater than or equal to `OPEN_MAX`; or the `cmd` argument is `F_GETLK`, `F_GETLK64`, `F_SETLK`, `F_SETLK64`, `F_SETLK64` and the data pointed to by `arg` is not valid; or `fd` refers to a file that does not support locking.

The `cmd` argument is `F_UNSHARE` and a reservation with this `f_id` for this process does not exist.

**EIO**
An I/O error occurred while reading from or writing to the file system.

**EMFILE**
The `cmd` argument is `F_DUPFD` and either `OPEN_MAX` file descriptors are currently open in the calling process, or no file descriptors greater than or equal to `arg` are available.

**ENOLCK**
The `cmd` argument is `F_SETLK`, `F_SETLK64`, `F_SETLK`, or `F_SETLK6464` and satisfying the lock or unlock request would result in the number of locked regions in the system exceeding a system-imposed limit.

**ENOLINK**
Either the `fd` argument is on a remote machine and the link to that machine is no longer active; or the `cmd` argument is `F_FREESP`, the file is on a remote machine, and the link to that machine is no longer active.

**EOVERFLOW**
One of the values to be returned cannot be represented correctly.

The `cmd` argument is `F_GETLK`, `F_SETLK`, or `F_SETLK64` and the smallest or, if `l_len` is non-zero, the largest, offset of any byte in the requested segment cannot be represented correctly in an object of type `off_t`. 
The `cmd` argument is `F_GETLK64`, `F_SETLK64`, or `F_SETLKW64` and the smallest or, if `l_len` is non-zero, the largest, offset of any byte in the requested segment cannot be represented correctly in an object of type `off64_t`.

The `fcntl()` function may fail if:

**EAGAIN**  
The `cmd` argument is `F_SETLK`, `F_SETLK64`, `F_SETLKW`, or `F_SETLKW64`, and the file is currently being mapped to virtual memory using `mmap(2)`.

**EDEADLK**  
The `cmd` argument is `F_SETLKW` or `F_SETLKW64`, the lock is blocked by some lock from another process and putting the calling process to sleep, waiting for that lock to become free would cause a deadlock.

The `cmd` argument is `F_FREESP`, mandatory record locking is enabled, `O_NDELAY` and `O_NONBLOCK` are clear and a deadlock condition was detected.

### ATTRIBUTES

See `attributes(5)` for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal Safe</td>
</tr>
</tbody>
</table>

### SEE ALSO

`lockd(1M)`, `chmod(2)`, `close(2)`, `creat(2)`, `dup(2)`, `exec(2)`, `fork(2)`, `mmap(2)`, `open(2)`, `pipe(2)`, `read(2)`, `sigaction(2)`, `write(2)`, `dup2(3C)`, `attributes(5)`, `fcntl(3HEAD)`

### Programming Interfaces Guide

### NOTES

In the past, the variable `errno` was set to `EACCES` rather than `EAGAIN` when a section of a file is already locked by another process. Therefore, portable application programs should expect and test for either value.

Advisory locks allow cooperating processes to perform consistent operations on files, but do not guarantee exclusive access. Files can be accessed without advisory locks, but inconsistencies may result. The network share locking protocol does not support the `f_deny` value of `F_COMPAT`. For network file systems, if `f_access` is `F_RDAcc`, `f_deny` is mapped to `F_RDDNY`. Otherwise, it is mapped to `F_RDONLY`.

To prevent possible file corruption, the system may reject `mmap()` requests for advisory locked files, or it may reject advisory locking requests for mapped files. Applications that require a file be both locked and mapped should lock the entire file (`l_start` and `l_len` both set to 0). If a file is mapped, the system may reject an unlock request, resulting in a lock that does not cover the entire file.
If the file server crashes and has to be rebooted, the lock manager (see `lockd(1M)`) attempts to recover all locks that were associated with that server. If a lock cannot be reclaimed, the process that held the lock is issued a `SIGLOST` signal.
fork(2)

NAME
fork, fork1 – create a new process

SYNOPSIS
#include <sys/types.h>
#include <unistd.h>

pid_t fork(void);
pid_t fork1(void);

DESCRIPTION
The fork() and fork1() functions create a new process. The new process (child
process) is an exact copy of the calling process (parent process). The child process
inherits the following attributes from the parent process:

- real user ID, real group ID, effective user ID, effective group ID
- environment
- open file descriptors
- close-on-exec flags (see exec(2))
- signal handling settings (that is, SIG_DFL, SIG_IGN, SIG_HOLD, function address)
- supplementary group IDs
- set-user-ID mode bit
- set-group-ID mode bit
- profiling on/off status
- nice value (see nice(2))
- scheduler class (see priocntl(2))
- all attached shared memory segments (see shmem(2))
- process group ID – memory mappings (see mmap(2))
- session ID (see exit(2))
- current working directory
- root directory
- file mode creation mask (see umask(2))
- resource limits (see getrlimit(2))
- controlling terminal
- saved user ID and group ID
- task ID and project ID
- processor bindings (see processor_bind(2))
- processor set bindings (see pset_bind(2))

Scheduling priority and any per-process scheduling parameters that are specific to a
given scheduling class may or may not be inherited according to the policy of that
particular class (see priocntl(2)). The child process differs from the parent process in
the following ways:
The child process has a unique process ID which does not match any active process group ID.

The child process has a different parent process ID (that is, the process ID of the parent process).

The child process has its own copy of the parent’s file descriptors and directory streams. Each of the child’s file descriptors shares a common file pointer with the corresponding file descriptor of the parent.

Each shared memory segment remains attached and the value of `shm_nattach` is incremented by 1.

All `semadj` values are cleared (see `semop(2)`).

Process locks, text locks, data locks, and other memory locks are not inherited by the child (see `plock(3C)` and `memcntl(2)`).

The child process’s `tms` structure is cleared: `tms_utime`, `stime`, `cutime`, and `cstime` are set to 0 (see `times(2)`).

The child processes resource utilizations are set to 0; see `getrlimit(2)`. The `it_value` and `it_interval` values for the `ITIMER_REAL` timer are reset to 0; see `getitimer(2)`.

The set of signals pending for the child process is initialized to the empty set.

Timers created by `timer_create(3RT)` are not inherited by the child process.

No asynchronous input or asynchronous output operations are inherited by the child.

Any preferred hardware address translation sizes (see `memcntl(2)`) are inherited by the child.

Record locks set by the parent process are not inherited by the child process (see `fcntl(2)`).

In applications that use the Solaris threads API rather than the POSIX threads API (applications linked with `-lthread` but not `-lpthread`), `fork()` duplicates in the child process all threads (see `thr_create(3THR)`) and LWPs in the parent process. The `fork1()` function duplicates only the calling thread (LWP) in the child process.

In applications that use the POSIX threads API rather than the Solaris threads API (applications linked with `-lpthread`, whether or not linked with `-lthread`), a call to `fork()` is like a call to `fork1()`, which replicates only the calling thread. There is no call that forks a child with all threads and LWPs duplicated in the child.

Note that if a program is linked with both libraries (`-lthread` and `-lpthread`), the POSIX semantic of `fork()` prevails.

If a Solaris threads application calls `fork1()` or a POSIX threads application calls `fork()`, and the child does more than simply call `exec()`, there is a possibility of deadlock occurring in the child. The application should use `pthread_atfork(3C)` to ensure safety with respect to this deadlock. Should there be any outstanding mutexes
throughout the process, the application should call pthread_atfork() to wait for and acquire those mutexes prior to calling fork() or fork1(). See "MT-Level of Libraries" on the attributes(5) manual page.

**RETURN VALUES**

Upon successful completion, fork() and fork1() return 0 to the child process and return the process ID of the child process to the parent process. Otherwise, (pid_t)-1 is returned to the parent process, no child process is created, and errno is set to indicate the error.

**ERRORS**

The fork() function will fail if:

- **EAGAIN** The system-imposed limit on the total number of processes under execution by a single user has been exceeded; or the total amount of system memory available is temporarily insufficient to duplicate this process.
- **ENOMEM** There is not enough swap space.

**ATTRIBUTES**

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>fork() is Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO**

alarm(2), exec(2), exit(2), fcntl(2), getitimer(2), getrlimit(2), memcntl(2), mmap(2), nice(2), priocntl(2), ptrace(2), semop(2), shmat(2), times(2), umask(2), wait(2), exit(3C), plock(3C), pthread_atfork(3C), signal(3C), system(3C), thr_create(3THR) timer_create(3RT), attributes(5), standards(5)

**NOTES**

An application should call _exit() rather than exit(3C) if it cannot execve(), since exit() will flush and close standard I/O channels and thereby corrupt the parent process’s standard I/O data structures. Using exit(3C) will flush buffered data twice. See exit(2).

The thread (or LWP) in the child that calls fork1() must not depend on any resources held by threads (or LWPs) that no longer exist in the child. In particular, locks held by these threads (or LWPs) will not be released.

In a multithreaded process, fork() or fork1() can cause blocking system calls to be interrupted and return with an EINTR error.

The fork() and fork1() functions suspend all threads in the process before proceeding. Threads that are executing in the kernel and are in an uninterruptible wait cannot be suspended immediately and therefore cause a delay before fork() and fork1() can complete. During this delay, since all other threads will have already been suspended, the process will appear “hung.”
fpathconf(2)

NAME    fpathconf, pathconf – get configurable pathname variables

SYNOPSIS #include <unistd.h>

long fpathconf(int fd, int name);
long pathconf(const char *path, int name);

DESCRIPTION The fpathconf() and pathconf() functions provide a method for the application to determine the current value of a configurable limit or option (variable) that is associated with a file or directory.

For pathconf(), the path argument points to the pathname of a file or directory.

For fpathconf(), the fd argument is an open file descriptor.

The name argument represents the variable to be queried relative to that file or directory. The variables in the following table come from <limits.h> or <unistd.h> and the symbolic constants, defined in <unistd.h>, are the corresponding values used for name:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Value of name</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>_PC_FILESIZEBITS</td>
<td>FILESIZEBITS</td>
<td>3,4</td>
</tr>
<tr>
<td>_PC_LINK_MAX</td>
<td>LINK_MAX</td>
<td>1</td>
</tr>
<tr>
<td>_PC_MAX_CANON</td>
<td>MAX_CANON</td>
<td>2</td>
</tr>
<tr>
<td>_PC_MAX_INPUT</td>
<td>MAX_INPUT</td>
<td>2</td>
</tr>
<tr>
<td>_PC_NAME_MAX</td>
<td>NAME_MAX</td>
<td>3,4</td>
</tr>
<tr>
<td>_PC_PATH_MAX</td>
<td>PATH_MAX</td>
<td>4,5</td>
</tr>
<tr>
<td>_PC_PIPE_BUF</td>
<td>PIPE_BUF</td>
<td>6</td>
</tr>
<tr>
<td>_PC_XATTR_ENABLED</td>
<td>XATTR_ENABLED</td>
<td>1</td>
</tr>
<tr>
<td>_PC_XATTR_EXISTS</td>
<td>XATTR_EXISTS</td>
<td>1</td>
</tr>
<tr>
<td>_PC_CHOWN_RESTRICTED</td>
<td>_POSIX_CHOWN_RESTRICTED</td>
<td>7</td>
</tr>
<tr>
<td>_PC_NO_TRUNC</td>
<td>_POSIX_NO_TRUNC</td>
<td>3,4</td>
</tr>
<tr>
<td>_PC_VDISABLE</td>
<td>_POSIX_VDISABLE</td>
<td>2</td>
</tr>
<tr>
<td>_PC_ASYNC_IO</td>
<td>_POSIX_ASYNC_IO</td>
<td>8</td>
</tr>
<tr>
<td>_PC_PRIO_IO</td>
<td>_POSIX_PRIO_IO</td>
<td>8</td>
</tr>
<tr>
<td>_PC_SYNC_IO</td>
<td>_POSIX_SYNC_IO</td>
<td>8</td>
</tr>
</tbody>
</table>

Notes:
1. If `path` or `fildes` refers to a directory, the value returned applies to the directory itself.

2. If `path` or `fildes` does not refer to a terminal file, it is unspecified whether an implementation supports an association of the variable name with the specified file.

3. If `path` or `fildes` refers to a directory, the value returned applies to filenames within the directory.

4. If `path` or `fildes` does not refer to a directory, it is unspecified whether an implementation supports an association of the variable name with the specified file.

5. If `path` or `fildes` refers to a directory, the value returned is the maximum length of a relative pathname when the specified directory is the working directory.

6. If `path` refers to a FIFO, or `fildes` refers to a pipe or FIFO, the value returned applies to the referenced object. If `path` or `fildes` refers to a directory, the value returned applies to any FIFO that exists or can be created within the directory. If `path` or `fildes` refers to any other type of file, it is unspecified whether an implementation supports an association of the variable name with the specified file.

7. If `path` or `fildes` refers to a directory, the value returned applies to any files, other than directories, that exist or can be created within the directory.

8. If `path` or `fildes` refers to a directory, it is unspecified whether an implementation supports an association of the variable name with the specified file.

If `name` is an invalid value, both `pathconf()` and `fpathconf()` return −1 and `errno` is set to indicate the error.

If the variable corresponding to `name` has no limit for the `path` or file descriptor, both `pathconf()` and `fpathconf()` return −1 without changing `errno`. If the implementation needs to use `path` to determine the value of `name` and the implementation does not support the association of `name` with the file specified by `path`, or if the process did not have appropriate privileges to query the file specified by `path`, or `path` does not exist, `pathconf()` returns −1 and `errno` is set to indicate the error.

If the implementation needs to use `fildes` to determine the value of `name` and the implementation does not support the association of `name` with the file specified by `fildes`, or if `fildes` is an invalid file descriptor, `fpathconf()` will return −1 and `errno` is set to indicate the error.

Otherwise `pathconf()` or `fpathconf()` returns the current variable value for the file or directory without changing `errno`. The value returned will not be more restrictive than the corresponding value available to the application when it was compiled with the implementation’s `<limits.h>` or `<unistd.h>`.

The `pathconf()` function will fail if:

**EINVAL** The value of `name` is not valid.
ELOOP
Too many symbolic links were encountered in resolving path.

The pathconf() function may fail if:

EACCES Search permission is denied for a component of the path prefix.

EINVAL The implementation does not support an association of the variable name with the specified file.

ENAMETOOLONG The length of the path argument exceeds PATH_MAX or a pathname component is longer than NAME_MAX.

ENAMETOOLONG Pathname resolution of a symbolic link produced an intermediate result whose length exceeds PATH_MAX.

ENOENT A component of path does not name an existing file or path is an empty string.

ENOTDIR A component of the path prefix is not a directory.

The fpathconf() function will fail if:

EINVAL The value of name is not valid.

The fpathconf() function may fail if:

EBADF The fildes argument is not a valid file descriptor.

EINVAL The implementation does not support an association of the variable name with the specified file.

ATTRIBUTES See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>fpathconf() is Standard; pathconf() is Stable</td>
</tr>
<tr>
<td>MT-Level</td>
<td>pathconf() is Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO sysconf(3C), limits(4), attributes(5), standards(5)
getacct(2)

NAME
getcct, putacct, wracct – get, put, or write extended accounting data

SYNOPSIS
#include <sys/exacct.h>

size_t getacct(idtype_t idtype, id_t id, void *buf, size_t bufsize);
int putacct(idtype_t idtype, id_t id, void *buf, size_t bufsize, int flags);
int wracct(idtype_t idtype, id_t id, int flags);

DESCRIPTION
These functions provide access to the extended accounting facility.

The getacct() function returns extended accounting buffers from the kernel for
currently executing tasks and processes. The resulting data buffer is a packed exact
object that can be unpacked using ea_unpack_object() (see
ea_pack_object(3EXACCT)) and subsequently manipulated using the functions of
the extended accounting library, libexacct(3LIB).

The putacct() function provides privileged processes the ability to tag accounting
records with additional data specific to that process. For instance, a queueing facility
might want to record to which queue a given task or process was submitted prior to
running. The flags argument determines whether the contents of buf should be treated
as raw data (EP_RAW) or as an embedded exact struct (EP_EXACCT_OBJECT). In
the case of EP_EXACCT_OBJECT, buf must be a packed exact object as returned by
ea_pack_object(3EXACCT). The use of an inappropriate flag or the inclusion of
corrupt exact data will likely corrupt the enclosing exact file.

The wracct() function requests the kernel to write, given its internal state of
resource usage, the appropriate data for the specified task or process. The flags field
determines whether a partial (EW_PARTIAL) or interval record (EW_INTERVAL) is
written.

These functions require root privilege, as they allow inquiry or reporting relevant to
system tasks and processes other than the invoking process. The putacct() and
wrracct() functions also cause the kernel to write records to the system’s extended
accounting files.

RETURN VALUES
The getacct() function returns the number of bytes required to represent the
extended accounting record for the requested system task or process. If bufsize exceeds
the returned size, buf will contain a valid accounting record buffer. If bufsize is less than
the return value, buf will contain the first bufsize bytes of the record. If bufsize is 0,
getacct() returns only the number of bytes required to represent the extended
accounting record. In the event of failure, −1 is returned and errno is set to indicate
the error.

The putacct() and wracct() functions return 0 if the record was successfully
written. Otherwise, −1 is returned and errno is set to indicate the error.

ERRORS
The getacct(), putacct(), and wracct() functions will fail if:
EINVAL
The idtype argument was not P_TASKID or P_PID.
getacct(2)

ENOSPC  The file system containing the extended accounting file is full. The 
wracct() or putacct() function will fail if the record size 
would exceed the amount of space remaining on the file system.

ENOTACTIVE  The extended accounting facility for the requested idtype_t is 
not active. Either putacct() attempted to write a task record 
when the task accounting file was unset, or getacct() attempted 
to retrieve accounting data for a process when extended process 
accounting was inactive.

EPERM  The invoking process lacks sufficient permission to perform the 
request operation.

ERSCH  The id argument does not refer to a presently active system task ID 
or process ID.

The putacct() and wracct() functions will fail if:

EINVAL  The flags argument is neither EW_PARTIAL nor EW_INTERVAL.

ATTRIBUTES  See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO  ea_pack_object(3EXACCT), libexacct(3LIB), attributes(5)
NAME
getaudit, setaudit, getaudit_addr, setaudit_addr – get and set process audit
information

SYNOPSIS
cc [ flag ... ] file ... -bsm -lsocket -lnsl -lintl [ library ... ]
#include <sys/param.h>
#include <bsm/audit.h>

int getaudit(struct auditinfo *info);
int setaudit(struct auditinfo *info);

int getaudit_addr(struct auditinfo_addr *info, int length);
int setaudit_addr(struct auditinfo_addr *info, int length);

DESCRIPTION
The getaudit() function gets the audit ID, the preselection mask, the terminal ID
and the audit session ID for the current process.

Note that getaudit() may fail and return an E2BIG errno if the address field in the
terminal ID is larger than 32 bits. In this case, getaudit_addr() should be used.

The setaudit() function sets the audit ID, the preselection mask, the terminal ID
and the audit session ID for the current process.

The getaudit_addr() function returns a variable length auditinfo_addr
structure that contains the audit ID, the preselection mask, the terminal ID, and the
audit session ID for the current process. The terminal ID contains a size field that
indicates the size of the network address.

The setaudit_addr() function sets the audit ID, the preselection mask, the terminal
ID, and the audit session ID for the current process. The values are taken from the
variable length struture auditinfo_addr. The terminal ID contains a size field that
indicates the size of the network address.

The auditinfo structure is used to pass the process audit information and contains
the following members:
auid_t ai_auid; /* audit user ID */
au_mask_t ai_mask; /* preselection mask */
au_tid_t ai_termid; /* terminal ID */
au_asid_t ai_asid; /* audit session ID */

The auditinfo_addr structure is used to pass the process audit information and
contains the following members:
auid_t ai_auid; /* audit user ID */
au_mask_t ai_mask; /* preselection mask */
au_tid_t ai_termid; /* terminal ID */
au_asid_t ai_asid; /* audit session ID */

RETURN VALUES
Upon successful completion, getaudit() and setaudit() return 0. Otherwise, −1
is returned and errno is set to indicate the error.

ERRORS
The getaudit() and setaudit() functions will fail if:
getaudit(2)

EFAULT  The info parameter points outside the process's allocated address space.

EPERM   The process's effective user ID is not superuser.

USAGE   Only processes with the effective user ID of the superuser can successfully execute these calls.

ATTRIBUTES See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Stable</td>
</tr>
<tr>
<td>MT-Level</td>
<td>MT-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO bsmconv(1M), audit(2), attributes(5)

NOTES  The functionality described in this man page is available only if the Basic Security Module (BSM) has been enabled. See bsmconv(1M) for more information.
NAME
getauid, setauid – get and set user audit identity

SYNOPSIS
cc [ flag ... ] file ... -bsm -lsocket -lnsl -lintl [ library ... ]
#include <sys/param.h>
#include <bsm/audit.h>

int getauid(au_id_t *auid);
int setauid(au_id_t *auid);

DESCRIPTION
The getauid() function returns the audit user ID for the current process. This value
is initially set at login time and inherited by all child processes. This value does not
change when the real/effective user IDs change, so it can be used to identify the
logged-in user even when running a setuid program. The audit user ID governs audit
decisions for a process.

The setauid() function sets the audit user ID for the current process.

RETURN VALUES
Upon successful completion, the getauid() function returns the audit user ID of the
current process on success. Otherwise, it returns -1 and sets errno to indicate the
error.

Upon successful completion the setauid() function returns 0. Otherwise, -1 is
returned and errno is set to indicate the error.

ERRORS
The getauid() and setauid() functions will fail if:

EFAULT The auid argument points to an invalid address.

EPERM The process’s effective user ID is not super-user.

USAGE
Only the super-user may successfully execute these calls.

SEE ALSO
bsmconv(1M), audit(2), getaudit(2)

NOTES
The functionality described in this man page is available only if the Basic Security
Module (BSM) has been enabled. See bsmconv(1M) for more information.

These system calls have been superseded by getaudit() and setaudit().
getcontext(2)

NAME  
getcontext, setcontext – get and set current user context

SYNOPSIS  
#include <ucontext.h>

int getcontext(ucontext_t *ucp);
int setcontext(const ucontext_t *ucp);

DESCRIPTION  
The getcontext() function initializes the structure pointed to by ucp to the current user context of the calling process. The ucontext_t type that ucp points to defines the user context and includes the contents of the calling process’ machine registers, the signal mask, and the current execution stack.

The setcontext() function restores the user context pointed to by ucp. A successful call to setcontext() does not return; program execution resumes at the point specified by the ucp argument passed to setcontext(). The ucp argument should be created either by a prior call to getcontext(), or by being passed as an argument to a signal handler. If the ucp argument was created with getcontext(), program execution continues as if the corresponding call of getcontext() had just returned. If the ucp argument was created with makecontext(3C), program execution continues with the function passed to makecontext(3C). When that function returns, the process continues as if after a call to setcontext() with the ucp argument that was input to makecontext(3C). If the ucp argument was passed to a signal handler, program execution continues with the program instruction following the instruction interrupted by the signal. If the uc_link member of the ucontext_t structure pointed to by the ucp argument is equal to 0, then this context is the main context, and the process will exit when this context returns. The effects of passing a ucp argument obtained from any other source are unspecified.

RETURN VALUES  
On successful completion, setcontext() does not return and getcontext() returns 0. Otherwise, −1 is returned.

ERRORS  
No errors are defined.

USAGE  
When a signal handler is executed, the current user context is saved and a new context is created. If the thread leaves the signal handler via longjmp(3UCB), then it is unspecified whether the context at the time of the corresponding setjmp(3UCB) call is restored and thus whether future calls to getcontext() will provide an accurate representation of the current context, since the context restored by longjmp(3UCB) may not contain all the information that setcontext() requires. Signal handlers should use siglongjmp(3C) instead.

Portable applications should not modify or access the uc_mcontext member of ucontext_t. A portable application cannot assume that context includes any process-wide static data, possibly including errno. Users manipulating contexts should take care to handle these explicitly when required.

SEE ALSO  
sigaction(2), sigaltstack(2), sigprocmask(2), bsd_signal(3C), makecontext(3C), setjmp(3UCB), sigsetjmp(3C), ucontext(3HEAD)
The `getdents()` function attempts to read `nbyte` bytes from the directory associated with the file descriptor `fildes` and to format them as file system independent directory entries in the buffer pointed to by `buf`. Since the file system independent directory entries are of variable lengths, in most cases the actual number of bytes returned will be less than `nbyte`. The file system independent directory entry is specified by the `dirent` structure. See `dirent(3HEAD)`.

On devices capable of seeking, `getdents()` starts at a position in the file given by the file pointer associated with `fildes`. Upon return from `getdents()`, the file pointer is incremented to point to the next directory entry.

Upon successful completion, a non-negative integer is returned indicating the number of bytes actually read. A return value of 0 indicates the end of the directory has been reached. Otherwise, -1 is returned and `errno` is set to indicate the error.

The `getdents()` function will fail if:

- **EBADF**: The `fildes` argument is not a valid descriptor open for reading.
- **EFAULT**: The `buf` argument points to an illegal address.
- **EINVAL**: The `nbyte` argument is not large enough for one directory entry.
- **EIO**: An I/O error occurred while accessing the file system.
- **ENOENT**: The current file pointer for the directory is not located at a valid entry.
- **ENOLINK**: The `fildes` argument points to a remote machine and the link to that machine is no longer active.
- **ENOTDIR**: The `fildes` argument is not a directory.
- **EOVERFLOW**: The value of the `dirent` structure member `d_ino` or `d_off` cannot be represented in an `ino_t` or `off_t`.

The `getdents()` function was developed to implement the `readdir(3C)` function and should not be used for other purposes.

The `getdents()` function has a transitional interface for 64-bit file offsets. See `lf64(5)`.

**SEE ALSO**

`readdir(3C), dirent(3HEAD), lf64(5)`
getgroups(2)

NAME
getgroups, setgroups – get or set supplementary group access list IDs

SYNOPSIS
#include <unistd.h>

int getgroups(int gidsetsize, gid_t *grouplist);
int setgroups(int ngroups, const gid_t *grouplist);

DESCRIPTION
The getgroups() function gets the current supplemental group access list of the
calling process and stores the result in the array of group IDs specified by grouplist.
This array has gidsetsize entries and must be large enough to contain the entire list.
This list cannot be larger than NGROUPS_MAX. If gidsetsize equals 0, getgroups() will
return the number of groups to which the calling process belongs without modifying
the array pointed to by grouplist.

The setgroups() function sets the supplementary group access list of the calling
process from the array of group IDs specified by grouplist. The number of entries is
specified by ngroups and can not be greater than NGROUPS_MAX.

RETURN VALUES
Upon successful completion, getgroups() returns the number of supplementary
group IDs set for the calling process and setgroups() returns 0. Otherwise, -1 is
returned and errno is set to indicate the error.

ERRORS
The getgroups() and setgroups() functions will fail if:

EFAULT
A referenced part of the array pointed to by grouplist is an illegal
address.

EINVAL
The value of gidsetsize is non-zero and less than the number of
supplementary group IDs set for the calling process.

The setgroups() function will fail if:

EINVAL
The value of ngroups is greater than NGROUPS_MAX.

EPERM
The effective user of the calling process is not super-user.

USAGE
Use of the setgroups() function requires superuser privileges.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
groups(1), chown(2), getuid(2), setuid(2), getgrnam(3C), initgroups(3C),
attributes(5)
getitimer, setitimer – get or set value of interval timer

#include <sys/time.h>

int getitimer(int which, struct itimerval *value);
int setitimer(int which, const struct itimerval *value, struct itimerval *ovalue);

The system provides each process with four interval timers, defined in sys/time.h. The getitimer() function stores the current value of the timer specified by which into the structure pointed to by value. The setitimer() function call sets the value of the timer specified by which to the value specified in the structure pointed to by value, and if ovalue is not NULL, stores the previous value of the timer in the structure pointed to by ovalue.

A timer value is defined by the itimerval structure (see gettimeofday(3C)) for the definition of timeval), which includes the following members:

struct timeval it_interval; /* timer interval */
struct timeval it_value; /* current value */

The it_value member indicates the time to the next timer expiration. The it_interval member specifies a value to be used in reloading it_value when the timer expires. Setting it_value to 0 disables a timer, regardless of the value of it_interval. Setting it_interval to 0 disables a timer after its next expiration (assuming it_value is non-zero).

Time values smaller than the resolution of the system clock are rounded up to the resolution of the system clock, except for ITIMER_REALPROF, whose values are rounded up to the resolution of the profiling clock. The four timers are as follows:

ITIMER_REAL
Decrments in real time. A SIGALRM signal is delivered when this timer expires.

ITIMER_VIRTUAL
Decrments in process virtual time. It runs only when the process is executing. A SIGVTALRM signal is delivered when it expires.

ITIMER_PROF
Decrments both in process virtual time and when the system is running on behalf of the process. It is designed to be used by interpreters in statistically profiling the execution of interpreted programs. Each time the ITIMER_PROF timer expires, the SIGPROF signal is delivered. Because this signal may interrupt in-progress functions, programs using this timer must be prepared to restart interrupted functions.

ITIMER_REALPROF
Decrments in real time. It is designed to be used for real-time profiling of multithreaded programs. Each time the ITIMER_REALPROF timer expires, one counter in a set of counters maintained by the system for each lightweight process (lwp) is incremented. The counter corresponds to the state of the lwp at the time of the timer tick. All lwps executing in user mode when the timer expires are
interrupted into system mode. When each lwp resumes execution in user mode, if any of the elements in its set of counters are non-zero, the SIGPROF signal is delivered to the lwp. The SIGPROF signal is delivered before any other signal except SIGKILL. This signal does not interrupt any in-progress function. A siginfo structure, defined in <sys/siginfo.h>, is associated with the delivery of the SIGPROF signal, and includes the following members:

- si_tstamp; /* high resolution timestamp */
- si_syscall; /* current syscall */
- si_nsysarg; /* number of syscall arguments */
- si_sysarg[ ]; /* actual syscall arguments */
- si_fault; /* last fault type */
- si_faddr; /* last fault address */
- si_mstate[ ]; /* ticks in each microstate */

The enumeration of microstates (indices into si_mstate) is defined in <sys/msacct.h>.

RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, -1 is returned and errno is set to indicate the error.

ERRORS
The getitimer() and setitimer() functions will fail if:

EINVAL The specified number of seconds is greater than 100,000,000, the number of microseconds is greater than or equal to 1,000,000, or the which argument is unrecognized.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>MT-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
alarm(2), gettimeofday(3C), sleep(3C), sysconf(3C), attributes(5), standards(5)

NOTES
The microseconds field should not be equal to or greater than one second.

The setitimer() function is independent of the alarm() function.

Do not use setitimer(ITIMER_REAL) with the sleep() routine. A sleep(3C) call wipes out knowledge of the user signal handler for SIGALRM.

The ITIMER_PROF and ITIMER_REALPROF timers deliver the same signal and have different semantics. They cannot be used together.

The granularity of the resolution of alarm time is platform-dependent.
getmsg, getpmsg – get next message off a stream

SYNOPSIS

```c
#include <stropts.h>

int getmsg(int filedes, struct strbuf *ctlptr, struct strbuf *dataptr, int *flagsp);
int getpmsg(int filedes, struct strbuf *ctlptr, struct strbuf *dataptr, int *bandp, int *flagsp);
```

DESCRIPTION

The `getmsg()` function retrieves the contents of a message (see `intro(2)`) located at the stream head read queue from a STREAMS file, and places the contents into user specified buffer(s). The message must contain either a data part, a control part, or both. The data and control parts of the message are placed into separate buffers, as described below. The semantics of each part is defined by the STREAMS module that generated the message.

The `getpmsg()` function behaved like `getmsg()`, but provides finer control over the priority of the messages received. Except where noted, all information pertaining to `getmsg()` also pertains to `getpmsg()`.

The `filedes` argument specifies a file descriptor referencing an open stream. The `ctlptr` and `dataptr` arguments each point to a `strbuf` structure, which contains the following members:

```c
int maxlen; /* maximum buffer length */
int len;    /* length of data */
char *buf;  /* ptr to buffer */
```

The `buf` member points to a buffer into which the data or control information is to be placed, and the `maxlen` member indicates the maximum number of bytes this buffer can hold. On return, the `len` member contains the number of bytes of data or control information actually received; 0 if there is a zero-length control or data part; or −1 if no data or control information is present in the message. The `flagsp` argument should point to an integer that indicates the type of message the user is able to receive, as described below.

The `ctlptr` argument holds the control part from the message and the `dataptr` argument holds the data part from the message. If `ctlptr` (or `dataptr`) is NULL or the `maxlen` member is −1, the control (or data) part of the message is not processed and is left on the stream head read queue. If `ctlptr` (or `dataptr`) is not NULL and there is no corresponding control (or data) part of the messages on the stream head read queue, `len` is set to −1. If the `maxlen` member is set to 0 and there is a zero-length control (or data) part, that zero-length part is removed from the read queue and `len` is set to 0. If the `maxlen` member is set to 0 and there are more than zero bytes of control (or data) information, that information is left on the read queue and `len` is set to 0. If the `maxlen` member in `ctlptr` or `dataptr` is less than, respectively, the control or data part of the message, `maxlen` bytes are retrieved. In this case, the remainder of the message is left on the stream head read queue and a non-zero return value is provided, as described below under `RETURN VALUES`.

```c
#include <stropts.h>
```
By default, `getmsg()` processes the first available message on the stream head read queue. A user may, however, choose to retrieve only high priority messages by setting the integer pointed to by `flagsp` to `RS_HIPRI`. In this case, `getmsg()` processes the next message only if it is a high priority message.

If the integer pointed to by `flagsp` is 0, `getmsg()` retrieves any message available on the stream head read queue. In this case, on return, the integer pointed to by `flagsp` will be set to `RS_HIPRI` if a high priority message was retrieved, or to 0 otherwise.

For `getpmsg()`, the `flagsp` argument points to a bitmask with the following mutually-exclusive flags defined: `MSG_HIPRI`, `MSG_BAND`, and `MSG_ANY`. Like `getmsg()`, `getpmsg()` processes the first available message on the stream head read queue. A user may choose to retrieve only high-priority messages by setting the integer pointed to by `flagsp` to `MSG_HIPRI` and the integer pointed to by `bandp` to 0. In this case, `getpmsg()` will only process the next message if it is a high-priority message. In a similar manner, a user may choose to retrieve a message from a particular priority band by setting the integer pointed to by `flagsp` to `MSG_BAND` and the integer pointed to by `bandp` to the priority band of interest. In this case, `getpmsg()` will only process the next message if it is in a priority band equal to, or greater than, the integer pointed to by `bandp`, or if it is a high-priority message. If a user just wants to get the first message off the queue, the integer pointed to by `flagsp` should be set to `MSG_ANY` and the integer pointed to by `bandp` should be set to 0. On return, if the message retrieved was a high-priority message, the integer pointed to by `flagsp` will be set to `MSG_HIPRI` and the integer pointed to by `bandp` will be set to 0. Otherwise, the integer pointed to by `flagsp` will be set to `MSG_BAND` and the integer pointed to by `bandp` will be set to the priority band of the message.

If `O_NDELAY` and `O_NONBLOCK` are clear, `getmsg()` blocks until a message of the type specified by `flagsp` is available on the stream head read queue. If `O_NDELAY` or `O_NONBLOCK` has been set and a message of the specified type is not present on the read queue, `getmsg()` fails and sets `errno` to `EAGAIN`.

If a hangup occurs on the stream from which messages are to be retrieved, `getmsg()` continues to operate normally, as described above, until the stream head read queue is empty. Thereafter, it returns 0 in the `len` member of `ctlptr` and `dataptr`.

Upon successful completion, a non-negative value is returned. A return value of 0 indicates that a full message was read successfully. A return value of `MORECTL` indicates that more control information is waiting for retrieval. A return value of `MOREDATA` indicates that more data are waiting for retrieval. A return value of `MORECTL` | `MOREDATA` indicates that both types of information remain. Subsequent `getmsg()` calls retrieve the remainder of the message. However, if a message of higher priority has been received by the stream head read queue, the next call to `getmsg()` will retrieve that higher priority message before retrieving the remainder of the previously received partial message.

The `getmsg()` and `getpmsg()` functions will fail if:

- `EAGAIN` The `O_NDELAY` or `O_NONBLOCK` flag is set and no messages are available.
<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>EBADF</td>
<td>The <code>fd</code>es argument is not a valid file descriptor open for reading.</td>
</tr>
<tr>
<td>EBADMSG</td>
<td>Queued message to be read is not valid for <code>getmsg</code>.</td>
</tr>
<tr>
<td>EFAULT</td>
<td>The <code>ctlptr</code>, <code>dataptr</code>, <code>bandp</code>, or <code>flagsp</code> argument points to an illegal address.</td>
</tr>
<tr>
<td>EINTR</td>
<td>A signal was caught during the execution of the <code>getmsg</code> function.</td>
</tr>
<tr>
<td>EINVAL</td>
<td>An illegal value was specified in <code>flagsp</code>, or the stream referenced by <code>fd</code>es is linked under a multiplexor.</td>
</tr>
<tr>
<td>ENOSTR</td>
<td>A stream is not associated with <code>fd</code>es.</td>
</tr>
</tbody>
</table>

The `getmsg()` function can also fail if a STREAMS error message had been received at the stream head before the call to `getmsg()`. The error returned is the value contained in the STREAMS error message.

**SEE ALSO**

`intro(2)`, `poll(2)`, `putmsg(2)`, `read(2)`, `write(2)`

*STREAMS Programming Guide*
getpid(2)

NAME
getpid, getpgrp, getppid, getpgid – get process, process group, and parent process IDs

SYNOPSIS
#include <unistd.h>

pid_t getpid(void);

pid_t getpgrp(void);

pid_t getppid(void);

pid_t getpgid(pid_t pid);

DESCRIPTION
The getpid() function returns the process ID of the calling process.
The getpgrp() function returns the process group ID of the calling process.
The getppid() function returns the parent process ID of the calling process.
The getpgid() function returns the process group ID of the process whose process ID is equal to pid, or the process group ID of the calling process, if pid is equal to 0.

RETURN VALUES
Upon successful completion, these functions return the process group ID. Otherwise, getpgid() returns (pid_t)-1 and sets errno to indicate the error.

ERRORS
The getpgid() function will fail if:

EPERM The process whose process ID is equal to pid is not in the same session as the calling process, and the implementation does not allow access to the process group ID of that process from the calling process.

ESRCH There is no process with a process ID equal to pid.

The getpgid() function may fail if:

EINVAL The value of the pid argument is invalid.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
intro(2), exec(2), fork(2), getsid(2), setpgid(2), setpgrp(2), setsid(2), signal(3C), attributes(5)
getrlimit, setrlimit – control maximum system resource consumption

SYNOPSIS
#include <sys/resource.h>

int getrlimit(int resource, struct rlimit *rlp);
int setrlimit(int resource, const struct rlimit *rlp);

DESCRIPTION
Limits on the consumption of a variety of system resources by a process and each process it creates may be obtained with the getrlimit() and set with setrlimit() functions.

Each call to either getrlimit() or setrlimit() identifies a specific resource to be operated upon as well as a resource limit. A resource limit is a pair of values: one specifying the current (soft) limit, the other a maximum (hard) limit. Soft limits may be changed by a process to any value that is less than or equal to the hard limit. A process may (irreversibly) lower its hard limit to any value that is greater than or equal to the soft limit. Only a process with an effective user ID of super-user can raise a hard limit. Both hard and soft limits can be changed in a single call to setrlimit() subject to the constraints described above. Limits may have an “infinite” value of RLIM_INFINITY. The rlp argument is a pointer to struct rlimit that includes the following members:

  rlim_t rlim_cur; /* current (soft) limit */
  rlim_t rlim_max; /* hard limit */

The type rlim_t is an arithmetic data type to which objects of type int, size_t, and off_t can be cast without loss of information.

The possible resources, their descriptions, and the actions taken when the current limit is exceeded are summarized as follows:

- **RLIMIT_CORE**: The maximum size of a core file in bytes that may be created by a process. A limit of 0 will prevent the creation of a core file. The writing of a core file will terminate at this size.

- **RLIMIT_CPU**: The maximum amount of CPU time in seconds used by a process. This is a soft limit only. The SIGXCPU signal is sent to the process. If the process is holding or ignoring SIGXCPU, the behavior is scheduling class defined.

- **RLIMIT_DATA**: The maximum size of a process’s heap in bytes. The brk(2) function will fail with errno set to ENOMEM.

- **RLIMITFSIZE**: The maximum size of a file in bytes that may be created by a process. A limit of 0 will prevent the creation of a file. The SIGXFSZ signal is sent to the process. If the process is holding or ignoring SIGXFSZ, continued attempts to increase the size of a file beyond the limit will fail with errno set to EFBIG.

- **RLIMIT_NOFILE**: One more than the maximum value that the system may assign to a newly created descriptor. This limit constrains the number of file descriptors that a process may create.
RLIMIT_STACK The maximum size of a process’s stack in bytes. The system will not automatically grow the stack beyond this limit.

Within a process, setrlimit() will increase the limit on the size of your stack, but will not move current memory segments to allow for that growth. To guarantee that the process stack can grow to the limit, the limit must be altered prior to the execution of the process in which the new stack size is to be used.

Within a multithreaded process, setrlimit() has no impact on the stack size limit for the calling thread if the calling thread is not the main thread. A call to setrlimit() for RLIMIT_STACK impacts only the main thread’s stack, and should be made only from the main thread, if at all.

The SIGSEGV signal is sent to the process. If the process is holding or ignoring SIGSEGV, or is catching SIGSEGV and has not made arrangements to use an alternate stack (see sigaltstack(2)), the disposition of SIGSEGV will be set to SIG_DFL before it is sent.

RLIMIT_VMEM The maximum size of a process’s mapped address space in bytes. If this limit is exceeded, the brk(2) and mmap(2) functions will fail with errno set to ENOMEM. In addition, the automatic stack growth will fail with the effects outlined above.

RLIMIT_AS This is the maximum size of a process’s total available memory, in bytes. If this limit is exceeded, the brk(2), malloc(3C), mmap(2) and sbrk(2) functions will fail with errno set to ENOMEM. In addition, the automatic stack growth will fail with the effects outlined above.

Because limit information is stored in the per-process information, the shell builtin ulimit command must directly execute this system call if it is to affect all future processes created by the shell.

The value of the current limit of the following resources affect these implementation defined parameters:

<table>
<thead>
<tr>
<th>Limit</th>
<th>Implementation Defined Constant</th>
</tr>
</thead>
<tbody>
<tr>
<td>RLIMIT_FSIZE</td>
<td>FCHR_MAX</td>
</tr>
<tr>
<td>RLIMIT_NOFILE</td>
<td>OPEN_MAX</td>
</tr>
</tbody>
</table>
When using the `getrlimit()` function, if a resource limit can be represented correctly in an object of type `rlim_t`, then its representation is returned; otherwise, if the value of the resource limit is equal to that of the corresponding saved hard limit, the value returned is `RLIM_SAVED_MAX`; otherwise the value returned is `RLIM_SAVED_CUR`.

When using the `setrlimit()` function, if the requested new limit is `RLIM_INFINITY`, the new limit will be “no limit”; otherwise if the requested new limit is `RLIM_SAVED_MAX`, the new limit will be the corresponding saved hard limit; otherwise, if the requested new limit is `RLIM_SAVED_CUR`, the new limit will be the corresponding saved soft limit; otherwise, the new limit will be the requested value. In addition, if the corresponding saved limit can be represented correctly in an object of type `rlim_t`, then it will be overwritten with the new limit.

The result of setting a limit to `RLIM_SAVED_MAX` or `RLIM_SAVED_CUR` is unspecified unless a previous call to `getrlimit()` returned that value as the soft or hard limit for the corresponding resource limit.

A limit whose value is greater than `RLIM_INFINITY` is permitted.

The `exec` family of functions also cause resource limits to be saved. See `exec`(2).

**RETURN VALUES**

Upon successful completion, `getrlimit()` and `setrlimit()` return 0. Otherwise, these functions return −1 and set `errno` to indicate the error.

**ERRORS**

The `getrlimit()` and `setrlimit()` functions will fail if:

- **EFAULT** The `rlp` argument points to an illegal address.
- **EINVAL** An invalid resource was specified; or in a `setrlimit()` call, the new `rlim_cur` exceeds the new `rlim_max`.
- **EPERM** The limit specified to `setrlimit()` would have raised the maximum limit value, and the effective user of the calling process is not super-user.

The `setrlimit()` function may fail if:

- **EINVAL** The limit specified cannot be lowered because current usage is already higher than the limit.

**USAGE**

The `getrlimit()` and `setrlimit()` functions have transitional interfaces for 64-bit file offsets. See `lfs64`(5).

The rlimit functionality is now provided by the more general resource control facility described on the `setrctl`(2) manual page. The actions associated with the resource limits described above are true at system boot, but an administrator can modify the local configuration to modify signal delivery or type. Application authors that utilize rlimts for the purposes of resource awareness should investigate the resource controls facility.
getrlimit(2)

SEE ALSO brk(2), exec(2), fork(2), open(2), setrlimit(2), ulimit(2),
getdtablesize(3C), malloc(3C), signal(3C), signal(3HEAD), sysconf(3C),
lf64(5)
### NAME
getsid – get process group ID of session leader

### SYNOPSIS
```c
#include <unistd.h>

pid_t getsid(pid_t pid);
```

### DESCRIPTION
The `getsid()` function obtains the process group ID of the process that is the session leader of the process specified by `pid`. If `pid` is `(pid_t) 0`, it specifies the calling process.

### RETURN VALUES
Upon successful completion, `getsid()` returns the process group ID of the session leader of the specified process. Otherwise, it returns `(pid_t)−1` and sets `errno` to indicate the error.

### ERRORS
The `getsid()` function will fail if:

- **EPERM** The process specified by `pid` is not in the same session as the calling process, and the implementation does not allow access to the process group ID of the session leader of that process from the calling process.

- **ESRCH** There is no process with a process ID equal to `pid`.

### SEE ALSO
`exec(2), fork(2), getpid(2), getpgid(2), setpgid(2), setsid(2)`
The `getuid()` function returns the real user ID of the calling process. The real user ID identifies the person who is logged in.

The `geteuid()` function returns the effective user ID of the calling process. The effective user ID gives the process various permissions during execution of “set-user-ID” mode processes which use `getuid()` to determine the real user ID of the process that invoked them.

The `getgid()` function returns the real group ID of the calling process.

The `getegid()` function returns the effective group ID of the calling process.

See `attributes(5)` for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

See also `intro(2), setuid(2), attributes(5)`
The `getustack()` function retrieves the address of per-LWP stack boundary information. The address is stored at the location pointed to by `spp`. If this address has not been defined using a previous call to `setustack()`, `NULL` is stored at the location pointed to by `spp`.

The `setustack()` function changes the address of the current thread’s stack boundary information to the value of `sp`.

Upon successful completion, these functions return 0. Otherwise, −1 is returned and `errno` is set to indicate the error.

These functions will fail if:

- `EFAULT` The `spp` or `sp` argument does not refer to a valid address.

Implementors of custom threading libraries should use `setustack()` to set the address of the stack bounds to in internal per-thread data structure.

See `attributes(5)` for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Evolving</td>
</tr>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO _stackGrow(3C), stack_getbounds(3C), stack_inbounds(3C), stack_setbounds(3C), stack_violation(3C), attributes(5)
ioctl(2)

NAME
ioctl – control device

SYNOPSIS
#include <unistd.h>
#include <stropts.h>

int ioctl(int fildes, int request, /* arg */ ...);

DESCRIPTION
The ioctl() function performs a variety of control functions on devices and STREAMS. For non-STREAMS files, the functions performed by this call are device-specific control functions. The request argument and an optional third argument with varying type are passed to the file designated by fildes and are interpreted by the device driver.

For STREAMS files, specific functions are performed by the ioctl() function as described in streamio(7I).

The fildes argument is an open file descriptor that refers to a device. The request argument selects the control function to be performed and depends on the device being addressed. The arg argument represents a third argument that has additional information that is needed by this specific device to perform the requested function. The data type of arg depends upon the particular control request, but it is either an int or a pointer to a device-specific data structure.

In addition to device-specific and STREAMS functions, generic functions are provided by more than one device driver (for example, the general terminal interface.) See termio(7I).

RETURN VALUES
Upon successful completion, the value returned depends upon the device control function, but must be a non-negative integer. Otherwise, −1 is returned and errno is set to indicate the error.

ERRORS
The ioctl() function will fail for any type of file if:

EBADF  The fildes argument is not a valid open file descriptor.
EINTR  A signal was caught during the execution of the ioctl() function.
EINVAL  The STREAM or multiplexer referenced by fildes is linked (directly or indirectly) downstream from a multiplexer.

The ioctl() function will also fail if the device driver detects an error. In this case, the error is passed through ioctl() without change to the caller. A particular driver might not have all of the following error cases. Under the following conditions, requests to device drivers may fail and set errno to indicate the error

EFAULT  The request argument requires a data transfer to or from a buffer pointed to by arg, but arg points to an illegal address.
EINVAL  The request or arg argument is not valid for this device.
EIO    Some physical I/O error has occurred.
ENOLINK  The fildes argument is on a remote machine and the link to that machine is no longer active.

ENOTTY  The fildes argument is not associated with a STREAMS device that accepts control functions.

ENXIO  The request and arg arguments are valid for this device driver, but the service requested can not be performed on this particular subdevice.

ENODEV  The fildes argument refers to a valid STREAMS device, but the corresponding device driver does not support the ioctl() function.

STREAMS errors are described in streamio(7I).

SEE ALSO  streamio(7I), termio(7I)
issetugid(2)

NAME
issetugid – determine if current executable is running setuid or setgid

SYNOPSIS
#include <unistd.h>

int issetugid(void);

DESCRIPTION
The issetugid() function enables library functions (in libtermlib, libc, or other libraries) to guarantee safe behavior when used in setuid or setgid programs. Some library functions might be passed insufficient information and not know whether the current program was started setuid or setgid because a higher level calling code might have made changes to the uid, euid, gid, or egid. These low-level library functions are therefore unable to determine if they are being run with elevated or normal privileges.

The issetugid() function should be used to determine if a path name returned from a getenv(3C) call can be used safely to open the specified file. It is often not safe to open such a file because the status of the effective uid is not known.

The result of a call to issetugid() is unaffected by calls to setuid(), setgid(), or other such calls. In case of a call to fork(2), the child process inherits the same status.

The status of issetugid() is affected only by execve() (see exec(2)). If a child process executes a new executable file, a new issetugid() status will be based on the existing process’s uid, euid, gid, and egid permissions and on the modes of the executable file. If the new executable file modes are setuid or setgid, or if the existing process is executing the new image with uid!=euid or gid!=egid, issetugid() will return 1 in the new process.

RETURN VALUES
The issetugid() function returns 1 if the process was made setuid or setgid as the result of the last or a previous call to execve(). Otherwise it returns 0.

ERRORS
The issetugid() function is always successful. No return value is reserved to indicate an error.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Evolving</td>
</tr>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
exec(2), fork(2), setuid(2), getenv(3C), attributes(5)
**NAME**
kill – send a signal to a process or a group of processes

**SYNOPSIS**
```c
#include <sys/types.h>
#include <signal.h>

int kill(pid_t pid, int sig);
```

**DESCRIPTION**
The `kill()` function sends a signal to a process or a group of processes. The process or group of processes to which the signal is to be sent is specified by `pid`. The signal that is to be sent is specified by `sig` and is either one from the list given in `signal(3HEAD)`, or 0. If `sig` is 0 (the null signal), error checking is performed but no signal is actually sent. This can be used to check the validity of `pid`.

The real or effective user ID of the sending process must match the real or saved (from one of functions in the `exec` family, see `exec(2)`) user ID of the receiving process unless the effective user ID of the sending process is superuser, (see `intro(2)`), or `sig` is `SIGCONT` and the sending process has the same session ID as the receiving process.

If `pid` is greater than 0, `sig` will be sent to the process whose process ID is equal to `pid`.

If `pid` is negative but not `(pid_t)−1`, `sig` will be sent to all processes whose process group ID is equal to the absolute value of `pid` and for which the process has permission to send a signal.

If `pid` is 0, `sig` will be sent to all processes excluding special processes (see `intro(2)`) whose process group ID is equal to the process group ID of the sender.

If `pid` is `(pid_t)−1` and the effective user ID of the sender is not super-user, `sig` will be sent to all processes excluding special processes whose real user ID is equal to the effective user ID of the sender.

If `pid` is `(pid_t)−1` and the effective user ID of the sender is super-user, `sig` will be sent to all processes excluding special processes.

**RETURN VALUES**
Upon successful completion, 0 is returned. Otherwise, −1 is returned, no signal is sent, and `errno` is set to indicate the error.

**ERRORS**
The `kill()` function will fail if:

- **EINVAL** The `sig` argument is not a valid signal number.
- **EPERM** The `sig` argument is `SIGKILL` and the `pid` argument is `(pid_t)1` (that is, the calling process does not have permission to send the signal to any of the processes specified by `pid`); or the effective user of the calling process does not match the real or saved user and is not super-user, and the calling process is not sending `SIGCONT` to a process that shares the same session ID.
- **ESRCH** No process or process group can be found corresponding to that specified by `pid`.

**USAGE**
The `sigsend(2)` function provides a more versatile way to send signals to processes.
kill(2)

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
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</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
kill(1), intro(2), exec(2), getpid(2), getsid(2), setpgid(2), sigaction(2), sigsend(2), signal(3C), attributes(5), signal(3HEAD)
#include <unistd.h>

int link(const char *existing, const char *new);

The `link()` function creates a new link (directory entry) for the existing file and increments its link count by one. The `existing` argument points to a path name naming an existing file. The `new` argument points to a pathname naming the new directory entry to be created.

To create hard links, both files must be on the same file system. Both the old and the new link share equal access and rights to the underlying object. The super-user may make multiple links to a directory. Unless the caller is the super-user, the file named by `existing` must not be a directory.

Upon successful completion, `link()` marks for update the `st_ctime` field of the file. Also, the `st_ctime` and `st_mtime` fields of the directory that contains the new entry are marked for update.

Upon successful completion, 0 is returned. Otherwise, -1 is returned, no link is created, and `errno` is set to indicate the error.

The `link()` function will fail if:

- `EACCES` A component of either path prefix denies search permission, or the requested link requires writing in a directory with a mode that denies write permission.
- `EDQUOT` The directory where the entry for the new link is being placed cannot be extended because the user’s quota of disk blocks on that file system has been exhausted.
- `EEXIST` The link named by `new` exists.
- `EFAULT` The `existing` or `new` argument points to an illegal address.
- `EINTR` A signal was caught during the execution of the `link()` function.
- `ELOOP` Too many symbolic links were encountered in translating `path`.
- `EMLINK` The maximum number of links to a file would be exceeded.
- `ENAMETOOLONG` The length of the `existing` or `new` argument exceeds `PATH_MAX`, or the length of a `existing` or `new` component exceeds `NAME_MAX` while `_POSIX_NO_TRUNC` is in effect.
The existing or new argument is a null pathname; a component of either path prefix does not exist; or the file named by existing does not exist.

The existing or new argument points to a remote machine and the link to that machine is no longer active.

The directory that would contain the link cannot be extended.

A component of either path prefix is not a directory.

The file named by existing is a directory and the effective user of the calling process is not super-user.

The requested link requires writing in a directory on a read-only file system.

The link named by new and the file named by existing are on different logical devices (file systems).

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

See also symlink(2), unlink(2), attributes(5)
NAME   llseek – move extended read/write file pointer

SYNOPSIS  #include <sys/types.h>
#include <unistd.h>

offset_t llseek(int fildes, offset_t offset, int whence);

DESCRIPTION  The llseek() function sets the 64-bit extended file pointer associated with the open file descriptor specified by fildes as follows:

- If whence is SEEK_SET, the pointer is set to offset bytes.
- If whence is SEEK_CUR, the pointer is set to its current location plus offset.
- If whence is SEEK_END, the pointer is set to the size of the file plus offset.

Although each file has a 64-bit file pointer associated with it, some existing file system types (such as tmpfs) do not support the full range of 64-bit offsets. In particular, on such file systems, non-device files remain limited to offsets of less than two gigabytes. Device drivers may support offsets of up to 1024 gigabytes for device special files.

Some devices are incapable of seeking. The value of the file pointer associated with such a device is undefined.

RETURN VALUES  Upon successful completion, llseek() returns the resulting pointer location as measured in bytes from the beginning of the file. Remote file descriptors are the only ones that allow negative file pointers. Otherwise, −1 is returned, the file pointer remains unchanged, and errno is set to indicate the error.

ERRORS  The llseek() function will fail if:

EBADF      The fildes argument is not an open file descriptor.

EINVAL     The whence argument is not SEEK_SET, SEEK_CUR, or SEEK_END; the offset argument is not a valid offset for this file system type; or the fildes argument is not a remote file descriptor and the resulting file pointer would be negative.

ESPIPE     The fildes argument is associated with a pipe or FIFO.

SEE ALSO  creat(2), dup(2), fcntl(2), lseek(2), open(2)
#include <sys/types.h>
#include <unistd.h>

off_t lseek(int fildes, off_t offset, int whence);

The `lseek()` function sets the file pointer associated with the open file descriptor specified by `fildes` as follows:

- If `whence` is `SEEK_SET`, the pointer is set to `offset` bytes.
- If `whence` is `SEEK_CUR`, the pointer is set to its current location plus `offset`.
- If `whence` is `SEEK_END`, the pointer is set to the size of the file plus `offset`.

The symbolic constants `SEEK_SET`, `SEEK_CUR`, and `SEEK_END` are defined in the header `<unistd.h>`.

Some devices are incapable of seeking. The value of the file pointer associated with such a device is undefined.

The `lseek()` function allows the file pointer to be set beyond the existing data in the file. If data are later written at this point, subsequent reads in the gap between the previous end of data and the newly written data will return bytes of value 0 until data are written into the gap.

If `fildes` is a remote file descriptor and `offset` is negative, `lseek()` returns the file pointer even if it is negative. The `lseek()` function will not, by itself, extend the size of a file.

Upon successful completion, the resulting offset, as measured in bytes from the beginning of the file, is returned. Otherwise, `(off_t)-1` is returned, the file offset remains unchanged, and `errno` is set to indicate the error.

The `lseek()` function will fail if:

- `EBADF` The `fildes` argument is not an open file descriptor.
- `EINVAL` The `whence` argument is not `SEEK_SET`, `SEEK_CUR`, or `SEEK_END`; or the `fildes` argument is not a remote file descriptor and the resulting file pointer would be negative.
- `EOVERFLOW` The resulting file offset would be a value which cannot be represented correctly in an object of type `off_t` for regular files.
- `ESPIPE` The `fildes` argument is associated with a pipe, a FIFO, or a socket.

In multithreaded applications, using `lseek()` in conjunction with a `read(2)` or `write(2)` call on a file descriptor shared by more than one thread is not an atomic operation. To ensure atomicity, use `pread()` or `pwrite()`.
ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
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<tbody>
<tr>
<td>MT-Level</td>
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</tbody>
</table>

SEE ALSO
creat(2), dup(2), fcntl(2), open(2), read(2), write(2), attributes(5), l64(5)
# NAME
\_lwp\_cond\_signal, \_lwp\_cond\_broadcast – signal a condition variable

## SYNOPSIS

```
#include <sys/lwp.h>

int \_lwp\_cond\_signal(lwp\_cond\_t *cvp);
int \_lwp\_cond\_broadcast(lwp\_cond\_t *cvp);
```

## DESCRIPTION

The \_lwp\_cond\_signal() function unblocks one LWP that is blocked on the LWP condition variable pointed to by \( cvp \).

The \_lwp\_cond\_broadcast() function unblocks all LWPs that are blocked on the LWP condition variable pointed to by \( cvp \).

If no LWPs are blocked on the LWP condition variable, then \_lwp\_cond\_signal() and \_lwp\_cond\_broadcast() have no effect.

Both functions should be called under the protection of the same LWP mutex lock that is used with the LWP condition variable being signaled. Otherwise, the condition variable may be signalled between the test of the associated condition and blocking in \_lwp\_cond\_wait(). This can cause an infinite wait.

## RETURN VALUES

Upon successful completion, 0 is returned. A non-zero value indicates an error.

## ERRORS

The \_lwp\_cond\_signal() and \_lwp\_cond\_broadcast() functions will fail if:

- \texttt{EINVAL} The \( cvp \) argument points to an invalid LWP condition variable.
- \texttt{EFAULT} The \( cvp \) argument points to an invalid address.

## SEE ALSO

- \_lwp\_cond\_wait(2), \_lwp\_mutex\_lock(2)
These functions are used to wait for the occurrence of a condition represented by an LWP condition variable. LWP condition variables must be initialized to 0 before use.

The _lwp_cond_wait() function atomically releases the LWP mutex pointed to by mp and causes the calling LWP to block on the LWP condition variable pointed to by cvp. The blocked LWP may be awakened by _lwp_cond_signal(2), _lwp_cond_broadcast(2), or when interrupted by delivery of a signal. Any change in value of a condition associated with the condition variable cannot be inferred by the return of _lwp_cond_wait() and any such condition must be re-evaluated.

The _lwp_cond_timedwait() function is similar to _lwp_cond_wait(), except that the calling LWP will not block past the time of day specified by abstime. If the time of day becomes greater than abstime, _lwp_cond_timedwait() returns with the error code ETIME.

The _lwp_cond_reltimedwait() function is similar to _lwp_cond_wait(), except that the calling LWP will not block past the relative time specified by reltime. If the time of day becomes greater than the starting time of day plus reltime, _lwp_cond_reltimedwait() returns with the error code ETIME.

The _lwp_cond_wait(), _lwp_cond_timedwait(), and _lwp_cond_reltimedwait() functions always return with the mutex locked and owned by the calling lightweight process.

RETURN VALUES
Upon successful completion, 0 is returned. A non-zero value indicates an error.

ERRORS
If any of the following conditions are detected, _lwp_cond_wait(), _lwp_cond_timedwait(), and _lwp_cond_reltimedwait() fail and return the corresponding value:

EINVAL
The cvp argument points to an invalid LWP condition variable or the mp argument points to an invalid LWP mutex.

EFAULT
The mp, cvp, or abstime argument points to an illegal address.

If any of the following conditions occur, _lwp_cond_wait(), _lwp_cond_timedwait(), and _lwp_cond_reltimedwait() fail and return the corresponding value:
_lwp_cond_wait(2)

EINTR The call was interrupted by a signal or fork(2).

If any of the following conditions occur, _lwp_cond_timedwait() and
_lwp_cond_reltimedwait() fail and return the corresponding value:

ETIME The time specified in abstime or reltime has passed.

EXAMPLES

EXAMPLE 1 Use the _lwp_cond_wait() function in a loop testing some condition.

The _lwp_cond_wait() function is normally used in a loop testing some condition,
as follows:

```c
lwp_mutex_t m;
lwp_cond_t cv;
int cond;
(void) _lwp_mutex_lock(&m);
while (cond == FALSE) {
    (void) _lwp_cond_wait(&cv, &m);
}
(void) _lwp_mutex_unlock(&m);
```

EXAMPLE 2 Use the _lwp_cond_timedwait() function in a loop testing some condition.

The _lwp_cond_timedwait() function is also normally used in a loop testing some
condition. It uses an absolute timeout value as follows:

```c
timestruc_t to;
lwp_mutex_t m;
lwp_cond_t cv;
int cond, err;
(void) _lwp_mutex_lock(&m);
to.tv_sec = time(NULL) + TIMEOUT;
to.tv_nsec = 0;
while (cond == FALSE) {
    err = _lwp_cond_timedwait(&cv, &m, &to);
    if (err == ETIME) {
        /* timeout, do something */
        break;
    }
}
(void) _lwp_mutex_unlock(&m);
```

This example sets a bound on the total wait time even though the
_lwp_cond_timedwait() may return several times due to the condition being
signalled or the wait being interrupted.

EXAMPLE 3 Use the _lwp_cond_reltimedwait() function in a loop testing some
condition.

The _lwp_cond_reltimedwait() function is also normally used in a loop testing
some condition. It uses a relative timeout value as follows:

```c
timestruc_t to;
lwp_mutex_t m;
```
EXAMPLE 3 Use the _lwp_cond_reltimedwait() function in a loop testing some condition.  

Continued

```c
lwp_cond_t cv;
int cond, err;
(void) _lwp_mutex_lock(&m);
while (cond == FALSE) {
    to.tv_sec = TIMEOUT;
    to.tv_nsec = 0;
    err = _lwp_cond_reltimedwait(&cv, &m, &t);
    if (err == ETIME) {
        /* timeout, do something */
        break;
    }
}
(void) _lwp_mutex_unlock(&m);
```

SEE ALSO

_lwp_cond_broadcast(2), _lwp_cond_signal(2), _lwp_kill(2),
_lwp_mutex_lock(2), fork(2), kill(2)
The \_lwp\_create() function adds a lightweight process (LWP) to the current process. The \context \_p argument specifies the initial signal mask, stack, and machine context (including the program counter and stack pointer) for the new LWP. The new LWP inherits the scheduling class and priority of the caller.

If \_lwp\_create() is successful and \new \_lwp \_p is not NULL, the ID of the new LWP is stored in the location pointed to by \new \_lwp\_p.

The flags argument specifies additional attributes for the new LWP. The value in flags is constructed by the bitwise inclusive OR operation of the following values:

- **LWP\_DETACHED**: The LWP is created detached.
- **LWP\_DAEMON**: The LWP is created as a daemon LWP.
- **LWP\_SUSPENDED**: The LWP is created suspended.

If LWP\_DETACHED or LWP\_DAEMON is specified, then the LWP is created in the detached state. Otherwise the LWP is created in the undetached state. The ID (and system resources) associated with a detached LWP can be automatically reclaimed when the LWP exits. The ID of an undetached LWP cannot be reclaimed until it exits and another LWP has reported its termination by way of \_lwp\_wait(2). This allows the waiting LWP to determine that the waited for LWP has terminated and to reclaim any process resources that it was using.

If LWP\_DAEMON is specified, then in addition to being created in the detached state, the LWP is created as a daemon LWP. Daemon LWPs do not interfere with the exit conditions for a process. A process will exit as though \_exit(0) \ had been called when the last non-daemon LWP calls \_lwp\_exit() (see \exit(2) and \_lwp\_exit(2)). Also, an LWP that is waiting in \_lwp\_wait(2) for any LWP to terminate will return EDEADLK when all remaining LWPs in the process are either daemon LWPs or other LWPs waiting in \_lwp\_wait().

If LWP\_SUSPENDED is specified, then the LWP is created in a suspended state. This allows the creator to change the LWP's inherited attributes before it starts to execute. The suspended LWP can only be resumed by way of \_lwp\_continue(2). If LWP\_SUSPENDED is not specified the LWP can begin to run immediately after it has been created.

Upon successful completion, 0 is returned. A non-zero value indicates an error.

If any of the following conditions are detected, \_lwp\_create() fails and returns the corresponding value:

- **EFAULT**: Either the context parameter or the new \_lwp parameter point to invalid addresses.
EAGAIN  A system limit is exceeded, (for example, too many LWPs were created for this real user ID).

EINVAL  The flags argument contains values other than those specified above.

EXAMPLES

EXAMPLE 1 How a stack is allocated to a new LWP.

This example shows how a stack is allocated to a new LWP. The _lwp_makecontext() function is used to set up the context parameter so that the new LWP begins executing a function.

```c
contextp = (ucontext_t *)malloc(sizeof(ucontext_t));
stackbase = malloc(stacksize);
_lwp_makecontext(contextp, func, arg, private, stackbase, stacksize);
sigprocmask(SIGSETMASK, NULL, &contextp->uc_sigmask);
error = _lwp_create(contextp, NULL, &new_lwp);
```

APPLICATIONS Applications should use bound threads rather than the _lwp_* functions (see thr_create(3THR)). Using LWPs directly is not advised because libraries are only safe to use with threads, not LWPs.

ATTRIBUTES See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
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</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
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<tr>
<td>MT-Level</td>
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</tbody>
</table>

SEE ALSO _lwp_cond_timedwait(2), _lwp_continue(2), _lwp_detach(2), _lwp_exit(2), _lwp_makecontext(2), _lwp_wait(2), alarm(2), exit(2), poll(2), signal(3HEAD), sleep(3C), thr_create(3THR), ucontext(3HEAD), attributes(5)

NOTES The _lwp_create() function is obsolete and will be removed in a future release.
The `_lwp_detach()` function marks the LWP specified by `target_lwp` as being a detached LWP. The effect is the same as if `target_lwp` had been created using the `LWP_DETACHED` flag (see `_lwp_create(2)`).

The `target_lwp` must be a non-detached LWP within the same process as the calling LWP.

Upon successful completion, 0 is returned. A non-zero value indicates an error.

If any of the following conditions occur, `_lwp_detach()` fails and returns the corresponding value:

- **EINVAL** The LWP with the ID specified by `target_lwp` is already detached.
- **ESRCH** No LWP with the ID specified by `target_lwp` can be found in the current process.

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
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</table>

The `_lwp_detach()` function is obsolete and will be removed in a future release.
The `_lwp_exit()` function causes the calling LWP to terminate. If it is the last non-daemon LWP in the process, the process exits with a status of 0 (see `exit(2)`).

If the LWP was created undetached, it is transformed into a “zombie LWP” that retains at least the LWP’s ID until it is waited for (see `_lwp_wait(2)`). Otherwise, its ID and system resources may be reclaimed immediately.

See attributes(5) for descriptions of the following attributes:

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</table>

The `_lwp_exit()` function is obsolete and will be removed in a future release.
The `_lwp_info()` function fills the `lwpinfo` structure pointed to by `buffer` with time-accounting information pertaining to the calling LWP. This call may be extended in the future to return other information to the `lwpinfo` structure as needed. The `lwpinfo` structure in `<sys/lwp.h>` includes the following members:

- `timestruc_t lwp_utime;`
- `timestruc_t lwp_stime;`

The `lwp_utime` member is the CPU time used while executing instructions in the user space of the calling LWP.

The `lwp_stime` member is the CPU time used by the system on behalf of the calling LWP.

Upon successful completion, `_lwp_info()` returns 0 and fills in the `lwpinfo` structure pointed to by `buffer`.

If the following condition is detected, `_lwp_info()` returns the corresponding value:

- `EFAULT` The `buffer` argument points to an illegal address.

Additionally, the `_lwp_info()` function will fail for 32-bit interfaces if:

- `EOVERFLOW` The size of the `tv_sec` member of the `timestruc_t` type pointed to by `lwp_utime` and `lwp_stime` is too small to contain the correct number of seconds.

See attributes(5) for descriptions of the following attributes:

<table>
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</table>

See also `times(2), attributes(5)`
NAME
_lwp_kill – send a signal to a LWP

SYNOPSIS
#include <sys/lwp.h>
#include <signal.h>

int _lwp_kill(lwpid_t target_lwp, int sig);

DESCRIPTION
The _lwp_kill() function sends a signal to the LWP specified by target_lwp. The signal that is to be sent is specified by sig and must be one from the list given in signal(3HEAD). If sig is 0 (the null signal), error checking is performed but no signal is actually sent. This can be used to check the validity of target_lwp.

The target_lwp must be an LWP within the same process as the calling LWP.

RETURN VALUES
Upon successful completion, 0 is returned. A non-zero value indicates an error.

ERRORS
If any of the following conditions occur, _lwp_kill() fails and returns the corresponding value:

EINVAL The sig argument is not a valid signal number.
ESRCH The target_lwp argument cannot be found in the current process.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

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SEE ALSO
kill(2), sigaction(2), sigprocmask(2), signal(3HEAD), attributes(5)
_lwp_makecontext

NAME
_lwp_makecontext – initialize an LWP context

SYNOPSIS
#include <sys/types.h>
#include <sys/lwp.h>
#include <ucontext.h>

void _lwp_makecontext(ucontext_t *ucp, void (*start_routine)(void *),
void *arg, void *private, caddr_t stack_base, size_t stack_size);

DESCRIPTION
The _lwp_makecontext() function initializes the user context structure pointed to
by ucp. The user context is defined by ucontext(3HEAD). The resulting user context
can be used by _lwp_create(2) for specifying the initial state of the new LWP. The
user context is set up to start executing the function start_routine with a single
argument, arg, and to call _lwp_exit(2) if start_routine returns. The new LWP will use
the storage starting at stack_base and continuing for stack_size bytes as an execution
stack. The initial value in LWP-private memory will be set to private (see
_lwp_setprivate(2)). The signal mask in the user context is not initialized.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
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<tbody>
<tr>
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</table>

SEE ALSO
_lwp_create(2), _lwp_exit(2), _lwp_setprivate(2), ucontext(3HEAD), attributes(5)

NOTES
The _lwp_makecontext() function is obsolete and will be removed in a future
release.
NAME

_lwp_mutex_lock, _lwp_mutex_unlock, _lwp_mutex_trylock – mutual exclusion

SYNOPSIS

#include <sys/lwp.h>

int _lwp_mutex_lock(lwp_mutex_t *mp);
int _lwp_mutex_trylock(lwp_mutex_t *mp);
int _lwp_mutex_unlock(lwp_mutex_t *mp);

DESCRIPTION

These functions serialize the execution of lightweight processes. They are useful for ensuring that only one lightweight process can execute a critical section of code at any one time (mutual exclusion). LWP mutexes must be initialized to 0 before use.

The _lwp_mutex_lock() function locks the LWP mutex pointed to by mp. If the mutex is already locked, the calling LWP blocks until the mutex becomes available. When _lwp_mutex_lock() returns, the mutex is locked and the calling LWP is the "owner".

The _lwp_mutex_trylock() function attempts to lock the mutex. If the mutex is already locked it returns with an error. If the mutex is unlocked, it is locked and _lwp_mutex_trylock() returns.

The _lwp_mutex_unlock() function unlocks a locked mutex. The mutex must be locked and the calling LWP must be the one that last locked the mutex (the owner). If any other LWPs are waiting for the mutex to become available, one of them is unblocked.

RETURN VALUES

Upon successful completion, 0 is returned. A non-zero value indicates an error.

ERRORS

If any of the following conditions are detected, _lwp_mutex_lock(), _lwp_mutex_trylock(), and _lwp_mutex_unlock() fail and return the corresponding value:

EINVAL

The mp argument points to an invalid LWP mutex.

EFAULT

The mp argument points to an illegal address.

If any of the following conditions occur, _lwp_mutex_trylock() fails and returns the corresponding value:

EBUSY

The mp argument points to a locked mutex.

SEE ALSO

intro(2), _lwp_cond_wait(2)
_lwp_self(2)

NAME       _lwp_self – get LWP identifier

SYNOPSIS   #include <sys/lwp.h>

            lwpid_t _lwp_self(void);

DESCRIPTION The _lwp_self() function returns the ID of the calling LWP.

ATTRIBUTES See attributes(5) for descriptions of the following attributes:

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</table>

SEE ALSO     _lwp_create(2), attributes(5)
# lwp_sema_wait(2)

**NAME**

_lwp_sema_wait, _lwp_sema_trywait, _lwp_sema_init, _lwp_sema_post – semaphore operations

**SYNOPSIS**

```c
#include <sys/lwp.h>

int _lwp_sema_wait(lwp_sema_t *sema);
int _lwp_sema_trywait(lwp_sema_t *sema);
int _lwp_sema_init(lwp_sema_t *sema, int count);
int _lwp_sema_post(lwp_sema_t *sema);
```

**DESCRIPTION**

Conceptually, a semaphore is an non-negative integer count that is atomically incremented and decremented. Typically this represents the number of resources available. The _lwp_sema_init() function initializes the count, _lwp_sema_post() atomically increments the count, and _lwp_sema_wait() waits for the count to become greater than 0 and then atomically decrements it.

LWP semaphores must be initialized before use. The _lwp_sema_init() function initializes the count associated with the LWP semaphore pointed to by *sema to count.

The _lwp_sema_wait() function blocks the calling LWP until the semaphore count becomes greater than 0 and then atomically decrements it.

The _lwp_sema_trywait() function atomically decrements the count if it is greater than zero. Otherwise it returns an error.

The _lwp_sema_post() function atomically increments the semaphore count. If there are any LWPs blocked on the semaphore, one is unblocked.

**RETURN VALUES**

Upon successful completion, 0 is returned. A non-zero value indicates an error.

**ERRORS**

The _lwp_sema_init(), _lwp_sema_trywait(), _lwp_sema_wait(), and _lwp_sema_post() functions will fail if:

- **EINVAL** The *sema argument points to an invalid semaphore.
- **EFAULT** The *sema argument points to an illegal address.

The _lwp_sema_wait() function will fail if:

- **EINTR** The function execution was interrupted by a signal or fork(2).

The _lwp_sema_trywait() function will fail if:

- **EBUSY** The function was called on a semaphore with a zero count.

The _lwp_sema_post() function will fail if:

- **EOVERFLOW** The value of the *sema argument exceeds SEM_VALUE_MAX.

**SEE ALSO**

fork(2)
NAME
_lwp_setprivate, _lwp_getprivate – set or get LWP specific storage

SYNOPSIS
#include <sys/lwp.h>

void _lwp_setprivate(void *buffer);
void *_lwp_getprivate(void);

DESCRIPTION
The _lwp_setprivate() function stores the value specified by buffer in LWP-private
memory that is unique to the calling LWP. This is typically used by thread library
implementations to maintain a pointer to information about the thread currently
running on the calling LWP.

The _lwp_getprivate() function returns the value stored in LWP-private memory.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

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</tbody>
</table>

SEE ALSO
_lwp_makecontext(2), attributes(5)

NOTES
The _lwp_setprivate() and _lwp_getprivate() functions are obsolete and will
be removed in a future release.
The `_lwp_suspend()` function immediately suspends the execution of the LWP specified by `target_lwp`. On successful return from `_lwp_suspend()`, `target_lwp` is no longer executing. Once a thread is suspended, subsequent calls to `_lwp_suspend()` have no affect.

The `_lwp_continue()` function resumes the execution of a suspended LWP. Once a suspended LWP is continued, subsequent calls to `_lwp_continue()` have no effect.

A suspended LWP will not be awakened by a signal. The signal stays pending until the execution of the LWP is resumed by `_lwp_continue()`.

Upon successful completion, 0 is returned. A non-zero value indicates an error.

If the following condition occurs, `_lwp_suspend()` and `_lwp_continue()` fail and return the corresponding value:

**ESRCH**

The `target_lwpid` argument cannot be found in the current process.

**ATTRIBUTES**

See `attributes(5)` for descriptions of the following attributes:

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</table>

**SEE ALSO**

`_lwp_create(2)`, `attributes(5)`
The `_lwp_wait()` function blocks the current LWP until the LWP specified by `wait_for` terminates. If the specified LWP terminated prior to the call to `_lwp_wait()`, `_lwp_wait()` returns immediately. If `wait_for` is zero, `_lwp_wait()` waits for any undetached LWP in the current process. If `wait_for` is not zero, it must specify an undetached LWP in the current process. If `departed_lwp` is not NULL, it points to a location where the ID of the exited LWP is stored (see `_lwp_exit(2)`).

When an LWP exits and there are one or more LWPs in the process waiting for this specific LWP to exit, one of the waiting LWPs is unblocked and it returns from `_lwp_wait()` successfully. Any other LWPs waiting for this same LWP to exit are also unblocked, but they return from `_lwp_wait()` with an error (ESRCH) indicating the waited-for LWP no longer exists. If there are no LWPs in the process waiting for this specific LWP to exit but there are one or more LWPs waiting for any LWP to exit, one of the waiting LWPs is unblocked and it returns from `_lwp_wait()` successfully.

If an LWP is waiting for any LWP to exit, it blocks until an undetached LWP for which no other LWP is waiting terminates, at which time it returns successfully, or until all other LWPs in the process are either daemon LWPs or LWPs waiting in `_lwp_wait()`, in which case it returns EDEADLK.

The ID of an LWP that has exited may be reused via `_lwp_create()` after the LWP has been successfully waited for.

Upon successful completion, 0 is returned. A non-zero value indicates an error.

If any of the following conditions occur, `_lwp_wait()` fails and returns the corresponding value:

- **EDEADLK**: A wait deadlock was detected, such as when an LWP attempts to wait for itself, or the calling LWP is waiting for any LWP to exit and only daemon LWPs or waiting LWPs exist in the process.
- **EINTR**: The `_lwp_wait()` function was interrupted by a signal.
- **EINVAL**: The LWP with the ID specified by `wait_for` is a detached LWP.
- **ESRCH**: No LWP with the ID specified by `wait_for` can be found in the current process.

See attributes(5) for descriptions of the following attributes:

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\_lwp\_wait(2)

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SEE ALSO  \_lwp\_create(2), \_lwp\_detach(2), \_lwp\_exit(2), attributes(5)

NOTES The \_lwp\_wait() function is obsolete and will be removed in a future release.
The `memcntl()` function allows the calling process to apply a variety of control operations over the address space identified by the mappings established for the address range \([addr, addr + len]\).

The `addr` argument must be a multiple of the pagesize as returned by `sysconf(3C)`. The scope of the control operations can be further defined with additional selection criteria (in the form of attributes) according to the bit pattern contained in `attr`.

The following attributes specify page mapping selection criteria:

- **SHARED**
  - Page is mapped shared.
- **PRIVATE**
  - Page is mapped private.

The following attributes specify page protection selection criteria. The selection criteria are constructed by a bitwise OR operation on the attribute bits and must match exactly.

- **PROT_READ**
  - Page can be read.
- **PROT_WRITE**
  - Page can be written.
- **PROT_EXEC**
  - Page can be executed.

The following criteria may also be specified:

- **PROC_TEXT**
  - Process text.
- **PROC_DATA**
  - Process data.

The `PROC_TEXT` attribute specifies all privately mapped segments with read and execute permission, and the `PROC_DATA` attribute specifies all privately mapped segments with write permission.

Selection criteria can be used to describe various abstract memory objects within the address space on which to operate. If an operation shall not be constrained by the selection criteria, `attr` must have the value 0.

The operation to be performed is identified by the argument `cmd`. The symbolic names for the operations are defined in `<sys/mman.h>` as follows:

- **MC_LOCK**
  - Lock in memory all pages in the range with attributes `attr`. A given page may be locked multiple times through different mappings; however, within a given mapping, page locks do not nest. Multiple lock operations on the same address in the same process will all be removed with a single unlock operation. A page locked in one
process and mapped in another (or visible through a different mapping in the locking process) is locked in memory as long as the locking process does neither an implicit nor explicit unlock operation. If a locked mapping is removed, or a page is deleted through file removal or truncation, an unlock operation is implicitly performed. If a writable MAP_PRIVATE page in the address range is changed, the lock will be transferred to the private page.

The arg argument is not used, but must be 0 to ensure compatibility with potential future enhancements.

**MC_LOCKAS**

Lock in memory all pages mapped by the address space with attributes attr. The addr and len arguments are not used, but must be NULL and 0 respectively, to ensure compatibility with potential future enhancements. The arg argument is a bit pattern built from the flags:

- **MCL_CURRENT** Lock current mappings.
- **MCL_FUTURE** Lock future mappings.

The value of arg determines whether the pages to be locked are those currently mapped by the address space, those that will be mapped in the future, or both. If MCL_FUTURE is specified, then all mappings subsequently added to the address space will be locked, provided sufficient memory is available.

**MC_SYNC**

Write to their backing storage locations all modified pages in the range with attributes attr. Optionally, invalidate cache copies. The backing storage for a modified MAP_SHARED mapping is the file the page is mapped to; the backing storage for a modified MAP_PRIVATE mapping is its swap area. The arg argument is a bit pattern built from the flags used to control the behavior of the operation:

- **MS_ASYNC** Perform asynchronous writes.
- **MS_SYNC** Perform synchronous writes.
- **MS_INVALIDATE** Invalidate mappings.

**MS_ASYNC** Return immediately once all write operations are scheduled; with **MS_SYNC** the function will not return until all write operations are completed.

**MS_INVALIDATE** Invalidate all cached copies of data in memory, so that further references to the pages will be obtained by the system from their backing storage locations. This operation should be used by applications that require a memory object to be in a known state.
**MC_UNLOCK**  Unlock all pages in the range with attributes `attr`. The `arg` argument is not used, but must be 0 to ensure compatibility with potential future enhancements.

**MC_UNLOCKAS**  Remove address space memory locks and locks on all pages in the address space with attributes `attr`. The `addr`, `len`, and `arg` arguments are not used, but must be `NULL`, 0 and 0, respectively, to ensure compatibility with potential future enhancements.

**MC_HAT_ADVISE**  Advise system how a region of user-mapped memory will be accessed. The `arg` argument is interpreted as a "`struct memcntl_mha *`". The following members are defined in a `struct memcntl_mha`:

- `uint_t mha_cmd;`
- `uint_t mha_flags;`
- `size_t mha_pagesize;`

The accepted values for `mha_cmd` are:

- `MHA_MPSIZE VA`
- `MHA_MPSIZE STACK`
- `MHA_MPSIZE BSSBRK`

The `mha_flags` member is reserved for future use and must always be set to 0. The `mha_pagesize` member must be a valid size as obtained from `getpagesize(3C)` or the constant value 0 to allow the system to choose an appropriate hardware address translation mapping size.

`MHA_MPSIZE VA` sets the preferred hardware address translation mapping size of the region of memory from `addr` to `addr + len`. Both `addr` and `len` must be aligned to an `mha_pagesize` boundary. The entire virtual address region from `addr` to `addr + len` must not have any holes. Permissions within each `mha_pagesize`-aligned portion of the region must be consistent. When a size of 0 is specified, the system selects an appropriate size based on the size and alignment of the memory region, type of processor, and other considerations.

`MHA_MPSIZE STACK` sets the preferred hardware address translation mapping size of the process main thread stack segment. The `addr` and `len` arguments must be `NULL` and 0, respectively.

`MHA_MPSIZE BSSBRK` sets the preferred hardware address translation mapping size of the process heap. The `addr` and `len` arguments must be `NULL` and 0, respectively. See the NOTES section of the `ppgsz(1)` manual page for additional information on process heap alignment.

The `attr` argument must be 0 for all `MC_HAT_ADVISE` operations.
The `mask` argument must be 0; it is reserved for future use.

Locks established with the lock operations are not inherited by a child process after `fork(2)`. The `memcntl()` function fails if it attempts to lock more memory than a system-specific limit.

Due to the potential impact on system resources, all operations except `MC_SYNC` are restricted to processes with superuser effective user ID.

**USAGE**

The `memcntl()` function subsumes the operations of `plock(3C)` and `mctl(3UCB)`.

`MC_HAT_ADVISE` is intended to improve performance of applications that use large amounts of memory on processors that support multiple hardware address translation mapping sizes; however, it should be used with care. Not all processors support all sizes with equal efficiency. Use of larger sizes may also introduce extra overhead that could reduce performance or available memory. Using large sizes for one application may reduce available resources for other applications and result in slower system wide performance.

**RETURN VALUES**

Upon successful completion, `memcntl()` returns 0; otherwise, it returns −1 and sets `errno` to indicate an error.

**ERRORS**

The `memcntl()` function will fail if:

- **EAGAIN** When the selection criteria match, some or all of the memory identified by the operation could not be locked when `MC_LOCK` or `MC_LOCKAS` was specified, some or all mappings in the address range `[addr, addr + len]` are locked for I/O when `MC_HAT_ADVISE` was specified, or the system has insufficient resources when `MC_HAT_ADVISE` was specified.

- **EBUSY** When the selection criteria match, some or all of the addresses in the range `[addr, addr + len]` are locked and `MC_SYNC` with the `MS_INVALIDATE` option was specified.

- **EINVAL** The `addr` argument specifies invalid selection criteria or is not a multiple of the page size as returned by `sysconf(3C)`; the `addr` and/or `len` argument does not have the value 0 when `MC_LOCKAS` or `MC_UNLOCKAS` is specified; the `arg` argument is not valid for the function specified; `mha_pagesize` or `mha_cmd` is invalid; or `MC_HAT_ADVISE` is specified and not all pages in the specified region have the same access permissions within the given size boundaries.

- **ENOMEM** When the selection criteria match, some or all of the addresses in the range `[addr, addr + len]` are invalid for the address space of a process or specify one or more pages which are not mapped.

- **EPERM** The process’s effective user ID is not superuser and `MC_LOCK`, `MC_LOCKAS`, `MC_UNLOCK`, or `MC_UNLOCKAS` was specified.
memcntl(2)

ATTRIBUTES

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>MT-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO

ppgsz(1), fork(2) mmap(2), mprotect(2), getpagesize(3C), mctl(3UCB), mlock(3C), mlockall(3C), msync(3C), plock(3C), sysconf(3C), attributes(5)
meminfo(2)

NAME
meminfo – provide information about memory

SYNOPSIS
#include <sys/types.h>
#include <sys/mman.h>

int meminfo(const uint64_t inaddr[], int addr_count, const uint_t info_req[], int info_count, uint64_t outdata[], uint_t validity[]);

PARAMETERS
inaddr array of input addresses; the maximum number of addresses that can be processed for each call is MAX_MEMINFO_CNT

addr_count number of addresses

info_req array of types of information requested

info_count number of pieces of information requested for each address in inaddr

outdata array into which results are placed; array size must be the product of info_req and addr_count

validity array of size addr_count containing bitwise result codes; 0th bit evaluates validity of corresponding input address, 1st bit validity of response to first member of info_req, and so on

DESCRIPTION
The meminfo() function provides information about virtual and physical memory particular to the calling process. The user or developer of performance utilities can use this information to analyze system memory allocations and develop a better understanding of the factors affecting application performance.

The caller of meminfo() can obtain the following types of information about both virtual and physical memory.

MEMINFO_VPHYSICAL physical address corresponding to virtual address

MEMINFO_VLGRP latency group of physical page corresponding to virtual address

MEMINFO_VPAGESIZE size of physical page corresponding to virtual address

MEMINFO_VREPLCNT number of replicated physical pages corresponding to specified virtual address

MEMINFO_VREPL | n nth physical replica of specified virtual address

MEMINFO_VREPL_LGRP | n lgrp of nth physical replica of specified virtual address

MEMINFO_PLGRP latency group of specified physical address

RETURN VALUES
Upon successful completion meminfo() returns 0. Otherwise −1 is returned and errno is set to indicate the error.

ERRORS
The meminfo() function will fail if:

EFAULT The area pointed to by outdata or validity could not be written, or the data pointed to by info_req or inaddr could not be read.
EINVAL  The value of info_count is greater than 31 or less than 1, or the value of addr_count is less than 1.

EXAMPLES

EXAMPLE 1 Print physical pages and page sizes corresponding to a set of virtual addresses.

The following example prints the physical pages and page sizes corresponding to a set of virtual addresses.

```c
void
print_info(void **addrvec, int how_many)
{
    static const uint_t info[] = {
        MEMINFO_VPHYSICAL,
        MEMINFO_VPAGESIZE
    },
    int info_num = sizeof (info) / sizeof (info[0]);
    int i;
    uint64_t *inaddr = alloca(sizeof (uint64_t) * how_many);
    uint64_t *outdata = alloca(sizeof (uint64_t) * how_many * info_num);
    uint_t *validity = alloca(sizeof (uint_t) * how_many);

    for (i = 0; i < how_many; i++)
        inaddr[i] = (uint64_t)addrvec[i];

    if (meminfo(inaddr, how_many, info, info_num, outdata, validity) < 0) {
        perror("meminfo");
        return;
    }

    for (i = 0; i < how_many; i++) {
        if ((validity[i] & 1) == 0)
            printf("address 0x%llx not part of address space\n", inaddr[i]);
        else if ((validity[i] & 2) == 0)
            printf("address 0x%llx has no physical page "
                "associated with it\n", inaddr[i]);
        else {
            char buff[80];
            if (((validity[i] & 4) == 0)
                strcpy(buff, "<Unknown>"),
            else
                sprintf(buff, "%lld",
                    outdata[i * info_num + 1]);
            printf("address 0x%llx is backed by physical "
                "page 0x%llx of size %s\n",
                    inaddr[i], outdata[i * info_num], buff);
    }
}
```
ATTRIBUTES

See attributes(5) for descriptions of the following attributes:

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</tr>
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<td>Interface Stability</td>
<td>Stable</td>
</tr>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO

mencnt1(2), mmap(2), gethomelgroup(3C), getpagesize(3C), madvise(3C), sysconf(3C), attributes(5)
mincore(2)

NAME
mincore – determine residency of memory pages

SYNOPSIS
#include <sys/types.h>

int mincore(caddr_t addr, size_t len, char *vec);

DESCRIPTION
The mincore() function determines the residency of the memory pages in the
address space covered by mappings in the range [addr, addr + len]. The status is
returned as a character-per-page in the character array referenced by *vec (which the
system assumes to be large enough to encompass all the pages in the address range).
The least significant bit of each character is set to 1 to indicate that the referenced page
is in primary memory, and to 0 to indicate that it is not. The settings of other bits in
each character are undefined and may contain other information in future
implementations.

Because the status of a page can change between the time mincore() checks and
returns the information, returned information might be outdated. Only locked pages
are guaranteed to remain in memory; see mlock(3C).

RETURN VALUES
Upon successful completion, mincore() returns 0. Otherwise, −1 is returned and
errno is set to indicate the error.

ERRORS
The mincore() function will fail if:

EFAULT    The vec argument points to an illegal address.
EINVAL    The addr argument is not a multiple of the page size as returned by
sysconf(3C), or the len argument has a value less than or equal to
0.
ENOMEM    Addresses in the range [addr, addr + len] are invalid for the address
space of a process or specify one or more pages which are not
mapped.

SEE ALSO
mmap(2), mlock(3C), sysconf(3C)
NAME  mkdir – make a directory
SYNOPSIS  
#include <sys/types.h>
#include <sys/stat.h>

int mkdir(const char *path, mode_t mode);

DESCRIPTION  The mkdir() function creates a new directory named by the path name pointed to by path. The mode of the new directory is initialized from mode (see chmod(2) for values of mode). The protection part of the mode argument is modified by the process’s file creation mask (see umask(2)).

The directory’s owner ID is set to the process’s effective user ID. The directory’s group ID is set to the process’s effective group ID, or if the S_ISGID bit is set in the parent directory, then the group ID of the directory is inherited from the parent. The S_ISGID bit of the new directory is inherited from the parent directory.

If path is a symbolic link, it is not followed.

The newly created directory is empty with the exception of entries for itself (.) and its parent directory (..).

Upon successful completion, mkdir() marks for update the st_atime, st_ctime and st_mtime fields of the directory. Also, the st_ctime and st_mtime fields of the directory that contains the new entry are marked for update.

RETURN VALUES  Upon successful completion, 0 is returned. Otherwise, -1 is returned, no directory is created, and errno is set to indicate the error.

ERRORS  The mkdir() function will fail if:

EACCES  Either a component of the path prefix denies search permission or write permission is denied on the parent directory of the directory to be created.

EDQUOT  The directory where the new file entry is being placed cannot be extended because the user’s quota of disk blocks on that file system has been exhausted; the new directory cannot be created because the user’s quota of disk blocks on that file system has been exhausted; or the user’s quota of inodes on the file system where the file is being created has been exhausted.

EEXIST  The named file already exists.

EFAULT  The path argument points to an illegal address.

EINVAL  An attempt was made to create an extended attribute that is a directory.

EIO  An I/O error has occurred while accessing the file system.
mkdir(2)

**ELOOP**
Too many symbolic links were encountered in translating *path*.

**EMLINK**
The maximum number of links to the parent directory would be exceeded.

**ENAMEETOOLONG**
The length of the *path* argument exceeds *PATH_MAX*, or the length of a *path* component exceeds *NAME_MAX* while _POSIX_NO_TRUNC is in effect.

**ENOENT**
A component of the path prefix does not exist or is a null pathname.

**ENOLINK**
The *path* argument points to a remote machine and the link to that machine is no longer active.

**ENOSPC**
No free space is available on the device containing the directory.

**ENOTDIR**
A component of the path prefix is not a directory.

**EROSFS**
The path prefix resides on a read-only file system.

**ATTRIBUTES**
See attributes(5) for descriptions of the following attributes:

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<td>Standard</td>
</tr>
<tr>
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</tr>
</tbody>
</table>

**SEE ALSO**
chmod(2), mknod(2), umask(2), stat(3HEAD), attributes(5)
mknod – make a directory, or a special or ordinary file

#include <sys/stat.h>

int mknod(const char *path, mode_t mode, dev_t dev);

The mknod() function creates a new file named by the path name pointed to by path. The file type and permissions of the new file are initialized from mode.

The file type is specified in mode by the S_IFMT bits, which must be set to one of the following values:

- S_IFIFO: fifo special
- S_IFCHR: character special
- S_IFDIR: directory
- S_IFBLK: block special
- S_IFREG: ordinary file

The file access permissions are specified in mode by the 0007777 bits, and may be constructed by a bitwise OR operation of the following values:

- S_ISUID: 04000 Set user ID on execution.
- S_ISGID: 020#0 Set group ID on execution if # is 7, 5, 3, or 1. Enable mandatory file/record locking if # is 6, 4, 2, or 0
- S_ISVTX: 01000 On directories, restricted deletion flag; on regular files on a UFS file system, do not cache flag.
- S_IRWXU: 00700 Read, write, execute by owner.
- S_IRUSR: 00400 Read by owner.
- S_IWUSR: 00200 Write by owner.
- S_IXUSR: 00100 Execute (search if a directory) by owner.
- S_IRWXG: 00070 Read, write, execute by group.
- S_IRGRP: 00040 Read by group.
- S_IWGRP: 00020 Write by group.
- S_IXGRP: 00010 Execute by group.
- S_IRWXO: 00007 Read, write, execute (search) by others.
- S_IROTH: 00004 Read by others.
- S_IWOTH: 00002 Write by others.
- S_IXOTH: 00001 Execute by others.
mknod(2)

The owner ID of the file is set to the effective user ID of the process. The group ID of the file is set to the effective group ID of the process. However, if the S_ISGID bit is set in the parent directory, then the group ID of the file is inherited from the parent. If the group ID of the new file does not match the effective group ID or one of the supplementary group IDs, the S_ISGID bit is cleared.

The access permission bits of mode are modified by the process’s file mode creation mask: all bits set in the process’s file mode creation mask are cleared (see umask(2)). If mode indicates a block or character special file, dev is a configuration-dependent specification of a character or block I/O device. If mode does not indicate a block special or character special device, dev is ignored. See makedev(3C).

If path is a symbolic link, it is not followed.

RETURN VALUES
Upon successful completion, mknod() returns 0. Otherwise, it returns −1, the new file is not created, and errno is set to indicate the error.

ERRORS
The mknod() function will fail if:

EACCES A component of the path prefix denies search permission, or write permission is denied on the parent directory.

EDQUOT The directory where the new file entry is being placed cannot be extended because the user’s quota of disk blocks on that file system has been exhausted, or the user’s quota of inodes on the file system where the file is being created has been exhausted.

EEXIST The named file exists.

EFAULT The path argument points to an illegal address.

EINTR A signal was caught during the execution of the mknod() function.

EINVAL An invalid argument exists.

EIO An I/O error occurred while accessing the file system.

ELOOP Too many symbolic links were encountered in translating path.

ENAME_TOO_LONG The length of the path argument exceeds PATH_MAX, or the length of a path component exceeds NAME_MAX while _POSIX_NO_TRUNC is in effect.

ENORESGENT A component of the path prefix specified by path does not name an existing directory or path is an empty string.

ENOLINK The path argument points to a remote machine and the link to that machine is no longer active.
The directory that would contain the new file cannot be extended or the file system is out of file allocation resources.

A component of the path prefix is not a directory.

The effective user of the calling process is not super-user.

The directory in which the file is to be created is located on a read-only file system.

The `mknod()` function may fail if:

| ENAMETOOLONG | Pathname resolution of a symbolic link produced an intermediate result whose length exceeds `PATH_MAX`. |

Normally, applications should use the `mkdir(2)` routine to make a directory, since the function `mknod()` may not establish directory entries for the directory itself (.) and the parent directory (..), and appropriate permissions are not required. Similarly, `mkfifo(3C)` should be used in place of `mknod()` in order to create FIFOs.

The `mknod()` function may be invoked only by a privileged user for file types other than FIFO special.

**SEE ALSO**

`chmod(2)`, `creat(2)`, `exec(2)`, `mkdir(2)`, `open(2)`, `stat(2)`, `umask(2)`, `makedev(3C)`, `mkfifo(3C)`, `stat(3HEAD)`
## NAME
mmap – map pages of memory

## SYNOPSIS
```c
#include <sys/mman.h>

void *mmap(void *addr, size_t len, int prot, int flags, int fildes, off_t off);
```

## DESCRIPTION
The `mmap()` function establishes a mapping between a process’s address space and a file or shared memory object. The format of the call is as follows:

```c
pa = mmap(addr, len, prot, flags, fildes, off);
```

The `mmap()` function establishes a mapping between the address space of the process at an address `pa` for `len` bytes to the memory object represented by the file descriptor `fildes` at offset `off` for `len` bytes. The value of `pa` is a function of the `addr` argument and values of `flags`, further described below. A successful `mmap()` call returns `pa` as its result. The address range starting at `pa` and continuing for `len` bytes will be legitimate for the possible (not necessarily current) address space of the process. The range of bytes starting at `off` and continuing for `len` bytes will be legitimate for the possible (not necessarily current) offsets in the file or shared memory object represented by `fildes`.

The `mmap()` function allows `[pa, pa + len)` to extend beyond the end of the object both at the time of the `mmap()` and while the mapping persists, such as when the file is created prior to the `mmap()` call and has no contents, or when the file is truncated. Any reference to addresses beyond the end of the object, however, will result in the delivery of a `SIGBUS` or `SIGSEGV` signal. The `mmap()` function cannot be used to implicitly extend the length of files.

The mapping established by `mmap()` replaces any previous mappings for those whole pages containing any part of the address space of the process starting at `pa` and continuing for `len` bytes.

If the size of the mapped file changes after the call to `mmap()` as a result of some other operation on the mapped file, the effect of references to portions of the mapped region that correspond to added or removed portions of the file is unspecified.

The `mmap()` function is supported for regular files and shared memory objects. Support for any other type of file is unspecified.

The `prot` argument determines whether read, write, execute, or some combination of accesses are permitted to the data being mapped. The `prot` argument should be either `PROT_NONE` or the bitwise inclusive OR of one or more of the other flags in the following table, defined in the header `<sys/mman.h>`.

| PROT_READ | Data can be read. |
| PROT_WRITE | Data can be written. |
| PROT_EXEC | Data can be executed. |
| PROT_NONE | Data cannot be accessed. |
If an implementation of `mmap()` for a specific platform cannot support the combination of access types specified by `prot`, the call to `mmap()` fails. An implementation may permit accesses other than those specified by `prot`; however, the implementation will not permit a write to succeed where `PROT_WRITE` has not been set or permit any access where `PROT_NONE` alone has been set. Each platform-specific implementation of `mmap()` supports the following values of `prot`: `PROT_NONE`, `PROT_READ`, `PROT_WRITE`, and the inclusive OR of `PROT_READ` and `PROT_WRITE`. On some platforms, the `PROT_WRITE` protection option is implemented as `PROT_READ|PROT_WRITE` and `PROT_EXEC` as `PROT_READ|PROT_EXEC`. The file descriptor `fd` is opened with read permission, regardless of the protection options specified. If `PROT_WRITE` is specified, the application must have opened the file descriptor `fd` with write permission unless `MAP_PRIVATE` is specified in the `flags` argument as described below.

The `flags` argument provides other information about the handling of the mapped data. The value of `flags` is the bitwise inclusive OR of these options, defined in `<sys/mman.h>`:

- `MAP_SHARED`: Changes are shared.
- `MAP_PRIVATE`: Changes are private.
- `MAP_FIXED`: Interpret `addr` exactly.
- `MAP_NORESERVE`: Do not reserve swap space.
- `MAP_ANON`: Map anonymous memory.
- `MAP_ALIGN`: Interpret `addr` as required alignment.

The `MAP_SHARED` and `MAP_PRIVATE` options describe the disposition of write references to the underlying object. If `MAP_SHARED` is specified, write references will change the memory object. If `MAP_PRIVATE` is specified, the initial write reference will create a private copy of the memory object page and redirect the mapping to the copy. The private copy is not created until the first write; until then, other users who have the object mapped `MAP_SHARED` can change the object. Either `MAP_SHARED` or `MAP_PRIVATE` must be specified, but not both. The mapping type is retained across `fork(2)`.

When `MAP_FIXED` is set in the `flags` argument, the system is informed that the value of `pa` must be `addr`, exactly. If `MAP_FIXED` is set, `mmap()` may return `(void *)-1` and set `errno` to `EINVAL`. If a `MAP_FIXED` request is successful, the mapping established by `mmap()` replaces any previous mappings for the process’s pages in the range `[pa, pa + len)`. The use of `MAP_FIXED` is discouraged, since it may prevent a system from making the most effective use of its resources.

When `MAP_FIXED` is set and the requested address is the same as previous mapping, the previous address is unmapped and the new mapping is created on top of the old one.
When `MAP_FIXED` is not set, the system uses `addr` to arrive at `pa`. The `pa` so chosen will be an area of the address space that the system deems suitable for a mapping of `len` bytes to the file. The `mmap()` function interprets an `addr` value of 0 as granting the system complete freedom in selecting `pa`, subject to constraints described below. A non-zero value of `addr` is taken to be a suggestion of a process address near which the mapping should be placed. When the system selects a value for `pa`, it will never place a mapping at address 0, nor will it replace any extant mapping, nor map into areas considered part of the potential data or stack “segments”.

When `MAP_ALIGN` is set, the system is informed that the alignment of `pa` must be the same as `addr`. The alignment value in `addr` must be 0 or some power of two multiple of page size as returned by `sysconf(3C)`. If `addr` is 0, the system will choose a suitable alignment.

The `MAP_NORESERVE` option specifies that no swap space be reserved for a mapping. Without this flag, the creation of a writable `MAP_PRIVATE` mapping reserves swap space equal to the size of the mapping; when the mapping is written into, the reserved space is employed to hold private copies of the data. A write into a `MAP_NORESERVE` mapping produces results which depend on the current availability of swap space in the system. If space is available, the write succeeds and a private copy of the written page is created; if space is not available, the write fails and a SIGBUS or SIGSEGV signal is delivered to the writing process. `MAP_NORESERVE` mappings are inherited across `fork()`. When `MAP_NORESERVE` mappings are inherited across `fork()`, swap space is reserved in the child for all private pages that currently exist in the parent; thereafter the child’s mapping behaves as described above.

When `MAP_ANON` is set in `flags`, and `fd` is set to -1, `mmap()` provides a direct path to return anonymous pages to the caller. This operation is equivalent to passing `mmap()` an open file descriptor on `/dev/zero` with `MAP_ANON` elided from the `flags` argument.

The `off` argument is constrained to be aligned and sized according to the value returned by `sysconf(3C)` when passed `SC_PAGESIZE` or `SC_PAGE_SIZE`. When `MAP_FIXED` is specified, the `addr` argument must also meet these constraints. The system performs mapping operations over whole pages. Thus, while the `len` argument need not meet a size or alignment constraint, the system will include, in any mapping operation, any partial page specified by the range `[pa, pa + len)`.

The system will always zero-fill any partial page at the end of an object. Further, the system will never write out any modified portions of the last page of an object which are beyond its end. References to whole pages following the end of an object will result in the delivery of a SIGBUS or SIGSEGV signal. SIGBUS signals may also be delivered on various file system conditions, including quota exceeded errors.

The `mmap()` function adds an extra reference to the file associated with the file descriptor `fd` which is not removed by a subsequent `close(2)` on that file descriptor. This reference is removed when there are no more mappings to the file by a call to the `munmap(2)` function.
The `st_atime` field of the mapped file may be marked for update at any time between the `mmap()` call and the corresponding `munmap()` call. The initial read or write reference to a mapped region will cause the file’s `st_atime` field to be marked for update if it has not already been marked for update.

The `st_ctime` and `st_mtime` fields of a file that is mapped with `MAP_SHARED` and `PROT_WRITE`, will be marked for update at some point in the interval between a write reference to the mapped region and the next call to `msync(3C)` with `MS_ASYNC` or `MS_SYNC` for that portion of the file by any process. If there is no such call, these fields may be marked for update at any time after a write reference if the underlying file is modified as a result.

If the process calls `mlockall(3C)` with the `MCL_FUTURE` flag, the pages mapped by all future calls to `mmap()` will be locked in memory. In this case, if not enough memory could be locked, `mmap()` fails and sets `errno` to `EAGAIN`.

Upon successful completion, the `mmap()` function returns the address at which the mapping was placed (`pBegin`); otherwise, it returns a value of `MAP_FAILED` and sets `errno` to indicate the error. The symbol `MAP_FAILED` is defined in the header `<sys/mman.h>`. No successful return from `mmap()` will return the value `MAP_FAILED`.

If `mmap()` fails for reasons other than `EACCES`, `EINVAL` or `ENOTSUP`, some of the mappings in the address range starting at `addr` and continuing for `len` bytes may have been unmapped.

The `mmap()` function will fail if:

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>EACCES</td>
<td>The file descriptor is not open for read, regardless of the protection specified; or <code>filedes</code> is not open for write and <code>PROT_WRITE</code> was specified for a <code>MAP_SHARED</code> type mapping.</td>
</tr>
<tr>
<td>EAGAIN</td>
<td>The mapping could not be locked in memory.</td>
</tr>
<tr>
<td>EBADF</td>
<td>The file descriptor is not open (and <code>MAP_ANON</code> was not specified).</td>
</tr>
<tr>
<td>EINVAL</td>
<td>The arguments <code>addr</code> (if <code>MAP_FIXED</code> was specified) or <code>off</code> are not multiples of the page size as returned by <code>sysconf()</code>.</td>
</tr>
<tr>
<td></td>
<td>The argument <code>addr</code> (if <code>MAP_ALIGN</code> was specified) is not 0 or some power of two multiple of page size as returned by <code>sysconf(3C)</code>.</td>
</tr>
<tr>
<td></td>
<td><code>MAP_FIXED</code> and <code>MAP_ALIGN</code> are both specified.</td>
</tr>
</tbody>
</table>

The field in `flags` is invalid (neither `MAP_PRIVATE` or `MAP_SHARED` is set).
mmap(2)

The argument \textit{len} has a value equal to 0.

\texttt{MAP\_ANON} was specified, but the file descriptor was not \texttt{-1}.

\texttt{EMFILE} The number of mapped regions would exceed an implementation-dependent limit (per process or per system).

\texttt{ENODEV} The \texttt{fd}es argument refers to an object for which \texttt{mmap()} is meaningless, such as a terminal.

\texttt{ENOMEM} The \texttt{MAP\_FIXED} option was specified and the range \texttt{[addr, addr + len]} exceeds that allowed for the address space of a process.

The \texttt{MAP\_FIXED} option was \textit{not} specified and there is insufficient room in the address space to effect the mapping.

The mapping could not be locked in memory, if required by \texttt{mlockall}(3C), because it would require more space than the system is able to supply.

The composite size of \texttt{len} plus the lengths obtained from all previous calls to \texttt{mmap()} exceeds \texttt{RLIMIT\_VMEM} (see \texttt{getrlimit}(2)).

\texttt{ENOTSUP} The system does not support the combination of accesses requested in the \texttt{prot} argument.

\texttt{ENXIO} Addresses in the range \texttt{[off, off + len]} are invalid for the object specified by \texttt{fd}es.

The \texttt{MAP\_FIXED} option was specified in \texttt{flags} and the combination of \texttt{addr}, \texttt{len} and \texttt{off} is invalid for the object specified by \texttt{fd}es.

\texttt{EOVERFLOW} The file is a regular file and the value of \texttt{off} plus \texttt{len} exceeds the offset maximum establish in the open file description associated with \texttt{fd}es.

The \texttt{mmap()} function may fail if:

\texttt{EAGAIN} The file to be mapped is already locked using advisory or mandatory record locking. See \texttt{fcntl(2)}.

\texttt{USAGE} Use of \texttt{mmap()} may reduce the amount of memory available to other memory allocation functions.

\texttt{MAP\_ALIGN} is useful to assure a properly aligned value of \texttt{pa} for subsequent use with \texttt{memcntl(2)} and the \texttt{MC\_HAT\_ADVISE} command. This is best used for large, long-lived, and heavily referenced regions. \texttt{MAP\_FIXED} and \texttt{MAP\_ALIGN} are always mutually-exclusive.
Use of MAP_FIXED may result in unspecified behavior in further use of brk(2), sbrk(2), malloc(3C), and shmat(2). The use of MAP_FIXED is discouraged, as it may prevent an implementation from making the most effective use of resources.

The application must ensure correct synchronization when using mmap() in conjunction with any other file access method, such as read(2) and write(2), standard input/output, and shmat(2).

The mmap() function has a transitional interface for 64-bit file offsets. See lf64(5).

The mmap() function allows access to resources using address space manipulations instead of the read()/write() interface. Once a file is mapped, all a process has to do to access it is use the data at the address to which the object was mapped.

Consider the following pseudo-code:

```c
fildes = open(...)
lseek(fildes, offset, whence)
read(fildes, buf, len)
/* use data in buf */
```

The following is a rewrite using mmap():

```c
fildes = open(...)
address = mmap((caddr_t) 0, len, (PROT_READ | PROT_WRITE),
                MAP_PRIVATE, fildes, offset)
/* use data at address */
```

## ATTRIBUTES

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Standard</td>
</tr>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

## SEE ALSO

close(2), exec(2), fcntl(2), fork(2), getrlimit(2), memcntl(2), mprotect(2), munmap(2), shmat(2), lockf(3C), mlockall(3C), msync(3C), plock(3C), sysconf(3C), attributes(5), lf64(5), standards(5), null(7D), zero(7D)
mount(2)

NAME
mount – mount a file system

SYNOPSIS
#include <sys/types.h>
#include <sys/mount.h>
#include <sys/mntent.h>

int mount(const char *spec, const char *dir, int mflag, char *fstype,
char *dataptr, int datalen, char *optptr, int optlen);

DESCRIPTION
The mount() function requests that a removable file system contained on the block
special file identified by spec be mounted on the directory identified by dir. The spec
and dir arguments are pointers to path names.

After a successful call to mount(), all references to the file dir refer to the root
directory on the mounted file system. The mounted file system is inserted into the
kernel list of all mounted file systems. This list can be examined through the mounted
file system table (see mnttab(4)).

The fstype argument is the file system type name. Standard file system names are
de
defined with the prefix MNTTYPE_ in <sys/mntent.h>.

The dataptr argument is 0 if no file system-specific data is to be passed; otherwise it
points to an area of size datalen that contains the file system-specific data for this
mount and the MS_DATA flag should be set.

If the MS_OPTIONSTR flag is set, then optptr points to a buffer containing the list of
options to be used for this mount. The optlen argument specifies the length of the
buffer. On completion of the mount() call, the options in effect for the mounted file
system are returned in this buffer. If MS OPTIONSTR is not specified, then the options
for this mount will not appear in the mounted file systems table.

The mflag argument is constructed by a bitwise-inclusive-OR of flags from the
following list, defined in <sys/mount.h>.

<table>
<thead>
<tr>
<th>Flag</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MS_DATA</td>
<td>The dataptr and datalen arguments describe a block of file system-specific binary data at address dataptr of length datalen. This is interpreted by file system-specific code within the operating system and its format depends on the file system type. If a particular file system type does not require this data, dataptr and datalen should both be 0.</td>
</tr>
<tr>
<td>MS_GLOBAL</td>
<td>Mount a file system globally if the system is configured and booted as part of a cluster (see clinfo(1M)).</td>
</tr>
<tr>
<td>MS_NOSUID</td>
<td>Prevent programs that are marked set-user-ID or set-group-ID from executing (see chmod(1)). It also causes open(2) to return ENXIO when attempting to open block or character special files.</td>
</tr>
<tr>
<td>MS_OPTIONSTR</td>
<td>The optptr and optlen arguments describe a character buffer at address optptr of size optlen. When calling mount(), the character buffer should contain a null-terminated string of options to be passed to the file system-specific code within the operating system.</td>
</tr>
</tbody>
</table>
On a successful return, the file system-specific code will return the list of options recognized. Unrecognized options are ignored. The format of the string is a list of option names separated by commas. Options that have values (rather than binary options such as suid or nosuid), are separated by "=" such as dev=2c4046c. Standard option names are defined in <sys/mntent.h>. Only strings defined in the "C" locale are supported. The maximum length option string that can be passed to or returned from a mount() call is defined by the MAX_MNTOPT_STR constant. The buffer should be long enough to contain more options than were passed in, as the state of any default options that were not passed in the input option string may also be returned in the recognized options list that is returned.

MS_OVERLAY  Allow the file system to be mounted over an existing file system mounted on dir, making the underlying file system inaccessible. If a mount is attempted on a pre-existing mount point without setting this flag, the mount will fail.

MS_RDONLY  Mount the file system for reading only. This flag should also be specified for file systems that are incapable of writing (for example, CDROM). Without this flag, writing is permitted according to individual file accessibility.

MS_REMOUNT  Remount a read-only file system as read-write.

RETURN VALUES  Upon successful completion, 0 is returned. Otherwise, -1 is returned and errno is set to indicate the error.

ERRORS  The mount() function will fail if:

EBUSY  The dir argument is currently mounted on, is someone’s current working directory, or is otherwise busy; the device associated with spec is currently mounted; or there are no more mount table entries.

EFAULT  The spec, dir, fstype, dataptr, or optptr argument points outside the allocated address space of the process.

EINVAL  The super block has an invalid magic number or the fstype is invalid.

ELOOP  Too many symbolic links were encountered in translating spec or dir.

ENAMETOOLONG  The length of the path argument exceeds PATH_MAX, or the length of a path component exceeds NAME_MAX while _POSIX_NO_TRUNC is in effect.

ENOENT  None of the named files exists or is a null pathname.
### mount(2)

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ENOLINK</td>
<td>The <code>path</code> argument points to a remote machine and the link to that machine is no longer active.</td>
</tr>
<tr>
<td>ENOSPC</td>
<td>The file system state in the super-block is not FsOKAY and <code>mflag</code> requests write permission.</td>
</tr>
<tr>
<td>ENOTBLK</td>
<td>The <code>spec</code> argument is not a block special device.</td>
</tr>
<tr>
<td>ENOTDIR</td>
<td>The <code>dir</code> argument is not a directory, or a component of a path prefix is not a directory.</td>
</tr>
<tr>
<td>ENOTSUP</td>
<td>A global mount is attempted (the <code>MS_GLOBAL</code> flag is set in <code>mflag</code>) on a machine which is not booted as a cluster or a local mount is attempted and <code>dir</code> is within a globally mounted file system.</td>
</tr>
<tr>
<td>ENXIO</td>
<td>The device associated with <code>spec</code> does not exist.</td>
</tr>
<tr>
<td>EOVERFLOW</td>
<td>The length of the option string to be returned in the <code>optptr</code> argument exceeds the size of the buffer specified by <code>optlen</code>.</td>
</tr>
<tr>
<td>EPERM</td>
<td>The effective user ID is not superuser.</td>
</tr>
<tr>
<td>EREMOTE</td>
<td>The <code>spec</code> argument is remote and cannot be mounted.</td>
</tr>
<tr>
<td>EROFS</td>
<td>The <code>spec</code> argument is write protected and <code>mflag</code> requests write permission.</td>
</tr>
</tbody>
</table>

#### Usage

The `mount()` function can be invoked only by processes with superuser privileges.

#### See Also

`mount(1M), umount(2), mnttab(4)`

#### Notes

`MS_OPTIONSTR`-type option strings should be used.

Some flag bits set file system options that can also be passed in an option string. Options are first set from the option string with the last setting of an option in the string determining the value to be set by the option string. Any options controlled by flags are then applied, overriding any value set by the option string.
**NAME**
mprotect – set protection of memory mapping

**SYNOPSIS**
```
#include <sys/mman.h>

int mprotect(void *addr, size_t len, int prot);
```

**DESCRIPTION**
The `mprotect()` function changes the access protections on the mappings specified by the range \([addr, addr + len]\), rounding `len` up to the next multiple of the page size as returned by `sysconf(3C)`, to be that specified by `prot`. Legitimate values for `prot` are the same as those permitted for `mmap(2)` and are defined in `<sys/mman.h>` as:

- `PROT_READ` /* page can be read */
- `PROT_WRITE` /* page can be written */
- `PROT_EXEC` /* page can be executed */
- `PROT_NONE` /* page can not be accessed */

When `mprotect()` fails for reasons other than `EINVAL`, the protections on some of the pages in the range \([addr, addr + len]\) may have been changed. If the error occurs on some page at `addr2`, then the protections of all whole pages in the range \([addr, addr2]\) will have been modified.

**RETURN VALUES**
Upon successful completion, `mprotect()` returns 0. Otherwise, it returns −1 and sets `errno` to indicate the error.

**ERRORS**
The `mprotect()` function will fail if:

- `EACCES` The `prot` argument specifies a protection that violates the access permission the process has to the underlying memory object.
- `EINVAL` The `len` argument has a value equal to 0, or `addr` is not a multiple of the page size as returned by `sysconf(3C)`.
- `ENOMEM` Addresses in the range \([addr, addr + len]\) are invalid for the address space of a process, or specify one or more pages which are not mapped.

The `mprotect()` function may fail if:

- `EAGAIN` The address range \([addr, addr + len]\) includes one or more pages that have been locked in memory and that were mapped `MAP_PRIVATE`; `prot` includes `PROT_WRITE`; and the system has insufficient resources to reserve memory for the private pages that may be created. These private pages may be created by store operations in the now-writable address range.

**SEE ALSO**
`mmap(2), plock(3C), mlock(3C), mlockall(3C), sysconf(3C)`
# msgctl(2)

## NAME
msgctl – message control operations

## SYNOPSIS
```c
#include <sys/msg.h>

int msgctl(int msqid, int cmd, struct msqid_ds *buf);
```

## DESCRIPTION
The `msgctl()` function provides a variety of message control operations as specified by `cmd`. The following commands are available:

- **IPC_STAT**: Place the current value of each member of the data structure associated with `msqid` into the structure pointed to by `buf`. The contents of this structure are defined in `intro(2)`.

- **IPC_SET**: Set the value of the following members of the data structure associated with `msqid` to the corresponding value found in the structure pointed to by `buf`:
  ```
  msg_perm.uid
  msg_perm.gid
  msg_perm.mode /* access permission bits only */
  msg_qbytes
  ```
  This `cmd` can only be executed by a process that has an effective user ID equal to either that of super-user, or to the value of `msg_perm.cuid` or `msg_perm.uid` in the data structure associated with `msqid`. Only super-user can raise the value of `msg_qbytes`.

- **IPC_RMID**: Remove the message queue identifier specified by `msqid` from the system and destroy the message queue and data structure associated with it. This `cmd` can only be executed by a process that has an effective user ID equal to either that of super-user, or to the value of `msg_perm.cuid` or `msg_perm.uid` in the data structure associated with `msqid`. The `buf` argument is ignored.

## RETURN VALUES
Upon successful completion, `msgctl()` returns 0. Otherwise, it returns −1 and sets `errno` to indicate the error.

## ERRORS
The `msgctl()` function will fail if:

- **EACCES**: The `cmd` argument is `IPC_STAT` and operation permission is denied to the calling process (see `intro(2)`).
- **EFAULT**: The `buf` argument points to an illegal address.
- **EINVAL**: The `msqid` argument is not a valid message queue identifier; or the `cmd` argument is not a valid command or is `IPC_SET` and `msg_perm.uid` or `msg_perm.gid` is not valid.
- **EPERM**: The `cmd` argument is `IPC_RMID` or `IPC_SET` and the effective user ID of the calling process is not super-user and is not equal to the value of `msg_perm.cuid` or `msg_perm.uid` in the data structure associated with `msqid`. 
The `cmd` argument is `IPC_SET`, an attempt is being made to increase to the value of `msg_qbytes`, and the effective user ID of the calling process is not super-user.

The `cmd` argument is `IPC_STAT` and `uid` or `gid` is too large to be stored in the structure pointed to by `buf`.

**SEE ALSO**

`intro(2), msgget(2), msgrcv(2), msgsnd(2)`
msgget(2)

NAME msgget – get message queue

SYNOPSIS #include <sys/msg.h>

    int msgget(key_t key, int msgflg);

DESCRIPTION The msgget() argument returns the message queue identifier associated with key.

A message queue identifier and associated message queue and data structure (see intro(2)) are created for key if one of the following are true:

- key is IPC_PRIVATE.
- key does not already have a message queue identifier associated with it, and (msgflg & IPC_CREAT) is true.

On creation, the data structure associated with the new message queue identifier is initialized as follows:

- msg_perm.cuid, msg_perm.uid, msg_perm.cgid, and msg_perm.gid are set to the effective user ID and effective group ID, respectively, of the calling process.
- The low-order 9 bits of msg_perm.mode are set to the low-order 9 bits of msgflg.
- msg_qnum, msg_lspid, msg_lrpid, msg_stime, and msg_rtime are set to 0.
- msg_ctime is set to the current time.
- msg_qbytes is set to the system limit.

RETURN VALUES Upon successful completion, a non-negative integer representing a message queue identifier is returned. Otherwise, -1 is returned and errno is set to indicate the error.

ERRORS The msgget() function will fail if:

EACCES A message queue identifier exists for key, but operation permission (see intro(2)) as specified by the low-order 9 bits of msgflg would not be granted.

EEXIST A message queue identifier exists for key but (msgflg & IPC_CREAT) and (msgflg & IPC_EXCL) are both true.

ENOENT A message queue identifier does not exist for key and (msgflg & IPC_CREAT) is false.

ENOSPC A message queue identifier is to be created but the system-imposed limit on the maximum number of allowed message queue identifiers system wide would be exceeded.

SEE ALSO intro(2), msgctl(2), msgrcv(2), msgsnd(2), ftok(3C)
NAME
msgid – discover all message queue identifiers

SYNOPSIS
#include <sys/msg.h>

int msgids(int *buf, uint_t nids, uint_t *pnids);

DESCRIPTION
The msgids() function copies all active message queue identifiers from the system into the user-defined buffer specified by buf, provided that the number of such identifiers is not greater than the number of integers the buffer can contain, as specified by nids. If the size of the buffer is insufficient to contain all of the active message queue identifiers in the system, none are copied.

Whether or not the size of the buffer is sufficient to contain all of them, the number of active message queue identifiers in the system is copied into the unsigned integer pointed to by pnids.

If nids is 0 or less than the number of active message queue identifiers in the system, buf is ignored.

RETURN VALUES
Upon successful completion, msgids() returns 0. Otherwise, -1 is returned and errno is set to indicate the error.

ERRORS
The msgids() function will fail if:
EFAULT The buf or pnids argument points to an illegal address.

USAGE
The msgids() function returns a snapshot of all the active message queue identifiers in the system. More may be added and some may be removed before they can be used by the caller.

EXAMPLES
EXAMPLE 1 msgids() example

This is sample C code indicating how to use the msgids() function (see msgsnap(2)):

```c
void examine_queues()
{
    int *ids = NULL;
    uint_t nids = 0;
    uint_t n;
    int i;

    for (;;) {
        if (msgids(ids, nids, &n) != 0) {
            perror("msgids");
            exit(1);
        }
        if (n <= nids) /* we got them all */
            break;
        /* we need a bigger buffer */
        ids = realloc(ids, (nids = n) * sizeof(int));
    }
    for (i = 0; i < n; i++)
```

System Calls 171
EXAMPLE 1  msgids( ) example (Continued)

        process_msgid(ids[i]);
        free(ids);
    }

ATTRIBUTES  See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO  ipcrm(1), ipcs(1), intro(2), msgctl(2), msgget(2), msgsnd(2), msgrcv(2),
           msgmsn(2), attributes(5)
**NAME**
msgrcv – message receive operation

**SYNOPSIS**
```c
#include <sys/msg.h>

ssize_t msgrcv(int msqid, void *msgp, size_t msgsz, long int msgtyp,
                 int msgflg);
```

**DESCRIPTION**
The `msgrcv()` function reads a message from the queue associated with the message queue identifier specified by `msqid` and places it in the user-defined buffer pointed to by `msgp`.

The `msgp` argument points to a user-defined buffer that must contain first a field of type `long int` that will specify the type of the message, and then a data portion that will hold the data bytes of the message. The structure below is an example of what this user-defined buffer might look like:

```c
struct mymsg {
    long int mtype; /* message type */
    char mtext[1]; /* message text */
}
```

The `mtype` member is the received message’s type as specified by the sending process.

The `mtext` member is the text of the message.

The `msgsz` argument specifies the size in bytes of `mtext`. The received message is truncated to `msgsz` bytes if it is larger than `msgsz` and `(msgflg&MSG_NOERROR)` is non-zero. The truncated part of the message is lost and no indication of the truncation is given to the calling process.

The `msgtyp` argument specifies the type of message requested as follows:

- If `msgtyp` is 0, the first message on the queue is received.
- If `msgtyp` is greater than 0, the first message of type `msgtyp` is received.
- If `msgtyp` is less than 0, the first message of the lowest type that is less than or equal to the absolute value of `msgtyp` is received.

The `msgflg` argument specifies which of the following actions is to be taken if a message of the desired type is not on the queue:

- If `(msgflg&IPC_NOWAIT)` is non-zero, the calling process will return immediately with a return value of `-1` and `errno` set to `ENOMSG`.
- If `(msgflg&IPC_NOWAIT)` is 0, the calling process will suspend execution until one of the following occurs:
  - A message of the desired type is placed on the queue.
  - The message queue identifier `msqid` is removed from the system (see `msgctl(2)`); when this occurs, `errno` is set equal to `EIDRM` and `-1` is returned.
  - The calling process receives a signal that is to be caught; in this case a message is not received and the calling process resumes execution in the manner prescribed in `sigaction(2)`.
msgrcv(2)

Upon successful completion, the following actions are taken with respect to the data structure associated with msqid (see intro(2)):

- msg_qnum is decremented by 1.
- msg_lrpid is set equal to the process ID of the calling process.
- msg_rtime is set equal to the current time.

RETURN VALUES

Upon successful completion, msgrcv() returns a value equal to the number of bytes actually placed into the buffer mtext. Otherwise, −1 is returned, no message is received, and errno is set to indicate the error.

ERRORS

The msgrcv() function will fail if:

- E2BIG The value of mtext is greater than msgsz and (msgflg&MSG_NOERROR) is 0.
- EACCES Operation permission is denied to the calling process. See intro(2).
- EIDRM The message queue identifier msqid is removed from the system.
- EINTR The msgrcv() function was interrupted by a signal.
- EINVAL The msqid argument is not a valid message queue identifier.
- ENOMSG The queue does not contain a message of the desired type and (msgflg&IPC_NOWAIT) is non-zero.

The msgrcv() function may fail if:

-EFAULT The msgp argument points to an illegal address.

USAGE

The value passed as the msgp argument should be converted to type void *.

SEE ALSO

intro(2), msgctl(2), msgget(2), msgsnd(2), sigaction(2)
The `msgsnap()` function reads all of the messages of type `msgtyp` from the queue associated with the message queue identifier specified by `msqid` and places them in the user-defined buffer pointed to by `buf`.

The `buf` argument points to a user-defined buffer that on return will contain first a buffer header structure:

```c
struct msgsnap_head {
    size_t msgsnap_size; /* bytes used/required in the buffer */
    size_t msgsnap_nmsg; /* number of messages in the buffer */
};
```

followed by `msgsnap_nmsg` messages, each of which starts with a message header:

```c
struct msgsnap_mhead {
    size_t msgsnap_mlen; /* number of bytes in the message */
    long msgsnap_mtype; /* message type */
};
```

and followed by `msgsnap_mlen` bytes containing the message contents.

Each subsequent message header is located at the first byte following the previous message contents, rounded up to a `sizeof(size_t)` boundary.

The `bufsz` argument specifies the size of `buf` in bytes. If `bufsz` is less than `sizeof(msgsnap_head)`, `msgsnap()` fails with `EINVAL`. If `bufsz` is insufficient to contain all of the requested messages, `msgsnap()` succeeds but returns with `msgsnap_nmsg` set to 0 and with `msgsnap_size` set to the required size of the buffer in bytes.

The `msgtyp` argument specifies the types of messages requested as follows:

- If `msgtyp` is 0, all of the messages on the queue are read.
- If `msgtyp` is greater than 0, all messages of type `msgtyp` are read.
- If `msgtyp` is less than 0, all messages with type less than or equal to the absolute value of `msgtyp` are read.

The `msgsnap()` function is a non-destructive operation. Upon completion, no changes are made to the data structures associated with `msqid`.

Upon successful completion, `msgsnap()` returns 0. Otherwise, −1 is returned and `errno` is set to indicate the error.

The `msgsnap()` function will fail if:

- `EACCES` Operation permission is denied to the calling process. See `intro(2)`.  

System Calls 175
**msgsnap(2)**

<table>
<thead>
<tr>
<th>EINVAL</th>
<th>The <em>msqid</em> argument is not a valid message queue identifier or the value of <em>bufsz</em> is less than sizeof(struct msgsnap_head).</th>
</tr>
</thead>
<tbody>
<tr>
<td>EFAULT</td>
<td>The <em>buf</em> argument points to an illegal address.</td>
</tr>
</tbody>
</table>

**USAGE**
The `msgsnap()` function returns a snapshot of messages on a message queue at one point in time. The queue contents can change immediately following return from `msgsnap()`.

**EXAMPLES**

**EXAMPLE 1** `msgsnap()` example

This is sample C code indicating how to use the `msgsnap` function (see `msgids(2)`).

```c
void process_msgid(int msqid)
{
    size_t bufsize;
    struct msgsnap_head *buf;
    struct msgsnap_mhead *mhead;
    int i;

    /* allocate a minimum-size buffer */
    buf = malloc(bufsize = sizeof(struct msgsnap_head));

    /* read all of the messages from the queue */
    for (;;) {
        if (msgsnap(msqid, buf, bufsize, 0) != 0) {
            perror("msgsnap");
            free(buf);
            return;
        }
        if (bufsize >= buf->msgsnap_size) /* we got them all */
            break;
        /* we need a bigger buffer */
        buf = realloc(buf, bufsize = buf->msgsnap_size);
    }

    /* process each message in the queue (there may be none) */
    mhead = (struct msgsnap_mhead *)(buf + 1); /* first message */
    for (i = 0; i < buf->msgsnap_nmsg; i++) {
        size_t mlen = mhead->msgsnap_mlen;

        /* process the message contents */
        process_message(mhead->msgsnap_mtype, (char *)(mhead+1), mlen);

        /* advance to the next message header */
        mhead = (struct msgsnap_mhead *)
            ((char *)mhead + sizeof(struct msgsnap_mhead) +
             (mlen + sizeof(size_t) - 1) & ~(sizeof(size_t) - 1));
    }

    free(buf);
}
```
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO**

ipcrm(1), ipcs(1), intro(2), msgctl(2), msgget(2), msgsids(2), msgrcv(2), msgsnd(2), attributes(5)
msgsnd(2)

NAME
msgsnd – message send operation

SYNOPSIS
#include <sys/msg.h>

int msgsnd(int msqid, const void *msgp, size_t msgsz, int msgflg);

DESCRIPTION
The msgsnd() function is used to send a message to the queue associated with the
message queue identifier specified by msqid.

The msgp argument points to a user-defined buffer that must contain first a field of
type long int that will specify the type of the message, and then a data portion that
will hold the data bytes of the message. The structure below is an example of what
this user-defined buffer might look like:

```
struct mymsg {
    long mtype; /* message type */
    char mtext[1]; /* message text */
}
```

The mtype member is a non-zero positive type long int that can be used by the
receiving process for message selection.

The mtext member is any text of length msgsz bytes. The msgsz argument can range
from 0 to a system-imposed maximum.

The msgflg argument specifies the action to be taken if one or more of the following are true:

- The number of bytes already on the queue is equal to msg_qbytes; see intro(2).
- The total number of messages on all queues system-wide is equal to the
  system-imposed limit.

These actions are as follows:

- If (msgflg & IPC_NOWAIT) is non-zero, the message will not be sent and the calling
  process will return immediately.
- If (msgflg & IPC_NOWAIT) is 0, the calling process will suspend execution until one
  of the following occurs:
    - The condition responsible for the suspension no longer exists, in which case the
      message is sent.
    - The message queue identifier msqid is removed from the system (see
      msgctl(2)); when this occurs, errno is set equal to EIDRM and −1 is returned.
    - The calling process receives a signal that is to be caught; in this case the
      message is not sent and the calling process resumes execution in the manner
      prescribed in sigaction(2).

Upon successful completion, the following actions are taken with respect to the data
structure associated with msqid (see intro(2)):

- msg_qnum is incremented by 1.
- msg_lspid is set equal to the process ID of the calling process.
msgsnd() is set equal to the current time.

**RETURN VALUES**
Upon successful completion, 0 is returned. Otherwise, −1 is returned, no message is sent, and errno is set to indicate the error.

**ERRORS**
The msgsnd() function will fail if:

- **EACCES** Operation permission is denied to the calling process. See intro(2).
- **EAGAIN** The message cannot be sent for one of the reasons cited above and (msgflags&IPC_NOWAIT) is non-zero.
- **EIDRM** The message queue identifier msgid is removed from the system.
- **EINTR** The msgsnd() function was interrupted by a signal.
- **EINVAL** The value of msqid is not a valid message queue identifier, or the value of mtype is less than 1; or the value of msgsz is less than 0 or greater than the system-imposed limit.

The msgsnd() function may fail if:

- **EFAULT** The msgp argument points to an illegal address.

**USAGE**
The value passed as the msgp argument should be converted to type void *.

**SEE ALSO**
intro(2), msgctl(2), msgget(2), msgrcv(2), sigaction(2)
munmap(2)

NAME    munmap – unmap pages of memory

SYNOPSIS #include <sys/mman.h>

    int munmap(void *addr, size_t len);

DESCRIPTION The munmap() function removes the mappings for pages in the range
[addr, addr + len), rounding the len argument up to the next multiple of the page size as
returned by sysconf(3C). If addr is not the address of a mapping established by a
prior call to mmap(2), the behavior is undefined. After a successful call to munmap() and
before any subsequent mapping of the unmapped pages, further references to
these pages will result in the delivery of a SIGBUS or SIGSEGV signal to the process.

The mmap(2) function often performs an implicit munmap().

RETURN VALUES Upon successful completion, munmap() returns 0; otherwise, it returns −1 and sets
errno to indicate an error.

ERRORS The munmap() function will fail if:

EINVAL The addr argument is not a multiple of the page size as returned by
sysconf(3C); addresses in the range [addr, addr + len) are outside
the valid range for the address space of a process; or the len
argument has a value less than or equal to 0.

SEE ALSO mmap(2), sysconf(3C)
nice(2)

NAME
nice – change priority of a process

SYNOPSIS
#include <unistd.h>

int nice(int incr);

DESCRIPTION
The nice() function allows a process to change its priority. The invoking process
must be in a scheduling class that supports the nice().

The nice() function adds the value of incr to the nice value of the calling process. A
process’s nice value is a non-negative number for which a greater positive value
results in lower CPU priority.

A maximum nice value of (2 * NZERO) – 1 and a minimum nice value of 0 are imposed
by the system. NZERO is defined in <limits.h> with a default value of 20. Requests
for values above or below these limits result in the nice value being set to the

The nice() function has no effect on the priority of processes or threads with

A process with superuser privileges can lower the nice value.

RETURN VALUES
Upon successful completion, nice() returns the new nice value minus NZERO.
Otherwise, −1 is returned, the process’s nice value is not changed, and errno is set to
 indicate the error.

ERRORS
The nice() function will fail if:

EINVAL The nice() function is called by a process in a scheduling class other than
time-sharing or fixed-priority.

EPERM The incr argument is negative or greater than 40 and the effective user ID of
the calling process is not superuser.

USAGE
The priocntl(2) function is a more general interface to scheduler functions.

Since −1 is a permissible return value in a successful situation, an application wishing
to check for error situations should set errno to 0, then call nice(), and if it returns
−1, check to see if errno is non-zero.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
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</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Standard</td>
</tr>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
nice(1), exec(2), priocntl(2), getpriority(3C), attributes(5), standards(5)
The `ntp_adjtime()` function adjusts the parameters used to discipline the local clock, according to the values in the struct `timex` pointed to by `tptr`. Before returning, it fills in the structure with the most recent values kept in the kernel.

The adjustment is effected in part by speeding up or slowing down the clock, as necessary, and in part by phase-locking onto a once-per second pulse (PPS) provided by a driver, if available.

```
struct timex {
    uint32_t modes; /* clock mode bits (w) */
    int32_t offset; /* time offset (us) (rw) */
    int32_t freq; /* frequency offset (scaled ppm) (rw) */
    int32_t maxerror; /* maximum error (us) (rw) */
    int32_t esterror; /* estimated error (us) (rw) */
    int32_t status; /* clock status bits (rw) */
    int32_t constant; /* pll time constant (rw) */
    int32_t precision; /* clock precision (us) (r) */
    int32_t tolerance; /* clock frequency tolerance (scaled ppm) (r) */
    int32_t ppsfreq; /* pps frequency (scaled ppm) (r) */
    int32_t jitter; /* pps jitter (us) (r) */
    int32_t shift; /* interval duration (s) (shift) (r) */
    int32_t stabil; /* pps stability (scaled ppm) (r) */
    int32_t jitcnt; /* jitter limit exceeded (r) */
    int32_t calcnt; /* calibration intervals (r) */
    int32_t errcnt; /* calibration errors (r) */
    int32_t stbcnt; /* stability limit exceeded (r) */
};
```

Upon successful completion, `ntp_adjtime()` returns the current clock state (see `<sys/timex.h>`). Otherwise, it returns −1 and sets `errno` to indicate the error.

The `ntp_adjtime()` function will fail if:

- **EFAULT** The `tptr` argument is an invalid pointer.
- **EINVAL** The `constant` member of the structure pointed to by `tptr` is less than 0 or greater than 30.
- **EPERM** The user is not super-user.

SEE ALSO

`xntpd(1M), ntp_gettime(2)`
**NAME**

ntp_gettime – get local clock values

**SYNOPSIS**

```c
#include <sys/timex.h>

int ntp_gettime(struct ntptimeval *tptr);
```

**DESCRIPTION**

The `ntp_gettime()` function reads the local clock value and dispersion, returning the information in `tptr`.

The `ntptimeval` structure contains the following members:

```c
struct ntptimeval {
    struct timeval time; /* current time (ro) */
    int32_t maxerror; /* maximum error (us) (ro) */
    int32_t esterror; /* estimated error (us) (ro) */
};
```

Upon successful completion, `ntp_gettime()` returns the current clock state (see `<sys/timex.h>`). Otherwise, it returns -1 and sets `errno` to indicate the error.

**ERRORS**

The `ntp_gettime()` function will fail if:

- **EFAULT** The `tptr` argument points to an invalid address.

The `ntp_gettime()` function will fail for 32-bit interfaces if:

- **EOVERFLOW** The size of the `time.tv_sec` member of the `ntptimeval` structure pointed to by `tptr` is too small to contain the correct number of seconds.

**SEE ALSO**
xntpd(1M), ntp_adjtime(2)
### NAME
open, openat – open a file

### SYNOPSIS
```
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>

int open(const char *path, int oflag, /* mode_t mode */ ...);
int openat(int fd, const char *path, int oflag, /* mode_t mode */ ...);
```

### DESCRIPTION
The `open()` function establishes the connection between a file and a file descriptor. It creates an open file description that refers to a file and a file descriptor that refers to that open file description. The file descriptor is used by other I/O functions to refer to that file. The `path` argument points to a pathname naming the file.

The `openat()` function is identical to the `open()` function except that the `path` argument is interpreted relative to the starting point implied by the `fd` argument. If the `fd` argument has the special value `AT_FDCWD`, a relative path argument will be resolved relative to the current working directory. If the `path` argument is absolute, the `fd` argument is ignored.

The `open()` function returns a file descriptor for the named file that is the lowest file descriptor not currently open for that process. The open file description is new, and therefore the file descriptor does not share it with any other process in the system. The `FD_CLOEXEC` file descriptor flag associated with the new file descriptor is cleared.

The file offset used to mark the current position within the file is set to the beginning of the file.

The file status flags and file access modes of the open file description are set according to the value of `oflag`. The `mode` argument is used only when `O_CREAT` is specified (see below.)

Values for `oflag` are constructed by a bitwise-inclusive-OR of flags from the following list, defined in `<fcntl.h>`. Applications must specify exactly one of the first three values (file access modes) below in the value of `oflag`:

- `O_RDONLY`: Open for reading only.
- `O_WRONLY`: Open for writing only.
- `O_RDWR`: Open for reading and writing. The result is undefined if this flag is applied to a FIFO.

Any combination of the following may be used:

- `O_APPEND`: If set, the file offset is set to the end of the file prior to each write.
- `O_CREAT`: Create the file if it does not exist. This flag requires that the `mode` argument be specified.
If the file exists, this flag has no effect except as noted under O_EXCL below. Otherwise, the file is created with the user ID of the file set to the effective user ID of the process. The group ID of the file is set to the effective group IDs of the process, or if the S_ISGID bit is set in the directory in which the file is being created, the file's group ID is set to the group ID of its parent directory. If the group ID of the new file does not match the effective group ID or one of the supplementary groups IDs, the S_ISGID bit is cleared. The access permission bits (see <sys/stat.h>) of the file mode are set to the value of mode, modified as follows (see creat(2)): a bitwise-AND is performed on the file-mode bits and the corresponding bits in the complement of the process's file mode creation mask. Thus, all bits set in the process's file mode creation mask (see umask(2)) are correspondingly cleared in the file's permission mask. The "save text image after execution bit" of the mode is cleared (see chmod(2)).

O_SYNC Write I/O operations on the file descriptor complete as defined by synchronized I/O integrity completion (see fcntl(3HEAD) definition of O_SYNC.) When bits other than the file permission bits are set, the effect is unspecified. The mode argument does not affect whether the file is open for reading, writing or for both.

O_DSYNC Write I/O operations on the file descriptor complete as defined by synchronized I/O data integrity completion.

O_EXCL If O_CREAT and O_EXCL are set, open() fails if the file exists. The check for the existence of the file and the creation of the file if it does not exist is atomic with respect to other processes executing open() naming the same filename in the same directory with O_EXCL and O_CREAT set. If O_CREAT is not set, the effect is undefined.

O_LARGEFILE If set, the offset maximum in the open file description is the largest value that can be represented correctly in an object of type off64_t.

O_NOCTTY If set and path identifies a terminal device, open() does not cause the terminal device to become the controlling terminal for the process.

O_NONBLOCK or O_NDELAY These flags may affect subsequent reads and writes (see read(2) and write(2)). If both O_NDELAY and O_NONBLOCK are set, O_NONBLOCK takes precedence.

When opening a FIFO with O_RDONLY or O_WRONLY set:

If O_NONBLOCK or O_NDELAY is set:
An open() for reading only returns without delay. An open() for writing only returns an error if no process currently has the file open for reading.

If O_NONBLOCK and O_NDELAY are clear:

An open() for reading only blocks until a process opens the file for writing. An open() for writing only blocks until a process opens the file for reading.

After both ends of a FIFO have been opened, there is no guarantee that further calls to open() O_RDONLY (O_WRONLY) will synchronize with later calls to open() O_WRONLY (O_RDONLY) until both ends of the FIFO have been closed by all readers and writers. Any data written into a FIFO will be lost if both ends of the FIFO are closed before the data is read.

When opening a block special or character special file that supports non-blocking opens:

If O_NONBLOCK or O_NDELAY is set:

The open() function returns without blocking for the device to be ready or available. Subsequent behavior of the device is device-specific.

If O_NONBLOCK and O_NDELAY are clear:

The open() function blocks until the device is ready or available before returning.

Otherwise, the behavior of O_NONBLOCK and O_NDELAY is unspecified.

O_RSYNC Read I/O operations on the file descriptor complete at the same level of integrity as specified by the O_DSYNC and O_SYNC flags. If both O_DSYNC and O_RSYNC are set in oflag, all I/O operations on the file descriptor complete as defined by synchronized I/O data integrity completion. If both O_SYNC and O_RSYNC are set in oflag, all I/O operations on the file descriptor complete as defined by synchronized I/O file integrity completion.

O_SYNC Write I/O operations on the file descriptor complete as defined by synchronized I/O file integrity completion.

O_TRUNC If the file exists and is a regular file, and the file is successfully opened O_RDWR or O_WRONLY, its length is truncated to 0 and the mode and owner are unchanged. It has no effect on FIFO special
files or terminal device files. Its effect on other file types is implementation-dependent. The result of using O_TRUNC with O_RDONLY is undefined.

O_XATTR If set in openat(), a relative path argument is interpreted as a reference to an extended attribute of the file associated with the supplied file descriptor. This flag therefore requires the presence of a legal path argument. If set in open(), the implied file descriptor is that for the current working directory. Extended attributes must be referenced with a relative path; providing an absolute path results in a normal file reference.

If O_CREAT is set and the file did not previously exist, upon successful completion, open() marks for update the st_atime, st_ctime, and st_mtime fields of the file and the st_ctime and st_mtime fields of the parent directory.

If O_TRUNC is set and the file did previously exist, upon successful completion, open() marks for update the st_ctime and st_mtime fields of the file.

If path refers to a STREAMS file, oflag may be constructed from O_NONBLOCK or O_NODELAY OR-ed with either O_RDONLY, O_WRONLY, or O_RDWR. Other flag values are not applicable to STREAMS devices and have no effect on them. The values O_NONBLOCK and O_NODELAY affect the operation of STREAMS drivers and certain functions (see read(2), getmsg(2), putmsg(2), and write(2)) applied to file descriptors associated with STREAMS files. For STREAMS drivers, the implementation of O_NONBLOCK and O_NODELAY is device-specific.

When open() is invoked to open a named stream, and the connld module (see connld(7M)) has been pushed on the pipe, open() blocks until the server process has issued an I_RECVFD ioctl() (see streamio(7I)) to receive the file descriptor.

If path names the master side of a pseudo-terminal device, then it is unspecified whether open() locks the slave side so that it cannot be opened. Portable applications must call unlockpt(3C) before opening the slave side.

If path is a symbolic link and O_CREAT and O_EXCL are set, the link is not followed.

Certain flag values can be set following open() as described in fcntl(2).

The largest value that can be represented correctly in an object of type off_t is established as the offset maximum in the open file description.

RETURN VALUES Upon successful completion, the open() function opens the file and return a non-negative integer representing the lowest numbered unused file descriptor. Otherwise, -1 is returned, errno is set to indicate the error, and no files are created or modified.

ERRORS The open() and openat() functions will fail if:
<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>EACCES</td>
<td>Search permission is denied on a component of the path prefix, or the file exists and the permissions specified by oflag are denied, or the file does not exist and write permission is denied for the parent directory of the file to be created, or O_TRUNC is specified and write permission is denied.</td>
</tr>
<tr>
<td>EBADF</td>
<td>The file descriptor provided to openat() is invalid.</td>
</tr>
<tr>
<td>EDQUOT</td>
<td>The file does not exist, O_CREAT is specified, and either the directory where the new file entry is being placed cannot be extended because the user's quota of disk blocks on that file system has been exhausted, or the user's quota of inodes on the file system where the file is being created has been exhausted.</td>
</tr>
<tr>
<td>EEXIST</td>
<td>The O_CREAT and O_EXCL flags are set, and the named file exists.</td>
</tr>
<tr>
<td>EINTR</td>
<td>A signal was caught during open().</td>
</tr>
<tr>
<td>EFAULT</td>
<td>The path argument points to an illegal address.</td>
</tr>
<tr>
<td>EINVAL</td>
<td>The system does not support synchronized I/O for this file, or the O_XATTR flag was supplied and the underlying file system does not support extended file attributes.</td>
</tr>
<tr>
<td>EIO</td>
<td>The path argument names a STREAMS file and a hangup or error occurred during the open().</td>
</tr>
<tr>
<td>EISDIR</td>
<td>The named file is a directory and oflag includes O_WRONLY or O_RDWR.</td>
</tr>
<tr>
<td>ELOOP</td>
<td>Too many symbolic links were encountered in resolving path.</td>
</tr>
<tr>
<td>EMFILE</td>
<td>OPEN_MAX file descriptors are currently open in the calling process.</td>
</tr>
<tr>
<td>EMULTIHOP</td>
<td>Components of path require hopping to multiple remote machines and the file system does not allow it.</td>
</tr>
<tr>
<td>ENAMETOOLONG</td>
<td>The length of the path argument exceeds PATH_MAX or a pathname component is longer than NAME_MAX.</td>
</tr>
<tr>
<td>ENFILE</td>
<td>The maximum allowable number of files is currently open in the system.</td>
</tr>
<tr>
<td>ENOENT</td>
<td>The O_CREAT flag is not set and the named file does not exist; or the O_CREAT flag is set and either the path prefix does not exist or the path argument points to an empty string.</td>
</tr>
<tr>
<td>ENOLINK</td>
<td>The path argument points to a remote machine, and the link to that machine is no longer active.</td>
</tr>
<tr>
<td>ENSR</td>
<td>The path argument names a STREAMS-based file and the system is unable to allocate a STREAM.</td>
</tr>
</tbody>
</table>
The directory or file system that would contain the new file cannot be expanded, the file does not exist, and O_CREAT is specified.

The device specified by path does not support the open operation.

A component of the path prefix is not a directory or a relative path was supplied to openat(), the O_XATTR flag was not supplied, and the file descriptor does not refer to a directory.

The O_NONBLOCK flag is set, the named file is a FIFO, the O_WRONLY flag is set, and no process has the file open for reading; or the named file is a character special or block special file and the device associated with this special file does not exist.

An attempt was made to open a path that corresponds to a AF_UNIX socket.

The named file is a regular file and either O_LARGEFILE is not set and the size of the file cannot be represented correctly in an object of type off_t or O_LARGEFILE is set and the size of the file cannot be represented correctly in an object of type off64_t.

The named file resides on a read-only file system and either O_WRONLY, O_RDWR, O_CREAT (if file does not exist), or O_TRUNC is set in the oflag argument.

The openat() function will fail if:

EBADF The filedes argument is not a valid open file descriptor or is not AT_FTCWD.

The open() function may fail if:

EAGAIN The path argument names the slave side of a pseudo-terminal device that is locked.

EINVAL The value of the oflag argument is not valid.

ENAMETOOLONG Pathname resolution of a symbolic link produced an intermediate result whose length exceeds PATH_MAX.

ENOMEM The path argument names a STREAMS file and the system is unable to allocate resources.

ETXTBSY The file is a pure procedure (shared text) file that is being executed and oflag is O_WRONLY or O_RDWR.

The open() function has a transitional interface for 64-bit file offsets. See l64(5).

Note that using open64() is equivalent to using open() with O_LARGEFILE set in oflag.
open(2)

**ATTRIBUTES**  See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td><code>open()</code> is Standard; <code>openat()</code> is Evolving</td>
</tr>
<tr>
<td>MT-Level</td>
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</tbody>
</table>

**SEE ALSO**  intro(2), chmod(2), close(2), creat(2), dup(2), exec(2), fcntl(2), getmsg(2), getrlimit(2), lseek(2), putmsg(2), read(2), stat(2), umask(2), write(2), attropen(3C), unlockpt(3C), attributes(5), fcntl(3HEAD), l64(5), stat(3HEAD), connld(7M), streamio(7I)

**NOTES**  Hierarchical Storage Management (HSM) file systems can sometimes cause long delays when opening a file, since HSM files must be recalled from secondary storage.
pause – suspeend process until signal

SYNOPSIS
#include <unistd.h>

int pause(void);

DESCRIPTION
The pause() function suspends the calling process until it receives a signal. The
signal must be one that is not currently set to be ignored by the calling process.

If the signal causes termination of the calling process, pause() does not return.

If the signal is caught by the calling process and control is returned from the
signal-catching function (see signal(3C)), the calling process resumes execution from
the point of suspension.

RETURN VALUES
Since pause() suspends thread execution indefinitely unless interrupted by a signal,
there is no successful completion return value. If interrupted, it returns -1 and sets
errno to indicate the error.

ERRORS
The pause() function will fail if:

EINTR A signal is caught by the calling process and control is returned
from the signal-catching function.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
alarm(2), kill(2), wait(2), signal(3C), attributes(5)
The pcsample() function provides CPU-use statistics by profiling the amount of CPU time expended by a program.

For profiling dynamically-linked programs and 64-bit programs, it is superior to the profil(2) function, which assumes that the entire program is contained in a small, contiguous segment of the address space, divides this segment into "bins", and on each clock tick increments the counter in the bin where the program is currently executing. With shared libraries creating discontinuous program segments spread throughout the address space, and with 64-bit address spaces so large that the size of "bins" would be measured in megabytes, the profil() function is of limited value.

The pcsample() function is passed an array samples containing nsamples pointer-sized elements. During program execution, the kernel samples the program counter of the process, storing unadulterated values in the array on each clock tick. The kernel stops writing to the array when it is full, which occurs after nsamples / HZ seconds of process virtual time. The HZ value is obtained by invoking the call sysconf(_SC_CLK_TCK). See sysconf(3C).

The sampling can be stopped by a subsequent call to pcsample() with the nsamples argument set to 0. Like profil(), sampling continues across a call to fork(2), but is disabled by a call to one of the exec family of functions (see exec(2)). It is also disabled if an update of the samples[] array causes a memory fault.

The pcsample() function always returns 0 the first time it is called. On subsequent calls, it returns the number of samples that were stored during the previous invocation. If nsamples is invalid, it returns −1 and sets errno to indicate the error.

The pcsample() function will fail if:

EINVAL The value of nsamples is not valid.

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
<tr>
<td>Interface Stability</td>
<td>Stable</td>
</tr>
</tbody>
</table>

SEE ALSO exec(2), fork(2), profil(2), sysconf(3C), attributes(5)
pipe – create an interprocess channel

#include <unistd.h>

int pipe(int fildes[2]);

The pipe() function creates an I/O mechanism called a pipe and returns two file descriptors, fildes[0] and fildes[1]. The files associated with fildes[0] and fildes[1] are streams and are both opened for reading and writing. The O_NDELAY and O_NONBLOCK flags are cleared.

A read from fildes[0] accesses the data written to fildes[1] on a first-in-first-out (FIFO) basis and a read from fildes[1] accesses the data written to fildes[0] also on a FIFO basis.

The FD_CLOEXEC flag will be clear on both file descriptors.

Upon successful completion pipe() marks for update the st_atime, st_ctime, and st_mtime fields of the pipe.

Upon successful completion, 0 is returned. Otherwise, −1 is returned and errno is set to indicate the error.

The pipe() function will fail if:

EMFILE There are OPEN_MAX−1 or more file descriptors currently open for this process.

ENOMEM A file table entry could not be allocated.

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

Since a pipe is bi-directional, there are two separate flows of data. Therefore, the size (st_size) returned by a call to fstat(2) with argument fildes[0] or fildes[1] is the number of bytes available for reading from fildes[0] or fildes[1] respectively. Previously, the size (st_size) returned by a call to fstat() with argument fildes[1] (the write-end) was the number of bytes available for reading from fildes[0] (the read-end).
poll(2)

NAME poll – input/output multiplexing

SYNOPSIS
#include <poll.h>

int poll(struct pollfd fds[], nfds_t nfds, int timeout);

DESCRIPTION
The poll() function provides applications with a mechanism for multiplexing
input/output over a set of file descriptors. For each member of the array pointed to by
fds, poll() examines the given file descriptor for the event(s) specified in events. The
number of pollfd structures in the fds array is specified by nfds. The poll() function identifies those file descriptors on which an application can read or write
data, or on which certain events have occurred.

The fds argument specifies the file descriptors to be examined and the events of
interest for each file descriptor. It is a pointer to an array with one member for each
open file descriptor of interest. The array’s members are pollfd structures, which
contain the following members:

int fd;    /* file descriptor */
short events; /* requested events */
short revents; /* returned events */

The fd member specifies an open file descriptor and the events and revents
members are bitmasks constructed by a logical OR operation of any combination of
the following event flags:

POLLIN      Data other than high priority data may be read without blocking.
            For STREAMS, this flag is set in revents even if the message is of
            zero length.

POLLRDNORM  Normal data (priority band equals 0) may be read without
            blocking. For STREAMS, this flag is set in revents even if the
            message is of zero length.

POLLRDBAND  Data from a non-zero priority band may be read without blocking.
            For STREAMS, this flag is set in revents even if the message is of
            zero length.

POLLPRI     High priority data may be received without blocking. For
            STREAMS, this flag is set in revents even if the message is of
            zero length.

POLLOUT     Normal data (priority band equals 0) may be written without
            blocking.

POLLOUT     The same as POLLOUT.

POLLRWRBAND Priority data (priority band > 0) may be written. This event only
            examines bands that have been written to at least once.

POLLOUT     An error has occurred on the device or stream. This flag is only
            valid in the revents bitmask; it is not used in the events
            member.
A hangup has occurred on the stream. This event and POLLOUT are mutually exclusive; a stream can never be writable if a hangup has occurred. However, this event and POLLIN, POLLRDNDM, POLLRDDBAND, or POLLPRI are not mutually exclusive. This flag is only valid in the revents bitmask; it is not used in the events member.

The specified fd value does not belong to an open file. This flag is only valid in the revents member; it is not used in the events member.

If the value fd is less than 0, events is ignored and revents is set to 0 in that entry on return from poll().

The results of the poll() query are stored in the revents member in the pollfd structure. Bits are set in the revents bitmask to indicate which of the requested events are true. If none are true, none of the specified bits are set in revents when the poll() call returns. The event flags POLLHUP, POLLERR, and POLLNVAL are always set in revents if the conditions they indicate are true; this occurs even though these flags were not present in events.

If none of the defined events have occurred on any selected file descriptor, poll() waits at least timeout milliseconds for an event to occur on any of the selected file descriptors. On a computer where millisecond timing accuracy is not available, timeout is rounded up to the nearest legal value available on that system. If the value timeout is 0, poll() returns immediately. If the value of timeout is −1, poll() blocks until a requested event occurs or until the call is interrupted. The poll() function is not affected by the O_NDELAY and O_NONBLOCK flags.

The poll() function supports regular files, terminal and pseudo-terminal devices, STREAMS-based files, FIFOs and pipes. The behavior of poll() on elements of fds that refer to other types of file is unspecified.

The poll() function supports sockets.

A file descriptor for a socket that is listening for connections will indicate that it is ready for reading, once connections are available. A file descriptor for a socket that is connecting asynchronously will indicate that it is ready for writing, once a connection has been established.

Regular files always poll() TRUE for reading and writing.

Upon successful completion, a non-negative value is returned. A positive value indicates the total number of file descriptors that has been selected (that is, file descriptors for which the revents member is non-zero). A value of 0 indicates that the call timed out and no file descriptors have been selected. Upon failure, −1 is returned and errno is set to indicate the error.

The poll() function will fail if:
### poll(2)

<table>
<thead>
<tr>
<th>Error</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>EAGAIN</td>
<td>Allocation of internal data structures failed, but the request may be attempted again.</td>
</tr>
<tr>
<td>EFAULT</td>
<td>Some argument points to an illegal address.</td>
</tr>
<tr>
<td>EINTR</td>
<td>A signal was caught during the <code>poll()</code> function.</td>
</tr>
<tr>
<td>EINVAL</td>
<td>The argument <code>nfds</code> is greater than <code>OPEN_MAX</code>, or one of the <code>fd</code> members refers to a STREAM or multiplexer that is linked (directly or indirectly) downstream from a multiplexer.</td>
</tr>
</tbody>
</table>

### SEE ALSO

- `intro(2)`, `getmsg(2)`, `getrlimit(2)`, `putmsg(2)`, `read(2)`, `write(2)`, `select(3C)`, `chpoll(9E)`

### STREAMS Programming Guide

### NOTES

Non-STREAMS drivers use `chpoll(9E)` to implement `poll()` on these devices.
The `p_online()` function changes or returns the operational status of processors. The state of the processor specified by the `processorid` argument is changed to the state represented by the `flag` argument.

Legal values for `flag` are `P_STATUS`, `P_ONLINE`, `P_OFFLINE`, and `P_NOINTR`.

When `flag` is `P_STATUS`, no processor status change occurs, but the current processor status is returned.

The `P_ONLINE`, `P_OFFLINE`, and `P_NOINTR` values for `flag` refer to valid processor states. A processor in the `P_ONLINE` state is allowed to process LWPs (lightweight processes) and perform system activities. The processor is also interruptible by I/O devices attached to the system.

A processor in the `P_OFFLINE` state is not allowed to process LWPs. The processor is as inactive as possible. If the hardware supports such a feature, the processor is not interruptible by attached I/O devices.

A processor in the `P_NOINTR` state is allowed to process LWPs, but it is not interruptible by attached I/O devices. Typically, interrupts, when they occur are routed to other processors in the system. Not all systems support putting a processor into the `P_NOINTR` state. At least one processor must always be available to service system clock interrupts.

Processor numbers are integers, greater than or equal to 0, and are defined by the hardware platform. Processor numbers are not necessarily contiguous, but “not too sparse.” Processor numbers should always be printed in decimal.

The maximum possible `processorid` value can be determined by calling `sysconf(_SC_CPUID_MAX)`. The list of valid processor numbers can be determined by calling `p_online()` with `processorid` values from 0 to the maximum returned by `sysconf(_SC_CPUID_MAX)`. The EINVAL error is returned for invalid processor numbers. See EXAMPLES below.

On successful completion, the value returned is the previous state of the processor, `P_ONLINE`, `P_OFFLINE`, `P_NOINTR`, or `P_POWEROFF`. Otherwise, −1 is returned and `errno` is set to indicate the error.

The `p_online()` function will fail if:

- EPERM: The effective user of the calling process is not super-user.
- EINVAL: A non-existent processor ID was specified or `flag` was invalid.
EBUSY The flag was P_OFFLINE and the specified processor is the only on-line processor, there are currently LWPs bound to the processor, or the processor performs some essential function that cannot be performed by another processor.

EBUSY The flag was P_NOINTR and the specified processor is the only interruptible processor in the system, or it handles interrupts that cannot be handled by another processor.

EBUSY The specified processor is powered off and cannot be powered on because some platform-specific resource is not available.

ENOTSUP The specified processor is powered off, and the platform does not support power on of individual processors.

EXAMPLES

EXAMPLE 1 List the legal processor numbers.

The following code sample will list the legal processor numbers:

```
#include <sys/unistd.h>
#include <sys/processor.h>
#include <sys/types.h>
#include <stdio.h>
#include <errno.h>

int main()
{
    processorid_t i, cpuid_max;
    cpuid_max = sysconf(_SC_CPUID_MAX);
    for (i = 0; i <= cpuid_max; i++) {
        if (p_online(i, P_STATUS) != -1)
            printf("processor %d present\n", i);
    }
    return (0);
}
```

ATTRIBUTES See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>MT-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO psradm(1M), psrinfo(1M), processor_bind(2), processor_info(2), pset_create(2), sysconf(3C), attributes(5)
priocntl – process scheduler control

SYNOPSIS

```c
#include <sys/types.h>
#include <sys/priocntl.h>
#include <sys/rtpriocntl.h>
#include <sys/tspriocntl.h>
#include <sys/fspriocntl.h>
#include <sys/fxpriocntl.h>

long priocntl(idtype_t idtype, id_t id, int cmd, /* arg */ ...);
```

DESCRIPTION

The `priocntl()` function provides for control over the scheduling of an active light weight process (LWP).

LWPs fall into distinct classes with a separate scheduling policy applied to each class. The classes currently supported are the realtime class, the time-sharing class, the fair-share class, and the fixed-priority class. The characteristics of these classes are described under the corresponding headings below.

The class attribute of an LWP is inherited across the `fork(2)` and `_lwp_create(2)` functions and the `exec` family of functions (see `exec(2)`). The `priocntl()` function can be used to dynamically change the class and other scheduling parameters associated with a running LWP or set of LWPs given the appropriate permissions as explained below.

In the default configuration, a runnable realtime LWP runs before any other LWP. Therefore, inappropriate use of realtime LWP can have a dramatic negative impact on system performance.

The `priocntl()` function provides an interface for specifying a process, set of processes, or an LWP to which the function applies. The `priocntlset(2)` function provides the same functions as `priocntl()`, but allows a more general interface for specifying the set of LWPs to which the function is to apply.

For `priocntl()`, the `idtype` and `id` arguments are used together to specify the set of LWPs. The interpretation of `id` depends on the value of `idtype`. The possible values for `idtype` and corresponding interpretations of `id` are as follows:

- **P_ALL**
  - The `priocntl()` function applies to all existing LWPs. The value of `id` is ignored. The permission restrictions described below still apply.

- **P_CID**
  - The `id` argument is a class ID (returned by the `priocntl()` `PC_GETCID` command as explained below). The `priocntl()` function applies to all LWPs in the specified class.

- **P_GID**
  - The `id` argument is a group ID. The `priocntl()` function applies to all LWPs with this effective group ID.

- **P_LWPID**
  - The `id` argument is an LWP ID. The `priocntl` function applies to the LWP with the specified ID within the calling process.
### priocntl(2)

<table>
<thead>
<tr>
<th>Argument</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>P_PGID</td>
<td>The id argument is a process group ID. The priocntl() function applies to all LWPs currently associated with processes in the specified process group.</td>
</tr>
<tr>
<td>P_PID</td>
<td>The id argument is a process ID specifying a single process. The priocntl() function applies to all LWPs currently associated with the specified process.</td>
</tr>
<tr>
<td>P_PPID</td>
<td>The id argument is a parent process ID. The priocntl() function applies to all LWPs currently associated with processes with the specified parent process ID.</td>
</tr>
<tr>
<td>P_PROJID</td>
<td>The id argument is a project ID. The priocntl() function applies to all LWPs with this project ID.</td>
</tr>
<tr>
<td>P_SID</td>
<td>The id argument is a session ID. The priocntl() function applies to all LWPs currently associated with processes in the specified session.</td>
</tr>
<tr>
<td>P_TASKID</td>
<td>The id argument is a task ID. The priocntl() function applies to all LWPs currently associated with processes in the specified task.</td>
</tr>
<tr>
<td>P_UID</td>
<td>The id argument is a user ID. The priocntl() function applies to all LWPs with this effective user ID.</td>
</tr>
</tbody>
</table>

An id value of P_MYID can be used in conjunction with the idtype value to specify the LWP ID, parent process ID, process group ID, session ID, task ID, class ID, user ID, group ID, or project ID of the calling LWP.

To change the scheduling parameters of an LWP (using the PC_SETPARMS or PC_SETXPARMS command as explained below), the real or effective user ID of the LWP calling priocntl() must match the real or effective user ID of the receiving LWP or the effective user ID of the calling LWP must be superuser. These are the minimum permission requirements enforced for all classes. An individual class might impose additional permissions requirements when setting LWPs to that class and/or when setting class-specific scheduling parameters.

A special SYS scheduling class exists for the purpose of scheduling the execution of certain special system processes (such as the swapper process). It is not possible to change the class of any LWP to SYS. In addition, any processes in the SYS class that are included in a specified set of processes are disregarded by priocntl(). For example, an idtype of P_UID and an id value of 0 would specify all processes with a user ID of 0 except processes in the SYS class and (if changing the parameters using PC_SETPARMS or PC_SETXPARMS) the init(1M) process.

The init process is a special case. For a priocntl() call to change the class or other scheduling parameters of the init process (process ID 1), it must be the only process specified by idtype and id. The init process can be assigned to any class configured on the system, but the time-sharing class is almost always the appropriate choice. (Other choices might be highly undesirable. See the System Administration Guide: Basic Administration for more information.)
The data type and value of \( \text{arg} \) are specific to the type of command specified by \textit{cmd}.

A \texttt{pcinfo_t} structure with the following members, defined in \texttt{<sys/priorcntl.h>}, is used by the \texttt{PC_GETCID} and \texttt{PC_GETCLINFO} commands.

\begin{verbatim}
  id_t pc_cid;  /* Class id */
  char pc_clname[PC_CLNMSZ]; /* Class name */
  int pc_clinfo[PC_CLINFOSZ]; /* Class information */
\end{verbatim}

The \texttt{pc_cid} member is a class ID returned by the \texttt{priocntl()} \texttt{PC_GETCID} command. The \texttt{pc_clname} member is a buffer of size \texttt{PC_CLNMSZ}, defined in \texttt{<sys/priorcntl.h>}, used to hold the class name: \texttt{RT} for realtime, \texttt{TS} for time-sharing, or \texttt{FX} for fixed-priority.

The \texttt{pc_clinfo} member is a buffer of size \texttt{PC_CLINFOSZ}, defined in \texttt{<sys/priorcntl.h>}, used to return data describing the attributes of a specific class. The format of this data is class-specific and is described under the appropriate heading (\texttt{REALTIME CLASS}, \texttt{TIME-SHARING CLASS}, or \texttt{FIXED-PRIORITY CLASS}) below.

A \texttt{pcparms_t} structure with the following members, defined in \texttt{<sys/priorcntl.h>}, is used by the \texttt{PC_SETPARMS} and \texttt{PC_GETPARMS} commands.

\begin{verbatim}
  id_t pc_cid;  /* LWP class */
  int pc_clparms[PC_CLPARMSZ]; /* Class-specific params */
\end{verbatim}

The \texttt{pc_cid} member is a class ID returned by the \texttt{priocntl()} \texttt{PC_GETCID} command. The special class ID \texttt{PC_CLNULL} can also be assigned to \texttt{pc_cid} when using the \texttt{PC_GETPARMS} command as explained below.

The \texttt{pc_clparms} buffer holds class-specific scheduling parameters. The format of this parameter data for a particular class is described under the appropriate heading below. \texttt{PC_CLPARMSZ} is the length of the \texttt{pc_clparms} buffer and is defined in \texttt{<sys/priorcntl.h>}.

The \texttt{PC_SETXPARMS} and \texttt{PC_GETXPARMS} commands exploit the \texttt{varargs} declaration of \texttt{priocntl()}. The argument following the command code is a class name: \texttt{RT} for realtime, \texttt{TS} for time-sharing, or \texttt{FX} for fixed-priority. The parameters after the class name build a chain of (key, value) pairs, where the key determines the meaning of the value within the pair. When using \texttt{PC_GETXPARMS}, the value associated with the key is always a pointer to a scheduling parameter. In contrast, when using \texttt{PC_SETXPARMS} the scheduling parameter is given as a direct value. A key value of 0 terminates the sequence and all further keys or values are ignored.

The \texttt{PC_SETXPARMS} and \texttt{PC_GETXPARMS} commands are more flexible than \texttt{PC_SETPARMS} and \texttt{PC_GETPARMS} and should replace \texttt{PC_SETPARMS} and \texttt{PC_GETPARMS} on a long-term basis.

\begin{center}
\textbf{COMMANDS} \hspace{2cm} Available \texttt{priocntl()} commands are:
\end{center}
<table>
<thead>
<tr>
<th>Command</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>PC_ADMIN</strong></td>
<td>This command provides functionality needed for the implementation of the <code>dispadmin(1M)</code> utility. It is not intended for general use by other applications.</td>
</tr>
<tr>
<td><strong>PC_DONICE</strong></td>
<td>Set or get nice value of the specified LWP(s) associated with the specified process(es). When this command is used with the <code>idtype</code> of <code>P_LWPID</code>, it sets the nice value of the LWP. The <code>arg</code> argument points to a structure of type <code>pcnice_t</code>. The <code>pc_val</code> member specifies the nice value and the <code>pc_op</code> specifies the type of the operation. When <code>pc_op</code> is set to <code>PC_GETNICE</code>, <code>priocntl()</code> sets the <code>pc_val</code> to the highest priority (lowest numerical value) pertaining to any of the specified LWPs. When <code>pc_op</code> is set to <code>PC_SETNICE</code>, <code>priocntl()</code> sets the nice value of all LWPs in the specified set to the value specified in <code>pc_val</code> member of <code>pcnice_t</code> structure.</td>
</tr>
<tr>
<td><strong>PC_GETCID</strong></td>
<td>Get class ID and class attributes for a specific class given the class name. The <code>idtype</code> and <code>id</code> arguments are ignored. If <code>arg</code> is non-null, it points to a structure of type <code>pcinfo_t</code>. The <code>pc_clname</code> buffer contains the name of the class whose attributes you are getting. On success, the class ID is returned in <code>pc_cid</code>, the class attributes are returned in the <code>pc_clinfo</code> buffer, and the <code>priocntl()</code> call returns the total number of classes configured in the system (including the <code>sys</code> class). If the class specified by <code>pc_clname</code> is invalid or is not currently configured, the <code>priocntl()</code> call returns <code>-1</code> with <code>errno</code> set to <code>EINVAL</code>. The format of the attribute data returned for a given class is defined in the <code>&lt;sys/rtpriocntl.h&gt;</code>, <code>&lt;sys/tspriocntl.h&gt;</code>, or <code>&lt;sys/fxpriocntl.h&gt;</code> header and described under the appropriate heading below. If <code>arg</code> is a null pointer, no attribute data is returned but the <code>priocntl()</code> call still returns the number of configured classes.</td>
</tr>
<tr>
<td><strong>PC_GETCLINFO</strong></td>
<td>Get class name and class attributes for a specific class given class ID. The <code>idtype</code> and <code>id</code> arguments are ignored. If <code>arg</code> is non-null, it points to a structure of type <code>pcinfo_t</code>. The <code>pc_cid</code> member is the class ID of the class whose attributes you are getting.</td>
</tr>
</tbody>
</table>
On success, the class name is returned in the `pc_clname` buffer, the class attributes are returned in the `pc_clinfo` buffer, and the `priocntl()` call returns the total number of classes configured in the system (including the `sys` class). The format of the attribute data returned for a given class is defined in the `<sys/rtpriocntl.h>`, `<sys/tspriocntl.h>`, or `<sys/fxpriocntl.h>` header and described under the appropriate heading below.

If `arg` is a null pointer, no attribute data is returned but the `priocntl()` call still returns the number of configured classes.

**PC_GETPARMS**

Get the class and/or class-specific scheduling parameters of an LWP. The `arg` member points to a structure of type `pcparms_t`.

If `pc_cid` specifies a configured class and a single LWP belonging to that class is specified by the `idtype` and `id` values or the `procset` structure, then the scheduling parameters of that LWP are returned in the `pc_clparms` buffer. If the LWP specified does not exist or does not belong to the specified class, the `priocntl()` call returns −1 with `errno` set to ESRCH.

If `pc_cid` specifies a configured class and a set of LWPs is specified, the scheduling parameters of one of the specified LWP belonging to the specified class are returned in the `pc_clparms` buffer and the `priocntl()` call returns the process ID of the selected LWP. The criteria for selecting an LWP to return in this case is class-dependent. If none of the specified LWPs exist or none of them belong to the specified class, the `priocntl()` call returns −1 with `errno` set to ESRCH.

If `pc_cid` is `PC_CLNULL` and a single LWP is specified, the class of the specified LWP is returned in `pc_cid` and its scheduling parameters are returned in the `pc_clparms` buffer.

**PC_GETXP ARMS**

Get the class or class-specific scheduling parameters of an LWP. The class name (first argument after `PC_GETXP ARMS`) specifies the class and the (key, value) pair sequence contains a pointer to the class-specific parameters. The keys and the types of the class-specific parameter data are described below and can also be found in the class-specific headers `<sys/rtpriocntl.h>`, `<sys/tspriocntl.h>`, and `<sys/fxpriocntl.h>`. If the specified class is a configured class and a single LWP belonging to that class is specified by the `idtype` and `id` values or the `procset` structure, then the scheduling parameters of that LWP are returned in the given (key, value) pair buffers. If the LWP specified does not exist or does not belong to the specified class, `priocntl()` returns −1 and `errno` is set to ESRCH.
If the class name specifies a configured class and a set of LWPs is given, the scheduling parameters of one of the specified LWPs belonging to the specified class are returned and the *priocntl()* call returns the process ID of the selected LWP. The criteria for selecting an LWP to return in this case is class-dependent. If none of the specified LWPs exist or none of them belong to the specified class, *priocntl()* returns −1 and *errno* is set to ESRCH.

If the class name is a null pointer, a single process or LWP is specified, and a (key, value) pair for a class name request is given, *priocntl()* fills the buffer pointed to by value with the class name of the specified process or LWP. The key for the class name request is PC_KY_CLNAME and the class name buffer should be declared as:

```c
char pc_clname[PC_CLNMSZ]; /* Class name */
```

### PC_SETPARMS

Set the class and class-specific scheduling parameters of the specified LWP(s) associated with the specified process(es). When this command is used with the *idtype* of P_LWPID, it will set the class and class-specific scheduling parameters of the LWP. The *arg* argument points to a structure of type pcparms_t. The pc_cid member specifies the class you are setting and the pc_clparms buffer contains the class-specific parameters you are setting. The format of the class-specific parameter data is defined in the `<sys/rtpriocntl.h>`, `<sys/tspriocntl.h>`, or `<sys/fxpriocntl.h>` header and described under the appropriate class heading below.

When setting parameters for a set of LWPs, *priocntl()* acts on the LWPs in the set in an implementation-specific order. If *priocntl()* encounters an error for one or more of the target processes, it might or might not continue through the set of LWPs, depending on the nature of the error. If the error is related to permissions (EPERM), *priocntl()* continues through the LWP set, resetting the parameters for all target LWPs for which the calling LWP has appropriate permissions. The *priocntl()* function then returns −1 with *errno* set to EPERM to indicate that the operation failed for one or more of the target LWPs. If *priocntl()* encounters an error other than permissions, it does not continue through the set of target LWPs but returns the error immediately.

### PC_SETXPARMS

Set the class and class-specific scheduling parameters of the specified LWP(s) associated with the specified process(es). When this command is used with P_LWPID as *idtype*, it will set the class and class-specific scheduling parameters of the LWP. The class name (first argument after PC_SETXPARMS) specifies the class to be changed and the following (key, value) pair sequence contains
the class-specific parameters to be changed. Only those (key,value) pairs whose scheduling behavior is to change must be specified. The keys and the types of the class-specific parameter data are described below and can also be found in the class-specific header files <sys/rtpiocntl.h>, <sys/tspriocntl.h>, and <sys/fxpriocntl.h>.

When setting parameters for a set of LWPs, priocntl() acts on the LWPs in the set in an implementation-specific order. If priocntl() encounters an error for one or more of the target processes, it might or might not continue through the set of LWPs, depending on the nature of the error. If the error is related to permissions (EPERM), priocntl() continues to reset the parameters for all target LWPs where the calling LWP has appropriate permissions. The priocntl() function returns −1 and errno is set to EPERM when the operation failed for one or more of the target LWPs. All errors other than EPERM result in an immediate termination of priocntl().

The realtime class provides a fixed priority preemptive scheduling policy for those LWPs requiring fast and deterministic response and absolute user/application control of scheduling priorities. If the realtime class is configured in the system, it should have exclusive control of the highest range of scheduling priorities on the system. This ensures that a runnable realtime LWP is given CPU service before any LWP belonging to any other class.

The realtime class has a range of realtime priority (rt_pri) values that can be assigned to an LWP within the class. Realtime priorities range from 0 to x, where the value of x is configurable and can be determined for a specific installation by using the priocntl() PC_GETCID or PC_GETCLINFO command.

The realtime scheduling policy is a fixed priority policy. The scheduling priority of a realtime LWP is never changed except as the result of an explicit request by the user/application to change the rt_pri value of the LWP.

For an LWP in the realtime class, the rt_pri value is, for all practical purposes, equivalent to the scheduling priority of the LWP. The rt_pri value completely determines the scheduling priority of a realtime LWP relative to other LWPs within its class. Numerically higher rt_pri values represent higher priorities. Since the realtime class controls the highest range of scheduling priorities in the system, it is guaranteed that the runnable realtime LWP with the highest rt_pri value is always selected to run before any other LWPs in the system.
In addition to providing control over priority, `priocntl()` provides for control over the length of the time quantum allotted to the LWP in the realtime class. The time quantum value specifies the maximum amount of time an LWP can run assuming that it does not complete or enter a resource or event wait state (sleep). If another LWP becomes runnable at a higher priority, the currently running LWP might be preempted before receiving its full time quantum.

The realtime quantum signal can be used for the notification of runaway realtime processes about the consumption of their time quantum. Those processes, which are monitored by the realtime time quantum signal, receive the configured signal in the event of time quantum expiration. The default value (0) of the time quantum signal will denote no signal delivery and a positive value will denote the delivery of the signal specified by the value. The realtime quantum signal can be set with the `priocntl() PC_SETXPARMS` command and displayed with the `priocntl() PC_GETXPARMS` command as explained below.

The system's process scheduler keeps the runnable realtime LWPs on a set of scheduling queues. There is a separate queue for each configured realtime priority and all realtime LWPs with a given `rt_pri` value are kept together on the appropriate queue. The LWPs on a given queue are ordered in FIFO order (that is, the LWP at the front of the queue has been waiting longest for service and receives the CPU first). Realtime LWPs that wake up after sleeping, LWPs that change to the realtime class from some other class, LWPs that have used their full time quantum, and runnable LWPs whose priority is reset by `priocntl()` are all placed at the back of the appropriate queue for their priority. An LWP that is preempted by a higher priority LWP remains at the front of the queue (with whatever time is remaining in its time quantum) and runs before any other LWP at this priority. Following a `fork(2)` or `_lwp_create(2)` function call by a realtime LWP, the parent LWP continues to run while the child LWP (which inherits its parent's `rt_pri` value) is placed at the back of the queue.

An `rtinfo_t` structure with the following members, defined in `<sys/rtpriocntl.h>`, defines the format used for the attribute data for the realtime class.

```c
short rt_maxpri; /* Maximum realtime priority */
```

The `priocntl()` `PC_GETCID` and `PC_GETCLINFO` commands return realtime class attributes in the `pc_clinfo` buffer in this format.

The `rt_maxpri` member specifies the configured maximum `rt_pri` value for the realtime class. If `rt_maxpri` is `x`, the valid realtime priorities range from 0 to `x`.

An `rtparms_t` structure with the following members, defined in `<sys/rtpriocntl.h>`, defines the format used to specify the realtime class-specific scheduling parameters of an LWP.

```c
short rt_pri; /* Real-Time priority */
uint_t rt_tqsecs; /* Seconds in time quantum */
int rt_tqnsecs; /* Additional nanoseconds in quantum */
```
When using the `priocntl()` `PC_SETPARMS` or `PC_GETPARMS` commands, if `pc_cid` specifies the realtime class, the data in the `pc_clparms` buffer are in this format.

These commands can be used to set the realtime priority to the specified value or get the current `rt_pri` value. Setting the `rt_pri` value of an LWP that is currently running or runnable (not sleeping) causes the LWP to be placed at the back of the scheduling queue for the specified priority. The LWP is placed at the back of the appropriate queue regardless of whether the priority being set is different from the previous `rt_pri` value of the LWP. A running LWP can voluntarily release the CPU and go to the back of the scheduling queue at the same priority by resetting its `rt_pri` value to its current realtime priority value. To change the time quantum of an LWP without setting the priority or affecting the LWP’s position on the queue, the `rt_pri` member should be set to the special value `RT_NOCHANGE`, defined in `<sys/rtpriocntl.h>`. Specifying `RT_NOCHANGE` when changing the class of an LWP to realtime from some other class results in the realtime priority being set to 0.

For the `priocntl()` `PC_GETPARMS` command, if `pc_cid` specifies the realtime class and more than one realtime LWP is specified, the scheduling parameters of the realtime LWP with the highest `rt_pri` value among the specified LWPs are returned and the LWP ID of this LWP is returned by the `priocntl()` call. If there is more than one LWP sharing the highest priority, the one returned is implementation-dependent.

The `rt_tqsecs` and `rt_tqnsecs` members are used for getting or setting the time quantum associated with an LWP or group of LWPs. `rt_tqsecs` is the number of seconds in the time quantum and `rt_tqnsecs` is the number of additional nanoseconds in the quantum. For example, setting `rt_tqsecs` to 2 and `rt_tqnsecs` to 500,000,000 (decimal) would result in a time quantum of two and one-half seconds. Specifying a value of 1,000,000,000 or greater in the `rt_tqnsecs` member results in an error return with `errno` set to `EINVAL`. Although the resolution of the `tq_nsec` member is very fine, the specified time quantum length is rounded up by the system to the next integral multiple of the system clock’s resolution. The maximum time quantum that can be specified is implementation-specific and equal to `INT_MAX` ticks. The `INT_MAX` value is defined in `<limits.h>`. Requesting a quantum greater than this maximum results in an error return with `errno` set to `ERANGE`, although infinite quantums can be requested using a special value as explained below. Requesting a time quantum of 0 by setting both `rt_tqsecs` and `rt_tqnsecs` to 0 results in an error return with `errno` set to `EINVAL`.

The `rt_tqnsecs` member can also be set to one of the following special values defined in `<sys/rtpriocntl.h>`, in which case the value of `rt_tqsecs` is ignored:

- **RT_TQINF**: Set an infinite time quantum.
- **RT_TQDEF**: Set the time quantum to the default for this priority (see `rt_dptbl(4)`).
- **RT_NOCHANGE**: Do not set the time quantum. This value is useful when you wish to change the realtime priority of an LWP without affecting the time quantum. Specifying this value when changing the class of an LWP to realtime from some other class results in the realtime priority being set to 0.
LWP to realtime from some other class is equivalent to specifying RT_TQDEF.

When using the priocntl() PC_SETXPARMS or PC_GETXPARMS commands, the first argument after the command code must be the class name of the realtime class ("RT"). The next arguments are formed as (key, value) pairs, terminated by a 0 key. The definition for the keys of the realtime class can be found in <sys/rtpriocntl.h>. A repeated specification of the same key results in an error return and errno set to EINVAL.

<table>
<thead>
<tr>
<th>Key</th>
<th>Value Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>RT_KY_PRI</td>
<td>pri_t</td>
<td>realtime priority</td>
</tr>
<tr>
<td>RT_KY_TQSECS</td>
<td>uint_t</td>
<td>seconds in time quantum</td>
</tr>
<tr>
<td>RT_KY_TQNSECS</td>
<td>int</td>
<td>nanoseconds in time quantum</td>
</tr>
<tr>
<td>RT_KY_TQSIG</td>
<td>int</td>
<td>realtime time quantum signal</td>
</tr>
</tbody>
</table>

When using the priocntl() PC_GETXPARMS command, the value associated with the key is always a pointer to a scheduling parameter of the value type shown in the table above. In contrast, when using the priocntl() PC_SETXPARMS command, the scheduling parameter is given as a direct value.

A priocntl() PC_SETXPARMS command with the class name ("RT") and without a following (key, value) pair will set or reset all realtime scheduling parameters of the target process(es) to their default values. Changing the class of an LWP to realtime from some other class causes the parameters to be set to their default values. The default realtime priority (RT_KY_PRI) is 0. A default time quantum (RT_TQDEF) is assigned to each priority class (see rt_dptbl(4)). The default realtime time quantum signal (RT_KY_TQSIG) is 0.

The value associated with RT_KY_TQSECS is the number of seconds in the time quantum. The value associated with RT_KY_TQNSECS is the number of nanoseconds in the quantum. Specifying a value of 1,000,000,000 or greater for the number of nanoseconds results in an error return and errno is set to EINVAL. The specified time quantum is rounded up by the system to the next integral multiple of the system clock’s resolution. The maximum time quantum that can be specified is implementation-specific and equal to INT_MAX ticks, defined in <limits.h>. Requesting a quantum greater than this maximum results in an error return and errno is set to ERANGE. If seconds (RT_KY_TQSECS) but no nanoseconds (RT_KY_TQNSECS) are supplied, the number of nanoseconds is set to 0. If nanoseconds (RT_KY_TQNSECS) but no seconds (RT_KY_TQSECS) are supplied, the number of seconds is set to 0. A time quantum of 0 (seconds and nanoseconds are 0) results in an error return with errno set toEINVAL. Special values for RT_KY_TQSECS are RT_TQINF and RT_TQDEF (as described above). The priocntl() command PC_SETXPARMS knows no special value RT_NOCHANGE.
To change the class of an LWP to realtime from any other class, the LWP invoking `priocntl()` must have superuser privileges. To change the priority or time quantum setting of a realtime LWP, the LWP invoking `priocntl()` must have superuser privileges or must itself be a realtime LWP whose real or effective user ID matches the real of effective user ID of the target LWP.

The realtime priority and time quantum are inherited across `fork(2)` and the `exec` family of functions. When using the time quantum signal with a user-defined signal handler across the `exec(2)` system call, the new image must install an appropriate user-defined signal handler before the time quantum expires. Otherwise, unpredictable behavior might result.

The time-sharing scheduling policy provides for a fair and effective allocation of the CPU resource among LWPs with varying CPU consumption characteristics. The objectives of the time-sharing policy are to provide good response time to interactive LWPs and good throughput to CPU-bound jobs, while providing a degree of user/application control over scheduling.

The time-sharing class has a range of time-sharing user priority (see `ts_upri` below) values that can be assigned to LWPs within the class. A `ts_upri` value of 0 is defined as the default base priority for the time-sharing class. User priorities range from $-x$ to $+x$ where the value of $x$ is configurable and can be determined for a specific installation by using the `priocntl()` `PC_GETCID` or `PC_GETCLINFO` command.

The purpose of the user priority is to provide some degree of user/application control over the scheduling of LWPs in the time-sharing class. Raising or lowering the `ts_upri` value of an LWP in the time-sharing class raises or lowers the scheduling priority of the LWP. It is not guaranteed, however, that an LWP with a higher `ts_upri` value will run before one with a lower `ts_upri` value, since the `ts_upri` value is just one factor used to determine the scheduling priority of a time-sharing LWP. The system can dynamically adjust the internal scheduling priority of a time-sharing LWP based on other factors such as recent CPU usage.

In addition to the system-wide limits on user priority (returned by the `PC_GETCID` and `PC_GETCLINFO` commands) there is a per LWP user priority limit (see `ts_uprilim` below) specifying the maximum `ts_upri` value that can be set for a given LWP. By default, `ts_uprilim` is 0.

A `tsinfo_t` structure with the following members, defined in `<sys/tspriocntl.h>`, defines the format used for the attribute data for the time-sharing class.

```c
short ts_maxupri; /* Limits of user priority range */
```

The `priocntl()` `PC_GETCID` and `PC_GETCLINFO` commands return time-sharing class attributes in the `pc_clinfo` buffer in this format.

The `ts_maxupri` member specifies the configured maximum user priority value for the time-sharing class. If `ts_maxupri` is $x$, the valid range for both user priorities and user priority limits is from $-x$ to $+x$. 

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A tsparms_t structure with the following members, defined in <sys/tspriocntl.h>, defines the format used to specify the time-sharing class-specific scheduling parameters of an LWP.

```c
short ts_uprilim; /* Time-Sharing user priority limit */
short ts_upri;    /* Time-Sharing user priority */
```

When using the priocntl() PC_SETPARMS or PC_GETPARMS commands, if pc_cid specifies the time-sharing class, the data in the pc_clparms buffer is in this format.

For the priocntl() PC_GETPARMS command, if pc_cid specifies the time-sharing class and more than one time-sharing LWP is specified, the scheduling parameters of the time-sharing LWP with the highest ts_upri value among the specified LWPs is returned and the LWP ID of this LWP is returned by the priocntl() call. If there is more than one LWP sharing the highest user priority, the one returned is implementation-dependent.

Any time-sharing LWP can lower its own ts_uprilim (or that of another LWP with the same user ID). Only a time-sharing LWP with superuser privileges can raise a ts_uprilim. When changing the class of an LWP to time-sharing from some other class, superuser privileges are required to set the initial ts_uprilim to a value greater than 0. Attempts by a non-superuser LWP to raise a ts_uprilim or set an initial ts_uprilim greater than 0 fail with a return value of -1 and errno set to EPERM.

Any time-sharing LWP can set its own ts_upri (or that of another LWP with the same user ID) to any value less than or equal to the LWP’s ts_uprilim. Attempts to set the ts_upri above the ts_uprilim (and/or set the ts_uprilim below the ts_upri) result in the ts_upri being set equal to the ts_uprilim.

Either of the ts_uprilim or ts_upri members can be set to the special value TS_NOCHANGE, defined in <sys/tspriocntl.h>, to set one of the values without affecting the other. Specifying TS_NOCHANGE for the ts_upri when the ts_uprilim is being set to a value below the current ts_upri causes the ts_upri to be set equal to the ts_uprilim being set. Specifying TS_NOCHANGE for a parameter when changing the class of an LWP to time-sharing (from some other class) causes the parameter to be set to a default value. The default value for the ts_uprilim is 0 and the default for the ts_upri is to set it equal to the ts_uprilim that is being set.

When using the priocntl() PC_SETXPARMS or PC_GETXPARMS commands, the first argument after the command code is the class name of the time-sharing class ("TS"). The next arguments are formed as (key, value) pairs, terminated by a 0 key. The definition for the keys of the time-sharing class can be found in <sys/tspriocntl.h>. A repeated specification of the same key results in an error return and errno set to EINVAL.
When using the `priocntl()` `PC_GETXPARMS` command, the value associated with the key is always a pointer to a scheduling parameter of the value type in the table above. In contrast, when using the `priocntl()` `PC_SETXPARMS` command, the scheduling parameter is given as a direct value.

A `priocntl()` `PC_SETXPARMS` command with the class name ("TS") and without a following (key, value) pair will set or reset all time-sharing scheduling parameters of the target process(es) to their default values. Changing the class of an LWP to time-sharing from some other class causes the parameters to be set to their default values. The default value for the user priority limit (TS_KY_UPRILIM) is 0. The default value for the user priority (TS_KY_UPRI) is equal to the user priority limit (TS_KY_UPRILIM) that is being set.

The `priocntl()` command `PC_SETXPARMS` knows no special value TS_NOCHANGE.

The time-sharing user priority and user priority limit are inherited across `fork()` and the `exec` family of functions.

The fair-share scheduling policy provides a fair allocation of CPU resources among projects, independent of the number of processes they contain. Projects are given "shares" to control their quota of CPU resources. See FSS(7) for more information about how to configure shares.

The fair share class supports the notion of per-LWP user priority (see fs_upri below) values for compatibility with the time-sharing scheduling class. An `fss_upri` value of 0 is defined as the default base priority for the fair-share class. User priorities range from -x to +x where the value of x is configurable and can be determined for a specific installation by using the `priocntl()` `PC_GETCID` or `PC_GETCLINFO` command.

The purpose of the user priority is to provide some degree of user/application control over the scheduling of LWPs in the fair-share class. Raising the `fss_upri` value of an LWP in the fair-share class tells the scheduler to give this LWP more CPU time slices, while lowering the `fss_upri` value tells the scheduler to give it less CPU slices. It is not guaranteed, however, that an LWP with a higher `fss_upri` value will run before one with a lower `fss_upri` value. This is because the `fss_upri` value is just one factor used to determine the scheduling priority of a fair-share LWP. The system can dynamically adjust the internal scheduling priority of a fair-share LWP based on other factors such as recent CPU usage. The fair-share scheduler attempts to provide an evenly graded effect across the whole range of user priority values.

---

<table>
<thead>
<tr>
<th>Key</th>
<th>Value Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>TS_KY_UPRILIM</td>
<td>pri_t</td>
<td>user priority limit</td>
</tr>
<tr>
<td>TS_KY_UPRI</td>
<td>pri_t</td>
<td>user priority</td>
</tr>
</tbody>
</table>
User priority values do not interfere with project shares. That is, changing a user priority value of a process does not have any effect on its project CPU entitlement, which is based on the number of shares it is allocated in comparison with other projects.

In addition to the system-wide limits on user priority (returned by the \texttt{PC_GETCID} and \texttt{PC_GETCLINFO} commands), there is a per-LWP user priority limit (see \texttt{fs_uprilim} below) that specifies the maximum \texttt{fss_upri} value that can be set for a given LWP. By default, \texttt{fss_uprilim} is 0.

A \texttt{fssinfo_t} structure with the following members, defined in \texttt{<sys/fsspriocntl.h>}, defines the format used for the attribute data for the fair-share class.

\begin{verbatim}
short fss_maxupri; /* Limits of user priority range */
\end{verbatim}

The \texttt{priocntl()} \texttt{PC_GETCID} and \texttt{PC_GETCLINFO} commands return fair-share class attributes in the \texttt{pc_clinfo} buffer in this format.

\texttt{fss_maxupri} specifies the configured maximum user priority value for the fair-share class. If \texttt{fss_maxupri} is \( x \), the valid range for both user priorities and user priority limits is from \(-x\) to \(+x\).

A \texttt{fssparms_t} structure with the following members, defined in \texttt{<sys/fsspriocntl.h>}, defines the format used to specify the fair-share class-specific scheduling parameters of an LWP.

\begin{verbatim}
short fss_uprilim; /* Fair-share user priority limit */
short fss_upri;   /* Fair-share user priority */
\end{verbatim}

When using the \texttt{priocntl()} \texttt{PC_SETPARMS} or \texttt{PC_GETPARMS} commands, if \texttt{pc_cid} specifies the fair-share class, the data in the \texttt{pc_clparms} buffer is in this format.

For the \texttt{priocntl()} \texttt{PC_GETPARMS} command, if \texttt{pc_cid} specifies the fair-share class and more than one fair-share LWP is specified, the scheduling parameters of the fair-share LWP with the highest \texttt{fs_upri} value among the specified LWPs is returned and the LWP ID of this LWP is returned by the \texttt{priocntl()} call. If there is more than one LWP sharing the highest user priority, the one returned is implementation-dependent.

Any fair-share LWP can lower its own \texttt{fss_uprilim} (or that of another LWP with the same user ID). Only a fair-share LWP with superuser privileges can raise an \texttt{fss_uprilim}. When changing the class of an LWP to fair-share from some other class, superuser privileges are required to set the initial \texttt{fss_uprilim} to a value greater than 0. Attempts by a non-superuser LWP to raise an \texttt{fs_uprilim} or set an initial \texttt{fs_uprilim} greater than 0 fail with a return value of \texttt{-1} and \texttt{errno} set to \texttt{EPERM}. 

\footnotesize
\begin{verbatim}
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\end{verbatim}
Any fair-share LWP can set its own fss_upri (or that of another LWP with the same user ID) to any value less than or equal to the LWP's fss_uprilim. Attempts to set the fss_upri above the fss_uprilim (and/or set the fss_uprilm below the fss_upri) result in the fss_upri being set equal to the fss_uprilim.

Either of the fss_uprilim or fss_upri members can be set to the special value FSS_NOCHANGE (defined in <sys/fsspriocntl.h>) to set one of the values without affecting the other. Specifying FSS_NOCHANGE for the fss_upri when the fss_uprilim is being set to a value below the current fss_upri causes the fss_upri to be set equal to the fss_uprilim being set. Specifying FSS_NOCHANGE for a parameter when changing the class of an LWP to fair-share (from some other class) causes the parameter to be set to a default value. The default value for the fss_uprilim is 0 and the default for the fss_upri is to set it equal to the fss_uprilim which is being set.

The fair-share user priority and user priority limit are inherited across fork() and the exec family of functions.

The fixed-priority class provides a fixed-priority preemptive scheduling policy for those LWPs requiring that the scheduling priorities do not get dynamically adjusted by the system and that the user/application have control of the scheduling priorities.

The fixed-priority class has a range of fixed-priority user priority (see fx_upri below) values that can be assigned to LWPs within the class. A fx_upri value of 0 is defined as the default base priority for the fixed-priority class. User priorities range from 0 to x where the value of x is configurable and can be determined for a specific installation by using the priocntl() PC_GETCID or PC_GETCLINFO command.

The purpose of the user priority is to provide user/application control over the scheduling of processes in the fixed-priority class. For processes in the fixed-priority class, the fx_upri value is, for all practical purposes, equivalent to the scheduling priority of the process. The fx_upri value completely determines the scheduling priority of a fixed-priority process relative to other processes within its class. Numerically higher fx_upri values represent higher priorities.

In addition to the system-wide limits on user priority (returned by the PC_GETCID and PC_GETCLINFO commands), there is a per-LWP user priority limit (see fx_uprilim below) that specifies the maximum fx_upri value that can be set for a given LWP. By default, fx_uprilim is 0.

A structure with the following member (defined in <sys/fxpriocntl.h>) defines the format used for the attribute data for the fixed-priority class.

```c
pri_t fx_maxupri; /* Maximum user priority */
```

The priocntl() PC_GETCID and PC_GETCLINFO commands return fixed-priority class attributes in the pc_clinfo buffer in this format.
The `fx_maxupri` member specifies the configured maximum user priority value for the fixed-priority class. If `fx_maxupri` is `x`, the valid range for both user priorities and user priority limits is from 0 to `x`.

A structure with the following members (defined in `<sysfxpriocntl.h>`) defines the format used to specify the fixed-priority class-specific scheduling parameters of an LWP.

```c
pri_t fx_upri; /* Fixed-priority user priority */
pri_t fx_uprilim; /* Fixed-priority user priority limit */
uint_t fx_tqsecs; /* seconds in time quantum */
int fx_tqnsecs; /* additional nanoseconds in time quantum */
```

When using the `priocntl()` `PC_SETPARMS` or `PC_GETPARMS` commands, if `pc_cid` specifies the fixed-priority class, the data in the `pc_clparms` buffer is in this format.

For the `priocntl()` `PC_GETPARMS` command, if `pc_cid` specifies the fixed-priority class and more than one fixed-priority LWP is specified, the scheduling parameters of the fixed-priority LWP with the highest `fx_upri` value among the specified LWPs is returned and the LWP ID of this LWP is returned by the `priocntl()` call. If there is more than one LWP sharing the highest user priority, the one returned is implementation-dependent.

Any fixed-priority LWP can lower its own `fx_uprilim` (or that of another LWP with the same user ID). Only a fixed-priority LWP with superuser privileges can raise a `fx_uprilim`. When changing the class of an LWP to fixed-priority from some other class, superuser privileges are required to set the initial `fx_uprilim` to a value greater than 0. Attempts by a non-superuser LWP to raise a `fx_uprilim` or set an initial `fx_uprilim` greater than 0 fail with a return value of -1 and `errno` set to `EPERM`.

Any fixed-priority LWP can set its own `fx_upri` (or that of another LWP with the same user ID) to any value less than or equal to the LWP's `fx_uprilim`. Attempts to set the `fx_upri` above the `fx_uprilim` (and/or set the `fx_uprilim` below the `fx_upri`) result in the `fx_upri` being set equal to the `fx_uprilim`.

Either of the `fx_uprilim` or `fx_upri` members can be set to the special value `FX_NOCHANGE` (defined in `<sysfxpriocntl.h>`) to set one of the values without affecting the other. Specifying `FX_NOCHANGE` for the `fx_upri` when the `fx_uprilim` is being set to a value below the current `fx_upri` causes the `fx_upri` to be set equal to the `fx_uprilim` being set. Specifying `FX_NOCHANGE` for a parameter when changing the class of an LWP to fixed-priority (from some other class) causes the parameter to be set to a default value. The default value for the `fx_uprilim` is 0 and the default for the `fx_upri` is to set it equal to the `fx_uprilim` that is being set. The default for time quantum is dependent on the `fx_upri` and on the system configuration; see `fx_dpthl(4)`.

The `fx_tqsecs` and `fx_tqnsecs` members are used for getting or setting the time quantum associated with an LWP or group of LWPs. `fx_tqsecs` is the number of seconds in the time quantum and `fx_tqnsecs` is the number of additional
nanoseconds in the quantum. For example, setting $fx\_tqsecs$ to 2 and $fx\_tqnsecs$ to 500,000,000 (decimal) would result in a time quantum of two and one-half seconds. Specifying a value of 1,000,000,000 or greater in the $fx\_tqnsecs$ member results in an error return with $errno$ set to EINVAL. Although the resolution of the $tq\_nsecs$ member is very fine, the specified time quantum length is rounded up by the system to the next integral multiple of the system clock’s resolution. The maximum time quantum that can be specified is implementation-specific and equal to $\text{INT\_MAX}$ ticks (defined in `<limits.h>`). Requesting a quantum greater than this maximum results in an error return with $errno$ set to ERANGE, although infinite quantums can be requested using a special value as explained below. Requesting a time quantum of 0 (setting both $fx\_tqsecs$ and $fx\_tqnsecs$ to 0) results in an error return with $errno$ set to EINVAL.

The $fx\_tqnsecs$ member can also be set to one of the following special values (defined in `<sys/fxpriocntl.h>`), in which case the value of $fx\_tqsecs$ is ignored:

- $FX\_TQINF$ Set an infinite time quantum.
- $FX\_TQDEF$ Set the time quantum to the default for this priority (see $fx\_dptbl$(4)).
- $FX\_NOCHANGE$ Do not set the time quantum. This value is useful in changing the user priority of an LWP without affecting the time quantum. Specifying this value when changing the class of an LWP to fixed-priority from some other class is equivalent to specifying $FX\_TQDEF$.

When using the `priocntl()` `PC_SETXPARMS` or `PC_GETXPARMS` commands, the first argument after the command code must be the class name of the fixed-priority class (FX). The next arguments are formed as (key, value) pairs, terminated by a 0 key. The definition for the keys of the fixed-priority class can be found in `<sys/fxpriocntl.h>`. A repeated specification of the same key results in an error return and $errno$ set to EINVAL.

<table>
<thead>
<tr>
<th>Key</th>
<th>Value Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>FX_KY_UPRILIM</td>
<td>pri_t</td>
<td>user priority limit</td>
</tr>
<tr>
<td>FX_KY_UPRI</td>
<td>pri_t</td>
<td>user priority</td>
</tr>
<tr>
<td>FX_KY_TQSECS</td>
<td>uint_t</td>
<td>seconds in time quantum</td>
</tr>
<tr>
<td>FX_KY_TQNSECS</td>
<td>int</td>
<td>nanoseconds in time quantum</td>
</tr>
</tbody>
</table>

When using the `priocntl()` `PC_GETXPARMS` command, the value associated with the key is always a pointer to a scheduling parameter of the value type shown in the table above. In contrast, when using the `priocntl()` `PC_SETXPARMS` command, the scheduling parameter is given as a direct value.
A `priocntl()` command with the class name (FX) and without a following (key, value) pair will set or reset all realtime scheduling parameters of the target process(es) to their default values. Changing the class of an LWP to fixed-priority from some other class causes the parameters to be set to their default values. The default value for the user priority limit (FX_KY_UPRILIM) is 0. The default value for the user priority (FX_KY_UPRI) is equal to the user priority limit (FX_KY_UPRILIM) that is being set. A default time quantum (FX_TQDEF) is assigned to each priority class (see `fx_dptbl(4)`).

The value associated with FX_KY_TQSECS is the number of seconds in the time quantum. The value associated with FX_KY_TQNSECS is the number of nanoseconds in the quantum. Specifying a value of 1,000,000,000 or greater for the number of nanoseconds results in an error return and `errno` is set to EINVAL. The specified time quantum is rounded up by the system to the next integral multiple of the system clock's resolution. The maximum time quantum that can be specified is implementation-specific and equal to INT_MAX ticks, defined in `<limits.h>`. Requesting a quantum greater than this maximum results in an error return and `errno` is set to ERANGE. If seconds (FX_KY_TQSECS) but no nanoseconds (FX_KY_TQNSECS) are supplied, the number of nanoseconds is set to 0. If nanoseconds (FX_KY_TQNSECS) but no seconds (FX_KY_TQSECS) are supplied, the number of seconds is set to 0. A time quantum of 0 (seconds and nanoseconds are 0) results in an error return with `errno` set to EINVAL. Special values for FX_KY_TQSECS are FX_TQINF and FX_TQDEF (as described above). The `priocntl()` command `PC_SETXPARMS` knows no special value FX_NOCHANGE.

The fixed-priority user priority and user priority limit are inherited across `fork(2)` and the exec family of functions (see `exec(2)`).

**RETURN VALUES**

Unless otherwise noted above, `priocntl()` returns 0 on success. On failure, `priocntl()` returns -1 and sets `errno` to indicate the error.

**ERRORS**

The `priocntl()` function will fail if:

- **EINVAL** An attempt to change the class of an LWP failed because of insufficient resources other than memory (for example, class-specific kernel data structures).
- **EFAULT** One of the arguments points to an illegal address.
- **EINVAL** The argument `cmd` was invalid, an invalid or unconfigured class was specified, or one of the parameters specified was invalid.
- **ENOMEM** An attempt to change the class of an LWP failed because of insufficient memory.
- **EPERM** The effective user of the calling LWP is not superuser.
- **ERANGE** The requested time quantum is out of range.
- **ESRCH** None of the specified LWPs exist.
SEE ALSO

priocntl(1), dispadmin(1M), init(1M), _lwp_create(2), exec(2), fork(2),
nice(2), priocntlset(2), fx_dptbl(4), rt_dptbl(4)

System Administration Guide: Basic Administration

Programming Interfaces Guide
priocntlset(2)

NAME
priocntlset – generalized process scheduler control

SYNOPSIS
#include <sys/types.h>
#include <sys/procset.h>
#include <sys/priocntl.h>
#include <sys/rtpriocntl.h>
#include <sys/tspriocntl.h>

long priocntlset(procset_t *psp, int cmd, /* arg */ ...);

DESCRIPTION
The priocntlset() function changes the scheduling properties of running processes. priocntlset() has the same functions as the priocntl() function, but a more general way of specifying the set of processes whose scheduling properties are to be changed.

cmd specifies the function to be performed. arg is a pointer to a structure whose type depends on cmd. See priocntl(2) for the valid values of cmd and the corresponding arg structures.

psp is a pointer to a procset structure, which priocntlset() uses to specify the set of processes whose scheduling properties are to be changed. The procset structure contains the following members:

idop_t   p_op;  /* operator connecting left/right sets */
idtype_t p_lidtype;  /* left set ID type */
id_t     p_lid;    /* left set ID */
idtype_t p_ridtype;  /* right set ID type */
id_t     p_rid;    /* right set ID */

The p_lidtype and p_lid members specify the ID type and ID of one (“left”) set of processes; the p_ridtype and p_rid members specify the ID type and ID of a second (“right”) set of processes. ID types and IDs are specified just as for the priocntl() function. The p_op member specifies the operation to be performed on the two sets of processes to get the set of processes the function is to apply to. The valid values for p_op and the processes they specify are:

POP_DIFF  Set difference: processes in left set and not in right set.
POP_AND   Set intersection: processes in both left and right sets.
POP_OR    Set union: processes in either left or right sets or both.
POP_XOR   Set exclusive-or: processes in left or right set but not in both.

The following macro, which is defined in <procset.h>, offers a convenient way to initialize a procset structure:

#define setprocset(psp, op, ltype, lid, rtype, rid) \
(psp)->p_op  = (op), \
(psp)->p_lidtype = (ltype), \
(psp)->p_lid = (lid), \
(psp)->p_ridtype = (rtype), \
(psp)->p_rid = (rid),
Unless otherwise noted above, `priocntlset()` returns 0 on success. Otherwise, it returns −1 and sets `errno` to indicate the error.

The `priocntlset()` function will fail if:

- **EAGAIN**: An attempt to change the class of a process failed because of insufficient resources other than memory (for example, class-specific kernel data structures).
- **EFAULT**: One of the arguments points to an illegal address.
- **EINVAL**: The argument `cmd` was invalid, an invalid or unconfigured class was specified, or one of the parameters specified was invalid.
- **ENOMEM**: An attempt to change the class of a process failed because of insufficient memory.
- **EPERM**: The effective user of the calling process is not super-user.
- **ERANGE**: The requested time quantum is out of range.
- **ESRCH**: None of the specified processes exist.

**SEE ALSO** `priocntl(1), priocntl(2)`
### processor_bind(2)

**NAME**
processor_bind – bind LWPs to a processor

**SYNOPSIS**
```c
#include <sys/types.h>
#include <sys/processor.h>
#include <sys/procset.h>

int processor_bind(idtype_t idtype, id_t id, processorid_t processorid, processorid_t *obind);
```

**DESCRIPTION**
The `processor_bind()` function binds the LWP (lightweight process) or set of LWPs specified by `idtype` and `id` to the processor specified by `processorid`. If `obind` is not NULL, this function also sets the `processorid_t` variable pointed to by `obind` to the previous binding of one of the specified LWPs, or to `PBIND_NONE` if the selected LWP was not bound.

If `idtype` is `P_PID`, the binding affects all LWPs of the process with process ID (PID) `id`.

If `idtype` is `P_LWPID`, the binding affects the LWP of the current process with LWP ID `id`.

If `idtype` is `P_TASKID`, the binding affects all LWPs of all processes with task ID `id`.

If `idtype` is `P_PROJID`, the binding affects all LWPs of all processes with project ID `id`.

If `id` is `P_MYID`, the specified LWP, process, task, or process is the current one.

If `processorid` is `PBIND_NONE`, the processor bindings of the specified LWPs are cleared.

If `processorid` is `PBIND_QUERY`, the processor bindings are not changed.

The effective user of the calling process must be superuser, or its real or effective user ID must match the real or effective user ID of the LWPs being bound. If the calling process does not have permission to change all of the specified LWPs, the bindings of the LWPs for which it does have permission will be changed even though an error is returned.

Processor bindings are inherited across `fork(2)` and `exec(2)`.

**RETURN VALUES**
Upon successful completion, 0 is returned. Otherwise, -1 is returned and `errno` is set to indicate the error.

**ERRORS**
The `processor_bind()` function will fail if:

- **EFAULT** The location pointed to by `obind` was not NULL and not writable by the user.
- **EINVAL** The specified processor is not on-line, or the `idtype` argument was not `P_PID`, `P_LWPID`, `P_PROJID`, or `P_TASKID`.
- **EPERM** The effective user of the calling process is not superuser, and its real or effective user ID does not match the real or effective user ID of one of the LWPs being bound.
ESRCH

No processes, LWPs, or tasks were found to match the criteria specified by idtype and id.

ATTRIBUTES

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Stable</td>
</tr>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO

psradm(1M), psrinfo(1M), exec(2), fork(2), p_online(2), pset_bind(2),
sysconf(3C), project(4)
processor_info(2)

NAME    processor_info – determine type and status of a processor

SYNOPSIS
#include <sys/types.h>
#include <sys/processor.h>

int processor_info(processorid_t processorid, processor_info_t *infop);

DESCRIPTION
The processor_info() function returns the status of the processor specified by
processorid in the processor_info_t structure pointed to by infop.

The structure processor_info_t contains the following members:

int pi_state;
char pi_processor_type[PI_TYPELEN];
char pi_fputypes[PI_FPUTYPE];
int pi_clock;

The pi_state member is the current state of the processor, either P_ONLINE,
P_OFFLINE, or P_POWEROFF.

The pi_processor_type member is a null-terminated ASCII string specifying the
type of the processor.

The pi_fputypes member is a null-terminated ASCII string containing the
comma-separated types of floating-point units (FPUs) attached to the processor. This
string will be empty if no FPU is attached.

The pi_clock member is the processor clock frequency rounded to the nearest
megahertz. It may be 0 if not known.

RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, –1 is returned and errno is set
to indicate the error.

ERRORS
The processor_info() function will fail if:

EINVAL    An non-existent processor ID was specified.
EFAULT     The processor_info_t structure pointed to by infop was not
            writable by the user.

SEE ALSO
psradm(1M), psrinfo(1M), p_online(2), sysconf(3C)

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#include <unistd.h>
void profil(unsigned short *buff, unsigned int bufsiz, unsigned int offset, unsigned int scale);

The profil() function provides CPU-use statistics by profiling the amount of CPU time expended by a program. The profil() function generates the statistics by creating an execution histogram for a current process. The histogram is defined for a specific region of program code to be profiled, and the identified region is logically broken up into a set of equal size subdivisions, each of which corresponds to a count in the histogram. With each clock tick, the current subdivision is identified and its corresponding histogram count is incremented. These counts establish a relative measure of how much time is being spent in each code subdivision. The resulting histogram counts for a profiled region can be used to identify those functions that consume a disproportionately high percentage of CPU time.

The buff argument is a buffer of bufsiz bytes in which the histogram counts are stored in an array of unsigned short int. Once one of the counts reaches 32767 (the size of a short int), profiling stops and no more data is collected.

The offset, scale, and bufsiz arguments specify the region to be profiled.

The offset argument is effectively the start address of the region to be profiled.

The scale argument is a contraction factor that indicates how much smaller the histogram buffer is than the region to be profiled. More precisely, scale is interpreted as an unsigned 16-bit fixed-point fraction with the decimal point implied on the left. Its value is the reciprocal of the number of bytes in a subdivision, per byte of histogram buffer. Since there are two bytes per histogram counter, the effective ratio of subdivision bytes per counter is one half the scale.

The values of scale are as follows:

- the maximum value of scale, 0xffff (approximately 1), maps subdivisions 2 bytes long to each counter.
- the minimum value of scale (for which profiling is performed), 0x0002 (1/32,768), maps subdivision 65,536 bytes long to each counter.
- the default value of scale (currently used by cc -gp), 0x4000, maps subdivisions 8 bytes long to each counter.

The values are used within the kernel as follows: when the process is interrupted for a clock tick, the value of offset is subtracted from the current value of the program counter (pc), and the remainder is multiplied by scale to derive a result. That result is used as an index into the histogram array to locate the cell to be incremented. Therefore, the cell count represents the number of times that the process was executing code in the subdivision associated with that cell when the process was interrupted.
The value of scale can be computed as \((RATIO \times 0.000001)\), where \(RATIO\) is the desired ratio of \(bufsiz\) to profiled region size, and has a value between 0 and 1. Qualitatively speaking, the closer \(RATIO\) is to 1, the higher the resolution of the profile information.

The value of \(bufsiz\) can be computed as \((size\_of\_region\_to\_be\_profiled \times RATIO)\).

Profiling is turned off by giving a scale value of 0 or 1, and is rendered ineffective by giving a \(bufsiz\) value of 0. Profiling is turned off when one of the \texttt{exec} family of functions (see \texttt{exec(2)}) is executed, but remains on in both child and parent processes after a \texttt{fork(2)}. Profiling is turned off if a buff update would cause a memory fault.

Usage

The \texttt{pcsample(2)} function should be used when profiling dynamically-linked programs and 64-bit programs.

See Also

exec(2), fork(2), pcsample(2), times(2), monitor(3C), prof(5)

Notes

In Solaris releases prior to 2.6, calling \texttt{profil()} in a multithreaded program would impact only the calling LWP; the profile state was not inherited at LWP creation time. To profile a multithreaded program with a global profile buffer, each thread needed to issue a call to \texttt{profil()} at threads start-up time, and each thread had to be a bound thread. This was cumbersome and did not easily support dynamically turning profiling on and off. In Solaris 2.6, the \texttt{profil()} system call for multithreaded processes has global impact — that is, a call to \texttt{profil()} impacts all LWPs/threads in the process. This may cause applications that depend on the previous per-LWP semantic to break, but it is expected to improve multithreaded programs that wish to turn profiling on and off dynamically at runtime.
pset_bind (2)

NAME

pset_bind – bind LWPs to a set of processors

SYNOPSIS

#include <sys/pset.h>

int pset_bind(psetid_t pset, idtype_t idtype, id_t id, psetid_t *opset);

DESCRIPTION

The pset_bind() function binds the LWP or set of LWPs specified by idtype and id to the processor set specified by pset. If obind is not NULL, pset_bind() sets the
psetid_t variable pointed to by opset to the previous processor set binding of one of the specified LWP, or to PS_NONE if the selected LWP was not bound.

If idtype is P_PID, the binding affects all LWPs of the process with process ID (PID) id.

If idtype is P_LWPID, the binding affects the LWP of the current process with LWP ID id.

If idtype is P_TASKID, the binding affects all LWPs of all processes with task ID id.

If idtype is P_PROJID, the binding affects all LWPs of all processes with project ID id.

If id is P_MYID, the specified LWP, process, task, or process is the current one.

If pset is PS_NONE, the processor set bindings of the specified LWPs are cleared.

If pset is PS_QUERY, the processor set bindings are not changed.

If pset is PS_MYID, the specified LWPs are bound to the same processor set as the caller. If the caller is not bound to a processor set, the processor set bindings are cleared.

The effective user of the calling process must be superuser, or its real or effective user ID must match the real or effective user ID of the LWPs being bound, or pset must be PS_QUERY. If the calling process does not have permission to change all of the specified LWPs, the bindings of the LWPs for which it does have permission will be changed even though an error is returned.

If the processor set type of pset is PS_PRIVATE (see pset_info(2)), the effective user of the calling process must be superuser.

LWPs that have been bound to a processor with processor_bind(2) may also be bound to a processor set if the processor is part of the processor set. If this occurs, the binding to the processor remains in effect. If the processor binding is later removed, the processor set binding becomes effective.

Processor set bindings are inherited across fork(2) and exec(2).

RETURN VALUES

Upon successful completion, 0 is returned. Otherwise, -1 is returned and errno is set to indicate the error.

ERRORS

The pset_bind() function will fail if:
EBUSY  One of the LWPs is bound to a processor, and the specified processor set does not include that processor.

EFAULT  The location pointed to by `opset` was not NULL and not writable by the user.

EINVAL  An invalid processor set ID was specified; or `idtype` was not `P_PID`, `P_LWPID`, `P_PROJID`, or `P_TASKID`.

EPERM  The effective user of the calling process is not superuser and either the real or effective user ID of the calling process does not match the real or effective user ID of one of the LWPs being bound, or the processor set from which one or more of the LWPs are being unbound has the `PSET_NOESCAPE` attribute set. See `pset_setattr(2)` for more information about processor set attributes.

ESRCH  No processes, LWPs, or tasks were found to match the criteria specified by `idtype` and `id`.

**ATTRIBUTES**

See `attributes(5)` for descriptions of the following attributes:

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</table>

**SEE ALSO**

`pbind(1M), pset(1M), exec(2), fork(2), processor_bind(2), pset_create(2), pset_info(2), pset_setattr(2), pset_getloadavg(3C), project(4), attributes(5)`
pset_create(2)

NAME
pset_create, pset_destroy, pset_assign – manage sets of processors

SYNOPSIS
#include <sys/pset.h>

int pset_create(psetid_t *newpset);
int pset_destroy(psetid_t pset);
int pset_assign(psetid_t pset, processorid_t cpu, psetid_t *opset);

DESCRIPTION
These functions control the creation and management of sets of processors. Processor sets allow a subset of the system’s processors to be set aside for exclusive use by specified LWPs and processes. The binding of LWPs and processes to processor sets is controlled by pset_bind(2).

The pset_create() function creates an empty processor set that contains no processors. On successful return, newpset will contain the ID of the new processor set.

The pset_destroy() function destroys the processor set pset, releasing its constituent processors and processes. If pset is PS_MYID, the processor set to which the caller is bound is destroyed.

The pset_assign() function assigns the processor cpu to the processor set pset. A processor that has been assigned to a processor set will run only LWPs and processes that have been explicitly bound to that processor set, unless another LWP requires a resource that is only available on that processor.

On successful return, if opset is non-null, opset will contain the processor set ID of the former processor set of the processor.

If pset is PS_NONE, pset_assign() releases processor cpu from its current processor set.

If pset is PS_QUERY, pset_assign() makes no change to processor sets, but returns the current processor set ID of processor cpu in opset.

If pset is PS_MYID, processor cpu is assigned to the processor set to which the caller belongs. If the caller does not belong to a processor set, processor cpu is released from its current processor set.

These functions are restricted to super-user use, except for pset_assign() when pset is PS_QUERY.

RETURN VALUES
Upon successful completion, these functions return 0. Otherwise, −1 is returned and errno is set to indicate the error.

ERRORS
These functions will fail if:

EBUSY The processor could not be moved to the specified processor set.

EFAULT The location pointed to by newpset was not writable by the user, or the location pointed to by opset was not NULL and not writable by the user.
The specified processor does not exist, the specified processor is not on-line, or an invalid processor set was specified.

There was insufficient space for pset_create to create a new processor set.

The effective user of the calling process is not super-user.

See attributes(5) for descriptions of the following attributes:

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Processors belonging to different processor sets of type PS_SYSTEM (see pset_info(2)) cannot be assigned to the same processor set of type PS_PRIVATE. If this is attempted, pset_assign() will fail and set errno to EINVAL.

Processors with LWPs bound to them using processor_bind(2) cannot be assigned to a new processor set. If this is attempted, pset_assign() will fail and set errno to EBUSY.
NAME pset_info – get information about a processor set

SYNOPSIS
#include <sys/pset.h>

int pset_info(psetid_t pset, int *type, uint_t *numcpus,
              processorid_t *cpulist);

DESCRIPTION
The pset_info() function returns information on the processor set pset.

If type is non-null, then on successful completion the type of the processor set will be stored in the location pointed to by type. Processor set types can have the following values:

- **PS_SYSTEM**: The processor set was created by the system. Processor sets of this type cannot be modified or removed by the user, but LWPs and processes can be bound to them using pset_bind(2).

- **PS_PRIVATE**: The processor set was created by pset_create(2) and can be modified by pset_assign(2) and removed by pset_destroy(2). LWPs and processes can also be bound to this processor set using pset_bind().

If numcpus is non-null, then on successful completion the number of processors in the processor set will be stored in the location pointed to by numcpus.

If numcpus and cpulist are both non-null, then cpulist points to a buffer where a list of processors assigned to the processor set is to be stored, and numcpus points to the maximum number of processor IDs the buffer can hold. On successful completion, the list of processors up to the maximum buffer size is stored in the buffer pointed to by cpulist.

If pset is **PS_NONE**, the list of processors not assigned to any processor set will be stored in the buffer pointed to by cpulist, and the number of such processors will be stored in the location pointed to by numcpus. The location pointed to by type will be set to **PS_NONE**.

If pset is **PS_MYID**, the processor list and number of processors returned will be those of the processor set to which the caller is bound. If the caller is not bound to a processor set, the result will be equivalent to setting pset to **PS_NONE**.

RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, -1 is returned and errno is set to indicate the error.

ERRORS
The pset_info() function will fail if:

- **EFAULT**: The location pointed to by type, numcpus, or cpulist was not null and not writable by the user.

- **EINVAL**: An invalid processor set ID was specified.
pset_info(2)

ATTRIBUTES

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Stable</td>
</tr>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO

psrinfo(1M), psrset(1M), processor_info(2), pset_assign(2), pset_bind(2), pset_create(2), pset_destroy(2), pset_getloadavg(3C), attributes(5)
NAME
pset_list – get list of processor sets

SYNOPSIS
#include <sys/pset.h>

int pset_list(psetid_t *psetlist, uint_t *numpsets);

DESCRIPTION
The pset_list function returns a list of processor sets in the system.

If numpsets is non-null, then on successful completion the number of processor sets in
the system will be stored in the location pointed to by numpsets.

If numpsets and psetlist are both non-null, then psetlist points to a buffer where a list
of processor sets in the system is to be stored, and numpsets points to the maximum
number of processor set IDs the buffer can hold. On successful completion, the list of
processor sets up to the maximum buffer size is stored in the buffer pointed to by
psetlist.

RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, -1 is returned and errno is set to
indicate the error.

ERRORS
The pset_list() function will fail if:
EFAULT The location pointed to by psetlist or numpsets was not null and not
writable by the user.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

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<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
pset(1M), processor_info(2), pset_bind(2), pset_create(2), pset_info(2),
pset_getloadavg(3C), attributes(5)
pset_setattr(2)

NAME
pset_setattr, pset_getattr – set or get processor set attributes

SYNOPSIS
#include <sys/pset.h>

int pset_setattr(psetid_t pset, uint_t attr);
int pset_getattr(psetid_t pset, uint_t *attr);

DESCRIPTION
The pset_setattr() function sets attributes of the processor set specified by pset. The bitmask of attributes to be set or cleared is specified by attr.

The pset_getattr function returns attributes of the processor set specified by pset. On successful return, attr will contain the bitmask of attributes for the specified processor set.

The value of the attr argument is the bitwise inclusive-OR of these attributes, defined in <sys/pset.h>:

PSET_NOESCAPE Unbinding of LWPs from the processor set with this attribute requires superuser privileges.

The binding of LWPs and processes to processor sets is controlled by pset_bind(2). When PSET_NOESCAPE attribute is cleared, a process calling pset_bind() can clear the processor set binding of any LWP whose real or effective user ID matches its own real of effective user ID. Setting PSET_NOESCAPE attribute forces pset_bind() to require superuser privileges for such an operation.

RETURN VALUES
Upon successful completion, these functions return 0. Otherwise, -1 is returned and errno is set to indicate the error.

ERRORS
These function will fail if:

EFAULT The location pointed to by attr was not writable by the user.
EINVAL An invalid processor set ID was specified.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

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<td>MT-Level</td>
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</tr>
</tbody>
</table>

SEE ALSO
psrset(1M), pset_bind(2), attributes(5)
ptrace – allows a parent process to control the execution of a child process

### SYNOPSIS

```c
#include <unistd.h>
#include <sys/types.h>

int ptrace(int request, pid_t pid, int addr, int data);
```

### DESCRIPTION

The `ptrace()` function allows a parent process to control the execution of a child process. Its primary use is for the implementation of breakpoint debugging. The child process behaves normally until it encounters a signal (see `signal(3HEAD)`), at which time it enters a stopped state and its parent is notified via the `wait(2)` function. When the child is in the stopped state, its parent can examine and modify its “core image” using `ptrace()`. Also, the parent can cause the child either to terminate or continue, with the possibility of ignoring the signal that caused it to stop.

The `request` argument determines the action to be taken by `ptrace()` and is one of the following:

<table>
<thead>
<tr>
<th>Request</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>This request must be issued by the child process if it is to be traced by its parent. It turns on the child’s trace flag that stipulates that the child should be left in a stopped state on receipt of a signal rather than the state specified by <code>func</code> (see <code>signal(3C)</code>). The <code>pid</code>, <code>addr</code>, and <code>data</code> arguments are ignored, and a return value is not defined for this request. Peculiar results ensue if the parent does not expect to trace the child.</td>
</tr>
<tr>
<td>1, 2</td>
<td>With these requests, the word at location <code>addr</code> in the address space of the child is returned to the parent process. If instruction and data space are separated, request 1 returns a word from instruction space, and request 2 returns a word from data space. If instruction and data space are not separated, either request 1 or request 2 may be used with equal results. The <code>data</code> argument is ignored. These two requests fail if <code>addr</code> is not the start address of a word, in which case −1 is returned to the parent process and the parent's <code>errno</code> is set to EIO.</td>
</tr>
<tr>
<td>3</td>
<td>With this request, the word at location <code>addr</code> in the child’s user area in the system’s address space (see <code>&lt;sys/user.h&gt;</code>) is returned to the parent process. The <code>data</code> argument is ignored. This request fails if <code>addr</code> is not the start address of a word or is outside the user area, in which case −1 is returned to the parent process and the parent’s <code>errno</code> is set to EIO.</td>
</tr>
<tr>
<td>4, 5</td>
<td>With these requests, the value given by the <code>data</code> argument is written into the address space of the child at location <code>addr</code>. If instruction and data space are separated, request 4 writes a word into instruction space, and request 5 writes a word into data space. If instruction and data space are not separated, either request 4 or request 5 may be used with equal results. On success, the value written into the address space of the child is returned to the parent process.</td>
</tr>
</tbody>
</table>

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ptrace(2)

the parent. These two requests fail if `addr` is not the start address of a word. On failure −1 is returned to the parent process and the parent’s `errno` is set to EIO.

6 With this request, a few entries in the child’s user area can be written. `data` gives the value that is to be written and `addr` is the location of the entry. The few entries that can be written are the general registers and the condition codes of the Processor Status Word.

7 This request causes the child to resume execution. If the `data` argument is 0, all pending signals including the one that caused the child to stop are canceled before it resumes execution. If the `data` argument is a valid signal number, the child resumes execution as if it had incurred that signal, and any other pending signals are canceled. The `addr` argument must be equal to 1 for this request. On success, the value of `data` is returned to the parent. This request fails if `data` is not 0 or a valid signal number, in which case −1 is returned to the parent process and the parent’s `errno` is set to EIO.

8 This request causes the child to terminate with the same consequences as `exit(2)`.

9 This request sets the trace bit in the Processor Status Word of the child and then executes the same steps as listed above for request 7. The trace bit causes an interrupt on completion of one machine instruction. This effectively allows single stepping of the child.

To forestall possible fraud, `ptrace()` inhibits the set-user-ID facility on subsequent calls to one of the `exec` family of functions (see `exec(2)`). If a traced process calls one of the `exec` functions, it stops before executing the first instruction of the new image showing signal SIGTRAP.

**ERRORS** The `ptrace()` function will fail if:

- EIO The `request` argument is an illegal number.
- EPERM The effective user of the calling process is not super-user.
- ESRCH The `pid` argument identifies a child that does not exist or has not executed a `ptrace()` call with request 0.

**USAGE** The `/proc` debugging interfaces should be used instead of `ptrace()`, which provides quite limited debugger support and is itself implemented using the `/proc` interfaces. There is no actual `ptrace()` system call in the kernel. See `proc(4)` for descriptions of the `/proc` debugging interfaces.

**ATTRIBUTES** See `attributes(5)` for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
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</table>
ptrace(2)

SEE ALSO exec(2), exit(2), wait(2), signal(3C), signal(3HEAD), attributes(5)
putmsg, putpmsg – send a message on a stream

#include <stropts.h>

int putmsg(int fildes, const struct strbuf *ctlptr, const struct strbuf *dataptr, int flags);

int putpmsg(int fildes, const struct strbuf *ctlptr, const struct strbuf *dataptr, int band, int flags);

The putmsg() function creates a message from user-specified buffer(s) and sends the message to a STREAMS file. The message may contain either a data part, a control part, or both. The data and control parts to be sent are distinguished by placement in separate buffers, as described below. The semantics of each part is defined by the STREAMS module that receives the message.

The putpmsg() function does the same thing as putmsg(), but provides the user the ability to send messages in different priority bands. Except where noted, all information pertaining to putmsg() also pertains to putpmsg().

The fildes argument specifies a file descriptor referencing an open stream. The ctlptr and dataptr arguments each point to a strbuf structure, which contains the following members:

```c
int maxlen; // not used here */
int len; // length of data */
void *buf; // ptr to buffer */
```

The ctlptr argument points to the structure describing the control part, if any, to be included in the message. The buf member in the strbuf structure points to the buffer where the control information resides, and the len member indicates the number of bytes to be sent. The maxlen member is not used in putmsg() (see getmsg(2)). In a similar manner, dataptr specifies the data, if any, to be included in the message. The flags argument indicates what type of message should be sent and is described later.

To send the data part of a message, dataptr must not be NULL, and the len member of dataptr must have a value of 0 or greater. To send the control part of a message, the corresponding values must be set for ctlptr. No data (control) part is sent if either dataptr (ctlptr) is NULL or the len member of dataptr (ctlptr) is negative.

For putmsg(), if a control part is specified, and flags is set to RS_HIPRI, a high priority message is sent. If no control part is specified, and flags is set to RS_HIPRI, putmsg() fails and sets errno to EINVAL. If flags is set to 0, a normal (non-priority) message is sent. If no control part and no data part are specified, and flags is set to 0, no message is sent, and 0 is returned.

The stream head guarantees that the control part of a message generated by putmsg() is at least 64 bytes in length.

For putpmsg(), the flags are different. The flags argument is a bitmask with the following mutually-exclusive flags defined: MSG_HIPRI and MSG_BAND. If flags is set to 0, putpmsg() fails and sets errno to EINVAL. If a control part is specified and flags
is set to MSG_HIPRI and band is set to 0, a high-priority message is sent. If flags is set to MSG_HIPRI and either no control part is specified or band is set to a non-zero value, putpmsg() fails and sets errno to EINVAL. If flags is set to MSG_BAND, then a message is sent in the priority band specified by band. If a control part and data part are not specified and flags is set to MSG_BAND, no message is sent and 0 is returned.

Normally, putmsg() will block if the stream write queue is full due to internal flow control conditions. For high-priority messages, putmsg() does not block on this condition. For other messages, putmsg() does not block when the write queue is full and O_NDELAY or O_NONBLOCK is set. Instead, it fails and sets errno to EAGAIN.

The putmsg() or putpmsg() function also blocks, unless prevented by lack of internal resources, waiting for the availability of message blocks in the stream, regardless of priority or whether O_NDELAY or O_NONBLOCK has been specified. No partial message is sent.

RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, −1 is returned and errno is set to indicate the error.

ERRORS
The putmsg() and putpmsg() functions will fail if:

- EAGAIN: A non-priority message was specified, the O_NDELAY or O_NONBLOCK flag is set and the stream write queue is full due to internal flow control conditions.
- EBADF: The fildes argument is not a valid file descriptor open for writing.
- EFAULT: The ctiptr or dataptr argument points to an illegal address.
- EINTR: A signal was caught during the execution of the putmsg() function.
- EINVAL: An undefined value was specified in flags; flags is set to RS_HIPRI and no control part was supplied; or the stream referenced by fildes is linked below a multiplexor.
- ENOSR: Buffers could not be allocated for the message that was to be created due to insufficient STREAMS memory resources.
- ENOSTR: The fildes argument is not associated with a STREAM.
- ENXIO: A hangup condition was generated downstream for the specified stream, or the other end of the pipe is closed.
- EPILE or EIO: The fildes argument refers to a STREAMS-based pipe and the other end of the pipe is closed. A SIGPIPE signal is generated for the calling process. This error condition occurs only with SUS-compliant applications. See standards(5).
- ERANGE: The size of the data part of the message does not fall within the range specified by the maximum and minimum packet sizes of the topmost stream module. This value is also returned if the control part of the message is larger than the maximum configured size of
putmsg(2)

the control part of a message, or if the data part of a message is larger than the maximum configured size of the data part of a message.

In addition, putmsg() and putpmsg() will fail if the STREAM head had processed an asynchronous error before the call. In this case, the value of errno does not reflect the result of putmsg() or putpmsg() but reflects the prior error.

The putpmsg() function will fail if:

EINVAL The flags argument is set to MSG_HIPRI and band is non-zero.

SEE ALSO

intro(2), getmsg(2), poll(2), read(2), write(2), standards(5)

STREAMS Programming Guide
The `read()` function attempts to read `nbyte` bytes from the file associated with the open file descriptor, `fildes`, into the buffer pointed to by `buf`.

If `nbyte` is 0, `read()` returns 0 and has no other results.

On files that support seeking (for example, a regular file), the `read()` starts at a position in the file given by the file offset associated with `fildes`. The file offset is incremented by the number of bytes actually read.

Files that do not support seeking (for example, terminals) always read from the current position. The value of a file offset associated with such a file is undefined.

If `fildes` refers to a socket, `read()` is equivalent to `recv(3SOCKET)` with no flags set.

No data transfer will occur past the current end-of-file. If the starting position is at or after the end-of-file, 0 will be returned. If the file refers to a device special file, the result of subsequent `read()` requests is implementation-dependent.

When attempting to read from a regular file with mandatory file/record locking set (see `chmod(2)`), and there is a write lock owned by another process on the segment of the file to be read:

- If `O_NDELAY` or `O_NONBLOCK` is set, `read()` returns -1 and sets `errno` to `EAGAIN`.
- If `O_NDELAY` and `O_NONBLOCK` are clear, `read()` sleeps until the blocking record lock is removed.

When attempting to read from an empty pipe (or FIFO):

- If no process has the pipe open for writing, `read()` returns 0 to indicate end-of-file.
- If some process has the pipe open for writing and `O_NDELAY` is set, `read()` returns 0.
- If some process has the pipe open for writing and `O_NONBLOCK` is set, `read()` returns -1 and sets `errno` to `EAGAIN`.
- If `O_NDELAY` and `O_NONBLOCK` are clear, `read()` blocks until data is written to the pipe or the pipe is closed by all processes that had opened the pipe for writing.
When attempting to read a file associated with a terminal that has no data currently available:

- If O_NDELAY is set, read() returns 0.
- If O_NONBLOCK is set, read() returns -1 and sets errno to EAGAIN.
- If O_NDELAY and O_NONBLOCK are clear, read() blocks until data become available.

When attempting to read a file associated with a socket or a stream that is not a pipe, a FIFO, or a terminal, and the file has no data currently available:

- If O_NDELAY or O_NONBLOCK is set, read() returns -1 and sets errno to EAGAIN.
- If O_NDELAY and O_NONBLOCK are clear, read() blocks until data becomes available.

The read() function reads data previously written to a file. If any portion of a regular file prior to the end-of-file has not been written, read() returns bytes with value 0. For example, lseek(2) allows the file offset to be set beyond the end of existing data in the file. If data is later written at this point, subsequent reads in the gap between the previous end of data and the newly written data will return bytes with value 0 until data is written into the gap.

For regular files, no data transfer will occur past the offset maximum established in the open file description associated with fildes.

Upon successful completion, where nbyte is greater than 0, read() will mark for update the st_atime field of the file, and return the number of bytes read. This number will never be greater than nbyte. The value returned may be less than nbyte if the number of bytes left in the file is less than nbyte, if the read() request was interrupted by a signal, or if the file is a pipe or FIFO or special file and has fewer than nbyte bytes immediately available for reading. For example, a read() from a file associated with a terminal may return one typed line of data.

If a read() is interrupted by a signal before it reads any data, it will return -1 with errno set to EINTR.

If a read() is interrupted by a signal after it has successfully read some data, it will return the number of bytes read.

A read() from a STREAMS file can read data in three different modes: byte-stream mode, message-nondiscard mode, and message-discard mode. The default is byte-stream mode. This can be changed using the I_SRDOPT ioctl(2) request, and can be tested with the I_GRDOPT ioctl(). In byte-stream mode, read() retrieves data from the STREAM until as many bytes as were requested are transferred, or until there is no more data to be retrieved. Byte-stream mode ignores message boundaries.
In STREAMS message-nondiscard mode, `read()` retrieves data until as many bytes as were requested are transferred, or until a message boundary is reached. If `read()` does not retrieve all the data in a message, the remaining data is left on the STREAM, and can be retrieved by the next `read()` call. Message-discard mode also retrieves data until as many bytes as were requested are transferred, or a message boundary is reached. However, unread data remaining in a message after the `read()` returns is discarded, and is not available for a subsequent `read()`, `readv()` or `getmsg()` call.

How `read()` handles zero-byte STREAMS messages is determined by the current read mode setting. In byte-stream mode, `read()` accepts data until it has read `nbyte` bytes, or until there is no more data to read, or until a zero-byte message block is encountered. The `read()` function then returns the number of bytes read, and places the zero-byte message block back on the STREAM to be retrieved by the next `read()`, `readv()` or `getmsg()` call. In message-nondiscard mode or message-discard mode, a zero-byte message returns 0 and the message is removed from the STREAM. When a zero-byte message is read as the first message on a STREAM, the message is removed from the STREAM and 0 is returned, regardless of the read mode.

A `read()` from a STREAMS file returns the data in the message at the front of the STREAM head read queue, regardless of the priority band of the message.

By default, STREAMs are in control-normal mode, in which a `read()` from a STREAMS file can only process messages that contain a data part but do not contain a control part. The `read()` function fails if a message containing a control part is encountered at the STREAM head. This default action can be changed by placing the STREAM in either control-data mode or control-discard mode with the `I_SRDPT ioctl()` command. In control-data mode, `read()` converts any control part to data and passes it to the application before passing any data part originally present in the same message. In control-discard mode, `read()` discards message control parts but returns to the process any data part in the message.

In addition, `read()` and `readv()` will fail if the STREAM head had processed an asynchronous error before the call. In this case, the value of `errno` does not reflect the result of `read()` or `readv()` but reflects the prior error. If a hangup occurs on the STREAM being read, `read()` continues to operate normally until the STREAM head read queue is empty. Thereafter, it returns 0.

`readv()` function is equivalent to `read()`, but places the input data into the `iov` buffers specified by the members of the `iov` array: `iov0`, `iov1`, ..., `iov[iovcnt-1]`. The `iovcnt` argument is valid if greater than 0 and less than or equal to `IOV_MAX`.

The `iov` structure contains the following members:

```c
    caddr_t     iov_base;
    int         iov_len;
```

Each `iov` entry specifies the base address and length of an area in memory where data should be placed. The `readv()` function always fills an area completely before proceeding to the next.
Upon successful completion, `readv()` marks for update the `st_atime` field of the file.

### `pread()`

The `pread()` function performs the same action as `read()`, except that it reads from a given position in the file without changing the file pointer. The first three arguments to `pread()` are the same as `read()` with the addition of a fourth argument `offset` for the desired position inside the file. `pread()` will read up to the maximum offset value that can be represented in an `off_t` for regular files. An attempt to perform a `pread()` on a file that is incapable of seeking results in an error.

### RETURN VALUES

Upon successful completion, `read()` and `readv()` return a non-negative integer indicating the number of bytes actually read. Otherwise, the functions return `-1` and set `errno` to indicate the error.

### ERRORS

The `read()`, `readv()`, and `pread()` functions will fail if:

- **EAGAIN** – Mandatory file/record locking was set, `O_NDELAY` or `O_NONBLOCK` was set, and there was a blocking record lock; total amount of system memory available when reading using raw I/O is temporarily insufficient; no data is waiting to be read on a file associated with a tty device and `O_NONBLOCK` was set; or no message is waiting to be read on a stream and `O_NDELAY` or `O_NONBLOCK` was set.

- **EBADF** – The `fd` argument is not a valid file descriptor open for reading.

- **EBADMSG** – Message waiting to be read on a stream is not a data message.

- **EDEADLK** – The read was going to go to sleep and cause a deadlock to occur.

- **EINTR** – A signal was caught during the read operation and no data was transferred.

- **EINVAL** – An attempt was made to read from a stream linked to a multiplexor.

- **EIO** – A physical I/O error has occurred, or the process is in a background process group and is attempting to read from its controlling terminal, and either the process is ignoring or blocking the `SIGTIN` signal or the process group of the process is orphaned.

- **EISDIR** – The `fd` argument refers to a directory on a file system type that does not support read operations on directories.

- **ENOLCK** – The system record lock table was full, so the `read()` or `readv()` could not go to sleep until the blocking record lock was removed.

- **ENOLINK** – The `fd` argument is on a remote machine and the link to that machine is no longer active.

- **ENXIO** – The device associated with `fd` is a block special or character special file and the value of the file pointer is out of range.
The `read()` and `pread()` functions will fail if:

- **EFAULT** The `buf` argument points to an illegal address.
- **EINVAL** The `nbyte` argument overflowed an `ssize_t`.

The `read()` and `readv()` functions will fail if:

- **EOVERFLOW** The file is a regular file, `nbyte` is greater than 0, the starting position is before the end-of-file, and the starting position is greater than or equal to the offset maximum established in the open file description associated with `fd`es.

The `readv()` function may fail if:

- **EFAULT** The `iov` argument points outside the allocated address space.
- **EINVAL** The `iovcnt` argument was less than or equal to 0 or greater than `{IOV_MAX}`. (See `intro(2)` for a definition of `{IOV_MAX}`).
- **EINVAL** One of the `iov_len` values in the `iov` array was negative, or the sum of the `iov_len` values in the `iov` array overflowed an `ssize_t`.

The `pread()` function will fail and the file pointer remain unchanged if:

- **ESPIPE** The `fd`es argument is associated with a pipe or FIFO.

**USAGE** The `pread()` function has a transitional interface for 64-bit file offsets. See `lfdso(5)`.

**ATTRIBUTES**

See `attributes(5)` for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td><code>read()</code> is Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO**

`intro(2)`, `chmod(2)`, `creat(2)`, `dup(2)`, `fcntl(2)`, `getmsg(2)`, `ioctl(2)`, `lseek(2)`, `open(2)`, `pipe(2)`, `recv(3SOCKET)`, `attributes(5)`, `lfdso(5)`, `streamio(7I)`, `termio(7I)`
### NAME
readlink – read the contents of a symbolic link

### SYNOPSIS
```c
#include <unistd.h>

int readlink(const char *path, char *buf, size_t bufsize);
```

### DESCRIPTION
The `readlink()` function places the contents of the symbolic link referred to by `path` in the buffer `buf` which has size `bufsize`. If the number of bytes in the symbolic link is less than `bufsize`, the contents of the remainder of `buf` are left unchanged. If the `buf` argument is not large enough to contain the link content, the first `bufsize` bytes are placed in `buf`.

### RETURN VALUES
Upon successful completion, `readlink()` returns the count of bytes placed in the buffer. Otherwise, it returns −1, leaves the buffer unchanged, and sets `errno` to indicate the error.

### ERRORS
The `readlink()` function will fail if:

- **EACCES** Search permission is denied for a component of the path prefix of `path`.
- **EFAULT** `path` or `buf` points to an illegal address.
- **EINVAL** The `path` argument names a file that is not a symbolic link.
- **EIO** An I/O error occurred while reading from the file system.
- **ENOENT** A component of `path` does not name an existing file or `path` is an empty string.
- **ELOOP** A loop exists in symbolic links encountered during resolution of the `path` argument.
- **ENAMETOOLONG** The length of `path` exceeds [PATH_MAX], or a pathname component is longer than [NAME_MAX] while `_POSIX_NO_TRUNC` is in effect.
- **ENOTDIR** A component of the path prefix is not a directory.
- **ENOSYS** The file system does not support symbolic links.

The `readlink()` function may fail if:

- **EACCES** Read permission is denied for the directory. This condition is reported.
- **ELOOP** More than [SYMLOOP_MAX] symbolic links were encountered in resolving `path`. This condition is reported.
- **ENAMETOOLONG** As a result of encountering a symbolic link in resolution of the `path` argument, the length of the substituted pathname string exceeded [PATH_MAX]. This condition is reported.

### USAGE
Portable applications should not assume that the returned contents of the symbolic link are null-terminated.
ATTRIBUTES: See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Standard</td>
</tr>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO: stat(2), symlink(2), attributes(5), standards(5)
rename(2)

NAME rename, renameat – change the name of a file

SYNOPSIS #include <stdio.h>

int rename(const char *old, const char *new);
int renameat(int fromfd, const char *old, int tofd, const char *new);

DESCRIPTION The rename() function changes the name of a file. The old argument points to the pathname of the file to be renamed. The new argument points to the new path name of the file.

The renameat() function renames an entry in a directory, possibly moving the entry into a different directory. See fattr(5). If the old argument is an absolute path, the fromfd is ignored. Otherwise it is resolved relative to the fromfd argument rather than the current working directory. Similarly, if the new argument is not absolute, it is resolved relative to the tofd argument. If either fromfd or tofd have the value AT_FDCWD, defined in <fcntl.h>, and their respective paths are relative, the path is resolved relative to the current working directory.

Current implementation restrictions will cause the renameat() function to return an error if an attempt is made to rename an extended attribute file to a regular (non-attribute) file, or to rename a regular file to an extended attribute file.

If old and new both refer to the same existing file, the rename() and renameat() functions return successfully and performs no other action.

If old points to the pathname of a file that is not a directory, new must not point to the pathname of a directory. If the link named by new exists, it will be removed and old will be renamed to new. In this case, a link named new must remain visible to other processes throughout the renaming operation and will refer to either the file referred to by new or the file referred to as old before the operation began.

If old points to the pathname of a directory, new must not point to the pathname of a file that is not a directory. If the directory named by new exists, it will be removed and old will be renamed to new. In this case, a link named new will exist throughout the renaming operation and will refer to either the file referred to by new or the file referred to as old before the operation began. Thus, if new names an existing directory, it must be an empty directory.

The new pathname must not contain a path prefix that names old. Write access permission is required for both the directory containing old and the directory containing new. If old points to the pathname of a directory, write access permission is required for the directory named by old, and, if it exists, the directory named by new.

If the directory containing old has the sticky bit set, at least one of the following conditions listed below must be true:

- the user must own old
- the user must own the directory containing old
- old must be writable by the user
the user must be a privileged user

If new exists, and the directory containing new is writable and has the sticky bit set, at least one of the following conditions must be true:

- the user must own new
- the user must own the directory containing new
- new must be writable by the user
- the user must be a privileged user

If the link named by new exists, the file’s link count becomes zero when it is removed, and no process has the file open, then the space occupied by the file will be freed and the file will no longer be accessible. If one or more processes have the file open when the last link is removed, the link will be removed before rename() or renameat() returns, but the removal of the file contents will be postponed until all references to the file have been closed.

Upon successful completion, the rename() and renameat() functions will mark for update the st_ctime and st_mtime fields of the parent directory of each file.

RETURN VALUES

Upon successful completion, 0 is returned. Otherwise, -1 is returned and errno is set to indicate an error.

ERRORS

The rename() function will fail if:

- **EACCES** A component of either path prefix denies search permission; one of the directories containing old and new denies write permissions; or write permission is denied by a directory pointed to by old or new.
- **EBUSY** The new argument is a directory and the mount point for a mounted file system.
- **EDQUOT** The directory where the new name entry is being placed cannot be extended because the user’s quota of disk blocks on that file system has been exhausted.
- **EEXIST** The link named by new is a directory containing entries other than "." (the directory itself) and ". . ." (the parent directory).
- **EFAULT** Either old or new references an invalid address.
- **EINVAL** The new argument directory pathname contains a path prefix that names the old directory, or an attempt was made to rename a regular file to an extended attribute or from an extended attribute to a regular file.
- **EISDIR** The new argument points to a directory but old points to a file that is not a directory.
- **ELOOP** Too many symbolic links were encountered in translating the pathname.
rename(2)

ENAMETOOLONG The length of old or new exceeds PATH_MAX, or a pathname component is longer than NAME_MAX while _POSIX_NO_TRUNC is in effect.

EMLINK The file named by old is a directory, and the link count of the parent directory of new would exceed LINK_MAX.

ENOENT The link named by old does not exist, or either old or new points to an empty string.

ENOSPC The directory that would contain new cannot be extended.

ENOTDIR A component of either path prefix is not a directory, or old names a directory and new names a nondirectory file, or tofd and dirfd in renameat() do not reference a directory.

EROF S The requested operation requires writing in a directory on a read-only file system.

EXDEV The links named by old and new are on different file systems.

EIO An I/O error occurred while making or updating a directory entry.

The renameat() functions will fail if:

ENOTSUP An attempt was made to rename a regular file as an attribute file or to rename an attribute file as a regular file.

ATTRIBUTES See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>rename() is Standard; renameat() is Evolving</td>
</tr>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO chmod(2), link(2), unlink(2), attributes(5), fsattr(5)

NOTES The system can deadlock if there is a loop in the file system graph. Such a loop can occur if there is an entry in directory a, a/name1, that is a hard link to directory b, and an entry in directory b, b/name2, that is a hard link to directory a. When such a loop exists and two separate processes attempt to rename a/name1 to b/name2 and b/name2 to a/name1, the system may deadlock attempting to lock both directories for modification. Use symbolic links instead of hard links for directories.
#include <unistd.h>

int resolvepath(const char *path, char *buf, size_t bufsiz);

The `resolvepath()` function fully resolves all symbolic links in the path name `path` into a resulting path name free of symbolic links and places the resulting path name in the buffer `buf` which has size `bufsiz`. The resulting path name names the same file or directory as the original path name. All “.” components are eliminated and every non-leading “..” component is eliminated together with its preceding directory component. If leading “..” components reach to the root directory, they are replaced by “/”. If the number of bytes in the resulting path name is less than `bufsiz`, the contents of the remainder of `buf` are unspecified.

Upon successful completion, `resolvepath()` returns the count of bytes placed in the buffer. Otherwise, it returns −1, leaves the buffer unchanged, and sets `errno` to indicate the error.

The `resolvepath()` function will fail if:

- **EACCES**: Search permission is denied for a component of the path prefix of `path` or for a path prefix component resulting from the resolution of a symbolic link.
- **EFAULT**: The `path` or `buf` argument points to an illegal address.
- **EIO**: An I/O error occurred while reading from the file system.
- **ENOENT**: The `path` argument is an empty string or a component of `path` or a path name component produced by resolving a symbolic link does not name an existing file.
- **ELOOP**: Too many symbolic links were encountered in resolving `path`.
- **ENAMETOOLONG**: The length of `path` exceeds `PATH_MAX`, or a path name component is longer than `NAME_MAX`. Path name resolution of a symbolic link produced an intermediate result whose length exceeds `PATH_MAX` or a component whose length exceeds `NAME_MAX`.
- **ENOTDIR**: A component of the path prefix of `path` or of a path prefix component resulting from the resolution of a symbolic link is not a directory.

No more than `PATH_MAX` bytes will be placed in the buffer. Applications should not assume that the returned contents of the buffer are null-terminated.

**SEE ALSO**

`readlink(2)`, `realpath(3C)`
rmdir(2)

NAME       rmdir – remove a directory
SYNOPSIS   #include <unistd.h>
            int rmdir(const char *path);

DESCRIPTION The rmdir() function removes the directory named by the path name pointed to by
        path. The directory must not have any entries other than “.” and “..”.

If the directory’s link count becomes zero and no process has the directory open, the
        space occupied by the directory is freed and the directory is no longer accessible. If
        one or more processes have the directory open when the last link is removed, the “.”
        and “..” entries, if present, are removed before rmdir() returns and no new entries
        may be created in the directory, but the directory is not removed until all references to
        the directory have been closed.

Upon successful completion rmdir() marks for update the st_ctime and
        st_mtime fields of the parent directory.

RETURN VALUES Upon successful completion, 0 is returned. Otherwise, −1 is returned, errno is set to
        indicate the error, and the named directory is not changed.

ERRORS      The rmdir() function will fail if:
EACCES      Search permission is denied for a component of the
        path prefix; write permission is denied on the directory
        containing the directory to be removed; the parent
        directory has the S_ISVTX variable set and is not
        owned by the user; the directory is not owned by the
        user and is not writable by the user; or the user is not a
        super-user.
EBUSY        The directory to be removed is the mount point for a
        mounted file system.
EEXIST       The directory contains entries other than those for “.”
        and “..”.
EFAULT       The path argument points to an illegal address.
EINVAL       The directory to be removed is the current directory, or
        the final component of path is “.”.
EIO           An I/O error occurred while accessing the file system.
ELOOP         Too many symbolic links were encountered in
        translating path.
ENAMETOOLONG The length of the path argument exceeds PATH_MAX, or
        the length of a path component exceeds NAME_MAX
        while _POSIX_NO_TRUNC is in effect.
ENOENT  The named directory does not exist or is the null pathname.
ENOLINK  The path argument points to a remote machine, and the connection to that machine is no longer active.
ENOTDIR  A component of the path prefix is not a directory.
EROFS   The directory entry to be removed is part of a read-only file system.

ATTRIBUTES

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
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<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO  mkdir(1), rm(1), mkdir(2), attributes(5)
The `semctl()` function provides a variety of semaphore control operations as specified by `cmd`. The fourth argument is optional, depending upon the operation requested. If required, it is of type `union semun`, which must be explicitly declared by the application program.

```
union semun {
    int val;
    struct semid_ds *buf;
    ushort_t *array;
} arg;
```

The permission required for a semaphore operation is given as `{token}`, where `token` is the type of permission needed. The types of permission are interpreted as follows:

- `00400` READ by user
- `00200` ALTER by user
- `00040` READ by group
- `00020` ALTER by group
- `00004` READ by others
- `00002` ALTER by others

See the Semaphore Operation Permissions subsection of the DEFINITIONS section of `intro(2)` for more information. The following semaphore operations as specified by `cmd` are executed with respect to the semaphore specified by `semid` and `semnum`.

- **GETVAL**
  - Return the value of `semval` (see `intro(2)`). `{READ}`

- **SETVAL**
  - Set the value of `semval` to `arg.val`. `{ALTER}`
  - When this command is successfully executed, the `semadj` value corresponding to the specified semaphore in all processes is cleared.

- **GETPID**
  - Return the value of `(int) sempid`. `{READ}`

- **GETNCNT**
  - Return the value of `semncnt`. `{READ}`

- **GETZCNT**
  - Return the value of `semzcnt`. `{READ}`

The following operations return and set, respectively, every `semval` in the set of semaphores.

- **GETALL**
  - Place `semvals` into array pointed to by `arg.array`. `{READ}`

- **SETALL**
  - Set `semvals` according to the array pointed to by `arg.array`. `{ALTER}`
  - When this cmd is successfully executed, the `semadj` values corresponding to each specified semaphore in all processes are cleared.
The following operations are also available.

**IPC_STAT**
Place the current value of each member of the data structure associated with `semid` into the structure pointed to by `arg.buf`. The contents of this structure are defined in `intro(2)`. [READ]

**IPC_SET**
Set the value of the following members of the data structure associated with `semid` to the corresponding value found in the structure pointed to by `arg.buf`:

```c
sem_perm.uid
sem_perm.gid
sem_perm.mode /* access permission bits only */
```

This command can be executed only by a process that has an effective user ID equal to either that of super-user, or to the value of `sem_perm.cuid` or `sem_perm.uid` in the data structure associated with `semid`.

**IPC_RMID**
Remove the semaphore identifier specified by `semid` from the system and destroy the set of semaphores and data structure associated with it. This command can only be executed by a process that has an effective user ID equal to either that of super-user, or to the value of `sem_perm.cuid` or `sem_perm.uid` in the data structure associated with `semid`.

**RETURN VALUES**
Upon successful completion, the value returned depends on `cmd` as follows:

- **GETVAL**
  - the value of `semval`
- **GETPID**
  - the value of `(int) sempid`
- **GETNCNT**
  - the value of `semncnt`
- **GETZCNT**
  - the value of `semzcnt`

All other successful completions return 0; otherwise, −1 is returned and `errno` is set to indicate the error.

**ERRORS**
The `semctl()` function will fail if:

- **EACCES**
  - Operation permission is denied to the calling process (see `intro(2)`).
- **EFAULT**
  - The source or target is not a valid address in the user process.
- **EINVAL**
  - The `semid` argument is not a valid semaphore identifier; the `semnum` argument is less than 0 or greater than `sem_nsems` − 1; or the `cmd` argument is not a valid command or is `IPC_SET` and `sem_perm.uid` or `sem_perm.gid` is not valid.
- **EPERM**
  - The `cmd` argument is equal to `IPC_RMID` or `IPC_SET` and the effective user of the calling process is not super-user, or `cmd` is
equal to the value of `sem_perm.cuid` or `sem_perm.uid` in the
data structure associated with `semid`.

**EOVERFLOW** The `cmd` argument is `IPC_STAT` and `uid` or `gid` is too large to be
stored in the structure pointed to by `arg.buf`.

**ERANGE** The `cmd` argument is `SETVAL` or `SETALL` and the value to which
`semval` is to be set is greater than the system imposed maximum.

**SEE ALSO** `ipcs(1)`, `intro(2)`, `semget(2)`, `semop(2)`
semget

- get set of semaphores

```c
#include <sys/types.h>
#include <sys/ipc.h>
#include <sys/sem.h>

int semget(key_t key, int nsems, int semflg);
```

**DESCRIPTION**

The semget() function returns the semaphore identifier associated with key.

A semaphore identifier and associated data structure and set containing nsems semaphores (see intro(2)) are created for key if one of the following is true:

- key is equal to IPC_PRIVATE.
- key does not already have a semaphore identifier associated with it, and (semflg&IPC_CREAT) is true.

On creation, the data structure associated with the new semaphore identifier is initialized as follows:

- sem_perm.cuid, sem_perm.uid, sem_perm.cgid, and sem_perm.gid are set equal to the effective user ID and effective group ID, respectively, of the calling process.
- The access permission bits of sem_perm.mode are set equal to the access permission bits of semflg.
- sem_nsems is set equal to the value of nsems.
- sem_otime is set equal to 0 and sem_ctime is set equal to the current time.

**RETURN VALUES**

Upon successful completion, a non-negative integer representing a semaphore identifier is returned. Otherwise, −1 is returned and errno is set to indicate the error.

**ERRORS**

The semget() function will fail if:

- EACCES A semaphore identifier exists for key, but operation permission (see intro(2)) as specified by the low-order 9 bits of semflg would not be granted.
- EEXIST A semaphore identifier exists for key but both (semflg&IPC_CREAT) and (semflg&IPC_EXCL) are both true.
- EINVAL The nsems argument is either less than or equal to 0 or greater than the system-imposed limit; or a semaphore identifier exists for key, but the number of semaphores in the set associated with it is less than nsems and nsems is not equal to 0.
- ENOENT A semaphore identifier does not exist for key and (semflg&IPC_CREAT) is false.
- ENOSPC A semaphore identifier is to be created but the system-imposed limit on the maximum number of allowed semaphores or semaphore identifiers system-wide would be exceeded.
SEE ALSO  ipcrm(1), ipcs(1), intro(2), semctl(2), semop(2), ftok(3C)
semids – discover all semaphore identifiers

#include <sys/sem.h>

int semids(int *buf, uint_t nids, uint_t *pnids);

The semids() function copies all active semaphore identifiers from the system into the user-defined buffer specified by buf, provided that the number of such identifiers is not greater than the number of integers the buffer can contain, as specified by nids. If the size of the buffer is insufficient to contain all of the active semaphore identifiers in the system, none are copied.

Whether or not the size of the buffer is sufficient to contain all of them, the number of active semaphore identifiers in the system is copied into the unsigned integer pointed to by pnids.

If nids is 0 or less than the number of active semaphore identifiers in the system, buf is ignored.

Upon successful completion, semids() returns 0. Otherwise, -1 is returned and errno is set to indicate the error.

The semids() function will fail if:

EFAULT The buf or pnids argument points to an illegal address.

The semids() function returns a snapshot of all the active semaphore identifiers in the system. More may be added and some may be removed before they can be used by the caller.

This is sample C code indicating how to use the semids() function.

```c
void examine_semids()
{
    int *ids = NULL;
    uint_t nids = 0;
    uint_t n;
    int i;

    for (;;) {
        if (semids(ids, nids, &n) != 0) {
            perror("semids");
            exit(1);
        }
        if (n <= nids) /* we got them all */
            break;
        /* we need a bigger buffer */
        ids = realloc(ids, (nids = n) * sizeof (int));
    }
    for (i = 0; i < n; i++)
```
EXAMPLE 1 semids() example  (Continued)

        process_semid(ids[i]);
        free(ids);
    }

ATTRIBUTES  See attributes(5) for descriptions of the following attributes:

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<tr>
<td>MT-Level</td>
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</tbody>
</table>

SEE ALSO  ipcrm(1), ipcs(1), intro(2), semctl(2), semget(2), semop(2), attributes(5)
The `semop()` function is used to perform atomically an array of semaphore operations on the set of semaphores associated with the semaphore identifier specified by `semid`. The `sops` argument is a pointer to the array of semaphore-operation structures. The `nsops` argument is the number of such structures in the array.

Each `sembuf` structure contains the following members:

- `short sem_num; /* semaphore number */`
- `short sem_op; /* semaphore operation */`
- `short sem_flg; /* operation flags */`

Each semaphore operation specified by `sem_op` is performed on the corresponding semaphore specified by `semid` and `sem_num`. The permission required for a semaphore operation is given as `{token}`, where `token` is the type of permission needed. The types of permission are interpreted as follows:

- 00400 READ by user
- 00200 ALTER by user
- 00040 READ by group
- 00020 ALTER by group
- 00004 READ by others
- 00002 ALTER by others

See the Semaphore Operation Permissions section of `intro(2)` for more information.

A process maintains a value, `semadj`, for each semaphore it modifies. This value contains the cumulative effect of operations the process has performed on an individual semaphore with the SEM_UNDO flag set (so that they can be undone if the process terminates unexpectedly). The value of `semadj` can affect the behavior of calls to `semop()`, `semtimedop()`, `exit()`, and `_exit()` (the latter two functions documented on `exit(2)`), but is otherwise unobservable. See below for details.

The `sem_op` member specifies one of three semaphore operations:

1. The `sem_op` member is a negative integer; {ALTER}
   - If `semval` (see `intro(2)`) is greater than or equal to the absolute value of `sem_op`, the absolute value of `sem_op` is subtracted from `semval`. Also, if `(sem_flg&SEM_UNDO)` is true, the absolute value of `sem_op` is added to the calling process’s `semadj` value (see `exit(2)`) for the specified semaphore.
semop(2)

- If semval is less than the absolute value of sem_op and (sem_flag & IPC_NOWAIT) is true, semop() returns immediately.
- If semval is less than the absolute value of sem_op and (sem_flag & IPC_NOWAIT) is false, semop() increments the semncnt associated with the specified semaphore and suspends execution of the calling process until one of the following conditions occur:
  - The value of semval becomes greater than or equal to the absolute value of sem_op. When this occurs, the value of semncnt associated with the specified semaphore is decremented, the absolute value of sem_op is subtracted from semval and, if (sem_flag & SEM_UNDO) is true, the absolute value of sem_op is added to the calling process's semadj value for the specified semaphore.
  - The semid for which the calling process is awaiting action is removed from the system (see semctl(2)). When this occurs, errno is set to EIDRM and -1 is returned.
  - The calling process receives a signal that is to be caught. When this occurs, the value of semncnt associated with the specified semaphore is decremented, and the calling process resumes execution in the manner prescribed in signal(3C).

2. The sem_op member is a positive integer; [ALTER]

   The value of sem_op is added to semval and, if (sem_flag & SEM_UNDO) is true, the value of sem_op is subtracted from the calling process's semadj value for the specified semaphore.

3. The sem_op member is 0; [READ]

   - If semval is 0, semop() returns immediately.
   - If semval is not equal to 0 and (sem_flag & IPC_NOWAIT) is true, semop() returns immediately.
   - If semval is not equal to 0 and (sem_flag & IPC_NOWAIT) is false, semop() increments the semncnt associated with the specified semaphore and suspends execution of the calling process until one of the following occurs:
     - The value of semval becomes 0, at which time the value of semncnt associated with the specified semaphore is set to 0 and all processes waiting on semval to become 0 are awakened.
     - The semid for which the calling process is awaiting action is removed from the system. When this occurs, errno is set to EIDRM and -1 is returned.
     - The calling process receives a signal that is to be caught. When this occurs, the value of semncnt associated with the specified semaphore is decremented, and the calling process resumes execution in the manner prescribed in signal(3C).

Upon successful completion, the value of sempid for each semaphore specified in the array pointed to by sops is set to the process ID of the calling process.
The `semtimedop()` function behaves as `semop()` except when it must suspend execution of the calling process to complete its operation. If `semtimedop()` must suspend the calling process after the time interval specified in `timeout` expires, or if the timeout expires while the process is suspended, `semtimedop()` returns with an error. If the `timespec` structure pointed to by `timeout` is zero-valued and `semtimedop()` needs to suspend the calling process to complete the requested operation(s), it returns immediately with an error. If `timeout` is the NULL pointer, the behavior of `semtimedop()` is identical to that of `semop()`.

**RETURN VALUES**

Upon successful completion, 0 is returned. Otherwise, −1 is returned and `errno` is set to indicate the error.

**ERRORS**

The `semop()` and `semtimedop()` functions will fail if:

- **E2BIG** The `nsops` argument is greater than the system-imposed maximum.
- **EACCES** Operation permission is denied to the calling process (see `intro(2)`).
- **EAGAIN** The operation would result in suspension of the calling process but (`sem_flg & IPC_NOWAIT`) is true.
- **EFAULT** The `sops` argument points to an illegal address.
- **EFBIG** The value of `sem_num` is less than 0 or greater than or equal to the number of semaphores in the set associated with `semid`.
- **EIDRM** A `semid` was removed from the system.
- **EINVAL** A signal was received.
- **EINVAL** The `semid` argument is not a valid semaphore identifier, or the number of individual semaphores for which the calling process requests a `SEM_UNDO` would exceed the limit.
- **ENOSPC** The limit on the number of individual processes requesting an `SEM_UNDO` would be exceeded.
- **ERANGE** An operation would cause a `semval` or a `semadj` value to overflow the system-imposed limit.

The `semtimedop()` function will fail if:

- **EAGAIN** The timeout expired before the requested operation could be completed.

The `semtimedop()` function will fail if one of the following is detected:

- **EFAULT** The `timeout` argument points to an illegal address.
- **EINVAL** The `timeout` argument specified a `tv_sec` or `tv_nsec` value less than 0, or a `tv_nsec` value greater than or equal to 1000 million.

**SEE ALSO**

ipcs(1), intro(2), exec(2), exit(2), fork(2), semctl(2), semget(2)
setpgid(2)

NAME
setpgid – set process group ID

SYNOPSIS
#include <sys/types.h>
#include <unistd.h>

int setpgid(pid_t pid, pid_t pgid);

DESCRIPTION
The setpgid() function sets the process group ID of the process with ID pid to pgid.

If pgid is equal to pid, the process becomes a process group leader. See intro(2) for more information on session leaders and process group leaders.

If pgid is not equal to pid, the process becomes a member of an existing process group.

If pid is equal to 0, the process ID of the calling process is used. If pgid is equal to 0, the process specified by pid becomes a process group leader.

RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, -1 is returned and errno is set to indicate the error.

ERRORS
The setpgid() function will fail if:

EACCES The pid argument matches the process ID of a child process of the calling process and the child process has successfully executed one of the exec family of functions (see exec(2)).

EINVAL The pgid argument is less than (pid_t) 0 or greater than or equal to PID_MAX, or the calling process has a controlling terminal that does not support job control.

EPERM The process indicated by the pid argument is a session leader.

EPERM The pid argument matches the process ID of a child process of the calling process and the child process is not in the same session as the calling process.

EPERM The pgid argument does not match the process ID of the process indicated by the pid argument, and there is no process with a process group ID that matches pgid in the same session as the calling process.

ESRCH The pid argument does not match the process ID of the calling process or of a child process of the calling process.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

ATTRIBUTE TYPE | ATTRIBUTE VALUE
---|---
MT-Level | Async-Signal-Safe

SEE ALSO
intro(2), exec(2), exit(2), fork(2), getpid(2), getsid(2), attributes(5)
**NAME**
setpgrp – set process group ID

**SYNOPSIS**
```c
#include <sys/types.h>
#include <unistd.h>

pid_t setpgrp(void);
```

**DESCRIPTION**
If the calling process is not already a session leader, the `setpgrp()` function makes it one by setting its process group ID and session ID to the value of its process ID, and releases its controlling terminal. See `intro(2)` for more information on process group IDs and session leaders.

**RETURN VALUES**
The `setpgrp()` function returns the value of the new process group ID.

**SEE ALSO**
`setpgrp(1)`, `intro(2)`, `exec(2)`, `fork(2)`, `getpid(2)`, `getsid(2)`, `kill(2)`, `signal(3C)`
The `setrctl()` and `getrctl()` functions provide interfaces for the modification and retrieval of resource control (rctl) values on active entities on the system, such as processes, tasks, or projects. All resource controls are unsigned 64-bit integers, however, a collection of flags are defined that modify which rctl value is to be set or retrieved.

Resource controls are restricted to three levels: basic controls that can be modified by the owner of the calling process, privileged controls that can be modified only by privileged callers, and system controls that are fixed for the duration of the operating system instance. Setting or retrieving each of these controls is performed by setting the privilege field of the resource control block to `RCTL_BASIC`, `RCTL_PRIVILEGED`, or `RCTL_SYSTEM` with `rctlblk_set_privilege()` (see `rctlblk_set_value(3C)`).

For limits on collective entities such as the task or project, the process ID of the calling process is associated with the resource control value. This ID is available by using `rctlblk_get_recipient_pid()` (see `rctlblk_set_value(3C)`). These values are visible only to that process and privileged processes within the collective.

The `getrctl()` function provides a mechanism for iterating through all of the established values on a resource control. The iteration is primed by calling `getrctl()` with `old_blk` set to `NULL`, a valid resource control block pointer in `new_blk`, and specifying `RCTL_FIRST` in the `flags` argument. Once a resource control block has been obtained, repeated calls to `getrctl()` with `RCTL_NEXT` in the `flags` argument and the obtained control in the `old_blk` argument will return the next resource control block in the sequence. The iteration reports the end of the sequence by failing and setting `errno` to `ENOENT`.

The `getrctl()` function allows the calling process to get the current usage of a controlled resource using `RCTL_USAGE` as the `flags` value. The current value of the resource usage is placed in the value field of the resource control block specified by `new_blk`. This value is obtained with `rctlblk_set_value()` (see `rctlblk_set_value(3C)`). All other members of the returned block are undefined and might be invalid.

The `setrctl()` function allows the creation, modification, or deletion of action-value pairs on a given resource control. When passed `RCTL_INSERT` as the `flag` value, `setrctl()` expects `new_blk` to contain a new action-value pair for insertion into the sequence. For `RCTL_DELETE`, the block indicated by `new_blk` is deleted from the sequence. For `RCTL_REPLACE`, the block matching `old_blk` is deleted and replaced by the block indicated by `new_blk`.
The kernel maintains a history of which resource control values have triggered for a particular entity, retrievable from a resource control block with the \texttt{rctlblk\_get\_firing\_time()} function (see \texttt{rctlblk\_set\_value(3C)}). The insertion or deletion of a resource control value at or below the currently enforced value might cause the currently enforced value to be reset. In the case of insertion, the newly inserted value becomes the actively enforced value. All higher values that have previously triggered will have their firing times zeroed. In the case of deletion of the currently enforced value, the next higher value becomes the actively enforced value.

The various resource control block properties are described on the \texttt{rctlblk\_set\_value(3C)} manual page.

Resource controls are inherited from the predecessor process or task. One of the \texttt{exec(2)} functions can modify the resource controls of a process by resetting their histories, as noted above for insertion or deletion operations.

**RETURN VALUES**

Upon successful completion, the \texttt{setrctl()} and \texttt{getrctl()} functions return 0. Otherwise they return -1 and set \texttt{errno} to indicate the error.

**ERRORS**

The \texttt{setrctl()} and \texttt{getrctl()} functions will fail if:

- **EFAULT**
  The \texttt{controlname}, \texttt{old\_blk}, or \texttt{new\_blk} argument points to an illegal address.

- **EINVAL**
  No rctl with the given name is known to the system.

- **ENOENT**
  No value beyond the given resource control block exists.

- **ESRCH**
  No value matching the given resource control block was found for any of RCTL\_NEXT, RCTL\_DELETE, or RCTL\_REPLACE.

- **ENOTSUPP**
  The resource control requested by RCTL\_USAGE does not support the usage operation.

The \texttt{setrctl()} function will fail if:

- **EACCESS**
  The rctl value specified cannot be changed by the current process.

- **EPERM**
  An attempt to set a system limit was attempted.

**EXAMPLES**

**EXAMPLE 1** Retrieve a rctl value.

Obtain the lowest enforced rctl value on the rctl limiting the number of LWPs in a task.

```c
#include <sys/types.h>
#include <rctl.h>
#include <stdio.h>

uint64_t value;
int cur_signal;
rctlblk_t *rblk;
...
```

---

**System Calls** 265
EXAMPLE 1 Retrieve a rctl value.  (Continued)

if ((rblk = malloc(rctlblk_size())) == NULL) {
    (void) fprintf(stderr, "malloc failed: %s
", strerror(errno);
    exit(1);
}

if (getrctl("task.max-lwps", NULL, rblk, RCTL_FIRST) == -1)
    (void) fprintf(stderr, "failed to get rctl: %s
", strerror(errno));
else
    (void) printf("task.max-lwps = %llu",
    rctlblk_get_value(rblk));

USAGE

Resource control blocks are matched on the value and privilege fields. Resource control operations act on the first matching resource control block. Multiple blocks of equal value and privilege will likely need to be entirely deleted and reinserted, rather than replaced, to have the correct outcome. Resource control blocks are sorted such that all blocks with the same value that lack the RCTL_LOCAL_DENY flag precede those having that flag set.

Only one RCPRIV_BASIC resource control value is permitted per process per control. Insertion of an RCPRIV_BASIC value will cause any existing RCPRIV_BASIC value owned by that process on the control to be deleted.

The resource control facility provides the backend implementation for both setrctl() / getrctl() and setrlimit() / getrlimit(). The facility behaves consistently when either of these interfaces is used exclusively; when using both interfaces, the caller must be aware of the ordering issues above, as well as the limit equivalencies described in the following paragraph.

The hard and soft process limits made available with setrlimit() and getrlimit() are mapped to the resource controls implementation. (New process resource controls will not be made available with the rlimit interface.) Because of the RCTL_INSERT and RCTL_DELETE operations, it is possible that the set of values defined on a resource control has more or fewer than the two values defined for an rlimit. In this case, the soft limit is the lowest priority resource control value with the RCTL_LOCAL_DENY flag set, and the hard limit is the resource control value with the lowest priority equal to or exceeding RCPRIV_PRIVILEGED with the RCTL_LOCAL_DENY flag set. If no identifiable soft limit exists on the resource control and setrlimit() is called, a new resource control value is created. If a resource control does not have the global RCTL_GLOBAL_LOWERABLE property set, its hard limit will not allow lowering by unprivileged callers.

ATTRIBUTES

See attributes(5) for descriptions of the following attributes:
**SETCTRL(2)**

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO**

getrlimit(2), errno(3C), rctlblk_set_value(3C), attributes(5)
The `setregid()` function is used to set the real and effective group IDs of the calling process. If `rgid` is −1, the real group ID is not changed; if `egid` is −1, the effective group ID is not changed. The real and effective group IDs may be set to different values in the same call.

If the effective user ID of the calling process is super-user, the real group ID and the effective group ID can be set to any legal value.

If the effective user ID of the calling process is not super-user, either the real group ID can be set to the saved set-group-ID from `execve(2)`, or the effective group ID can either be set to the saved set-group-ID or the real group ID.

In either case, if the real group ID is being changed (that is, if `rgid` is not −1), or the effective group ID is being changed to a value not equal to the real group ID, the saved set-group-ID is set equal to the new effective group ID.

Upon successful completion, 0 is returned. Otherwise, −1 is returned, `errno` is set to indicate the error, and neither of the group IDs will be changed.

The `setregid()` function will fail if:

**EINVAL** The value of `rgid` or `egid` is less than 0 or greater than `UID_MAX` (defined in `<limits.h>`).

**EPERM** The calling process’s effective UID is not the super-user and a change other than changing the real group ID to the saved set-group-ID or changing the effective group ID to the real group ID or the saved group ID, was specified.

If a set-group-ID process sets its effective group ID to its real group ID, it can still set its effective group ID back to the saved set-group-ID.
NAME | setreuid – set real and effective user IDs

SYNOPSIS | #include <unistd.h>

int setreuid(uid_t ruid, uid_t euid);

DESCRIPTION | The setreuid() function is used to set the real and effective user IDs of the calling process. If ruid is −1, the real user ID is not changed; if euid is −1, the effective user ID is not changed. The real and effective user IDs may be set to different values in the same call.

If the effective user ID of the calling process is super-user, the real user ID and the effective user ID can be set to any legal value.

If the effective user ID of the calling process is not super-user, either the real user ID can be set to the effective user ID, or the effective user ID can either be set to the saved set-user ID from execve() (see exec(2)) or the real user ID.

In either case, if the real user ID is being changed (that is, if ruid is not −1), or the effective user ID is being changed to a value not equal to the real user ID, the saved set-user ID is set equal to the new effective user ID.

RETURN VALUES | Upon successful completion, 0 is returned. Otherwise, −1 is returned, errno is set to indicate the error, and neither of the user IDs will be changed.

ERRORS | The setreuid() function will fail if:

EINVAL | The value of ruid or euid is less than 0 or greater than UID_MAX (defined in <limits.h>).

EPERM | The calling process’s effective user ID is not the super-user and a change other than changing the real user ID to the effective user ID, or changing the effective user ID to the real user ID or the saved set-user ID, was specified.

USAGE | If a set-user-ID process sets its effective user ID to its real user ID, it can still set its effective user ID back to the saved set-user ID.

SEE ALSO | exec(2), getuid(2), setregid(2), setuid(2)
setsid(2)

NAME
setsid – create session and set process group ID

SYNOPSIS
#include <sys/types.h>
#include <unistd.h>

pid_t setsid(void);

DESCRIPTION
The setsid() function creates a new session, if the calling process is not a process
group leader. Upon return the calling process will be the session leader of this new
session, will be the process group leader of a new process group, and will have no
controlling terminal. The process group ID of the calling process will be set equal to
the process ID of the calling process. The calling process will be the only process in the
new process group and the only process in the new session.

RETURN VALUES
Upon successful completion, setsid() returns the value of the process group ID of
the calling process. Otherwise it returns (pid_t)−1 and sets errno to indicate the
error.

ERRORS
The setsid() function will fail if:

EPERM
The calling process is already a process group leader, or the
process group ID of a process other than the calling process
matches the process ID of the calling process.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
getsid(2), setpgid(2), setpgrp(2), attributes(5)

WARNINGS
A call to setsid() by a process that is a process group leader will fail. A process can
become a process group leader by being the last member of a pipeline started by a job
control shell. Thus, a process that expects to be part of a pipeline, and that calls
setsid(), should always first fork; the parent should exit and the child should call
setsid(). This will ensure that the calling process will work reliably when started by
both job control shells and non-job control shells.
NAME    settaskid, gettaskid, getprojid – set or get task or project IDs

SYNOPSIS #include <sys/types.h>
#include <sys/task.h>
#include <unistd.h>

taskid_t settaskid(projid_t project, int flags);
taskid_t gettaskid(void);
projid_t getprojid(void);

DESCRIPTION The settaskid() function makes a request of the system to assign a new task ID to
the calling process, changing the associated project ID to that specified. The calling
process must have superuser privileges to perform this operation. The flags argument
should be either TASK_NORMAL for a regular task, or TASK_FINAL, which disallows
subsequent settaskid() calls by the created task.

The gettaskid() function returns the task ID of the calling process.

The getprojid() function returns the project ID of the calling process.

RETURN VALUES Upon successful completion, these functions return the appropriate task or project ID.
Otherwise, −1 is returned and errno is set to indicate the error.

ERRORS The settaskid() function will fail if:
EACCES The invoking task was created with the TASK_FINAL flag.
EPERM The effective user of the calling process is not superuser.
EINVAL The given project ID is not within the valid project ID range.

ATTRIBUTES See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO setsid(2), project(4), attributes(5)
The `setuid()` function sets the real user ID, effective user ID, and saved user ID of the calling process. The `setgid()` function sets the real group ID, effective group ID, and saved group ID of the calling process. The `setegid()` and `seteuid()` functions set the effective group and user IDs respectively for the calling process. See `intro(2)` for more information on real, effective, and saved user and group IDs.

At login time, the real user ID, effective user ID, and saved user ID of the login process are set to the login ID of the user responsible for the creation of the process. The same is true for the real, effective, and saved group IDs; they are set to the group ID of the user responsible for the creation of the process.

When a process calls one of the `exec` family of functions (see `exec(2)`) to execute a file (program), the user and/or group identifiers associated with the process can change. If the file executed is a set-user-ID file, the effective and saved user IDs of the process are set to the owner of the file executed. If the file executed is a set-group-ID file, the effective and saved group IDs of the process are set to the group of the file executed. If the file executed is not a set-user-ID or set-group-ID file, the effective user ID, saved user ID, effective group ID, and saved group ID are not changed.

If the effective user ID of the process calling `setuid()` is the super-user, the real, effective, and saved user IDs are set to the `uid` argument.

If the effective user ID of the calling process is not the super-user, but `uid` is either the real user ID or the saved user ID of the calling process, the effective user ID is set to `uid`.

If the effective user ID of the process calling `setgid()` is the super-user, the real, effective, and saved group IDs are set to the `gid` argument.

If the effective user ID of the calling process is not the super-user, but `gid` is either the real group ID or the saved group ID of the calling process, the effective group ID is set to `gid`.

Upon successful completion, 0 is returned. Otherwise, -1 is returned and `errno` is set to indicate the error.

The `setuid()` and `setgid()` functions will fail if:

- `EINVAL` The value of `uid` or `gid` is out of range.
EPERM

For setuid() and seteuid() the effective user of the calling process is not super-user, and the uid argument does not match either the real or saved user IDs. For setgid() and setegid() the effective user of the calling process is not the super-user, and the gid argument does not match either the real or saved group IDs.

ATTRIBUTES

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>setuid() and setgid() and Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
intro(2), exec(2), getgroups(2), getuid(2), stat(3HEAD), attributes(5)
shmctl(2)

NAME
shmctl – shared memory control operations

SYNOPSIS
#include <sys/types.h>
#include <sys/ipc.h>
#include <sys/shm.h>

int shmctl(int shmid, int cmd, struct shmid_ds *buf);

DESCRIPTION
The shmctl() function provides a variety of shared memory control operations as
specified by cmd. The permission required for a shared memory control operation is
given as {token}, where token is the type of permission needed. The types of permission
are interpreted as follows:

<table>
<thead>
<tr>
<th>Token</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>00400</td>
<td>READ by user</td>
</tr>
<tr>
<td>00200</td>
<td>WRITE by user</td>
</tr>
<tr>
<td>00040</td>
<td>READ by group</td>
</tr>
<tr>
<td>00020</td>
<td>WRITE by group</td>
</tr>
<tr>
<td>00004</td>
<td>READ by others</td>
</tr>
<tr>
<td>00002</td>
<td>WRITE by others</td>
</tr>
</tbody>
</table>

See the Shared Memory Operation Permissions section of intro(2) for more information.

The following operations require the specified tokens:

<table>
<thead>
<tr>
<th>Operation</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>IPC_STAT</td>
<td>Place the current value of each member of the data structure associated with shmid into the structure pointed to by buf. The contents of this structure are defined in intro(2). {READ}</td>
</tr>
</tbody>
</table>
| IPC_SET   | Set the value of the following members of the data structure associated with shmid to the corresponding value found in the structure pointed to by buf:  

\[
\begin{align*}
\text{shm_perm.uid} \\
\text{shm_perm.gid} \\
\text{shm_perm.mode} & \text{ /* access permission bits only */} 
\end{align*}
\]

This command can be executed only by a process that has an effective user ID equal to that of super-user, or to the value of \text{shm_perm.cuid} or \text{shm_perm.uid} in the data structure associated with shmid.

<table>
<thead>
<tr>
<th>Operation</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>IPC_RMID</td>
<td>Remove the shared memory identifier specified by shmid from the system and destroy the shared memory segment and data structure associated with it. This command can be executed only by a process that has an effective user ID equal to that of super-user, or to the value of \text{shm_perm.cuid} or \text{shm_perm.uid} in the data structure associated with shmid.</td>
</tr>
<tr>
<td>SHM_LOCK</td>
<td>Lock the shared memory segment specified by shmid in memory. This command can be executed only by a process that has an effective user ID equal to super-user.</td>
</tr>
<tr>
<td>SHM_UNLOCK</td>
<td>Unlock the shared memory segment specified by shmid. This command can be executed only by a process that has an effective user ID equal to super-user.</td>
</tr>
</tbody>
</table>
Shared memory segments must be explicitly removed after the last reference to them has been removed.

RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, −1 is returned and errno is set to indicate the error.

ERRORS
The shmctl() function will fail if:

- **EACCESS** The *cmd* argument is equal to IPC_STAT and {READ} operation permission is denied to the calling process.
- **EFAULT** The *buf* argument points to an illegal address.
- **EINVAL** The *shmid* argument is not a valid shared memory identifier; or the *cmd* argument is not a valid command or is IPC_SET and *shm_perm.uid* or *shm_perm.gid* is not valid.
- **ENOMEM** The *cmd* argument is equal to SHM_LOCK and there is not enough memory.
- **EOVERFLOW** The *cmd* argument is equal to IPC_STAT and *uid* or *gid* is too large to be stored in the structure pointed to by *buf*.
- **EPERM** The *cmd* argument is equal to IPC_RMID or IPC_SET and the effective user ID of the calling process is not super-user and it is not equal to the value of *shm_perm.cuid* or *shm_perm.uid* in the data structure associated with *shmid*.
- **EPERM** The *cmd* argument is equal to SHM_LOCK or SHM_UNLOCK and the effective user ID of the calling process is not equal to that of super-user.

SEE ALSO
ipcs(1), intro(2), shmget(2), shmap(2)
# shmget

## NAME

shmget – get shared memory segment identifier

## SYNOPSIS

```c
#include <sys/types.h>
#include <sys/ipc.h>
#include <sys/shm.h>

int shmget(key_t key, size_t size, int shmflag);
```

## DESCRIPTION

The `shmget()` function returns the shared memory identifier associated with `key`.

A shared memory identifier and associated data structure and shared memory segment of at least `size` bytes (see `intro(2)`) are created for `key` if one of the following are true:

- The `key` argument is equal to `IPC_PRIVATE`.
- The `key` argument does not already have a shared memory identifier associated with it, and (`shmflag` & `IPC_CREAT`) is true.

Upon creation, the data structure associated with the new shared memory identifier is initialized as follows:

- The values of `shm_perm.cuid`, `shm_perm.uid`, `shm_perm.cgid`, and `shm_perm.gid` are set equal to the effective user ID and effective group ID, respectively, of the calling process.
- The access permission bits of `shm_perm.mode` are set equal to the access permission bits of `shmflag`. `shm_segsz` is set equal to the value of `size`.
- The values of `shm_lpid`, `shm_nattch`, `shm_atime`, and `shm_dtime` are set equal to 0.
- The `shm_ctime` is set equal to the current time.

Shared memory segments must be explicitly removed after the last reference to them has been removed.

## RETURN VALUES

Upon successful completion, a non-negative integer representing a shared memory identifier is returned. Otherwise, -1 is returned and `errno` is set to indicate the error.

## ERRORS

The `shmget()` function will fail if:

- `EACCES` A shared memory identifier exists for `key` but operation permission (see `intro(2)`) as specified by the low-order 9 bits of `shmflag` would not be granted.
- `EEXIST` A shared memory identifier exists for `key` but both (`shmflag` & `IPC_CREATE`) and (`shmflag` & `IPC_EXCL`) are true.
- `EINVAL` The `size` argument is less than the system-imposed minimum or greater than the system-imposed maximum.
- `EINVAL` A shared memory identifier exists for `key` but the size of the segment associated with it is less than `size` and `size` is not equal to 0.
<table>
<thead>
<tr>
<th>ENOENT</th>
<th>A shared memory identifier does not exist for key and $(shmfg&amp;IPC_CREATE)$ is false.</th>
</tr>
</thead>
<tbody>
<tr>
<td>ENOMEM</td>
<td>A shared memory identifier and associated shared memory segment are to be created but the amount of available memory is not sufficient to fill the request.</td>
</tr>
<tr>
<td>ENOSPC</td>
<td>A shared memory identifier is to be created but the system-imposed limit on the maximum number of allowed shared memory identifiers system-wide would be exceeded.</td>
</tr>
</tbody>
</table>

**SEE ALSO** intro(2), shmctl(2), shmop(2), ftok(3C)
The `shmids()` function copies all active shared memory identifiers from the system into the user-defined buffer specified by `buf`, provided that the number of such identifiers is not greater than the number of integers the buffer can contain, as specified by `nids`. If the size of the buffer is insufficient to contain all of the active shared memory identifiers in the system, none are copied.

Whether or not the size of the buffer is sufficient to contain all of them, the number of active shared memory identifiers in the system is copied into the unsigned integer pointed to by `pnids`.

If `nids` is 0 or less than the number of active shared memory identifiers in the system, `buf` is ignored.

Upon successful completion, `shmids()` returns 0. Otherwise, –1 is returned and `errno` is set to indicate the error.

The `shmids()` function will fail if:

- **EFAULT** The `buf` or `pnids` argument points to an illegal address.

The `shmids()` function returns a snapshot of all the active shared memory identifiers in the system. More may be added and some may be removed before they can be used by the caller.

This is sample C code indicating how to use the `shmids()` function.

```c
void examine_shmids()
{
    int *ids = NULL;
    uint_t nids = 0;
    uint_t n;
    int i;

    for (; ;) {
        if (shmids(ids, nids, &n) != 0) {
            perror("shmids");
            exit(1);
        }
        if (n <= nids) /* we got them all */
            break;
        /* we need a bigger buffer */
        ids = realloc(ids, (nids = n) * sizeof (int));
    }
    for (i = 0; i < n; i++)
```

NAME | shmids – discover all shared memory identifiers

SYNOPSIS | #include <sys/shm.h>

```c
int shmids(int *buf, uint_t nids, uint_t *pnids);
```

DESCRIPTION | The `shmids()` function copies all active shared memory identifiers from the system into the user-defined buffer specified by `buf`, provided that the number of such identifiers is not greater than the number of integers the buffer can contain, as specified by `nids`. If the size of the buffer is insufficient to contain all of the active shared memory identifiers in the system, none are copied.

Whether or not the size of the buffer is sufficient to contain all of them, the number of active shared memory identifiers in the system is copied into the unsigned integer pointed to by `pnids`.

If `nids` is 0 or less than the number of active shared memory identifiers in the system, `buf` is ignored.

RETURN VALUES | Upon successful completion, `shmids()` returns 0. Otherwise, –1 is returned and `errno` is set to indicate the error.

ERRORS | The `shmids()` function will fail if:

- **EFAULT** The `buf` or `pnids` argument points to an illegal address.

USAGE | The `shmids()` function returns a snapshot of all the active shared memory identifiers in the system. More may be added and some may be removed before they can be used by the caller.

EXAMPLES | **EXAMPLE 1** `shmids()` example

This is sample C code indicating how to use the `shmids()` function.

```c
void examine_shmids()
{
    int *ids = NULL;
    uint_t nids = 0;
    uint_t n;
    int i;

    for (; ;) {
        if (shmids(ids, nids, &n) != 0) {
            perror("shmids");
            exit(1);
        }
        if (n <= nids) /* we got them all */
            break;
        /* we need a bigger buffer */
        ids = realloc(ids, (nids = n) * sizeof (int));
    }
    for (i = 0; i < n; i++)
```
EXAMPLE 1  

    shmids() example  (Continued)

    process_shmids(ids[i]);
    free(ids);
}

ATTRIBUTES

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO  
ipcrm(1), ipcs(1), intrc(2), shmct1(2), shmmget(2), shmoop(2), attributes(5)
The shmat() function attaches the shared memory segment associated with the
shared memory identifier specified by shmid to the data segment of the calling process.

The permission required for a shared memory control operation is given as `{token}`, where
`token` is the type of permission needed. The types of permission are interpreted
as follows:

00400 READ by user
00200 WRITE by user
00040 READ by group
00020 WRITE by group
00004 READ by others
00002 WRITE by others

See the Shared Memory Operation Permissions section of intro(2) for more information.

When `(shmflg&SHM_SHARE_MMU)` is true, virtual memory resources in addition to
shared memory itself are shared among processes that use the same shared memory.

When `(shmflg&SHM_PAGEABLE)` is true, virtual memory resources are shared and the
dynamic shared memory (DISM) framework is created. The dynamic shared memory
can be resized dynamically within the specified size in shmget(2). The DISM shared
memory is pageable unless it is locked.

The shared memory segment is attached to the data segment of the calling process at
the address specified based on one of the following criteria:

- If `shmaddr` is equal to `(void *) 0`, the segment is attached to the first available
  address as selected by the system.
- If `shmaddr` is equal to `(void *) 0` and `(shmflg&SHM_SHARE_MMU)` or
  `(shmflg&SHM_PAGEABLE)` is true, then the segment is attached to the first available
  suitably aligned address. When `(shmflg&SHM_SHARE_MMU)` or
  `(shmflg&SHM_PAGEABLE)` is set, however, the permission given by `shmget()`
determines whether the segment is attached for reading or reading and writing.
- If `shmaddr` is not equal to `(void *) 0` and `(shmflg&SHM_RND)` is true, the segment
  is attached to the address given by `(shmaddr - (shmaddr modulus SHMLBA))`.
- If `shmaddr` is not equal to `(void *) 0` and `(shmflg&SHM_RND)` is false, the segment
  is attached to the address given by `shmaddr`.
- The segment is attached for reading if `(shmflg&SHM_RDONLY)` is true [READ],
otherwise it is attached for reading and writing [READ/WRITE].
The `shmdt()` function detaches from the calling process’s data segment the shared memory segment located at the address specified by `shmaddr`. If the application is standard-conforming (see `standards(5)`), the `shmaddr` argument is of type `const void *`. Otherwise it is of type `char *`.

Shared memory segments must be explicitly removed after the last reference to them has been removed.

**RETURN VALUES**

Upon successful completion, `shmat()` returns the data segment start address of the attached shared memory segment; `shmdt()` returns 0. Otherwise, −1 is returned, the shared memory segment is not attached, and `errno` is set to indicate the error.

**ERRORS**

The `shmat()` function will fail if:

- **EACCES** Operation permission is denied to the calling process (see `intro(2)`).
- **EINVAL** The `shmid` argument is not a valid shared memory identifier.
- **EINVAL** The `shmaddr` argument is not equal to 0, and the value of `(shmaddr - (shmaddr modulus SHMLBA))` is an illegal address.
- **EINVAL** The `shmaddr` argument is not equal to 0, is an illegal address, and (`shmflag&SHM_RND`) is false.
- **EINVAL** The `shmaddr` argument is not equal to 0, is not properly aligned, and (`shmflag&SHM_SHARE_MMU`) is true.
- **EINVAL** `SHM_SHARE_MMU` is not supported in certain architectures.
- **EMFILE** The number of shared memory segments attached to the calling process would exceed the system-imposed limit.
- **ENOMEM** The available data space is not large enough to accommodate the shared memory segment.

The `shmdt()` function will fail if:

- **EINVAL** The `shmaddr` argument is not the data segment start address of a shared memory segment.

**SEE ALSO**

`intro(2), exec(2), exit(2), fork(2), shmat(2), shmget(2), standards(5)`
### NAME

sigaction – detailed signal management

### SYNOPSIS

```c
#include <signal.h>

int sigaction(int sig, const struct sigaction *act, struct sigaction *oact);
```

### DESCRIPTION

The `sigaction()` function allows the calling process to examine or specify the action to be taken on delivery of a specific signal. See `signal(3HEAD)` for an explanation of general signal concepts.

The `sig` argument specifies the signal and can be assigned any of the signals specified in `signal(3HEAD)` except SIGKILL and SIGSTOP. In a multithreaded process, `sig` cannot be SIGWAITING, SIGCANCEL, or SIGLWP.

If the argument `act` is not NULL, it points to a structure specifying the new action to be taken when delivering `sig`. If the argument `oact` is not NULL, it points to a structure where the action previously associated with `sig` is to be stored on return from `sigaction()`.

The `sigaction` structure includes the following members:

- `void (*sa_handler)();`
- `void (*sa_sigaction)(int, siginfo_t *, void *);`
- `sigset_t sa_mask;`
- `int sa_flags;`

The storage occupied by `sa_handler` and `sa_sigaction` may overlap, and a standard-conforming application (see `standards(5)`) must not use both simultaneously.

The `sa_handler` member identifies the action to be associated with the specified signal, if the SA_SIGINFO flag (see below) is cleared in the `sa_flags` field of the `sigaction` structure. It may take any of the values specified in `signal(3HEAD)` or that of a user specified signal handler. If the SA_SIGINFO flag is set in the `sa_flags` field, the `sa_sigaction` field specifies a signal-catching function.

The `sa_mask` member specifies a set of signals to be blocked while the signal handler is active. On entry to the signal handler, that set of signals is added to the set of signals already being blocked when the signal is delivered. In addition, the signal that caused the handler to be executed will also be blocked, unless the SA_NODEFER flag has been specified. SIGSTOP and SIGKILL cannot be blocked (the system silently enforces this restriction).

The `sa_flags` member specifies a set of flags used to modify the delivery of the signal. It is formed by a logical OR of any of the following values:

- **SA_ONSTACK**: If set and the signal is caught, and if the thread that is chosen to processes a delivered signal has an alternate signal stack declared with `sigaltstack(2)`, then it will process the signal on that stack. Otherwise, the signal is delivered on the thread’s normal stack.
If set and the signal is caught, the disposition of the signal is reset to `SIG_DFL` and the signal will not be blocked on entry to the signal handler (SIGILL, SIGTRAP, and SIGPWR cannot be automatically reset when delivered; the system silently enforces this restriction).

If set and the signal is caught, the signal will not be automatically blocked by the kernel while it is being caught.

If set and the signal is caught, functions that are interrupted by the execution of this signal’s handler are transparently restarted by the system, namely `fcntl(2)`, `ioctl(2)`, `wait(2)`, `waitid(2)`, and the following functions on slow devices like terminals: `getmsg()` and `getpmsg()` (see `getmsg(2)`); `putmsg()` and `putpmsg()` (see `putmsg(2)`); `pread()`, `read()`, and `readv()` (see `read(2)`); `pwrite()`, `write()`, and `writev()` (see `write(2)`); `recv()`, `recvfrom()`, and `recvmsg()` (see `recv(3SOCKET)`); and `send()`, `sendto()`, and `sendmsg()` (see `send(3SOCKET)`). Otherwise, the function returns an `EINTR` error.

If cleared and the signal is caught, `sig` is passed as the only argument to the signal-catching function. If set and the signal is caught, two additional arguments are passed to the signal-catching function. If the second argument is not equal to `NULL`, it points to a `siginfo_t` structure containing the reason why the signal was generated (see `siginfo(3HEAD)`); the third argument points to a `ucontext_t` structure containing the receiving process’s context when the signal was delivered (see `ucontext(3HEAD)`).

If set and `sig` equals `SIGCHLD`, the system will not create zombie processes when children of the calling process exit. If the calling process subsequently issues a `wait(2)`, it blocks until all of the calling process’s child processes terminate, and then returns −1 with `errno` set to `ECHILD`.

If set and `sig` equals `SIGCHLD`, `SIGCHLD` will not be sent to the calling process when its child processes stop or continue.

Upon successful completion, 0 is returned. Otherwise, −1 is returned, `errno` is set to indicate the error, and no new signal handler is installed.

The `sigaction()` function will fail if:

- **EINVAL**: The value of the `sig` argument is not a valid signal number or is equal to `SIGKILL` or `SIGSTOP`. In addition, if in a multithreaded process, it is equal to `SIGWAITING`, `SIGCANCEL`, or `SIGLWP`.

**RETURN VALUES**

**ERRORS**
sigaction(2)

ATTRIBUTES

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Standard</td>
</tr>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO

kill(1), intro(2), exit(2), fcntl(2), getmsg(2), ioctl(2), kill(2), pause(2),
putmsg(2), read(2), sigaltstack(2), sigprocmask(2), sigsend(2),
sigsuspend(2), wait(2), waitid(2), write(2), recv(3SOCKET), send(3SOCKET),
siginfo(3HEAD), signal(3C), signal(3HEAD), sigsetops(3C),
thr_create(3THR), ucontext(3HEAD), attributes(5), standards(5)

NOTES

The handler routine can be declared:

```c
void handler (int sig, siginfo_t *sip, ucontext_t *ucp);
```

The `sig` argument is the signal number. The `sip` argument is a pointer (to space on the stack) to a `siginfo_t` structure, which provides additional detail about the delivery of the signal. The `ucp` argument is a pointer (again to space on the stack) to a `ucontext_t` structure (defined in `<sys/ucontext.h>`) which contains the context from before the signal. It is not recommended that `ucp` be used by the handler to restore the context from before the signal delivery.
**NAME**
sigaltstack – set or get signal alternate stack context

**SYNOPSIS**
#include <signal.h>

int sigaltstack(const stack_t *ss, stack_t *oss);

**DESCRIPTION**
The `sigaltstack()` function allows a thread to define and examine the state of an alternate stack area on which signals are processed. If `ss` is non-zero, it specifies a pointer to and the size of a stack area on which to deliver signals, and informs the system whether the thread is currently executing on that stack. When a signal’s action indicates its handler should execute on the alternate signal stack (specified with a `sigaction(2)` call), the system checks whether the thread chosen to execute the signal handler is currently executing on that stack. If the thread is not currently executing on the signal stack, the system arranges a switch to the alternate signal stack for the duration of the signal handler’s execution.

The `stack_t` structure includes the following members:

```c
int *ss_sp
long ss_size
int ss_flags
```

If `ss` is not `NULL`, it points to a structure specifying the alternate signal stack that will take effect upon successful return from `sigaltstack()`. The `ss_sp` and `ss_size` members specify the new base and size of the stack, which is automatically adjusted for direction of growth and alignment. The `ss_flags` member specifies the new stack state and may be set to the following:

- **SS_DISABLE** The stack is to be disabled and `ss_sp` and `ss_size` are ignored. If `SS_DISABLE` is not set, the stack will be enabled.

If `oss` is not `NULL`, it points to a structure specifying the alternate signal stack that was in effect prior to the call to `sigaltstack()`. The `ss_sp` and `ss_size` members specify the base and size of that stack. The `ss_flags` member specifies the stack’s state, and may contain the following values:

- **SS_ONSTACK** The thread is currently executing on the alternate signal stack. Attempts to modify the alternate signal stack while the thread is executing on it will fail.
- **SS_DISABLE** The alternate signal stack is currently disabled.

**RETURN VALUES**
Upon successful completion, `0` is return. Otherwise, `-1` is returned and `errno` is set to indicate the error.

**ERRORS**
The `sigaltstack()` function will fail if:

- **EFAULT** The `ss` or `oss` argument points to an illegal address.
- **EINVAL** The `ss` argument is not a null pointer, and the `ss_flags` member pointed to by `ss` contains flags other than `SS_DISABLE`.
- **ENOMEM** The size of the alternate stack area is less than `MINSIGSTKSZ`.

System Calls 285
EPERM
An attempt was made to modify an active stack.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
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</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
getcontext(2), mmap(2), sigaction(2), ucontext(3HEAD)

NOTES
The value SIGSTKSZ is defined to be the number of bytes that would be used to cover
the usual case when allocating an alternate stack area. The value MINSIGSTKSZ is
defined to be the minimum stack size for a signal handler. In computing an alternate
stack size, a program should add that amount to its stack requirements to allow for the
operating system overhead.

The following code fragment is typically used to allocate an alternate stack with an
adjacent red zone (an unmapped page) to guard against stack overflow, as with
default stacks:

```c
#include <signal.h>
#include <sys/mman.h>

stack_t sigstk;
sigstk.ss_sp = mmap(NULL, SIGSTKSZ, PROT_READ | PROT_WRITE,
                     MAP_PRIVATE | MAP_ANON, -1, 0);
if (sigstk.ss_sp == MAP_FAILED)
    /* error return */;
sigstk.ss_size = SIGSTKSZ;
sigstk.ss_flags = 0;
if (sigaltstack(&sigstk, NULL) < 0)
    perror("sigaltstack");
```
sigpending – examine signals that are blocked and pending

#include <signal.h>

int sigpending(sigset_t *set);

The sigpending() function retrieves those signals that have been sent to the calling process but are being blocked from delivery by the calling process’s signal mask. The signals are stored in the space pointed to by the set argument.

Upon successful completion, 0 is returned. Otherwise, −1 is returned and errno is set to indicate the error.

The sigpending() function will fail if:

EFAULT The set argument points to an illegal address.

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

sigaction(2), sigprocmask(2), sigsetops(3C), attributes(5)
The `sigprocmask()` function is used to examine and/or change the caller’s signal mask. If the value is `SIG_BLOCK`, the set pointed to by the `set` argument is added to the current signal mask. If the value is `SIG_UNBLOCK`, the set pointed by the `set` argument is removed from the current signal mask. If the value is `SIG_SETMASK`, the current signal mask is replaced by the set pointed to by the `set` argument. If the `oset` argument is not `NULL`, the previous mask is stored in the space pointed to by `oset`. If the value of the `set` argument is `NULL`, the value `how` is not significant and the caller’s signal mask is unchanged; thus, the call can be used to inquire about currently blocked signals.

If there are any pending unblocked signals after the call to `sigprocmask()`, at least one of those signals will be delivered before the call to `sigprocmask()` returns.

It is not possible to block those signals that cannot be ignored this restriction is silently imposed by the system. See `sigaction(2)`.

If `sigprocmask()` fails, the caller’s signal mask is not changed.

Upon successful completion, 0 is returned. Otherwise, -1 is returned and `errno` is set to indicate the error.

The `sigprocmask()` function will fail if:

- `EFAULT` The `set` or `oset` argument points to an illegal address.
- `EINVAL` The value of the `how` argument is not equal to one of the defined values.

See `attributes(5)` for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

In a multithreaded program, the call to `sigprocmask()` impacts only the calling thread’s signal mask and is therefore identical to a call to `thr_sigsetmask(3THR)`.

Signals that are generated synchronously should not be masked. If such a signal is blocked and delivered, the receiving process is killed.
NAME
sigsend, sigsendset – send a signal to a process or a group of processes

SYNOPSIS

#include <signal.h>

int sigsend(idtype_t idtype, id_t id, int sig);
int sigsendset(procset_t *psp, int sig);

DESCRIPTION

The sigsend() function sends a signal to the process or group of processes specified by id and idtype. The signal to be sent is specified by sig and is either 0 or one of the values listed in signal(3HEAD). If sig is 0 (the null signal), error checking is performed but no signal is actually sent. This value can be used to check the validity of id and idtype.

The real or effective user ID of the sending process must match the real or saved user ID of the receiving process, unless the effective user ID of the sending process is super-user, or sig is SIGCONT and the sending process has the same session ID as the receiving process.

If idtype is P_PID, sig is sent to the process with process ID id.

If idtype is P_PGID, sig is sent to all processes with process group ID id.

If idtype is P_SID, sig is sent to all processes with session ID id.

If idtype is P_TASKID, sig is sent to all processes with task ID id.

If idtype is P_UID, sig is sent to any process with effective user ID id.

If idtype is P_GID, sig is sent to any process with effective group ID id.

If idtype is P_PROJID, sig is sent to any process with project ID id.

If idtype is P_CID, sig is sent to any process with scheduler class ID id (see priocntl(2)).

If idtype is P_ALL, sig is sent to all processes and id is ignored.

If id is P_MYID, the value of id is taken from the calling process.

The process with a process ID of 0 is always excluded. The process with a process ID of 1 is excluded unless idtype is equal to P_PID.

The sigsendset() function provides an alternate interface for sending signals to sets of processes. This function sends signals to the set of processes specified by psp. psp is a pointer to a structure of type procset_t, defined in <sys/procset.h>, which includes the following members:

idop_t p_op;
idtype_t p_ididtype;
id_t p_id;
idtype_t p_idtype;
id_t p_rid;
The `p_lidtype` and `p_lid` members specify the ID type and ID of one ("left") set of processes; the `p_ridtype` and `p_rid` members specify the ID type and ID of a second ("right") set of processes. ID types and IDs are specified just as for the `idtype` and `id` arguments to `sigsend()`. The `p_op` member specifies the operation to be performed on the two sets of processes to get the set of processes the function is to apply to. The valid values for `p_op` and the processes they specify are:

- **POp_Diff**: Set difference: processes in left set and not in right set.
- **POp_And**: Set intersection: processes in both left and right sets.
- **POp_Or**: Set union: processes in either left or right set or both.
- **POp_Xor**: Set exclusive-or: processes in left or right set but not in both.

**Return Values**: Upon successful completion, 0 is return. Otherwise, −1 is returned and `errno` is set to indicate the error.

**Errors**: The `sigsend()` and `sigsendset()` functions will fail if:

- **EINVAL**: The `sig` argument is not a valid signal number, or the `idtype` argument is not a valid idtype field.
- **EINVAL**: The `sig` argument is `SIGKILL`, `idtype` is `P_PID` and `id` is 1 (proc1).
- **EPERM**: The effective user of the calling process is not superuser and its real or effective user ID does not match the real or effective user ID of the receiving process, and the calling process is not sending `SIGCONT` to a process that shares the same session.
- **ESRCH**: No process can be found corresponding to that specified by `id` and `idtype`.

The `sigsendset()` function will fail if:

- **EFAULT**: The `psp` argument points to an illegal address.

**See Also**: `kill()`, `getpid()`, `kill()`, `priocntl()`, `signal()`, `signal(3C)`, `signal(3HEAD)`
sighandle – install a signal mask and suspend caller until signal

#include <signal.h>

int sigsuspend(const sigset_t *set);

The `sigsuspend()` function replaces the caller’s signal mask with the set of signals pointed to by the `set` argument and suspends the caller until delivery of a signal whose action is either to execute a signal catching function or to terminate the process.

If the action is to terminate the process, `sigsuspend()` does not return. If the action is to execute a signal catching function, `sigsuspend()` returns after the signal catching function returns. On return, the signal mask is restored to the set that existed before the call to `sigsuspend()`.

It is not possible to block those signals that cannot be ignored (see `signal(3HEAD)`); this restriction is silently imposed by the system.

Since `sigsuspend()` suspends the caller’s execution indefinitely, there is no successful completion return value. On failure, it returns `-1` and sets `errno` to indicate the error.

The `sigsuspend()` function will fail if:

- **EFAULT** The `set` argument points to an illegal address.
- **EINTR** A signal was caught by the caller and control was returned from the signal catching function.

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO** `sigaction(2), sigprocmask(2), sigwait(2), signal(3C), signal(3HEAD), sigsetops(3C), attributes(5)`

**NOTES**

If the caller specifies more than one unblocked signal in the mask to `sigsuspend()`, more than one signal might be processed before the call to `sigsuspend()` returns.

While the caller is executing the signal handler that interrupted its call to `sigsuspend()`, its signal mask is the one passed to `sigsuspend()`, modified as usual by the signal mask specification in the signal’s `sigaction(2)` parameters. The caller’s signal mask is not restored to its previous value until the caller returns from all the signal handlers that interrupted `sigsuspend()`.
The **sigwait()** function selects a signal in `set` that is pending on the calling thread (see `thr_create(3THR)` and `pthread_create(3THR)`). If no signal in `set` is pending, then `sigwait()` blocks until a signal in `set` becomes pending. The selected signal is cleared from the set of signals pending on the calling thread and the number of the signal is returned, or in the POSIX version (see `standards(5)`) placed in `sig`.

The selection of a signal in `set` is independent of the signal mask of the calling thread. This means a thread can synchronously wait for signals that are being blocked by the signal mask of the calling thread. To ensure that only the caller receives the signals defined in `set`, all threads should have signals in `set` masked including the calling thread.

If `sigwait()` is called on an ignored signal, then the occurrence of the signal will be ignored, unless `sigaction()` changes the disposition. If more than one thread waits for the same signal, only one is unblocked when the signal arrives.

Upon successful completion, the default version of `sigwait()` returns a signal number; the POSIX version returns 0 and stores the received signal number at the location pointed to by `sig`. Otherwise, −1 is returned and `errno` is set to indicate an error.

The `sigwait()` function will fail if:

- **EFAULT** The `set` argument points to an invalid address.
- **EINTR** The wait was interrupted by an unblocked, caught signal.
- **EINVAL** The `set` argument contains an unsupported signal number.

The following sample C code creates a thread to handle the receipt of a signal. More specifically, it catches the asynchronously generated signal, **SIGINT**.

```c
#include <pthread.h> #include <signal.h>

/* compile with -D_POSIX_PTHREAD_SEMANTICS switch; */
/* required by sigwait() */
/* sigint thread handles delivery of signal. uses sigwait() to wait */
/* for SIGINT signal. */
#include <pthread.h>
```
EXAMPLE 1 Creating a thread to handle receipt of a signal  (Continued)

#include <stdlib.h>
#include <stdio.h>
#include <string.h>
#include <unistd.h>
#include <signal.h>
#include <synch.h>

static void *threadTwo(void *);
static void *threadThree(void *);
static void *sigint(void *);

sigset_t signalSet;

void *
main(void)
{
    pthread_t t;
    pthread_t t2;
    pthread_t t3;

    sigfillset(&signalSet);
    /*
     * Block signals in initial thread. New threads will
     * inherit this signal mask.
     */
    pthread_sigmask(SIG_BLOCK, &signalSet, NULL);

    printf("Creating threads\n");

    pthread_create(&t, NULL, sigint, NULL);
    pthread_create(&t2, NULL, threadTwo, NULL);
    pthread_create(&t3, NULL, threadThree, NULL);

    printf("##################\n");
    printf("press CTRL-C to deliver SIGINT to sigint thread\n");
    printf("##################\n");

    pthread_exit((void *)0);
}

static void *
threadTwo(void *arg)
{
    printf("hello world, from threadTwo [tid: %d]\n", 
           pthread_self());
    printf("threadTwo [tid: %d] is now complete and exiting\n", 
           pthread_self());
    pthread_exit((void *)0);
}

static void *
threadThree(void *arg)
{
    printf("hello world, from threadThree [tid: %d]\n", 
           pthread_self());
    printf("threadThree [tid: %d] is now complete and exiting\n", 
           pthread_self());
    pthread_exit((void *)0);
}

System Calls  293
EXAMPLE 1 Creating a thread to handle receipt of a signal (Continued)

```c
    printf("threadThree [tid: %d} is now complete and exiting\n", pthread_self( ));
    pthread_exit((void *)0);
}

void *
    sigint(void *arg)
{
    int sig;
    int err;

    printf("thread sigint [tid: %d] awaiting SIGINT\n", pthread_self( ));

    /*
    /* use POSIX sigwait( ) -- 2 args: signal set, signum
    */
    err = sigwait ( &signalSet, &sig );

    /* test for SIGINT; could catch other signals */
    if (err || sig != SIGINT)
        abort( );

    printf("\nSIGINT signal %d caught by sigint thread [tid: %d]\n", sig, pthread_self( ));
    pthread_exit((void *)0);
}
```

ATTRIBUTES

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO

sigaction(2), sigpending(2), sigprocmask(2), sigsuspend(2),
    pthread_create(3THR), pthread_sigmask(3THR), signal(3HEAD),
    thr_create(3THR), thr_sigsetmask(3THR), standards(5)

NOTES

The sigwait() function cannot be used to wait for signals that cannot be caught (see sigaction(2)). This restriction is silently imposed by the system.

Solaris 2.4 and earlier releases provided a sigwait() facility as specified in POSIX.1c Draft 6. The final POSIX.1c standard changed the interface as described above. Support for the Draft 6 interface is provided for compatibility only and may not be supported in future releases. New applications and libraries should use the POSIX standard interface.
__sparc_utrap_install(2)

NAME
__sparc_utrap_install – install a SPARC V9 user trap handler

SYNOPSIS
#include <sys/utrap.h>

int __sparc_utrap_install(utrap_entry_t type, utrap_handler_t
    new_precise, utrap_handler_t new_deferred, utrap_handler_t
    *old_precise, utrap_handler_t *old_deferred);

DESCRIPTION
The __sparc_utrap_install() function establishes new_precise and new_deferred user trap handlers as the new values for the specified type and returns the existing user trap handler values in *old_precise and *old_deferred in a single atomic operation. A new handler address of NULL means no user handler of that type will be installed. A new handler address of UTH_NOCHANGE means that the user handler for that type should not be changed. An old handler pointer of NULL means that the user is not interested in the old handler address.

A precise trap is caused by a specific instruction and occurs before any program-visible state has been changed by this instruction. When a precise trap occurs, the program counter (PC) saved in the Trap Program Counter (TPC) register points to the instruction that induced the trap; all instructions prior to this trapping instruction have been executed. The next program counter (nPC) saved in the Trap Next Program Counter (TnPC) register points to the next instruction following the trapping instruction, which has not yet been executed. A deferred trap is also caused by a particular instruction, but unlike a precise trap, a deferred trap may occur after the program-visible state has been changed. See the SPARC Architecture Manual, Version 9 for further information on precise and deferred traps.

The list that follows contains hardware traps and their corresponding user trap types. User trap types marked with a plus-sign (+) are required and must be provided by all ABI-conforming implementations. The others may not be present on every implementation; an attempt to install a user trap handler for those conditions will return EINVAL. User trap types marked with an asterisk (*) are implemented as precise traps only.

<table>
<thead>
<tr>
<th>Trap Name</th>
<th>User Trap Type (utrap_entry_t)</th>
</tr>
</thead>
<tbody>
<tr>
<td>illegal_instruction</td>
<td>UT_ILLTRAP_INSTRUCTION +* or UT_ILLEGAL_INSTRUCTION</td>
</tr>
<tr>
<td>fp_disabled</td>
<td>UT_FP_DISABLED +*</td>
</tr>
<tr>
<td>fp_exception_ieee_754</td>
<td>UT_FP_EXCEPTION_IEEE_754 +</td>
</tr>
<tr>
<td>fp_exception_other</td>
<td>UT_FP_EXCEPTION_OTHER</td>
</tr>
<tr>
<td>tag_overflow</td>
<td>UT_TAG_OVERFLOW +*</td>
</tr>
<tr>
<td>division_by_zero</td>
<td>UT_DIVISION_BY_ZERO +</td>
</tr>
<tr>
<td>mem_address_not_aligned</td>
<td>UT_MEM_ADDRESS_NOT_ALIGNED +</td>
</tr>
</tbody>
</table>
The following explanations are provided for those user trap types that are not self-explanatory.

**UT_ILLTRAP_INSTRUCTION**
This trap is raised by user execution of the ILLTRAP INSTRUCTION. It is always precise.

**UT_ILLEGAL_INSTRUCTION**
This trap will be raised by the execution of otherwise undefined opcodes. It is implementation-dependent as to what opcodes raise this trap; the ABI only specifies the interface. The trap may be precise or deferred.

**UT_PRIVILEGED_OPCODE**
All opcodes declared to be privileged in SPARC V9 will raise this trap. It is implementation-dependent whether other opcodes will raise it as well; the ABI only specifies the interface.

**UT_DATA_EXCEPTION, UT_INSTRUCTION_EXCEPTION**
No valid user mapping can be made to this address, for a data or instruction access, respectively.

**UT_DATA_PROTECTION, UT_INSTRUCTION_PROTECTION**
A valid mapping exists, and user privilege to it exists, but the type of access (read, write, or execute) is denied, for a data or instruction access, respectively.

**UT_DATA_ERROR, UT_INSTRUCTION_ERROR**
A valid mapping exists, and both user privilege and the type of access are allowed, but an unrecoverable error occurred in attempting the access, for a data or instruction access, respectively. %11 will contain either BUS_ADDRERR or BUS_OBJERR.
This trap is raised when an application issues a floating point instruction (including load or store) and the SPARC V9 Floating Point Registers State (FPRS) FEF bit is 0. If a user handler is installed for this trap, it will be given control. Otherwise the system will set FEF to one and retry the instruction.

For all traps, the handler executes in a new register window, where the **in** registers are the **out** registers of the previous frame and have the value they contained at the time of the trap, similar to a normal subroutine call after the `save` instruction. The **global** registers (including the special registers `%ccr, %asi, and %y`) and the **floating-point** registers have their values from the time of the trap. The stack pointer register `%sp` plus the BIAS will point to a properly-aligned 128-byte register save area; if the handler needs scratch space, it should decrement the stack pointer to obtain it. If the handler needs access to the previous frame’s **in** registers or **local** registers, it should execute a `FLUSHW` instruction, and then access them off of the frame pointer. If the handler calls an ABI-conforming function, it must set the `%asi` register to `ASI_PRIMARY_NOFAULT` before the call.

On entry to a precise user trap handler `%l6` contains the `%pc` and `%l7` contains the `%npc` at the time of the trap. To return from a handler and reexecute the trapped instruction, the handler would execute:

```assembly
jmp %l6, %g0 ! Trapped PC supplied to user trap handler
return %l7 ! Trapped nPC supplied to user trap handler
```

To return from a handler and skip the trapped instruction, the handler would execute:

```assembly
jmp %l7, %g0 ! Trapped nPC supplied to user trap handler
return %l7+4 ! Trapped nPC + 4
```

On entry to a deferred trap handler `%o0` contains the address of the instruction that caused the trap and `%o1` contains the actual instruction (right-justified, zero-extended), if the information is available. Otherwise `%o0` contains the value −1 and `%o1` is undefined. Additional information may be made available for certain cases of deferred traps, as indicated in the following table.

<table>
<thead>
<tr>
<th>Instructions</th>
<th>Additional Information</th>
</tr>
</thead>
<tbody>
<tr>
<td>LD-type (LDSTUB)</td>
<td><code>%o2</code> contains the effective address (`rs1 + rs2</td>
</tr>
<tr>
<td>ST-type (CAS, SWAP)</td>
<td><code>%o2</code> contains the effective address (`rs1 + rs2</td>
</tr>
<tr>
<td>Integer arithmetic</td>
<td><code>%o2</code> contains the <code>rs1</code> value. <code>%o3</code> contains the `rs2</td>
</tr>
<tr>
<td>Floating-point arithmetic</td>
<td><code>%o2</code> contains the address of <code>rs1</code> value. <code>%o3</code> contains the address of <code>rs2</code> value.</td>
</tr>
<tr>
<td>Control-transfer</td>
<td><code>%o2</code> contains the target address (`rs1 + rs2</td>
</tr>
</tbody>
</table>
Asynchronous data errors 

%o2 contains the address that caused the error. %o3 contains the effective ASI, if available, else −1.

To return from a deferred trap, the trap handler issues:

```
ta 68 !ST_RETURN_FROM_DEFERRED_TRAP
```

The following pseudo-code explains how the operating system dispatches traps:

```c
if (precise trap) {
    if (precise_handler) {
        invoke(precise_handler);
        /* not reached */
    } else {
        convert_to_signal(precise_trap);
    }
} else if (deferred_trap) {
    invoke(deferred_handler);
    /* not reached */
} else {
    convert_to_signal(deferred_trap);
}
```

User trap handlers must preserve all registers except the locals (%l0-7) and the outs (%o0-7), that is, %i0-7, %g1-7, %d0-d6, %asi, %fsr, %fprs, %ccr, and %y, except to the extent that modifying the registers is part of the desired functionality of the handler. For example, the handler for UT_FP_DISABLED may load floating-point registers.

**RETURN VALUES**

Upon successful completion, 0 is returned. Otherwise, a non-zero value is returned and errno is set to indicate the error.

**ERRORS**

The `__sparc_utrap_install()` function will fail if:

**EINVAL**

The type argument is not a supported user trap type; the new user trap handler address is not word aligned; the old user trap handler address cannot be returned; or the user program is not a 64-bit executable.

**EXAMPLES**

**EXAMPLE 1** A sample program using the `__sparc_utrap_install()` function.

The `__sparc_utrap_install()` function is normally used by user programs that wish to provide their own tailored exception handlers as a faster alternative to signal(3C), or to handle exceptions that are not directly supported by the signal() interface, such as fp_disabled.

```c
extern void *fpdis_trap_handler();
utrap_handler_t new_precise = (utrap_handler_t)fpdis_trap_handler;
double d;
```
EXAMPLE 1  A sample program using the __sparc_utrap_install() function.
(Continued)

```c
int err;
err = __sparc_utrap_install(UT_FP_DISABLED, new_precise,
          UTH_NOCHANGE, NULL, NULL);
if (err == EINVAL) {
    /* unexpected error, do something */
    exit (1);
}
d = 1.0e-300;
ENTRY(fpdis_trap_handler)
wr  %g0, FPRS_FEF, %fprs
jmpl  %l6, %g0
return  %l7
SET_SIZE(fpdis_trap_handler)
```

This example turns on bit 2, FEF, in the Floating-Point Registers State (FPRS) Register, after a floating-point instruction causes an fp_disabled trap. (Note that this example simulates part of the default system behavior; programs do not need such a handler. The example is for illustrative purposes only.)

ATTRIBUTES  See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>MT-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO  signal(3C), attributes(5)

SPARC Architecture Manual, Version 9

Manufacturer’s processor chip user manuals

NOTES  The Exceptions and Interrupt Descriptions section of the SPARC V9 manual documents which hardware traps are mandatory or optional, and whether they can be implemented as precise or deferred traps, or both. The manufacturer’s processor chip user manuals describe the details of the traps supported for the specific processor implementation.
stat(2)

NAME
stat, lstat, fstat, fstatat – get file status

SYNOPSIS
#include <sys/types.h>
#include <sys/stat.h>

int stat(const char *path, struct stat *buf);
int lstat(const char *path, struct stat *buf);
int fstat(int fildes, struct stat *buf);
int fstatat(int fildes, const char *path, struct stat *buf, int flag);

DESCRIPTION
The stat() function obtains information about the file pointed to by path. Read,
write, or execute permission of the named file is not required, but all directories listed
in the path name leading to the file must be searchable.

The lstat() function obtains file attributes similar to stat(), except when the
named file is a symbolic link; in that case lstat() returns information about the link,
while stat() returns information about the file the link references.

The fstat() function obtains information about an open file known by the file
descriptor fildes, obtained from a successful open(2), creat(2), dup(2), fcntl(2), or
pipe(2) function.

The fstatat() function obtains file attributes similar to the stat(), lstat(), and
fstat() functions. If the path argument is a relative path, it is resolved relative to the
fildes argument rather than the current working directory. If path is absolute, the fildes
argument is unused. If the fildes argument has the special value AT_FDCWD, defined in
<fcntl.h>, relative paths are resolved from the current working directory. If the flag
argument is AT_SYMLINK_NOFOLLOW, defined in <fcntl.h>, the function behaves
like lstat() and does not automatically follow symbolic links. See fsattr(5).

The buf argument is a pointer to a stat structure into which information is placed
concerning the file. A stat structure includes the following members:

mode_t st_mode; /* File mode (see mknod(2)) */
ino_t st_ino; /* Inode number */
dev_t st_dev; /* ID of device containing */
/* a directory entry for this file */
dev_t st_rdev; /* ID of device */
/* This entry is defined only for */
/* char special or block special files */
nlink_t st_nlink; /* Number of links */
uid_t st_uid; /* User ID of the file’s owner */
gid_t st_gid; /* Group ID of the file’s group */
off_t st_size; /* File size in bytes */
time_t st_atime; /* Time of last access */
time_t st_mtime; /* Time of last data modification */
time_t st_ctime; /* Time of last file status change */
/* Times measured in seconds since */
/* 00:00:00 UTC, Jan. 1, 1970 */
long st_blksize; /* Preferred I/O block size */
blkcnt_t st_blocks; /* Number of 512 byte blocks allocated*/
Descriptions of structure members are as follows:

<table>
<thead>
<tr>
<th>Member</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>st_mode</td>
<td>The mode of the file as described in mknod(2). In addition to the modes described in mknod(2), the mode of a file can also be S_IFLNK if the file is a symbolic link. S_IFLNK can be returned either by lstat() or by fstat() when the AT_SYMLINK_NOFOLLOW flag is set.</td>
</tr>
<tr>
<td>st_ino</td>
<td>This field uniquely identifies the file in a given file system. The pair st_ino and st_dev uniquely identifies regular files.</td>
</tr>
<tr>
<td>st_dev</td>
<td>This field uniquely identifies the file system that contains the file. Its value may be used as input to the ustat() function to determine more information about this file system. No other meaning is associated with this value.</td>
</tr>
<tr>
<td>st_rdev</td>
<td>This field should be used only by administrative commands. It is valid only for block special or character special files and only has meaning on the system where the file was configured.</td>
</tr>
<tr>
<td>st_nlink</td>
<td>This field should be used only by administrative commands.</td>
</tr>
<tr>
<td>st_uid</td>
<td>The user ID of the file’s owner.</td>
</tr>
<tr>
<td>st_gid</td>
<td>The group ID of the file’s group.</td>
</tr>
<tr>
<td>st_size</td>
<td>For regular files, this is the address of the end of the file. For block special or character special, this is not defined. See also pipe(2).</td>
</tr>
<tr>
<td>st_atime</td>
<td>Time when file data was last accessed. Changed by the following functions: creat(), mknod(), pipe(), utime(2), and read(2).</td>
</tr>
<tr>
<td>st_mtime</td>
<td>Time when data was last modified. Changed by the following functions: creat(), mknod(), pipe(), utime(), and write(2).</td>
</tr>
<tr>
<td>st_ctime</td>
<td>Time when file status was last changed. Changed by the following functions: chmod(), chown(), creat(), link(2), mknod(), pipe(), unlink(2), utime(), and write().</td>
</tr>
<tr>
<td>st_blksize</td>
<td>A hint as to the &quot;best&quot; unit size for I/O operations. This field is not defined for block special or character special files.</td>
</tr>
<tr>
<td>st_blocks</td>
<td>The total number of physical blocks of size 512 bytes actually allocated on disk. This field is not defined for block special or character special files.</td>
</tr>
</tbody>
</table>

RETURN VALUES

Upon successful completion, 0 is returned. Otherwise, −1 is returned and errno is set to indicate the error.

ERRORS

The stat(), fstat(), lstat(), and fstatat() functions will fail if:

EOVERFLOW The file size in bytes or the number of blocks allocated to the file or the file serial number cannot be represented correctly in the structure pointed to by buf.
The `stat()`, `lstat()`, and `fstatat()` functions will fail if:

- **EACCES**: Search permission is denied for a component of the path prefix.
- **EFAULT**: The `buf` or `path` argument points to an illegal address.
- **EINTR**: A signal was caught during the execution of the `stat()` or `lstat()` function.
- **ELOOP**: Too many symbolic links were encountered in translating `path`.
- **ENAMETOOLONG**: The length of the `path` argument exceeds `PATH_MAX`, or the length of a `path` component exceeds `NAME_MAX` while `_POSIX_NO_TRUNC` is in effect.
- **ENOENT**: The named file does not exist or is the null pathname.
- **ENOLINK**: The `path` argument points to a remote machine and the link to that machine is no longer active.
- **ENOTDIR**: A component of the path prefix is not a directory, or the `fildes` argument does not refer to a valid directory when given a non-null relative path.
- **EOVERFLOW**: A component is too large to store in the structure pointed to by `buf`.

The `fstat()` and `fstatat()` functions will fail if:

- **EBADF**: The `fildes` argument is not a valid open file descriptor. Note that in `fstatat()` the `fildes` argument may also have the valid value of `AT_FDCWD`.
- **EFAULT**: The `buf` argument points to an illegal address.
- **EINTR**: A signal was caught during the execution of the `fstat()` function.
- **ENOLINK**: The `fildes` argument points to a remote machine and the link to that machine is no longer active.
- **EOVERFLOW**: A component is too large to store in the structure pointed to by `buf`.

**USAGE**

The `stat()`, `fstat()`, and `lstat()` functions have transitional interfaces for 64-bit file offsets. See `lfs64(5)`.

**ATTRIBUTES**

See `attributes(5)` for descriptions of the following attributes:
**SEE ALSO**

`chmod(2), chown(2), creat(2), link(2), mknod(2), pipe(2), read(2), time(2), unlink(2), utime(2), write(2), fattach(3C), stat(3HEAD), attributes(5), fsattr(5), lf64(5)`

**NOTES**

If `chmod(2)` is used to change the file group owner permissions on a file with ACL entries, both the file group owner permissions and the ACL mask are changed to the new permissions. The new ACL mask permissions might change the effective permissions for additional users and groups who have ACL entries on the file.
statvfs(2)

NAME     statvfs, fstatvfs – get file system information

SYNOPSIS  
```c
#include <sys/types.h>
#include <sys/statvfs.h>

int statvfs(const char *path, struct statvfs *buf);
int fstatvfs(int fildes, struct statvfs *buf);
```

DESCRIPTION The `statvfs()` function returns a “generic superblock” describing a file system; it can be used to acquire information about mounted file systems. The `buf` argument is a pointer to a structure (described below) that is filled by the function.

The `path` argument should name a file that resides on that file system. The file system type is known to the operating system. Read, write, or execute permission for the named file is not required, but all directories listed in the path name leading to the file must be searchable.

The `statvfs` structure pointed to by `buf` includes the following members:

- `u_long f_bsize;` /* preferred file system block size */
- `u_long f_frsize;` /* fundamental filesystem block
  (size if supported) */
- `fsblkcnt_t f_blocks;` /* total # of blocks on file system
  in units of f_frsize */
- `fsblkcnt_t f_bfree;` /* total # of free blocks */
- `fsblkcnt_t f_bavail;` /* # of free blocks avail to
  non-super-user */
- `fsfilcnt_t f_files;` /* total # of file nodes (inodes) */
- `fsfilcnt_t f_ffree;` /* total # of free file nodes */
- `fsfilcnt_t f_favail;` /* # of inodes avail to
  non-super-user*/
- `u_long f_fsid;` /* file system id (dev for now) */
- `char f_basetype[FSTYPSZ];` /* target fs type name,
  null-terminated */
- `u_long f_flag;` /* bit mask of flags */
- `u_long f_namemax;` /* maximum file name length */
- `char f_fstr[32];` /* file system specific string */
- `u_long f_filler[16];` /* reserved for future expansion */

The `f_basetype` member contains a null-terminated FSType name of the mounted target.

The following values can be returned in the `f_flag` field:

- `ST_RDONLY 0x01 /* read-only file system */`
- `ST_NOSUID 0x02 /* does not support setuid/setgid semantics */`
- `ST_NOTRUNC 0x04 /* does not truncate file names longer than
  NAME_MAX */`

The `fstatvfs()` function is similar to `statvfs()`, except that the file named by `path` in `statvfs()` is instead identified by an open file descriptor `fildes` obtained from a successful `open(2)`, `creat(2)`, `dup(2)`, `fcntl(2)`, or `pipe(2)` function call.

RETURN VALUES Upon successful completion, 0 is returned. Otherwise, –1 is returned and `errno` is set to indicate the error.
The `statvfs()` and `fstatvfs()` functions will fail if:

<table>
<thead>
<tr>
<th>Error</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>EOVERFLOW</td>
<td>One of the values to be returned cannot be represented correctly in the structure pointed to by <code>buf</code>.</td>
</tr>
</tbody>
</table>

The `statvfs()` function will fail if:

<table>
<thead>
<tr>
<th>Error</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>EACCES</td>
<td>Search permission is denied on a component of the path prefix.</td>
</tr>
<tr>
<td>EFAULT</td>
<td>The <code>path</code> or <code>buf</code> argument points to an illegal address.</td>
</tr>
<tr>
<td>EINTR</td>
<td>A signal was caught during the execution of the <code>statvfs()</code> function.</td>
</tr>
<tr>
<td>EIO</td>
<td>An I/O error occurred while reading the file system.</td>
</tr>
<tr>
<td>ELOOP</td>
<td>Too many symbolic links were encountered in translating <code>path</code>.</td>
</tr>
<tr>
<td>ENAMETOOLONG</td>
<td>The length of a <code>path</code> component exceeds <code>NAME_MAX</code> characters, or the length of <code>path</code> exceeds <code>PATH_MAX</code> characters.</td>
</tr>
<tr>
<td>ENOENT</td>
<td>Either a component of the path prefix or the file referred to by <code>path</code> does not exist.</td>
</tr>
<tr>
<td>ENOLINK</td>
<td>The <code>path</code> argument points to a remote machine and the link to that machine is no longer active.</td>
</tr>
<tr>
<td>ENOTDIR</td>
<td>A component of the path prefix of <code>path</code> is not a directory.</td>
</tr>
</tbody>
</table>

The `fstatvfs()` function will fail if:

<table>
<thead>
<tr>
<th>Error</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>EBADF</td>
<td>The <code>fdes</code> argument is not an open file descriptor.</td>
</tr>
<tr>
<td>EFAULT</td>
<td>The <code>buf</code> argument points to an illegal address.</td>
</tr>
<tr>
<td>EINTR</td>
<td>A signal was caught during the execution of the <code>fstatvfs()</code> function.</td>
</tr>
<tr>
<td>EIO</td>
<td>An I/O error occurred while reading the file system.</td>
</tr>
</tbody>
</table>

**USAGE**

The `statvfs()` and `fstatvfs()` functions have transitional interfaces for 64-bit file offsets. See `lfs64(5)`.

**SEE ALSO**

`chmod(2)`, `chown(2)`, `creat(2)`, `dup(2)`, `fcntl(2)`, `link(2)`, `mknod(2)`, `open(2)`, `pipe(2)`, `read(2)`, `time(2)`, `unlink(2)`, `utime(2)`, `write(2)`, `lfs64(5)`

**BUGS**

The values returned for `f_files`, `f_ffree`, and `f_favail` may not be valid for NFS mounted file systems.
### NAME
stime – set system time and date

### SYNOPSIS
```
#include <unistd.h>

int stime(const time_t *tp);
```

### DESCRIPTION
The `stime()` function sets the system’s idea of the time and date. The *tp* argument points to the value of time as measured in seconds from 00:00:00 UTC January 1, 1970.

### RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, −1 is returned and *errno* is set to indicate the error.

### ERRORS
The `stime()` function will fail if:

- **EINVAL** The *tp* argument points to an invalid (negative) time value.
- **EPERM** The effective user of the calling process is not super-user.

### SEE ALSO
- `time(2)`
The `swapctl()` function adds, deletes, or returns information about swap resources. `cmd` specifies one of the following options contained in `<sys/swap.h>`:

- **SC_ADD** /* add a resource for swapping */
- **SC_LIST** /* list the resources for swapping */
- **SC_REMOVE** /* remove a resource for swapping */
- **SC_GETNSWP** /* return number of swap resources */

When **SC_ADD** or **SC_REMOVE** is specified, `arg` is a pointer to a `swapres` structure containing the following members:

- `char *sr_name;` /* pathname of resource */
- `off_t sr_start;` /* offset to start of swap area */
- `off_t sr_length;` /* length of swap area */

The `sr_start` and `sr_length` members are specified in 512-byte blocks. A swap resource can only be removed by specifying the same values for the `sr_start` and `sr_length` members as were specified when it was added. Swap resources need not be removed in the order in which they were added.

When **SC_LIST** is specified, `arg` is a pointer to a `swaptable` structure containing the following members:

- `int swt_n;` /* number of swapents following */
- `struct swapent swt_ent[];` /* array of swt_n swapents */

A `swapent` structure contains the following members:

- `char *ste_path;` /* name of the swap file */
- `off_t ste_start;` /* starting block for swapping */
- `off_t ste_length;` /* length of swap area */
- `long ste_pages;` /* number of pages for swapping */
- `long ste_free;` /* number of ste_pages free */
- `long ste_flags;` /* ST_INDEL bit set if swap file */
  /* is now being deleted */

The **SC_LIST** function causes `swapctl()` to return at most `swt_n` entries. The return value of `swapctl()` is the number actually returned. The ST_INDEL bit is turned on in `ste_flags` if the swap file is in the process of being deleted.

When **SC_GETNSWP** is specified, `swapctl()` returns as its value the number of swap resources in use. `arg` is ignored for this operation.

The **SC_ADD** and **SC_REMOVE** functions will fail if calling process does not have appropriate privileges.
Upon successful completion, the function `swapctl()` returns a value of 0 for `SC_ADD` or `SC_REMOVE`, the number of `struct swapent` entries actually returned for `SC_LIST`, or the number of swap resources in use for `SC_GETNSWP`. Upon failure, the function `swapctl()` returns a value of -1 and sets `errno` to indicate an error.

Under the following conditions, the function `swapctl()` fails and sets `errno` to:

- **EEXIST** Part of the range specified by `sr_start` and `sr_length` is already being used for swapping on the specified resource (`SC_ADD`).
- **EFAULT** Either `arg`, `sr_name`, or `ste_path` points to an illegal address.
- **EINVAL** The specified function value is not valid, the path specified is not a swap resource (`SC_REMOVE`), part of the range specified by `sr_start` and `sr_length` lies outside the resource specified (`SC_ADD`), or the specified swap area is less than one page (`SC_ADD`).
- **EISDIR** The path specified for `SC_ADD` is a directory.
- **ELOOP** Too many symbolic links were encountered in translating the pathname provided to `SC_ADD` or `SC_REMOVE`.
- **ENAMETOOLONG** The length of a component of the path specified for `SC_ADD` or `SC_REMOVE` exceeds `NAME_MAX` characters or the length of the path exceeds `PATH_MAX` characters and `_POSIX_NO_TRUNC` is in effect.
- **ENOENT** The pathname specified for `SC_ADD` or `SC_REMOVE` does not exist.
- **ENOMEM** An insufficient number of `struct swapent` structures were provided to `SC_LIST`, or there were insufficient system storage resources available during an `SC_ADD` or `SC_REMOVE`, or the system would not have enough swap space after an `SC_REMOVE`.
- **ENOSYS** The pathname specified for `SC_ADD` or `SC_REMOVE` is not a file or block special device.
- **ENOTDIR** Pathname provided to `SC_ADD` or `SC_REMOVE` contained a component in the path prefix that was not a directory.
- **EPERM** The effective user of the calling process is not super-user.
- **EROFS** The pathname specified for `SC_ADD` is a read-only file system.
Additionally, the swapctl() function will fail for 32-bit interfaces if:

EOVERFLOW The amount of swap space configured on the machine is too large to be represented by a 32-bit quantity.

EXAMPLE 1 The usage of the SC_GETNSWP and SC_LIST commands.

The following example demonstrates the usage of the SC_GETNSWP and SC_LIST commands.

```c
#include <sys/stat.h>
#include <sys/swap.h>
#include <stdio.h>

#define MAXSTRSIZE 80

int main(argc, argv)
    int argc;
    char *argv[];
{
    swaptbl_t *s;
    int i, n, num;
    char *strtab; /* string table for path names */

    again:
    if ((num = swapctl(SC_GETNSWP, 0)) == -1) {
        perror("swapctl: GETNSWP");
        exit(1);
    }
    if (num == 0) {
        fprintf(stderr, "No Swap Devices Configured\n");
        exit(2);
    }
    /* allocate swaptable for num+1 entries */
    if ((s = (swaptbl_t *)
        malloc(num * sizeof(swapent_t) +
               sizeof(struct swaptable))) ==
        (void *) 0) {
        fprintf(stderr, "Malloc Failed\n");
        exit(3);
    }
    /* allocate num+1 string holders */
    if ((strtab = (char *)
        malloc((num + 1) * MAXSTRSIZE)) == (void *) 0) {
        fprintf(stderr, "Malloc Failed\n");
        exit(3);
    }
    /* initialize string pointers */
    for (i = 0; i < (num + 1); i++) {
        s->swt_ent[i].ste_path = strtab + (i * MAXSTRSIZE);
    }
    s->swt_n = num + 1;
    if ((n = swapctl(SC_LIST, s)) < 0) {
        perror("swapctl");
        exit(1);
    }
}
```

EXAMPLES
EXAMPLE 1 The usage of the SC_GETNSWP and SC_LIST commands. (Continued)

```c
}  
    if (n > num) { /* more were added */
        free(s);
        free(strtab);
        goto again;
    }
    for (i = 0; i < n; i++)
        printf("%s %ld\n",
            s->swt_ent[i].ste_path, s->swt_ent[i].ste_pages);
```
**NAME**
symlink – make a symbolic link to a file

**SYNOPSIS**
```
#include <unistd.h>

int symlink(const char *name1, const char *name2);
```

**DESCRIPTION**
The `symlink()` function creates a symbolic link `name2` to the file `name1`. Either name may be an arbitrary pathname, the files need not be on the same file system, and `name1` may be nonexistent.

The file to which the symbolic link points is used when an `open(2)` operation is performed on the link. A `stat()` operation performed on a symbolic link returns the linked-to file, while an `lstat()` operation returns information about the link itself. See `stat(2)`. Unexpected results may occur when a symbolic link is made to a directory. To avoid confusion in applications, the `readlink(2)` call can be used to read the contents of a symbolic link.

**RETURN VALUES**
Upon successful completion, 0 is returned. Otherwise, −1 is returned, `errno` is set to indicate the error, and the symbolic link is not made.

**ERRORS**
The `symlink()` function will fail if:

- **EACCES**: Search permission is denied for a component of the path prefix of `name2`.
- **EDQUOT**: The directory where the entry for the new symbolic link is being placed cannot be extended because the user's quota of disk blocks on that file system has been exhausted; the new symbolic link cannot be created because the user's quota of disk blocks on that file system has been exhausted; or the user's quota of inodes on the file system where the file is being created has been exhausted.
- **EEXIST**: The file referred to by `name2` already exists.
- **EFAULT**: The `name1` or `name2` argument points to an illegal address.
- **EIO**: An I/O error occurs while reading from or writing to the file system.
- **ELOOP**: Too many symbolic links are encountered in translating `name2`.
- **ENAMETOOLONG**: The length of the `name2` argument exceeds `PATH_MAX`, or the length of a `name2` component exceeds `NAME_MAX` while `_POSIX_NO_TRUNC` is in effect.
- **ENOENT**: A component of the path prefix of `name2` does not exist.
- **ENOSPC**: The directory in which the entry for the new symbolic link is being placed cannot be extended because no space is left on the file system containing the directory; the new symbolic link cannot be created because no space is left on the file system which will contain the link; or there are no free inodes on the file system on which the file is being created.
The file system does not support symbolic links

A component of the path prefix of `name2` is not a directory.

The file `name2` would reside on a read-only file system.

**ATTRIBUTES**

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Standard</td>
</tr>
</tbody>
</table>

**SEE ALSO**

cp(1), link(2), open(2), readlink(2), stat(2), unlink(2), attributes(5)
NAME
sync – update super block

SYNOPSIS
#include <unistd.h>

void sync(void);

DESCRIPTION
The sync() function writes all information in memory that should be on disk, including modified super blocks, modified inodes, and delayed block I/O.

Unlike fsync(3C), which completes the writing before it returns, sync() schedules but does not necessarily complete the writing before returning.

USAGE
The sync() function should be used by applications that examine a file system, such as fsck(1M), and df(1M), and is mandatory before rebooting.

SEE ALSO
df(1M), fsck(1M), fsync(3C)
sysfs(2)

NAME
sysfs – get file system type information

SYNOPSIS
#include <sys/fstyp.h>
#include <sys/fsid.h>

int sysfs(int opcode, const char *fsname);
int sysfs(int opcode, int fs_index, char *buf);
int sysfs(int opcode);

DESCRIPTION
The sysfs() function returns information about the file system types configured in the system. The number of arguments accepted by sysfs() depends on the opcode argument, which can take the following values:

GETFSIND
Translate fsname, a null-terminated file-system type identifier, into a file-system type index.

GETFSTYP
Translate fs_index, a file-system type index, into a null-terminated file-system type identifier and write it into the buffer pointed to by buf, which must be at least of size FSTYPSZ as defined in <sys/fstyp.h>.

GETNFSTYP
Return the total number of file system types configured in the system.

RETURN VALUES
Upon successful completion, the value returned depends upon the opcode argument as follows:

GETFSIND the file-system type index
GETFSTYP 0
GETNFSTYP the number of file system types configured

Otherwise, -1 is returned and errno is set to indicate the error.

ERRORS
The sysfs() function will fail if:

EFAULT The buf or fsname argument points to an illegal address.
EINVAL The fsname argument points to an invalid file-system identifier; the fs_index argument is 0 or invalid; or the opcode argument is invalid.
sysinfo(2)

NAME
    sysinfo – get and set system information strings

SYNOPSIS
    #include <sys/systeminfo.h>

    long sysinfo(int command, char *buf, long count);

DESCRIPTION
    The sysinfo() function copies information relating to the operating system on
    which the process is executing into the buffer pointed to by buf. It can also set certain
    information where appropriate commands are available. The count parameter indicates
    the size of the buffer.

    The POSIX P1003.1 interface (see standards(5)) sysconf(3C) provides a similar
    class of configuration information, but returns an integer rather than a string.

    The values for command are as follows:

    SI_SYSNAME
        Copy into the array pointed to by buf the string that
        would be returned by uname(2) in the sysname field. This is the name of the
        implementation of the operating system, for example, SunOS or UTS.

    SI_HOSTNAME
        Copy into the array pointed to by buf a string that
        names the present host machine. This is the string that
        would be returned by uname(2) in the nodename field. This hostname or
        nodename is often the name the machine is known by locally. The hostname is the name
        of this machine as a node in some network. Different networks may have different names
        for the node, but presenting the nodename to the appropriate network directory or
        name-to-address mapping service should produce a transport end point address. The name
        may not be fully qualified. Internet host names may be up to 256 bytes in length (plus
        the terminating null).

    SI_SET_HOSTNAME
        Copy the null-terminated contents of the array pointed
        to by buf into the string maintained by the kernel
        whose value will be returned by succeeding calls to
        sysinfo() with the command SI_HOSTNAME. This command requires that the effective-user-id
        be super-user.

    SI_RELEASE
        Copy into the array pointed to by buf the string that
        would be returned by uname(2) in the release field. Typical values might be 5.2 or 4.1.

    SI_VERSION
        Copy into the array pointed to by buf the string that
        would be returned by uname(2) in the version field. The syntax and semantics of this string
        are defined by the system provider.
### sysinfo(2)

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>SI_MACHINE</td>
<td>Copy into the array pointed to by <code>buf</code> the string that would be returned by <code>uname(2)</code> in the <code>machine</code> field, for example, <code>sun4u</code>.</td>
</tr>
<tr>
<td>SI_ARCHITECTURE</td>
<td>Copy into the array pointed to by <code>buf</code> a string describing the basic instruction set architecture of the current system, for example, <code>sparc</code>, <code>mc68030</code>, <code>m32100</code>, or <code>i386</code>. These names may not match predefined names in the C language compilation system.</td>
</tr>
<tr>
<td>SI_ISALIST</td>
<td>Copy into the array pointed to by <code>buf</code> the names of the variant instruction set architectures executable on the current system. The names are space-separated and are ordered in the sense of best performance. That is, earlier-named instruction sets may contain more instructions than later-named instruction sets; a program that is compiled for an earlier-named instruction set will most likely run faster on this machine than the same program compiled for a later-named instruction set. Programs compiled for an instruction set that does not appear in the list will most likely experience performance degradation or not run at all on this machine. The instruction set names known to the system are listed in <code>isalist(5)</code>; these names may or may not match predefined names or compiler options in the C language compilation system.</td>
</tr>
<tr>
<td>SI_PLATFORM</td>
<td>Copy into the array pointed to by <code>buf</code> a string describing the specific model of the hardware platform, for example, <code>SUNW,Sun_4_75</code>, <code>SUNW,SPARCsystem-600</code>, or <code>i86pc</code>.</td>
</tr>
<tr>
<td>SI_HW_PROVIDER</td>
<td>Copies the name of the hardware manufacturer into the array pointed to by <code>buf</code>.</td>
</tr>
<tr>
<td>SI_HW_SERIAL</td>
<td>Copy into the array pointed to by <code>buf</code> a string which is the ASCII representation of the hardware-specific serial number of the physical machine on which the function is executed. Note that this may be implemented in Read-Only Memory, using software constants set when building the operating system, or by other means, and may contain non-numeric characters. It is anticipated that manufacturers will not issue the same “serial number” to more than one physical machine. The pair of strings returned by <code>SI_HW_PROVIDER</code> and <code>sysinfo(2)</code> will consist of the names of the machine and the serial number.</td>
</tr>
</tbody>
</table>
SI_HW_SERIAL is likely to be unique across all vendor’s SVR4 implementations.

SI_SRPC_DOMAIN Copies the Secure Remote Procedure Call domain name into the array pointed to by buf.

SI_SET_SRPC_DOMAIN Set the string to be returned by sysinfo() with the SI_SRPC_DOMAIN command to the value contained in the array pointed to by buf. This command requires that the effective-user-id be super-user.

SI_DHCP_CACHE Copy into the array pointed to by buf an ASCII string consisting of the ASCII hexadecimal encoding of the name of the interface configured by boot(1M) followed by the DHCPACK reply from the server. This command is intended for use only by the dhcpagent(1M) DHCP client daemon for the purpose of adopting the DHCP maintenance of the interface configured by boot.

Upon successful completion, the value returned indicates the buffer size in bytes required to hold the complete value and the terminating null character. If this value is no greater than the value passed in count, the entire string was copied. If this value is greater than count, the string copied into buf has been truncated to count −1 bytes plus a terminating null character.

Otherwise, −1 is returned and errno is set to indicate the error.

The sysinfo() function will fail if:

EFAULT The buf argument does not point to a valid address.

EINVAL The data for a SET command exceeds the limits established by the implementation.

EPERM The effective user of the calling process is not super-user.

In many cases there is no corresponding programming interface to set these values; such strings are typically settable only by the system administrator modifying entries in the /etc/system directory or the code provided by the particular OEM reading a serial number or code out of read-only memory, or hard-coded in the version of the operating system.

A good estimation for count is 257, which is likely to cover all strings returned by this interface in typical installations.

boot(1M), dhcpagent(1M), uname(2), gethostid(3C), gethostname(3C), sysconf(3C), isalist(5), standards(5)
time() function returns the value of time in seconds since 00:00:00 UTC, January 1, 1970.

If \( tloc \) is non-zero, the return value is also stored in the location to which \( tloc \) points. If \( tloc \) points to an illegal address, \( \text{time()} \) fails and its actions are undefined.

Upon successful completion, \( \text{time()} \) returns the value of time. Otherwise, \( (\text{time_t})-1 \) is returned and \( \text{errno} \) is set to indicate the error.

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

See also \( \text{stime(2)}, \text{ctime(3C)}, \text{attributes(5)} \)
times(2)

NAME times – get process and child process times

SYNOPSIS #include <sys/times.h>
#include <limits.h>
clock_t times(struct tms *buffer);

DESCRIPTION The times() function fills the tms structure pointed to by buffer with time-accounting information. The tms structure, defined in <sys/times.h>, contains the following members:

clock_t tms_utime;
clock_t tms_stime;
clock_t tms_cutime;
clock_t tms_cstime;

All times are reported in clock ticks. The specific value for a clock tick is defined by the variable CLK_TCK, found in the header <limits.h>.

The times of a terminated child process are included in the tms_cutime and tms_cstime members of the parent when wait(2) or waitpid(2) returns the process ID of this terminated child. If a child process has not waited for its children, their times will not be included in its times.

The tms_utime member is the CPU time used while executing instructions in the user space of the calling process.

The tms_stime member is the CPU time used by the system on behalf of the calling process.

The tms_cutime member is the sum of the tms_utime and the tms_cutime of the child processes.

The tms_cstime member is the sum of the tms_stime and the tms_cstime of the child processes.

RETURN VALUES Upon successful completion, times() returns the elapsed real time, in clock ticks, since an arbitrary point in the past (for example, system start-up time). This point does not change from one invocation of times() within the process to another. The return value may overflow the possible range of type clock_t. If times() fails, (clock_t)-1 is returned and errno is set to indicate the error.

ERRORS The times() function will fail if:

EFAULT The buffer argument points to an illegal address.

ATTRIBUTES See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>
#include <sys/uadmin.h>

int uadmin(int cmd, int fcn, uintptr_t mdep);

The uadmin() function provides control for basic administrative functions. This function is tightly coupled to the system administrative procedures and is not intended for general use. The argument mdep is provided for machine-dependent use and is not defined here.

As specified by cmd, the following commands are available:

A_SHUTDOWN The system is shut down. All user processes are killed, the buffer cache is flushed, and the root file system is unmounted. The action to be taken after the system has been shut down is specified by fcn. The functions are generic; the hardware capabilities vary on specific machines.

  AD_HALT Halt the processor(s).
  AD_POWEROFF Halt the processor(s) and turn off the power.
  AD_BOOT Reboot the system, using the kernel file.
  AD_IBOOT Interactive reboot; user is prompted for bootable program name.

A_REBOOT The system stops immediately without any further processing. The action to be taken next is specified by fcn as above.

A_DUMP The system is forced to panic immediately without any further processing and a crash dump is written to the dump device (see dumpadm(1M)). The action to be taken next is specified by fcn as above.

A_REMOUNT The root file system is mounted again after having been fixed. This should be used only during the startup process.

A_FREEZE Suspend the whole system. The system state is preserved in the state file. The following three subcommands are available.

  AD_COMPRESS Save the system state to the state file with compression of data.
  AD_CHECK Check if your system supports suspend and resume. Without performing a system suspend/resume, this command checks if this feature is currently available on your system.
  AD_FORCE Force AD_COMPRESS even when threads of user applications are not suspendable.

RETURN VALUES Upon successful completion, the value returned depends on cmd as follows:
uadmin(2)

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>A_SHUTDOWN</td>
<td>Never returns.</td>
</tr>
<tr>
<td>A_REBOOT</td>
<td>Never returns.</td>
</tr>
<tr>
<td>A_FREEZE</td>
<td>0 upon resume.</td>
</tr>
<tr>
<td>A_REMOUNT</td>
<td>0.</td>
</tr>
</tbody>
</table>

Otherwise, −1 is returned and errno is set to indicate the error.

**ERRORS**

The uadmin() function will fail if:

- **EPERM**: The effective user of the calling process is not super-user.
- **ENOMEM**: Suspend/resume ran out of physical memory.
- **ENOSPC**: Suspend/resume could not allocate enough space on the root file system to store system information.
- **ENOTSUP**: Suspend/resume not supported on this platform.
- **ENXIO**: Unable to successfully suspend system.
- **EBUSY**: Suspend already in progress.

**SEE ALSO**
dumpadm(1M), kernel(1M), uadmin(1M)
ulimit – get and set process limits

SYNOPSIS
#include <ulimit.h>

long ulimit(int cmd, /* newlimit */...);

DESCRIPTION
The ulimit() function provides for control over process limits. It is effective in limiting the growth of regular files. Pipes are limited to PIPE_MAX bytes.

The cmd values, defined in <ulimit.h>, include:

UL_GETFSIZE Return the soft file size limit of the process. The limit is in units of 512-byte blocks and is inherited by child processes. Files of any size can be read. The return value is the integer part of the soft file size limit divided by 512. If the result cannot be represented as a long int, the result is unspecified.

UL_SETFSIZE Set the hard and soft file size limits for output operations of the process to the value of the second argument, taken as a long int. Any process may decrease its own hard limit, but only a process with appropriate privileges may increase the limit. The new file size limit is returned. The hard and soft file size limits are set to the specified value multiplied by 512. If the result would overflow an rlimit_t, the actual value set is unspecified.

UL_GMEMLIM Get the maximum possible break value (see brk(2)).

UL_GDESLIM Get the current value of the maximum number of open files per process configured in the system.

RETURN VALUES
Upon successful completion, ulimit() returns the value of the requested limit. Otherwise, -1 is returned, the limit is not changed, and errno is set to indicate the error.

ERRORS
The ulimit() function will fail if:

EINVAL The cmd argument is not valid.

EPERM A process not having appropriate privileges attempts to increase its file size limit.

USAGE
Since all return values are permissible in a successful situation, an application wishing to check for error situations should set errno to 0, then call ulimit(), and if it returns -1, check if errno is non-zero.

The getrlimit() and setrlimit() functions provide a more general interface for controlling process limits, and are preferred over ulimit(). See getrlimit(2).

SEE ALSO
brk(2), getrlimit(2), write(2)
NAME
umask – set and get file creation mask

SYNOPSIS
#include <sys/types.h>
#include <sys/stat.h>

mode_t umask(mode_t cmask);

DESCRIPTION
The umask() function sets the process’s file mode creation mask to cmask and returns
the previous value of the mask. Only the access permission bits of cmask and the file
mode creation mask are used. The mask is inherited by child processes. See intro(2)
for more information on masks.

RETURN VALUES
The previous value of the file mode creation mask is returned.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
mkdir(1), sh(1), intro(2), chmod(2), creat(2), mknod(2), open(2), stat(3HEAD),
attributes(5)
umount, umount2 – umount a file system

SYNOPSIS
#include <sys/mount.h>

int umount(const char *file);
int umount2(const char *file, int mflag);

DESCRIPTION
The umount() function requests that a previously mounted file system contained on a block special device or directory be unmounted. The file argument is a pointer to the absolute pathname of the file system to be unmounted. After unmounting the file system, the directory upon which the file system was mounted reverts to its ordinary interpretation.

The umount2() function is identical to umount(), with the additional capability of unmounting file systems even if there are open files active. The mflag argument must contain one of the following values:

0 Perform a normal unmount that is equivalent to umount(). The umount2() function returns EBUSY if there are open files active within the file system to be unmounted.

MS_FORCE Unmount the file system, even if there are open files active. A forced unmount may resort in loss of data, so it should be used only when a regular unmount is unsuccessful. The umount2() function returns ENOTSUP if the specified file systems does not support MS_FORCE. Currently only nfs- and ufs-type file systems support MS_FORCE.

RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, −1 is returned and errno is set to indicate the error.

ERRORS
The umount() and umount2() functions will fail if:

EBUSY A file on file is busy.
EFAULT The file pointed to by file points to an illegal address.
EINVAL The file pointed to by file is not mounted.
ENOENT The file pointed to by file does not exist.
ELOOP Too many symbolic links were encountered in translating the path pointed to by file.
ENAMETOOLONG The length of the file argument exceeds PATH_MAX, or the length of a file component exceeds NAME_MAX while _POSIX_NO_TRUNC is in effect.
ENOLINK The file pointed to by file is on a remote machine and the link to that machine is no longer active.
ENOTBLK The file pointed to by file is not a block special device.
EPERM The process’s effective user ID is not superuser.
The file pointed to by *file* is remote.

The `umount2()` function will fail if:

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>EREMOTE</td>
<td>The file pointed to by <em>file</em> does not support this operation.</td>
</tr>
</tbody>
</table>

**USAGE**

The `umount()` and `umount2()` functions may be invoked only by the superuser. Because it provides greater functionality, the `umount2()` function is preferred.

**SEE ALSO**

`mount(2)`

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NAME  uname – get name of current operating system

SYNOPSIS  
```
#include <sys/utsname.h>

int uname(struct utsname *name);
```

DESCRIPTION  The `uname()` function stores information identifying the current operating system in the structure pointed to by `name`.

The `uname()` function uses the `utsname` structure, defined in `<sys/utsname.h>`, whose members include:

```
char sysname[SYS_NMLN];
char nodename[SYS_NMLN];
char release[SYS_NMLN];
char version[SYS_NMLN];
char machine[SYS_NMLN];
```

The `uname()` function returns a null-terminated character string naming the current operating system in the character array `sysname`. Similarly, the `nodename` member contains the name by which the system is known on a communications network. The `release` and `version` members further identify the operating system. The `machine` member contains a standard name that identifies the hardware on which the operating system is running.

RETURN VALUES  Upon successful completion, a non-negative value is returned. Otherwise, –1 is returned and `errno` is set to indicate the error.

ERRORS  The `uname()` function will fail if:

- `EFAULT`  The `name` argument points to an illegal address.

ATTRIBUTES  See `attributes(5)` for descriptions of the following attributes:

```
<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>
```

SEE ALSO  `uname(1), sysinfo(2), sysconf(3C), attributes(5)`
unlink(2)

NAME
unlink, unlinkat – remove directory entry

SYNOPSIS
#include <unistd.h>

int unlink(const char *path);

int unlinkat(int dirfd, const char *path, int flag);

DESCRIPTION
The unlink() function removes a link to a file. If path names a symbolic link, unlink() removes the symbolic link named by path and does not affect any file or directory named by the contents of the symbolic link. Otherwise, unlink() removes the link named by the pathname pointed to by path and decrements the link count of the file referenced by the link.

The unlinkat() function also removes a link to a file. See fsattr(5). If the flag argument is 0, the behavior of unlinkat() is the same as unlink() except in the processing of its path argument. If path is absolute, unlinkat() behaves the same as unlink() and the dirfd argument is unused. If path is relative and dirfd has the value AT_FDCWD, defined in <fcntl.h>, unlinkat() also behaves the same as unlink(). Otherwise, path is resolved relative to the directory referenced by the dirfd argument.

If the flag argument is set to the value AT_REMOVEDIR, defined in <fcntl.h>, unlinkat() behaves the same as rmdir(2) except in the processing of the path argument as described above.

When the file’s link count becomes 0 and no process has the file open, the space occupied by the file will be freed and the file is no longer accessible. If one or more processes have the file open when the last link is removed, the link is removed before unlink() or unlinkat() returns, but the removal of the file contents is postponed until all references to the file are closed.

The path argument must not name a directory unless the process has appropriate privileges and the implementation supports using unlink() and unlinkat() on directories.

Upon successful completion, unlink() and unlinkat() will mark for update the st_ctime and st_mtime fields of the parent directory. If the file’s link count is not 0, the st_ctime field of the file will be marked for update.

RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, -1 is returned, errno is set to indicate the error, and the file is not unlinked.

ERRORS
The unlink() and unlinkat() functions will fail if:

EACCES Search permission is denied for a component of the path prefix; write permission is denied on the directory containing the link to be removed; the parent directory has the sticky bit set and the file is not writable by the user; or the user does not own the parent directory and the user does not own the file.
EBUSY
The entry to be unlinked is the mount point for a mounted file system.

EFAULT
The path argument points to an illegal address.

EINTR
A signal was caught during the execution of the unlink() function.

ELOOP
Too many symbolic links were encountered in translating path.

ENAMETOOLONG
The length of the path argument exceeds PATH_MAX, or the length of a path component exceeds NAME_MAX while _POSIX_NO_TRUNC is in effect.

ENOENT
The named file does not exist or is a null pathname.

ENOLINK
The path argument points to a remote machine and the link to that machine is no longer active.

ENOTDIR
A component of the path prefix is not a directory or the provided directory descriptor for unlinkat() is not AT_FDCWD or does not reference a directory.

EPERM
The named file is a directory and the effective user of the calling process is not superuser.

EROFS
The directory entry to be unlinked is part of a read-only file system.

The unlink() and unlinkat() functions may fail if:

ENAMETOOLONG
Pathname resolution of a symbolic link produced an intermediate result whose length exceeds PATH_MAX.

ETXTBSY
The entry to be unlinked is the last directory entry to a pure procedure (shared text) file that is being executed.

USAGE
Applications should use rmdir(2) to remove a directory.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
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</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>unlink() is Standard; unlinkat() is Evolving</td>
</tr>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

SEE ALSO
rm(1), close(2), link(2), open(2), rmdir(2), remove(3C), attributes(5), fsattr(5)
# ustat(2)

## NAME
ustat – get file system statistics

## SYNOPSIS
```
#include <sys/types.h>
#include <ustat.h>

int ustat(dev_t dev, struct ustat *buf);
```

## DESCRIPTION
The `ustat()` function returns information about a mounted file system. The `dev` argument is a device number identifying a device containing a mounted file system (see `makedev(3C)`). The `buf` argument is a pointer to a `ustat` structure that includes the following members:

- `daddr_t f_tfree; /* Total free blocks */`
- `ino_t f_tinode; /* Number of free inodes */`
- `char f_fname[6]; /* Filsys name */`
- `char f_fpack[6]; /* Filsys pack name */`

The `f_fname` and `f_fpack` members may not contain significant information on all systems; in this case, these members will contain the null character as the first character.

## RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, −1 is returned and `errno` is set to indicate the error.

## ERRORS
The `ustat()` function will fail if:

- **ECOMM** The `dev` argument is on a remote machine and the link to that machine is no longer active.
- **EFAULT** The `buf` argument points to an illegal address.
- **EINTR** A signal was caught during the execution of the `ustat()` function.
- **EINVAL** The `dev` argument is not the device number of a device containing a mounted file system.
- **ENOLINK** The `dev` argument refers to a device on a remote machine and the link to that machine is no longer active.
- **EOVERFLOW** One of the values returned cannot be represented in the structure pointed to by `buf`.

## USAGE
The `statvfs(2)` function should be used in favor of `ustat()`.

## SEE ALSO
`stat(2), statvfs(2), makedev(3C), lfcntl(5)`

## BUGS
The NFS revision 2 protocol does not permit the number of free files to be provided to the client; therefore, when `ustat()` has completed on an NFS file system, `f_tinode` is always −1.
NAME
utime – set file access and modification times

SYNOPSIS
#include <sys/types.h>
#include <utime.h>

int utime(const char *path, const struct utimbuf *times);

DESCRIPTION
The utime() function sets the access and modification times of the file pointed to by
path, and causes the time of the last file status change (st_ctime) to be updated.

If times is NULL, the access and modification times of the file are set to the current time.
A process must be the owner of the file or have write permission to use utime() in
this manner.

If times is not NULL, times is interpreted as a pointer to a utimbuf structure (defined in
<utime.h>) and the access and modification times are set to the values contained in
the designated structure. Only the owner of the file or the super-user may use
utime() in this manner.

The utimbuf structure contains the following members:

    time_t actime; /* access time */
    time_t modtime; /* modification time */

The times contained in the members of the utimbuf structure are measured in
seconds since 00:00:00 UTC, January 1, 1970.

RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, −1 is returned and errno is set
to indicate the error.

ERRORS
The utime() function will fail if:

EACCESS    Search permission is denied by a component of the path prefix; or
the effective user ID of the process is not superuser and not the
owner of the file, write permission is denied for the file, and times
is NULL.

EFAULT     The path argument points to an illegal address.

EINTR      A signal was caught during the execution of the utime() function.

EIO        An I/O error occurred while reading from or writing to the file
system.

ELOOP      Too many symbolic links were encountered in translating path.

ENAMETOOLONG The length of the path argument exceeds PATH_MAX, or the length
of a path component exceeds NAME_MAX while _POSIX_NO_TRUNC
is in effect.

ENOENT     The named file does not exist or is a null pathname.
ENOLINK: The *path* argument points to a remote machine and the link to that machine is no longer active.

ENOTDIR: A component of the *path* prefix is not a directory.

EPERM: The effective user of the calling process is not super-user and not the owner of the file, and *times* is not NULL.

EROF: The file system containing the file is mounted read-only.

**ATTRIBUTES**

See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Standard</td>
</tr>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO** stat(2), attributes(5)
utimes, futimesat – set file access and modification times

SYNOPSIS
#include <sys/time.h>

int utimes(const char *path, const struct timeval times[2]);
int futimesat(int fildes, const char *path, const struct timeval times[2]);

DESCRIPTION
The utimes() function sets the access and modification times of the file pointed to by the path argument to the value of the times argument. It allows time specifications accurate to the microsecond.

The futimesat() function also sets access and modification times. See fsattr(5). If path is a relative path name, however, futimesat() resolves the path relative to the fildes argument rather than the current working directory. If fildes is set to AT_FDCWD, defined in <fcntl.h>, futimesat() resolves the path relative to the current working directory. If path is a null pointer, futimesat() sets the access and modification times on the file referenced by fildes. The fildes argument is ignored even when futimesat() is provided with an absolute path.

The times argument is an array of timeval structures. The first array member represents the date and time of last access, and the second member represents the date and time of last modification. The times in the timeval structure are measured in seconds and microseconds since the Epoch, although rounding toward the nearest second may occur.

If the times argument is a null pointer, the access and modification times of the file are set to the current time. The effective user ID of the process must be the same as the owner of the file, or must have write access to the file or super-user privileges to use this call in this manner. Upon completion, utimes() will mark the time of the last file status change, st_ctime, for update.

RETURN VALUES
Upon successful completion, 0 is returned. Otherwise, –1 is returned, errno is set to indicate the error, and the file times will not be affected.

ERRORS
The utimes() and futimesat() functions will fail if:

EACCES Search permission is denied by a component of the path prefix; or the times argument is a null pointer and the effective user ID of the process does not match the owner of the file and write access is denied.

EFAULT The path or times argument points to an illegal address. For futimesat(), path might have the value NULL if the fildes argument refers to a valid open file descriptor.

EINVAL A signal was caught during the execution of the utimes() function.
utimes(2)

EINVAL The number of microseconds specified in one or both of the timeval structures pointed to by times was greater than or equal to 1,000,000 or less than 0.

EIO An I/O error occurred while reading from or writing to the file system.

ELOOP Too many symbolic links were encountered in resolving path.

ENAMETOOLONG The length of the path argument exceeds [PATH_MAX] or a pathname component is longer than [NAME_MAX].

ENOLINK The path argument points to a remote machine and the link to that machine is no longer active.

ENOENT A component of path does not name an existing file or path is an empty string.

ENOTDIR A component of the path prefix is not a directory or the path argument is relative and the fildes argument is not AT_FDCWD or does not refer to a valid directory.

EPERM The times argument is not a null pointer and the calling process’s effective user ID has write access to the file but does not match the owner of the file and the calling process does not have the appropriate privileges.

EROFS The file system containing the file is read-only.

The utimes() and futimesat() functions may fail if:

ENAMETOOLONG Path name resolution of a symbolic link produced an intermediate result whose length exceeds [PATH_MAX].

ATTRIBUTES See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>utimes() is Standard; futimesat() is Evolving</td>
</tr>
</tbody>
</table>

SEE ALSO stat(2), attributes(5), fsattr(5)
NAME
vfork – spawn new process in a virtual memory efficient way

SYNOPSIS
#include <unistd.h>

pid_t vfork(void);

DESCRIPTION
The vfork() function creates new processes without fully copying the address space
of the old process. This function is useful in instances where the purpose of a fork(2)
operation would be to create a new system context for an execve() operation (see
exec(2)).

Unlike with the fork() function, the child process borrows the parent’s memory and
thread of control until a call to execve() or an exit (either abnormally or by a call to
_exit() (see exit(2)). The parent process is suspended while the child is using its
resources.

In a multithreaded application, vfork() borrows only the thread of control that
called vfork() in the parent; that is, the child contains only one thread. The use of
vfork() in multithreaded applications, however, is not advised.

The vfork() function can normally be used the same way as fork(). The procedure
that called vfork(), however, should not return while running in the child’s context,
since the eventual return from vfork() would be to a stack frame that no longer
exists. The _exit() function should be used in favor of exit(3C) if unable to
perform an execve() operation, since exit() will flush and close standard I/O
channels, and thereby corrupt the parent process's standard I/O data structures. The
_exit() function should be used even with fork() to avoid flushing the buffered
data twice.

RETURN VALUES
Upon successful completion, vfork() returns 0 to the child process and returns the
process ID of the child process to the parent process. Otherwise, −1 is returned to the
parent process, no child process is created, and errno is set to indicate the error.

ERRORS
The vfork() function will fail if:

EAGAIN The system-imposed limit on the total number of processes under
execution (either system-quality or by a single user) would be
exceeded. This limit is determined when the system is generated.

ENOMEM There is insufficient swap space for the new process.

ATTRIBUTES
See attributes(5) for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Unsafe</td>
</tr>
</tbody>
</table>

SEE ALSO exec(2), exit(2), fork(2), ioctl(2), wait(2), exit(3C), attributes(5)
vfork(2)

NOTES  The use of `vfork()` for any purpose other than as a prelude to an immediate call to a function from the `exec` family or to `_exit()` is not advised.

To avoid a possible deadlock situation, processes that are children in the middle of a `vfork()` are never sent `SIGTTOU` or `SIGTTIN` signals; rather, output or ioctl's are allowed and input attempts result in an EOF indication.

On some systems, the implementation of `vfork()` causes the parent to inherit register values from the child. This can create problems for certain optimizing compilers if `<unistd.h>` is not included in the source calling `vfork()`.
vhangup(2)

NAME
vhangup – virtually “hangup” the current controlling terminal

SYNOPSIS
#include <unistd.h>

void vhangup(void);

DESCRIPTION
The vhangup() function is used by the initialization process init(1M) (among
others) to ensure that users are given “clean” terminals at login by revoking access of
the previous users’ processes to the terminal. To effect this, vhangup() searches the
system tables for references to the controlling terminal of the invoking process and
revokes access permissions on each instance of the terminal that it finds. Further
attempts to access the terminal by the affected processes will yield I/O errors (EBADF
or EIO). A SIGHUP (hangup signal) is sent to the process group of the controlling
terminal.

SEE ALSO
init(1M)

BUGS
Access to the controlling terminal using /dev/tty is still possible.

This call should be replaced by an automatic mechanism that takes place on process
exit.

System Calls  337
The `wait()` function will suspend execution of the calling thread until status information for one of its terminated child processes is available, or until delivery of a signal whose action is either to execute a signal-catching function or to terminate the process. If more than one thread is suspended in `wait()` or `waitpid(2)` awaiting termination of the same process, exactly one thread will return the process status at the time of the target process termination. If status information is available prior to the call to `wait()`, return will be immediate.

If `wait()` returns because the status of a child process is available, it returns the process ID of the child process. If the calling process specified a non-zero value for `stat_loc`, the status of the child process is stored in the location pointed to by `stat_loc`. That status may be evaluated with the macros described on the `wstat(3XFN)` manual page.

In the following, `status` is the object pointed to by `stat_loc`:

- If the child process stopped, the high order 8 bits of `status` will contain the number of the signal that caused the process to stop and the low order 8 bits will be set equal to `WSTOPFLG`.
- If the child process terminated due to an `_exit()` call, the low order 8 bits of `status` will be 0 and the high order 8 bits will contain the low order 8 bits of the argument that the child process passed to `_exit()`; see `exit(2)`.
- If the child process terminated due to a signal, the high order 8 bits of `status` will be 0 and the low order 8 bits will contain the number of the signal that caused the termination. In addition, if `WCOREFLG` is set, a “core image” will have been produced; see `signal(3HEAD)` and `wstat(3XFN)`.

If the calling process has SA_NOCLDWAIT set or has SIGCHLD set to SIG_IGN, and the process has no unwaited children that were transformed into zombie processes, it will block until all of its children terminate, and `wait()` will fail and set `errno` to `ECHILD`.

If a parent process terminates without waiting for its child processes to terminate, the parent process ID of each child process is set to 1, with the initialization process inheriting the child processes; see `intro(2)`.

When `wait()` returns due to a terminated child process, the process ID of the child is returned to the calling process. Otherwise, -1 is returned and `errno` is set to indicate the error.

The `wait()` function will fail if:

- `ECHILD` The calling process has no existing unwaited-for child processes.
wait(2)

**EINTR**  The function was interrupted by a signal.

**USAGE**  Since `wait()` blocks on a stopped child, a calling process wishing to see the return results of such a call should use `waitid(2)` or `waitpid(2)` instead of `wait()`.

**ATTRIBUTES**  See `attributes(5)` for descriptions of the following attributes:

<table>
<thead>
<tr>
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</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO**  `intro(2), exec(2), exit(2), fork(2), pause(2), ptrace(2), waitid(2), waitpid(2), signal(3C), attributes(5), signal(3HEAD), wstat(3XFN)`
waitid(2)

NAME
waitid – wait for child process to change state

SYNOPSIS
#include <wait.h>

int waitid(idtype_t idtype, id_t id, siginfo_t *infop, int options);

DESCRIPTION
The waitid() function suspends the calling process until one of its child processes
changes state. It records the current state of a child in the structure pointed to by infop.
It returns immediately if a child process changed state prior to the call.

The idtype and id arguments specify which children waitid() is to wait for, as
follows:
- If idtype is P_PID, waitid() waits for the child with a process ID equal to
  (pid_t) id.
- If idtype is P_PGID, waitid() waits for any child with a process group ID equal to
  (pid_t) id.
- If idtype is P_ALL, waitid() waits for any child and id is ignored.

The options argument is used to specify which state changes waitid() is to wait for.
It is formed by bitwise OR operation of any of the following flags:
WCONTINUED Return the status for any child that was stopped and has been
  continued.
WEXITED Wait for process(es) to exit.
WNOHANG Return immediately.
WNOWAIT Keep the process in a waitable state.
WSTOPPED Wait for and return the process status of any child that has stopped
  upon receipt of a signal.
WTRAPPED Wait for traced process(es) to become trapped or reach a
  breakpoint (see ptrace(2)).

The infop argument must point to a siginfo_t structure, as defined in
siginfo(3HEAD). If waitid() returns because a child process was found that
satisfies the conditions indicated by the arguments idtype and options, then the
structure pointed to by infop will be filled by the system with the status of the process.
The si_signo member will always be equal to SIGCHLD.

RETURN VALUES
If waitid() returns due to a change of state of one of its children and WNOHANG was
not used, 0 is returned. Otherwise, −1 is returned and errno is set to indicate the
error. If WNOHANG was used, 0 can be returned (indicating no error); however, no
children may have changed state if info->si_pid is 0.

ERRORS
The waitid() function will fail if:
ECHILD The set of processes specified by idtype and id does not contain any
  unwaited processes.
EFAULT
The infop argument points to an illegal address.

EINTR
The waitid() function was interrupted due to the receipt of a signal by the calling process.

EINVAL
An invalid value was specified for options, or idtype and id specify an invalid set of processes.

**USAGE**
With idtype equal to P_ALL and options equal to WEXITED | WTRAPPED, waitid() is equivalent to wait(2).

**SEE ALSO**
intro(2), exec(2), exit(2), fork(2), pause(2), ptrace(2), sigaction(2), wait(2), signal(3C), siginfo(3HEAD)
waitpid – wait for child process to change state

#include <sys/types.h>
#include <sys/wait.h>

pid_t waitpid(pid_t pid, int *stat_loc, int options);

The waitpid() function will suspend execution of the calling thread until status information for one of its terminated child processes is available, or until delivery of a signal whose action is either to execute a signal-catching function or to terminate the process. If more than one thread is suspended in waitpid() or wait(2) awaiting termination of the same process, exactly one thread will return the process status at the time of the target process termination. If status information is available prior to the call to waitpid(), return will be immediate.

The pid argument specifies a set of child processes for which status is requested, as follows:

- If pid is equal to (pid_t)−1, status is requested for any child process.
- If pid is greater than (pid_t)0, it specifies the process ID of the child process for which status is requested.
- If pid is equal to (pid_t)0 status is requested for any child process whose process group ID is equal to that of the calling process.
- If pid is less than (pid_t)−1, status is requested for any child process whose process group ID is equal to the absolute value of pid.

If the calling process has SA_NOCLDWAIT set or has SIGCHLD set to SIG_IGN and the process has no unwaited children that were transformed into zombie processes, it will block until all of its children terminate, and waitpid() will fail and set errno to ECHILD.

If waitpid() returns because the status of a child process is available, then that status may be evaluated with the macros defined by wstat(3XFN). If the calling process had specified a non-zero value of stat_loc, the status of the child process will be stored in the location pointed to by stat_loc.

The options argument is constructed from the bitwise inclusive OR of zero or more of the following flags, defined in the header <sys/wait.h>:

- WCONTINUED: The status of any continued child process specified by pid, whose status has not been reported since it continued, is also reported to the calling process.
- WNOHANG: waitpid() will not suspend execution of the calling process if status is not immediately available for one of the child processes specified by pid.
- WNOWAIT: Keep the process whose status is returned in stat_loc in a waitable state. The process may be waited for again with identical results.
The status of any child processes specified by *pid* that are stopped, and whose status has not yet been reported since they stopped, is also reported to the calling process.

If `waitpid()` returns because the status of a child process is available, it returns a value equal to the process ID of the child process for which status is reported. If `waitpid()` returns due to the delivery of a signal to the calling process, −1 is returned and `errno` is set to EINTR. If `waitpid()` was invoked with `WNOHANG` set in `options`, it has at least one child process specified by `pid` for which status is not available, and status is not available for any process specified by `pid`, then 0 is returned. Otherwise, −1 is returned and `errno` is set to indicate the error.

The `waitpid()` function will fail if:

- **ECHILD**: The process or process group specified by `pid` does not exist or is not a child of the calling process or can never be in the states specified by `options`.
- **EINTR**: The `waitpid()` function was interrupted due to the receipt of a signal sent by the calling process.
- **EINVAL**: An invalid value was specified for `options`.

With `options` equal to 0 and `pid` equal to `(pid_t)−1`, `waitpid()` is identical to `wait(2)`.

See `attributes(5)` for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>MT-Level</td>
<td>Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO** `intro(2), exec(2), exit(2), fork(2), pause(2), ptrace(2), sigaction(2), wait(2), signal(3C), attributes(5), siginfo(3HEAD), wstat(3XFN)`
write(2)

NAME  
write, pwrite, writev – write on a file

SYNOPSIS

```c
#include <unistd.h>

ssize_t write(int fd, const void *buf, size_t nbyte);

ssize_t pwrite(int fd, const void *buf, size_t nbyte, off_t offset);

#include <sys/uio.h>

ssize_t writev(int fd, const struct iovec *iov, int iovcnt);
```

DESCRIPTION

The `write()` function attempts to write `nbyte` bytes from the buffer pointed to by `buf` to the file associated with the open file descriptor, `fd`. If `nbyte` is 0, `write()` will return 0 and have no other results if the file is a regular file; otherwise, the results are unspecified.

On a regular file or other file capable of seeking, the actual writing of data proceeds from the position in the file indicated by the file offset associated with `fd`. Before successful return from `write()`, the file offset is incremented by the number of bytes actually written. On a regular file, if this incremented file offset is greater than the length of the file, the length of the file will be set to this file offset.

If the `O_SYNC` bit has been set, `write` I/O operations on the file descriptor complete as defined by synchronized I/O file integrity completion.

If `fd` refers to a socket, `write()` is equivalent to `send(3SOCKET)` with no flags set.

On a file not capable of seeking, writing always takes place starting at the current position. The value of a file offset associated with such a device is undefined.

If the `O_APPEND` flag of the file status flags is set, the file offset will be set to the end of the file prior to each write and no intervening file modification operation will occur between changing the file offset and the write operation.

For regular files, no data transfer will occur past the offset maximum established in the open file description with `fd`.

A `write()` to a regular file is blocked if mandatory file/record locking is set (see `chmod(2)`), and there is a record lock owned by another process on the segment of the file to be written:

- If `O_NDELAY` or `O_NONBLOCK` is set, `write()` returns −1 and sets `errno` to `EAGAIN`.
- If `O_NDELAY` and `O_NONBLOCK` are clear, `write()` sleeps until all blocking locks are removed or the `write()` is terminated by a signal.

If a `write()` requests that more bytes be written than there is room for—for example, if the write would exceed the process file size limit (see `getrlimit(2)` and `ulimit(2)`), the system file size limit, or the free space on the device—only as many
bytes as there is room for will be written. For example, suppose there is space for 20
bytes more in a file before reaching a limit. A write() of 512-bytes returns 20. The
next write() of a non-zero number of bytes gives a failure return (except as noted for
pipes and FIFO below).

If write() is interrupted by a signal before it writes any data, it will return −1 with
errno set to EINTR.

If write() is interrupted by a signal after it successfully writes some data, it will
return the number of bytes written.

After a write() to a regular file has successfully returned:
- Any successful read(2) from each byte position in the file that was modified by
  that write will return the data specified by the write() for that position until such
  byte positions are again modified.
- Any subsequent successful write() to the same byte position in the file will
  overwrite that file data.

Write requests to a pipe or FIFO are handled the same as a regular file with the
following exceptions:
- There is no file offset associated with a pipe, hence each write request appends to
  the end of the pipe.
- Write requests of \{PIPE_BUF\} bytes or less are guaranteed not to be interleaved
  with data from other processes doing writes on the same pipe. Writes of greater
  than \{PIPE_BUF\} bytes may have data interleaved, on arbitrary boundaries, with
  writes by other processes, whether or not the O_NONBLOCK or O_NDELAY flags are
  set.
- If O_NONBLOCK and O_NDELAY are clear, a write request may cause the process to
  block, but on normal completion it returns nbyte.
- If O_NONBLOCK and O_NDELAY are set, write() does not block the process. If a
  write() request for PIPE_BUF or fewer bytes succeeds completely write() returns
  nbyte. Otherwise, if O_NONBLOCK is set, it returns −1 and sets errno to
  EAGAIN or if O_NDELAY is set, it returns 0. A write() request for greater than
  \{PIPE_BUF\} bytes transfers what it can and returns the number of bytes written
  or it transfers no data and, if O_NONBLOCK is set, returns −1 with errno set to
  EAGAIN or if O_NDELAY is set, it returns 0. Finally, if a request is greater than
  PIPE_BUF bytes and all data previously written to the pipe has been read,
  write() transfers at least PIPE_BUF bytes.

When attempting to write to a file descriptor (other than a pipe, a FIFO, a socket, or a
STREAM) that supports nonblocking writes and cannot accept the data immediately:
- If O_NONBLOCK and O_NDELAY are clear, write() blocks until the data can be
  accepted.
- If O_NONBLOCK or O_NDELAY is set, write() does not block the process. If some
data can be written without blocking the process, write() writes what it can and
returns the number of bytes written. Otherwise, if O_NONBLOCK is set, it returns −1.
and sets \texttt{errno} to \texttt{EAGAIN} or if \texttt{O_NDELAY} is set, it returns 0.

Upon successful completion, where \texttt{nbyte} is greater than 0, \texttt{write()} will mark for update the \texttt{st_ctime} and \texttt{st_mtime} fields of the file, and if the file is a regular file, the \texttt{S_ISUID} and \texttt{S_ISGID} bits of the file mode may be cleared.

For STREAMS files (see \texttt{intro(2) and streamio(7I)}), the operation of \texttt{write()} is determined by the values of the minimum and maximum \texttt{nbyte} range ("packet size") accepted by the STREAM. These values are contained in the topmost STREAM module, and can not be set or tested from user level. If \texttt{nbyte} falls within the packet size range, \texttt{nbyte} bytes are written. If \texttt{nbyte} does not fall within the range and the minimum packet size value is zero, \texttt{write()} breaks the buffer into maximum packet size segments prior to sending the data downstream (the last segment may be smaller than the maximum packet size). If \texttt{nbyte} does not fall within the range and the minimum value is non-zero, \texttt{write()} fails and sets \texttt{errno} to \texttt{ERANGE}. Writing a zero-length buffer (\texttt{nbyte} is zero) to a STREAMS device sends a zero length message with zero returned. However, writing a zero-length buffer to a pipe or FIFO sends no message and zero is returned. The user program may issue the \texttt{I_SWOPT ioctl(2)} to enable zero-length messages to be sent across the pipe or FIFO (see \texttt{streamio(7I)}).

When writing to a STREAM, data messages are created with a priority band of zero. When writing to a socket or to a STREAM that is not a pipe or a FIFO:

- If \texttt{O_NDELAY} and \texttt{O_NONBLOCK} are not set, and the STREAM cannot accept data (the STREAM write queue is full due to internal flow control conditions), \texttt{write()} blocks until data can be accepted.
- If \texttt{O_NDELAY} or \texttt{O_NONBLOCK} is set and the STREAM cannot accept data, \texttt{write()} returns -1 and sets \texttt{errno} to \texttt{EAGAIN}.
- If \texttt{O_NDELAY} or \texttt{O_NONBLOCK} is set and part of the buffer has already been written when a condition occurs in which the STREAM cannot accept additional data, \texttt{write()} terminates and returns the number of bytes written.

The \texttt{write()} and \texttt{writev()} functions will fail if the STREAM head had processed an asynchronous error before the call. In this case, the value of \texttt{errno} does not reflect the result of \texttt{write()} or \texttt{writev()} but reflects the prior error.

\textbf{pwrite()}

The \texttt{pwrite()} function performs the same action as \texttt{write()}, except that it writes into a given position without changing the file pointer. The first three arguments to \texttt{pwrite()} are the same as \texttt{write()} with the addition of a fourth argument \texttt{offset} for the desired position inside the file.

\textbf{writev()}

The \texttt{writev()} function performs the same action as \texttt{write()}, but gathers the output data from the \texttt{iovlen} buffers specified by the members of the \texttt{iov} array: \texttt{iov[0]}, \texttt{iov[1]}, ..., \texttt{iov[iovlen - 1]}. The \texttt{iovlen} buffer is valid if greater than 0 and less than or equal to \texttt{(IOV_MAX)}. See \texttt{intro(2)} for a definition of \texttt{(IOV_MAX)}.

The \texttt{iovec} structure contains the following members:

\begin{verbatim}
  caddr_t  iov_base;
  int      iov_len;
\end{verbatim}
Each `iovec` entry specifies the base address and length of an area in memory from which data should be written. The `writev()` function always writes all data from an area before proceeding to the next.

If `fd` refers to a regular file and all of the `iov_len` members in the array pointed to by `iov` are 0, `writev()` will return 0 and have no other effect. For other file types, the behavior is unspecified.

If the sum of the `iov_len` values is greater than `SSIZE_MAX`, the operation fails and no data is transferred.

RETURN VALUES
Upon successful completion, `write()` returns the number of bytes actually written to the file associated with `fd`. This number is never greater than `nbyte`. Otherwise, −1 is returned, the file-pointer remains unchanged, and `errno` is set to indicate the error.

Upon successful completion, `writev()` returns the number of bytes actually written. Otherwise, it returns −1, the file-pointer remains unchanged, and `errno` is set to indicate an error.

ERRORS
The `write()`, `pwrite()`, and `writev()` functions will fail if:

- **EAGAIN**: Mandatory file/record locking is set, `O_NDELAY` or `O_NONBLOCK` is set, and there is a blocking record lock; an attempt is made to write to a STREAM that can not accept data with the `O_NDELAY` or `O_NONBLOCK` flag set; or a write to a pipe or FIFO of `PIPE_BUF` bytes or less is requested and less than `nbytes` of free space is available.

- **EBADF**: The `fd` argument is not a valid file descriptor open for writing.

- **EDEADLK**: The write was going to go to sleep and cause a deadlock situation to occur.

- **EDQUOT**: The user’s quota of disk blocks on the file system containing the file has been exhausted.

- **EFBIG**: An attempt is made to write a file that exceeds the process’s file size limit or the maximum file size (see `getrlimit(2)` and `ulimit(2)`).

- **EFBIG**: The file is a regular file, `nbyte` is greater than 0, and the starting position is greater than or equal to the offset maximum established in the file description associated with `fd`.

- **EINVAL**: A signal was caught during the write operation and no data was transferred.

- **EIO**: The process is in the background and is attempting to write to its controlling terminal whose `TOSTOP` flag is set, or the process is neither ignoring nor blocking `SIGTTOU` signals and the process group of the process is orphaned.
### write(2)

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>ENOLCK</strong></td>
<td>Enforced record locking was enabled and {LOCK_MAX} regions are already locked in the system, or the system record lock table was full and the write could not go to sleep until the blocking record lock was removed.</td>
</tr>
<tr>
<td><strong>ENOLINK</strong></td>
<td>The <code>fildes</code> argument is on a remote machine and the link to that machine is no longer active.</td>
</tr>
<tr>
<td><strong>ENOSPC</strong></td>
<td>During a write to an ordinary file, there is no free space left on the device.</td>
</tr>
<tr>
<td><strong>ENOSR</strong></td>
<td>An attempt is made to write to a STREAMS with insufficient STREAMS memory resources available in the system.</td>
</tr>
<tr>
<td><strong>ENXIO</strong></td>
<td>A hangup occurred on the STREAM being written to.</td>
</tr>
<tr>
<td><strong>EPIPE</strong></td>
<td>An attempt is made to write to a pipe or a FIFO that is not open for reading by any process, or that has only one end open (or to a file descriptor created by <code>socket(3SOCKET)</code>, using type <code>SOCK_STREAM</code> that is no longer connected to a peer endpoint). A SIGPIPE signal will also be sent to the process. The process dies unless special provisions were taken to catch or ignore the signal.</td>
</tr>
<tr>
<td><strong>ERANGE</strong></td>
<td>The transfer request size was outside the range supported by the STREAMS file associated with <code>fildes</code>.</td>
</tr>
</tbody>
</table>

The `write()` and `pwrite()` functions will fail if:

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>EFAULT</strong></td>
<td>The <code>buf</code> argument points to an illegal address.</td>
</tr>
<tr>
<td><strong>EINVAL</strong></td>
<td>The <code>nbyte</code> argument overflowed an <code>ssize_t</code>.</td>
</tr>
</tbody>
</table>

The `pwrite()` function fails and the file pointer remains unchanged if:

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>ESPIPE</strong></td>
<td>The <code>fildes</code> argument is associated with a pipe or FIFO.</td>
</tr>
</tbody>
</table>

The `write()` and `writev()` functions may fail if:

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>EINVAL</strong></td>
<td>The STREAM or multiplexer referenced by <code>fildes</code> is linked (directly or indirectly) downstream from a multiplexer.</td>
</tr>
<tr>
<td><strong>ENXIO</strong></td>
<td>A request was made of a non-existent device, or the request was outside the capabilities of the device.</td>
</tr>
</tbody>
</table>

A write to a STREAMS file may fail if an error message has been received at the STREAM head. In this case, `errno` is set to the value included in the error message.

The `writev()` function may fail if:

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Description</th>
</tr>
</thead>
</table>
| **EINVAL** | The `iovcnt` argument was less than or equal to 0 or greater than `{IOV_MAX}; one of the `iov_len` values in the `iov` array was negative; or the sum of the `iov_len` values in the `iov` array...
overflowed an `ssize_t`.

**USAGE**
The `pwrite()` function has a transitional interface for 64-bit file offsets. See `lf64(5)`.

**ATTRIBUTES**
See `attributes(5)` for descriptions of the following attributes:

<table>
<thead>
<tr>
<th>ATTRIBUTE TYPE</th>
<th>ATTRIBUTE VALUE</th>
</tr>
</thead>
<tbody>
<tr>
<td>Interface Stability</td>
<td>Standard</td>
</tr>
<tr>
<td>MT-Level</td>
<td><code>write()</code> is Async-Signal-Safe</td>
</tr>
</tbody>
</table>

**SEE ALSO**
`intro(2), chmod(2), creat(2), dup(2), fcntl(2), getrlimit(2), ioctl(2), lseek(2), open(2), pipe(2), ulimit(2), send(3SOCKET), socket(3SOCKET), attributes(5), lf64(5), streamio(7I)`
yield(2)

<table>
<thead>
<tr>
<th>NAME</th>
<th>yield – yield execution to another lightweight process</th>
</tr>
</thead>
<tbody>
<tr>
<td>SYNOPSIS</td>
<td><code>#include &lt;unistd.h&gt;</code></td>
</tr>
<tr>
<td></td>
<td><code>void yield(void);</code></td>
</tr>
<tr>
<td>DESCRIPTION</td>
<td>The <code>yield()</code> function causes the current lightweight process to yield its execution in favor of another lightweight process with the same or greater priority.</td>
</tr>
<tr>
<td>SEE ALSO</td>
<td><code>thr_yield(3THR)</code></td>
</tr>
</tbody>
</table>

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